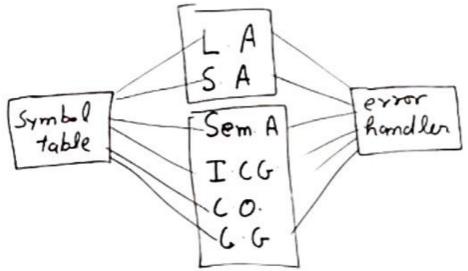
Symbol Table.

Symbol table 18 a data structure created and common maintained by compiler in order to store Information about variables, function, class, object ele



Formate of Symbol table. Compiler was following type Information in symbol table. 1. Data tipe, 3. Name, 3 Scobe 4. Address 5. other 4 tribule.

eg static int a;

SNI	Name	Tybe	Attribute
1.	a	int	Static
1	6	flogi	+ -

Symbol table Representation

University Academ

Teaching[Training]Informative

I Fixed length 2. Vanable length

Eramble int calculate; int sum; int a, b,

Nome	Type.
a E clullate	e Int
sium IIII	int
a	int
	int
plinifi	\ .,,
-1401	•

longh	tope
10	int
4	int
2	121+
2	Int
	10

calculate \$ 84 m \$ 9 9 15 15 17

Symbol Table

Operation on Symbol table.

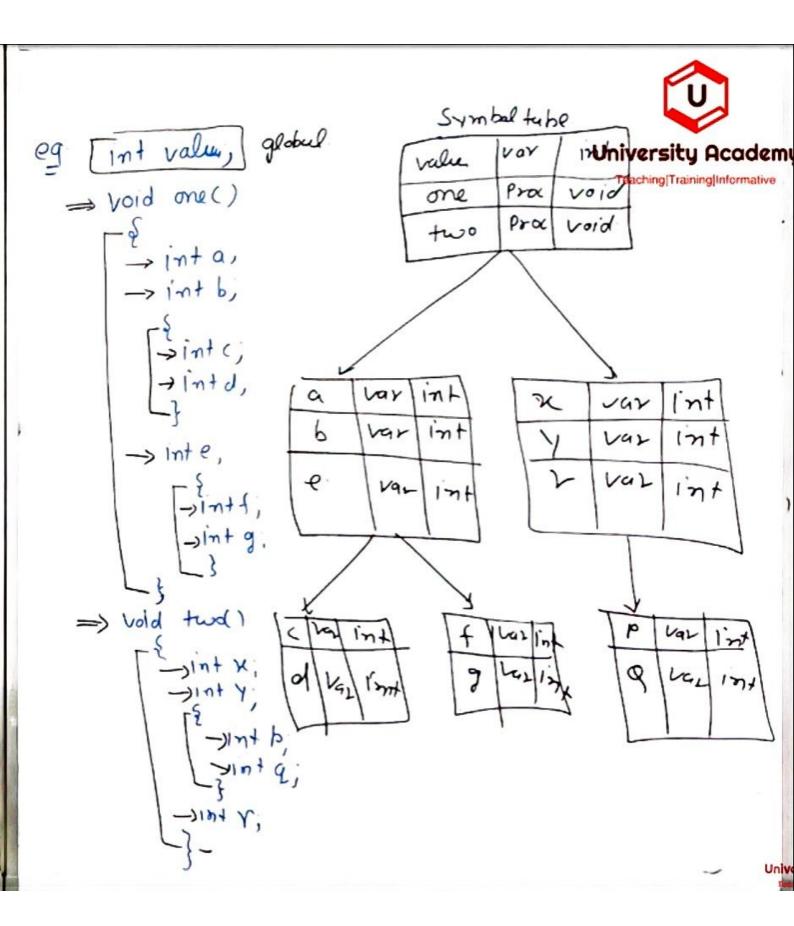
Nam Pape.

Insert(): insert (a, inf)

2 Lookub(): Lookub(symbol):- Lookub(a)

3. Delete(): delete(Simbul): delete(a)

4 Scope mgmt: local and Global Vanil.



Symbol Table

Implimentation of symbol table.

1. <u>linear list</u>: Linear list is simplest way to implement symbol table. An Array used to stor information

Name	Attribute	searching
id,	info	Searching
idz	infoz	
	\	
idn	intoz	7
s		name name
	id,	ida info, ida info,

2. Self organizing list this symbol table Imblementus, wing linked list. A link field added to each record

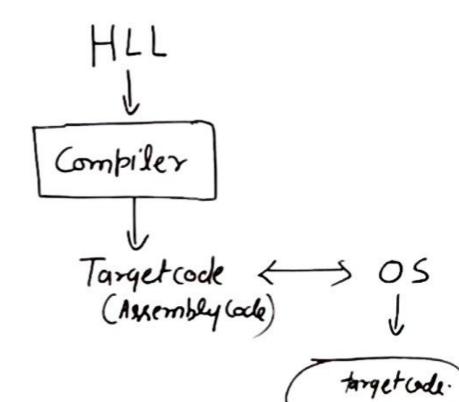
Multin and	Name	Attribute	link	
Γ	id,	into	\	7)
	ldz	intoz		1
	1d3	intoz		Y
	Ida	inlor	-	
Available				

4. Hash table. hashing is Most powerfull Implementations
Technique in symbol table. in hashing scheme two
table are maintained.

