

# Lecture 4.1

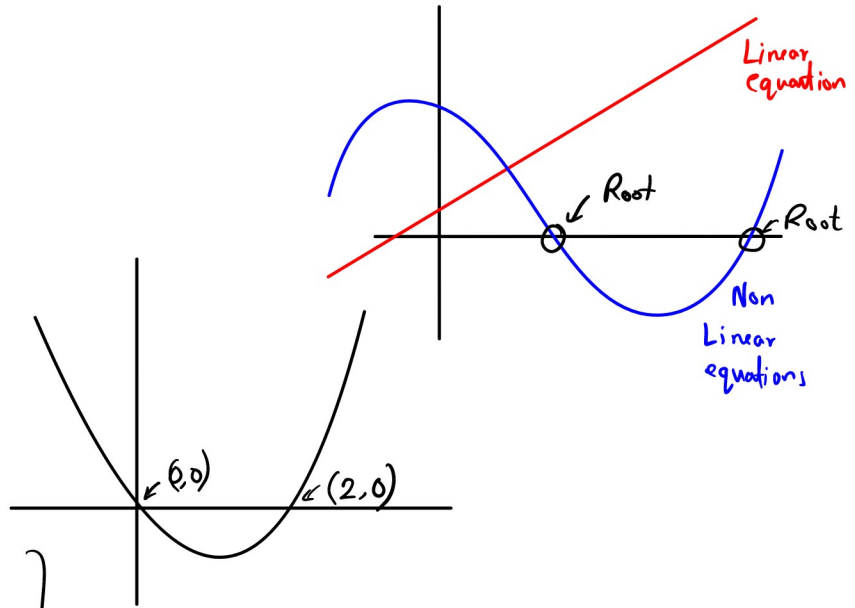
## Root finding for non linear equations:

$$f(x) = x^2 - 2x = 0$$

$$\Rightarrow x(x-2) = 0$$

$$\Rightarrow x = 0, 2$$

$$\text{Roots} \rightarrow x^* = 0, 2$$



# Polynomials degree  $\geq 2$

# Trigonometric functions (sin, cos)

# exponential " ( $e^x$ )

#  $\left(\frac{1}{x} - 2\right)$

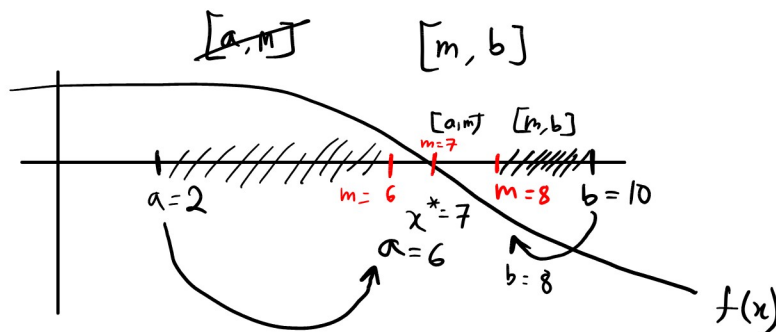
} Non linear function

## Bisection Root Finding Algorithm

Interval Bisection method

Given,

$f(x) = \sim$ ,  $[a, b]$ , find  $x^*$



$$f(m) = f(7) = 0$$

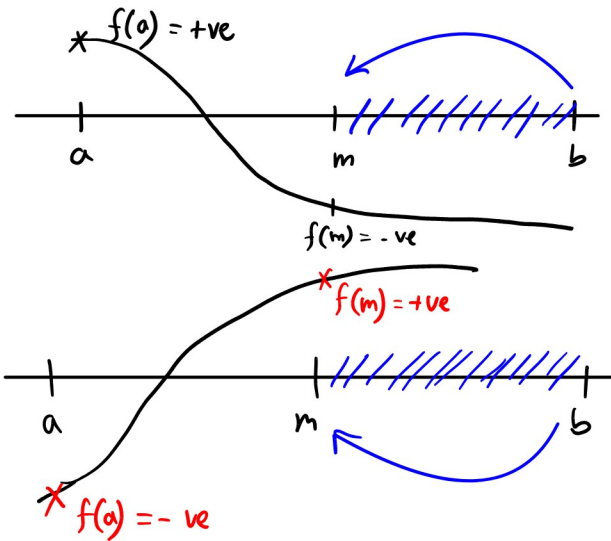
$m=7$  is the root.

