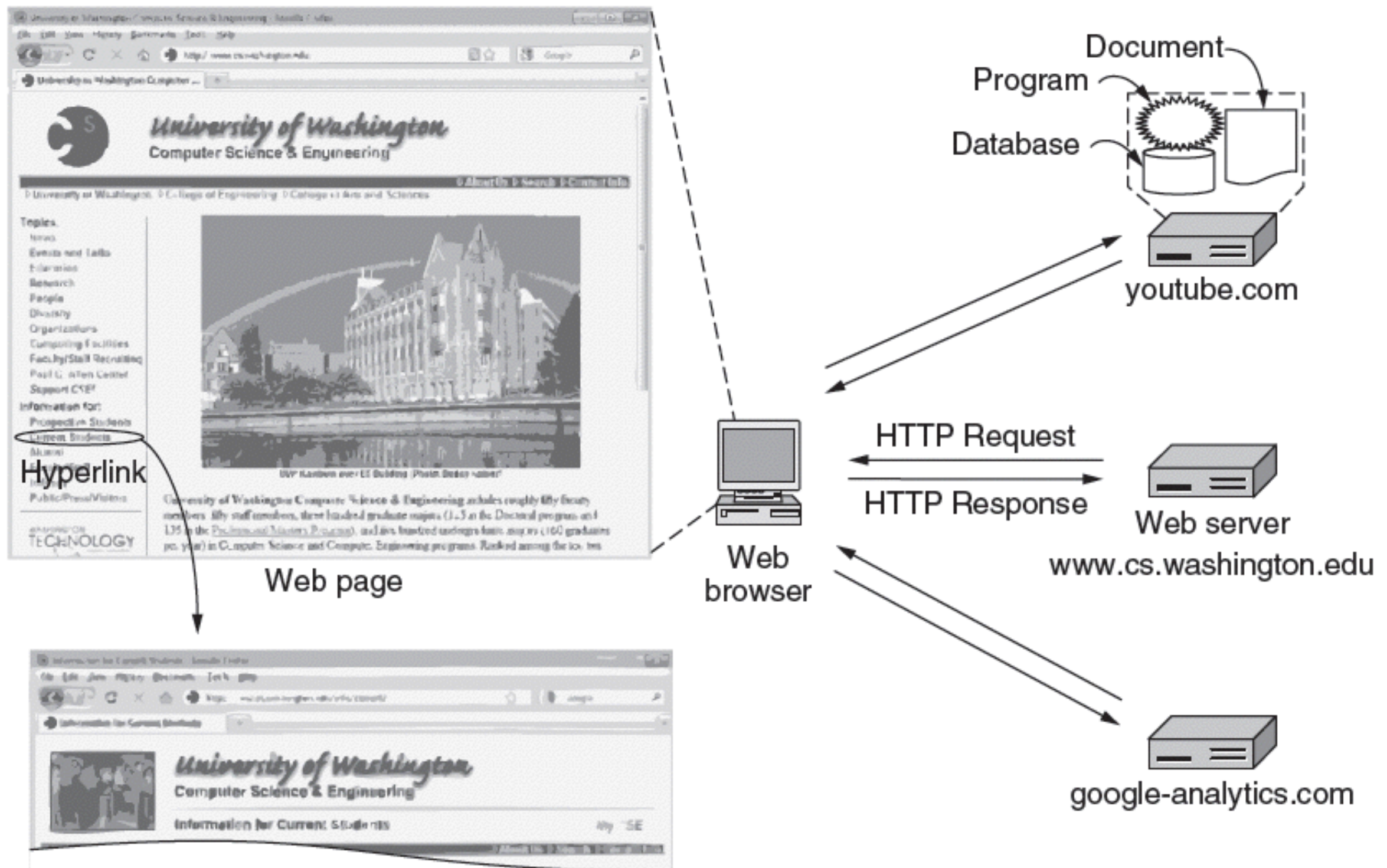


# Couche application

HTTP

*Computer Networks. Tanenbaum*  
*Computer Networking. Kurose&Ross*

# Http: Principles



# Web et HTTP

- ❖ Une *page web* contient des *objets*
- ❖ Objet : fichier HTML, images JPEG, applet, fichiers audio,...
- ❖ Une page web page consiste en un fichier de base *HTML* contenant des objets référencés
- ❖ Chaque objet est adressable par une *URL* (*Uniform Resource Locator*)

`www.someschool.edu/someDept/pic.gif`

host name

path name

# URL:

- Example: <http://www.phdcomics.com/comics.php>

Protocol

Server

Page on server

Our  
focus →

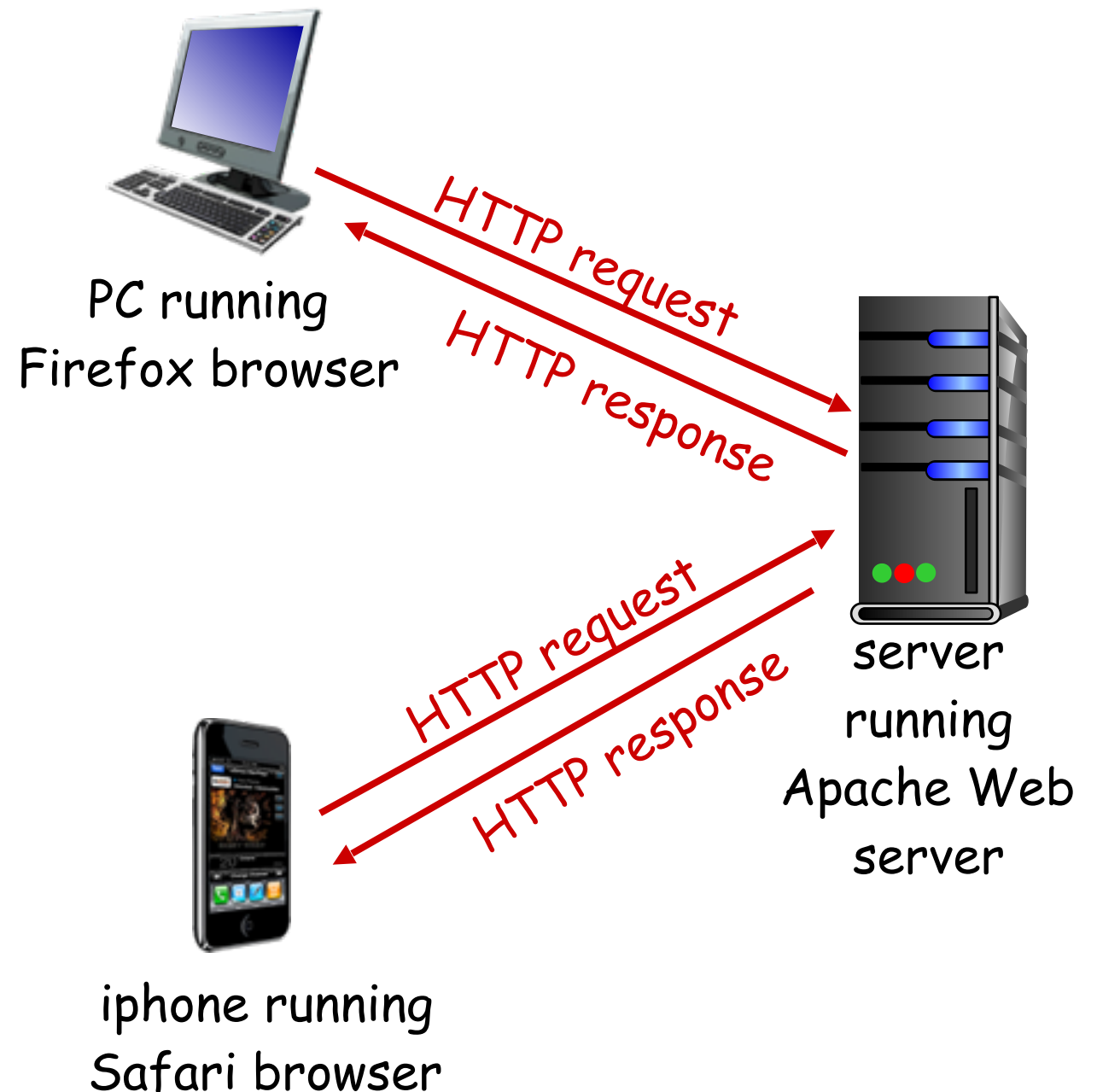
Name	Used for	Example
http	Hypertext (HTML)	<a href="http://www.ee.uwa.edu/~rob/">http://www.ee.uwa.edu/~rob/</a>
https	Hypertext with security	<a href="https://www.bank.com/accounts/">https://www.bank.com/accounts/</a>
ftp	FTP	<a href="ftp://ftp.cs.vu.nl/pub/minix/README">ftp://ftp.cs.vu.nl/pub/minix/README</a>
file	Local file	<a href="file:///usr/suzanne/prog.c">file:///usr/suzanne/prog.c</a>
mailto	Sending email	<a href="mailto:JohnUser@acm.org">mailto:JohnUser@acm.org</a>
rtsp	Streaming media	<a href="rtsp://youtube.com/montypython.mpg">rtsp://youtube.com/montypython.mpg</a>
sip	Multimedia calls	<a href="sip:eve@adversary.com">sip:eve@adversary.com</a>
about	Browser information	<a href="about:plugins">about:plugins</a>

