# FARA: Future-aware Ranking Algorithm for Fairness Optimization

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## **Abstract**

Ranking systems are the key components of modern Information Retrieval (IR) applications, such as search engines and recommender systems. Besides the ranking relevance to users, the exposure fairness to item providers has also been considered an important factor in ranking optimization. Many fair ranking algorithms have been proposed to jointly optimize both ranking relevance and fairness. However, we find that most existing fair ranking methods adopt greedy algorithms that only optimize rankings for the next immediate session or request. As shown in this paper, such a myopic paradigm could limit the upper bound of ranking optimization and lead to suboptimal performance in the long term.

To this end, we propose FARA, a novel Future-Aware Ranking Algorithm for ranking relevance and fairness optimization. Instead of greedily optimizing rankings for the next immediate session, FARA plans ahead by jointly optimizing multiple ranklists together and saves them for future sessions. Particularly, FARA first uses the Taylor expansion to investigate how future ranklists will influence the overall fairness of the system. Then, based on the analysis of the Taylor expansion, FARA adopts a two-phase optimization algorithm where we first solve an optimal future exposure planning problem and then construct the optimal ranklists according to the optimal future exposure planning. Theoretically, we show that FARA is optimal for ranking relevance and fairness joint optimization. Empirically, our extensive experiments on three semi-synthesized datasets show that FARA is efficient, effective, and can deliver significantly better ranking performance compared to state-of-the-art fair ranking methods.

#### **Keywords**

Fair Ranking, Position Bias, Exposure, Exposure Fairness

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Conference acronym 'XX, June 03–05, 2018, Woodstock, NY
© 2018 Association for Computing Machinery.
ACM ISBN 978-1-4503-XXXX-X/18/06...\$15.00
https://doi.org/XXXXXXXXXXXXXXX

# 1 INTRODUCTION

Ranking systems are one of the important cornerstones of information retrieval (IR). Existing ranking systems are usually constructed to optimize ranking relevance with the Probability Ranking Principle (PRP) [35] where items of greater likely relevance should be ranked higher. The PRP is a user-centered ranking strategy that helps save users energy and time since users could satisfy their needs with the top-ranked items [19]. However, recent research has shown that, besides users, item providers also draw utility from ranking systems, and the PRP could result in severe unfairness for item providers [4, 38]. Particularly, the PRP always assigns a few top items with high-rank positions. Those top items usually get the majority of exposure while other items rarely get exposure, although other items might still be relevant [20, 30]. The unbalanced exposure leads to unfair opportunities and unfair economic gains for item providers. Such unfairness will eventually force unfairly treated providers to leave the system, and fewer options will be left for users [50]. Therefore, IR researchers have argued that ranking relevance and fairness are both important for modern ranking systems [38, 39]. Many fair ranking algorithms have been proposed to optimize both of them jointly [30, 54].

However, existing fair algorithms are mostly greedy algorithms and could only deliver suboptimal ranking performance in the long run. In particular, existing fair ranking algorithms [4, 25, 27, 38, 47] usually behave greedily to sequentially produce the locally optimal ranklist for the next immediate session without being aware of the influence of future sessions (see more discussion in §2). The unawareness could lead to unmitigated ranking conflict between relevance and fairness optimization. For example, imagine a case that there are in total 3 items in consideration, item A, item B, and item C, where item A is the most relevant one and item C is the least relevant one. Ranklist [A, B, C] is the ranklist to maximize ranking relevance. We now consider a scenario where item C is severely unfairly treated in history. To optimize exposure fairness, we need to allocate item C more exposure by boosting item C to a higher position. However, If we try to greedily boost item C within the next immediate session, it is highly likely that item C will be boosted to the first rank to get the maximum exposure and the result ranklist is [C, A, B]. However, Ranklist [C, A, B] is of poor

<sup>&</sup>lt;sup>1</sup>In this paper, we define a session as a query issued by a user.

ranking relevance due to the ranking conflict that the least relevant item (item *C*) is put on the most important rank (the first rank).

Intuitively, the ranking conflict can be smoothed if we plan ahead and jointly optimize multiple future sessions' ranklists together instead of greedily optimizing the next immediate session. For example, the multiple ranklists after joint optimization can be [[A, C, B], [A, C, B], ...], where item C is smoothly boosted in multiple ranklists and the most relevant item, i.e., item A, is still ranked the highest. Based on the above idea of planning for future, we propose FARA, a novel Future-Aware Ranking Algorithm for relevance and fairness optimization. FARA precomputes and jointly optimizes multiple ranklists together and save them for future use. Particularly, to be able to plan for the future, FARA first uses the Taylor expansion to investigate how future ranklists will influence fairness. Then, based on the influence, FARA uses a two-phase optimization to jointly optimizes multiple ranklists together for future use. In phase 1, we solve an exposure planning problem and get the optimal future planning for item exposure. In phase 2, we construct the optimal ranklists according to the optimal future planning for item exposure. Theoretically, we show that FARA is optimal for ranking relevance and fairness joint optimization. Empirically, we conduct experiments on three semi-synthesized datasets. The experiment results show that FARA is efficient, effective, and significantly outperforms existing state-of-the-art fair ranking algorithms in terms of fairness and relevance joint optimization.

#### 2 RELATED WORK

In this section, we brief the related works on fair ranking algorithms.

Ranking Fairness: Due to the importance of rankings for providers (sellers, job seekers, content creators, etc.) [9, 18, 38, 52, 53], there has been growing interest in ranking fairness for providers [5, 13, 16, 21, 26, 33, 41, 41, 47]. However, the definitions of ranking fairness vary a lot in the existing literature, and there exists no universal definition. At a high level, existing fairness definitions can be grouped into probability-based fairness and exposure-based fairness [30, 54]. Probability-based fairness [3, 6, 14] usually requires a minimum number or proportion of protected (e.g., race, gender) items to be distributed evenly across a ranklist. However, only considering the number or proportion of items in a ranklist neglects the fact that different ranks usually have different importance. To address this, exposure-based fairness [4, 8, 12, 38, 48] assigns values to each ranking position based on the expected user attention or click probability. Exposure-based fairness argues that total exposure is a limited resource for a ranking system and advocates for fair distribution of exposure among items to ensure fairness for item providers [30]. In this paper, we limit our discussion of fair ranking algorithms within the scope of exposure-based fairness.

Fair Ranking Algorithms: Recently, a few ranking algorithms [11, 16, 24, 30, 36, 41, 54] have been proposed to achieve exposure-based fairness. In this work, we classify them as **open-loop** algorithms or **feedback-loop** algorithms depending on whether historically generated ranklists are used to correct ranking scores. For open-loop fair algorithms [16, 27, 37–39, 42, 46], each item usually has a static and fixed ranking score once the ranking model is optimized. Then ranklists are stochastically sampled for each session according to the static ranking scores. Various techniques have been

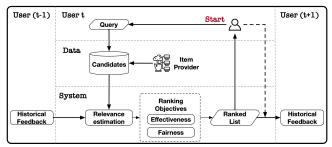


Figure 1: Ranking Service Workflow.

used to optimize the static ranking model, such as linear programming [16, 38], policy gradient [39], differentiable PL model optimization [27]. However, given the fact that ranking scores are static, open-loop algorithms are usually not robust. For example, if unfairness already happened in some historical sessions by some external disturbances, there is no way to correct the unfairness automatically in later sessions since the ranking score is fixed. The disturbances can be some online experiments, sponsored results, etc. To improve ranking robustness, feedback-loop algorithms [4, 25, 48, 51] dynamically take historical ranklists as input to correct items' scores. For example, Morik et al. [25] proposes to use a proportional controller to boost ranking scores of historically unfairly treated items. Openloop and feedback-loop fair algorithms are greedy algorithms since their ranklist optimization for the next immediate session does not consider the influence of future sessions' ranklists. Being greedy algorithms, they usually get stuck at the local optimum and fail to reach optimal ranking performance in the long run [10].

