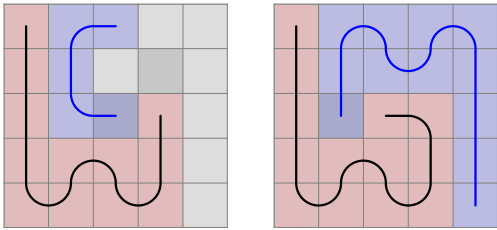
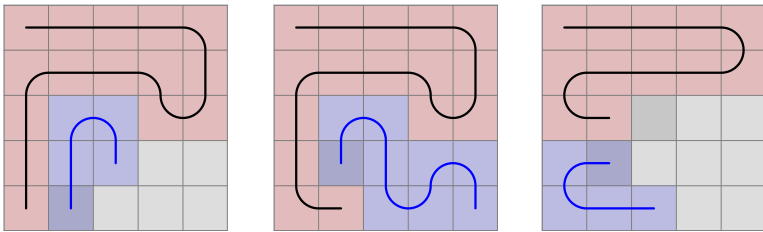


- A 2 x 2 empty area next to the live end that is walled by three sides (2-2-2 long) will have a future line going through along the walls. In this example, the far end is already extended by one step as it had only one option to move.



- Future line extension when we step on a future line: The far can be extended if it was 2 distance away from the near end. It can now fill the C-shape.



The same goes with 1 x- and y-distance. A C-Shape is not always created in this case.

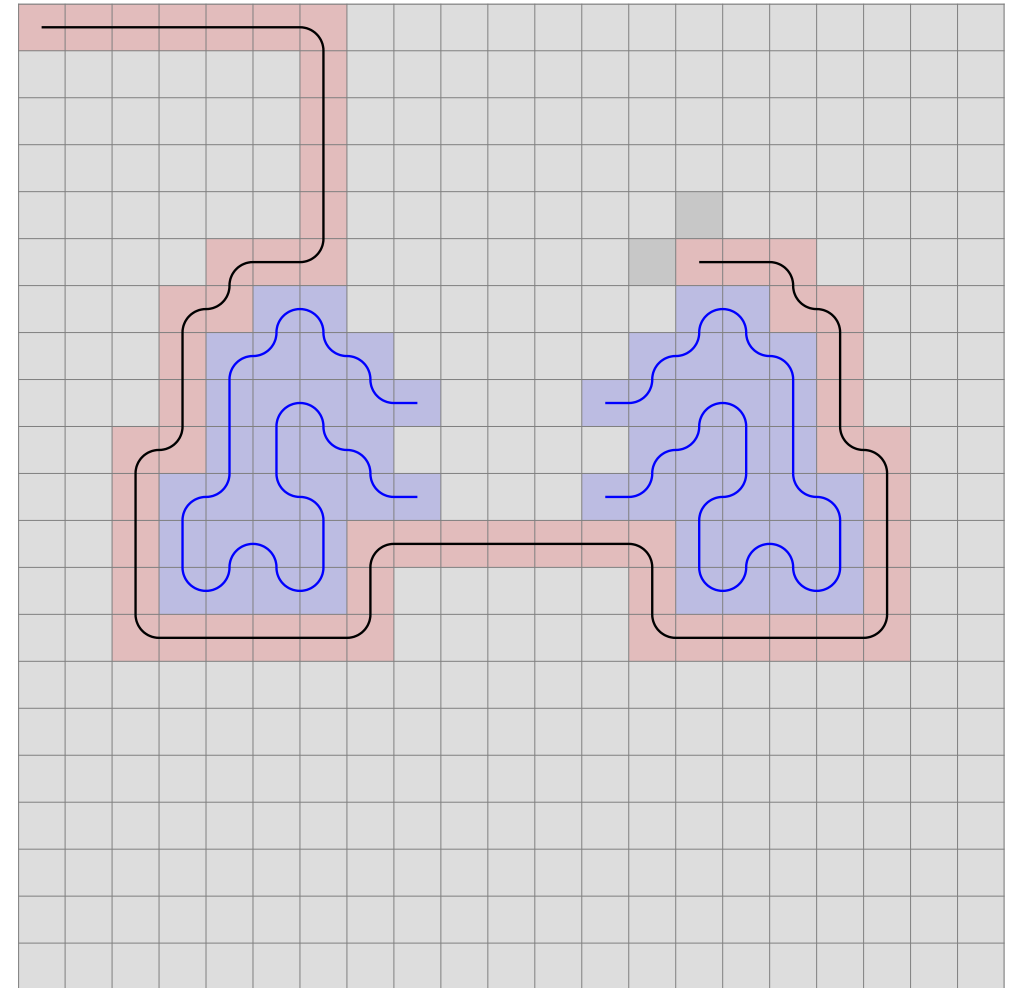
## The one-way labyrinth algorithm

([github.com/fodorbalint/OneWayLabyrinth](https://github.com/fodorbalint/OneWayLabyrinth))

This research aims to solve the following problem:

"Draw a line that goes through an  $n \times n$  grid (where  $n$  is an odd number), passing through each field once. The line has to start from the field at the upper left corner (1,1) and end at (n,n). At any time it is allowed to move left, right, up or down, and it has to randomly choose between the available fields."

At first sight it may look easy. But look at the following example:

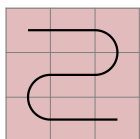


Based on the black line's movement, blue fragments were drawn to indicate a path we have to go through in the future in order to fill the board.  
Do you see why the situation is impossible from now on?

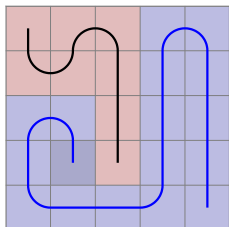
The question is, is there a single rule or a set of rules that will guarantee you can draw a labyrinth of any size? Or do the rules get infinitely complex?

To assist with the research, I have written a computer program. In the beginning, I let it run on a 21 x 21 field, and whenever I noticed a trouble, I coded the solution into it. While you can discover many patterns this way, they will be random and do not help in gaining a fundamental understanding. At one point you will find things get too complex, and you are still far from solving the 21 x 21 board. That's where a gradual approach comes in.

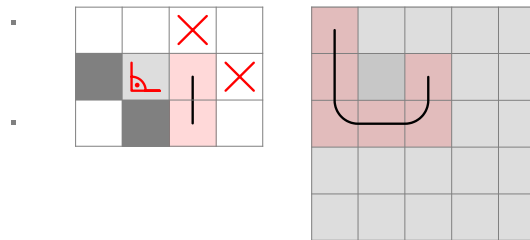
A 3 x 3 area can only be filled in two ways, like this and mirrored:



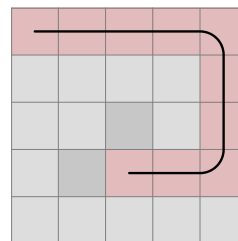
The 5 x 5 requires much more consideration. Whenever it is possible to draw future lines, the program has to be able to do it. The future lines can not only extend at each step but connect too.



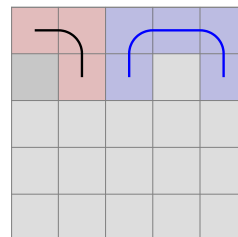
By August 21, 2023 all 5 x 5 scenarios were discovered. The number of walkthroughs are 104.  
Here are the things to consider on a grid of this size:



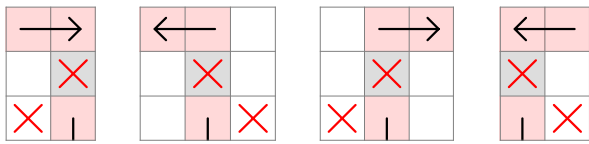
- A single field next to the live end that is walled from two other sides (either by the border or the line) needs to be filled in the next step. I call it C-shape. The pattern is both mirrored and rotated, so that the empty field is straight ahead. To qualify for this rule, the empty field cannot be the end corner. If there is a C-shape, we don't need to check other rules.



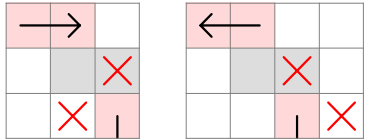
- Movement near the edge: In the example, we cannot step left (3,5), since the (2,5) field is empty.



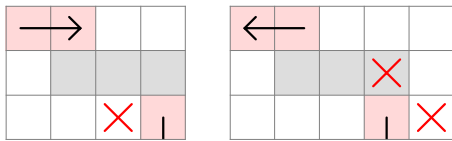
- A 2 x 3 empty area next to the live end that is walled by three sides (2-3-2 long) will have a future line going through along the walls. At the wall next to the main line, its direction is the opposite of the main line, meaning it will go from (3,2) upwards whereas the main line just took a step downwards. How the middle field will be filled is not yet known. Either the near end (the one the main line will go through first) or the far end can fill it.



The gray square means empty field. When the field 2 to straight is taken, its left or right side will be taken too.

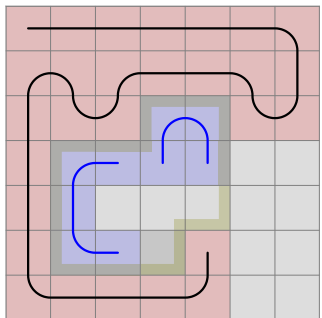


These will only be checked if one of the above 4 situations were not present. (They have to be mirrored, too.)

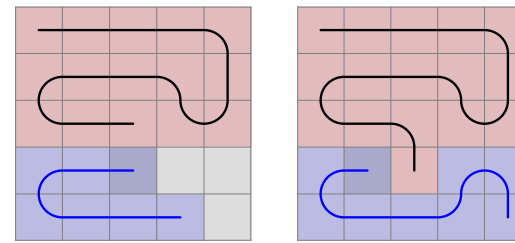


Likewise, these will be not be checked if the previous rules were true.

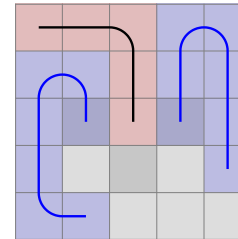
And when none of the 1-distance situations are valid, we check for 2-distance.



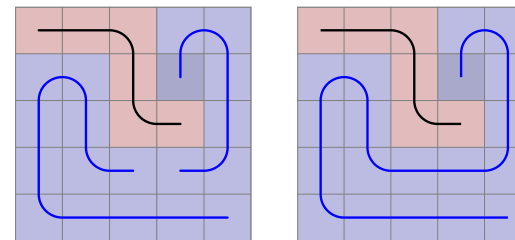
Impair areas can now happen inside the grid, not just on the edge, and the following rules have to be applied:



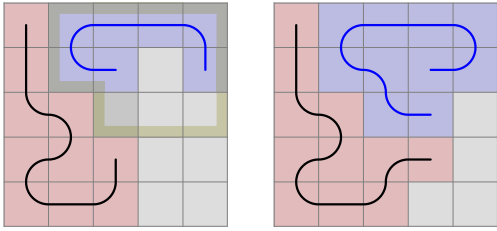
If the far end was near the end corner, it has to choose the other empty field.



- Future line extension when stepping away: If there was a near end where the main line was in the previous step, it now may have only one choice to move, so it can be extended.



- Future line connection: In this case, the line being stepped on extends until the far end has two options. (When the end corner is one of them, it has to be removed.) Then, the line on the left extends and now has no other option than to connect to the line on the right.

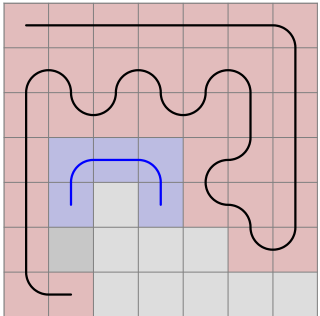


- When we are two distance away from the edge, we need to check if stepping towards it is possible.

It is because if we do so, an enclosed area is created, with one way to go out of it. If that area has an impair amount of cells, it cannot be filled, so we cannot take that step.

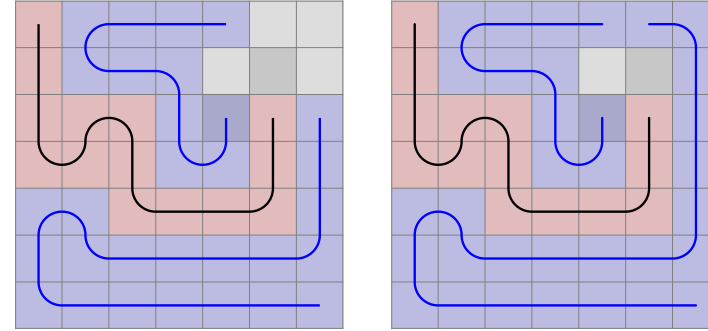
The explanation is simple: Imagine if the table was a chess board. In order to step from white to black, you would need to take an impair amount of steps - the color changes at every step. Here, the entry of the area would be (4,3) and the exit (5,3). An impair amount of steps means pair amount of cells.

In the example, you can also say that we cannot step right, because there is a future line start 2 to straight and an end 2 to straight and 2 to right. On 7 x 7, there will be examples where this is the rule we have to apply, because area counting is not getting triggered:

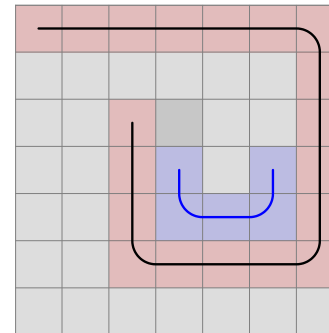


But let's start with the simpler rules:

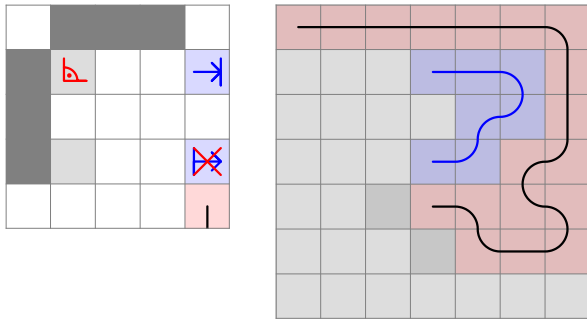
- Future line extension: When a near end is at 2 distance left or right from the live end, it will fill the field between them if the live end steps elsewhere. That's what happened in the 5 x 5 example above before the line failed.



In other situations, there is a 1-thin future line next to the live end that can be extended if its far end is at the corner. Though disabling this rule does not affect the total amount of walkthroughs on a 7 x 7 grid, I chose to include it in the project on the basis that if a future line can be extended, we should do it. It can make a considerable difference. The left picture is without the rule, the right is with it.



- Just like moving near the edge, we need to disable some fields if we are approaching an older section of the main line. In order to determine on which side the enclosed area is created, we need to examine the direction of the line at the connection point.



- And these are the remaining size-specific rules. Future 2 x 2 Start End, Future 2 x 3 Start End and Future 3 x 3 Start End.

The program, in fast mode, can run through approximately 100 cases per second, depending on your computer speed. This enables us to discover all 7 x 7 walkthroughs, which is 111 712.

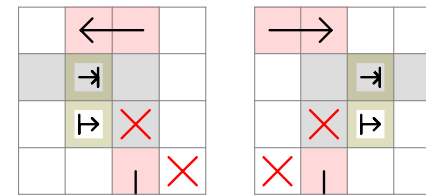
It is equal to what is described in the Online Encyclopedia of Integer Series (Number of simple Hamiltonian paths connecting opposite corners of a  $2n+1 \times 2n+1$  grid).

As the sizes grow, it will be impossible to run through all cases with one computer in a reasonable time. In order to discover the patterns, we need to run the program randomly.

Is it possible to develop an algorithm that works for all sizes? The edge-related and area-counting rules are universal, but the size-specific rules get more and more complex. Can you define them with one statement?

I have made statistics about how many random walkthroughs you can complete on different grids using the 7 x 7-specific and the universal rules before running into an error. Based on 1000 attempts, here are the results:

9: 19.5  
11: 5.7  
13: 2.6  
15: 1.2  
17: 0.7  
19: 0.4  
21: 0.2

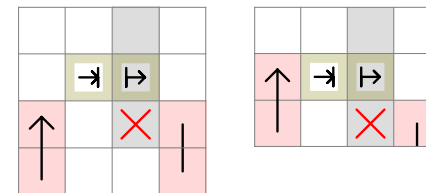


The procedure is similar to the the straight 2-distance rule. The only difference is that we count the area starting and ending at the marked fields. In the first, the direction of the circle is left, in the second right.

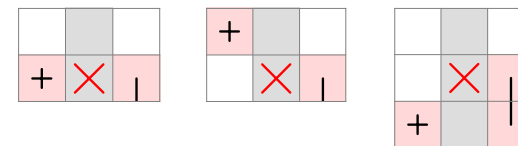
Besides mirroring them, we also have to rotate them both counter-clockwise and clockwise.

But we do not need 12 of such rules. Taking the first, the live end cannot come from the left, because the area parity was already checked in the previous step, and now we just added 2 fields to it. It can come from the right, and then there is naturally only one field we might have to disable.

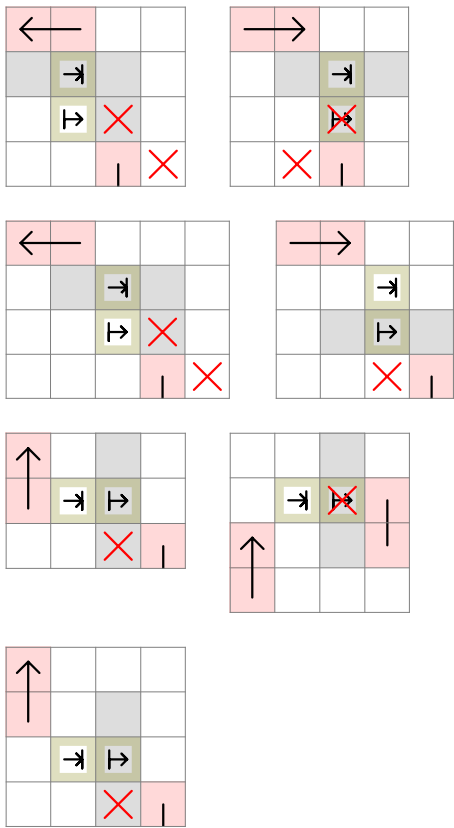
Here are the representations of the two scenarios for the left side:



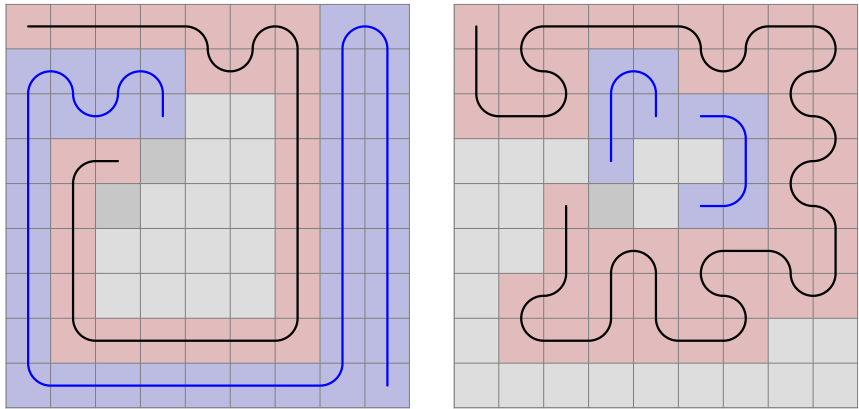
Similarly to the straight rules, these will only apply if there is no wall 2 distance to the left or right. Let's construct these preconditions.



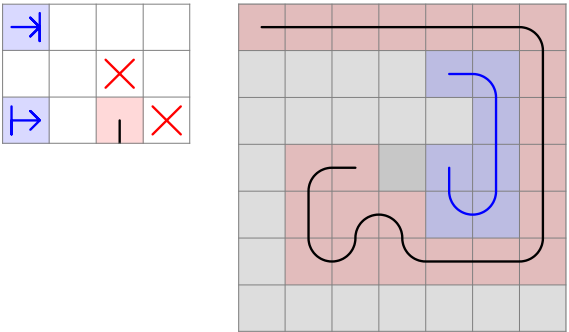
We are not finished. Did you notice the example above is not covered by these rules? We have to move the taken fields 1 and 2 steps to the side, both in straight and side direction.



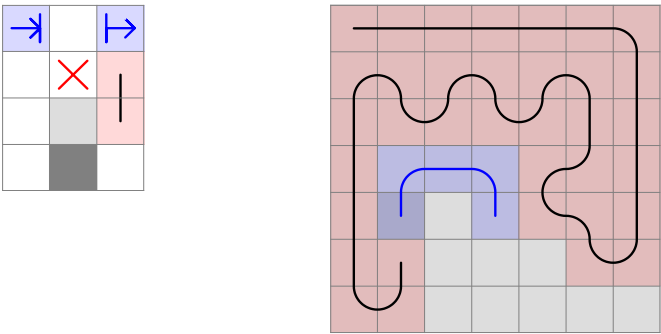
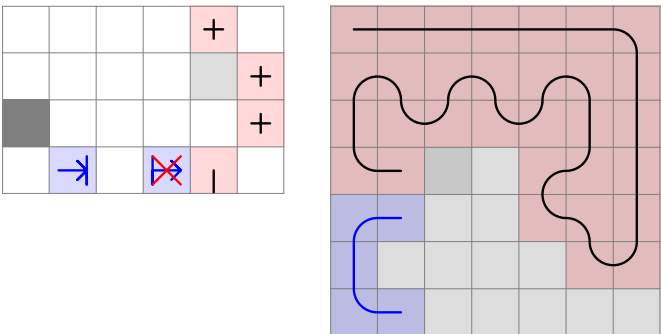
When any of the straight 2-distance rules are present, we don't need to check the side rules or the area created with the border. This is not entirely proven, but take these 9 x 9 examples:

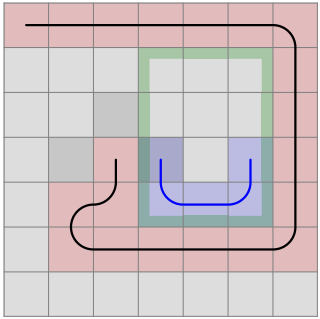


And these are the rest of the rules:



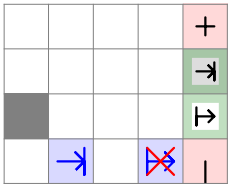
- This is what I started the 7 x 7 introduction with. I will call it Future L.





