Another Cone for Integration

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Abstract

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1. Introduction

In [2] we considered the problem of integration and the cone of integrands

$$C_{\tau} := \{ f \in \mathcal{V}^1 : \operatorname{Var}(f') \le \tau \| f' - f(1) + f(0) \|_1 \}, \tag{1}$$

where the total variation and the \mathcal{L}_p norms are defined as

$$\operatorname{Var}(f) := \sup_{\substack{n \in \mathbb{N} \\ 0 = x_0 < x_1 < \dots < x_n = 1}} \sum_{i=1}^n |f(x_i) - f(x_{i-1})|,$$

$$\|f\|_p := \left\{ \left[\int_0^1 |f(x)|^p \, \mathrm{d}x \right]^{1/p}, \quad 1 \le p < \infty, \right.$$

$$\sup_{\substack{0 \le x \le 1 \\ 0 \le x \le 1}} |f(x)|, \qquad p = \infty,$$

We derived an algorithm [2, Algorithm 4] that was guaranteed for integrands in \mathcal{C}_{τ} . In this note we consider another algorithm and other cones.

First we recall some notation and results from [2]. For all $n \in \mathcal{I} := \{0, 2, 3, \ldots\}$ we have the linear spline. By convention $A_0(f) = 0$, and for n > 0,

$$x_{i,n} := x_i := \frac{i-1}{n-1}, \qquad i = 1, \dots, n,$$
 (2a)

$$A_n(f)(x) := (n-1) \left[f(x_i)(x_{i+1} - x) + f(x_{i+1})(x - x_i) \right]$$
for $x_i \le x \le x_{i+1}$. (2b)

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The cost of each function value is one and so the cost of A_n is n. The dependence of the nodes, x_i on n is often suppressed for simplicity. Integrating the linear spline gives us the trapezoidal rule based on n-1 trapezoids:

$$T_n(f) := \int_0^1 A_n(f) \, \mathrm{d}x = \frac{1}{2n-2} [f(x_1) + 2f(x_2) + \dots + 2f(x_{n-1}) + f(x_n)]$$

for $n \geq 2$ and $T_0(f) = 0$.

The error of the trapezoidal rule has the following upper bound [1, (7.15)]:

$$\left| \int_0^1 f(x) \, dx - T_n(f) \right| \le \frac{\operatorname{Var}(f')}{8(n-1)^2} \qquad n \in \mathcal{I} \setminus \{0\}. \tag{3}$$

For any $n \in \mathcal{I}$, let $\mathcal{J}_n = \{m \in \mathbb{N} : (n-1)/(m-1) \in \mathbb{N}\}$. This means that T_n integrates exactly any function that is a linear spline using m nodes for $m \in \mathcal{J}_n$. This implies that

$$\left| \int_{0}^{1} f(x) dx - T_{n}(f) \right| = \left| \int_{0}^{1} [f(x) - A_{m}(f)(x)] dx - T_{n}(f - A_{m}(f)) \right|$$

$$\leq \frac{\operatorname{Var}(f' - A_{m}(f)')}{8(n-1)^{2}} \quad \forall m \in \mathcal{J}_{n}, \ n \in \mathcal{I} \setminus \{0\}.$$
 (4)

The variation of the first derivative of f is bounded below by the variation of the first derivative of the linear spline of f. For all $f \in \mathcal{V}^1$ it follows that

$$\operatorname{Var}(f') \ge F_n(f) := \operatorname{Var}(A_n(f)')$$

$$= \begin{cases} 0, & n = 0, 2, \\ (n-1) \sum_{i=1}^{n-2} |f(x_i) - 2f(x_{i+1}) + f(x_{i+2})|, & n \ge 3. \end{cases}$$
(5)

Also note that

$$F_m(f) = F_m(A_n(f)) \le \operatorname{Var}(A_n(f)') = F_n(f) \qquad \forall m \in \mathcal{J}_n.$$
 (6)

The bound in further implies that

$$F_n(f) \le \operatorname{Var}(f') \le \operatorname{Var}(f' - A_m(f)') + \operatorname{Var}(A_m(f))$$

$$= \operatorname{Var}(f' - A_m(f)') + F_m(f). \quad (7)$$

Another useful fact from ??? is that

$$||f' - f(1) + f(0)||_1 \le \widetilde{F}_n(f) + \widetilde{h}(n)\operatorname{Var}(f') \qquad \forall f \in \mathcal{V}^1$$
(8)