Common URL protocols

# HTTP overview

## HTTP: hypertext transfer protocol

- ❖ Web's application layer protocol
- ❖ client/server model
  - **client**: browser that requests, receives, (using HTTP protocol) and "displays" Web objects
  - **server**: Web server sends (using HTTP protocol) objects in response to requests



# HTTP overview

---

## *uses TCP:*

- ❖ client initiates TCP connection (creates socket) to server, port 80
- ❖ server accepts TCP connection from client
- ❖ HTTP messages (application-layer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- ❖ TCP connection closed

## *HTTP is “stateless”*

- ❖ server maintains no information about past client requests

### protocols that maintain “state” are complex!

- v past history (state) must be maintained
- v if server/client crashes, their views of “state” may be inconsistent, must be reconciled



# Overview

Steps a client (browser) takes to follow a hyperlink:

- Determine the protocol (HTTP)
- Ask DNS for the IP address of server
- Make a TCP connection to server
- Send request for the page; server sends it back
- Fetch other URLs as needed to display the page
- Close idle TCP connections

Steps a server takes to serve pages:

- Accept a TCP connection from client
- Get page request and map it to a resource (e.g., file name)
- Get the resource (e.g., file from disk)
- Send contents of the resource to the client.
- Release idle TCP connections

# HTTP connections

## *non-persistent HTTP*

- ❖ at most one object sent over TCP connection
  - connection then closed
- ❖ downloading multiple objects required multiple connections

## *persistent HTTP*

- ❖ multiple objects can be sent over single TCP connection between client, server



# Non-persistent HTTP

suppose user enters URL:

`www.someSchool.edu/someDepartment/home.index`

(contains text,  
references to 10  
jpeg images)

1a. HTTP client initiates TCP connection to HTTP server (process) at `www.someSchool.edu` on port 80

1b. HTTP server at host `www.someSchool.edu` waiting for TCP connection at port 80. "accepts" connection, notifying client

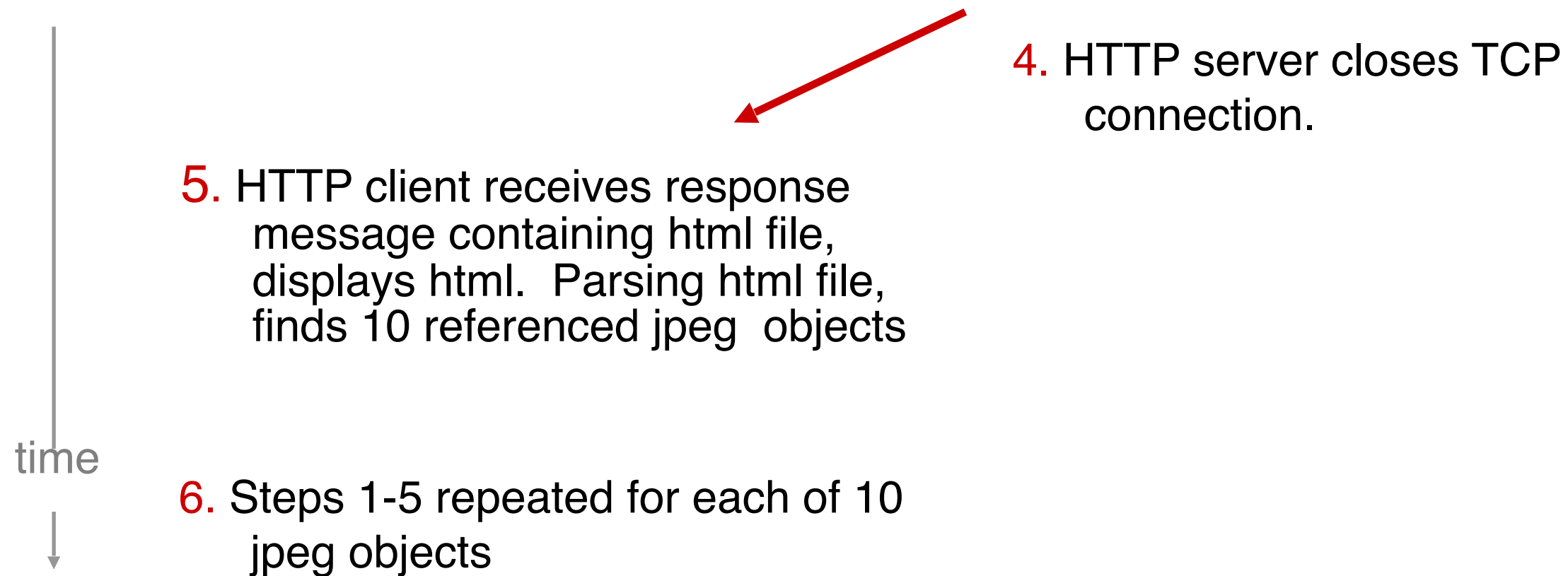
2. HTTP client sends HTTP *request message* (containing URL) into TCP connection socket. Message indicates that client wants object `someDepartment/home.index`