It would be a mistake to disable the right field.

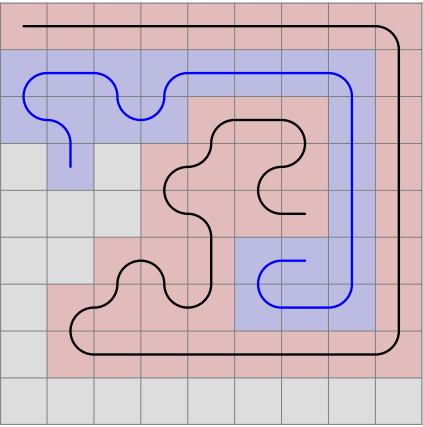
So we need to check if an enclosed has been created on that side, but counting the area is unnecessary. Nevertheless, we can represent the rule this way, setting the circle direction to right:



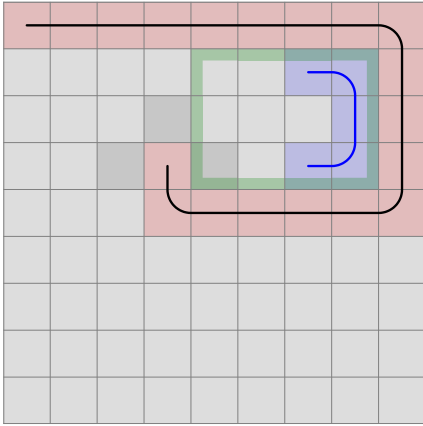
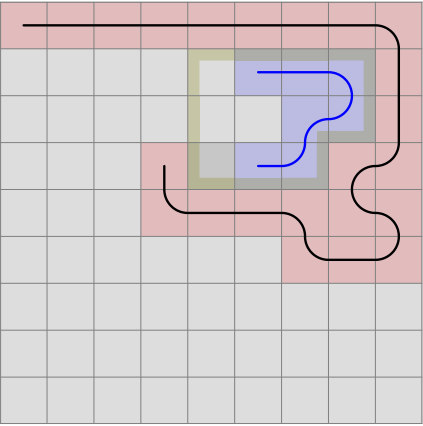
The code generator will examine if the count area start and end fields are 1 or 2 distance apart. In the first case, it will only determine in which direction the taken field straight ahead is going to, and if it is right, the forbidden field will take effect.

You may ask, why that field is "taken", not "taken or border". From what I found through some examples, if that field is border, the enclosed area on the right is impair, so the line cannot step in the other direction anyway. But it needs further examination.

To discover 9-specific patterns, I run the program keeping it left as long as the time to get to the first error is too big. After that, I will run it randomly. The first 13 826 walkthroughs are completed before we encounter a situation. It is similar to the last one we discovered on 7 x 7:



Let's simplify the pattern. Which will be impossible to fill?

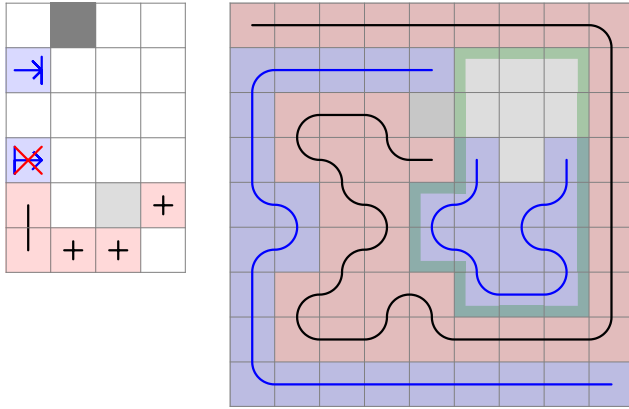


It is the picture on the left. Since the yellow-bordered area is impair, adding the (4,2) (4,3) (4,4) fields will be pair. We enter the area at (4,4), so we will exit at (4,3). Now we enter the 3 x 3 area in the top left corner at its side, (3,3) and will exit at (2,4). The results is two C-shapes on each side:





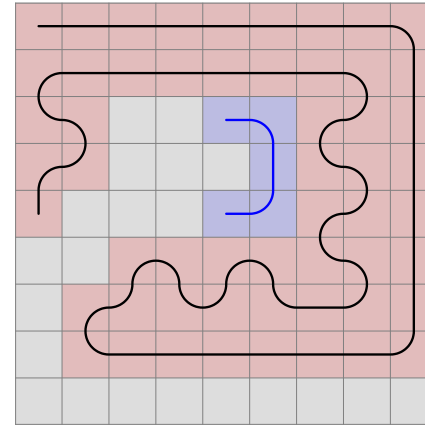
One certain situation reveals the incorrectness of the 7-rules when it comes to a 9-grid. In the following example, when I apply a rule rotated, it will disable a field that would otherwise be viable.



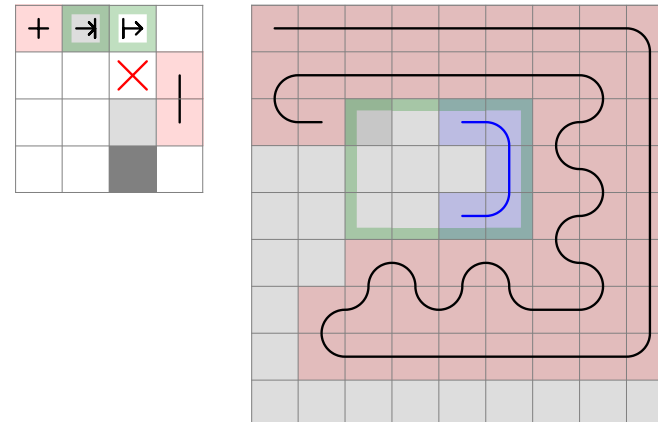
Rotating was not necessary to start with on 7 x 7, because no such situation occurred.

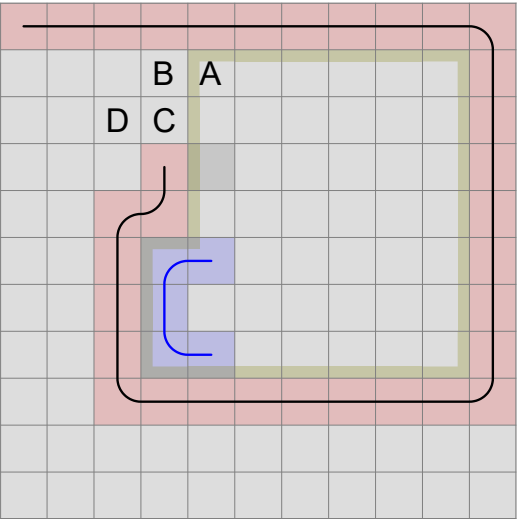
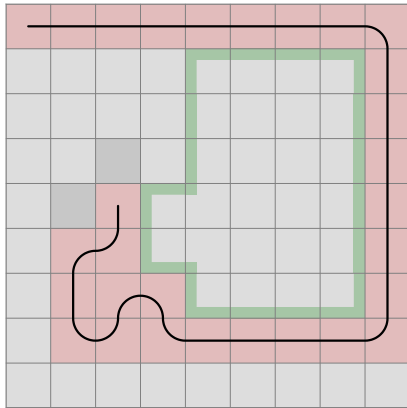
We can see that defining a rule with future line starts and ends does not tell us on which side the future line was created. That is the side that contains the enclosed area. We need to therefore replace such rules with area counting, which we actually already did, with the exception of Future L. Here the future line couldn't have been created on the other side, because that's the side the live end is at right now. And area counting is not always possible, like in this situation:

The next error, at 14 004 has something to do with how I defined the universal rules of approaching an older section of the line, it needs to be reworked in light of the C-shape the main line can create with the border.



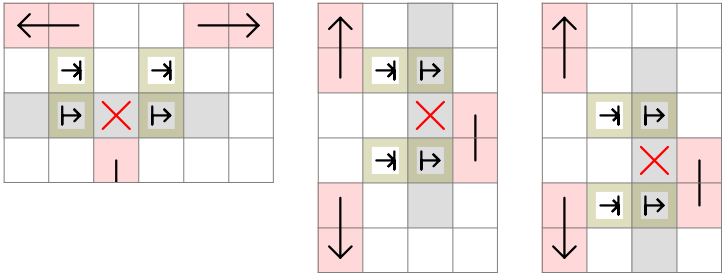
We need to take a few steps back, and then we can create the rule. It is similar to the universal 2-distance rule on the side, it just checks the field 2 behind and 1 to the side too. Even though the area counted is pair, now stepping to the right is disabled.



[illegible]

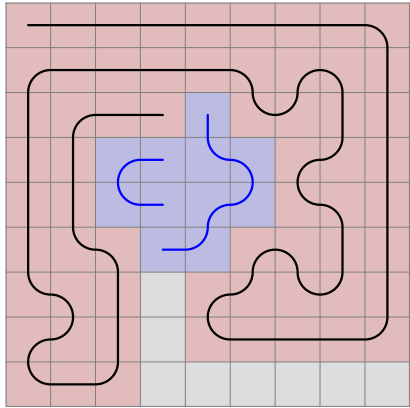
18 19

A 10x10 grid with a red path and blue obstacles. The red path starts at (0,0), goes right to (2,0), down to (2,2), right to (3,2), down to (3,4), right to (4,4), down to (4,6), right to (5,6), down to (5,8), right to (6,8), and finally down to (6,10). Blue obstacles are located at (0,2)-(1,3), (2,2)-(3,3), (3,2)-(4,3), (4,2)-(5,3), (5,2)-(6,3), (5,4)-(6,5), and (6,4)-(7,5).

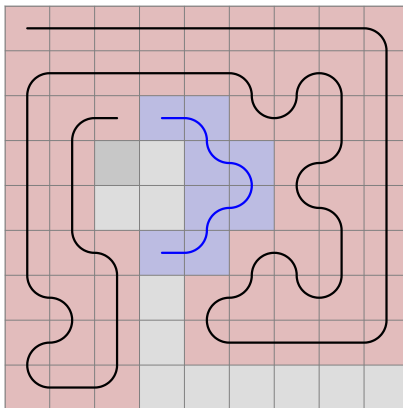
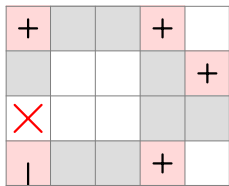


Any of the far straight rules (straight, mid across and across as I call them, depending on the horizontal distance of the obstacle) on the left side can be combined with any of those on the right side when the enclosed area is going to the same direction - left for left side and right for right side. And the same is true when the pattern is rotated to the left or right side. As far as programming concerned, it just needed a rework of the universal rules, we didn't need to make completely new ones.

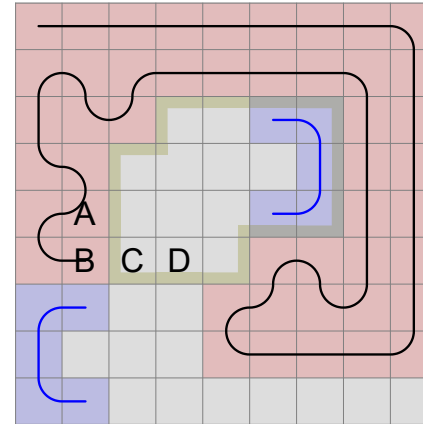
As we run the program further, we will discover this at 227 200:



Intuitively, we can draw up the square, and let's mark the exit as well. There can be loops on the upper, lower and right side, they have no importance when tracing it back to the live end. There is only one way to go through.

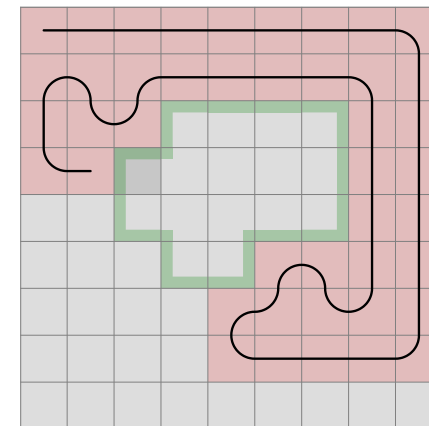
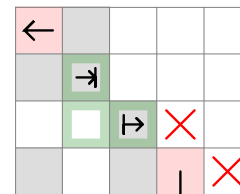


233 810 will look like:

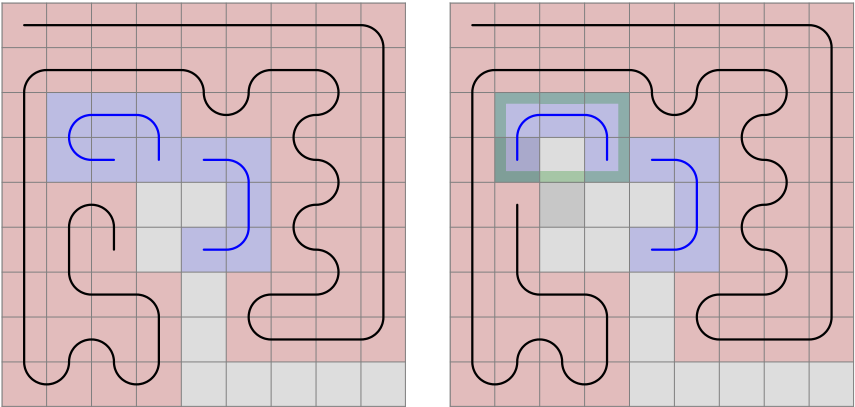


Once we step to A, it is unavoidable to get to B before entering the outlined area. It is because we can only reach B from the left or the bottom. The area is impair, therefore we cannot complete it starting in C and ending in D.

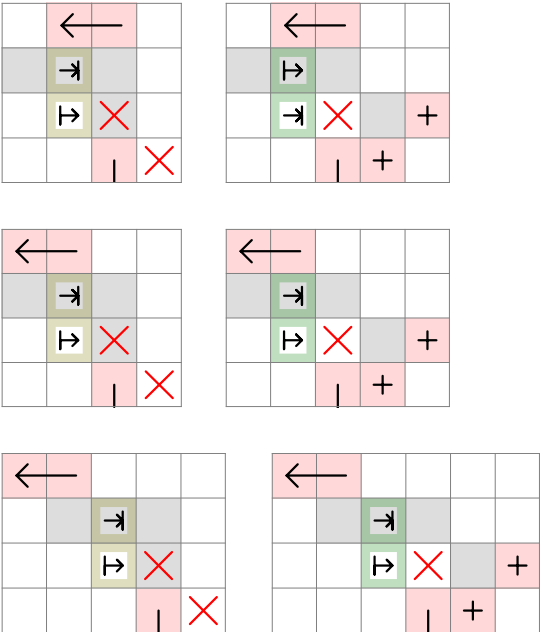
If we omit the C field from the area, the area becomes pair. It is clear that the start and end field being across each other, a pair amount of fields cannot be filled. We must therefore enter the area now.



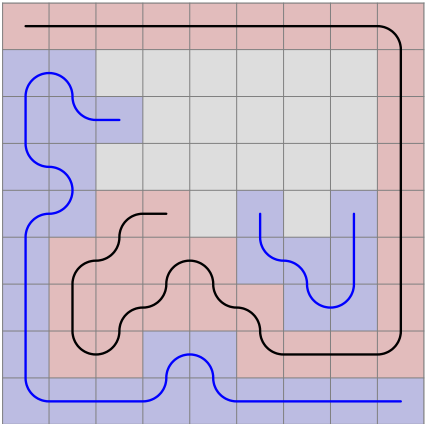
The same concept we encounter at 635 301, only the C-shape is created when we enter an area, on the other side of it.



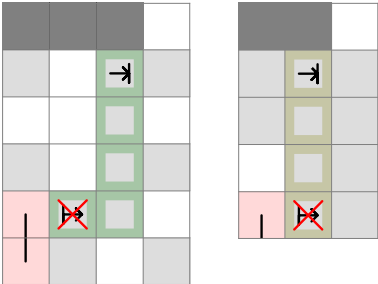
We have seen this in the third 9 x 9 rule. There the taken field next to the exit was in middle across position, and now it is across. And we also need to think about an obstacle straight ahead. Here are the original universal rules and their modifications.  
Straight, circle direction left:



At 349 215, we find this:

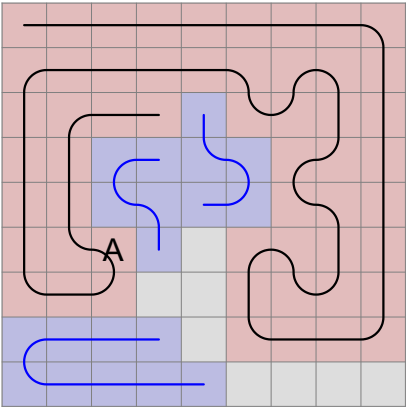


Though a double C-shape has been created in backwards direction, it indicates that the area on the right cannot be filled either.  
We have made a similar rule previously. Now we need to simplify it.

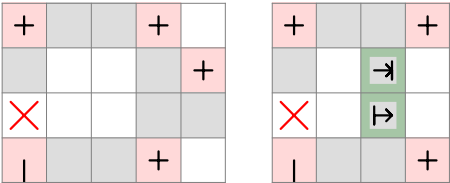


The area now has to be impair for the right direction to be forbidden.  
Essentially, we just added the three extra fields to the pair area.

478 361 is similar to what we have seen before, only now there is a 2-wide path to exit the area:

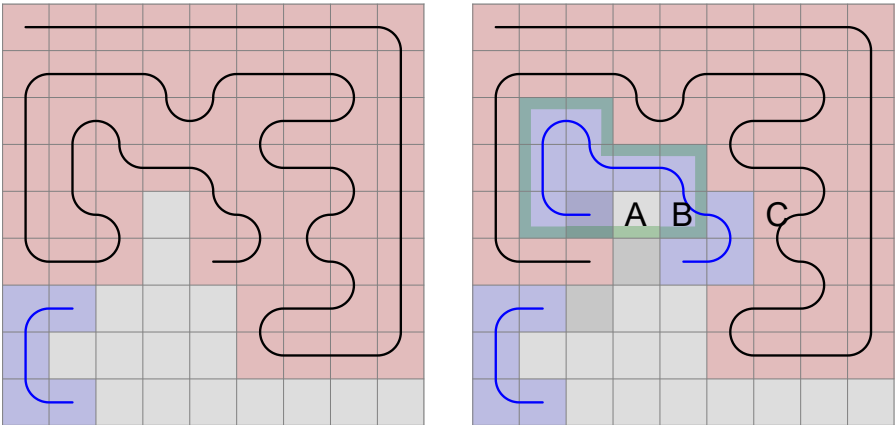


We have to mark where the area has been created in another way.

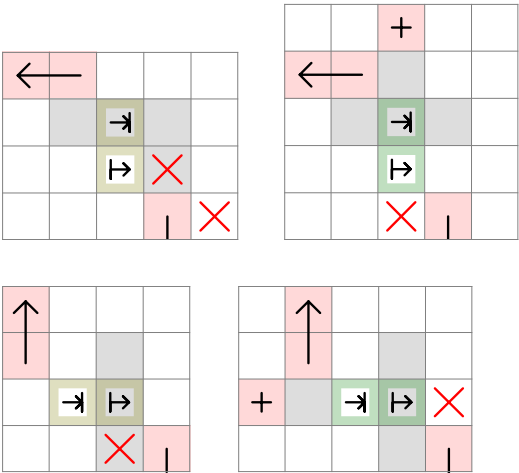


The taken field in the upper right corner is now checked for direction, but it is not enough. It can go upwards, and the exit of the area can still be on the bottom edge, just look at the example and imagine the live end was at A with the pattern already drawn. (On 11 x 11, it is possible to draw it.) In order to establish an enclosed area, we must not encounter the bottom-right corner of the grid when walking along the edge of it.

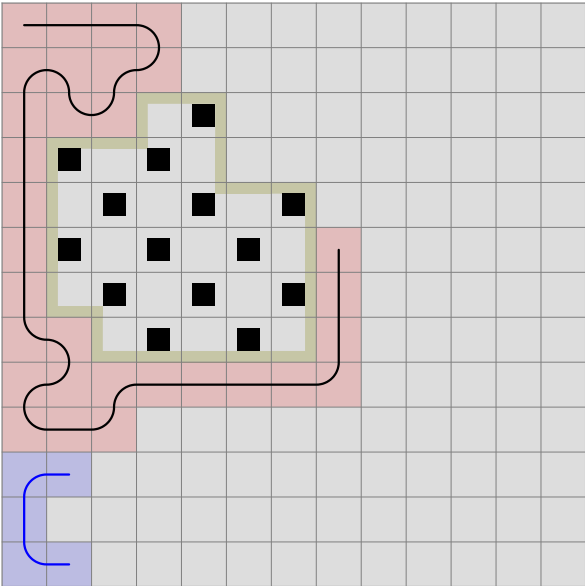
626 071 is:



With the marked area being pair, if we enter the area by stepping left, we will exit at A. But we can only get there from B; if we entered from the top, nothing would fill B, and we cannot enter and exit it after we left the area - subtracting 1 from the area would make it impair, so then we couldn't have exited at A. The taken field C creates a C-shape, which we need to step into from B. The universal far across rule have to be extended. By default, we disable the option to step straight or right if the counted area is impair. When it is pair, we need to disable the left field.

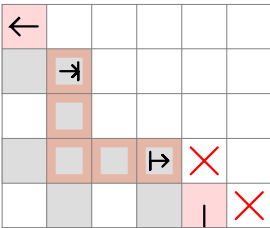


Let's mark the original example as a checkerboard.

