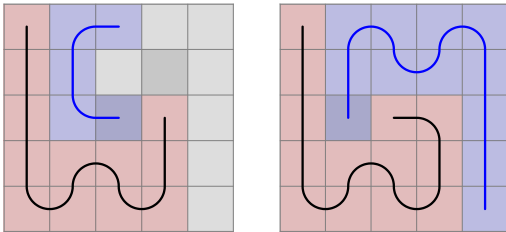
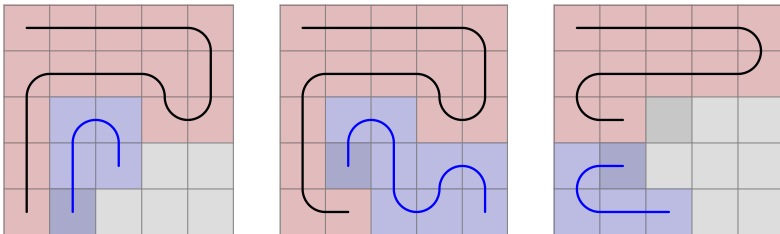


- A 2 x 2 empty area next to the live end that is walled by three sides (2-2-2 long) will have a future line going through along the walls. In this example, the far end is already extended by one step as it had only one option to move.



- Future line extension when we step on a future line: The far can be extended if it was 2 distance away from the near end. It can now fill the C-shape.



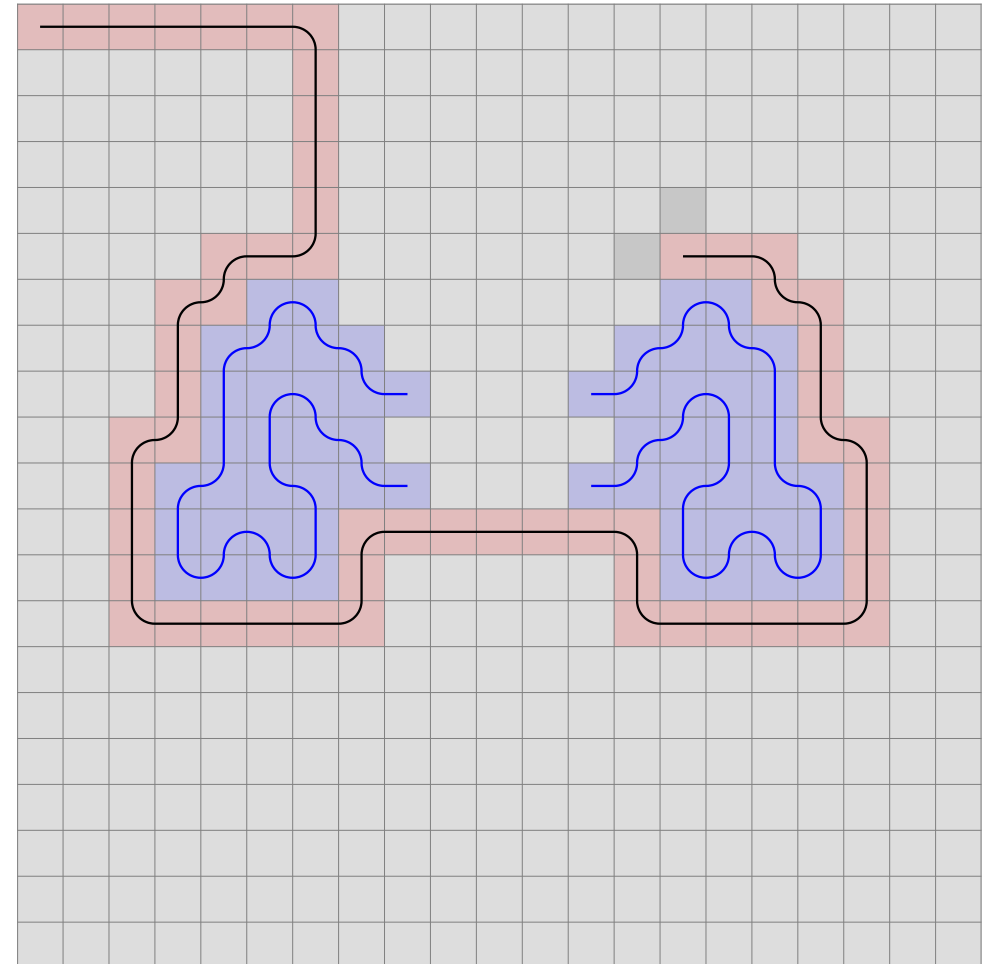
The same goes with 1 x- and y-distance. A C-Shape is not always created in this case.

## The one-way labyrinth algorithm

This research aims to solve the following problem:

"Draw a line that goes through an  $n \times n$  grid (where  $n$  is an odd number), passing through each field once. The line has to start from the field at the upper left corner (1,1) and end at (n,n). At any time it is allowed to move left, right, up or down, and it has to randomly choose between the available fields."

At first sight it may look easy. But look at the following example:

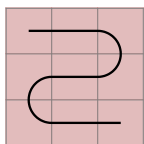


Based on the black line's movement, blue fragments were drawn to indicate a path we have to go through in the future in order to fill the board.  
Do you see why the situation is impossible from now on?

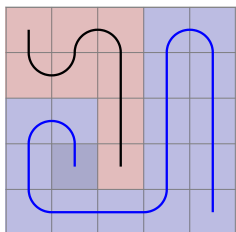
The question is, is there a single rule or a set of rules that will guarantee you can draw a labyrinth of any size? Or do the rules get infinitely complex?

To assist with the research, I have written a computer program. In the beginning, I let it run on a 21 x 21 field, and whenever I noticed a trouble, I coded the solution into it. While you can discover many patterns this way, they will be random and do not help in gaining a fundamental understanding. At one point you will find things get too complex, and you are still far from solving the 21 x 21 board.  
That's where a gradual approach comes in.

A 3 x 3 area can only be filled in two ways, like this and mirrored:

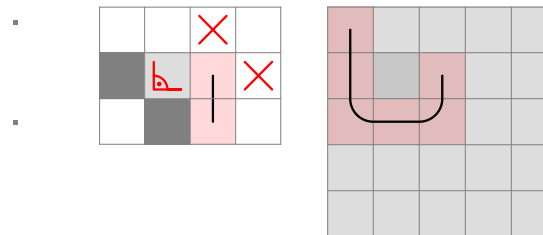


The 5 x 5 requires much more consideration. Whenever it is possible to draw future lines, the program has to be able to do it. The future lines can not only extend at each step but connect too.

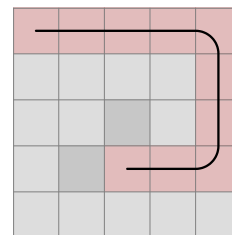


By August 21, 2023 all 5 x 5 scenarios were discovered. The number of walkthroughs are 104.

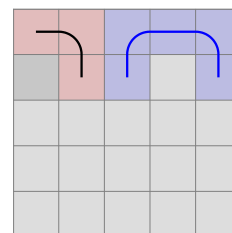
Here are the things to consider on a grid of this size:



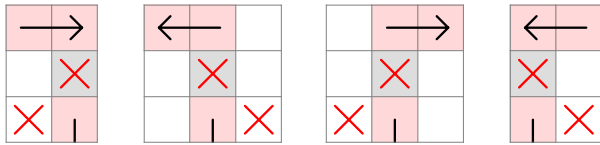
- A single field next to the live end that is walled from two other sides (either by the border or the line) needs to be filled in the next step. I call it C-shape. The pattern is both mirrored and rotated, so that the empty field is straight ahead. To qualify for this rule, the empty field cannot be the end corner. If there is a C-shape, we don't need to check other rules.



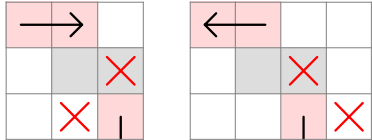
- Movement near the edge: In the example, we cannot step left (3,5), since the (2,5) field is empty.



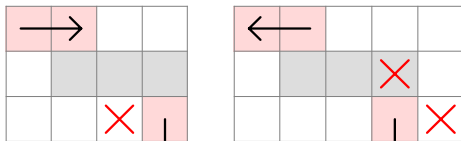
- A 2 x 3 empty area next to the live end that is walled by three sides (2-3-2 long) will have a future line going through along the walls. At the wall next to the main line, its direction is the opposite of the main line, meaning it will go from (3,2) upwards whereas the main line just took a step downwards. How the middle field will be filled is not yet known. Either the near end (the one the main line will go through first) or the far end can fill it.



The gray square means empty field. When the field 2 to straight is taken, its left or right side will be taken too.

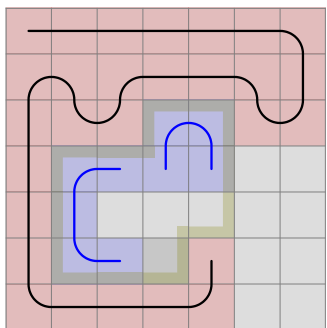


These will only be checked if one of the above 4 situations were not present. (They have to be mirrored, too.)

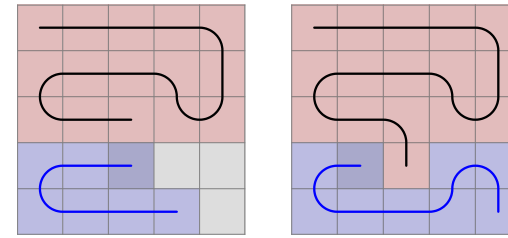


Likewise, these will not be checked if the previous rules were true.

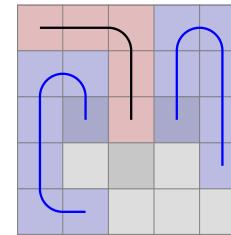
And when none of the 1-distance situations are valid, we check for 2-distance.



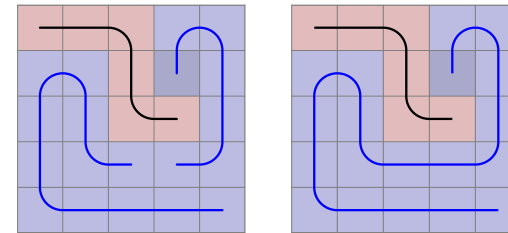
Impair areas can now happen inside the grid, not just on the edge, and the following rules have to be applied:



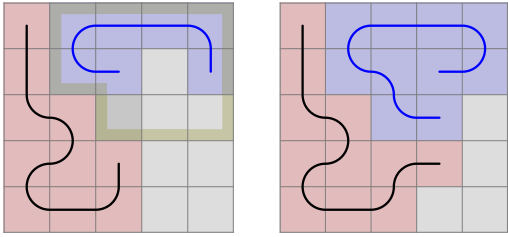
If the far end was near the end corner, it has to choose the other empty field.



- Future line extension when stepping away: If there was a near end where the main line was in the previous step, it now may have only one choice to move, so it can be extended.



