

## FLA (Fall 2022) – Assignment 2

Name: \_\_\_\_\_ Dept: \_\_\_\_\_

Grade: \_\_\_\_\_ ID: \_\_\_\_\_

**Due: Nov. 15, 2022**

### Problem 1

Give context free grammars that generate the following languages, and give a brief description of the functionality of each variable in your grammars (in natural language).

- a.  $\{w \in \{a, b\}^* \mid w = w^R\}$
- b.  $\{w \in \{a, b\}^* \mid |b|_w = 2|a|_w\}$ , where  $|x|_w$  denotes the number of occurrences of  $x$  in string  $w$ .
- c.  $\{a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i + j = k\}$

## Problem 2

Consider the following context free grammar:  $G = (\{ S \}, \{ x, +, (, ) \}, P, S)$ , where  $P$  consists of:

$$S \rightarrow S + S \mid (S) \mid x$$

Is grammar  $G$  ambiguous? If it is, give an example string that has more than one parse tree according to  $G$ , and give an unambiguous grammar that generates the same language. If not, prove your conclusion.

### Problem 3

Consider the following context free grammar:  $G = (\{ S, B \}, \{ 0, 1 \}, P, S)$ , where  $P$  consists of:

$$S \rightarrow BSA \mid A$$

$$A \rightarrow 0A \mid \epsilon$$

$$B \rightarrow B10 \mid \epsilon$$

- a. For the string 1010000, give its parse tree and rightmost derivation according to  $G$ .
- b. Provide a nondeterministic PDA  $P$  that accepts the language  $L(G)$  by empty stack.

## Problem 4

Begin with the grammar:

$$S \rightarrow BSA \mid A$$

$$A \rightarrow aA \mid \epsilon$$

$$B \rightarrow Bba \mid \epsilon$$

1. Eliminate  $\epsilon$ -productions.
2. Eliminate any unit productions in the resulting grammar of (1.).
3. Put the resulting grammar of (2.) into Chomsky normal form.

## Problem 5

Given grammar  $G$ :

$$S \rightarrow aSb \mid bSa \mid \epsilon$$

Please use CYK algorithm to decide whether string  $aabbabb$  belongs to  $L(G)$ .

## Problem 6

Prove that each of these languages is not context free.

- a.  $L = \{ 0^{i^2} \mid i \geq 0 \}$ .
- b.  $L = \{ 0^i 1^j \mid i^2 \geq j \}$ .
- c.  $L = \{ \omega \in \{0, 1, 2\}^* \mid \omega \text{ has equal number of 0's, 1's and 2's} \}$ .
- d. **(Bonus)**  $L = \{ \omega \omega \omega \mid \omega \in \{0, 1\}^* \}$ .

## Problem 7

For any context-free language  $L$  and any regular language  $R$ , answer each of the following statements **True** or **False**. If your answer is **True**, give an explanation. If your answer is **False**, give a counterexample.

- a.  $L - R$  is context-free.
- b.  $R - L$  is context-free.
- c.  $S(L) = \{w | \exists v \in \Sigma^*. vw \in L\}$ .  $S(L)$  is context-free.
- d. **(Bonus)**  $H(L) = \{w | \exists v \in \Sigma^*. vw \in L \wedge |v| = |w|\}$ .  $H(L)$  is context-free. **(Hint: intersection with a regular language.)**

## Problem 8

We define an operation  $\bowtie$  for language  $L$  and  $R$  to be

$$L \bowtie R = \{w \mid w = x_1y_1x_2y_2 \cdots x_ny_n \text{ for some } n, \text{ where } x_1x_2 \cdots x_n \in L \text{ and } y_1y_2 \cdots y_n \in R, \text{ each } x_i, y_i \in \Sigma^*\}$$

- a. Show that if  $L$  is context-free and  $R$  is regular, then  $L \bowtie R$  is context-free.
- b. Show that the class of CFL is not closed under  $\bowtie$  operation.