

A formalization of forcing and the consistency of the failure of the continuum hypothesis

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Abstract

We describe a formalization of forcing using Boolean-valued models in the Lean 3 theorem prover, including the fundamental theorem of forcing and a deep embedding of first-order logic with a Boolean-valued soundness theorem. As an application of our framework, we specialize our construction to a Boolean completion of the Cohen poset and formally verify in the resulting model the failure of the continuum hypothesis.

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Introduction

The continuum hypothesis states that there are no sets strictly larger than the countable natural numbers and strictly smaller than the uncountable real numbers. It was introduced by Cantor in 1878 [?] and was the very first problem on Hilbert’s list of twenty-three outstanding problems in mathematics. Gödel proved in 1938 [?] that the continuum hypothesis was consistent with ZFC, and later conjectured that the continuum hypothesis is independent of ZFC, i.e. neither provable nor disprovable from the ZFC axioms. In 1963, Paul Cohen developed *forcing* [?], which allowed him to prove the consistency of the negation of the continuum hypothesis, and therefore complete the independence proof. For this work, which marked the beginning of modern set theory, he was awarded a Fields medal—the only one to ever be awarded for a work in mathematical logic.

In this paper we discuss the formalization of a Boolean-valued model of set theory where the continuum hypothesis fails. The work we describe is part of the Flypitch project, which aims to formalize the independence of the continuum hypothesis. Our results mark a major milestone towards that goal.

Our formalization is written in the Lean 3 theorem prover. Lean is an interactive proof assistant under active development at Microsoft Research [?, ?]. It implements the Calculus

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of Inductive Constructions and has a similar metatheory to Coq, adding definitional proof irrelevance, quotient types, and a noncomputable choice principle. There is a well-known encoding of ZFC into dependent type theory with CIC, due to Aczel [?, ?] and Werner [?], which has been implemented in Lean’s mathematical library `mathlib`² by Mario Carneiro. The fact that Lean’s metatheory is powerful enough to encode a model of ZFC already allows us to perform metatheoretic arguments about ZFC which were unavailable to e.g. Paulson [?], who went to extreme lengths to circumvent them inside Isabelle/ZF. While this was a formidable task, we content ourselves with treating ZFC as a mathematical object of study, and freely ignore the restrictions which it would impose as a foundation for the metatheory.

Indeed, our formalization makes as much use of the expressiveness of Lean’s dependent type theory as possible, using constructions which are impossible or unwieldy to encode in HOL, much less ZF: Lean’s ordinals and cardinals, which are defined as equivalence classes of well-ordered types, live one universe level up and play a crucial role in the forcing argument; the models of set theory we construct require as input an entire universe of types; our encoding of first-order logic uses parametrized inductive types to equate type-correctness with well-formedness, eliminating the need for separate well-formedness proofs.

The method of forcing with Boolean-valued models was developed by Solovay and Scott (and independently, Vopěnka) in ’65-’66 [?, ?] as a simplification of Cohen’s method. Some of these simplifications were incorporated by Shoenfield [?] into a general theory of forcing using partial orders, and it is in this form that forcing is usually practiced. While both approaches have essentially the same mathematical content (see e.g. the discussion in Kunen [?] or Jech [?]), there are several reasons why we chose Boolean-valued models for our formalization:

- **Modularity.** The theory of forcing with Boolean-valued models cleanly splits into several components (a general theory of Boolean-valued semantics for first-order logic, a library for calculations inside complete Boolean algebras, the construction of Boolean-valued models of set theory, and the specifics of the forcing argument itself) which could be formalized in parallel and then recombined.
- **Directness.** For the purposes of an independence proof, the Boolean-valued soundness theorem eliminates the need to produce a two-valued model. This approach also bypasses any requirement for the reflection theorem/Löwenheim-Skolem theorems, Mostowski collapse, countable transitive models, or genericity considerations for filters.
- **Novelty and reusability.** As far as we were able to tell, the Boolean-valued approach to forcing has never been formalized. Furthermore, while for the purposes of an independence proof, forcing with Boolean-valued models and forcing with countable transitive models accomplish the same thing, a general library for Boolean-valued semantics of a deeply embedded logic could be used for formal verification applications outside of set theory, e.g. to formalize the Boolean-valued semantics of the stochastic λ -calculus [?].
- **Amenability to structural induction.** As with Coq, Lean is able to encode extremely complex objects and reason about their specifications using inductive types. However, the user must be careful to choose the encoding so that properties they wish to reason about are accessible by structural induction, which is the most natural mode of reasoning in the proof assistant. After observing (1) that the Aczel-Werner encoding of ZFC as an inductive type is essentially a special case of the recursive *name* construction from forcing (c.f. Section 3), and (2) that the automatically-generated induction principle for

² <https://github.com/leanprover-community/mathlib>

that inductive type *is* \in -induction, it is easy to see that this encoding can be modified to produce a Boolean-valued model of set theory where, again, \in -induction comes for free.

In Section 1 we outline the method of Boolean-valued models and sketch the forcing argument. Section 2 discusses a deep embedding of first-order logic, including a proof system, Boolean-valued semantics, and the Boolean-valued soundness theorem. Section 3 discusses our construction of Boolean-valued models of set theory. Section 4 describes the formalization of the forcing argument and the construction of a suitable Boolean algebra \mathbb{B} for forcing $\neg\text{CH}$. Section 5 describes the formalization of some transfinite combinatorics. We conclude with a reflection on our formalization and an indication of future work.

1 Outline of the proof

ZFC is a collection of first-order sentences in the language of a single binary predicate relation $\{\in\}$, used to axiomatize set theory. The continuum hypothesis can be written in this fashion as a first-order sentence CH. A proof of CH is a finite list of deductions starting from ZFC and ending at CH. The soundness theorem says that provability implies satisfiability, i.e. if $\text{ZFC} \vdash \text{CH}$, then CH interpreted in any model of ZFC is true. Taking the contrapositive, we can demonstrate the unprovability (equivalently, the consistency of the negation) of CH by exhibiting a single model where CH is not true.

A model of a first-order theory T in a language L is in particular a way of assigning true or false in a coherent way to sentences in L . Modulo provable equivalence, the sentences form a Boolean algebra and “coherent” means the assignment is a Boolean algebra homomorphism (so \wedge becomes meet, \vee becomes join, \forall becomes an indexed infimum, etc.) into $\mathbf{2} = \{\text{true}, \text{false}\}$. The soundness theorem ensures that this homomorphism v sends a proof $\phi \vdash \psi$ to an inequality $v(\phi) \leq v(\psi)$. But $\mathbf{2}$ does not really play a special role in this scheme, and may be replaced by any complete Boolean algebra \mathbb{B} , where the top and bottom elements \top, \perp take the place of true and false. It is straightforward to extend this analogy to a \mathbb{B} -valued semantics for first-order logic, and in this generality, the soundness theorem now says that for any such \mathbb{B} , if $\text{ZFC} \vdash \text{CH}$, then for any \mathbb{B} -valued structure where all the axioms of ZFC have truth-value \top , CH does also. Then as before, to demonstrate the consistency of the negation of CH it suffices to find just one \mathbb{B} and a single \mathbb{B} -valued model where CH is not “true”.

