# Soundness of reset workflow nets

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LiCS 2024 Tallin

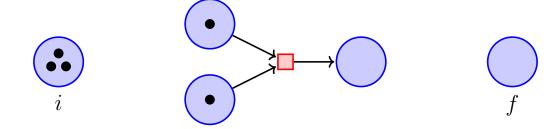
Petri nets for business processes [van der Aalst, 1997]

• Processes are initialised and finalised

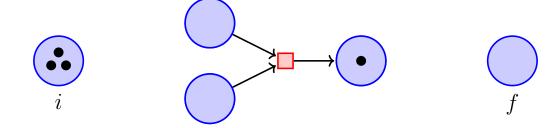




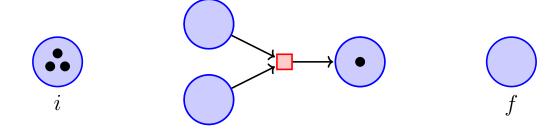
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- There are some transitions moving the tokens



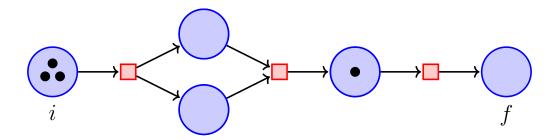
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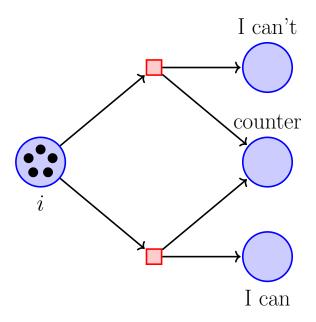
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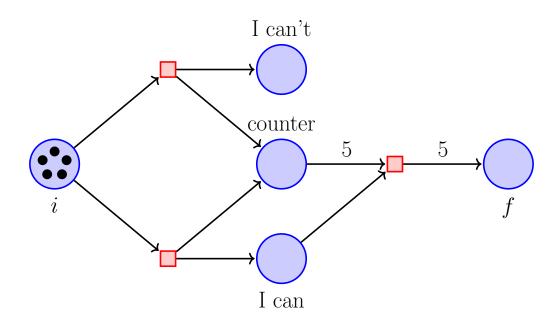


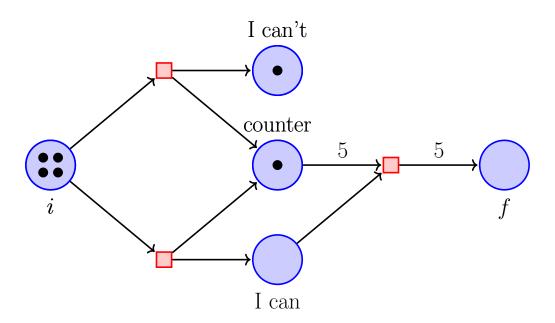
- Processes are initialised and finalised
- There are some transitions moving the tokens You cannot: add tokens to i and remove tokens from f
- $\bullet$  Everything is on a path from i to f



