**Q1**.A directed graph G is semi-connected if for every pair of nodes  $u, v \in V$  (G), there is a path from u to v or a path from v to u. Give an efficient algorithm in pseudo code which, given G, determines whether G is semi-connected or not. Prove the correctness and analyze the running time of your algorithm.

#### Code:

```
IsSemiconnected(G) //G (v1, v2, ....., vn) is a directed graph
Using Depth-first-search find all strongly connected components G1, G2, ..., Gk;
V[V1, V2, ....., Vk]; //Vi is the vertex set of Gi
Topological-sorting( V [V1,V2, ....., Vk] );
flag = true; //flag is a bool type variable
for i = 1 up to k do
    If edge_exist(V[i], V[i+1]) = false then
        flag = false;
return flag;
```

#### Prove:

If a directed graph G is semi-connected, then there exists a path that passes all vertex in G. If we have G1, ....., Gk as strongly connected components in G, we know that the graph is semi-connected for G1, ....., Gk. If there exists a path in G1, ....., Gk, such that contains at least one vertex of every graph in G1, ....., Gk, we ca conclude that the directed G is semi-connected since all sub graph are semi-connected, and are connected with each other.

#### ♦ Time analysis:

Depth-first-search find all strongly connected components G1, G2, ..., Gk costs O(V+E), Topological-sorting costs O(V+E), loop costs O(V). Thus the total time cost is T(n) = O(V+E).

**Q2**. Prove that if all edge weights of a graph are positive, then any subset of edges that connects all vertices and has minimum total weight must be a tree. Give an example to show that the same conclusion does not follow if we allow some weights to be non-positive.

### ◆ Prove:

Suppose we have a graph, vertexes still connected after removing the edge form the cycle. We call the new graph G'. The weight of G' is equal to weight of G minus weight of the removed edge. If G' is a tree which means there is no cycle. Adding any edge between any two vertex will increases the total weight and forming cycle, which contradicts the definition of a tree. Thus then any subset of edges that connects all vertices and has minimum total weight must be a tree.

If we allow negative weight edges in graph G. Suppose G is a tree containing only one cycle which consist of edges e1 and e2 between two vertex v1 and v2. Assume weight of e1 and e2 are negative, removing any edges between e1 and e2 would increase the total weight, which contradicts the minimum weight constrains.

# ◆ Running time report of **Prim's algorithm**:

# [1]. Prim's minimum spanning tree algorithm using adjacent list

	E	3*n	N^1.5	(n-1)/2
V				
100		1744800	2243300	1342700
		ns	ns	ns
200		2547100	2614400	1831400
		ns	ns	ns
400		3171600	3527000	2074600
		ns	ns	ns
800		3977300	5267200	3015900
		ns	ns	ns

# [2]. Prim's minimum spanning tree algorithm using adjacent matrix

	Е	3*n	N^1.5	(n-1)/2
V				
100		17607100	2243300	16723500
		ns	ns	ns
200		30210400	29648400	36947700
		ns	ns	ns
400		145349200	147127400	120842600
		ns	ns	ns
800		881566000	882888100	888191100
		ns	ns	ns

# ◆ Running time report of **creating input graphs**:

# [1]. Procedure of creating graph by ${\bf adjacent\; list\;}$

	Е	3*n	N^1.5	(n-1)/2
V				
100		927400	1210400	609400
		ns	ns	ns
200		1040100	1601500	743800
		ns	ns	ns
400		1432600	2377800	1071900
		ns	ns	ns
800		1548400	4898000	1215100
		ns	ns	ns

```
else{
matrix[i][j] = INT_MAX;
    \begin{split} & \text{int remain} = \text{E-V+1;} \\ & \text{for(int k = 0; k<remain; k++)} \\ & \text{int src = (int)(0+Math.random()*(V-1-0+0));} \\ & \text{int des = (int)(0+Math.random()*(V-1-0+0));} \\ & \text{int weight = (int)(1+Math.random()*(50-1+1));} \end{split}
         while (matrix[src][des] != INT_MAX){
    src = (int)(0+Math.random()*(V-1-0+0));
    des = (int)(0+Math.random()*(V-1-0+0));
         matrix[src][des] = weight;
     return matrix;
```