$$m = \frac{a+b}{2} = \frac{2+10}{2} = 6$$

$$m = \frac{a+b}{2} = \frac{6+10}{2} = 8$$

$$m = \frac{a+b}{2} = \frac{6+8}{2} = 7$$

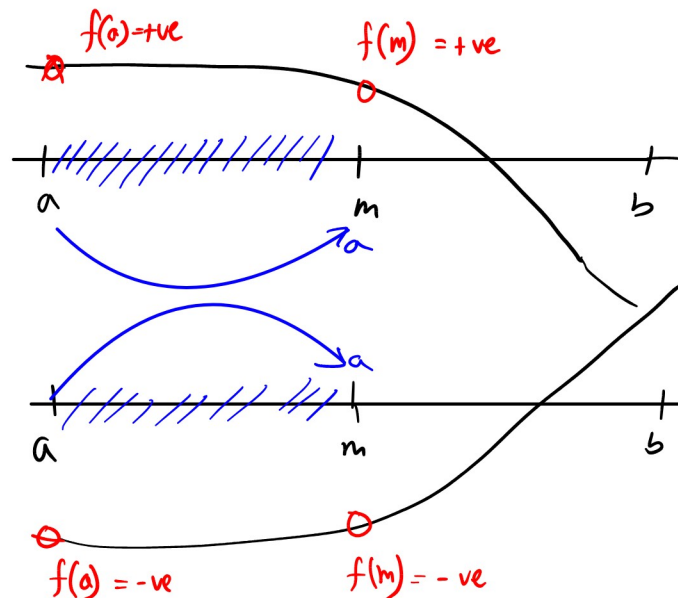
Case 1 : Root exists in  $[a, m]$



$f(a)$  and  $f(m)$   
has different sign.

$f(a) \times f(m) = -ve$   
means root exist in  
 $[a, m]$ ,  $b = m$

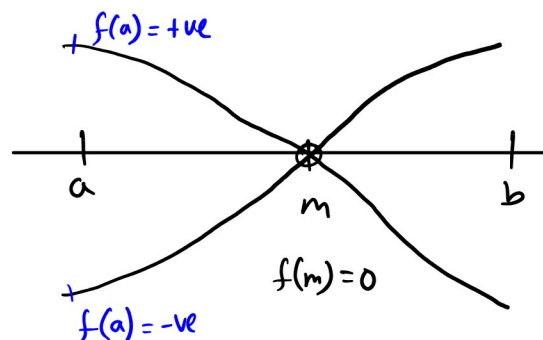
Case 2: Root doesn't exist in  $[a, m]$



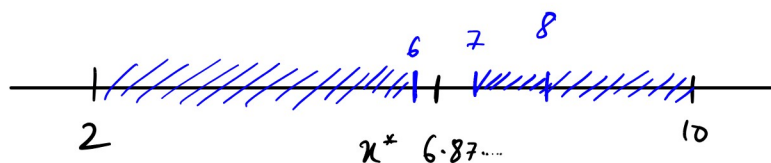
$f(a)$  &  $f(m)$  have  
similar sign.

$f(a) \times f(m) = +ve$   
root doesn't exist in  
 $[a, m]$ ,  $a = m$

Case 3:



if  $f(a) \times f(m) = 0$ ,  
then  $f(m) = 0$   
 $\therefore m$  is the root.



There is a limit for iteration. We do this until we find actual root,

or we reach machine epsilon  $\epsilon$ .  
 $\rightarrow$  Stopping Criteria.

