Implicit Type Conversions

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- What are Implicit Type Conversions (ITC)?
- ITC out there
- What do we want?
- Immediate concerns
- TrIC
- The approach taken
- In-depth ideas

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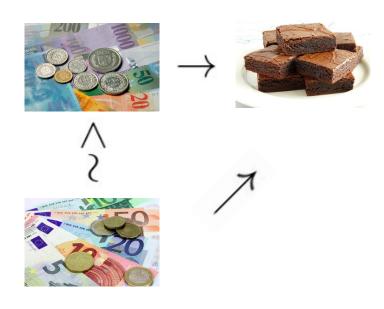
The idea

Base Types:









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ITC in Scala

```
implicit def conversion (s:String):Int = 3
def function (i:Int):String = "Woohoo!"
println(function("Oops"))
```

RESULT

Woohoo!

Polymorphic and constrained ITCs

```
implicit def conversion [A] (s:List[A]):A = a.head
```

```
implicit def conv (y:List[Int])(implicit cond:(Int => Char)):Char= cond(y.head)
```

```
implicit def conversion1 (s:String):Char = 'c'
implicit def conversion2 (c:Char): Int = 3
def function (i:Int):String = "Woohoo!"
println(function("Oops"))
ERROR
ScalaFiddle.scala:6: error: type mismatch;
 found : lang.this.String("Oops")
 required: scala.this.Int
   println(function("Oops"))
```

```
implicit def conversion (c:Char): Int = 3
def function (c:Char):String =
c match {
  case 3 => "Int 3"
 case 97 => "Int 97"
 case 'a' => "Char 'a'"
 case => "Other"
println(function('a'))
```

RESULT

Int 97

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This are the features we want to support

- Transparent extension (including type inference & type classes)
- Polymorphism
- Constrained Conversions
- User-friendliness
- Transitivity
- Local Scoping

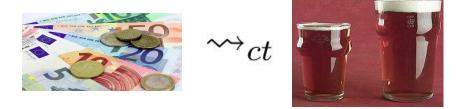
Transitivity



Local Scoping

- What are Implicit Type Conversions (ITC)?
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 - Ambiguity
- TrIC
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Ambiguity #1







Implicit Environment

Ambiguity #2





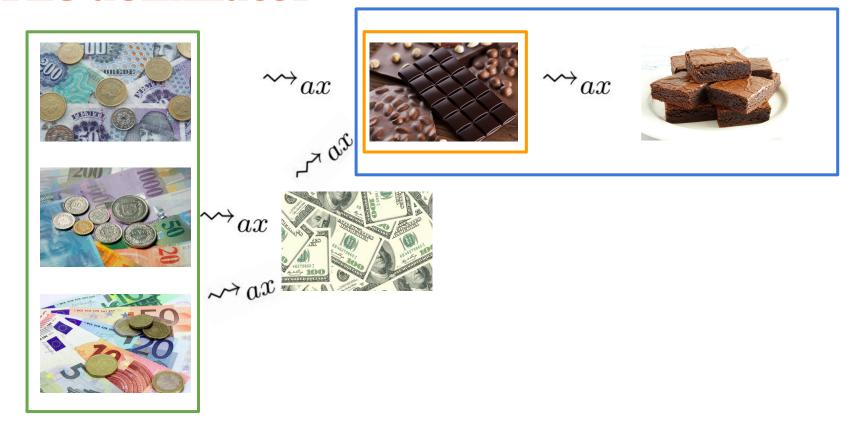
Attitude towards ambiguity

If there is no <u>obviously better</u> way to fix the program, don't.

