

Automatic Exploit Generation

Given a program, find bugs and generate exploits

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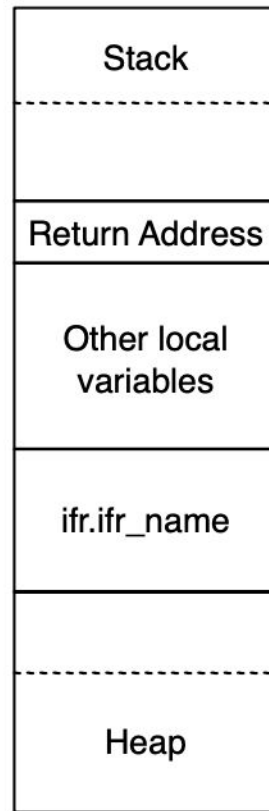
Agenda

- WarmUp
- Demo
- Overview of AEG
- Challenges
- Formal Verification
- Design
- Evaluation
- Discussion
- Conclusion

WarmUp

```
1 int main(int argc, char **argv) {
2     int skfd;           /* generic raw socket desc.    */
3     if(argc == 2)
4         print_info(skfd, argv[1], NULL, 0);
5     ...
6 static int print_info(int skfd, char *ifname, char *args[], int count)
7     {
8         struct wireless_info info;
9         int rc;
10        rc = get_info(skfd, ifname, &info);
11    ...
12 static int get_info(int skfd, char *ifname, struct wireless_info * info
13    ) {
14        struct iwreq wrq;
15        if(iw_get_ext(skfd, ifname, SIOCGIWNAME, &wrq) < 0) {
16            struct ifreq ifr;
17            strcpy(ifr.ifr_name, ifname); /* buffer overflow */
18        ...
19    }
```

Code snippet from iwconfig

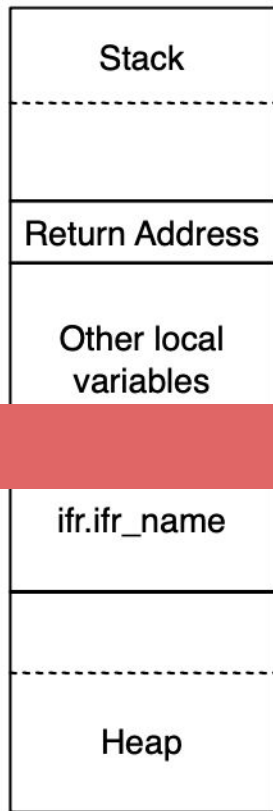


Memory

WarmUp

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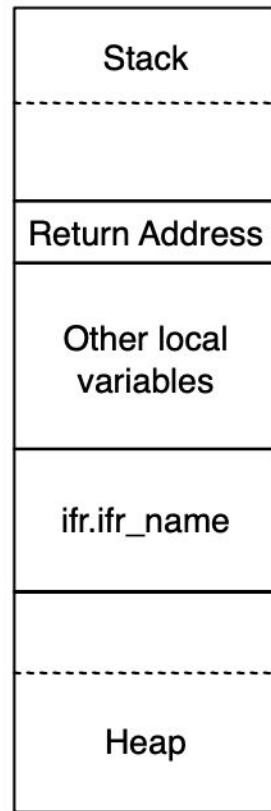


Memory

WarmUp

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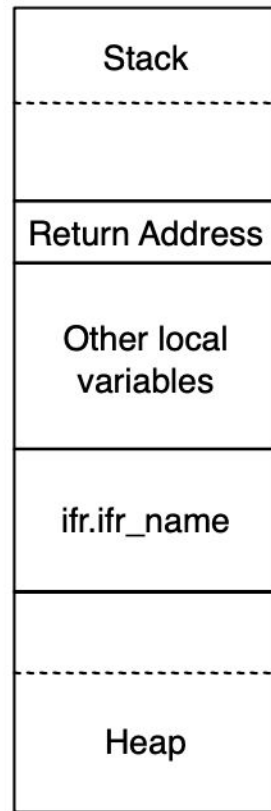


Memory

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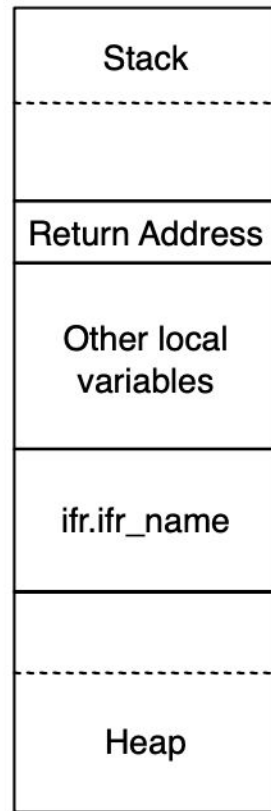


Memory

WarmUp

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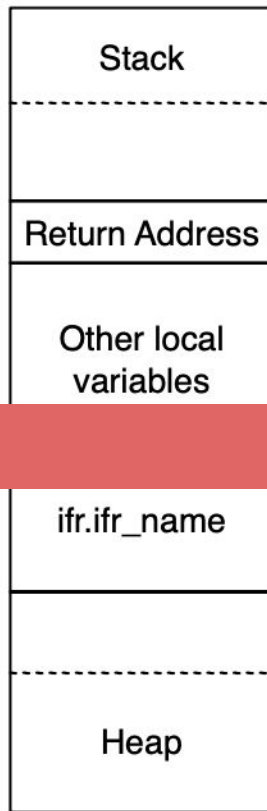


Memory

WarmUp

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```

Note that this `ifr.ifr_name` is a 32 bytes array



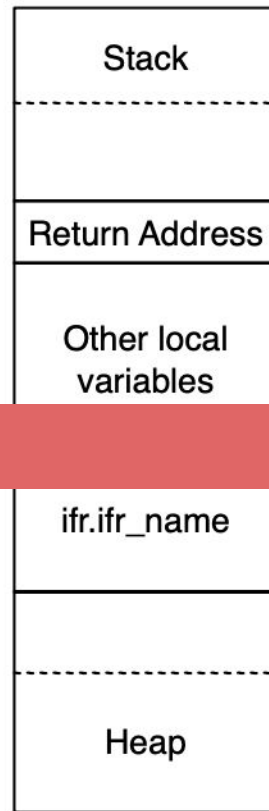
Memory

Code snippet from iwconfig

WarmUp

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```

What if we put a significant a long input to this program?



