Fred J. Hickernell

Room E1-208, Department of Applied Mathematics, Illinois Institute of Technology, 10 W. 32nd St., Chicago, IL 60616

Lluís Antoni Jiménez Rugama

Room E1-120, Department of Applied Mathematics, Illinois Institute of Technology, 10 W. 32nd St., Chicago, IL 60616

Abstract

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1. Bases and Node Sets

1.1. Group-Like Structures

Consider the half open d-dimensional unit cube, $\mathcal{X} := [0,1)^d$, on which the functions of interest are to be defined. Define \mathcal{X} to be a field with the additive operation $\oplus : \mathcal{X} \times \mathcal{X} \to \mathcal{X}$, $\mathbf{x} \oplus \mathbf{y} = \mathbf{x} + \mathbf{y} \pmod{1}$. Indeed, (\mathcal{X}, \oplus) is an Abelian group. Here $\mathbf{0}$ is the additive identity. The unique additive inverse of \mathbf{x} is $\ominus \mathbf{x} := \mathbf{1} - \mathbf{x}$, and $\mathbf{x} \ominus \mathbf{t}$ means $\mathbf{x} \oplus (\ominus \mathbf{t})$. Moreover, such a set \mathcal{X} is also a vector space under the field \mathbb{Z} and the multiplicative operation is seen by means of \oplus :

$$ax := \underbrace{x \oplus \cdots \oplus x}_{a \text{ times}} \ \forall a \in \mathbb{N}, \qquad ax := \underbrace{\ominus x \ominus \cdots \ominus x}_{-a \text{ times}} \ \forall a \in \mathbb{Z} \backslash \mathbb{N}_0.$$

The set $\mathbb{K} := \mathbb{Z}^d$ is used to index series expressions for the integrands. This is a field with the natural sum and multiplication. Similarly to \mathcal{X} , it is also a vector space under \mathbb{Z} .

Now, define the bilinear operation $\langle \cdot, \cdot \rangle : \mathbb{K} \times \mathcal{X} \to \mathcal{X}$,

$$\langle \boldsymbol{k}, \boldsymbol{x} \rangle = \boldsymbol{k}^T \boldsymbol{y} \pmod{1}.$$
 (1a)

For all $t, x \in \mathcal{X}$, $k, l \in \mathbb{K}$, and $a \in \mathbb{Z}$, it follows that

$$\langle \boldsymbol{k}, \boldsymbol{0} \rangle = \langle \boldsymbol{0}, \boldsymbol{x} \rangle = 0,$$
 (1b)

$$\langle \boldsymbol{k}, a\boldsymbol{x} \oplus \boldsymbol{t} \rangle = a \langle \boldsymbol{k}, \boldsymbol{x} \rangle + \langle \boldsymbol{k}, \boldsymbol{t} \rangle \pmod{1}$$
 (1c)

$$\langle a\mathbf{k} + \mathbf{l}, \mathbf{x} \rangle = a \langle \mathbf{k}, \mathbf{x} \rangle + \langle \mathbf{l}, \mathbf{x} \rangle \pmod{1},$$
 (1d)

$$\langle \boldsymbol{k}, \boldsymbol{x} \rangle = 0 \ \forall \boldsymbol{k} \in \mathbb{K} \Longrightarrow \boldsymbol{x} = \boldsymbol{0}.$$
 (1e)

1.2. Sequences, Nets, and Dual Nets

Suppose that there exists a sequence of points in \mathcal{X} , denoted $\mathcal{P}_{\infty} = \{t_i\}_{i=0}^{\infty}$. Any $\mathcal{P}_m := \{t_i\}_{i=0}^{b^m-1}$ dotted with \oplus is an Abelian subgroup of \mathcal{P}_{∞} . They are called *nets* and all are nested, i.e. $\{0\} = \mathcal{P}_0 \subseteq \cdots \subseteq \mathcal{P}_m \subseteq \cdots \subseteq \mathcal{P}_{\infty}$. Furthermore, \mathcal{P}_{∞} is assumed to satisfy the following properties:

$$\{t_1, t_b, t_{b^2}, \ldots\}$$
 are linearly independent, (2a)

$$b\mathbf{t}_{b^m} = \mathbf{t}_{b^{m-1}},\tag{2b}$$

$$\boldsymbol{t}_{i} = \sum_{\ell=0}^{\infty} i_{\ell} \boldsymbol{t}_{b^{\ell}}, \quad \text{where } \vec{i} = (i_{0}, i_{1}, i_{2}, \ldots) \in \mathbb{F}_{b}^{\infty},$$
 (2c)

$$\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle = 0 \ \forall i \in \mathbb{N}_0 \Longrightarrow \boldsymbol{k} = \boldsymbol{0}.$$
 (2d)