#### 3 BACKGROUND AND PRIOR KNOWLEDGE

In this section, we provide readers with background knowledge of the paper, including the workflow of ranking services, biased user feedback, and ranking metrics. A summary of notations we use throughout the paper can be found in Table 1.

**Ranking Services Workflow:** In Figure 1, we take web search as an example to detail the ranking service workflow. At time step t, a ranking session starts when user t issues a query q. For query q, there exist candidate items provided by item providers. With the query and candidate items, the ranking system first estimates each item's relevance and then constructs a ranklist of candidate items by optimizing certain ranking objectives. For example, fairness and ranking relevance can be adopted here. After ranking optimization, the ranking system presents the ranklist to users and collects users' feedback (e.g., clicks) which can be further utilized to update the relevance estimator.

**Partial and Biased Feedback:** As shown in the ranking service workflow in Fig. 1, users' feedback is usually used to update the relevance estimation. However, users' feedback is usually a noisy and biased indicator of relevance because users only provide meaningful feedback for items that users have examined. If we consider user clicks as the main signal for user feedback, then we have

$$c = \begin{cases} r, & \text{if } e = 1\\ 0, & \text{otherwise} \end{cases} \tag{1}$$

where e, r, c are binary random variables indicating whether an item is examined, perceived as relevant, and clicked, respectively.

Table 1: A summary of notations in this paper.

1 D()	F
d, q, D(q)	For a query $q$ , $D(q)$ is the set of candidates items. $d \in$
	D(q) is an item.
e, r, c	All are binary random variables indicating whether an
	item $d$ is examined ( $e = 1$ ), perceived as relevant ( $r = 1$ )
	and clicked $(c = 1)$ by a user respectively.
$R, P_i, E, \pi$	R = P(r = 1 d), is the probability of an item $d$ perceived
	as relevant. $P_i = P(e = 1   rank(d   \pi) = i)$ is the exami-
	nation probability of item $d$ when it is put in $i^{th}$ rank in
	a ranklist $\pi$ . $E$ is item's accumulated examination proba-
	bility (see Eq.7).
$k_s, k_c$	Users will stop examining items lower than rank $k_s$ due to
	selection bias (see Eq. 3). $k_c$ is the cutoff prefix to evaluate
	Cum-NDCG and $k_c \leq k_s$ .

According to the well-established User Examination Hypothesis [34], we model users' click probability as

$$P(c = 1) = P(e = 1)P(r = 1)$$
(2)

For the rest of the paper, we use R = P(r = 1) to simplify the notation. Although there exist several types of biases in the examination probability P(e = 1), we focus on two most important ones: positional bias and selection bias.

**Positional Bias** [7]: Examination probability is decided by the rank (also called position) and drops along ranks. Particularly, the examining probability is denoted as  $P_{rank(d|\pi)}$ , where  $rank(d|\pi)$  is item d's rank in ranklist  $\pi$ .

**Selection Bias** [28, 29]: This bias exists when not all of the items are selected to be shown to users, or some lists are so long that no users will examine every item in them. Assuming the items ranked lower than rank  $k_s$  will not be examined [28], we model this as:

$$P(e=1|d,\pi) = \begin{cases} P_{rank(d|\pi)}, & \text{if } rank(d|\pi) \le k_s \\ 0, & \text{otherwise} \end{cases}$$
 (3)

**Ranking Utility Measurement:** Here we introduce the evaluation of ranking performance from both the user side and the provider side, which in later sections will guide the ranking optimization.

The User-side Utility: User-side utility measures a ranking system's ability to put relevant items on higher ranks. A popular user-side utility measurement is DCG [17]. Specifically, given a query q and a ranklist  $\pi$ ,  $DCG@k_c$  is defined as,

$$DCG@k_c(\pi) = \sum_{i=1}^{k_c} R(\pi[i], q) \lambda_i = \sum_{i=1}^{k_c} R(\pi[i], q) P_i$$
 (4)

where  $\pi[i]$  indicates the  $i^{th}$  ranked item in ranklist  $\pi$ ,  $R(\pi[i], q)$  indicates  $\pi[i]$ 's relevance to query q, cutoff  $k_c$  indicates the prefix we want to evaluate,  $\lambda_i$  indicates the weight we put on  $i^{th}$  rank.  $\lambda_i$  is usually monotonically decreases as i increases, e.g.,  $\lambda_i$  is usually set to  $\frac{1}{\log_2(i+1)}$  [17]. In this paper, following previous works in [38], we set  $\lambda_i$  as the examination probability  $P_i$  at the  $i^{th}$  rank when computing DCG. Furthermore, based on DCG, we can measure multiple ranklists with the cumulative NDCG,

eff.@
$$k_c = cum$$
-NDCG@ $k_c(q, t) = \sum_{\tau=1}^{t} \gamma^{t-\tau} \frac{DCG@k_c(\pi_{\tau})}{DCG@k_c(\pi^*)}$  (5)

where  $0 \le \gamma \le 1$  is a constant discount factor and t is the current time step.  $\pi^*$  is the ideal ranklist constructed by ranking items

according to true relevance and we use  $DCG@k_c(\pi^*)$  to normalize  $DCG@k_c(\pi_\tau)$ . By ignoring  $\gamma$ , we can get the average-NDCG as,

$$aver-NDCG@k_{c} = \frac{\sum_{\tau=1}^{t} DCG@k_{c}(\pi_{\tau})}{tDCG@k_{c}(\pi^{*})}$$

$$= \frac{\sum_{d \in D(q)} R(d) \left(\sum_{\tau=1}^{t} \sum_{j=1}^{k_{c}} P_{j} \mathbb{1}_{\pi_{i}[j]==d}\right)}{tDCG@k_{c}(\pi^{*})}$$

$$= \frac{\sum_{d \in D(q)} R(d) E^{t}@k_{c}(d)}{tDCG@k_{c}(\pi^{*})}$$
(6)

where  $E^t@k_c(d)$  is the cumulative exposure at top  $k_c$  ranks,

$$E^{t}@k_{c}(d) = \sum_{\tau=1}^{t} \sum_{j=1}^{k_{c}} P_{j} \mathbb{1}_{\pi_{i}[j] = d}$$
(7)

where  $\pi_i[j]$  indicates the  $j^{th}$  item in ranklist  $\pi_i$ ,  $\mathbb{1}$  is an indicator function which means we only accumulate item d's exposure. To simplify notations, we use  $E^t(d)$  to denote  $E^t@k_s(d)$  and eff. for  $eff.@k_s$ , where  $k_s$  is the ranklist length introduced in Eq. 3.