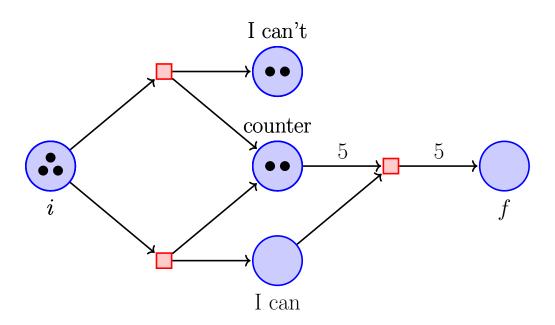






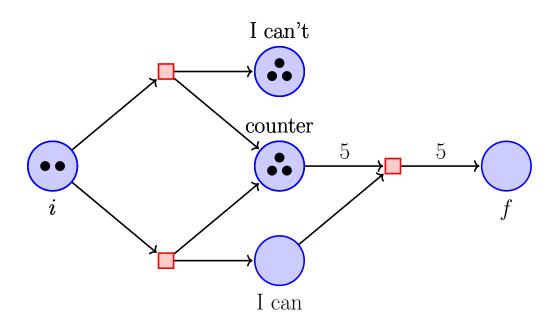


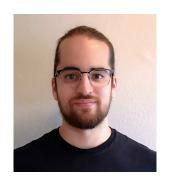






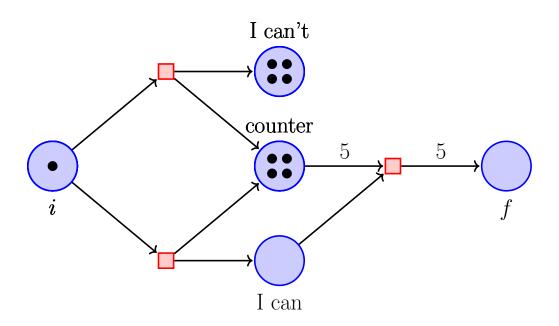










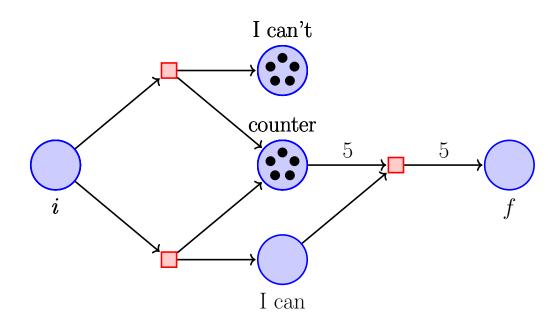














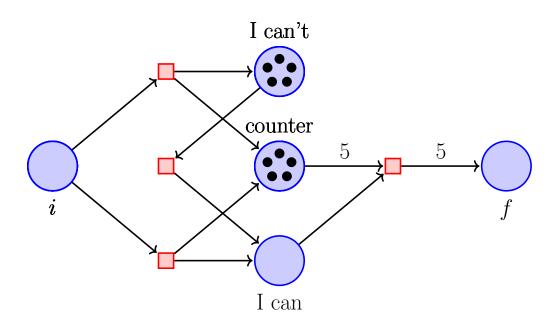








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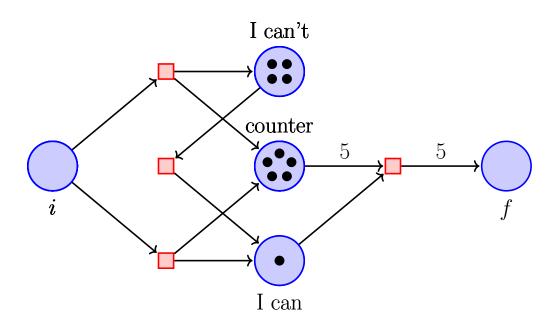














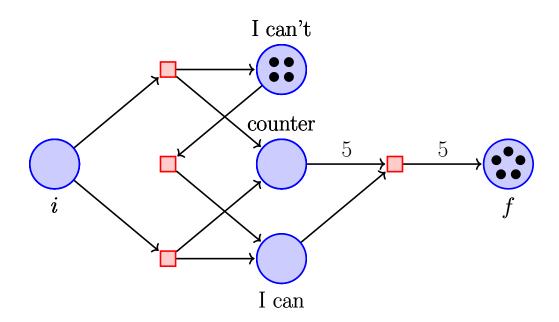








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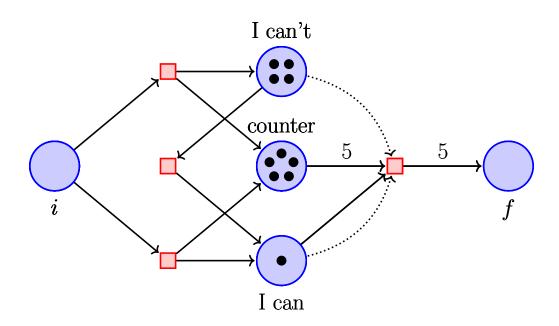








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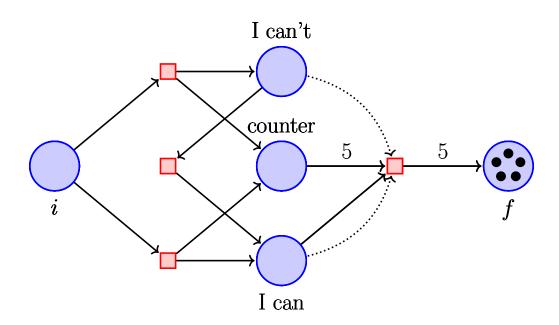








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(In Ackermann time)

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• What about workflow nets with a pushdown stack?