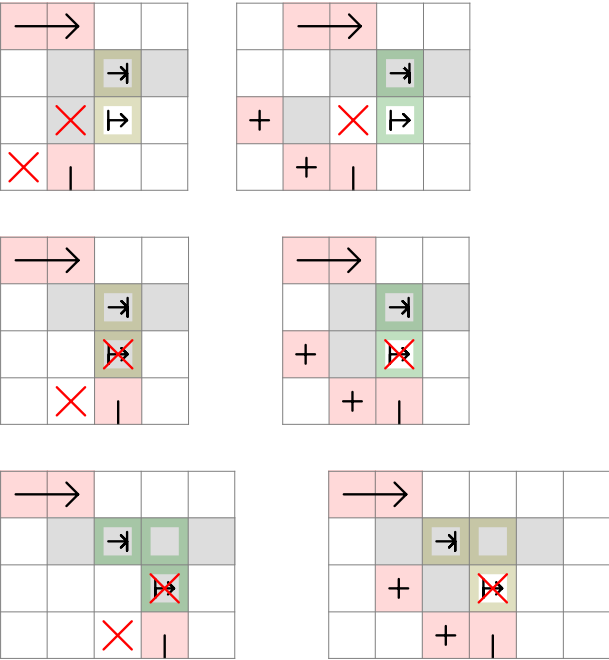


We enter at a black field and exit at black too, so the number of black fields should be one more than the number of white fields. Here there are 14 black fields and 15 white. That's why the area cannot be filled. The up and right directions need to be disabled, so we can only step left.

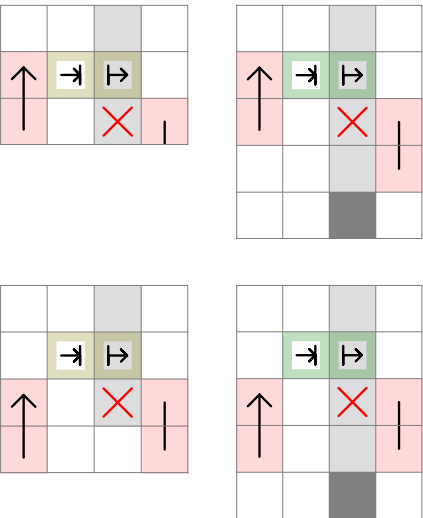
This is the rule representation. The reddish arealine now means the arealine is impair, and we know that the entry and exit points are the arealine start and end fields.

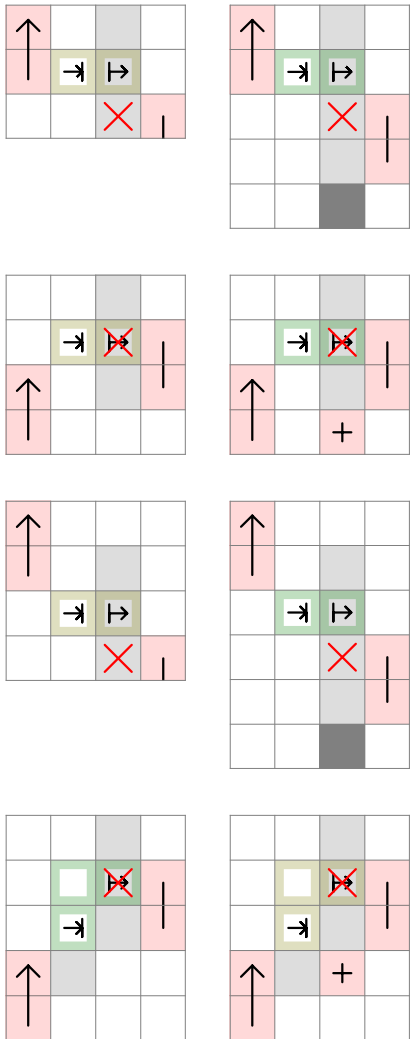


Circle direction right:



Side, with taken fields above and below:



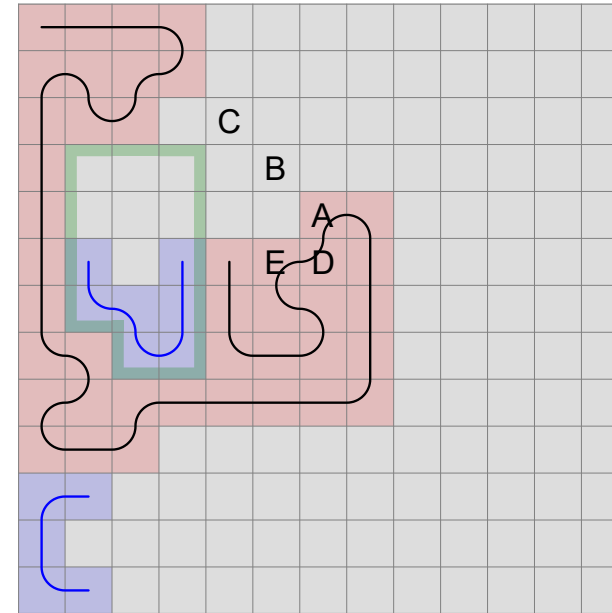


I have made some changes by adding some empty fields in side positions, so they are the same as the straight rules, just rotated.

Also, I have added the side across down rule and changed the straight across rule accordingly. Not only fields next to each other can define an area, they can be across too. In that case, if the area originally was marked impair, now it has to be pair.

Notice that in side rules, when the taken field that would create the C-shape is below the obstacle creating the area, it can be a border field too. We have seen an example of that previously.

Now what if both the start and end C-conditions are true? We can construct this on 13 x 13:



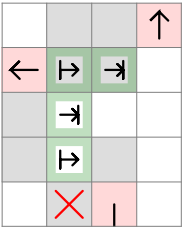
Several walkthrough attempts will leave you thinking why you cannot fill the area once obstacle responsible for the start C-shape is created (A). The area enclosed by A, B and C is pair. So when you enter it at A or B (obviously C is not a possibility), in order to exit at C, you need to leave out an impair amount of fields from the area. In case of entering at B, you cannot leave out A, but when you enter at A, you can leave out B, and no more than that. Now the area will be impair.

The minimal area would be stepping left from A, left again, up and up to get to C. You have covered 5 fields.

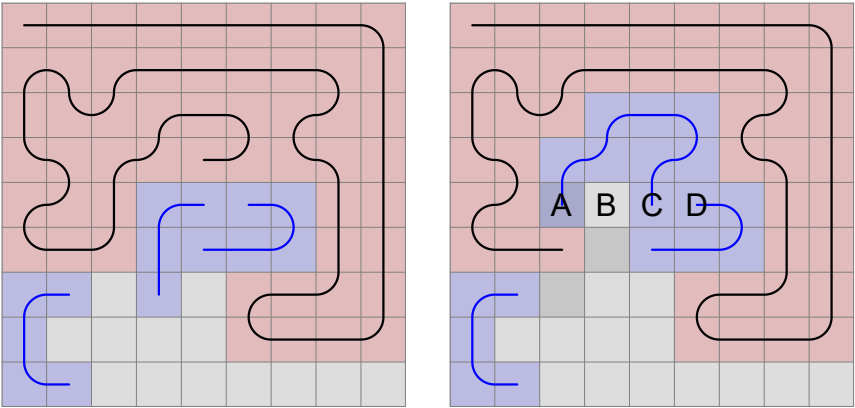
In order to make a walkthroughable area, you would need to extend it by pairs of fields next to each other, like D and E. One will be filled at a pair amount of steps, the other at an impair amount.



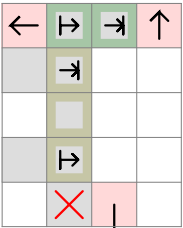
The rule will now look like this:



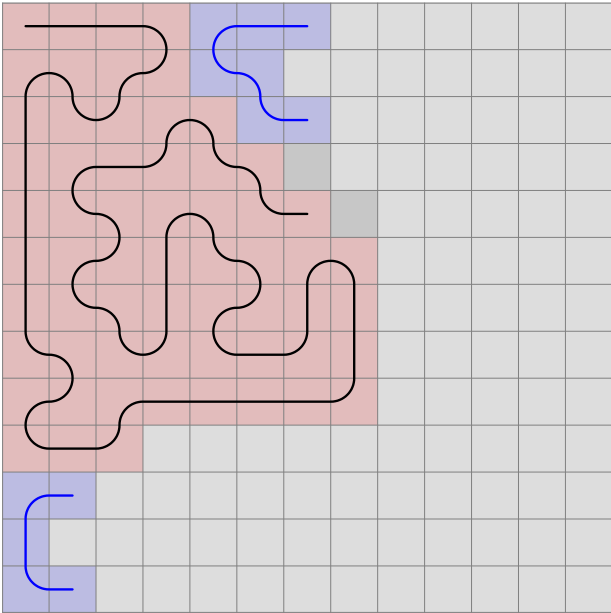
740 039 is a slight modification.



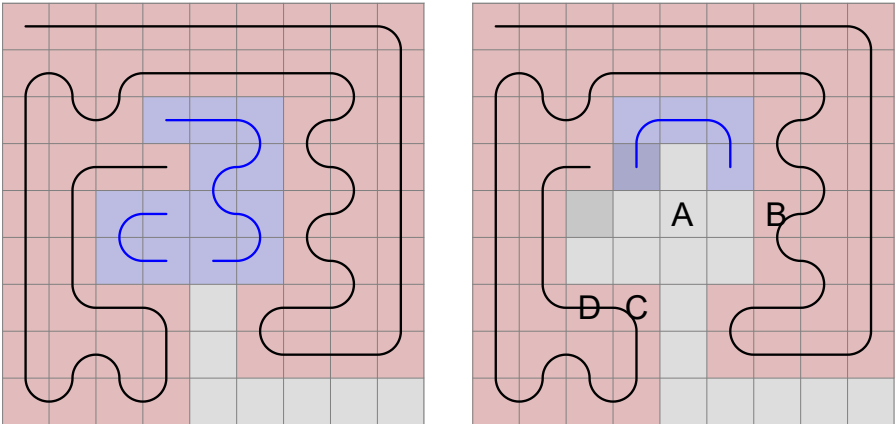
The only difference is, that the obstacle is 3-distance away. With the area being impair, if we enter at A, we must exit at C. What if we omit D from the area? Then the area will be pair, so we must exit at B, and the only way to get there is from C. And if D is included, we can only step to C from there. Either way, we step away from the area beyond D, so the rule will be:



And now the walkthrough is possible.

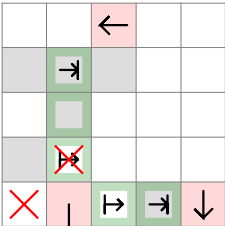


Continuing the 9 x 9 program, we get this at 641 019:

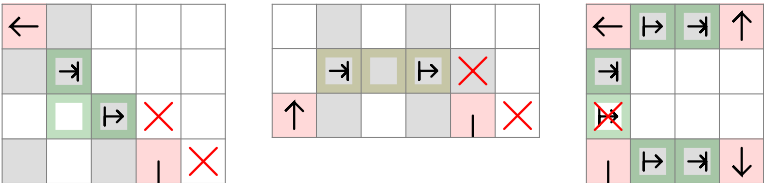


If you enter the pair area straight ahead, you will exit at A, and you need to turn towards B because of the C-shape. Now you cannot go in and out of the area enclosed by C, and the situation would be the same if that obstacle was in D.

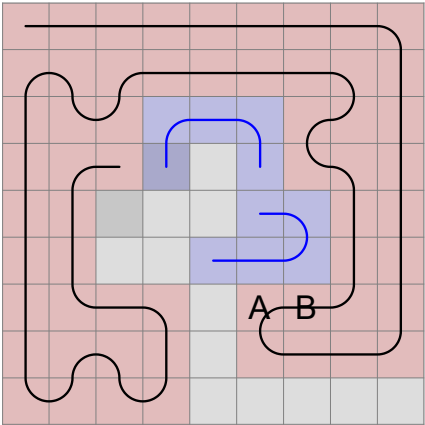
To mark the two areas, each one has to be given a directional obstacle next to the count area end field. In this case, it represents a taken field, but we don't go wrong if we include the border as well.



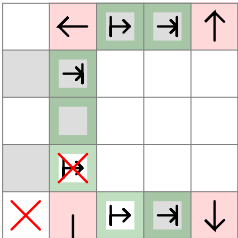
And with this marking system, we can correct the rules previously made. All rules featuring future line start and end fields have to be rewritten to start with.  
 So we get the 2-distance across rule, the straight 3-distance rule to prevent a double C-shape, and the square constellation with 3 areas. All of them are rotated clockwise or counter-clockwise.



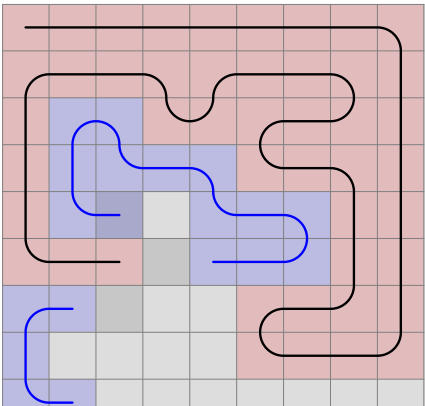
Let's return to the last example and make a modification:



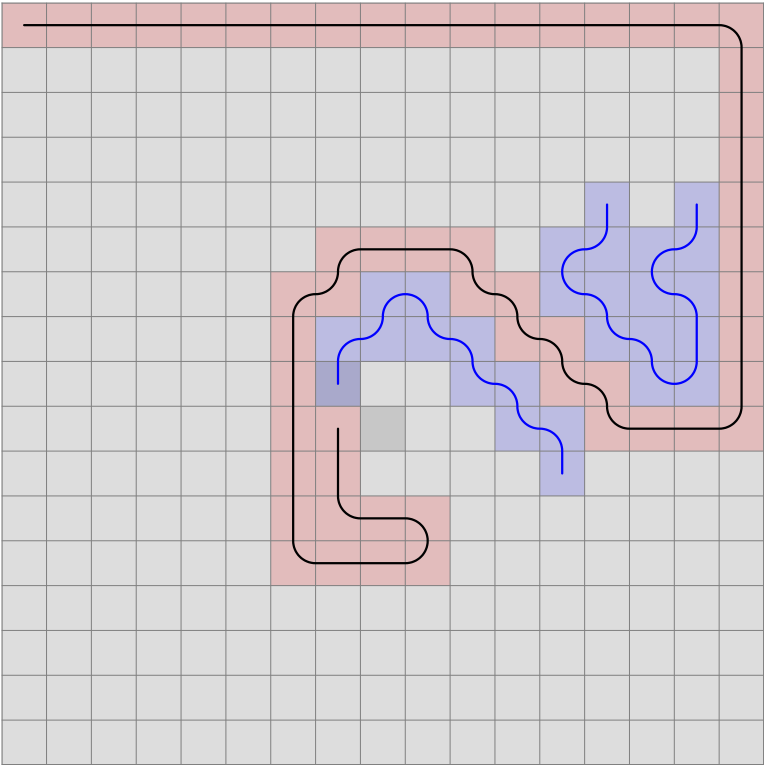
The field previously marked with B is now empty. But we still need to step in that direction, due to the area enclosed by A, which obstacle could as well be in B.  
 The rule will be now symmetrical. It is similar to the square obstacle pattern.



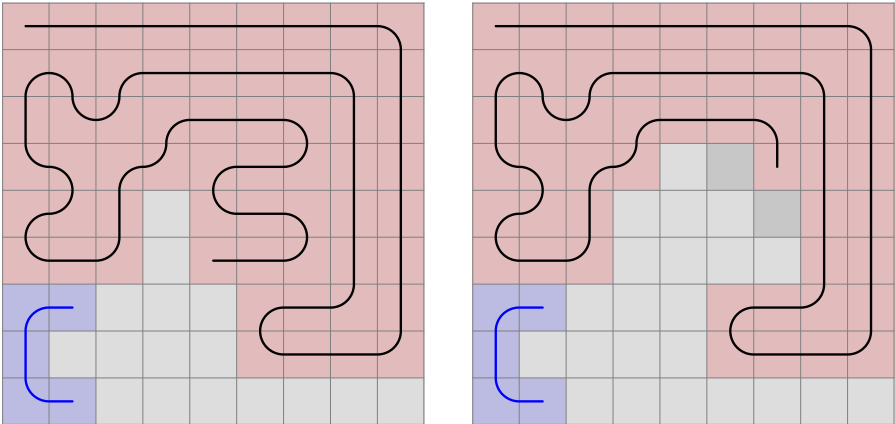
The same concept we encounter at 725 325. We have seen this previously, just with C-shape, not an area.



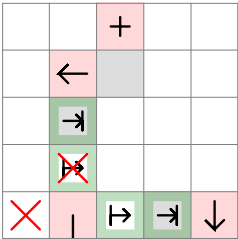
On 17 x 17, we can construct a situation where the obstacle across the stair is 2 behind and 2 to right. As the table size increases, the stair-obstacle narrowing can move infinite distance away from the live end. That's why it is important to group these rules as one.



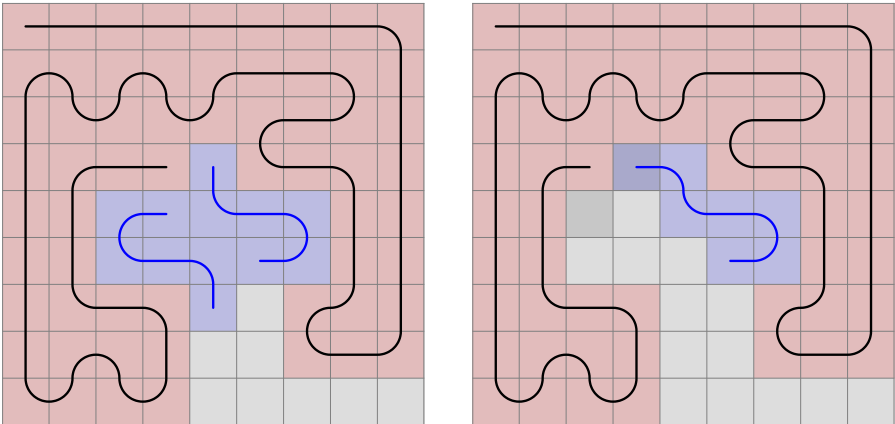
811 808:



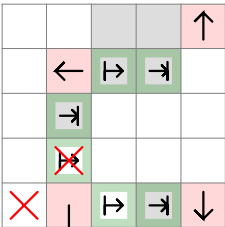
Recognize it is a variation of the square obstacle pattern where instead of an area, there is a C-shape at the rule's upper edge.



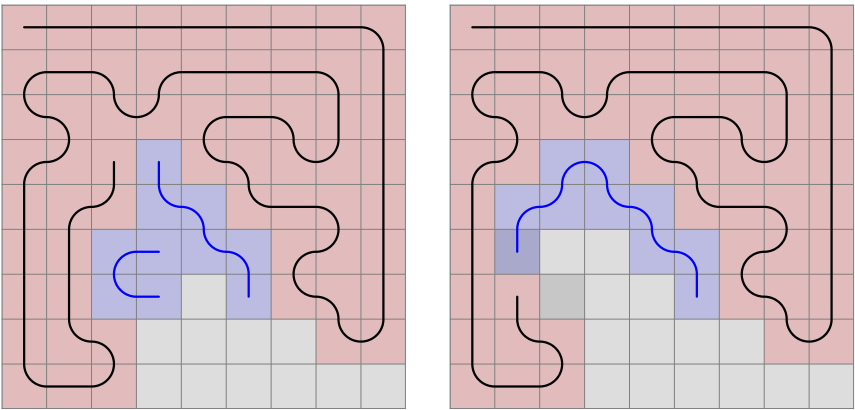
1 261 580:



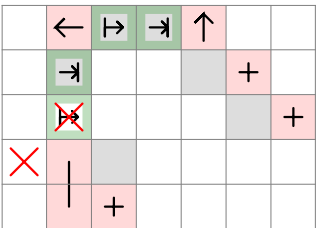
Again, same pattern with area. The upper obstacle is now moved, but it will satisfy the previous examples too. The rule replaces the old one.



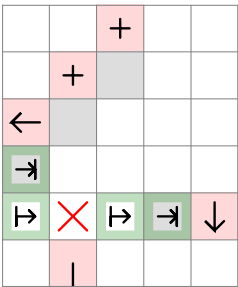
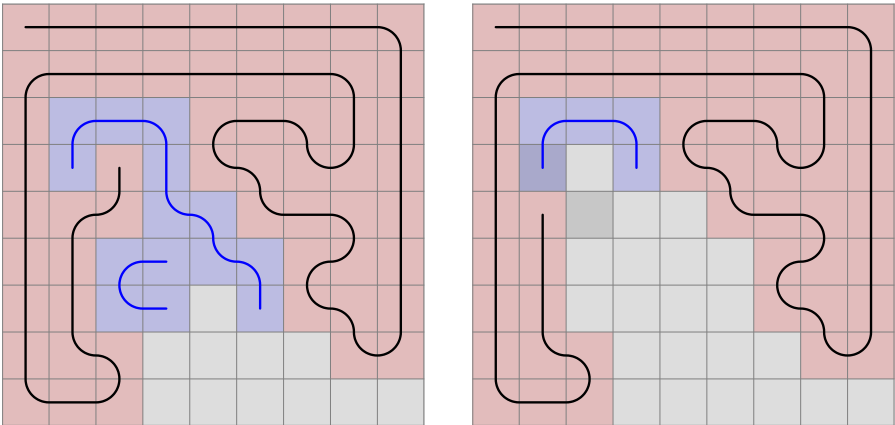
2 022 337 is getting stuck because of the stair-shaped walls that force the future line along them. Therefore, an area is created with only one field to go in and out of it. What is the solution?



