

DTM - Deterministic Turing Machine
 NTM - Non-deterministic Turing Machine
 PT - Probabilistic Turing Machine
 PDTM - Deterministic Turing Machine *running in polynomial time*
 PNTM - Non-deterministic Turing Machine *running in polynomial time*
 PPT - Probabilistic Turing Machine *running in polynomial time*
n - Input length

Time Complexity Classes

$$\begin{aligned}
 \text{DTIME}(f) &= \{A \in \Sigma^* \mid \exists \text{ DTM } T : T \text{ decides } A \wedge T \text{ runs in } \mathcal{O}(f(n)) \} \\
 \text{NTIME}(f) &= \{A \in \Sigma^* \mid \exists \text{ NTM } T : T \text{ decides } A \wedge T \text{ runs in } \mathcal{O}(f(n)) \}
 \end{aligned}$$

$$\begin{aligned}
 P &= \bigcup_{k \geq 1} \text{DTIME}(n^k) & NP &= \bigcup_{k \geq 1} \text{NTIME}(n^k) \\
 E &= \bigcup_{k \geq 1} \text{DTIME}(2^{kn}) & NE &= \bigcup_{k \geq 1} \text{NTIME}(2^{kn}) \\
 EXP &= \bigcup_{k \geq 1} \text{DTIME}(2^{n^k}) & NEXP &= \bigcup_{k \geq 1} \text{NTIME}(2^{n^k})
 \end{aligned}$$

Space Complexity Classes

$$\begin{aligned}
 \text{DSPACE}(s) &= \{A \in \Sigma^* \mid \exists \text{ DTM } T : T \text{ decides } A \wedge T \text{ uses } \mathcal{O}(s(n)) \text{ space} \} \\
 \text{NSPACE}(s) &= \{A \in \Sigma^* \mid \exists \text{ NTM } T : T \text{ decides } A \wedge T \text{ uses } \mathcal{O}(s(n)) \text{ space} \}
 \end{aligned}$$

$$\begin{aligned}
 \text{LOGSPACE} &= L = & \text{DSPACE}(\log n) \\
 \text{NLOGSPACE} &= NL = & \text{NSPACE}(\log n) \\
 \text{LINSPACE} &= & \text{DSPACE}(n)
 \end{aligned}$$

$$\begin{aligned}
PSPACE &= \bigcup_{k \geq 1} DSPACE(n^k) &= & NPSPACE = \bigcup_{k \geq 1} NSPACE(n^k) \\
ESPACE &= \bigcup_{k \geq 1} DSPACE(2^{kn}) &= & NESPACE = \bigcup_{k \geq 1} NSPACE(2^{kn}) \\
EXPSPACE &= \bigcup_{k \geq 1} DSPACE(2^{n^k}) &= & NEXPSPACE = \bigcup_{k \geq 1} NSPACE(2^{n^k})
\end{aligned}$$

Polynomial Hierarchy

$$\begin{aligned}
\Delta_0^P &= \Sigma_0^P = \Pi_0^P = P \\
\Delta_{k+1}^P &= P^{\Sigma_k^P} & \Delta_1^P &= P & \Delta_2^P &= P^{NP} \\
\Sigma_{k+1}^P &= NP^{\Sigma_k^P} & \Sigma_1^P &= NP & \Sigma_2^P &= NP^{NP} \\
\Pi_{k+1}^P &= co-NP^{\Sigma_k^P} & \Pi_1^P &= co-NP & \Pi_2^P &= co-NP^{NP}
\end{aligned}$$

$$PH = \bigcup_{k \geq 0} \Delta_k^P = \bigcup_{k \geq 0} \Sigma_k^P = \bigcup_{k \geq 0} \Pi_k^P$$

Probabilistic Complexity Classes

$$\begin{aligned} P &= \{A \in \Sigma^* \mid \exists \text{PDTM } T : \forall x \in A : P(T(x) = 1) = 1 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ NP &= \{A \in \Sigma^* \mid \exists \text{PNTM } T : \forall x \in A : P(T(x) = 1) > 0 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ PP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) \leq 0.5 \} \\ BPP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 + \epsilon \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) \leq 0.5 - \epsilon \} \\ RP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ ZPP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad P(T(x) = ?) < 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \wedge \\ &\quad P(T(x) = ?) < 0.5 \} \end{aligned}$$