Memory

Code snippet from iwconfig

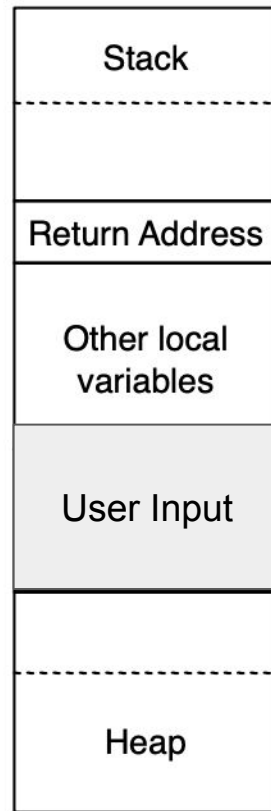
Trigger buffer overflow !

How would you exploit it ?

WarmUp

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Code snippet from iwconfig

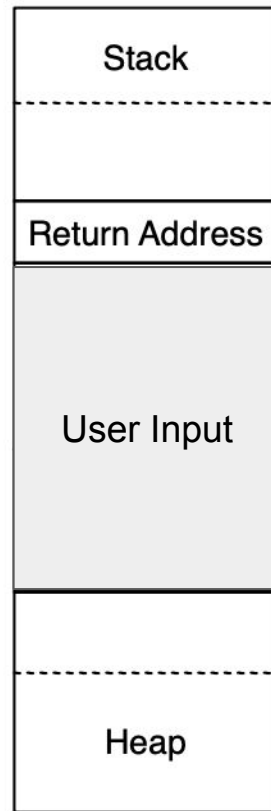


Memory

WarmUp

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Code snippet from iwconfig

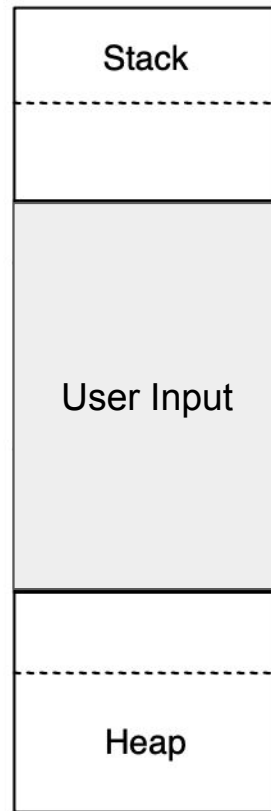


Memory

WarmUp

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20 }
```

Code snippet from iwconfig



Memory

Execve bin/sh !

Demo !


