# OSVVM's Co-simulation Framework

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#### 1 Overview

The co-simulation features of OSVVM are a compliment to the Address Bus Model Independent Transactions features, as outlined in [1] and the Streaming Model Independent features (see [3]). It is highly recommended that these documents are read before this one, as it will be assumed that the reader is familiar with the transactions associated with that feature.

A typical address bus transaction test bench might look something like the diagram below for an AXI4 environment, which is based on the example test bench in AXI4/Axi4/testbench/TbAxi4Memory.vhd. with the Axi4Memory VC standing in for a unit under test component.

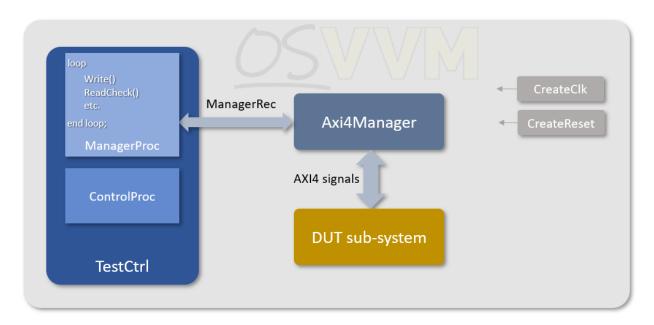


Figure 1: Address Bus Transaction Test Environment

In this example the test has a manger process using calls to to the address bus procedures, within an operation loop, to generate transaction via the ManagerRec which the Axi4Manager VC component translates into specific AXI4 signalling. In this case the target is another OSVVM VC, which is a memory component but would normally be the DUT component or sub-system. A subordinate process inspects arriving traffic and performs checks. In some scenarios the subordinate process generates the responses but, in this case, the Axi4Memory VC generates these internally. The advantage of this system is that the transactions are abstracted away from the bus specific model with re-usable components. If the Axi4Manager VC is replaced with another address bus VC, then the manager code remains unchanged and can drive the new system

The aim of the co-simulation features is to extend this architecture (and similarly for streaming) to allow the generation of the transactions to be performed within a C/C++ domain. It extends the calls to the address bus procedures to calls to methods within a C++ class. This opens up new possibilities not availabe in a pure VHDL domains. Some of these are listed below:

- Writing of tests in C++
  - o Test development environment extended to software engineers

- Running C++ models and code on simulated hardware
  - o Instruction set simulators or cycle accurate processor models
  - System models with a mix of software models and logic IP
- Connection to an external program via TCP/IP sockets
  - Where model might not be callable directly from the C++ code

With co-simulation of processor and system models, the possibility of developing software alongside the logic IP exists before silicon is available, but using both that actual hard and soft components together and not facsimiles, reducing risk.

The co-simulation features replace the transaction calls to TestCtrl's ManagerProc with a connection to the co-simulation interface. The diagram below shows the updated framework for address bus model independent AXI4 manager VC.

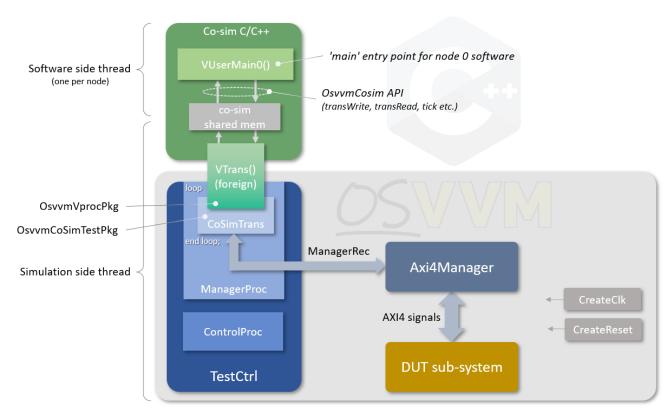


Figure 2: Co-simulation Environment Address Bus Example

In place of calls directly to the address bus transaction procedures, a CoSimTrans procedure is provided to replace these and drive ManagerRec instead. This hides details away for calls to the the underlying cosimulation code, with calls to foreign procedures that communicate with user program. The code, at startup will call a function (VUserMain0 for 'node 0', for example) that the user provides with all their code and sub-code called from there (much like main()). The user code now has access to the transaction API provided by the co-simulation code, and can generate read and write transactions in the OSVVM domain.

A complementary setup for subordinate responses is also available. Here the co-simulation features replace the transaction calls to TestCtrl's SubordinateProc with a connection to the co-simulation interface. The diagram below shows the updated framework for address bus model independent AXI4 subordinate VC.

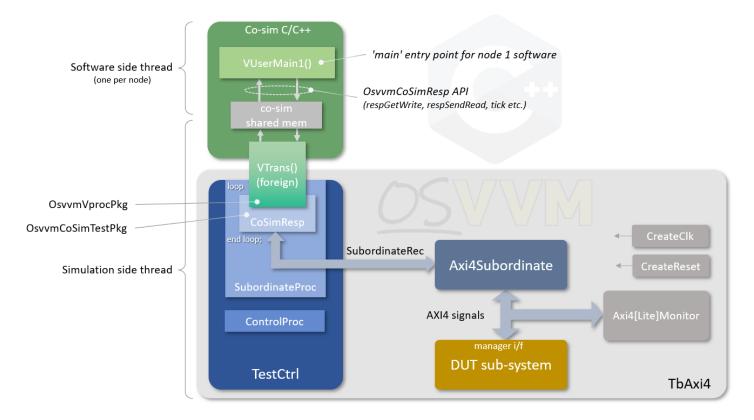


Figure 3: Co-simulation Environment Address Bus Responder Example

In place of calls directly to the address bus transaction procedures, a CoSimTransResp procedure is provided to replace these and drive SubordinateRec instead. This hides details away for calls to the the underlying co-simulation code, with calls to foreign procedures that communicate with user program. The code, at startup will call a function (VUserMain1 for 'node 1', for example) that the user provides with all their code and sub-code called from there. The user code now has access to the transaction API provided by the co-simulation code, and can send read responses and ftech received write transactions in the OSVVM domain. The complexity of the user code to model other system features is now unlimited, and can direct address bus accesses towards the simulation, and components modelled there via the API as appropriate.

A similar setup is found for the stream model independent setup. The diagram below shows an example test setup using the xMii MAC VC for Ethernet (with a setup for the xMii PHY VC being very similar):

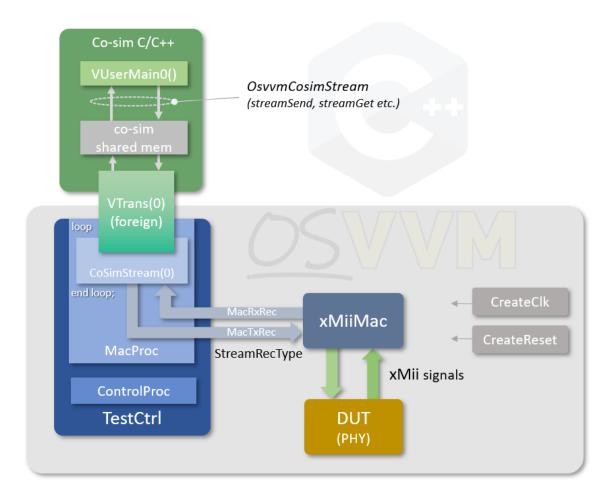


Figure 4: Co-simulation Environment Stream Example

The main features, then, of the co-simulation C++ code are as follows:

- Ability to support multiple, independent Address Bus Interfaces (such as AXI, Avalon etc.) for word and burst transfers
- Multiple nodes to drive separate AddressBusRecType transaction signals from multiple processes
- Ability to handle interrupts in the software via callbacks
- Ability to advance simulation time (without a transaction)

In the rest of this document will be detailed the usage of the co-simulation features to drive address bus transactions from the VHDL/OSVVM side and the C++ side.

# 2 Prerequisites for Windows

To run the co-simulation features on Windows requires the use of MSYS2 and the gcc toolchains for 64-and 32-bit compilation, as well as make and python3. The list below provides a link to the MSYS2 installer. Once installed, the required gcc modules can be installed as shown for the latest MSYS2 at time of writing (recommend updating existing installations before proceeding).

• msys2

- mingw-w64 32- and 64-bit gcc toolchains, with required non-included libraries
- From an MSYS2 MINGW64 shell:

```
    pacman -S mingw-w64-x86_64-gcc
    pacman -S mingw-w64-i686-gcc
    pacman -S mingw-w64-x86_64-dlfcn
    pacman -S mingw-w64-i686-dlfcn
```

- Utility functions
- From an MSYS2 MINGW64 shell:

```
    pacman -S make
    pacman -S python3
    pacman -S rlwrap (for GHDL and NVC)
    pacman -S mingw-w64-x86_64-tclib (for GHDL and NVC)
```

Note that **ModelSim** require running from the MSYS2 MINGW32 shell, as it is a 32-bit executable and the co-simulation code must be compiled for 32-bits, whereas all the other simulators require using the MSYS32 MINGW64 shell for 64-bit compilation.

If **GHDL** is required to be used, the mingw-w64-x86\_64-ghd1-llvm package must be installed with pacman and should be run from the MSYS2 MINGW64 shell. If must also be ensured that tclsh is available (pacman -S tcl).

For Active-HDL or Riviera-PRO to use the local MSYS2 environment and toochain the tool installation's mingw folder should be renamed (e.g., to \_mingw) to prevent interference. Alternatively PATH can be setupt to point to <path to Aldec tool>/mingw/bin to use the bundled toolchain and the LD\_LIBRARY\_PATH updated with <path to Aldec tool>/mingw/x86\_64-w64-mingw32/lib.

