

Persistent Data Structures & Finger Trees

Fun Club / Haskell User Group – Dec 13th, 2012
Frederic Kettelhoit

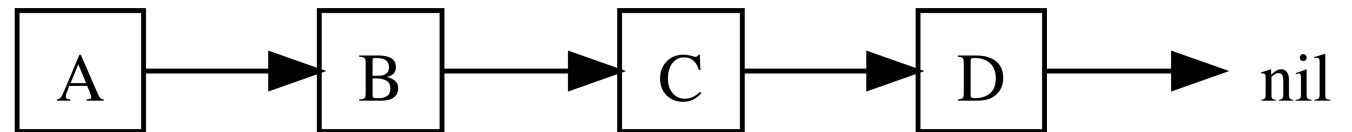
Functional Data Structures

- Functional languages emphasize **immutability**
 - We cannot just update data structures in place
 - Return a **new version** after every change
 - We could copy the whole data structure every time
 - Very inefficient

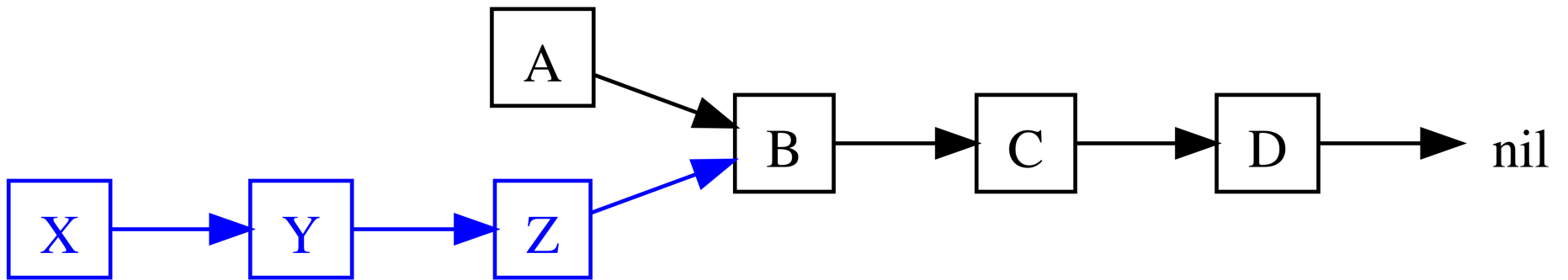
Persistent Data Structures

- When modified: preserve previous versions
 - immutable, always return **new versions**
- **Re-use** unchanged parts of the old version
- Example: (Singly) linked lists

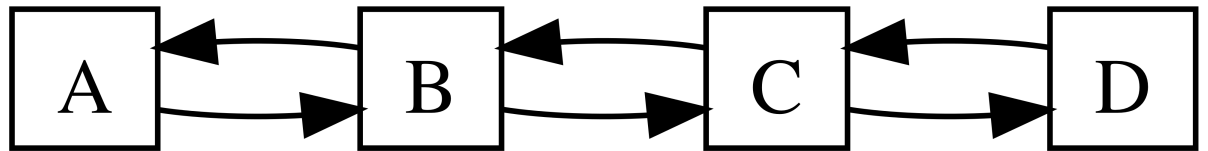
Linked Lists



Linked Lists

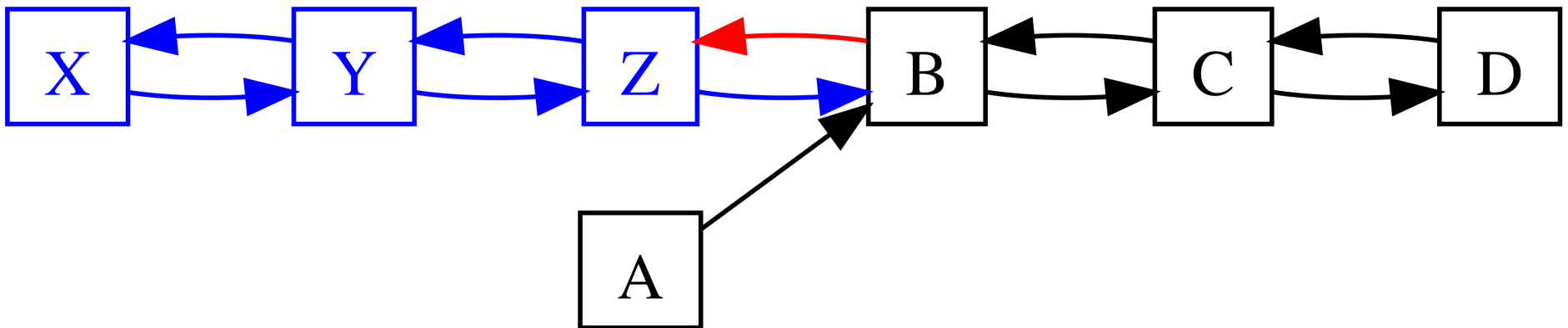


Doubly Linked List



Persistent?

Doubly Linked List



Not persistent!

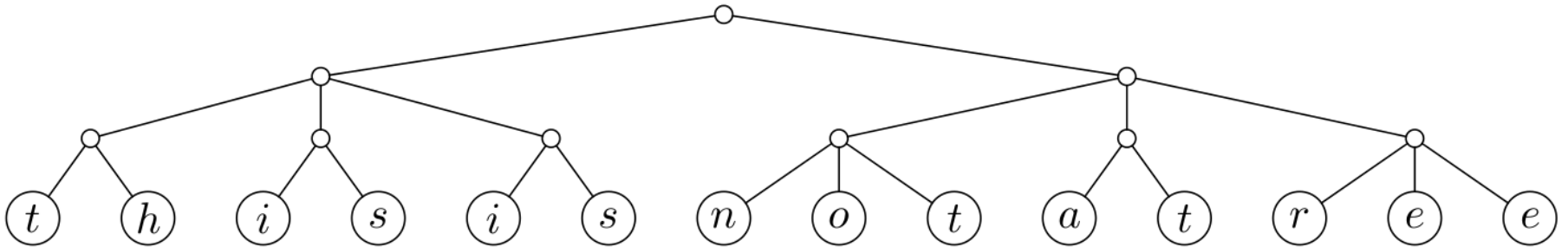
Finger Trees

- Many properties of doubly linked lists
 - cons/head/tail on both sides (deque)
- Customizable through monoids, can act as:
 - Random Access Sequence, Max-Priority Queue, Sorted Set, ...
- Good upper bounds for most operations
- Persistent

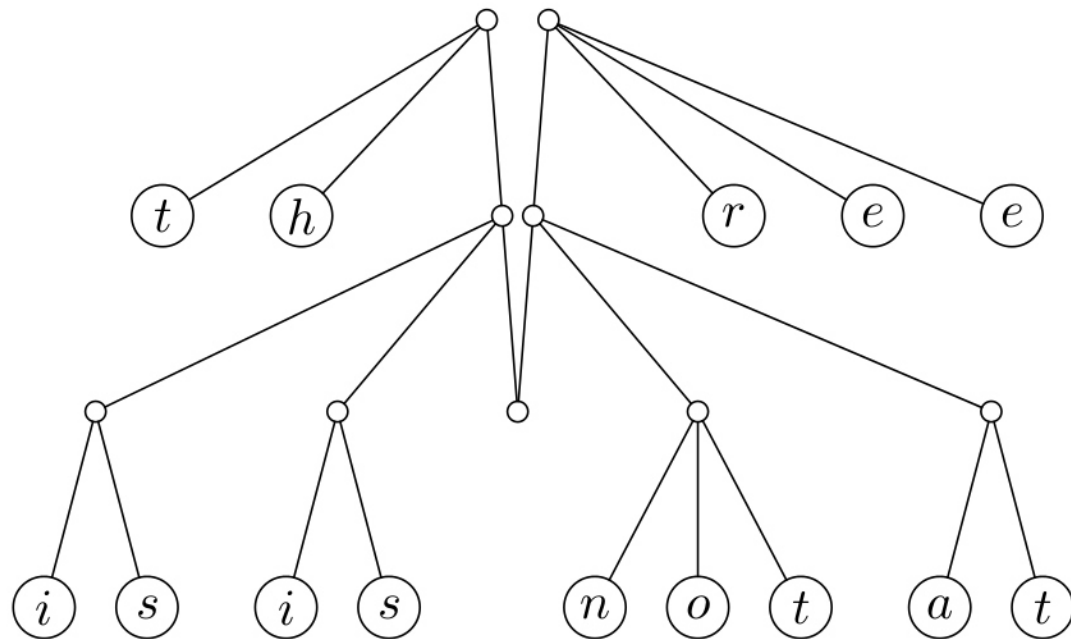
Upper Bounds

- head/tail in amortized time **$O(1)$**
 - left and right
- cons/snoc in amortized time **$O(1)$**
- concat in time **$O(\log (\min m n))$**
- split in time **$O(\log (\min m n))$**

2-3 Tree



2-3 Finger Tree

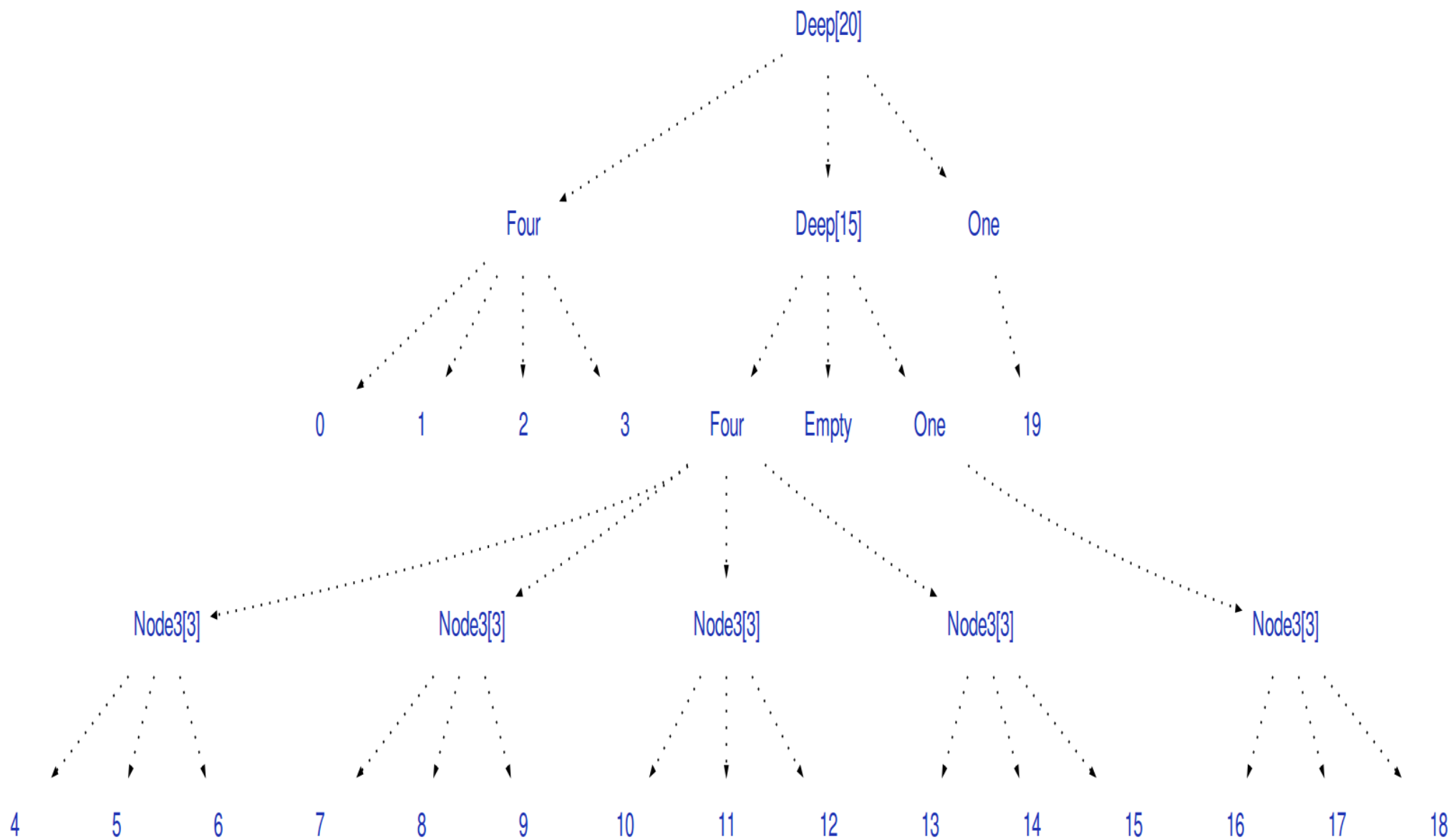


Finger Trees in Haskell

```
data FingerTree a = Empty
                  | Single a
                  | Deep (Digit a) (FingerTree (Node a)) (Digit a)
```

```
type Digit a = One    a
             | Two    a a
             | Three  a a a
             | Four   a a a a
```