```
root
ls
lsakefile
access.log
aeg.sh
aeg_#-data
aeg_#-data-stat
aeg_stdin
aeg_stdin-stat
error.log
filesize
klee-last
klee-out-#
portno
recvinfo
runtime_info
server.conf
serverd
serverd.bc
stype
tapfile
[]
```

**AEG automatically got a root shell
exploit in under a second !**

What is an Exploit ?

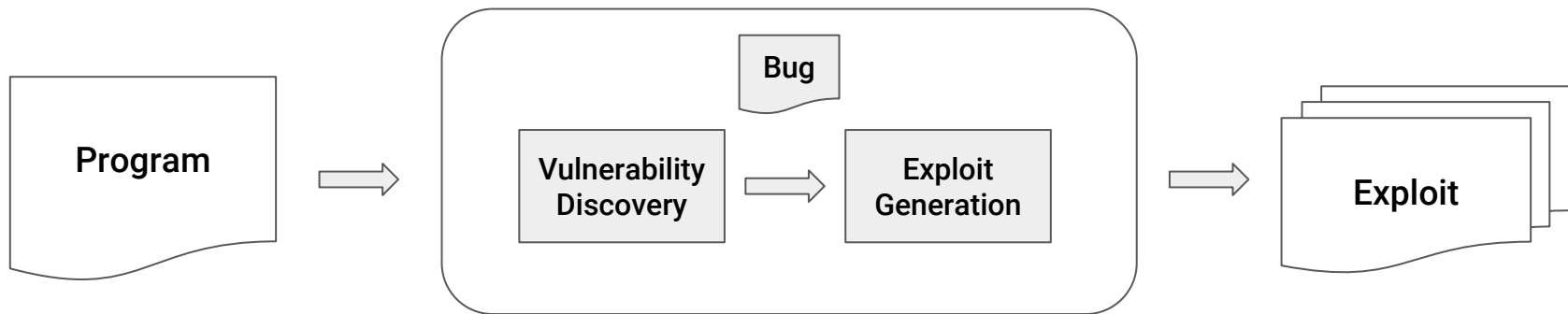
Hacking Community

Control Flow Hijack

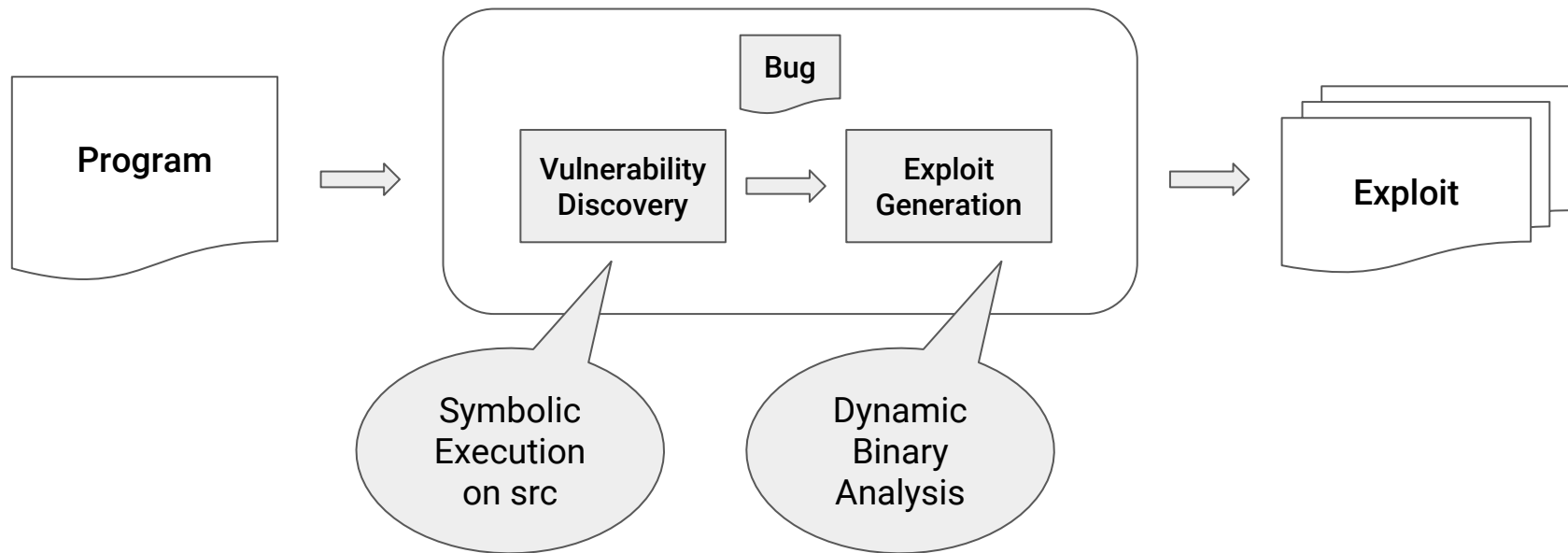
Overview



Overview



Overview



Challenges

- State Space Explosion
- Path Selection
- Environment Modelling
- Mixed Analysis challenge
- Exploit Verification

Challenges

- State Space Explosion
 - Challenge: There are possible infinite paths that AEG needs to explore
 - Solution: Preconditioned symbolic execution
- Path Selection
- Environment Modelling
- Mixed Analysis challenge
- Exploit Verification

Challenges

- State Space Explosion
- Path Selection
 - Challenge: AEG must select the path which should be explore firstly
 - Solution: Path prioritization technique
- Environment Modelling
- Mixed Analysis challenge
- Exploit Verification

Challenges

- State Space Explosion
- Path Selection
- Environment Modelling
 - Challenge: Make accurate analysis
 - Solution: Model environment IO behavior
- Mixed Analysis challenge
- Exploit Verification

Challenges

- State Space Explosion
- Path Selection
- Environment Modelling
- Mixed Analysis challenge
 - Challenge: Perform a mix of binary and source level analysis to scale large program
 - Solution: Combine DBA and SRC analysis
- Exploit Verification

Challenges

- State Space Explosion
- Path Selection
- Environment Modelling
- Mixed Analysis challenge
- Exploit Verification
 - Challenge: Is it success ?
 - Solution: Execute `'/bin/sh'`

How can we describe problems?

**Treat it as a formal verification
problem !**

Formal Verification