- Future line connection: In this case, the line being stepped on extends until the far end has two options. (When the end corner is one of them, it has to be removed.) Then, the line on the left extends and now has no other option than to connect to the line on the right.

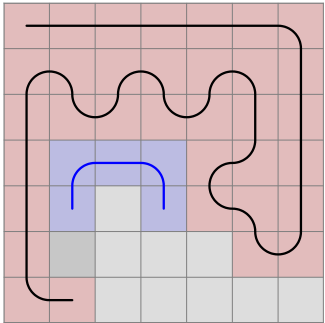


- When we are two distance away from the edge, we need to check if stepping towards it is possible.

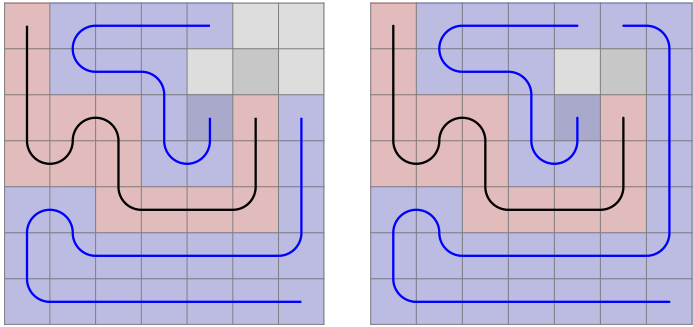
It is because if we do so, an enclosed area is created, with one way to go out of it. If that area has an impair amount of cells, it cannot be filled, so we cannot take that step.

The explanation is simple: Imagine if the table was a chess board. In order to step from white to black, you would need to take an impair amount of steps - the color changes at every step. Here, the entry of the area would be (4,3) and the exit (5,3). An impair amount of steps means pair amount of cells.

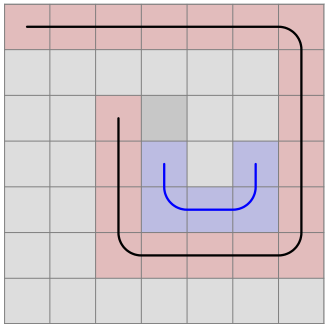
In the example, you can also say that we cannot step right, because there is a future line start 2 to straight and an end 2 to straight and 2 to right. On 7 x 7, there will be examples where this is the rule we have to apply, because area counting is not getting triggered:



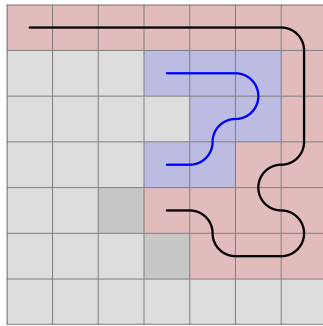
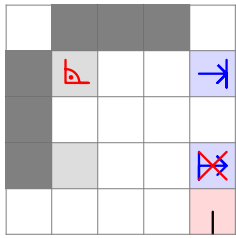
- But let's start with the simpler rules:
- - Future line extension: When a near end is at 2 distance left or right from the live end, it will fill the field between them if the live end steps elsewhere. That's what happened in the 5 x 5 example above before the line failed.



In other situations, there is a 1-thin future line next to the live end that can be extended if its far end is at the corner. Though disabling this rule does not affect the total amount of walkthroughs on a 7 x 7 grid, I chose to include it in the project on the basis that if a future line can be extended, we should do it. It can make a considerable difference. The left picture is without the rule, the right is with it.



- - Just like moving near the edge, we need to disable some fields if we are approaching an older section of the main line. In order to determine on which side the enclosed area is created, we need to examine the direction of the line at the connection point.



- And these are the remaining size-specific rules. Future 2 x 2 Start End, Future 2 x 3 Start End and Future 3 x 3 Start End.

The program, in fast mode, can run through approximately 100 cases per second, depending on your computer speed. This enables us to discover all 7 x 7 walkthroughs, which is 111 712.

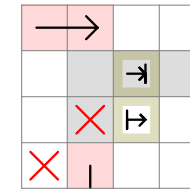
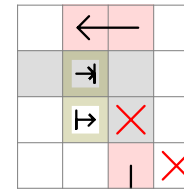
It is equal to what is described in the Online Encyclopedia of Integer Series (Number of simple Hamiltonian paths connecting opposite corners of a  $2n+1 \times 2n+1$  grid).

As the sizes grow, it will be impossible to run through all cases with one computer in a reasonable time. In order to discover the patterns, we need to run the program randomly.

Is it possible to develop an algorithm that works for all sizes? The edge-related and area-counting rules are universal, but the size-specific rules get more and more complex. Can you define them with one statement?

I have made statistics about how many random walkthroughs you can complete on different grids using the 7 x 7-specific and the universal rules before running into an error. Based on 1000 attempts, here are the results:

9: 19.5  
11: 5.7  
13: 2.6  
15: 1.2  
17: 0.7  
19: 0.4  
21: 0.2

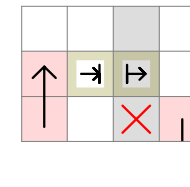
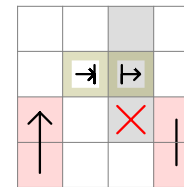


The procedure is similar to the the straight 2-distance rule. The only difference is that we count the area starting and ending at the marked fields. In the first, the direction of the circle is left, in the second right.

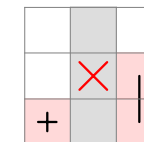
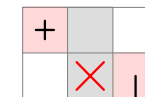
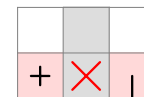
Besides mirroring them, we also have to rotate them both counter-clockwise and clockwise.

But we do not need 12 of such rules. Taking the first, the live end cannot come from the left, because the area parity was already checked in the previous step, and now we just added 2 fields to it. It can come from the right, and then there is naturally only one field we might have to disable.

Here are the representations of the two scenarios for the left side:

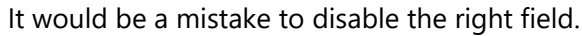


Similarly to the straight rules, these will only apply if there is no wall 2 distance to the left or right. Let's construct these preconditions.



We are not finished. Did you notice the example above is not covered by these rules? We have to move the taken fields 1 and 2 steps to the side, both in straight and side direction.

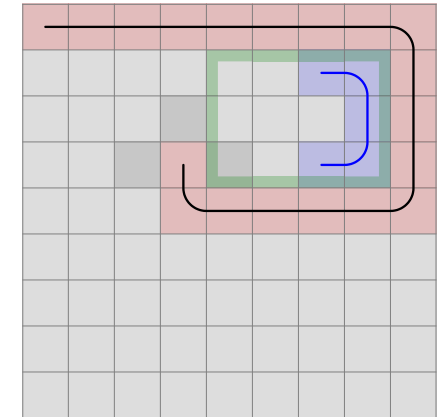




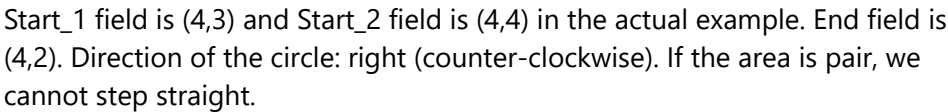
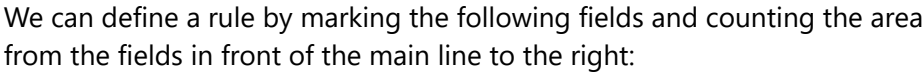
|  |   |  |   |   |
|--|---|--|---|---|
|  |   |  |   | + |
|  |   |  |   | → |
|  |   |  |   | → |
|  | → |  | → | + |

You may ask, why that field is "taken", not "taken or border". From what I found through some examples, if that field is border, the enclosed area on the right is impair, so the line cannot step in the other direction anyway. But it needs further examination.

A 10x10 grid with red and blue shaded cells. A black line starts at (0,0) and ends at (9,9), following a path that avoids blue cells. A blue line starts at (0,0) and ends at (9,9), following a path that avoids red cells.



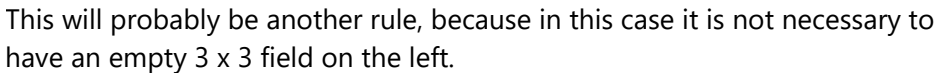
16      13



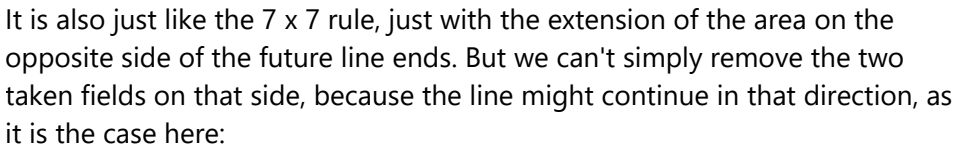
When generating code from the drawing, we have to check on which side the enclosed area was created. Here, we want it to be on the right side, so there are two cases to look at:

- The taken or border field beyond the end field is a taken field. In this case, if the field to its left is taken, its index must be lower. If the field to the right is taken, its index must be higher.
- It is the border. Add together the x- and y-coordinates to get a value. A higher value is closer to the end corner. Here, we compare the border field straight ahead and on its left, and we want the first-mentioned to be the smaller.

I have applied this rule rotated clockwise (besides mirroring it, of course), so that the live end can both come from the bottom and the right. But it can also come from the left in this example:



Now let's run the program further up to number 13 992:

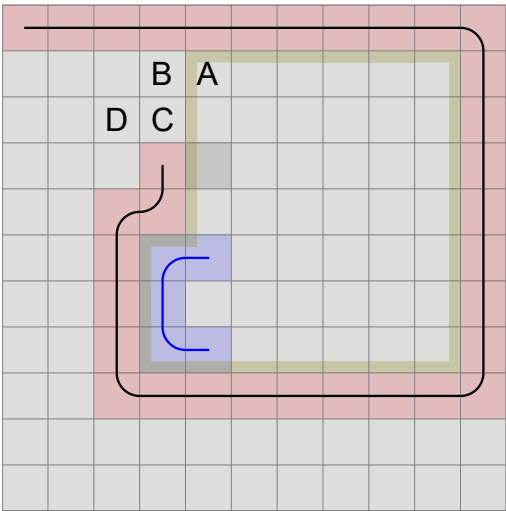
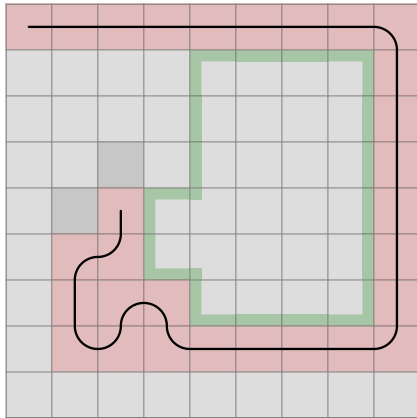




A 5x5 grid with various symbols and a path. The symbols are: a grey square at (1,1), a blue arrow pointing right at (2,1), a blue square with a red 'X' at (3,1), a black vertical line at (4,1), a black plus sign at (4,5), and two black plus signs at (5,2) and (5,3). A path is shown with a blue line starting at (1,5), going right to (5,5), then left to (5,1), then up to (2,1), and finally right to (2,2). A black line starts at (2,2), goes right to (2,3), then down to (3,3), then left to (3,2), then up to (3,1), then right to (4,1), then down to (4,2), then right to (4,3), then down to (5,3), then right to (5,4), then up to (4,4), then right to (4,5), then up to (3,5), then left to (3,4), then up to (2,4), then right to (2,5), then up to (1,5), and finally right to (2,5).

