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ACM-ICPC World Finals 2022

2024-04-08

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## Contest (1)

template.cpp 14 lines

```
#include <bits/stdc++.h>
using namespace std;

#define rep(i, a, b) for(int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
typedef long long ll;
typedef pair<int, int> pii;
typedef vector<int> vi;
```

```
int main() {
    cin.tie(0); ios::sync_with_stdio(0);
    cin.exceptions(cin.failbit);
}
```

.bashrc 2 lines

```
alias c='g++ -Wall -Wconversion -Wfatal-errors -g -std=c++17 \
-fsanitize=undefined,address'
```

hash.sh 3 lines

```
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed.
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | dos2unix |
md5sum | cut -c-6
```

## Mathematics (2)

### 2.1 Equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by  $x = -b/2a$ .

$$\begin{aligned} ax + by &= e & \Rightarrow & x = \frac{ed - bf}{ad - bc} \\ cx + dy &= f & & y = \frac{af - ec}{ad - bc} \end{aligned}$$

Generally for  $Ax = b$ ,

$$x_i = \frac{\det A'_i}{\det A}$$

where  $A'_i$  is  $A$  with the  $i$ 'th column replaced by  $b$ .

### 2.2 Recurrences

If  $a_n = c_1 a_{n-1} + \dots + c_k a_{n-k}$ , and  $r_1, \dots, r_k$  are distinct roots of  $x^k - c_1 x^{k-1} - \dots - c_k$ , there are  $d_1, \dots, d_k$  s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots  $r$  become polynomial factors, e.g.

$$a_n = (d_1 n + d_2) r^n.$$

### 2.3 Trigonometry

$$\begin{aligned} \tan(v + w) &= \frac{\tan v + \tan w}{1 - \tan v \tan w} \\ \sin v + \sin w &= 2 \sin \frac{v + w}{2} \cos \frac{v - w}{2} \\ \cos v + \cos w &= 2 \cos \frac{v + w}{2} \cos \frac{v - w}{2} \end{aligned}$$

$$(V + W) \tan(v - w)/2 = (V - W) \tan(v + w)/2$$

where  $V, W$  are lengths of sides opposite angles  $v, w$ .

$$a \cos x + b \sin x = r \cos(x - \phi)$$

$$a \sin x + b \cos x = r \sin(x + \phi)$$

where  $r = \sqrt{a^2 + b^2}$ ,  $\phi = \text{atan2}(b, a)$ .

## 2.4 Geometry

### 2.4.1 Triangles

Side lengths:  $a, b, c$

$$\text{Semiperimeter: } p = \frac{a + b + c}{2}$$

$$\text{Area: } A = \sqrt{p(p - a)(p - b)(p - c)}$$

$$\text{Circumradius: } R = \frac{abc}{4A}$$

$$\text{Inradius: } r = \frac{A}{p}$$

Length of median (divides triangle into two equal-area triangles):

$$m_a = \frac{1}{2} \sqrt{2b^2 + 2c^2 - a^2}$$

Length of bisector (divides angles in two):

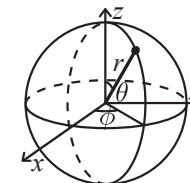
$$s_a = \sqrt{bc \left[ 1 - \left( \frac{a}{b + c} \right)^2 \right]}$$

$$\text{Law of sines: } \frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$$

$$\text{Law of cosines: } a^2 = b^2 + c^2 - 2bc \cos \alpha$$

$$\text{Law of tangents: } \frac{a + b}{a - b} = \frac{\tan \frac{\alpha + \beta}{2}}{\tan \frac{\alpha - \beta}{2}}$$

### 2.4.2 Spherical coordinates



$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \arccos(z / \sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \text{atan2}(y, x) \end{aligned}$$

### 2.4.3 Quadrilaterals

With side lengths  $a, b, c, d$ , diagonals  $e, f$ , diagonals angle  $\theta$ , area  $A$  and magic flux  $F = b^2 + d^2 - a^2 - c^2$ :

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is  $180^\circ$ ,  $ef = ac + bd$ , and  $A = \sqrt{(p - a)(p - b)(p - c)(p - d)}$ .

2.5 Derivatives/Integrals

$$\frac{d}{dx} \arcsin x = \frac{1}{\sqrt{1-x^2}}$$
$$\frac{d}{dx} \tan x = 1 + \tan^2 x$$
$$\int \tan ax = -\frac{\ln|\cos ax|}{a}$$
$$\int e^{-x^2} = \frac{\sqrt{\pi}}{2}\text{erf}(x)$$

$$\frac{d}{dx} \arccos x = -\frac{1}{\sqrt{1-x^2}}$$
$$\frac{d}{dx} \arctan x = \frac{1}{1+x^2}$$
$$\int x \sin ax = \frac{\sin ax - ax \cos ax}{a^2}$$
$$\int x e^{ax} dx = \frac{e^{ax}}{a^2}(ax-1)$$

Integration by parts:

$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx$$

$$\int x^n dx = \frac{1}{n+1} x^{n+1}$$

$$\int \frac{1}{x} dx = \ln|x|$$

$$\int u dv = uv - \int v du$$

$$\int \frac{1}{ax+b} dx = \frac{1}{a} \ln|ax+b|$$

$$\int \frac{1}{(x+a)^2} dx = -\frac{1}{x+a}$$

$$\int (x+a)^n dx = \frac{(x+a)^{n+1}}{n+1}, n \neq -1$$

$$\int x(x+a)^n dx = \frac{(x+a)^{n+1}((n+1)x-a)}{(n+1)(n+2)}$$

$$\int \frac{1}{1+x^2} dx = \tan^{-1} x$$

$$\int \frac{1}{a^2+x^2} dx = \frac{1}{a} \tan^{-1} \frac{x}{a}$$

$$\int \frac{x}{a^2+x^2} dx = \frac{1}{2} \ln|a^2+x^2|$$

$$\int \frac{x^2}{a^2+x^2} dx = x - a \tan^{-1} \frac{x}{a}$$

$$\int \frac{x^3}{a^2+x^2} dx = \frac{1}{2}x^2 - \frac{1}{2}a^2 \ln|a^2+x^2|$$

$$\int \frac{1}{ax^2+bx+c} dx = \frac{2}{\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}}$$

$$\int \frac{1}{(x+a)(x+b)} dx = \frac{1}{b-a} \ln \frac{a+x}{b+x}, \quad a \neq b$$

$$\int \frac{x}{(x+a)^2} dx = \frac{a}{a+x} + \ln|a+x|$$

$$\int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln|ax^2+bx+c| - \frac{b}{a\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}}$$

$$\int \sqrt{x-a} dx = \frac{2}{3}(x-a)^{3/2}$$

$$\int \frac{1}{\sqrt{x \pm a}} dx = 2\sqrt{x \pm a}$$

$$\int \frac{1}{\sqrt{a-x}} dx = -2\sqrt{a-x}$$

$$\int x\sqrt{x-a} dx = \frac{2}{3}a(x-a)^{3/2} + \frac{2}{5}(x-a)^{5/2}$$

$$\int \sqrt{ax+bd} dx = \left(\frac{2b}{3a} + \frac{2x}{3}\right) \sqrt{ax+b}$$

$$\int (ax+b)^{3/2} dx = \frac{2}{5a} (ax+b)^{5/2}$$

$$\int \frac{x}{\sqrt{x \pm a}} dx = \frac{2}{3}(x \mp 2a)\sqrt{x \pm a}$$

$$\int \sqrt{\frac{x}{a-x}} dx = -\sqrt{x(a-x)} - a \tan^{-1} \frac{\sqrt{x(a-x)}}{x-a}$$

$$\int \sqrt{\frac{x}{a+x}} dx = \sqrt{x(a+x)} - a \ln\left[\sqrt{x} + \sqrt{x+a}\right]$$

$$\int x\sqrt{ax+bd} dx = \frac{2}{15a^2}(-2b^2+abx+3a^2x^2)\sqrt{ax+b}$$

$$\int \sqrt{x(ax+b)} dx = \frac{1}{4a^{3/2}} \left[ (2ax+b)\sqrt{ax(ax+b)} - b^2 \ln\left|a\sqrt{x} + \sqrt{a(ax+b)}\right| \right]$$

$$\int \sqrt{x^3(ax+b)} dx = \left[ \frac{b}{12a} - \frac{b^2}{8a^2x} + \frac{x}{3} \right] \sqrt{x^3(ax+b)} + \frac{b^3}{8a^{5/2}} \ln\left|a\sqrt{x} + \sqrt{a(ax+b)}\right|$$

$$\int \sqrt{x^2 \pm a^2} dx = \frac{1}{2}x\sqrt{x^2 \pm a^2} \pm \frac{1}{2}a^2 \ln\left|x + \sqrt{x^2 \pm a^2}\right|$$

$$\int \sqrt{a^2-x^2} dx = \frac{1}{2}x\sqrt{a^2-x^2} + \frac{1}{2}a^2 \tan^{-1} \frac{x}{\sqrt{a^2-x^2}}$$

$$\int x\sqrt{x^2 \pm a^2} dx = \frac{1}{3}\left(x^2 \pm a^2\right)^{3/2}$$

$$\int \frac{1}{\sqrt{x^2 \pm a^2}} dx = \ln\left|x + \sqrt{x^2 \pm a^2}\right|$$

$$\int \frac{1}{\sqrt{a^2-x^2}} dx = \sin^{-1} \frac{x}{a}$$

$$\int \frac{x}{\sqrt{x^2 \pm a^2}} dx = \sqrt{x^2 \pm a^2}$$

$$\int \frac{x}{\sqrt{a^2-x^2}} dx = -\sqrt{a^2-x^2}$$

$$\int \frac{x^2}{\sqrt{x^2 \pm a^2}} dx = \frac{1}{2}x\sqrt{x^2 \pm a^2} \mp \frac{1}{2}a^2 \ln\left|x + \sqrt{x^2 \pm a^2}\right|$$

$$\int \sqrt{ax^2+bx+cd} dx = \frac{b+2ax}{4a} \sqrt{ax^2+bx+c} + \frac{4ac-b^2}{8a^{3/2}} \ln\left|2ax+b+2\sqrt{a(ax^2+bx+c)}\right|$$

$$\int x\sqrt{ax^2+bx+c} = \frac{1}{48a^{5/2}} \left( 2\sqrt{a}\sqrt{ax^2+bx+c} \times \left(-3b^2+2abx+8a(c+ax^2)\right) + 3(b^3-4abc) \ln\left|b+2ax+2\sqrt{a}\sqrt{ax^2+bx+c}\right| \right)$$

$$\int \frac{1}{\sqrt{ax^2+bx+c}} dx = \frac{1}{\sqrt{a}} \ln\left|2ax+b+2\sqrt{a(ax^2+bx+c)}\right|$$

$$\int \frac{x}{\sqrt{ax^2+bx+c}} dx = \frac{1}{a} \sqrt{ax^2+bx+c} - \frac{b}{2a^{3/2}} \ln\left|2ax+b+2\sqrt{a(ax^2+bx+c)}\right|$$

$$\int \frac{dx}{(a^2+x^2)^{3/2}} = \frac{x}{a^2\sqrt{a^2+x^2}}$$

$$\int \ln ax dx = x \ln ax - x$$

$$\int \frac{\ln ax}{x} dx = \frac{1}{2} (\ln ax)^2$$

$$\int \ln(ax+b) dx = \left(x + \frac{b}{a}\right) \ln(ax+b) - x, a \neq 0$$

$$\int \ln(x^2+a^2) \, \mathrm{d}x = x \ln(x^2+a^2) + 2a \tan^{-1} \frac{x}{a} - 2x$$

$$\int \ln(x^2-a^2) \, \mathrm{d}x = x \ln(x^2-a^2) + a \ln \frac{x+a}{x-a} - 2x$$

$$\int \ln\left(ax^2+bx+c\right) dx = \frac{1}{a} \sqrt{4ac-b^2} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}} - 2x + \left(\frac{b}{2a} + x\right) \ln\left(ax^2+bx+c\right)$$

$$\int x \ln(ax+b) dx = \frac{bx}{2a} - \frac{1}{4}x^2 + \frac{1}{2} \left( x^2 - \frac{b^2}{a^2} \right) \ln(ax+b) \quad (2.48)$$

$$\int x \ln(a^2 - b^2 x^2) dx = -\frac{1}{2}x^2 + \frac{1}{2} \left( x^2 - \frac{a^2}{b^2} \right) \ln(a^2 - b^2 x^2) \quad (2.49)$$

$$\int e^{ax} dx = \frac{1}{a} e^{ax} \quad (2.50)$$

$$\int \sqrt{x} e^{ax} dx = \frac{1}{a} \sqrt{x} e^{ax} + \frac{i\sqrt{\pi}}{2a^{3/2}} \operatorname{erf}(i\sqrt{ax}),$$

where  $\operatorname{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-t^2} dt$  (2.51)

$$\int x e^x dx = (x-1)e^x \quad (2.52)$$

$$\int x e^{ax} dx = \left( \frac{x}{a} - \frac{1}{a^2} \right) e^{ax} \quad (2.53)$$

$$\int x^2 e^x dx = (x^2 - 2x + 2) e^x \quad (2.54)$$

$$\int x^2 e^{ax} dx = \left( \frac{x^2}{a} - \frac{2x}{a^2} + \frac{2}{a^3} \right) e^{ax} \quad (2.55)$$

$$\int x^3 e^x dx = (x^3 - 3x^2 + 6x - 6) e^x \quad (2.56)$$

$$\int x^n e^{ax} dx = \frac{x^n e^{ax}}{a} - \frac{n}{a} \int x^{n-1} e^{ax} dx \quad (2.57)$$

$$\int x^n e^{ax} dx = \frac{(-1)^n}{a^{n+1}} \Gamma[1+n, -ax],$$

where  $\Gamma(a, x) = \int_x^\infty t^{a-1} e^{-t} dt$  (2.58)

$$\int e^{ax^2} dx = -\frac{i\sqrt{\pi}}{2\sqrt{a}} \operatorname{erf}(i\sqrt{a}x) \quad (2.59)$$

$$\int e^{-ax^2} dx = \frac{\sqrt{\pi}}{2\sqrt{a}} \operatorname{erf}(x\sqrt{a}) \quad (2.60)$$

$$\int x e^{-ax^2} dx = -\frac{1}{2a} e^{-ax^2} \quad (2.61)$$

$$\int x^2 e^{-ax^2} dx = \frac{1}{4} \sqrt{\frac{\pi}{a^3}} \operatorname{erf}(x\sqrt{a}) - \frac{x}{2a} e^{-ax^2} \quad (2.62)$$

$$\int \sin ax dx = -\frac{1}{a} \cos ax \quad (2.63)$$

$$\int \sin^2 ax dx = \frac{x}{2} - \frac{\sin 2ax}{4a} \quad (2.64)$$

$$\int \sin^3 ax dx = -\frac{3 \cos ax}{4a} + \frac{\cos 3ax}{12a} \quad (2.65)$$

$$\int \cos ax dx = \frac{1}{a} \sin ax \quad (2.66)$$

$$\int \cos^2 ax dx = \frac{x}{2} + \frac{\sin 2ax}{4a} \quad (2.67)$$

$$\int \cos^3 ax dx = \frac{3 \sin ax}{4a} + \frac{\sin 3ax}{12a} \quad (2.68)$$

$$\int \cos ax \sin bxdx = \frac{\cos[(a-b)x]}{2(a-b)} - \frac{\cos[(a+b)x]}{2(a+b)}, a \neq b \quad (2.69)$$

$$\int \sin^2 ax \cos bxdx = -\frac{\sin[(2a-b)x]}{4(2a-b)} + \frac{\sin bx}{2b} - \frac{\sin[(2a+b)x]}{4(2a+b)} \quad (2.70)$$

$$\int \sin^2 x \cos x dx = \frac{1}{3} \sin^3 x \quad (2.71)$$

$$\int \cos^2 ax \sin bxdx = \frac{\cos[(2a-b)x]}{4(2a-b)} - \frac{\cos bx}{2b} - \frac{\cos[(2a+b)x]}{4(2a+b)} \quad (2.72)$$

$$\int \cos^2 ax \sin ax dx = -\frac{1}{3a} \cos^3 ax \quad (2.73)$$

$$\int \sin^2 ax \cos^2 bxdx = \frac{x}{4} - \frac{\sin 2ax}{8a} - \frac{\sin[2(a-b)x]}{16(a-b)} + \frac{\sin 2bx}{8b} - \frac{\sin[2(a+b)x]}{16(a+b)} \quad (2.74)$$

$$\int \sin^2 ax \cos^2 ax dx = \frac{x}{8} - \frac{\sin 4ax}{32a} \quad (2.75)$$

$$\int \tan ax dx = -\frac{1}{a} \ln \cos ax \quad (2.76)$$

$$\int \tan^2 ax dx = -x + \frac{1}{a} \tan ax \quad (2.77)$$

$$\int \tan^3 ax dx = \frac{1}{a} \ln \cos ax + \frac{1}{2a} \sec^2 ax \quad (2.78)$$

$$\int \sec x dx = \ln |\sec x + \tan x| = 2 \tanh^{-1} \left( \tan \frac{x}{2} \right) \quad (2.79)$$

$$\int \sec^2 ax dx = \frac{1}{a} \tan ax \quad (2.80)$$

$$\int \sec^3 x dx = \frac{1}{2} \sec x \tan x + \frac{1}{2} \ln |\sec x + \tan x| \quad (2.81)$$

$$\int \sec x \tan x dx = \sec x \quad (2.82)$$

$$\int \sec^2 x \tan x dx = \frac{1}{2} \sec^2 x \quad (2.83)$$

$$\int \sec^n x \tan x dx = \frac{1}{n} \sec^n x, n \neq 0 \quad (2.84)$$

$$\int \csc x dx = \ln \left| \tan \frac{x}{2} \right| = \ln |\csc x - \cot x| + C \quad (2.85)$$

$$\int \csc^2 ax dx = -\frac{1}{a} \cot ax \quad (2.86)$$

$$\int \csc^3 x dx = -\frac{1}{2} \cot x \csc x + \frac{1}{2} \ln |\csc x - \cot x| \quad (2.87)$$

$$\int \csc^n x \cot x dx = -\frac{1}{n} \csc^n x, n \neq 0 \quad (2.88)$$

$$\int \sec x \csc x dx = \ln |\tan x| \quad (2.89)$$

$$\int x \cos x dx = \cos x + x \sin x \quad (2.90)$$

$$\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{x}{a} \sin ax \quad (2.91)$$

$$\int x^2 \cos x dx = 2x \cos x + (x^2 - 2) \sin x \quad (2.92)$$

$$\int x^2 \cos ax dx = \frac{2x \cos ax}{a^2} + \frac{a^2 x^2 - 2}{a^3} \sin ax \quad (2.93)$$

$$\int x^n \cos x dx = -\frac{1}{2}(i)^{n+1} [\Gamma(n+1, -ix) + (-1)^n \Gamma(n+1, ix)] \quad (2.94)$$

$$\int x^n \cos ax dx = \frac{1}{2}(ia)^{1-n} [(-1)^n \Gamma(n+1, -iax) - \Gamma(n+1, iax)] \quad (2.95)$$

$$\int x \sin x dx = -x \cos x + \sin x \quad (2.96)$$

$$\int x \sin ax dx = -\frac{x \cos ax}{a} + \frac{\sin ax}{a^2} \quad (2.97)$$

$$\int x^2 \sin x dx = (2 - x^2) \cos x + 2x \sin x \quad (2.98)$$

$$\int x^2 \sin ax dx = \frac{2 - a^2 x^2}{a^3} \cos ax + \frac{2x \sin ax}{a^2} \quad (2.99)$$

$$\int x^n \sin x dx = -\frac{1}{2} (i)^n [\Gamma(n+1, -ix) - (-1)^n \Gamma(n+1, -ix)] \tag{2.100}$$

$$\int e^x \sin x dx = \frac{1}{2} e^x (\sin x - \cos x) \tag{2.101}$$

$$\int e^{bx} \sin ax dx = \frac{1}{a^2 + b^2} e^{bx} (b \sin ax - a \cos ax) \tag{2.102}$$

$$\int e^x \cos x dx = \frac{1}{2} e^x (\sin x + \cos x) \tag{2.103}$$

$$\int e^{bx} \cos ax dx = \frac{1}{a^2 + b^2} e^{bx} (a \sin ax + b \cos ax) \tag{2.104}$$

$$\int x e^x \sin x dx = \frac{1}{2} e^x (\cos x - x \cos x + x \sin x) \tag{2.105}$$

$$\int x e^x \cos x dx = \frac{1}{2} e^x (x \cos x - \sin x + x \sin x) \tag{2.106}$$

2.6 Sums

$$c^a + c^{a+1} + \dots + c^b = \frac{c^{b+1} - c^a}{c - 1}, c \neq 1$$

$$\begin{aligned} \sum_{i=1}^n i^1 &= \frac{n(n+1)}{2} = \frac{1}{2} n^2 + \frac{1}{2} n \\ \sum_{i=1}^n i^2 &= \frac{n(n+1)(2n+1)}{6} = \frac{1}{3} n^3 + \frac{1}{2} n^2 + \frac{1}{6} n \\ \sum_{i=1}^n i^3 &= \left[ \frac{n(n+1)}{2} \right]^2 = \frac{1}{4} n^4 + \frac{1}{2} n^3 + \frac{1}{4} n^2 \\ \sum_{i=1}^n i^4 &= \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30} \\ &= \frac{1}{5} n^5 + \frac{1}{2} n^4 + \frac{1}{3} n^3 - \frac{1}{30} n \\ \sum_{i=1}^n i^5 &= \frac{n^2(n+1)^2(2n^2+2n-1)}{12} \\ &= \frac{1}{6} n^6 + \frac{1}{2} n^5 + \frac{5}{12} n^4 - \frac{1}{12} n^2 \\ \sum_{i=1}^n i^6 &= \frac{n(n+1)(2n+1)(3n^4+6n^3-3n+1)}{42} \\ &= \frac{1}{7} n^7 + \frac{1}{2} n^6 + \frac{1}{2} n^5 - \frac{1}{6} n^3 + \frac{1}{42} n \end{aligned}$$

2.7 Series

$$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \leq 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \leq x \leq 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

Polynomial Tricks

Newton’s Method:  $F_{t+1}(x) \equiv F_t(x) - \frac{G(F_t(x))}{G'(F_t(x))}$

Lagrange inversion:  $n[x^n]F(x)^k = k[x^{-k}]G(x)^{-n}$  (where  $F(G(x)) = x$ ) or  $[x^n]F(x)^k = [x^{-k-1}]G(x)^{-n-1}G'$

With  $H$ :  $[x^n]H(F(x)) = 1/n[x^{n-1}]H'(x)(x/G(x))^n$  or  $[x^n]H(F(x)) = [x^n]H(x)(x/G(x))^{n+1}G'(x)$

Bernoulli

Let  $S_k(x) = \sum_{i=0}^{x-1} i^k$ .

$$S_k(x) = \frac{1}{k+1} \sum_{g=0}^k C_{k+1}^g B_g x^{k+1-g} = k! \sum_{g=1}^{k+1} x^g / g! \times B_{k+1-g} / (k+1-g)!.$$

$$\sum_{i=0}^\infty \frac{B_i}{i!} x^i = \frac{1}{\sum_{i=0}^\infty \frac{x^i}{(i+1)!}}$$

If you need  $\sum_{i=0}^x i^k$  add one to  $B_1$ .

$$\sum_{i=1}^n i^m = \frac{1}{m+1} \sum_{k=0}^m C(m+1, k) B_k n^{m+1-k} \quad (m > 0)$$

2.8 Probability theory

Let  $X$  be a discrete random variable with probability  $p_X(x)$  of assuming the value  $x$ . It will then have an expected value (mean)  $\mu = \mathbb{E}(X) = \sum_x x p_X(x)$  and variance  $\sigma^2 = V(X) = \mathbb{E}(X^2) - (\mathbb{E}(X))^2 = \sum_x (x - \mathbb{E}(X))^2 p_X(x)$  where  $\sigma$  is the standard deviation. If  $X$  is instead continuous it will have a probability density function  $f_X(x)$  and the sums above will instead be integrals with  $p_X(x)$  replaced by  $f_X(x)$ .

