Standard Code Library

FLself

SCUT

August 26, 2022

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一切的开始

一些宏定义

● 需要 C++11

```
// #define DEBUG
   // #define InTerminal
   // #define std cpp17
   #include<bits/stdc++.h>
   #define int long long
   #define PII std::pair<int, int>
   #define VI std::vector<int>
   #define VPII std::vector<std::pair<int, int> >
   #define VVI std::vector<std::vector<int> >
   #define ALL(a) (a).begin(), (a).end()
   #define SIZ(a) ((int)(a).size())
11
   #define FOR(i, l, r) for (int i = (l); i \le (r); ++i)
   #define REP(i, r, l) for (int i = (r); i \ge (l); --i)
13
   #define lowbit(x) ((x) & (-(x)))
   #define lbpos(x) (__builtin_ctz(x))
15
   #define hbpos(x) (31 - __builtin_clz(x))
16
17
    template<typename S, typename T> std::istream &operator>>(std::istream &is, std::pair<S, T> &pp) { is >> pp.first >>
18

    pp.second; return is; }

     \textbf{template} < \textbf{typename S, typename T} > \textbf{std}::ostream \& \textbf{operator} << (\textbf{std}::ostream \& \textbf{os, std}::pair < \textbf{S, T} > \textbf{pp}) \  \  \{ \  \  os << \  \  " (" < < \  \  ostream \  \  \  \  ) \} 
19
     template<typename S, std::size_t _siz> std::istream &operator>>(std::istream &is, std::array<S, _siz> &arr) { for
    template<typename S, std::size_t _siz> std::ostream &operator<<(std::ostream &os, std::array<S, _siz> arr) { os <</pre>
     \textbf{template} < \textbf{typename T} > \textbf{std}:: \textbf{istream \& operator} >> (\textbf{std}:: \textbf{istream \& is, std}:: \textbf{vector} < \textbf{T} > \textbf{\&vec}) ~ \{ \textbf{ for (auto \& x: vec) is } >> \textbf{x; } \}
22

→ return is: }

    template<typename T> std::ostream &operator<<(std::ostream &os, const std::vector<T> &vec) { os << '{'}; for (auto
23
     ⇔ &x: vec) os << x << ", "; return os << "}";}</pre>
    #ifdef std cpp17
24
    template < class Tuple, std::size_t... Is> void print_tuple_impl(std::ostream &os, const Tuple &t,
    ⇔ std::index_sequence<Is...>) { ((os << (Is == 0? "" : ", ") << std::get<Is>(t)), ...); }
    template <class... Args> std::ostream &operator << (std::ostream &os, const std::tuple <Args...> &t) { os << "(";

    print_tuple_impl(os, t, std::index_sequence_for<Args...>{}); return os << ")"; }
</pre>
    #endif
27
    #ifdef DEBUG
   #ifdef InTerminal
29
    #define dbg(x...) do { std::cerr << "\033[32;1m" << #x << " -> "; <math>err(x); } while (0)
    void err() { std::cerr << "\033[39;0m" << std::endl; }</pre>
31
32
    #define dbg(x...) do { std::cerr << \#x << " -> "; err(x); } while (0)
   void err() { std::cerr << std::endl; }</pre>
34
    template<typename T, typename... A>
36
    void err(T a, A... x) { std::cerr << a << ' '; err(x...); }</pre>
37
    #else
38
    #define dbg(...)
39
    #endif
41
42
    using namespace std;
    const int maxn = 2e5 + 3;
43
    const int INF = 0x3f3f3f3f3f3f3f3f3f3f3;
44
    const int mod = 998244353;
    mt19937 RD(time(0));
46
47
48
49
50
   void solv() {
51
52
53
        return ;
54
   }
55
56
    signed main() {
        // freopen("./data.in", "r", stdin);
```

```
std::ios::sync_with_stdio(false), std::cin.tie(0), std::cout.tie(0);
58
59
        int beg__TT = clock();
60
        signed _ttt;
61
62
        cin >> _ttt;
63
        while(_ttt--)
64
65
            solv():
66
        #ifdef DEBUG
67
        std::cerr << "use : " << (clock() - beg__TT) << "ms\n";
68
        #endif
70
        return 0;
   }
71
    数据结构
    ST 表
       一维
    class Sparcetable {
        vector<vector<int> > st;
        int siz;
        bool MX_flg = 0;
        inline int renew(int x, int y) {
            if (MX_flg) return max(x, y);
            return min(x, y);
    public:
        // 注意 bhpos(0) 返回-1
10
        bool (*comp)(int, int);
11
        Sparcetable():siz(maxn) {st.resize(hbpos(maxn - 1) + 1, std::vector<int> (maxn));}
        Sparcetable(const std::vector<int>& a, bool _MX_flg = 1): siz(a.size()), MX_flg(_MX_flg) {
13
14
            int n = a.size();
            st.resize(hbpos(n) + 1, vector < int > (n + 1));
15
            for (int i = 1; i <= n; ++i) st[0][i] = a[i];</pre>
16
            for (int i = 1; i <= hbpos(siz); ++i) {</pre>
17
                 for (int j = 1; j + (1 << i) <= siz + 1; ++j) {
18
19
                     st[i][j] = renew(st[i - 1][j], st[i - 1][j + (1 << (i - 1))]);
                }
20
            }
21
22
        int query(int l, int r) {
23
            int len = hbpos(r - l + 1);
24
            return renew(st[len][l], st[len][r - (1 << len) + 1]);</pre>
25
27
   };
       二维
    int f[10][10][maxn][maxn];
    #define _highbit(x) (31 - __builtin_clz(x))
    inline int calc(int x, int y, int xx, int yy, int p, int q) {
            \max(f[p][q][x][y], f[p][q][xx - (1 << p) + 1][yy - (1 << q) + 1]),
            \max(f[p][q][xx - (1 << p) + 1][y], f[p][q][x][yy - (1 << q) + 1])
        );
   }
8
    void init() {
        for (int x = 0; x <= _highbit(n); ++x)</pre>
10
        for (int y = 0; y <= _highbit(m); ++y)</pre>
11
            for (int i = 0; i \le n - (1 \le x); ++i)
12
            for (int j = 0; j \le m - (1 << y); ++j) {
13
                 if (!x && !y) { f[x][y][i][j] = a[i][j]; continue; }
14
15
                 f[x][y][i][j] = calc(
16
                     i, j,
                     i + (1 << x) - 1, j + (1 << y) - 1,
17
```

max(x - 1, 0), max(y - 1, 0)