The Provider-side Utility (Fairness): As items' rankings can have significant effects on their providers' profit, it is important to create a fair ranking environment. To evaluate whether exposure is fairly allocated to users, we use the negative exposure disparity between item pairs as fairness measurement [27],

$$unfair.(q,t) = \frac{1}{n(n-1)} \sum_{d_x \in D(q)} \sum_{d_y \in D(q)} \left( E^t(d_x) R(d_y) - E^t(d_y) R(d_x) \right)^2$$

$$fair.(q,t) = -unfair.(q,t) \tag{8}$$

$$(9)$$

where D(q) is the set of candidate items that will construct the ranklist of query q and n=|D(q)|. The intuition of the above fairness measurement is that the optimal fairness can be achieved when items exposure is proportional to their relevance, i.e.,  $\forall d_X, d_Y \in D(q), \frac{E^t(d_X)}{R(d_X)} = \frac{E^t(d_Y)}{R(d_Y)}$ . In other words, exposure fairness means we should let items of similar relevance get similar exposure. In this paper, we choose the exposure fairness evaluation proposed by Oosterhuis [27] instead of the original evaluation proposed in [38]. The reason for the choice is that the fairness evaluation in [38] needs to divide the exposure of an item by its relevance, i.e.,  $\frac{E^t(d)}{R(d)}$ , which has zero denominator problem when item d is irrelevant and R(d) is near zero. This paper uses average unfairness across different queries to evaluate a ranking algorithm. We refer to the average unfairness as the unfairness tolerance.

## 4 PROPOSED METHOD

Most existing fair algorithms are greedy algorithms, i.e., they sequentially construct the locally optimal ranklist for the next intermediate session. They usually fail to optimize the construction procedure if we expect multiple sessions will come for the same query in the future. To mitigate this gap and reach a global optimal for a query, we propose to (i) plan ahead and precompute multiple ranklists for future use (§ 4.1) and (ii) jointly optimize those ranklists together to maximize both fairness and ranking relevance (§ 4.2 & § 4.3). We hypothesize that the joint optimization of multiple ranklists together can construct better ranklists than sequentially greedily optimizing one ranklist at

each time step. This hypothesis is verified by both the theoretical analysis in § 5 and the empirical results in § 6.2.

# 4.1 Future-aware Ranking Objective

We first propose a ranking fairness objective to plan and optimize the future  $\Delta T$  ranklists for a query q. Specifically, when we are at time step t+1, the objective is to pre-compute the optimal ranklists  $\mathcal{B}^* = [\pi_{t+1}, ..., \pi_{t+\Delta T}]$  that can maximize the marginal fairness  $\Delta \widehat{fair}(q, t, t + \Delta T)$ ,

$$\mathcal{B}^* = \underset{\mathcal{B} = [\pi_{t+1}, \dots, \pi_{t+\Delta T}]}{\operatorname{arg max}} \quad \Delta \widehat{fair.}(q, t, t + \Delta T)$$
 (10)

 $\widehat{\Delta fair}(q,t,t+\Delta T)=\widehat{fair}(q,t+\Delta T)-\widehat{fair}(q,t),$  and  $\widehat{fair}$  is the estimated fairness.  $\widehat{fair}$  is calculated with Eq. 9 by substituting true relevance R with the estimated relevance, denoted as  $\widehat{R}$ , since true relevance is mostly not available during optimization. Here maximizing the marginal fairness  $\widehat{fair}(q,t,t+\Delta T)$  is equivalent to maximizing final fairness  $\widehat{fair}(q,t+\Delta T)$  at time step  $t+\Delta T$ . The reason for the equivalence is that ranklists  $\mathcal{B}=[\pi_{t+1},...,\pi_{t+\Delta T}]$  can only influence the marginal fairness from timestep t+1 to timestep  $t+\Delta T$  rather than fairness before timestep t. Here, fairness evaluation in Eq. 9 is a direct objective in our method.

To the best of our knowledge, there is no trivial algorithm to get the optimal ranklists  $\mathcal{B}^*$  due to the discontinuity of ranking problem [22]. One example of discontinuity is that increasing an item's ranking score may not change output ranklist, and fairness stays the same unless the increased score can surpass another item's score, and fairness will experience a sudden change. To alleviate the discontinuity of  $\mathcal{B}^* \leftarrow \operatorname{argmax}_B \Delta \widehat{fair}$ , we propose a novel two-phase solution path by introducing a continuous variable,  $\Delta E$ ,

$$\mathcal{B}^* \xleftarrow{\text{Vertical Allocation}} \Delta E^*(d) \ \forall d \in D(q) \xleftarrow{\text{argmax}_{\Delta E}} \Delta \widehat{fair}. \quad (11)$$

where  $\Delta E(d)$ , also referred to as the planning exposure, is the marginal (or incremental) exposure we plan to assign to item d within the next  $\Delta T$  timesteps.  $\Delta E^*(d)$  is the optimal marginal exposure. We found that introducing  $\Delta E$  helps to effectively maximize  $\Delta \widehat{fair}$ . in the phase 1 of our solution. In phase 2, we construct the optimal ranklists  $\mathcal{B}^*$  by allocating the optimal exposure  $\Delta E^*$  to each item with a vertical allocation method (more details are in §4.3).

To get  $\Delta E^*$ , we carry out a Taylor series expansion to investigate how future exposure will influence the fairness objective,

$$\begin{split} \widehat{\Delta fair.}(q,t,t+\Delta T) &= \sum_{d \in D(q)} \frac{\widehat{\partial fair.}}{\partial E(d)} \Delta E(d) \\ &+ \frac{1}{2} \sum_{d_x \in D(q)} \sum_{d_y \in D(q)} \frac{\partial^2 \widehat{fair.}}{\partial E(d_x) \partial E(d_y)} \Delta E(d_x) \Delta E(d_y) \\ &= \sum_{d \in D(q)} \widehat{G}(d) \Delta E(d) \\ &- \frac{1}{2} \sum_{d_x \in D(q)} \sum_{d_y \in D(q)} \widehat{H}(d_x,d_y) \Delta E(d_x) \Delta E(d_y) \\ &= \widehat{\vec{G}} \cdot \Delta \vec{E} - \frac{1}{2} \Delta \vec{E}^T \cdot \widehat{H} \cdot \Delta \vec{E} \end{split}$$

$$(12)$$

where  $\widehat{G}$  and  $\widehat{H}$  are the first and the second order derivative, i.e., the gradient vector and the Hessian matrix, respectively,

$$\begin{split} \widehat{G}(d) &= \frac{4}{n(n-1)} \left( \widehat{R}(d) \sum_{l} E^{t}(l) \widehat{R}(l) - E^{t}(d) \sum_{h} \widehat{R}^{2}(h) \right) \\ \widehat{H}(d_{x}, d_{y}) &= \frac{4}{n(n-1)} \left( (\sum_{d \in D(a)} \widehat{R}^{2}(d)) \mathbb{1}_{x=y} - \widehat{R}(d_{x}) \widehat{R}(d_{y}) \right) \end{split}$$

By observing Eq. 8, we could derive two facts about the above second-order expansion in Eq. 12. (i) The above second-order expansion is not an approximation, but equality since (un)fairness in Eq. 8 is defined as a polynomial of E with a degree of two, and its derivative of order higher than two is zero. (ii) Since (un)fairness in Eq. 8 is defined as a sum of squares, the second order derivative  $\widehat{H}$  is semi-definite. As an equality, Eq. 12 allows us to **correctly estimate future fairness given marginal exposure**  $\Delta E$  **even when we consider a long-term future (large**  $\Delta T$  **and**  $\Delta E$ ). Based on the correct future fairness estimation, it is possible to find **the optimal marginal exposure planning, denoted as**  $\Delta E^*$ , that can maximize future fairness. Since  $\widehat{H}$  is semi-definite, Quadratic Programming-based (QP) optimization is valid to find  $\Delta E^*$ . We give the specific QP problem formulation to find  $\Delta E^*$  in § 4.2 and leave constructing optimal ranklists  $\mathcal{B}^*$  from  $\Delta E^*$  in § 4.3.