This is where forcing comes in. Given a universe V of set theory which contains a Boolean algebra \mathbb{B} , one constructs in analogy to the cumulative hierarchy a new \mathbb{B} -valued universe $V^{\mathbb{B}}$ of set theory, where the powerset operation is replaced by taking functions into \mathbb{B} . Thus, the structure of \mathbb{B} informs the decisions made by $V^{\mathbb{B}}$ about what subsets, hence functions, exist among the members of $V^{\mathbb{B}}$; the real challenge lies in selecting a suitable \mathbb{B} and reasoning about how its structure affects the structure of $V^{\mathbb{B}}$. While $V^{\mathbb{B}}$ may vary wildly depending on the choice of \mathbb{B} , the original universe V always embeds into $V^{\mathbb{B}}$ via an operation $x \mapsto \check{x}$, and while the passage of x to \check{x} may not always preserve its original properties, properties which can be specified with only bounded quantification are preserved; in particular, $V^{\mathbb{B}}$ thinks $\check{\mathbb{N}}$ is \mathbb{N} .

To force the negation of the continuum hypothesis, we use the Boolean algebra $\mathbb{B} := \text{RO}(2^{\aleph_2 \times \mathbb{N}})$ of regular opens of the Cantor space $2^{\aleph_2 \times \mathbb{N}}$. For each $\nu \in \aleph_2$, we associate the \mathbb{B} -valued characteristic function $\chi_\nu : \mathbb{N} \rightarrow \mathbb{B}$ by $n \mapsto \{f \mid f(\nu, n) = 1\}$. This induces what $V^{\mathbb{B}}$ thinks is a new subset $\check{\chi}_\nu \subseteq \mathbb{N}$, called a *Cohen real*, and furthermore, simultaneously performing this construction on all $\nu : \aleph_2$ induces what $V^{\mathbb{B}}$ thinks is a function from $\check{\aleph}_2 \rightarrow \mathcal{P}(\mathbb{N})$. After showing that $V^{\mathbb{B}}$ thinks this function is injective, to finish the proof it suffices to show that $x \mapsto \check{x}$ preserves cardinal inequalities, as then we will have squeezed $\check{\aleph}_1$

properly between \mathbb{N} and $\mathcal{P}(\mathbb{N})$. This is really the technical heart of the matter, and relies on a combinatorial property of \mathbb{B} called the *countable chain condition* (CCC), the proof of which requires a detailed combinatorial analysis of the basis of the product topology for $2^{\mathbb{N}_2 \times \mathbb{N}}$; we handle this with a general result in transfinite combinatorics called the *Δ -system lemma*.

So far we have mentioned nothing about how this argument, which is wholly set-theoretic, is to be interpreted inside type theory. To do this, it was important for us to separate the mathematical content from the metamathematical content of the argument. While our objective is only to produce some model of ZFC satisfying certain properties, traditional presentations of forcing are careful to stay within the foundations of ZFC, emphasizing that all arguments may be performed internally inside a model of ZFC, etc., and it is not immediately clear what parts of the argument use that set-theoretic foundation in an essential way and require modification in the passage to type theory. As we will see, our formalization clarifies some of these questions.

Sources

Our strategy for constructing a Boolean-valued model in which the continuum hypothesis fails is a synthesis of the proofs in the textbooks of Bell ([?], Chapter 2) and Manin ([?], Chapter 8). For the Δ -system lemma, we follow Kunen ([?], Chapters 1 and 5).

Viewing the formalization

The code blocks in this paper were taken directly from our formalization, but for the sake of formatting and readability, we sometimes omit or modify universe levels, type ascriptions, and casts. We refer the interested reader to our source code.³ The forcing argument for the negation of CH is located in `forcing.lean`. In a Lean-aware editor such as Emacs, the user is encouraged start at the theorem `neg_CH` and jump backwards to trace the dependencies of the proof. (TODO(jesse) maybe change this to point at the “summary” file)

2 First-order logic

The starting point for first-order logic is a *language* of relation and function symbols. We represent a language as a pair of \mathbb{N} -indexed families of types, each of which is to be thought of as the collection of relation (resp. function) symbols stratified by arity:

```
structure Language : Type (u+1) :=
  (functions :  $\mathbb{N} \rightarrow \text{Type } u$ ) (relations :  $\mathbb{N} \rightarrow \text{Type } u$ )
```

2.1 (Pre)terms, (pre)formulas

The main novelty of our implementation of first-order logic is the use of *partially applied* terms and formulas, encoded in a parametrized inductive type where the \mathbb{N} parameter measures the difference between the arity and the number of applications. The benefit of this is that it is impossible to produce an ill-formed term or formula, because type-correctness is equivalent to well-formedness. This eliminates the need for separate well-formedness proofs.

Fix a language L . We define the type of **preterms** as follows:

³ <https://github.com/flypitch/flypitch>

```

171
172 inductive preterm : ℕ → Type u
173 | var {} : ∀ (k : ℕ), preterm 0
174 | func : ∀ {l : ℕ} (f : L.functions l), preterm l
175 | app : ∀ {l : ℕ} (t : preterm (l + 1)) (s : preterm 0), preterm l
176

```

We use de Bruijn indices to avoid variable shadowing. A member of `preterm n` is a partially applied term. If applied to `n` terms, it becomes a term. Every element of `preterm L 0` is a well-formed term. We use this encoding to avoid mutual or nested inductive types, since those are not too convenient to work with in Lean.

The type of `preformulas` is defined similarly:

```

182
183 inductive preformula : ℕ → Type u
184 | falsum {} : preformula 0 -- notation '\bot'
185 | equal (t1 t2 : term L) : preformula 0 -- notation '\simeq'
186 | rel {l : ℕ} (R : L.relations l) : preformula l
187 | apprel {l : ℕ} (f : preformula (l + 1)) (t : term L) : preformula l
188 | imp (f1 f2 : preformula 0) : preformula 0 -- notation '⇒'
189 | all (f : preformula 0) : preformula 0 -- notation '∀'
190 -- ¬ f := f ⇒ ⊥, notation '~f'
191 -- ∃ f := ~ ∀' ~f, notation '∃' f'
192

```

A member of `preformula n` is a partially applied formula. If applied to `n` terms, it becomes a formula. Implication is the only binary connective. Since we use classical logic, we can define the other connectives from implication and `falsum`. Similarly, universal quantification is our only quantifier.