18

);

```
}
20
21
    inline int get_max(int x, int y, int xx, int yy) {
22
23
        return calc(x, y, xx, yy, _highbit(xx - x + 1), _highbit(yy - y + 1));
    Fenwick Tree(树状数组)
       一维
    template<typename T>
1
    class FenwickT {
        int n:
        vector<T> tr;
    public:
5
        FenwickT(int siz): tr(siz), n(siz) {}
        FenwickT(int siz, T ini): tr(siz, ini), n(siz) {}
        void add(int p, T x) {
            for (int i = p; i < n; i += i & (-i)) tr[i] += x;</pre>
10
11
        T query(int p) {
12
            T ret = T(0);
            for (int i = p; i > 0; i -= i & (-i)) ret += tr[i];
13
14
            return ret;
15
16
        T range_sum(int l, int r) {
            return (query(r) - query(l - 1));
17
18
   };
    线段树
    class SegT{
1
        struct Pt{
            int val, ls, rs, sl;
3
            Pt():val(0), ls(0), rs(0), sl(0) {}
            Pt(int v, int l, int r, int s): val(v), ls(l), rs(r), sl(s) {}
            ~Pt() {}
        int renew(int x, int y) {
            return x + y;
10
        int renewsl(int x, int y) {
11
12
            return x + y;
        }
13
14
        int renewqj(int x, int y, int l, int r) {
            return x + y * (r - l + 1);
15
16
    public:
17
        std::vector<Pt> tr;
18
        int cnt = 0; // 结点数目
19
        SegT() {cnt = 1; tr.emplace_back();} // 建一棵空树
20
21
        // 修改
22
23
        void add(int l, int r, int p, int ql, int qr, int x) {
24
            if (ql <= l && r <= qr) {
                tr[p].val = renewqj(tr[p].val, x, l, r);
25
                tr[p].sl = renewsl(tr[p].sl, x);
27
                return;
28
            int mid = l + (r - l) / 2;
29
            if (tr[p].sl) {
30
                if (!tr[p].ls) tr[p].ls = cnt++;
                if (!tr[p].rs) tr[p].rs = cnt++;
32
                if (cnt >= tr.size()) tr.resize(cnt + 1);
33
34
                int ls = tr[p].ls, rs = tr[p].rs;
               tr[ls].val = renewqj(tr[ls].val, tr[p].sl, l, mid), tr[rs].val = renewqj(tr[rs].val, tr[p].sl, mid + 1, r);
35
                tr[ls].sl = renewsl(tr[ls].sl, tr[p].sl), tr[rs].sl = renewsl(tr[rs].sl, tr[p].sl);
                tr[p].sl = 0;
37
            }
```

```
if (ql <= mid) {
39
40
                if (!tr[p].ls) tr[p].ls = cnt++;
                if (cnt >= tr.size()) tr.resize(cnt + 1);
41
42
                add(l, mid, tr[p].ls, ql, qr, x);
43
            if (qr > mid) {
44
                if (!tr[p].rs) tr[p].rs = cnt++;
45
                if (cnt >= tr.size()) tr.resize(cnt + 1);
46
                add(mid + 1, r, tr[p].rs, ql, qr, x);
47
48
            tr[p].val = renew(tr[p].ls? tr[tr[p].ls].val: 0, tr[p].rs? tr[tr[p].rs].val: 0);
49
50
        };
51
52
        int query(int l, int r, int p, int ql, int qr) {
53
            if (ql <= l && r <= qr) return tr[p].val;</pre>
54
            int mid = l + (r - l) / 2;
55
            if (tr[p].sl) {
56
                if (!tr[p].ls) tr[p].ls = cnt++;
                if (!tr[p].rs) tr[p].rs = cnt++;
58
59
                if (cnt >= tr.size()) tr.resize(cnt + 1);
60
                int ls = tr[p].ls, rs = tr[p].rs;
               tr[ls].val = renewqj(tr[ls].val, tr[p].sl, l, mid), tr[rs].val = renewqj(tr[rs].val, tr[p].sl, mid + 1, r);
61
                tr[ls].sl = renewsl(tr[ls].sl, tr[p].sl), tr[rs].sl = renewsl(tr[rs].sl, tr[p].sl);
                tr[p].sl = 0;
63
64
            int ret = 0;
65
            if (ql <= mid) {
66
67
                if (!tr[p].ls) tr[p].ls = cnt++;
                if (cnt >= tr.size()) tr.resize(cnt + 1);
68
                ret = renew(ret, query(l, mid, tr[p].ls, ql, qr));
69
70
            if (qr > mid) {
71
72
                if (!tr[p].rs) tr[p].rs = cnt++;
                if (cnt >= tr.size()) tr.resize(cnt + 1);
73
                ret = renew(ret, query(mid + 1, r, tr[p].rs, ql, qr));
74
75
            tr[p].val = renew(tr[p].ls? tr[tr[p].ls].val: 0, tr[p].rs? tr[tr[p].rs].val: 0);
76
77
            return ret;
78
        };
   };
```

数学

数论

数论整数

```
constexpr int P = 998244353;
    // assume -P \le x \le 2P
    int norm(int x) {
        // x %= P;
        if (x < 0) \{ x += P; \}
        if (x >= P) { x -= P; }
        return x;
    }
    template<typename E>
10
    E power(E n, int k) {
        E ret = E(1);
11
12
        while (k) {
            if (k & 1) ret *= n;
13
            n \star = n;
            k >>= 1;
15
        } return ret;
16
17
    struct Z {
18
        int x;
        Z(int x = 0) : x(norm(x)) \{ \}
20
        int val() const { return x; }
21
        Z operator-() const { return Z(norm(P - x)); }
22
```

```
Z inv() const { assert(x != 0); return power(*this, P - 2); }
23
24
                 Z & operator *= (const Z & rhs) { x = (long long)(x) * rhs.x % P; return *this; }
                 Z &operator+=(const Z &rhs) { x = norm(x + rhs.x); return *this; }
25
                 Z &operator-=(const Z &rhs) { x = norm(x - rhs.x); return *this; }
26
27
                 Z &operator/=(const Z &rhs) { return *this *= rhs.inv(); }
                 friend Z operator*(const Z &lhs, const Z &rhs) { Z res = lhs; res *= rhs; return res; }
28
                 friend Z operator+(const Z &lhs, const Z &rhs) { Z res = lhs; res += rhs; return res;
29
                 friend Z operator-(const Z &lhs, const Z &rhs) { Z res = lhs; res -= rhs; return res; }
30
                 friend Z operator/(const Z &lhs, const Z &rhs) { Z res = lhs; res /= rhs; return res; }
31
                 friend std::istream &operator>>(std::istream &is, Z &a) { long long v; is >> v; a = Z(v); return is; }
32
                 friend std::ostream &operator<<(std::ostream &os, const Z &a) { return os << a.val(); }</pre>
33
       };
        拉格朗日差值法
        欧几里得
                • 扩展欧几里得
        int exgcd(int a, int b, int& x, int& y) {
2
                 if (a == 0) {
                        x = 0, y = 1;
                         return b;
                 int ret = exgcd(b % a, a, x, y), xx = x;
                 x = y - b / a * x;
                y = xx;
                 return ret;
        类欧几里得
               • f(a,b,c,n) = \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor: 当 a \geq c or b \geq c 时, f(a,b,c,n) = (\frac{a}{c})n(n+1)/2 + (\frac{b}{c})(n+1) + f(a \bmod c, b \bmod c, c, n); 否则 f(a,b,c,n) = nm - f(c,c-b-1,a,m-1) 。
               • g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor : \exists a \geq c \text{ or } b \geq c \text{ fl}, g(a,b,c,n) = (\frac{a}{c})n(n+1)(2n+1)/6 + (\frac{b}{c})n(n+1)/2 + (\frac{b}{c
                    g(a \bmod c, b \bmod c, c, n); 否则 g(a, b, c, n) = \frac{1}{2}(n(n+1)m - f(c, c-b-1, a, m-1) - h(c, c-b-1, a, m-1))。
               • h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2: \exists a \geq c \text{ or } b \geq c \text{ pl}, h(a,b,c,n) = (\frac{a}{c})^2 n(n+1)(2n+1)/6 + (\frac{b}{c})^2 (n+1) + (\frac{b}{c})^2 (n+1)
                    (\frac{a}{c})(\frac{b}{c})n(n+1)+h(a \bmod c,b \bmod c,c,n)+2(\frac{a}{c})g(a \bmod c,b \bmod c,c,n)+2(\frac{b}{c})f(a \bmod c,b \bmod c,c,n);否则
                    h(a,b,c,n) = nm(m+1) - 2g(c,c-b-1,a,m-1) - 2f(c,c-b-1,a,m-1) - f(a,b,c,n)
        template<typename T>
        T lintp(const vector<int>& x, const vector<T>& y, int k) {
                 T ans = 0:
                 for (int i = 0; i < x.size(); ++i) {</pre>
                         if (k == x[i])
                                return y[i];
                         T u = 1, v = 1;
                         for (int j = 0; j < x.size(); ++j) {</pre>
                                 if (i == j) continue;
                                 u *= (k - x[j]);
10
11
                                 v *= (x[i] - x[j]);
12
                         ans += y[i] * u / v;
13
15
                 return ans:
        中国剩余定理
        int CRT(std::vector<int> a, std::vector<int> r) {
                long long n = 1, ans = 0, k = a.size();
                 for (int i = 0; i < k; ++i) n = n * r[i];</pre>
                 for (int i = 0; i < k; ++i) {</pre>
                         long long m = n / r[i], b, y;
                         exgcd(m, r[i], b, y);
                        ans = (ans + a[i] * m % n * b % n + n) % n;
                }
```

```
9 return ans;
10 }
```