Formal verification is the act of proving or disproving the correctness of intended algorithms underlying a system with respect to a certain formal specification or property, using formal methods of mathematics. - Wikipedia

Typically, in security field, we would use it prove the system whether it is safe.

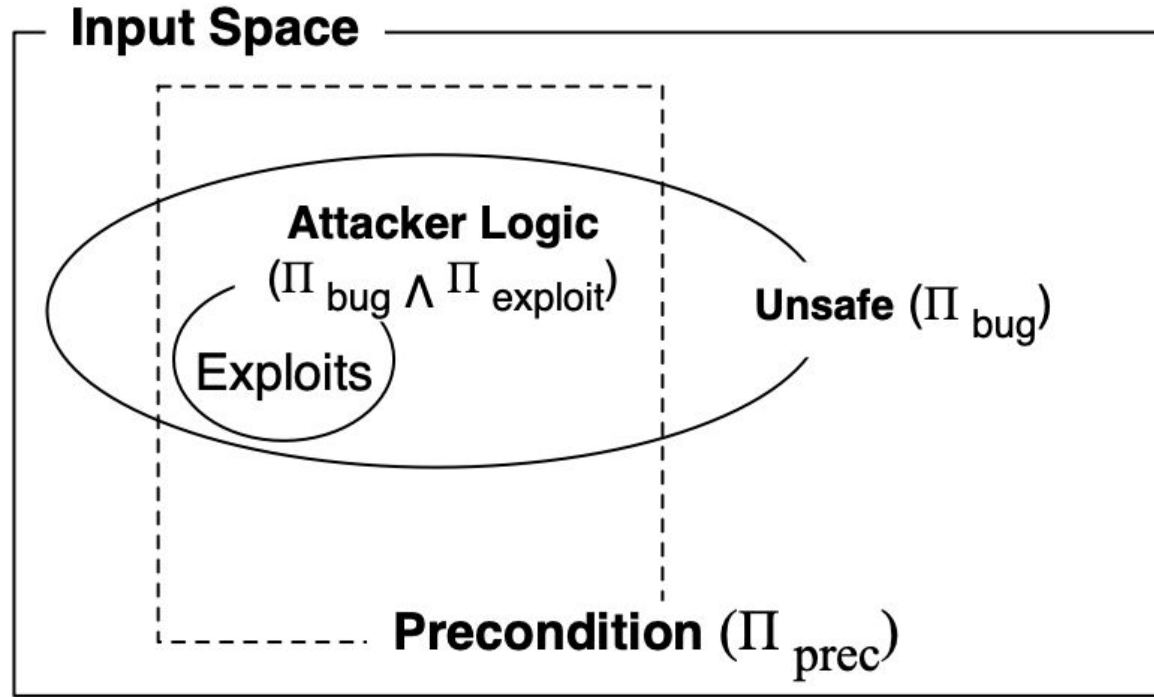
**Verify the program whether it is
exploitable**

Formal Verification

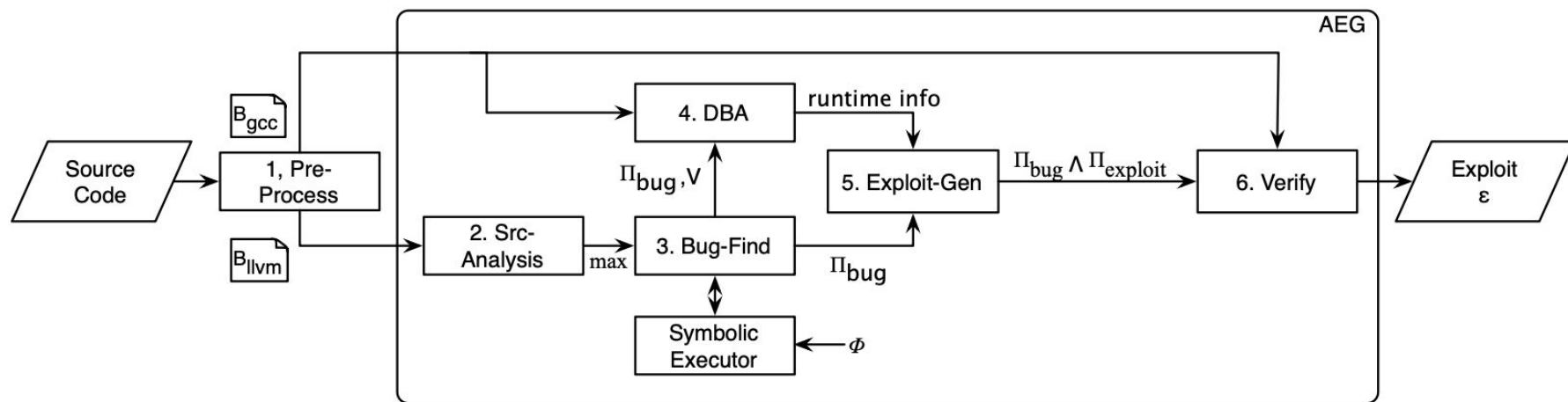
- Predicates

- Unsafe Path Predicate (Π_{bug})
 - Represents a path predicate of an execution that violates safe property φ
- Exploit Predicate (Π_{exploit})
 - Represents attacker's logic
- Preconditioned symbolic parameter (Π_{prec})
 - Represents the narrow state

Formal Verification



Design



Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$

Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
 - A compiled binary $Bgcc$ give it to AEG to generate exploits
 - A LLVM bytecode $Bllvm$ will be used in bug finding infrastructure
- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
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Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
- Src-Analysis: $Bllvm \rightarrow max$
 - Analyze source code to generate max size of input length
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$

Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
 - Π_{bug} contains path predicates
 - V contains source-level information about detected vulnerabilities
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$

Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
 - Analyze *Bgcc* with a concrete buggy input and output runtime Information
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$

Design

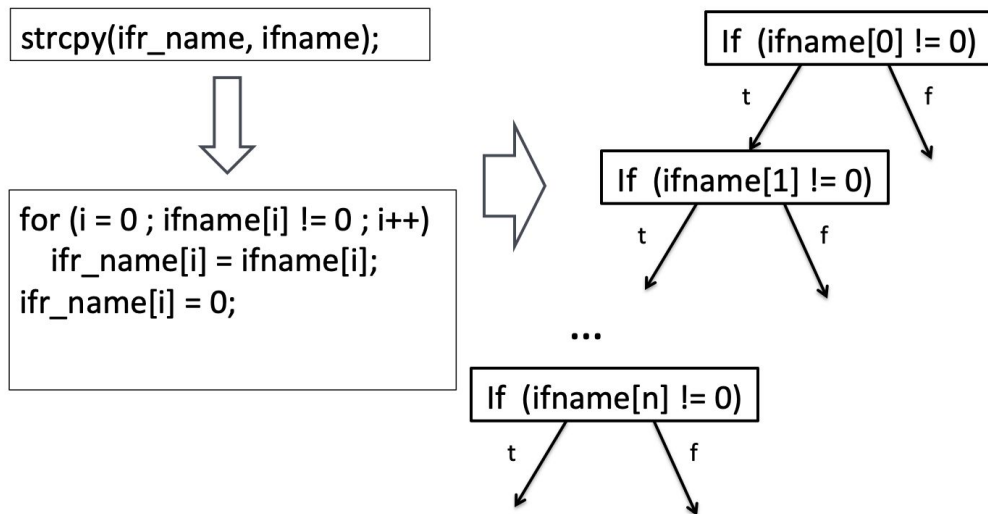
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- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
 - Constructs a formula for a exploit
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$

Design

- Pre-Process: $\text{src} \rightarrow (Bgcc, Bllvm)$