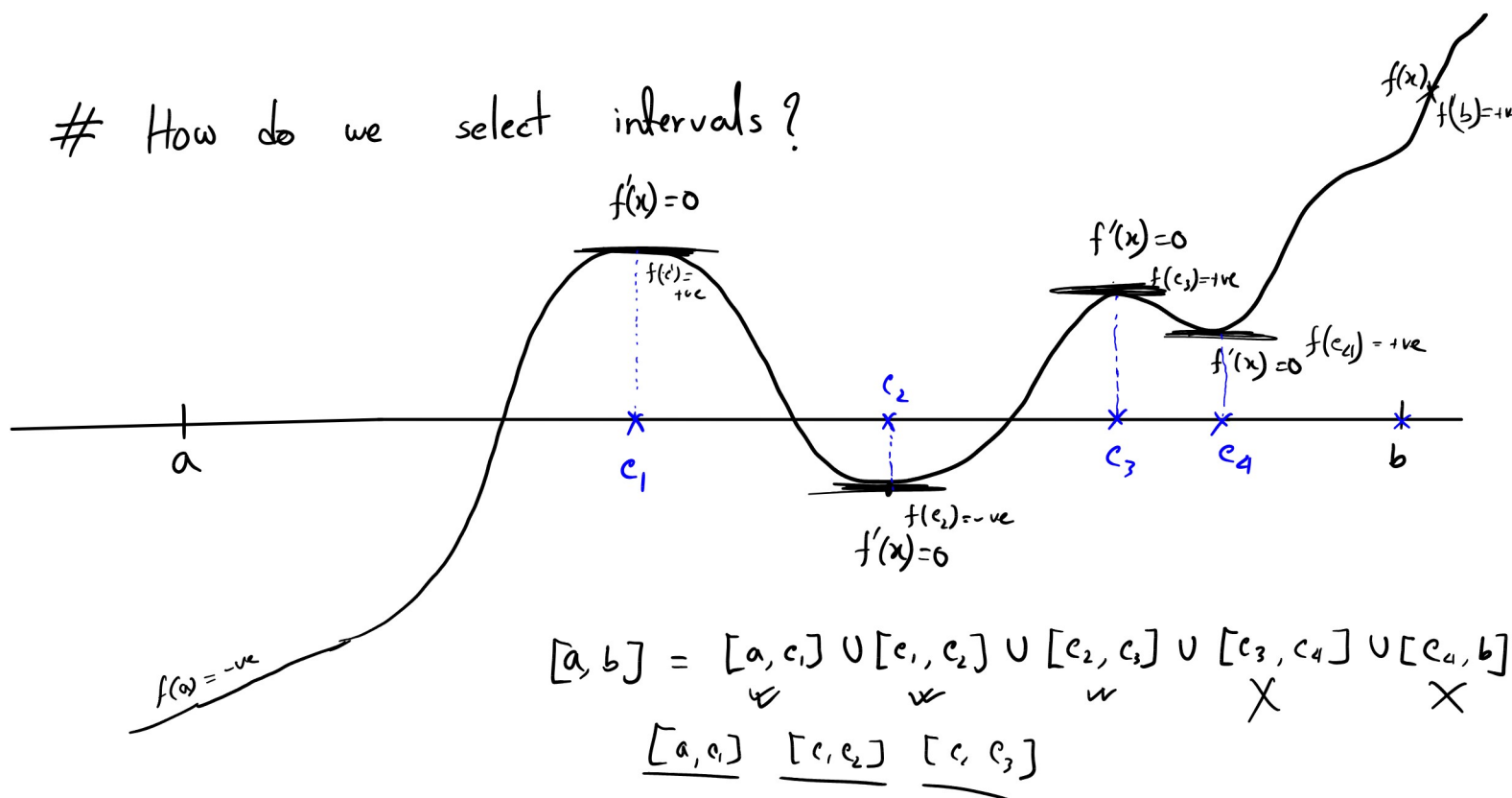
# Using Bisection method, find the root of

$f(x) = x^3 - 7x^2 + 14x - 6$  in interval  $[1, 3.2]$ . Your solution must be accurate within  $0.05$   
 $\epsilon$

Iteration	a	b	$m = \frac{a+b}{2}$	$f(a)$	$f(m)$	Root exist in $[a, m]$	new interval
0	1	3.2	2.1	2	1.79	Doesn't exist	$[m, b] = [2.1, 3.2]$
1	2.1	3.2	2.65	1.79	0.55	Doesn't exist	$[m, b] = [2.65, 3.2]$
2	2.65	3.2	2.925	0.55	0.086	Doesn't exist	$[m, b] = [2.925, 3.2]$
3	2.925	3.2	3.0625	0.086	-0.054	Exists	$[a, m] = [2.925, 3.0625]$
4	2.925	3.0625	2.99375	0.086	0.006	$\rightarrow < \epsilon$ 0.05	

$\therefore$  Root  $x^* = 2.99375$

# How do we select intervals?



# how many iteration is required for given intervals & epsilon values to reach the root?

$$[a, b], \quad \epsilon = 10^{-2}$$

$$\left. \begin{array}{l} \text{number of iteration} \\ \text{required} \end{array} \right\} k \geq \frac{\log(|b_0 - a_0|) - \log(\epsilon)}{\log(2)} - 1$$

Don't include if accuracy is specified.

# Find the number of iteration required to get the root with error bound of machine epsilon,  $\epsilon = 1.1 \times 10^{-6}$  in interval  $[1.5, 3]$   
 $a_0 \quad b_0$

$$k \geq \frac{\log(|b_0 - a_0|) - \log(\epsilon)}{\log(2)} - 1$$

$$k \geq \frac{\log(|3 - 1.5|) - \log(1.1 \times 10^{-6})}{\log(2)} - 1$$

$$k \geq 19.379$$

$$k \geq 20$$

$\therefore$  number of iteration required is 20.