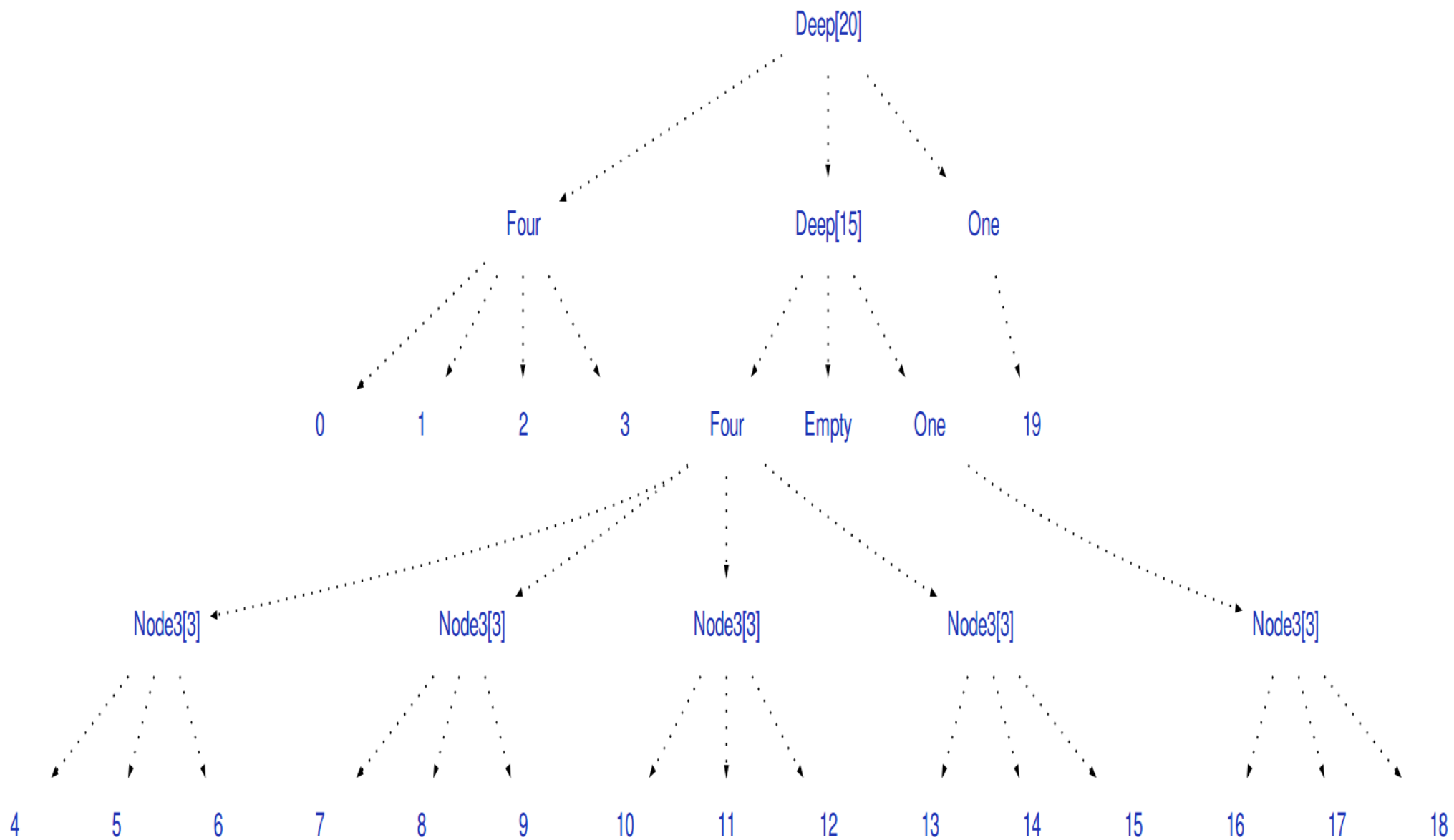
```
data Node a = Node2 a a
            | Node3 a a a
```

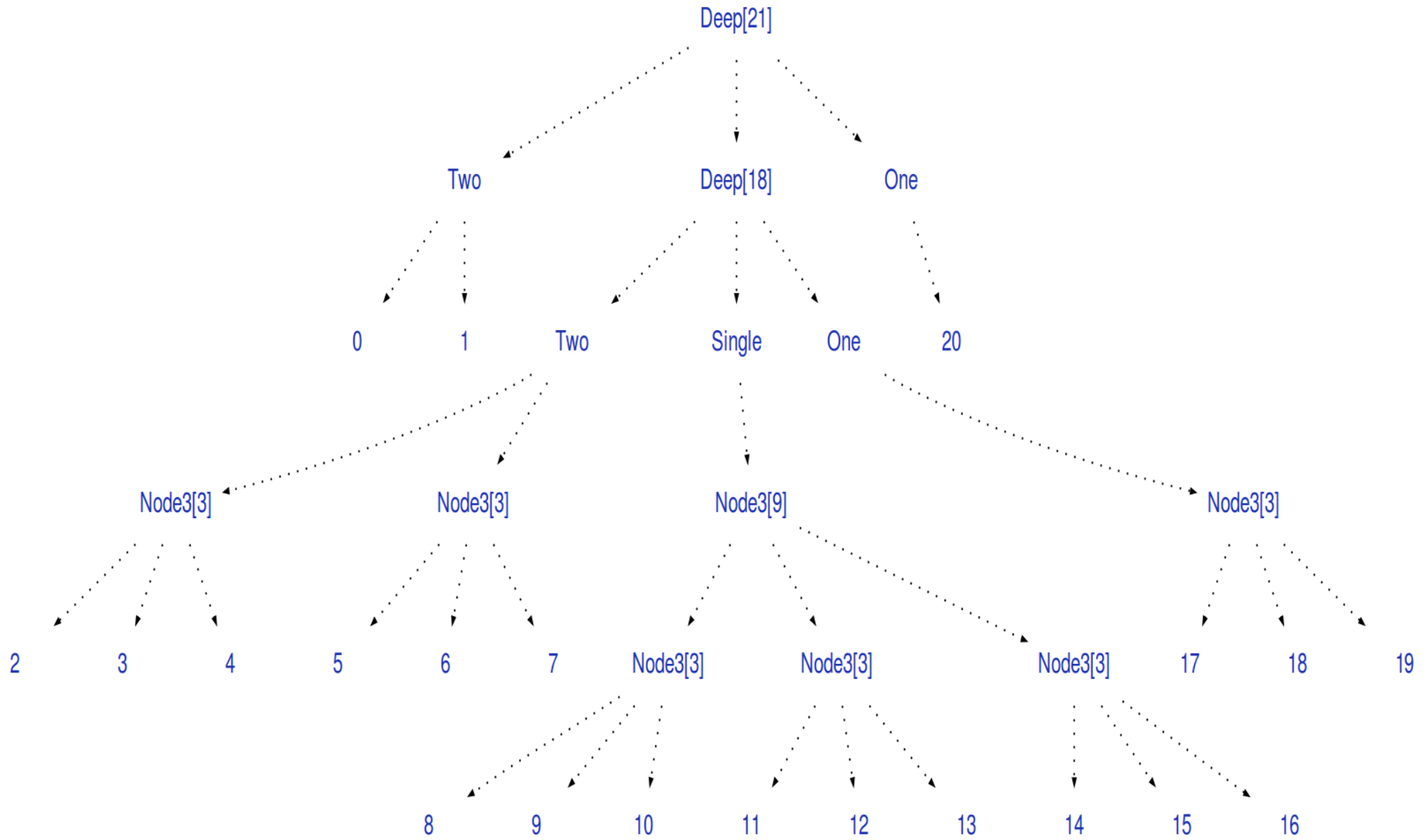


cons

```
cons :: a -> FingerTree a -> FingerTree a
cons a Empty                    = Single a
cons a (Single b)              = Deep (One a) Empty (One b)
cons a (Deep (One b) m sf)     = Deep (Two a b) m sf
cons a (Deep (Two b c) m sf)   = Deep (Three a b c) m sf
cons a (Deep (Three b c d) m sf) = Deep (Four a b c d) m sf
cons a (Deep (Four b c d e) m sf)
    = Deep (Two a b) (cons (Node3 c d e) m) sf
```

- $O(\log n)$ worst case, $O(1)$ amortized
- One, Two Three: Safe
- Four: Dangerous
- *(for the other side everything is simply swapped)*





concat

```
app3 :: FingerTree a -> [a] -> FingerTree a -> FingerTree a
app3 Empty ts xs      = cons' ts xs
app3 xs ts empty      = snoc' xs ts
app3 (Single x) ts xs = cons x (cons' ts xs)
app3 xs ts (Single x) = snoc (snoc' xs ts) x
app3 (Deep pr1 m1 sf1) ts (Deep pr2 m2 sf2)
    = Deep pr1 (app3 m1 (nodes sf1 ++ ts ++ pr2) m2) sf2
```

```
nodes :: [a] -> [Node a]
nodes [a, b] = [Node2 a b]
nodes [a, b, c] = [Node3 a b c]
nodes [a, b, c, d] = [Node2 a b, Node2 c d]
nodes (a : b : c : xs) = Node3 a b c : nodes xs
```

```
concat :: FingerTree a -> FingerTree a -> FingerTree a
concat xs ys = app3 xs [] ys
```

- concat in Zeit $O(\log (\min m n))$

Customization using Monoids

- A Monoid is a structure with 2 operations

```
mempty  :: a
mappend :: a -> a -> a
```

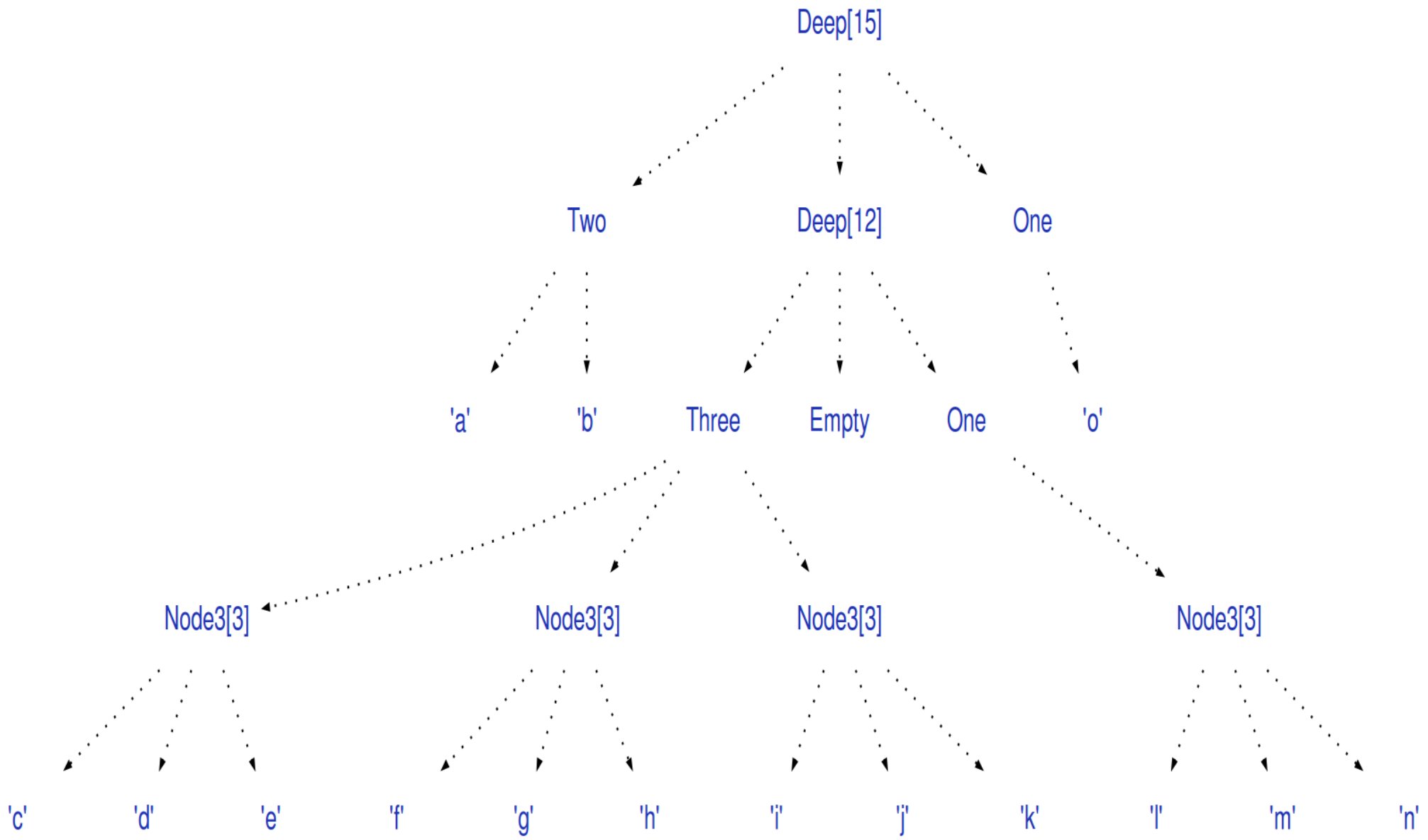
- Monoid Law:

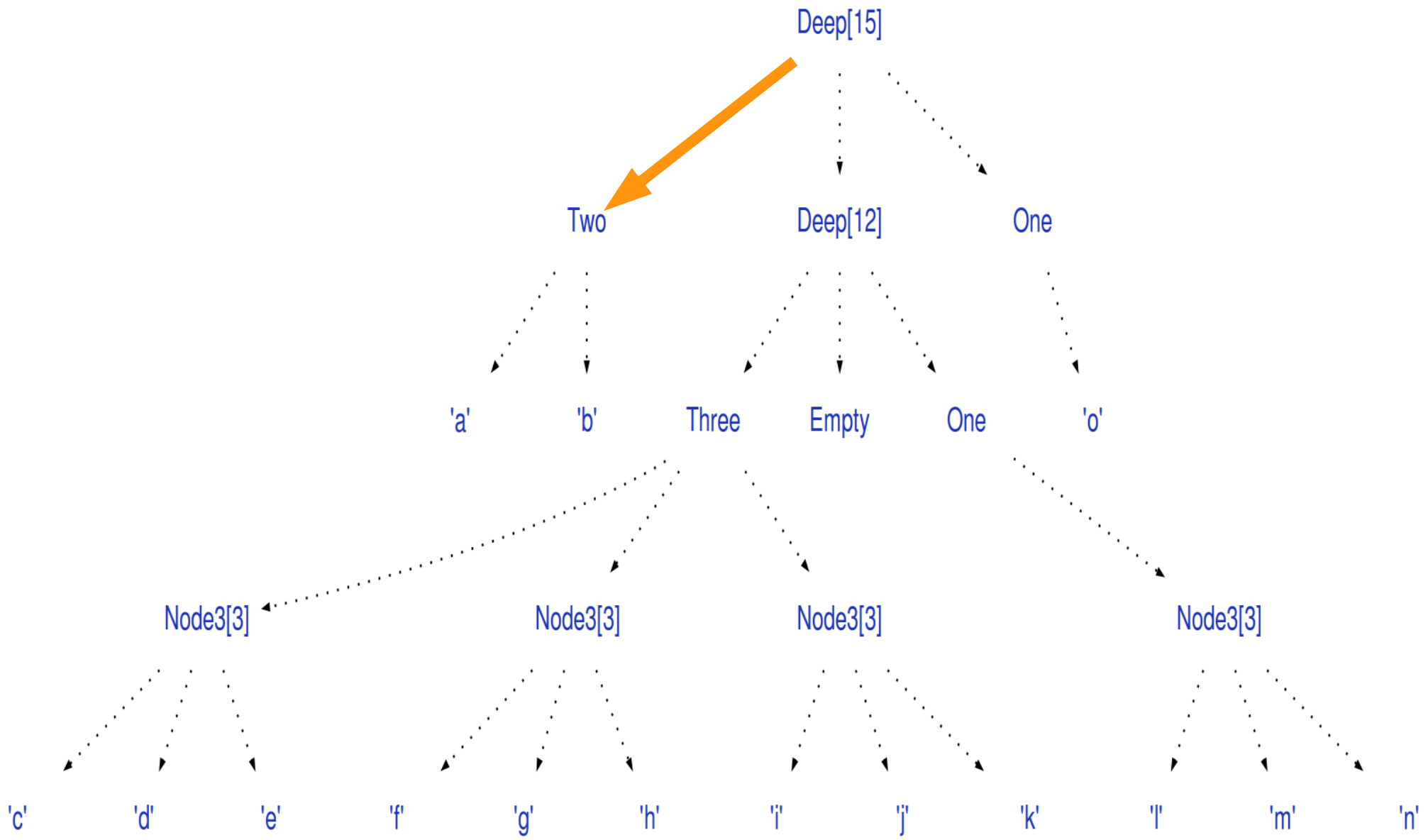
```
mappend mempty x = x
mappend x mempty = x
mappend x (mappend y z) = mappend (mappend x y) z
```

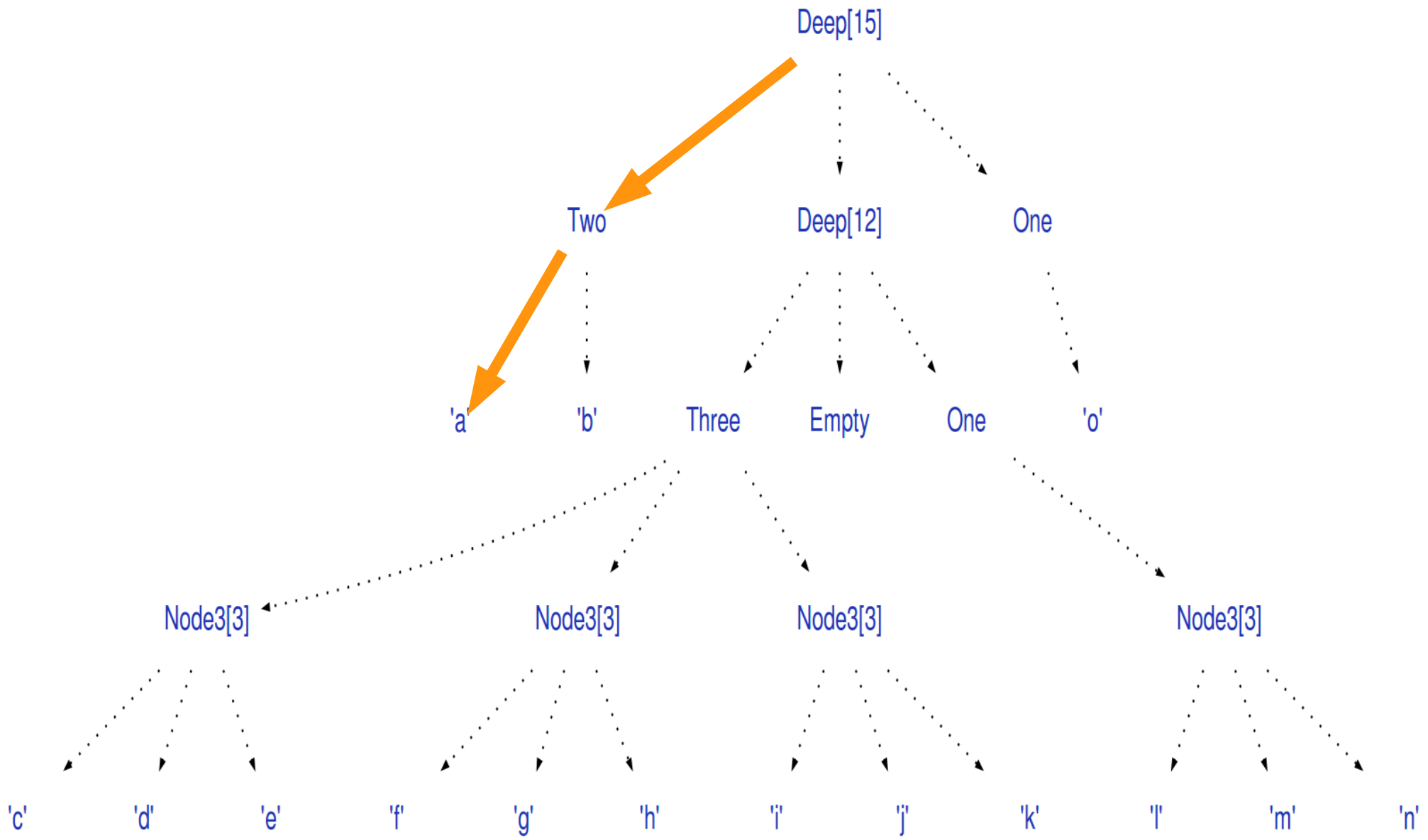
- And for the non-Category Theorists:
 - (Monoid = binary associative operation + identity)
 - For example: + and 0
- Finger Trees can be customized using Monoids

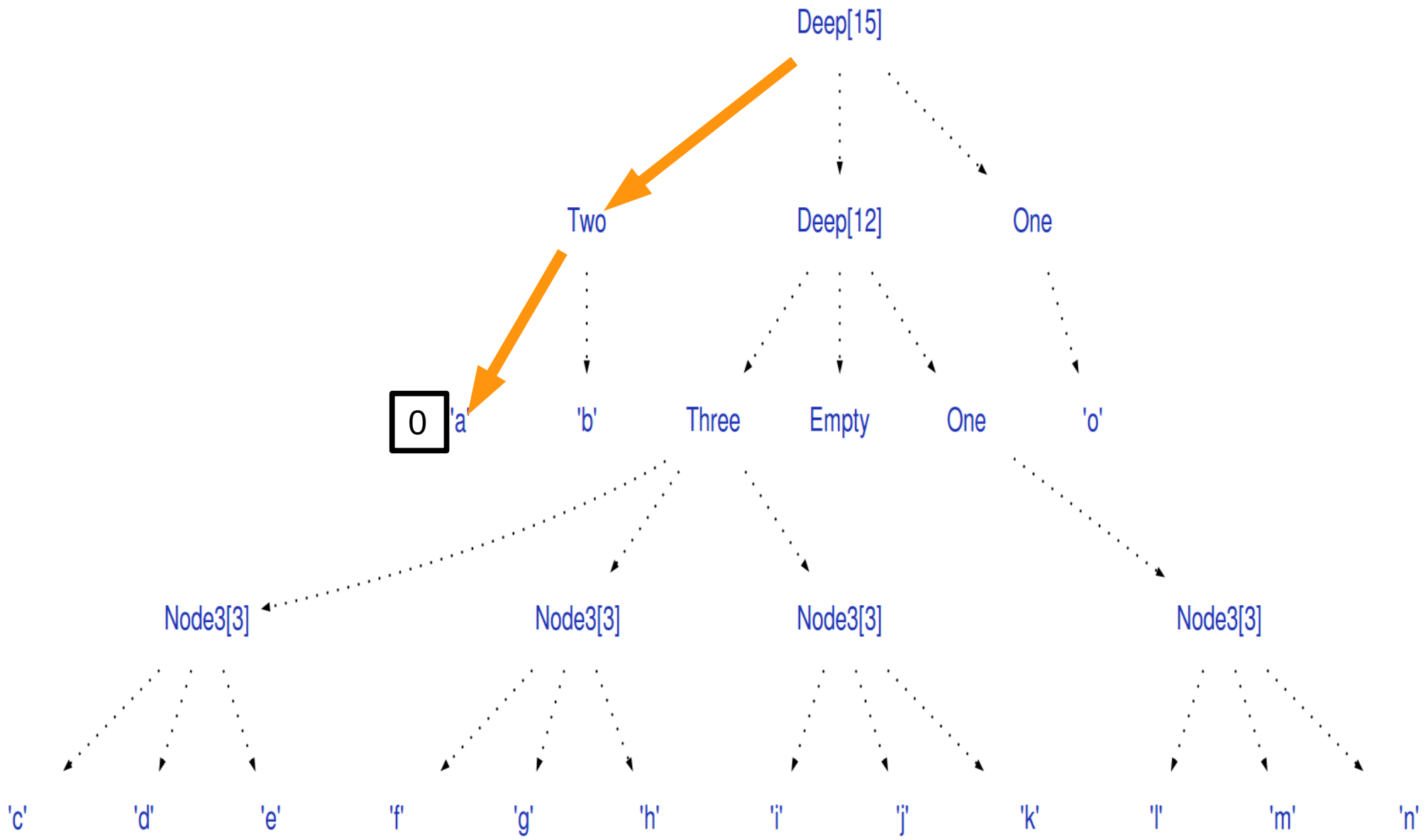
Customization using Monoids

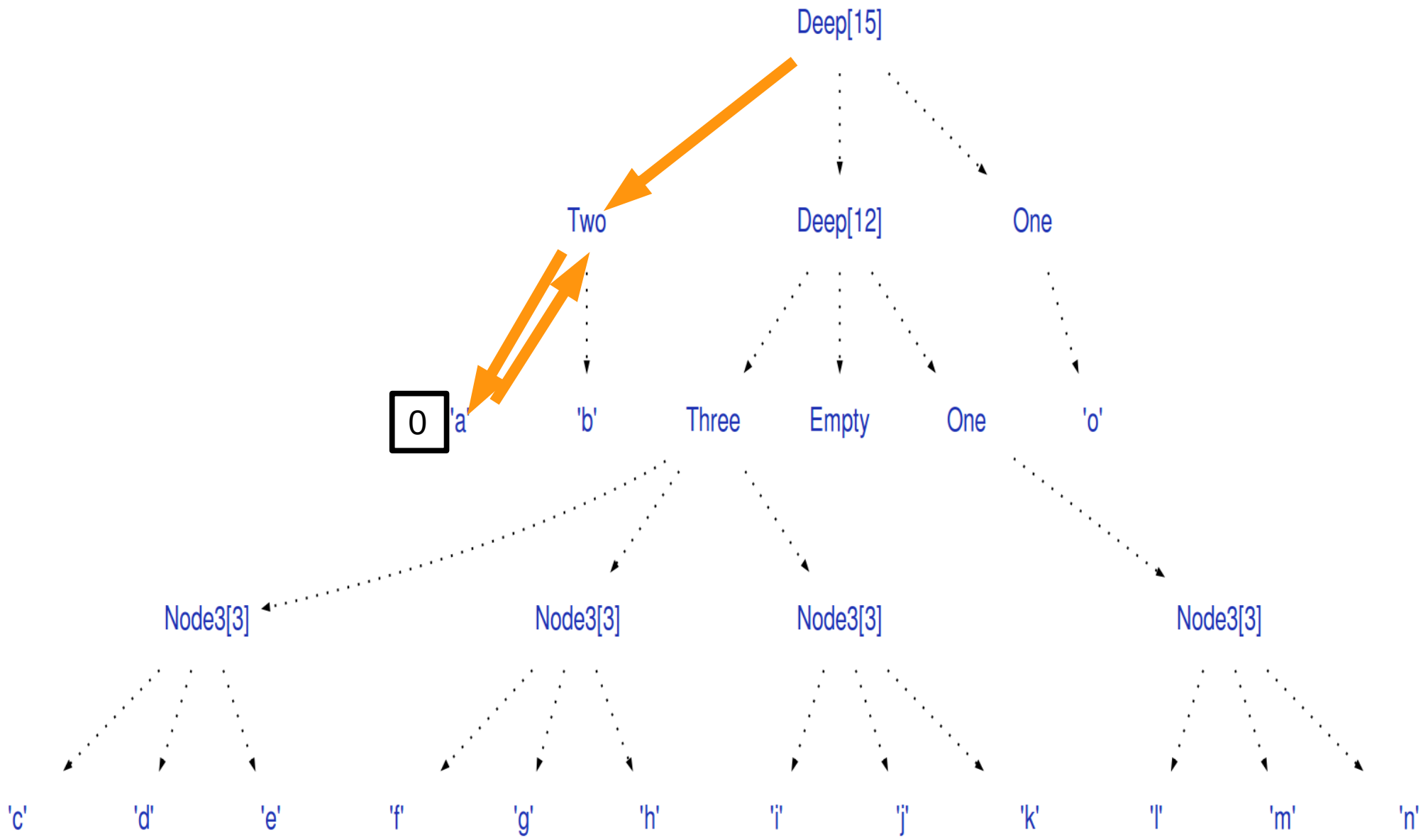
- Every node contains an additional Monoid annotation with the value:
 - Empty = Identity of Monoid
 - Single $a = ||a||$
 - Deep $pr\ m\ sf = \text{mappend } (||pr||\ ||m||)\ ||sf||$
- Example: Size Monoid
 - Empty = 0
 - Single $a = ||a|| = (1 \text{ for primitive elements})$
 - Deep $pr\ m\ sf = ||pr|| + ||m|| + ||sf||$ (number of elements in the tree)