3. HTTP server receives request message, forms *response message* containing requested object, and sends message into its socket

time

Carole Delporte

# Non-persistent HTTP (cont.)

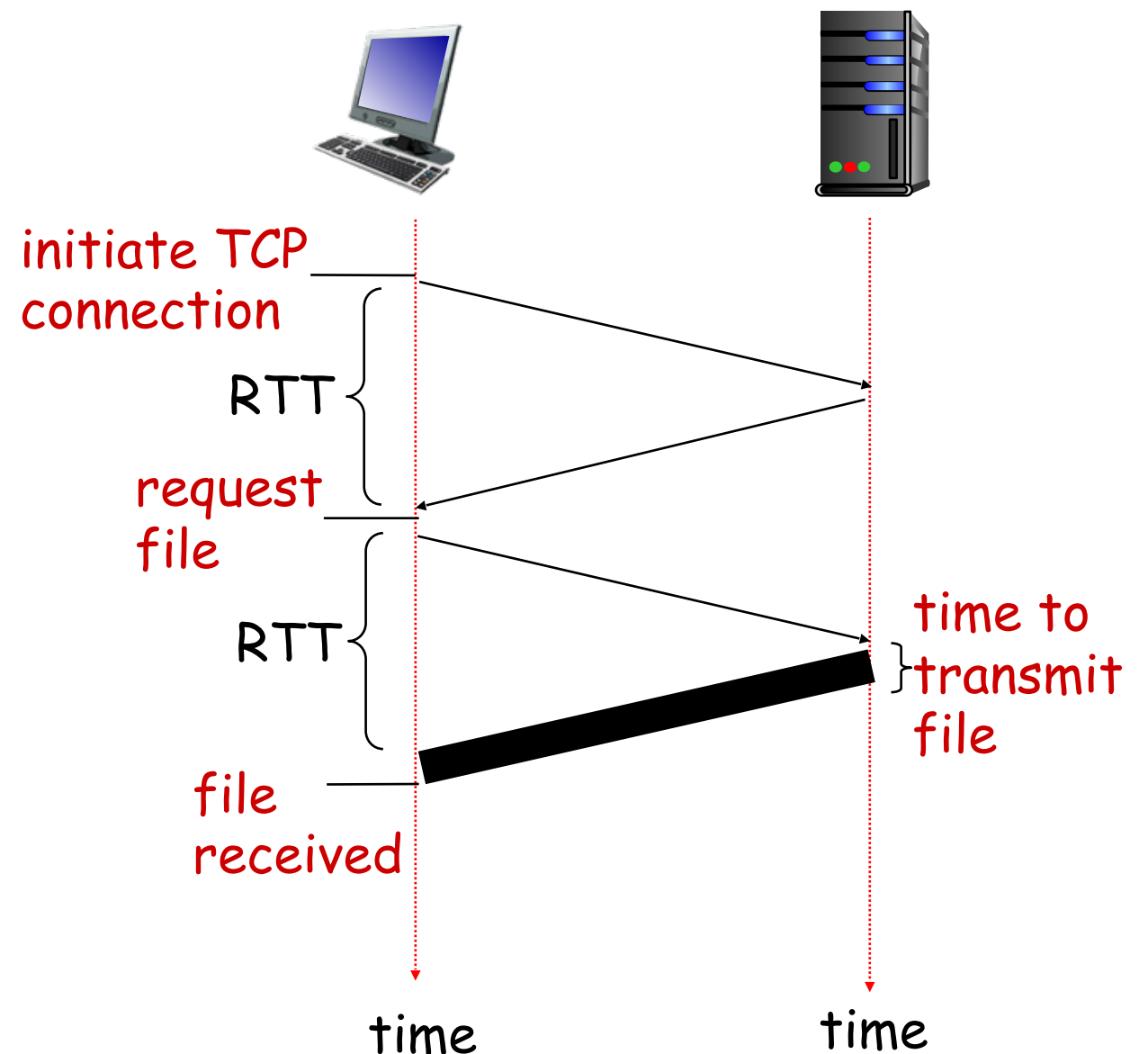


# Non-persistent HTTP: response time

**RTT (definition):** time for a small packet to travel from client to server and back

**HTTP response time:**

- ❖ one RTT to initiate TCP connection
- ❖ one RTT for HTTP request and first few bytes of HTTP response to return
- ❖ file transmission time
- ❖ non-persistent HTTP response time =  $2\text{RTT} + \text{file transmission time}$



# Persistent HTTP

## *non-persistent HTTP issues:*

- ❖ requires 2 RTTs per object
- ❖ OS overhead for *each* TCP connection
- ❖ browsers often open parallel TCP connections to fetch referenced objects

## *persistent HTTP:*

- ❖ server leaves connection open after sending response
- ❖ subsequent HTTP messages between same client/server sent over open connection
- ❖ client sends requests as soon as it encounters a referenced object
- ❖ as little as one RTT for all the referenced objects

# HTTP request message

- ❖ two types of HTTP messages: *request, response*
- ❖ **HTTP request message:**
  - ASCII (human-readable format)

request line  
(GET, POST,  
HEAD commands)

header  
lines

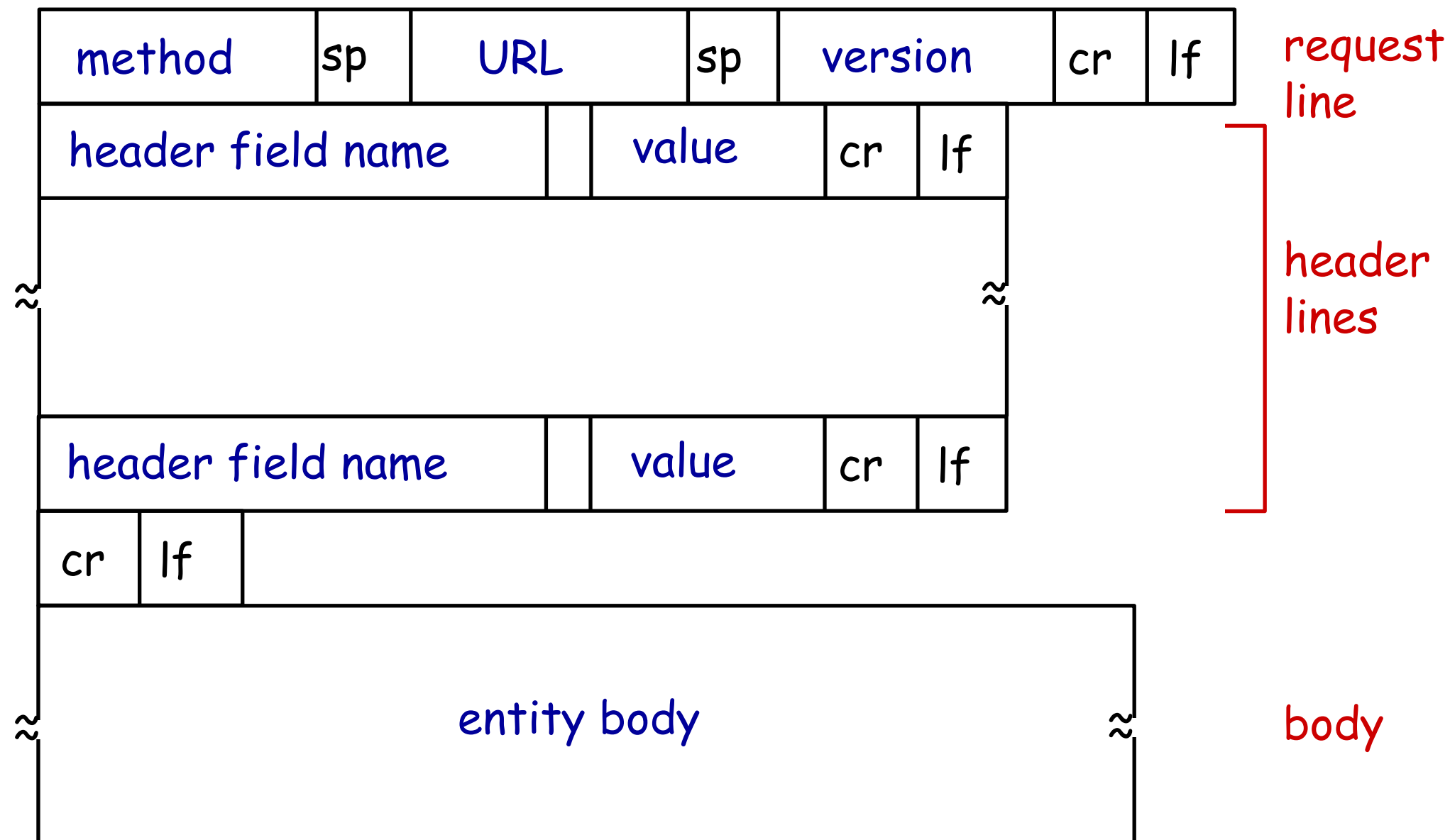
carriage return,  
line feed at start  
of line indicates  
end of header lines

```
GET /index.html HTTP/1.1\r\n
Host: www-net.cs.umass.edu\r\n
User-Agent: Firefox/3.6.10\r\n
Accept: text/html,application/xhtml+xml\r\n
Accept-Language: en-us,en;q=0.5\r\n
Accept-Encoding: gzip,deflate\r\n
Accept-Charset: ISO-8859-1,utf-8;q=0.7\r\n
Keep-Alive: 115\r\n
Connection: keep-alive\r\n
\r\n
```