We can see that defining a rule with future line starts and ends does not tell us on which side the future line was created. That is the side that contains the enclosed area. We need to therefore replace such rules with area counting, which we actually already did, with the exception of Future L. Here the future line couldn't have been created on the other side, because that's the side the live end is at right now. And area counting is not always possible, like in this situation:

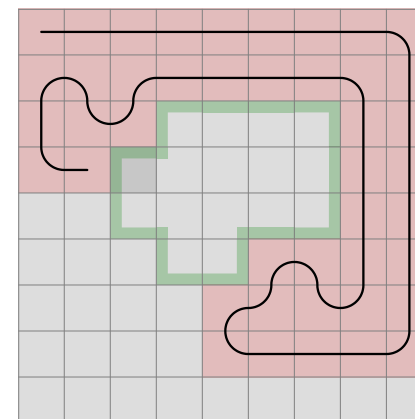
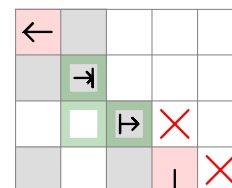
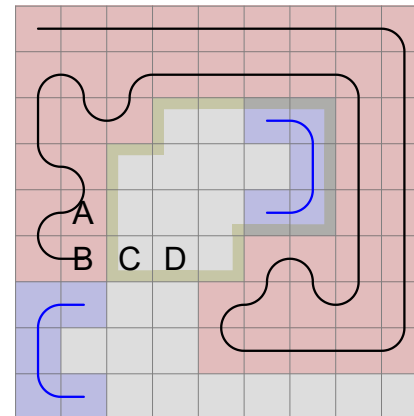
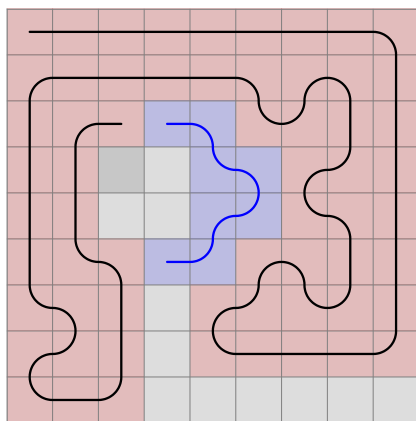
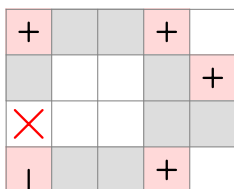
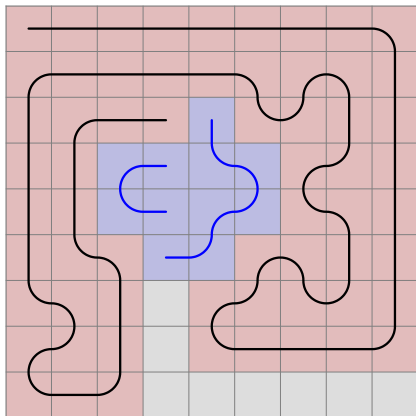
The diagram illustrates the process of learning a policy for a robot in a grid world. The left part shows a 4x4 grid with a red 'X' in the bottom-right cell and a black line indicating a path. The right part shows a 4x4 grid with a green square in the center, a blue square in the bottom-right cell, and a black line indicating a path.



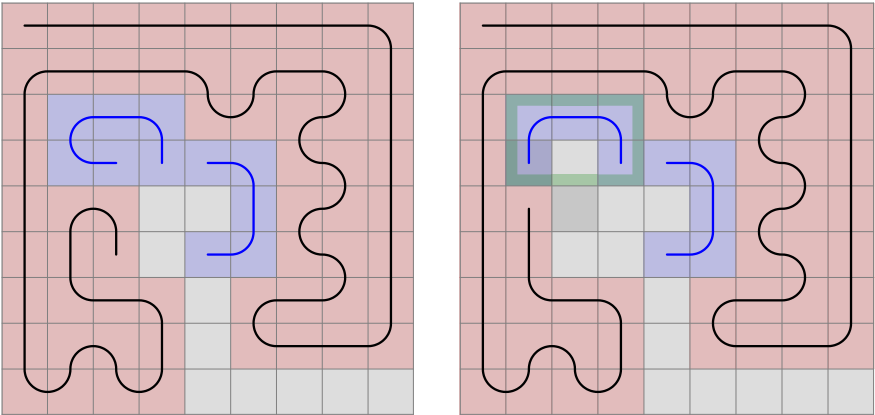
18 19

A 10x10 grid with a black path and blue obstacles. The path starts at (0,0), goes right to (9,0), then down to (9,9), then left to (0,9), and finally up to (0,0). There are blue obstacles at (0,1), (1,1), (2,1), (3,1), (4,1), (5,1), (6,1), (7,1), (8,1), (9,1), (0,2), (1,2), (2,2), (3,2), (4,2), (5,2), (6,2), (7,2), (8,2), (9,2), (0,3), (1,3), (2,3), (3,3), (4,3), (5,3), (6,3), (7,3), (8,3), (9,3), (0,4), (1,4), (2,4), (3,4), (4,4), (5,4), (6,4), (7,4), (8,4), (9,4), (0,5), (1,5), (2,5), (3,5), (4,5), (5,5), (6,5), (7,5), (8,5), (9,5), (0,6), (1,6), (2,6), (3,6), (4,6), (5,6), (6,6), (7,6), (8,6), (9,6), (0,7), (1,7), (2,7), (3,7), (4,7), (5,7), (6,7), (7,7), (8,7), (9,7), (0,8), (1,8), (2,8), (3,8), (4,8), (5,8), (6,8), (7,8), (8,8), (9,8), (0,9), (1,9), (2,9), (3,9), (4,9), (5,9), (6,9), (7,9), (8,9), (9,9). The path is black and the obstacles are blue.

As far as programming concerned, it just needed a rework of the universal rules, we didn't need to make completely new ones.

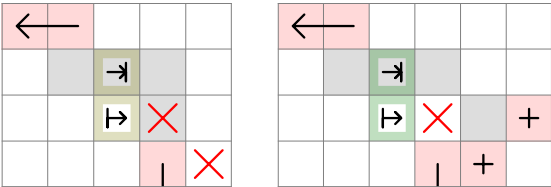
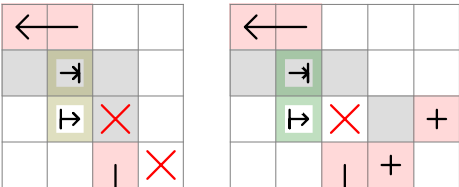
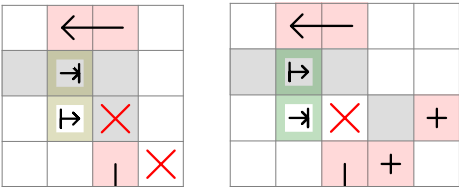


The same concept we encounter at 635 301, only the C-shape is created when we enter an area, on the other side of it.

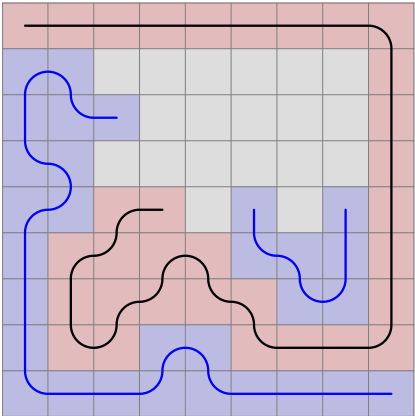


We have seen this in the third 9 x 9 rule. There the taken field next to the exit was in middle across position, and now it is across. And we also need to think about an obstacle straight ahead. Here are the original universal rules and their modifications.

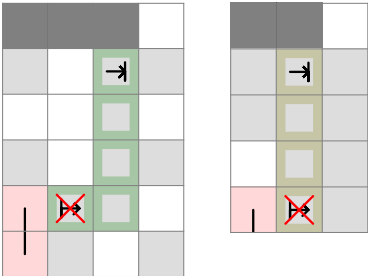
Straight, circle direction left:



At 349 215, we find this:



Though a double C-shape has been created in backwards direction, it indicates that the area on the right cannot be filled either. We have made a similar rule previously. Now we need to simplify it.

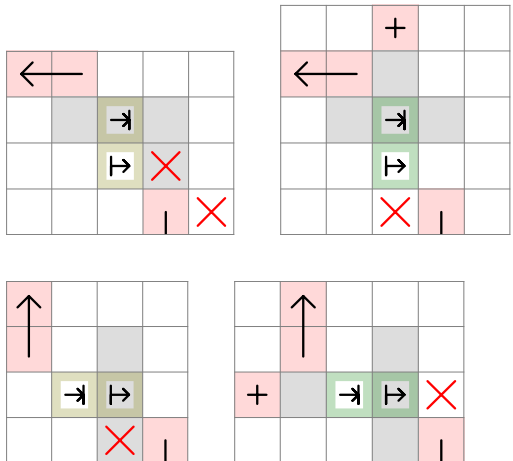


The area now has to be impair for the right direction to be forbidden. Essentially, we just added the three extra fields to the pair area.

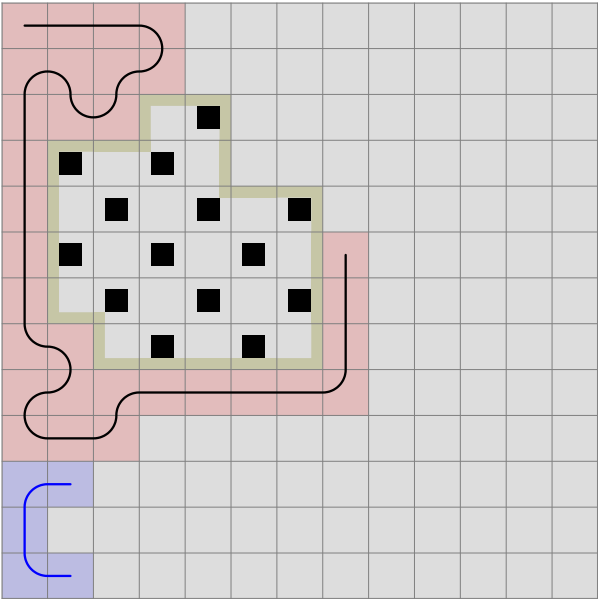
The diagram shows a 5x5 grid world environment. The grid is composed of 25 cells. A black start point 'A' is located at the bottom-left corner (row 5, column 1). A red goal point 'B' is located at the top-right corner (row 1, column 5). The grid contains several obstacles represented by black outlines: a large obstacle on the left side (rows 2-4, columns 1-3), a central obstacle (rows 2-4, columns 4-5), and a large obstacle on the right side (rows 2-4, columns 6-8). The grid is also divided into colored regions: blue regions (rows 2-4, columns 1-3; rows 2-4, columns 4-5; rows 5, columns 1-3) and red regions (rows 1, columns 4-5; rows 2-4, columns 6-8; rows 5, columns 4-5). A green path is shown starting from 'A' and ending at 'B'.