Expectation is linear:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

For independent  $X$  and  $Y$ ,

$$V(aX + bY) = a^2 V(X) + b^2 V(Y).$$

2.8.1 Discrete distributions

Binomial distribution

The number of successes in  $n$  independent yes/no experiments, each which yields success with probability  $p$  is  $\text{Bin}(n, p)$ ,  $n = 1, 2, \dots$ ,  $0 \leq p \leq 1$ .

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \sigma^2 = np(1-p)$$

$\text{Bin}(n, p)$  is approximately  $\text{Po}(np)$  for small  $p$ .

First success distribution

The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability  $p$  is  $\text{Fs}(p)$ ,  $0 \leq p \leq 1$ .

$$p(k) = p(1-p)^{k-1}, k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution

The number of events occurring in a fixed period of time  $t$  if these events occur with a known average rate  $\kappa$  and independently of the time since the last event is  $\text{Po}(\lambda)$ ,  $\lambda = t\kappa$ .

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$

$$\mu = \lambda, \sigma^2 = \lambda$$

2.8.2 Continuous distributions

Uniform distribution

If the probability density function is constant between  $a$  and  $b$  and 0 elsewhere it is  $\text{U}(a, b)$ ,  $a < b$ .

$$f(x) = \begin{cases} \frac{1}{b-a} & a < x < b \\ 0 & \text{otherwise} \end{cases}$$

$$\mu = \frac{a+b}{2}, \sigma^2 = \frac{(b-a)^2}{12}$$

Exponential distribution

The time between events in a Poisson process is  $\text{Exp}(\lambda)$ ,  $\lambda > 0$ .

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \geq 0 \\ 0 & x < 0 \end{cases}$$

$$\mu = \frac{1}{\lambda}, \sigma^2 = \frac{1}{\lambda^2}$$

Normal distribution

Most real random values with mean  $\mu$  and variance  $\sigma^2$  are well described by  $\mathcal{N}(\mu, \sigma^2)$ ,  $\sigma > 0$ .

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

If  $X_1 \sim \mathcal{N}(\mu_1, \sigma_1^2)$  and  $X_2 \sim \mathcal{N}(\mu_2, \sigma_2^2)$  then

$$aX_1 + bX_2 + c \sim \mathcal{N}(\mu_1 + \mu_2 + c, a^2\sigma_1^2 + b^2\sigma_2^2)$$

2.9 Markov chains

A *Markov chain* is a discrete random process with the property that the next state depends only on the current state. Let  $X_1, X_2, \dots$  be a sequence of random variables generated by the Markov process. Then there is a transition matrix  $\mathbf{P} = (p_{ij})$ , with  $p_{ij} = \Pr(X_n = i | X_{n-1} = j)$ , and  $\mathbf{p}^{(n)} = \mathbf{P}^n \mathbf{p}^{(0)}$  is the probability distribution for  $X_n$  (i.e.,  $p_i^{(n)} = \Pr(X_n = i)$ ), where  $\mathbf{p}^{(0)}$  is the initial distribution.

$\pi$  is a stationary distribution if  $\pi = \pi \mathbf{P}$ . If the Markov chain is *irreducible* (it is possible to get to any state from any state), then  $\pi_i = \frac{1}{\mathbb{E}(T_i)}$  where  $\mathbb{E}(T_i)$  is the expected time between two visits in state  $i$ .  $\pi_j / \pi_i$  is the expected number of visits in state  $j$  between two visits in state  $i$ .

For a connected, undirected and non-bipartite graph, where the transition probability is uniform among all neighbors,  $\pi_i$  is proportional to node  $i$ 's degree.

A Markov chain is *ergodic* if the asymptotic distribution is independent of the initial distribution. A finite Markov chain is ergodic iff it is irreducible and *aperiodic* (i.e., the gcd of cycle lengths is 1).  $\lim_{k \rightarrow \infty} \mathbf{P}^k = \mathbf{1}\pi$ .

A Markov chain is an A-chain if the states can be partitioned into two sets **A** and **G**, such that all states in **A** are absorbing ( $p_{ii} = 1$ ), and all states in **G** leads to an absorbing state in **A**. The probability for absorption in state  $i \in \mathbf{A}$ , when the initial state is  $j$ , is  $a_{ij} = p_{ij} + \sum_{k \in \mathbf{G}} a_{ik} p_{kj}$ . The expected time until absorption, when the initial state is  $i$ , is  $t_i = 1 + \sum_{k \in \mathbf{G}} p_{ki} t_k$ .

2.10 Misc  
Surreal Number

For real  $x < y$ : if contains integer, integer closest to 0. Otherwise the one with smallest denominator which is a power of 2.

| $L/R$        | 0                                         | *                            | $\uparrow$                            | $\downarrow$                              |
|--------------|-------------------------------------------|------------------------------|---------------------------------------|-------------------------------------------|
| 0            | $\ast = \{0 \mid 0\}$                     | $\uparrow = \{0 \mid \ast\}$ | $\uparrow \ast = \{0 \mid \uparrow\}$ | $\ast = \{0 \mid 0\}$                     |
| *            | $\downarrow = \{\ast \mid 0\}$            | $0 = \{\mid\}$               | $0 = \{\mid\}$                        | $\{\ast \mid \downarrow\}$                |
| $\uparrow$   | $\ast = \{0 \mid 0\}$                     | $\{\uparrow \mid \ast\}$     | $\uparrow \ast = \{0 \mid \uparrow\}$ | $\ast = \{0 \mid 0\}$                     |
| $\downarrow$ | $\downarrow \ast = \{\downarrow \mid 0\}$ | $0 = \{\mid\}$               | $0 = \{\mid\}$                        | $\downarrow \ast = \{\downarrow \mid 0\}$ |

Data structures (3)

```
Rope.cpp
Description: "ext/rope" contains a data structure that is a generalization of a string.
<bits/stdc++.h>, <ext/rope>, std::, _gnu_cxx:: 59d0df, 47 lines
// rope: push_back(x), insert(pos,x), erase(pos,x),
// replace(pos,x), substr(pos,x), copy(x), at[x], +=, mutable
const int maxn=1e5+10;
rope<char> *his[maxn];
int n,d[maxn];
inline void updata(int x){
```

```
while(x<=n) d[x]++, x+=x&-x;
}
inline int get(int x){
int res=0;
while(x) res+=d[x], x-=x&-x;
return res;
}
void deb(rope<char> s){
for(int i=0;i<s.length();i++)
cout<<s[i];puts("");
}
int main(){
cin>>n;
his[0]=new rope<char>();
for(int i=1;i<=n;i++){
his[i]=new rope<char>(*his[i-1]);
// deb(*his[i]);
char opt; cin>>opt;
if(opt=='T'){
cin>>opt;
his[i]->push_back(opt);
updata(i);
}else if(opt=='U'){
updata(i);
int x,l=1,r=i,mid,now=get(i);
cin>>x;
while(l<r){
mid=(l+r)>>1;
if(now-get(mid)>x) l=mid+1;
else r=mid;
}
his[i]=his[l-1];
}else if(opt=='Q'){
int x; cin>>x; --x;
putchar(his[i]->at(x));
putchar('\n');
}
}
return 0;
}
```

```
OrderStatisticTree.h
Description: A set (not multiset!) with support for finding the n'th element, and finding the index of an element (number of smaller elements). To get a map, change null_type.
Time: O(log N)
_gnu_pbds:: 4dcf35, 15 lines
#include <bits/extc++.h>

template<class T>
using Tree = tree<T, null_type, less<T>, rb_tree_tag, tree_order_statistics_node_update>;

void example() {
Tree<int> t, t2; t.insert(8);
auto it = t.insert(10).first;
assert(it == t.lower_bound(9));
assert(t.order_of_key(10) == 1);
assert(t.order_of_key(11) == 2);
assert(*t.find_by_order(0) == 8);
t.join(t2); // assuming T< T2 or T> T2, merge t2 into t
}
```

```
DSU.cpp
Description: Union find set.
0782da, 27 lines
struct DSU {
vector<int> sz, parent;
```

```
void make_set(int v) {
parent[v] = v;
sz[v] = 1;
}
int find_set(int v) {
if (v == parent[v])
return v;
return find_set(parent[v]);
}
void union_sets(int a, int b) {
a = find_set(a);
b = find_set(b);
if (a != b) {
if (sz[a] < sz[b])
swap(a, b);
parent[b] = a;
sz[a] += sz[b];
}
}
DSU(int n) {
parent.resize(n);
sz.resize(n);
for (int i = 0; i<n; i++) make_set(i);
}
```

```
HashTable.cpp
Description: A hash table. TODO: test if it's faster than e.g. cuckoo. 895b04, 16 lines
#define MOD 1000007
struct my_hash_table{
int h[MOD],nxt[SZ],vv[SZ],ec; ll fl[SZ];
int &operator[](ll v){
for(int e=h[v%MOD];e;e=nxt[e])if(fl[e]==v)return vv[e];
++ec,fl[ec]=v,nxt[ec]=h[v%MOD],h[v%MOD]=ec;
vv[ec]=0; return vv[ec];
}
bool count(int v){
for(int e=h[v%MOD];e;e=nxt[e])if(fl[e]==v)return true;
return false;
}
void clear() {
for(int i=1;i<=ec;++i) h[fl[i]%MOD]=0;ec=0;
}
};
```

```
LinkCutTree.cpp
Description: Link cut tree. 6d9141, 57 lines
int ch[SZ][2],fa[SZ],sum[SZ],vv[SZ],mx[SZ];
bool rev[SZ];
bool top(int x) {return !(ch[fa[x]][0]==x||ch[fa[x]][1]==x);}
void pd(int x) {
if(!rev[x]) return;
rev[x]=0;
rev[ch[x][0]]^=1;
rev[ch[x][1]]^=1;
swap(ch[x][0],ch[x][1]);
}
void upd(int x) {
sum[x]=sum[ch[x][0]]+sum[ch[x][1]]+vv[x];
if(vv[mx[ch[x][0]]]>vv[mx[ch[x][1]]]
&&vv[mx[ch[x][0]]]>vv[x]) mx[x]=mx[ch[x][0]];
else if(vv[mx[ch[x][1]]]>vv[x]) mx[x]=mx[ch[x][1]];
else mx[x]=x;
}
void rot(int x) {
if(top(x)) return;
int y=fa[x],c=ch[y][0]==x;
```

```
int f=fa[y];
if(!top(y)) ch[f][ch[f][1]==y]=x;
if(ch[x][c]) fa[ch[x][c]]=y;
ch[y][!c]=ch[x][c];
ch[x][c]=y; fa[x]=f; fa[y]=x;
upd(y); upd(x);
}
int ss[SZ],sn;
void splay(int x) {
    sn=0;
    for(int c=x;;c=fa[c]) {
        ss[++sn]=c;
        if(top(c)) break;
    }
    while(sn) pd(ss[sn--]);
    while(!top(x)) {
        int y=fa[x];
        if(!top(y)) {
            if(ch[fa[y]][0]==y^ch[y][0]==x) rot(x);
            else rot(y);
        }
        rot(x);
    }
}
void access(int x) {
    for(int c=0;x;c=x,x=fa[x]) splay(x), ch[x][1]=c, upd(x);
}
void makeroot(int x) {access(x); splay(x); rev[x]^=1;}
void link(int a,int b) {makeroot(a); fa[a]=b;}
void cut(int a,int b) {makeroot(a); access(b); splay(b); ch[b][0]=fa[a]=0;}
int findroot(int x) {
    access(x); splay(x);
    int lc=x;
    while(ch[lc][0]) lc=ch[lc][0];
    splay(lc); return lc;
}
int getrd(int a,int b) {makeroot(a); access(b); splay(b);
    return b;}
```

278193, 85 lines

Treap.cpp
 Description: Treap.

```
pd(x);
if(sz[ch[x][0]]>=s)
    b=x, split(ch[x][0],a,ch[x][0],s), upd(x);
else
    a=x, split(ch[x][1],ch[x][1],b,s-sz[ch[x][0]]-1), upd(x);
}
}
void merge(int& ad,int a,int b) {
    if(a&&b) {
        if((int)mtl1()% (sz[a]+sz[b])<sz[a]) {
            pd(ad=a); merge(ch[a][1],ch[a][1],b);
        }
        else {
            pd(ad=b); merge(ch[b][0],a,ch[b][0]);
        }
        upd(ad);
    }
    else ad=a|b;
}
void edit(int l,int r,ll v) {
    int a,b,c;
    split(root,a,b,l-1);
    split(b,b,c,r-l+1);
    tag[b]+=v;
    sum[b]+=v*sz[b];
    merge(a,a,b);
    merge(root,a,c);
}
ll query(int l,int r) {
    int a,b,c;
    split(root,a,b,l-1);
    split(b,b,c,r-l+1);
    ll ans=sum[b];
    merge(a,a,b);
    merge(root,a,c);
    return ans;
}
int n,ns[SZ];
int ins(int l,int r) {
    if(l>r) return 0;
    if(l==r) {
        int ad; addnode(ad,ns[l]); return ad;
    }
    int mid=(l+r)>>1,lc=ins(l,mid),rc=ins(mid+1,r);
    merge(lc,lc,rc); return lc;
}
void init() {
    addnode(root,0);
    merge(root,ins(1,n),root);
}
int main() {
    scanf("%d",&n);
    for(int i=1;i<=n;i++) scanf("%d",&ns[i]);
    init();
    while(n--) {
        int o,l,r,c;
        scanf("%d%d%d%d",&o,&l,&r,&c);
        if(!o) edit(l,r,c);
        else ++c,printf("%lld\n", (query(l,r)%c+c)%c);
    }
}
```

Sec1c7, 30 lines

LineContainer.h
 Description: Container where you can add lines of the form  $kx+m$ , and query maximum values at points  $x$ . Useful for dynamic programming (“convex hull trick”).
 Time:  $\mathcal{O}(\log N)$

struct Line {

```
mutable ll k, m, p;
bool operator<(const Line& o) const { return k < o.k; }
bool operator<(ll x) const { return p < x; }
};

struct LineContainer : multiset<Line, less<>> {
    // (for doubles, use inf = 1/.0, div(a,b) = a/b)
    static const ll inf = LLONG_MAX;
    ll div(ll a, ll b) { // floored division
        return a / b - ((a ^ b) < 0 && a % b); }
    bool isect(iterator x, iterator y) {
        if (y == end()) return x->p = inf, 0;
        if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
        else x->p = div(y->m - x->m, x->k - y->k);
        return x->p >= y->p;
    }
    void add(ll k, ll m) {
        auto z = insert({k, m, 0}), y = z++, x = y;
        while (isect(y, z)) z = erase(z);
        if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
        while ((y = x) != begin() && (--x)->p >= y->p)
            isect(x, erase(y));
    }
    ll query(ll x) {
        assert(!empty());
        auto l = *lower_bound(x);
        return l.k * x + l.m;
    }
};

RMQ.h
Description: Range Minimum Queries on an array. Returns min(V[a], V[a + 1], ... V[b - 1]) in constant time.
Usage: RMQ rmq(values);
rmq.query(inclusive, exclusive);
Time:  $\mathcal{O}(|V| \log |V| + Q)$ 
510c32, 16 lines

template<class T>
struct RMQ {
    vector<vector<T>> jmp;
    RMQ(const vector<T>& V) : jmp(1, V) {
        for (int pw = 1, k = 1; pw * 2 <= sz(V); pw *= 2, ++k) {
            jmp.emplace_back(sz(V) - pw * 2 + 1);
            rep(j, 0, sz(jmp[k]))
                jmp[k][j] = min(jmp[k - 1][j], jmp[k - 1][j + pw]);
        }
    }
    T query(int a, int b) {
        assert(a < b); // or return inf if a == b
        int dep = 31 - __builtin_clz(b - a);
        return min(jmp[dep][a], jmp[dep][b - (1 << dep)]);
    }
};

Numerical (4)

4.1 Polynomials and recurrences
PolyInterpolate.h
Description: Given  $n$  points  $(x[i], y[i])$ , computes an  $n-1$ -degree polynomial  $p$  that passes through them:  $p(x) = a[0] * x^0 + \dots + a[n - 1] * x^{n-1}$ . For numerical precision, pick  $x[k] = c * \cos(k / (n - 1) * \pi)$ ,  $k = 0 \dots n - 1$ .
Time:  $\mathcal{O}(n^2)$ 
vd = vector<double>
285367, 12 lines

vd interpolate(vd x, vd y, int n) {
    vd res(n), temp(n);
    rep(k, 0, n-1) rep(i, k+1, n)
        y[i] = (y[i] - y[k]) / (x[i] - x[k]);
}
```

```
double last = 0; temp[0] = 1;
rep(k,0,n) rep(i,0,n) {
    res[i] += y[k] * temp[i];
    swap(last, temp[i]);
    temp[i] -= last * x[k];
}
return res;
}
```

LinearRecurrence.cpp

**Description:** Linear recurrence toolkit.

ll = long long, std::420684, 60 lines

```
const int MOD=1e9+7;
ll qp(ll a,ll b) {
    ll x=1; a%=MOD;
    while(b) {
        if(b&1) x=x*a%MOD;
        a=a*a%MOD; b>>=1;
    }
    return x;
}
namespace linear_seq {
inline vector<int> BM(vector<int> x) {
    vector<int> ls,cur;
    int pn=0,lf,ld;
    for(int i=0;i<int(x.size());++i) {
        ll t=-x[i]%MOD;
        for(int j=0;j<int(cur.size());++j)
            t=(t+x[i-j-1]*(ll)cur[j])%MOD;
        if(!t) continue;
        if(!cur.size())
            {cur.resize(i+1); lf=i; ld=t; continue;}
        ll k=-t*qp(ld,MOD-2)%MOD;
        vector<int> c(i-lf-1); c.pb(-k);
        for(int j=0;j<int(ls.size());++j) c.pb(ls[j]*k%MOD);
        if(c.size()<cur.size()) c.resize(cur.size());
        for(int j=0;j<int(cur.size());++j)
            c[j]=(c[j]+cur[j])%MOD;
        if(i-lf+(int)ls.size()>=(int)cur.size())
            ls=cur,lf=i,ld=t;
        cur=c;
    }
    vector<int>&o=cur;
    for(int i=0;i<int(o.size());++i)
        o[i]=(o[i]%MOD+MOD)%MOD;
    return o;
}
int N; ll a[SZ],h[SZ],t_[SZ],s[SZ],t[SZ];
inline void mull(ll*p,ll*q) {
    for(int i=0;i<N+N;++i) t_[i]=0;
    for(int i=0;i<N;++i) if(p[i])
        for(int j=0;j<N;++j)
            t_[i+j]=(t_[i+j]+p[i]*q[j])%MOD;
    for(int i=N+N-1;i>=N;--i) if(t_[i])
        for(int j=N-1;~j;--j)
            t_[i-j-1]=(t_[i-j-1]+t_[i]*h[j])%MOD;
    for(int i=0;i<N;++i) p[i]=t_[i];
}
inline ll calc(ll K) {
    for(int i=N;~i;--i) s[i]=t[i]=0;
    s[0]=1; if(N!=1) t[1]=1; else t[0]=h[0];
    for(;K;mull(t,t),K>>=1) if(K&1) mull(s,t); ll su=0;
    for(int i=0;i<N;++i) su=(su+s[i]*a[i])%MOD;
    return (su%MOD+MOD)%MOD;
}
inline int gao(vector<int> x,ll n) {
    if(n<int(x.size())) return x[n];
    vector<int> v=BM(x); N=v.size(); if(!N) return 0;
    for(int i=0;i<N;++i) h[i]=v[i],a[i]=x[i];
}
```

```
return calc(n);
}
}
```

4.2 Optimization

GoldenSectionSearch.h  
**Description:** Finds the argument minimizing the function  $f$  in the interval  $[a,b]$  assuming  $f$  is unimodal on the interval, i.e. has only one local minimum and no local maximum. The maximum error in the result is  $\epsilon$  $\epsilon$ s. Works equally well for maximization with a small change in the code. See Ternary-Search.h in the Various chapter for a discrete version.  
**Usage:** double func(double x) { return 4+x+.3\*x\*x; }  
double xmin = gss(-1000,1000,func);  
**Time:**  $\mathcal{O}(\log((b-a)/\epsilon))$

31d45b, 14 lines

```
double gss(double a, double b, double (*f)(double)) {
    double r = (sqrt(5)-1)/2, eps = 1e-7;
    double x1 = b - r*(b-a), x2 = a + r*(b-a);
    double f1 = f(x1), f2 = f(x2);
    while (b-a > eps)
        if (f1 < f2) { //change to > to find maximum
            b = x2; x2 = x1; f2 = f1;
            x1 = b - r*(b-a); f1 = f(x1);
        } else {
            a = x1; x1 = x2; f1 = f2;
            x2 = a + r*(b-a); f2 = f(x2);
        }
    return a;
}
```