2. New Cone, New Algorithm

Let $\overline{\mathcal{I}}$ be some non-empty subset of $\{(\ell, m, n) \in \{2, 3, \ldots\}^3 : \ell \in \mathcal{J}_m, m \in \mathcal{J}_n\}$, and let $\overline{\tau} : \overline{\mathcal{I}} \to (0, \infty)$ be some given function. The new cone considered

here is defined as

$$\overline{\mathcal{C}}_{\overline{\tau}} := \{ f \in \mathcal{V}^1 : \operatorname{Var}(f' - A_m(f)') \le F_n(f - A_m(f)) + \overline{\tau}(\ell, m, n) \| f' - A_\ell(f)' \|_1 \}$$

$$\forall (\ell, m, n) \in \overline{\mathcal{I}} \}. \quad (9)$$

Note that for all $(\ell, m, n) \in \overline{\mathcal{I}}$

$$||f' - A_{\ell}(f)'||_{1} = ||f' - A_{m}(f)' + A_{m}(f)' - A_{\ell}(f)'||_{1}$$

$$\leq ||f' - A_{m}(f)'||_{1} + \widetilde{F}_{m}(f - A_{\ell}(f))$$

$$\leq \widetilde{F}_{m}(f - A_{m}(f)) + \widetilde{h}(m) \operatorname{Var}(f' - A_{m}(f)') + \widetilde{F}_{m}(f - A_{\ell}(f))$$

$$= \widetilde{h}(m) \operatorname{Var}(f' - A_{m}(f)') + \widetilde{F}_{m}(f - A_{\ell}(f)) \quad \text{by (8)}$$

For all $f \in \overline{\mathcal{C}}_{\overline{\tau}}$ it then follows that

$$\operatorname{Var}(f' - A_m(f)') \leq F_n(f - A_m(f)) + \overline{\tau}(\ell, m, n) \widetilde{F}_{\ell}(f)$$

$$\leq F_n(f - A_m(f)) + \overline{\tau}(\ell, m, n) [\widetilde{h}(m) \operatorname{Var}(f' - A_m(f)') + \widetilde{F}_m(f - A_{\ell}(f))]$$

$$\operatorname{Var}(f' - A_m(f)') \leq \frac{F_n(f - A_m(f)) + \overline{\tau}(\ell, m, n) \widetilde{F}_m(f - A_{\ell}(f))}{1 - \overline{\tau}(\ell, m, n) \widetilde{h}(m)},$$

$$\operatorname{provided} \overline{\tau}(\ell, m, n) \widetilde{h}(m) < 1.$$

Here $\hat{\tau}: \widehat{\mathcal{I}} \to [0, \infty)$ is some specified function that defines the cone, and $\widehat{\mathcal{I}} = \{N_{\min}, N_{\min} + 1, \ldots\}$, where N_{\min} is some integer no smaller than 3.

Algorithm 1 (New Cone Adaptive Univariate Integration). Let the sequence of algorithms $\{T_n\}_{n\in\mathcal{I}}$ $\{F_n\}_{n\in\mathcal{I}}$, and $\widehat{C}_{\hat{\tau}}$ be as described above. Set i=1, and let $n_1=N_{\min}$. For any error tolerance ε and input function f, do the following:

Step 1. Bound Var(f') and check for convergence. Compute $F_{n_i}(f)$ in (5). Check whether n_i is large enough to satisfy the error tolerance, i.e.

$$\hat{\tau}(n_i)F_{n_i}(f) \le 8(n_i - 1)^2 \varepsilon.$$

If this is true, then return $T_{n_i}(f)$ and terminate the algorithm.

Step 2. Increase the number of trapezoids. If the above condition is false, choose $n_{i+1} = 2n_i$, increment i, and go to Step 1.

3. New Cone, New Algorithm

The new cone considered here is defined as

$$\widehat{\mathcal{C}}_{\widehat{\tau}} := \{ f \in \mathcal{V}^1 : \min_{m \in \mathcal{I}_n} \operatorname{Var}(f' - A_m(f)') \le \widehat{\tau}(n) F_n(f) \ \forall n \in \widehat{\mathcal{I}} \}, \tag{10}$$

Here $\hat{\tau}: \widehat{\mathcal{I}} \to [0, \infty)$ is some specified function that defines the cone, and $\widehat{\mathcal{I}} = \{N_{\min}, N_{\min} + 1, \ldots\}$, where N_{\min} is some integer no smaller than 3.

