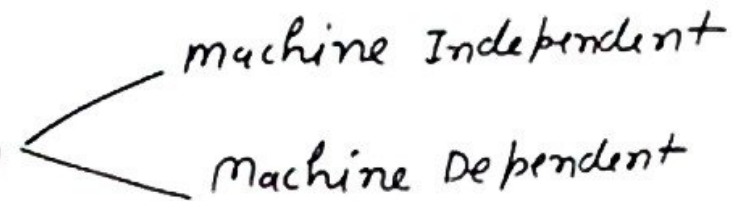


CODE OPTIMIZATION

The Code optimization is required to produce an efficient code. The improvement over Intermediate Code by transformation is called optimization.

Code optimization 
machine Independent
Machine Dependent

there are three Criteria that we applied optimization.

- (1) must preserve meaning of program.
- (2) speedup program by measurable amount of time.
- (ii) must worth the effort.

Term used in Code optimization.

- 1. Basic Blocks :- Break Intermediate Code into Blocks
- 2. flow Graph \leftrightarrow Graphical Representation of Basic Block

Ex: Convert following program into Three Address Code and
 Identify the Basic Blocks.



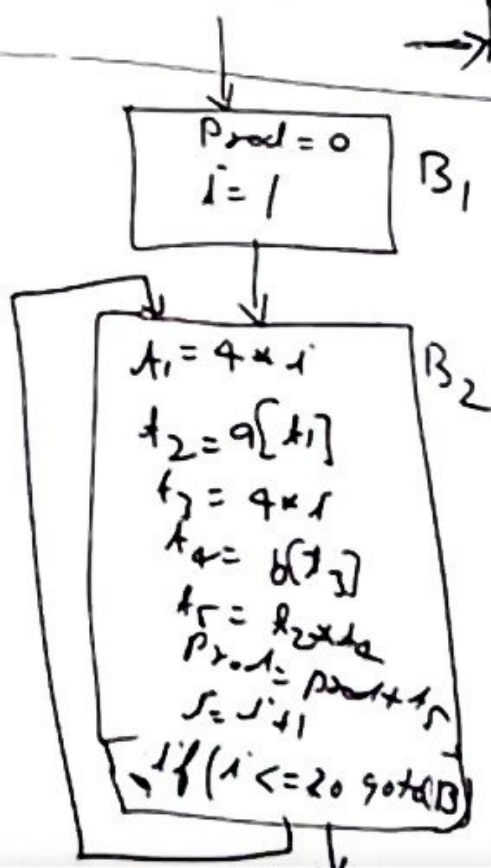
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```

begin
  Prod = 0;
  i = 1;
do
  begin
    Prod = Prod + a[i] * b[i]
    i = i + 1;
  end
while i <= 20
end
  
```