	Е	3*n	N^1.5	(n-1)/2
V				
100		315200	402600	332000
		ns	ns	ns
200		1076500	1299800	1028500
		ns	ns	ns
400		3108400	4046600	2954600
		ns	ns	ns
800		6171500	7713700	6373800
		ns	ns	ns

# Code:

```
public class prims {
static class Mynode {
     int dest;
     int weight;
Mynode(int a, int b)
       dest = a;
weight = b;
  static class Graph {
     // List of adjacent nodes of a given vertex LinkedList<Mynode>[] adj;
     // Constructor
     Graph(int e)
    {
V = e;
'' = n
       adj = new LinkedList[V];
for (int o = 0; o < V; o++)
          adj[o] = new LinkedList<>();
    int vertex;
   class comparator implements Comparator<node> {
    യverride public int compare(node node0, node node1) {
        return node0.key - node1.key;
   static void addEdge(Graph graph, int src, int dest, int weight)
     Mynode node0 = new Mynode(dest, weight);
```

```
Mynode node = new Mynode(src, weight);
graph.adj[src].addLast(node0);
     graph.adj[dest].addLast(node);
   void prims_mst(Graph graph)
     Boolean[] mstset = new Boolean[graph.V];
node[] e = new node[graph.V];
     int[] parent = new int[graph.V];
     for (int o = 0; o < graph.V; o++)
e[o] = new node();
     for (int o = 0; o < graph.V; o++) {
    mstset[o] = false;
        e[o].key = Integer.MAX_VALUE;
        e[o].vertex = o;
        parent[o] = -1;
     mstset[0] = true;
     e[0].key = 0;
     TreeSet<node> queue = new TreeSet<node>(new comparator());
     for (int o = 0; o < graph.V; o++)
        queue.add(e[o]);
     while (!queue.isEmpty()) {
        node node0 = queue.pollFirst();
        mstset[node0.vertex] = true;
        for (Mynode iterator : graph.adj[node0.vertex]) {
           if (mstset[iterator.dest] == false) {
             (mstsetj(terator.dest) == raise) {
    f(e[iterator.dest].key > iterator.weight) {
        queue.remove(e[iterator.dest]);
        eliterator.dest], key = iterator.weight;
        queue.add(e[iterator.dest]);
        parent[iterator.dest] = node0.vertex;
}
     //
//
//
   public static void fillTheGraph(Graph G, int e)
     for (int i = 0; i<G.V-1; i++){
       // 1 to 50 random weight
addEdge(G, i, i+1, (int)(1+Math.random()*(50-1+1)));
     int remain = e - G.V + 1;
     //Adjacent Matrix
  static int INT_MAX = Integer.MAX_VALUE;
   static boolean isValidEdge(int u, int v, boolean[] inMST)
     if (u == v)
return false;
     if (inMST[u] == false && inMST[v] == false)
return false;
else if (inMST[u] == true && inMST[v] == true)
       return false;
  public static int [][] fillMatrix(int V, int E){
     int[][] matrix = new int [V][V];
for(int i=0; i<V; i++){
       for(int j=0; j<V; j++){
    if(j == i+1){
        int weight = (int)(1+Math.random()*(50-1+1));
             matrix[i][j] = weight;
//System.out.println(matrix[i][j]);
             matrix[i][j] = INT_MAX;
```

```
for(int k = 0; k<remain; k++){
  int src = (int)(0+Math.random()*(V-1-0+0));
  int des = (int)(0+Math.random()*(V-1-0+0));
     int \ weight = (int)(1+Math.random()*(50-1+1));\\
     while (matrix[src][des] != INT_MAX){
    src = (int)(0+Math.random()*(V-1-0+0));
    des = (int)(0+Math.random()*(V-1-0+0));
     matrix[src][des] = weight;
  return matrix;
static void primMST(int cost[][], int V)
  boolean []inMST = new boolean[V];
  inMST[0] = true;
  int edge_count = 0, mincost = 0;
while (edge_count < V - 1)
     int min = INT_MAX, a = -1, b = -1; for (int i = 0; i < V; i++)
        for (int j = 0; j < V; j++)
        {
    if (cost[i][j] < min)
               if (isValidEdge(i, j, inMST))
                 min = cost[i][j];
//System.out.println("Mini = "+min);
                 a = i;
b = j;
      }
     if (a != -1 && b != -1)
        edge count++;
        euge_count+,
//system.outprintf("Edge %d:(%d, %d) cost: %d \n",edge_count, a, b, min);
mincost = mincost + min;
inMST[b] = inMST[a] = true;
  , System.out.printf("\n Minimum cost = %d \n", mincost);
public static void main(String[] args)
  int V = 100;
  Graph graph = new Graph(V);
fillTheGraph(graph, E);
  long begin1, end1;
long time1;
  prims e = new prims();
  begin1 = System.nanoTime();
  e.prims_mst(graph);
end1 = System.nanoTime();
  time1 = end1 - begin1;
  System.out.println(time1+" ns"+"\n");
  int[][] matrix = fillMatrix(V, E);
  long begin2, end2;
long time2;
begin2 = System.nanoTime();
  primMST(matrix, V);
end2 = System.nanoTime();
time2 = end2 - begin2;
  System.out.println(time2+" ns"+"\n");
```