carriage return character  
line-feed character

Carole Delporte

# HTTP request message: general format





# HTTP

## Headers:

Function	Example Headers
Browser capabilities (client → server)	User-Agent, Accept, Accept-Charset, Accept-Encoding, Accept-Language
Caching related (mixed directions)	If-Modified-Since, If-None-Match, Date, Last-Modified, Expires, Cache-Control, ETag
Browser context (client → server)	Cookie, Referer, Authorization, Host
Content delivery (server → client)	Content-Encoding, Content-Length, Content-Type, Content-Language, Content-Range, Set-Cookie

# Uploading form input

## POST method:

- ❖ web page often includes form input
- ❖ input is uploaded to server in entity body

## URL method:

- ❖ uses GET method
- ❖ input is uploaded in URL field of request line:

`www.somesite.com/animalsearch?monkeys&banana`

# Method types

## HTTP/1.0:

- ❖ GET
- ❖ POST
- ❖ HEAD
  - asks server to leave requested object out of response

## HTTP/1.1:

- ❖ GET, POST, HEAD
- ❖ PUT
  - uploads file in entity body to path specified in URL field
- ❖ DELETE
  - deletes file specified in the URL field

# HTTP

## Request methods.

Fetch a page →

Used to send  
input data to a  
server program →

Method	Description
GET	Read a Web page
HEAD	Read a Web page's header
POST	Append to a Web page
PUT	Store a Web page
DELETE	Remove the Web page
TRACE	Echo the incoming request
CONNECT	Connect through a proxy
OPTIONS	Query options for a page

# HTTP response message

status line  
(protocol  
status code  
status phrase)

header  
lines

data, e.g.,  
requested  
HTML file

```
HTTP/1.1 200 OK\r\n
Date: Sun, 26 Sep 2010 20:09:20 GMT\r\n
Server: Apache/2.0.52 (CentOS)\r\n
Last-Modified: Tue, 30 Oct 2007 17:00:02 GMT\r\n
ETag: "17dc6-a5c-bf716880"\r\n
Accept-Ranges: bytes\r\n
Content-Length: 2652\r\n
Keep-Alive: timeout=10, max=100\r\n
Connection: Keep-Alive\r\n
Content-Type: text/html; charset=ISO-8859-1\r\n
\r\n
data data data data data ...
```

# HTTP response status codes

v status code appears in 1st line in server-to-client response message.

v some sample codes:

**200 OK**

- request succeeded, requested object later in this msg

**301 Moved Permanently**

- requested object moved, new location specified later in this msg  
(Location:)

**400 Bad Request**

- request msg not understood by server

**404 Not Found**

- requested document not found on this server

**505 HTTP Version Not Supported**



# HTTP

Response codes tell the client how the request fared:

Code	Meaning	Examples
1xx	Information	100 = server agrees to handle client's request
2xx	Success	200 = request succeeded; 204 = no content present
3xx	Redirection	301 = page moved; 304 = cached page still valid
4xx	Client error	403 = forbidden page; 404 = page not found
5xx	Server error	500 = internal server error; 503 = try again later

# Trying out HTTP (client side) for yourself

## 1. Telnet to your favorite Web server:

```
telnet www.irif.fr 80
```

opens TCP connection to port 80  
(default HTTP server port) at cis.poly.edu.  
anything typed in sent  
to port 80 at cis.poly.edu

## 2. type in a GET HTTP request:

```
GET /bla HTTP/1.1
```

```
Host: www.irif.fr
```

by typing this in (hit carriage  
return twice), you send  
this minimal (but complete)  
GET request to HTTP server

## 3. look at response message sent by HTTP server!

```
$ telnet www.irif.fr 80
```

```
Trying 81.194.27.176...
```

```
Connected to www.irif.fr.
```

```
Escape character is '^['.
```

```
GET /blaa HTTP/1.1
```

```
Host: www.irif.fr
```

```
HTTP/1.1 301 Moved Permanently
```

```
Date: Mon, 10 Oct 2015 10:32:29 GMT
```

```
Server: Apache/2.4.23 ( FreeBSD) PHP/5.6.24 OpenSSL/0.9.8sz-freebsd
```

```
Location: https://www.irif.fr/bla
```

```
Content-Length: 231
```

```
Content-Type: text/html; charset=iso-8859-1
```

```
<!DOCTYPE html PUBLIC « -//IETF//DTD HTML 2.0 EN » >
```

```
<html><head>.....
```

