

Chapter 4 - Solutions

Arithmetic Sum:

$$\begin{aligned} \text{first element: } \sum_{i=1}^1 i &= \frac{(1)(2)}{2} = 1, & \text{assume: } \sum_{i=1}^{n-1} i &= \frac{(n-1)(n)}{2} \\ \Rightarrow \sum_{i=1}^n i &= n + \sum_{i=1}^{n-1} i = n + \frac{(n-1)(n)}{2} = \frac{n^2 + n}{2} = \frac{(n)(n+1)}{2} \end{aligned}$$

Geometric Sum:

$$\begin{aligned} S_n &= \sum_{i=0}^n r^i = 1 + r + \dots + r^n \\ rS_n &= r + r^2 + \dots + r^{n+1} \Rightarrow S_n - rS_n = 1 - r^{n+1} \\ S_n &= \frac{1 - r^{n+1}}{1 - r} \end{aligned}$$

Reinforcement

$$4.3. \quad 2n^3 \geq 40n^2 \Rightarrow n \geq 20$$

$$4.8. \quad 2^{10}, 2^{\log n}, 3n + 100 \log n, 4n, n \log n, 4n \log n + 2n, n^2 + 10n, n^3, 2^n$$

$$4.9. \quad O(n) \quad 4.10. \quad O(n) \quad 4.11. \quad O(n^2) \quad 4.12. \quad O(n) \quad 4.13. \quad O(n^3)$$

4.16.

$$\begin{aligned} d(n) \in O(f(n)) &\Rightarrow d(n) \leq c_1 f(n), \quad \forall n \geq n_1 \\ e(n) \in O(g(n)) &\Rightarrow e(n) \leq c_2 g(n), \quad \forall n \geq n_2 \\ c_0 &= \max(c_1, c_2), \quad n_0 = \max(n_1, n_2) \end{aligned}$$

$$\begin{aligned} d(n) + e(n) &\leq c_1 f(n) + c_2 g(n) \leq c_0 f(n) + c_0 g(n) = c_0 (f(n) + g(n)), \quad \forall n \geq n_0 \\ &\Rightarrow d(n) + e(n) \in O(f(n) + g(n)) \end{aligned}$$

Creativity

$$4.45. \quad \text{Sum for all elements in the range } [0, n-1] \Rightarrow \text{sum} = \sum_{i=0}^{n-1} i = \frac{(n-1)n}{2}$$

public static int find_missing_number (**int**[] arr)

```
{
    int arr_total = 0, n=arr.length;
    int sum = n*(n-1)/2;
    for ( int i = 0; i < n; i++)
        arr_total += arr[i];
    return (sum - arr_total);
}
```