#### 3 Running the Co-simulation Examples

See the steps in the OSVVM Overview guide for running the OSVVM demos. Do the steps shown in Figure 5 in your simulator. StartUp.tcl needs to be sourced each time you start the simulator. See the Script User Guide for additional details for Aldec's ActiveHDL and GHDL.

```
cd sim
source ../OsvvmLibraries/Scripts/StartUp.tcl
build ../OsvvmLibraries/OsvvmLibraries.pro
build ../OsvvmLibraries/CoSim/RunAllTests.pro
```

Figure 5. Compiling and Running Co-simulation tests

# 4 Co-simulation Supported Platforms

Below is a list of simulators and the tested support for them for the co-simulation features of OSVVM.

Vendor	Simulator	Linux	Windows
Siemens	ModelSim	✓	✓
Siemens	QuestaSim	✓	✓
Open Source	GHDL	✓	✓
Aldec	Active-HDL	N/A	✓
Aldec	Riviera-PRO	✓	✓
Open Source	NVC	✓	✓

**Table 1: Supported Simulators** 

#### 5 OSVVM Co-simulation Address Transaction Procedure

To add address bus transaction co-simulation to the VHDL test logic in OSVVM, procedures are provided in OsvvmTestCoSimPkg. The procedures stand in for address bus transaction generation logic such as that demonstrated, for example, in the AXI4/Axi4/TestCases/TbAxi4\_MemoryReadWrite.vhd test. The co-simulation procedures use the standard OSVVM test environments driving the AddressBusRecType records, usually from a manager process, to instigate bus reads and writes on, for instance, AXI4 or Axi4Lite busses, as dictated by the test bench used. The instigators of the transactions are then from C++ programs compiled and run with the simulation instead of direct calls to the OSVVM Address Bus Transaction procedures [1]. The co-simulation demonstration programs use the TbAxi4Memory.vhd test benches fro Axi4 and Axi4Lite.

Before any calls to co-simulated code can be made, the code must be initialised for *all* nodes to be used. Separate programs can be run concurrently if needed (though not usual), and each must have a unique node ID. This is done using the CoSimInit procedure, as shown below.

```
procedure CoSimInit (

variable NodeNum : in integer := 0
);
```

The procedure would normally be called from within the ControlProc of the test case, or possibly from within the MangerProc if only a single node used, before any calls made to other co-simulation procedures.

To generate address bus transactions from the co-simulation software repeated calls to a procedure CoSimTrans must be made from a process (in a loop for example) that drive an AddressBusRecType signal. This would normally be from a test case's MangerProc, but multiple driving processes are allowed and each can make calls to CoSimTrans, so long as each uses a unique mode ID and that node has been initialised with a call to CoSimInit (see, for example,

CoSim/TbInterrupt/TestCases\_Interrupt/TbAb\_InterruptCosim1.vhd). The CoSimTrans procedure is defined below:

```
procedure CoSimTrans (

signal ManagerRec : inout AddressBusRecType;
variable Done : inout integer;
variable Error : inout integer;
variable IntReq : in integer := 0;
variable NodeNum : in integer := 0
);
```

The main argument to the procedure is the ManagerRec that is used to send a transaction request to the address bus transaction VC, such as an AXI4 manager VC. The Done, and Error arguments indicate status from the co-simulation software. These must be connected to local variables in the calling process in order for state to be preserved between calls. The calling procedure can choose to use these signals or not, but that inidcate state that is useful in controlling the simulation from state that the software has indicated.

The Done argument is an indication that the co-simulation software has completed its run, when non-zero, and will no longer generate output. The calling process may use this to halt testing, but the status is an indication that the software is complete, rather than a request to halt. In systems with multiple nodes, a test may wish only to stop when all nodes have indicated that they are done.

The co-simulation software can also have internal errors and self-check failures. It communicates this by setting the Error argument to a non-zero value. The calling process can use this to assert a failure and even stop at the first instance this error state.

The IntReq input argument is used to communicate interrupt status to the co-simulation software. This boolean, when true, will call a function in the co-simulation software if one has been registered as an interrupt callback function (see the next section). The callback function is called for every call to CoSimTrans that the input is true, so if the source of the interrupt in the test bench is a level driven interrupt, the calling process should only set the IntReq true when the state changes and the callback to set/clear interrupt state to the software on alternate calls. An edge driven interrupt would need the IntReq true just for cycles on its active edge. For environments that do not use the interrupt callback features, the argument can be left off and the default false value used.

The NodeNum argument is the node ID. When only one co-simulation process present, this can be left off and the node defaults to 0. Where multiple processes are generating transactions each must use a unique NodeNum.

This, then, is all that is required from the VHDL/OSVVM environment side to be able to generate transactions for both word/sub-word and burst address bus transactions.

#### 6 OSVVM Address Bus Transaction C++ API

The user side C++ code has access to an API in order to generate address bus transactions, allow simulation time to advance and to generate status. It can also get interrupt status from the simulation if present. A simple example for a program is given in Figure 7. For each node used in the OSVVM test bench, there is expected to be a user written function named VUserMainn, where n is the number matching the node ID used (e.g. VUserMain0). This is like the main() of a C/C++ program for the node, or perhaps more like WinMain() for Windows non-console programs. An example definition is shown below:

The source code for the user code would normally be gathered into a single sub-directory for a given test. This might include just a single source file (e.g. VUserMain0.cpp) but can have sub-files to any complexity and would also include the code for all active nodes. All the source code then has access to the co-simulation API.

#### 6.1 OsvvmCosim Class Constructor

The API is defined in the OsvvmCosim.h header in CoSim/code. This is a C++ wrapper class to abstract away the lower level calls to the co-simulation infrastructure software. A constructor defines which node the object is attached to.

```
OsvvmCosim (
   int nodeIn = 0,
   std::string test_name = ""
);
```

Any user file can construct the API object. This class is a wrapper and does not hold state (except the node number) and so sub-programs can create their own API objects (with the correct node for the top level program) to gain access to the API methods.

In order to identify the different co-simulation tests, that may all uses the same underlying test bench, the software can set a test name when constructing an OsvvmCosim object using the test\_name argument. By default, the argument is set to "" and the test name will not be set. Any other value will be used to set the test name in the OSVVM environment.

# 6.2 Advancing Time

Advancing time without generating transactions is done with the tick() method as defined below.

```
int tick (
   const int ticks,
   const bool done = false,
   const bool error = false
);
```

The method is called with a ticks parameter to indicate how many calls to the CoSimTrans VHDL procedure should be made without generating a transaction. If the calling process does not add any additional clock waits in the loop calling CoSimTrans, then the ticks will be clock cycles. The optional done and error inputs send status to the equivalent arguments on CoSimTrans as explained in the previous section.

The transaction methods are the methods that are used to generate the address bus activity.

#### 6.3 Transaction methods

The basic available write and read transaction methods are listed below for both single word transactions and burst transactions.

```
uint<nn> t transWrite (
______
 const uint<nn>_t addr,
const uint<nn>_t data,
 const int
       prot = 0 );
______
uint<nn> t transWriteAsync (
______
 const uint<nn>_t addr,
 const uint<nn>_t data,
 const int prot = 0 );
______
uint<nn> t transWriteAndRead (
-----
 const uint<nn>_t addr,
 const uint<nn>_t wdata,
 const uint<nn>_t *rdata,
 const int prot = 0 );
```

```
uint<nn> t transWriteAndReadAsync (
______
 const uint<nn>_t addr,
 const uint<nn> t wdata,
         prot = 0 );
 const int
______
void transRead (
______
 const uint<nn>_t addr,
 const uint<nn>_t *data,
 const int
           prot = 0);
______
bool transReadPoll (
______
 const uint<nn>_t addr,
 const uint<nn>_t *data,
 const int
          idx,
 const int bitval,
const int waittime
const int prot = 0
          waittime = 10,
);
______
void transBurstWrite (
______
 const uint<nn>_t addr,
 const int
          bytesize,
 const int
          prot = 0);
void transBurstWriteAsync (
 const uint<nn>_t addr,
    uint8 t *data,
 const int
           bytesize,
          prot = 0);
 const int
void transBurstRead (
______
 const uint<nn>_t addr,
    uint8_t
          *data,
 const int
           bytesize,
 const int
           prot = 0);
```

These transaction methods are overloaded for various sizes of address and data , using the stdint.h definitions, where uint<nn>\_t indicates different uint sizes. For address parameters this is uint32\_t or

uint64\_t. For the data parameters of transRead and transWrite this is one of uint8\_t, uint16\_t, uint32\_t or uint64\_t (if address type is also uint64\_t). Note that using uint64\_t can only be done if the target test bench is configured for a 64-bit bus architecture.

The single word methods, transRead and transWrite, and its non-blocking asynchronous equivalent transWriteAsync, take an address, data and optional prot (protection) arguments, with the transRead having a pointer to the variable for returning the read data. The transWrite returns a value which is a future-proofing for returning read data in the same transaction as a write, but is not yet supported. The transRead method does not return a value and a pointer is used to return the data as this is how the method selects the size of the read transaction (byte, half-word etc.) based on the argument's size. The transWriteAndRead method does a write to the specified address with the wdata value, and returns the old value in the rdata argument. The transWriteAndReadAsync method is simlar, but the read data is not returned at the call. The read transaction is instigated, however, and the data can be fetched with methods such as transReadData. The transReadPoll method will poll a memory location specified by the addr argument, at intervals defined by the waittime argument, until the data read has its bit indexed at idx matching the value in bitval. The value read on matching is returned in data.