Algorithm 2 (New Cone Adaptive Univariate Integration). Let the sequence of algorithms $\{T_n\}_{n\in\mathcal{I}}$ $\{F_n\}_{n\in\mathcal{I}}$, and $\widehat{C}_{\hat{\tau}}$ be as described above. Set i=1, and let $n_1=N_{\min}$. For any error tolerance ε and input function f, do the following:

Step 1. Bound Var(f') and check for convergence. Compute $F_{n_i}(f)$ in (5). Check whether n_i is large enough to satisfy the error tolerance, i.e.

$$\hat{\tau}(n_i)F_{n_i}(f) \le 8(n_i - 1)^2 \varepsilon.$$

If this is true, then return $T_{n_i}(f)$ and terminate the algorithm.

Step 2. Increase the number of trapezoids. If the above condition is false, choose $n_{i+1} = 2n_i$, increment i, and go to Step 1.

4. The New Cone's Relationship to Other Cones

The cone defined in (10) makes Algorithm 2 work. In this section we show that it contains and is contained in other cones that might be more intuitive. One family of cones of interest is defined by replacing $F_n(f)$ by Var(f') in (10):

$$\widetilde{\mathcal{C}}_{\tilde{\tau}} := \{ f \in \mathcal{V}^1 : \operatorname{Var}(f' - A_n(f)') \le \tilde{\tau}(n) \operatorname{Var}(f'), \ n \in \mathcal{I} \}, \tag{11}$$

where $\tilde{\tau}: \mathcal{I} \to [0,2]$ is non-increasing. Another family of cones is related to (1) and is defined as

$$C_{\overline{\tau}} := \{ f \in \mathcal{V}^1 : \text{Var}(f' - A_n(f)') \le \overline{\tau}(n) \| f' - A_n(f)' \|_1, \ n \in \mathcal{I} \}, \tag{12}$$

where $\overline{\tau}: \mathcal{I} \to [0, \infty]$. Under this definition \mathcal{C}_{τ} corresponds to defining $\overline{\tau}(2) = \tau$, $\overline{\tau}(n) = \infty$ for n > 2.

To facilitate the comparison of $\widehat{C}_{\hat{\tau}}$, $\widetilde{C}_{\tilde{\tau}}$, and $C_{\overline{\tau}}$ we note several inequalities. For all $f \in \mathcal{V}^1$,

$$Var(f') \le Var(f' - A_n(f)') + Var(A_n(f)') = Var(f' - A_n(f)') + F_n(f),$$
 (13)

$$\operatorname{Var}(f' - A_n(f)') \le \operatorname{Var}(f') + \operatorname{Var}(A_n(f')) = \operatorname{Var}(f') + F_n(f). \tag{14}$$

prove the following lemma. From (5) and (14) it follows that

$$\operatorname{Var}(f' - A_n(f)') \le 2 \operatorname{Var}(f') \quad \forall f \in \mathcal{V}^1,$$
 (15)

which is why $\tilde{\tau}(n) \leq 2$ for all n. Moreover, if $\tilde{\tau}(n) = 2$ for all $n \in \mathcal{I}$, then $\widetilde{\mathcal{C}}_{\tilde{\tau}} = \mathcal{V}^1$.

Theorem 1. Given the function $\hat{\tau}: \widehat{\mathcal{I}} \to [0, \infty)$, suppose that $\tilde{\tau}_j: \mathcal{I} \to [0, 2]$, j = 1, 2 satisfy the inequality

$$\tilde{\tau}_1(n) \le \frac{\hat{\tau}(n)}{1 + \hat{\tau}(n)} \le \min(2, \hat{\tau}(n)) \le \tilde{\tau}_2(n) \quad \forall n \in \hat{\mathcal{I}}.$$
 (16)

It follows that $\widetilde{\mathcal{C}}_{\tilde{\tau}_1} \leq \widehat{\mathcal{C}}_{\hat{\tau}} \subseteq \widetilde{\mathcal{C}}_{\tilde{\tau}_2}$.