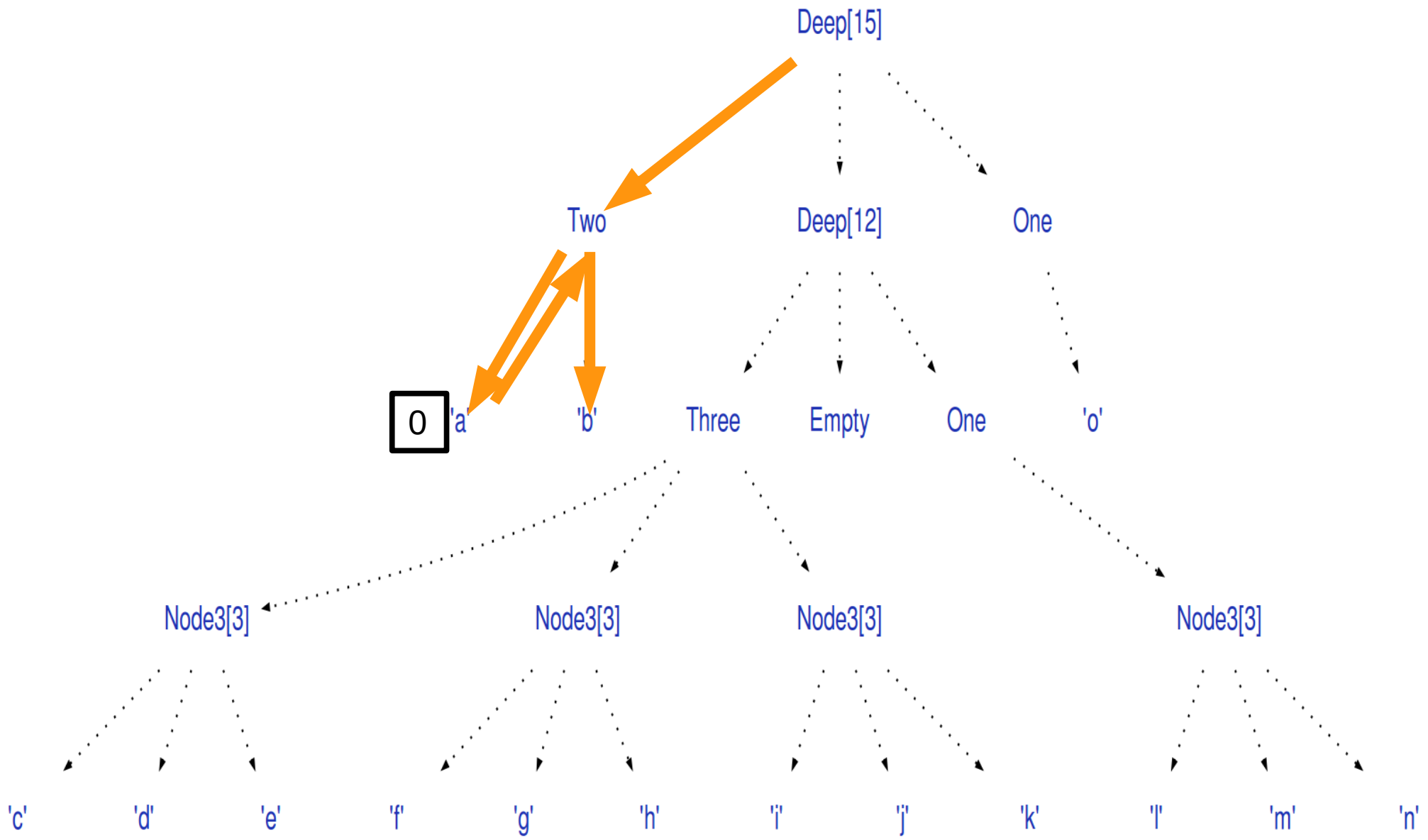


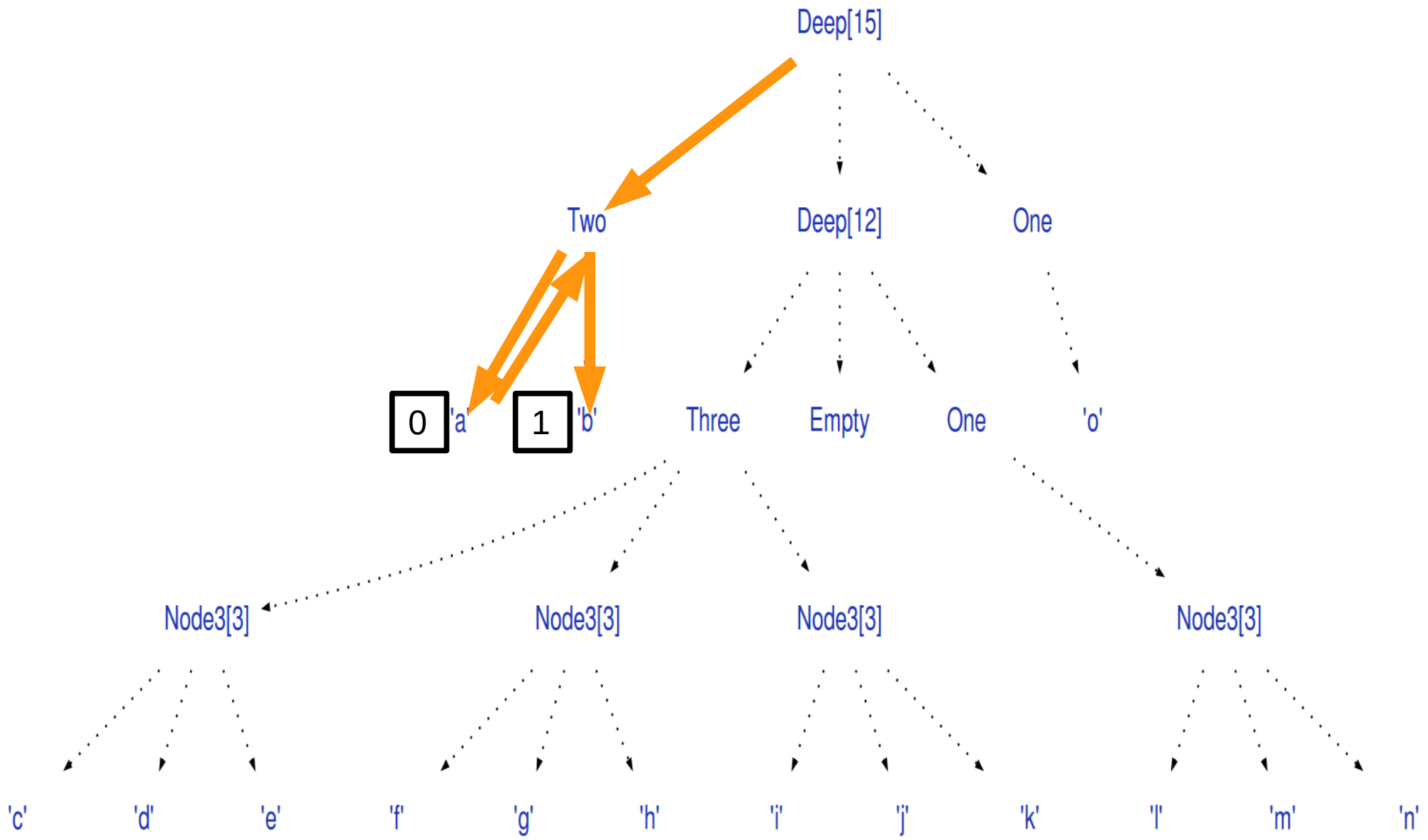


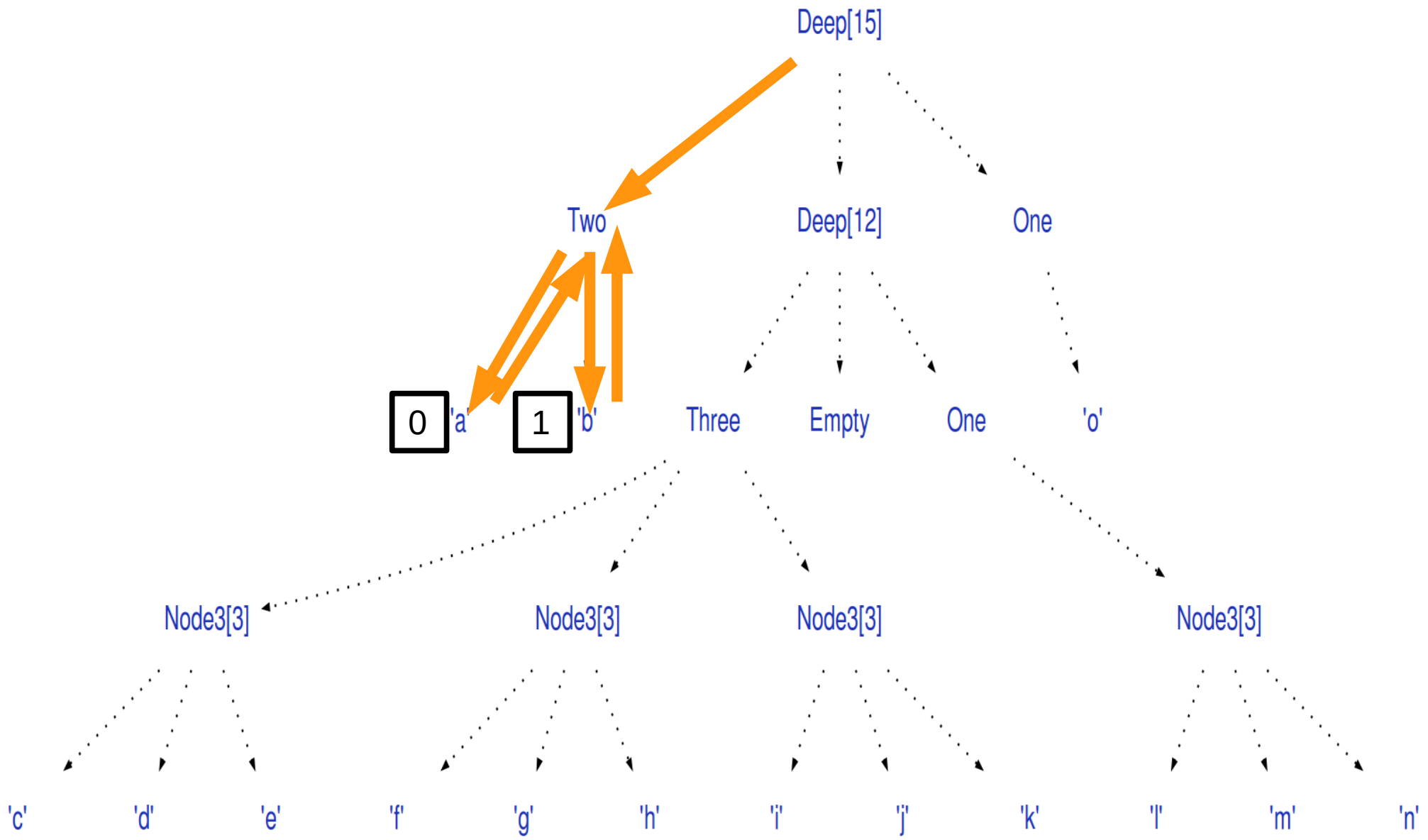


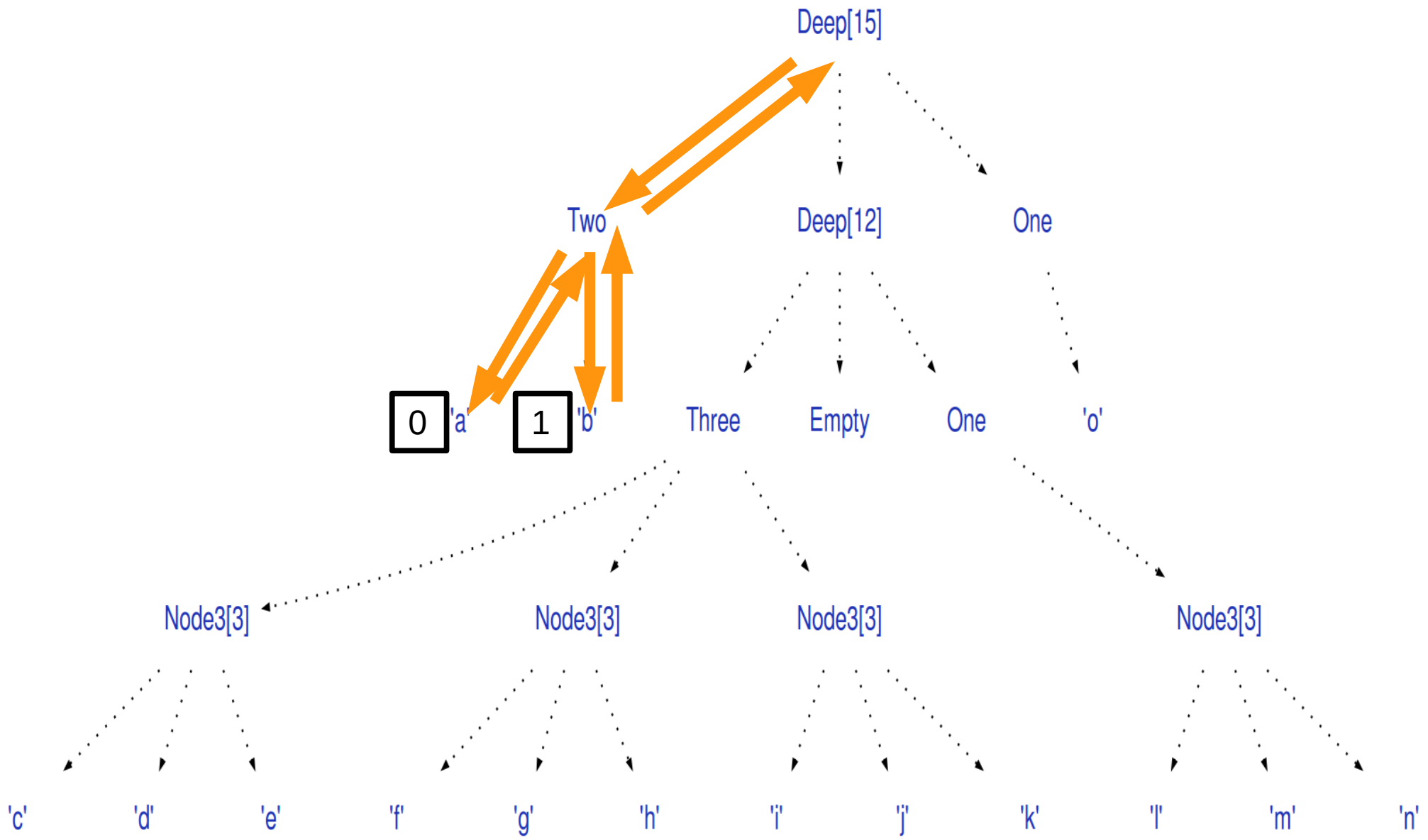


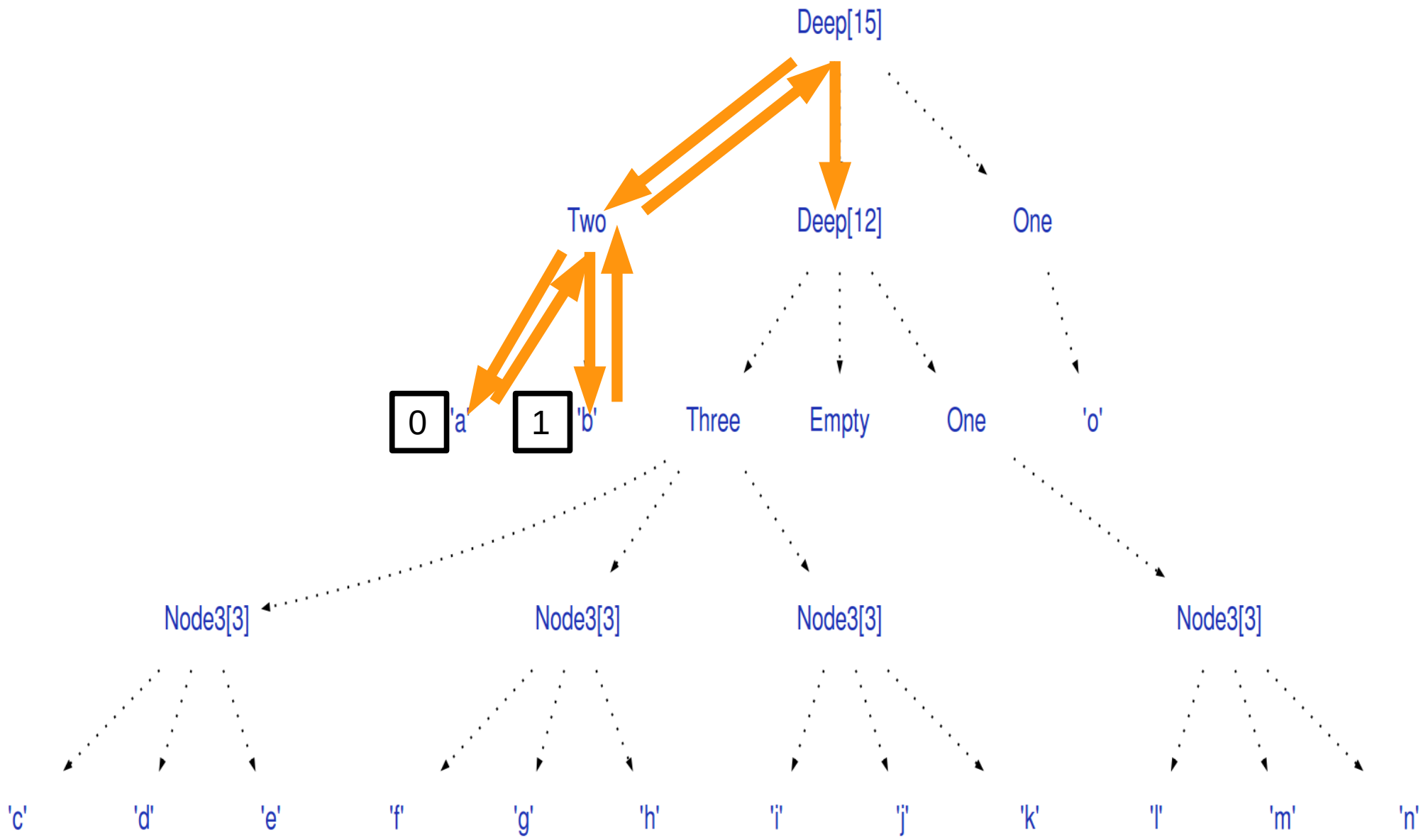


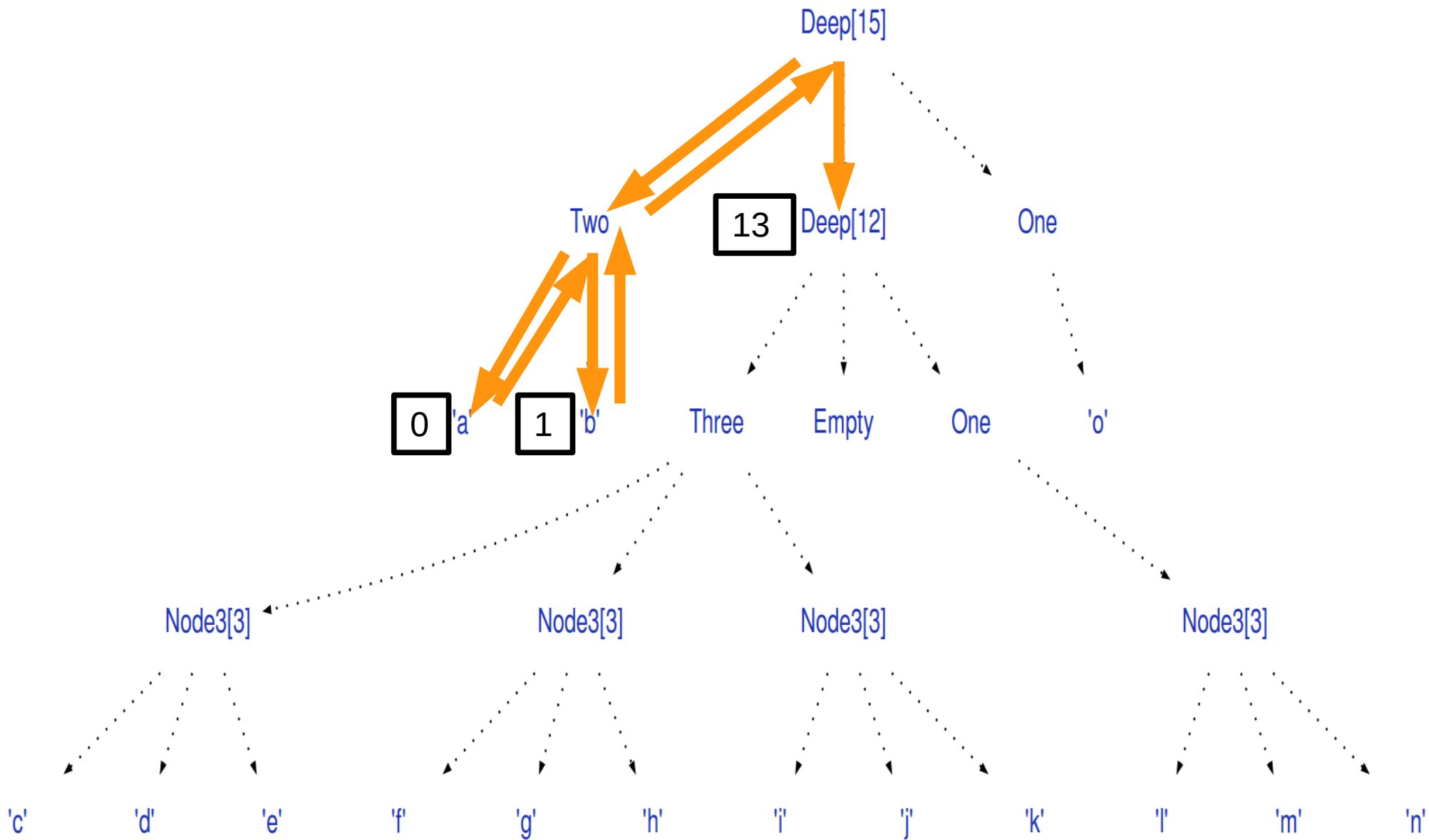


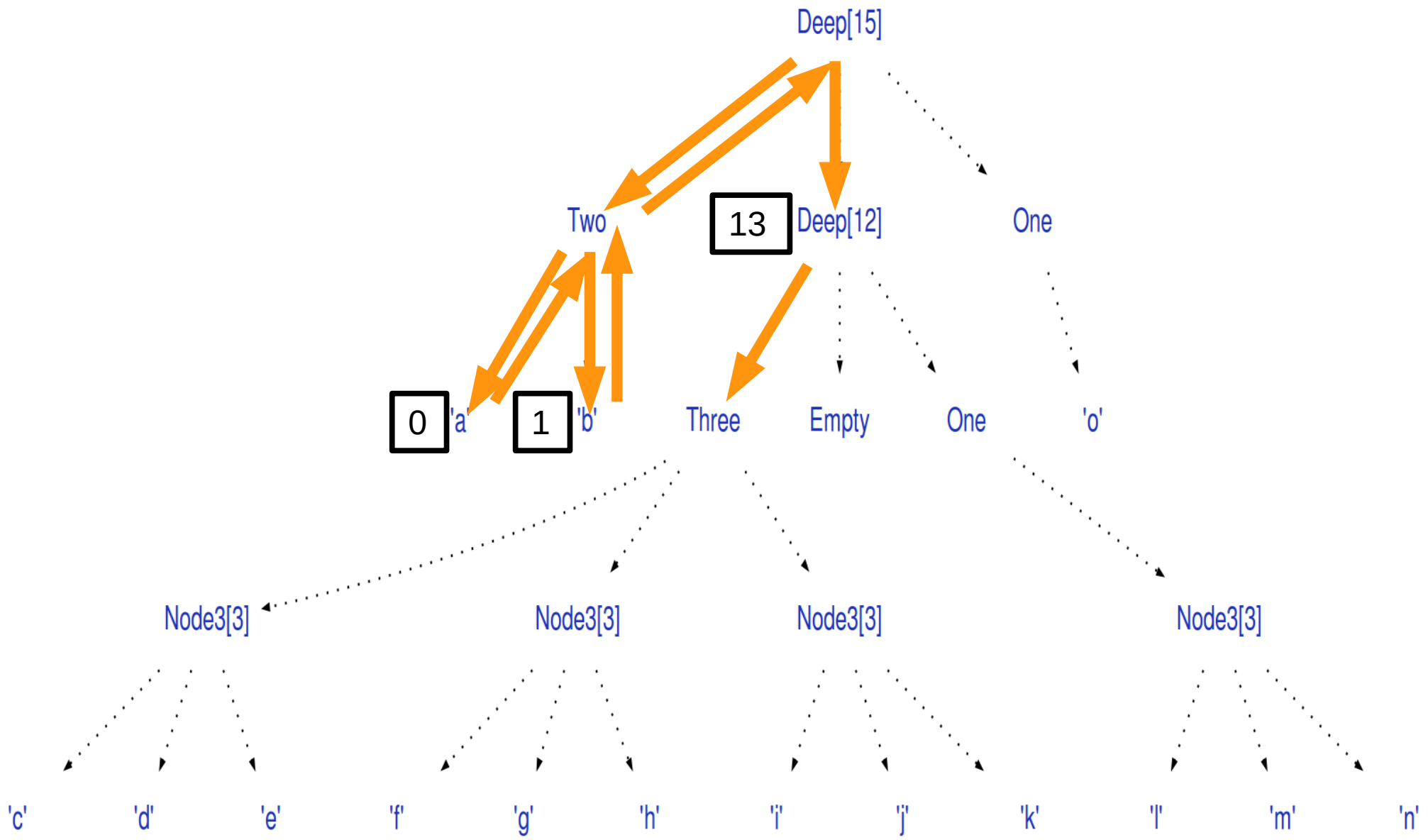