The dominator



The dominator



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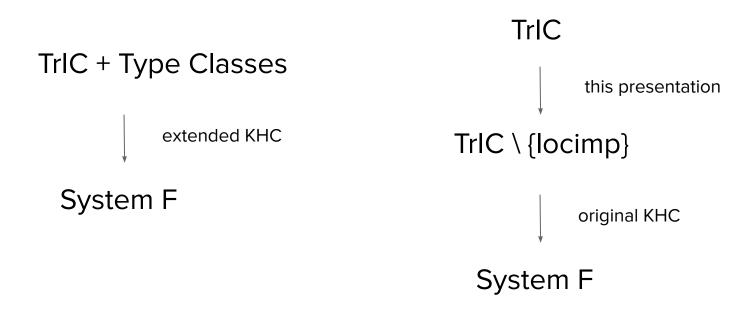
TrIC - syntax

e ::=

```
data EUR = K eur Float
data CHF = K chf Float
data DKK = K dkk Float
data USD = K usd Float
data Wallet (a :: *) = K wallet a a
(locimp \x. case x of K dkk y -> K eur (0.13*y) : DKK ~> EUR in
(locimp \x. case x of K chf y -> K eur (0.92*y) : CHF ~> EUR in
(locimp \x. case x of K eur y -> K usd (1.1*y) : EUR ~> USD in
(K_wallet (K_chf 5.25)
          (locimp \x.Pi : Int ~> Float in (K dkk 1)))
)))
```

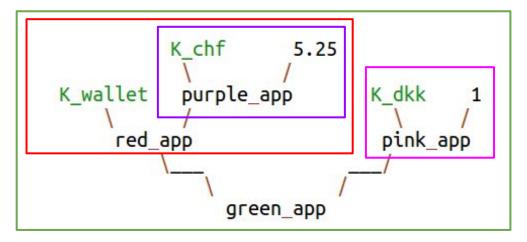
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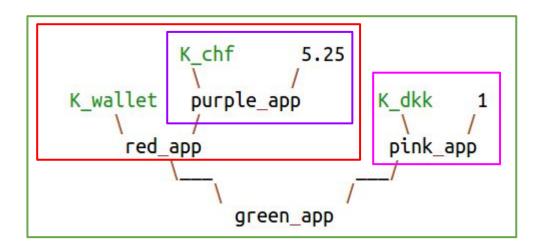
KHC-dialect = TrIC \ {locimp} + Type Classes





```
tm1 : X
                tm2 : Y
    (tm1 tm2): idk ; X ~ (Y -> idk)
                 tm1 : X
                 tm2 : Y
(tm1 tm2): idk; X ~ (arg -> idk); Y ~> arg
```





Equality constraints

Conversion constraints

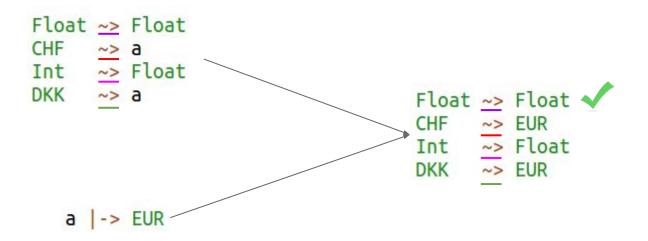
Equality constraints

```
(Float -> CHF ) ~ (ty_conv_from_Float -> ty_purple)
(a -> a -> Wallet a) ~ (ty_conv_from_ty_purple -> ty_red
ty_conv_from_Float
                         |-> Float
      ty purple
                         -> CHF
      ty_conv_from_ty_purple
                         |-> (a -> Wallet a)
      ty red
      ty conv from Int |-> Float
      ty pink
                        -> DKK
      ty conv from ty pink
                         -> Wallet a
      ty green
```

```
Float ~> ty_conv_from_Float
ty puple -> ty conv from ty purple
ty_pink ~> ty_conv_from_ty_pink
                                                  Float ~> Float
                                                  CHF
                                                       ~> Float
ty_conv_from_Float
                      -> Float
                                                  Int
                                                  DKK
                                                       ~> a
ty purple
                      -> CHF
ty conv from ty purple
                      -> a
ty red
                      -> (a -> Wallet a)
ty conv from Int
                      -> Float
ty pink
                      -> DKK
ty conv from ty pink
                      -> a
                      -> Wallet a
ty green
```