- Src-Analysis: $Bllvm \rightarrow max$
- Bug-Find: $(Bllvm, \varphi, max) \rightarrow (\Pi_{bug}, V)$
- DBA: $(Bgcc, (\Pi_{bug}, V)) \rightarrow R$
- Exploit-Gen: $(\Pi_{bug}, R) \rightarrow \Pi_{bug} \wedge \Pi_{exploit}$
- Verify: $(Bgcc, \Pi_{bug} \wedge \Pi_{exploit}) \rightarrow \{\epsilon, \perp\}$
 - If there is a solution then outputs ϵ

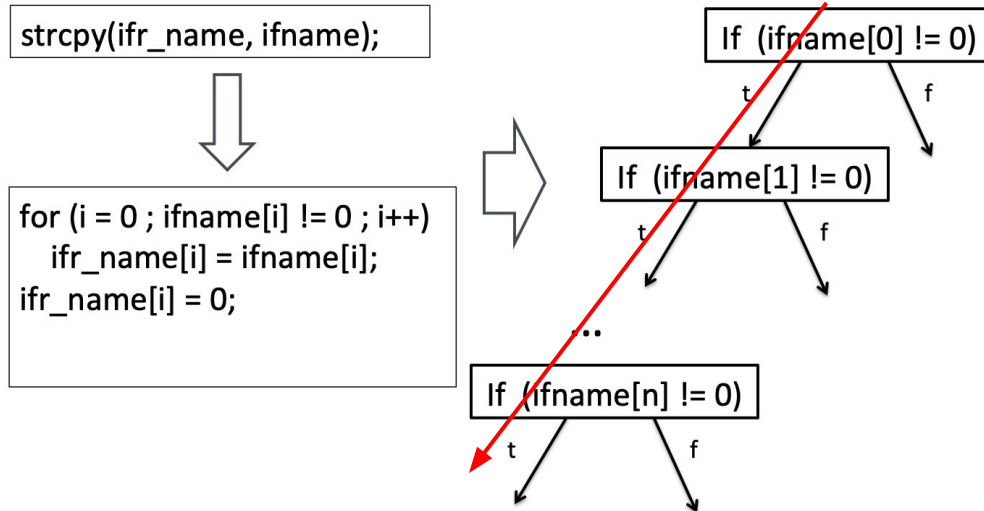
Design (Bug-Find)

- Traditional symbolic execution
 - Try to cover all paths
 - Slow to find exploitable bugs



Design (Bug-Find)

- Preconditioned Symbolic Execution
 - Decrease the state space
 - Find bugs efficiently



Design (Bug-Find)

- Preconditions
 - None
 - No precondition
 - Known Length
 - Max length of input
 - Known Prefix
 - Constrains the input
 - Ex. PNG utility
 - Concolic Execution
 - Specify all inputs have a specific value

**Not all paths are likely to be
exploitable**

Design (Bug-Find)

- Path Prioritization
 - Which paths should explore first ?
- Buggy-Path-First
 - Paths contain bugs are more likely to be exploitable
 - Prioritize buggy paths higher
- Loop Exhaustion
 - Run a high priority interpreter as many times as possible

Design (DBA)

- Dynamic Binary Analysis
 - Test exploitability of buggy path
- Inputs
 - *Bgcc*
 - Path constraints
 - Names of vuln functions and buffers
- Output
 - Address to overwrite
 - Starting address AEG write to
 - Additional constraints that describe stack memory

Design (Exploit-Gen)

- Generate an exploit
 - Determines which class of attacks
 - Return-to-stack Exploit
 - Return-to-libc Exploit
 - Shellcode
 - Size
 - Position of shellcode

Design (Verify)

- Checks the shell whether it has been spawned
 - Success: A shell
 - Failure: Nothing

Evaluation

- Take popular advisories
 - CVE
 - OSVDB
 - EDB

Evaluation

	Program	Ver.	Exploit Type	Vulnerable Input src	Gen. Time (sec.)	Executable Lines of Code	Advisory ID.
None	aeon	0.2a	Local Stack	Env. Var.	3.8	3392	CVE-2005-1019