# User-server state: cookies

many Web sites use cookies

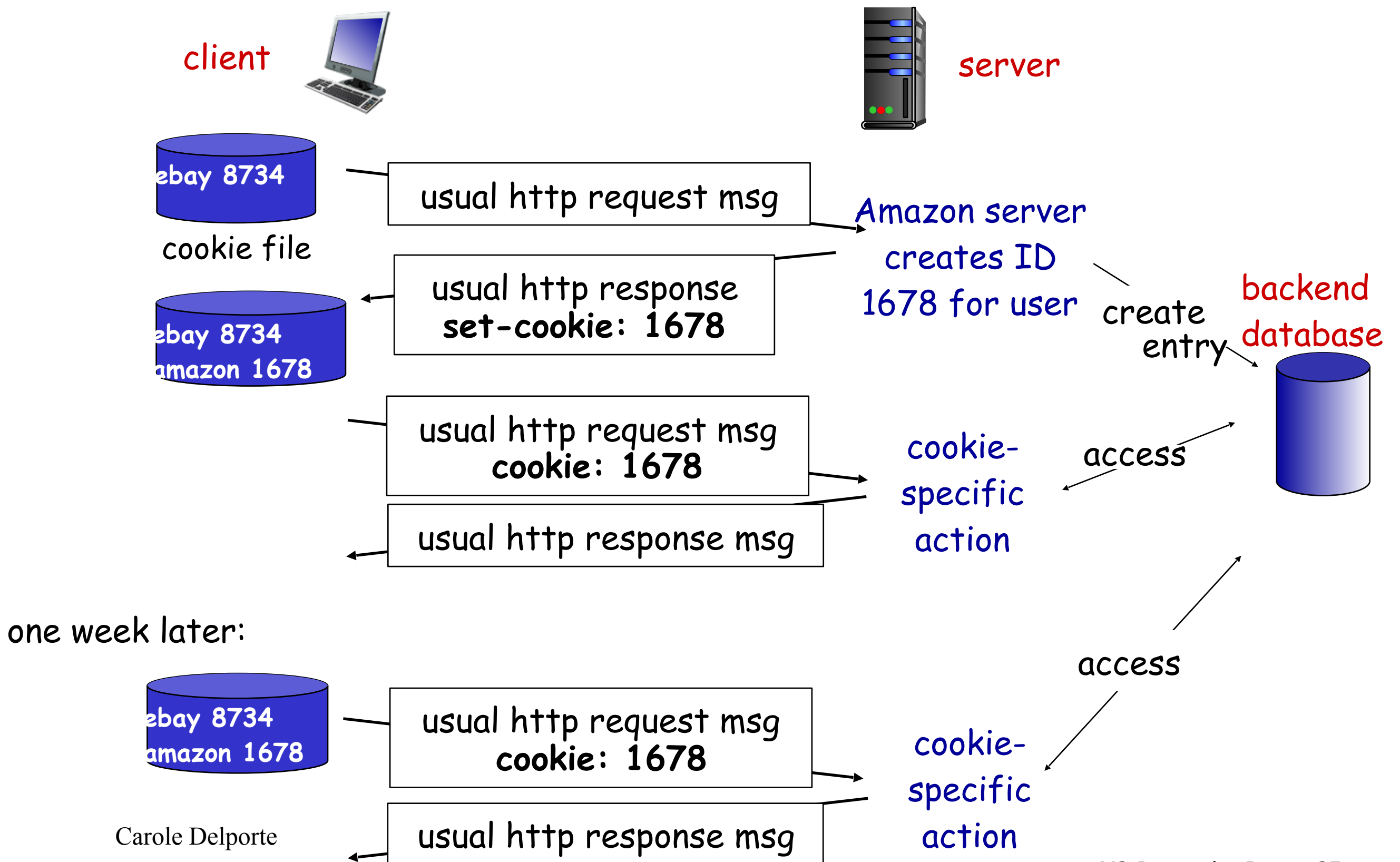
## *four components:*

- 1) cookie header line of HTTP *response* message
- 2) cookie header line in next HTTP *request* message
- 3) cookie file kept on user's host, managed by user's browser
- 4) back-end database at Web site

## *example:*

- ❖ Susan always access Internet from PC
- ❖ visits specific e-commerce site for first time
- ❖ when initial HTTP requests arrives at site, site creates:
  - unique ID
  - entry in backend database for ID

# Cookies: keeping “state” (cont.)



# Cookies (continued)

*what cookies can be used for:*

- ❖ authorization
- ❖ shopping carts
- ❖ recommendations
- ❖ user session state (Web e-mail)

*how to keep “state”:*

- v protocol endpoints: maintain state at sender/receiver over multiple transactions
- v cookies: http messages carry state

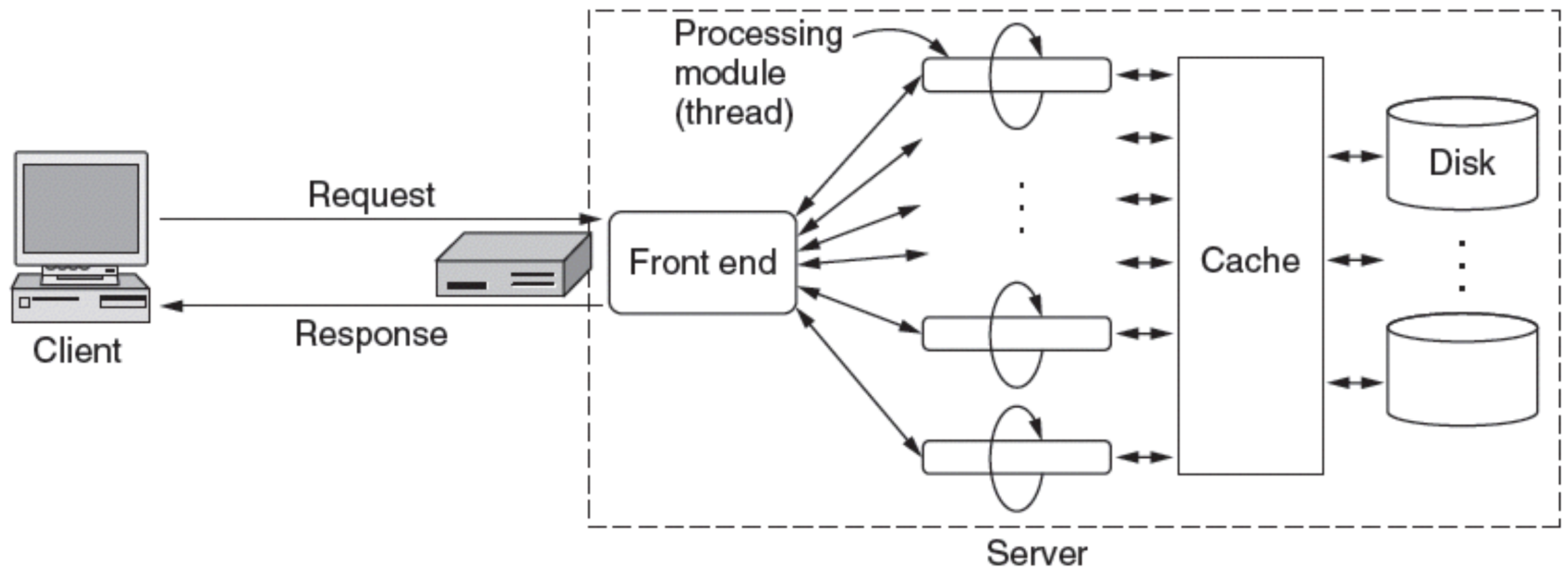
— aside —  
*cookies and privacy.*

- v cookies permit sites to learn a lot about you
- v you may supply name and e-mail to sites



# Caching

- To scale performance, Web servers can use:
- Caching, multiple threads, and a front end