各种筛

筛法	场景	效率
Min-25 筛 PN 筛 杜教筛	$f(p)$ 是一个关于 p 的多项式, $f(p^c)$ 能快速求 找一个好求前缀和的积性函数 $g()$ 在 p 处 $f(p)=g(p)$ 找一个 $g()$ 使得 $f*g()$ 好求前缀和	$O(\frac{n^{\frac{3}{4}}}{\log n})$ $O(\sqrt{n})$ $O(n^{\frac{2}{3}})$

Min-25 筛

$$\begin{split} F_k(n) &= \sum_{i=2}^n [p_k \leq \mathrm{lpf}(i)] f(i) \\ &= \sum_{k \leq i} \sum_{\substack{c \geq 1 \\ p_i^2 \leq n}} f\left(p_i^c\right) \left([c > 1] + F_{i+1}\left(n/p_i^c\right)\right) + \sum_{k \leq i} f(p_i) \\ &= \sum_{k \leq i} \sum_{\substack{c \geq 1 \\ p_i^2 \leq n}} f\left(p_i^c\right) \left([c > 1] + F_{i+1}\left(n/p_i^c\right)\right) + F_{\mathrm{prime}}(n) - F_{\mathrm{prime}}(p_{k-1}) \\ &= \sum_{k \leq i} \sum_{\substack{c \geq 1 \\ p_i^2 \leq n}} \int_{\substack{c \geq 1 \\ p_i^2 \leq n}} \left(f\left(p_i^c\right) F_{i+1}\left(n/p_i^c\right) + f\left(p_i^{c+1}\right)\right) + F_{\mathrm{prime}}(n) - F_{\mathrm{prime}}(p_{k-1}) \\ &G_k(n) = G_{k-1}(n) - \left[p_k^2 \leq n\right] g(p_k) (G_{k-1}(n/p_k) - G_{k-1}(p_{k-1})) \end{split}$$

```
namespace Min_25{
        bool INI = 0;
2
        int M = 0, sqn = 0;
        std::vector<int> primes;
        Z (*fp)(int, int); // f(p^c)
        Z fp_single(int p, int c) {
            return power(Z(p), c);
        Z sgp(int n, int c) {
            if (c == 0) return Z(n);
            if (c == 1) return Z(n+1) * n / 2;
11
            if (c == 2) return Z(2*n+1) * (n+1) * n / 6;
12
            // if (c == 3) return Z(n) * n * (n+1) * (n+1) / 4;
13
            // if (c == 4) return Z(n) * (n+1) * (2*n+1) * (3*n*n%P + 3*n - 1) / 30;
14
        struct Gpoly{
16
17
            int M, sqn, a;
            std::vector<Z> F_prime1, F_prime2;
18
            Gpoly(int scale, int a, int c): M(scale), a(a) {
19
                sqn = sqrt(M) + 3;
                F_prime1.resize(sqn), F_prime2.resize(sqn);
21
22
                int st = M, ed = 1;
23
                for (int i = 1; st >= ed; st = min(M / ++i, st - 1)) {
24
25
                     (*this)[st] = sgp(st, c) - 1;
26
                for (int j = 1; j < primes.size(); ++j) {</pre>
                     st = M, ed = primes[j] * primes[j];
28
                     for (int i = 1; st >= ed; st = min(M / ++i, st - 1)) {
29
                         (*this)[st] -= fp_single(primes[j], c) * ((*this)[st/primes[j]] - (*this)[primes[j-1]]);
                     }
31
                }
32
33
            void seta(int aa) {a = aa;}
34
            Z & operator [](int idx) {
35
                if (idx < sqn) return F_prime1[idx];</pre>
36
                return F_prime2[M/idx];
```

```
}
38
39
        };
40
        std::vector<Gpoly> Gp;
41
        Z G(int x){
42
            Z res = 0;
43
             for (auto &g: Gp) res += g[x] * g.a;
44
            return res;
45
46
        void init(int scale) {
47
            if (M < scale) INI = 0;</pre>
48
49
            M = scale, sqn = sqrt(M) + 5;
50
            if (!INI) {
51
52
                 INI = 1;
                 std::vector<signed> vis(sqn);
53
54
                 primes.resize(1, 1);
                 for (int i = 2; i < sqn; ++i) {</pre>
55
                     if (!vis[i]) primes.push_back(i);
57
                     for (int j = 1; j < primes.size(); ++j) {</pre>
                         auto prm = primes[j];
58
                         if (prm * i >= sqn) break;
59
                         vis[i * prm] = 1;
60
                          if (i % prm == 0) break;
                     }
62
                 }
63
64
            }
65
            Gp.clear();
67
            Gp.emplace_back(scale, N, 0);
68
            // Gp.emplace_back(scale, a, c); // a*p^c
69
70
        Z seive(int n, int k) {
            Z res = G(n) - G(primes[k-1]);
72
             for (int i = k; i < primes.size() && primes[i] * primes[i] <= n; ++i) {</pre>
73
                 Z fpj = fp(primes[i], 1), fpj1;
74
                 for (int j = 1, pc = primes[i]; pc * primes[i] <= n; ++j, pc *= primes[i], fpj = fpj1) {
75
76
                     fpj1 = fp(primes[i], j + 1);
                     res += fpj \star seive(n/pc, i + 1) + fpj1;
77
78
            }
79
            return res;
80
81
        Z S_f(int n) {
82
83
            return seive(n, 1) + 1;
84
85
        void seta(const vector<int>& a) {for (int i = 0; i < Gp.size() && i < a.size(); ++i) Gp[i].seta(a[i]);}</pre>
   }
86
    PN 筛 (Powerful Number)
    // f = (g*h); f(p) = g(p) + h(p) [with f(p) = g(p)] -> h[p] = 0.
1
2
    namespace PowerfulNumber{
        bool INI = 0;
        std::vector<int> primes;
        std::vector<std::vector<Z> > h;
5
        int M;
        Z (*fp)(int, int); // f(p^c)
        Z (*gp)(int, int); // g(p^{c})
        Z (*S_g)(int);
                             // preffix sum of g()
10
        // if lack of a formula of h(i, j), you need to set f() and g() beforehand
11
        void init(int scale) {
12
            if (M < scale) INI = 0;</pre>
13
            M = scale;
            int n = std::sqrt(M)+10;
15
16
             if (!INI) {
                 INI = 1;
17
                 primes.resize(0);
18
                 std::vector<signed> vis(n);
```

```
for (int i = 2; i < n; ++i) {</pre>
20
21
                      if (!vis[i]) primes.push_back(i);
                     for (auto prm: primes) {
22
                          if (prm * i >= n) break;
23
                          vis[i * prm] = 1;
24
                          if (i % prm == 0) break;
25
26
27
                 h.resize(primes.size(), std::vector<Z>((int)(log2(M))+1));
28
            }
29
30
31
             // get the function h() (with f() and g() set or with formula of h())
            for (int i = 0; i < h.size(); ++i) {</pre>
32
                 int pp = primes[i] * primes[i];
33
34
                 if (pp > M) break;
                 h[i][0] = 1;
35
36
                 for (int j = 2; pp <= M && j < h[i].size(); ++j, pp *= primes[i]) {</pre>
                     Z sgh = 0;
37
                     for (int k = 1, xp = primes[i]; k <= j; ++k, xp *= primes[i]) {</pre>
38
39
                          sgh += gp(primes[i], xp) * h[i][j-k];
40
41
                     h[i][j] = fp(pp, j) - sgh;