	iwconfig	V.26	Local Stack	Arguments	1.5	11314	CVE-2003-0947
	glftpd	1.24	Local Stack	Arguments	2.3	6893	OSVDB-ID#16373
	ncompress	4.2.4	Local Stack	Arguments	12.3	3198	CVE-2001-1413
Length	htget (processURL)	0.93	Local Stack	Arguments	57.2	3832	CVE-2004-0852
	htget (HOME)	0.93	Local Stack	Env. Var	1.2	3832	Zero-day
	expect (DOTDIR)	5.43	Local Stack	Env. Var	187.6	458404	Zero-day
	expect (HOME)	5.43	Local Stack	Env. Var	186.7	458404	OSVDB-ID#60979
	socat	1.4	Local Format	Arguments	3.2	35799	CVE-2004-1484
	tipxd	1.1.1	Local Format	Arguments	1.5	7244	OSVDB-ID#12346
Prefix	aspell	0.50.5	Local Stack	Local File	15.2	550	CVE-2004-0548
	exim	4.41	Local Stack	Arguments	33.8	241856	EDB-ID#796
	xserver	0.1a	Remote Stack	Sockets	31.9	1077	CVE-2007-3957
	rsync	2.5.7	Local Stack	Env. Var	19.7	67744	CVE-2004-2093
	xmail	1.21	Local Stack	Local File	1276.0	1766	CVE-2005-2943
Concolic	corehttp	0.5.3	Remote Stack	Sockets	83.6	4873	CVE-2007-4060
Average Generation Time & Executable Lines of Code					114.6	56784	

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	tipxd	1.1.1	Local Format	Arguments	1.5	7244	OSVDB-ID#12346
Prefix	aspell	0.50.5	Local Stack	Local File	15.2	550	CVE-2004-0548
	exim	4.41	Local Stack	Arguments	33.8	241856	EDB-ID#796
	xserver	0.1a	Remote Stack	Sockets	31.9	1077	CVE-2007-3957
	rsync	2.5.7	Local Stack	Env. Var	19.7	67744	CVE-2004-2093
	xmail	1.21	Local Stack	Local File	1276.0	1766	CVE-2005-2943
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Format String

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Buffer
overflow

Evaluation

	Program	Ver.	Exploit Type	Vulnerable Input src	Gen. Time (sec.)	Executable Lines of Code	Advisory ID.
None	aeon	0.2a	Local Stack	Env. Var.	3.8	3392	CVE-2005-1019
	iwconfig	V.26	Local Stack	Arguments	1.5	11314	CVE-2003-0947
	glftpd	1.24	Local Stack	Arguments	2.3	6893	OSVDB-ID#16373
	ncompress	4.2.4	Local Stack	Arguments	12.3	3198	CVE-2001-1413