Note that from (1) together with (2b) it follows,

$$\langle \boldsymbol{k}, \boldsymbol{t}_{b^{m-1}} \rangle = \langle b\boldsymbol{k}, \boldsymbol{t}_{b^m} \rangle$$
 (3)

One example is the extensible rank-1 lattices [1]. For $\mathcal{P}_m := \left\{ \boldsymbol{z} \frac{n}{b^m}, \ n \in \mathbb{F}_{b^m} \right\}$ the nested structure detailed above is well defined and the vector \boldsymbol{z} can be seen in each coordinate as an infinite digit integer. In addition, for every \mathcal{P}_m we can find a generator. If we want $\boldsymbol{t}_{b^{m-1}} = \boldsymbol{z} \frac{j_m}{b^m}$ to the the generator of \mathcal{P}_m , it only suffices to verify that $\gcd(j_m, b^m) = 1$ with $j_m \in \mathbb{F}_{b^m}$. This vector $\boldsymbol{j} = (j_0, j_1, \ldots)$ will define the choice of generators for all subgroups of the sequence \mathcal{P}_{∞} . In order to satisfy (2c), note that the order of the sequence \mathcal{P}_{∞} given the generators must be the Sobol order. Equation (2b) also gives us another condition on j_m 's: $b^{m-1} \mid j_m - j_{m-1} \to j_m = j_{m-1} + b^{m-1}, \forall m \in \mathbb{N}$.

For $m \in \mathbb{N}_0$ define the dual net corresponding to \mathcal{P}_m as

$$\mathcal{P}_m^{\perp} = \{ \boldsymbol{k} \in \mathbb{K} : \langle \boldsymbol{k}, \boldsymbol{t}_i \rangle = 0, \ i = 0, \dots, b^m - 1 \}$$
$$= \{ \boldsymbol{k} \in \mathbb{K} : \langle \boldsymbol{k}, \boldsymbol{t}_{h^{\ell}} \rangle = 0, \ \ell = 0, \dots, m - 1 \}.$$

By this definition $\mathcal{P}_0^{\perp} = \mathbb{K}$. The properties of the bilinear transform, (1), imply that the dual net \mathcal{P}_m^{\perp} is a subgroup, and even a subspace, of the dual net $\mathcal{P}_{\ell}^{\perp}$ for all $\ell = 0, \ldots, m-1$.

The next goal is to define the map $\hat{\boldsymbol{\nu}}: \mathbb{K} \to \mathbb{F}_b^{\infty}$, and $\tilde{\nu}_m: \mathbb{K} \to \mathbb{F}_{b^m}$ that facilitates the calculation of the discrete Fourier transform introduced below.

Definition 1. For every $k \in \mathbb{K}$, let

$$\hat{\boldsymbol{\nu}}(\boldsymbol{k}) = (\hat{\nu}_0(\boldsymbol{k}), \hat{\nu}_1(\boldsymbol{k}), \hat{\nu}_2(\boldsymbol{k}), \ldots), \tag{4a}$$

$$\hat{\nu}_0(\mathbf{k}) = b \langle \mathbf{k}, \mathbf{t}_1 \rangle, \qquad \hat{\nu}_m(\mathbf{k}) = b \langle \mathbf{k}, \mathbf{t}_{b^m} \rangle - \langle \mathbf{k}, \mathbf{t}_{b^{m-1}} \rangle, \quad m \in \mathbb{N},$$
 (4b)

$$\tilde{\nu}_m(\mathbf{k}) = \sum_{\ell=0}^{m-1} \hat{\nu}_\ell(\mathbf{k}) b^\ell, \quad m \in \mathbb{N}.$$
(4c)

These maps have certain desirable properties.