```

→ 1. Prod = 0
   2. i = 1
→ 3. t1 = 4 * i
   4. t2 = a[t1]
   5. t3 = 4 * i
   6. t4 = b[t3]
   7. t5 = t2 * t4
   8. Prod = Prod + t5
   9. i = i + 1
→ 10. if i <= 20 goto (3)
  
```



CODE OPTIMIZATION

Local

Global

1. Function preserving transformation.

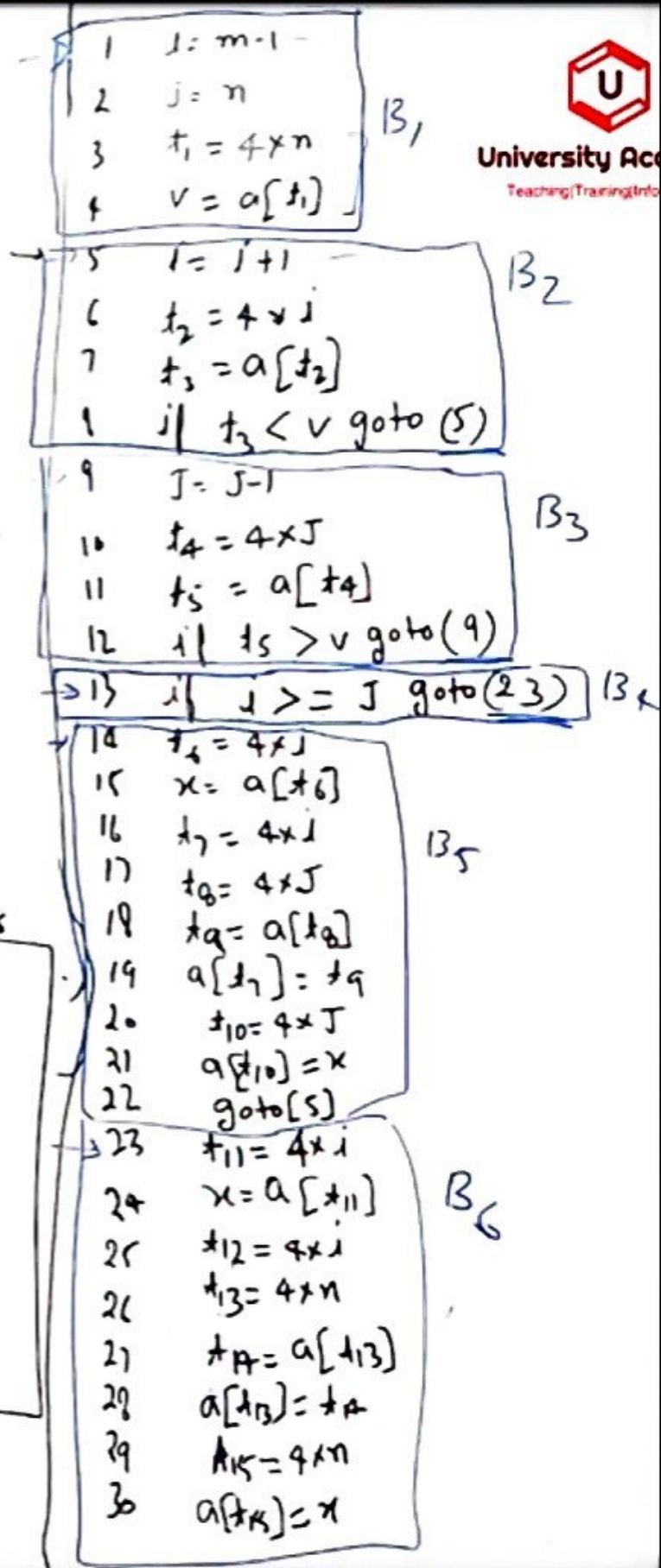
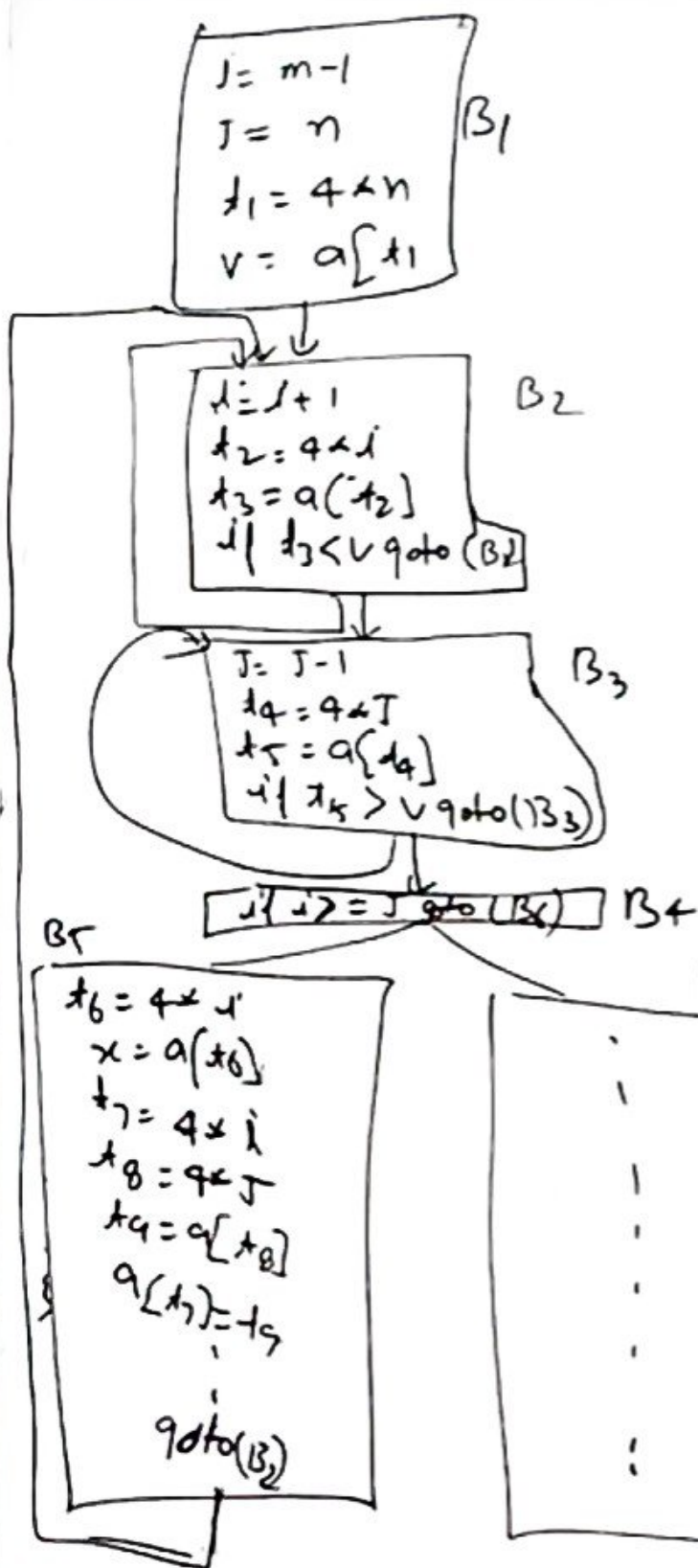
- (a) Common Sub expression elimination
- (b) Copy propagation
- (c) dead code elimination
- (d) Constant folding

2. Loop Optimization.

