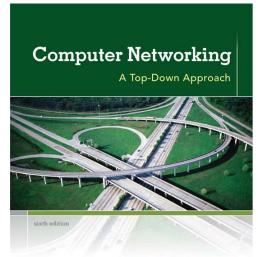
# Chapter 6 Wireless and Mobile Networks



KUROSE ROSS

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Networking: A Top
Down Approach
6th edition
Jim Kurose, Keith Ross
Addison-Wesley

March 2012

### Ch. 6: Wireless and Mobile Networks

### **Background:**

- # wireless (mobile) phone subscribers now exceeds # wired phone subscribers (5-to-I)!
- # wireless Internet-connected devices equals # wireline Internet-connected devices
  - laptops, Internet-enabled phones promise anytime untethered Internet access
- two important (but different) challenges
  - wireless: communication over wireless link
  - mobility: handling the mobile user who changes point of attachment to network

# Chapter 6 outline

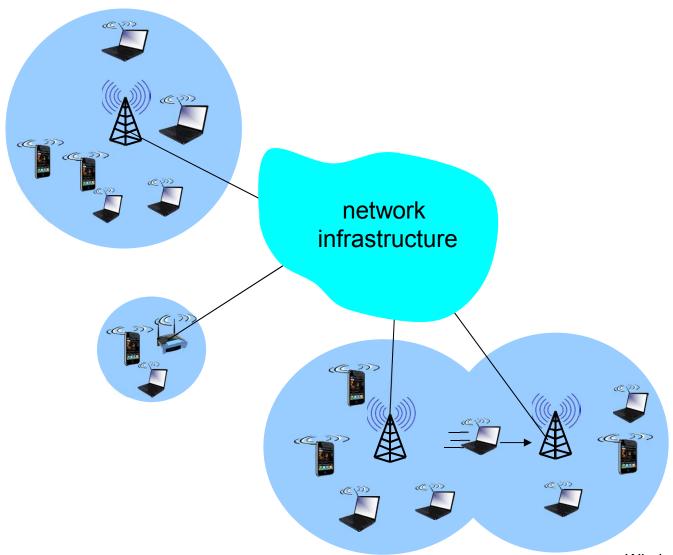
#### 6. I Introduction

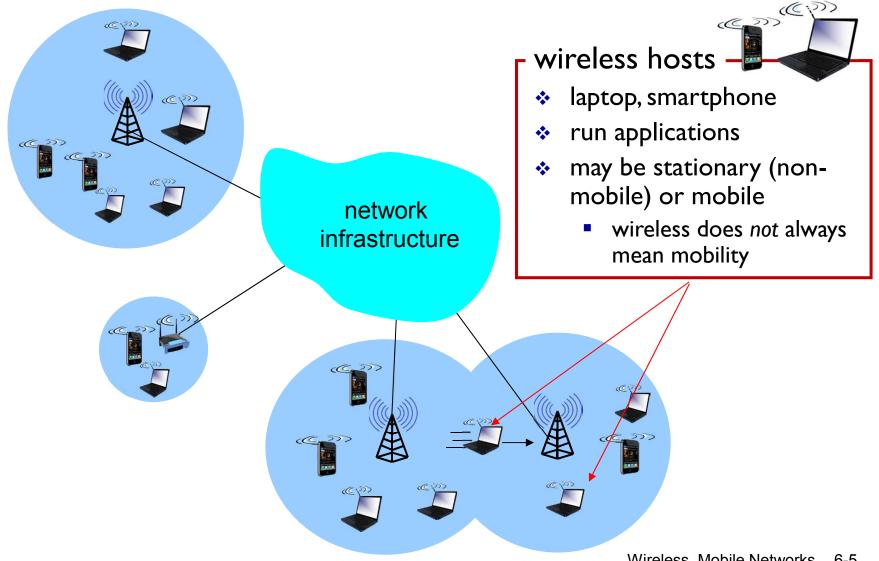
#### Wireless

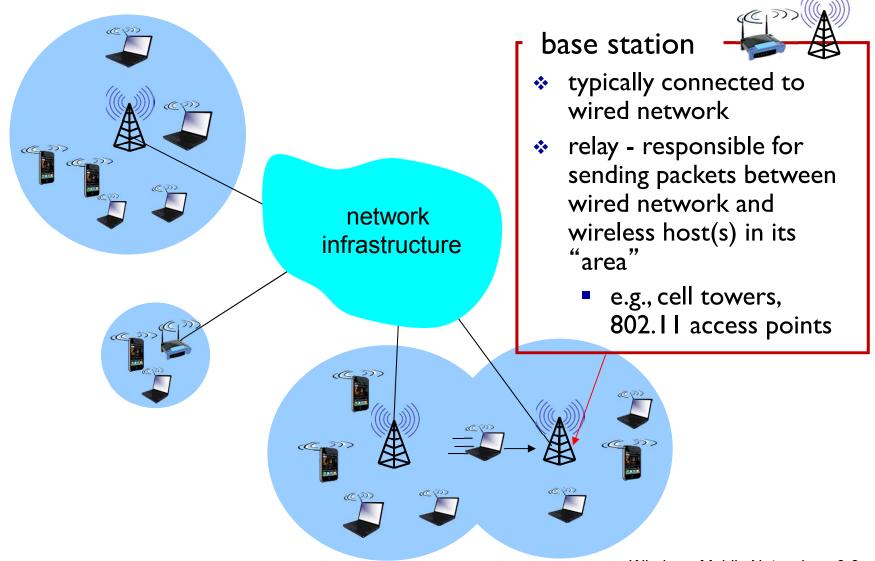
- 6.2 Wireless links, characteristics
  - CDMA
- 6.3 IEEE 802.11 wireless LANs ("Wi-Fi")
- 6.4 Cellular Internet Access
  - architecture
  - standards (e.g., GSM)

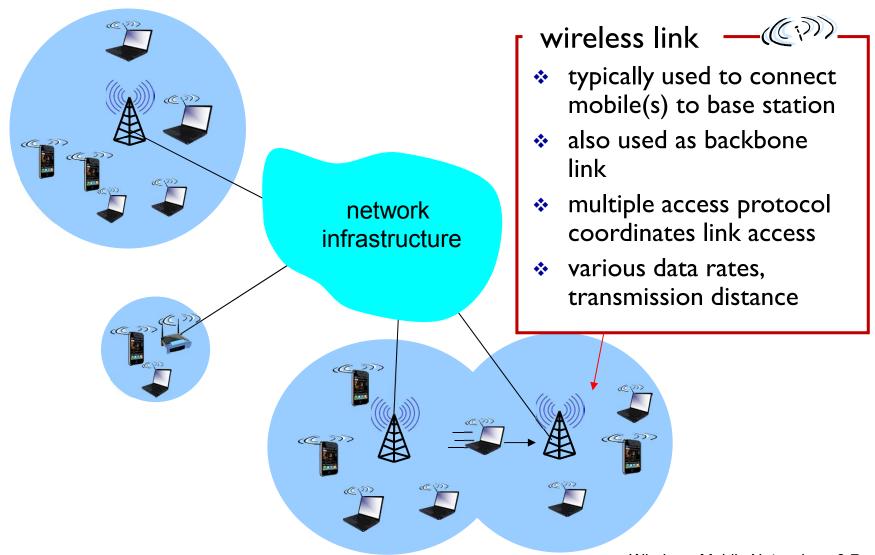
### **Mobility**

- 6.5 Principles: addressing and routing to mobile users
- 6.6 Mobile IP
- 6.7 Handling mobility in cellular networks
- 6.8 Mobility and higher-layer protocols
- 6.9 Summary