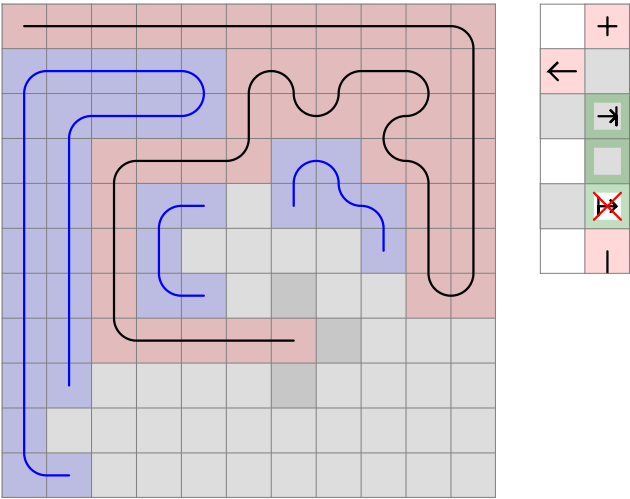
Though not as universal as we want it to be, this will solve this specific situation:



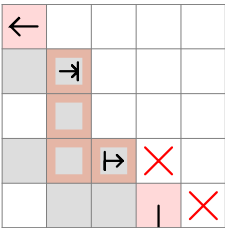
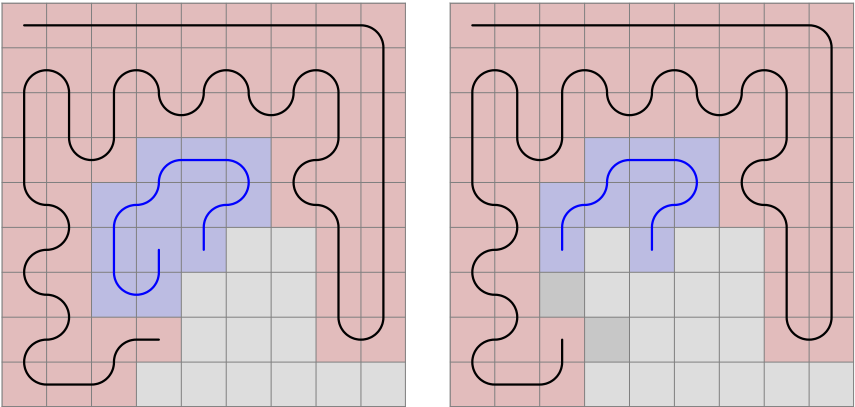
And soon, at 2 022 773 we encounter a similar one:



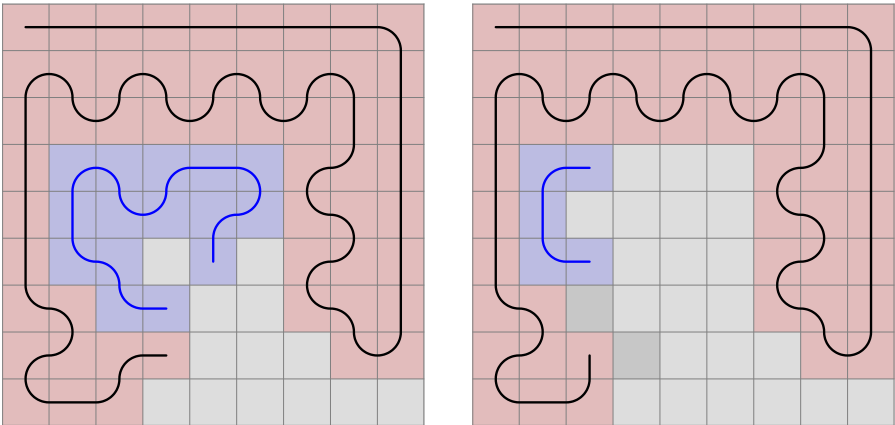
From our experience, the area can be substituted with C-shape.



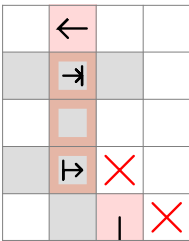
3 225 432 is a variation of the impair area imbalance rules we have seen before.



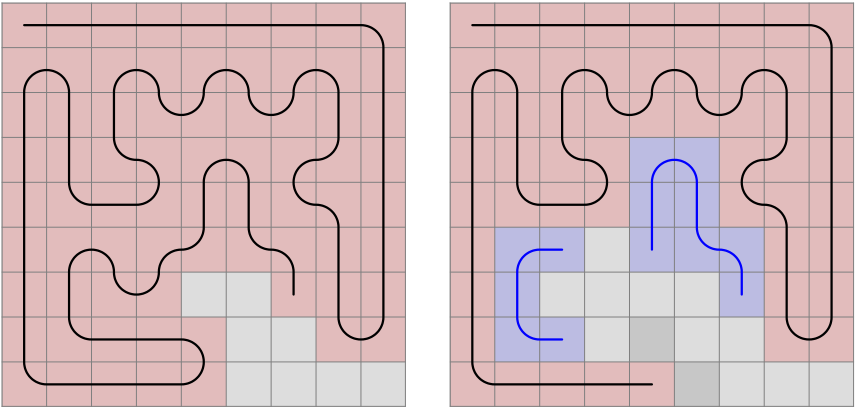
We have all the tools to handle 2 034 575.



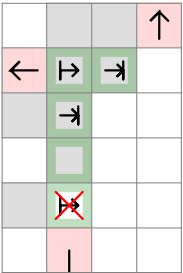
It is an impair area where the number of the starting field's color is less than the other color.



Next stop is at 3 224 847.

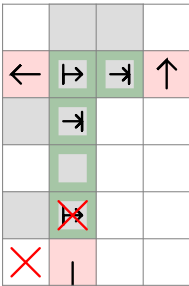


A pair area is created with the obstacle 3 distance away, so if we step into it, we will exit at the middle, but because of an area, we cannot step there.

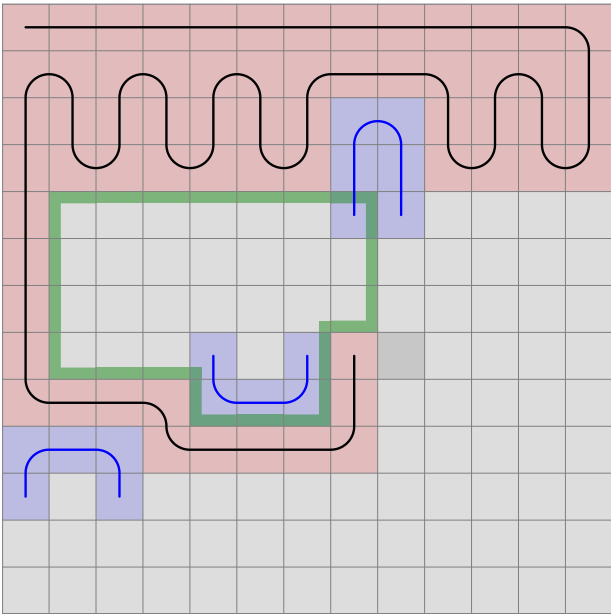


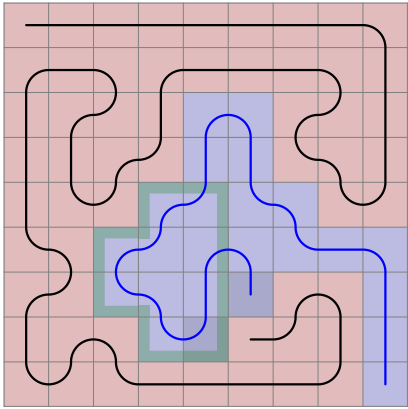
Beware of disabling the left field. If the count area end field is excluded from the area, the area will be impair, thus we will exit at the count area start field, coming from the middle.

But if the obstacle in the upper right corner is moved down, even this will be impossible.



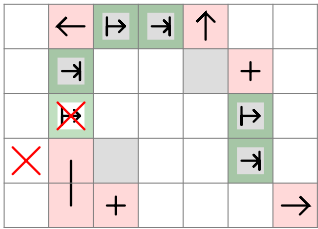
We can recreate this example on 13 x 13.



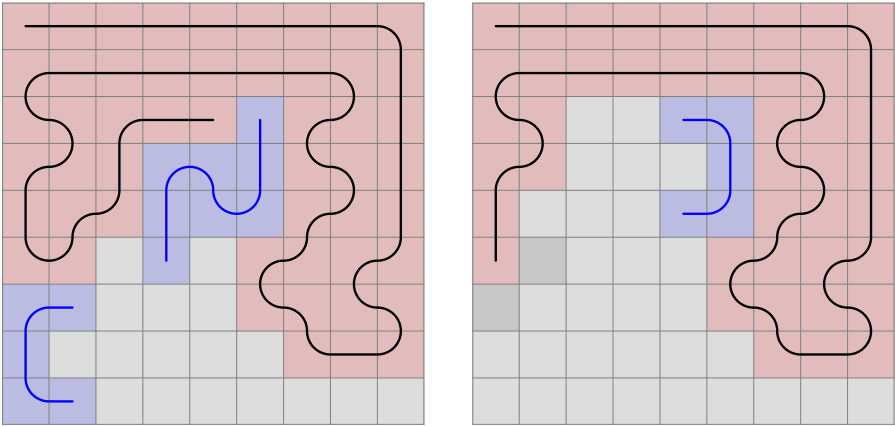


It is not the only thing. So far, when we entered a future line, the program just followed it without checking the possibilities for the next step. This behaviour needs to be changed too. Future lines are no longer needed, and we should restart the 9 x 9 walkthroughs.

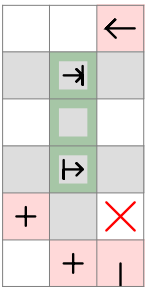
For now, here is the solution to this and the next cases:



8 076 012 builds upon the existing rule where C-shapes are created on both sides if we enter an impair area.



Here, a C-shape at the start would force the line to enter the area.



In the rule, I introduced a new field that indicates the entry point; this has always been the start field until now.

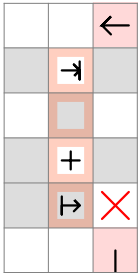


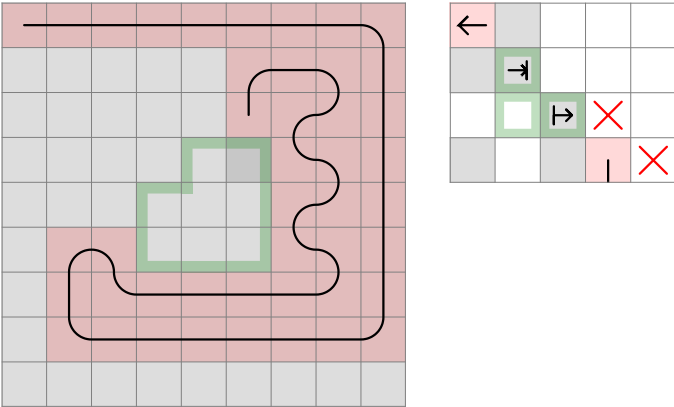
Figure 1 consists of two side-by-side diagrams illustrating a 2D grid environment for a pathfinding problem. Both diagrams feature a 10x10 grid with a red background. A black line outlines a complex shape, possibly representing a building or a boundary. In the left diagram, a blue path starts at the bottom left, moves right, then up, then right again, and finally up to the top right. A green obstacle is located in the center of the grid. In the right diagram, the blue path starts at the bottom left, moves right, then up, then right again, and finally up to the top right. The green obstacle is located in the center of the grid, but its shape is different from the one in the left diagram.

The diagram shows a 10x10 grid. A complex black boundary encloses most of the area, leaving a small rectangular opening at the bottom center. The interior is divided into red and blue regions. A blue path starts from the right edge, enters through the bottom opening, moves left, then up and left, ending in a large central blue region.

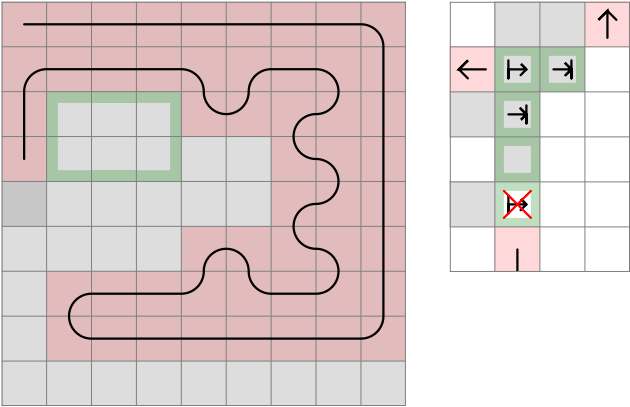
46 47



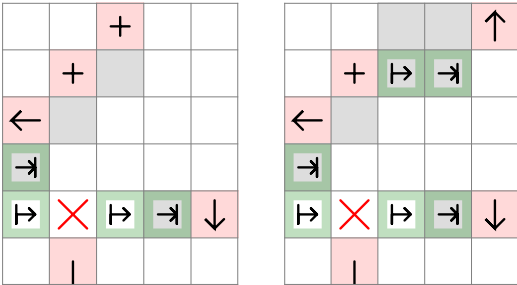
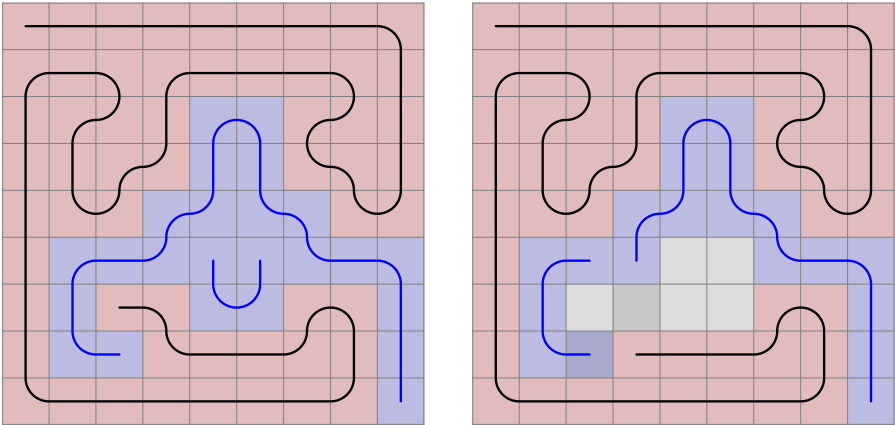
9 121, Count Area 2 Across



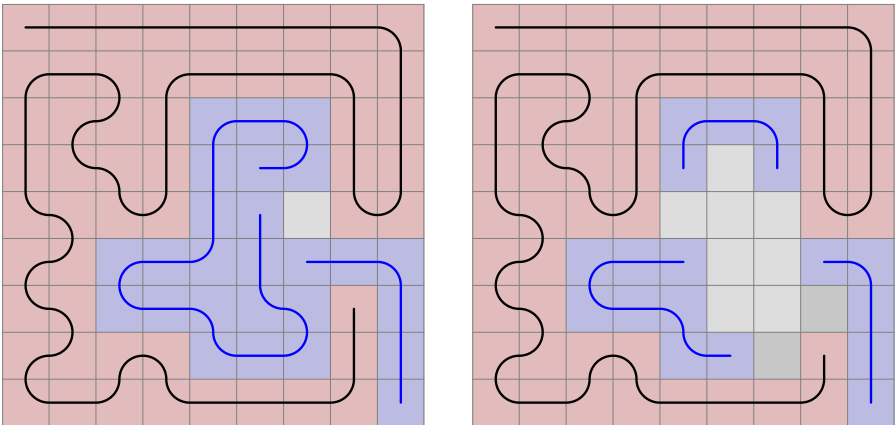
22 328, Straight Mid Across 3 End Area



19 718 148 is a slight modification of 2 022 773 where there is an area instead of a C-shape straight ahead.



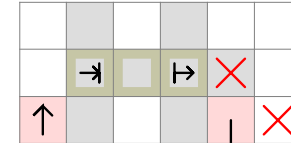
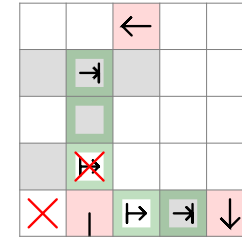
We encounter a new constellation of 3 areas in 23 310 321 where the exit is next to the live end.

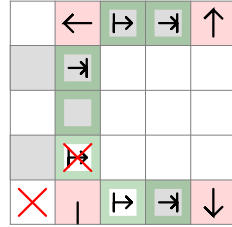


←	→		→	↑	
→					
→				×	
←	→		→	↓	×

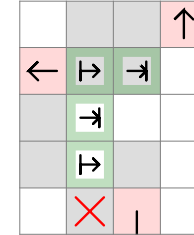
In the following section I list the 9 x 9 rules in chronological order. The patterns are not introduced when they are first recognized, but when they are first needed, meaning that they disable fields that the other rules don't. And the disabled fields have to be empty.

Still, the number of completed walkthroughs before the appearance of the rule may not be the same as the number of those before getting stuck in the lack of that rule. If the rule disables a field right to a possible field, the left branch would run through first.

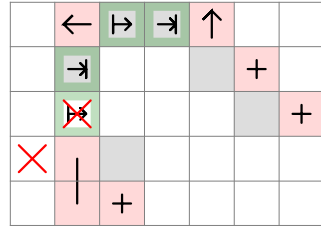


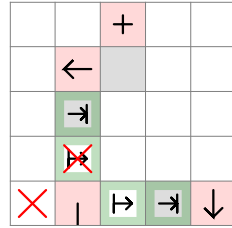
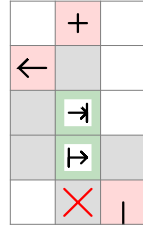


A 10x10 grid with a red region and a black boundary. The red region is defined by the following cells (row, column): (1,1), (1,2), (1,3), (1,4), (1,5), (1,6), (1,7), (1,8), (1,9), (1,10), (2,1), (2,2), (2,3), (2,4), (2,5), (2,6), (2,7), (2,8), (2,9), (2,10), (3,1), (3,2), (3,3), (3,4), (3,5), (3,6), (3,7), (3,8), (3,9), (3,10), (4,1), (4,2), (4,3), (4,4), (4,5), (4,6), (4,7), (4,8), (4,9), (4,10), (5,1), (5,2), (5,3), (5,4), (5,5), (5,6), (5,7), (5,8), (5,9), (5,10), (6,1), (6,2), (6,3), (6,4), (6,5), (6,6), (6,7), (6,8), (6,9), (6,10), (7,1), (7,2), (7,3), (7,4), (7,5), (7,6), (7,7), (7,8), (7,9), (7,10), (8,1), (8,2), (8,3), (8,4), (8,5), (8,6), (8,7), (8,8), (8,9), (8,10), (9,1), (9,2), (9,3), (9,4), (9,5), (9,6), (9,7), (9,8), (9,9), (9,10), (10,1), (10,2), (10,3), (10,4), (10,5), (10,6), (10,7), (10,8), (10,9), (10,10). The black boundary is a closed curve that starts at (1,1), goes right to (1,10), down to (10,10), left to (10,1), and up to (1,1). There are several indentations and protrusions: a small indentation at (2,1), a large protrusion at (3,1) extending to (3,10), a small indentation at (4,1), a large protrusion at (5,1) extending to (5,10), a small indentation at (6,1), a large protrusion at (7,1) extending to (7,10), a small indentation at (8,1), a large protrusion at (9,1) extending to (9,10), a small indentation at (10,1), and a large protrusion at (11,1) extending to (11,10).

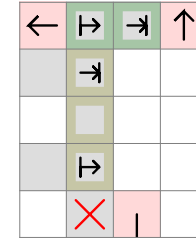
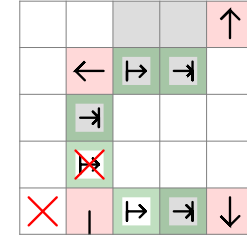


A 10x10 grid with a black boundary. The grid contains a complex shape defined by black lines. The interior of the shape is divided into white and light gray cells. A green rectangle highlights a 2x1 area of white cells in the lower-left quadrant.

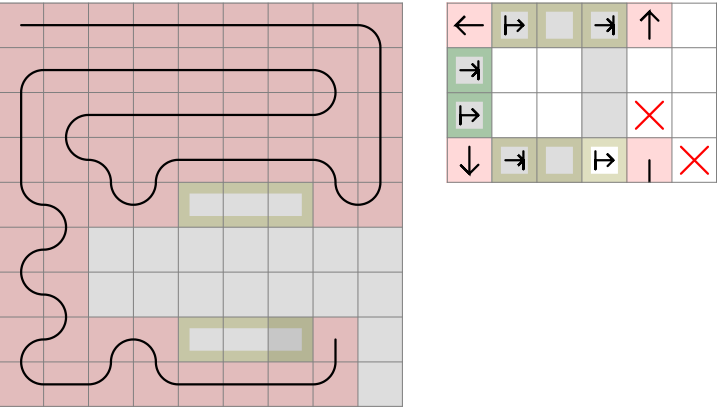




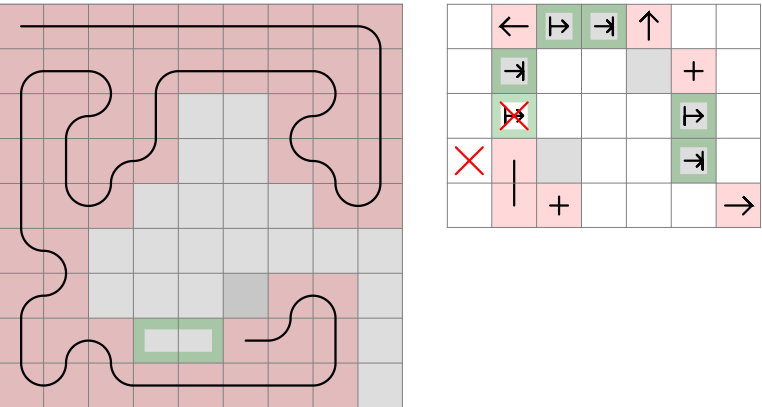
The figure shows a 10x10 grid with red and light blue cells. A black line forms a complex shape. Two green rectangles highlight specific areas: one on a red cell at (3,4) and one on a light blue cell at (4,3).