(a) Code motion

(b) Induction variable.

(c) Reduction in strength.



CODE OPTIMIZATION

Local

Global

1. Function Preserving transformation.

(a) Common Sub expression elimination. (b) Copy propagation.

(c) dead code elimination.

(d) Constant folding

2. Loop Optimization.

(a) Code motion.

(b) Induction Variable.

(c) Reduction in strength.

Common Sub expression Elimination:

$x = t_3$
 $a[t_2] = t_5$
 $a[t_4] = t_3$
 goto B_2

$a[t_2] = t_5$
 $a[t_4] = t_3$
 goto B_2

B_5

$x = t_3$
 $a[t_2] = t_5$
 $a[t_4] = t_3$
 goto B_2

B_5

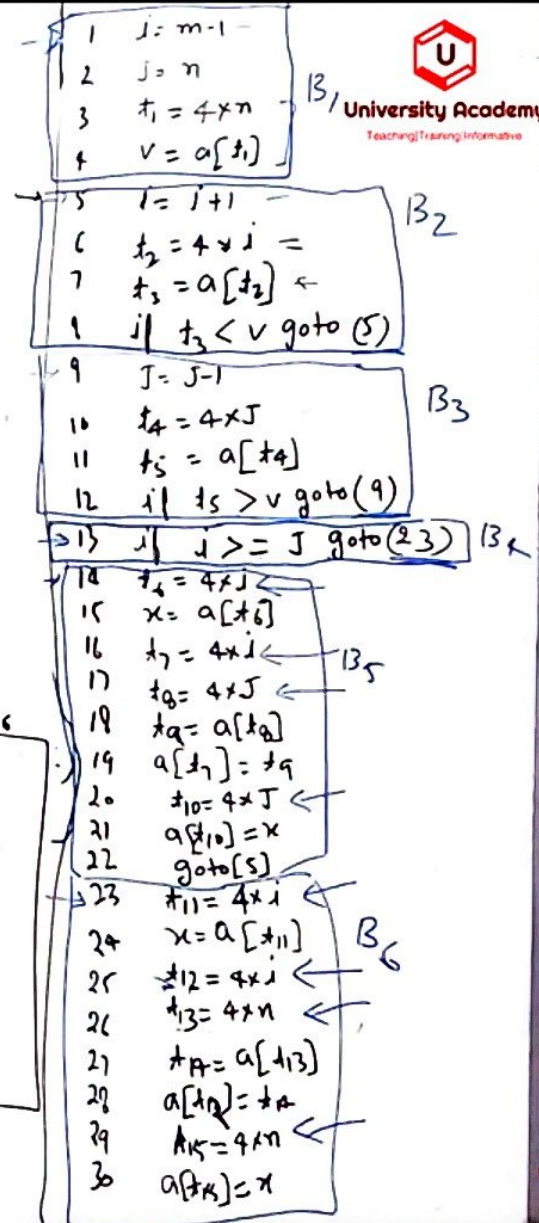
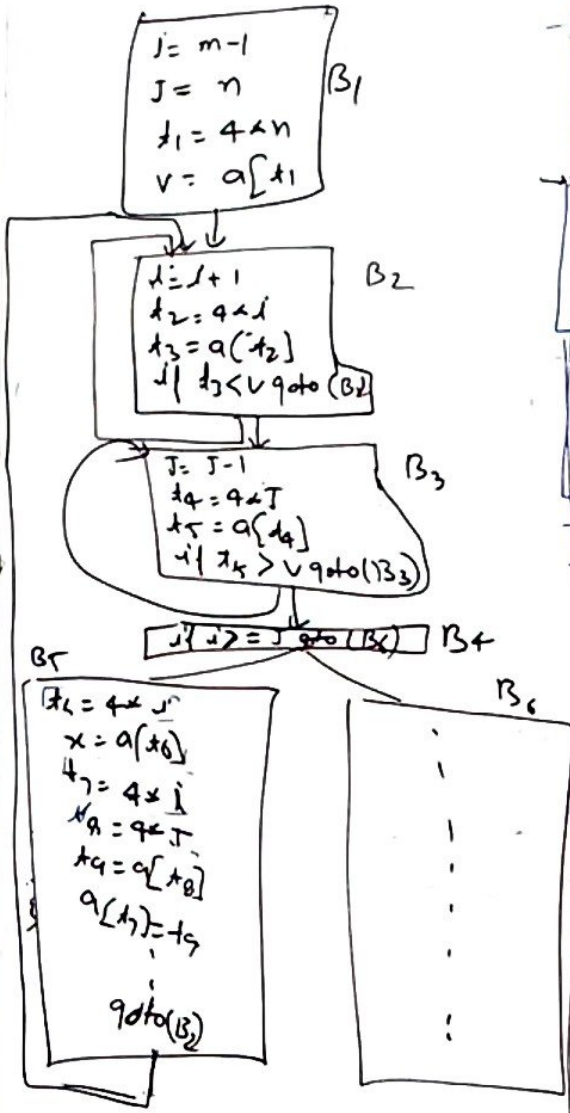
$t_{11} = 4 \times i$
 $x = a[t_{11}]$
 $t_{13} = 4 \times n$
 $t_{14} = a[t_{13}]$
 $a[t_{11}] = t_{14}$
 $a[t_{13}] = x$

B_6

$x = a[t_2]$
 $t_9 = a[t_4]$
 $a[t_2] = t_9$
 $a[t_4] = x$
 goto B_2

$x = t_3$
 $t_9 = t_5$
 $a[t_2] = t_9$
 $a[t_4] = x$
 goto B_2

B_5



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Loop Unrolling.