## ◆ Structure of optimal solution

Dx(y): the least cost from x to y

Goal: each node notifies neighbors only when its Dv changes, then neighbors notify their neighbors

- Case 1: source node is the destination, x = y, Dx(y) = 0.
- Case 2: source node differs from destination, x = y,  $Dx(y) = min(v)\{c(x, v) + Dv(y)\}$ .

## ◆ Bellman equation:

```
Dx(y) = \{ 0 , x = y

min(v)\{ c(x, v) + Dv(y) \} , x != y }
```

#### ◆ Code

```
Distance_vector_algorithm(G, i, s, v) //G is a graph represented by adjacent matrix
   If s is equal to v then
      return 0;
   w;       //w is neighbor nodes of s
   while there are neighbor node w not visited do
      Path = mini{ opt(i-1, w1, v), opt(i-2, w2, v), ....., opt(i-k, wk, v) };
   return path + cost(s, wk);
```

#### ◆ Time analysis

From time-to-time, each node sends its own distance vector estimate to neighbors, which takes O(V). When a node x receives new DV estimate from any neighbor v, it saves v's distance vector and it updates its own DV, which takes  $O(V^3)$ . Thus the time complexity is  $T(n) = O(V^3)$ .

## Q5.

## • Running time report of **Dijkstra's algorithm**:

### [1]. Dijkstra's algorithm using adjacent list

	E	3*n	N^1.5	(n-1)/2
V				
100		1368200	19771500	293600
		ns	ns	ns
200		2460600	77584400	365700
		ns	ns	ns
400		2533700	355401400	499000
		ns	ns	ns
800		3502100	2929802100	832000
		ns	ns	ns

## [2]. Dijkstra's algorithm using adjacent matrix

	Ε	3*n	N^1.5	(n-1)/2
V				
100		17607100	2243300	16723500
		ns	ns	ns
200		30210400	29648400	36947700
		ns	ns	ns
400		145349200	147127400	120842600
		ns	ns	ns
800		881566000	882888100	888191100
		ns	ns	ns

## Running time report of creating input graphs:

#### [1]. Procedure of creating graph by adjacent list public static void fillTheGraph(List<List<Node>> adj, int V, int E) {

```
for (int i = 0; i < V - 1; i++) {
    int weight = (int) (1 + Math.random() * (20 - 1 + 1));
    adj.get(i).add(new Node(i + 1, weight));
}
int remain = E - V + 1;
for (int j = 0; | < remain; j + ) {
    int weight = (int) (1 + Math.random() * (20 - 1 + 1));
    adj.get((int) (0 + Math.random() * (V - 1 - 0 + 0))).add(new Node((int) (0 + Math.random() * (V - 1 - 0 + 0))), weight));
}
```