Our proof system is a natural deduction calculus, and all rules are motivated to work well with backwards-reasoning:

```

199
200
201 inductive prf : set (formula L) → formula L → Type u
202 | axm {Γ A} (h : A ∈ Γ) : prf Γ A
203 | impI {Γ} {A B} (h : prf (insert A Γ) B) : prf Γ (A ⇒ B)
204 | impE {Γ} {A} {B} (h1 : prf Γ (A ⇒ B)) (h2 : prf Γ A) : prf Γ B
205 | falsumE {Γ} {A} (h : prf (insert ~A Γ) ⊥) : prf Γ A
206 | allI {Γ A} (h : prf Γ A) : prf Γ (∀' A)
207 | allE2 {Γ} A t (h : prf Γ (∀' A)) : prf Γ (A[t // 0])
208 | ref (Γ t) : prf Γ (t ≈ t)
209 | subst2 {Γ} (s t f) (h1 : prf Γ (s ≈ t)) (h2 : prf Γ (f[s // 0])) :
210 prf Γ (f[t // 0])
211

```

A member of `prf Γ A` is a proof tree encoding a derivation of `A` from `Γ`. Note that `prf` is `Type`- instead of `Prop`-valued, so different members of `prf Γ A` are not definitionally equal.

2.2 Completeness

As part of our formalization of first-order logic, we completed a verification of the Gödel completeness theorem. Although our present development of forcing did not require it, we anticipate that it will be useful later to e.g. prove the downward Löwenheim-Skolem theorem for extracting countable transitive models. Like soundness, it also serves as a proof-of-concept and stress-test of our chosen encoding of first-order logic.

For our formalization, we chose the Henkin-style approach of constructing a canonical term model. In order to perform the argument, which normally involves modifying the language “in place” to iteratively add new constant symbols, we had to adapt it to type theory. Since our languages are represented by pairs of indexed types instead of sets, we cannot really modify them in-place with new constant symbols. Instead, at each step of the construction, we must construct an entirely new language in which the previous one embeds, and in the limit we must compute a directed colimit of types instead of a union. This construction induces similar constructions on terms and formulas, and completing the argument requires reasoning with all of them. As a result of our design decisions, only a few arguments required anything more than straightforward case-analysis and structural induction. The final statement makes no restrictions on the cardinality of the language.

2.3 Boolean-valued semantics for first-order logic

A **complete Boolean algebra** is a type \mathbb{B} equipped with the structure of a Boolean algebra and additionally operations Inf and Sup (which we write as \sqcap and \sqcup) returning the infimum and supremum of an arbitrary collection of members of \mathbb{B} . For more details on complete Boolean algebras, we refer the reader to the textbook of Halmos-Givant [?].

► **Definition 1.** Fix a language L and a type \mathbb{B} with the structure of a complete Boolean algebra. A \mathbb{B} -valued structure is an instance of the following **structure**:

```

structure bStructure :=
  (carrier : Type u)
  (fun_map : ∀{n}, L.functions n → vector carrier n → carrier)
  (rel_map : ∀{n}, L.relations n → vector carrier n →  $\mathbb{B}$ )
  (eq : carrier → carrier →  $\mathbb{B}$ )
  (eq_refl : ∀ x, eq x x =  $\top$ )
  (eq_symm : ∀ x y, eq x y = eq y x)
  (eq_trans : ∀{x} y {z}, eq x y  $\sqcap$  eq y z ≤ eq x z)
  (fun_congr : ∀{n} (f : L.functions n) (x y : vector carrier n),
     $\sqcap(\text{map2 eq } x \ y) \leq \text{eq } (\text{fun\_map } f \ x) (\text{fun\_map } f \ y)$ )
  (rel_congr : ∀{n} (R : L.relations n) (x y : vector carrier n),
     $\sqcap(\text{map2 eq } x \ y) \sqcap \text{rel\_map } R \ x \leq \text{rel\_map } R \ y$ )

```

Above, “ $\sqcap(\text{map2 eq } x \ y)$ ” means “for all i , fold the binary inf \sqcap across the list whose i th entry is eq applied to $x[i]$ and $y[i]$ ”.

Note that Boolean-valued equality is not really an equivalence relation, but “ \mathbb{B} thinks it is”. One complication which then arises in Boolean-valued semantics is keeping track of the congruence lemmas for formulas. However, as part of the soundness theorem shows, once these extensionality proofs are provided for the basic symbols in the language, they extend by structural induction to all formulas.

2.4 The soundness theorem

A soundness theorem says that a proof tree may be replayed to produce an actual proof in the object of truth-values. When the object of truth-values is **Prop**, this says that a proof tree compiles to a proof term. When the object of truth-values is a Boolean algebra, this says that the proof tree becomes an internal implication from the interpretation of the context to the interpretation of the conclusion:

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```
lemma boolean_soundness {Γ : set (formula L)} {A : formula L}
  (H : Γ ⊢ A) : ∀ M, (⊓ γ ∈ Γ, M[γ]) ≤ M[A]
```

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As a result of our design decisions, the proofs of both the ordinary and Boolean-valued soundness theorems are straightforward structural inductions.

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3 Constructing Boolean-valued models of set theory

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Throughout this section, we fix a universe level u , a type $\mathbb{B} : \text{Type } u$ and an instance of a complete Boolean algebra structure on \mathbb{B} .

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In set theory (see e.g. Jech [?] or Bell [?]), Boolean-valued models are obtained by imitating the construction of the von Neumann cumulative hierarchy via a transfinite recursion where iterations of the powerset operation (taking functions into $\mathbf{2} = \{\text{true}, \text{false}\}$) are replaced by iterations of the “ \mathbb{B} -valued powerset operation” (taking functions into \mathbb{B}).

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Since this construction by transfinite recursion does not easily translate into type theory, our construction of Boolean-valued models of set theory is instead a variation on a well-known encoding originally due to Aczel [?, ?, ?]. This encoding was adapted by Werner [?] to encode ZFC into Coq, whose metatheory is close to that of Lean. Werner’s construction was re-implemented in Lean’s `mathlib` by Carneiro as part of [?]. In this approach, one takes a universe of types `Type u` as the starting point and then imitates the cumulative hierarchy by constructing the inductive type

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```
inductive pSet : Type (u+1)
| mk (α : Type u) (A : α → pSet) : pSet
```

290

291

The Aczel-Werner encoding is closely related to the recursive definition of *names*, which is used in forcing to construct forcing extensions:

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► **Definition 2.** Let P be a partial order (which one thinks of as a collection of forcing conditions). A P -name is a collection of pairs (y, p) where y is a P -name and $p : P$.

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If P consists of only one element, then a P -name is specified by essentially the same information as a member of the inductive type `pSet` above. Conversely, specializing P to an arbitrary complete Boolean algebra \mathbb{B} , we generalize the definition of `pSet.mk` so that elements are recursively assigned Boolean truth-values:

```
inductive bSet (B : Type u) [complete_boolean_algebra B] : Type (u+1)
| mk (α : Type u) (A : α → bSet) (B : α → B) : bSet
```

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Thus `bSet B` is the type of \mathbb{B} -names, and will be the underlying type of our Boolean-valued model of set theory. For convenience, if $x : \text{bSet } \mathbb{B}$ and $x := \langle \alpha, A, B \rangle$, we put $x.\text{type} := \alpha$, $x.\text{func} := A$, $x.\text{bval} := B$.