26 27

The universal far across rule have to be extended. By default, we disable the option to step straight or right if the counted area is impair. When it is pair, we need to disable the left field.

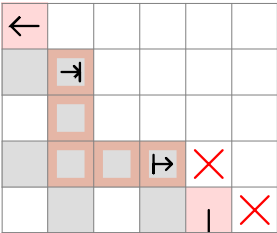


Let's mark the original example as a checkerboard.

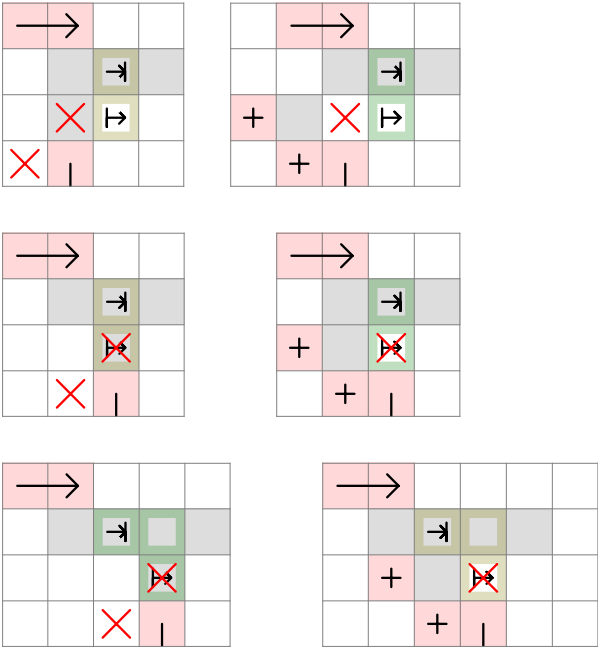


We enter at a black field and exit at black too, so the number of black fields should be one more than the number of white fields. Here there are 14 black fields and 15 white. That's why the area cannot be filled. The up and right directions need to be disabled, so we can only step left.

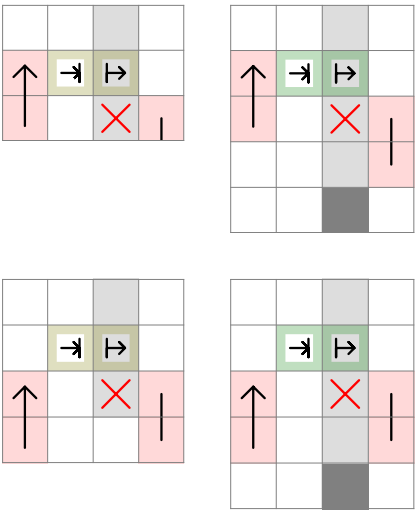
This is the rule representation. The reddish arealine now means the arealine is impair, and we know that the entry and exit points are the arealine start and end fields.

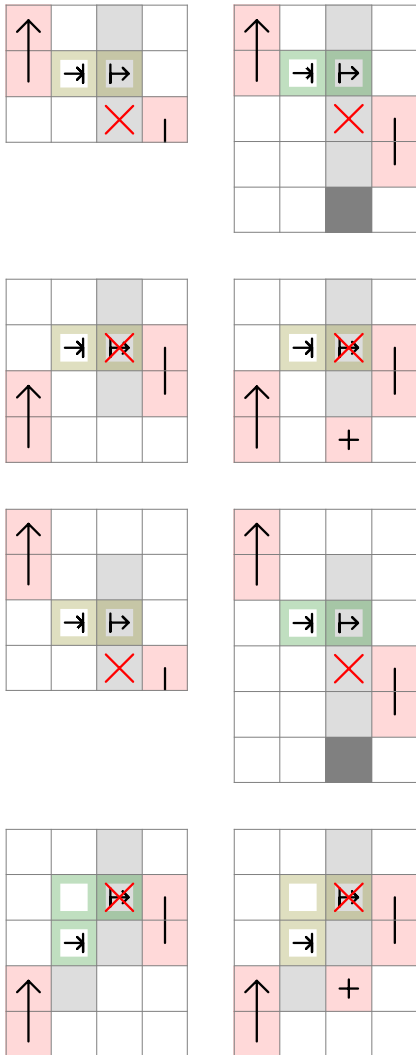


Circle direction right:



Side, with taken fields above and below:



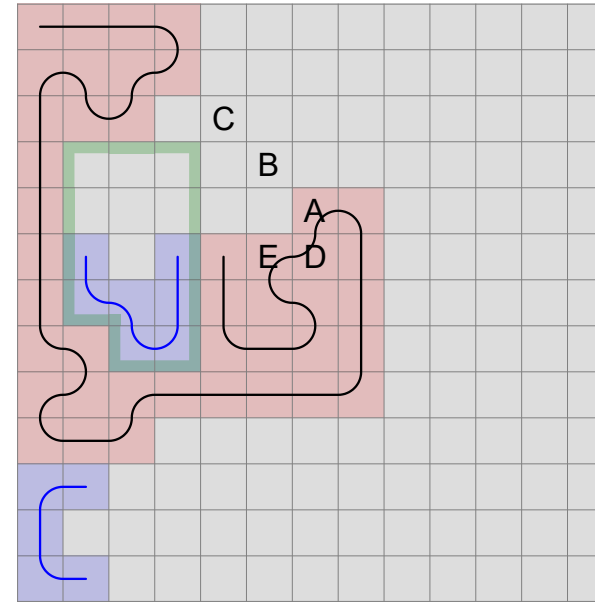


I have made some changes by adding some empty fields in side positions, so they are the same as the straight rules, just rotated.

Also, I have added the side across down rule and changed the straight across rule accordingly. Not only fields next to each other can define an area, they can be across too. In that case, if the area originally was marked impair, now it has to be pair.

Notice that in side rules, when the taken field that would create the C-shape is below the obstacle creating the area, it can be a border field too. We have seen an example of that previously.

Now what if both the start and end C-conditions are true? We can construct this on 13 x 13:



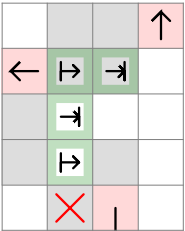
Several walkthrough attempts will leave you thinking why you cannot fill the area once obstacle responsible for the start C-shape is created (A). The area enclosed by A, B and C is pair. So when you enter it at A or B (obviously C is not a possibility), in order to exit at C, you need to leave out an impair amount of fields from the area. In case of entering at B, you cannot leave out A, but when you enter at A, you can leave out B, and no more than that. Now the area will be impair.

The minimal area would be stepping left from A, left again, up and up to get to C. You have covered 5 fields.

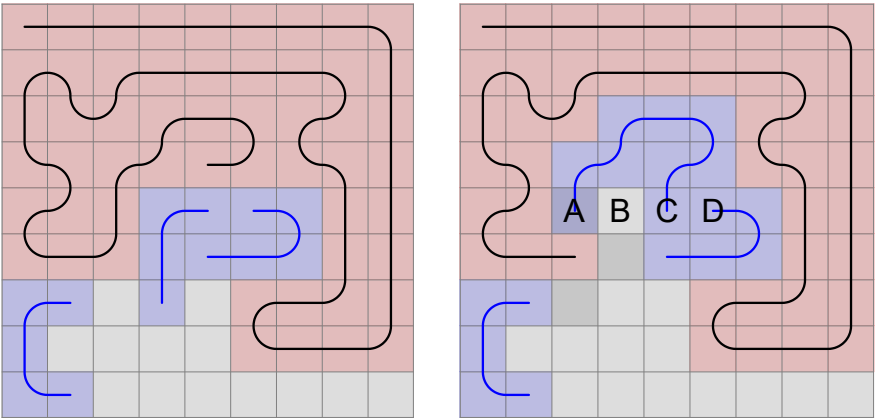
In order to make a walkthroughable area, you would need to extend it by pairs of fields next to each other, like D and E. One will be filled at a pair amount of steps, the other at an impair amount.



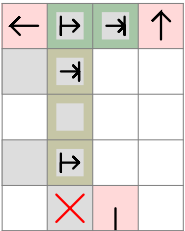
The rule will now look like this:



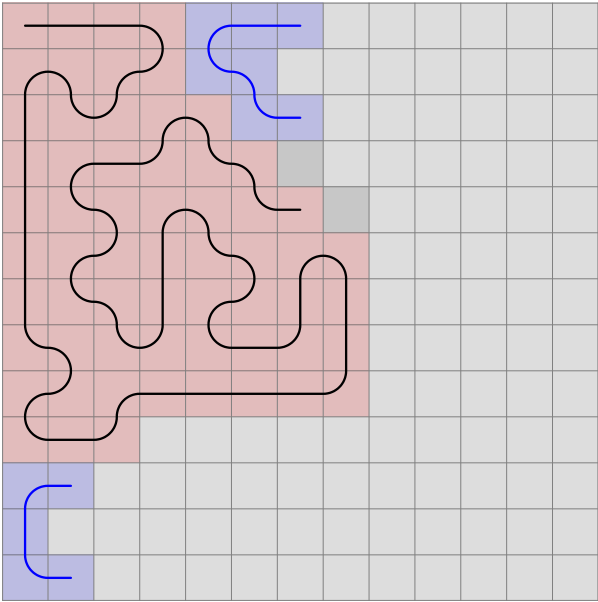
740 039 is a slight modification.



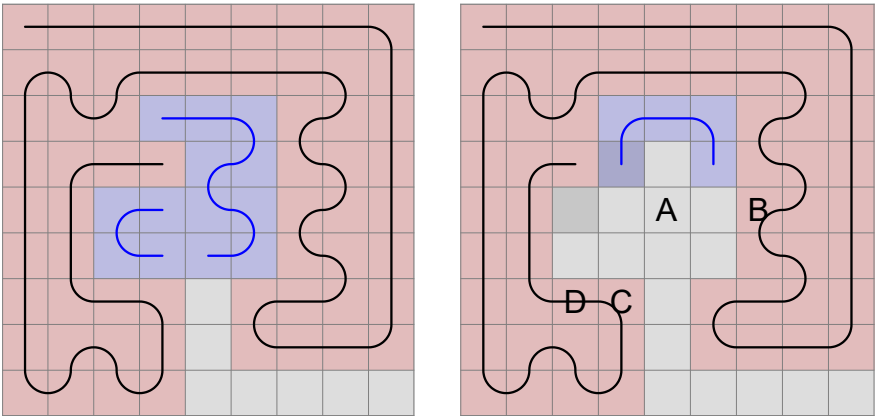
The only difference is, that the obstacle is 3-distance away. With the area being impair, if we enter at A, we must exit at C. What if we omit D from the area? Then the area will be pair, so we must exit at B, and the only way to get there is from C. And if D is included, we can only step to C from there. Either way, we step away from the area beyond D, so the rule will be:



And now the walkthrough is possible.

