Universität Tübingen

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Functional Programming

Notes from

 ${\rm Philipp~Moers} \\ < {\rm soziflip~funnycharacter~gmail~dot~com} >$

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Abstract

This is just the product of me taking notes on the lecture. Nothing official. If you find mistakes or have got any questions, please feel free to contact me. Cheers!

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»A programming language is a medium for expressing ideas (not to get a computer to perform operations) and only incidentally for machines to execute.«	-
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0.1. Links

Site: http://db.inf.uni-tuebingen.de/teaching/FunctionalProgrammingSS2014

Ilias: http://goo.gl/rlqbkK

0.2. Literature

• Lipovača:

Learn You a Haskell for Great Good No Starch Press 2011,

http://learnyouahaskell.com

• O'Sullivan, Steward, Goerzen:

Real World Haskell

O'Reilly 2010

http://book.realworldhaskell.org

• Haskell 2010 Report,

http://www.haskell.org/onlinereport/haskell2010

1. Introduction

Computational model in Functional Programming: **reduction** (replace expression to values)

In Functional Programming, expressions are formed by applying functions to values.

- 1. Functions as in math: $x = y \Rightarrow f(x) = f(y)$
- 2. Functions are values (just like numbers, text $\dots)$

	Functional	Imperative
program construction	function application and composition	statement sequencing
execution	reduction (expression evaluation)	state changes
semantics	lambda calculus	complex (denotational)

Example

 $n \in \mathbb{N}, n \geq 2$ is a prime number iff the set of non-trivial factors is empty:

```
n \text{ is prime} \Leftrightarrow \{\ m \mid m \in \{2, \dots, n-1\},\ n \mod m = 0\} = \emptyset
```

Listing 1.1: isPrime.hs

```
isPrime :: Integer -> Bool
isPrime n = factors == []
where
factors = [ m | m <- [2..n-1], n 'mod' m == 0 ]

main :: IO ()
main = do
let n = 42
print (isPrime n)</pre>
```

2. Haskell Ramp-Up

(Read \equiv as 'denotes the same value as')

- Apply f to value e: f e (juxtaposition, 'apply', binary operator _, Haskell speak: infixL 10 _)
- \cup has max precedence (10): $f e_1 + e_2 \equiv (f e_1) + e_2$
- \cup associates to the left: g f e \equiv (g f) e (' (g f)' is a function)
- Function composition:
 - (g.f) e \equiv g (fe) (. is something like mathematical \circ 'after')
 - Alternative 'apply'-operator \$ (lowest precedence, associates to the right, infixR \emptyset \$):

$$g \ f \ f \ e \equiv g \ f \ (f \ e) \equiv g \ (f \ e)$$

- Prefix application of binary infix operator \otimes : (\otimes) e_1 e_2 \equiv e_1 \otimes e_2
- Infix application of binary function f: e_1 'f' $e_2 \equiv f e_1 e_2$:
 - * 1 'elem' [1,2,3] $(1 \in \{1,2,3\})$
 - * n 'mod' m
 - * ...
- User defined operators, built from symbols

3. Values and Types

Any Haskell expression e has a type t (e:: t) that is determined at compile time. The **type assigmnent**:: is either given explicitly or inferred by the compiler.

3.1. Base Types

Type	Description	values
Int	fixed-prec. integer	0, 1, (-42)
Integer	arbitrary prec. integer	10^100
Float, Double	single/double floating point (IEEE)	0.1, 1e02
Char	Unicode character	'x', '\t', '△', '\8710'
Bool	Boolean	True, False
()	Unit	()

3.2. Type Constructors

- Build new types from existing types
- Let a, b ... denote arbitrary types (type variables)

Type	Description	values
(a, b)	pairs of values of type a, b	(1, True) :: (Int, Bool)
$(\mathbf{a}_1,\mathbf{a}_2,\ldots\mathbf{a}_n)$	n-tuples	
[a]	list of values of type a	[True, False] :: [Bool], []::[a]
Maybe a	optional value of type a	Just 42 :: Maybe Int
		Nothing :: Maybe a
Either a b	choice	Left 'x' :: Either Char b
		Right pi :: Either a Double
IO a	I/O actions that return	print 42 :: IO ()
	a value of type a	
a -> b	functions from a to b	isLetter :: Char -> Bool

3.3. Currying

- Recall: e_1 ++ e_2 \equiv (++) e_1 e_2
- (++) e_1 $e_2 \equiv ((++)$ $e_1)$ e_2
- Function application happens one argument at a time. (Currying, Haskell B. Curry)
- Type of n-ary function is $a_1 \rightarrow a_2 \rightarrow \dots a_n \rightarrow b$
- Type fun -> associates to the right, read above type as $a_1 -> (a_2 -> (\dots (a_n -> b)))$
- Enables Partial Application