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3.1 Boolean-valued equality and membership

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In `pSet`, equivalence of sets is defined by structural recursion as follows: two sets x and y are equivalent if and only if for every $w \in x$, there exists a $w' \in y$ such that w is equivalent to w' , and vice-versa. Analogously, by translating quantifiers and connectives into operations on \mathbb{B} , Boolean-valued equality is defined in the same way:


```

310 def bv_eq : ∀ (x y : bSet ℤ), ℤ
311 | ⟨α, A, B⟩ ⟨α', A', B'⟩ :=
312   (⊓ a : α, B a ⇒ ⊓ a', B' a' ⊓ bv_eq (A a) (A' a')) ⊓
313   (⊓ a' : α', B' a' ⇒ ⊓ a, B a ⊓ bv_eq (A a) (A' a'))
314
315

```

316 We abbreviate `bv_eq` with the infix operator `=B`. With equality in place, it is easy to
 317 define membership, by translating “ x is a member of y if and only if there exists a $w \in y$
 318 such that $x = w$.” As with equality, we denote \mathbb{B} -valued membership with `∈B`.

```

319
320
321 def mem : bSet ℤ → bSet ℤ → ℤ
322 | a (mk α' A' B') := ⊓ a', B' a' ⊓ a =B A' a'
323

```

3.2 Automation and metaprogramming for reasoning in \mathbb{B}

324 As Scott stresses in [?], “A main point ... is that the well-known algebraic characterizations
 325 of [complete Heyting algebras] and [complete Boolean algebras] exactly mimic the rules of
 326 deduction in the respective logics. . .” Indeed, that is really why the Boolean-valued soundness
 327 theorem is true. One thinks of the \leq symbol in an inequality of Boolean truth-values as a
 328 turnstile in a proof state: the conjunctands on the left as a list of assumptions in context, and
 329 the quantity on the right as the goal. For example, given $a \ b : \mathbb{B}$, the identity $(a \Rightarrow b) \sqcap a \leq b$
 330 could be proven by unfolding the definition of material implication, but it is really just modus
 331 ponens; similarly, given an indexed family $a : I \rightarrow \mathbb{B}$, $\sqcup i, a \ i \leq b \leftrightarrow \forall i, a \ i \leq b$
 332 is just \exists -elimination.

333 Difficulties arise when the statements to be proved become only slightly more complicated.
 334 Consider the following example, which should be “**by assumption**”:

```

335
336
337 ∀ a b c d e f g : ℤ, (d ⊓ e) ⊓ (f ⊓ g ⊓ ((b ⊓ a) ⊓ c)) ≤ a
338

```

339 or slightly less trivially, the following example where the goal is attainable by “just applying
 340 a hypothesis to an assumption”

```

341
342 ∀ a b c d : ℤ, (a ⇒ b) ⊓ c ⊓ (d ⊓ a) ≤ b
343

```

344 There are three ways to deal with goals like these, which approximately describe the
 345 evolution of our approach. First, one can try using the basic lemmas in `mathlib`, using
 346 the simplifier to normalize expressions, and performing clever rewrites with the deduction
 347 theorem.⁴ Second, one can take the LCF-style approach and expand the library of lemmas
 348 with increasingly sophisticated derived inference rules.

349 Third, one can make the following observation:

350 ► **Lemma 3** (Yoneda lemma for posets). *Let (P, \leq) be a partially ordered set. Let $a \ b : P$.
 351 Then $a \leq b$ if and only if $\forall \Gamma : P, \Gamma \leq a \rightarrow \Gamma \leq b$.*

352 This is a consequence of the Yoneda lemma for partially ordered sets, and its proof is utterly
 353 trivial. However, one side of the equivalence is much easier for Lean to reason with. Take
 354 the example which should have been “**by assumption**”. The following proof, in which the
 355 user navigates down the binary tree of nested \sqcap s, will work:

⁴ The deduction theorem in a Boolean algebra says that for all a, b and c , $a \sqcap b \leq c \iff a \leq b \Rightarrow c$.


```

356 example {a b c d e f g :  $\mathbb{B}$ } : (d  $\sqcap$  e)  $\sqcap$  (f  $\sqcap$  g  $\sqcap$  ((b  $\sqcap$  a)  $\sqcap$  c))  $\leq$  a :=
357 by {apply inf_le_right_of_le, apply inf_le_right_of_le,
358     apply inf_le_left_of_le, apply inf_le_right_of_le, refl}
359
360

```

But if we use the right-hand side of Lemma 3 instead, then after some preprocessing, `assumption` will literally work:

```

363
364
365 example {a b c d e f g :  $\mathbb{B}$ } : (d  $\sqcap$  e)  $\sqcap$  (f  $\sqcap$  g  $\sqcap$  ((b  $\sqcap$  a)  $\sqcap$  c))  $\leq$  a :=
366 by {apply poset_yoneda, intros  $\Gamma$  H,
367     simp only [le_inf_iff] at H /- turns H into conjunction -/,
368     repeat{auto_cases} /- splits conjunction -/,
369     assumption}
370 /- Goal state before 'assumption':
371 [...]
372 H_right_right_left_left :  $\Gamma \leq b$ ,
373 H_right_right_left_right :  $\Gamma \leq a$ 
374  $\vdash \Gamma \leq a$  -/
375

```

A key feature of Lean is that it is its own metalanguage, allowing for seamless in-line definitions of custom tactics. This feature was an invaluable asset, as it allowed the rapid development of a custom tactic library for simulating natural-deduction style proofs inside \mathbb{B} after applying Lemma 3. The preprocessing steps before the call to `assumption` in the previous example are subsumed into a single tactic. Boolean-valued versions of natural deduction rules like \vee/\wedge -elimination, instantiation of existentials, implication introduction, and even basic automation were easy to write. The result is that the user is able to pretend, with absolute rigor, that they are simply writing proofs in first-order logic while calculations in the complete Boolean algebra are being performed under the hood.

One use-case where automation is crucial is context-specialization. For example, suppose that after preprocessing with `poset_yoneda`, the goal is $\Gamma \leq a \implies b$, and one would like to “introduce the implication”, adding $\Gamma \leq a$ to context and reducing the goal to $\Gamma \leq b$. This is impossible as stated. Rather, the deduction theorem lets us rewrite the goal to $\Gamma \sqcap a \leq b$, and now we may add $\Gamma \sqcap a \leq a$. So we may introduce the implication after all, but at the cost of specializing the context Γ to the smaller context $\Gamma' := \Gamma \sqcap a$. But now, in order for the user to continue the pretense that they are merely doing first-order logic, this change of variables must be propagated to the rest of the assumptions which may still be of the form $\Gamma \leq _$ —which is extremely tedious to do by hand, but easy to automate.