IntegrateAdaptive.h

**Description:** Fast integration using an adaptive Simpson's rule.  
**Usage:** double sphereVolume = quad(-1, 1, [](double x) {  
return quad(-1, 1, [&](double y) {  
return quad(-1, 1, [&](double z) {  
return x\*x + y\*y + z\*z < 1; }));});};};  
d = double50bf66, 14 lines

#define S(a,b) (f(a) + 4\*f((a+b) / 2) + f(b)) \* (b-a) / 6

```
template <class F>
d rec(F& f, d a, d b, d eps, d S) {
    d c = (a + b) / 2;
    d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;
    if (abs(T - S) <= 15 * eps || b - a < 1e-10)
        return T + (T - S) / 15;
    return rec(f, a, c, eps / 2, S1) + rec(f, c, b, eps / 2, S2);
}
template<class F>
d quad(d a, d b, F f, d eps = 1e-8) {
    return rec(f, a, b, eps, S(a, b));
}
```

Simplex.cpp

**Description:** Simplex. Maximize  $\sum_j c_j x_j$  subject to  $\sum_j a_{i,j} x_j \leq b_i$  and  $x_j \geq 0$ .

f6ff59, 27 lines

```
const int N=25;const double eps=1e-7;
int n,m,op,lid[N],cid[N];double a[N][N],b[N];
void pivot(int l,int e){
    int i,j;double r=a[l][e];a[l][e]=1;for(i=0;i<=m;i++)a[l][i]/=r;swap(lid[l],cid[e]);
    for(i=0;i<=n;i++)if(i!=l)for(r=a[i][e],a[i][e]=j=0;j<=m;j++)a[i][j]-=r*a[l][j];
}
int main(){
    int i,j,l,e;double t,tt;scanf("%d%d%d",&m,&n,&op);
    for(i=1;i<=m;i++)scanf("%lf",&a[0][i]),cid[i]=i;
}
```

```
for(i=1;i<=n;scanf("%lf",&a[i++][0]))for(j=1;j<=m;j++)scanf("%lf",&a[i][j]);
while(true){
    for(l=0,i=1,t=-eps;i<=n;i++)if(a[i][0]<t)t=a[l=i][0];if(!l)break;
    for(e=0,i=1,t=-eps;i<=m;i++)if(a[l][i]<t&&(!e||(rand()&1)))t=a[l][e=i];
    if(!e)puts("Infeasible"),exit(0);pivot(l,e);
}
while(true){
    for(e=0,i=1,t=eps;i<=m;i++)if(a[0][i]>t)t=a[0][e=i];if(!e)break;
    for(l=0,i=1;i<=n;i++)if(a[i][e]>eps&&((tt=a[i][0]/a[i][e])<t||!l))t=tt,l=i;
    if(!l)puts("Unbounded"),exit(0);pivot(l,e);
}
printf("%.10lf\n",-a[0][0]);
if(op){
    for(i=1;i<=n;i++)if(lid[i])b[lid[i]]=a[i][0];
    for(i=1;i<=m;i++)printf("%.10lf%c",b[i]," \n"[i==m]);
}
return 0;
}
```

4.3 Matrices

Determinant.h

**Description:** Calculates determinant of a matrix. Destroys the matrix.  
**Time:**  $\mathcal{O}(N^3)$

bd5cec, 15 lines

```
double det(vector<vector<double>>& a) {
    int n = sz(a); double res = 1;
    rep(i,0,n) {
        int b = i;
        rep(j,i+1,n) if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
        if (i != b) swap(a[i], a[b]), res *= -1;
        res *= a[i][i];
        if (res == 0) return 0;
        rep(j,i+1,n) {
            double v = a[j][i] / a[i][i];
            if (v != 0) rep(k,i+1,n) a[j][k] -= v * a[i][k];
        }
    }
    return res;
}
```

SolveLinear.h

**Description:** Solves  $A * x = b$ . If there are multiple solutions, an arbitrary one is returned. Returns rank, or -1 if no solutions. Data in  $A$  and  $b$  is lost.  
**Time:**  $\mathcal{O}(n^2m)$

vd = vector<double>daa98b, 37 lines

```
const double eps = 1e-12;

int solveLinear(vector<vd>& A, vd& b, vd& x) {
    int n = sz(A), m = sz(x), rank = 0, br, bc;
    if (n) assert(sz(A[0]) == m);
    vi col(m); iota(all(col), 0);

    rep(i,0,n) {
        double v, bv = 0;
        rep(r,i,n) rep(c,i,m)
            if ((v = fabs(A[r][c])) > bv)
                br = r, bc = c, bv = v;
        if (bv <= eps) {
            rep(j,i,n) if (fabs(b[j]) > eps) return -1;
            break;
        }
        swap(A[i], A[br]);
        swap(b[i], b[br]);
    }
}
```



```
swap(col[i], col[bc]);
rep(j,0,n) swap(A[j][i], A[j][bc]);
bv = 1/A[i][i];
rep(j,i+1,n) {
    double fac = A[j][i] * bv;
    b[j] -= fac * b[i];
    rep(k,i+1,m) A[j][k] -= fac*A[i][k];
}
rank++;
}

x.assign(m, 0);
for (int i = rank; i--;) {
    b[i] /= A[i][i];
    x[col[i]] = b[i];
    rep(j,0,i) b[j] -= A[j][i] * b[i];
}
return rank; // (multiple solutions if rank < m)
}
```

SolveLinear2.h

**Description:** To get all uniquely determined values of  $x$  back from SolveLinear, make the following changes:

```
"SolveLinear.h" 08e495, 7 lines

rep(j,0,n) if (j != i) // instead of rep(j,i+1,n)
// ... then at the end:
x.assign(m, undefined);
rep(i,0,rank) {
    rep(j,rank,m) if (fabs(A[i][j]) > eps) goto fail;
    x[col[i]] = b[i] / A[i][i];
fail:; }
```

SolveLinearBinary.h

**Description:** Solves  $Ax = b$  over  $\mathbb{F}_2$ . If there are multiple solutions, one is returned arbitrarily. Returns rank, or -1 if no solutions. Destroys  $A$  and  $b$ .

```
"Time:  $\mathcal{O}(n^2m)$ "
bs = bitset<1000> 26d73e, 32 lines

int solveLinear(vector<bs>& A, vi& b, bs& x, int m) {
    int n = sz(A), rank = 0, br;
    assert(m <= sz(x));
    vi col(m); iota(all(col), 0);
    rep(i,0,n) {
        for (br=i; br<n; ++br) if (A[br].any()) break;
        if (br == n) {
            rep(j,i,n) if(b[j]) return -1;
            break;
        }
        int bc = (int)A[br]._Find_next(i-1);
        swap(A[i], A[br]);
        swap(b[i], b[br]);
        swap(col[i], col[bc]);
        rep(j,0,n) if (A[j][i] != A[j][bc]) {
            A[j].flip(i); A[j].flip(bc);
        }
        rep(j,i+1,n) if (A[j][i]) {
            b[j] ^= b[i];
            A[j] ^= A[i];
        }
        rank++;
    }
}
```

```
x = bs();
for (int i = rank; i--;) {
    if (!b[i]) continue;
    x[col[i]] = 1;
    rep(j,0,i) b[j] ^= A[j][i];
}
return rank; // (multiple solutions if rank < m)
```

}  
**MatrixInverse.h**  
**Description:** Invert matrix  $A$ . Returns rank; result is stored in  $A$  unless singular (rank < n). Can easily be extended to prime moduli; for prime powers, repeatedly set  $A^{-1} = A^{-1}(2I - AA^{-1}) \pmod{p^k}$  where  $A^{-1}$  starts as the inverse of  $A \pmod{p}$ , and  $k$  is doubled in each step.  
**Time:**  $\mathcal{O}(n^3)$

```
ebfff6, 35 lines

int matInv(vector<vector<double>>& A) {
    int n = sz(A); vi col(n);
    vector<vector<double>> tmp(n, vector<double>(n));
    rep(i,0,n) tmp[i][i] = 1, col[i] = i;

    rep(i,0,n) {
        int r = i, c = i;
        rep(j,i,n) rep(k,i,n)
            if (fabs(A[j][k]) > fabs(A[r][c]))
                r = j, c = k;
        if (fabs(A[r][c]) < 1e-12) return i;
        A[i].swap(A[r]); tmp[i].swap(tmp[r]);
        rep(j,0,n)
            swap(A[j][i], A[j][c]), swap(tmp[j][i], tmp[j][c]);
        swap(col[i], col[c]);
        double v = A[i][i];
        rep(j,i+1,n) {
            double f = A[j][i] / v;
            A[j][i] = 0;
            rep(k,i+1,n) A[j][k] -= f*A[i][k];
            rep(k,0,n) tmp[j][k] -= f*tmp[i][k];
        }
        rep(j,i+1,n) A[i][j] /= v;
        rep(j,0,n) tmp[i][j] /= v;
        A[i][i] = 1;
    }

    for (int i = n-1; i > 0; --i) rep(j,0,i) {
        double v = A[j][i];
        rep(k,0,n) tmp[j][k] -= v*tmp[i][k];
    }

    rep(i,0,n) rep(j,0,n) A[col[i]][col[j]] = tmp[i][j];
    return n;
}
```

Tridiagonal.h

**Description:**  $x = \text{tridiagonal}(d, p, q, b)$  solves the equation system

$$\begin{pmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_{n-1} \end{pmatrix} = \begin{pmatrix} d_0 & p_0 & 0 & 0 & \cdots & 0 \\ q_0 & d_1 & p_1 & 0 & \cdots & 0 \\ 0 & q_1 & d_2 & p_2 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & q_{n-3} & d_{n-2} & p_{n-2} \\ 0 & 0 & \cdots & 0 & q_{n-2} & d_{n-1} \end{pmatrix} \begin{pmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_{n-1} \end{pmatrix}.$$

This is useful for solving problems on the type

$$a_i = b_i a_{i-1} + c_i a_{i+1} + d_i, \quad 1 \leq i \leq n,$$

where  $a_0, a_{n+1}, b_i, c_i$  and  $d_i$  are known.  $a$  can then be obtained from

$$\{a_i\} = \text{tridiagonal}(\{1, -1, -1, \dots, -1, 1\}, \{0, c_1, c_2, \dots, c_n\}, \{b_1, b_2, \dots, b_n, 0\}, \{a_0, d_1, d_2, \dots, d_n, a_{n+1}\}).$$

Fails if the solution is not unique.  
If  $|d_i| > |p_i| + |q_{i-1}|$  for all  $i$ , or  $|d_i| > |p_{i-1}| + |q_i|$ , or the matrix is positive definite, the algorithm is numerically stable and neither  $\text{tr}$  nor the check for  $\text{diag}[i] == 0$  is needed.

```
"Time:  $\mathcal{O}(N)$ "
T = double d284f8, 25 lines
```

```
vector<T> tridiagonal(vector<T> diag, const vector<T>& super,
    const vector<T>& sub, vector<T> b) {
    int n = sz(b); vi tr(n);
    rep(i,0,n-1) {
        if (abs(diag[i]) < 1e-9 * abs(super[i])) { // diag[i] == 0
            b[i+1] -= b[i] * diag[i+1] / super[i];
            if (i+2 < n) b[i+2] -= b[i] * sub[i+1] / super[i];
            diag[i+1] = sub[i]; tr[++i] = 1;
        } else {
            diag[i+1] -= super[i]*sub[i]/diag[i];
            b[i+1] -= b[i]*sub[i]/diag[i];
        }
    }
    for (int i = n; i--;) {
        if (tr[i]) {
            swap(b[i], b[i-1]);
            diag[i-1] = diag[i];
            b[i] /= super[i-1];
        } else {
            b[i] /= diag[i];
            if (i) b[i-1] -= b[i]*super[i-1];
        }
    }
    return b;
}
```

4.4 Fourier transforms

FastFourierTransform.h

**Description:**  $\text{fft}(a)$  computes  $\hat{f}(k) = \sum_a a[x] \exp(2\pi i \cdot kx/N)$  for all  $k$ .  $N$  must be a power of 2. Useful for convolution:  $\text{conv}(a, b) = c$ , where  $c[x] = \sum a[i]b[x-i]$ . For convolution of complex numbers or more than two vectors: FFT, multiply pointwise, divide by  $n$ , reverse(start+1, end), FFT back. Rounding is safe if  $(\sum a_i^2 + \sum b_i^2) \log_2 N < 9 \cdot 10^{14}$  (in practice  $10^{16}$ ; higher for random inputs). Otherwise, use NTT/FFTMod.  
**Time:**  $\mathcal{O}(N \log N)$  with  $N = |A| + |B|$  ( $\sim 1s$  for  $N = 2^{22}$ )

```
C = complex<double>, vd = vector<double> 22c834, 33 lines

void fft(vector<C>& a) {
    int n = sz(a), L = 31 - __builtin_clz(n);
    static vector<complex<long double>> R(2, 1);
    static vector<C> rt(2, 1); // (^ 10% faster if double)
    for (static int k = 2; k < n; k *= 2) {
        R.resize(n); rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k);
        rep(i,k,2*k) rt[i] = R[i] = i&1 ? R[i/2] * x : R[i/2];
    }
    vi rev(n);
    rep(i,0,n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
    rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int k = 1; k < n; k *= 2)
        for (int i = 0; i < n; i += 2 * k) rep(j,0,k) {
            C z = rt[j+k] * a[i+j+k]; // (25% faster if hand-rolled)
            a[i + j + k] = a[i + j] - z;
            a[i + j] += z;
        }
}

vd conv(const vd& a, const vd& b) {
    if (a.empty() || b.empty()) return {};
    vd res(sz(a) + sz(b) - 1);
    int L = 32 - __builtin_clz(sz(res)), n = 1 << L;
    vector<C> in(n), out(n);
    copy(all(a), begin(in));
    rep(i,0,sz(b)) in[i].imag(b[i]);
    fft(in);
    for (C& x : in) x *= x;
    rep(i,0,n) out[i] = in[-i & (n - 1)] - conj(in[i]);
    fft(out);
    rep(i,0,sz(res)) res[i] = imag(out[i]) / (4 * n);
    return res;
}
```

```
}

// ...
}
```

### FastFourierTransformMod.h

**Description:** Higher precision FFT, can be used for convolutions modulo arbitrary integers as long as  $N \log_2 N \cdot \text{mod} < 8.6 \cdot 10^{14}$  (in practice  $10^{16}$  or higher). Inputs must be in  $[0, \text{mod})$ .

**Time:**  $\mathcal{O}(N \log N)$ , where  $N = |A| + |B|$  (twice as slow as NTT or FFT)  
"FastFourierTransform.h", vi = vector<ll> c1f2f1, 21 lines

```
template<int M> vl convMod(const vl &a, const vl &b) {
    if (a.empty() || b.empty()) return {};
    vl res(sz(a) + sz(b) - 1);
    int B=32-__builtin_clz(sz(res)), n=1<<B, cut=int(sqrt(M));
    vector<C> L(n), R(n), outs(n), outl(n);
    rep(i,0,sz(a)) L[i] = C((int)a[i] / cut, (int)a[i] % cut);
    rep(i,0,sz(b)) R[i] = C((int)b[i] / cut, (int)b[i] % cut);
    fft(L), fft(R);
    rep(i,0,n) {
        int j = -i & (n - 1);
        outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
        outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / 1i;
    }
    fft(outl), fft(outs);
    rep(i,0,sz(res)) {
        ll av = ll(real(outl[i])+.5), cv = ll(imag(outs[i])+.5);
        ll bv = ll(imag(outl[i])+.5) + ll(real(outs[i])+.5);
        res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
    }
    return res;
}
```

### FastSubsetTransform.h

**Description:** Transform to a basis with fast convolutions of the form  $c[z] = \sum_{z=x \oplus y} a[x] \cdot b[y]$ , where  $\oplus$  is one of AND, OR, XOR. The size of  $a$  must be a power of two.

**Time:**  $\mathcal{O}(N \log N)$

```
void FST(vi& a, bool inv) {
    for (int n = sz(a), step = 1; step < n; step *= 2) {
        for (int i = 0; i < n; i += 2 * step) rep(j,i,i+step) {
            int &u = a[j], &v = a[j + step]; tie(u, v) =
                inv ? pii(v - u, u) : pii(v, u + v); // AND
                inv ? pii(v, u - v) : pii(u + v, u); // OR
                pii(u + v, u - v); // XOR
        }
    }
    if (inv) for (int& x : a) x /= sz(a); // XOR only
}

vi conv(vi a, vi b) {
    FST(a, 0); FST(b, 0);
    rep(i,0,sz(a)) a[i] *= b[i];
    FST(a, 1); return a;
}
```

## Number theory (5)

### 5.1 Modular arithmetic

#### ModularArithmetic.h

**Description:** Operators for modular arithmetic. You need to set mod to some number first and then you can use the structure.

```
"euclid.h" 35bfea, 18 lines

const ll mod = 17; // change to something else
struct Mod {
    ll x;
    Mod(ll xx) : x(xx) {}
    Mod operator+(Mod b) { return Mod((x + b.x) % mod); }
    Mod operator-(Mod b) { return Mod((x - b.x + mod) % mod); }
```

```
Mod operator*(Mod b) { return Mod((x * b.x) % mod); }
Mod operator/(Mod b) { return *this * invert(b); }
Mod invert(Mod a) {
    ll x, y, g = euclid(a.x, mod, x, y);
    assert(g == 1); return Mod((x + mod) % mod);
}
Mod operator^(ll e) {
    if (!e) return Mod(1);
    Mod r = *this ^ (e / 2); r = r * r;
    return e&1 ? *this * r : r;
}
};
```

#### ModInverse.h

**Description:** Pre-computation of modular inverses. Assumes  $\text{LIM} \leq \text{mod}$  and that mod is a prime.

```
6f684f, 3 lines

const ll mod = 10000000007, LIM = 200000;
ll* inv = new ll[LIM] - 1; inv[1] = 1;
rep(i,2,LIM) inv[i] = mod - (mod / i) * inv[mod % i] % mod;
```

#### ModPow.h

```
b83e45, 7 lines

const ll mod = 10000000007; // faster if const
ll modpow(ll b, ll e) {
    ll ans = 1;
    for (; e; b = b * b % mod, e /= 2)
        if (e & 1) ans = ans * b % mod;
    return ans;
}
```

#### ModLog.h

**Description:** Returns the smallest  $x > 0$  s.t.  $a^x = b \pmod m$ , or  $-1$  if no such  $x$  exists. modLog(a,1,m) can be used to calculate the order of  $a$ .

**Time:**  $\mathcal{O}(\sqrt{m})$

```
c040b8, 11 lines

ll modLog(ll a, ll b, ll m) {
    ll n = (ll) sqrt(m) + 1, e = 1, f = 1, j = 1;
    unordered_map<ll, ll> A;
    while (j <= n && (e = f * e * a % m) != b % m)
        A[e * b % m] = j++;
    if (e == b % m) return j;
    if (__gcd(m, e) == __gcd(m, b))
        rep(i,2,n+2) if (A.count(e = e * f % m))
            return n * i - A[e];
    return -1;
}
```

#### ModMulLL.h

**Description:** Calculate  $a \cdot b \bmod c$  (or  $a^b \bmod c$ ) for  $0 \leq a, b \leq c \leq 7.2 \cdot 10^{18}$ . Requires x87 80-bit floats (e.g. GCC). Only valid for  $< 2^{52} \approx 4.5 \cdot 10^{15}$  if long double is 64-bit.

**Time:**  $\mathcal{O}(1)$  for modmul,  $\mathcal{O}(\log b)$  for modpow

```
fabd71, 10 lines

ull = unsigned long long
ull modmul(ull a, ull b, ull M) {
    ll ret = a * b - M * ull(1.L / M * a * b);
    return ret + M * (ret < 0) - M * (ret >= (ll)M);
}
ull modpow(ull b, ull e, ull mod) {
    ull ans = 1;
    for (; e; b = modmul(b, b, mod), e /= 2)
        if (e & 1) ans = modmul(ans, b, mod);
    return ans;
}
```

#### ModSqrt.h

**Description:** Tonelli-Shanks algorithm for modular square roots. Finds  $x$  s.t.  $x^2 = a \pmod p$  ( $-x$  gives the other solution).