The burst transactions use byte buffers to send and return data. The calling code provides buffer space to a maximum of 4Kbytes preloaded with byte data for writes, or updated with byte data on reads. If no data is specified, then the burst write will use whatever has been pushed to the write fifo using transBurstPushData (see later). The bytesize parameter indicates the transfer size. The co-simulation code will not go beyond 4Kbytes when reading or writing to the buffers, so it is recommended to use 4Kbyte buffers in all cases to avoid overrun. A non-blocking version of transBurstWrite is provided as transBurstWriteAsync. In addition, the burst writes can have predetermined patterns as a source of data, namely incrementing or random, with a data byte provided to indicate the first word to use for the pattern. The methods are shown below:

```
void transBurstWriteIncrement (

const uint<nn>_t addr,
  const uint8_t firstbyte,
  const int bytesize,
  const int prot = 0);

void transBurstWriteIncrementAsync (

const uint<nn>_t addr,
  const uint8_t firstbyte,
  const int bytesize,
  const int prot = 0);
```

```
void transBurstWriteRandom (
    const uint<nn>_t addr,
    const uint8_t firstbyte,
    const int bytesize,
    const int prot = 0);

void transBurstWriteRandomAsync (
    const uint<nn>_t addr,
    const uint8_t firstbyte,
    const int bytesize,
    const int bytesize,
    const int prot = 0);
```

# 6.3.1 Split Transactions

For VCs that support it, methods are provided for split transactions—i.e., sending address and data seperately, at different times. The methods to support this are listed below:

```
______
void transWriteAdressAsync (
______
const uint<nn>_t addr,
const int prot = 0);
______
uint<nn>_t transWriteDataAsync (
______
const uint<nn>_t data,
const uint32_t bytelane,
const int
       prot = 0 );
______
void transReadAdressAsync (
______
const uint<nn>_t addr,
const int prot = 0);
void transReadData (
-----
const uint<nn>_t *data
);
```

```
bool transTryReadData (
------
const uint<nn>_t *data
);
```

The transTryReadData is the same as transReadData, but will return if no data available, rather than block. It will return false if no data available and true if data was availabe, which is then returned in the data buffer.

# 6.3.2 Push and Pop Burst Data

Data can be pushed and pooped from the transaction fifos independently of transaction method calls for use in split transactions etc. In addition, the fifos can be filled with predetermined patterns. The methods for this are listed below.

```
______
void transBurstPushData (
______
const uint8 t *data,
const int bytesize);
______
void transBurstPushIncrement (
______
const uint8_t data,
const int bytesize);
void transBurstPushRandom (
_____
const uint8_t data,
const int bytesize);
______
void transBurstPopData (
______
const uint8_t *data,
const int bytesize);
```

# 6.3.3 Read Checks

Methods are provided that do check on data as they are read for eith words or bursts. This can be using provided data, using data already pushed to the Fifo, or predetermined patterns. The provided methods for word transactions are listed below:

The first method does a blocking read transaction, and compares the read value with that supplied as the data argument. The second method does the same check, but with already received data from a previous read. The last method is the same as the second but is non-blocking, returning false if no word available and true if data was available and was checked.

The burst transactions also have check methods, with methods that instigate a transaction and check, and methods that check previously received data. In addition, the data can be checked agains a supplied set of data, or against a predetermined pattern. The provided methods are shown below:

```
void transBurstReadCheckData (
    const uint<nn>_t addr,
    const uint8_t *data,
    const int prot = 0);

void transBurstCheckData (
    const uint8_t *data,
    const int prot = 0);

void transBurstReadCheckIncrement (
    const uint
void transBurstReadCheckIncrement (
    const uint
void transBurstReadCheckIncrement (
    const uint
void transBurstReadCheckIncrement (
```

```
void transBurstReadCheckRandom (
    const uint<nn>_t addr,
    const uint8_t data,
    const int prot = 0);

void transBurstCheckDataRandom (
    const uint8_t data,
    const int prot = 0);
```

#### 6.4 Directive Methods

Some directive methods are provided for transaction statistics and synchronization, as listed below:

```
int transGetTransactionCount (
______
void
);
int transGetWriteTransactionCount (
______
void
);
______
int transGetReadTransactionCount (
______
void
);
void transWaitForTransaction (
______
void
);
______
void transWaitForWriteTransaction (
void
);
```

```
void transWaitForReadTransaction (
    void
);
```

#### 6.5 Interrupt Callback

User code can register a callback function which will be called by the simulation every time the IntReq input to the CoSimTrans procedure changes state in the OSVVM test process. To register an interrupt callback function the following method is used:

#### 6.6 Additional Methods

As well as the methods outlined above a couple of other methods are provided to iad usage of the class. These are shown below:

The getNodeNumber method simply returns the node number associated with the instantiation object of the class. As stated before, mutiple object are allowed for access to a given node. To aid in managing which node a particular object has been associated with, this method gives access to that state.

The waitForSim method is for running in GHDL when in callable simulation mode. For all other modes, this call does nothing. When active, a call to this method blocks until the simulation is in a suitable state to have transaction methods called to instigate traffic on the bus. For more details see section 12.6, which also has example code using this method.

# 6.7 OsvvmCosimInt Interrupt Support Class

The OsvvmCosimInt class, derived from the OsvvmClass, adds additional programming support for handling interrupts for new test code that has no other means of processing interrupts such as when

using a software processor model. This inherits all the methods of the parent class, as described in the sections above, including the regInterruptCB() method.

The type of callback function, as defined by pVUserInt\_t, is int <myfunc>(int). When the callback is executed the int argument contatins the current state of the IntReq input to the CoSimTrans procedure. Example code can be found in OsvvmLibraries/CoSim/test/interruptCB. In addition it provided methods for registering up to 32 interrupt service routine (ISR) functions:

Each function registered to a valid level (from 0 to 31) will be called if the interrupt level indicates that it is active, it is enabled and no higher priority is active. The priority has 0 as the highest and 31 as the lowest. Various methods are provided for controlling whether an interrupt is enabled or not. A master interrupt can eb enabled or disabled to globally control all interrupts.

```
______
void enableMasterInterrupt (
-----
 void
);
______
void disableMasterInterrupt (
______
 void
);
Each individual level also has an enable and disable method for individual control of each ISR:
______
void enableIsr (
-----
 const int level
);
______
void disableIsr (
 const int level
);
In order to update interrupt state, a method is provided to pass in the current interrupt request status:
-----
void updateIntReq (
______
 const int intReq
```

);

This method would normally be used in conjunction with a callback function registered with the regInterruptCB() method, which is called whenever the interrupt request input to CoSimTrans changes. This function would have access to the OsvvmCosimInt object and call its updateIntReq() method to pass in the interrupt request status.

#### 7 OSVVM Co-simulation Address Bus Responder Procedure

As a complimentary feature to the address bus model independent transaction features detailed in the last two sections an address bus transaction responder is also provided. This is used, for example, when testing IP with a manager interface and the test environment is required to generate the responses to read and write requests. From the VHDL side, the responder re-uses the CoSimInit procedure (see section 5) to initialise a node for a responder co-simulation user software. The supplied node number must be unique for each instance that will be used, and from any other co-simulation elements for transactions and streams).

To generate address bus transactions responser from the co-simulation software repeated calls to a procedure CoSimResp must be made from a process (in a loop for example) that drive an AddressBusRecType signal. This would normally be from a test case's SubordinateProc, but multiple driving processes are allowed and each can make calls to CoSimResp, so long as each uses a unique mode ID and that node has been initialised with a call to CoSimInit (see, for example, CoSim/TestCases/TbAb\_Responder.vhd). The CoSimResp procedure is defined below:

```
procedure CoSimResp (
    signal SubordinateRec : inout AddressBusRecType ;
    variable Done : inout integer ;
    variable Error : inout integer ;
    variable NodeNum : in integer := 1
) ;
```

The main argument to the procedure is the SubordinateRec that is used to send transaction responses to the address bus transaction VC, such as an AXI4 manager VC, or to an AXI4 manager interface on a DUT. The Done, and Error arguments indicate status from the co-simulation software. These must be connected to local variables in the calling process in order for state to be preserved between calls. The calling procedure can choose to use these signals or not, but they indicate state that is useful in controlling the simulation from that indicated from software.

The Done argument is an indication that the co-simulation software has completed its run, when non-zero, and will no longer generate output. The calling process may use this to halt testing, but the status is an indication that the software is complete, rather than a request to halt. In systems with multiple nodes, a test may wish only to stop when all nodes have indicated that they are done.

The co-simulation software can also have internal errors and self-check failures. It communicates this by setting the Error argument to a non-zero value. The calling process can use this to assert a failure and even stop at the first instance this error state.

The NodeNum argument is the node ID. When only one co-simulation process present, this can be left off and the node defaults to 1 for the responder. Where multiple processes are generating transactions each must use a unique NodeNum.

#### 8 OSVVM Address Bus Responder C++ API

The user side C++ code has access to an API in order to generate address bus reponses, allow simulation time to advance and to generate status. For each node used in the OSVVM test bench, there is expected to be a user written function named VUserMainn, where n is the number matching the node ID used (e.g. VUserMain1, with a node number of 1 being the default for a responder). This is like the main() of a C/C++ program for the node, or perhaps more like WinMain() for Windows non-console programs. An example definition is shown below:

The source code for the user code would normally be gathered into a single sub-directory for a given test. This might include just a single source file (e.g. VUserMain1.cpp) but can have sub-files to any complexity and would also include the code for all active nodes. All the source code then has access to the co-simulation API.