# 4.2 Phase 1: Future Exposure Planning

When giving the QP problem formulation, we noticed that existing ranking fairness optimization usually considers two settings: (i) **the post-processing setting** [4, 38, 39] where relevance is assumed to be known or well estimated in advance; and (ii) **the online setting** [25, 48] where fairness is optimized while relevance is still being learned. To consider both settings, we first illustrate the QP problem formulation in the post-processing setting in § 4.2.1 and then extend it to work in the online setting in § 4.2.2.

4.2.1 The post-processing setting . To get the optimal exposure planning  $\Delta E^*$ , we propose the following QP formulation with  $\Delta E(d) \ \forall d \in D(q)$  as decision variables,

$$\max_{\Delta F} \quad \Delta \widehat{fair}.(q, t, t + \Delta T) \tag{13a}$$

s.t. 
$$\sum_{d \in D(q)} \Delta E(d) = \sum_{i=1}^{\Delta T} \sum_{j=1}^{k_s} P_j$$
 (13b)

$$\sum_{d \in D(q)} \Delta E(d) \widehat{R}(d) \ge (1 - \alpha) \sum_{i=1}^{\Delta T} \sum_{j=1}^{k_s} P_j \widehat{R}(d_{order_j})$$
 (13c)

$$\Delta E(d) \ge 0, \forall d \in D(q)$$
 (13d)

$$\Delta E(d) \le \sum_{i=1}^{\Delta T} P_1, \forall d \in D(q)$$
 (13e)

where  $k_s$  is the length of ranklists,  $P_j$  is the examining probability at rank j. Eq.13b indicates that the sum of items' marginal exposure should equal the sum of the  $\Delta T$  ranklists' exposure. Eq.13c indicates the NDCG requirement, where  $\widehat{R}(d_{order_j})$  means the  $j^{th}$  largest relevance according to items' relevance estimation,  $(1-\alpha)$  indicates the minimum aver-NDCG@ $k_s$  requirement we want to guarantee (refer to Eq. 6). In the post-processing setting, relevance is assumed to be given or already well-estimated prior to ranking

optimization. Eq.13d indicates that an item's marginal exposure  $\Delta E(d)$  should be no less than 0. In Eq.13e, an item's marginal exposure  $\Delta E(d)$  should be more than the accumulation of the first rank's exposure in the  $\Delta T$  ranklists because an item should be unique in a ranklist and there are  $\Delta T$  ranklists under consideration. Here we assume that the first rank's exposure is the greatest and exposure drops from the top rank to the lower ranks.

4.2.2 The online setting. In the online setting, ranklists are optimized while relevance is still being learned. How to actively explore items and get more accurate relevance for ranking optimization is essential. Yang et al. [49] show that a more accurate relevance estimation for an item can be achieved by exposing an item more because more exposure leads to more interaction with users. Based on this, we do explorations by setting a minimum exposure requirement for items and propose the following QP formulation,

$$\max_{\Delta E} \quad \Delta \widehat{fair.}(q, t, t + \Delta T) - \beta \sum_{d \in D(q)} s(d)$$
 (14a)

s.t. 
$$Eq. (13b, 13c, 13d, 13e)$$
 (14b)

$$s(d) \ge 0, \forall d \in D(q)$$
 (14c)

$$s(d) + \Delta E(d) + E^{t}(d) \ge E_{min}, \forall d \in D(q)$$
 (14d)

where  $\beta$  indicates the importance of exploration,  $s(d) \forall d \in D(q)$  are slack variables to encourage exploration. Both s(d) and  $E^t(d) \forall d \in$ D(q) are decision variables in the setting. In Eq. 14d,  $E_{min}$  is the minimum exposure requirement,  $E^{t}(d)$  is item d's exposure at time step t, and  $\Delta E(d)$  is the marginal exposure we plan to allocate to item d within the next  $\Delta T$  steps. s(d) can be interpreted as the additional exposure still needed to satisfy the minimum exposure requirement after the next  $\Delta T$  steps. When  $E^t(d) \geq E_{min}$  for item d, i.e., minimum exposure requirement is already satisfied for item d, it is straightforward that s(d) will be 0 and will not contribute to the ranking objective in Eq. 14a. When  $E^{t}(d) < E_{min}$ , i.e. item d does not meet the minimum exposure requirement, the objective in Eq. 14a will try to minimize s(d), in other words,  $\Delta E((d))$  will be boosted in order to satisfy the constraint in Eq. 14d. With more exposure, item d will be explored more. In this paper, we refer to the introduction of s as **Exploration**. And we treat  $\beta$  in Eq. 14a and  $E_{min}$  as hyper-parameters to control the degree of exploration.

As quadratic programming has been well studied, there are many available existing solvers. In this paper, we use quadratic programming library qpsolvers<sup>2</sup> within python to solve Equation 13 and Equation 14 to get the optimal exposure planning  $\Delta E^*$ .

# 4.3 Phase 2: Ranklists Construction

Following the solution path in Eq. 11, the next step is to construct the optimal ranklists  $\mathcal{B}^*$  as  $\Delta E^*$  has been solved in Phase 1. To construct ranklists  $\mathcal{B}^*$ , we should let each item get exactly its optimal exposure  $\Delta E^*$  within ranklists  $\mathcal{B}^*$ . However, we find that solution for  $\mathcal{B}^*$  is not unique, and many solutions can accurately allocate each item the optimal exposure planning  $\Delta E^*$  (see Theorem 5.2). Given the same exposure planning  $\Delta E^*$ , those solutions share the same fairness and aver-NDCG@ $k_s$ (see Theorem§ 5.4), where  $k_s$  is the length of a ranklist. In this paper, besides accurately allocating

#### Algorithm 1: Vertical Allocation

- 1 **Input**: The optimal exposure planning  $\Delta E$ , the number of planning sessions to consider  $\Delta T$ , the ranked list length  $k_s$ , relevance estimation  $\hat{R}$ :
- 2 **Initialize:** ranking lists  $\mathcal{B}^*(i) \leftarrow [\varnothing]$  for  $i \in range(\Delta T)$ , set  $\tilde{E}(d) \leftarrow 0 \quad \forall d \in D ;$ s for  $rnk \in [1, 2, ..., k_s]$  do **for** sess ∈  $[1, 2, ..., \Delta T]$  **do**  $Set1 \leftarrow \{d | \Delta E(d) - \tilde{E}(d) \ge P_{rnk} \};$  $Set2 \leftarrow \{d|d \notin \mathcal{B}^*(sess)\};$ 6 if  $Set1 \cap Set2 = \emptyset$  then  $Candidates \leftarrow Set2$  $Candidates \leftarrow Set1 \cap Set2$ 10  $d^* \leftarrow \operatorname{argmax} \widehat{R}(d)$ ; 11  $d \in Candidates$  $\mathcal{B}^*(sess).append(d^*);$ 12  $\tilde{E}(d^*) \leftarrow \tilde{E}(d^*) + P_{rnk};$ 13

14 Output:  $\mathcal{B}^*$ ;

the optimal exposure  $\Delta E^*$ , we additionally aim to find the optimal unique ranklists  $\mathcal{B}^*$  that can optimize **all top ranks' effectiveness**, i.e., aver-NDCG@ $k_c$ ,  $\forall k_c \leq k_s$ , instead of just aver-NDCG@ $k_s$ . Optimizing top ranks' effectiveness is important since users usually pay more attention to top ranks.

