For each $e \in E$, define $\{\{e\}\} \subseteq M$ such that:

1. $\forall s \in M, s \in \{\{s\}\}$.
2. $\forall s, s' \in M$ if $s' \in \{\{s\}\}$ then $s' = s$.
3. if $e \notin W$ then $\{\{e\}\} = \emptyset$.
4. Special case 1, the Zugzwang paper stream:
   i. $e \supseteq s \Rightarrow s \in \{\{e\}\}$.
   ii. $e \subseteq s \Rightarrow s \in \{\{e\}\}$.
5. Special case 2, Levenshtein distance:
   i. We know that $x_0 = \min\{d_L(s, s') : s, s' \in M, s' \neq s\} > 1$. But need to prove.
   ii. $\{\{e\}\}_x = \{s \in M : d_L(e, s) \leq x\}$.
   iii. For $x = 0$ we get a distribution with all mass in $M$; what “others” are doing when define a probability on $M$.
   iv. We can set

$$\{\{e\}\}' = \{s \in M : d_L(e, s) = \min d_L(e, M)\}$$

   and then consider two cases:

   - If $\{\{e\}\}'$ has more than one element, $\{\{e\}\} = \emptyset$, or not.
6. Is there any relation with the Hamming distance and code correction?