	Е	3*n	N^1.5	(n-1)/2
V				
100		1000500	1267700	728400
		ns	ns	ns
200		1082200	1873300	882500
		ns	ns	ns
400		1398800	3064900	1054500
		ns	ns	ns
800		1555400	5734600	1180800
		ns	ns	ns

```
erse{
    matrix[i][j] = INT_MAX;
}
    \begin{split} & \text{int remain} = \text{E-V+1;} \\ & \text{for(int k = 0; k<remain; k++)} \{ \\ & \text{int src} = (& \text{int})(0+& \text{Math.random}()^*(V-1-0+0));} \\ & \text{int des} = (& \text{int})(0+& \text{Math.random}()^*(V-1-0+0));} \\ & \text{int weight} = (& \text{int})(1+& \text{Math.random}()^*(5-0+1));} \end{split}
          while (matrix[src][des] != INT_MAX){
    src = (int)(0+Math.random()*(V-1-0+0));
    des = (int)(0+Math.random()*(V-1-0+0));
          matrix[src][des] = weight;
return matrix;
```

	Е	3*n	N^1.5	(n-1)/2
V				
100		352100	448700	367600
		ns	ns	ns
200		1350600	1429100	1173600
		ns	ns	ns
400		3111700	3618700	499000
		ns	ns	ns
800		7334400	8793300	6459800
		ns	ns	ns

```
Code
class Graph_pq {
  int dist[];
  Set<Integer> visited;
PriorityQueue<Node> pqueue;
  int V;
List<List<Node>> adj_list;
 public Graph_pq(int V) {
    this.V = V;
    dist = new int[V];
    visited = new HashSet<Integer>();
    pqueue = new PriorityQueue<Node>(V, new Node());
  public void algo_dijkstra(List<List<Node>> adj_list, int src_vertex) {
    this.adj_list = adj_list;
      for (int i = 0; i < V; i++)
dist[i] = Integer.MAX_VALUE;
      pqueue.add(new Node(src_vertex, 0));
      dist[src_vertex] = 0;
while (visited.size() != V) {
         int u = pqueue.remove().node;
    visited.add(u);
graph_adjacentNodes(u);
}
  private void graph_adjacentNodes(int u) {
      int edgeDistance = -1;
int newDistance = -1;
      for (int i = 0; i < adj_list.get(u).size(); i++) {
    Node v = adj_list.get(u).get(i);
         if (!visited.contains(v.node)) {
             edgeDistance = v.cost;
newDistance = dist[u] + edgeDistance;
            if (newDistance < dist[v.node])
  dist[v.node] = newDistance;</pre>
             pqueue.add(new Node(v.node, dist[v.node]));
  public static void fillTheGraph(List<List<Node>> adj, int V, int E) {
      for (int i = 0; i < V - 1; i++) {
    int weight = (int) (1 + Math.random() * (20 - 1 + 1));
    adj.get(i).add(new Node(i + 1, weight));
      //-----
                               -----Adiacent Matrix----
  private static char[] points;
private static int[][] arc;
private static final int INF = Integer.MAX_VALUE;
  private static void dijkstra(int start) {
  boolean[] flag = new boolean[points.length];
  int[] distence = arc[start];
  int[] prev = new int[points.length];
      for (int i = 0; i < points.length; i++) {
    prev[i] = start;
}</pre>
      flag[start] = true;
prev[start] = start;
Stack path = new Stack();
int currentSmall = 0;
      for (int j = 1; j < points.length; j++) {
  int minDistence = INF;
  for (int i = 0; i < points.length; i++) {</pre>
            if (!flag[i] && distence[i] < minDistence) {
                minDistence = distence[i];
currentSmall = i;
          flag[currentSmall] = true;
         \label{eq:formula} \begin{aligned} &\text{for (int } i = 0; \ i < points.length; \ i++) \ \{ \\ &\text{if (!flag[i])} \end{aligned}
                 int\ temp = (arc[currentSmall][i] == INF\ ?\ INF\ : (minDistence + arc[currentSmall][i]));
          distence[i] = temp;
prev[i] = currentSmall;
}
```

```
int j = prev[i];
while (j != start) {
  path.push(points[j]);
  j = prev[j];
                   J
System.out.print("shortestDistenct("+ points[start] + "," + points[i] + ") = " + distence[i] + " ");
System.out.print("path = " + points[start] + " ");
while (lpath.isEmpty()) {
System.out.print(path.pop() + " ");
                   System.out.println(points[i]);
  public static int [][] fillMatrix(int V, int E){