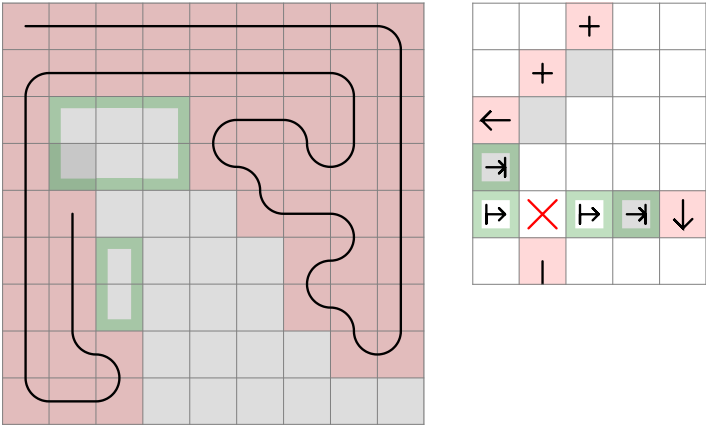
18 665 383, Triple Area Exit Down



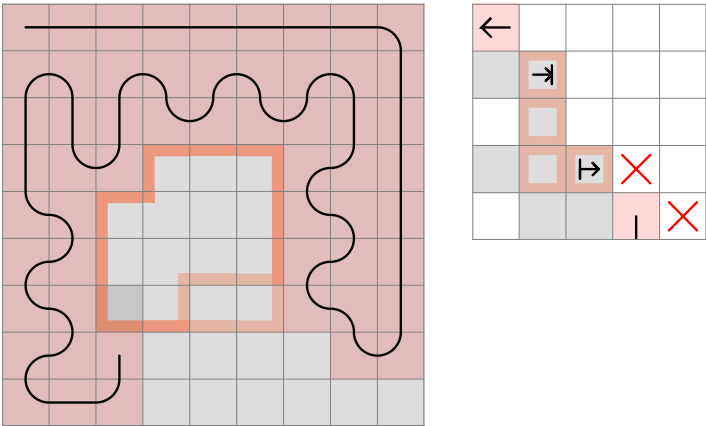
19 720 122, Triple Area Stair



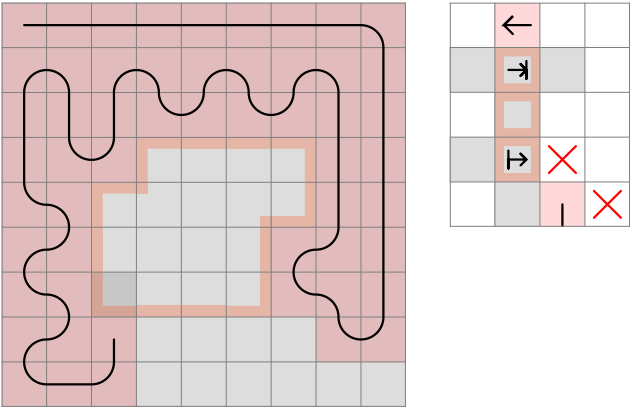
2 023 198, Double Area Stair 2



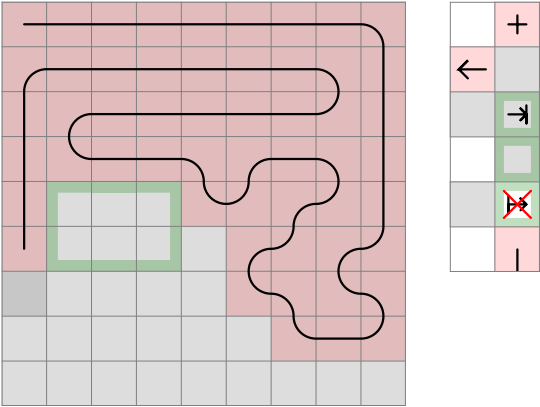
2 034 435, Mid Mid Across 3 Determined (and Mid Across 3 Impair Determined)



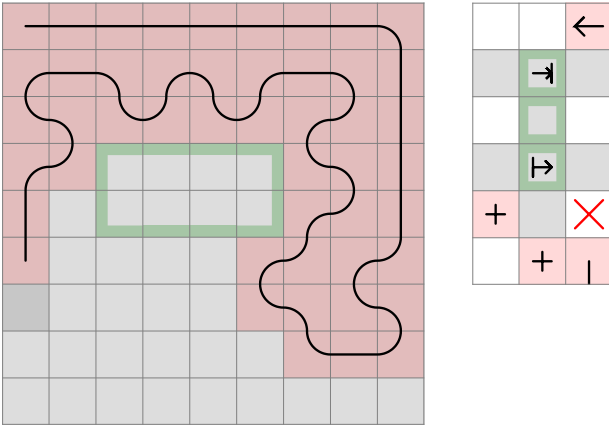
2 059 934, Mid Across 3 Impair Determined



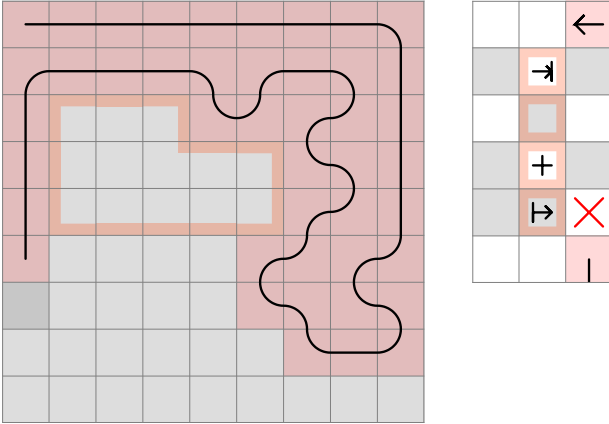
8 076 202, Straight Mid Across 3 End C



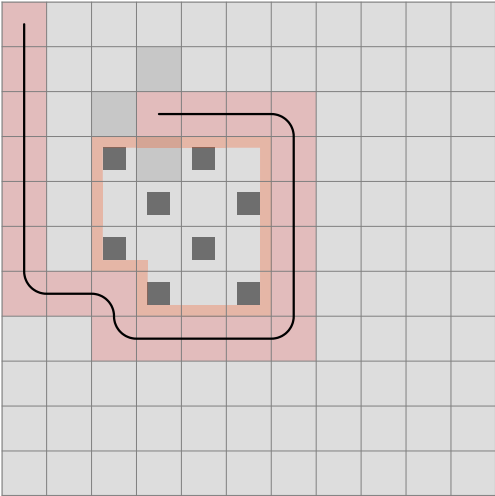
8 076 706, Double C-Shape Start C



8 076 762, Double C-Shape Determined

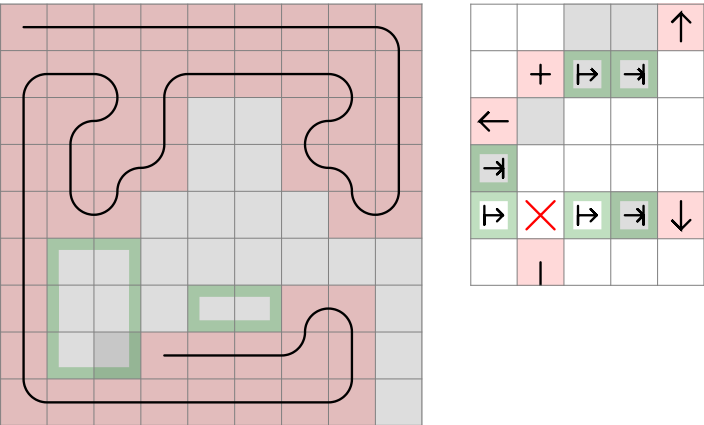


Here, the area is still pair, but there are 12 black fields and 10 white. To fill it, two lines of an impair length would be needed, each starting and finishing on a black field. Now, it is possible to enter and exit the black field in the upper left corner of the area, but we cannot do it with the field 2 below. And there are no more black fields on the boundary of the area. Filling it is therefore impossible.

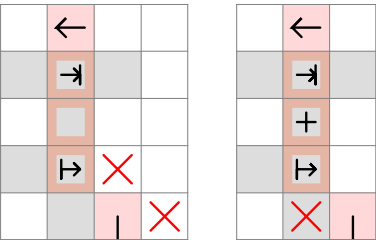
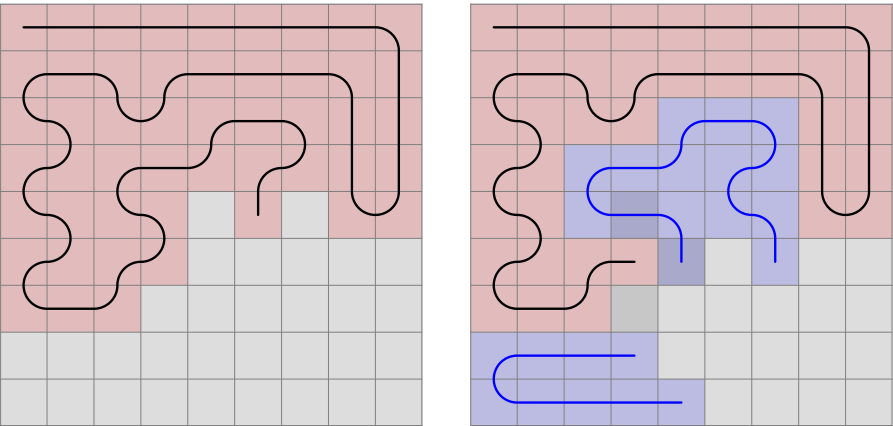


When there are one more black field than white (8 and 7 in this case), the area is impair. If we enter now, we can exit at the count area end (to make a line of pair length), and the corner black can be filled later. We can also enter at the corner black later to exit at the count area end.

19 720 614, Double Area Stair Area



23 350 320 is new, but it shows similarity to the Mid Across 3 Impair Determined rule. As the double C-shape reveals, it is about pair/impair field imbalance.



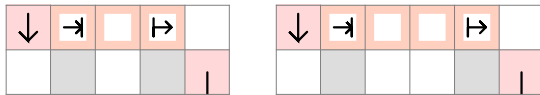
If we step left in the impair area, we can only exit at the count area middle field, but there is one less field of that type than the other.

And no fields can be omitted from the area for entry and exit later.

When the count area start field + middle field is omitted (subtracting a pair amount of cells from the area), the possible exit is the count area end field, which has a different parity than the field to the left.

When the count area middle + end field is omitted, the possible exit is the count area start field, which has again different parity.

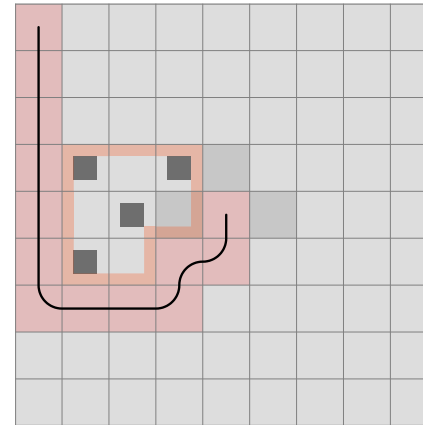
By now, we are able to group some rules and even solve the original 21 x 21 example. Previously, we have covered all of the cases where an obstacle is 2 distance away from the live end. Let's examine distances of 3, 4 and so on in this constellation:



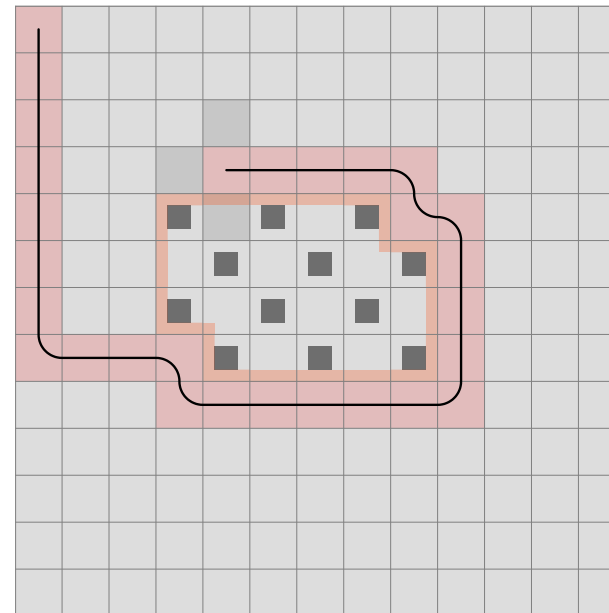
Besides the count area fields being empty, here I also assume that the left field and the field below the obstacle is also empty. In the future, this might have to be changed and new constellations added.

Knowing the difference between the number of pair and impair fields, we can make a decision in some cases.

Let's start with this:

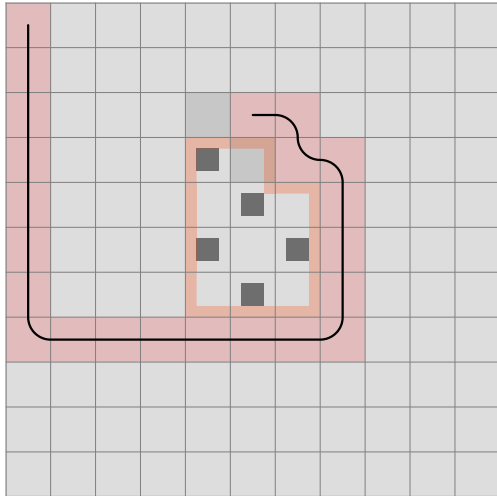


The distance is 3, the area is pair, and the number of pair and impair fields are the same. We can either enter the area now (stepping left) or later (up or right). Either way, we can start and end the area on a different color, so nothing should be disabled.

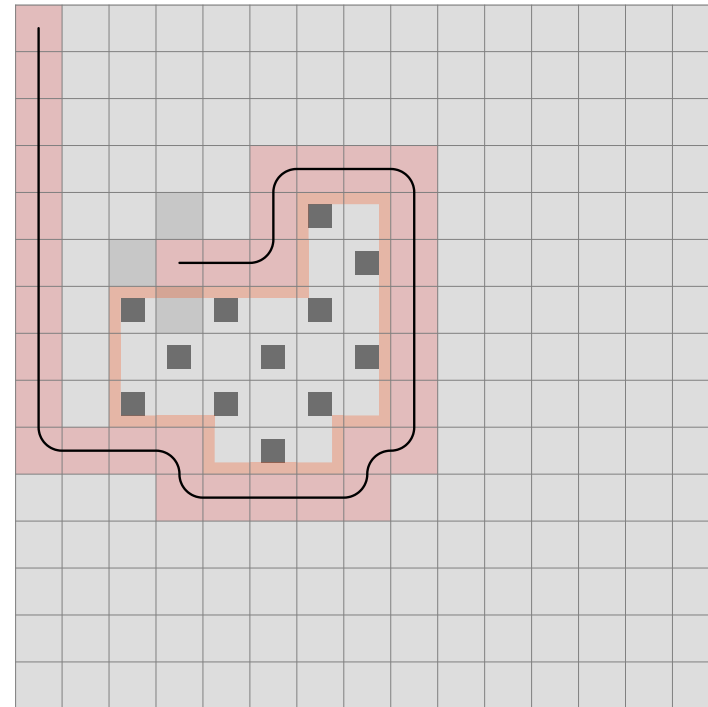




If  $\text{black} = \text{white} - 1$ :



A line that starts and ends on white can be drawn, no matter if we enter now or later. If we enter now, the next field has to be the corner black, and then there will be a line between the two white fields on the boundary.



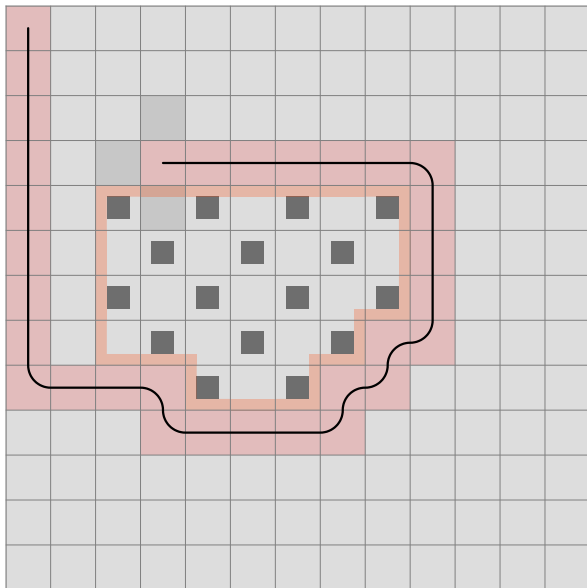
Here, there are one more white field than black. To fill the area, 1 line of impair length would be needed that starts and ends on white, but if the enter the area now, the corner black could not have been filled. Now there are 2 less black fields available for the rest of the area, but making two lines starting and ending on white is impossible, there are not that many white fields on the boundary.

Needless to say that we cannot enter later either, there is only one white field on the boundary.

If there are even fewer black fields relative to the white, the situation is the same.

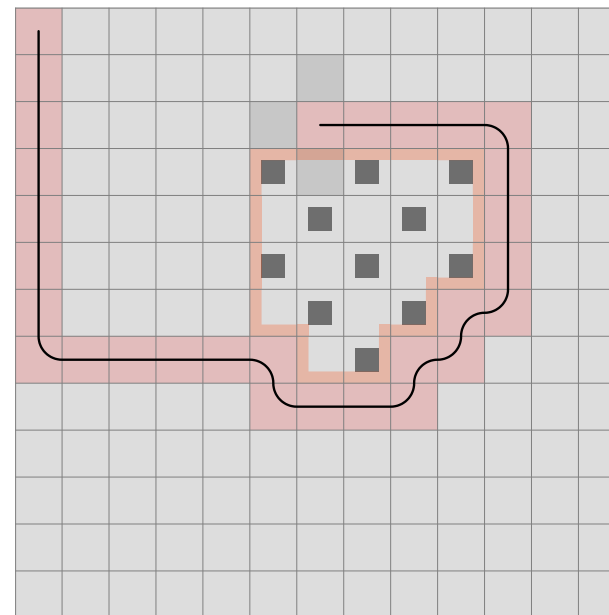
At 4 distance (and in any case), we still don't have a problem filling an area that has equal number of black and white fields.

When the number of black fields are two more than the whites (16 and 14):



Two black to black lines would be needed to fill the area. Apart from the black in the corner, there is only one black field on the boundary.

If black = white + 1:



If we enter now by stepping left, there will have to be one more line even if we exit on black. That line has to go from black to black, so it can only be the corner field. Have we exited at the other black field (the third on the boundary), either the second or the fourth could not have been filled.

**This direction therefore has to be disabled.**

We can enter later without problems to start at the corner field, then the second, go inside the area and end at the third.

**2) Pair distance, impair black and white: 6, 10 etc.**

If we enter now and exit at black,  $(B-1) / 2$  black to black lines can be drawn.

If we exit at white,  $(W+1) / 2$  white to white lines can be drawn. (We have to move to the corner black and exit at the first white during the first line)

If we enter later,  $(B+1) / 2$  and  $(W-1) / 2$  black and white lines are possible, respectively.

The rule is double:

- If the number of black fields is greater or equal than the number of whites plus  $(B+1) / 2$ , we cannot enter now.
- If the number of white fields is greater or equal than the number of blacks plus  $(W+1) / 2$ , we have to enter now.

**3) Impair distance, impair black and pair white: 5, 9 etc.**

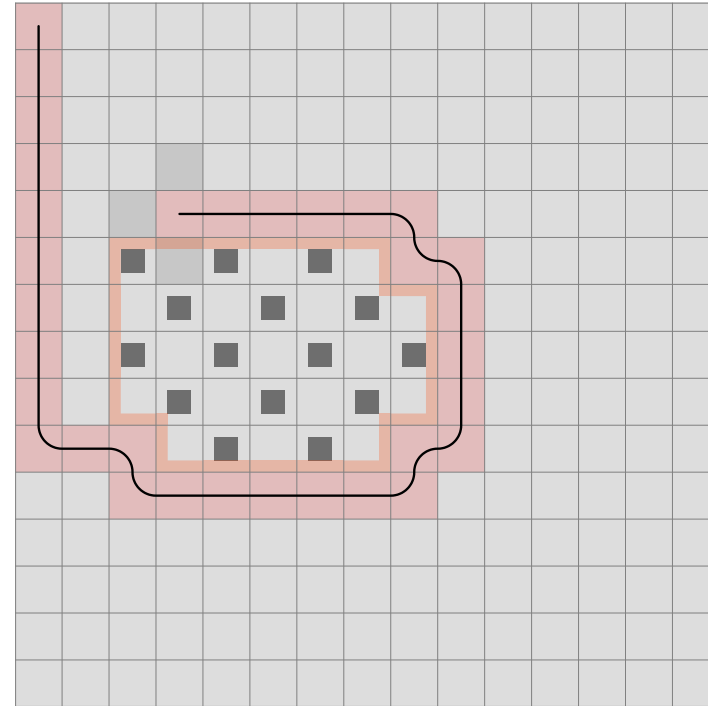
If we enter now and exit at black,  $(B-1) / 2$  black lines are possible.

If we exit at white,  $W/2$  white lines are possible.

If we enter later, the number of black lines can be up to  $(B+1)/2$ , the whites  $W/2$

Single rule: When the difference is at least  $(B+1)/2$ , we cannot enter now.

If black = white - 2:



Two white to white lines would be needed, but there are only 3 white fields on the boundary, and none is the corner.