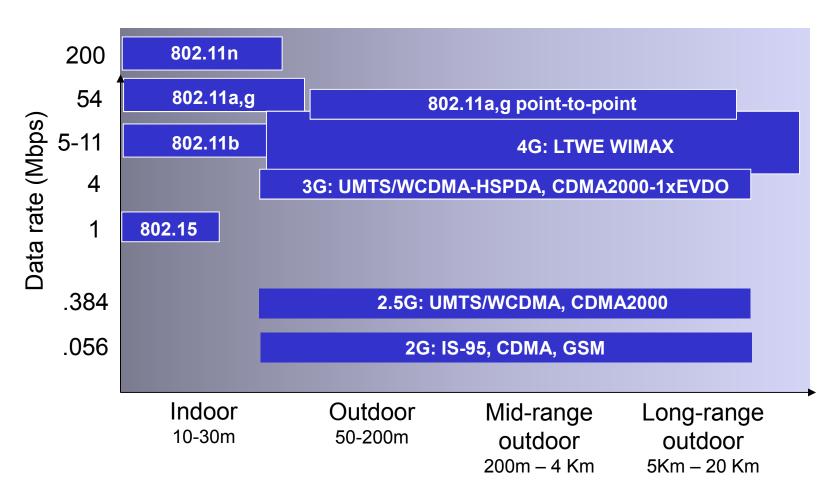


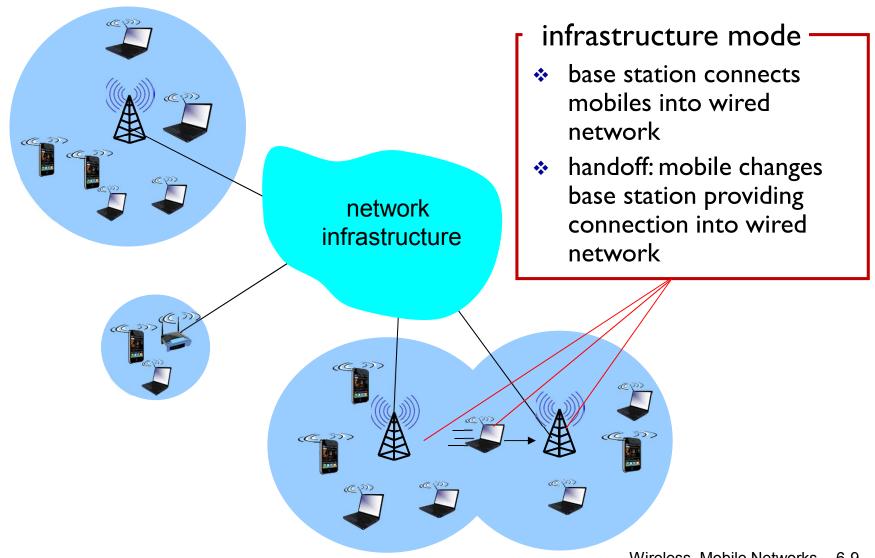


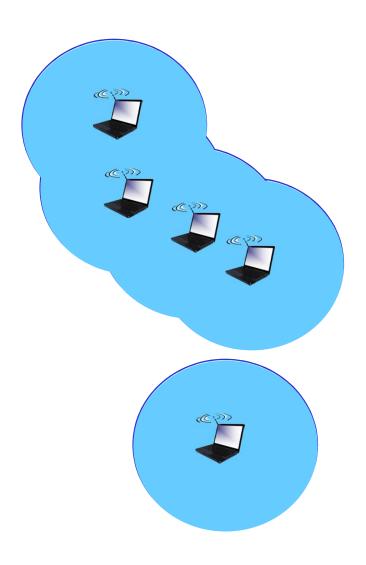




### Characteristics of selected wireless links







#### ad hoc mode

- no base stations
- nodes can only transmit to other nodes within link coverage
- nodes organize themselves into a network: route among themselves

# Wireless network taxonomy

	single hop	multiple hops
infrastructure (e.g., APs)	host connects to base station (WiFi, WiMAX, cellular) which connects to larger Internet	host may have to relay through several wireless nodes to connect to larger Internet: mesh net
no infrastructure	no base station, no connection to larger Internet (Bluetooth, ad hoc nets)	no base station, no connection to larger Internet. May have to relay to reach other given wireless node MANET, VANET

### Wireless Link Characteristics (I)

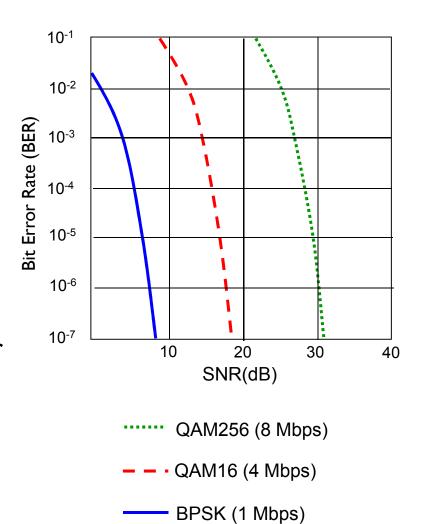
important differences from wired link ....

- decreased signal strength: radio signal attenuates as it propagates through matter (path loss)
- interference from other sources: standardized wireless network frequencies (e.g., 2.4 GHz) shared by other devices (e.g., phone); devices (motors) interfere as well
- multipath propagation: radio signal reflects off objects ground, arriving at destination at slightly different times

.... make communication across (even a point to point) wireless link much more "difficult"

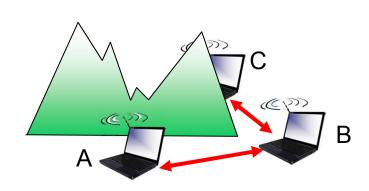
### Wireless Link Characteristics (2)

- SNR: signal-to-noise ratio
  - larger SNR easier to extract signal from noise (a "good thing")
- SNR versus BER tradeoffs
  - given physical layer: increase power -> increase SNR->decrease BER
  - given SNR: choose physical layer that meets BER requirement, giving highest throughput
    - SNR may change with mobility: dynamically adapt physical layer (modulation technique, rate)



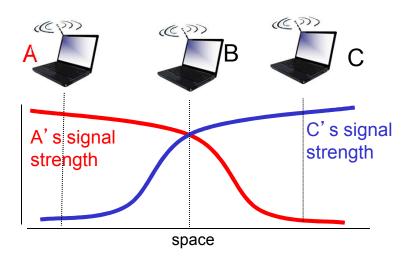
### Wireless network characteristics

Multiple wireless senders and receivers create additional problems (beyond multiple access):



#### Hidden terminal problem

- B,A hear each other
- ❖ B, C hear each other
- A, C can not hear each other means A, C unaware of their interference at B

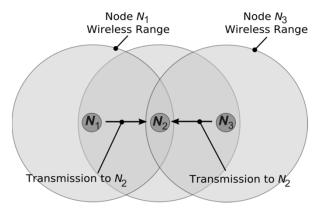


#### Signal attenuation:

- ❖ B,A hear each other
- B, C hear each other
- A, C cannot hear each other interfering at B

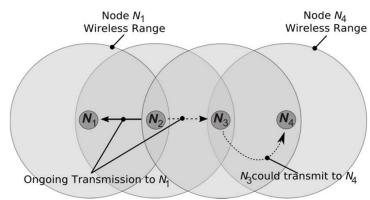
### Wireless network characteristics

Sentir o canal no transmissor não fornece informação acerca do canal no receptor



#### Hidden node problem

N1 e N3 não se escutam mutuamente devido a obstáculos ou atenuação: os seus pacotes colidem em N2



### **Exposed node problem**

N1 e N4 poderiam ser receptores simultâneos mas os respectivos emissores N2 e N3 estão em zona de

alcance -

Dois principais problemas

Problema menos grave que o anterior →reduz a utilização →menos estudado

Wireless, Mobile Networks 6-16

### IEEE 802.11 Wireless LAN

#### 802.11b

- 2.4-5 GHz unlicensed spectrum
- up to 11 Mbps
- direct sequence spread spectrum (DSSS) in physical layer
  - all hosts use same chipping code