         ublic static int []] fillMatrix(int V, int E){
int[]] matrix = new int [V][V];
for(int i=0; i<V; i++){
    for(int j=0; i<V; i++){
        if(j ==i+1 | j ==i+1){
            int weight = (int)(1+Math.random()*(50-1+1);
            matrix[j][j] = weight;
            matrix[j][j] = weight;
            //System.out.println(matrix[i][j]);
        }
}</pre>
                                matrix[i][j] = INF;
          \begin{split} & \text{int remain} = \text{E-V+1;} \\ & \text{for(int } k = 0; \ k < \text{remain; } k + + \} \{ \\ & \text{int } \text{src} = (\text{int})(0 + \text{Math.random}()^*(V - 1 - 0 + 0));} \\ & \text{int } \text{des} = (\text{int})(0 + \text{Math.random}()^*(V - 1 - 0 + 0));} \\ & \text{int } \text{weight} = (\text{int})(1 + \text{Math.random}()^*(5 - 0 + 1));} \end{split}
                  while (matrix[src][des] != INF){
    src = (int)(0+Math.random()*(V-1-0+0));
    des = (int)(0+Math.random()*(V-1-0+0));
                   matrix[src][des] = weight;
return matrix;
  public static void main(String arg[]) {
          long begin1, end1;
long time1;
          long begin2, end2;
long time2;
long begin3, end3;
long time3;
          int V = 800:
          int E = (V-1)/2;
int source = 0;
          List<List<Node>> adj_list = new ArrayList<List<Node>>();
          for (int i = 0; i < V; i++) {
    List<Node> item = new ArrayList<Node>();
    adj_list.add(item);
          begin1 = System.nanoTime();
fillTheGraph(adj_list, V, E);
end1 = System.nanoTime();
time1 = end1 - begin1;
          System.out.println("Fill the list: "+time1);
           begin2 = System.nanoTime();
arc = fillMatrix(V, E);
          end2 = System.nanoTime();
time2 = end2 - begin2;
           System.out.println("Fill the matrix: "+time2);\\
          // call Dijkstra's algo method
Graph_pq dpq = new Graph_pq(V);
          begin3 = System.nanoTime();
dpq.algo_dijkstra(adj_list, source);
end3 = System.nanoTime();
time3 = end3 - begin3;
       // Print the shortest path from source node to all the nodes System.out.println("The shorted path from source node to other nodes:"); System.out.println("Source\t\t\t\" + "Node#\t\t" + "Distance"); for (int i=0, i<0, i<0
          arc = fillMatrix(V, E);
```

```
dijkstra(1);
 // Node class
static class Node implements Comparator<Node> {
   public int node;
   public int cost;
   public Node() {
   public Node(int node, int cost) {
  this.node = node;
  this.cost = cost;
   @Override
public int compare(Node node1, Node node2) {
  if (node1.cost < node2.cost)
    return -1;</pre>
    if (node1.cost > node2.cost)
return 1;
return 0;
Q6.
EXTEND-SHORTEST-PATHS(\prod, L, W)
       n = L.row;
                            // L' is a n \times n matrix
       L' = l'[i, j];
       \prod' = \pi'[i, j];
                         // ∏' is a n × n matrix
       for i = 1 to n do
          for j = 1 to n do
                 I'[i, j] = INFINITE; \pi'[i, j] = NIL;
                 for k = 1 to n do
                     if I[i, k] + I[k, j] < I[i, j] then
                        I[i, j] = I[i, k] + I[k, j];
                            if k = j then \pi'[i, j] = k;
                            else \pi'[i, j] = \pi[i, j];
       return (∏', L');
SLOW-ALL-PAIRS-SHORTEST-PATHS(W)
       n = W.rows;
       L(1) = W;
       \prod(1) = \pi[i, j](1);
       if there is an edge from i to j then
          \pi[i, j](1) = i;
       else \pi[i, j](1) = NIL;
       for m = 2 to n - 1 do
          \prod(m), L(m) = EXTEND-SHORTEST-PATHS(\prod(m-1), L(m-1), W);
       return (\prod(n-1), L(n-1));
Q7.
Faster-All-Pairs-Shortest-Paths(W)
   matrix = W; r = 1; // matrix is 2 x 2 matrix
   while r < n - 1 do
       matrix =Extend-Shortest-Paths(matrix, matrix);
       r = 2r;
   return matrix.
```