                     // h[i][j] = Z(j-1) * (pp * primes[i] % P - pp % P);
42
                     // h[i][j] = Z(-pp) / (j * (j-1));
43
                     // [a formula of h(i, j)], better faster than log, this example can be optimized to O(1).
44
45
                 }
            }
46
        }
47
48
        // assistance func to get the sum
49
        Z PN_sieve(int n, int flr, Z hd) {
50
            Z res = S_g(n)*hd;
51
            for(int i = flr+1; i < primes.size(); ++i) {</pre>
52
53
                 int prm=primes[i], k=1;
                 int val=n/prm, pk=prm;
54
55
                 if(val < prm) break;</pre>
56
                 while(val >= prm) {
57
58
                     val /= prm;
                     pk *= prm;
59
60
                     ++k;
61
                     res += PN_sieve(val, i, hd*h[i][k]);
62
63
                 }
64
65
            return res;
66
        // func to get the sum
        Z getsumf(int n) {
68
69
            return PN_sieve(n, -1, 1);
70
   }
71
    杜教筛
    // g(1)S(n) = \sum_{i=1}^{n} (f*g)(i) - \sum_{i=2}^{n} g(i)S(n/i)
    namespace dujiaoshai{
2
        vector<Z> smallS;
3
4
        map<int, Z> bigS;
        void Set_smallS(int);
        Z (*S_fg)(int);
        Z (*S_g)(int);
8
        Z getS(int n) {
            if (n < smallS.size()) return smallS[n];</pre>
10
            else if (bigS.count(n)) return bigS[n];
            Z res = S_fg(n);
12
13
            for (int l = 2, r; l <= n; l = r + 1) {
                 r = n / (n / l);
14
                 res -= (S_g(r) - S_g(l-1)) * getS(n/l);
15
            }
```

```
// res /= S_g(1);
17
18
             return (bigS[n] = res);
19
    }
20
    void dujiaoshai::Set_smallS(int siz) {
22
        int n = pow(siz, 0.67);
23
        smallS.assign(n, 0);
24
        smallS[1] = 1;
25
26
        vector<signed> vis(n), primes;
        for (int i = 2; i < n; ++i) {</pre>
27
28
             if (!vis[i]) {
29
                 primes.push_back(i);
                 smallS[i] = i - 1;
30
31
             for (auto &prm: primes) {
32
33
                 if (i * prm >= n) break;
                 vis[i * prm] = 1;
34
                 if (i % prm == 0) {
35
                     smallS[i * prm] = smallS[i] * prm;
36
37
                     break;
                 }
38
39
                 smallS[i * prm] = smallS[i] * (prm - 1);
             }
            smallS[i] = smallS[i-1] + smallS[i]*i;
41
42
    }
43
    多项式
        • Poly with NTT
    std::vector<int> rev;
1
    std::vector<Z> roots{0, 1};
2
    void dft(std::vector<Z> &a) {
        int n = a.size();
5
        if ((int)(rev.size()) != n) {
6
             int k = __builtin_ctz(n) - 1;
             rev.resize(n);
             for (int i = 0; i < n; i++) {</pre>
                 rev[i] = rev[i >> 1] >> 1 | (i & 1) << k;
10
             }
11
        }
12
13
14
        for (int i = 0; i < n; i++) {</pre>
             if (rev[i] < i) {
15
                 std::swap(a[i], a[rev[i]]);
16
             }
17
18
        if ((int)(roots.size()) < n) {</pre>
             int k = __builtin_ctz(roots.size());
20
             roots.resize(n);
21
             while ((1 << k) < n) {
22
                 Z = power(Z(3), (P - 1) >> (k + 1));
23
                 for (int i = 1 << (k - 1); i < (1 << k); i++) {
                     roots[2 * i] = roots[i];
25
26
                     roots[2 * i + 1] = roots[i] * e;
                 }
27
                 k++;
28
            }
29
30
        for (int k = 1; k < n; k *= 2) {</pre>
31
             for (int i = 0; i < n; i += 2 * k) {
32
                 for (int j = 0; j < k; j++) {
33
                     Z u = a[i + j];
34
                     Z v = a[i + j + k] * roots[k + j];
35
36
                     a[i + j] = u + v;
                     a[i + j + k] = u - v;
37
                 }
38
            }
39
```

```
40
41
    }
    void idft(std::vector<Z> &a) {
42
43
         int n = a.size();
         std::reverse(a.begin() + 1, a.end());
44
         dft(a);
45
         Z inv = (1 - P) / n;
46
         for (int i = 0; i < n; i++) {
47
             a[i] *= inv;
48
49
    }
50
51
    struct Poly {
         std::vector<Z> a;
52
         Poly() {}
53
         Poly(const std::vector<Z> &a) : a(a) {}
54
         Poly(const std::initializer_list<Z> &a) : a(a) {}
55
56
         int size() const {
57
             return a.size();
         void resize(int n) {
59
             a.resize(n);
60
61
         Z operator[](int idx) const {
62
             if (idx < size()) {
                 return a[idx];
64
65
             } else {
66
                 return 0;
             }
67
68
         Z & operator[](int idx) {
69
             return a[idx];
70
71
         Poly mulxk(int k) const {
72
73
             auto b = a;
             b.insert(b.begin(), k, 0);
74
75
             return Poly(b);
76
         Poly modxk(int k) const {
77
78
             k = std::min(k, size());
             return Poly(std::vector<Z>(a.begin(), a.begin() + k));
79
80
         Poly divxk(int k) const {
81
             if (size() <= k) {
82
83
                 return Poly();
84
85
             return Poly(std::vector<Z>(a.begin() + k, a.end()));
86
         friend Poly operator+(const Poly &a, const Poly &b) {
             std::vector<Z> res(std::max(a.size(), b.size()));
88
89
             for (int i = 0; i < (int)(res.size()); i++) {</pre>
90
                 res[i] = a[i] + b[i];
             }
91
             return Poly(res);
93
94
         friend Poly operator-(const Poly &a, const Poly &b) {
             std::vector<Z> res(std::max(a.size(), b.size()));
95
             for (int i = 0; i < (int)(res.size()); i++) {</pre>
96
97
                 res[i] = a[i] - b[i];
98
             return Poly(res);
99
100
         friend Poly operator*(Poly a, Poly b) {
101
102
             if (a.size() == 0 || b.size() == 0) {
                 return Poly();
103
104
             int sz = 1, tot = a.size() + b.size() - 1;
105
             while (sz < tot) {</pre>
106
107
                 sz *= 2;
108
             a.a.resize(sz);
             b.a.resize(sz);
110