# Caching...

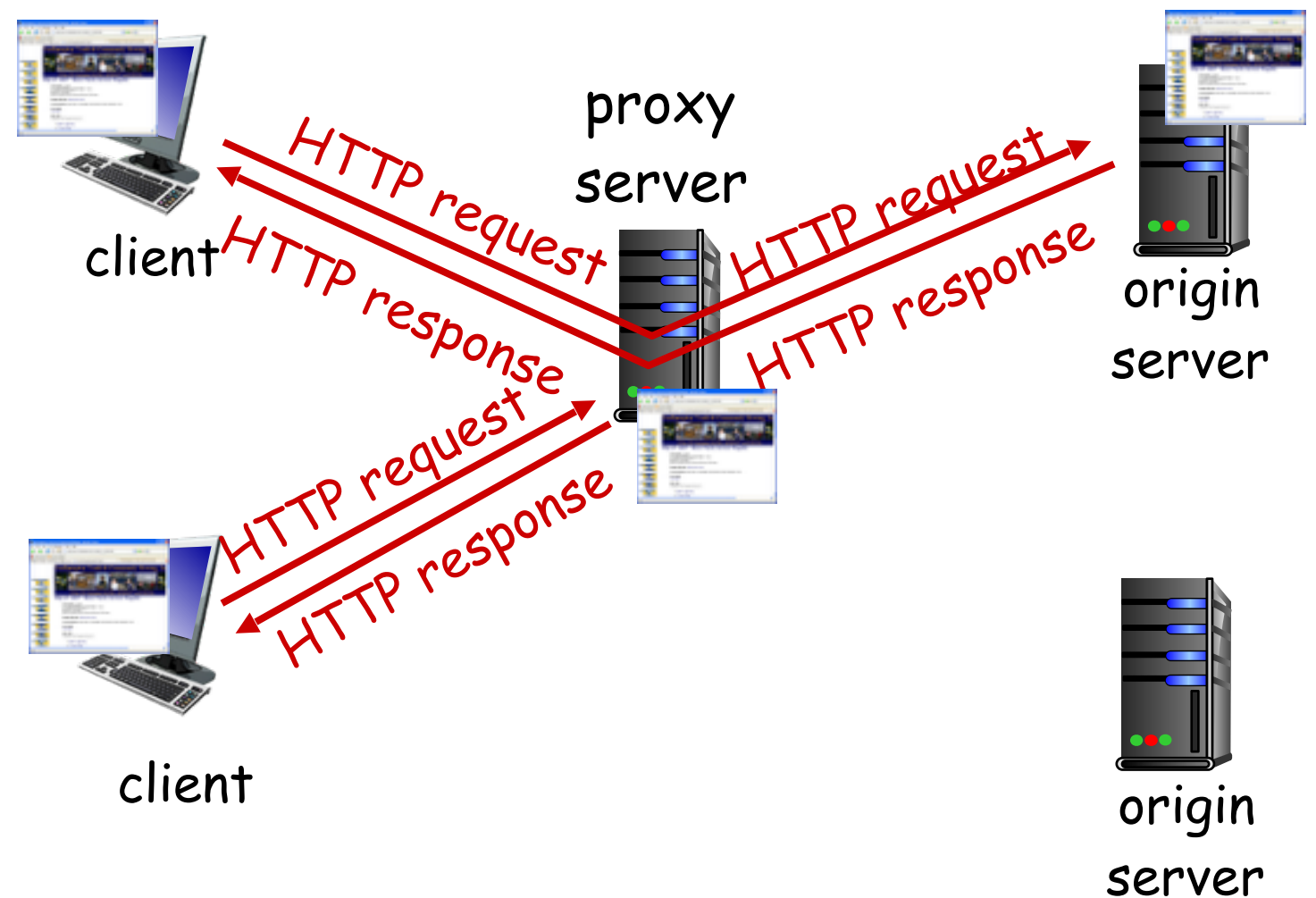
## Server steps, revisited:

- Resolve name of Web page requested
- Perform access control on the Web page
- Check the cache
- Fetch requested page from disk or run program
- Determine the rest of the response
- Return the response to the client
- Make an entry in the server log

# Web caches (proxy server)

*goal:* satisfy client request without involving origin server

- ❖ user sets browser: Web accesses via cache
- ❖ browser sends all HTTP requests to cache
  - object in cache: cache returns object
  - else cache requests object from origin server, then returns object to client



# More about Web caching

- ❖ cache acts as both client and server
  - server for original requesting client
  - client to origin server
- ❖ typically cache is installed by ISP (university, company, residential ISP)

## *why Web caching?*

- ❖ reduce response time for client request
- ❖ reduce traffic on an institution's access link
- ❖ Internet dense with caches: enables “poor” content providers to effectively deliver content (so too does P2P file sharing)

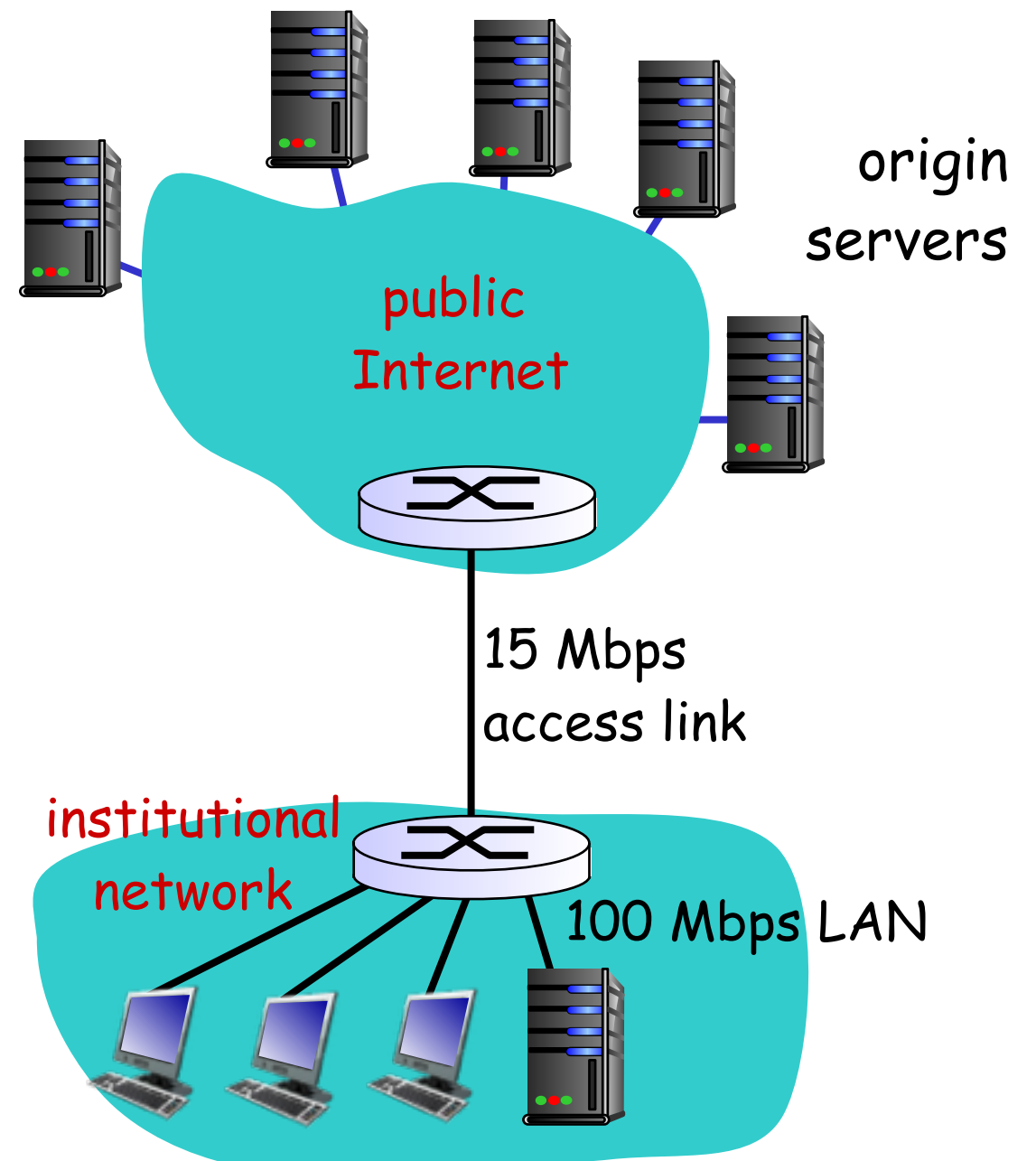
# Caching example:

## *assumptions:*

- v avg object size: 1 Mbits
- v avg request rate from browsers to origin servers: 15 request/sec
- v avg RTT from institutional router to any origin server: 2 sec
- v access link rate: 15 Mbps

## *consequences:*

- v Traffic intensity on the LAN : 15%
- v access link utilization = 100% <sup>problem!</sup>
- v total delay = Internet delay + access delay + LAN delay  
= 2 sec + minutes + millisecs



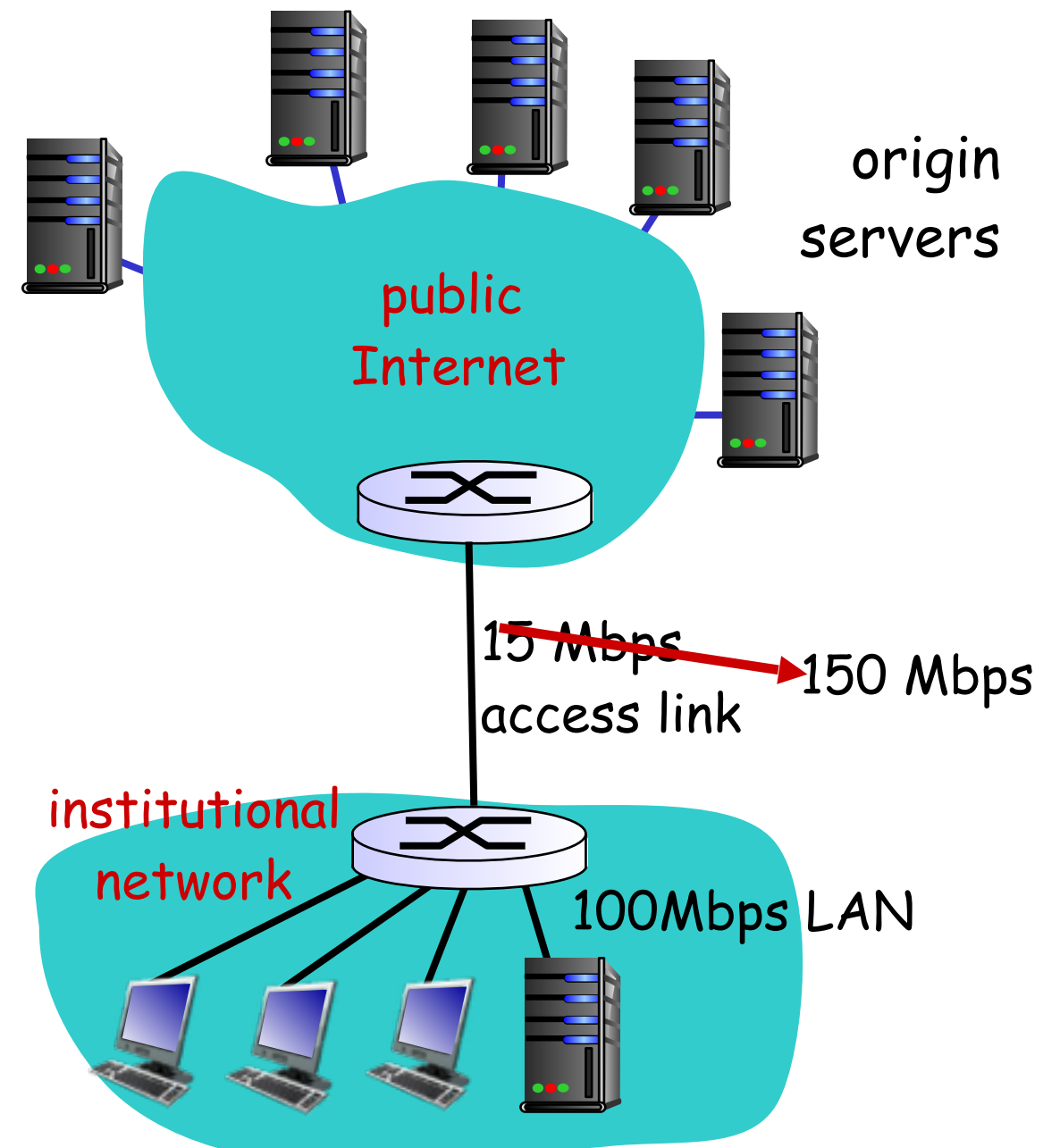
# Caching example: fatter access link

## *assumptions:*

- v avg object size: 1 Mbits
- v avg request rate from browsers to origin servers: 15/sec
- v RTT from institutional router to any origin server: 2 sec
- v access link rate: 15 Mbps

## *consequences:*

- v LAN utilization: 15%
- v access link utilization = 100%
- v total delay = Internet delay + LAN delay  
= 2 sec + minutes + msecs



**Cost:** increased access link speed (not cheap!)



# Caching example: install local cache

## *assumptions:*

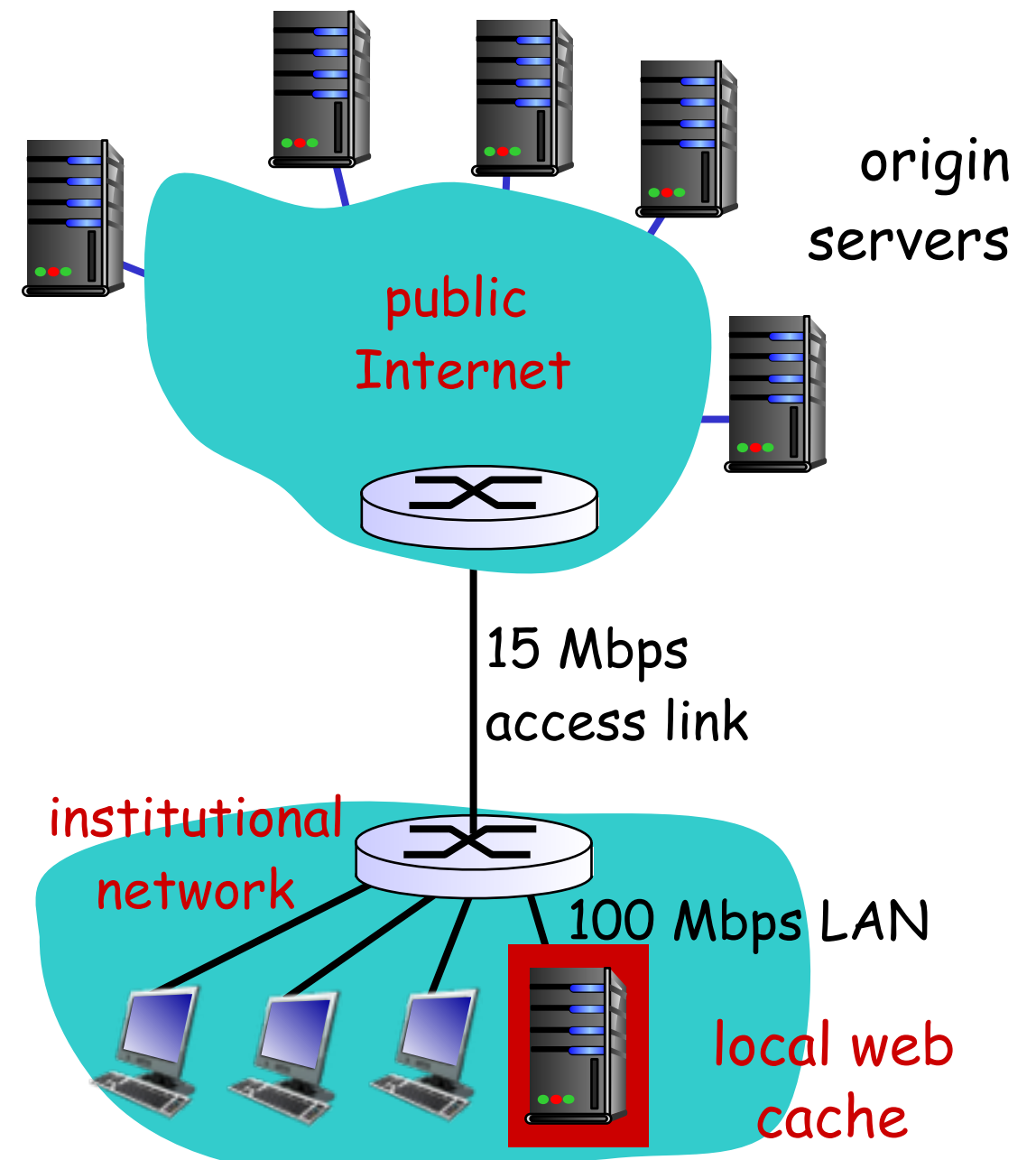
- v avg object size: 1 Mbits
- v avg request rate from browsers to origin servers: 15/sec
- v RTT from institutional router to any origin server: 2 sec
- v access link rate: 15 Mbps

