

# FLA (Fall 2022) – Assignment 1

Name: \_\_\_\_\_ Dept: \_\_\_\_\_

Grade: \_\_\_\_\_ ID: \_\_\_\_\_

**Due: Oct. 18, 2022**

## Problem 1

Provide DFAs and REs of the following languages. In all parts, the alphabet  $\Sigma = \{0, 1\}$  and  $|v|_\omega$  means the number of substring v occurrences in string  $\omega$ .

- a.  $\{\omega \mid |01|_\omega \geq 2\}$
- b.  $\{\omega \mid \omega \text{ contains at least one } 1 \text{ and two } 0\}$
- c.  $\{\omega \mid \omega \text{ has odd length and ends with } 1, \text{ or, has even length and ends with } 0\}$
- d.  $\{\omega \mid |01|_\omega \bmod 2 \equiv 0\}$  (DFA only)

**Solution.**

## Problem 2

Let  $R = (a + b)^*a(b + c)^*b(a + c)^*$ .

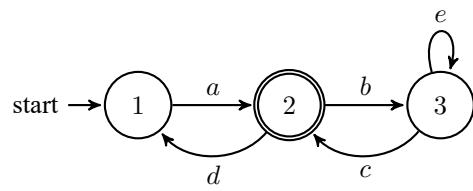
- a. Convert  $R$  to an  $\epsilon$ -NFA.
- b. Convert the  $\epsilon$ -NFA to a DFA by subset construction.

**Solution.**

### Problem 3

Give a DFA as figure below, please give the regular expression for the following  $R_{ij}^k$ , and try to simplify the expressions as much as possible.

- a. All the REs  $R_{ij}^0$
- b. All the REs  $R_{ij}^1$
- c. All the REs  $R_{ij}^2$
- d. The RE for this DFA



**Solution.**

## Problem 4

Prove that the following languages are not regular. You may use the pumping lemma and the closure properties of the class of regular languages.

- a.  $\{\omega\omega \mid \omega \in \{0, 1\}^*\}$
- b.  $\{\omega \mid |0|_\omega \geq |1|_\omega, \omega \in \{0, 1\}^*\}$
- c.  $\{\omega \mid |0|_\omega + |1|_\omega = n^2, n \in N\}$
- d.  $\{0^x \mid x \text{ is a prime number}\}$

**Proof.**

## **Problem 5**

Let a "restricted DFA" be any DFA that contains exactly one accepting state. Are restricted DFAs strictly less powerful than DFAs in terms of language expressing? If so, give an example regular language and prove that it cannot be decided by any restricted DFA. If not, prove the equivalence between DFAs and restricted DFAs.

**Proof.**

## Problem 6

Let's define a new binary operator  $\star$  for languages. Let  $L_1, L_2 \in \Sigma^*$ , where  $\Sigma$  is an alphabet. We define

$$L_1 \star L_2 = \{w | (w \in L_1 \vee w \in L_2) \wedge \neg(w \in L_1 \wedge w \in L_2)\}.$$

Is the class of regular languages closed under  $\star$ ? If so, give a sound proof. If not, give an counterexample.

## Problem 7

Prove or disprove the following statement:

- a. Every regular language has a regular proper subset.
- b.  $A, B$  are two languages over alphabet  $\Sigma$ . If neither  $A$  nor  $B$  is regular, then  $A \cup B$  is also not regular.
- c. If  $A$  and  $B$  are not regular languages and  $C$  is a language such that  $A \subseteq C \subseteq B$ , then  $C$  is not regular.
- d. An infinite regular language  $L$  can be split into two infinite disjoint, nonempty, regular subsets  $L_1, L_2$ . (**Hint: use the pumping lemma.**)

**Solution.**

## **Problem 8 (Bonus.)**

- a. For a language  $L$ , let  $A(L)$  be the language  $\{wv \mid \exists v \in \Sigma. vw \in L\}$ . Show that if  $L$  is regular, so is  $A(L)$ .
- b. For a language  $L$ , let  $B(L)$  be the language  $\{wv \mid \exists v \in \Sigma^*. vw \in L\}$ . Show that if  $L$  is regular, so is  $B(L)$ .

**Solution.**