Lemma 1. The following is true for the maps defined in Definition 1:

- a) $\hat{\boldsymbol{\nu}}(\mathbf{0}) = \mathbf{0}$ and $\tilde{\nu}_m(\mathbf{0}) = 0$ for all $m \in \mathbb{N}$.
- b) $\hat{\nu}_m(\mathbf{k}) \in \{0, \dots, b-1\}$ and $\tilde{\nu}_m(\mathbf{k}) \in \{0, \dots, b^m 1\}$ for all $m \in \mathbb{N}_0$. c) for all $m \in \mathbb{N}_0$ and all $\mathbf{\nu} \in \mathbb{F}_b^m$ there exist a unique $\mathbf{k} \in \mathbb{K}$ with $\hat{\mathbf{\nu}}(\mathbf{k}) = \mathbf{k}$ $(\nu_0,\ldots,\nu_{m-1},\ldots).$
- d) for any $m \in \mathbb{N}_0$, $i \in \{0, ..., b^m 1\}$, $\tilde{\nu}_m(\mathbf{k}) = \nu = (\nu_0, \nu_1, ...)$, and $\vec{i} = (i_0, i_1, \ldots), it follows that$

$$\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle = \sum_{\ell=0}^{m-1} i_{\ell} [\nu \pmod{b^{(\ell+1)}}] b^{-(\ell+1)} \pmod{1}$$
 (5)

Proof. a) Directly from definition.

- b) Using (2c) and by construction, $\hat{\nu}_0(\mathbf{k}) \in \{0, \dots, b-1\}$ and $\hat{\nu}_m(\mathbf{k}) \in (-1, b)$. Using the assumption (2b), $\hat{\nu}_m(\mathbf{k}) \pmod{1} = \mathbf{k}^T b \mathbf{t}_{b^m} \pmod{1} - \mathbf{k}^T \mathbf{t}_{b^{m-1}}$ $(\text{mod } 1) = 0. \text{ Then, } \hat{\nu}_m(\mathbf{k}) \in (-1, b) \cap \mathbb{Z} = \{0, \dots, b-1\}, \forall m \in \mathbb{N}_0.$
- c) For injectivity, we prove that $\hat{\boldsymbol{\nu}}(\boldsymbol{k}) = \hat{\boldsymbol{\nu}}(\boldsymbol{l}) \Rightarrow \boldsymbol{k} = \boldsymbol{l}$. If $\hat{\boldsymbol{\nu}}(\boldsymbol{k}) = \hat{\boldsymbol{\nu}}(\boldsymbol{l})$, $\hat{\nu}_m(\boldsymbol{k}) = \hat{\boldsymbol{\nu}}(\boldsymbol{l})$ $\hat{\nu}_m(\mathbf{l}), \forall m \in \mathbb{N}_0$. In particular for m = 0, this implies $\langle \mathbf{k}, \mathbf{t}_1 \rangle - \langle \mathbf{l}, \mathbf{t}_1 \rangle = 0$. Assume now that $\langle \mathbf{k}, \mathbf{t}_{b^m} \rangle - \langle \mathbf{l}, \mathbf{t}_{b^m} \rangle = 0$. Since $\hat{\nu}_{m+1}(\mathbf{k}) - \hat{\nu}_{m+1}(\mathbf{l}) = 0$, then $\langle \boldsymbol{k}, t_{b^{m+1}} \rangle - \langle \boldsymbol{l}, t_{b^{m+1}} \rangle = 0$. By induction $\langle \boldsymbol{k}, t_{b^m} \rangle - \langle \boldsymbol{l}, t_{b^m} \rangle = \langle \boldsymbol{k} - \boldsymbol{l}, t_{b^m} \rangle = 0$ for all $m \in \mathbb{N}_0$. Thus, by (2d), $\mathbf{k} = \mathbf{l}$.

For the surjection, by (1e) there exists k such that $\hat{\nu}_0(k) = \nu \neq 0$. Furthermore due to the property (1d), $\hat{\nu}_0(a\mathbf{k}) = a\nu \pmod{b}$ and recalling the Lagrange's Theorem, any element ν different of the identity generates the group \mathbb{F}_b . Therefore, a = 1, ..., b - 1 gives us any element of \mathbb{F}_b . Now, for any $l \leq m \in \mathbb{N}$ and using (3),

$$\hat{\nu}_{l}(\mathbf{k} + b^{m}\mathbf{a}) = \begin{cases} b\langle \mathbf{k}, \mathbf{t}_{b^{l}} \rangle - \langle \mathbf{k}, \mathbf{t}_{b^{l-1}} \rangle + b\langle \mathbf{a}, \mathbf{t}_{1} \rangle \pmod{b} & \text{if } l = m, \\ b\langle \mathbf{k}, \mathbf{t}_{b^{l}} \rangle - \langle \mathbf{k}, \mathbf{t}_{b^{l-1}} \rangle & \text{if } l < m \end{cases}$$

$$= \begin{cases} \hat{\nu}_{l}(\mathbf{k}) + \hat{\nu}_{0}(\mathbf{a}) \pmod{b} & \text{if } l = m, \\ \hat{\nu}_{l}(\mathbf{k}) & \text{if } l < m \end{cases}$$