#### 8.1 OsvvmCosim Class Constructor

The API is defined in the OsvvmCosimResp.h header in CoSim/code. This is a C++ wrapper class to abstract away the lower level calls to the co-simulation infrastructure software. A constructor defines which node the object is attached to.

```
OsvvmCosimResp (
    int nodeIn = 0,
    std::string test_name = ""
);
```

Any user file can construct the API object. This class is a wrapper and does not hold state (except the node number) and so sub-programs can create their own API objects (with the correct node for the top level program) to gain access to the API methods.

In order to identify the different co-simulation tests, that may all uses the same underlying test bench, the software can set a test name when constructing an OsvvmCosim object using the test\_name

argument. By default, the argument is set to "" and the test name will not be set. Any other value will be used to set the test name in the OSVVM environment.

# 8.2 Advancing Time

Advancing time without generating transactions is done with the tick() method as defined below.

```
int tick (
    const int ticks,
    const bool done = false,
    const bool error = false
);
```

The method is called with a ticks parameter to indicate how many calls to the CoSimTrans VHDL procedure should be made without generating a transaction. If the calling process does not add any additional clock waits in the loop calling CoSimTrans, then the ticks will be clock cycles. The optional done and error inputs send status to the equivalent arguments on CoSimTrans as explained in the previous section.

#### 8.3 Response methods

The response methods are the methods that are used to generate the address bus response activity. The basic available write and read transaction methods are listed below:

```
______
void respGetWrite (
_____
uint<nn>_t *addr,
uint<nn>_t *data
);
______
bool respTryGetWrite (
_____
uint<nn>_t *addr,
uint<nn>_t *data
);
------
void respSendRead (
-----
uint<nn>_t *addr,
uint<nn>_t data
);
______
```

```
bool respTrySendRead (
------
uint<nn>_t *addr,
uint<nn>_t data
);
```

These reponse methods are overloaded for various sizes of address and data, using the stdint.h definitions, where uint<nn>\_t indicates different uint sizes. For address parameters this is uint32\_t or uint64\_t. For the data parameters of transRead and transWrite this is one of uint8\_t, uint16\_t, uint32\_t or uint64\_t (if address type is also uint64\_t). Note that using uint64\_t can only be done if the target test bench is configured for a 64-bit bus architecture.

The methods respGetWrite and respSendRead, and their non-blocking asynchronous equivalents respTryGetWrite, and respTrySendData, take and address pointer argument, which returns the address associated with the received transaction. The write methods take a data pointer argument in which the written data is returned. The read methods take a data input argument to send as the read response data. The respGetWrite and respSendRead methods are blocking, and wil not return until the transaction is completed. For writes, this means that a whole write transaction has been received, and for reads, that the repsonse data has been sent. The try methods return a boolean status. If the response can't be processed it will return false. For writes, this means no write data is available. For reads, this means that no read address has been received.

# 8.3.1 Split Response Methods

For VCs that support it, methods are provided for split transactions—i.e., fetching (or sending, where applicable) address and data seperately, at different times. The methods to support this are listed below.

```
void respTryGetWriteData (
______
uint<nn>_t *data
______
void respGetReadAddress (
______
uint<nn>_t *addr
);
______
void respTryGetReadAddress (
_____
uint<nn>_t *addr
);
______
void respSendReadData (
        -----
uint<nn>_t data
);
bool respSendReadDataAsync (
______
uint<nn>_t data
);
```

There are methods for fetching both read and write received addresses, with respGetWriteAddress and respGetReadAddress being block, waiting until an address is available. The non-blocking versions (respTryGetWriteAddress and respTryGetReadAddress) return false if no address is available. Retrieving write data is done with respGetWriteData (blocking) and respTryGetWriteData (non-blocking, returning false if no write data available). To send the response data, respSendReadData (blocking until a read address is available) and respSendReadData (non-blocking). This last does return the status of whether a read address was availabe at the point the methods was called, but the data is placed in the respond fifo regardless.

# 8.4 Directive Methods

Some directive methods are provided for transaction statistics and synchronization, as listed below:

```
int respGetTransactionCount (
______
void
);
______
int respGetWriteTransactionCount (
void
);
______
int respGetReadTransactionCount (
______
void
);
______
void respWaitForTransaction (
______
void
);
_____
void respWaitForWriteTransaction (
void
);
void respWaitForReadTransaction (
______
void
);
```

#### 8.5 Additional Methods

As well as the methods outlined above a couple of other methods are provided to iad usage of the class. These are shown below:

```
void
);
```

The getNodeNumber method simply returns the node number associated with the instantiation object of the class. As stated before, mutiple object are allowed for access to a given node. To aid in managing which node a particular object has been associated with, this method gives access to that state.

The waitForSim method is for running in GHDL when in callable simulation mode. For all other modes, this call does nothing. When active, a call to this method blocks until the simulation is in a suitable state to have transaction methods called to instigate traffic on the bus. For more details see section 12.6, which also has example code using this method.

#### 9 OSVVM Co-simulation Stream Transaction Procedure

To enable streaming transaction co-simulation in the VHDL test logic of OSVVM, procedures are provided in OsvvmTestCoSimPkg. The procedures stand in for streaming transaction generation logic such as that demonstrated, for example, in the Ethernet/TestStandAlone/Tb\_xMii.vhd test. The co-simulation procedures use the standard OSVVM test environments driving the StreamRecType records, usually from a manager process, to instigate bus reads and writes on, for instance, Ethernet or UART interconnects, as dictated by the test bench used. The instigators of the transactions are then from C++ programs compiled and run with the simulation instead of direct calls to the OSVVM Streaming Transaction procedures [3].

Just as for the adress bus transactions, before any calls to co-simulated code can be made, the code must be initialised for *all* nodes to be used. Separate programs can be run concurrently if needed (though not usual), and each must have a unique node ID. This is done using the CoSimInit procedure, as shown below.

To generate address bus transactions from the co-simulation software repeated calls to a procedure CoSimStream must be made from a process (in a loop for example) that drive a StreamRecType signal. This would normally be from a test case's MangerProc, but multiple driving processes are allowed and each can make calls to CoSimStream, so long as each uses a unique mode ID and that node has been initialised with a call to CoSimInit. The CoSimStream procedure is defined below:

```
procedure CoSimStream (

signal TxRec : inout StreamRecType;
signal RxRec : inout StreamRecType;
variable Done : inout integer;
variable Error : inout integer;
variable NodeNum : in integer := 0);
```

The main arguments to the procedure are the TxRec and RxRec that are used to send transaction requests and receive transactions to and from a streaming transaction VC, such as an Ethernet VC. The Done, and Error arguments indicate status from the co-simulation software. These must be connected to local variables in the calling process in order for state to be preserved between calls. The calling procedure can choose to use these signals or not, but that inidcate state that is useful in controlling the simulation from state that the software has indicated.

The Done argument is an indication that the co-simulation software has completed its run, when non-zero, and will no longer generate output. The calling process may use this to halt testing, but the status is an indication that the software is complete, rather than a request to halt. In systems with multiple nodes, a test may wish only to stop when all nodes have indicated that they are done.

#### 10 OSVVM Streaming C++ API

Like for the Address Bus Transaction, the streaming features have a C++ API in order to generate streaming transactions, allow simulation time to advance and to generate status. User code is also entered via a VUserMain*n* function, with *n* being the node number (defaulting to 0).

#### 10.1 OsvvmCosimStream Class Constructor

The Streaming API is provided via a class, OsvvmCosimStream, defined in OsvvmCosimStream.h, which can be found in CoSim/code. This is a C++ wrapper class to abstract away the lower level calls to the cosimulation infrastructure software. A constructor defines which node the object is attached to.

Any user file can construct the API object. This class is a wrapper and does not hold state (except the node number) and so sub-programs can create their own API objects (with the correct node for the top level program) to gain access to the API methods. In order to identify the different co-simulation tests, that may all uses the same underlying test bench, the software can set a test name when constructing an OsvvmCosim object using the test\_name argument. By default, the argument is set to "" and the test name will not be set. Any other value will be used to set the test name in the OSVVM environment.

#### 10.2 Advancing Time

Advancing time without generating transactions is done with the tick() method as defined below.

```
int tick (
const int ticks,
const bool done = false,
const bool error = false );
```

The method is called with a ticks parameter to indicate how many calls to the CoSimStream VHDL procedure should be made without generating a transaction. If the calling process does not add any additional clock waits in the loop calling CoSimStream, then the ticks will be clock cycles. The optional done and error inputs send status to the equivalent arguments on CoSimStream as explained in the previous section. The transaction methods are the methods that are used to generate the streaming activity.

#### 10.3 Transaction Methods

The basic streaming transaction methods are shown below:

```
uint<nn>_t streamSend (
______
const uint<nn>_t data,
const int param = 1
);
uint<nn> t streamSendAsync (
-----
const uint<nn>_t data,
const int param = 1
);
void streamGet (
______
 uint<nn>_t *data
);
______
void streamGet (
______
uint<nn>_t *data,
 const int *status
);
void streamBurstSend (
______
   uint8_t *data,
const int bytesize
);
```

```
void streamBurstSend (
______
const int bytesize
void streamBurstSendAsync (
______
   uint8_t *data,
const int bytesize
);
______
void streamBurstSendAsync (
______
 const int
        bytesize
);
void streamBurstGet (
______
   uint8_t *data,
const int bytesize
);
```

The streamSend and streamGet transaction methods are overloaded for various sizes of data, using the stdint.h definitions, where uint<nn>\_t indicates different uint sizes. For the data parameter, this is one of uint8\_t, uint16\_t, uint32\_t or uint64\_t. There are also non-bloking variants of the of the transSend and transBurstSend methods; transSendAsync and transBurstSendAsync. In addition the transBurstSend and transBurstSendAsync can have a data buffer pointer argument or not. When supplied the buffer data is pushed to the transaction fifo before the transaction is instigated. When no data buffer argument is given, the data already in the fifo is used, loaded previously via the push methods (see below).