Without finding concrete examples, at 5 distance I only draw the boundary and go through the different possibilities.



black = white + 2

If we enter now and finish at black, only two black fields remain. Drawing two black to black lines is not possible, so this direction has to be disabled. We can enter later, go through the corner and draw another black to black line.

$$\text{black} = \text{white} + 3$$

There is not enough black fields for three black to black lines.

$$\text{black} = \text{white} + 1$$

We can enter now, finish on a black field and draw another black to black line. Or enter later.

$$\text{black} = \text{white} - 1$$

A white to white line is possible in either case.

$$\text{black} = \text{white} - 2$$

We cannot enter now, because even if we end on a white field, only one white remains, and that is not the corner. And cannot enter later either.

The same procedure applies at 6 distance.



### **black = white + 2**

The number of black fields is the same as previously, so entering now is not possible.

black = white + 3 is impossible.

black = white + 1 and black = white - 1 is possible.

### **black = white - 2**

Entering now is possible if we step upwards and exit at the neighbouring white field. Two white fields remains.

" But we cannot enter later and make two white to white lines using three white fields.

" From now on, we can distinguish between four cases:

**1) Pair distance, pair black and white** (indicated by B and W): 4, 8, 12 etc. distance

If we enter now and exit on black, B-1 black field remains on the border. B-1 is impair, and the corner alone can make a black to black line. Hence  $(B-2) / 2 + 1 = B / 2$  black to black line would be possible, but...

- if the first line finishes at the corner black, the opportunity for it to make a single line is lost, and the remaining B-1 black fields cannot make  $B / 2$  black lines.

- if it finishes on any other black field, on one side there will be a section that has white fields on both ends (below the greenish fields were taken by the first line)



" A black to black line cannot have three white fields on the border (unless more black fields were used up).

So there can be at most  $B / 2 - 1$  black to black lines.

If we enter now and exit on white, W-1 white fields remain, so  $1 + (W-2) / 2 = W / 2$  white to white lines is possible.

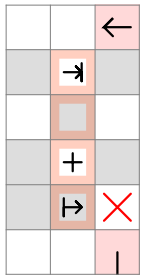
If we enter later, the number of possible black to black lines is at most  $B / 2$ , and the white ones  $W / 2$ .

Our rule will look like this: If the number of black fields in the area is greater or equal than the number of white fields plus  $B / 2$ , we cannot enter now.

←			
	■		
	■		
	■		
	■		
	■	■	
	■	■	↓

The figure consists of four 5x5 grids, each representing a step in a pathfinding process. The start cell is marked with a black arrow, and the goal cell is marked with a red arrow. Obstacles are represented by gray cells. The path is highlighted in red. Red 'X' marks indicate cells that are not part of the current path.

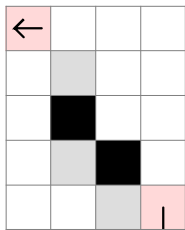
- Grid 1:** The path starts at (1,1) and goes down to (4,1). The goal is at (4,5). Obstacles are at (1,2), (2,2), (3,2), and (4,2). Red 'X' marks are at (4,3) and (4,4).
- Grid 2:** The path starts at (1,1) and goes down to (3,1). The goal is at (3,5). Obstacles are at (1,2), (2,2), (3,2), and (4,2). Red 'X' marks are at (3,3) and (3,4).
- Grid 3:** The path starts at (1,1) and goes down to (3,1). The goal is at (3,5). Obstacles are at (1,2), (2,2), (3,2), and (4,2). Red 'X' marks are at (3,3) and (3,4).
- Grid 4:** The path starts at (1,1) and goes down to (3,1). The goal is at (3,5). Obstacles are at (1,2), (2,2), (3,2), and (4,2). Red 'X' marks are at (3,3) and (3,4).



As a reminder, they indicate impair areas, and in the first two case if the count area start field (black) type is not 1 field more in the area, we cannot enter later.

In the fourth, the field marked with + (white) is the type that needs to be 1 more than black, otherwise stepping forward is forbidden.

I will take the second rule under examination as it contains both a horizontal and vertical offset.



If we enter now:

- We can exit at the end white, so 1 white line is possible.
- We can exit at the black farthest away and then make a line using the black field closest. 1 black line is possible.

I will mark it like this: 1W -> 1B

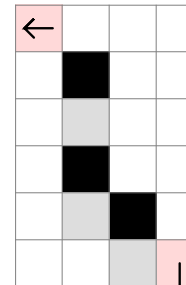
If we enter later:

- Having two black fields, 1 black line is possible.
- There is only one white field. It sits in a corner, but the two black fields will give 1 black line. 0 white line is possible.

0W -> 1B

- So if there are 1 more white fields in the area than black (1W), we cannot enter later.

- Now let's increase the vertical distance.



If we enter now, 1 white line to 2 black lines are possible. There are 2 black fields on a corner.

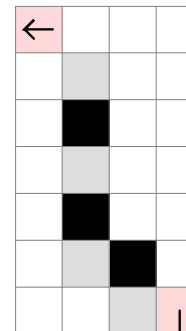
1W -> 2B

Later, 0 white line and 2 black lines can be drawn.

- 0W -> 2B

Conclusion: in case of 1W, we cannot enter later.

5 distance:

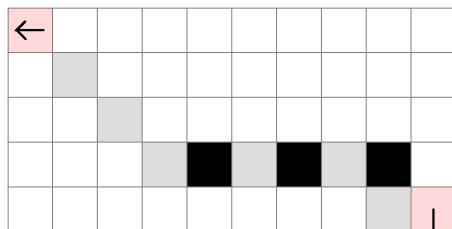


Now: 2W -> 1B

Later: 1W -> 2B

- 2W: cannot enter later
- 2B: cannot enter now

But we are not finished, we still need to examine the distance of 8.



Now: 4W -> 1B

After entry, we need to move up to fill the corner black and exit at the first white field.

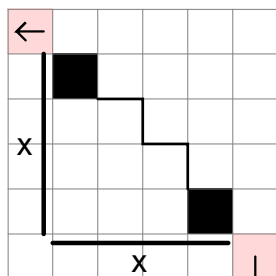
At 4 distance, only 2W was possible, but now we can exit at the first white and fill the area when entering at the second and exiting at the third.

Later: 3W -> 1B

After this practice, let's calculate how many white and black lines we can draw when we have an obstacle x and y distance away.

There are three cases to look at.

1. Equal horizontal and vertical distance



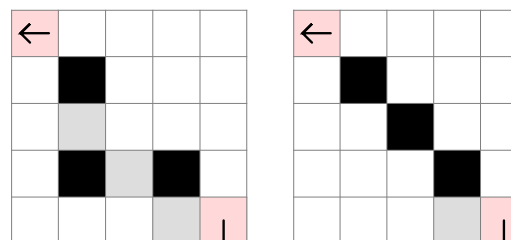
There are an x number of black fields on the area boundary.

In case of 7, there are 4 fields added to the 3-distance example. We would expect that in case of 2W, we cannot enter later, but now, 2W is possible even when entering later, because one line will be the corner white, and the other can go between the other two white fields, taking up all the blacks along the way.

From now on, increasing the numbers by 1 for every 4 distance increase will work.

The next thing to do is the horizontal increase.

3 distance:



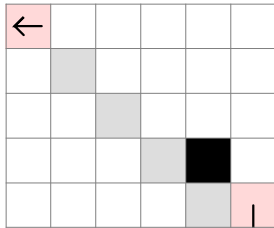
The picture on the left is the representation we have used so far. However, we cannot exit at the black in the middle, and when we exit at one of the whites, the other is only accessible for immediate entry. Therefore, I will add the extra field. If there is a taken field anywhere ahead acting as the obstacle, we can get to it by drawing a straight line and a stair.

Now: 0W -> 2B

Later: 1B -> 2B

If we used all 3 black corners for a separate line, we would not enter the area.

4 distance:

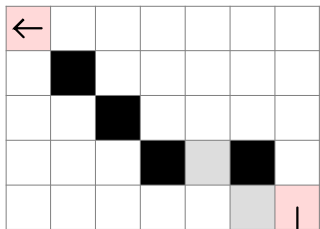


Now: 2W -> 0B

If we exited at the nearest white after entering, we have either not entered the area or not filled the black field.

Later: 2W -> 0B

5 distance:

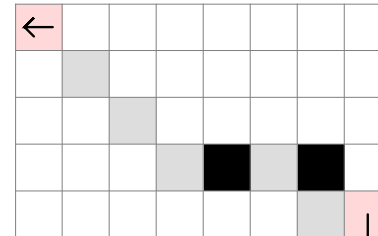


Now: 0W -> 3B

If we reserved the 3 black corners, the only way to end the first line on black is to move downwards after entry.

Later: 0W -> 3B

6 distance:

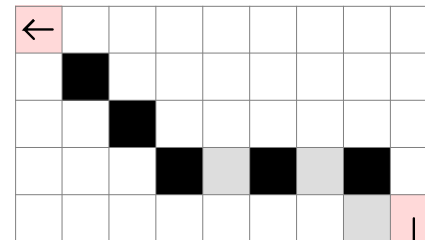


Now: 3W -> 1B

Similarly, we need to move down after entry in order to finish at the second black field, leaving the first for itself.

Later: 3W -> 1B

7 distance:



Now: 1W -> 3B

Later: 0W -> 4B

There cannot be one white line, because out of the first three black fields only 2 would be filled.

Notice that as we added 4 distance to 3, now both the white and the black end of the ranges have increased by one.



Later: Because of three inline white fields, one white line is possible. For blacks, we can use all the corners and two out of the three inline blacks.  
 $1W \rightarrow (x+1)B$

We don't need more examples.

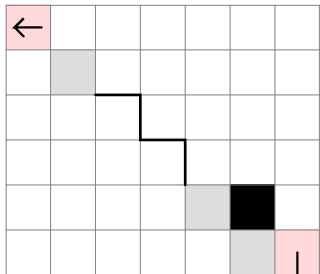
For  $2n$  added fields, there will be  $n$  amount of inline white fields.

If we enter now,  $(n+1 - (n+1) \% 2) / 2$  white lines can be drawn if  $n > 1$ .  
 The black line count is all the corners minus one for finishing the first line, plus one for each remaining pair.  
 $x + (n-1 - (n-1) \% 2) / 2$

For entering later, the white line count is  $(n - n \% 2) / 2$ , and the black line count is  $x + (n - n \% 2) / 2$ .

**If the added distance is impair:**

$2n + 1 = 1, n = 0$ :



Now: If we end the first line at the first white, we could not have filled the corner black and the area simultaneously.  $xW$  is therefore not possible, but  $(x-1)W$  is.

And since there is only one black field, the black line count will be 0.  
 $(x-1)W \rightarrow 0B$

$x = 1$ :  
 Now:  $0W \rightarrow 0B$   
 We cannot enter later.

$x = 2$ :  
 Now:  $0W \rightarrow 1B$   
 Later:  $1B$

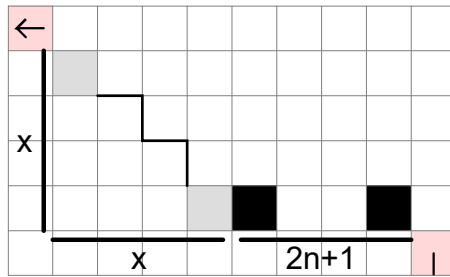
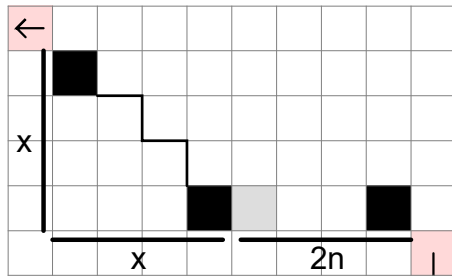
$x = 3$ :  
 Now:  $0W \rightarrow 2B$   
 Later:  $1B \rightarrow 2B$

$x = 4$ :  
 Now:  $0W \rightarrow 3B$   
 Later:  $1B \rightarrow 3B$

$x = n$ :  
 Now:  $0W \rightarrow (n-1)B$   
 Later:  $1B \rightarrow (n-1)B$   
 There is an  $x$  amount of corner blacks, but we need to enter the area as well.

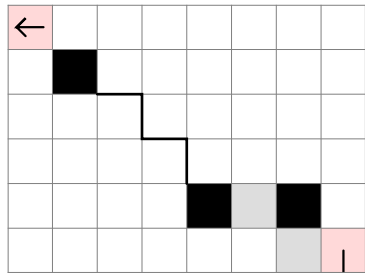
Conclusion: if the white and black fields in the area are equal, we cannot enter later.

## 2. Larger horizontal distance



**If the added distance is pair:**

$2n = 2$ :



Now: Can we make 1W? No, because that would require going through the first black field to end at the first white, while filling the other parts of the area too.

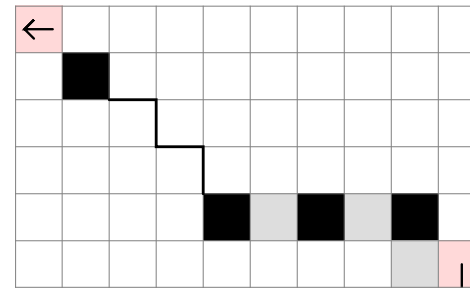
What about the black line count? Without the horizontal addition,  $(x-1)B$  was possible. Can we now make  $x$  amount? Yes, by stepping down and ending the line at the first white field and the second black, having filled everything except the corner blacks.

$0W \rightarrow xB$

Later: We can use all corner blacks, and the line starting at the first white and ending at the second black will fill the area.

$0W \rightarrow xB$

$2n = 4$ :



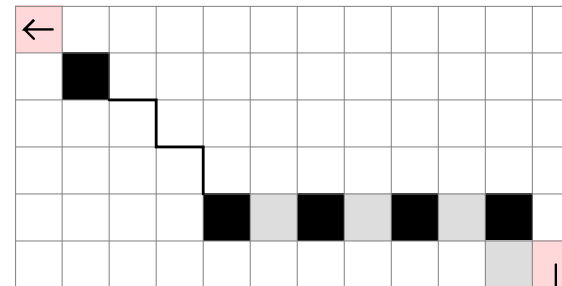
Now: 1W is now possible by ending at the second white, and everything can be filled on the way, while  $(x+1)B$  is not possible. Two inline blacks remain after taking off all the corner blacks, and we have to use one of them to exit the first line.

$1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

The line connecting the two inline blacks can fill the rest of the area.

$2n = 6$ :

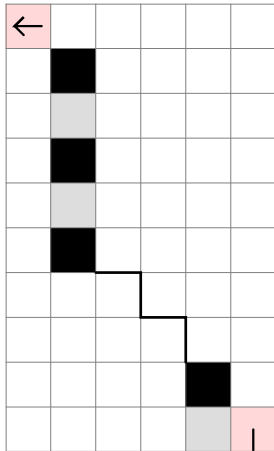


Now: If the first line goes through the first black and exits at the first white, the second line can go between the remaining two whites.

As far as the blacks concerned, if we exit at the second black, we can use the third and fourth for a line, plus the  $x$  amount of corners.

$2W \rightarrow (x+1)B$

$2n = 4$ :



Now: There cannot be more than  $1W$ . And the black line count is unchanged too. One of the two inline blacks will be used for completing the first line, and one remains plus  $x$  amount of corner.

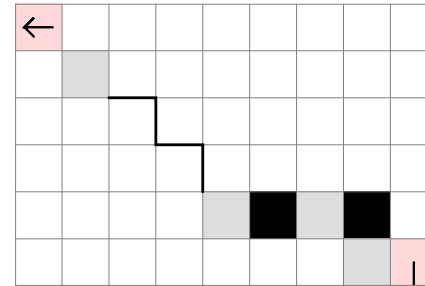
$1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

Later: All the corner whites plus the neutral line makes  $(x-1)W$ . The black line count is still 0. The black field is a corner, but it will be counterbalanced by at least one white to white line.

$(x-1)W \rightarrow 0B$

$2n + 1 = 3, n = 1$ :



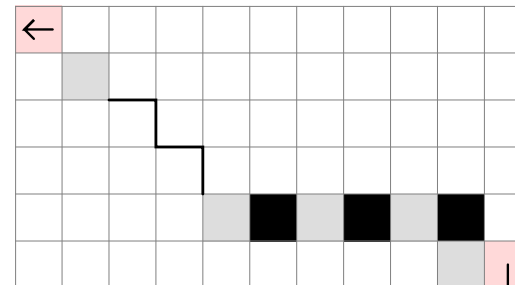
Now: Aside the corner whites,  $1W$  is possible by ending at the second white, just like in the  $2n = 4$  case.

The black line count is now 1, but we need to step downwards and finish at the first white and second black in order to have the corner black available.

$xW \rightarrow 1B$

Later:  $xW \rightarrow 1B$

$2n + 1 = 5, n = 2$ :



Now: We can step up and finish at the first white, because line connecting the remaining two whites can fill the area.

$(x+1)W \rightarrow 1B$

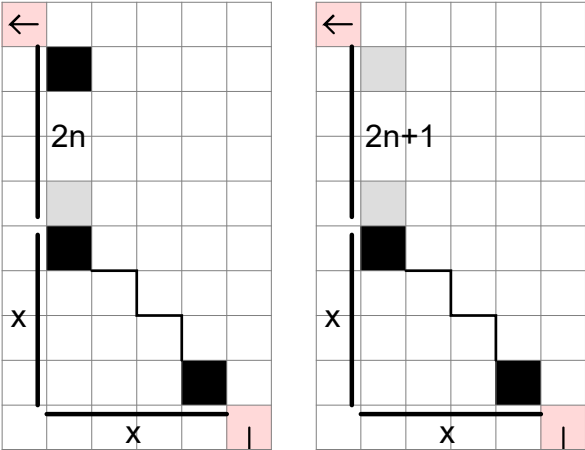
Later:  $xW \rightarrow 2B$

From now on, the calculations are the following:

If we enter now,  $x + (n - n \% 2) / 2$  white lines can be drawn if  $n > 0$ .  
The number of black fields is  $n+1$ , and when we use up one to finish the first line,  $n$  amount remains, one of which is a corner. Add one to make pairs, and the formula will be  $(n+1 - (n+1) \% 2) / 2$ .

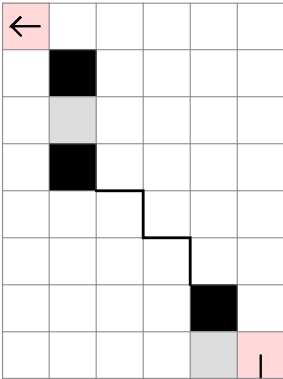
When entering later, the white line count is  $x-1 + (n+1 - (n+1) \% 2) / 2$ , and the black line count is  $(n+2 - (n+2) \% 2) / 2$  if  $n > 0$ .

3. Larger vertical distance



If the added distance is pair:

$2n = 2$ :

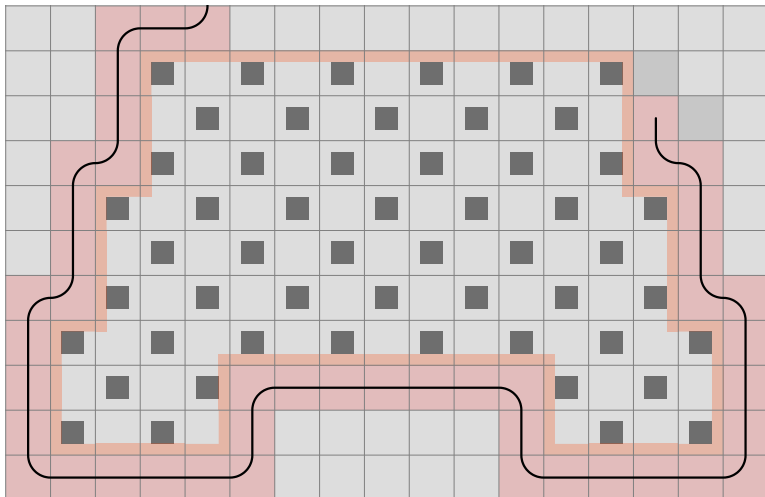
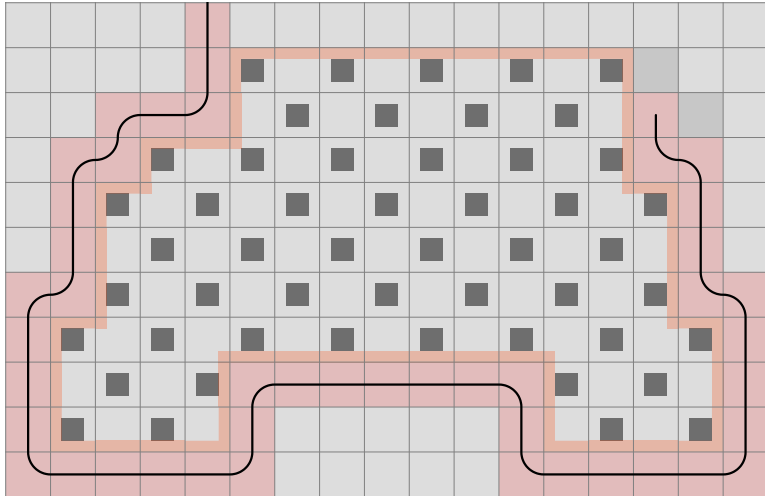


Now: It is possible to exit at the white field, having filled everything. The previous step must have been the last black.  
When counting the blacks, we exit the first line at the first black field above the stair, and  $x$  amount of corner blacks will remain.  
 $1W \rightarrow xB$

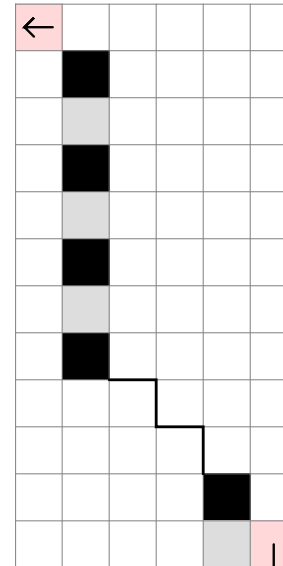
Later: Same as the  $2n = 2$  case previously.  
 $0W \rightarrow xB$

Now the area can be counted. And after this, we increase the vertical distance by one and stop / mirror sides when a field at one horizontal distance is taken or is the border.