Therefore, for all $m \in \mathbb{N}_0$ and all $(\nu_0, \dots, \nu_m) \in \mathbb{F}_b^{m+1}$ there exist a $k \in \mathbb{K}$ with $\hat{\nu}_l(\mathbf{k}) = \nu_l$, $l = 0, \dots, m$. This means $\hat{\boldsymbol{\nu}}(\mathbf{k})$ is bijective.

d) Follows by applying (1c) and Definition 1:

$$\langle \boldsymbol{k}, \boldsymbol{t}_{i} \rangle = \left\langle \boldsymbol{k}, \sum_{\ell=0}^{m-1} i_{\ell} \boldsymbol{t}_{b^{\ell}} \right\rangle = \sum_{\ell=0}^{m-1} i_{\ell} \langle \boldsymbol{k}, \boldsymbol{t}_{b^{\ell}} \rangle \pmod{1}$$

$$= \sum_{\ell=0}^{m-1} i_{\ell} \sum_{j=0}^{\ell} \nu_{j} b^{j-(\ell+1)} \pmod{1}$$

$$= \sum_{\ell=0}^{m-1} i_{\ell} [\nu \pmod{b^{(\ell+1)}}] b^{-(\ell+1)} \pmod{1}.$$

1.3. Fourier Series and Discrete Transforms

The integrands are assumed to belong to some subset of $\mathcal{L}_2([0,1)^d)$, the space of square integrable functions. The \mathcal{L}_2 inner product is defined as

$$\langle f, g \rangle_2 = \int_{[0,1)^d} f(\boldsymbol{x}) \overline{g(\boldsymbol{x})} \, \mathrm{d}\boldsymbol{x}.$$

Let $\{\varphi(\cdot, \mathbf{k}) \in \mathcal{L}_2([0, 1)^d) : \mathbf{k} \in \mathbb{K}\}$ be the complete orthonormal Walsh function basis for $\mathcal{L}_2([0, 1)^d)$, i.e.,

$$\varphi(\boldsymbol{x}, \boldsymbol{k}) = e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{x} \rangle/b}, \qquad \boldsymbol{k} \in \mathbb{K}, \ \boldsymbol{x} \in [0, 1)^d.$$

Then any function in \mathcal{L}_2 may be written in series form as

$$f(\boldsymbol{x}) = \sum_{\boldsymbol{k} \in \mathbb{K}} \hat{f}(\boldsymbol{k}) \varphi(\boldsymbol{x}, \boldsymbol{k}), \text{ where } \hat{f}(\boldsymbol{k}) = \langle f, \varphi(\cdot, \boldsymbol{k}) \rangle_2,$$
 (6)

and the inner product of two functions in \mathcal{L}_2 is the ℓ_2 inner product of their series coefficients:

$$\langle f,g\rangle_2 = \sum_{\boldsymbol{k}\in\mathbb{K}} \hat{f}(\boldsymbol{k})\overline{\hat{g}(\boldsymbol{k})} =: \left\langle \left(\hat{f}(\boldsymbol{k})\right)_{\boldsymbol{k}\in\mathbb{K}}, \left(\hat{g}(\boldsymbol{k})\right)_{\boldsymbol{k}\in\mathbb{K}}\right\rangle_2.$$

For all $k \in \mathbb{K}$ and $x \in \mathcal{P}$, it follows that

$$0 = \frac{1}{b^m} \sum_{i=0}^{b^m - 1} \left[\varphi(\boldsymbol{t}_i, \boldsymbol{k}) - \varphi(\boldsymbol{t}_i \oplus \boldsymbol{x}, \boldsymbol{k}) \right] = \frac{1}{b^m} \sum_{i=0}^{b^m - 1} \left[e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle} - e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \oplus \boldsymbol{x} \rangle} \right]$$
$$= \frac{1}{b^m} \sum_{i=0}^{b^m - 1} \left[e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle} - e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle + \langle \boldsymbol{k}, \boldsymbol{x} \rangle} \right] \quad \text{by (1c)}$$
$$= \left[1 - e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{x} \rangle} \right] \frac{1}{b^m} \sum_{i=0}^{b^m - 1} e^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle}.$$

