## galois

## **Background: Sequence Comprehensions**

## What is a sequence comprehension?

Cryptol uses a sequence comprehension to manage state changes during a computation. A simple example of a sequence comprehension is this:

```
Cryptol> [x | x <- [1...] | y <- [1..10]] [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
```

The outer [ ] means that a sequence is being constructed. The leftmost x, which is to the left of the leftmost |, means that values of x are being appended to the sequence. Those values come from the construct x <- [1...] which means x values are taken from an infinite sequence. The variable y controls how many of those values x will take. Variable y gets a value every time x gets a value due to the rightmost |. Thus, when x gets 1, y gets 1, when x gets 2, y gets 2, and so on. But y gets only up to 10. At that point computation stops, but not before the output sequence grew from empty to [1] to [1,2], to [1,2,3], up to [1,2,3,4,5,6,7,8,9,10] due to the 10 x values that were appended to the sequence. Now replace the rightmost | with , and observe the output of this:

```
Cryptol> [ (x,y) | x <- [10..12] , y <- [1..4]]
[(10, 1), (10, 2), (10, 3), (10, 4), (11, 1), (11, 2), (11, 3),
(11, 4), (12, 1), (12, 2), (12, 3), (12, 4)]
```

The tuple (x,y) is appended to the sequence, per iteration, but x and y do not take values synchronously, rather as follows: values for x are taken from the sequence [10,11,12] but for each value of x all the values for y are taken from the sequence [1,2,3,4] and paired with a static x as many times as there are y values. A tuple (x,y) is appended for each such pairing of x and y. Another simple example:

```
Cryptol> [ (x,y,z) | x <- [10,11] , y <- [20,21], z <- [30,31]] [(10, 20, 30), (10, 20, 31), (10, 21, 30), (10, 21, 31), (11, 20, 30), (11, 20, 31), (11, 21, 30), (11, 21, 31)]
```

The rightmost | can be added to control the number of tuples that is appended to the sequence, for example like this.

```
Cryptol> [ (x,y,z) \mid x <- [10,11] , y <- [20,21] \mid z <- [30,31,32] [(10, 20, 30), (10, 21, 31), (11, 20, 32)]
```

Now things get interesting. The elements that are appended to a sequence are typically states of a computation. But next states typically depend on current or previous states. The following is a simple illustration of how this is managed with sequence comprehensions.

```
Cryptol> let s = [1]\#[x+y \mid x <-s \mid y <- [1]\#s]
Cryptol> take `{10} s
[1, 2, 3, 5, 8, 13, 21, 34, 55, 89]
```

Sequence s begins with just one element, namely 1, due to [1]#[x+y | ...] and [x+y | ...] is initially empty. Next, the first x+y is appended to s. The first x is 1 which is the only

element existing for s at this point. The first y comes from the sequence [1]#s which starts off as [1]#[1] which is [1,1]. So, the first y is 1 and the first x+y is 2 which is appended to s. For the second x+y, s is now [1,2]. Variable x has already taken 1 from s so now it takes 2. Variable y is now taking elements from [1,1,2] since 2 has been appended to s. Since y already took the first 1 it now takes the second 1. Thus x+y is equal to 2 (from x) + 1 (from y) = 3 which is appended to s. On the next iteration s = [1,2,3], x gets the 3, y gets the 2 and the sum 3+2=5 is appended to s. On the next iteration s = [1,2,3], x gets 5, y gets 3 and 8 is appended to s. This is an infinite computation but no time is spent making the computation because the elements in s are computed only on demand. In this case, take `{10} s demands that the first 10 of the elements of s be computed and displayed. It should be evident that the above comprehension computes all Fibonacci numbers. To find, say, the  $30^{th}$  number in the sequence do this: Cryptol> s@29 (the  $0^{th}$  element is in the first position of s so the  $29^{th}$  element of the sequence is in the  $30^{th}$  position).

Now consider a classic problem in logic: prove that if the number of pigeon holes is less than the number of pigeons, and pigeons are distributed over the holes, then at least one hole has more than one pigeon. First consider a representation of the problem. Let a sequence ph represent a collection of pigeon holes. Each element of ph will be a number of pigeons in the hole corresponding to that element. The total number of pigeons in holes can be calculated with this:

Now write a comprehension that returns True if and only if all elements of ph are no greater than 1:

```
Cryptol> :s base=2  
Cryptol> let nh = [True]\#[ (x \le 1) / y | x \le ph | y \le nh ]  
Cryptol> nh  
0b1111110000  
Cryptol> let pigeons0K = nh ! 0  
Cryptol> pigeons0K  
False
```

Sequence nh keeps getting True appended to is as long as x, the number of pigeons in the current hole, is no greater than 1. The appended True comes from  $(x \le 1) /$  True where the True is the previous sequence member of nh. If  $(x \le 1)$  is False, the  $(x \le 1) /$  True is False and False is appended to nh. From then on, no matter what  $(x \le 1)$  is,  $(x \le 1) /$  False is False and False will be appended to nh. If nh ! 0 is False then some hole has more than 1 pigeon. Otherwise, all holes have at most one pigeon.

The number of pigeon holes is:

```
Cryptol> let nholes = length ph
Cryptol> nholes
```

The property of interest is this:

```
Cryptol> if (npigeons > nholes) then pigeonsOK == False else True
```

which is the same as saying (npigeons > nholes) implies pigeonsOK is False. In a subsequent unit the above will be rewritten so that this property may be proved.