#### 802.11a

- 5-6 GHz range
- up to 54 Mbps

#### 802.11g

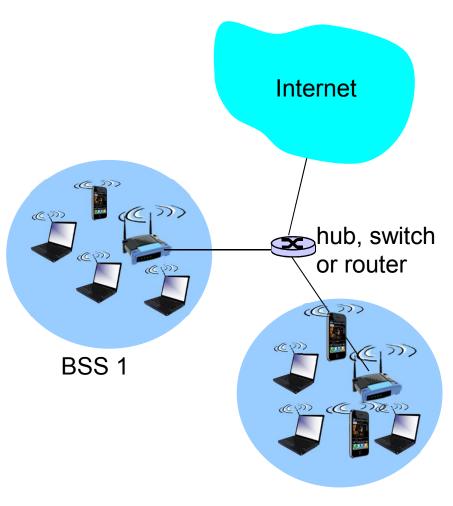
- 2.4-5 GHz range
- up to 54 Mbps

#### 802. I In: multiple antennae

- 2.4-5 GHz range
- up to 200 Mbps

- all use CSMA/CA for multiple access
- all have base-station and ad-hoc network versions

### 802.11 LAN architecture



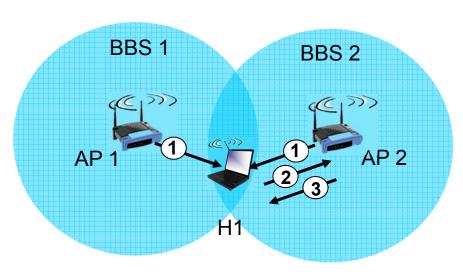
BSS 2

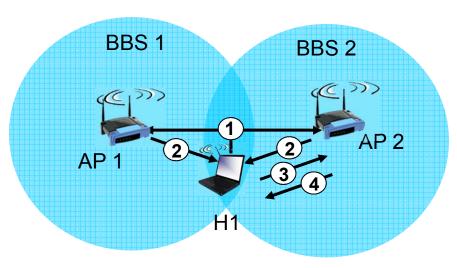
- wireless host communicates with base station
  - base station = access point (AP)
- \* Basic Service Set (BSS) (aka "cell") in infrastructure mode contains:
  - wireless hosts
  - access point (AP): base station
  - ad hoc mode: hosts only

### 802.11: Channels, association

- \* 802.11b: 2.4GHz-2.485GHz spectrum divided into 11 channels at different frequencies
  - AP admin chooses frequency for AP
  - interference possible: channel can be same as that chosen by neighboring AP!
- \* host: must associate with an AP
  - scans channels, listening for beacon frames containing AP's name (SSID) and MAC address
  - selects AP to associate with
  - may perform authentication
  - will typically run DHCP to get IP address in AP's subnet

# 802. I I: passive/active scanning





#### passive scanning:

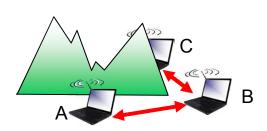
- (I) beacon frames sent from APs
- (2) association Request frame sent: HI to selected AP
- (3) association Response frame sent from selected AP to HI

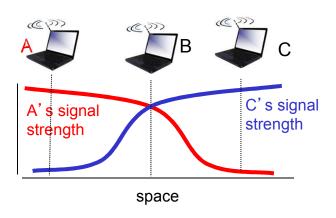
#### active scanning:

- (1) Probe Request frame broadcast from H1
- (2) Probe Response frames sent from APs
- (3) Association Request frame sent: H1 to selected AP
- (4) Association Response frame sent from selected AP to H1

### IEEE 802.11: multiple access

- avoid collisions: 2<sup>+</sup> nodes transmitting at same time
- \* 802.11: CSMA sense before transmitting
  - don't collide with ongoing transmission by other node
- \* 802.11: no collision detection!
  - difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
  - can't sense all collisions in any case: hidden terminal, fading
  - goal: avoid collisions: CSMA/C(ollision)A(voidance)





### IEEE 802.11 MAC Protocol: CSMA/CA

#### 802.11 sender

**DIFS**: **DCF** (Distributed Coordination Function) Inter-Frame Sequence

SIFS: Short Inter-Frame Sequence

1 if sense channel idle for **DIFS** then transmit entire frame (no CD)

2 if sense channel busy then

start random backoff time

timer counts down while channel idle

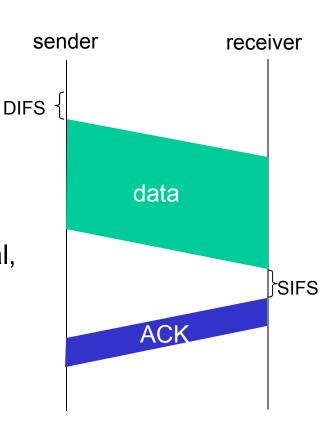
transmit when timer expires

if no ACK, increase random backoff interval, repeat 2

#### <u>802.11 receiver</u>

- if frame received OK

return ACK after **SIFS** (ACK needed due to hidden terminal problem)

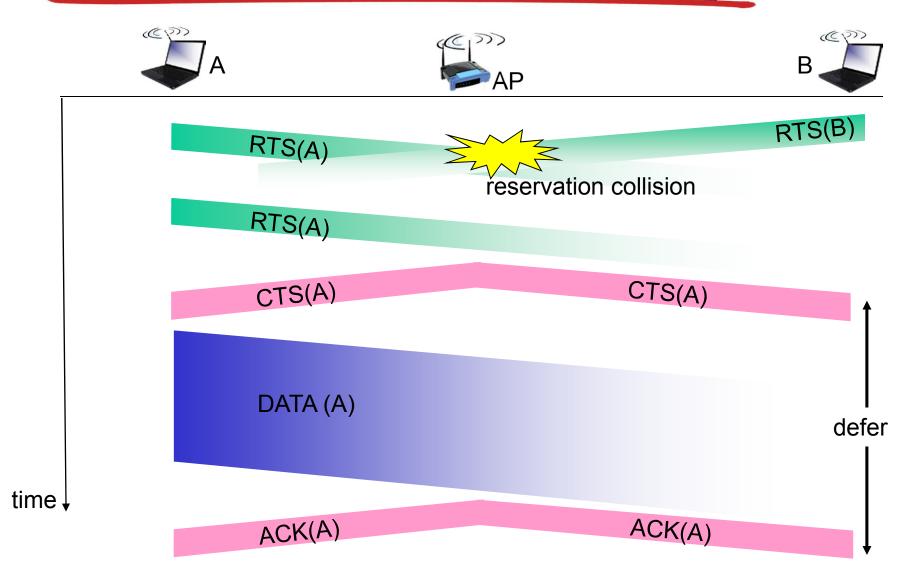


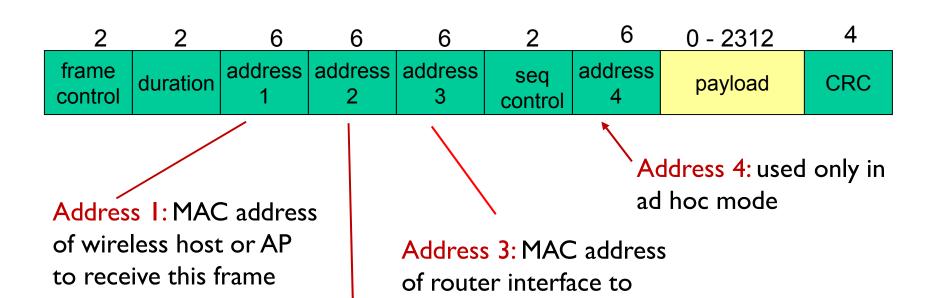
# Avoiding collisions (more)

- idea: allow sender to "reserve" channel rather than random access of data frames: avoid collisions of long data frames
- sender first transmits small request-to-send (RTS) packets to BS using CSMA
  - RTSs may still collide with each other (but they' re short)
- BS broadcasts clear-to-send CTS in response to RTS
- CTS heard by all nodes
  - sender transmits data frame
  - other stations defer transmissions

avoid data frame collisions completely using small reservation packets!