**Time:**  $\mathcal{O}(\log^2 p)$  worst case,  $\mathcal{O}(\log p)$  for most  $p$

```
"ModPow.h" 19a793, 24 lines

ll sqrt(ll a, ll p) {
    a %= p; if (a < 0) a += p;
    if (a == 0) return 0;
    assert(modpow(a, (p-1)/2, p) == 1); // else no solution
    if (p % 4 == 3) return modpow(a, (p+1)/4, p);
    // a^(n+3)/8 or 2^(n+3)/8 * 2^(n-1)/4 works if p % 8 == 5
    ll s = p - 1, n = 2;
    int r = 0, m;
    while (s % 2 == 0)
        ++r, s /= 2;
    while (modpow(n, (p - 1) / 2, p) != p - 1) ++n;
    ll x = modpow(a, (s + 1) / 2, p);
    ll b = modpow(a, s, p), g = modpow(n, s, p);
    for (; r = m) {
        ll t = b;
        for (m = 0; m < r && t != 1; ++m)
            t = t * t % p;
        if (m == 0) return x;
        ll gs = modpow(g, 1LL << (r - m - 1), p);
        g = gs * gs % p;
        x = x * gs % p;
        b = b * g % p;
    }
}
```

#### Euclidean.cpp

**Description:** Floor sum.

```
9d1ecd, 20 lines

ll = long long
ll flr(ll x, ll y) { // floor(x / y). need y > 0
    ll res = x / y;
    if (res * y > x) res -= 1;
    return res;
}
ll f(ll a, ll b, ll c, ll n) { // sum_{0<=i<=n} flr((ai+b)/c).
    need c > 0
    ll pls = (a / c) * n * (n + 1) / 2;
    a %= c;
    if (a < 0) pls -= n * (n + 1) / 2, a += c;
    pls += (n + 1) * (b / c);
    b %= c;
    if (b < 0) pls -= (n + 1), b += c;
    if (a == 0) {
        ll ans = flr(b, c) * (n + 1);
        return ans + pls;
    }
    ll m = flr(a * n + b, c);
    ll res = n * m - f(c, c - b - 1, a, m - 1);
    return res + pls;
}
```

## 5.2 Primality

#### FastEratosthenes.h

**Description:** Prime sieve for generating all primes smaller than LIM.

**Time:** LIM=1e9  $\approx$  1.5s

```
6b2912, 20 lines

const int LIM = 1e6;
bitset<LIM> isPrime;
vi eratosthenes() {
    const int S = (int)round(sqrt(LIM)), R = LIM / 2;
    vi pr = {2}, sieve(S+1); pr.reserve(int(LIM/log(LIM)*1.1));
    vector<pii> cp;
    for (int i = 3; i <= S; i += 2) if (!sieve[i]) {
        cp.push_back({i, i * i / 2});
        for (int j = i * i; j <= S; j += 2 * i) sieve[j] = 1;
    }
    for (int L = 1; L <= R; L += S) {
```

```
array<bool, S> block{};
for (auto &p, idx] : cp)
    for (int i=idx; i < S+L; idx = (i+=p)) block[i-L] = 1;
rep(i,0,min(S, R - L))
    if (!block[i]) pr.push_back((L + i) * 2 + 1);
}
for (int i : pr) isPrime[i] = 1;
return pr;
}
```

MillerRabin.h

**Description:** Deterministic Miller-Rabin primality test. Guaranteed to work for numbers up to  $7 \cdot 10^{18}$ ; for larger numbers, use Python and extend A randomly.  
**Time:** 7 times the complexity of  $a^b \bmod c$ .

```
"ModMulLL.h" 60dcd1, 12 lines
bool isPrime(ull n) {
    if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;
    ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022},
    s = __builtin_ctzll(n-1), d = n >> s;
    for (ull a : A) { // ^ count trailing zeroes
        ull p = modpow(a%n, d, n), i = s;
        while (p != 1 && p != n - 1 && a % n && i--)
            p = modmul(p, p, n);
        if (p != n-1 && i != s) return 0;
    }
    return 1;
}
```

Factor.h

**Description:** Pollard-rho randomized factorization algorithm. Returns prime factors of a number, in arbitrary order (e.g. 2299 -> {11, 19, 11}).  
**Time:**  $\mathcal{O}(n^{1/4})$ , less for numbers with small factors.

```
"ModMulLL.h", "MillerRabin.h" a33cf6, 18 lines
ull pollard(ull n) {
    auto f = [n](ull x) { return modmul(x, x, n) + 1; };
    ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
    while (t++ % 40 || __gcd(prd, n) == 1) {
        if (x == y) x = ++i, y = f(x);
        if ((q = modmul(prd, max(x,y) - min(x,y), n))) prd = q;
        x = f(x), y = f(f(y));
    }
    return __gcd(prd, n);
}
vector<ull> factor(ull n) {
    if (n == 1) return {};
    if (isPrime(n)) return {n};
    ull x = pollard(n);
    auto l = factor(x), r = factor(n / x);
    l.insert(l.end(), all(r));
    return l;
}
```

Min25SieveDFS.cpp

**Description:** Min25 sieve with DFS.

```
<bits/stdc++.h>, std::, us = unsigned long long, ll = long long 83c4ef, 105 lines
const us MOD=1e9+7;
#define SS 233333
ll n,c0[SS],c1[SS],b0[SS],b1[SS];
#define MU 6000000
int S,ps[MU/10],pn=0;
int U=MU;
#define V 40
unsigned fc[V][MU];
inline ull F(ull x,us g) {
    if(g<V&&x<U) return fc[g][x];
```

```
if (x<=1||ps[g]>x) return 0;
ull ans=((x>S)?b1[n/x]:c1[x])-c1[ps[g-1]]%MOD;
for (us j=g; j<=pn&&ps[j]*(ll)ps[j]<=x;++j) {
    ull cn=x/ps[j],ce=ps[j]*(ll)ps[j];
    for (us e=1;cn>ps[j];++e,cn/=ps[j],ce*=ps[j])
        ans+=F(cn,j+1)*(ps[j]^e)+(ps[j]^(e+1)),ans%=MOD;
}
return ans%MOD;
}
int mf[MU];
unsigned f[MU];
int main() {
    // n=1e11;
    cin>>n;
    S=sqrtl(n);
    U=min(U,int(S*30));
    for (int i=1;i<U;++i) f[i]=1;
    for (int i=2;i<U;++i) {
        if (!mf[i]) {
            f[i]=i^1; ps[++pn]=i; mf[i]=pn;
        }
        for (int j=1;j<=pn&&i*ps[j]<U;++j) {
            int w=i*ps[j]; mf[w]=j;
            if (i%ps[j]==0) {
                int p=ps[j];
                int o=w,c=0;
                while (o%p==0) o/=p,++c;
                f[w]=f[o]*(p^c);
                break;
            }
            else f[w]=f[i]*(ps[j]^1);
        }
    }
    for (int j=0;j<V;++j) {
        unsigned*fcj=fc[j];
        for (int i=2;i<U;++i) {
            if (mf[i]>=j) {
                fcj[i]=fcj[i-1]+f[i];
                (fcj[i]>=MOD)?(fcj[i]-=MOD):0;
            }
            else fcj[i]=fcj[i-1];
        }
    }
    pn=0;
    for (int i=1;i<=S;++i) {
        ll t=(n/i)%MOD; b0[i]=t;
        b1[i]=t*(t+1)/2%MOD; c0[i]=i;
        c1[i]=i*(ll)(i+1)/2%MOD;
    }
    for (int i=2;i<=S;++i) {
        if (c0[i]==c0[i-1]) continue; //not a prime
        ll x0=c0[i-1],x1=c1[i-1]%MOD,r=(ll)i*i; ps[++pn]=i;
        int u=min((ll)S,n/(i*(ll)i)),uu=min(u,S/i);
        for (int j=1;j<=uu;++j)
            (b1[j]==(b1[j*i]-x1)*i)%=MOD,
            b0[j]-=b0[j*i]-x0;
        ll t=n/i;
        if (t<=2147483647) {
            int tt=t;
            for (int j=uu+1;j<=u;++j)
                b1[j]-=(c1[tt/j]-x1)*i,
                b0[j]-=c0[tt/j]-x0;
        }
        else {
            for (int j=uu+1;j<=u;++j)
                (b1[j]==(c1[t/j]-x1)*i)%=MOD,
                b0[j]-=c0[t/j]-x0;
        }
    }
```

```
for (int j=S;j>=r;--j)
    c1[j]-=(c1[j/i]-x1)*i,
    c0[j]-=c0[j/i]-x0;
}
for (int i=1;i<=S;++i) {
    c1[i]=MOD;
    b1[i]=MOD;
    c1[i]-=c0[i];
    b1[i]-=b0[i];
    if (i>=2) c1[i]+=2;
    if (n>=2LL*i) b1[i]+=2;
    c1[i]=(c1[i]%MOD+MOD)%MOD;
    b1[i]=(b1[i]%MOD+MOD)%MOD;
}
ps[pn+1]=S+1;
vector<ll> u;
for (ll i=1;i<=n;) {
    ll d=n/i,r=n/d; i=r+1;
    u.pb(((1+F(d,1))%MOD+MOD)%MOD);
}
sort(u.begin(),u.end());
u.erase(unique(u.begin(),u.end()),u.end());
ll ans=0;
for (auto c:u) ans^=c;
cout<<ans<<"\n";
}
```

5.3 Divisibility

euclid.h

**Description:** Finds two integers  $x$  and  $y$ , such that  $ax + by = \gcd(a, b)$ . If you just need gcd, use the built in `__gcd` instead. If  $a$  and  $b$  are coprime, then  $x$  is the inverse of  $a \pmod b$ .

```
33ba8f, 5 lines
ll euclid(ll a, ll b, ll &x, ll &y) {
    if (!b) return x = 1, y = 0, a;
    ll d = euclid(b, a % b, y, x);
    return y -= a/b * x, d;
}
```

MagicEuclid.cpp

**Description:** Magic euclid. Calculate  $\sum_{x=1}^L A^x B^{(Px+R)/Q}$ .

```
<bits/stdc++.h>, std::, ll = __int128 73f8b3, 58 lines
const int mod=998244353;
long long p,q,r,l; int n;
inline int add(int x,int y){return x+y<mod?x+y:x+y-mod;}
struct matrix{
    int a[21][21];
    matrix() {memset(a,0,sizeof(a));}
    matrix operator +(const matrix &b){
        matrix c;
        for (int i=1;i<=n;i++)
            for (int j=1;j<=n;j++)
                c.a[i][j]=add(a[i][j],b.a[i][j]);
        return c;
    }
    matrix operator *(const matrix &b){
        matrix c;
        for (int i=1;i<=n;i++){
            for (int k=1;k<=n;k++){
                if (!a[i][k]) continue;
                for (int j=1;j<=n;j++)
                    c.a[i][j]=add(c.a[i][j],lll*a[i][k]*b.a[k][j]%mod);
            }
        }
        return c;
    }
}a,b,I,ey;
struct node{
```

```
matrix u,r,s;
node operator +(const node &k){
    return node{u*k.u,r*k.r,s+r*k.s*u};
}
node operator *(ll n){
    node x=node{I,I,ey},y=node{u,r,s};
    while(n){
        if(n&1) x=x+y;
        y=y+y,n>>=1;
    }
    return x;
}
}
ans;
inline node solve(ll p,ll q,ll b,ll l,node u,node r){
    if(!l) return node{I,I,ey};
    if(p>=q) return solve(p*q,q,b,l,u,u*(p/q)+r);
    ll cnt=(l*p+b)/q; if(!cnt) return r*1;
    return r*((q-b-1)/p)+u+solve(q,p,(q-b-1)%p,cnt-1,r,u)+r*(1-(
        cnt*q-b-1)/p);
}
int main(){
    cin>>p>>q>>r>>l>>n;
    for(int i=1;i<=n;i++){
        for(int j=1;j<=n;j++) cin>>a.a[i][j];
        for(int i=1;i<=n;i++){
            for(int j=1;j<=n;j++) cin>>b.a[i][j];
        }
        for(int i=1;i<=n;i++) I.a[i][i]=1;
        ans=node{b,I,ey}*r/q+solve(p,q,r*q,l,node{b,I,ey},node{I,a,
            a});
        for(int i=1;i<=n;i++){
            for(int j=1;j<=n;j++) cout<<ans.s.a[i][j]<<' ';
            cout<<endl;
        }
    }
}
```

CRT.h

**Description:** Chinese Remainder Theorem.

crt(a, m, b, n) computes  $x$  such that  $x \equiv a \pmod{m}$ ,  $x \equiv b \pmod{n}$ . If  $|a| < m$  and  $|b| < n$ ,  $x$  will obey  $0 \leq x < \text{lcm}(m,n)$ . Assumes  $mn < 2^{62}$ .

**Time:**  $\log(n)$

```
"euclid.h" 04d93a, 7 lines

ll crt(ll a, ll m, ll b, ll n) {
    if (n > m) swap(a, b), swap(m, n);
    ll x, y, g = euclid(m, n, x, y);
    assert((a - b) % g == 0); // else no solution
    x = (b - a) % n * x % n / g * m + a;
    return x < 0 ? x + m*n/g : x;
}
```

5.4 Fractions

ContinuedFractions.h

**Description:** Given  $N$  and a real number  $x \geq 0$ , finds the closest rational approximation  $p/q$  with  $p, q \leq N$ . It will obey  $|p/q - x| \leq 1/qN$ .

For consecutive convergents,  $p_{k+1}q_k - q_{k+1}p_k = (-1)^k$ . ( $p_k/q_k$  alternates between  $> x$  and  $< x$ .) If  $x$  is rational,  $y$  eventually becomes  $\infty$ ; if  $x$  is the root of a degree 2 polynomial the  $a$ 's eventually become cyclic.

**Time:**  $\mathcal{O}(\log N)$

```
dd6c5e, 21 lines

typedef double d; // for N ~ 1e7; long double for N ~ 1e9
pair<ll, ll> approximate(d x, ll N) {
    ll LP = 0, LQ = 1, P = 1, Q = 0, inf = LLONG_MAX; d y = x;
    for (;;) {
        ll lim = min(P ? (N-LP) / P : inf, Q ? (N-LQ) / Q : inf),
            a = (ll)floor(y), b = min(a, lim),
            NP = b*P + LP, NQ = b*Q + LQ;
        if (a > b) {
            // If b > a/2, we have a semi-convergent that gives us a
            // better approximation; if b = a/2, we *may* have one.
```

```
// Return {P, Q} here for a more canonical approximation.
return (abs(x - (d)NP / (d)NQ) < abs(x - (d)P / (d)Q)) ?
    make_pair(NP, NQ) : make_pair(P, Q);
}
if (abs(y = 1/(y - (d)a)) > 3*N) {
    return {NP, NQ};
}
LP = P; P = NP;
LQ = Q; Q = NQ;
}
```

FracBinarySearch.h

**Description:** Given  $f$  and  $N$ , finds the smallest fraction  $p/q \in [0, 1]$  such that  $f(p/q)$  is true, and  $p, q \leq N$ . You may want to throw an exception from  $f$  if it finds an exact solution, in which case  $N$  can be removed.

**Usage:** fracBS({}(Frac f) { return f.p>=3\*f.q; }, 10); // {1,3}

**Time:**  $\mathcal{O}(\log(N))$

```
27ab3e, 25 lines

struct Frac { ll p, q; };

template<class F>
Frac fracBS(F f, ll N) {
    bool dir = 1, A = 1, B = 1;
    Frac lo{0, 1}, hi{1, 1}; // Set hi to 1/0 to search (0, N]
    if (f(lo)) return lo;
    assert(f(hi));
    while (A || B) {
        ll adv = 0, step = 1; // move hi if dir, else lo
        for (int si = 0; step; (step *= 2) >= si) {
            adv += step;
            Frac mid{lo.p * adv + hi.p, lo.q * adv + hi.q};
            if (abs(mid.p) > N || mid.q > N || dir == !f(mid)) {
                adv -= step; si = 2;
            }
        }
        hi.p += lo.p * adv;
        hi.q += lo.q * adv;
        dir = !dir;
        swap(lo, hi);
        A = B; B = !adv;
    }
    return dir ? hi : lo;
}
```

FarreySequence.cpp

**Description:** Use the Farrey sequence to find fraction between a/b and c/d with minimum denominator.

```
<cstdio>, <algorithm>, std::, ll = long long, pair<ll,ll>P = d18c2a, 24 lines

ll a,b,c,d,t;
ll gcd(ll a,ll b){return b?gcd(b,a%b):a;}
P cal(ll a,ll b,ll c,ll d){
    ll x=a/b+1;
    if(x*d<c) return P(x,1);
    if(!a) return P(1,d/c+1);
    if(a<=b&&c<=d){
        P t=cal(d,c,b,a);
        swap(t.first,t.second);
        return t;
    }
    x=a/b;
    P t=cal(a-b*x,b,c-d*x,d);
    t.first+=t.second*x;
    return t;
}
int main(){
    while(~scanf("%lld%lld%lld%lld",&a,&b,&c,&d)){
        t=gcd(a,b),a/=t,b/=t;
        t=gcd(c,d),c/=t,d/=t;
```

```
P p=cal(a,b,c,d);
printf("%lld/%lld\n",p.first,p.second);
}
}
```

PellEquation.py

```
28 lines

# Pell equation solver
# TODO: figure out if it's working
n=int(raw_input())
j=1
while j*j<n:
    j+=1
if j*j==n:
    print j,1
if j*j>n:
    p=[0 for i in range(0,1001)]
    q=[0 for i in range(0,1001)]
    a=[0 for i in range(0,1001)]
    g=[0 for i in range(0,1001)]
    h=[0 for i in range(0,1001)]
    p[1]=q[0]=h[1]=1
    p[0]=q[1]=g[1]=0
    a[2]=j-1
    i=2
    while 1:
        g[i]=-g[i-1]+a[i]*h[i-1]
        h[i]=(n-g[i]*g[i])/h[i-1]
        a[i+1]=(g[i]+a[2])/h[i]
        p[i]=a[i]*p[i-1]+p[i-2]
        q[i]=a[i]*q[i-1]+q[i-2]
        if(p[i]*p[i]-n*q[i]*q[i]==1):
            print p[i],q[i]
            break
        i+=1
```

5.5 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), \quad b = k \cdot (2mn), \quad c = k \cdot (m^2 + n^2),$$

with  $m > n > 0$ ,  $k > 0$ ,  $m \perp n$ , and either  $m$  or  $n$  even.

Combinatorial (6)

6.1 Permutations

6.1.1 Factorial

IntPerm.h

**Description:** Permutation -> integer conversion. (Not order preserving.)

Integer -> permutation can use a lookup table.

**Time:**  $\mathcal{O}(n)$

```
044568, 6 lines

int permToInt(vi& v) {
    int use = 0, i = 0, r = 0;
    for(int x:v) r = r * ++i + __builtin_popcount(use & -(1<<x)),
        use |= 1 << x; // (note: minus, not ~!)
    return r;
}
```

6.1.2 Cycles

Let  $g_S(n)$  be the number of  $n$ -permutations whose cycle lengths all belong to the set  $S$ . Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp \left( \sum_{n \in S} \frac{x^n}{n} \right)$$

6.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

6.2 Partitions and subsets

6.2.1 Partition function

Number of ways of writing  $n$  as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \quad p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

|        |   |   |   |   |   |   |    |    |    |    |     |            |            |
|--------|---|---|---|---|---|---|----|----|----|----|-----|------------|------------|
| $n$    | 0 | 1 | 2 | 3 | 4 | 5 | 6  | 7  | 8  | 9  | 20  | 50         | 100        |
| $p(n)$ | 1 | 1 | 2 | 3 | 5 | 7 | 11 | 15 | 22 | 30 | 627 | $\sim 2e5$ | $\sim 2e8$ |

6.2.2 Lucas’ Theorem

Let  $n, m$  be non-negative integers and  $p$  a prime. Write  $n = n_k p^k + \dots + n_1 p + n_0$  and  $m = m_k p^k + \dots + m_1 p + m_0$ . Then  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ .

6.3 General purpose numbers

6.3.1 Bernoulli numbers

EGF of Bernoulli numbers is  $B(t) = \frac{t}{e^t - 1}$  (FFT-able).

$$B[0, \dots] = [1, -\tfrac{1}{2}, \tfrac{1}{6}, 0, -\tfrac{1}{30}, 0, \tfrac{1}{42}, \dots]$$

Sums of powers:

$$\sum_{i=1}^n n^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\begin{aligned} \sum_{i=m}^\infty f(i) &= \int_m^\infty f(x) dx - \sum_{k=1}^\infty \frac{B_k}{k!} f^{(k-1)}(m) \\ &\approx \int_m^\infty f(x) dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m)) \end{aligned}$$

6.3.2 Stirling numbers of the first kind

Number of permutations on  $n$  items with  $k$  cycles.

$$c(n, k) = c(n-1, k-1) + (n-1)c(n-1, k), \quad c(0, 0) = 1$$
$$\sum_{k=0}^n c(n, k) x^k = x(x+1) \dots (x+n-1)$$

$$c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1$$
$$c(n, 2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$$

PushRelabel MinCostFlowAnton

6.3.3 Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly  $k$  elements are greater than the previous element.  $k$   $j$ :s s.t.  $\pi(j) > \pi(j+1)$ ,  $k+1$   $j$ :s s.t.  $\pi(j) \geq j$ ,  $k$   $j$ :s s.t.  $\pi(j) > j$ .

$$E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$$

$$E(n, 0) = E(n, n-1) = 1$$

$$E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$$

6.3.4 Stirling numbers of the second kind

Partitions of  $n$  distinct elements into exactly  $k$  groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

6.3.5 Bell numbers

Total number of partitions of  $n$  distinct elements.  $B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$  For  $p$  prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

6.3.6 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \quad C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \quad C_{n+1} = \sum C_i C_{n-i}$$

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$$

- sub-diagonal monotone paths in an  $n \times n$  grid.
- strings with  $n$  pairs of parenthesis, correctly nested.
- binary trees with with  $n+1$  leaves (0 or 2 children).
- ordered trees with  $n+1$  vertices.
- ways a convex polygon with  $n+2$  sides can be cut into triangles by connecting vertices with straight lines.
- permutations of  $[n]$  with no 3-term increasing subseq.

Graph (7)

7.1 Network flow

PushRelabel.h

**Description:** Push-relabel using the highest label selection rule and the gap heuristic. Quite fast in practice. To obtain the actual flow, look at positive values only.

**Time:**  $\mathcal{O}(V^2\sqrt{E})$

```
0ae1d4, 46 lines
struct PushRelabel {
    struct Edge {
        int dest, back;
        ll f, c;
    };
    vector<vector<Edge*>> g;
    vector<ll> ec;
    vector<Edge*> cur;
    vector<vi> hs; vi H;
    PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}
    void addEdge(int s, int t, ll cap, ll rcap=0) {
        if (s == t) return;
        g[s].push_back({t, sz(g[t]), 0, cap});
        g[t].push_back({s, sz(g[s])-1, 0, rcap});
    }
    void addFlow(Edge& e, ll f) {
        Edge &back = g[e.dest][e.back];
        if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
        e.f += f; e.c -= f; ec[e.dest] += f;
        back.f -= f; back.c += f; ec[back.dest] -= f;
    }
    ll calc(int s, int t) {
        int v = sz(g); H[s] = v; ec[t] = 1;
        vi co(2*v); co[0] = v-1;
        rep(i, 0, v) cur[i] = g[i].data();
        for (Edge& e : g[s]) addFlow(e, e.c);

        for (int hi = 0;;) {
            while (hs[hi].empty()) if (!hi--) return -ec[s];
            int u = hs[hi].back(); hs[hi].pop_back();
            while (ec[u] > 0) // discharge u
                if (cur[u] == g[u].data() + sz(g[u])) {
                    H[u] = 1e9;
                    for (Edge& e : g[u]) if (e.c && H[u] > H[e.dest]+1)
                        H[u] = H[e.dest]+1, cur[u] = &e;
                    if (++co[H[u]], !--co[hi] && hi < v)
                        rep(i, 0, v) if (hi < H[i] && H[i] < v)
                            --co[H[i]], H[i] = v + 1;
                    hi = H[u];
                } else if (cur[u]->c && H[u] == H[cur[u]->dest]+1)
                    addFlow(*cur[u], min(ec[u], cur[u]->c));
                else ++cur[u];
        }
        bool leftOfMinCut(int a) { return H[a] >= sz(g); }
    };
};
```

MinCostFlowAnton.cpp

**Description:** Simple min cost flow.

```
9e97e2, 67 lines
struct Edge {
    int from, to, capacity, cost;
};
vector<vector<int*>> adj, cost, capacity;
const int INF = 1e9;
void shortest_paths(int n, int v0, vector<int*>& d, vector<int*>& p) {
    p.assign(n, INF);
    d[v0] = 0;
    vector<bool> inq(n, false);
    queue<int> q;
    q.push(v0);
    p.assign(n, -1);
    while (!q.empty()) {
```



```
int u = q.front();
q.pop();
inq[u] = false;
for (int v : adj[u]) {
    if (capacity[u][v] > 0 && d[v] > d[u] + cost[u][v]) {
        d[v] = d[u] + cost[u][v];
        p[v] = u;
        if (!inq[v]) {
            inq[v] = true;
            q.push(v);
        }
    }
}
}

int min_cost_flow(int N, vector<Edge> edges, int K, int s, int t) {
    adj.assign(N, vector<int>());
    cost.assign(N, vector<int>(N, 0));
    capacity.assign(N, vector<int>(N, 0));
    for (Edge e : edges) {
        adj[e.from].push_back(e.to);
        adj[e.to].push_back(e.from);
        cost[e.from][e.to] = e.cost;
        cost[e.to][e.from] = -e.cost;
        capacity[e.from][e.to] = e.capacity;
    }
    int flow = 0;
    int cost = 0;
    vector<int> d, p;
    while (flow < K) {
        shortest_paths(N, s, d, p);
        if (d[t] == INF) break;
        // find max flow on that path
        int f = K - flow;
        int cur = t;
        while (cur != s) {
            f = min(f, capacity[p[cur]][cur]);
            cur = p[cur];
        }
        // apply flow
        flow += f;
        cost += f * d[t];
        cur = t;
        while (cur != s) {
            capacity[p[cur]][cur] -= f;
            capacity[cur][p[cur]] += f;
            cur = p[cur];
        }
    }
    if (flow < K)
        return -1;
    else
        return cost;
}
```

Dinic.cpp

Description: Dinic.