By this equality it follows that the average of a basis function, $\varphi(\cdot, \mathbf{k})$, over the points in a node set is either one or zero, depending on whether \mathbf{k} is in the dual set or not.

$$\frac{1}{b^m} \sum_{i=0}^{b^m-1} \mathrm{e}^{2\pi\sqrt{-1}\langle \boldsymbol{k}, \boldsymbol{t}_i \rangle} = \mathbb{1}_{\mathcal{P}_m^{\perp}}(\boldsymbol{k}) = \begin{cases} 1, & \boldsymbol{k} \in \mathcal{P}_m^{\perp} \\ 0, & \boldsymbol{k} \in \mathbb{K} \backslash \mathcal{P}_m^{\perp}. \end{cases}$$

Given the sequence $\{t_i\}_{i=0}^{\infty}$, one may also define a shifted sequence $\{x_i = t_i \oplus \Delta\}_{i=0}^{\infty}$, where $\Delta \in [0,1)^d$. Define the discrete transform of a function, f,

over the shifted net as

$$\tilde{f}_{m}(\mathbf{k}) := \frac{1}{b^{m}} \sum_{i=0}^{b^{m}-1} e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{x}_{i} \rangle} f(\mathbf{x}_{i}) \tag{7}$$

$$= \frac{1}{b^{m}} \sum_{i=0}^{b^{m}-1} \left[e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{x}_{i} \rangle} \sum_{\mathbf{l} \in \mathbb{K}} \hat{f}(\mathbf{l}) e^{2\pi\sqrt{-1}\langle \mathbf{l}, \mathbf{x}_{i} \rangle} \right]$$

$$= \sum_{\mathbf{l} \in \mathbb{K}} \hat{f}(\mathbf{l}) \frac{1}{b^{m}} \sum_{i=0}^{b^{m}-1} e^{2\pi\sqrt{-1}\langle \mathbf{l} - \mathbf{k}, \mathbf{x}_{i} \rangle}$$

$$= \sum_{\mathbf{l} \in \mathbb{K}} \hat{f}(\mathbf{l}) e^{2\pi\sqrt{-1}\langle \mathbf{l} - \mathbf{k}, \Delta \rangle} \frac{1}{b^{m}} \sum_{i=0}^{b^{m}-1} e^{2\pi\sqrt{-1}\langle \mathbf{l} - \mathbf{k}, \mathbf{t}_{i} \rangle}$$

$$= \sum_{\mathbf{l} \in \mathbb{K}} \hat{f}(\mathbf{l}) e^{2\pi\sqrt{-1}\langle \mathbf{l} - \mathbf{k}, \Delta \rangle} \mathbb{1}_{\mathcal{P}_{m}^{\perp}}(\mathbf{l} - \mathbf{k})$$

$$= \sum_{\mathbf{l} \in \mathcal{P}_{m}^{\perp}} \hat{f}(\mathbf{k} + \mathbf{l}) e^{2\pi\sqrt{-1}\langle \mathbf{l}, \Delta \rangle}$$

$$= \hat{f}(\mathbf{k}) + \sum_{\mathbf{l} \in \mathcal{P}_{m}^{\perp}} \hat{f}(\mathbf{k} + \mathbf{l}) e^{2\pi\sqrt{-1}\langle \mathbf{l}, \Delta \rangle}, \quad \forall \mathbf{k} \in \mathbb{K}.$$
(8)

It is seen here that the discrete transform $\tilde{f}_m(\mathbf{k})$ is equal to the integral transform $\hat{f}(\mathbf{k})$, defined in (6), plus the aliasing terms corresponding to $\hat{f}(\mathbf{l})$ where $\mathbf{l} - \mathbf{k} \in \mathcal{P}_m^{\perp} \setminus \mathbf{0}$.