# Find the number of iteration required to get the root with the accuracy of  $1.1 \times 10^{-6}$  in interval  $[1.5, 3]$   
 $a_0 \quad b_0$

$$k \geq \frac{\log(|b_0 - a_0|) - \log(1.1 \times 10^{-6})}{\log(2)}$$

$$k \geq 20.379$$

$$k \geq 21$$

Bisection number of iteration formula derivation:

For interval,  $a_0, b_0$ :

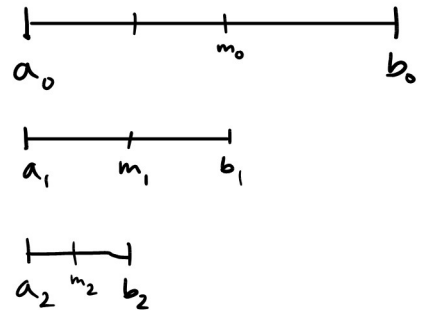
$$|b_1 - a_1| = \frac{|b_0 - a_0|}{2}$$

$$|b_2 - a_2| = \frac{|b_1 - a_1|}{2} = \frac{|b_0 - a_0|}{2^2}$$

$$|b_3 - a_3| = \frac{|b_2 - a_2|}{2} = \frac{|b_0 - a_0|}{2^3}$$

$\vdots$

$$|b_k - a_k| = \frac{|b_0 - a_0|}{2^k} \quad \text{--- ①}$$



After k iteration,

$$\text{Actual Error} = |m_k - x^*|$$

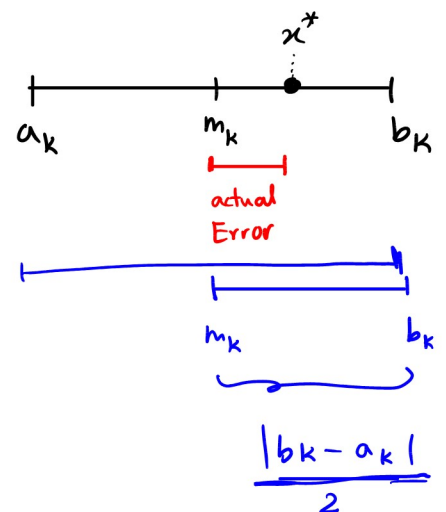
Intuitively,

$$\text{Actual Error} \leq \frac{|b_k - a_k|}{2}$$

$$\Rightarrow |m_k - x^*| \leq \frac{|b_k - a_k|}{2}$$

$$\Rightarrow |m_k - x^*| \leq \frac{|b_0 - a_0|}{2^k \cdot 2}$$

[ from equation 1 ]



$$\Rightarrow |m_k - x^*| \leq \frac{|b_0 - a_0|}{2^{k+1}}, \quad \text{for maximum possible error, } |m_k - x^*|_{\max} = \frac{|b_0 - a_0|}{2^{k+1}} \quad \text{--- (1)}$$

Since, we stop our iteration when actual error is less than given epsilon,

$$\text{Actual Error} \leq \epsilon$$

$$\Rightarrow |m_k - x^*| \leq \epsilon$$

$$\Rightarrow \frac{|b_0 - a_0|}{2^{k+1}} \leq \epsilon \quad [\text{from equation 2}]$$

$$\Rightarrow \frac{|b_0 - a_0|}{\epsilon} \leq 2^{k+1}$$

$$\Rightarrow \log\left(\frac{|b_0 - a_0|}{\epsilon}\right) \leq \log(2^{k+1})$$

$$[\log\left(\frac{x}{y}\right) = \log(x) - \log(y)]$$

$$[\log(x^m) = m \log(x)]$$

$$\Rightarrow \log(|b_0 - a_0|) - \log(\epsilon) \leq (k+1) \log(2)$$

$$\Rightarrow k+1 \geq \frac{\log(|b_0 - a_0|) - \log(\epsilon)}{\log(2)}$$

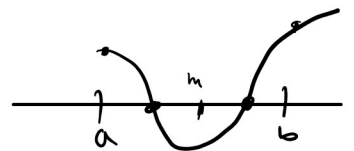
$$\therefore k \geq \frac{\log(|b_0 - a_0|) - \log(\epsilon)}{\log(2)} - 1$$

[showed]

Disadvantages of Bisection:

# Slow convergence.

# Can't find Multiple roots in single interval.



Advantages: Even though it is slow, it is more robust and guaranteed.