Proof. First suppose that $f \in \widetilde{\mathcal{C}}_{\tilde{\tau}_1}$ where $\tilde{\tau}_1$ satisfies inequality (16). It then follows that

$$Var(f' - A_n(f)') = (1 + \hat{\tau}(n)) Var(f' - A_n(f)') - \hat{\tau}(n) Var(f' - A_n(f)')$$

$$\leq [1 + \hat{\tau}(n)] \tilde{\tau}_1(n) Var(f') - \hat{\tau}(n) Var(f' - A_n(f)') \quad \text{by (11)}$$

$$\leq \hat{\tau}(n) [Var(f') - Var(f' - A_n(f)')] \quad \text{by (16)}$$

$$\leq \hat{\tau}(n) F_n(f) \quad \text{by (13)}.$$

Thus, $\widetilde{C}_{\tilde{\tau}_1} \leq \widehat{C}_{\hat{\tau}}$. Now suppose that $f \in \widehat{C}_{\hat{\tau}}$. It follows by (5) and (15) that $f \in \widetilde{C}_{\tilde{\tau}_2}$.

To prove the relationship between the cones defined in (10) and (12) the following bound is needed.

Lemma 1. For all $n \in \mathcal{I}$ and all $f \in \mathcal{V}^1$ it follows that

$$F_n(f) \le 2(n-1) \|f' - A_n(f)'\|_1$$
 (17)

Proof. For all $f \in \mathcal{V}^1$ we use the triangle inequality:

$$F_{n}(f) = (n-1) \sum_{i=1}^{n-2} |f(x_{i}) - 2f(x_{i+1}) + f(x_{i+2})|$$

$$\leq (n-1) \sum_{i=1}^{n-2} |f(x_{i}) - f(x_{i+1}) + \frac{f(1) - f(0)}{n-1}|$$

$$+ (n-1) \sum_{i=1}^{n-2} |-f(x_{i+1}) + f(x_{i+2}) - \frac{f(1) - f(0)}{n-1}|$$

$$\leq 2(n-1) \sum_{i=1}^{n-1} |f(x_{i+1}) - f(x_{i}) - \frac{f(1) - f(0)}{n-1}|$$

$$= 2(n-1) ||A_{n}(f)' - f(1) + f(0)||_{1}$$

$$\leq 2(n-1) ||f' - f(1) + f(0)||_{1}.$$

Theorem 2. Given the function $\hat{\tau}: \widehat{\mathcal{I}} \to [0, \infty)$, suppose that $\overline{\tau}_j: \mathcal{I} \to [0, \infty)$, j = 1, 2 satisfy the inequality

$$\min\{2(n-1)\hat{\tau}(n): n \in \widehat{\mathcal{I}}\}??\overline{\tau}_{?}(n)$$
(18)

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It follows that $C_{\overline{\tau}_1} \leq \widehat{C}_{\hat{\tau}} \subseteq C_{\overline{\tau}_2}$.

Proof. For all $f \in \widehat{\mathcal{C}}_{\hat{\tau}}$ it follows from (17) that

$$\operatorname{Var}(f') \le \hat{\tau}(n) F_n(f) = 2(n-1)\hat{\tau}(n) \|f' - f(1) + f(0)\|_1 \quad \forall n \ge N_{\min}$$

Applying the definition of τ completes the proof.

Now we define a cone that is contained in $\widehat{\mathcal{C}}_{\hat{\tau}}$. Let

$$\widetilde{C}_{\tilde{\tau}} := \{ f \in \mathcal{V}^1 : \operatorname{Var}(f') \le \tilde{\tau}(n) \| f' - f(1) + f(0) \|_1 \ \forall n \ge 3 \}, \tag{19}$$

Theorem 3. For any non-increasing $\hat{\tau}: \mathcal{I} \to (1, \infty)$, let

$$\tau = \min\{2(n-1)\hat{\tau}(n) : n \ge N_{\min}\}.$$

It follows that $\widehat{\mathcal{C}}_{\hat{\tau}} \subseteq \mathcal{C}_{\tau}$.

Proof. For all $f \in \widehat{\mathcal{C}}_{\hat{\tau}}$ it follows from (17) that

$$\operatorname{Var}(f') \le \hat{\tau}(n) F_n(f) = 2(n-1)\hat{\tau}(n) \|f' - f(1) + f(0)\|_1 \quad \forall n \ge N_{\min}.$$

Applying the definition of τ completes the proof.

References

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