14ea98, 40 lines

```
int N,S,T,M=1; //M=1 is important!!
int fst[SZ],vb[SZ],nxt[SZ],vc[SZ];
void ad_de_(int a,int b,int c)
{++M;nxt[M]=fst[a];fst[a]=M;vb[M]=b;vc[M]=c;}
void ad_de(int a,int b,int c){ad_de_(a,b,c);ad_de_(b,a,0);}
int d[SZ],ff[SZ];
bool bfs() {
    static int qs[SZ]; int h=0,t=0;
    memset(d,-1,sizeof(int)*(N+1));
```

```
qs[t++]=S; d[S]=0;
while(h<t) {
    int x=qs[h++];
    for esb(x,e,b) {
        if((~d[b])||!vc[e]) continue;
        d[b]=d[x]+1; qs[t++]=b;
    }
}
return d[T]!=-1;
}
ll dfs(int x,ll u) {
    if(x==T||!u) return u;
    ll f=0;
    for(int&e=fst[x];e=nxt[e]) {
        int b=vb[e];
        if(d[b]!=d[x]+1) continue;
        ll s=dfs(vb[e],min(u-f,(ll)vc[e]));
        f+=s; vc[e]-=s; vc[e^1]+=s;
        if(f==u) break;
    }
    if(!f) d[x]=-1;
    return f;
}
ll dinic() {
    ll ans=0;
    memcpy(ff,fst,sizeof(int)*(N+1));
    while(bfs())
        ans+=dfs(S,1e18),
        memcpy(fst,ff,sizeof(int)*(N+1));
    return ans;
}
```

DinicAnton.cpp

Description: Dinic.

b8d428, 75 lines

```
struct FlowEdge {
    int v, u;
    long long cap, flow = 0;
    FlowEdge(int v, int u, long long cap) : v(v), u(u), cap(cap) {}
};

struct Dinic {
    const long long flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector<vector<int>> adj;
    int n, m = 0;
    int s, t;
    vector<int> level, ptr;
    queue<int> q;
    Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
        ptr.resize(n);
    }
    void add_edge(int v, int u, long long cap) {
        edges.emplace_back(v, u, cap);
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
        adj[u].push_back(m + 1);
        m += 2;
    }
    bool bfs() {
        while (!q.empty()) {
            int v = q.front();
            q.pop();
            for (int id : adj[v]) {
                if (edges[id].cap - edges[id].flow < 1)
                    continue;
                if (level[edges[id].u] != -1)
```

```
                    continue;
                    level[edges[id].u] = level[v] + 1;
                    q.push(edges[id].u);
            }
        }
        return level[t] != -1;
    }
    long long dfs(int v, long long pushed) {
        if (pushed == 0)
            return 0;
        if (v == t)
            return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++) {
            int id = adj[v][cid];
            int u = edges[id].u;
            if (level[v] + 1 != level[u] || edges[id].cap - edges[id].flow < 1)
                continue;
            long long tr = dfs(u, min(pushed, edges[id].cap - edges[id].flow));
            if (tr == 0)
                continue;
            edges[id].flow += tr;
            edges[id ^ 1].flow -= tr;
            return tr;
        }
        return 0;
    }
    long long flow() {
        long long f = 0;
        while (true) {
            fill(level.begin(), level.end(), -1);
            level[s] = 0;
            q.push(s);
            if (!bfs())
                break;
            fill(ptr.begin(), ptr.end(), 0);
            while (long long pushed = dfs(s, flow_inf)) {
                f += pushed;
            }
        }
        return f;
    }
};
```

MinCut.h

Description: After running max-flow, the left side of a min-cut from s to t is given by all vertices reachable from s, only traversing edges with positive residual capacity.

GlobalMinCut.h

Description: Find a global minimum cut in an undirected graph, as represented by an adjacency matrix.

Time:  $\mathcal{O}(V^3)$

8b0e19, 21 lines

```
pair<int, vi> globalMinCut(vector<vi> mat) {
    pair<int, vi> best = {INT_MAX, {}};
    int n = sz(mat);
    vector<vi> co(n);
    rep(i,0,n) co[i] = {i};
    rep(ph,1,n) {
        vi w = mat[0];
        size_t s = 0, t = 0;
        rep(it,0,n-ph) { //  $\mathcal{O}(V^2) \rightarrow \mathcal{O}(E \log V)$  with prio. queue
            w[t] = INT_MIN;
            s = t, t = max_element(all(w)) - w.begin();
            rep(i,0,n) w[i] += mat[t][i];
        }
```

```

if(x<=n) return q.push_back(x);
for(int i=0; i<flower[x].size(); i++) q.push(flower[x][i]);
}
void set_st(int x,int b){
    st[x]=b;
    if(x<=n) return;
    for(int i=0; i<flower[x].size(); ++i) set_st(flower[x][i],b);
}
int get_pr(int b,int xr){
    int pr=find(flower[b].begin(),flower[b].end(),xr)-flower[b].begin();
    if(pr%2==1){
        reverse(flower[b].begin()+1,flower[b].end());
        return (int) flower[b].size()-pr;
    }
    else return pr;
}
void set_match(int u,int v){
    match[u]=g[u][v].v;
    if(u<=n) return;
    Edge e=g[u][v];
    int xr=flower_from[u][e.u],pr=get_pr(u,xr);
    for(int i=0; i<pr; ++i) set_match(flower[u][i],flower[u][i^1])
        ;
    set_match(xr,v);
    rotate(flower[u].begin(),flower[u].begin()+pr,flower[u].end())
        );
}
void augment(int u,int v){
    int xnv=st[match[u]];
    set_match(u,v);
    if(!xnv) return;
    set_match(xnv,st[pa[xnv]]);
    augment(st[pa[xnv]],xnv);
}
int get_lca(int u,int v){
    static int t=0;
    for(++t; u||v; swap(u,v)){
        if(u==0) continue;
        if(vis[u]==t) return u;
        vis[u]=t;
        u=st[match[u]];
        if(u) u=st[pa[u]];
    }
    return 0;
}
void add_blossom(int u,int lca,int v){
    int b=n+1;
    while(b<=n_x&&st[b]) ++b;
    if(b>n_x) ++n_x;
    lab[b]=0,S[b]=0;
    match[b]=match[lca];
    flower[b].clear();
    flower[b].push_back(lca);
    for(int x=u,y; x!=lca; x=st[pa[y]])
        flower[b].push_back(x),flower[b].push_back(y=st[match[x]]),
            q.push(y);
    reverse(flower[b].begin()+1,flower[b].end());
    for(int x=v,y; x!=lca; x=st[pa[y]])
        flower[b].push_back(x),flower[b].push_back(y=st[match[x]]),
            q.push(y);
    set_st(b,b);
    for(int x=1; x<=n_x; ++x) g[b][x].w=g[x][b].w=0;
    for(int x=1; x<=n; ++x) flower_from[b][x]=0;
    for(int i=0; i<flower[b].size(); ++i){
        int xs=flower[b][i];
        for(int x=1; x<=n_x; ++x)
            if(g[b][x].w==0||DIST(g[xs][x])<DIST(g[b][x]))
                g[b][x]=g[xs][x],g[x][b]=g[x][xs];
    }
}

```

```
        for(int x=1; x<=n; ++x)
            if(flower_from[xs][x]) flower_from[b][x]=xs;
    }
    set_slack(b);
}
void expand_blossom(int b)    // S[b] == 1
{
    for(int i=0; i<flower[b].size(); ++i)
        set_st(flower[b][i], flower[b][i]);
    int xr=flower_from[b][g[b][pa[b]].u], pr=get_pr(b, xr);
    for(int i=0; i<pr; i+=2){
        int xs=flower[b][i], xns=flower[b][i+1];
        pa[xs]=g[xns][xs].u;
        S[xs]=1, S[xns]=0;
        slack[xs]=0, set_slack(xns);
        q_push(xns);
    }
    S[xr]=1, pa[xr]=pa[b];
    for(int i=pr+1; i<flower[b].size(); ++i){
        int xs=flower[b][i];
        S[xs]=-1, set_slack(xs);
    }
    st[b]=0;
}
bool on_found_Edge(const Edge &e){
    int u=st[e.u], v=st[e.v];
    if(S[v]==-1){
        pa[v]=e.u, S[v]=1;
        int nu=st[match[v]];
        slack[v]=slack[nu]=0;
        S[nu]=0, q_push(nu);
    }
    else if(S[v]==0){
        int lca=get_lca(u, v);
        if(!lca) return augment(u, v), augment(v, u), 1;
        else add_blossom(u, lca, v);
    }
    return 0;
}
bool matching(){
    fill(S, S+n_x+1, -1), fill(slack, slack+n_x+1, 0);
    q.clear();
    for(int x=1; x<=n_x; ++x)
        if(st[x]==x&&!match[x]) pa[x]=0, S[x]=0, q_push(x);
    if(q.empty()) return 0;
    for(;;){
        while(q.size()){
            int u=q.front();
            q.pop_front();
            if(S[st[u]]==1) continue;
            for(int v=1; v<=n; ++v)
                if(g[u][v].w>0&&st[u]!=st[v]){
                    if(DIST(g[u][v])==0){
                        if(on_found_Edge(g[u][v])) return 1;
                    }
                    else update_slack(u, st[v]);
                }
        }
        int d=INF;
        for(int b=n+1; b<=n_x; ++b)
            if(st[b]==b&&S[b]==1) d=min(d, lab[b]/2);
        for(int x=1; x<=n_x; ++x)
            if(st[x]==x&&slack[x]){
                if(S[x]==-1) d=min(d, DIST(g[slack[x]][x]));
                else if(S[x]==0) d=min(d, DIST(g[slack[x]][x])/2);
            }
        for(int u=1; u<=n; ++u){
            if(S[st[u]]==0){
                if(lab[u]<=d) return 0;
            }
        }
    }
}
```

```
        lab[u]-=d;
    }
    else if(S[st[u]]==1) lab[u]+=d;
}
for(int b=n+1; b<=n_x; ++b)
    if(st[b]==b){
        if(S[st[b]]==0) lab[b]+=d*2;
        else if(S[st[b]]==1) lab[b]-=d*2;
    }
q.clear();
for(int x=1; x<=n_x; ++x)
    if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&DIST(g[slack[x]][x])==0)
        if(on_found_Edge(g[slack[x]][x])) return 1;
for(int b=n+1; b<=n_x; ++b)
    if(st[b]==b&&S[b]==1&&lab[b]==0) expand_blossom(b);
}
return 0;
}
pair<ll, int> weight_blossom(){
    fill(match, match+n+1, 0);
    n_x=n;
    int n_matches=0;
    ll tot_weight=0;
    for(int u=0; u<=n; ++u) st[u]=u, flower[u].clear();
    int w_max=0;
    for(int u=1; u<=n; ++u)
        for(int v=1; v<=n; ++v){
            flower_from[u][v]=(u==v?u:0);
            w_max=max(w_max, g[u][v].w);
        }
    for(int u=1; u<=n; ++u) lab[u]=w_max;
    while(matching()) ++n_matches;
    for(int u=1; u<=n; ++u)
        if(match[u]&&match[u]<u)
            tot_weight+=g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
}
int main(){
    cin>>n>>m;
    for(int u=1; u<=n; ++u)
        for(int v=1; v<=n; ++v)
            g[u][v]=Edge {u, v, 0};
    for(int i=0, u, v, w; i<m; ++i){
        cin>>u>>v>>w;
        g[u][v].w=g[v][u].w=w;
    }
    cout<<weight_blossom().first<<'\\n';
    for(int u=1; u<=n; ++u) cout<<match[u]<<' ';
}

hopcroftKarp.h
Description: Fast bipartite matching algorithm. Graph  $g$  should be a list of neighbors of the left partition, and  $btoa$  should be a vector full of -1's of the same size as the right partition. Returns the size of the matching.  $btoa[i]$  will be the match for vertex  $i$  on the right side, or  $-1$  if it's not matched.
Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);
Time:  $\mathcal{O}(\sqrt{VE})$ 
```

```
bool dfs(int a, int L, vector<vi>& g, vi& btoa, vi& A, vi& B) {
    if (A[a] != L) return 0;
    A[a] = -1;
    for (int b : g[a]) if (B[b] == L + 1) {
        B[b] = 0;
        if (btoa[b] == -1 || dfs(btoa[b], L + 1, g, btoa, A, B))
            return btoa[b] = a, 1;
    }
    return 0;
}
```

```
int hopcroftKarp(vector<vi>& g, vi& btoa) {
    int res = 0;
    vi A(g.size()), B(btoa.size()), cur, next;
    for (;;) {
        fill(all(A), 0);
        fill(all(B), 0);
        cur.clear();
        for (int a : btoa) if (a != -1) A[a] = -1;
        rep(a, 0, sz(g)) if (A[a] == 0) cur.push_back(a);
        for (int lay = 1;; lay++) {
            bool islast = 0;
            next.clear();
            for (int a : cur) for (int b : g[a]) {
                if (btoa[b] == -1) {
                    B[b] = lay;
                    islast = 1;
                }
                else if (btoa[b] != a && !B[b]) {
                    B[b] = lay;
                    next.push_back(btoa[b]);
                }
            }
            if (islast) break;
            if (next.empty()) return res;
            for (int a : next) A[a] = lay;
            cur.swap(next);
        }
        rep(a, 0, sz(g))
            res += dfs(a, 0, g, btoa, A, B);
    }
}
```

DFSMatching.h

**Description:** Simple bipartite matching algorithm. Graph  $g$  should be a list of neighbors of the left partition, and  $btoa$  should be a vector full of -1's of the same size as the right partition. Returns the size of the matching.  $btoa[i]$  will be the match for vertex  $i$  on the right side, or  $-1$  if it's not matched.

**Usage:** vi btoa(m, -1); dfsMatching(g, btoa);

**Time:**  $\mathcal{O}(VE)$

```
bool find(int j, vector<vi>& g, vi& btoa, vi& vis) {
    if (btoa[j] == -1) return 1;
    vis[j] = 1; int di = btoa[j];
    for (int e : g[di])
        if (!vis[e] && find(e, g, btoa, vis)) {
            btoa[e] = di;
            return 1;
        }
    return 0;
}
int dfsMatching(vector<vi>& g, vi& btoa) {
    vi vis;
    rep(i, 0, sz(g)) {
        vis.assign(sz(btoa), 0);
        for (int j : g[i])
            if (find(j, g, btoa, vis)) {
                btoa[j] = i;
                break;
            }
    }
    return sz(btoa) - (int)count(all(btoa), -1);
}
```

MinimumVertexCover.h

**Description:** Finds a minimum vertex cover in a bipartite graph. The size is the same as the size of a maximum matching, and the complement is a maximum independent set.

"DFSMatching.h"



```
vi cover(vector<vi>& g, int n, int m) {
    vi match(m, -1);
    int res = dfsMatching(g, match);
    vector<bool> lfound(n, true), seen(m);
    for (int it : match) if (it != -1) lfound[it] = false;
    vi q, cover;
    rep(i,0,n) if (lfound[i]) q.push_back(i);
    while (!q.empty()) {
        int i = q.back(); q.pop_back();
        lfound[i] = 1;
        for (int e : g[i]) if (!seen[e] && match[e] != -1) {
            seen[e] = true;
            q.push_back(match[e]);
        }
    }
    rep(i,0,n) if (!lfound[i]) cover.push_back(i);
    rep(i,0,m) if (seen[i]) cover.push_back(n+i);
    assert(sz(cover) == res);
    return cover;
}
```

**WeightedMatching.h**  
**Description:** Given a weighted bipartite graph, matches every node on the left with a node on the right such that no nodes are in two matchings and the sum of the edge weights is minimal. Takes cost[N][M], where cost[i][j] = cost for L[i] to be matched with R[j] and returns (min cost, match), where L[i] is matched with R[match[i]]. Negate costs for max cost. Requires  $N \leq M$ .  
**Time:**  $\mathcal{O}(N^2M)$

1e0fe9, 31 lines

```
pair<int, vi> hungarian(const vector<vi> &a) {
    if (a.empty()) return {0, {}};
    int n = sz(a) + 1, m = sz(a[0]) + 1;
    vi u(n), v(m), p(m), ans(n - 1);
    rep(i,1,n) {
        p[0] = i;
        int j0 = 0; // add "dummy" worker 0
        vi dist(m, INT_MAX), pre(m, -1);
        vector<bool> done(m + 1);
        do { // dijkstra
            done[j0] = true;
            int i0 = p[j0], j1, delta = INT_MAX;
            rep(j,1,m) if (!done[j]) {
                auto cur = a[i0 - 1][j - 1] - u[i0] - v[j];
                if (cur < dist[j]) dist[j] = cur, pre[j] = j0;
                if (dist[j] < delta) delta = dist[j], j1 = j;
            }
            rep(j,0,m) {
                if (done[j]) u[p[j]] += delta, v[j] -= delta;
                else dist[j] -= delta;
            }
            j0 = j1;
        } while (p[j0]);
        while (j0) { // update alternating path
            int j1 = pre[j0];
            p[j0] = p[j1], j0 = j1;
        }
    }
    rep(j,1,m) if (p[j]) ans[p[j] - 1] = j - 1;
    return {-v[0], ans}; // min cost
}
```

**GeneralMatching.h**  
**Description:** Matching for general graphs. Fails with probability  $N/mod$ .  
**Time:**  $\mathcal{O}(N^3)$

../numerical/MatrixInverse-mod.hcb1912, 40 lines

```
vector<pii> generalMatching(int N, vector<pii>& ed) {
    vector<vector<ll>> mat(N, vector<ll>(N)), A;
    for (pii pa : ed) {
```

```
int a = pa.first, b = pa.second, r = rand() % mod;
mat[a][b] = r, mat[b][a] = (mod - r) % mod;
}

int r = matInv(A = mat), M = 2*N - r, fi, fj;
assert(r % 2 == 0);

if (M != N) do {
    mat.resize(M, vector<ll>(M));
    rep(i,0,N) {
        mat[i].resize(M);
        rep(j,N,M) {
            int r = rand() % mod;
            mat[i][j] = r, mat[j][i] = (mod - r) % mod;
        }
    }
} while (matInv(A = mat) != M);

vi has(M, 1); vector<pii> ret;
rep(it,0,M/2) {
    rep(i,0,M) if (has[i])
        rep(j,i+1,M) if (A[i][j] && mat[i][j]) {
            fi = i; fj = j; goto done;
        }
    assert(0); done:
    if (fj < N) ret.emplace_back(fi, fj);
    has[fi] = has[fj] = 0;
    rep(sw,0,2) {
        ll a = modpow(A[fi][fj], mod-2);
        rep(i,0,M) if (has[i] && A[i][fj]) {
            ll b = A[i][fj] * a % mod;
            rep(j,0,M) A[i][j] = (A[i][j] - A[fi][j] * b) % mod;
        }
        swap(fi, fj);
    }
}
return ret;
}
```

**MinimumAverageCycle.cpp**  
**Description:** Find the directed cycle with minimum average.

d343b9, 19 lines

```
const int N=3005,M=10005;
const double inf=1e18;
int n,m,i,j,u[M],v[M];double w[M],f[N][N],ans=1e9,now,tmp;
inline void up(double&a,double b){a>b?(a=b):0;}
int main(){
    scanf("%d%d",&n,&m);
    for(i=1;i<=m;i++)scanf("%d%d%lf",&u[i],&v[i],&w[i]);
    for(i=1;i<=n;i++)for(j=1;j<=n;j++)f[i][j]=inf;
    for(i=0;i<n;i++)for(j=1;j<=m;j++)up(f[i+1][v[j]],f[i][u[j]]+w[j]);
    for(i=1;i<=n;i++)if(f[n][i]<inf/2){
        now=-1e9;
        for(j=0;j<n;j++)if(f[j][i]<inf/2){
            tmp=1.0*(f[n][i]-f[j][i])/(n-j);
            if(now<tmp)now=tmp;
        }
        up(ans,now);
    }
    printf("%.8f",ans);
}
```

**7.3 DFS algorithms**  
**DominatorTree.cpp**  
**Description:** Building dominator tree.

9b5715, 41 lines

```
<cstdio>
// Assuming S can reach all graph (flowgraph), dom tree with
source S
```

```
// dfn[x]: x's DFS order
// id[x]: x-th in DFS order
// gd[x]: list of childs for x-th in DFS order
// idom[x]: parent in dom tree for x-th in DFS order
// sd[x]: (semi-)dominator for x-th in DFS order
// id[idom[dfn[x]]]: closest dominator of x
const int N=5010,M=200010;
int n,m,i,x,y,q[N],ans;
int g1[N],g2[N],gd[N],v[M*3+N],nxt[M*3+N],ed;
int cnt,dfn[N],id[N],fa[N],f[N],mn[N],sd[N],idom[N];
void add(int*x,int x,int y){v[++ed]=y;nxt[ed]=g[x];g[x]=ed;}
int F(int x){
    if(f[x]==x)return x;
    int y=F(f[x]);
    if(sd[mn[x]]>sd[mn[f[x]]])mn[x]=mn[f[x]];
    return f[x]=y;
}

void dfs(int x){
    id[dfn[x]++]=x;
    for(int i=g1[x];i;i=nxt[i])if(!dfn[v[i]])dfs(v[i]),fa[dfn[v[i]]]=dfn[x];
}

void tarjan(int S){
    int i,j,k,x;
    for(cnt=0,i=1;i<=n;i++)gd[i]=dfn[i]=id[i]=fa[i]=idom[i]=0,f[i]=sd[i]=mn[i]=i;
    dfs(S);
    for(i=n;i>1;i--){
        for(j=g2[id[i]];j;j=nxt[j])F(k=dfn[v[j]]),sd[i]=sd[i]<sd[mn[k]]?sd[i]:sd[mn[k]];
        add(gd,sd[i],i);
        for(j=gd[f[i]]=x=fa[i];j;j=nxt[j])F(k=v[j]),idom[k]=sd[mn[k]]<x?mn[k]:x;
        gd[x]=0;
    }
    for(i=2;i<=n;add(gd,idom[i],i),i++)if(idom[i]!=sd[i])idom[i]=idom[idom[i]];
}

int main(){
    while(~scanf("%d%d",&n,&m)){
        for(ed=0,i=1;i<=n;i++)g1[i]=g2[i]=0;
        while(m--){scanf("%d%d",&x,&y),add(g1,x,y),add(g2,y,x);
            tarjan(1);
        }
    }
}
```

**SCC.h**  
**Description:** Finds strongly connected components in a directed graph. If vertices  $u, v$  belong to the same component, we can reach  $u$  from  $v$  and vice versa.  
**Usage:** scc(graph, [&](vi& v) { ... }) visits all components in reverse topological order. comp[i] holds the component index of a node (a component only has edges to components with lower index). ncomps will contain the number of components.  
**Time:**  $\mathcal{O}(E + V)$

76b5c9, 24 lines

```
vi val, comp, z, cont;
int Time, ncomps;
template<class G, class F> int dfs(int j, G& g, F& f) {
    int low = val[j] = ++Time, x; z.push_back(j);
    for (auto e : g[j]) if (comp[e] < 0)
        low = min(low, val[e] ?: dfs(e,g,f));

    if (low == val[j]) {
        do {
            x = z.back(); z.pop_back();
            comp[x] = ncomps;
            cont.push_back(x);
        } while (x != j);
    }
```

```
f(cont); cont.clear();
ncomps++;
}
return val[j] = low;
}
template<class G, class F> void scc(G& g, F f) {
    int n = sz(g);
    val.assign(n, 0); comp.assign(n, -1);
    Time = ncomps = 0;
    rep(i,0,n) if (comp[i] < 0) dfs(i, g, f);
}
```

BiconnectedComponents.h

**Description:** Finds all biconnected components in an undirected graph, and runs a callback for the edges in each. In a biconnected component there are at least two distinct paths between any two nodes. Note that a node can be in several components. An edge which is not in a component is a bridge, i.e., not part of any cycle.