```
int i = 1
while (i < 100)
{
    a[i] = b[i];
    i++;
}
```



```
int i = 1
while (i < 100)
{
    a[i] = b[i];
    i++;
    a[i] = b[i];
    i++;
}
```

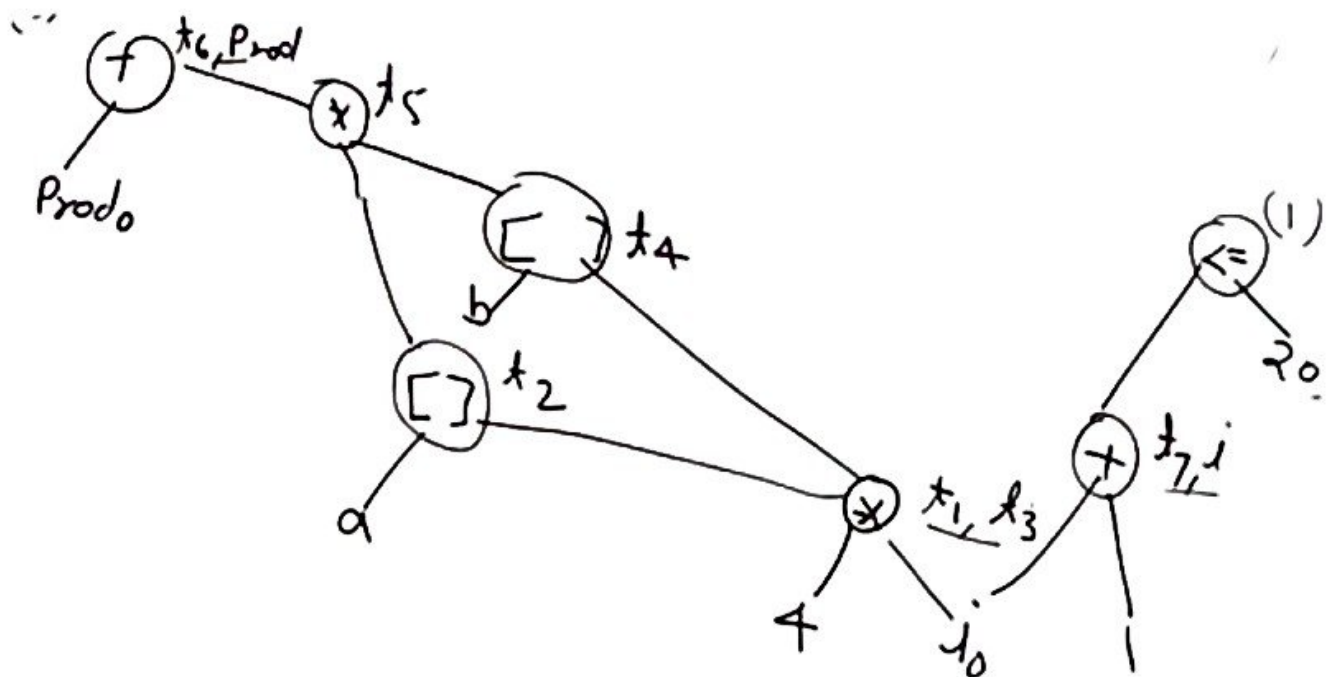
Loop fusion or Loop tumbing.