```

```
dft(a.a);
111
112
              dft(b.a);
              for (int i = 0; i < sz; ++i) {
113
                  a.a[i] = a[i] * b[i];
114
              idft(a.a);
116
              a.resize(tot);
117
              return a:
118
119
120
         friend Poly operator*(Z a, Poly b) {
              for (int i = 0; i < (int)(b.size()); i++) {</pre>
121
122
                  b[i] *= a;
              }
123
              return b;
124
125
         friend Poly operator*(Poly a, Z b) {
126
127
              for (int i = 0; i < (int)(a.size()); i++) {</pre>
                  a[i] *= b;
128
129
130
              return a;
131
132
         Poly &operator+=(Poly b) {
              return (*this) = (*this) + b;
133
134
         Poly &operator-=(Poly b) {
135
              return (*this) = (*this) - b;
136
137
         Poly &operator *= (Poly b) {
138
139
              return (*this) = (*this) * b;
140
         Poly deriv() const {
141
              if (a.empty()) {
142
                  return Poly();
143
144
              std::vector<Z> res(size() - 1);
145
              for (int i = 0; i < size() - 1; ++i) {</pre>
146
                  res[i] = (i + 1) * a[i + 1];
147
148
149
              return Poly(res);
150
151
         Poly integr() const {
              std::vector<Z> res(size() + 1);
152
              for (int i = 0; i < size(); ++i) {</pre>
153
                  res[i + 1] = a[i] / (i + 1);
154
155
156
              return Poly(res);
157
158
         Poly inv(int m) const {
              Poly x{a[0].inv()};
159
              int k = 1;
160
              while (k < m) {
161
                  k *= 2;
162
                  x = (x * (Poly{2} - modxk(k) * x)).modxk(k);
              }
164
165
              return x.modxk(m);
166
         Poly log(int m) const {
167
168
              return (deriv() * inv(m)).integr().modxk(m);
169
         Poly exp(int m) const {
170
171
              Poly x{1};
              int k = 1;
172
173
              while (k < m) {</pre>
                  k *= 2;
174
175
                  x = (x * (Poly{1} - x.log(k) + modxk(k))).modxk(k);
176
177
              return x.modxk(m);
178
         Poly pow(int k, int m) const {
179
              int i = 0;
              while (i < size() && a[i].val() == 0) {</pre>
181
```

```
i++;
182
183
             if (i == size() || 1LL * i * k >= m) {
184
                 return Poly(std::vector<Z>(m));
185
             Z v = a[i];
187
             auto f = divxk(i) * v.inv();
188
             return (f.log(m - i * k) * k).exp(m - i * k).mulxk(i * k) * power(v, k);
189
190
191
         Poly sqrt(int m) const {
             Poly x{1};
192
             int k = 1;
193
             while (k < m) {
194
                 k *= 2;
195
                 x = (x + (modxk(k) * x.inv(k)).modxk(k)) * ((P + 1) / 2);
196
197
198
             return x.modxk(m);
199
200
         Poly mulT(Poly b) const {
             if (b.size() == 0) {
201
                 return Poly();
202
203
             int n = b.size();
204
             std::reverse(b.a.begin(), b.a.end());
205
             return ((*this) * b).divxk(n - 1);
206
207
208
         std::vector<Z> eval(std::vector<Z> x) const {
             if (size() == 0) {
209
                 return std::vector<Z>(x.size(), 0);
211
             const int n = std::max((int)(x.size()), size());
212
             std::vector<Poly> q(4 * n);
213
             std::vector<Z> ans(x.size());
214
215
             x.resize(n);
             std::function<void(int, int, int)> build = [&](int p, int l, int r) {
216
                  if (r - l == 1) {
217
                     q[p] = Poly{1, -x[l]};
218
                 } else {
219
                      int m = (l + r) / 2;
220
                      build(2 * p, l, m);
221
222
                      build(2 * p + 1, m, r);
                      q[p] = q[2 * p] * q[2 * p + 1];
223
                 }
224
225
             };
             build(1, 0, n);
226
227
             std::function<void(int, int, int, const Poly &)> work = [&](int p, int l, int r, const Poly &num) {
                 if (r - l == 1) {
228
                      if (l < (int)(ans.size())) {
229
230
                          ans[l] = num[0];
                     }
231
                 } else {
232
                     int m = (l + r) / 2;
233
                      work(2 * p, l, m, num.mulT(q[2 * p + 1]).modxk(m - l));
234
                      work(2 \star p + 1, m, r, num.mulT(q[2 \star p]).modxk(r - m));
235
236
237
             work(1, 0, n, mulT(q[1].inv(n)));
238
239
             return ans;
240
         }
    };
241
        • FFT
    namespace FFT { // n_ 是初始的数组长度,不一定为 2 的幂次; n 是 init 之后的长度,保证为 2 的幂次长度
1
         const double PI = acos(-1);
2
         int rev[1 << 20];</pre>
         int init(int n_) {
4
             int step = 0, n = 1;
             for (; n < n_; n <<= 1) ++step;</pre>
             for (int i = 1; i < n; ++i) {</pre>
7
                  rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (step - 1));
             }
```

```
return n:
10
11
       void FFT(complex<double> a[], int n, int f) {
12
           for (int i = 0; i < n; ++i) {</pre>
13
               if (i < rev[i]) std::swap(a[i], a[rev[i]]);</pre>
15
           for (int h = 2; h <= n; h <<= 1) {</pre>
16
               complex<double> wn(cos(f * 2 * PI / h), sin(f * 2 * PI / h));
17
               for (int i = 0; i < n; i += h) {
18
                   complex<double> w(1, 0), u;
19
                   for (int j = i, k = h >> 1; j < i + k; ++j) {
20
                      u = a[j + k] * w;
21
                       a[j + k] = a[j] - u;
22
                       a[j] = a[j] + u;
23
                       w = w * wn;
24
                   }
25
               }
27
           if (f == -1) {
               for (int i = 0; i < n; ++i) {</pre>
29
                   a[i] = {a[i].real() / n, 0};
30
31
           }
32
       34
35
           int n = init(n_);
36
           FFT(a, n, 1);
           FFT(b, n, 1);
37
           for (int i = 0; i < n; ++i) a[i] *= b[i];</pre>
           FFT(a, n, -1);
39
40
   }
41
```