**Usage:** int eid = 0; ed.resize(N);  
for each edge (a,b) {  
ed[a].emplace\_back(b, eid);  
ed[b].emplace\_back(a, eid++); }  
bicomps([&](const vi& edgelist) {...});

**Time:**  $\mathcal{O}(E + V)$

2965e5, 33 lines

```
vi num, st;
vector<vector<pii>> ed;
int Time;
template<class F>
int dfs(int at, int par, F& f) {
    int me = num[at] = ++Time, e, y, top = me;
    for (auto pa : ed[at]) if (pa.second != par) {
        tie(y, e) = pa;
        if (num[y]) {
            top = min(top, num[y]);
            if (num[y] < me)
                st.push_back(e);
        } else {
            int si = sz(st);
            int up = dfs(y, e, f);
            top = min(top, up);
            if (up == me) {
                st.push_back(e);
                f(vi(st.begin() + si, st.end()));
                st.resize(si);
            }
            else if (up < me) st.push_back(e);
            else { /* e is a bridge */ }
        }
    }
    return top;
}

template<class F>
void bicomps(F f) {
    num.assign(sz(ed), 0);
    rep(i,0,sz(ed)) if (!num[i]) dfs(i, -1, f);
}
```

EulerWalk.h

**Description:** Eulerian undirected/directed path/cycle algorithm. Input should be a vector of (dest, global edge index), where for undirected graphs, forward/backward edges have the same index. Returns a list of nodes in the Eulerian path/cycle with src at both start and end, or empty list if no cycle/path exists. To get edge indices back, add .second to s and ret.

**Time:**  $\mathcal{O}(V + E)$

780b64, 15 lines

```
vi eulerWalk(vector<vector<pii>>& gr, int nedges, int src=0){
    int n = sz(gr);
```

```
vi D(n), its(n), eu(nedges), ret, s = {src};
D[src]++; // to allow Euler paths, not just cycles
while (!s.empty()) {
    int x = s.back(), y, e, &it = its[x], end = sz(gr[x]);
    if (it == end){ ret.push_back(x); s.pop_back(); continue; }
    tie(y, e) = gr[x][it++];
    if (!eu[e]) {
        D[x]--, D[y]++;
        eu[e] = 1; s.push_back(y);
    }
}
for (int x : D) if (x < 0 || sz(ret) != nedges+1) return {};
return {ret.rbegin(), ret.rend()};
}
```

7.4 Coloring

ChordalGraphOrder.cpp

**Description:** Ordering for chordal graph.

66d260, 8 lines

```
int i,j,k,col[N],ans;bool vis[N];
int main(){
    for(i=n;i;i--){
        for(k=0,j=1;j<=n;j++)if(!vis[j]&&col[j]>=col[k])k=j; //
            optimize this
        for(vis[k]=1,j=g[k];j;j=nxt[j])if(!vis[v[j]])if(++col[v[j]]>ans)ans=col[v[j]];
    }
    printf("%d",ans+1);
}
```

EdgeColoring.h

**Description:** Given a simple, undirected graph with max degree  $D$ , computes a  $(D + 1)$ -coloring of the edges such that no neighboring edges share a color. ( $D$ -coloring is NP-hard, but can be done for bipartite graphs by repeated matchings of max-degree nodes.)

**Time:**  $\mathcal{O}(NM)$

e210e2, 31 lines

```
vi edgeColoring(int N, vector<pii> eds) {
    vi cc(N + 1), ret(sz(eds)), fan(N), free(N), loc;
    for (pii e : eds) ++cc[e.first], ++cc[e.second];
    int u, v, ncols = *max_element(all(cc)) + 1;
    vector<vi> adj(N, vi(ncols, -1));
    for (pii e : eds) {
        tie(u, v) = e;
        fan[0] = v;
        loc.assign(ncols, 0);
        int at = u, end = u, d, c = free[u], ind = 0, i = 0;
        while (d = free[v], !loc[d] && (v = adj[u][d]) != -1)
            loc[d] = ++ind, cc[ind] = d, fan[ind] = v;
        cc[loc[d]] = c;
        for (int cd = d; at != -1; cd ^= c ^ d, at = adj[at][cd])
            swap(adj[at][cd], adj[end = at][cd ^ c ^ d]);
        while (adj[fan[i]][d] != -1) {
            int left = fan[i], right = fan[++i], e = cc[i];
            adj[u][e] = left;
            adj[left][e] = u;
            adj[right][e] = -1;
            free[right] = e;
        }
        adj[u][d] = fan[i];
        adj[fan[i]][d] = u;
        for (int y : {fan[0], u, end})
            for (int& z = free[y] = 0; adj[y][z] != -1; z++);
    }
    rep(i,0,sz(eds))
        for (tie(u, v) = eds[i]; adj[u][ret[i]] != v;) ++ret[i];
    return ret;
}
```

7.5 Heuristics

MaximalCliques.h

**Description:** Runs a callback for all maximal cliques in a graph (given as a symmetric bitset matrix; self-edges not allowed). Callback is given a bitset representing the maximal clique.

**Time:**  $\mathcal{O}\left(3^{n/3}\right)$ , much faster for sparse graphs

b63b43, 11 lines

```
B = bitset<128>
template<class F>
void cliques(vector<B>& eds, F f, B P = ~B(), B X={}, B R={}) {
    if (!P.any()) { if (!X.any()) f(R); return; }
    auto q = (P | X)._Find_first();
    auto cands = P & ~eds[q];
    rep(i,0,sz(eds)) if (cands[i]) {
        R[i] = 1;
        cliques(eds, f, P & eds[i], X & eds[i], R);
        R[i] = P[i] = 0; X[i] = 1;
    }
}
```

7.6 Trees

CompressTree.h

**Description:** Given a rooted tree and a subset  $S$  of nodes, compute the minimal subtree that contains all the nodes by adding all (at most  $|S| - 1$ ) pairwise LCA's and compressing edges. Returns a list of (par, orig\_index) representing a tree rooted at 0. The root points to itself.

**Time:**  $\mathcal{O}(|S| \log |S|)$

83c9a2, 20 lines

```
"LCA.h", vpi = vector<pair<int, int>>
vpi compressTree(LCA& lca, const vi& subset) {
    static vi rev; rev.resize(sz(lca.time));
    vi li = subset, &T = lca.time;
    auto cmp = [&](int a, int b) { return T[a] < T[b]; };
    sort(all(li), cmp);
    int m = sz(li)-1;
    rep(i,0,m) {
        int a = li[i], b = li[i+1];
        li.push_back(lca.lca(a, b));
    }
    sort(all(li), cmp);
    li.erase(unique(all(li)), li.end());
    rep(i,0,sz(li)) rev[li[i]] = i;
    vpi ret = {pii(0, li[0])};
    rep(i,0,sz(li)-1) {
        int a = li[i], b = li[i+1];
        ret.emplace_back(rev[lca.lca(a, b)], b);
    }
    return ret;
}
```

DirectedMST.h

**Description:** Finds a minimum spanning tree/arborescence of a directed graph, given a root node. If no MST exists, returns -1.

**Time:**  $\mathcal{O}(E \log V)$

39e620, 60 lines

```
"../data-structures/UnionFindRollback.h"
struct Edge { int a, b; ll w; };
struct Node {
    Edge key;
    Node *l, *r;
    ll delta;
    void prop() {
        key.w += delta;
        if (l) l->delta += delta;
        if (r) r->delta += delta;
        delta = 0;
    }
    Edge top() { prop(); return key; }
};
```

```

Node *merge(Node *a, Node *b) {
    if (!a || !b) return a ? b;
    a->prop(), b->prop();
    if (a->key.w > b->key.w) swap(a, b);
    swap(a->l, (a->r = merge(b, a->r)));
    return a;
}

void pop(Node*& a) { a->prop(); a = merge(a->l, a->r); }

pair<ll, vi> dmst(int n, int r, vector<Edge*& g) {
    RollbackUF uf(n);
    vector<Node*> heap(n);
    for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e});
    ll res = 0;
    vi seen(n, -1), path(n), par(n);
    seen[r] = r;
    vector<Edge> Q(n), in(n, {-1,-1}), comp;
    deque<tuple<int, int, vector<Edge*>>> cycs;
    rep(s,0,n) {
        int u = s, qi = 0, w;
        while (seen[u] < 0) {
            if (!heap[u]) return {-1,{};};
            Edge e = heap[u]->top();
            heap[u]->delta -= e.w, pop(heap[u]);
            Q[qi] = e, path[qi++] = u, seen[u] = s;
            res += e.w, u = uf.find(e.a);
            if (seen[u] == s) {
                Node* cyc = 0;
                int end = qi, time = uf.time();
                do cyc = merge(cyc, heap[w = path[--qi]]);
                while (uf.join(u, w));
                u = uf.find(u), heap[u] = cyc, seen[u] = -1;
                cycs.push_front({u, time, {&Q[qi], &Q[end]}});
            }
        }
        rep(i,0,qi) in[uf.find(Q[i].b)] = Q[i];
    }

    for (auto& [u,t,comp] : cycs) { // restore sol (optional)
        uf.rollback(t);
        Edge inEdge = in[u];
        for (auto& e : comp) in[uf.find(e.b)] = e;
        in[uf.find(inEdge.b)] = inEdge;
    }
    rep(i,0,n) par[i] = in[i].a;
    return {res, par};
}

```

## 7.7 Math

### 7.7.1 Number of Spanning Trees

Create an  $N \times N$  matrix  $\text{mat}$ , and for each edge  $a \rightarrow b \in G$ , do  $\text{mat}[a][b]--$ ,  $\text{mat}[b][b]++$  (and  $\text{mat}[b][a]--$ ,  $\text{mat}[a][a]++$  if  $G$  is undirected). Remove the  $i$ th row and column and take the determinant; this yields the number of directed spanning trees rooted at  $i$  (if  $G$  is undirected, remove any row/column).

### 7.7.2 Erdős–Gallai theorem

A simple graph with node degrees  $d_1 \geq \dots \geq d_n$  exists iff  $d_1 + \dots + d_n$  is even and for every  $k = 1 \dots n$ ,

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k).$$

## Geometry (8)

### 2DGeometry.cpp

**Description:** Basics of 2D geometry

```

<bits/stdc++.h>, std::, ld = long double
83be95, 128 lines

const ld PI = acos(-1);
const ld EPS = 1e-12;
int sgn(ld x) {
    if (x > EPS) return 1;
    else if (x < -EPS) return -1;
    return 0;
}

struct point {
    ld x, y;
    point(ld x = 0, ld y = 0) : x(x), y(y) {}
    point operator+(const point& o) const {
        return point(x + o.x, y + o.y);
    }
    point operator-(const point& o) const {
        return point(x - o.x, y - o.y);
    }
    point operator*(const ld& o) const {
        return point(x * o, y * o);
    }
    point operator/(const ld& o) const {
        return point(x / o, y / o);
    }
    ld operator*(const point& o) const {
        return x * o.y - y * o.x;
    }
    ld operator^(const point& o) const {
        return x * o.x + y * o.y;
    }
    point rotate() const {
        return point(-y, x);
    }
    point unit() const {
        ld w = abs();
        return point(x / w, y / w);
    }
    int get() const {
        return sgn(y) == 1 || (sgn(y) == 0 && sgn(x) == -1);
    }
    ld abs() const {
        return sqrt(x * x + y * y);
    }
};

struct line {
    point p0, p1;
    line(point p0, point p1) : p0(p0), p1(p1) {}
    bool contain(point p) const {
        return sgn((p1 - p0) * (p - p0)) == 1;
    }
    line move(ld d) const {
        point delta = (p1 - p0).rotate().unit() * d;
        return line(p0 + delta, p1 + delta);
    }
};

struct pointID {
    point p;
    int id;
    pointID(point p, int id) : p(p), id(id) {}
    bool operator<(const pointID& o) const {
        return make_pair(p.x, p.y) < make_pair(o.p.x, o.p.y);
    }
};

```

```

bool compare_angle(point p0, point p1) {
    return p0.get() < p1.get() || (p0.get() == p1.get() && sgn(p0 * p1) > 0);
}

point line_intersection(point p0, point p1, point p2, point p3) {
    ld foo = (p0 - p2) * (p3 - p2);
    ld bar = (p3 - p2) * (p1 - p2);
    return (p0 * bar + p1 * foo) / (foo + bar);
}

point line_intersection(line l0, line l1) {
    return line_intersection(l0.p0, l0.p1, l1.p0, l1.p1);
}

bool parallel(line l0, line l1) {
    return sgn((l0.p1 - l0.p0) * (l1.p1 - l1.p0)) == 0;
}

bool same_direction(line l0, line l1) {
    return parallel(l0, l1) && sgn((l0.p1 - l0.p0) ^ (l1.p1 - l1.p0)) == 1;
}

bool operator<(const line& l0, const line& l1) {
    if (same_direction(l0, l1))
        return l1.contain(l0.p0);
    else
        return compare_angle(l0.p1 - l0.p0, l1.p1 - l1.p0);
}

bool check(line l0, line l1, line l2) {
    return l2.contain(line_intersection(l0, l1));
}

vector<line> halfplane_intersection(vector<line> l) {
    sort(l.begin(), l.end());
    deque<line> q;
    for (int i = 0; i < (int) l.size(); ++i) {
        if (i && same_direction(l[i], l[i - 1]))
            continue;
        while ((int) q.size() > 1 && !check(q[q.size() - 2], q[q.size() - 1], l[i]))
            q.pop_back();
        while ((int) q.size() > 1 && !check(q[l], q[0], l[i]))
            q.pop_front();
        q.push_back(l[i]);
    }
    while ((int) q.size() > 2 && !check(q[q.size() - 2], q[q.size() - 1], q[0]))
        q.pop_back();
    while ((int) q.size() > 2 && !check(q[l], q[0], q[q.size() - 1]))
        q.pop_front();
    if ((int) q.size() == 2)
        return vector<line>();
    vector<line> res;
    for (int i = 0; i < (int) q.size(); ++i)
        res.push_back(q[i]);
    return res;
}

vector<pointID> convex_hull(vector<pointID> pts) {
    sort(pts.begin(), pts.end());
    if (pts.size() <= 1) return pts;
    vector<pointID> ch;
    for (int rot = 0; rot < 2; ++rot) {
        int lim = ch.size();
        for (pointID p : pts) {
            while ((int) ch.size() > lim + 1 && (ch[ch.size() - 1].p - ch[ch.size() - 2].p) * (p.p - ch[ch.size() - 2].p) <= 0)
                ch.pop_back();
            ch.push_back(p);
        }
        ch.pop_back();
    }
}

```

```
reverse(pts.begin(), pts.end());
}
return ch;
}
```

## 8.1 Geometric primitives

**Point.h**  
**Description:** Class to handle points in the plane. T can be e.g. double or long long. (Avoid int.)

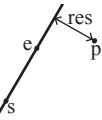
47ec0a, 28 lines

```
template <class T> int sgn(T x) { return (x > 0) - (x < 0); }
template<class T>
struct Point {
    typedef Point P;
    T x, y;
    explicit Point(T x=0, T y=0) : x(x), y(y) {}
    bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y); }
    bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }
    P operator+(P p) const { return P(x+p.x, y+p.y); }
    P operator-(P p) const { return P(x-p.x, y-p.y); }
    P operator*(T d) const { return P(x*d, y*d); }
    P operator/(T d) const { return P(x/d, y/d); }
    T dot(P p) const { return x*p.x + y*p.y; }
    T cross(P p) const { return x*p.y - y*p.x; }
    T cross(P a, P b) const { return (a-*this).cross(b-*this); }
    T dist2() const { return x*x + y*y; }
    double dist() const { return sqrt((double)dist2()); }
    // angle to x-axis in interval [-pi, pi]
    double angle() const { return atan2(y, x); }
    P unit() const { return *this/dist(); } // makes dist()=1
    P perp() const { return P(-y, x); } // rotates +90 degrees
    P normal() const { return perp().unit(); }
    // returns point rotated 'a' radians ccw around the origin
    P rotate(double a) const {
        return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
    friend ostream& operator<<(ostream& os, P p) {
        return os << "(" << p.x << ", " << p.y << ")"; }
};
```

lineDistance.h

Description:

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-negative distance. For Point3D, call .dist on the result of the cross product.



f6bf6b, 4 lines

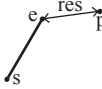
```
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
    return (double) (b-a).cross(p-a)/(b-a).dist();
}
```

SegmentDistance.h

Description:

Returns the shortest distance between point p and the line segment from point s to e.

Usage: Point<double> a, b(2,2), p(1,1); bool onSegment = segDist(a,b,p) < 1e-10;



ae751a, 5 lines

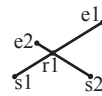
```
double segDist(P& s, P& e, P& p) {
    if (s==e) return (p-s).dist();
    auto d = (e-s).dist2(), t = min(d,max(.0, (p-s).dot(e-s)));
    return ((p-s)*d-(e-s)*t).dist()/d;
}
```

SegmentIntersection.h

Description:

If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.

Usage: vector<P> inter = segInter(s1,e1,s2,e2); if (sz(inter)==1) cout << "segments intersect at " << inter[0] << endl;



9d57f2, 13 lines

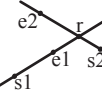
```
template<class P> vector<P> segInter(P a, P b, P c, P d) {
    auto oa = c.cross(d, a), ob = c.cross(d, b),
        oc = a.cross(b, c), od = a.cross(b, d);
    // Checks if intersection is single non-endpoint point.
    if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
        return {(a * ob - b * oa) / (ob - oa)};
    set<P> s;
    if (onSegment(c, d, a)) s.insert(a);
    if (onSegment(c, d, b)) s.insert(b);
    if (onSegment(a, b, c)) s.insert(c);
    if (onSegment(a, b, d)) s.insert(d);
    return {all(s)};
}
```

lineIntersection.h

Description:

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists {-1, (0,0)} is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.

Usage: auto res = lineInter(s1,e1,s2,e2); if (res.first == 1) cout << "intersection point at " << res.second << endl;



a01f81, 8 lines

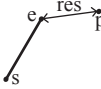
```
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
    auto d = (e1 - s1).cross(e2 - s2);
    if (d == 0) // if parallel
        return {(s1.cross(e1, s2) == 0), P(0, 0)};
    auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
    return {1, (s1 * p + e1 * q) / d};
}
```

sideOf.h

Description:

Returns where p is as seen from s towards e. 1/0/-1 ⇔ left/on line/right. If the optional argument eps is given 0 is returned if p is within distance eps from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

Usage: bool left = sideOf(p1,p2,q)==1;



3af81c, 9 lines

```
template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }

template<class P>
int sideOf(const P& s, const P& e, const P& p, double eps) {
    auto a = (e-s).cross(p-s);
    double l = (e-s).dist()*eps;
    return (a > l) - (a < -l);
}
```

OnSegment.h

Description:

Returns true iff p lies on the line segment from s to e. Use (segDist(s,e,p) <= epsilon) instead when using Point<double>.

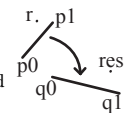
c597e8, 3 lines

```
template<class P> bool onSegment(P s, P e, P p) {
    return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
}
```

linearTransformation.h

Description:

Apply the linear transformation (translation, rotation and scaling) which takes line p0-p1 to line q0-q1 to point r.