Inspired by [50], we propose a vertical exposure allocation method in Algorithm 1 to construct the optimal  $\mathcal{B}^*$  based on  $\Delta E^*$ . Algorithm 1 is theoretically justified to accurately allocate exposure  $\Delta E^*$  (see Theorem 5.1) and can construct optimal  $\mathcal{B}^*$  for aver-NDCG@ $k_c \ \forall k_c \le k_s$  (see Theorem§ 5.3). In Algorithm 1,  $\mathcal{B}^*$  is the generated  $\Delta T$  ranklists and  $\mathcal{B}^*(sess, rnk)$  denote the  $rnk^{th}$  rank of the  $sess^{th}$  ranklist. Algorithm 1 adopts a vertical way to fill out  $\mathcal{B}^*$  which follows the order,  $\mathcal{B}^*(sess, rnk)$ ,  $\mathcal{B}^*(sess + 1, rnk)$ ,  $\mathcal{B}^*(sess + 2, rnk), \dots, \mathcal{B}^*(\Delta T, rnk), (1, rnk + 1), (2, rnk + 1), \dots$  To fill out  $\mathcal{B}^*(sess, rnk)$ , the proposed vertical allocation first generates a feasible candidate set, i.e., Candidates, which contains items that are not selected for this session before  $(d \notin B(sess))$  but still have planned exposure left  $(\Delta E(d) - \tilde{E}(d) \ge P_{rnk})$ . Here, examination probability  $P_{rnk}$  serves as the margin,  $\tilde{E}(d)$  stores the actual exposure item d receives. Within Candidates, our algorithm selects the most relevant item from the candidate set to fill out  $\mathcal{B}^*(sess, rnk)$ .

Although inspired by the vertical method in [50], the proposed vertical allocation is different from it. Yang et al. [50] focus on a certain share of exposure to be guaranteed and have a complicated 3-step procedure, i.e., allocation, appending, and resorting, which cannot be used to allocate  $\Delta E^*$  for our problem. However, the proposed vertical allocation uses up all  $\Delta E^*$ , and the allocation procedure is less complicated and more straightforward.

# 4.4 FARA: Future-Aware Ranking Algorithm

Combining *Phase 1* and *Phase 2*, we propose a future-aware ranking algorithm for fairness optimization, **FARA**, detailed in Algorithm 2. FARA serves users in an online manner as explained in Fig. 1 where we pre-compute  $\mathcal{B}^*$ , the next  $\Delta T$  ranklists, and randomly pop out

<sup>&</sup>lt;sup>2</sup>https://pypi.org/project/qpsolvers/

#### Algorithm 2: FARA: Future-aware Ranking Algorithm

- 1 **Input**: The number of planning sessions to consider  $\Delta T$ , fairness-effectiveness tradeoff parameter  $\alpha$ . And in the online setting, we need to additionally give the exploration parameters  $\beta$  and  $E_{min}$  (see Eq. 14);
- 2 **Initialize** time step  $t \leftarrow 0$ , initialize an empty dictionary  $\mathbb{B} = \{\}$  to store ranked lists, items' exposure  $E \leftarrow 0$  and cumulative click  $cumC \leftarrow 0$ ;

```
3 while True do
        t \leftarrow t + 1;
4
        A user issues a query q_t;
        if q_t \notin \mathbb{B} then
 6
         \mathbb{B}[q_t] = [], i.e., add an empty list
 7
        if \mathbb{B}[q_t] is empty then
             Get \Delta E^* by solving Eq. 13 (post-processing) or
              Eq. 14 (online);
             Get \mathcal{B}^* with Algorithm 1;
10
             Randomly shuffle \mathcal{B};
11
            \mathbb{B}[q_t] \leftarrow \mathcal{B}^*;
12
        Pop out a ranking list from \mathbb{B}[q_t] and present it;
13
        Update cumulative click cumC and items' exposure E;
14
        Update the relevance estimation \hat{R} via Eq. 15;
15
```

one ranklist from  $\mathcal{B}^*$  when needed. When  $\mathcal{B}^*$  is used up and empty, We will re-compute  $\mathcal{B}^*$  for future  $\Delta T$  timesteps.  $\mathbb{B}[q]$  is used to store  $\mathcal{B}^*$  for query q. Besides, FARA does not depend on any specific relevance estimation model, therefore, can be seamlessly integrated into most existing ranking applications. In this paper, we follow works by [49] to use the following unbiased estimator of relevance,

$$\widehat{R}(d) = \frac{cumC^{t}(d)}{E^{t}(d)}$$
(15)

where  $cumC^t(d) = \sum_{i=1}^t \sum_{j=1}^{k_s} C_{i,j} \mathbbm{1}_{\pi_i[j]==d}$  is the cumulative clicks<sup>3</sup>. Moreover, it is worth noting that the relevance estimator can be replaced with other relevance estimators as well.

# Theoretical Analysis

Here we provide some theoretical analysis for FARA.

Theorem 5.1. Algorithm 1 can theoretically guarantee  $\Delta E(d)$  –  $\tilde{E}(d) \leq P_{k_s}$  for at least  $|D| - k_s$  items.

PROOF. Here we discuss the exposure allocation error bounds in Phase 2 of FARA, i.e.,  $|E(d) - \Delta E(d)|$ , in Algorithm 1. We noticed that there are two possible scenarios of lines 6-12 in Algorithm 1:

Scenario 1: There exists a  $(rnk^*, sess^*)$  pair where  $Set1 \cap Set2 = \emptyset$ . If this happens,  $(k_s, sess^*)$  will also have  $Set1 \cap Set2 = \emptyset$  since the size of Set1 and Set2 monotonically decrease for lower rank of the same session. As Set2 is the set of unselected items for a session, we know that there are at least  $|D| - k_s$  items in Set2, i.e.,  $|Set2| \ge |D| - k_s$ . If  $Set1 \cap Set2 = \emptyset$ , those  $|D| - k_s$  items are not in Set1. In other words, there are at least  $|D| - k_s$  items that satisfy  $\Delta E(d) - \tilde{E}(d) < P_{k_s}.$ 

*Scenario 2: Set*  $1 \cap Set 2 \neq \emptyset \ \forall (rnk, sess)$  *pair.* Due to the margin  $P_{rnk}$  in line 5 of Algorithm 1,  $\Delta E(d) \geq E(d) \ \forall d \in D$  should always hold if  $Set1 \cap Set2 = \emptyset$  never happens. By considering  $\Delta E(d) \ge$  $E(d) \ \forall d \in D \ \text{and the identity} \ \sum_{d \in D(q)} \Delta E(d) \equiv \sum_{d \in D(q)} E(d), \text{ we}$ would know that  $\Delta E(d) = \tilde{E}(d) \ \forall d \in D$ , which means exposures are perfectly allocated according to  $\Delta E(d)$ .

Combing the two scenarios, the vertical allocation in Algorithm 1 can theoretically guarantee  $\Delta E(d) - E(d) \leq P_{k_s}$  for at least  $|D| - k_s$ items. Since  $\Delta E(d) >> P_{k_s}$  and  $|D| >> k_s$  when using FARA, we still claim that Algorithm 1 correctly allocates exposure.

Theorem 5.2. Ranklist solutions to allocate  $\Delta E^*$  are not unique.

PROOF. The discussion in Sec. 5.1 is independent of how we select item  $d^*$  from the candidates set, and it does not need to be the most relevant one as we do in line 11 of Algorithm 1. So we could know that many ranklists can accurately allocate exposure, and the solution is not unique.

THEOREM 5.3. FARA can reach the optimal ranklists solution of Equation 6 with exposure planning.