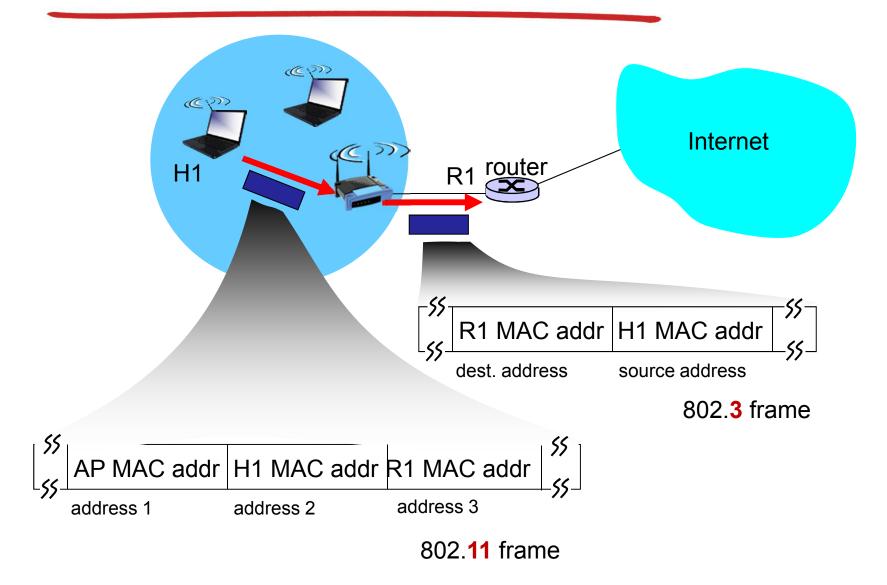
### Collision Avoidance: RTS-CTS exchange





which AP is attached

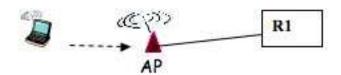
Address 2: MAC address of wireless host or AP transmitting this frame



- The addresses semantics depends on the frame type and directionality
- ❖ addrI MAC address of the receiver (varies);
- ❖ addr2 MAC address of the sender

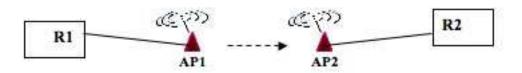
DS: Distribution System

toDS	fromDS	addr1	addr2	addr3	addr4	obs.
0	0	DA	SA	BSSID	02	ad hoc
0	1	DA	BSSID	SA	100	do AP
1	0	BSSID	SA	DA	39	para AP
1	1	RA	TA	DA	SA	dentro DS



#### toDS=1, fromDS=0

A1 (RA) = BSSID = MAC AP A2 (TA) = SA = MAC STA A3 (DA) = MAC R1

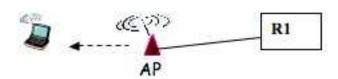


#### toDS=1, fromDS=1

A1 (RA) = MAC AP2

A2 (TA) = MACAP1A3 (DA) = MACR2

A4(SA) = MACRI



#### toDS=0, fromDS=1

A1 (RA) = DA = MAC STA A2 (TA) = BSSID = MAC AP A3 (SA) = MAC RI DA - Destination Address - receptor final

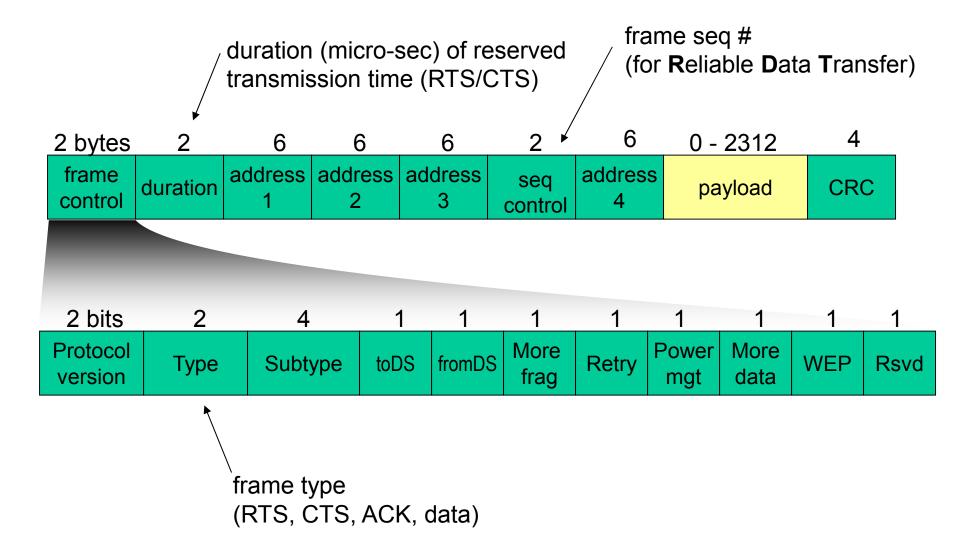
SA - Source Address - origem da transmissão

RA - Receiver Address - estação wireless que deve processar a trama wireless STA -> RA=DA wired node -> RA=MAC AP; DA=router

TA - Transmitter Address - interface wireless que transmitiu a trama

BSSID - MAC da interface wireless do AP (Infrastructure networks); aleatório BSSID (Ad-hoc networks)

### 802.11 frame: more



# 802.11 frame types and subtypes

ANSI/IEEE Std 802.11, 1999 Edition

LOCAL AND METROPOLITAN AREA NETWORKS: WIRELESS LAN

Type 00 - Management frames e.g. Beacon, Association request, Probe request, etc.

Type 01 - Control frames e.g. RTS, CTS, ACK, etc.

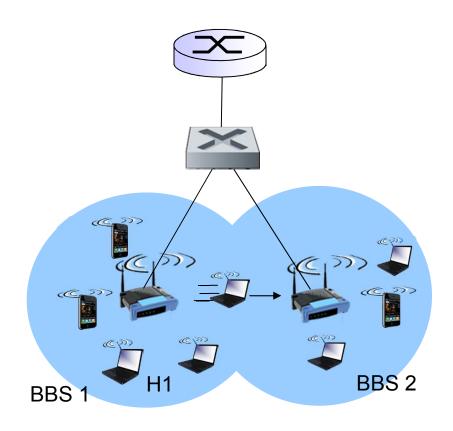
Type 10 - Data frames e.g. Data

Table 1-Valid type and subtype combinations

Type value b3 b2	Type description	Subtype value b7 b6 b5 b4	Subtype description
00	Management	0000	Association request
00	Management	0001	Association response
00	Management	0010	Reassociation request
00	Management	0011	Reassociation response
00	Management	0100	Probe request
00	Management	0101	Probe response
00	Management	0110-0111	Reserved
00	Management	1000	Beacon
00	Management	1001	Announcement traffic indication message (ATIM)
00	Management	1010	Disassociation
00	Management	1011	Authentication
00	Management	1100	Deauthentication
00	Management	1101-1111	Reserved
01	Control	0000-1001	Reserved
01	Control	1010	Power Save (PS)-Poll
01	Control	1011	Request To Send (RTS)
01	Control	1100	Clear To Send (CTS)
01	Control	1101	Acknowledgment (ACK)
01	Control	1110	Contention-Free (CF)-End
01	Control	1111	CF-End + CF-Ack
10	Data	0000	Data
10	Data	0001	Data + CF-Ack
10	Data	0010	Data + CF-Poll
10	Data	0011	Data + CF-Ack + CF-Poll
10	Data	0100	Null function (no data)
10	Data	0101	CF-Ack (no data)
10	Data	0110	CF-Poll (no data)
10	Data	0111	CF-Ack + CF-Poll (no data)
10	Data	1000-1111	Reserved
11	Reserved	0000-1111	Reserved

### 802. I I: mobility within same subnet

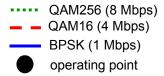
- H1 remains in same IP subnet: IP address can remain same
- switch: which AP is associated with H1?
  - self-learning (Ch. 5):
     switch will see frame
     from H1 and
     "remember" which
     switch port can be used
     to reach H1

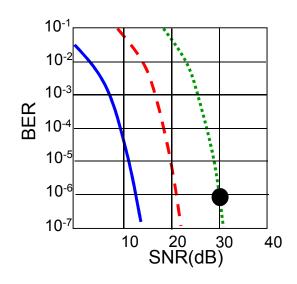


### 802. I I: advanced capabilities

### Rate adaptation

base station, mobile
 dynamically change
 transmission rate
 (physical layer modulation
 technique) as mobile
 moves, SNR varies





- 1. SNR decreases, BER increase as node moves away from base station
- 2. When BER becomes too high, switch to lower transmission rate but with lower BER