3.3 Check-names

From the definitions of `pSet` and `bSet`, one immediately sees that there is a canonical map `check : pSet \rightarrow bSet \mathbb{B}` , defined by

```

397
398 def check : pSet  $\rightarrow$  bSet  $\mathbb{B}$ 
399 |  $\langle \alpha, A \rangle$  :=  $\langle \alpha, \lambda a, \text{check } (A \ a), \lambda a, \top \rangle$ 
400

```

That is, `check` takes a `pSet` and recursively attaches the Boolean truth-value \top to all elements. We call members of the image of `check` *check-names*,⁵ after the usual diacritic

⁵ This terminology is standard, c.f. [?, ?].

notation \check{x} for `check (x : pSet)`. These are also known as *canonical names*, as they are the canonical representation of standard two-valued sets inside a Boolean-valued model of set theory. We were pleased to discover Lean’s support for custom notation allowed us to declare the Unicode modifier character \checkmark as a postfix operator for `check`.

3.4 The fundamental theorem of forcing

The fundamental theorem of forcing for Boolean-valued models [?] states that for any complete Boolean algebra B , V^B is a Boolean-valued model of ZFC. Since, in type theory, a type universe `Type u` takes the place of the standard universe V , the analogous statement in our setting is that for every complete Boolean algebra \mathbb{B} , `bSet \mathbb{B}` is a Boolean-valued model of ZFC.

Bell [?] gives an extremely detailed account of the verification of the ZFC axioms, and we faithfully followed his presentation for this part of the formalization. Most of it is routine. We describe some aspects of `bSet \mathbb{B}` which are revealed by this verification.

The axiom of infinity

$\omega : \text{bSet } \mathbb{B}$ is $\check{\omega}$. ω is defined in `pSet` to be the collection of all finite von Neumann ordinals, which are defined by induction on \mathbb{N} . While it is easy to show $\check{\omega}$ satisfies the axiom of infinity

```
def axiom_of_infinity_spec (u : bSet  $\mathbb{B}$ ) :  $\mathbb{B}$  :=
  ( $\emptyset \in^{\mathbb{B}} u$ )  $\cap$  ( $\bigcap i\_x, \bigcup i\_y, (u.\text{func } i\_x \in^{\mathbb{B}} u.\text{func } i\_y)$ )
```

it can furthermore be shown to satisfy the universal property of ω , which says that ω is a subset of any set which contains \emptyset and is closed under the successor operation $x \mapsto x \cup \{x\}$.

The axiom of powerset

► **Definition 4.** Fix a \mathbb{B} -valued set $x = \langle \alpha, A, b \rangle$. Let $\chi : \alpha \rightarrow \mathbb{B}$ be a function. The subset of x associated to χ is a \mathbb{B} -valued set defined as follows:

```
def set_of_indicator {x} ( $\chi : x.\text{type} \rightarrow \mathbb{B}$ ) :=  $\langle x.\text{type}, x.\text{func}, \chi \rangle$ 
```

The **powerset** $\mathcal{P}(x)$ of x is defined to be the following \mathbb{B} -valued set, whose underlying type is the type of all functions $x.\text{type} \rightarrow \mathbb{B}$:

```
def bv_powerset (u : bSet  $\mathbb{B}$ ) : bSet  $\mathbb{B}$  :=
   $\langle u.\text{type} \rightarrow \mathbb{B}, \lambda f, \text{set\_of\_indicator } f, \lambda f, \text{set\_of\_indicator } f \subseteq^{\mathbb{B}} u \rangle$ 
```

The axiom of choice

Following Bell, we verified Zorn’s lemma, which is provably equivalent over ZF to the axiom of choice. As is the case with `pSet`, establishing the axiom of choice requires the use of a choice principle from the metatheory. This was the most involved part of our verification of the fundamental theorem of forcing, and relies on the technical tool of *mixtures*, which allow sequences of \mathbb{B} -valued sets to be “averaged” into new ones, and the *maximum principle*, which allows existentially quantified statements to be instantiated without changing their truth-value.

445 The smallness of \mathbb{B}

446 Before ending this section, we remark that the “smallness” (or more precisely, the fact that \mathbb{B}
 447 lives in the same universe of types out of which `bSet` \mathbb{B} is being built), is essential in making
 448 `bSet` \mathbb{B} a model of ZFC. It is required for extracting the witness needed for the maximum
 449 principle, and is also required to even define the powerset operation, because the underlying
 450 type of the powerset is the function type of all maps into \mathbb{B} .

451 4 Forcing

452 4.1 Representing Lean’s ordinals inside `pSet` and `bSet`

453 The treatment of ordinals in `mathlib` associates a class of ordinals to every type universe,
 454 defined as isomorphism classes of well-ordered types, and includes interfaces for both well-
 455 founded and transfinite recursion. Lean’s ordinals may be represented inside `pSet` by defining
 456 a map `ordinal.mk : ordinal → pSet` via transfinite recursion; it is nothing more than
 457 the von Neumann definition of ordinals. In pseudocode,

```
458 def ordinal.mk : ordinal → pSet
459 | 0 := ∅
460 | succ ξ := pSet.succ (ordinal.mk ξ) -- (mk ξ ∪ {mk ξ})
461 | is_limit ξ := ⋃ η < ξ, (ordinal.mk η)
```

464 Composing by `check (??)` yields a map `check ∘ ordinal.mk : ordinal → bSet` \mathbb{B} . (We
 465 could just as well defined `ordinal.mk' : ordinal → bSet` \mathbb{B} analogously to `ordinal.mk`
 466 such that `ordinal.mk' = check ∘ ordinal.mk`; the point is that there is a link between
 467 the metatheory’s notion of size and order with that of the forcing extension.)

468 Cardinals in Lean are defined separately from ordinals as bijective equivalence classes
 469 of types, but are canonically represented by ordinals which is are not bijective with any
 470 predecessor. We let `aleph : ordinal → ordinal` index these representatives. For the rest
 471 of this section, unadorned alephs (e.g. “ \aleph_2 ”) will mean either an ordinal of the form `aleph ξ`
 472 or a choice of representative from the isomorphism class of well-ordered types, and checked
 473 alephs (e.g. “ \aleph_2^\sim ”) will mean the `check ∘ ordinal.mk` of that ordinal.

474 4.2 The Cohen poset and the regular open algebra

475 Forcing with partial orders and forcing with complete Boolean algebras are related by the
 476 fact that every poset of forcing conditions can be embedded into a complete Boolean algebra
 477 as a dense suborder. This will be the case for our forcing argument: our Boolean algebra
 478 is the algebra of regular opens on $2^{\aleph_2 \times \mathbb{N}}$. The poset of forcing condition typically used for
 479 Cohen forcing embeds in this Boolean algebra as a dense suborder.

480 ► **Definition 5.** The **Cohen poset** for adding \aleph_2 -many Cohen reals is the collection of all
 481 finite partial functions $\aleph_2 \times \mathbb{N} \rightarrow 2$, ordered by reverse inclusion.