## *consequences:*

- v LAN utilization: 15%
- v access link utilization = 100%
- v total delay = Internet delay + access delay + LAN delay  
= 2 sec + minutes + usecs

*How to compute link utilization, delay?*

**Cost:** web cache (cheap!)



# Caching example: install local cache

## *Calculating access link utilization, delay with cache:*

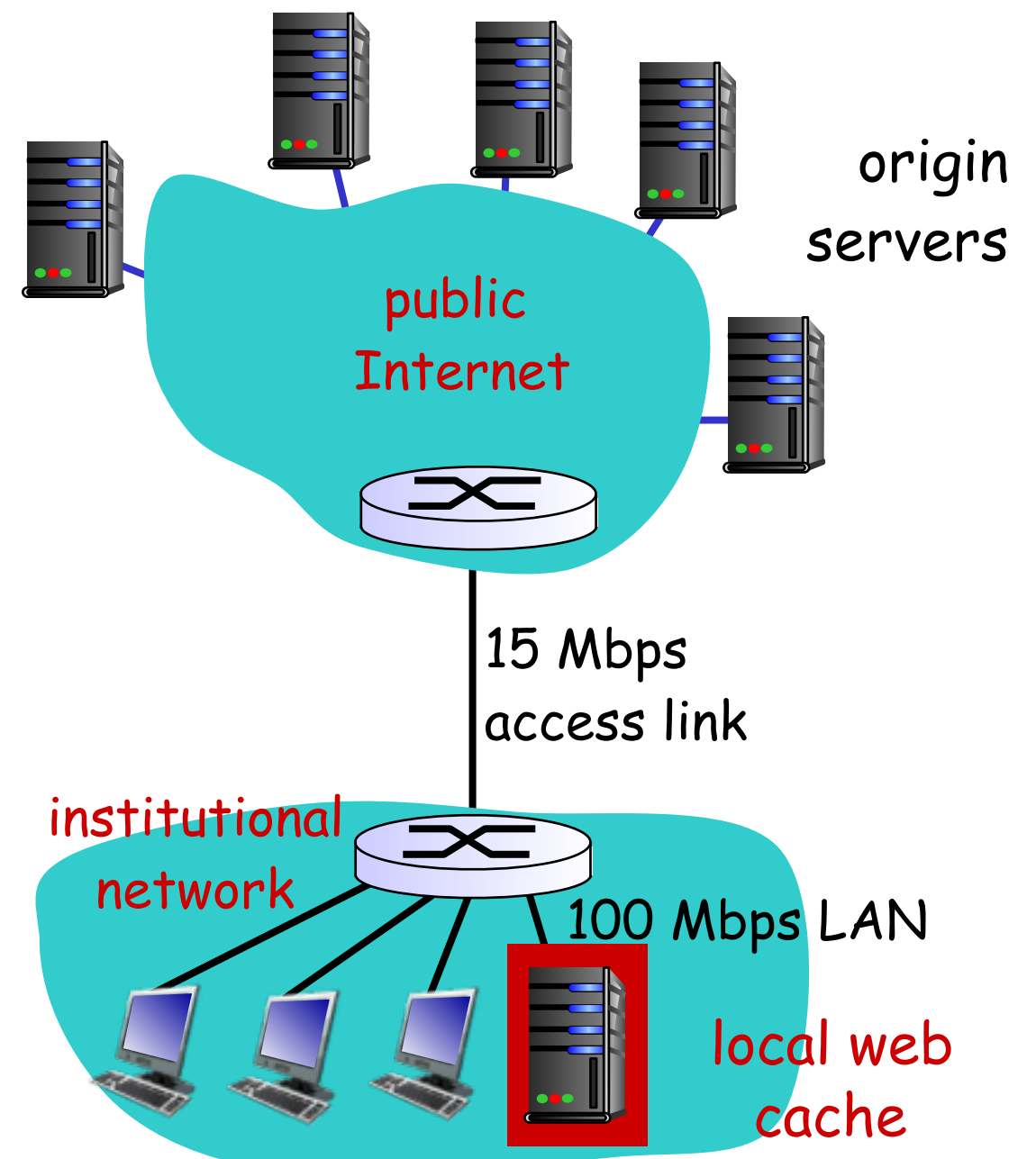
- ❖ suppose cache hit rate is 0.4
  - 40% requests satisfied at cache, 60% requests satisfied at origin

### v access link utilization:

- § 60% of requests use access link
- § access link utilization = 60%

### v total delay

- §  $= 0.6 * (\text{delay from origin servers}) + 0.4 * (\text{delay when satisfied at cache})$
- §  $= 0.6 (2.01) + 0.4 (\sim \text{msecs})$
- §  $= \sim 1.2 \text{ secs}$
- § less than with 150 Mbps link (and cheaper too!)



# Conditional GET

- ❖ **Goal:** don't send object if cache has up-to-date cached version
  - no object transmission delay
  - lower link utilization
- ❖ **cache:** specify date of cached copy in HTTP request  
If-modified-since: <date>
- ❖ **server:** response contains no object if cached copy is up-to-date:  
HTTP/1.0 304 Not Modified

client



server



HTTP request msg  
If-modified-since: <date>

object  
not  
modified  
before  
<date>

HTTP response  
HTTP/1.0  
304 Not Modified



HTTP request msg  
If-modified-since: <date>

object  
modified  
after  
<date>

HTTP response  
HTTP/1.0 200 OK  
<data>