# 802. I I: advanced capabilities

#### power management

- node-to-AP: "I am going to sleep until next beacon frame"
  - AP knows not to transmit frames to this node
  - node wakes up before next beacon frame
- beacon frame: contains list of mobiles with APto-mobile frames waiting to be sent
  - node will stay awake if AP-to-mobile frames to be sent; otherwise sleep again until next beacon frame

### Chapter 6 outline

6. I Introduction

#### Wireless

- 6.2 Wireless links, characteristics
  - CDMA
- 6.3 IEEE 802.11 wireless LANs ("Wi-Fi")
- 6.4 Cellular Internet access
  - architecture
  - standards (e.g., GSM)

#### **Mobility**

- 6.5 Principles: addressing and routing to mobile users
- 6.6 Mobile IP
- 6.7 Handling mobility in cellular networks
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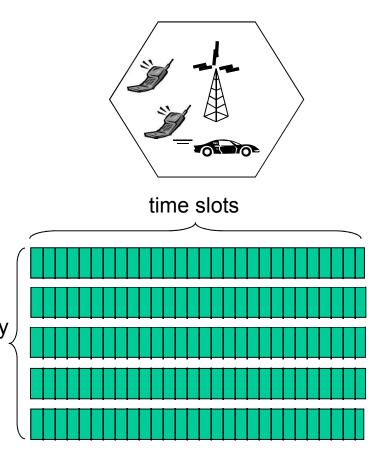
### Components of cellular network architecture

#### **MSC** connects cells to wired tel. net. manages call setup (more later!) handles mobility (more later!) cell covers geographical region Mobile base station (BS) Switching analogous to 802.11 AP Center Public telephone \* mobile users attach to network network through BS air-interface: physical Mobile and link layer protocol **Switching** Center between mobile and BS wired network

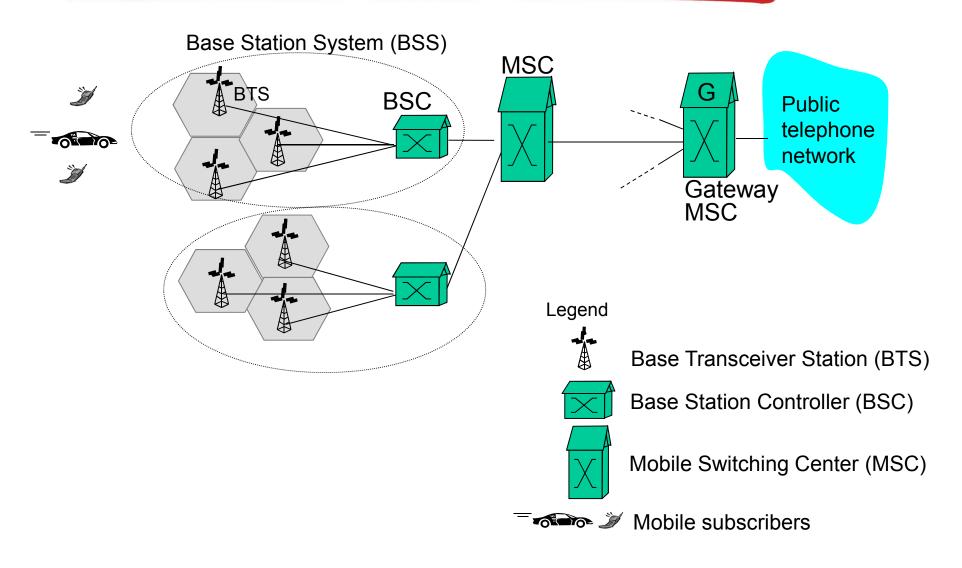
### Cellular networks: the first hop

Two techniques for sharing mobile-to-BS radio spectrum

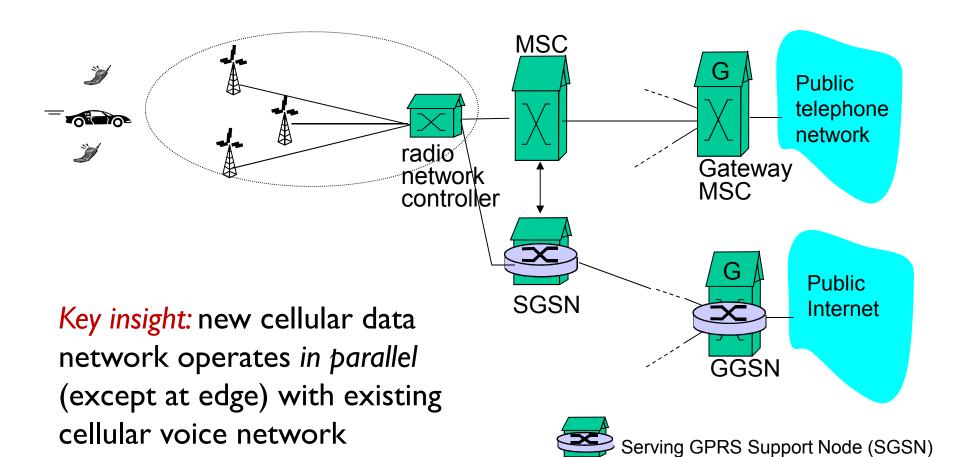
- combined FDMA/TDMA: divide spectrum in frequency channels, divide each channel into time slots
- CDMA: code division multiple access
  frequency
  bands



### 2G (voice) network architecture



## 3G (voice+data) network architecture

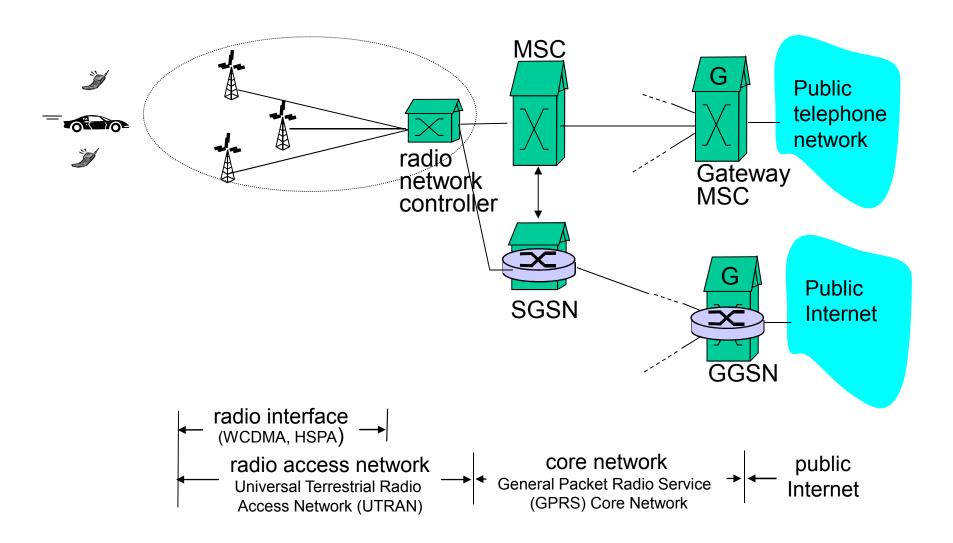


- voice network unchanged in core
- data network operates in parallel



Gateway GPRS Support Node (GGSN)

## 3G (voice+data) network architecture



## Chapter 6 outline

6. I Introduction

#### Wireless

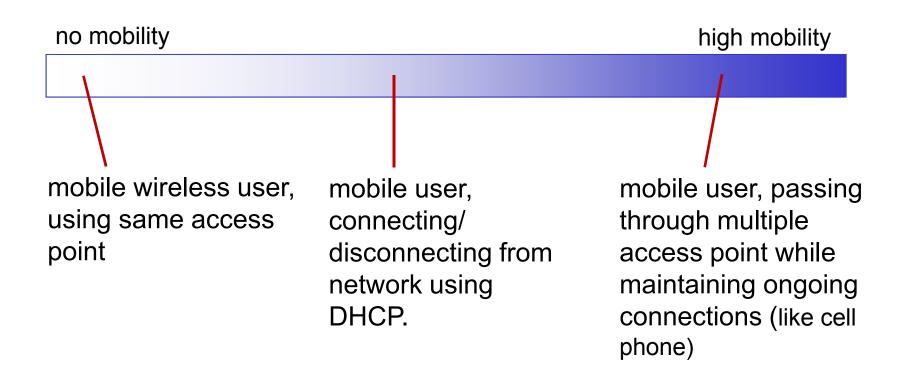
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### Mobility