3.4. Defining Values (and thus functions)

• |= binds names to values. Names must not start with A-Z (Haskell style: camelCase)

- Define constant (0-ary function) c. Value of c is value of expression e.
 c = e
- Define n-ary function f with arguments x_i . f may occur in e.

```
f x_1 x_2 \dots x_n = e
```

- A Haskell program is a set of bindings.
- Good style: give type assignments for top-level (global) bindings:

3.4.1. Guards

Guards are conditional expressions (something like 'switch' in Java). They are a lot more readable and more powerful than if ...then ...else

Guards are introduced by | :

Guards (q_i) are expressions of type Bool, evaluated top to bottom.

Listing 3.1: factorial.hs

```
1 -- Compute n!
2 fac :: Integer -> Integer
3 fac n = if n <= 1 then 1 else n * fac (n - 1)
4
5 -- A reformulation using guards
6 fac' :: Integer -> Integer
7 fac' n | n <= 1 = 1
8 | otherwise = n * fac' (n - 1)</pre>
```

```
main :: IO ()
main = print $ (fac 10, fac' 10)
```

3.4.2. Local Definitions

1. Where bindings: local definitions visible in the entire rhs of a definition.

Listing 3.2: power.hs

```
- Efficient power computation, basic idea: x^2k = (x^2 + x^2)
     ^2) ^k
2
  power :: Double -> Integer -> Double
  power x k \mid k == 1
             | even k
                         = power (x * x) (halve k)
5
             | otherwise = x * power (x * x) (halve k)
6
    where
7
      even n = n 'mod' 2 == 0
8
      halve n = n 'div' 2
10
  main :: IO ()
 main = print $ power 2 16
```

2. Let expressions: local definitions visible inside one expression.

```
let g_1 = ...

g_2 = ...

...

g_o

in e
```

3.4.3. Lists

- Recursive definitions:
 - 1. [] is a list (nil), type [] :: [a]
 - 2. x:xs is a list, if x::a, xs::[a] (x is head, xs is tail)
- Notation: $3:(2:(1:[])) \equiv 3:2:1:[] \equiv [3,2,1] \equiv 3:[2,1]$
- Law: \forall xs :: [a] : $(xs \neq [])$ head xs : tail xs == xs

3.4.4. Pattern Matching

• The idiomatic Haskell way to define a function by cases:

Pattern	Matches If	Bindings in e_r
constant c	$\mathrm{x}_i == \mathrm{c}$	
variable v	always	$v \equiv x_i$
wildcard _	always	
tuple $(p_1, \ldots p_m)$	components of x_i match patterns p	
$[] \hspace{1cm} \mathrm{x}_i == []$		
$(p_1 : p_2)$ head x_i matches p_1 , tail x_i matches p_2		

Listing 3.3: tally.hs

```
-- Equivalent definitions of sum (over lists of integers)
  -- (1) Conditional expression
  sum' :: [Integer] -> Integer
  sum' xs = if xs == [] then 0 else head xs + <math>sum' (tail xs)
  -- (2) Guards
  sum'' :: [Integer] -> Integer
  sum'' xs | xs == [] = 0
            | otherwise = head xs + sum'' (tail xs)
10
11
  -- (3) Pattern matching
12
  sum''' :: [Integer] -> Integer
  sum''' [] = 0
  sum''' (x:xs) = x + sum''' xs
16
17
  main :: IO ()
  main = print $ (sum' [1..100], sum'' [1..100], sum'''
     [1..100])
```

Listing 3.4: take.hs

```
-- Finite prefix of a list
```

```
take' :: Integer -> [a] -> [a]

take' 0 _ = []

take' _ [] = []

take' n (x:xs) = x : take' (n - 1) xs

main :: IO ()
main = print $ take' 20 [1,3..]
```

Listing 3.5: mergesort.hs

```
-- Mergesort
  mergeSort :: (a -> a -> Bool) -> [a] -> [a]
  mergeSort (<<<) [] = []
  mergeSort (<<<) [x] = [x]
  mergeSort (<<<) xs = merge (<<<) (mergeSort (<<<) ls) (
     mergeSort (<<<) rs)</pre>
    where
       (ls, rs) = splitAt (length xs 'div' 2) xs
      merge :: (a -> a -> Bool) -> [a] -> [a] -> [a]
10
      merge (<<<) []
                          УS
                                     = ys
      merge (<<<) xs
                          []
12
      merge (<<<) (x:xs) (y:ys)
13
         | x <<< y = x:merge (<<<) xs (y:ys)
14
        | otherwise = y:merge (<<<) (x:xs) ys</pre>
 main :: IO ()
  main = print $ mergeSort (>) [1,3..19]
```