1.4. Computation of the Discrete Transform

The discrete transform defined in (7) may also be expressed as

$$\tilde{f}_{m}(\mathbf{k}) = \frac{1}{b^{m}} \sum_{i=0}^{b^{m}-1} e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{t}_{i} \oplus \mathbf{\Delta} \rangle} f(\mathbf{t}_{i} \oplus \mathbf{\Delta})$$

$$= \frac{e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{\Delta} \rangle}}{b^{m}} \sum_{i=0}^{b^{m}-1} e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{t}_{i} \rangle} f(\mathbf{t}_{i} \oplus \mathbf{\Delta}).$$

Letting $y_i = f(t_i \oplus \Delta)$,

$$Y_{m,0}(i_0,\ldots,i_{m-1})=y_i, \qquad i=i_0+i_1b+\cdots+i_{m-1}b^{m-1},$$

and invoking Lemma 1, for any $\mathbf{k} \in \mathbb{K}$ with $\tilde{\nu}_m(\mathbf{k}) = \nu = \nu_0 + \nu_1 b + \dots + \nu_{m-1} b^{m-1}$ one may write

$$\begin{split} \tilde{f}_{m}(\boldsymbol{k}) &= \mathrm{e}^{-2\pi\sqrt{-1}\langle\boldsymbol{k},\boldsymbol{\Delta}\rangle}Y_{m,m}(\nu_{0},\ldots,\nu_{m-1}), \\ Y_{m,m}(\nu_{0},\ldots,\nu_{m-1}) \\ &:= \frac{1}{b^{m}}\sum_{i=0}^{b^{m}-1}\mathrm{e}^{-2\pi\sqrt{-1}\langle\boldsymbol{k},\boldsymbol{t}_{i}\rangle}y_{i} \\ &= \frac{1}{b^{m}}\sum_{i_{m-1}=0}^{b-1}\cdots\sum_{i_{0}=0}^{b-1}\mathrm{e}^{-2\pi\sqrt{-1}\sum_{\ell=0}^{m-1}i_{\ell}[\nu\pmod{b^{(\ell+1)}}]b^{-(\ell+1)}}Y_{m,0}(i_{0},\ldots,i_{m-1}) \\ &= \frac{1}{b}\sum_{i_{m-1}=0}^{b-1}\mathrm{e}^{-2\pi\sqrt{-1}i_{m-1}[\nu\pmod{b^{m}}]b^{-m}}\cdots \\ &\qquad \qquad \frac{1}{b}\sum_{i_{0}=0}^{b-1}\mathrm{e}^{-2\pi\sqrt{-1}i_{0}[\nu\pmod{b}]b^{-1}}Y_{m,0}(i_{0},\ldots,i_{m-1}) \end{split}$$

This sum can be computed recursively:

$$Y_{m,\ell+1}(\nu_0,\ldots,\nu_{\ell},i_{\ell+1},\ldots,i_m)$$

$$= \frac{1}{b} \sum_{i_{\ell}=0}^{b-1} e^{-2\pi\sqrt{-1}i_{\ell}[\nu \pmod{b^{(\ell+1)}}]b^{-(\ell+1)}} Y_{m,\ell}(\nu_0,\ldots,\nu_{\ell-1},i_{\ell},\ldots,i_m)$$

In light of this development we define $\mathring{f}_m(\nu) = Y_{m,m}(\nu_0,\dots,\nu_{m-1})$ for $\nu=0,\dots,b^m-1$. Then

$$\tilde{f}(\mathbf{k}) = e^{-2\pi\sqrt{-1}\langle \mathbf{k}, \mathbf{\Delta} \rangle} \mathring{f}_m(\tilde{\nu}(\mathbf{k})).$$

2. Error Estimation and an Automatic Algorithm

2.1. Wavenumber Map

Now we are going to map the non-negative numbers into the space of all wavenumbers using the dual sets. For every $\kappa \in \mathbb{N}_0$, we assign a wavenumber $\tilde{\boldsymbol{k}}(\kappa) \in \mathbb{K}$ iteratively according to the following constraints:

- i) k(0) = 0:
- ii) For any $\lambda, m \in \mathbb{N}_0$ and $\kappa = 0, \dots, b^m 1$, it follows that $\tilde{\nu}_m(\tilde{\boldsymbol{k}}(\kappa)) = \tilde{\nu}_m(\tilde{\boldsymbol{k}}(\kappa + \lambda b^m))$.

This last condition implies that $\tilde{\mathbf{k}}(\kappa) - \tilde{\mathbf{k}}(\kappa + \lambda b^m) \in \mathcal{P}_m^{\perp}$.