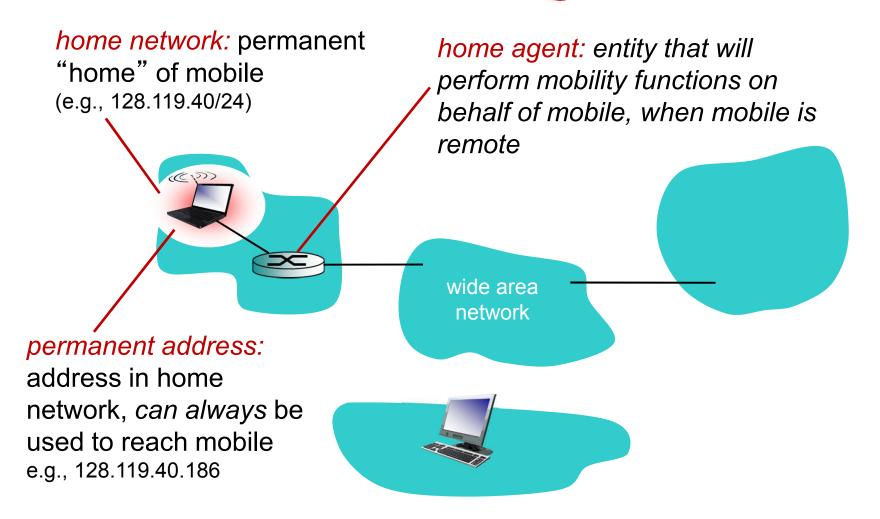
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## What is mobility?

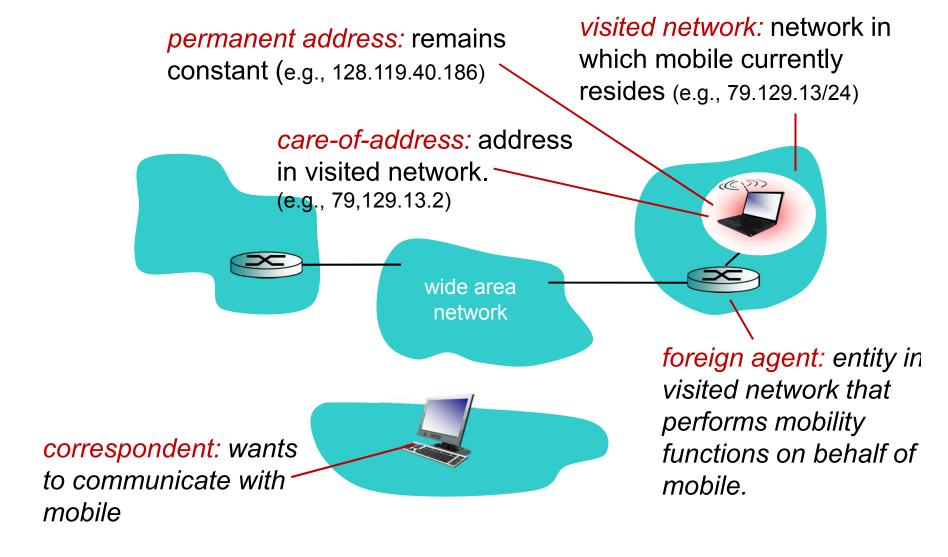
spectrum of mobility, from the network perspective:



## Mobility: vocabulary



# Mobility: more vocabulary



## How do you contact a mobile friend:

Consider friend frequently changing addresses, how do you find her?

search all phone books?

- call her parents?
- expect her to let you know where he/she is?

I wonder where Alice moved to?



## Mobility: approaches

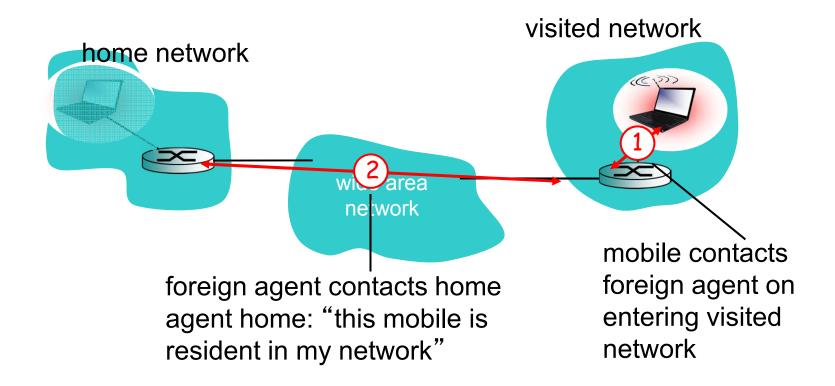
- let routing handle it: routers advertise permanent address of mobile-nodes-in-residence via usual routing table exchange.
  - routing tables indicate where each mobile located
  - no changes to end-systems
- let end-systems handle it:
  - indirect routing: communication from correspondent to mobile goes through home agent, then forwarded to remote
  - direct routing: correspondent gets foreign address of mobile, sends directly to mobile

## Mobility: approaches

- vertise permanent address of let routing handle it: r mobile-nodes-in-r sual routing table exchange. not
  - routing tables

- each mobile located
- no changes to mobiles
- let end-systems handle it.
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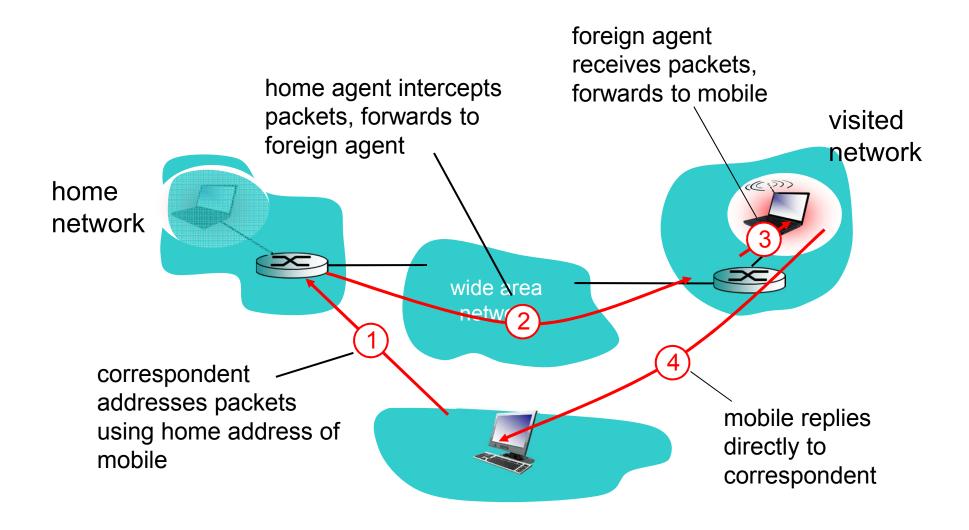
### Mobility: registration



#### end result:

- foreign agent knows about mobile
- home agent knows location of mobile

### Mobility via indirect routing



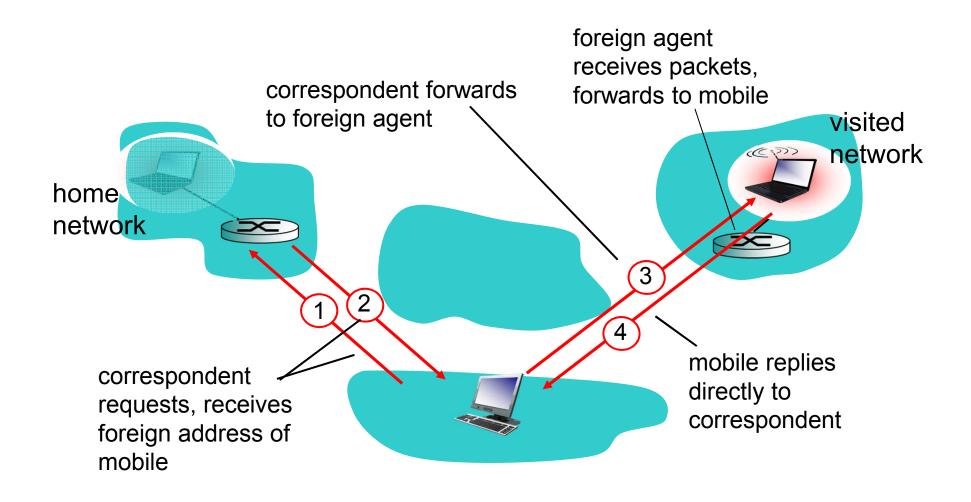
### Indirect Routing: comments

- mobile uses two addresses:
  - permanent address: used by correspondent (hence mobile location is transparent to correspondent)
  - care-of-address: used by home agent to forward datagrams to mobile
- foreign agent functions may be done by mobile itself
- triangle routing: correspondent-home-network
  - mobile
  - inefficient when correspondent, mobile are in same network