Compare these two cases:



$2n = 6$ :



Now:  $2W \rightarrow (x+1)B$

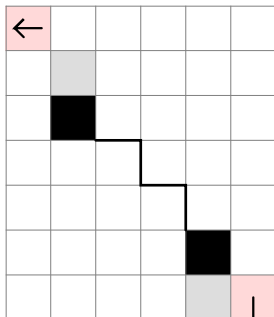
Later:  $1W \rightarrow (x+1)B$

When entering now, the general formula will be  $(n+1 - (n+1) \% 2) / 2$  white and  $x + (n-1 - (n-1) \% 2) / 2$  black.

The later case is the same as previously,  $(n - n \% 2) / 2$  for white and  $x + (n - n \% 2) / 2$  for black.

## If the added distance is impair:

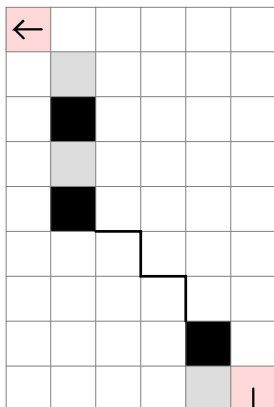
$$2n + 1 = 1, n = 0:$$



Now: 1W -> (x-1)B

Later: (see horizontal case) 0W -> (x-1)B

$$2n + 1 = 3, n = 1:$$



Now: 2W -> (x-1)B

Later: 1W -> xB

- Now: The corner white always gives 1. Then make pairs with the remaining inline whites plus the white field we are stepping first.

$$1 + (n+1 - (n+1) \% 2) / 2$$

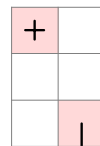
- For the blacks, one of the inline black fields will be taken by the first line. We can then make pairs with the remaining inline blacks and add the corners.

$$x - 1 + (n - n \% 2) / 2.$$

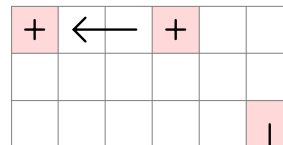
Later: The white line count is  $1 + (n - n \% 2) / 2$  if  $n > 0$ , and the black line count is  $x - 1 + (n+1 - (n+1) \% 2) / 2$ .

Next, we will look at the corner discovery algorithm.

Starting with 1 horizontal and 2 vertical distance, we check if that field is taken.

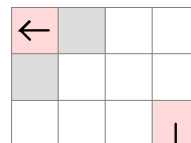


- If so, we mirror sides and start the algorithm on the right side.
- If not, we increase the horizontal distance by one until we find a taken field or run into the border.

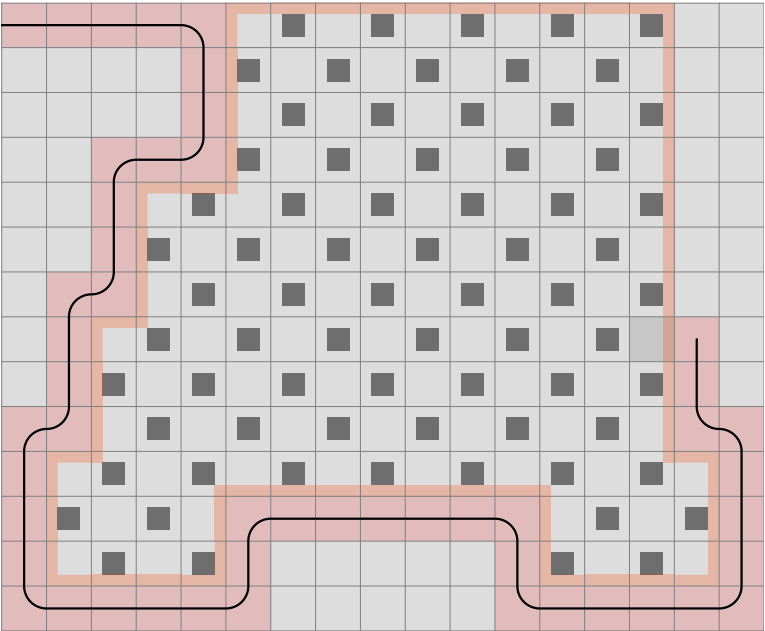


If it is a border field, we increase the vertical distance and start with 1 horizontal distance again.

Otherwise, we check if the bottom field is free, and by comparing the index of the corner field with the field above, we can determine if the line is going down and left, so the enclosed area is on the side we want.

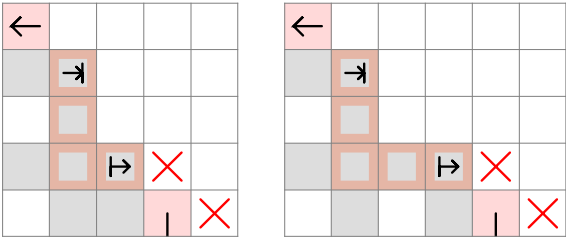


And as we step back, we find the point where the line should move in another direction.

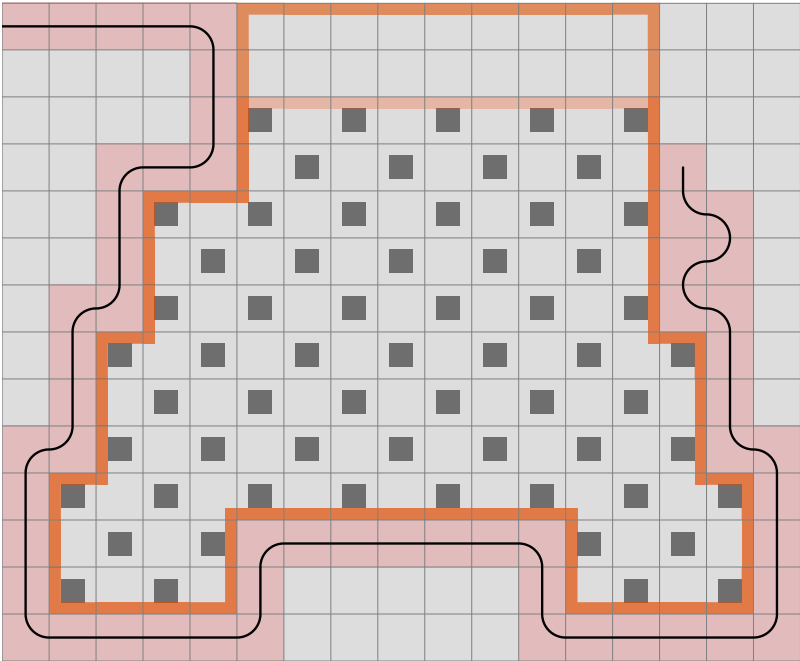


The only difference is the added 2x2 area. To the first, we apply the straight-to-side algorithm, while at the second, we have a corner that defines the area, but essentially, the procedure is the same.

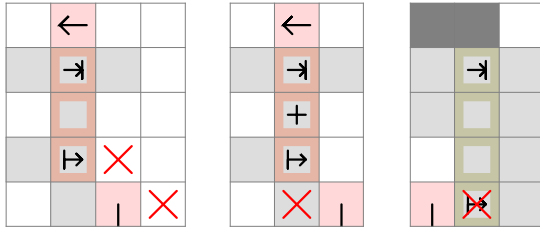
Having the universal algorithm, these two size-specific rules can be deleted:



As we continue the case above, soon we will discover a deficiency which actually has been visible all along.



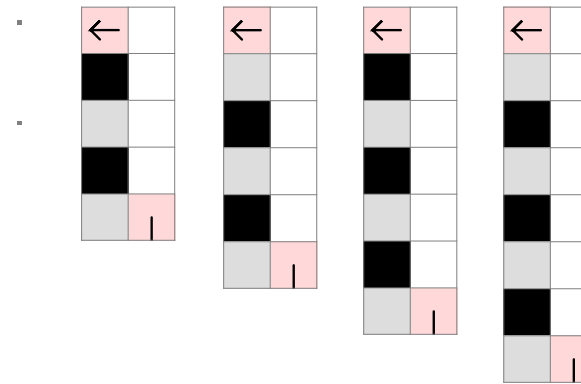
With the border on top, now we have no option to move.  
The following rules are active, in addition to the universal one that disables the left field:



There are three more black fields in the area than white, so there is not enough vertical distance for entering and exiting that many times.  
The straight-to-side algorithm has to be rotated upwards, so we get this:



Let's look at the distances from 3 to 6 to have an example of each case of modulo 4.



$D \text{ (distance)} \% 4 = 3$ :

Now:  $1W \rightarrow 0B$ ,  $2W \rightarrow 1B$  etc. =  $(D+1)/4 W \rightarrow (D-3)/4 B$   
Later:  $0W \rightarrow 1B$ ,  $1W \rightarrow 2B$  etc. =  $(D-3)/4 W \rightarrow (D+1)/4 B$   
(double rule)

$D \% 4 = 0$ :

Now:  $1W \rightarrow 0B$ ,  $2W \rightarrow 1B$  etc. =  $D/4 W \rightarrow D/4 - 1 B$   
Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $D/4 W \rightarrow D/4 B$   
(single rule)

$D \% 4 = 1$ :

Now:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-1)/4 W \rightarrow (D-1)/4 B$   
Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-1)/4 W \rightarrow (D-1)/4 B$   
(no rule)

$D \% 4 = 2$ :

Now:  $2W \rightarrow 1B$ ,  $3W \rightarrow 2B$  etc. =  $(D+2)/4 W \rightarrow (D-2)/4 B$   
Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-2)/4 W \rightarrow (D-2)/4 B$   
(single rule)



The figure consists of two side-by-side 10x10 grids. The left grid shows a 5x5 orange square representing a kernel applied to a 5x5 black square representing the input image. The right grid shows the result of applying the kernel, where the output is a blurred version of the input, with the 5x5 black square now appearing as a 5x5 gray square.

- 0 vertical distance small area
- x horizontal and y vertical distance small area
- 0 horizontal distance small area
- 0 horizontal distance big area

The diagram shows a complex, irregularly shaped region on a grid. The region is defined by a black outline and filled with a light red color. The region has several protrusions and indentations. A large rectangular area at the bottom is outlined in orange and contains a grid of small black squares.

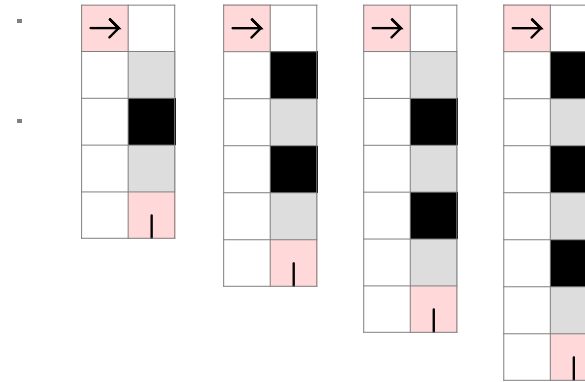
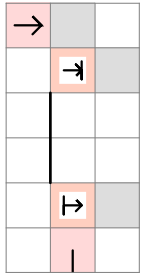
			←
	→		
	+		
	→	×	
			.

The number of black fields is one more than the white. Without the size-specific rule, we could step straight. Let me remind you that this rule is based upon the Double C-Shape. If we stepped straight and then into the area, we would come out in the middle, creating two C-shapes.

What this actually means is that there would be two fields of the same color that cannot be filled simultaneously.

If we now extended the area to include the 4 fields straight ahead, there would be still one more black than white, and by stepping straight, to a white field, it is clear that a black to black line cannot be drawn.

An extension of the universal rule is necessary to include cases of a "big" area where the obstacle is on the other side of the live end.



$D \% 4 = 3$ :

Now:  $1W \rightarrow 0B = (D+1)/4$   $W \rightarrow (D-3)/4$   $B$

Later:  $1W \rightarrow 0B = (D+1)/4$   $W \rightarrow (D-3)/4$   $B$   
(no rule)

$D \% 4 = 0$ :

Now:  $1W \rightarrow 0B = D/4$   $W \rightarrow D/4 - 1$   $B$

Later:  $1W \rightarrow 1B = D/4$   $W \rightarrow D/4$   $B$   
(single rule)

$D \% 4 = 1$ :

Now:  $2W \rightarrow 0B = (D+3)/4$   $W \rightarrow (D-5)/4$   $B$

Later:  $1W \rightarrow 1B = (D-1)/4$   $W \rightarrow (D-1)/4$   $B$   
(double rule)

$D \% 4 = 2$ :

Now:  $2W \rightarrow 1B = (D+2)/4$   $W \rightarrow (D-2)/4$   $B$

Later:  $1W \rightarrow 1B = (D-2)/4$   $W \rightarrow (D-2)/4$   $B$   
(single rule)

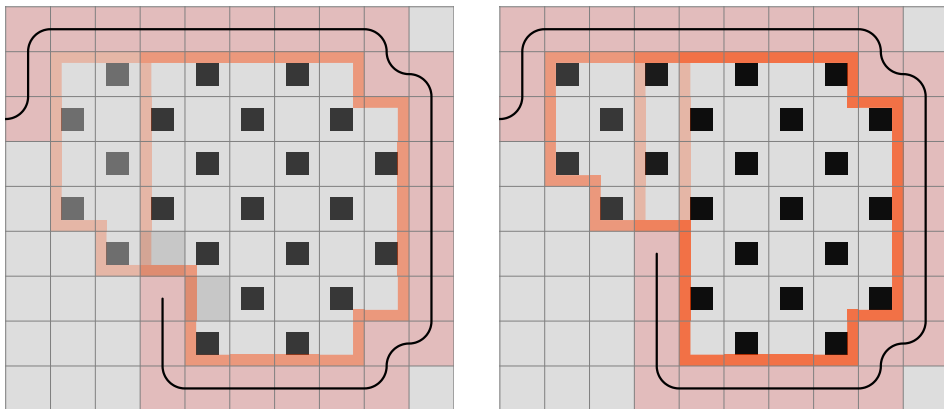
General:

Now: If  $n > 0$ , we can use all corner blacks after exiting the first line.

$(n+2 - (n+2) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$  if  $n > 0$ .

Later:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$  if  $n > 0$ .

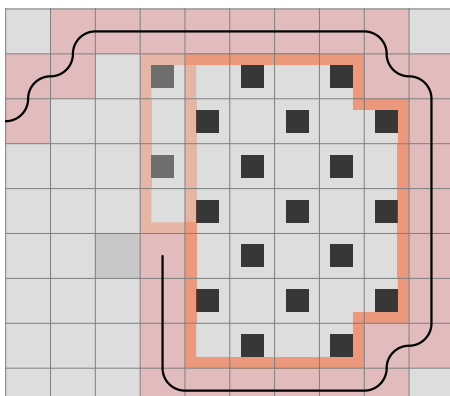
While creating a case to verify the newly created rule set, I have encountered this:



It is not possible to continue after stepping upwards.

But where is the missing part?

On the left, the largest area contains 1 more white field than black, and on the right it is 2 more black. It would be possible if the upper left corner was filled, like this:



I will now take the case of  $x$  horizontal and  $y$  vertical distance big area.

All representations are the same as in the small area case, only the area is on the other side.

So I will just summarize the black and white limits here.

### 1. Equal horizontal and vertical distance

$x = 2$ :

Now:  $0W \rightarrow 1B$

Later:  $1B$

$x = 3$ :

Now:  $0W \rightarrow 2B$

Later:  $1B \rightarrow 2B$

$x = n$ :

Now:  $0W \rightarrow (n-1)B$

Later:  $1B \rightarrow (n-1)B$

Exactly the same as with the small area.

### 2. Larger horizontal distance

#### If the added distance is pair:

$2n = 2$ :

Now:  $1W$  is possible. The entry field and the first white is at least 2 distance from each other, the whole area can be filled between them.

$1W \rightarrow xB$

Later: All corner blacks can be used. The line between the first black and first white will fill the area.

$0W \rightarrow xB$

General:

The only difference is the  $n = 1$  case. Otherwise, the number of inline and corner fields are the same.

Now:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

**If the added distance is impair:**

$2n + 1 = 1, n = 0:$

Now:  $xW$  is possible.

$xW \rightarrow 0B$

Later: Same values as previously.

$(x-1)W \rightarrow 0B$

$2n + 1 = 3, n = 1:$

Note that in the small area case, there were  $x-1$  corner whites and 1 corner black. Now there are  $x$  amount of corner whites and 2 inline blacks instead of 1.

Now:  $xW \rightarrow 0B$

Later:  $xW \rightarrow 1B$

$2n + 1 = 5, n = 2:$

Now:  $(x+1)W \rightarrow 1B$

Later:  $(x+1)W \rightarrow 1B$

General:

Now:  $x + (n+1 - (n+1) \% 2) / 2 W \rightarrow (n - n \% 2) / 2 B$

Later:  $x + (n - n \% 2) / 2 W$  if  $n > 0 \rightarrow (n+1 - (n+1) \% 2) / 2 B$

3. Larger vertical distance

**If the added distance is pair:**

We will find it is the same as the small area.

$2n = 2:$

Now:  $1W \rightarrow xB$

Later:  $0W \rightarrow xB$

$2n = 4:$

Now:  $1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

General:

Now:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

**If the added distance is impair:**

$2n + 1 = 1, n = 0:$

Now:  $1W \rightarrow (x-1)B$

Later:  $0W \rightarrow (x-1)B$

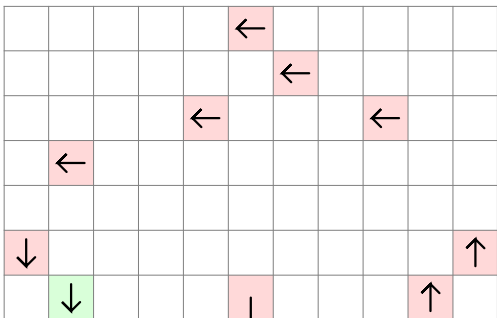
$2n + 1 = 3, n = 1:$

Here comes the change again, due to having one more corner black field and one less corner white field than in the small area case.

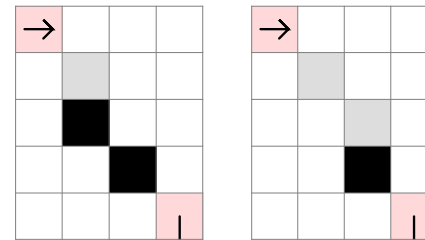
Now:  $1W \rightarrow xB$

Later:  $1W \rightarrow xB$

Now when there is a C-Shape, it is possible to step there, and no other rules will be checked. But C-Shape is just one of the nine cases of area checking.

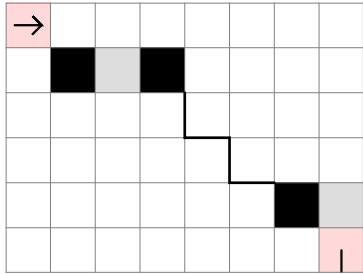


One thing is sure, we have been using the small area representations when defining the rule set, which does not give us the minimal area in this case. See the difference:



Notice that these are the same as the small area patterns, things are just mirrored, so that the previous horizontal expansion is now vertical.

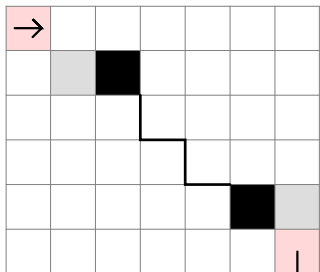
n = 1



Now:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

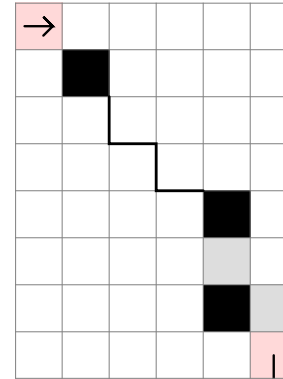
n = 0



Now:  $1 + (n+1 - (n+1) \% 2) / 2 W \rightarrow x - 1 + (n - n \% 2) / 2 B$

Later:  $1 + (n - n \% 2) / 2 W$  if  $n > 0$  (0 if  $n = 0$ )  $\rightarrow x - 1 + (n+1 - (n+1) \% 2) / 2 B$

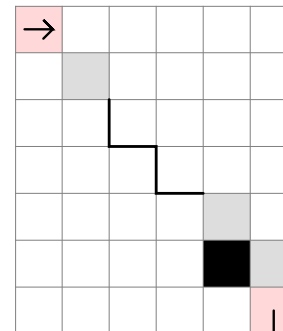
n = 1



Now:  $(n+1 - (n+1) \% 2) / 2 W$  if  $n > 1$  (0 if  $n = 1$ )  $\rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

n = 0



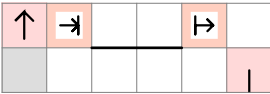
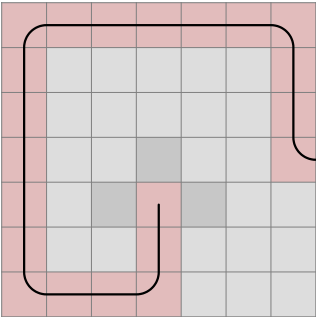
Now:  $x - 1 + (n+2 - (n+2) \% 2) / 2 W$  if  $n > 0$  ( $x - 1$  if  $n = 0$ )  $\rightarrow (n+1 - (n+1) \% 2) / 2 B$

Later:  $x - 1 + (n+1 - (n+1) \% 2) / 2 W \rightarrow (n+2 - (n+2) \% 2) / 2 B$  if  $n > 0$  (0 if  $n = 0$ )

The difference between stepping up and right still remains in these vertical expansion cases. We didn't have to deal with it at the small area, because the line could not step backwards.