组合数学

卡特兰数

可以将问题划分为两个子问题的问题,满足递推式:

$$H_n = \begin{cases} \sum_{i=1}^{n} H_{i-1} H_{n-i} & n \ge 2, n \in \mathbb{N}_{+} \\ 1 & n = 0, 1 \end{cases}$$

递推式:
$$H_n=rac{H_{n-1}(4n-2)}{n+1}$$
 通项式: $H_n={2n\choose n}-{2n\choose {n-1}}$

斯特林数

将 n 个不相同的元素划分为 k 个互不区分的集合, $\left\{ \begin{matrix} n \\ k \end{matrix} \right\}$,也记作 S(n,k) .

递推式: $S(n,k)=S(n-1,k)+k\times S(n-1,k)$ 通项公式: $\binom{n}{m}=\sum_{i=0}^m\frac{(-1)^{m-i}i^n}{i!(m-i)!}$

第二类斯特林数

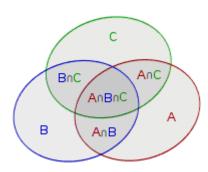
```
};
13
14
15
    struct Stirling2_oneline {
16
17
        std::vector<Z> S;
        int n;
18
        // S[n][k] = \sum_{i=0}^{n} {i^n (-1)^{m-i}}{i! (m-i)!}
19
        // that is convolution of f(x) = x^n / (x!) and g(x) = (-1)^x / (x!)
20
        Stirling2_oneline(int _n): n(_n) {
21
             vector\langle Z \rangle f(n+1), g(n+1), ifact(n+1);
             g[0] = 1;
23
24
             Z fact = 1;
25
             for (int i = 1; i <= n; ++i) fact *= i;</pre>
26
             ifact[n] = fact.inv();
27
             for (int i = n-1; i >= 1; --i) {
28
                 ifact[i] = ifact[i+1] * (i+1);
30
             for (int i = 1; i <= n; ++i) {</pre>
32
                 f[i] = power(Z(i), n) * ifact[i];
33
                 g[i] = ifact[i] * (i % 2? -1: 1);
35
             S = (Poly(f) * Poly(g)).modxk(n+1).a;
37
38
        Z operator [](int idx) {return S[idx];}
39
    };
```

分拆数

将n分成k个部分的分拆,称为k部分拆,记作p(n,k)

递推式: p(n,k) = p(n-1,k-1) + p(n-k,k)

容斥原理



$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |C \cap A| + |A \cap B \cap C|$$

进一步推广:

$$\begin{split} \left|\bigcup_{i=1}^n S_i\right| &= \sum_i |S_i| - \sum_{i < j} |S_i \cap S_j| + \sum_{i < j < k} |S_i \cap S_j \cap S_k| - \cdots \\ &+ (-1)^{m-1} \sum_{a_i < a_{i+1}} \left|\bigcap_{i=1}^m S_{a_i}\right| + \cdots + (-1)^{n-1} |S_1 \cap \cdots \cap S_n| \end{split}$$

经典例子: 不定方程非负整数解

线性代数

线性基

```
struct LinearBase{
        std::vector<int> LBP;
        LinearBase(): LBP(64) {}
        LinearBase(int _n): LBP(_n) {}
4
        void insert(int x) {
            for (int i = (int)LBP.size() - 1; i >= 0; --i) {
                if ((x >> i) & 1) {
                     if (!LBP[i]) {LBP[i] = x; break;}
10
                     else x ^= LBP[i];
11
                else continue;
12
            }
13
14
15
        bool query(int x) {
            for (int i = (int)LBP.size() - 1; i >= 0; --i) {
16
17
                if ((x >> i) & 1) x ^= LBP[i];
            }
18
19
            return (x == 0);
20
        int MaxSum() {
21
22
            int ret = 0;
            for (int i = (int)LBP.size() - 1; i >= 0; --i) {
23
24
                if ((ret >> i) & 1) continue;
                else ret ^= LBP[i];
25
26
27
            return ret;
28
   };
```

博弈论

Sprague-Grundy Theorem (SG 定理), 将多个游戏的 SG 值异或起来就能表示组合起来的胜负状态. 可以归纳法来证明.