The streamSend and streamSendAsync optional param argument is used by some VCs. For example, the UART VC uses this to inject parity errors and other conditions. The default value is for no errors. The streamGet also has an optional second argument, status, which can be used to get status returned from a streaming VC, if it supports this. For example, the UART VC can indicate receiving a a transaction with a parity error.

There are also pattern generating versions of the burst send methods. These can generate patterns in the fifo as either incrementing or rendom, with a data argument specifying the first byte of the pattern. The methods are listed below.

```
void streamBurstSendIncrement (
______
   uint8_t data,
const int bytesize,
const int param = 1
);
______
void streamBurstSendIncrementAsync (
______
   uint8 t data,
const int bytesize,
const int param = 1
);
______
void streamBurstSendRandom (
______
   uint8_t data,
const int bytesize,
const int param = 1
);
______
void streamBurstSendRandomAsync (
______
   uint8_t data,
const int bytesize,
const int param = 1
);
```

The streaming burst transactions use byte buffers to send and return data. The calling code provides buffer space to a maximum of 4Kbytes preloaded with byte data for writes, or updated with byte data on reads. The bytesize parameter indicates the transfer size. The co-simulation code will not go beyond 4Kbytes when reading or writing to the buffers, so it is recommended to use 4Kbyte buffers in all cases to avoid overrun.

#### 10.3.1 Push and Pop Methods

Various methods are provided to push and pop data from the transaction fifos, both Tx and Rx. The push methods are either for transmission and are pushed to the Tx fifo, or are for comparison with received data, when they push to the Rx fifo (these are the push to check methods). There are also methods to load the fifos with predetermined patterns, namely increment and random petterns, with the first byte specified as a parameter. The methods are listed below:

```
______
void streamBurstPushData (
______
   uint8_t *data,
const int
    bytesize
);
______
void streamBurstPushCheckData (
______
   uint8_t *data,
const int
     bytesize
);
______
void streamBurstPushIncrement (
______
   uint8_t data,
const int bytesize
);
______
void streamBurstPushCheckIncrement (
______
   uint8_t data,
      bytesize
const int
);
_____
void streamBurstPushRandom (
______
   uint8_t data,
const int bytesize
);
_____
void streamBurstPushCheckRandom (
______
   uint8_t data,
const int bytesize
);
______
void streamBurstPopData (
______
   uint8_t *data,
const int bytesize
);
```

#### 10.3.2 Check Methods

A set of methods are provided to check data word and burst data in the receive fifo, and these are listed below.

```
_____
void streamCheck (
______
 const uint<nn>_t data,
const int param = 0
);
______
void streamBurstCheck (
_____
   uin8_t *data
const int bytesize,
const int param = 1
);
void streamBurstCheck (
______
const int bytesize,
const int param = 1
);
void streamBurstCheckIncrement (
______
    uin8_t data
const int bytesize,
const int param = 1
        param = 1
);
______
void streamBurstCheckRandom (
______
    uin8_t data
const int bytesize,
const int param = 1
);
```

# 10.3.3 Non-blocking Get and Check

Methods are provided which are non-blocking get and check operations. These return false if no data is available, otherwise they return true and return or check the data. These are listed below.

```
-----
bool streamTryGet (
______
uint<nn>_t *data
_____
bool streamTryGet (
______
uint<nn>_t *data,
int
    *status
);
______
bool streamTryCheck (
______
uint<nn>_t *data,
const int param = 0
);
______
bool streamBurstTryGet (
______
uint<nn>_t *data,
const int bytesize,
const int param = 1
);
______
bool streamBurstTryGet (
-----
const int bytesize,
const int param = 1
);
______
bool streamBurstTryCheck (
______
uint<nn>_t *data,
const int bytesize,
const int param = 1
);
______
bool streamBurstTryCheck (
______
const int bytesize,
const int param = 1
);
```

```
______
bool streamBurstTryIncrement (
______
uint8 t
      data,
const int bytesize,
const int param = 1
);
______
bool streamBurstTryRandom (
_____
uint8 t
      data,
const int bytesize,
const int param = 1
);
```

The class also has getNodeNumber and waitForSim methods as for the OsvvmClass, with the same functionality (see section 6.6).

#### 10.3.4 Directive Methods

Some additional directive methods are provided for statistics and flow synchronisation, as listed below:

```
int streamGetRxTransactionCount (
______
void
);
______
int streamGetTxTransactionCount (
______
void
);
______
void streamWaitForRxTransaction (
______
void
);
______
void streamWaitForTxTransaction (
______
void
);
```

# 11 Compiling Co-simulation Code

To compile the co-simulation files, both OSVVM and C++, some additional steps are required over a pure OSVVM environment. Some Tcl procedures are available to compile the C++ code which are gathered into CoSim/Scripts/MakeVproc.tcl. Ultimately, these scripts call the makefile in the CoSim directory, setting parameters as appropriate. This builds the co-simulation code into a shared object called VProc.so, and the user code is compiled to VUser.so.

For normal OSVVM test case code, this might use RunTest to analyse and then run the test code. For a co-simulation test additional arguments must be given to compile the C++ code. The MkVproc TCL function is used and has two arguments:

The first argument is a path to the directory where the user source code is located, containing all the user code for the particular test. An optional libname argument allows for libraries within CoSim/lib to be linked with the user code by tagging a -1:libname> argument. An example .pro script to compile a co-simulation test is shown below, where the script is located in AXI4/Axi4/TestCases:

```
MkVproc $::osvvm::OsvvmCoSimDirectory/tests/usercode_size simulate TbAxi4_CoSim
```

Figure 6: Compiling Co-simulation Code

In this example, the test case, TbAxi4\_Cosim, is run with arguments to build the software to run with it. The MkVproc call specified that the top level test directory is CoSim (using the global \$::osvvm::OsvvmCoSimDirectory), and that the particular test code to be compiled is in tests/usercode\_size, under the top level test directory. The output of the make process, instigated by MkVproc, will be in the directory from which the .pro script was called, which would normally be outside of the OsvvmLibraries directory. The co-simulation shared objects, VProc.so and VUser.so, along with the intermediate compiled files and static libraries are all output to the running directory. With a call to MkVproc a clean compile is done. That is so all the build files of any previous compilation are deleted first in order to ensure no clashes or pick up of old files in the event of a compilation error. A MkVprocNoClean procedure is provided to skip the cleaning step for diagnostic purposes but would not be used in normal operations.

# 12 Usage Examples

In this section will be described a set of usage cases for the co-simulation address bus features. Each has a demonstration test case which may be referred to as a reference for each of the case.

### 12.1 C++ test code for word and burst transactions

At its simplest, the co-simulation features allow address bus transaction tests to be written in C++. The user code provides a VUserMainO function, instantiates an OsvvmCosim object and then has access to all the API features described above. The diagram in Figure 2 shows this test setup.

Demo tests for both word/sub-word and burst transactions exist, with the C++ source as:

- CoSim/tests/usercode\_size/VUserMain0.cpp
- CoSim/tests/usercode\_burst/VUserMain0.cpp

Both of these programs can be run on the TbAxiMemory.vhd test bench for both Axi4Lite and Axi4. The Axi4Lite and Axi4 test case VHDL for the tests are:

- AXI4/Axi4Lite/testbench/TbAxi4\_CoSim.vhd
- AXI4/Axi4/TestCases/TbAxi4 CoSim.vhd

These VHDL test case files are common to most of the relevant tests as it is just different C++ programs that need to be run on them. A simple example of a program using the word/sub-word API calls is show below.

```
#include "OsvvmCosim.h"
extern "C" void VUserMain0() {
   bool
                          error = false;
   uint32 t
                          wdata = 0;
   OsvvmCosim
                          cosim(0);
   for (int loop = 0; loop < 4; loop++)
    {
        uint32_t addr = 0x10000000;
       uint32_t rdata;
       for (int idx = 0; idx < 4; idx++) {
            cosim.transWrite(addr, wdata + idx);
            addr += 4;
        }
       addr = 0x10000000;
        for (int idx = 0; idx < 4; idx++) {
            cosim.transRead(addr, &rdata);
            if (rdata != (wdata + idx)) {
                error = true;
                break;
            }
            addr += 4;
        }
       wdata += 0x10;
    }
   // Flag to the simulation we're finished, after 10 more iterations
   cosim.tick(10, true, error);
   // If ever got this far then sleep forever
   SLEEPFOREVER;
```

Figure 7: Example Co-simulation Test Program

### 12.2 RISC-V instruction set simulator

The first case above gives full access to the API to write new test code to drive the address bus transactions. A more useful ability might be to run a complex C++ model and have this access logic components in the simulation. The test code in CoSim/tests/iss hooks up a RISC-V instruction set simulator. This is a pre-existing open-source model available on github which models a RISC-V processor to the RV32GC+Zicsr standard, with machine level privileges. The model can be called from an external function and provides a means to register a callback function to be called whenever the processor does a load or store access. This callback function can choose to process the access or hand it back to the ISS. The model can also be connected to a gdb debugger as a remote target over TCP/IP and thus an IDE such as Eclipse. The test environment now looks like the following:

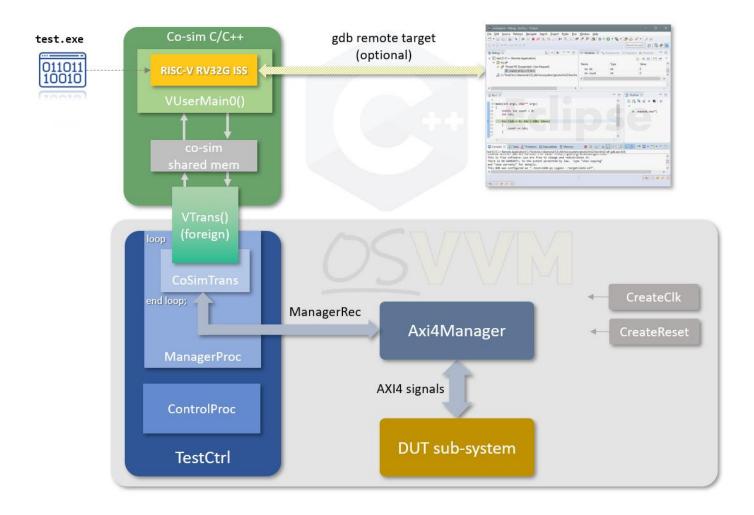


Figure 8: ISS model Co-Simulation Environment

In this test the driving of the transactions now come from within the model itself, with a pre-compiled RISC-V program (test.exe, in the test directory) that is one of RISC-V International's unit test—to test the **sb** (store byte) instruction. The VUserMain0 function is now just a means to instantiate the model, load a program (if not loaded over gdb) and run it. Binary libraries for the RISC-V ISS model are provided in CoSim/lib and the headers in CoSim/include.

A simplifed VUserMain0 program is shown below.

```
#include "OsvvmCosim.h"
#include "rv32.h"
#include "rv32 cpu gdb.h"
extern "C" void VUserMain0()
   OsvvmCosim cosim(node);
   rv32i_cfg_s cfg; // ISS config structure
   bool error = false;
   rv32* pCpu = new rv32(); // ISS object
   // Register memory access callback
   pCpu->register_ext_mem_callback(memcosim);
   // Load a program and run the ISS model
   if (!pCpu->read elf("test.exe")) {
       pCpu->run(cfg);
       error = check_exit_status(pCpu);
   }
   else
       error = true;
   // Clean up
   delete pCpu;
   // Flag to sim that the test is finished
   cosim.tick(10, true, error);
   // Let the simulation free run
   SLEEPFOREVER;
```

Figure 9: Main Program with RISC-V ISS

The registered call back function (memcosim in the example above) is called each load and store from the model and has parameters for byte address data and type. A simplified version of the callback is shown below

```
int memcosim (const uint32 t byte addr, uint32 t &data,
            const int type, const rv32i_time_t time)
 OsvvmCosim cosim(node);
 int cycle count = 5;
 uint8 t rdata8;
 uint16 t rdata16;
 uint32 t rdata32;
 switch (type)
   case MEM WR ACCESS BYTE : cosim.transWrite(byte addr, (uint8 t)data);
     break;
   case MEM WR ACCESS HWORD: cosim.transWrite(byte addr, (uint16 t)data);
     break;
   case MEM WR ACCESS WORD: cosim.transWrite(byte addr, (uint32 t)data);
     break;
   case MEM WR ACCESS INSTR: cosim.transWrite(byte addr, (uint32 t)data);
     break:
   case MEM RD ACCESS BYTE: cosim.transRead(byte addr, &rdata8); data=rdata8;
     break;
   case MEM RD ACCESS HWORD: cosim.transRead(byte addr, &rdata16); data=rdata16;
     break;
   case MEM RD ACCESS WORD: cosim.transRead(byte addr, &rdata32); data=rdata32;
     break:
   case MEM RD ACCESS INSTR: cosim.transRead(byte addr, &rdata32); data=rdata32;
     break;
   default: cycle count = RV32I EXT MEM NOT PROCESSED;
     break
 return cycle_count;
```

Figure 10: ISS Callback Function

This example serves to demonstrate the simplicity of connecting an existing C++ model environment to the OSVVM co-simulation features. The test bench uses the TbAxi4Memory VC as a target, with the program loaded and then run from that memory, but this just stands in for simulated DUT logic, either a peripheral or a sub-system, say. The callback might be enhanced to do some address decoding to only forward to OSVVM relevant accesses, and hand back those out of range for the rest of the model to handle.

The method described here relies on having a model that is callable from the VUserMain0 program. The simulators, in general, are executables and are not callable from external programs. The OSVVM cosimulation code provides the means to have a 'free running' program that can call API methods to initiate activity in the logic simulation but, ultimately, this was all initiated from the simulator process. In order to have a program, separate from the simulator, drive the transactions another method is required.

### 12.3 Connecting to External Programs

If the modelling system to be hooked to OSVVM is, itself, an executable and can't be called from the co-simulation code, then a means is needed to communicate between the external model and the OSVVM co-simulation program. In CoSim/code is some source code that defines an OsvvmCosimSkt class which acts as a TCP/IP socket server. It defines a protocol for sending word or sub-word read or write

transactions, sending back responses to a connected client. The protocol is based on gdb's remote target protocol, but the class is defined in such a way as to allow a parsing and response methods to be overwritten in a derived class to alter or extend that protocol, so long as the basic structure is a start-of-packet (SOP) byte, data bytes, end-of-packet byte, followed by a fixed number of bytes (which can be 0). The current protocol is restricted to word and sub-word transactions and has no support for time advancement or interrupts, but this is easily added by extending the class to enhance the protocol and add burst transfers, time, and interrupt support. It serves as a demonstration and, of course, a user can write their own socket code. The test environment now looks like that shown in the figure below:

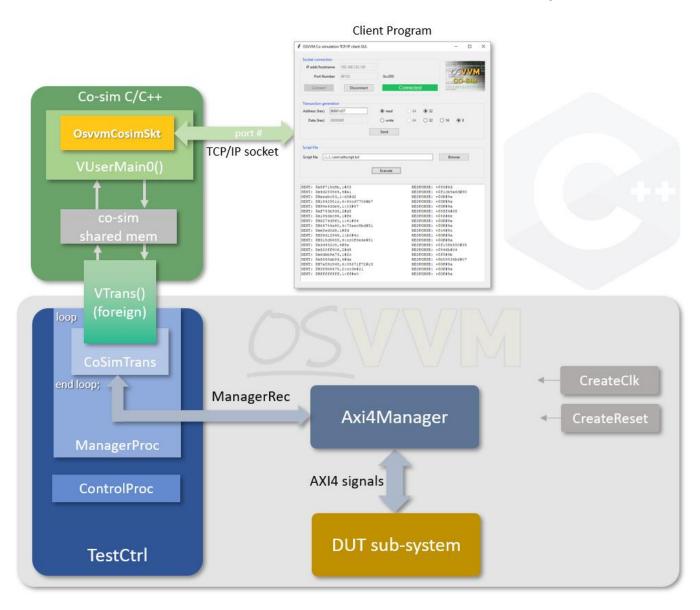


Figure 11: TCP/IP Socket Based Test Environment

The CoSim/tests/socket test code makes use of this class to open up a TCP/IP port and listen for an external connection. The VUserMain0 program is shown below:

```
extern "C" void VUserMain0()
{
   OsvvmCosim cosim(node);
   OsvvmCosimSkt skt(node);
   bool error = false;
   if (skt.ProcessPkts() != OsvvmCosimSkt::OSVVM COSIM OK)
   {
       fprintf(stderr, "***ERROR: socket exited with bad status\n");
       error = true;
   }
   else
   {
       printf("DONE\n");
   }
   // Flag to the simulation we're finished, after 10 more iterations
   cosim.tick(10, true, error);
   SLEEPFOREVER:
```

Figure 12: Use of OsvvmCosimSkt Class

Standing in for an external program, a python script is provided in CoSim/Scripts. A batch version, used in running the demo tests is called client\_batch.py. A GUI version (client\_gui.py) can be used to interact directly with the simulation. When the simulation is started the simulation will wait on connection to a port (by default 49152—which is 0xC000). The GUI can then connect to this port and send commands and receive responses. The GUI is shown as the client in Figure 12 above. As well as sending and receiving data manually, the GUI can read a script, and the ISS demo test generates a script as it runs that can be used. There is also a sktscript.txt file in the CoSim/tests/socket directory, used in the demo tests.

## 12.4 InterruptHandler VC usage

In the 2022.11 release of OSVVM an Interrupt Handler VC component was added to allow for testing of the interrupt generation features of a DUT (see [2] for details). The test environment adds to the basic environment by having an additional InterruptProc process that can generate transactions in the same manner as a MangerProc process. The InterruptHandler arbitrates between the two sources to forward transactions to the Axi4Manager VC based on the state of the DUT's interrupt request line. To use the handler with co-simulation code, the multi-node capabilities are used, where the manager process uses node 0 (the case in the previous examples, and the interrupt process uses node 1. Each node will call its respective top level functions—VUserMain0 and VUserMain1—each running as separate programs. This arrangement is shown in the diagram below.