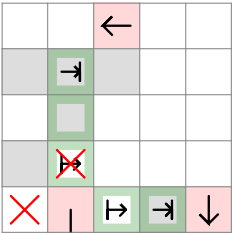


Continuing the 9 x 9 program, we get this at 641 019:

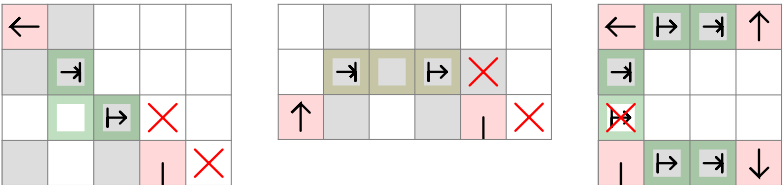


If you enter the pair area straight ahead, you will exit at A, and you need to turn towards B because of the C-shape. Now you cannot go in and out of the area enclosed by C, and the situation would be the same if that obstacle was in D.

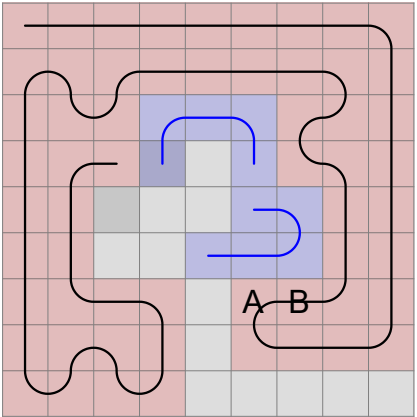
To mark the two areas, each one has to be given a directional obstacle next to the count area end field. In this case, it represents a taken field, but we don't go wrong if we include the border as well.



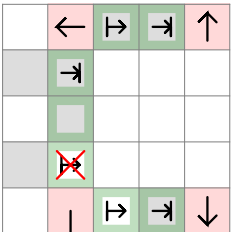
And with this marking system, we can correct the rules previously made. All rules featuring future line start and end fields have to be rewritten to start with.  
So we get the 2-distance across rule, the straight 3-distance rule to prevent a double C-shape, and the square constellation with 3 areas. All of them are rotated clockwise or counter-clockwise.



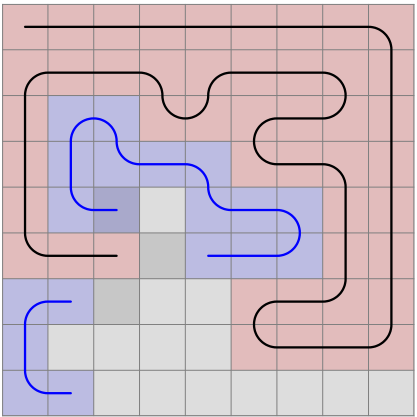
Let's return to the last example and make a modification:



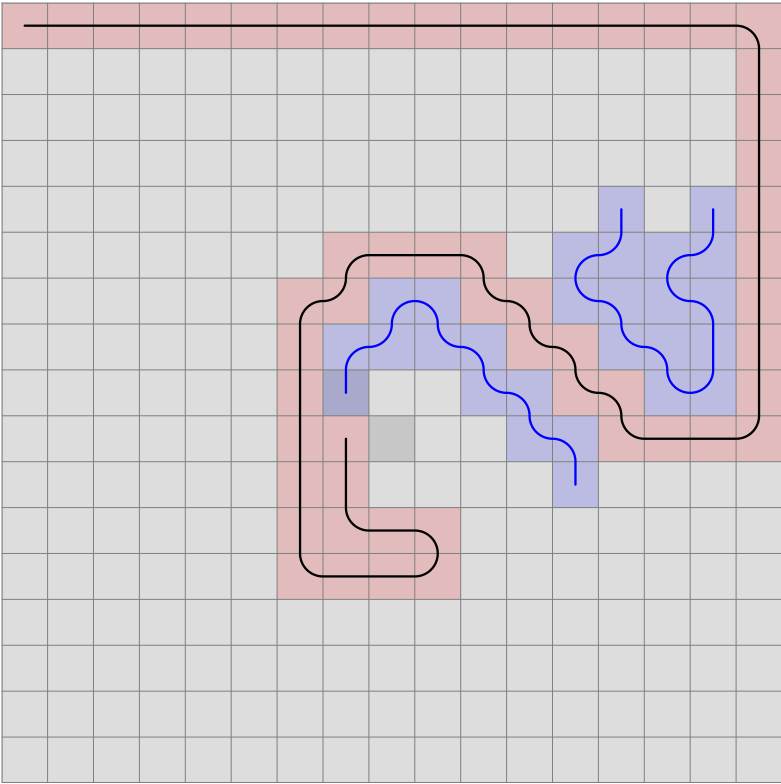
The field previously marked with B is now empty. But we still need to step in that direction, due to the area enclosed by A, which obstacle could as well be in B.  
The rule will be now symmetrical. It is similar to the square obstacle pattern.



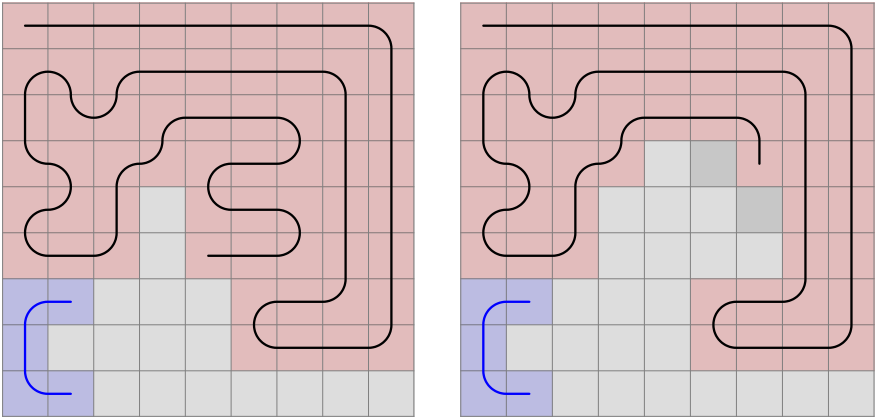
The same concept we encounter at 725 325. We have seen this previously, just with C-shape, not an area.



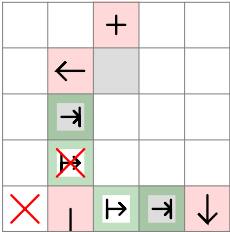
On 17 x 17, we can construct a situation where the obstacle across the stair is 2 behind and 2 to right. As the table size increases, the stair-obstacle narrowing can move infinite distance away from the live end. That's why it is important to group these rules as one.



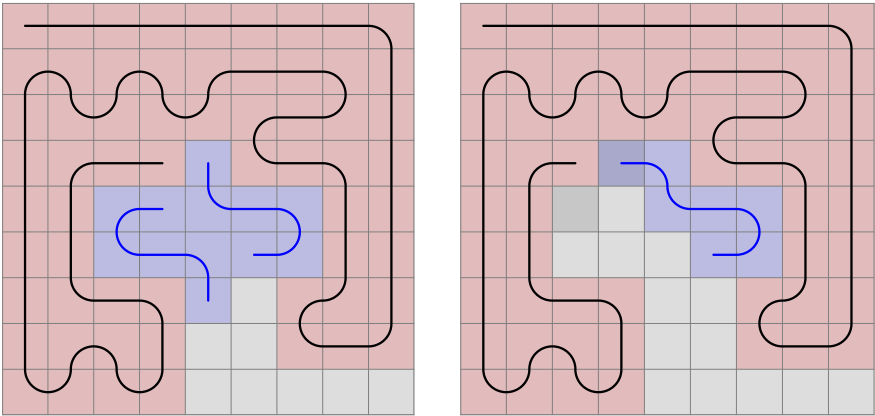
811 808:



Recognize it is a variation of the square obstacle pattern where instead of an area, there is a C-shape at the rule's upper edge.

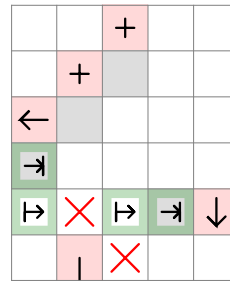


1 261 580:

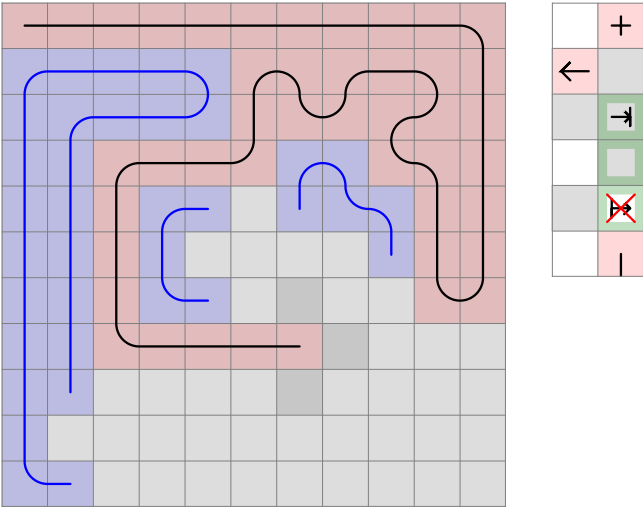


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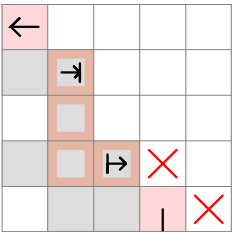
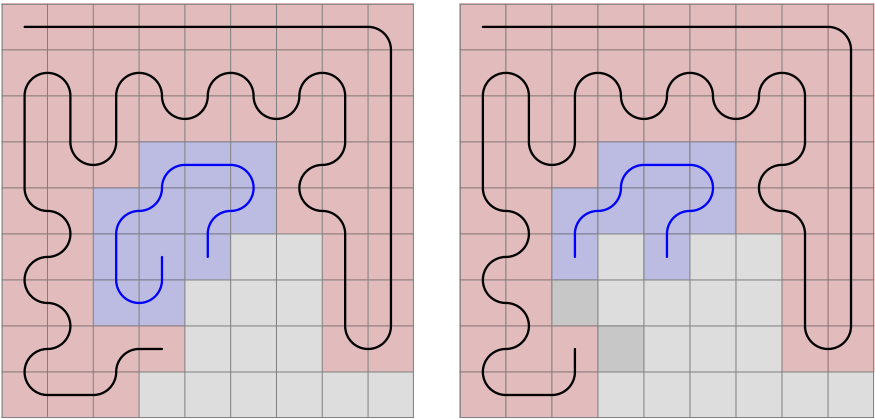
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| ✗ |   |   |   |   | + |
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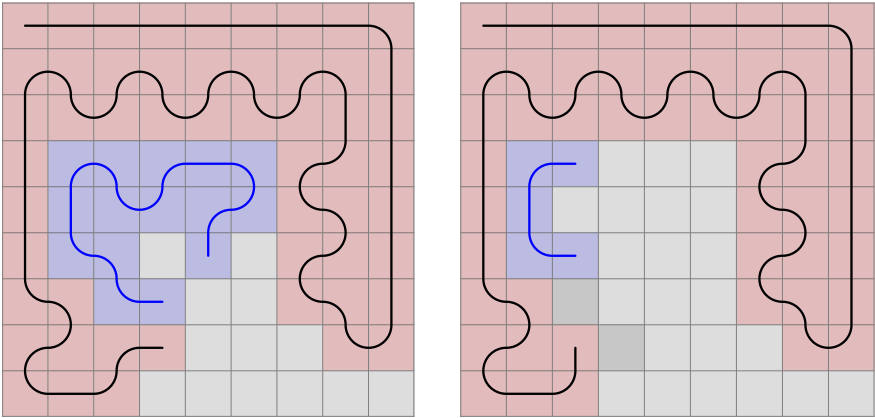
From our experience, the area can be substituted with C-shape.