PROOF. Here we provide theoretical proof that vertical allocation, i.e., phase 2, can optimize effectiveness (aver-NDCG) when exposure planning  $\Delta E$  is given. Specifically, maximizing aver-NDCG@ $k_c$  in Eq. 6 is equivalent to

$$\max \sum_{d \in D(q)} R(d) E@k_c(d)$$
 (16a)

$$\max \sum_{d \in D(q)} R(d)E@k_c(d) \tag{16a}$$
 s.t. 
$$\sum_{d \in D(q)} E@k_c(d) = Const. \tag{16b}$$

$$0 \le E@k_c(d) \le \Delta E(d) \ \forall d \in D$$
 (16c)

where normalization is ignored in Eq. 16a, the sum of top ranks exposure should be a constant in Eq. 16b, and the top ranks exposure should be less than the total exposure planning in Eq. 16c. According to Rearrangement Inequality [15], it is straightforward to know that aver-NDCG@ $k_c$ , i.e., Eq. 16a, can be optimized by letting item of greater relevance R get more exposure at top ranks, i.e., greater  $E@k_c$ . In other words, we should prioritize letting items of greater relevance R fulfill their exposure planning  $\Delta E$  at top  $k_c$  ranks since  $E@k_c$  is bounded in  $[0, \Delta E]$ . By assuming that aver-NDCG at higher ranks is more important [35], we should maximize aver-NDCG@ $k_c$ before maximizing aver-NDCG@ $(k_c + 1)$ ,  $\forall 1 \le k_c < k_s$ . As we maximize aver-NDCG from top to lower ranks, it is straightforward that the optimal way is to follow a greedy selection strategy to let an item of greater relevance R fulfill its exposure planning  $\Delta E$  at its highest possible ranks. In Algorithm 1, the proposed vertical allocation exactly follows the above greedy selection strategy to let item of greater relevance R (line 11 in Algorithm 1) fulfill its exposure planning  $\Delta E$  at the highest possible ranks (setting rank loops as the outer loop in 3 and line 4 of Algorithm 1). So it can reach optimal effectiveness at the top ranks. For a two-phase algorithm, FARA's exposure planning  $\Delta E^*$  is optimal by solving the QP problem in phase 1, and the final constructed ranklists are also optimal in phase 2.

Theorem 5.4. Effectiveness and fairness are fixed when  $\Delta E$  is fixed.

<sup>&</sup>lt;sup>3</sup>Given the page limit, we skip the proof of unbiasedness for the above estimator and refer interested readers to [49]

Table 2: Datasets statistics. For each dataset, the table below shows the number of queries, the average number of docs for each query, and the relevance annotation y's range.

Datasets	#Queries	#Aver. Docs per Query	y's range
MQ2008	800	20	0 - 2
MSLR-10k	10k	122	0 - 4
Istella-S	33k	103	0 - 4

PROOF. Given the same exposure planning  $\Delta E$ , effectiveness (aver-NDCG@ $k_s$ ) and fairness are fixed since we can substitute exposure planning  $\Delta E(d)$  for  $E^t @k_c(d)$  in Eq. 6 and substituting  $\Delta E$  for E in Eq. 8, respectively. In other words, for any ranklist  $\mathcal{B}^*$ , as long as exposure planning  $\Delta E$  can be accurately allocated in  $\mathcal{B}^*$ , the effectiveness and fairness are fixed.

#### **6 EXPERIMENTS**

# 6.1 Experimental setup

To evaluate our method, we conduct semi-synthetic experiments. We cover the experimental settings in this section. To facilitate reproducibility, we release our code  $^4$ .

**Datasets:** In this work, we use three public Learning-to-Rank (LTR) datasets: MQ2008 [31], MSLR10k<sup>5</sup> and Istella-S [23]. Datasets' statistics are shown in Table 2. MQ2008 has a three-level relevance judgment (from 0 to 2). MSLR10k and Istella-S have a five-level relevance judgment (from 0 to 4). Queries in each dataset are already divided into training, validation, and test partitions according to a 60%-20%-20% scheme. In this work, we mainly focus on comparison within the LTR tasks. However, the proposed method can be adapted to recommendation tasks by considering personal relevance, which we leave for future studies.

Baselines: In this paper, we compare the following methods,

- **TopK**: Sort items according to  $\widehat{R}(d)$
- RandomK: Randomly rank items.
- FairCo [25]: Fair ranking algorithm based on a proportional controller.  $\alpha \in [0.0, 1000.0]$
- MCFair [51]: Fair ranking algorithm directly uses gradient as the ranking score.  $\alpha \in [0.0, 1000.0]$
- ILP [4]: Fair ranking algorithm based on Integer Linear Programming (ILP). $\alpha \in [0.0, 1.0]$
- LP [38]: Fair ranking algorithm based on Linear Programming (LP).α ∈ [0.0, 1000.0]
- MMF [48]. Similar to FairCo but focus on top ranks fairness.  $\alpha \in [0.0, 1.0]$
- **PLFair** [27]. A fair ranking algorithm based on Placket-Luce optimization.  $\alpha \in [0.0, 1.0]$
- **FARA-Horiz.** (ours): A variant of FARA. Compared to FARA, we switch line 3 and line 4 in Algorithm 1 to first iterate the sessions and then iterate the ranks. We refer to the iterations as the **horizontal allocation** paradigm.  $\alpha \in [0.0, 1.0]$
- FARA (ours). The proposed fair ranking algorithm.  $\alpha \in [0.0, 1.0]$  Among the above ranking algorithm, TopK and RandomK are unfair algorithms, while the others are fair algorithms. While all the

fair ranking algorithms aim to maximize effectiveness and fairness, FARA and FARA-Horiz. differ from others by taking a joint optimization across multiple ranklists rather than a traditional greedy optimization approach. For fair ranking algorithms, there exists a tradeoff parameter  $\alpha$ , similar to  $\alpha$  in Eq. 13, to balance effectiveness and fairness. For fair algorithms, the greater  $\alpha$  is, the more we care about fairness while potentially sacrificing more effectiveness. For example, when increasing  $\alpha$  in Eq. 13c, FARA can maximize fairness with less effectiveness constraint. For different fair algorithms,  $\alpha$ lies in different ranges. For FairCo, MCFair, LP,  $\alpha$  are originally within  $[0.0, +\infty]$ , and we adopt  $\alpha \in [0.0, 1000.0]$  which is enough according to our experiments. For ILP, MMF, PLFair, FARA-Horiz. and FARA,  $\alpha \in [0.0, 1.0]$ . Although the vertical allocation in Algorithm 1 was inspired by [50], [50] cannot be used as a baseline because [50] works with offline ranking services where all user queries are known in advance. However, in this paper, we consider the online services depicted in Figure 1.

**Ranking Service Simulation:** Following the workflow in Figure 1, at each time step, a simulated user will issue a query q, which is randomly sampled from the training, validation, or test partition. Corresponding to the query q, a ranking algorithm will construct a ranklist  $\pi$  of candidate items and present it to the simulated user. To collect users' feedback for the ranked list  $\pi$ , we need to simulate relevance and examination (see Eq.2). Same as [2], the relevance probabilities of each document-query pair (d,q) are simulated with their relevance judgement y as

$$P(r = 1|d, q) = \epsilon + (1 - \epsilon) \frac{2^{y} - 1}{2^{y_{max}} - 1}$$

where  $y_{max}$  is the maximum value of relevance judgement y, i.e., 2 or 4 depending on the datasets. Besides relevance, following [25, 29], we simulate users' examination probability as,

$$P(e=1|d,\pi) = \begin{cases} \frac{1}{\log_2(rank(d|\pi)+1)}, & \text{if } rank(d|\pi) \le k_s \\ 0, & \text{otherwise} \end{cases}$$

For simplicity, we only simulate users' examination behavior on top ranks, and we set  $k_s$  to 5 throughout the experiments (refer to Eq. 3 for more details of  $k_s$ ). With  $P(r=1|d,q,\pi)$  and  $P(e=1|d,\pi)$ , we sample clicks with Equation 2. The advantage of the simulation is that it allows us to do online experiments on a large scale while still being easy to reproduce by researchers without access to live ranking systems [29]. For simplicity, same as existing works [25, 29, 48, 49], we assume that users' examination  $P(e=1|d,\pi)$  is known in experiment since many existing works [1, 2, 32, 45] have been proposed to estimate it. Due to different data sizes, we simulate 200k steps for MQ2008 and 4M steps for MSLR10k and Istella-S.