482 In the formalization, the Cohen poset is represented as a `structure` with three fields:

```
483 structure C : Type :=
484   (ins : finset (ℵ₂.type × ℕ))
485   (out : finset (ℵ₂.type × ℕ))
486   (H : ins ∩ out = ∅)
```

That is, we identify a finite partial function f with the triple $\langle f.\text{ins}, f.\text{out}, f.H \rangle$, where $f.\text{ins}$ is the preimage of $\{1\}$, $f.\text{out}$ is the preimage of $\{0\}$, and $f.H$ ensures well-definedness. While f is usually defined as a finite partial function, we found that in practice f is really only needed to give a finite partial specification of a subset of $\aleph_2 \times \mathbb{N}$ (i.e. a finite set $f.\text{ins}$ which *must* be in the subset, and a finite set $f.\text{out}$ which *must not* be in the subset), and chose this representation to make that information immediately accessible.

► **Definition 6.** Let X be a topological space, and for any open set U , let U^\perp denote the complement of the closure of U . The **regular open algebra** of a topological space X , written $\text{RO}(X)$, is the collection of all open sets U such that $U = (U^\perp)^\perp$, equipped with the structure of a complete Boolean algebra, with $x \sqcap y := x \cap y$, $x \sqcup y := ((x \sqcup y)^\perp)^\perp$, $\neg x := x^\perp$, and $\bigsqcup x_i := ((\bigcup x_i)^\perp)^\perp$.

The Boolean algebra which we will use for forcing $\neg\text{CH}$ is $\text{RO}(2^{\aleph_2 \times \mathbb{N}})$. Unless stated otherwise, for the rest of this section, we put $\mathbb{B} := \text{RO}(2^{\aleph_2 \times \mathbb{N}})$.

► **Definition 7.** We define the **canonical embedding** of the Cohen poset into \mathbb{B} as follows:

```
def  $\iota : \mathcal{C} \rightarrow \mathbb{B} := \lambda p, \{S \mid p.\text{ins} \subseteq S \wedge p.\text{out} \subseteq - S\}$ 
```

That is, we send each $c : \mathcal{C}$ all the subsets which satisfy the specification given by c . This is a clopen set, hence regular. Crucially, this embedding is *dense*:

```
lemma  $\mathcal{C\_dense} \{b : \mathbb{B}\} (H : \perp < b) : \exists p : \mathcal{C}, \iota p \leq b$ 
```

Recalling that \leq in \mathbb{B} is subset-inclusion, we see that this is essentially because the image of $\iota : \mathcal{C} \rightarrow \mathbb{B}$ is the standard basis for the product topology. Our chosen encoding of the Cohen poset also made it easier to perform this identification when formalizing this proof.

4.3 Adding \aleph_2 -many distinct Cohen reals

As we saw in Definition 4, for any \mathbb{B} -valued set x , characteristic functions into \mathbb{B} from the underlying type of x determine \mathbb{B} -valued subsets of x . While the ingredients \aleph_2 and \mathbb{N} for \mathbb{B} are types and thus external to $\text{bSet } \mathbb{B}$, they are represented nonetheless inside $\text{bSet } \mathbb{B}$ by their check-names $\check{\aleph}_2$ and $\check{\mathbb{N}}$, and in fact \aleph_2 is $\check{\aleph}_2.\text{type}$ and \mathbb{N} is $\check{\mathbb{N}}.\text{type}$. Given our specific choice of \mathbb{B} , this will allow us to construct an \aleph_2 -indexed family of distinct subsets of $\check{\mathbb{N}}$, which we can then convert into an injective function from $\check{\aleph}_2$ to $\check{\mathbb{N}}$, *inside* $\text{bSet } \mathbb{B}$.

► **Definition 8.** Let $\nu : \aleph_2$. For any $n : \mathbb{N}$, the collection of all subsets of $\aleph_2 \times \mathbb{N}$ which contain (ν, n) is a regular open of $2^{\aleph_2 \times \mathbb{N}}$, called the **principal open** $\mathbf{P}_{(\nu, n)}$ over (ν, n) .

► **Definition 9.** Let $\nu : \aleph_2$. We associate to ν the \mathbb{B} -valued characteristic function $\chi_\nu : \mathbb{N} \rightarrow \mathbb{B}$ defined by $\chi_\nu(n) := \mathbf{P}_{(\nu, n)}$. In light of our previous observations, we see that each χ_ν induces a new \mathbb{B} -valued subset $\widetilde{\chi}_\nu \subseteq \check{\mathbb{N}}$. We call $\widetilde{\chi}_\nu$ a **Cohen real**.

This gives us an \aleph_2 -indexed family of Cohen reals. Converting this data into an injective function from $\check{\aleph}_2$ to \mathbb{N} inside $\text{bSet } \mathbb{B}$ requires some care. One must check that $\nu \mapsto \widetilde{\chi}_\nu$ is externally injective, and this is where the characterization of the Cohen poset as a dense subset of \mathbb{B} (and moving back and forth between this representation and the definition as finite partial functions) comes in. Furthermore, one has to develop machinery similar to that for the powerset operation to convert an external injective function $\mathbf{x.type} \rightarrow \text{bSet } \mathbb{B}$ to a \mathbb{B} -valued set which $\text{bSet } \mathbb{B}$ thinks is a injective function, while maintaining conditions on the intended codomain. Our custom tactics and automation for reasoning inside \mathbb{B} made this latter task significantly easier than it would have been otherwise. We refer the interested reader to our formalization for details.

4.4 Preservation of cardinal inequalities

So far, we have shown for $\mathbb{B} = \text{RO}(2^{\aleph_2 \times \aleph})$ that $\mathbf{bSet} \ \mathbb{B}$ thinks \aleph_2^\vee is smaller than $\mathcal{P}(\aleph)$. Although Lean believes there is a strict inequality of cardinals $\aleph_0 < \aleph_1 < \aleph_2$, in general we can only deduce that their representations inside $\mathbf{bSet} \ \mathbb{B}$ are subsets of each other: $\top \leq \aleph_0^\vee \subseteq^{\mathbb{B}} \aleph_1^\vee \subseteq^{\mathbb{B}} \aleph_2^\vee$. To finish negating CH, it suffices to show that $\mathbf{bSet} \ \mathbb{B}$ thinks \aleph_0^\vee is strictly smaller than \aleph_1^\vee , and that $\mathbf{bSet} \ \mathbb{B}$ thinks \aleph_1^\vee is strictly smaller than \aleph_2^\vee . That is, we want that the passage from \aleph_i to \aleph_i^\vee preserves cardinal inequalities.

► **Definition 10.** For our purposes, “ X is strictly smaller than Y ” means “there exists no function f such that for every $y \in Y$, there exists an $x \in X$ such that $(x, y) \in f$ ”. Thus, “ X is strictly smaller than Y ” translates to the Boolean truth-value

$$\neg(\bigsqcup f, (\text{is_func } f) \sqcap \bigsqcup y, y \in^{\mathbb{B}} Y \implies \bigsqcup x, x \in^{\mathbb{B}} X \sqcap (x, y) \in^{\mathbb{B}} f).$$

The condition on an arbitrary \mathbb{B} which ensures the preservation of cardinal inequalities is the *countable chain condition*.