Length	htget (processURL)	0.93	Local Stack	Arguments	57.2	3832	CVE-2004-0852
	htget (HOME)	0.93	Local Stack	Env. Var	1.2	3832	Zero-day
	expect (DOTDIR)	5.43	Local Stack	Env. Var	187.6	458404	Zero-day
	expect (HOME)	5.43	Local Stack	Env. Var	186.7	458404	OSVDB-ID#60979
	socat	1.4	Local Format	Arguments	3.2	35799	CVE-2004-1484
	tipxd	1.1.1	Local Format	Arguments	1.5	7244	OSVDB-ID#12346
Prefix	aspell	0.50.5	Local Stack	Local File	15.2	550	CVE-2004-0548
	exim	4.41	Local Stack	Arguments	33.8	241856	EDB-ID#796
	xserver	0.1a	Remote Stack	Sockets	31.9	1077	CVE-2007-3957
	rsync	2.5.7	Local Stack	Env. Var	19.7	67744	CVE-2004-2093
	xmail	1.21	Local Stack	Local File	1276.0	1766	CVE-2005-2943
Concolic	corehttp	0.5.3	Remote Stack	Sockets	83.6	4873	CVE-2007-4060
Average Generation Time & Executable Lines of Code					114.6	56784	

Remote
Attack

Evaluation

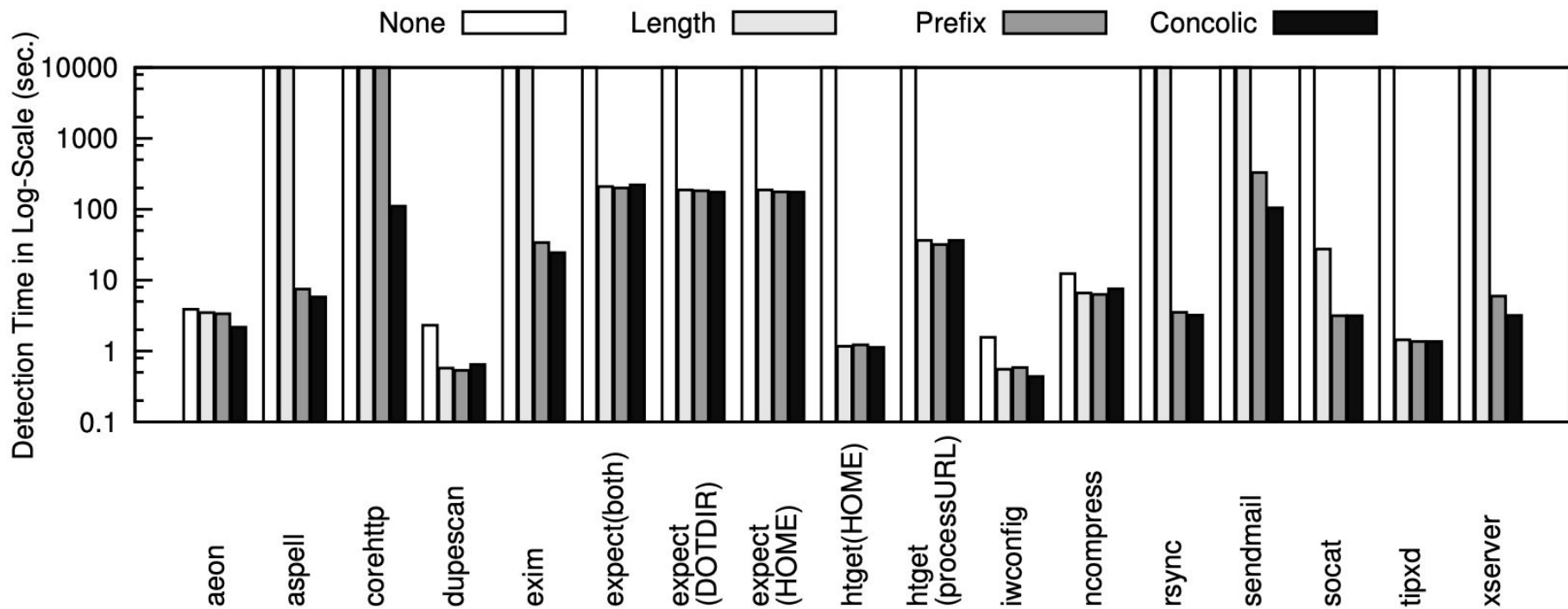
	Program	Ver.	Exploit Type	Vulnerable Input src	Gen. Time (sec.)	Executable Lines of Code	Advisory ID.
None	aeon	0.2a	Local Stack	Env. Var.	3.8	3392	CVE-2005-1019
	iwconfig	V.26	Local Stack	Arguments	1.5	11314	CVE-2003-0947
	glftpd	1.24	Local Stack	Arguments	2.3	6893	OSVDB-ID#16373
	ncompress	4.2.4	Local Stack	Arguments	12.3	3198	CVE-2001-1413
Length	htget (processURL)	0.93	Local Stack	Arguments	57.2	3832	CVE-2004-0852
	htget (HOME)	0.93	Local Stack	Env. Var	1.2	3832	Zero-day
	expect (DOTDIR)	5.43	Local Stack	Env. Var	187.6	458404	Zero-day
	expect (HOME)	5.43	Local Stack	Env. Var	186.7	458404	OSVDB-ID#60979
	socat	1.4	Local Format	Arguments	3.2	35799	CVE-2004-1484
	tipxd	1.1.1	Local Format	Arguments	1.5	7244	OSVDB-ID#12346
D.C	aspell	0.50.5	Local Stack	Local File	15.2	550	CVE-2004-0548
	exim	4.41	Local Stack	Arguments	33.8	241856	EDB-ID#796
	rsync	2.5.7	Local Stack	Env. Var	19.7	67744	CVE-2004-2093
	xmail	1.21	Local Stack	Local File	1276.0	1766	CVE-2005-2943
Average Generation Time & Executable Lines of Code					114.6	56784	

Local Attack