096877, 5 lines

```
Point.h", P = Point<double>
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
    P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
    return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2();
}
```

Angle.h

Description:

A class for ordering angles (as represented by int points and a number of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

Usage: vector<Angle> v = {w[0], w[0].t360() ...}; // sorted int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) ++j; } // sweeps j such that (j-i) represents the number of positively oriented triangles with vertices at 0 and i

0f0602, 35 lines

```
struct Angle {
    int x, y;
    int t;
    Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
    Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
    int half() const {
        assert(x || y);
        return y < 0 || (y == 0 && x < 0);
    }
    Angle t90() const { return {-y, x, t + (half() && x >= 0)}; }
    Angle t180() const { return {-x, -y, t + half()}; }
    Angle t360() const { return {x, y, t + 1}; }
};
bool operator<(Angle a, Angle b) {
    // add a.dist2() and b.dist2() to also compare distances
    return make_tuple(a.t, a.half(), a.y * (ll)b.x <
        make_tuple(b.t, b.half(), a.x * (ll)b.y);
}

// Given two points, this calculates the smallest angle between
// them, i.e., the angle that covers the defined line segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
    if (b < a) swap(a, b);
    return (b < a.t180() ?
        make_pair(a, b) : make_pair(b, a.t360()));
}
Angle operator+(Angle a, Angle b) { // point a + vector b
    Angle r(a.x + b.x, a.y + b.y, a.t);
    if (a.t180() < r) r.t--;
    return r.t180() < a ? r.t360() : r;
}
Angle angleDiff(Angle a, Angle b) { // angle b - angle a
    int tu = b.t - a.t; a.t = b.t;
    return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)};
}
```

## 8.2 Circles

### CircleIntersection.h

**Description:** Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

```
"Point.h", P = Point<double> c64785, 10 lines

bool circleInter(P a,P b,double r1,double r2,pair<P, P>* out) {
    if (a == b) { assert(r1 != r2); return false; }
    P vec = b - a;
    double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
        p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
    if (sum*sum < d2 || dif*dif > d2) return false;
    P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2) / d2);
    *out = {mid + per, mid - per};
    return true;
}
```

### CircleTangents.h

**Description:** Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents – 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and .second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0.

```
"Point.h" b0153d, 13 lines

template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, double r2) {
    P d = c2 - c1;
    double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
    if (d2 == 0 || h2 < 0) return {};
    vector<pair<P, P>> out;
    for (double sign : {-1, 1}) {
        P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
        out.push_back({c1 + v * r1, c2 + v * r2});
    }
    if (h2 == 0) out.pop_back();
    return out;
}
```

### CirclePolygonIntersection.h

**Description:** Returns the area of the intersection of a circle with a ccw polygon.

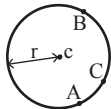
```
"Time: O(n)
".../content/geometry/Point.h", P = Point<double> 11e475, 18 lines

#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
    auto tri = [&](P p, P q) {
        auto r2 = r * r / 2;
        P d = q - p;
        auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2();
        auto det = a * a - b;
        if (det <= 0) return arg(p, q) * r2;
        auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
        if (t < 0 || 1 <= s) return arg(p, q) * r2;
        P u = p + d * s, v = p + d * t;
        return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
    };
    auto sum = 0.0;
    rep(i,0,sz(ps))
        sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
    return sum;
}
```

### circumcircle.h

#### Description:

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



```
"Point.h", P = Point<double> c98102, 8 lines

double ccRadius(const P& A, const P& B, const P& C) {
    return (B-A).dist()* (C-B).dist()* (A-C).dist() /
        abs((B-A).cross(C-A))/2;
}

P ccCenter(const P& A, const P& B, const P& C) {
    P b = C-A, c = B-A;
    return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
}
```

### MinimumEnclosingCircle.h

**Description:** Computes the minimum circle that encloses a set of points.

```
"Time: expected O(n)
"circumcircle.h" 09dd0a, 17 lines

pair<P, double> mec(vector<P> ps) {
    shuffle(all(ps), mt19937(time(0)));
    P o = ps[0];
    double r = 0, EPS = 1 + 1e-8;
    rep(i,0,sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
        o = ps[i], r = 0;
        rep(j,0,i) if ((o - ps[j]).dist() > r * EPS) {
            o = (ps[i] + ps[j]) / 2;
            r = (o - ps[i]).dist();
            rep(k,0,j) if ((o - ps[k]).dist() > r * EPS) {
                o = ccCenter(ps[i], ps[j], ps[k]);
                r = (o - ps[i]).dist();
            }
        }
    }
    return {o, r};
}
```

## 8.3 Polygons

### InsidePolygon.h

**Description:** Returns true if p lies within the polygon. If strict is true, it returns false for points on the boundary. The algorithm uses products in intermediate steps so watch out for overflow.

**Usage:** vector<P> v = {P{4,4}, P{1,2}, P{2,1}}; bool in = inPolygon(v, P{3, 3}, false);

```
"Time: O(n)
"Point.h", "OnSegment.h", "SegmentDistance.h" 2bf504, 11 lines

template<class P>
bool inPolygon(vector<P> &p, P a, bool strict = true) {
    int cnt = 0, n = sz(p);
    rep(i,0,n) {
        P q = p[(i + 1) % n];
        if (onSegment(p[i], q, a)) return !strict;
        //or: if (segDist(p[i], q, a) <= eps) return !strict;
        cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
    }
    return cnt;
}
```

### PolygonArea.h

**Description:** Returns twice the signed area of a polygon. Clockwise enumeration gives negative area. Watch out for overflow if using int as T!

```
"Point.h" f12300, 6 lines

template<class T>
T polygonArea2(vector<Point<T>>& v) {
    T a = v.back().cross(v[0]);
    rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
}
```

```
return a;
}
```

### PolygonCenter.h

**Description:** Returns the center of mass for a polygon.

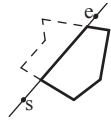
```
"Time: O(n)
"Point.h", P = Point<double> 0d0d84, 8 lines

P polygonCenter(const vector<P>& v) {
    P res(0, 0); double A = 0;
    for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
        res = res + (v[i] + v[j]) * v[j].cross(v[i]);
        A += v[j].cross(v[i]);
    }
    return res / A / 3;
}
```

### PolygonCut.h

#### Description:

Returns a vector with the vertices of a polygon with everything to the left of the line going from s to e cut away.



```
"Usage: vector<P> p = ...;
p = polygonCut(p, P(0,0), P(1,0));
"Point.h", "lineIntersection.h", P = Point<double> af7b53, 12 lines

vector<P> polygonCut(const vector<P>& poly, P s, P e) {
    vector<P> res;
    rep(i,0,sz(poly)) {
        P cur = poly[i], prev = i ? poly[i-1] : poly.back();
        bool side = s.cross(e, cur) < 0;
        if (side != (s.cross(e, prev) < 0))
            res.push_back(lineInter(s, e, cur, prev).second);
        if (side)
            res.push_back(cur);
    }
    return res;
}
```

### ConvexHull.h

#### Description:

Returns a vector of the points of the convex hull in counter-clockwise order. Points on the edge of the hull between two other points are not considered part of the hull.



```
"Time: O(n log n)
"Point.h", P = Point<ll> ec85f8, 12 lines

vector<P> convexHull(vector<P> pts) {
    if (sz(pts) <= 1) return pts;
    sort(all(pts));
    vector<P> h(sz(pts)+1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(all(pts)))
        for (P p : pts) {
            while (t >= s + 2 && h[t-2].cross(h[t-1], p) <= 0) t--;
            h[t++] = p;
        }
    return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};
}
```

### HullDiameter.h

**Description:** Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

```
"Time: O(n)
"Point.h", P = Point<ll> 5f726b, 11 lines

array<P, 2> hullDiameter(vector<P> S) {
    int n = sz(S), j = n < 2 ? 0 : 1;
    pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
    rep(i,0,j)
        for (;;) j = (j + 1) % n) {
            res = max(res, {{S[i] - S[j]}.dist2(), {S[i], S[j]}});
        }
```



```
        if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
            break;
    }
    return res.second;
}
```

PointInsideHull.h

**Description:** Determine whether a point t lies inside a convex hull (CCW order, with no collinear points). Returns true if point lies within the hull. If strict is true, points on the boundary aren't included.

**Time:**  $\mathcal{O}(\log N)$

"Point.h", "sideOf.h", "OnSegment.h", P = Point<ll> c74639, 12 lines

```
bool inHull(const vector<P>& l, P p, bool strict = true) {
    int a = 1, b = sz(l) - 1, r = !strict;
    if (sz(l) < 3) return r && onSegment(l[0], l.back(), p);
    if (sideOf(l[0], l[a], l[b]) > 0) swap(a, b);
    if (sideOf(l[0], l[a], p) >= r || sideOf(l[0], l[b], p) <= -r)
        return false;
    while (abs(a - b) > 1) {
        int c = (a + b) / 2;
        (sideOf(l[0], l[c], p) > 0 ? b : a) = c;
    }
    return sgn(l[a].cross(l[b], p)) < r;
}
```

LineHullIntersection.h

**Description:** Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:  $\bullet(-1, -1)$  if no collision,  $\bullet(i, -1)$  if touching the corner  $i$ ,  $\bullet(i, i)$  if along side  $(i, i + 1)$ ,  $\bullet(i, j)$  if crossing sides  $(i, i + 1)$  and  $(j, j + 1)$ . In the last case, if a corner  $i$  is crossed, this is treated as happening on side  $(i, i + 1)$ . The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

**Time:**  $\mathcal{O}(\log n)$

"Point.h" 7cf45b, 39 lines

```
#define cmp(i, j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
template <class P> int extrVertex(vector<P>& poly, P dir) {
    int n = sz(poly), lo = 0, hi = n;
    if (extr(0)) return 0;
    while (lo + 1 < hi) {
        int m = (lo + hi) / 2;
        if (extr(m)) return m;
        int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
        (ls < ms || (ls == ms && ls == cmp(lo, m)) ? hi : lo) = m;
    }
    return lo;
}

#define cmpL(i) sgn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
    int endA = extrVertex(poly, (a - b).perp());
    int endB = extrVertex(poly, (b - a).perp());
    if (cmpL(endA) < 0 || cmpL(endB) > 0)
        return {-1, -1};
    array<int, 2> res;
    rep(i, 0, 2) {
        int lo = endB, hi = endA, n = sz(poly);
        while ((lo + 1) % n != hi) {
            int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
            (cmpL(m) == cmpL(endB) ? lo : hi) = m;
        }
        res[i] = (lo + !cmpL(hi)) % n;
        swap(endA, endB);
    }
    if (res[0] == res[1]) return {res[0], -1};
}
```

```
    if (!cmpL(res[0]) && !cmpL(res[1]))
        switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
            case 0: return {res[0], res[0]};
            case 2: return {res[1], res[1]};
        }
    return res;
}
```

8.4 Misc. Point Set Problems

FastDelaunay.h

**Description:** Fast Delaunay triangulation. Each circumcircle contains none of the input points. There must be no duplicate points. If all points are on a line, no triangles will be returned. Should work for doubles as well, though there may be precision issues in 'circ'. Returns triangles in order  $\{t[0][0], t[0][1], t[0][2], t[1][0], \dots\}$ , all counter-clockwise.

**Time:**  $\mathcal{O}(n \log n)$

"Point.h", P = Point<ll>, Q = struct Quad\* 1846cb, 86 lines

```
typedef __int128_t ll1; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point

struct Quad {
    Q rot, o; P p = arb; bool mark;
    P& F() { return r()->p; }
    Q& r() { return rot->rot; }
    Q prev() { return rot->o->rot; }
    Q next() { return r()->prev(); }
} *H;

bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    ll1 p2 = p.dist2(), A = a.dist2()-p2,
        B = b.dist2()-p2, C = c.dist2()-p2;
    return p.cross(a,b)*C + p.cross(b,c)*A + p.cross(c,a)*B > 0;
}

Q makeEdge(P orig, P dest) {
    Q r = H ? H : new Quad(new Quad{new Quad{0}}});
    H = r->o; r->r()->r() = r;
    rep(i, 0, 4) r = r->rot, r->p = arb, r->o = i & 1 ? r : r->r();
    r->p = orig; r->F() = dest;
    return r;
}

void splice(Q a, Q b) {
    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
}

Q connect(Q a, Q b) {
    Q q = makeEdge(a->F(), b->p);
    splice(q, a->next());
    splice(q->r(), b);
    return q;
}
```

```
pair<Q, Q> rec(const vector<P>& s) {
    if (sz(s) <= 3) {
        Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
        if (sz(s) == 2) return {a, a->r()};
        splice(a->r(), b);
        auto side = s[0].cross(s[1], s[2]);
        Q c = side ? connect(b, a) : 0;
        return {side < 0 ? c->r() : a, side < 0 ? c : b->r()};
    }

#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
    Q A, B, ra, rb;
    int half = sz(s) / 2;
    tie(ra, A) = rec({all(s) - half});
    tie(B, rb) = rec({sz(s) - half + all(s)});
    while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
        (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
}
```

```
Q base = connect(B->r(), A);
if (A->p == ra->p) ra = base->r();
if (B->p == rb->p) rb = base;

#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) { \
        Q t = e->dir; \
        splice(e, e->prev()); \
        splice(e->r(), e->r()->prev()); \
        e->o = H; H = e; e = t; \
    }
for (;;) {
    DEL(LC, base->r(), o); DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
        base = connect(RC, base->r());
    else
        base = connect(base->r(), LC->r());
}
return {ra, rb};
}
```

```
vector<P> triangulate(vector<P> pts) {
    sort(all(pts)); assert(unique(all(pts)) == pts.end());
    if (sz(pts) < 2) return {};
    Q e = rec(pts).first;
    vector<Q> q = {e};
    int qi = 0;
    while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
    q.push_back(c->r()); c = c->next(); } while (c != e); }
    ADD; pts.clear();
    while (qi < sz(q) if (!(e = q[qi++])->mark) ADD;
    return pts;
}
```

MinkowskiSum.cpp

**Description:** A random(bad) minkowski sum template found online.

```
<bits/stdc++.h>, std::, ld = long double, ll = long long, co = complex<ll> abbc5b,
69 lines

#define X real()
#define Y imag()
#define F first
#define S second
ll ccw(co a, co b, co c) {
    return ((c-a)*conj(b-a)).Y;
}

int ar(co x) {
    if (x.Y>=0&&x.X<0) return 1;
    if (x.X>=0&&x.Y>0) return 2;
    if (x.Y<=0&&x.X>0) return 3;
    return 4;
}

bool cp(pair<co, pair<int, int> > p1, pair<co, pair<int, int> >
    p2) {
    if (ar(p1.F)!=ar(p2.F))
        return ar(p1.F)<ar(p2.F);
    assert((ccw({0, 0}, p1.F, p2.F)==0)==(ccw({0, 0}, p2.F, p1.F)
        ==0));
    if (ccw({0, 0}, p1.F, p2.F)==0)
        return p1.S>p2.S;
    return ccw({0, 0}, p2.F, p1.F)>0;
}

vector<co> minkowski(vector<co>& a, vector<co>& b) {
    int n=a.size();
    int m=b.size();
    if (n==0) return b;
    if (m==0) return a;
    if (n==1) {
```

```
vector<co> ret(m);
for (int i=0;i<m;i++) ret[i]=b[i]+a[0];
return ret;
}
if (m==1) {
vector<co> ret(n);
for (int i=0;i<n;i++) ret[i]=a[i]+b[0];
return ret;
}
vector<pair<co, pair<int, int> > > pp;
int f1=0,f2=0;
for (int i=0;i<n;i++) {
if (ccw(a[(i-1+n)%n], a[i], a[(i+1)%n])!=0) {
f1=i; break;
}
}
for (int i=0;i<n;i++)
pp.push_back({a[(i+1+f1)%n]-a[(i+f1)%n], {1, i}});
for (int i=0;i<m;i++) {
if (ccw(b[(i-1+m)%m], b[i], b[(i+1)%m])!=0) {
f2=i;break;
}
}
for (int i=0;i<m;i++)
pp.push_back({b[(i+1+f2)%m]-b[(i+f2)%m], {2, i}});
sort(pp.rbegin(), pp.rend(), cp);
co s={0, 0},ad={0, 0};
for (int i=0;i<(int)pp.size();i++) {
s+=pp[i].F;
if (pp[i].S.F!=pp[i+1].S.F) {
if (pp[i].S.F==1) ad=a[(pp[i].S.S+1+f1)%n]+b[(pp[i+1].S.S
+f2)%m];
else ad=b[(pp[i].S.S+1+f2)%m]+a[(pp[i+1].S.S+f1)%n];
ad-=s;break;
}
}
s=ad;
vector<co> ret(pp.size());
for (int i=0;i<(int)pp.size();i++) {
ret[i]=s;s+=pp[i].F;
}
return ret;
}
```

### 8.5 3D PolyhedronVolume.h

**Description:** Magic formula for the volume of a polyhedron. Faces should point outwards.

```
template<class V, class L>
double signedPolyVolume(const V& p, const L& trilst) {
double v = 0;
for (auto i : trilst) v += p[i.a].cross(p[i.b]).dot(p[i.c]);
return v / 6;
}
```

### Point3D.h

**Description:** Class to handle points in 3D space. T can be e.g. double or long long.

```
template<class T> struct Point3D {
typedef Point3D P;
typedef const P& R;
T x, y, z;
explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {}
bool operator<(R p) const {
return tie(x, y, z) < tie(p.x, p.y, p.z); }
bool operator==(R p) const {
return tie(x, y, z) == tie(p.x, p.y, p.z); }
```

```
P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
P operator*(T d) const { return P(x*d, y*d, z*d); }
P operator/(T d) const { return P(x/d, y/d, z/d); }
T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
P cross(R p) const {
return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
}
T dist2() const { return x*x + y*y + z*z; }
double dist() const { return sqrt((double)dist2()); }
//Azimuthal angle (longitude) to x-axis in interval [-pi, pi]
double phi() const { return atan2(y, x); }
//Zenith angle (latitude) to the z-axis in interval [0, pi]
double theta() const { return atan2(sqrt(x*x+y*y),z); }
P unit() const { return *this/(T)dist(); } //makes dist()==1
//returns unit vector normal to *this and p
P normal(P p) const { return cross(p).unit(); }
//rotated point rotated 'angle' radians ccw around axis
P rotate(double angle, P axis) const {
double s = sin(angle), c = cos(angle); P u = axis.unit();
return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
}
};
```

### 3dHull.h

**Description:** Computes all faces of the 3-dimension hull of a point set. \*No four points must be coplanar\*, or else random results will be returned. All faces will point outwards.

```
Time:  $\mathcal{O}(n^2)$ 
"Point3D.h", P3 = Point3D<double> b8fb93, 47 lines
struct PR {
void ins(int x) { (a == -1 ? a : b) = x; }
void rem(int x) { (a == x ? a : b) = -1; }
int cnt() { return (a != -1) + (b != -1); }
int a, b;
};
```

```
struct F { P3 q; int a, b, c; };
```

```
vector<F> hull3d(const vector<P3>& A) {
assert(sz(A) >= 4);
vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
#define E(x,y) E[f.x][f.y]
vector<F> FS;
auto mf = [&](int i, int j, int k, int l) {
P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
if (q.dot(A[l]) > q.dot(A[i]))
q = q * -1;
F f{q, i, j, k};
E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
FS.push_back(f);
};
rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)
mf(i, j, k, 6 - i - j - k);
```

```
rep(i,4,sz(A)) {
rep(j,0,sz(FS)) {
F f = FS[j];
if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
E(a,b).rem(f.c);
E(a,c).rem(f.b);
E(b,c).rem(f.a);
swap(FS[j--], FS.back());
FS.pop_back();
}
}
int nw = sz(FS);
rep(j,0,nw) {
F f = FS[j];
```

```
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.c);
C(a, b, c); C(a, c, b); C(b, c, a);
}
}
for (F& it : FS) if ((A[it.b] - A[it.a]).cross(
A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);
return FS;
};
```

### sphericalDistance.h

**Description:** Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 ( $\phi_1$ ) and f2 ( $\phi_2$ ) from x axis and zenith angles (latitude) t1 ( $\theta_1$ ) and t2 ( $\theta_2$ ) from z axis (0 = north pole). All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx\*radius is then the difference between the two points in the x direction and d\*radius is the total distance between the points.

```
611f07, 8 lines
double sphericalDistance(double f1, double t1,
double f2, double t2, double radius) {
double dx = sin(t2)*cos(f2) - sin(t1)*cos(f1);
double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
double dz = cos(t2) - cos(t1);
double d = sqrt(dx*dx + dy*dy + dz*dz);
return radius*2*asin(d/2);
}
```

### 3DConvexHull.cpp

**Description:** 3d convex hull in  $O(n^2 \log n)$ . First 4 points should not be coplanar.

```
f6e03c, 56 lines
struct v3 {
ld x,y,z;
v3() : x(0), y(0), z(0) {}
v3(ld xx, ld yy, ld zz) : x(xx), y(yy), z(zz) {}
};
v3 operator+(v3 a, v3 b) {
return v3(a.x+b.x, a.y+b.y, a.z+b.z);
}
v3 operator-(v3 a, v3 b) {
return v3(a.x-b.x, a.y-b.y, a.z-b.z);
}
v3 operator*(v3 a, v3 b) {
return v3(a.y*b.z-a.z*b.y,a.z*b.x-a.x*b.z,a.x*b.y-a.y*b.x);
}
ld operator^(v3 a, v3 b) {
return a.x*b.x+a.y*b.y+a.z*b.z;
}
bool onSameHalfSpace(v3 a, v3 b, v3 p1, v3 p2, v3 p3) {
v3 hlp = (p2-p1)*(p3-p1);
return (hlp^(a-p1))*(hlp^(b-p1))>=0;
}
int n;
v3 pts[1010];
map<pair<int,int>,int>,int> convHull;
void toggleHull(int i1, int i2, int i3, int ref) {
if (convHull.count({i1,i2,i3})) {
convHull.erase({i1,i2,i3});
} else convHull[{i1,i2,i3}]=ref;
}
void makeHull() {
convHull[{0,1,2}]=3;
convHull[{0,1,3}]=2;
convHull[{0,2,3}]=1;
convHull[{1,2,3}]=0;
for (int i=4;i<n;i++) {
vector<pair<pair<int,int>,pair<int,int>>> toChange;
for (auto hullFace : convHull) {
```

```
int i1=hullFace.F.F.F;
int i2=hullFace.F.F.S;
int i3=hullFace.F.S;
v3 pt1=pts[i1];
v3 pt2=pts[i2];
v3 pt3=pts[i3];
v3 ref=pts[hullFace.S];
if (!onSameHalfSpace(pts[i],ref,pt1,pt2,pt3)) {
    toChange.push_back({{i1,i2},{i,i3}});
    toChange.push_back({{i1,i3},{i,i2}});
    toChange.push_back({{i2,i3},{i,i1}});
    toChange.push_back({{i1,i2},{i3,i}});
}
}
for (auto diff : toChange) {
    toggleHull(diff.F.F,diff.F.S,diff.S.F, diff.S.S);
}
}
```

Strings (9)

Zfunc.h

**Description:** z[x] computes the length of the longest common prefix of s[i] and s, except z[0] = 0. (abacaba -> 0010301)

**Time:**  $\mathcal{O}(n)$

ee09e2, 12 lines

```
vi Z(const string& S) {
    vi z(sz(S));
    int l = -1, r = -1;
    rep(i,l,sz(S)) {
        z[i] = i >= r ? 0 : min(r - i, z[i - l]);
        while (i + z[i] < sz(S) && S[i + z[i]] == S[z[i]])
            z[i]++;
        if (i + z[i] > r)
            l = i, r = i + z[i];
    }
    return z;
}
```

Manacher.h

**Description:** For each position in a string, computes p[0][i] = half length of longest even palindrome around pos i, p[1][i] = longest odd (half rounded down).