► **Definition 11.** We say that \mathbb{B} has the **countable chain condition** (CCC) if every antichain $\mathcal{A} : I \rightarrow \mathbb{B}$ (i.e. an indexed collection of elements $\mathcal{A} := \{a_i\}$ such that whenever $i \neq j, a_i \sqcap a_j = \perp$) has a countable image.

We sketch the argument that CCC implies the preservation of cardinal inequalities. The proof is by contraposition. Let κ_1 and κ_2 be cardinals such that $\kappa_1 < \kappa_2$, and suppose that \aleph_1^\vee is not strictly smaller than \aleph_2^\vee . Then there exists some $f : \mathbf{bSet} \ \mathbb{B}$ and some $\Gamma > \perp$ such that $\Gamma \leq (\text{is_func } f) \sqcap \bigsqcup y, y \in^{\mathbb{B}} \aleph_1^\vee \implies \bigsqcup x, x \in^{\mathbb{B}} \aleph_2^\vee \sqcap (x, y) \in^{\mathbb{B}} f$. Then one can show:

```
lemma AE_of_check_larger_than_check :
  ∀ β < κ₂, ∃ η < κ₁, ⊥ < (is_func f) ∩ (η̃, β̃) ∈B f
```

The name of this lemma emphasizes that what was happened here is that, given this f and the assumption that it satisfies some $\forall\text{-}\exists$ formula inside $\mathbf{bSet} \ \mathbb{B}$, we are able to extract, by virtue of \aleph_1^\vee and \aleph_2^\vee being check-names, a $\forall\text{-}\exists$ statement in the *metatheory*. Using Lean’s choice principle, we can then convert this $\forall\text{-}\exists$ statement into a function $g : \kappa_2 \rightarrow \kappa_1$, such that for every $\beta, \perp < (\text{is_func } f) \sqcap (g(\beta)^\vee, \beta^\vee) \in^{\mathbb{B}} f$. Since $\kappa_2 > \kappa_1$, it follows from the infinite pigeonhole principle that there exists some $\eta < \kappa_1$ such that the $g^{-1}(\{\eta\})$ is uncountable. Define $\mathcal{A} : g^{-1}(\{\eta\}) \rightarrow \mathbb{B}$ by $\mathcal{A}(\beta) := (\text{is_func } f) \sqcap (g(\beta)^\vee, \beta^\vee) \in^{\mathbb{B}} f$. This is an uncountable antichain because if $\beta_1 \neq \beta_2$, then the well-definedness part of $\text{is_func } f$ ensures that, since $g(\beta_1) = g(\beta_2)$, the truth-value $\beta_1^\vee = f(g(\beta_1)) \neq^{\mathbb{B}} f(g(\beta_2)) = \beta_2^\vee$ is \perp .

Thus, conditional on showing that $\mathbb{B} = \text{RO}(2^{\aleph_2 \times \aleph})$ has the CCC, we now have that cardinal inequalities are preserved in $\mathbf{bSet} \ \mathbb{B}$. Combining this with the injection $\aleph_2^\vee \rightarrow \mathcal{P}(\aleph)$, we obtain:

```
theorem neg_CH : ⊤ = (ℵ < (ℵ₁)̃ ∩ (ℵ₁)̃ < (ℵ₂)̃ ∩ (ℵ₂)̃ ≤ P(ℵ))
```

The arguments sketched in subsection 4.3 and subsection 4.4 form the heart of the forcing argument. Their proofs involve taking objects in $\text{Type } u$ and $\mathbf{bSet} \ \mathbb{B}$, constructing corresponding objects on the other side, and reasoning about them in ordinary and \mathbb{B} -valued logic simultaneously to determine cardinalities in $\mathbf{bSet} \ \mathbb{B}$. We have omitted many details from our discussion, but of course, all the proofs have been formally verified.

4.5 The unprovability of CH

We conclude this section by briefly describing how the previous results may be converted into a formal proof of the unprovability of CH. We work in a conservative expansion ZFC' of ZFC with an expanded language $L_{\text{ZFC}'}$ with symbols for pairing, union, powerset, and ω . We define ZFC' to be precisely the ZFC axioms which were verified in the fundamental theorem of forcing, along with specifications for the new function symbols. CH can then be written as a deeply-embedded $L_{\text{ZFC}'}$ sentence

```
def CH : sentence L_ZFC' := ¬ ∃' ∃' (ω < &1) ∧ (&1 < &0) ∧ (&0 ≤ P(ω))
```

where $<$ and \leq are abbreviations with the same meaning as in the previous section. Then proving $\text{bSet } \mathbb{B} \models \text{ZFC}' + \neg\text{CH}$ is a straightforward matter of checking that sentences are interpreted correctly as Boolean truth values which we have already proved to be \top . Applying the contrapositive of the Boolean-valued soundness theorem yields the result.

5 Transfinite combinatorics and the countable chain condition

What remains now is to prove that $\text{RO}(2^{\aleph_2 \times \aleph})$ has the CCC. There are several ways forward, but we chose the most general one, anticipating its usefulness in future formalizations of set theory and forcing.

5.1 The Δ -system lemma

► **Definition 12.** A Δ -system is... Lorem ipsum dolor sit amet, consectetur adipiscing elit. Donec molestie rutrum sapien et sagittis. Curabitur varius egestas tortor. Sed faucibus tincidunt felis, et tincidunt erat consequat eu. Praesent ullamcorper interdum ex in ornare. Vestibulum quam sem, molestie ac aliquam sit amet, efficitur non nunc. Suspendisse rutrum metus vitae ligula sagittis.

```
def is_delta_system {α : Type u} (A : set (set α)) :=
  ∃(root : set α), ∀{x y}, x ∈ A → y ∈ A → x ≠ y → x ∩ y = root
```

TODO(floris) Lorem ipsum dolor sit amet, consectetur adipiscing elit. Donec molestie rutrum sapien et sagittis. Curabitur varius egestas tortor. Sed faucibus tincidunt felis, et tincidunt erat consequat eu. Praesent ullamcorper interdum ex in ornare. Vestibulum quam sem, molestie ac aliquam sit amet, efficitur non nunc. Suspendisse rutrum metus vitae ligula sagittis, commodo lacinia metus ultricies. Sed ultricies fringilla magna ac luctus. Vivamus dolor mauris, vulputate in tempus dictum, faucibus non eros. Praesent aliquet, justo et tristique interdum, tellus sapien tempus sem, eu vestibulum nisl sem at ligula. Aenean commodo mauris nec leo pellentesque porttitor. Class aptent taciti sociosqu ad litora torquent per conubia nostra, per inceptos himenaeos. Duis interdum eget nulla ac molestie. Proin vel sem auctor, porttitor velit nec, cursus arcu.

```
theorem delta_system_lemma {α : Type u} {κ : cardinal}
  (hκ : cardinal.omega ≤ κ) {θ} (hκθ : κ < θ) (hθ : is_regular θ)
  (hθ_le : ∀(c < θ), c ^< κ < θ) (A : set (set α))
  (hA : θ ≤ mk A) (h2A : ∀{s : set α} (h : s ∈ A), mk s < κ) :
  ∃(B ⊆ A), mk B = θ ∧ is_delta_system B :=
```