### Indirect routing: moving between networks

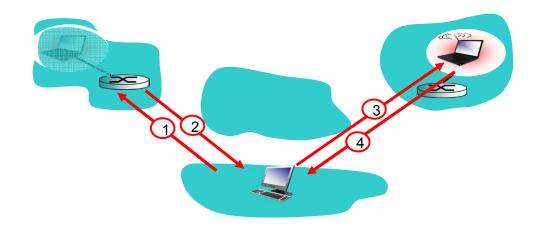
- suppose mobile user moves to another network
  - registers with new foreign agent
  - new foreign agent registers with home agent
  - home agent update care-of-address for mobile
  - packets continue to be forwarded to mobile (but with new care-of-address)
- mobility, changing foreign networks transparent: on going connections can be maintained!

## Mobility via direct routing



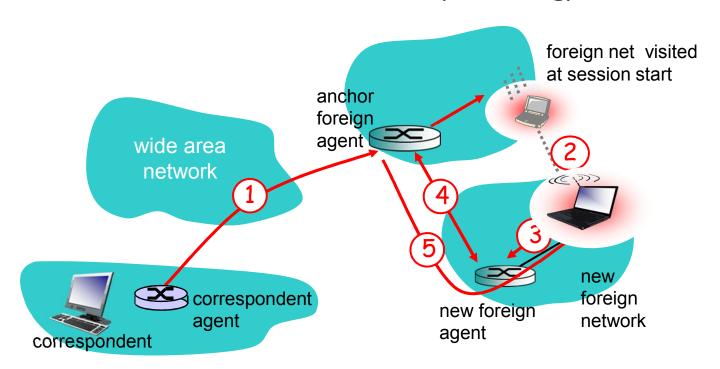
## Mobility via direct routing: comments

- overcome triangle routing problem
- non-transparent to correspondent: correspondent must get care-of-address from home agent
  - what if mobile changes visited network?



### Accommodating mobility with direct routing

- anchor foreign agent: FA in first visited network
- data always routed first to anchor FA
- when mobile moves: new FA arranges to have data forwarded from old FA (chaining)



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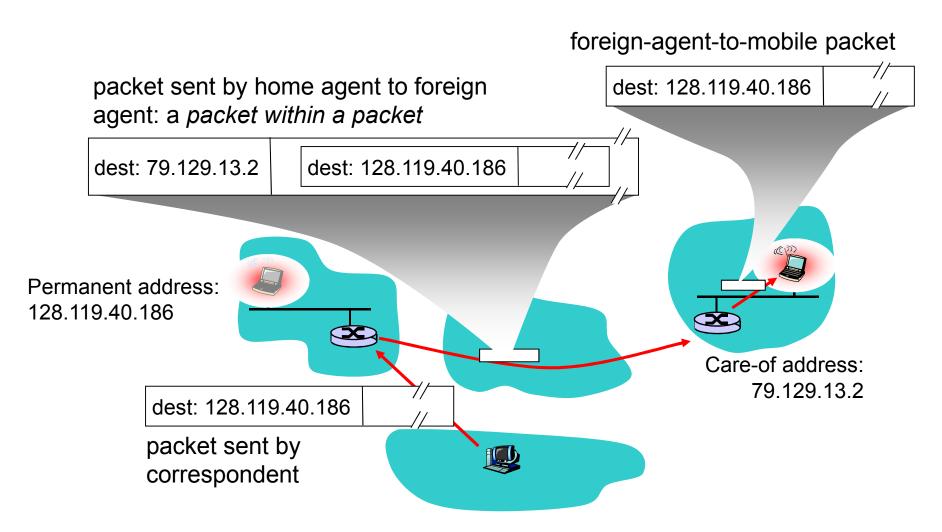
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# Mobile IP

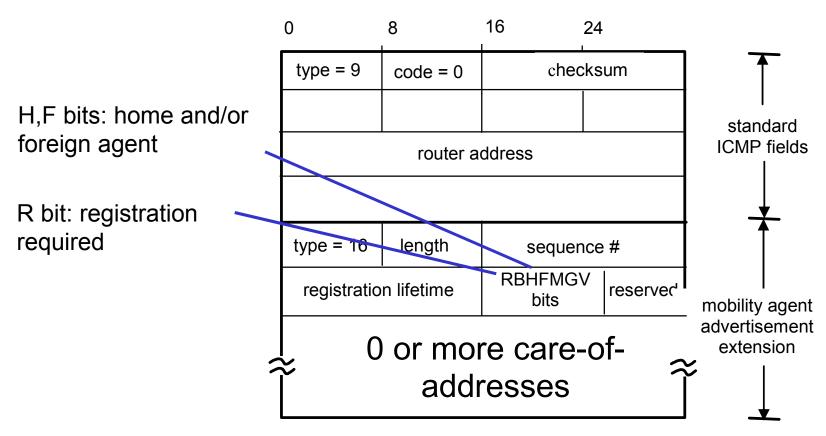
- \* RFC 3344
- has many features we've seen:
  - home agents, foreign agents, foreign-agent registration, care-of-addresses, encapsulation (packet-within-apacket)
- three components to standard:
  - indirect routing of datagrams
  - agent discovery
  - registration with home agent

## Mobile IP: indirect routing

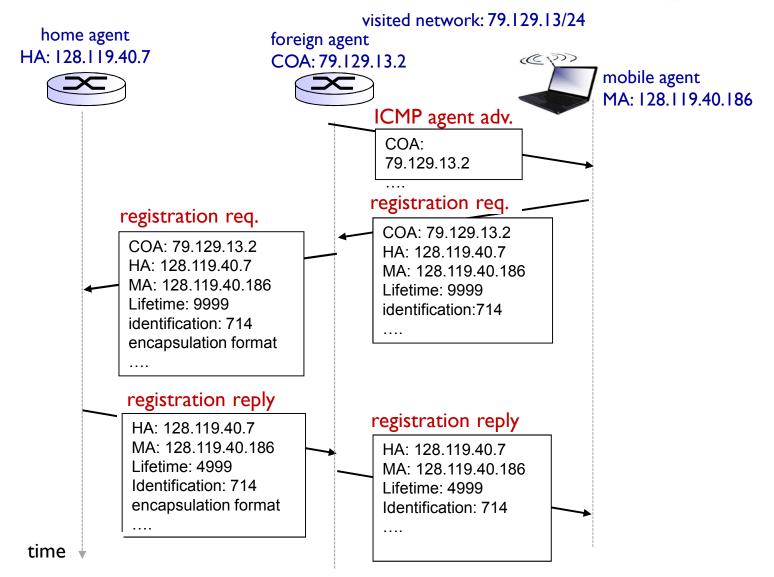


# Mobile IP: agent discovery

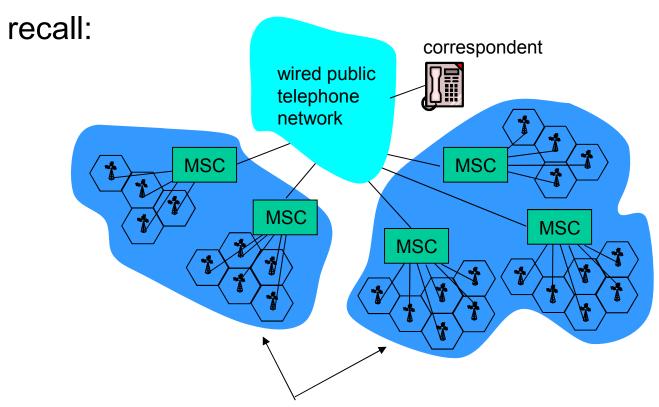
agent advertisement: foreign/home agents advertise service by broadcasting ICMP messages (typefield = 9)