```
for i = 1 to n do
    for j = 1 to m do
        a[j] = 10;
```

CODE OPTIMIZATION

Application of DAG

1. Determining Common subexpression. and eliminate it.
2. Determining which Name. are used inside but not Valued outside the Block.
3. Determining which statement of Block could have their computed value but used outside the Block.



CODE OPTIMIZATION

Application of DAG

1. Determining Common subexpression and eliminate it.
2. Determining which Name are used inside but evaluated outside the Block.
3. Determining which statement of Block could have their computed value but used outside the Block.

1. $a = b * c$
2. $d = b$
3. $e = d * c$
4. $b = e$
5. $f = b + c$
6. $g = f + d$

Ex Block of three Address Code is given below.

$t_1, t_2, t_3, t_4, t_5, t_6, t_7$ University Academy
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1. $t_1 = 4 * i$
2. $t_2 = a[t_1]$
3. $t_3 = 4 * i$ X
4. $t_4 = b[t_3]$
5. $t_5 = t_2 * t_4$
6. $t_6 = prod + t_5$
7. $prod = t_6$
8. $t_7 = i + 1$
9. $i = t_7$
10. if $i \leq 20$ goto (1)

B₁

1. $t_1 = 4 * i$
2. $t_2 = a[t_1]$
3. $t_4 = b[t_1]$
4. $t_5 = t_2 * t_4$
5. $prod = prod + t_5$
6. $i = i + 1$
7. if $i \leq 20$ goto (1)

CODE GENERATION

Code generation is the final activity of compiler.
code generation is the process of creating Assembly / machine language. There are some properties of code generation

- (i) Correctness
- (ii) High Quality
- (iii) Efficient use of resource of target machine.
- (iv) Quick code generation.

Absolute Machine Code: Fixed location in memory and immediately executed. Useful for small program

Relocatable Machine Code: Not a fixed location, code can be placed wherever linker finds the room in RAM.
Useful for commercial compilers

Intermediate code generation
or
Code optimization.

Intermediate code:

- (i) Three Address code
- (ii) Quadruple
- (iii) Triples
- (iv) Postfix Notation
- (v) Syntax tree / DAG

Symbol
table

Code generation

error
handler

Target code / Assembly code.

Assembler

Machine code

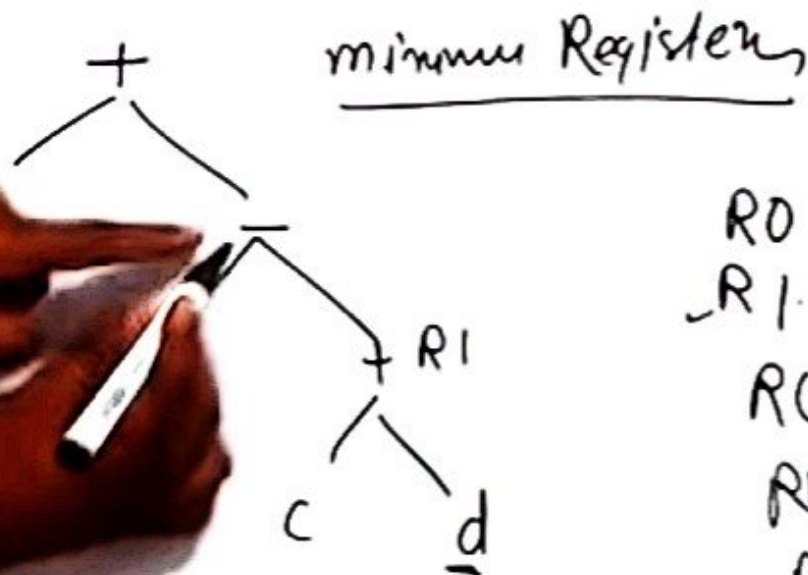
- (i) Absolute Machine code
- (ii) Relocatable Machine code



CODE GENERATION

Issue in code generation

1. Input to Code generation.
2. Target program.
3. Memory Management
4. Instruction selection
5. Register Allocation.



$R0 \leftarrow a$

$R1 \leftarrow b$

$R0 \leftarrow R0 - R1$

$R1 \leftarrow c$

$R2 \leftarrow d$

$R1 \leftarrow R2$ Swipe again

$R2 \leftarrow e$

$R2 - R1$

CODE GENERATION

PEEPHOLE OPTIMIZATION

A statement by statement code generation strategy generate target code that contain redundant instruction. to optimize such target code we use Peephole optimization technique.

1. Redundant Instruction

(i) mov R0, a
(ii) mov a, R0 } eliminate (ii) instruction.

(2) Flow of control optimization

```
goto test
-----
test goto done
-----
done
```

⇒

```
goto done
-----
test goto done
-----
done
```

.in



③ Algebraic simplification.

$x = x + 0$
or
 $x = x * 1$ } eliminate such instruction

④ Reduction strength.

Addition is cheaper than multiplication

Subtraction is cheaper than division.

⑤ Machine idioms.

$x = x + 1$ → Remove and use Auto Increment.

$x = x - 1$ → Remove and use Auto decrement.