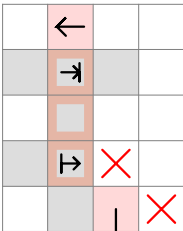
3 225 432 is a variation of the impair area imbalance rules we have seen before.



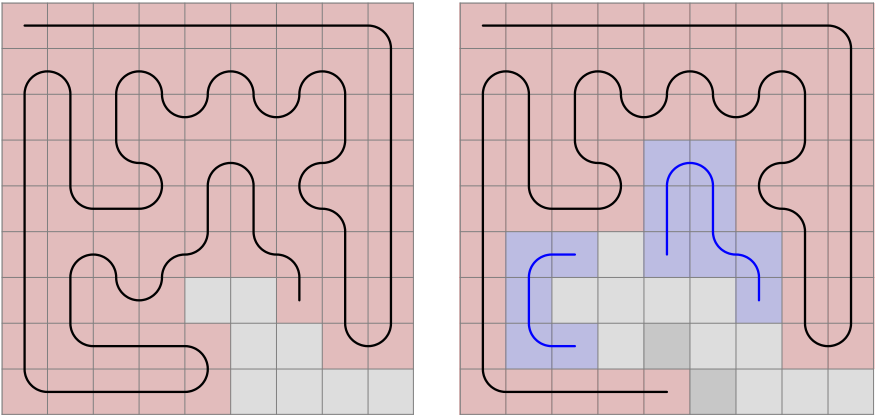
We have all the tools to handle 2 034 575.



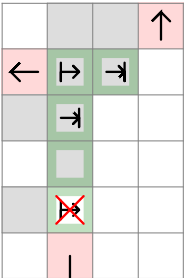
It is an impair area where the number of the starting field's color is less than the other color.



Next stop is at 3 224 847.

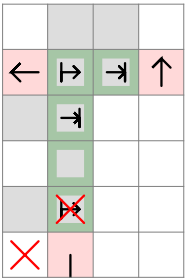


A pair area is created with the obstacle 3 distance away, so if we step into it, we will exit at the middle, but because of an area, we cannot step there.

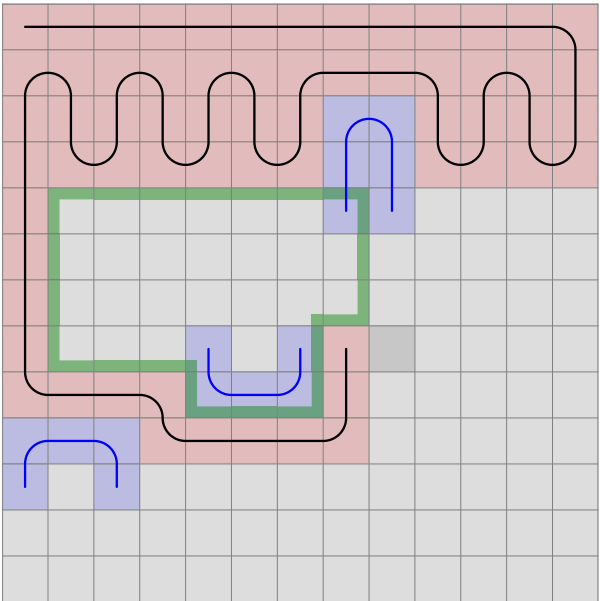


Beware of disabling the left field. If the count area end field is excluded from the area, the area will be impair, thus we will exit at the count area start field, coming from the middle.

But if the obstacle in the upper right corner is moved down, even this will be impossible.

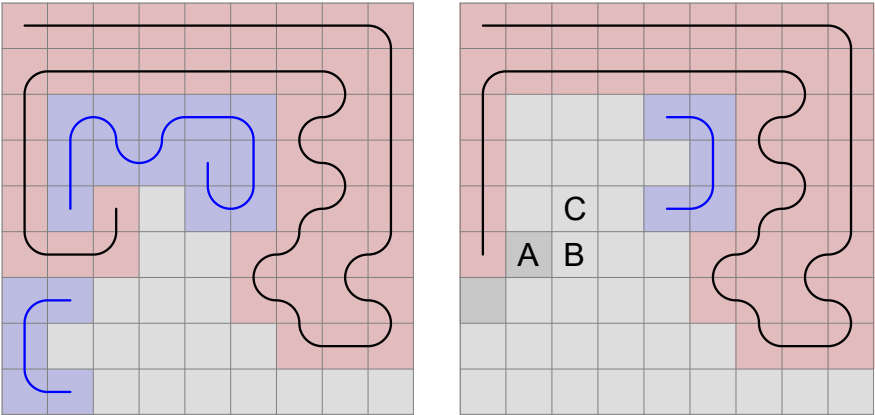


We can recreate this example on 13 x 13.

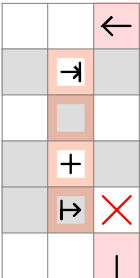




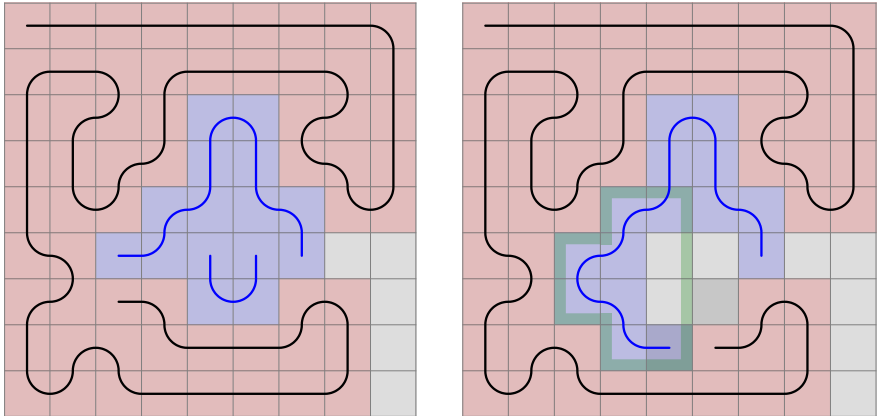
Soon we get a similar situation, only here it is the imbalance of pair and impair fields that is to blame.



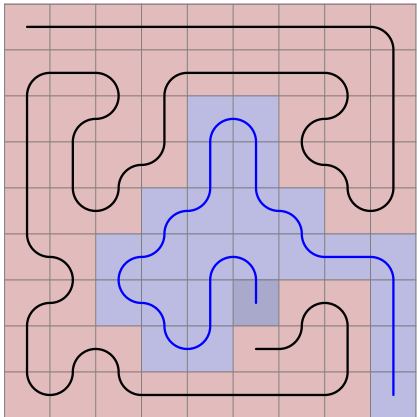
If we step to A, we cannot step left and therefore must continue to B (or right). From B, the only possibility is C, but the 5 x 3 area is not just impair, there is less of the C-parity field than the other.  
In the rule, I introduced a new field that indicates the entry point; this has always been the start field until now.



At 19 717 655 the program stops.



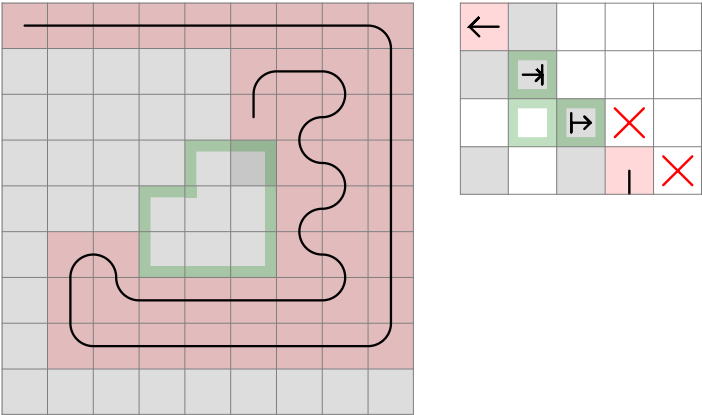
Obviously, we cannot step straight, but had we extended the future line until the end corner, the situation would not have occurred and we would have just got this:



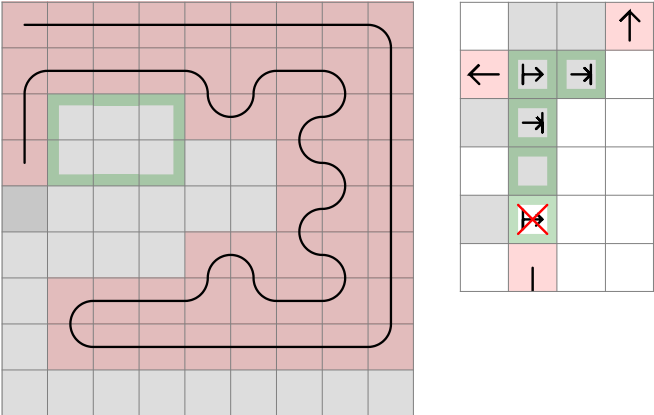
Though the algorithm including the reliance on the future lines is just as solvable, we miss patterns and therefore narrow the spectrum of the discoverable rules. We would eventually discover the patterns as we increase the table, but why not gain the most out of the 9 x 9 study? From now on, future lines are treated as a visible aid, but they do not play a role in deciding which fields are available for the next move. When a possible field is within the body of a future line, the program should stop.