This wavenumber map allows us to introduce a shorthand notation that facilitates the later analysis for $\kappa \in \mathbb{N}_0$ and $m \in \mathbb{N}$:

$$\hat{f}_{\kappa} = \hat{f}(\tilde{\boldsymbol{k}}(\kappa)),$$

$$\tilde{f}_{m,\kappa} = \tilde{f}_{m}(\tilde{\boldsymbol{k}}(\kappa)) = e^{-2\pi\sqrt{-1}\langle \tilde{\boldsymbol{k}}(\kappa), \boldsymbol{\Delta} \rangle} \mathring{f}_{m}(\tilde{\boldsymbol{\nu}}_{m}(\tilde{\boldsymbol{k}}(\kappa)))$$

$$= e^{-2\pi\sqrt{-1}\langle \tilde{\boldsymbol{k}}(\kappa), \boldsymbol{\Delta} \rangle} \mathring{f}_{m,\kappa},$$

where $\mathring{f}_{m,\kappa} := \mathring{f}_m(\tilde{\nu}_m(\tilde{k}(\kappa)))$. According to (8), it follows that

$$\tilde{f}_{m,\kappa} = \hat{f}_{\kappa} + \sum_{\lambda=1}^{\infty} \hat{f}_{\kappa+\lambda b^m} e^{2\pi\sqrt{-1}\left\langle \tilde{\mathbf{k}}(\kappa+\lambda b^m) - \tilde{\mathbf{k}}(\kappa), \mathbf{\Delta} \right\rangle}.$$
 (9)

We want to use $\tilde{f}_{m,\kappa}$ to estimate \hat{f}_{κ} if m is signficantly larger than $\lfloor \log_b(\kappa) \rfloor$.

2.2. Sums of Series Coefficients and Their Bounds

Consider the following sums of the series coefficients defined for $\ell, m \in \mathbb{N}_0$, $\ell \leq m$:

$$S(m) = \sum_{\kappa = \lfloor b^{m-1} \rfloor}^{b^{m}-1} |\hat{f}_{\kappa}|, \qquad \hat{S}(\ell, m) = \sum_{\kappa = \lfloor b^{\ell-1} \rfloor}^{b^{\ell}-1} \sum_{\lambda = 1}^{\infty} |\hat{f}_{\kappa + \lambda b^{m}}|,$$

$$\check{S}(m) = \hat{S}(0, m) + \dots + \hat{S}(m, m) = \sum_{\kappa = b^{m}}^{\infty} |\hat{f}_{\kappa}|,$$

$$\tilde{S}(\ell, m) = \sum_{\kappa = \lfloor b^{\ell-1} \rfloor}^{b^{\ell}-1} |\tilde{f}_{m, \kappa}| = \sum_{\kappa = \lfloor b^{\ell-1} \rfloor}^{b^{\ell}-1} |\mathring{f}_{m, \kappa}|.$$

The first three kinds of sums, $S(\cdot)$, $\hat{S}(\cdot, \cdot)$, and $\check{S}(\cdot)$, which involve the true series coefficients, cannot be observed, but the last one, $\tilde{S}(\cdot, \cdot)$, which involves the discrete transform coefficients, can easily be observed.

We now make critical assumptions that $\widehat{S}(\ell,m)$ and $\widecheck{S}(m)$ can be bounded above in terms of $S(\ell)$, provided that ℓ is large enough. Let $\ell, m \in \mathbb{N}_0$ with $\ell \leq m$, and fix $\ell_* \in \mathbb{N}$. It is assumed that their exist known, non-negative valued functions $\widehat{\omega}$ and $\widecheck{\omega}$ with $\lim_{m\to\infty} \widecheck{\omega}(m) = 0$ such that

$$\widehat{S}(\ell,m) \leqslant \widehat{\omega}(m-\ell)\widecheck{S}(m) \quad \forall \ell, \qquad \widecheck{S}(m) \leqslant \widecheck{\omega}(m-\ell)S(\ell) \quad \forall \ell_* \leqslant \ell.$$
 (10)

By the definition of $\check{S}(m)$, the choice $\widehat{\omega}(m) := 1$ for all m is always guaranteed to work. However, one might also consider choosing $\widehat{\omega}(m) = Cb^{-m}$ for some C. The reason for enforcing the second assumption only for $\ell \geqslant \ell_*$ is that for small ℓ , one might have a coincidentally small $S(\ell)$, since it only involves b^{ℓ} coefficients, while $\check{S}(m)$ is large.