**Experiment Settings:** We noticed that LP and ILP methods are proposed in **the post-processing setting**, where relevance is already known or well estimated in advance. However, in most real-world settings, ranking optimization and relevance learning are carried out at the same time, which we refer to as **the online setting**. To give a comprehensive comparison, we evaluate ranking methods in both settings with the online services depicted in Figure 1. In the post-processing setting, all the ranking methods in Section 6.1 are based on true relevance R, and FARA will set β as 0. In the online setting, all the ranking methods in Section 6.1 are based on the

<sup>&</sup>lt;sup>4</sup>https://github.com/Taosheng-ty/WSDM23-MCFair.git

<sup>&</sup>lt;sup>5</sup>https://www.microsoft.com/en-us/research/project/mslr/

relevance estimation  $\widehat{R}$  in Eq. 15 to perform ranking optimization. FARA set  $\beta$  to 1 and  $E_{min}=10$  unless otherwise explicitly specified, as they work well across all of our experiments.

**Evaluation:** We use the cum-NDCG (cNDCG) in Eq. 5 with  $\gamma = 0.995$  (same  $\gamma$  adopted in [43, 44]) to evaluate the effectiveness at different cutoffs,  $1 \le k_c \le 5$ . Aside from effectiveness, unfairness defined in Eq. 8 is used for unfairness measurement. We run each experiment five times and report the average evaluation performance on the test partition. We use the Fisher randomization test [40] with p < 0.05 to do significant tests. Due to the time cost (see Table 3), we do not run ILP and LP on the larger datasets, MSLR10k and Istella-S, and the performances are not available (NA).

# 6.2 Results and Analysis

In this section, we first compare the ranking relevance performance given different degrees of fairness requirements. Then we compare fair algorithms' fairness upper bound, i.e., maximum fairness. Finally, we dive deep into our method FARA to offer more insights into FARA's supremacy.

6.2.1 Can FARA reach a better balance between fairness and effectiveness? In Figure 2, we compare ranking methods' effectivenessfairness balance given different fairness requirements. To generate the balance curves in Figure 2, we incrementally sample  $\alpha$  from the minimum value to the maximum value within  $\alpha$ 's ranges indicated in Section 6.1. For each method, twenty  $\alpha$  are sampled with the step size as  $(\alpha_{max} - \alpha_{min})/20$ . After sampling, we perform ranking simulation experiments for each  $\alpha$  to get a (cNDCG, unfairness) pair. Then we connect different  $\alpha$ 's (cNDCG), unfairness) pair to form a curve for each method respectively in Figure 2. All the curves start from the top right to the bottom left as  $\alpha$  increases, which means there exists a tradeoff between fairness and effectiveness (cNDCG). The reason behind this tradeoff is that requiring more fairness will bring more constraints on optimizing effectiveness. Since TopK and RandomK do not have trade-off parameters, both of them only have one single pair of (cNDCG, unfairness), and their performances are shown as single points in Figure 2.

In Figure 2, our methods FARA and FARA-Horiz. outperform all other fair methods since our methods reach the best cNDCG given the same unfairness tolerance. And FARA's supremacy is consistent in both post-processing and online settings. All fair ranking algorithms are effective fair ranking algorithms since they all show the tradeoff, i.e., higher cNDCG when increasing the unfairness tolerance. For unfair algorithms, TopK performs differently in postprocessing and online settings. In the post-processing setting, TopK reaches the highest cNDCG since relevance is known, and ranking relevance is the only consideration. However, in the online setting, TopK can not reaches the highest cNDCG. We think the drop in cNDCG is that TopK naively trusts the relevance estimation without any exploration when optimizing ranking relevance. However, fair algorithms are shown to be robust to the online setting since they mostly can reach better cNDCG than Topk when increasing unfairness tolerance. We think the reason for the robustness is that fair algorithms usually rerank items for different sessions to optimize fairness, and such reranking brings explorations.

6.2.2 What is the fairness upper bound that FARA can reach? In Table 3, lower unfairness means higher fairness capacity and fairness

upper bound, i.e., the maximum possible fairness one algorithm can reach. Fair ranking algorithms, including FairCo, LP, FARA-Horiz. and FARA, have similar fairness capacity and outperform unfair ranking algorithms in terms of unfairness. The success of FARA-Horiz, and FARA validates the proposed quadratic programming formulation can optimize fairness. For other fair ranking algorithms, ILP and MMF, and PLFair show inferior fairness capacity. As for the possible reason, ILP uses the integer linear programming method, which may not be effective in optimizing fairness. MMF actually follows a slightly different definition of fairness which require fairness at any cutoff should be fair, which is more strict than the definition we use in this paper. As for PLFair, PLFair tries to learn the ranking score that optimizes fairness based on the feature representation (the exact setting in original paper [27]). However, the feature representation is initially designed for relevance which makes PLFair suboptimal for fairness optimization.

6.2.3 How is FARA's effectiveness at different cutoffs? In Table 3, we show cNDCG at different cutoffs. Although FairCo, LP, FARA-Horiz. and FARA have similar fairness capacities, FARA significantly outperforms those fair algorithms for cNDCG@1 and cNDCG@3 on all three datasets. Compared to FARA-Horiz., shown in Table 3, FARA still significantly outperforms FARA-Horiz. at top ranks, which shows the necessity of vertical allocation.

6.2.4 How is FARA's time efficiency? Besides fairness and effectiveness optimization, we also empirically compare the time efficiency. In Table 4, ILP and LP are really time-consuming, especially on large datasets, MSLR10k and Istella-S. Compared to ILP and LP, FARA is more than 1000× time efficient on MSLR10k and Istella-S, although all three of them are programming-based methods. There are two reasons behind FARA's time efficiency. The first one is that FARA has a much fewer number of decision variables since FARA only has O(n) decision variables, while ILP and LP have  $O(n^2)$ decision variables. The second one is that FARA does not need to solve quadratic programming for every time step. By solving quadratic programming once, we can get  $\Delta T$  ranklists used for future  $\Delta T$  sessions. FARA can reach comparable time efficiency with non-programming-based algorithms like TopK, RandomK, and FairCo. Compared with those non-programming-based algorithms, the slightly additional time cost of FARA is acceptable given FARA' superior ranking performance in Tab. 3 and in Fig. 2.

6.2.5 How does  $\Delta T$  influence FARA? In Figure 3, we show the results of cNDCG and unfairness by varying the value of  $\Delta T$ . With greater  $\Delta T$ , we see a clear boost of cNDCG for FARA, while such a boost does not happen for FARA-Horiz. We know that the proposed vertical allocation is the key reason to have better-ranking relevance when we get the optimal exposure planning  $\Delta E^*$ , and we theoretically analyze the reason in § 5.3. Besides, as we increase  $\Delta T$ , unfairness does not vary much, and its value stays close to the minimum unfairness we can achieve in Table 3. We think the reason for the steady value of unfairness is that FARA already reaches the upper limit of fairness when  $\Delta T$  is small, and it is hard to improve when we increase  $\Delta T$ .