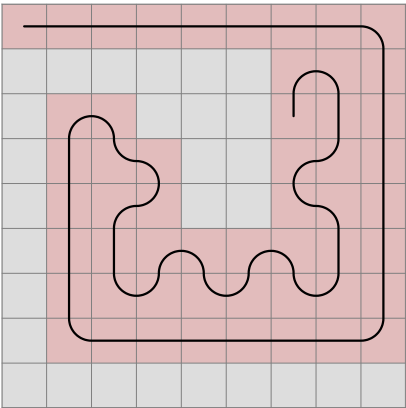
It just means we have to remove the Now W conditions for stepping right.

And here is an example of the x horizontal, 0 vertical distance big area corner. None of the existing cases cover it. The area is impair, so we shouldn't be able to step right.

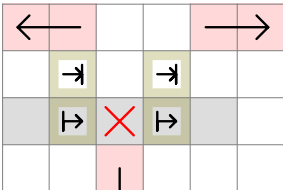


When we re-run the program on 9 x 9, it will be obvious that covering all the above 9 cases of area checking still does not solve everything.

If we don't apply the first size-specific rule, we get stuck here:



Remember we have made rules previously that take both sides into account, like this:

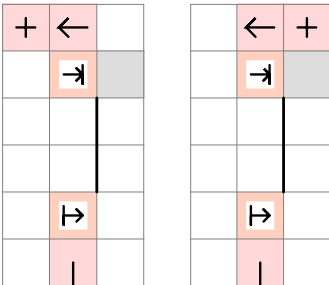


To combine all cases where after the next step, on the left side only the left field will be possible, and on the right, only the right, is not feasible.

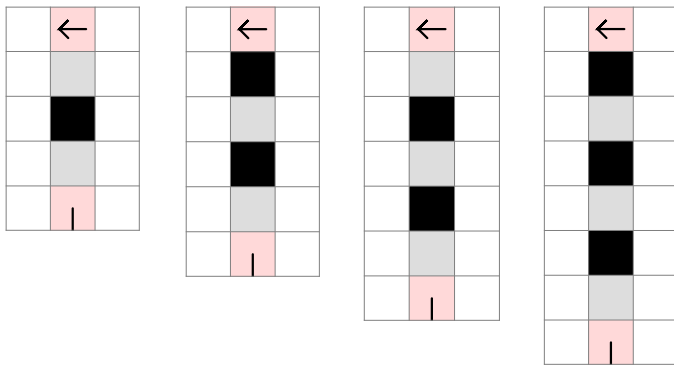
Let's say instead that if one of the possible fields will lead to an impossible situation, we cannot step there. It still results in a usable algorithm.

(An algorithm is usable if the number of operations are on the order of  $n^2$  where  $n$  indicates the size of the table.  
An unusable algorithm would require a number of operations on the order of  $2^n$  or more: that is the random path where you just guess the next step until you get an impossible situation, but the crossroad might have been as far back as almost  $n \times n$  steps.)

We are still missing the straight obstacle case:



Besides all the fields straight ahead, the gray field has to be empty in order to apply the rule.



Notice the area is the same as with the case where the obstacle is one field to the right, but there is a difference:  
 Look at the 3-distance case. Since the obstacle is straight ahead, if we enter now by stepping left, we cannot exit at the black field, because then one of the whites could not have been filled.  
 The area can only be filled if it is 1W. (This is the case we previously called Double C-Shape.) Adding 4 extra distance does not make the problem disappear: if we exit at the first black and the previous field was the first white, now we get the 5-distance case with 3 white fields and 2 black. A black to black line is therefore not possible.

These are the new values:

$D \% 4 = 3$ :

Now:  $1W \rightarrow -1B = (D+1)/4$   $W \rightarrow (D-7)/4$  B

Later:  $1W \rightarrow 0B = (D+1)/4$   $W \rightarrow (D-3)/4$  B  
 (single rule)

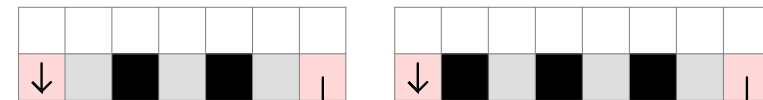
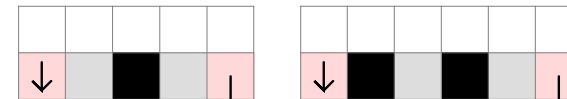
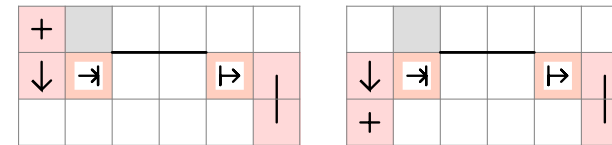
$D \% 4 = 5$ :

Now:  $2W \rightarrow 0B = (D+3)/4$   $W \rightarrow (D-5)/4$  B

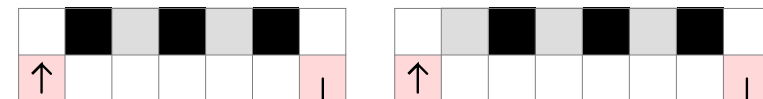
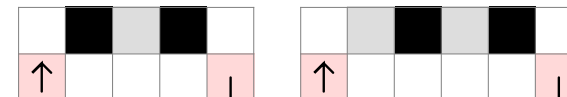
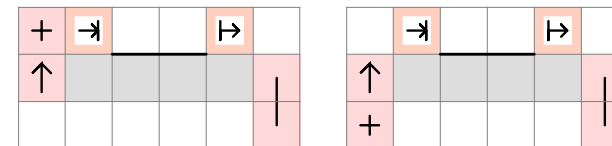
Later:  $1W \rightarrow 0B = (D-1)/4$   $W \rightarrow (D-5)/4$  B  
 (single rule)

Also, pay attention to the 2-distance case. If we enter now by stepping left, 1W is possible. If we step straight, it is 0W. When the distance is 6, there is no difference, because if the step straight, we can exit the area immediately.

Side straight cases, small area:



Big area:





The figure consists of two 10x10 grids. In both, red cells represent obstacles and light blue cells represent free space. A black line indicates a path from the top-left corner (0,0) to the bottom-right corner (9,9).

**Left Grid:** The path starts at (0,0), goes right to (4,0), then down to (4,1), right to (5,1), up to (5,2), right to (6,2), down to (6,3), right to (7,3), up to (7,4), right to (8,4), down to (8,5), right to (9,5), and finally down to (9,9).

**Right Grid:** The path starts at (0,0), goes right to (4,0), then down to (4,1), right to (5,1), up to (5,2), right to (6,2), down to (6,3), right to (7,3), up to (7,4), right to (8,4), down to (8,5), right to (9,5), and finally down to (9,9). This path is identical to the one in the left grid.

The four grids illustrate the step-by-step solution of a 5-disk Tower of Hanoi puzzle. The grids are arranged in two rows of two. Each grid is a 5x5 table representing the state of the puzzle. The top row shows the initial state and the first three moves. The bottom row shows the final state and the last three moves. The goal state is the final state, which is the goal state.

**Grid 1 (Top Left):** Initial state. Disk 1 is on column 1, row 5. Disk 2 is on column 2, row 4. Disk 3 is on column 3, row 3. Disk 4 is on column 4, row 2. Disk 5 is on column 5, row 1. The goal state is reached.

**Grid 2 (Top Right):** State after move 1. Disk 1 is on column 2, row 5. Disk 2 is on column 3, row 4. Disk 3 is on column 4, row 3. Disk 4 is on column 5, row 2. Disk 5 is on column 1, row 1. The goal state is reached.

**Grid 3 (Bottom Left):** State after move 2. Disk 1 is on column 3, row 5. Disk 2 is on column 4, row 4. Disk 3 is on column 5, row 3. Disk 4 is on column 1, row 2. Disk 5 is on column 2, row 1. The goal state is reached.

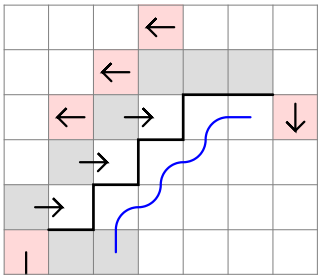
**Grid 4 (Bottom Right):** State after move 3. Disk 1 is on column 4, row 5. Disk 2 is on column 5, row 4. Disk 3 is on column 1, row 3. Disk 4 is on column 2, row 2. Disk 5 is on column 3, row 1. The goal state is reached.

The diagram shows a 6x6 grid world. Red squares represent obstacles, forming a large L-shape on the left and top-left. Green squares represent goal areas or targets. A black path starts at the bottom-left, moves right, then up, then right again, ending at the top-right. There are two green squares along the path: one at (3, 4) and another at (4, 5). There are also two green squares off the path: one at (2, 5) and another at (3, 6).

What happens here is that the  $3 \times 2$  area has a certain exit point, the middle, because the area is pair. Then it will create a C-Shape on one side and an area on the other side with a 1-wide gap.

The figure consists of three 5x5 grids illustrating the step-by-step construction of a Young diagram for the partition (4, 3, 1).  
 Grid 1: The bottom-left cell (row 5, column 1) contains a vertical bar. The cell at (row 5, column 5) is pink and contains a downward arrow.  
 Grid 2: A new cell is added to the top row (row 1, column 4), which is grey. The cell at (row 5, column 5) remains pink with a downward arrow.  
 Grid 3: A new cell is added to the second row (row 2, column 3), which is grey. The cell at (row 5, column 5) remains pink with a downward arrow.

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If an area is created with either of the directional obstacles, it cannot be filled.

The algorithm first checks if the two fields to the right are empty. Then, starting from the field straight ahead, it takes each field in the row until an obstacle is encountered. If the horizontal distance of the obstacle is vertical distance + 3, the area is of the desired shape and the 3 empty fields on the top are checked.

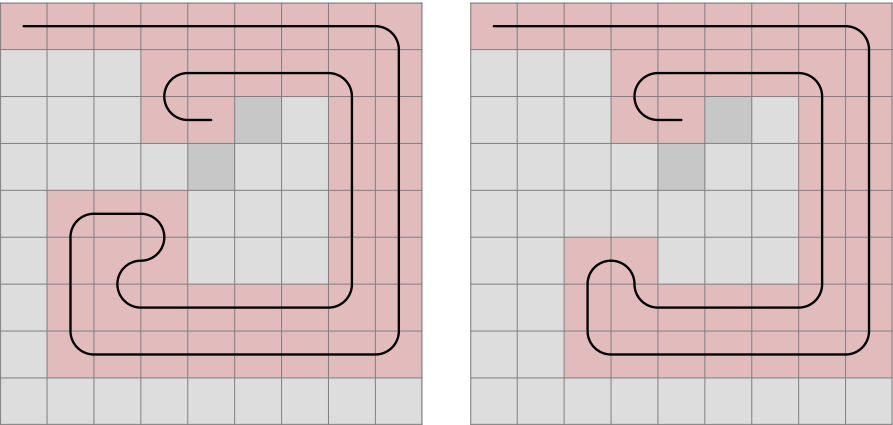
If the number of black and white fields are equal, the program checks at each corner point on the border if there is a mid across obstacle going to the left, creating a small area on that side.

The whole process is then repeated, increasing the vertical distance, until a non-empty field is found at  $n$  horizontal and  $n + 1$  vertical distance.

The fields on the border (blue line) are of course also checked for being empty.

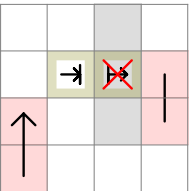
You could make the process more bulletproof by including the across obstacles at the corner points or the straight obstacle at the top, but is it necessary? It is not proven until we have a case for them.

- But let's continue the program. There are still single area cases we haven't thought about yet. (Previously we had rules for them, but not in the new system.)



- Obviously, at 1 distance we cannot step right, but neither can we at 2 distance if the area is impair. It is because stepping straight allows for 1W, while stepping right only 0W is possible.

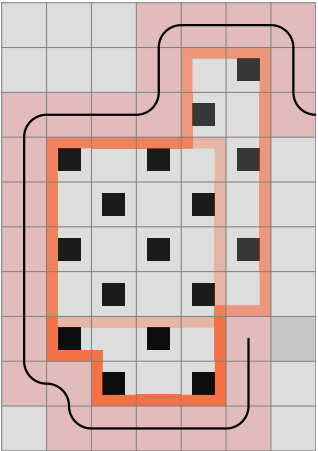
- Previously, we represented it like this:



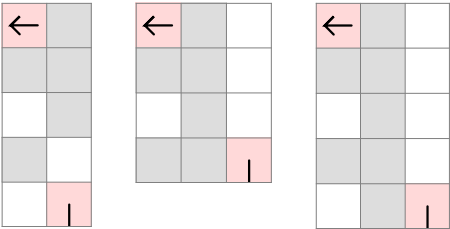
Reviewing those 2-distance rules, we can see that many of them has a double area (one of which is a C-shape). They cannot be solved with the single area patterns. It is best to re-enable the whole set even if it is a repetition in other cases.



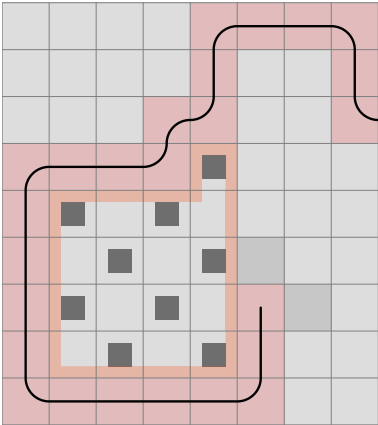
So far, we simply assumed that it is the opposite, just because the area is impair, but such case does not exist. Because of the single area rule with the obstacle, the step to get there is disabled. In the following example the desired area is 1B. So is the area defined by the obstacle.



When checking the conditions for the first area, these are the fields that need to be empty:



Notice that the first two cases can exist simultaneously. Then we examine the smaller area (second case); the obstacle in the first will create a C-shape with the exit of that area and the leftwards step will be disabled by the single area rule:

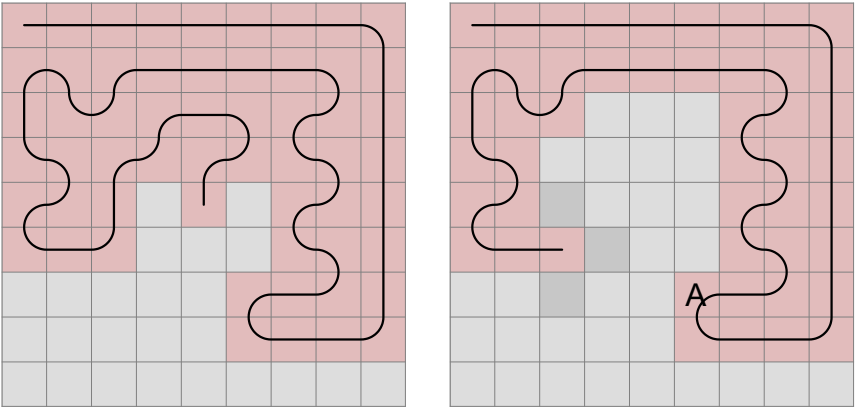


Can the border line be longer?

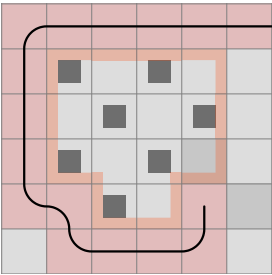


If we look at 4 distance: It is true that if the area is 1W and we step forward, only one white field remains on the border, but we can also exit immediately, leaving a 3-long border line and a 0W area. Then, the area can be entered at the second white and exited at the last black.

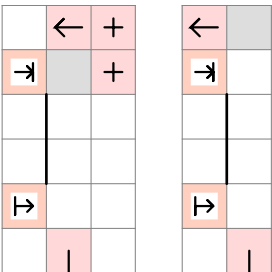
We get stuck at 641 027 in a case that was previously solved with Double C-Shape.



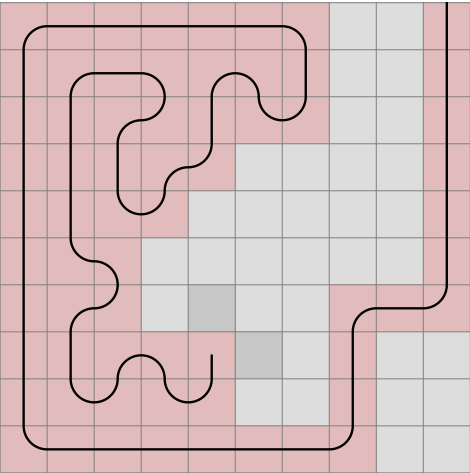
If we didn't have the taken field marked with A, the straight obstacle rule would rightfully disable the left field.



But, since we want to keep the field next to the furthest border field empty, we need to apply a rule for the smaller area. This will be the same as in the 0 horizontal distance small area case.

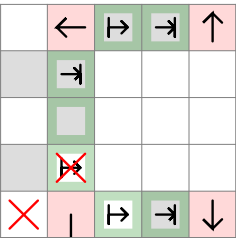


We can see that the mechanism is not limited to a few cases, but it can go on indefinitely. If after exiting the area, there are repeated C-shapes on the left side (a stair shape), the right side obstacle may be far away, as in this example:



A sequence is made from the start area and the close obstacle cases, and this can be programmed.

But let's build the program step by step, based on the discovered 9 x 9 rules. The second case, Triple Area, uses an area where the obstacle is 3 distance away.

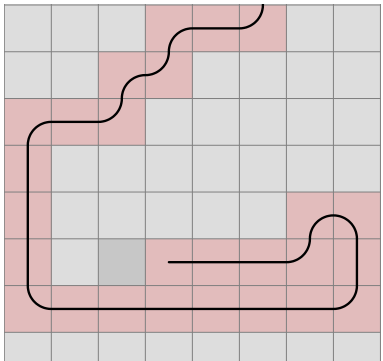
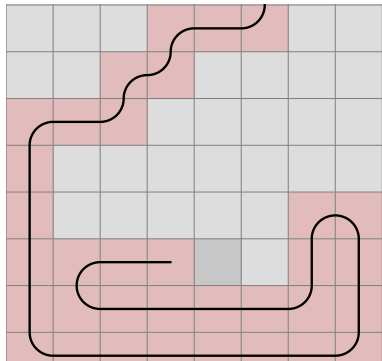


If the area is pair, the exit field will be straight in the middle, and the next step is one to the right. That's where we can check for close obstacles on either side. Stepping to the left is already disabled in the single area straight obstacle rule. Otherwise, the pattern is rotated, like in the previous case.

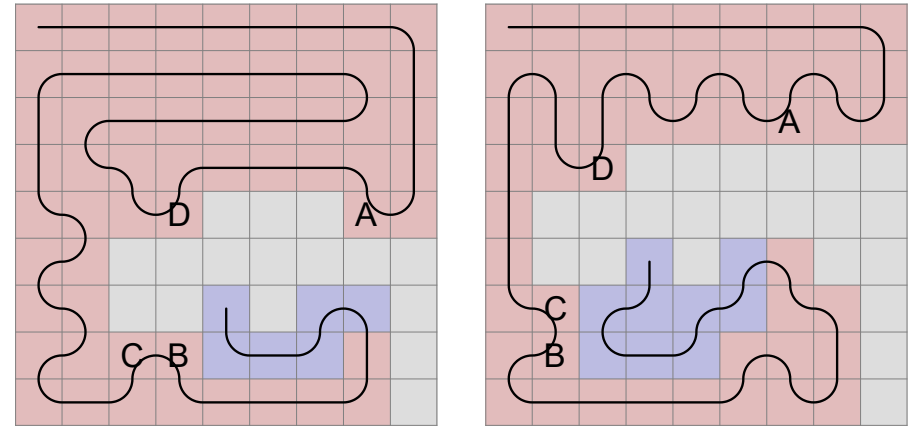
	←	↔	→	↑		
	→	A	B		+	
	<del>↔</del>			C		+
<del>×</del>						
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122 123

		+		
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←				
↖				
→	×	→	↖	↓
	↓			

[illegible]

Now, take a look at the following two cases. The first is the well-known Triple Area Exit Down, at over 18 million, while the other comes at around 51 million.



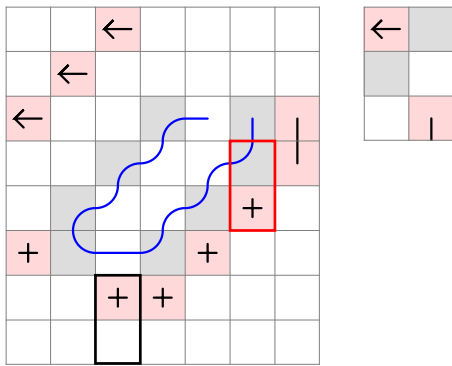
In the first, we can only step left. In the second, we cannot step left. It is easy to see that the pattern can be longer. If there is a stair shape downwards with 3 fields at the bottom, there will be a stair backwards, which conflicts with the obstacle on the left (D).

But there are more things that need to be present:

- There should be an obstacle at 2 distance straight ahead to start with.
- If the B and C weren't taken, the area could be filled.

What is the simplest algorithm to apply?

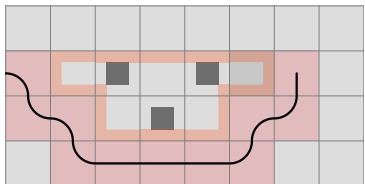
In the following, I check these fields as empty, taken or directional taken.



The area enclosed with red line is repeated downwards left, until both of the fields are taken.

For each return step, I check a mid across directional field. Their right and down field need to be empty as well.

Why is C-shape checking to the left unnecessary? It is already taken care of.



To accomodate the Triple Area Exit Down case, let's first check the presence of an obstacle straight ahead at 3 vertical distance. Stepping one forward, we can apply the same procedure.