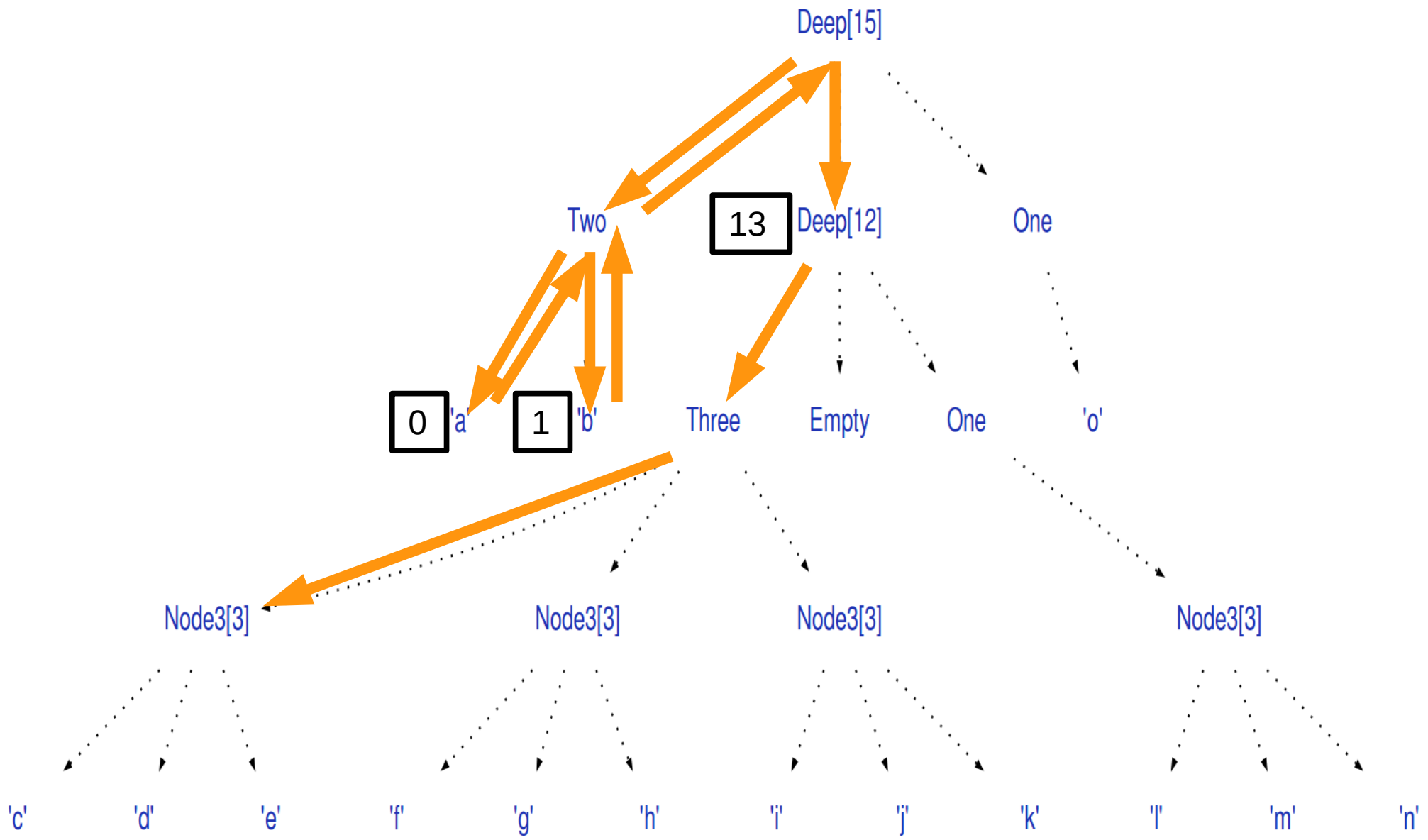


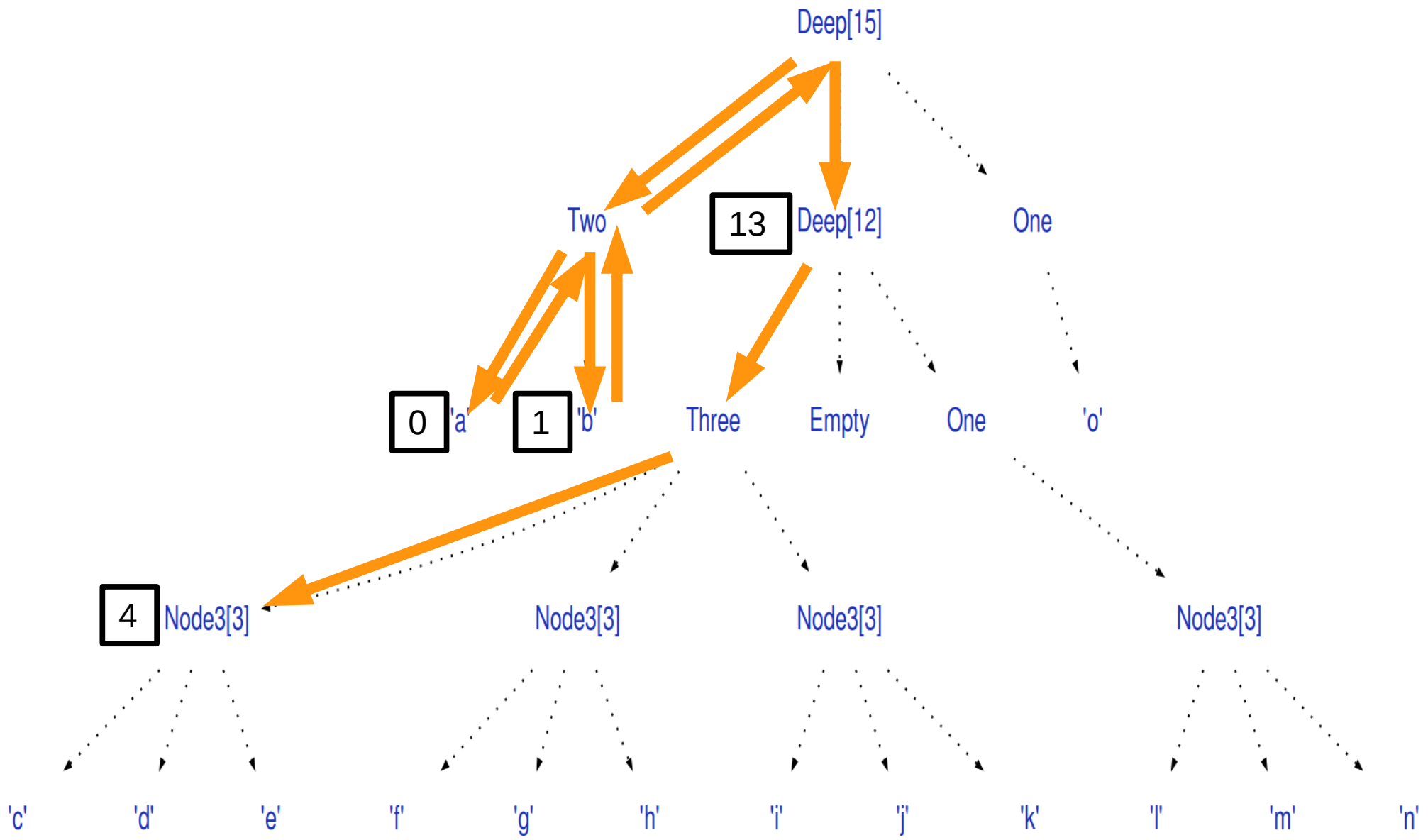


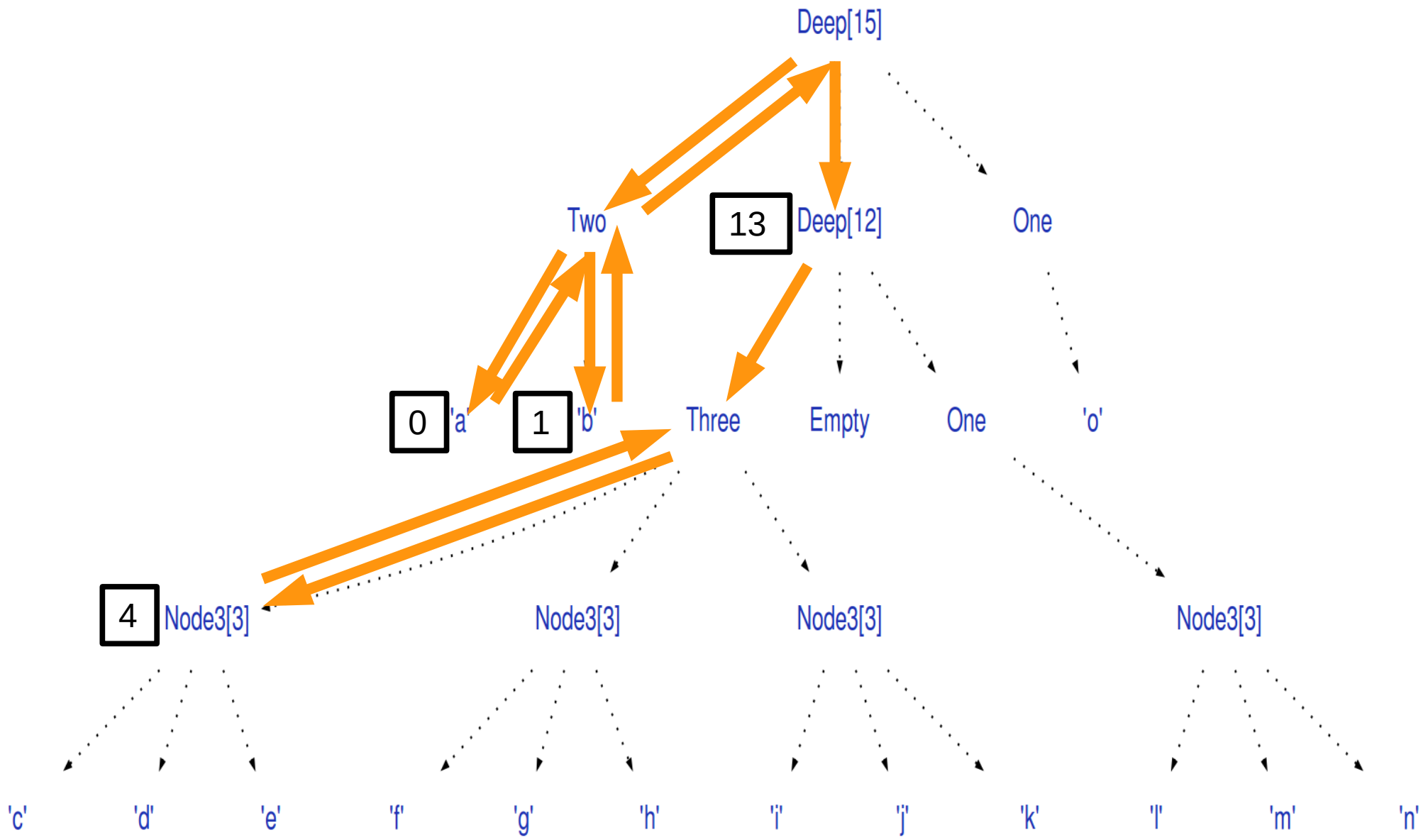


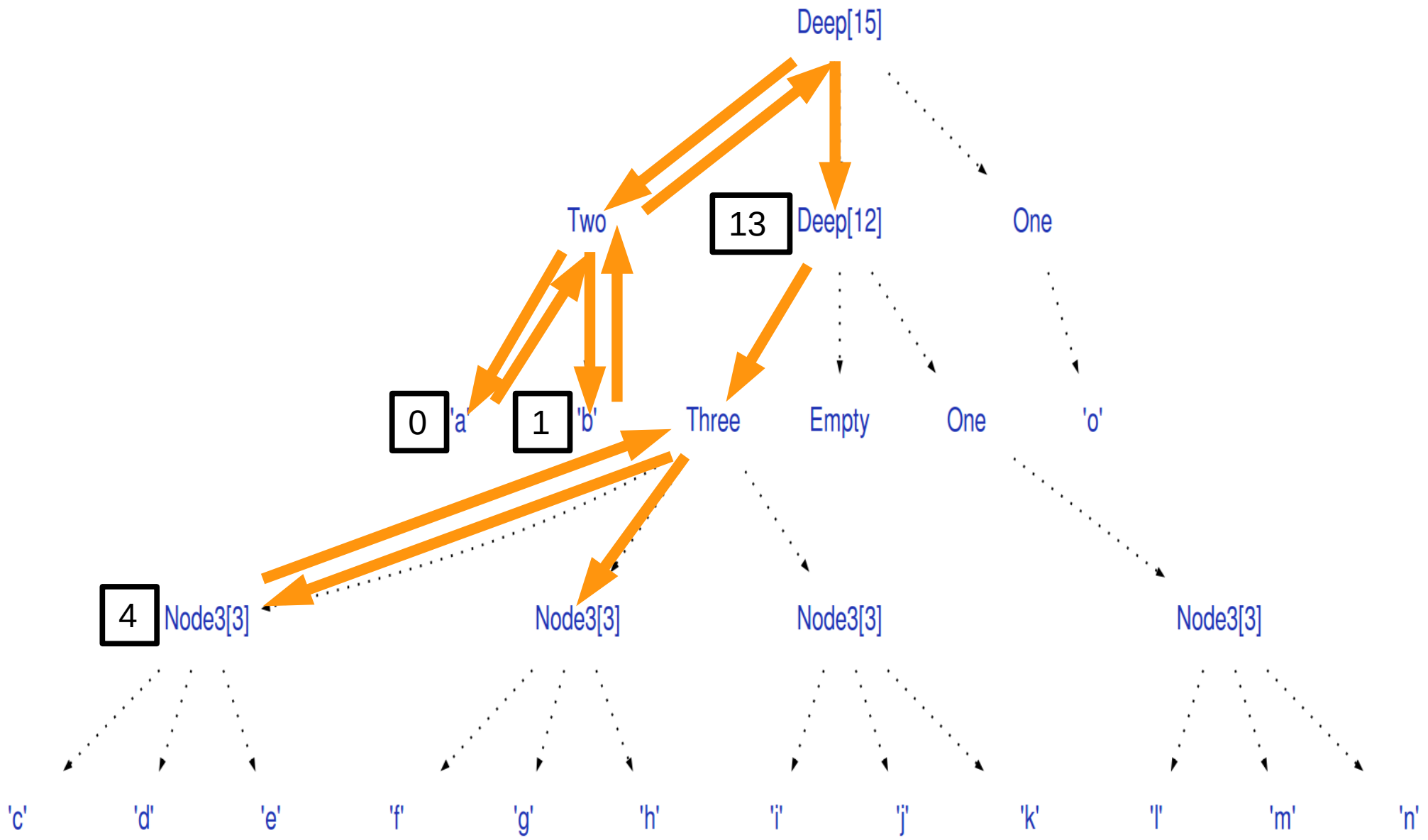


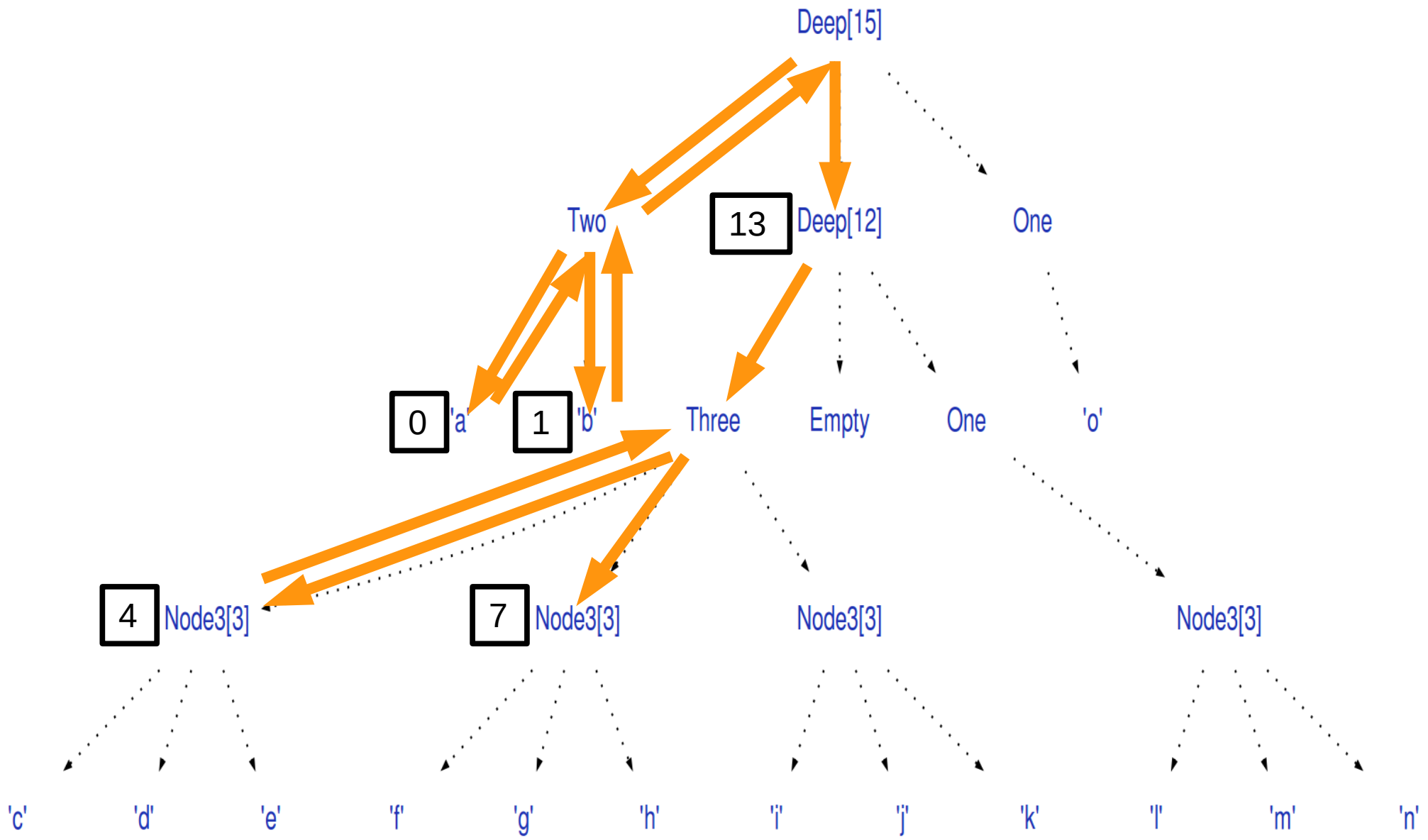


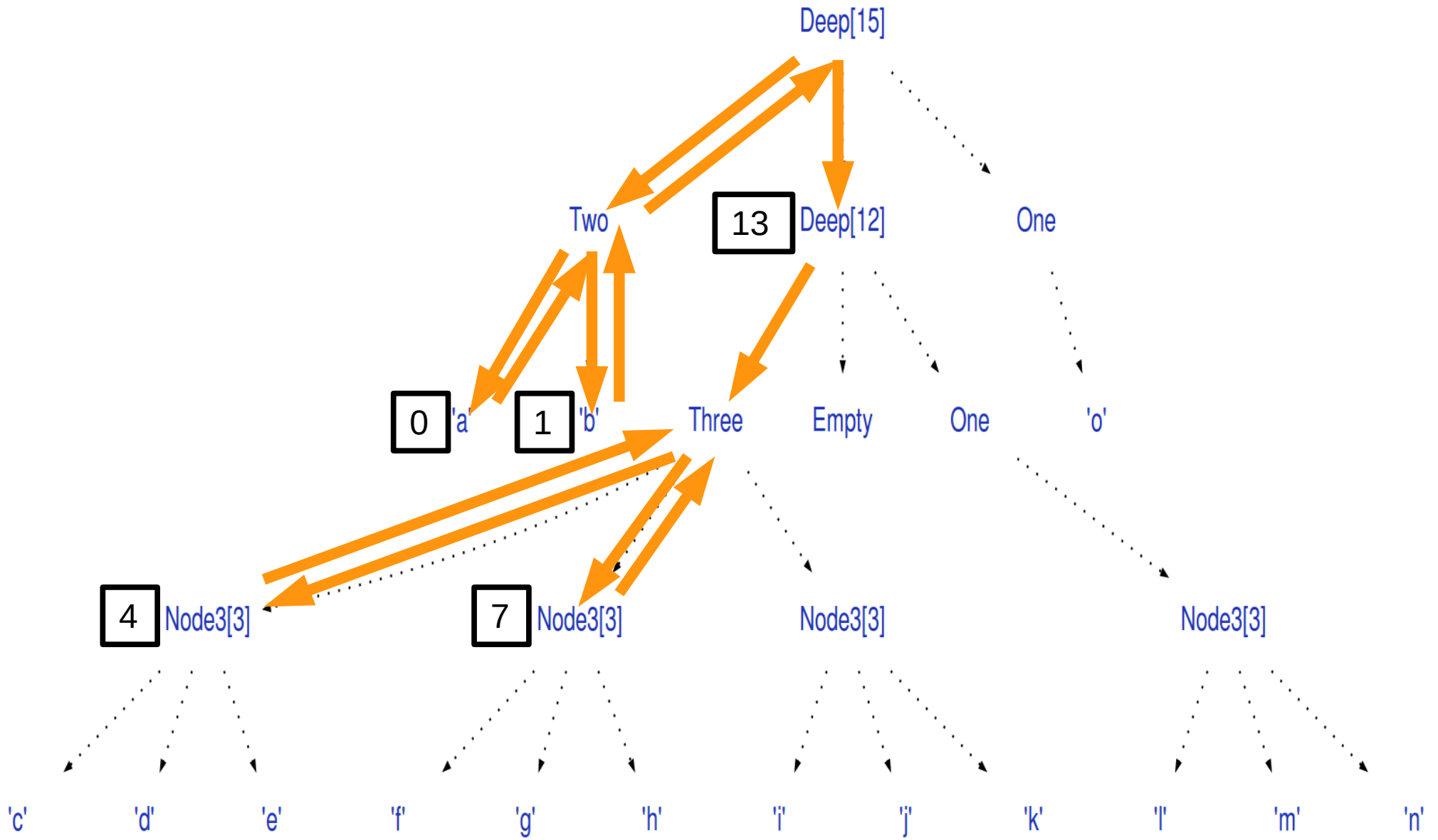


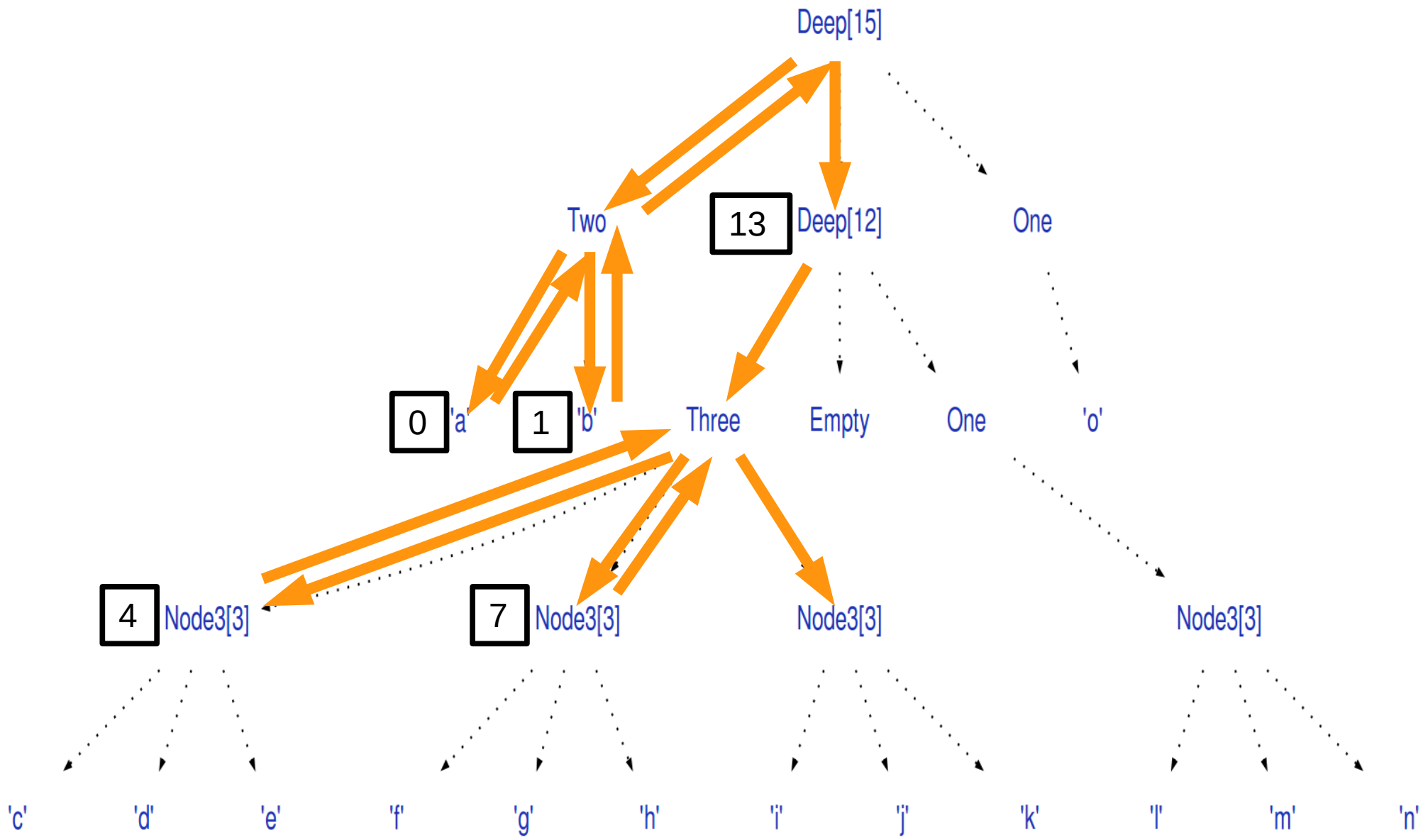