6.2.6 How does exploration influence FARA in the online setting? To study how the exploration part (the slack variables *s* in Eq. 14a) influences FARA, we did an ablation study for FARA with or without exploration. Due to the page limit, we only show the ablation

Table 3: Comparison of cNDCG@(1,3,5) and unfairness tolerance in the post-processing setting. Significant improvements or degradations with respect to FairCo are indicated with +/-. Within fair algorithms, the best performance with statistical significance is bolded and underlined. Here,  $\alpha$  is set to the maximum value (see Sec. 6.1 for  $\alpha$ 's range) for each fair algorithm respectively, and the unfairness here means the least unfairness a fair algorithm has to tolerate. In other words, lower unfairness in this table means higher fairness capacity or upper bound. ILP and LP are NA for MSLR10k and Istella-S due to time costs (refer to Table 4). LP, FairCo, FARA-Horiz, and FARA reach similar cNDCG@5 and unfairness, and there is no significantly better one. Similar cNDCG@5 and unfairness are expected according to Theorem 5.4. However, FARA significantly outperforms all other fair algorithms in cNDCG@1 and cNDCG@3. Results are rounded to one decimal place.

									_			
Methods	MSLR-10k			Istella-S			MQ2008					
	cNDCG@1	cNDCG@3	cNDCG@5	unfair.	@1	@3	@5	unfair.	@1	@3	@5	unfair.
TopK	200.0+	200.0+	200.0+	4165.0 -	200.0+	200.0+	200.0 <sup>+</sup>	310.1	200.0+	200.0+	200.0+	86001.1 -
Randomk	68.0-	$74.7^{-}$	79.7	119.0 -	30.2	35.7 -	41.1 -	56.7 -	74.0 -	95.7	$114.6^{-}$	$104632.2^-$
PLFair	68.2 -	74.8 -	79.9	119.6 -	31.8 -	36.2	41.5 -	54.6 -	79.2	99.3	117.2 -	101245.1
MMF	84.4	92.8	99.9	8.0	62.2	68.8	80.1	6.6	132.8	162.3	$172.5^{-}$	$20688.7^{-}$
ILP	NA	NA	NA	NA	NA	NA	NA	NA	185.6 <sup>+</sup>	183.8	186.7	19916.3 -
LP	NA	NA	NA	NA	NA	NA	NA	NA	188.4 <sup>+</sup>	187.4 <sup>+</sup>	187.9	9425.7
MCFair	114.8+	102.7+	101.0	0.0	113.7 +	85.25 +	81.34	0.4	193.5 <sup>+</sup>	186.0 <sup>+</sup>	186.6	9113.7
FairCo	85.5	93.7	100.8	0.0	63.3	69.9	80.4	0.5	179.0	182.0	187.4	9382.0
FARA-Horiz.(Ours)	90.7+	96.1 <sup>+</sup>	100.4	0.0	78.5 <sup>+</sup>	79.8 <sup>+</sup>	82.6	0.9	187.3 <sup>+</sup>	186.1 <sup>+</sup>	187.0	9125.9
FARA(Ours)	<u>129.0</u> <sup>+</sup>	<u>107.0</u> +	99.7	0.0	<u>135.5</u> <sup>+</sup>	<u>89.1</u> +	82.6	0.9	<u>196.3</u> +	<u>190.9</u> +	186.8	9129.9
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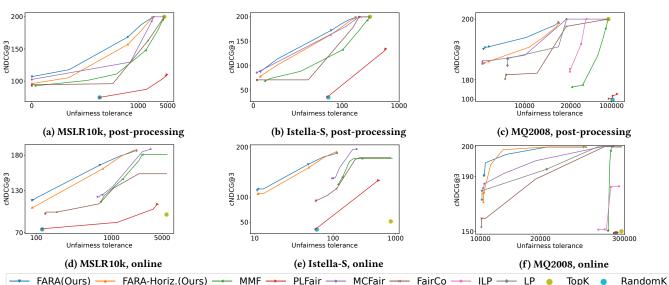


Figure 2: c-NDCG vs. unfairness tolerance (Eq. 8) in the post-processing setting and the online setting. Given the same unfairness, the higher curves or points lie, the better their performances are. Our methods FARA and FARA-Horiz. lie higher than all fair baselines in all figures. ILP and LP are unavailable for MSLR10k and Istella-S due to time costs (refer to Table 4).

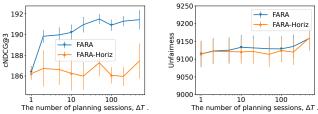


Figure 3: The numbers of planning session  $\Delta T$ 's influence on FARA in the post-processing setting on MQ2008.  $\alpha$  is set as 1.

results on the larger dataset, i.e., MSLR10k and Istella-S, in Figure 4. The advantage of exploration is two-folded based on Figure 4. Firstly, FARA lies higher than FARA-w/o-Exp. in the figure, which

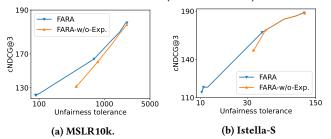


Figure 4: Ablation study of exploration in the online setting. The higher curves lie, the better their performances are.

suggests exploration leads to a better effectiveness-fairness balance. Secondly, FARA has a smaller lower bound of unfairness tolerance,

Table 4: The average time (seconds per 1k ranklists) cost with standard deviations in parentheses. Since ILP and LP are time-consuming on large datasets, the time costs on MSLR-10k and Istella-S are estimated by only running 1k steps instead of the total simulation steps indicated in Sec. 6.1.

Algorithma	Datasets						
Algorithms	MSLR10k	Istella	MQ2008				
TopK	$0.65_{(0.14)}$	$0.50_{(0.00)}$	$0.55_{(0.10)}$				
Randomk	$0.63_{(0.12)}$	$0.57_{(0.04)}$	$0.59_{(0.14)}$				
PLFair	2.24(0.04)	3.11 <sub>(0.07)</sub>	1.77 <sub>(0.04)</sub>				
MMF	$8.01_{(0.39)}$	$6.57_{(0.23)}$	$1.82_{(0.28)}$				
ILP	1208.90(85.80)	1102.30 <sub>(75.20)</sub>	$19.70_{(1.29)}$				
LP	≥10 days	≥10 days	$2.09_{(0.48)}$				
MCFair	$0.724_{(0.016)}$	$0.660_{(0.025)}$	$0.567_{(0.035)}$				
FairCo	$0.73_{(0.04)}$	$0.71_{(0.03)}$	$0.70_{(0.12)}$				
FARA-Horiz.(Ours)	$1.00_{(0.17)}$	$0.86_{(0.07)}$	$0.97_{(0.22)}$				
FARA(Ours)	$0.91_{(0.07)}$	$0.91_{(0.00)}$	$0.97_{(0.30)}$				

which implies exploration enables FARA to have a higher fairness capacity and can meet a more strict fairness requirement.

## 7 Conclusions

In this work, we study the problem of ranking relevance and fairness joint optimization. We found that existing fair ranking algorithms are usually greedy algorithms that sequentially produce a locally optimal ranklist for each session. To reach a global optimum, we propose FARA, a future-aware ranking algorithm for fairness, which optimizes multiple future sessions' ranklists together. With extensive experiments on semi-synthesized datasets, FARA achieves better performance compared to existing fair ranking algorithms.

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