# Mobile IP: registration example



### Components of cellular network architecture

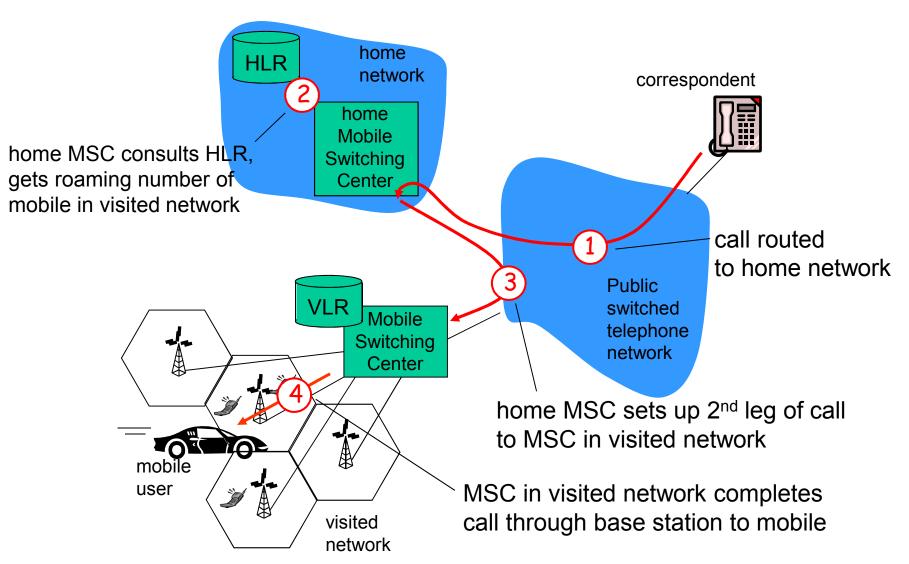


different cellular networks, operated by different providers

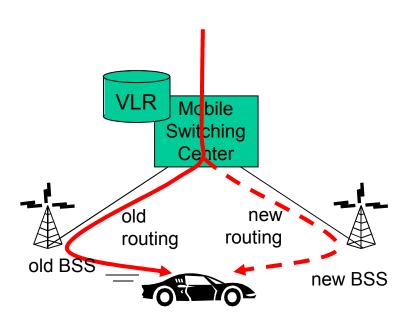
## Handling mobility in cellular networks

- home network: network of cellular provider you subscribe to (e.g., Sprint PCS, Verizon)
  - home location register (HLR): database in home network containing permanent cell phone #, profile information (services, preferences, billing), information about current location (could be in another network)
- visited network: network in which mobile currently resides
  - visitor location register (VLR): database with entry for each user currently in network
  - could be home network

## GSM: indirect routing to mobile

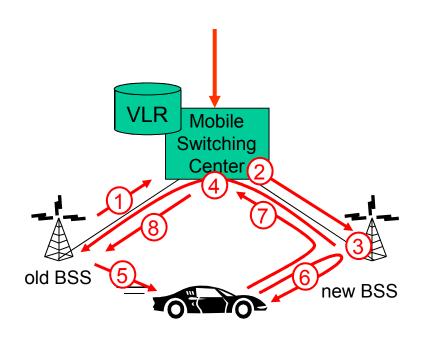


### GSM: handoff with common MSC



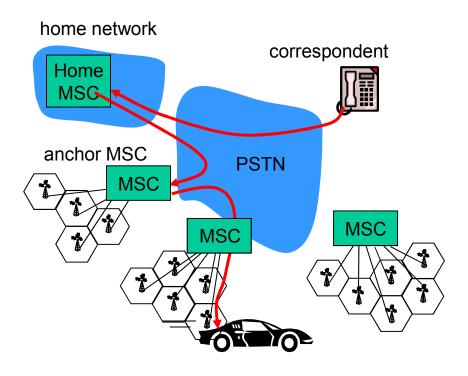
- handoff goal: route call via new base station (without interruption)
- reasons for handoff:
  - stronger signal to/from new BSS (continuing connectivity, less battery drain)
  - load balance: free up channel in current BSS
  - GSM doesnt mandate why to perform handoff (policy), only how (mechanism)
- handoff initiated by old BSS

### GSM: handoff with common MSC



- 1. old BSS informs MSC of impending handoff, provides list of 1<sup>+</sup> new BSSs
- 2. MSC sets up path (allocates resources) to new BSS
- 3. new BSS allocates radio channel for use by mobile
- 4. new BSS signals MSC, old BSS: ready
- 5. old BSS tells mobile: perform handoff to new BSS
- 6. mobile, new BSS signal to activate new channel
- 7. mobile signals via new BSS to MSC: handoff complete. MSC reroutes call
- 8 MSC-old-BSS resources released

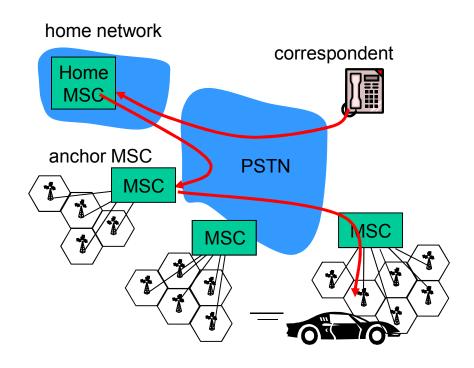
## GSM: handoff between MSCs



(a) before handoff

- anchor MSC: first MSC visited during call
  - call remains routed through anchor MSC
- new MSCs add on to end of MSC chain as mobile moves to new MSC
- optional path minimization step to shorten multi-MSC chain

## GSM: handoff between MSCs



(b) after handoff

- anchor MSC: first MSC visited during call
  - call remains routed through anchor MSC
- new MSCs add on to end of MSC chain as mobile moves to new MSC
- optional path minimization step to shorten multi-MSC chain

## Mobility: GSM versus Mobile IP

GSM element	Comment on GSM element Mo	obile IP element
Home system	Network to which mobile user's permanent phone number belongs	Home network
Gateway Mobile Switching Center, or "home MSC". Home Location Register (HLR)	Home MSC: point of contact to obtain routable address of mobile user. HLR: database in home system containing permanent phone number, profile information, current location of mobile user, subscription information	Home agent
Visited System	Network other than home system where mobile user is currently residing	Visited network
Visited Mobile services Switching Center. Visitor Location Record (VLR)	Visited MSC: responsible for setting up calls to/from mobile nodes in cells associated with MSC. VLR: temporary database entry in visited system, containing subscription information for each visiting mobile user	Foreign agent
Mobile Station Roaming Number (MSRN), or "roaming number"	Routable address for telephone call segment between home MSC and visited MSC, visible to neither the mobile nor the correspondent.	Care-of- address

### Wireless, mobility: impact on higher layer protocols

- ❖ logically, impact should be minimal ...
  - best effort service model remains unchanged
  - TCP and UDP can (and do) run over wireless, mobile
- ... but performance-wise:
  - packet loss/delay due to bit-errors (discarded packets, delays for link-layer retransmissions), and handoff
  - TCP interprets loss as congestion, will decrease congestion window un-necessarily
  - delay impairments for real-time traffic
  - limited bandwidth of wireless links

## Chapter 6 summary

### Wireless

- wireless links:
  - capacity, distance
  - channel impairments
  - CDMA
- ❖ IEEE 802.11 ("Wi-Fi")
  - CSMA/CA reflects wireless channel characteristics
- cellular access
  - architecture
  - standards (e.g., GSM, 3G, 4G LTE)

### Mobility

- principles: addressing, routing to mobile users
  - home, visited networks
  - direct, indirect routing
  - care-of-addresses
- case studies
  - mobile IP
  - mobility in GSM
- impact on higher-layer protocols