Under this assumption, for $\ell, m \in \mathbb{N}$, $\ell_* \leq \ell \leq m$, it is possible to bound the sum of the true coefficients, $S(\ell)$, in terms of the observed sum of the discrete

coefficients, $\widetilde{S}(\ell, m)$, as follows:

$$\begin{split} S(\ell) &= \sum_{\kappa = b^{\ell-1}}^{b^{\ell}-1} \left| \hat{f}_{\kappa} \right| = \sum_{\kappa = b^{\ell-1}}^{b^{\ell}-1} \left| \tilde{f}_{m,\kappa} - \sum_{\lambda = 1}^{\infty} \hat{f}_{\kappa + \lambda b^{m}} \mathrm{e}^{2\pi \sqrt{-1} \left\langle \tilde{\mathbf{k}}(\kappa + \lambda b^{m}) \ominus \tilde{\mathbf{k}}(\kappa), \mathbf{\Delta} \right\rangle / b} \right| \\ &\leqslant \sum_{\kappa = b^{\ell-1}}^{b^{\ell}-1} \left| \tilde{f}_{m,\kappa} \right| + \sum_{\kappa = b^{\ell-1}}^{b^{\ell}-1} \sum_{\lambda = 1}^{\infty} \left| \hat{f}_{\kappa + \lambda b^{m}} \right| = \widetilde{S}(\ell,m) + \widehat{S}(\ell,m) \\ &\leqslant \widetilde{S}(\ell,m) + \widehat{\omega}(m-\ell) \widecheck{\omega}(m-\ell) S(\ell) \\ \\ S(\ell) \leqslant \frac{\widetilde{S}(\ell,m)}{1 - \widehat{\omega}(m-\ell) \widecheck{\omega}(m-\ell)} \quad \text{provided that } \widehat{\omega}(m-\ell) < 1. \end{split}$$

Using this upper bound, one can then conservatively bound the error of integration using the shifted node set. For for $\ell, m \in \mathbb{N}$, $\ell_* \leq \ell \leq m$, it follows that

$$\left| \int_{[0,1)^d} f(\boldsymbol{x}) \, d\boldsymbol{x} - \frac{1}{b^m} \sum_{i=0}^{b^m - 1} f(\boldsymbol{x}_i) \right|$$

$$= \left| \hat{f}(\boldsymbol{0}) - \tilde{f}_m(\boldsymbol{0}) \right| = \left| \hat{f}_0 - \tilde{f}_{m,0} \right| = \left| \sum_{\lambda=1}^{\infty} \hat{f}_{\lambda b^m} e^{2\pi \sqrt{-1} \left\langle \tilde{\boldsymbol{k}}(\lambda b^m), \boldsymbol{\Delta} \right\rangle} \right|$$

$$\leq \sum_{\lambda=1}^{\infty} \left| \hat{f}_{\lambda b^m} \right| = \hat{S}(0, m) \leq \hat{\omega}(m) \check{S}(m) \leq \hat{\omega}(m) \check{\omega}(m - \ell) S(\ell)$$

$$\leq \frac{\tilde{S}(\ell, m) \hat{\omega}(m) \check{\omega}(m - \ell)}{1 - \hat{\omega}(m - \ell) \check{\omega}(m - \ell)}.$$

This error bound suggests the following algorithm. Choose $r \in \mathbb{N}$ such that $\widehat{\omega}(r)\widecheck{\omega}(r) < 1$. For $j \in \mathbb{N}$ define

$$\ell_j = j + \ell_* - 1, \qquad m_j = j + \ell_* + r - 1, \qquad \mathfrak{C} = \frac{\widecheck{\omega}(r)}{1 - \widehat{\omega}(r)\widecheck{\omega}(r)}.$$

Define $\ell_j = \ell_* + j - 1$ and $m_j = \ell_j + r$. Given a tolerance ε , and an integrand f, do the following: for $j = 1, 2, \ldots$ check whether

$$\mathfrak{C}\widehat{\omega}(m_j)\widetilde{S}(\ell_j,m_j) \leqslant \varepsilon.$$

If so, we're done. If not, increment j by one and repeat.

Given $\widehat{\omega}$, $\widecheck{\omega}$, and r, one can compute \mathfrak{C} . Alternatively, given \mathfrak{C} , $\widehat{\omega}$, and r, one can compute $\widecheck{\omega}(r)$:

$$\mathfrak{C} = \frac{\widecheck{\omega}(r)}{1 - \widehat{\omega}(r)\widecheck{\omega}(r)} \Longleftrightarrow \widecheck{\omega}(r) = \frac{\mathfrak{C}}{1 + \mathfrak{C}\widehat{\omega}(r)}.$$

 F. J. Hickernell and H. Niederreiter, "The existence of good extensible rank-1 lattices," J. Complexity, vol. 19, pp. 286–300, 2003.