Now consider a comprehension that reverses the elements of a sequence lst:

```
Cryptol> let z = [lst] \# [[b]\#(tail (a >> 1)) | a <- z | b <- lst ] Cryptol> let rev = z ! 0
```

The following shows the value of z given input sequence lst = [16, 42, 23, 77]

```
Cryptol> z [[16,42,23,77],[16,16,42,23],[42,16,16,42],[23,42,16,16],[77,23,42,16]] Cryptol> z ! 0 [77,23,42,16]
```

So, the next sequence to be appended to z begins with the next element of the input sequence lst (from b <- lst), and is followed by the tail of the previous sequence of z (from a <- z). After all elements are pulled from lst the last sequence of z is the reverse sequence of lst. Then z! 0 shows that last sequence.

The following creates a sequence of 16 numbers, beginning with the first one, x, as input, doubled, then followed by numbers that are double the previous one.

```
Cryptol> let x = 3
Cryptol> let iv = [ 2*lr | lr <- [x] # iv | k <- [1 .. 16]]
Cryptol> iv
[6,12,24,48,96,192,384,768,1536,3072,6144,12288,24576,49152,98304,196608]
```

The following is an infinite version of the above

```
Cryptol> let x = 2
Cryptol> let iv = [2*lr | lr <- [x]#iv | k <- [1...]]
Cryptol> iv
[4, 8, 16, 32, 64, ...]
Cryptol> let res = take `{12} iv
[4, 8, 16, 32, 64, 128, 256, 512, 1024, 2048, 4096, 8192]
```

Consider the following comprehension (but see below warning before running it):

```
Cryptol> let ys = [ x + y | x <- ys | y <- [0] \# ys ]
Cryptol> :t ys
ys : {n, a} (Ring a, Literal 0 a) => [n]a
```

So, from the type, ys exists but it has no elements

```
Cryptol> ys == []
True
```

This is because x never gets a value from x < -ys.

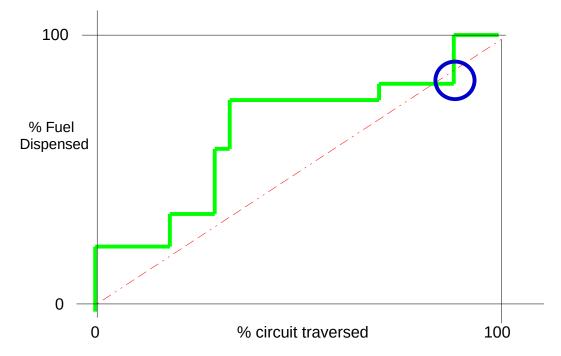
**Warning:** computing ys may cause all of memory to be consumed eventually, then a crash, so kill cryptol as soon as you can in case you want to try the above.

Finally, consider a problem of supplying fuel to a race car traversing a race circuit. There are a number of fueling stations located at various points around the circuit. Suppose the amount of fuel at each fueling station is different but the total fuel around the circuit is exactly what is needed by the race car to make one complete circuit. Find the point on the circuit from which a race car, with an empty tank, may make one complete circuit without running out of fuel by tanking up at every fueling station in the circuit. Assume the tank is large enough to hold enough fuel to complete the circuit. The location of each fueling station is specified by a sequence of increasing ratios, each of which states the percentage of the total circuit distance that the fueling station, corresponding to a sequence element, is from the starting line. Here is an example:

The highest ratio in loc must be no greater than (ratio 1 1).

The amount of fuel a station contains is stated as a percentage of the fuel needed to complete one circuit. This information is contained in a sequence where the station at position x corresponds to the location at position x in the location sequence. Here is an example:

The sum of the numbers in gas must be 1.



The red line in the above Figure depicts the consumption of fuel by the race car over the entire circuit. The green line depicts the fuel that has been dispensed to the race car. The vertical jumps state the fuel supplied by a fueling station at the station's position on the circuit. As long as the green line is above the red line, the race car has fuel and may continue on the circuit. If the green line crosses the red line, the race car has run out of fuel. From the above, it is clear that moving the red line to the right as far as it will go while intersecting the green

line and keeping the same slope, the red line will finish moving at the station that is the solution to the problem. In the above Figure that is the station indicated inside the blue circle. Let fuel\_seen be the percentage of fuel dispensed to the race car so far at the arrival of a fuel station. Let track\_covered be the percentage of the track covered by the race car at that location. If track\_covered - fuel\_seen is negative, there is a fuel deficit and the race car will run out of fuel. Creating a sequence comprehension one can find the deficit at each station and, considering each station in order, if that deficit is the greatest seen so far then record that station as the current correct station to start from. Thus, the following structure, a tuple, will be updated by the comprehension: (fuel\_seen, deficit, station). The station will be identified by its position in the input sequences p and g. Initially, the tuple is (0,0,-1) where the -1 just indicates a lack of station candidate so far and the two 0s represent the origin of the race car in the circuit. Begin building the comprehension like this:

**Note:** before cutting and pasting the above three lets to Cryptol, make each let statement just one line by concatenating all lines. Then take a look at s:

```
Cryptol> s
[((ratio 0 1), (ratio 0 1), -1), ((ratio 1 10), (ratio 1 5), 0),
  ((ratio 23 100), (ratio 1 5), 0), ((ratio 33 100), (ratio 11 50), 2),
  ((ratio 12 25), (ratio 11 50), 2), ((ratio 71 100), (ratio 23 100), 4),
  ((ratio 39 50), (ratio 23 100), 4), ((ratio 19 20), (ratio 23 100), 4),
  ((ratio 1 1), (ratio 23 100), 4)]
```

So, the station to start from is at (s!0).2 and the location of that station is loc@(s!0).2.

```
Cryptol> (s!0).2
4
Cryptol> loc@(s!0).2
(ratio 71 100)
```