**Time:**  $\mathcal{O}(N)$

e7ad79, 13 lines

```
array<vi, 2> manacher(const string& s) {
    int n = sz(s);
    array<vi,2> p = {vi(n+1), vi(n)};
    rep(z,0,2) for (int i=0,l=0,r=0; i < n; i++) {
        int t = r-i+!z;
        if (i<r) p[z][i] = min(t, p[z][l+t]);
        int L = i-p[z][i], R = i+p[z][i]-!z;
        while (L>=1 && R+1<n && s[L-1] == s[R+1])
            p[z][i]++, L--, R++;
        if (R>r) l=L, r=R;
    }
    return p;
}
```

MinRotation.h

**Description:** Finds the lexicographically smallest rotation of a string.

**Usage:** rotate(v.begin(), v.begin()+minRotation(v), v.end());

**Time:**  $\mathcal{O}(N)$

d07a42, 8 lines

```
int minRotation(string s) {
    int a=0, N=sz(s); s += s;
    rep(b,0,N) rep(k,0,N) {
```

```
        if (a+k == b || s[a+k] < s[b+k]) {b += max(0, k-1); break;}
        if (s[a+k] > s[b+k]) { a = b; break; }
    }
    return a;
}
```

LyndonRuns.cpp

**Description:** Compute lyndon and runs. TODO: check whether it works

<bits/stdc++.h>, std::, ull = uint64\_t3cb092, 70 lines

```
const ull base = 131;
const int N = 1e6 + 5, mod = 1e9 + 7;
ull pw[N], h[N], pw2[N], h2[N];
char s[N];
int n, ly[N], stk[N], tp, tot;
void init_hash(const char *s, int n) {
    pw[0] = 1;
    for(int i=1; i<=n; i++) {
        h[i] = (h[i-1]*base+s[i])%mod;
        pw[i] = pw[i-1]*base%mod;
    }
}
ull get(int l, int r) { return (h[r] - h[l-1]*pw[r-l+1]%mod+mod)%mod; }
int lcp(int x, int y) {
    int l = 1, r = n - max(x, y) + 1, ans = 0;
    while(l<=r) {
        int mid = (l+r)>>1;
        if(get(x, x+mid-1)==get(y, y+mid-1)) l = mid + 1, ans = mid;
        else r = mid - 1;
    }
    return ans;
}
int lcs(int x, int y) {
    int l = 1, r = min(x, y), ans = 0;
    while(l<=r) {
        int mid = (l+r)>>1;
        if(get(x-mid+1, x)==get(y-mid+1, y)) l = mid + 1, ans = mid;
        else r = mid - 1;
    }
    return ans;
}
struct runs {
    int l, r, p;
    bool operator < (const runs &oth) {
        return l==oth.l ? (r==oth.r ? p<oth.p : r<oth.r) : l<oth.l;
    }
}r[N<<1];
void check(int x, int y) {
    int L = lcs(x-1, y-1), R = lcp(x, y);
    if(L+R>=y-x) r[++tot] = {x-L, y+R-1, y-x};
}
bool cmp(int x, int y) {
    int len = lcp(x, y);
    return s[x+len]<s[y+len];
}
void lyndon(bool op) {
    ly[n] = n, stk[0] = n + 1, stk[tp=1] = n;
    for(int i=n-1; i>=1; i--) {
        while(tp && cmp(i, stk[tp])==op) --tp;
        stk[++tp] = i;
        ly[i] = stk[tp-1] - 1;
    }
}
int main() {
    scanf("%s", s+1);
    n = strlen(s+1);
```

```
init_hash(s, n);
for(int op=0; op<=1; op++) {
    lyndon(op);
    for(int i=1; i<n; i++) check(i, ly[i]+1);
}
sort(r+1, r+tot+1);
int cnt = 0;
for(int i=1; i<=tot; i++) {
    if(r[i].l!=r[i-1].l||r[i].r!=r[i-1].r) r[++cnt] = r[i];
}
printf("%d\n", cnt);
for(int i=1; i<=cnt; i++) printf("%d %d %d\n", r[i].l, r[i].r, r[i].p);
return 0;
}
```

PalindromeTree.cpp

**Description:** Palindromic tree. Odd and even roots etc.

653931, 27 lines

```
struct PTree {
    int ch[SZ][26],len[SZ],fail[SZ],cnt[SZ],s[SZ],cl,an,lst;
    int addn(int l) {len[an]=l; return an++;}
    PTree() {
        cl=an=lst=0;
        memset(ch,0,sizeof(ch));
        addn(0); addn(-1);
        fail[0]=1; s[0]=-233;
    }
    int gfail(int x,int l) {
        while(s[l-len[x]-1]!=s[l]) x=fail[x];
        return x;
    }
    void add(int c) {
        s[++cl]=c;
        int cp=gfail(lst,cl);
        if(!ch[cp][c]) {
            int nn=addn(len[cp]+2);
            fail[nn]=ch[ghfail(fail[cp],cl)][c];
            ch[cp][c]=nn;
        }
        cnt[lst=ch[cp][c]]++;
    }
    void getcnt() {
        for(int i=an-1;i>=2;i--) cnt[fail[i]]+=cnt[i];
    }
}pt;
```

SuffixAM.cpp

**Description:** Suffix automaton. To build upon a Trie use BFS.

35c15b, 37 lines

```
int N,ch[SZ][26],fa[SZ],ml[SZ];
void init() {N=0; fa[0]=-1; ml[0]=0;}
int ins(int p,int c) {
    int x=ch[p][c],y=0;
    if(!x) {
        x=++N;ml[x]=ml[p]+1;y=1;
        for(;p!=-1&&!ch[p][c];p=fa[p]) ch[p][c]=x;
    }
    if(p==-1) fa[x]=0;
    else if(ml[ch[p][c]]==ml[p]+1) y?(fa[x]=ch[p][c]):0;
    else {
        int q=++N,cc=ch[p][c];ml[q]=ml[p]+1;
        for(int i=0;i<26;i++) ch[q][i]=ch[cc][i];
        fa[q]=fa[cc],fa[cc]=fa[x]=q;
        for(;p!=-1&&ch[p][c]==cc;p=fa[p]) ch[p][c]=q;
        if(!y) return q;
    }
    return x;
}
```



```
int qs[SZ],od[SZ]; //od[0]..od[N]
void topo() {
    for(int i=0;i<=N;i++) qs[i]=0;
    for(int i=0;i<=N;i++) qs[ml[i]]++;
    for(int i=1;i<=N;+i) qs[i]+=qs[i-1];
    for(int i=N;i>=0;i--) od[--qs[ml[i]]]=i;
}
char s[SZ]; int ri[SZ];
int main(){
    init(); scanf("%s",s+1); int w=0;
    for(int j=1;s[j];j++) ++ri[w=ins(w,s[j]-'a')];
    topo(); ll ans=0;
    for(int i=N;i;--i) {
        int x=od[i]; ri[fa[x]]+=ri[x];
        if(ri[x]!=1) ans=max(ans,(ll)ri[x]*ml[x]);
    }
    printf("%lld\n",ans);
}
```

### SuffixArray.cpp

**Description:** Suffix array, 0-indexed. lcp(a,b) equals minimum h between rk[a] and rk[b].

ebc8b2, 41 lines

```
int sa[SZ],rk[SZ],su[SZ],ork[SZ];
char s[SZ]; int n,h[SZ];
bool diff(int a,int b,int g)
{return ork[a]!=ork[b]||ork[a+g]!=ork[b+g];}
void build_sa() {
    int G=max(n,200);
    for(int i=0;i<=G;+i) su[i]=0;
    for(int i=0;i<n;+i) ++su[rk[i]=s[i]];
    for(int i=1;i<=G;+i) su[i]+=su[i-1];
    for(int i=0;i<n;+i) sa[--su[rk[i]]]=i;
    for(int g=1;g<=n;g<=<=1) {
        for(int i=0;i<=G;+i) su[i]=0;
        static int ts[SZ]; int tn=0;
        for(int i=n-1;i>=0;--i) {
            if(sa[i]>=g) ts[+tn]=sa[i]-g;
            ork[i]=rk[i];
        }
        for(int i=n-g;i<n;+i) ts[+tn]=i;
        for(int i=1;i<=tn;+i) ++su[rk[ts[i]]];
        for(int i=1;i<=G;+i) su[i]+=su[i-1];
        for(int i=1;i<=tn;+i)
            sa[--su[rk[ts[i]]]]=ts[i];
        int t=0;
        for(int i=0;i<n;+i) {
            if(i&&diff(sa[i-1],sa[i],g)) ++t;
            rk[sa[i]]=t;
        }
    }
    for(int i=0;i<n;+i) rk[sa[i]]=i;
    int g=0;
    for(int i=0;i<n;+i) {
        g=max(g-1,0);
        if(rk[i]==n-1) {
            h[rk[i]]=0;
            continue;
        }
        int A=i,B=sa[rk[i]+1];
        while(s[A+g]==s[B+g]) ++g;
        h[rk[i]]=g;
    }
}
```

### BasicSubstringStructure.cpp

**Description:** Basic substring structure. The triangle is partitioned into substairs and each group is of several identical substairs. Total perimeter of substairs is linear.

b57e91, 50 lines

```
template<int NN,int Z>
class SAM{
    static constexpr int N=NN*2;
public:
    int OO,tot,n,Lst,ch[N][Z],fa[N],len[N],pre[N],posl[N],posr[
        N],occ[N],lis[N],tcnt[N],tmp[N],_ffa[_lg(N)+1][N],ffa
        [N][_lg(N)+1]; vector<int> son[N];
    SAM(int OOO=0) : OO(OOO) { tot=Lst=1; }
    inline void push_back(int c){
        int np=++tot,p=Lst; Lst=np; occ[np]=1; len[np]=++n; pre
        [n]=np; posr[np]=n;
        while(p&&!ch[p][c]) ch[p][c]=np,p=fa[p];; if(!p) return
        fa[np]=1,void();
        int v=ch[p][c]; if(len[v]==len[p]+1) return fa[np]=v,
        void();
        int nv=++tot; memcpy(ch[nv],ch[v],sizeof(ch[nv])); fa[
        nv]=fa[v]; posr[nv]=posr[v]; len[nv]=len[p]+1; fa[
        v]=fa[np]=nv;
        while(p&&ch[p][c]==v) ch[p][c]=nv,p=fa[p];
    }
    inline void build(int oo=0){
        For(u,2,tot) _ffa[0][u]=fa[u];; For(i,1,_lg(n)) For(u
        ,2,tot) _ffa[i][u]=_ffa[i-1][_ffa[i-1][u]];
        For(u,1,tot) For(i,0,_lg(n)) ffa[u][i]=_ffa[i][u];
        For(u,2,tot) son[fa[u]].push_back(u);;
        function<void(int)> dfs; dfs = [&](int u) { for(int v:
        son[u]) dfs(v),posr[u]=max(posr[u],posr[v]); };
        if(oo) dfs(1);
        For(u,2,tot) posl[u]=posr[u]-len[u]+1;
        For(u,1,tot) tcnt[posl[u]]++;; For(i,1,n) tcnt[i]+=tcnt
        [i-1];; For(u,1,tot) tmp[tcnt[posl[u]]--]=u;; if(
        oo) reverse(tmp+1,tmp+1+tot);
        For(i,1,n) tcnt[i]=0;; For(i,1,tot) tcnt[len[tmp[i]
        ]]]++;; For(i,1,n) tcnt[i]+=tcnt[i-1];; Rev(i,tot
        ,1) lis[tcnt[len[tmp[i]]--]=tmp[i];
        Rev(i,tot,1) { int u=lis[i]; for(int v:son[u]) occ[u]+=
        occ[v]; }
    }
    inline int find(int l,int r) const { int u=pre[r]; Rev(i,
        _lg(n),0) if(len[ffa[u][i]]>=r-l+1) u=ffa[u][i];;
        return u; }
};
template<int NN,int Z>
class BasicStringStructure{
public:
    static constexpr int N=NN*2; struct Node { int u,r; };
    using pii = pair<int,int>; struct Stair { int posl,
        posr; vector<Node> X,Y; } st[N];
    int n,stcnt,num0[N],num1[N]; SAM<NN,Z> S0,S1;
    BasicStringStructure<NN,Z>() : S0(0), S1(1) {}
    inline void build(int _n,int* s){
        n=_n; For(i,1,n) S0.push_back(s[i]);; Rev(i,n,1) S1.
        push_back(s[i]);; fprintf(stderr,"T1:%ld\n",clock
        ()); S0.build(0); S1.build(1);
        stcnt=0; Rev(i,S0.tot,2){
            int u=S0.lis[i],l=S0.posl[u],r=S0.posr[u],l2=1+S0.
            len[u]-S0.len[S0.fa[u]]-1,v=0; For(c,0,Z-1) if
            (S0.ch[u][c]&&S0.occ[S0.ch[u][c]]==S0.occ[u])
            { v=S0.ch[u][c]; break; }
            if(!v) { num0[u]=++stcnt; st[stcnt].posl=1; st[
            stcnt].posr=r; st[stcnt].X.push_back(Node{u,l2
            }); }
            else { st[num0[u]=num0[v]].X.push_back(Node{u,l2});
            }
        }
    }
};
```

```

    }
    stcnt=0; Rev(i,S1.tot,2){
        int u=S1.lis[i],l=n-S1.posr[u]+1,r=n-S1.posl[u]+1,
        r2=r-(S1.len[u]-S1.len[S1.fa[u]])+1,v=0; For(c
        ,0,Z-1) if(S1.ch[u][c]&&S1.occ[S1.ch[u][c]]==
        S1.occ[u]) { v=S1.ch[u][c]; break; }
        if(!v) { num1[u]=++stcnt; assume(st[stcnt].posl==1
        &&st[stcnt].posr==r); st[stcnt].Y.push_back(
        Node{u,r2}); }
        else { st[num1[u]=num1[v]].Y.push_back(Node{u,r2});
        }
    }
    // For(i,1,stcnt){
    //     printf("Stair #%d (%d,%d): X = ",i,st[i].posl,st
    //         [i].posr);
    //     for(auto x:st[i].X) printf("%d[%d,%d] ",x.u,x.l,
    //         x.r);
    //     printf(" Y = "); for(auto y:st[i].Y) printf("%d
    //         [%d,%d] ",y.u,y.l,y.r);; puts("");
    // }
    }
};
```

## Various (10)

### 10.1 Intervals

#### IntervalContainer.h

**Description:** Add and remove intervals from a set of disjoint intervals. Will merge the added interval with any overlapping intervals in the set when adding. Intervals are [inclusive, exclusive).

**Time:**  $\mathcal{O}(\log N)$

edce47, 23 lines

```
set<pii>::iterator addInterval(set<pii>& is, int L, int R) {
    if (L == R) return is.end();
    auto it = is.lower_bound({L, R}), before = it;
    while (it != is.end() && it->first <= R) {
        R = max(R, it->second);
        before = it = is.erase(it);
    }
    if (it != is.begin() && (--it)->second >= L) {
        L = min(L, it->first);
        R = max(R, it->second);
        is.erase(it);
    }
    return is.insert(before, {L,R});
}
```

```
void removeInterval(set<pii>& is, int L, int R) {
    if (L == R) return;
    auto it = addInterval(is, L, R);
    auto r2 = it->second;
    if (it->first == L) is.erase(it);
    else (int&)it->second = L;
    if (R != r2) is.emplace(R, r2);
}
```

#### IntervalCover.h

**Description:** Compute indices of smallest set of intervals covering another interval. Intervals should be [inclusive, exclusive). To support [inclusive, inclusive], change (A) to add || R.empty(). Returns empty set on failure (or if G is empty).

**Time:**  $\mathcal{O}(N \log N)$

9e9d8d, 19 lines

```
template<class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
    vi S(sz(I)), R;
    iota(all(S), 0);
```

```
sort(all(S), [&](int a, int b) { return I[a] < I[b]; });
T cur = G.first;
int at = 0;
while (cur < G.second) { //(A)
    pair<T, int> mx = make_pair(cur, -1);
    while (at < sz(I) && I[S[at]].first <= cur) {
        mx = max(mx, make_pair(I[S[at]].second, S[at]));
        at++;
    }
    if (mx.second == -1) return {};
    cur = mx.first;
    R.push_back(mx.second);
}
return R;
}
```

10.2 Misc. algorithms

ExpressionEvaluation.cpp

Description: Expression evaluation sample.

2b0f18, 78 lines

```
bool delim(char c) {return c == ' ';}
bool is_op(char c) {
    return c == '+' || c == '-' || c == '*' || c == '/';
}
bool is_unary(char c) {
    return c == '+' || c == '-';
}
int priority (char op) {
    if (op < 0) // unary operator
        return 3;
    if (op == '+' || op == '-')
        return 1;
    if (op == '*' || op == '/')
        return 2;
    return -1;
}
void process_op(stack<int>& st, char op) {
    if (op < 0) {
        int l = st.top(); st.pop();
        switch (-op) {
            case '+': st.push(l); break;
            case '-': st.push(-l); break;
        }
    } else {
        int r = st.top(); st.pop();
        int l = st.top(); st.pop();
        switch (op) {
            case '+': st.push(l + r); break;
            case '-': st.push(l - r); break;
            case '*': st.push(l * r); break;
            case '/': st.push(l / r); break;
        }
    }
}
int evaluate(string s) {
    stack<int> st;
    stack<char> op;
    bool may_be_unary = true;
    for (int i = 0; i < (int)s.size(); i++) {
        if (delim(s[i])) continue;
        if (s[i] == '(') {
            op.push('('); may_be_unary = true;
        } else if (s[i] == ')') {
            while (op.top() != '(') {
                process_op(st, op.top());
                op.pop();
            }
            op.pop(); may_be_unary = false;
        } else if (is_op(s[i])) {
```

ExpressionEvaluation FastKnapsack Datetime

```
char cur_op = s[i];
if (may_be_unary && is_unary(cur_op))
    cur_op = -cur_op;
while (!op.empty() && (
    (cur_op >= 0 && priority(op.top()) >= priority(
        cur_op)) ||
    (cur_op < 0 && priority(op.top()) > priority(cur_op
    ))) {
    process_op(st, op.top()); op.pop();
}
op.push(cur_op);
may_be_unary = true;
} else {
    int number = 0;
    while (i < (int)s.size() && isalnum(s[i]))
        number = number * 10 + s[i++] - '0';
    --i;
    st.push(number);
    may_be_unary = false;
}
}
while (!op.empty()) {
    process_op(st, op.top()); op.pop();
}
return st.top();
}
int32_t main() {
    cout << evaluate("12-3-(3 - 4)");
    return 0;
}
```

FastKnapsack.h

Description: Given N non-negative integer weights w and a non-negative target t, computes the maximum S <= t such that S is the sum of some subset of the weights.

Time: O(N max(w\_i))

b20ecc, 16 lines

```
int knapsack(vi w, int t) {
    int a = 0, b = 0, x;
    while (b < sz(w) && a + w[b] <= t) a += w[b++];
    if (b == sz(w)) return a;
    int m = *max_element(all(w));
    vi u, v(2*m, -1);
    v[a+m-t] = b;
    rep(i,b,sz(w)) {
        u = v;
        rep(x,0,m) v[x+w[i]] = max(v[x+w[i]], u[x]);
        for (x = 2*m; --x > m;) rep(j, max(0,u[x]), v[x])
            v[x-w[j]] = max(v[x-w[j]], j);
    }
    for (a = t; v[a+m-t] < 0; a--);
    return a;
}
```

Datetime.py

12 lines

```
# python datetime example
import datetime
d = datetime.datetime(2022, 12, 25)
print((d.year, d.month, d.day, d.hour, d.minute, d.second, d.
    microsecond))
d2 = datetime.datetime.fromtimestamp(1326244364)
dt = datetime.timedelta(days = 3, hours = 34, minutes = 56,
    seconds = 234, microseconds=1000001)
d = d + dt
print((d.year, d.month, d.day, d.hour, d.minute, d.second, d.
    microsecond))
```

```
print(d-d2) #<class 'datetime.timedelta'>, 4005 days,
    14:47:11.000001
print(d.strftime("%Y/%m/%d, %H:%M:%S"))
d3 = datetime.datetime.strptime("25 December, 2022", "%d %B, %Y
    ")
print(["Sunday", "Monday", "Tuesday", "Wednesday", "Thursday",
    "Friday", "Saturday"][d3.weekday()])
```

ZellerCongruence

$w = (\lfloor \frac{c}{4} \rfloor - 2c + y + \lfloor \frac{y}{4} \rfloor + \lfloor \frac{13(m+1)}{5} \rfloor + d - 1) \bmod 7$

w: 0 sunday, 1 monday, 2 tuesday, 6 saturday. c: first two digits of year. y: last two digits of year. m: month where  $3 \leq m \leq 14$ : change mon 1, 2 to mon 13, 14 in the previous year! d: day. Modular result may be negative.

10.3 Debugging tricks

- signal(SIGSEGV, [](int) { \_Exit(0); }); converts segfaults into Wrong Answers. Similarly one can catch SIGABRT (assertion failures) and SIGFPE (zero divisions). \_GLIBCXX\_DEBUG failures generate SIGABRT (or SIGSEGV on gcc 5.4.0 apparently).
- feenableexcept(29); kills the program on NaNs (1), 0-divs (4), infinities (8) and denormals (16).

10.4 Optimization tricks

\_\_builtin\_ia32\_ldmxcsr(40896); disables denormals (which make floats 20x slower near their minimum value).

10.4.1 Bit hacks

- x & -x is the least bit in x.
- for (int x = m; x; ) { --x &= m; ... } loops over all subset masks of m (except m itself).
- c = x&-x, r = x+c; ((r^x) >> 2)/c | r is the next number after x with the same number of bits set.
- rep(b,0,K) rep(i,0,(1 << K)) if (i & 1 << b) D[i] += D[i^(1 << b)]; computes all sums of subsets.

10.4.2 Pragmas

- #pragma GCC optimize ("Ofast") will make GCC auto-vectorize loops and optimizes floating points better.
- #pragma GCC target ("avx2") can double performance of vectorized code, but causes crashes on old machines.
- #pragma GCC optimize ("trapv") kills the program on integer overflows (but is really slow).