图论

LCA

● 倍增

```
void dfs(int u, int fa) {
        pa[u][0] = fa; dep[u] = dep[fa] + 1;
        for (int i = 1; i < SP; ++i) pa[u][i] = pa[pa[u][i - 1]][i - 1];</pre>
        for (int& v: G[u]) {
            if (v == fa) continue;
            dfs(v, u);
        }
   }
    int lca(int u, int v) {
        if (dep[u] < dep[v]) swap(u, v);</pre>
11
        int t = dep[u] - dep[v];
12
        for (int i = 0; i < SP; ++i) if (t & (1 << i)) u = pa[u][i];
13
        for (int i = SP-1; i >= 0; --i) {
14
15
            int uu = pa[u][i], vv = pa[v][i];
            if (uu != vv) { u = uu; v = vv; }
16
        }
17
        return u == v ? u : pa[u][0];
18
   }
19
       • 树链剖分
    namespace HLD{
        const int N = 1e5+3;
2
        std::vector<std::pair<int, int> > g[N]; // 原本的树
3
        int idx[N]; // idx[x] -> 点 x 的 dfs 序号
```

```
int ridx[N];// ridx[id] -> dfs 序为 id 的点号
        int dep[N]; // dep[x] -> 点 x 的深度
6
        int siz[N]; // siz[x] -> 以点 x 为根的树的大小
        int top[N]; // top[x] -> 点 x 所在重链的顶点
        int fa[N]; // fa[x] -> 点 x 的父亲结点
        int son[N]; // son[x] -> 点 x 的重儿子
10
        int clk = 0;// clk -> dfs 序号
11
12
        void init() {
13
14
            for (int i = 0; i <= clk; ++i)</pre>
               g[i].clear();
15
16
            clk = 0;
17
        int lca(int u, int v) {
18
19
            while (top[u] != top[v]) {
                if (dep[top[u]] < dep[top[v]]) std::swap(u, v);</pre>
20
21
                u = fa[top[u]];
22
23
            if (dep[u] > dep[v]) std::swap(u, v);
24
            return u;
25
        void HLDinit(int rt) {
26
            std::function<void(int, int)> predfs1 = [&](int u, int f) {
27
                siz[u] = 1;
                int maxs = son[u] = -1;
29
                for (auto &[to, w]: g[u]) {
30
31
                    if (to == f) continue;
                    fa[to] = u, dep[to] = dep[u] + 1;
32
                    predfs1(to, u);
                    siz[u] += siz[to];
34
                    if (maxs == -1 || siz[to] > siz[maxs]) maxs = to;
35
                }
36
            };
37
            std::function<void(int, int)> predfs2 = [&](int u, int tp) {
                top[u] = tp;
39
                idx[u] = ++clk, ridx[clk] = u;
40
                if (son[u] != -1) {
41
                    predfs2(son[u], tp);
42
43
                for (auto \&[to, w]: g[u]) {
44
45
                    if (to != fa[u] && to != son[u]) predfs2(to, to);
46
47
            };
48
            predfs1(rt, -1);
            predfs2(rt, rt);
49
50
51
        void add(int u, int v, int w = 0) {
            g[u].emplace_back(v, w);
53
54
            g[v].emplace_back(u, w);
55
   }
56
    连通性
    tarjan
    const int N = 1e5 + 3;
   std::vector<int> g[N];
    int dfn[N], dfnx = 0; // dfs 序
    int low[N]; // 最高点
    int col[N], colx = 0; // 连通块标号
    int stk[N], stktp = 0; // 当前链
    bool instk[N];
    void tarjan(int u, int f) {
        dfn[u] = low[u] = ++dfnx;
        stk[++stktp] = u;
10
11
        for (auto to: g[u]) {
            // if (to == f) continue;
12
            if (dfn[to]) {if (instk[to]) low[u] = std::min(dfn[to], low[u]);}
13
            else tarjan(to, u), low[u] = std::min(low[u], low[to]);
14
```

```
15
16
        if (low[u] == dfn[u]) {
17
            ++colx;
            while (stk[stktp] != u) instk[stk[stktp]] = 0, col[stk[stktp--]] = colx;
18
            instk[stk[stktp]] = 0, col[stk[stktp--]] = colx;
        }
20
   }
21
    树
    虚树
          用栈处理形式上与笛卡尔树类似,按顺序后处理右链部分
   namespace VirtualTree{
1
        const int N = 3e5+3;
2
        std::vector<std::pair<int, int> > g[N]; // 原本的树
        int idx[N]; // idx[x] -> 点 x 的 dfs 序号
        int ridx[N];// ridx[id] -> dfs 序为 id 的点号
        int dep[N]; // dep[x] -> 点 x 的深度
        int siz[N]; // siz[x] \rightarrow 以点 x 为根的树的大小
        int top[N]; // top[x] -> 点 x 所在重链的顶点
        int fa[N]; // fa[x] -> 点 x 的父亲结点
        int son[N]; // son[x] -> 点 x 的重儿子
10
        int clk = 0;// clk -> dfs 序号
11
12
        void init() {
13
            for (int i = 0; i <= clk; ++i)</pre>
14
                g[i].clear();
15
            clk = 0;
16
17
        int lca(int u, int v) {
18
            while (top[u] != top[v]) {
                if (dep[top[u]] < dep[top[v]]) std::swap(u, v);</pre>
20
                u = fa[top[u]];
21
22
            if (dep[u] > dep[v]) std::swap(u, v);
23
24
            return u;
25
26
        void HLDinit(int rt) {
            std::function<void(int, int)> predfs1 = [&](int u, int f) {
27
                siz[u] = 1;
28
                int maxs = son[u] = -1;
                for (auto \&[to, w]: g[u]) {
30
                    if (to == f) continue;
31
                    fa[to] = u, dep[to] = dep[u] + 1;
32
                    predfs1(to, u);
33
34
                    siz[u] += siz[to];
                    if (maxs == -1 || siz[to] > siz[maxs]) maxs = to;
35
37
            };
            std::function<void(int, int)> predfs2 = [&](int u, int tp) {
38
39
                top[u] = tp;
                idx[u] = ++clk, ridx[clk] = u;
40
41
                if (son[u] != -1) {
                    predfs2(son[u], tp);
42
43
                for (auto \&[to, w]: g[u]) {
44
45
                    if (to != fa[u] && to != son[u]) predfs2(to, to);
46
            };
47
            predfs1(rt, -1);
            predfs2(rt, rt);
49
50
51
        int stk[N]; // 栈
52
53
        bool isl[N];
        std::vector<int> vg[N]; // 虚树
54
        std::vector<int> lst; // 当前虚树中的点
55
        void buildvt(std::vector<int> pt, int rt = 0) {
56
            for (auto x: lst) vg[x].clear(), isl[x] = 0;
```

```
lst.clear();
58
59
             sort(pt.begin(), pt.end(), [](int u, int v) {return idx[u] < idx[v];});</pre>
60
61
             int tp = 0;
             stk[++tp] = rt; lst.push_back(rt);
62
             for (auto u: pt) {
63
                 if (u != rt) {
64
                     int ll = lca(u, stk[tp]);
65
                     if (ll != stk[tp]) {
66
67
                          while (idx[ll] < idx[stk[tp-1]])</pre>
                              vg[stk[tp-1]].push\_back(stk[tp]), --tp;
68
                          vg[ll].push_back(stk[tp]);
70
                          if (idx[ll] > idx[stk[tp-1]]) stk[tp] = ll, lst.push_back(ll);
                          else --tp;
71
72
                     }
                     stk[++tp] = u, lst.push_back(u);
73
74
                 isl[u] = 1;
75
            while (tp > 1) {
77
                 vg[stk[tp-1]].push_back(stk[tp]);
78
79
80
             }
        }
82
83
        void add(int u, int v, int w = 0) {
84
             g[u].emplace_back(v, w);
            g[v].emplace_back(u, w);
85
    }
87
    网络流
    Dinic
    class Dinic{
1
        \textbf{struct } \; \textbf{E} \{
             int to, cp;
3
             E(int t, int c): to(t), cp(c) {}
        };
5
        int n, m, s, t;
        std::vector<E> edges;
        std::vector<std::vector<int> > G;
        int *d, *cur;
10
11
        bool BFS() {
12
             memset(d, 0, sizeof(int) * (n));
13
14
             std::queue<int> Q;
            Q.push(s); d[s] = 1;
15
16
             while (!Q.empty()) {
                 int x = Q.front(); Q.pop();
17
                 for (auto &i: G[x]) {
18
19
                     E &e = edges[i];
                     if (!d[e.to] && e.cp > 0) {
20
                          d[e.to] = d[x] + 1;
21
                          Q.push(e.to);
22
                     }
23
                 }
24
25
26
             return d[t];
27
28
        int DFS(int u, int cp) {
29
30
             if (u == t || !cp) return cp;
             int tmp = cp, f;
31
             for (int& i = cur[u]; i < G[u].size(); ++i) {</pre>
32
33
                 E& e = edges[G[u][i]];
                 if (d[u] + 1 == d[e.to]) {
34
                     f = DFS(e.to, std::min(cp, e.cp));
35
                     e.cp -= f;
```

```
edges[G[u][i] ^ 1].cp += f;
37
38
                       cp -= f;
                       if (!cp) break;
39
                  }
40
             }
             return tmp - cp;
42
43
44
    public:
45
         Dinic(\textbf{int} \ nn): \ n(nn), \ G(nn), \ m(0) \ \{d = \textbf{new int}[nn], \ cur = \textbf{new int}[nn]; \}
46
         ~Dinic() {delete[] d; delete[] cur;}
47
48
         void add(int u, int v, int cap) {
49
             edges.emplace_back(v, cap);
50
51
             edges.emplace_back(u, 0);
             G[u].push_back(m++);
52
53
             G[v].push_back(m++);
54
         int go(int ss, int tt) {
56
57
             s = ss, t = tt;
58
             int flow = 0;
             while (BFS()) {
59
                  memset(cur, 0, sizeof(int) * n);
                  flow += DFS(s, INF);
61
62
63
             return flow;
         }
64
    };
```