629 Lorem ipsum dolor sit amet, consectetur adipiscing elit. Donec molestie rutrum sapien
 630 et sagittis. Curabitur varius egestas tortor. Sed faucibus tincidunt felis, et tincidunt erat
 631 consequat eu. Praesent ullamcorper interdum ex in ornare. Vestibulum quam sem, molestie
 632 ac aliquam sit amet, efficitur non nunc. Suspendisse rutrum metus vitae ligula sagittis,
 633 commodo lacinia metus ultricies. Sed ultricies fringilla magna ac luctus. Vivamus dolor
 634 mauris, vulputate in tempus dictum, faucibus non eros. Praesent aliquet, justo et tristique
 635 interdum, tellus sapien tempus sem, eu vestibulum nisl sem at ligula. Aenean commodo
 636 mauris nec leo pellentesque porttitor.

637 5.2 $RO(2^{\mathbb{N}_2 \times \mathbb{N}})$ has the countable chain condition

638 TODO(floris) Lorem ipsum dolor sit amet, consectetur adipiscing elit. Donec molestie
 639 rutrum sapien et sagittis. Curabitur varius egestas tortor. Sed faucibus tincidunt felis,
 640 et tincidunt erat consequat eu. Praesent ullamcorper interdum ex in ornare. Vestibulum
 641 quam sem, molestie ac aliquam sit amet, efficitur non nunc. Suspendisse rutrum metus
 642 vitae ligula sagittis, commodo lacinia metus ultricies. Sed ultricies fringilla magna ac luctus.
 643 Vivamus dolor mauris, vulputate in tempus dictum, faucibus non eros. Praesent aliquet,
 644 justo et tristique interdum, tellus sapien tempus sem, eu vestibulum nisl sem at ligula.
 645 Aenean commodo mauris nec leo pellentesque porttitor. Class aptent taciti sociosqu ad litora
 646 torquent per conubia nostra, per inceptos himenaeos. Duis interdum eget nulla ac molestie.
 647 Proin vel sem auctor, porttitor velit nec, cursus arcu.

```
648
649
650 theorem B_CCC : CCC regular_open_algebra :=
651 begin
652   simp[Suspendisse rutrum metus vitae ligula sagittis,
653     commodo lacinia metus ultricies,
654     Sed ultricies fringilla magna ac luctus,
655     Vivamus dolor mauris, vulputate in tempus dictum];
656   tidy
657 end
658
```

659 6 Related work

660 TODO - Note that the usefulness poset Yoneda lemma was discovered separately by Mario
 661 for making proofs in Metamath easier.

662 7 Conclusions and future work

663 Reflections on the proof

664 As our formalization has shown, for the purposes of a consistency proof, one can perform
 665 forcing entirely outside of the set-theoretic foundations in which forcing is usually presented.
 666 There is no need to work inside an ambient model of set theory, or to even have a ground
 667 model of set theory over which one constructs a forcing extension. Instead, the recursive
 668 *name* construction (generalizing the Aczel-Werner encoding) applied to a universe of types is
 669 the key. The type universe, with its classical two-valued metatheory and its own notion of
 670 ordinals, takes the place of the standard universe of sets. These external ordinals are then
 671 represented in the internal ordinals of the forcing extension by indexing the construction

of von Neumann ordinals. With a clever choice of forcing conditions \mathbb{B} , one can make this representation of ordinals send externally distinct cardinals to internally distinct cardinals, and then force an uncountable cardinal beneath $\mathcal{P}(\mathbb{N})$.

In particular, `pSet`, being only another special case of the construction which produces `bSet` \mathbb{B} , is no longer a prerequisite for working with `bSet` \mathbb{B} , but merely a convenient tool for organizing the check-names—this is the only role it played in the proof. The check-names themselves were actually not necessary either: as we remarked, the canonical map `ordinal` \rightarrow `bSet` \mathbb{B} can be defined without reference to them. However, since in all of our sources, `pSet` additionally played the role of the universe of types, and an interface for it was readily available in `mathlib`, we started our formalization by following the usual arguments, implementing these simplifications as we became aware of them.

Lessons learned

- Originally, we thought set-theoretic arguments involving transfinite/ordinal induction, which are ubiquitous, would be difficult to implement. In practice, Lean’s tools for well-founded recursion and the robust interface to ordinals in `mathlib` made the implementation of such arguments painless.
- Definitions and lemmas should be stated as generally as possible. This maximizes reusability, minimizes redundancy, and by exposing only the information required to complete the proof, improves the performance of automation.
- One should invest early in domain-specific automation. The formalization of the fundamental theorem was completed using only the first two strategies outlined in subsection 3.2; the calculations, while tedious, were recorded in our sources and it seemed easier to follow them. If we had followed through on the observations around Lemma 3 and developed the custom tactic library earlier, we would have saved a significant amount of time.

Towards a formal proof of the independence of the continuum hypothesis

The work we have described in this paper was undertaken as part of the Flypitch project, which aims to produce a formal proof of the independence of the continuum hypothesis. As such, the obvious next goal is a formalization of the consistency of the continuum hypothesis.

As indicated at the start of section 1, this means a formal proof of the independence of CH means showing that CH and \neg CH cannot be proved from the ZFC axioms. Although our work includes a formal proof of the unprovability of a version of CH from a version of the ZFC axioms in a conservative extension of the language of ZFC, verifying this was easy because we chose our axiomatization of ZFC to match exactly what we had proven about our model.

What is more interesting is formalizing the equivalence of various common formulations of ZFC and CH, so that a skeptical user may verify that their preferred version of CH is unprovable from their preferred version of ZFC. This would require formalizations of the conservativity of commonly-used extensions of ZFC, and of the equivalence of the various ways to say that one set is strictly smaller than another. The completeness theorem will be useful for this, because it constructs deeply-embedded proofs from ordinary proofs carried out in an arbitrary model; even the proof of the completeness theorem required certain conservativity results, which shows that we are equipped to handle such tasks.

Although the stated goal of our project is to achieve a formal proof of the independence of the continuum hypothesis, we are also interested in developing a reusable library for set theory and mathematical logic. We have completed a formalization of forcing, but we are nowhere near completing a library for forcing which a set theorist could use to verify their

research. Just as, more than 50 years ago, Cohen's proof marked only the beginning of modern set theory, a formal proof of the independence of the continuum hypothesis will mark only the beginning of formalized set theory.

8 References

TODO