

Complexity Classes

DTM - Deterministic Turing Machine

NTM - Non-deterministic Turing Machine

PT - Probabilistic Turing Machine

PDTM - Deterministic Turing Machine *running in polynomial time*

PNTM - Non-deterministic Turing Machine *running in polynomial time*

PPT - Probabilistic Turing Machine *running in polynomial time*

n - Input length

Time Complexity Classes

$\text{DTIME}(f) = \{A \in \Sigma^* \mid \exists \text{DTM } T : T \text{ decides } A \wedge T \text{ runs in } \mathcal{O}(f(n)) \}$

$\text{NTIME}(f) = \{A \in \Sigma^* \mid \exists \text{NTM } T : T \text{ decides } A \wedge T \text{ runs in } \mathcal{O}(f(n)) \}$

$$P = \bigcup_{k \geq 1} \text{DTIME}(n^k)$$

$$NP = \bigcup_{k \geq 1} \text{NTIME}(n^k)$$

$$E = \bigcup_{k \geq 1} \text{DTIME}(2^{kn})$$

$$NE = \bigcup_{k \geq 1} \text{NTIME}(2^{kn})$$

$$EXP = \bigcup_{k \geq 1} \text{DTIME}(2^{n^k})$$

$$NEXP = \bigcup_{k \geq 1} \text{NTIME}(2^{n^k})$$

Space Complexity Classes

$$\begin{aligned} \text{DSPACE}(s) &= \{ A \in \Sigma^* \mid \exists \text{ DTM } T : T \text{ decides } A \wedge T \text{ uses } \mathcal{O}(s(n)) \text{ space} \} \\ \text{NSPACE}(s) &= \{ A \in \Sigma^* \mid \exists \text{ NTM } T : T \text{ decides } A \wedge T \text{ uses } \mathcal{O}(f(n)) \text{ space} \} \end{aligned}$$

$$\begin{aligned} \text{LOGSPACE} &= L = & \text{DSPACE}(\log n) \\ \text{NLOGSPACE} &= \text{NL} = & \text{NSPACE}(\log n) \\ \text{LSPACE} &= & \text{DSPACE}(n) \end{aligned}$$

$$\begin{aligned} \text{PSPACE} &= \bigcup_{k \geq 1} \text{DSPACE}(n^k) = & \text{NPSPACE} &= \bigcup_{k \geq 1} \text{NSPACE}(n^k) \\ \text{EXPSPACE} &= \bigcup_{k \geq 1} \text{DSPACE}(2^{n^k}) = & \text{NEXPSPACE} &= \bigcup_{k \geq 1} \text{NSPACE}(2^{n^k}) \end{aligned}$$

Polynomial Hierarchy

$$\begin{aligned} \Delta_0^P &= \Sigma_0^P = \Pi_0^P = P \\ \Delta_{k+1}^P &= P^{\Sigma_k^P} & \Delta_1^P &= P & \Delta_2^P &= P^{NP} \\ \Sigma_{k+1}^P &= NP^{\Sigma_k^P} & \Sigma_1^P &= NP & \Sigma_2^P &= NP^{NP} \\ \Pi_{k+1}^P &= \text{co-NP}^{\Sigma_k^P} & \Pi_1^P &= \text{co-NP} & \Pi_2^P &= \text{co-NP}^{NP} \end{aligned}$$

$$PH = \bigcup_{k \geq 0} \Delta_k^P = \bigcup_{k \geq 0} \Sigma_k^P = \bigcup_{k \geq 0} \Pi_k^P$$

Probabilistic Complexity Classes

$$\begin{aligned} P &= \{A \in \Sigma^* \mid \exists \text{PDTM } T : \forall x \in A : P(T(x) = 1) = 1 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ NP &= \{A \in \Sigma^* \mid \exists \text{PNTM } T : \forall x \in A : P(T(x) = 1) > 0 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ PP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) \leq 0.5 \} \\ BPP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 + \epsilon \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) \leq 0.5 - \epsilon \} \\ RP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \} \\ ZPP &= \{A \in \Sigma^* \mid \exists \text{PPT } T : \forall x \in A : P(T(x) = 1) > 0.5 \wedge \\ &\quad P(T(x) = 0) = 0 \wedge \\ &\quad P(T(x) = ?) \leq 0.5 \wedge \\ &\quad \forall x \notin A : P(T(x) = 1) = 0 \wedge \\ &\quad P(T(x) = 0) > 0.5 \wedge \\ &\quad P(T(x) = ?) \leq 0.5 \} \end{aligned}$$

Complexity Classes Relations

Time Hierarchy Theorem

$$P \subset E \subset EXP$$

Borodin-Trakhtenbrot Gap-Theorem

\forall total computable function g with $g(n) \geq n$:
 \exists total computable function $f : DTIME(f) = DTIME(g \circ f)$

Space Hierarchy Theorem

$$L \subseteq NL \subset PSPACE \subset EXPSPACE$$

$$L \subseteq NL \subseteq P \subseteq NP \subseteq PH \subseteq PSPACE \subseteq EXP$$

Savitch

$$NSPACE(s(n)) \subseteq DSPACE(s(n)^2)$$

(compare to $NTIME(f(n)) \subseteq DTIME(2^{O(f(n))})$)

Immerman-Szelepcsnyi

$$s(n) \geq \log(n) \Rightarrow NSPACE(s(n)) = co-NSPACE(s(n))$$

Probabilistic World

$$\begin{array}{lll} P \subseteq ZPP \subseteq & RP \subseteq BPP & \subseteq PP \subseteq PSPACE \\ P \subseteq ZPP \subseteq & RP \subseteq NP & \subseteq PP \subseteq PSPACE \\ P \subseteq ZPP \subseteq & co-RP \subseteq co-NP & \subseteq PP \subseteq PSPACE \end{array}$$

$$RP \cap co-RP = ZPP$$

Misc

$$DIP = PV = NP \subseteq PH \subseteq PSPACE$$

$$P = NP \iff PH = P$$

$$NP \neq co-NP \Rightarrow P \neq NP$$