上下界只需要在建图的时候把上下界之差的出入不平衡的调整一下

最大流 == 最小割

计算几何

二维几何: 点与向量

```
const double EPS = 1e-9;
    inline int sign(double a) { return a < -EPS ? -1 : a > EPS; }
    inline int cmp(double a, double b){ return sign(a-b); }
5
    struct P {
        double x, y;
        P() {}
        P(\textbf{double } \_x, \ \textbf{double } \_y) \ : \ x(\_x), \ y(\_y) \ \{\}
        P operator+(P p) { return {x + p.x, y + p.y}; }
10
11
        P operator-(P p) { return \{x - p.x, y - p.y\}; \}
        P operator*(double d) { return \{x \, * \, d, \, y \, * \, d\}; }
12
         P operator/(double d) { return {x / d, y / d}; }
13
14
15
         bool operator<(P p) const {</pre>
16
             int c = cmp(x, p.x);
             if (c) return c == -1;
17
18
             return cmp(y, p.y) == -1;
        }
19
20
         bool operator==(P o) const{
21
             return cmp(x,o.x) == 0 \&\& cmp(y,o.y) == 0;
22
23
24
         double dot(P p) { return x * p.x + y * p.y; }
25
         double det(P p) { return x * p.y - y * p.x; }
26
27
28
         double distTo(P p) { return (*this-p).abs(); }
         double alpha() { return atan2(y, x); }
29
         void read() { std::cin >> x >>y; }
         void write() {std::cout << "(" << x << "," << y << ")" << std::endl;}</pre>
31
         double abs() { return sqrt(abs2());}
```

```
double abs2() { return x * x + y * y; }
33
34
        P rot90() { return P(-y,x);}
        P unit() { return *this/abs(); }
35
        int quad() const { return sign(y) == 1 \mid \mid (sign(y) == 0 \&\& sign(x) >= 0); }
36
37
        P rot(double an) { return \{x*\cos(an)-y*\sin(an),x*\sin(an) + y*\cos(an)\}; }
   };
38
39
    struct L{ //ps[0] -> ps[1]
40
        P ps[2];
41
42
        P& operator[](int i) { return ps[i]; }
        P dir() { return ps[1] - ps[0]; }
43
44
        L (P a, P b) {
            ps[0]=a;
45
            ps[1]=b;
46
47
        bool include(P p) \{ return  sign((ps[1] - ps[0]).det(p - ps[0])) > 0; \}
48
        L push(){ // push eps outward
            const double eps = 1e-8;
50
            P delta = (ps[1] - ps[0]).rot90().unit() * eps;
            return {ps[0] + delta, ps[1] + delta};
52
53
        }
   };
55
    #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
    字符串
    KMP
          可以尝试从 fail 树的角度理解
    bool kmp(string &s, string &p) {
        int n = p.size();
        vector<int> nex(n + 1);
        for (int i = 1, l = 0; i < n; ++i) {</pre>
            while (| && p[|] != p[i]) {| = nex[|];}
            if (p[l] != p[i]) nex[i] = 0;
            else nex[i] = ++l;
        n = s.size();
        for (int i = 0, now = 0; i < n; ++i) {
            while (now \&\& s[i] != p[now]) now = nex[now - 1];
            if (s[i] == p[now]) ++now;
12
            if (now == p.size()) return true;
        }
14
        return false;
   }
16
    后缀自动机
    struct SuffixAutomaton {
1
        static constexpr int ALPHABET_SIZE = 26;
        int N = 1e5;
3
        struct Node {
            int len;
            int link;
            // int next[ALPHABET_SIZE];
            std::vector<int> next;
            Node() : len(0), link(0), next(ALPHABET_SIZE) {}
        };
10
11
        std::vector<Node> t;
        int cntNodes;
12
13
        int extend(int p, int c) {
14
            if (t[p].next[c]) {
                int q = t[p].next[c];
15
                if (t[q].len == t[p].len + 1)
16
17
                    return q;
                int r = ++cntNodes;
18
```

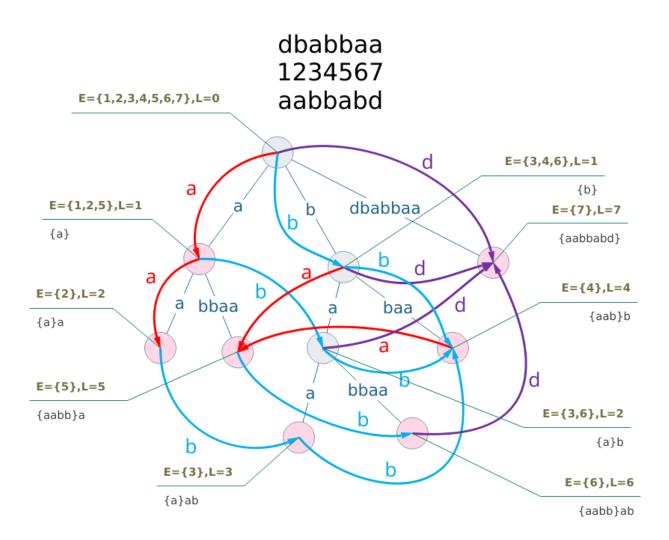


Figure 1: 后缀自动机图解

```
t[r].len = t[p].len + 1;
19
20
                t[r].link = t[q].link;
                // std::copy(t[q].next, t[q].next + ALPHABET_SIZE, t[r].next);
21
                t[r].next = t[q].next;
22
                t[q].link = r;
                while (t[p].next[c] == q) {
24
                     t[p].next[c] = r;
25
                     p = t[p].link;
26
                }
27
28
                return r;
29
            int cur = ++cntNodes;
            t[cur].len = t[p].len + 1;
31
            while (!t[p].next[c]) {
32
                t[p].next[c] = cur;
33
                p = t[p].link;
34
35
            t[cur].link = extend(p, c);
36
37
            return cur;
        }
38
        SuffixAutomaton(int N_{-}): t(2 * N_{-}), N(N_{-}) {
39
40
            cntNodes = 1;
            // std::fill(t[0].next, t[0].next + ALPHABET_SIZE, 1);
41
42
            t[0].next.assign(ALPHABET_SIZE, 1);
43
            t[0].len = -1;
44
        SuffixAutomaton(string s): t(2 * s.size()), N(s.size()) {
45
            cntNodes = 1;
46
47
            t[0].next.assign(ALPHABET_SIZE, 1);
            t[0].len = -1;
48
49
            int p = 1;
50
51
            for (auto ch: s) {
52
                p = extend(p, ch-'a');
            }
53
54
    };
55
    杂项
    STL
       copy
    template <class InputIterator, class OutputIterator>
      OutputIterator copy (InputIterator first, InputIterator last, OutputIterator result);
```