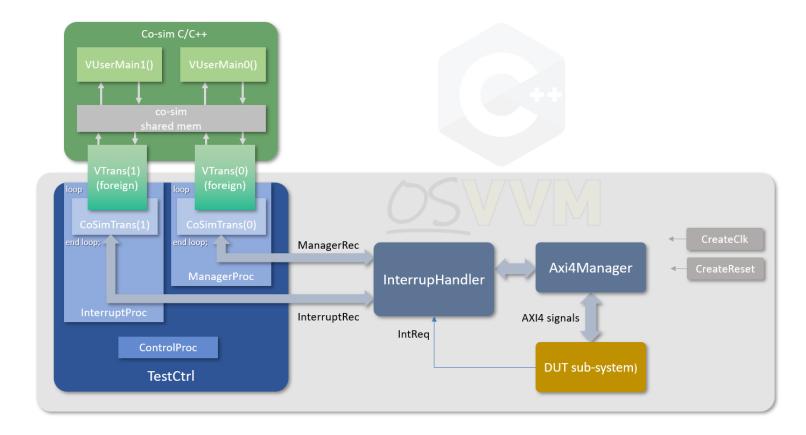


Figure 13: Co-simulation with InterruptHandler VC

A test example is given in the code in CoSim/tests/interrupt and runs on the TbAxi4Memory.vhd test bench in Common/TbInterrupt/testbench. The VHDL test case is TbAxi4\_InterruptCosim1.vhd under Common/TbInterrupt/TestCases. The two VUserMain programs running simply emulate the functionality in the TbAxi4\_Interrupt1.vhd, just using calls to the API to instigate the transactions.

## 12.5 Interrupt callback

The interrupt callback method of handling interrupts make use of a VHDL global signal, gIntReq, defined in the InterruptGlobalSignalPkg. The CoSimInterruptHandler reflects the state of the IntReq input and sets the global signal accordingly. The manager process of the test case can then access this signal to set the interrupt input of CoSimTrans as required. In the case of the TbAxi4\_InterruptCosim3.vhd test case, the input to CoSimTrans is set to 1 whenever the gIntReq changes state. The test code in CoSim/tests/interruptIss registers an interrupt callback function with the co-simulation software and this function toggles a local static variable, IntReq, each time it is called. The test uses the RISC-V ISS model, and another function is registered as a callback from this model to inspect IntReq and return the status to the model. The associated RISC-V test program has the exception vectors set up and enables interrupts before running code to write to a given address to initiate an external interrupt, which the test bench monitors for and sets the IntReq output from the TestCtrl\_e test bench component in the logic simulation. The RISC-V interrupt routine clears the interrupt (writes 0 to the same address that set the interrupt signal) and sets a RISC-V register to say it has been called, which the main test program then checks for, flagging a pass/fail condition.

Co-sim C/C++ state update available to main program

VUserMainO()

Co-sim shared mem sha

This setup is shown in the diagram below:

Figure 14: Interrupt Callback Environment

This test, then, demonstrates the connection between an interrupt signal in logic simulation causing an external interrupt initiating a jump to exception code, within a software processor system model.

## 12.6 Using GHDL callable Simulation

The default behaviour of the co-simulation code is to call one or more user programs, running in threads, with entry points at VUserMainx (one for each active node). The user test code is called from there and can be any code complexity, modelling whole systems if required. The assumption here is that the user's code, modelling their system, is callable from external code. This is true of the example with the RISC-V ISS, which is instantiated and run from a VUserMain program. This may not always be possible for pre-existing modelling systems, where the tool being used is an executable in its own right, and custom models are added to this environment and are called from there.

The supported GHDL simulator has a documented feature whereby the simulator can be called from an external program via the <a href="mailto:ghdl main">ghdl main</a> function. As documented, this is meant to be called from user code compiled as a shared object and loaded as such as for foreign language routines. Again, since a user's model may not be in the form of a shared object, a different model is needed to make calling the simulation possible. GHDL's two steps for compiling code is to first analyse all the VHDL files and then,

secondly, elaborate the top level. The GHDL tool must be compiled for an LLVM backend (though gcc might work, but not tried) and not mcode (see <a href="here">here</a>). The first step results in a set of objects for each component in a respective library. For OSVVM, all these files are in subdirectories of VHDL\_LIBS. The elaboration step results in an object in the compile directory with the form e~<name of top level>.o, for example e~tbaxi\_cosim.o. It also results in an executable which is the simulation that can be run, but this is not yet suitable. These objects in VHDL\_LIBS can all be gathered into a single static library, along with the relevant objects from GHDL for the normal VHDL libraries (e.g. <root to GHDL libs>/std/08/\*.o) to be used in compiling with the main code.

The OSVVM co-simulation code itself must be compiled slightly differently so that it does not attempt to start the threads containing the <code>VUserMainx</code> code. This is done by compiling with <code>-DDISABLE\_VUSERMAIN\_THREAD</code> in the <code>gcc/g++</code> flags. This will still create the usual <code>VProc.so</code> and <code>VUser.so</code> shared object files. It also prevents the user code from setting the <code>done</code> argument to <code>true</code> in any call to the <code>tick</code> method. A call to this method with <code>done</code> set to <code>true</code> will block forever as it will terminate the simulation. The co-simulation software expects a response for each API call from the simulator that is no longer forthcoming, as the simulation has terminated. When the code is run as a thread from the simulator, this is not an issue, as the simulator terminating is the end of the simulation and the program exits. When called from an external program, it must not hang on this call, and external means are required to call the <code>tick</code> method with <code>done</code> set to <code>true</code> (in code which hasn't been compiled with <code>-DDISABLE VUSERMAIN THREAD</code>).

The two co-simulation shared objects, along with the new static library of component object files and the elaboration object can now be compiled with the user's main program to replace the original executable with one that now has user code that can call the GHDL simulator and use the co-simulation API directly. However, the simulator must be started first via the call to <code>ghdl\_main</code>. This is a blocking call and won't exit until the simulation has finished. Therefore this must be run in a separate thread. Once started, and the simulation has initialised the co-simulation code, the user code can then make calls to the co-simulation API, via an <code>OSvvmCosimClass</code>. The diagram below summarises this environment

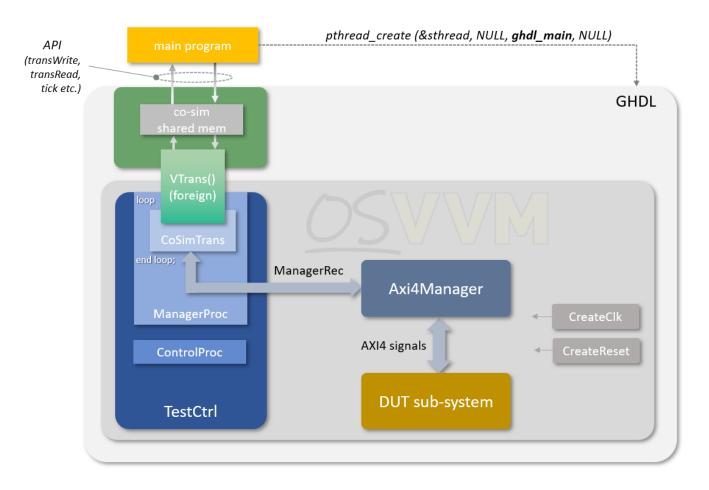


Figure 15: Callable GHDL test environment

The user code must take care of initialising the simulator, running <code>ghdl\_main</code> in a thread, and waiting for it to be ready for API accesses. When the simulation finishes, the simulation must be stopped with a call to the <code>tick</code> method with the <code>done</code> argument set to <code>true</code>. As mentioned above, such a call to this method will block forever, and so this needs to be called fro a separate thread. Whan called, the <code>ghdl\_main</code> function will return and the thread it's running in will terminate. The main program can wait on this thread in order to ensure the simulation has, in fact, finished, using a <code>pthread\_join</code> call. The abbreviated code fragment below summarises the use of these features.

```
#include <pthread.h>
#include "OsvvmCosim.h"
extern "C" int ghdl main (int argc, char **argv);
extern "C" int VUserMain0 ();
static int node = 0;
static int large;
static char** largv;
// -----
int run sim (int dummy) {
   int status = ghdl main(largc, largv);
   return status;
}
// -----
int stop sim thread (int dummy) {
 OsvvmCosim cosim(node);
 // This call will hang because the simulation will not return
 // a respoinse since it was stopped.
 cosim.tick(1, true, false);
 return 0;
// -----
void stop_sim(void) {
   pthread t stop thread;
   pthread create(&stop thread, NULL, (pThreadFunc t)stop sim thread,
                (void *)((long long)node));
// ======== M A I N =========
int main (int argc, char **argv) {
   OsvvmCosim cosim(node);
   pthread t sim thread;
   // Export arguments for use by thread code
   largc = argc; largv = argv;
   // Create a thread for calling GHDL simulator
   pthread create(&sim thread, NULL, (pThreadFunc t)run sim, (void *)((long long)node));
   \ensuremath{//} Allow the simulation's co-simulation code to initialise and be ready
   // for the first API call.
   cosim.waitForSim();
   // Do some tests, calling the VUserMainO code directly.
   VUserMain0();
   // Post a done status to the simulator.
   stop sim();
   // Wait for the simulator to exit and the thread to complete.
   pthread_join(sim_thread, NULL);
   return 0;
```

The steps necessary to set up the GHDL callable simulation are summarised below.