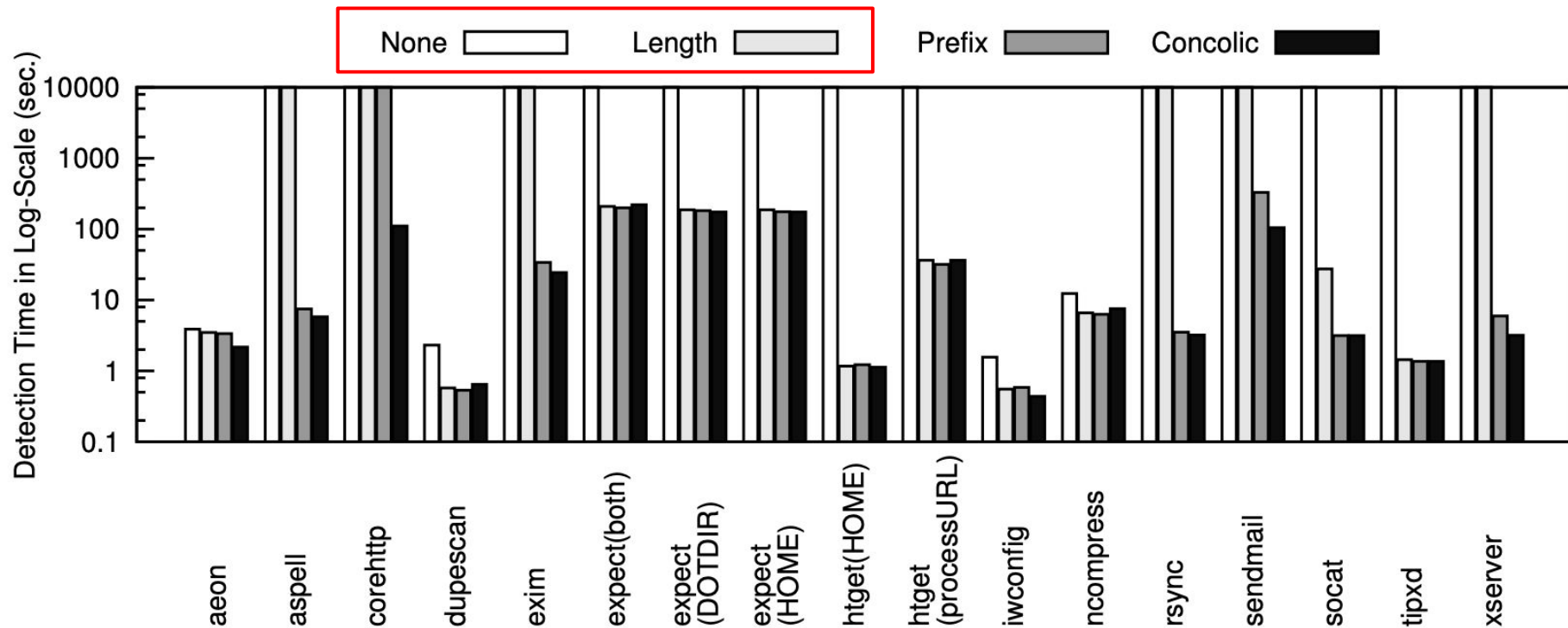
Evaluation

	Program	Ver.	Exploit Type	Vulnerable Input src	Gen. Time (sec.)	Executable Lines of Code	Advisory ID.
None	aeon	0.2a	Local Stack	Env. Var.	3.8	3392	CVE-2005-1019
	iwconfig	V.26	Local Stack	Arguments	1.5	11314	CVE-2003-0947
	glftpd	1.24	Local Stack	Arguments	2.3	6893	OSVDB-ID#16373
	ncompress	4.2.4	Local Stack	Arguments	12.3	3198	CVE-2001-1413
Length	htget (processURL)	0.93	Local Stack	Arguments	57.2	3832	CVE-2004-0852
	htget (HOME)	0.93	Local Stack	Env. Var	1.2	3832	Zero-day
	expect (DOTDIR)	5.43	Local Stack	Env. Var	187.6	458404	Zero-day
	expect (HOME)	5.43	Local Stack	Env. Var	186.7	458404	OSVDB-ID#60979
	socat	1.4	Local Format	Arguments	3.2	35799	CVE-2004-1484
	tipxd	1.1.1	Local Format	Arguments	1.5	7244	OSVDB-ID#12346
Prefix	aspell	0.50.5	Local Stack	Local File	15.2	550	CVE-2004-0548
	exim	4.41	Local Stack	Arguments	33.8	241856	EDB-ID#796
	xserver	0.1a	Remote Stack	Sockets	31.9	1077	CVE-2007-3957
	rsync	2.5.7	Local Stack	Env. Var	19.7	67744	CVE-2004-2093
	xmail	1.21	Local Stack	Local File	1276.0	1766	CVE-2005-2943
Concolic	corehttp	0.5.3	Remote Stack	Sockets	83.6	4873	CVE-2007-4060
Average Generation Time & Executable Lines of Code					114.6	56784	

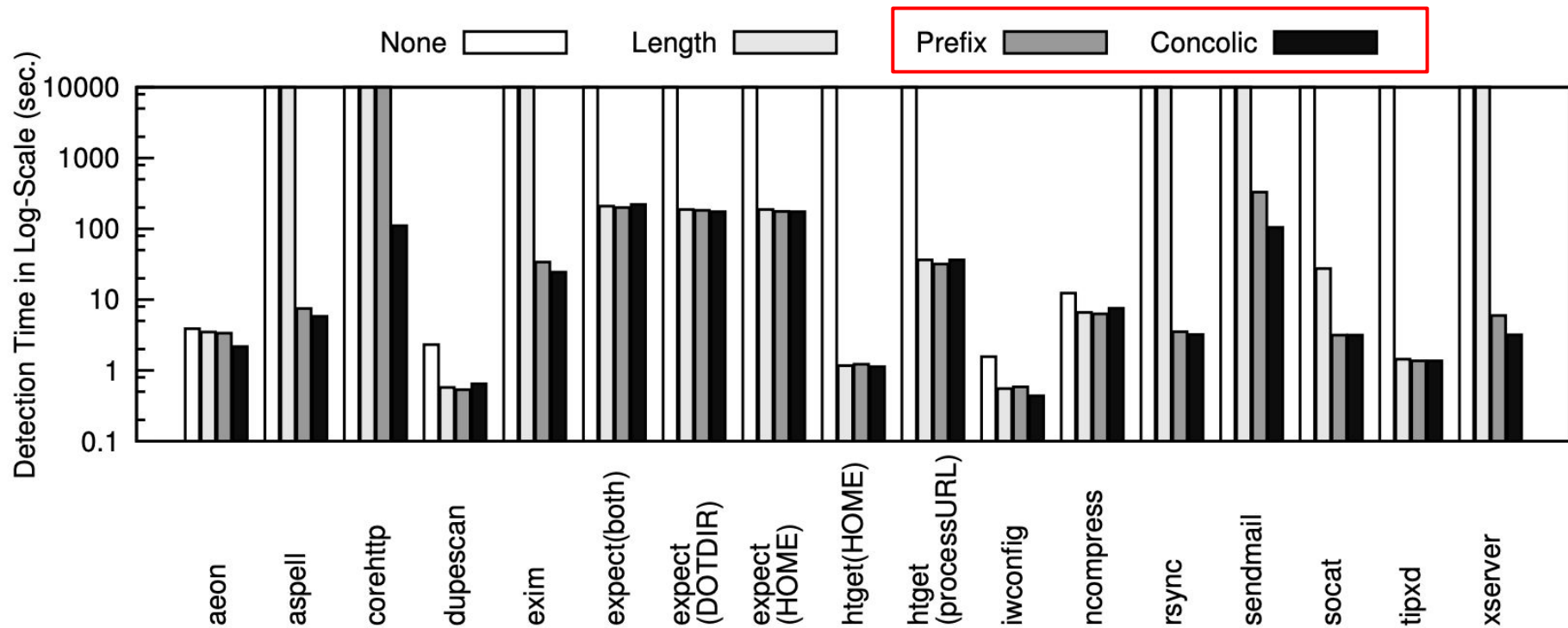
Evaluation



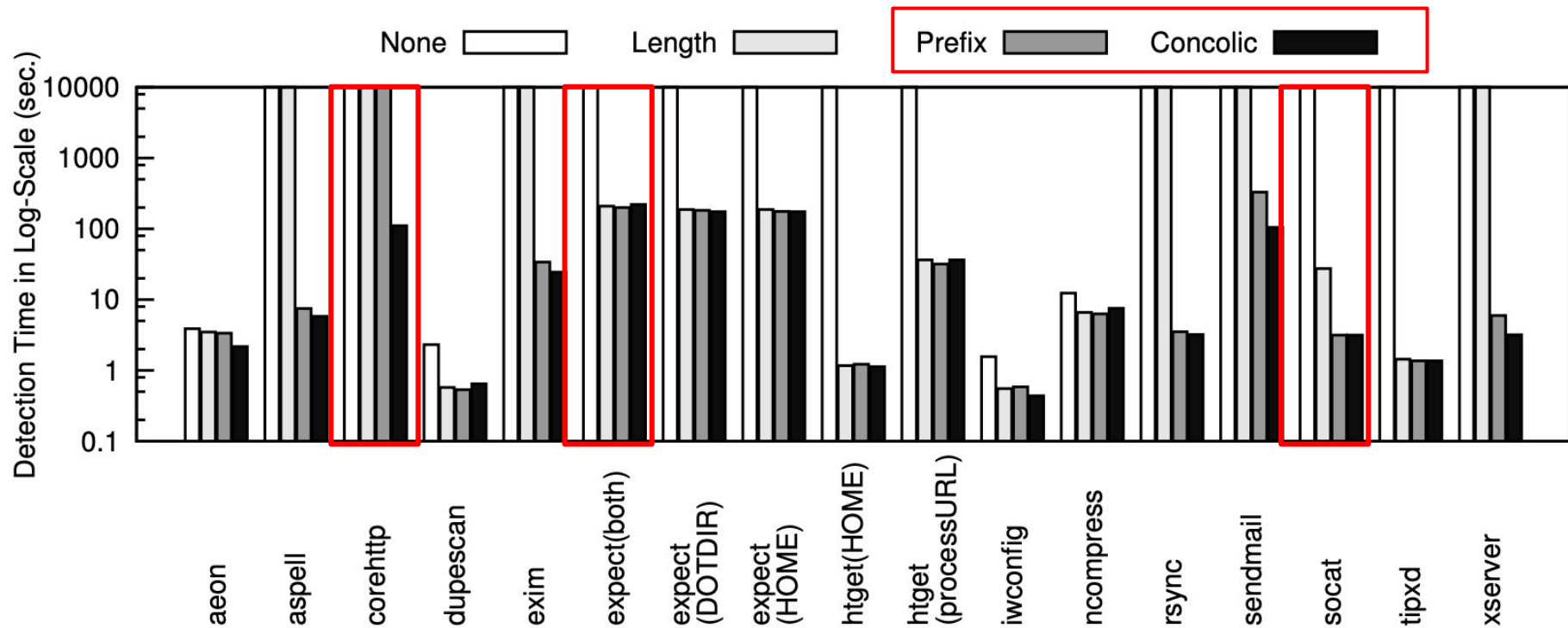
Evaluation



Evaluation



Evaluation

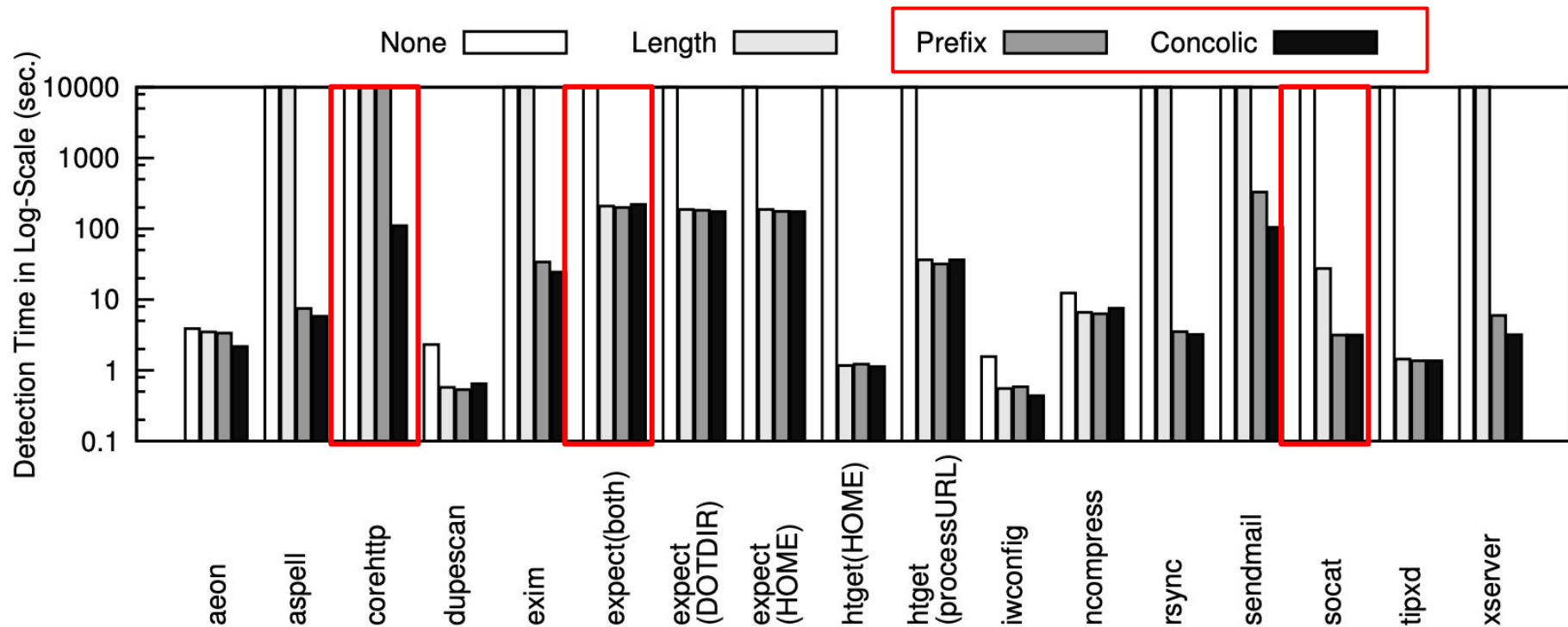


AEF in smpCTF

Solved one of problems less 10 mins

Discussion

What do you think about the causes make this happen ?



Futurework

- Heap Exploitation
- Portable Exploits
 - AEG only generates an exploit for a GNU compiled binary

Conclusion

- First fully automatic end-to-end approach for exploit generation
- Implementation of AEG
- Analyzed 14 open source projects
- Generated 16 control flow hijack exploits (2 zero day)
- Find and identify exploitable bugs
 - Developed Preconditioned symbolic execution
 - Developed Path prioritization algorithms

End!