Compute the dominator of the conversion constraints involving each type variable

Implicit environment associated with both \cong and \cong :



```
CHF ~> EUR
```

```
CHF CHF Implicit Environment
```

```
data EUR = K eur Float
data CHF = K chf Float
data DKK = K dkk Float
data USD = K usd Float
data Wallet (a :: *) = K wallet a a
(locimp \x. case x of K dkk y -> K eur (0.13*y) : DKK ~> EUR in
(locimp \x. case x of K chf y -> K eur (0.92*y) : CHF ~> EUR in
(locimp \x. case x of K eur y -> K usd (1.1*y) : EUR \x USD in
(K wallet (K chf 5.25)
          (locimp \x.Pi : Int ~> Float in (K dkk 1)))
)))
```

```
K_wallet (K_chf 5.25)

K wallet ((\x. case x of K chf y -> K eur (0.92*y)) (K chf 5.25))
```

```
data EUR = K eur Float
data CHF = K chf Float
data DKK = K dkk Float
data USD = K usd Float
data Wallet (a :: *) = K wallet a a
(locimp \x. case x of K_dkk y -> K_eur (0.13*y) : DKK \xspace EUR in
(locimp \x. case x of K_chf y -> K_eur (0.92*y) : CHF ~> EUR in
(locimp \x. case x of K eur y -> K usd (1.1*y) : EUR ~> USD in
(K wallet (K chf 5.25)
          (locimp \x.Pi : Int ~> Float in (K dkk 1)))
)))
```

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Randomly choose a variable and compute the dominator?

```
data Wallet (a :: *) = K_wallet a a
            (K wallet
              (K wallet (K chf 5.25) (K dkk 2.5))
              (K wallet (K eur 6.5) (K usd 0.25)))
                          DKK ~,
CHF ~> a
DKK ~> a
                                         Wallet a ~> c
                          CHF ~> a
EUR ~> b
USD ~> b
                                          Wallet b
Wallet a ~> c
Wallet b ~> c
```

When is an axiom applicable?



```
JSD -> EUR -> DKK

Wallet a -> Wallet EUR

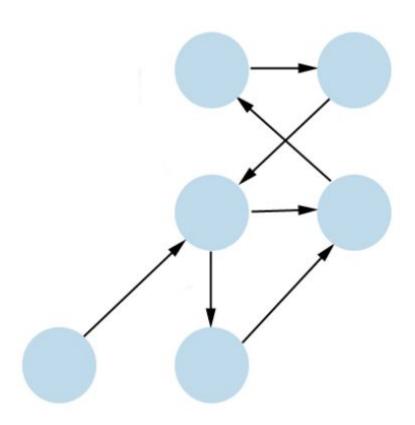
j:a -> b => Wallet a -> Wallet b

j:a -> EUK => Wallet a -> Wallet EUK

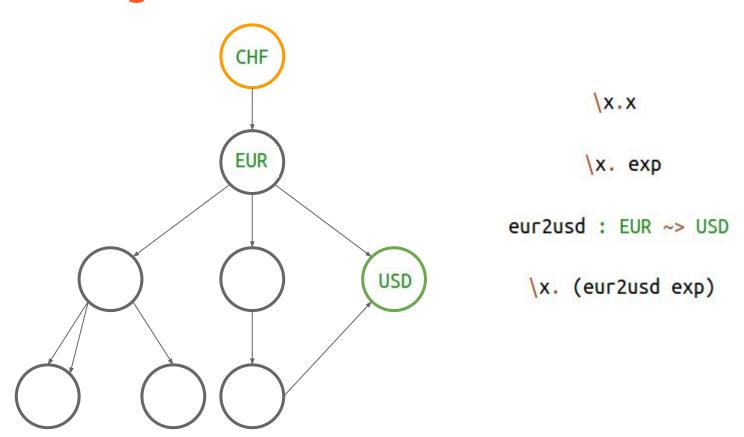
Implicit Environment
```



Paths



Constructing Paths



Type Classes

```
data DKK = K_dkk Float
data CHF = K_chf Float
data EUR = K_eur Float
data Wallet (a :: *) = K_wallet a a

class Total a :: * where
total :: a -> EUR

instance Total (Wallet EUR) where
total = \x. case x of K_wallet (K_eur x1) (K_eur x2) -> K_eur (x1+x2)

(locimp \x. case x of K_dkk y -> K_eur (0.13*y) : DKK ~> EUR in
total (locimp \x.case x of K_chf y -> K_eur (0.92*y) : CHF ~> EUR in K_wallet (K_dkk 6.3) (K_chf 3.14)))
```