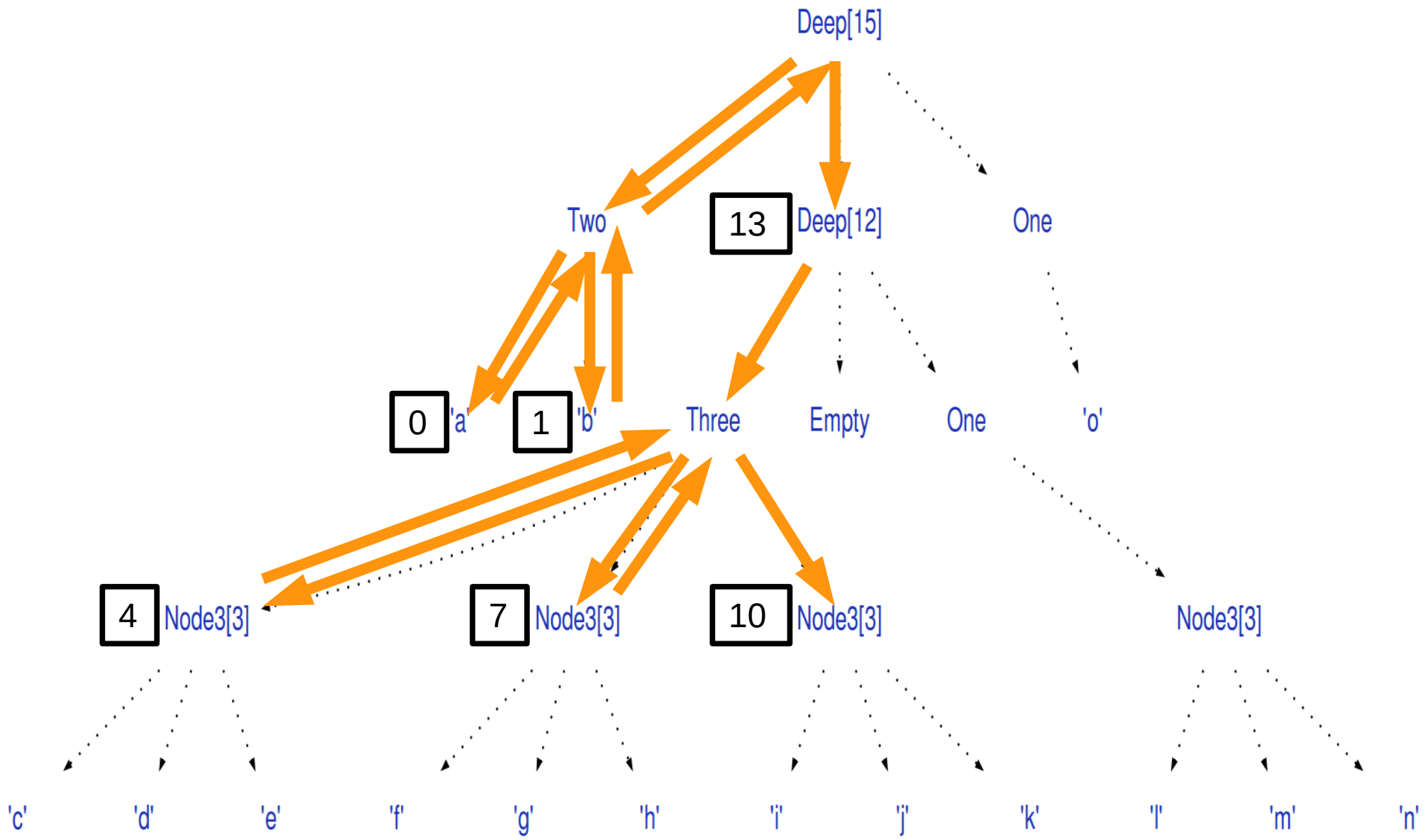


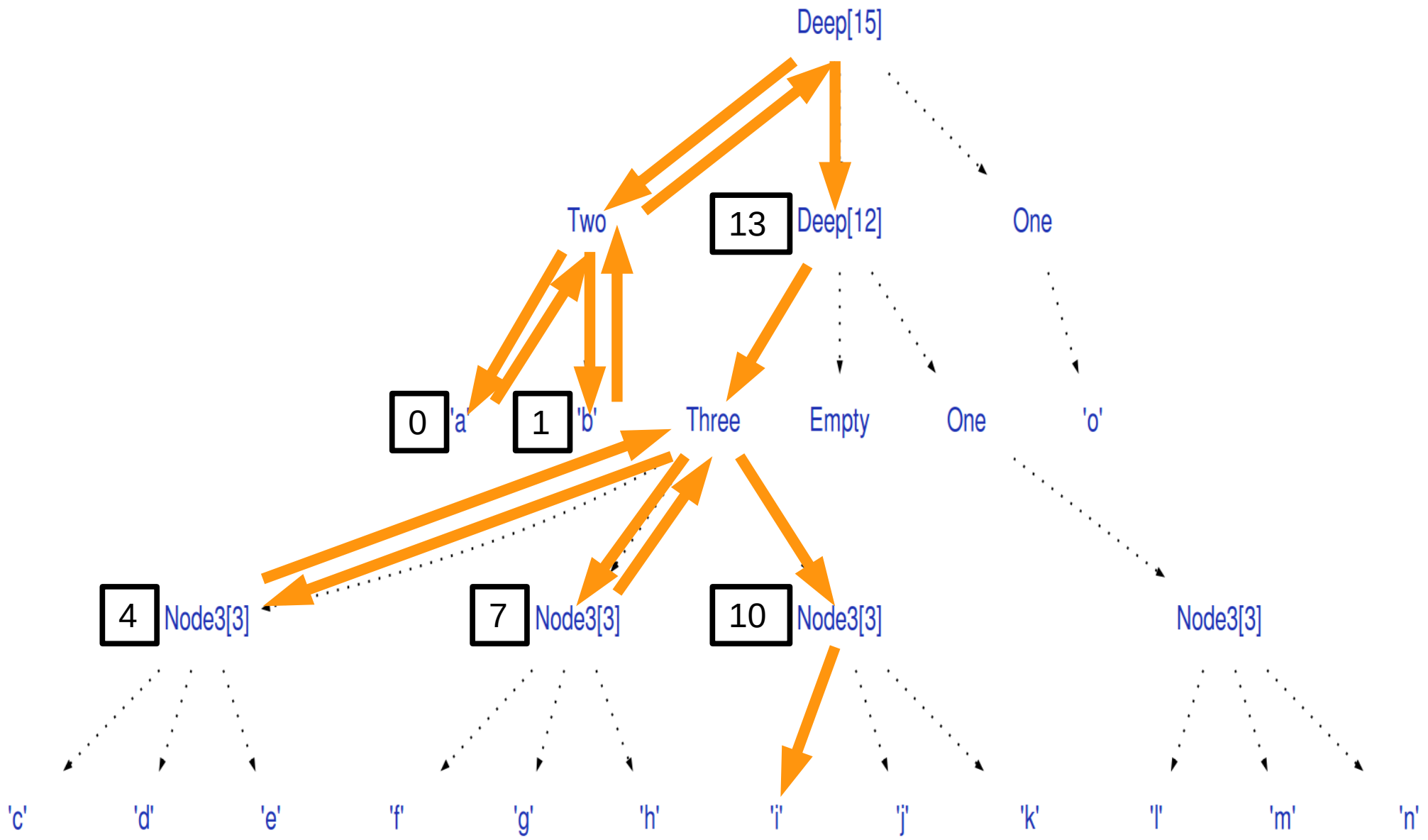


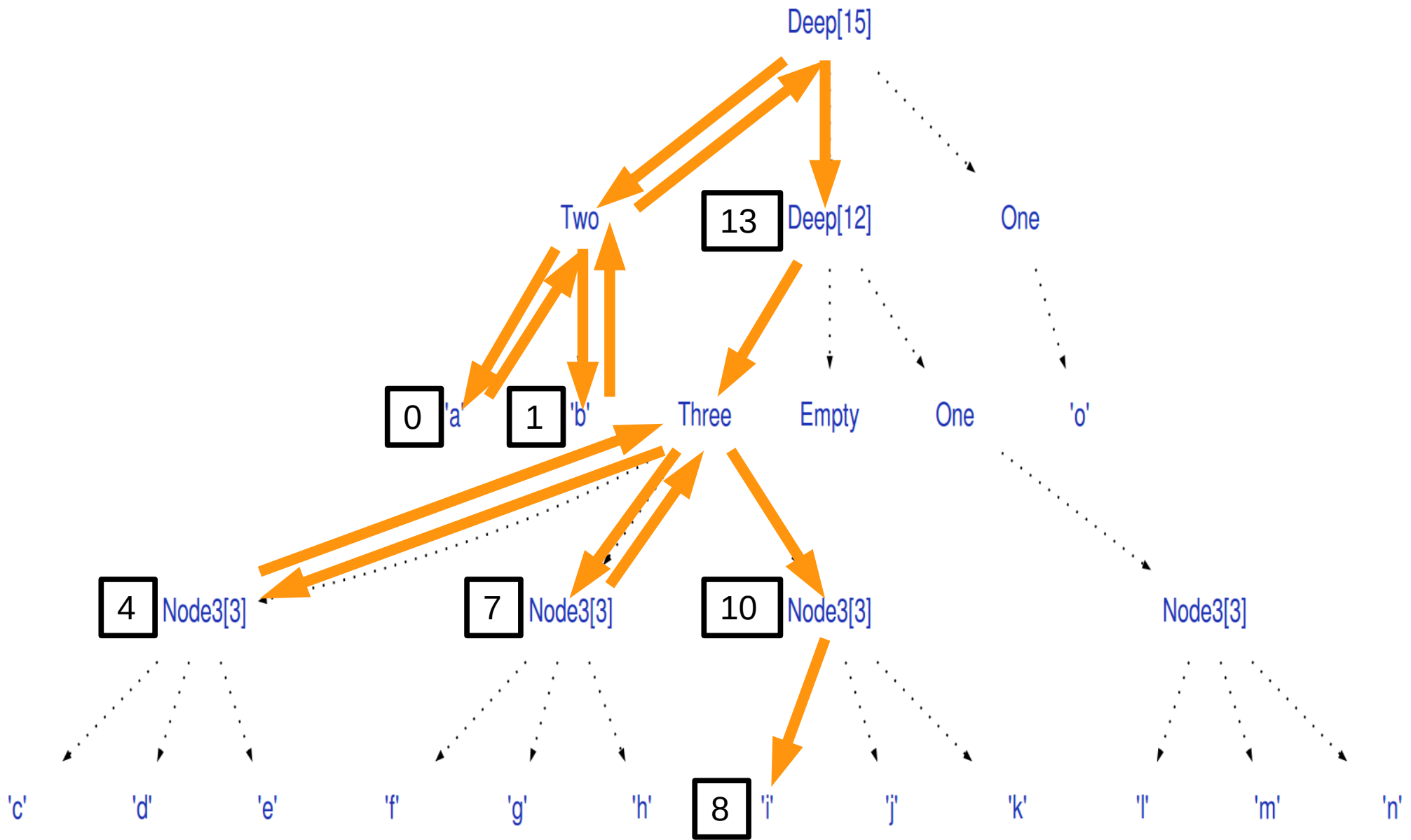












Finger Trees in Comparison

	Time (in ns) per operation randomly selected from			
	$\triangleleft, view_L$ (stack)	$\triangleright, view_L$ (queue)	$\triangleleft, view_L,$ $\triangleright, view_R$ (deque)	<i>index</i>
Bankers queue	51	147	—	—
Bankers deque	56	229	75	—
Catenable deque	78	215	100	—
Skew binary random access list	44	—	—	295
Finger tree	67	106	89	—
Finger tree with sizes	74	128	94	482

Table 1. *Comparing persistent sequence implementations*

Resources

- The original paper:
 - <http://www.soi.city.ac.uk/~ross/papers/FingerTree.html>
- Finger Trees in Haskell:
 - <http://www.haskell.org/ghc/docs/latest/html/libraries/containers/Data-Sequence.html>
- Finger Trees in Clojure:
 - <https://github.com/clojure/data.finger-tree>