- Analyse VHDL files: ghdl -a <list of files>
- Elaborate top level: ghdl -e <top level>
- Compile co-sim code: make -C <path to OsvvmLibraries/CoSim> [options]
  - o The USRFLAGS must be set to include -DDISABLE VUSERMAIN THREAD
- Compile user code: gcc -c main.c -o main.o
- Gather analysed file objects, user code, and GHDL library objects into a library:

```
o ar crs libtb.a main.o
  `find VHDL_LIBS -name *.o`
  `find <GHDL rootdir>/std/08; -name *.o`
  `find <GHDL rootdir>/ieee/08; -name *.o`
```

- Generate a new executable:

-L<root>/lib/gcc/x86\_64-mingw32/12.2.0/adalib

The above steps, not including the analysis of the VHDL files are encapsulated in the procedure MkVprocGhdlMain, which has the same arguments as MkVproc described in section Error! Reference s

ource not found. This makes use of a make file, makefile.ghdl, in OsvvmLibraries/CoSim. This make file is customisable with user overridable variables, as documented in the file's header, so can be used for user code outside of the OsvvmLibraries directory structure.

### 13 Summary

This document defines the features of the co-simulation capabilities of OSVVM. A sub-set of the Address Bus Model Independent Transaction and Stream Model Independent interfaces are exposed within a C++ environment to allow reads and writes of words or bursts, driving an AddressBusRecType or StreamRecType signal to initiate transactions on one of the pre-existing Address Bus or streaming virtual components. A single point of entry from VHDL is given with the CoSimTrans or CoSimStream procedure (for a given node), abstracting away the details of the programming interfaces of the different simulators, and keeping the OSVVM environment intact. Support for interrupts using the InterruptHandler VC component is also given for two usage models.

With the domain crossed from VHDL to C++ the possibilities for modelling systems that are partially simulated in C++ and partially simulated in HDL in a logic simulator is demonstrated. Two methods were described (with examples) with a model's code called directly from user co-simulation code, or a TCP/IP socket link is established to drive the simulator from an external program.

## 13.1 Summary of OsvvmCosim Transaction Methods

The table below summarises the transaction methods of the OsvvmCosim class. All of these methods are inherited (and some overloaded) by the OsvvmCosimInt interrupt class.

method	function
transWrite	Blocking write byte, hword, word or dword
transWriteAsync	Non-blocking write byte, hword, word or dword
transWriteAddressAsync	Non-blocking write address (split transaction)
transWriteDataAsync	Non-blocking write byte, hword, word or byte (split transaction)
transWriteAndRead	Blocking write and read byte, hword, word or dword
transWriteAndReadAsync	Non-blocking write and read byte, hword, word or dword
transBurstWrite	Blocking burst write
transBurstWriteAsync	Non-blocking burst write
transBurstWriteIncrement	Blocking burst write of incrmenting pattern
transBurstWriteIncrementAsync	Non-blocking burst write of incrmenting pattern
transBurstWriteRandom	Blocking burst write of incrmenting pattern
transBurstWriteRandomAsync	Non-blocking burst write of incrmenting pattern
transRead	Blocking read byte, hword, word or dword
transReadAddressAsync	Non-blocking read address (split transaction)
transReadData	Blocking data only read of byte, hword, word or dword (split transaction)
transTryReadData	Non-blocking data only read of byte, hword, word or dword (split transaction)
transReadCheck	Blocking read and check of byte, hword, word or dword
transReadDataCheck	Blocking data only read and check of byte, hword, word or dword (split transaction)
transTryReadDataCheck	Non-blocking data only read and check of byte, hword, word or dword (split transaction)
transBurstRead	Blocking burst read
transBurstReadCheckData	Blocking burst read and check of data (split transaction)

transBurstReadCheckIncrement	Blocking burst read and check data incrementing pattern(split transaction)
transBurstReadCheckRandom	Blocking burst read and check random pattern (split transaction)
transBurstCheckData	Blocking check of received data
transBurstCheckIncrement	Blocking check of received incrementing pattern
transBurstCheckRandom	Blocking check of received random pattern
transPushData	Blocking push of data to fifo for transmission or check
transPushIncrement	Blocking push of incrementing pattern to fifo for transmission or check
transPushRandom	Blocking push of random pattern to fifo for traansmission or check
transPopData	Blocking pop of data from receive fifo
transGetTransactionCount	Get the number of interactions over the interface
transGetWriteTransactionCount	Get the number of write transactions over the interface
transGetReadTransactionCount	Get the number of read transactions over the interface
transWaitForTransaction	Wait for a transaction
transWaitForWriteTransaction	Wait for a write transaction
transWaitForReadTransaction	Wait for a read transaction

## 13.2 Summary of OsvvmCosimResp Response Methods

The table below summarises the response methods of the  ${\tt OsvvmCosimResp}$  class.

method	function
respGetWrite	Get write transaction address and data
respTryGetWrite	Get write transaction address and data or return false if none
respGetWriteAddress	Get write address
respTryGetWriteAddress	Get write address or return false if none

respGetWriteData	Get write data
respTryGetWriteData	Get write data or return false if none
respSendRead	Send read response for received transaction
respTrySendRead	Send read response for received transaction or return false if none
respGetReadAddress	Get read address
respTryGetReadAddress	Get read address or return false if none
respSendReadData	Send response data for received transaction
respSedReadDataAsync	Send response data for received transaction asynchronously
respWaitForTransaction	Wait for a transaction
respWaitForWriteTransaction	Wait for a write transaction
respWaitForReadTransaction	Wait for a read transaction
respGetTransactionCount	Get the number of interactions over the interface
respGetWriteTransactionCount	Get the number of write transactions over the interface
respGetReadTransactionCount	Get the number of read transactions over the interface

# 13.3 Summary of OsvvmCosimStream Transaction Methods

The table below summarises the transaction methods of the OsvvmCosimStream class.

method	function
streamSend	Blocking send of byte, hword, word or dword
streamSendAsync	Non-blocking send of byte, hword, word or dword
streamBurstSend	Blocking send of burst data
streamBurstSendAsync	Non-blocking send of burst data
streamBurstSendIncrement	Blocking send of burst incrementing pattern
steramBurstSendIncrementAsync	Non-blocking send of burst incrementing pattern

streamBurstSendRandom	Blocking send of burst random pattern
streamBurstSendRandomAsync	Non-blocking send of burst random pattern
streamGet	Blocking get of byte, hword, word or dword
streamGetBurstGet	Blocking get of burst data
streamCheck	Blocking get and check of byte, hword, word or dword
streamBurstCheck	Blocking get and check of burst data
streamBurstCheckIncrement	Blocking get and check of burst incrementing pattern
streamBurstCheckRandom	Blocking get and check of burst random pattern
streamTryGet	Non-blocking get of byte, hword, word or dword
streamTryCheck	Non-blocking get and check of byte, hword, word or dword
streamBurstTryGet	Non-blocking get of burst data
streamBurstTryCheck	Non-blocking get and check of burst data
streamBurstTryIncrement	Non-blocking get and check of burst incrementing pattern
streamBurstTryRandom	Non-blocking get and check of burst random pattern
streamBurstPushData	Push data to fifo for transmission
streamBurstPushCheckData	Push data to fifo for checking against received data
streamBurstPushIncrement	Push incrementing pattern to fifo for transmission
streamBurstPushCheckIncrement	Push incrementing pattern to fifo for checking against received data
streamBurstPushRandom	Push random pattern to fifo for transmission
streamBurstPushCheckRandom	Push random pattern to fifo for checking against received data
streamGetRxTransactionCount	Get number of completed transactions for transmit interface
streamGetTxTransactionCount	Get number of completed transactions for receive interface
streamWaitForRxTransaction	Wait for a transaction completion on transmit interface
streamWaitForTxTransaction	Wait for a transaction completion on receive interface

### 14 About the OSVVM

The OSVVM utility and verification component libraries were developed and are maintained by Jim Lewis of SynthWorks VHDL Training and the co-simulation libraries/features were developed and are maintained by Simon Southwell. These libraries evolved from methodology and packages developed for SynthWorks' VHDL Testbenches and verification class.

Please support OSVVM by purchasing your VHDL training from SynthWorks.

### 15 About the Authors and Contributors

## 15.1 About the Co-simulation Contributor – Simon Southwell

Simon Southwell has thirty plus years of embedded software, ASIC and FPGA design experience in fields that include high performance computing, wireless, cellular modem (LTE) and processor system modelling. Mr. Southwell has developed and successfully deployed co-simulation techniques at a variety of companies using his own open-source IP.

Mr. Southwell is currently developing open source IP in areas such as RISC-V and co-simulation, is actively mentoring undergraduate and new graduate engineers in digital systems design and is also collaborating on the development of OSVVM.

If you find bugs the co-simulation packages or would like to request enhancements, Mr. Southwell can be reached at simon.southwell@gmail.com.

### 15.2 About the Author - Jim Lewis

Jim Lewis, the founder of SynthWorks, has thirty plus years of design, teaching, and problem solving experience. In addition to working as a Principal Trainer for SynthWorks, Mr. Lewis has done ASIC and FPGA design, custom model development, and consulting.

Mr. Lewis is chair of the IEEE 1076 VHDL Working Group (VASG) and is the primary developer of the Open Source VHDL Verification Methodology (OSVVM.org) packages. Neither of these activities generate revenue. Please support our volunteer efforts by buying your VHDL training from SynthWorks.

If you find bugs these packages or would like to request enhancements, you can reach me at jim@synthworks.com.

#### 16 References

- [1] Address Bus Model Independent Transactions User Guide
- [2] Interrupt Handler User Guide
- [3] Stream Model Independent Transaction User Guide