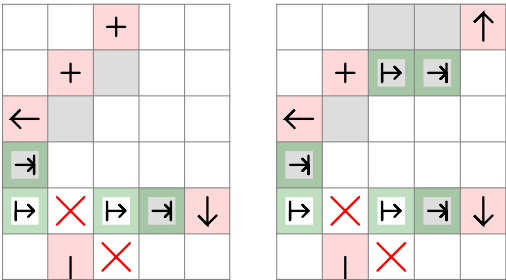
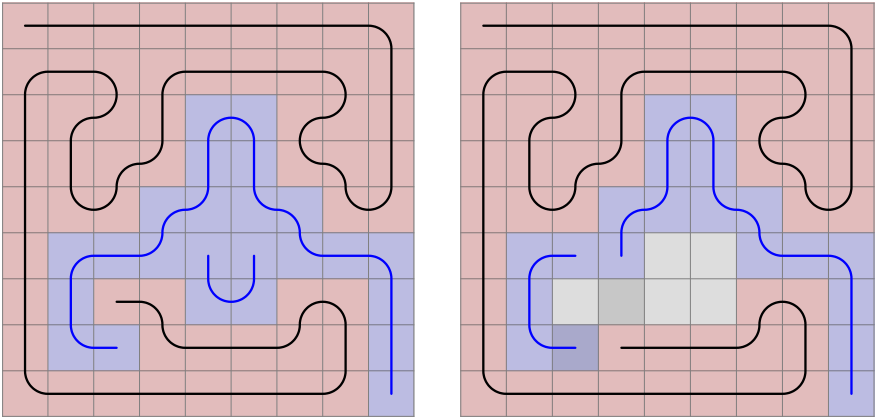
9 121, Count Area 2 Across



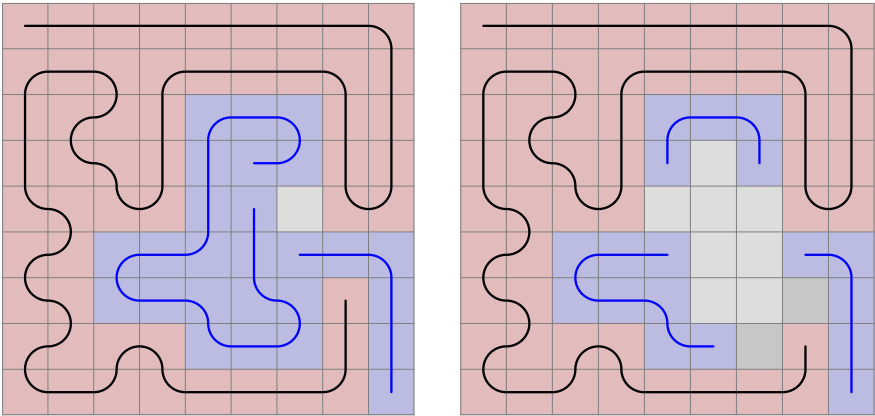
22 328, Straight Mid Across 3 End Area



19 718 148 is a slight modification of 2 022 773 where there is an area instead of a C-shape straight ahead.



We encounter a new constellation of 3 areas in 23 310 321 where the exit is next to the live end.

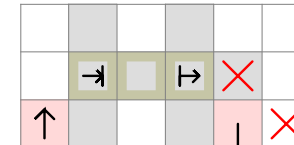
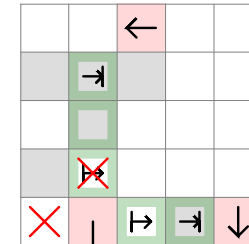


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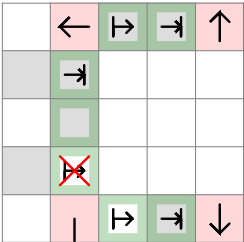
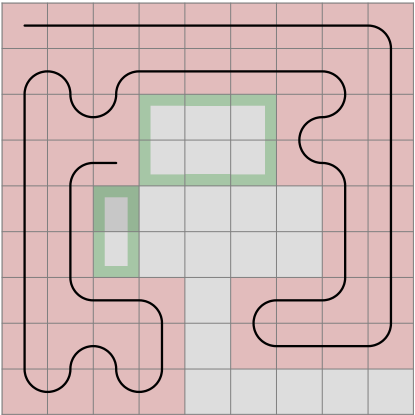
In the following section I list the 9 x 9 rules in chronological order. The patterns are not introduced when they are first recognized, but when they are first needed, meaning that they disable fields that the other rules don't. And the disabled fields have to be empty.

Still, the number of completed walkthroughs before the appearance of the rule may not be the same as the number of those before getting stuck in the lack of that rule. If the rule disables a field right to a possible field, the left branch would run through first.

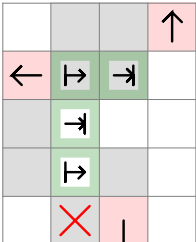
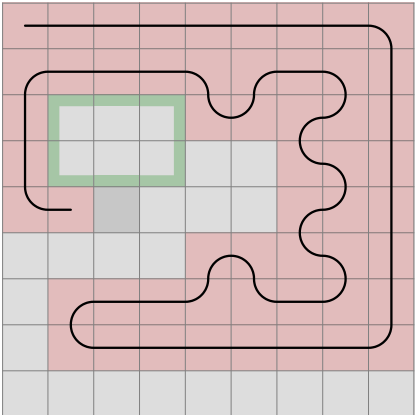
The diagram illustrates a 10x10 grid with a red-shaded region on the left side, representing a 'red region'. A black outline of a complex shape is drawn on the grid. Two green rectangles are placed within the grid, one of which is partially inside the red region. This setup is used to demonstrate the relationship between the red region and the black outline.



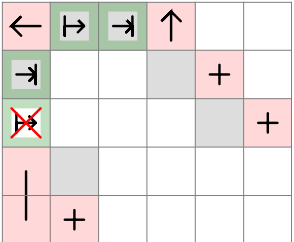
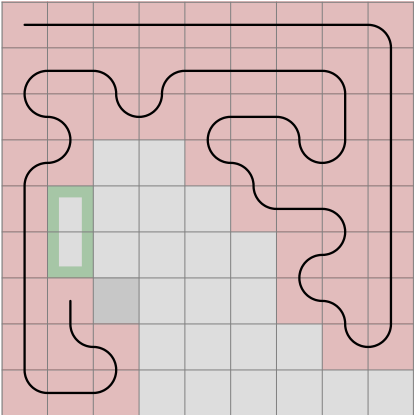
740 363, Triple Area



22 328, Straight Across End Area



2 022 763, Double Area Stair

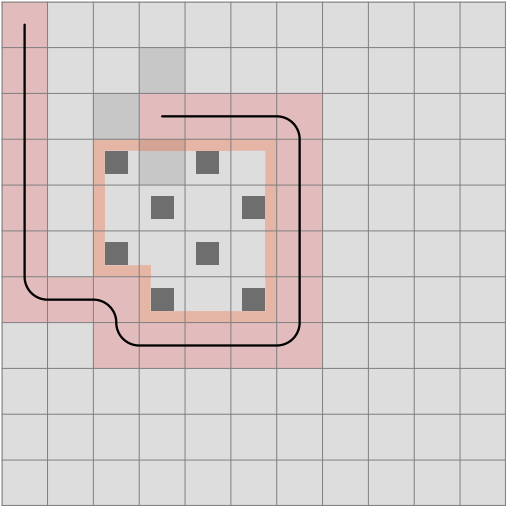






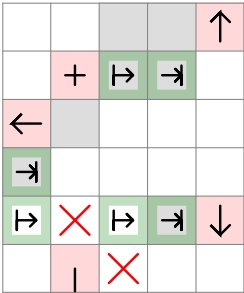
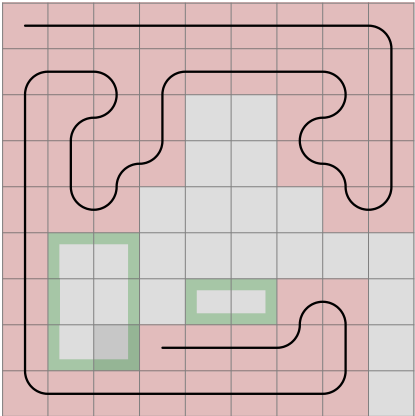


Here, the area is still pair, but there are 12 black fields and 10 white. To fill it, two lines of an impair length would be needed, each starting and finishing on a black field. Now, it is possible to enter and exit the black field in the upper left corner of the area, but we cannot do it with the field 2 below. And there are no more black fields on the boundary of the area. Filling it is therefore impossible.

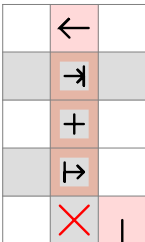
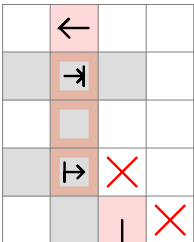
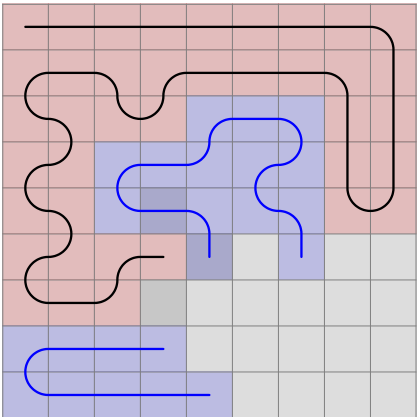
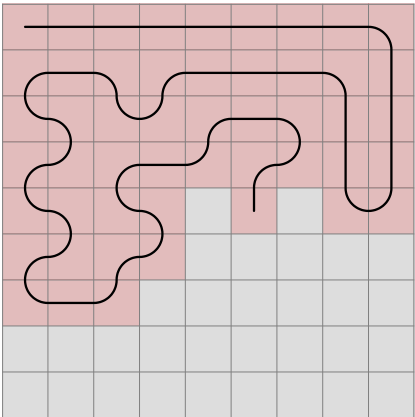


When there are one more black field than white (8 and 7 in this case), the area is impair. If we enter now, we can exit at the count area end (to make a line of pair length), and the corner black can be filled later. We can also enter at the corner black later to exit at the count area end.

19 720 614, Double Area Stair Area



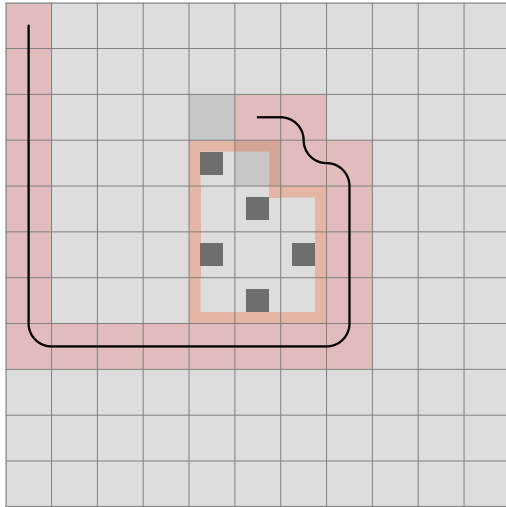
23 350 320 is new, but it shows similarity to the Mid Across 3 Impair Determined rule. As the double C-shape reveals, it is about pair/impair field imbalance.