~	hash table	Nam	Allin	sley	
Ø	Sum	Sum	Int	The state of the s	\
-1 -2			jnt		
		5	int		
ika	ava	Avail.	flaut		
4(0)		funct h(nomu)	returns	integer hele o	ķ

Unive

Run time Storage Administration.



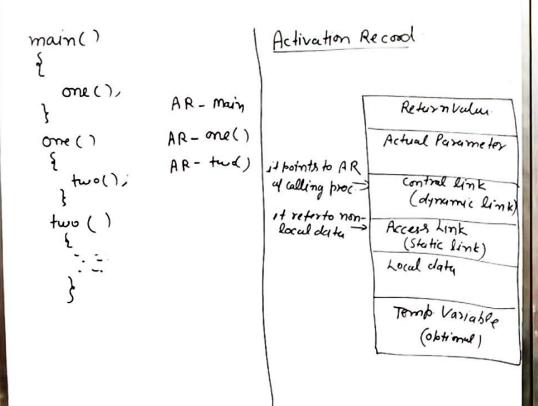
- 1 Implementation of Asimple Stack Allocation scheme.
- me mory/ Storage
- @ Implementation of Black structured lang.



higher end. dymanuc Heap memosy Allocation -> Tob Stack Conviant, and global variable compile time. Static / Global dal code Area Fixed size & hold target cale lower end

Run time Storage Administration.

1) Simple stack Allocation scheme:



main ()

{
 int f;
 f = fact(3);
 }
 int fact (int n)

 {
 il(n==1)
 return 1;
 elu
 return(n*fact(n-1));
 }

	,	U
return value	Hobs	rsity Acade
Parameter		bing[Training]informati
dymanuc link	/ •	
returnitules		
Parameter	2	
dynamic wink	•	
returnibuly)
Paramietr	3	AR for
dynamic link	•	Jack
Ycturnvulu		AR for
local	t	main

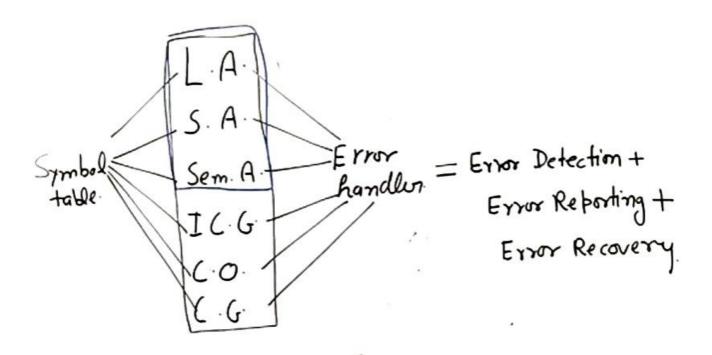
Run time Storage Admunistration.

(2) Implementation of Block-Structured Language.

Storage. Non-local local (handling by scope Informally (handling by AR) dynamic Static (Non-Block (Block structure Structure Storage) 1 local data: Storage Hon - local date: test() int a, 6; Eint x, y; { int c, d;

Static scope rule. (Lexical scope): eq. C, ADA, PASCAL scop-test()

Error Detection and Recovery



Combile time run time.

Levicul Huse Syntactic Phase Semantic Phase error.

Recovery method	Lexical bhan	Syntoctic Phase Error	Sertimiversity Ac
Paric Mode Phrase level Error prodution	× ×	~	X X X
Clobal broduction	*	X	^ ~,

(iii) Ummertihed string or comment

et. Void main()

{
int a, @, 1; variable declaration */

a=10;

|mintf("%d", a); \$
}

Error Detection and Recovery

Syntactic Phase Error.

- (i) missing parenthesis eq. printf("hello";
- missing operator a+b,c
- Misspelled Kennord Swifth (ch)
- Colon in place of semicolon a=1:x
- Extra Blank space . /x comment x g/
- [1) wing symbol tubles:

- 3 Error production. Add exta grammar production and supput make on Augmented grammar and parex university Acode
- (4) global correction. The parson examine the whole program and tries to find out closest much for it which is error free. due to high space and time complexity It Is not implemented practically.

Semantic Phaser error:

- int flow (1) Incompatible type of operands.
- Mil) Not matching actual argument with formul argument.