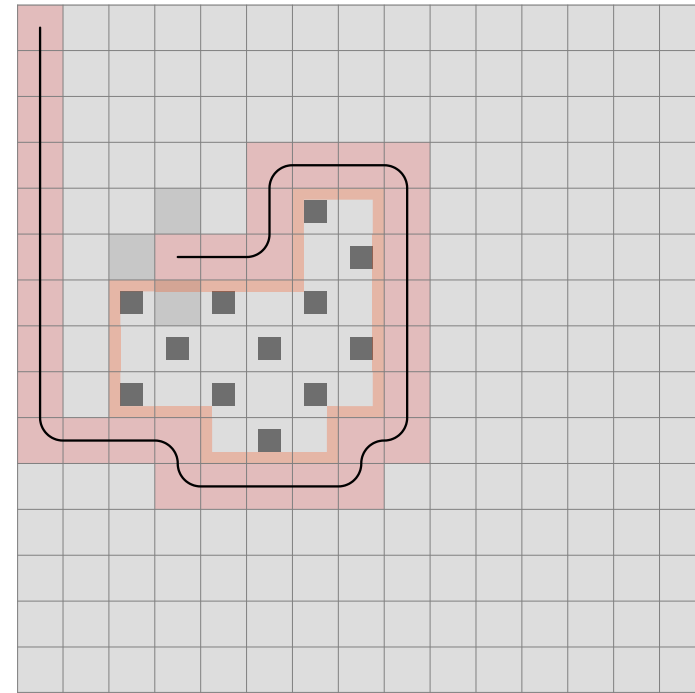




If  $\text{black} = \text{white} - 1$ :



A line that starts and ends on white can be drawn, no matter if we enter now or later. If we enter now, the next field has to be the corner black, and then there will be a line between the two white fields on the boundary.



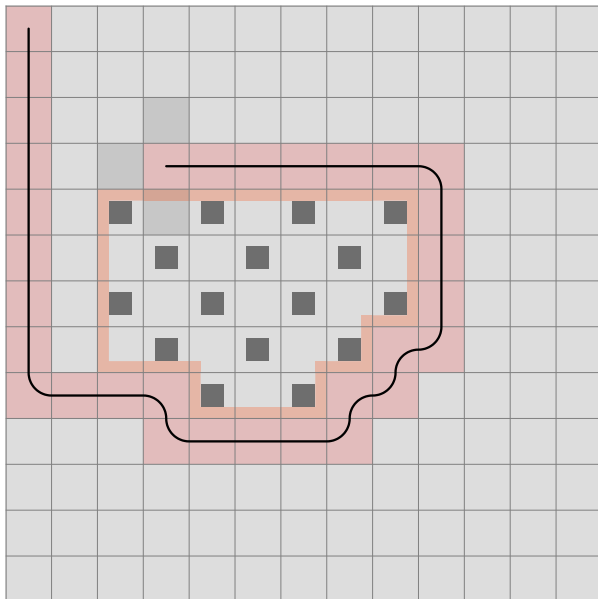
Here, there are one more white field than black. To fill the area, 1 line of impair length would be needed that starts and ends on white, but if the enter the area now, the corner black could not have been filled. Now there are 2 less black fields available for the rest of the area, but making two lines starting and ending on white is impossible, there are not that many white fields on the boundary.

Needless to say that we cannot enter later either, there is only one white field on the boundary.

If there are even fewer black fields relative to the white, the situation is the same.

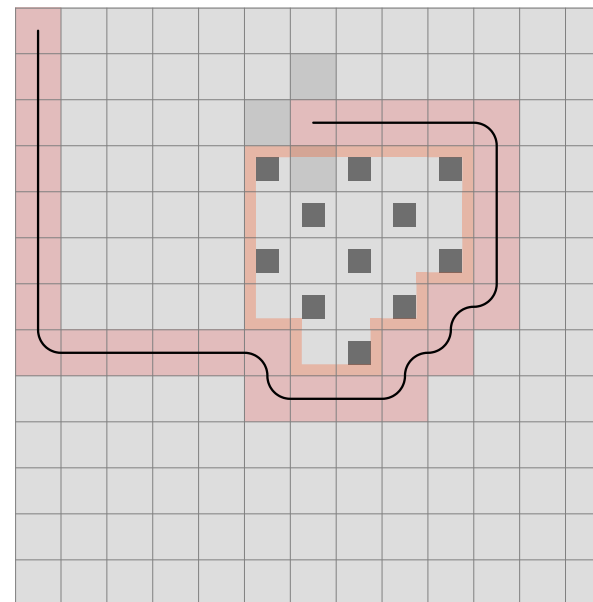
At 4 distance (and in any case), we still don't have a problem filling an area that has equal number of black and white fields.

When the number of black fields are two more than the whites (16 and 14):



Two black to black lines would be needed to fill the area. Apart from the black in the corner, there is only one black field on the boundary.

If black = white + 1:



If we enter now by stepping left, there will have to be one more line even if we exit on black. That line has to go from black to black, so it can only be the corner field. Have we exited at the other black field (the third on the boundary), either the second or the fourth could not have been filled.

**This direction therefore has to be disabled.**

We can enter later without problems to start at the corner field, then the second, go inside the area and end at the third.

## 2) Pair distance, impair black and white: 6, 10 etc.

If we enter now and exit at black,  $(B-1) / 2$  black to black lines can be drawn.

If we exit at white,  $(W+1) / 2$  white to white lines can be drawn. (We have to move to the corner black and exit at the first white during the first line)

If we enter later,  $(B+1) / 2$  and  $(W-1) / 2$  black and white lines are possible, respectively.

The rule is double:

- If the number of black fields is greater or equal than the number of whites plus  $(B+1) / 2$ , we cannot enter now.
- If the number of white fields is greater or equal than the number of blacks plus  $(W+1) / 2$ , we have to enter now.

## 3) Impair distance, impair black and pair white: 5, 9 etc.

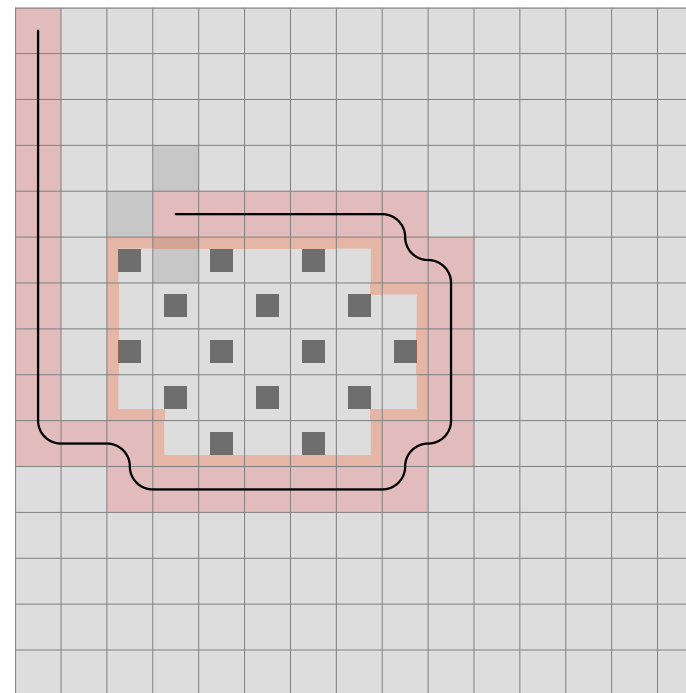
If we enter now and exit at black,  $(B-1) / 2$  black lines are possible.

If we exit at white,  $W/2$  white lines are possible.

If we enter later, the number of black lines can be up to  $(B+1)/2$ , the whites  $W/2$

Single rule: When the difference is at least  $(B+1)/2$ , we cannot enter now.

If black = white - 2:



Two white to white lines would be needed, but there are only 3 white fields on the boundary, and none is the corner.

Without finding concrete examples, at 5 distance I only draw the boundary and go through the different possibilities.



## black = white + 2

If we enter now and finish at black, only two black fields remain. Drawing two black to black lines is not possible, so this direction has to be disabled.

We can enter later, go through the corner and draw another black to black line.

$$\text{black} = \text{white} + 3$$

There is not enough black fields for three black to black lines.

$$\text{black} = \text{white} + 1$$

We can enter now, finish on a black field and draw another black to black line.  
Or enter later.

$$\text{black} = \text{white} - 1$$

A white to white line is possible in either case.

$$\text{black} = \text{white} - 2$$

We cannot enter now, because even if we end on a white field, only one white remains, and that is not the corner. And cannot enter later either.

The same procedure applies at 6 distance.



### **black = white + 2**

The number of black fields is the same as previously, so entering now is not possible.

black = white + 3 is impossible.

black = white + 1 and black = white - 1 is possible.

### **black = white - 2**

Entering now is possible if we step upwards and exit at the neighbouring white field. Two white fields remains.

But we cannot enter later and make two white to white lines using three white fields.

From now on, we can distinguish between four cases:

**1) Pair distance, pair black and white** (indicated by B and W): 4, 8, 12 etc. distance

If we enter now and exit on black, B-1 black field remains on the border. B-1 is impair, and the corner alone can make a black to black line. Hence  $(B-2) / 2 + 1 = B / 2$  black to black line would be possible, but...

- if the first line finishes at the corner black, the opportunity for it to make a single line is lost, and the remaining B-1 black fields cannot make  $B / 2$  black lines.

- if it finishes on any other black field, on one side there will be a section that has white fields on both ends (below the greenish fields were taken by the first line)



A black to black line cannot have three white fields on the border (unless more black fields were used up).

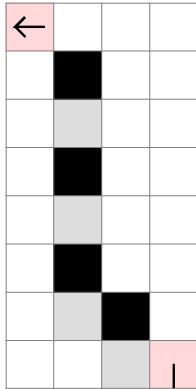
So there can be at most  $B / 2 - 1$  black to black lines.

If we enter now and exit on white, W-1 white fields remain, so  $1 + (W-2) / 2 = W / 2$  white to white lines is possible.

If we enter later, the number of possible black to black lines is at most  $B / 2$ , and the white ones  $W / 2$ .

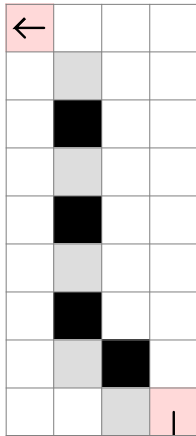
Our rule will look like this: If the number of black fields in the area is greater or equal than the number of white fields plus  $B / 2$ , we cannot enter now.

6 distance:



Now: 1W -> 2B  
 Later: 1W -> 3B  
 3B: cannot enter now

7 distance:



Now: 2W -> 2B  
 Later: 2W -> 2B  
 No rule.

#### 4) Impair distance, pair black and impair white: 3, 7 etc.

If we enter now and exit at black,  $B / 2$  black lines are possible.

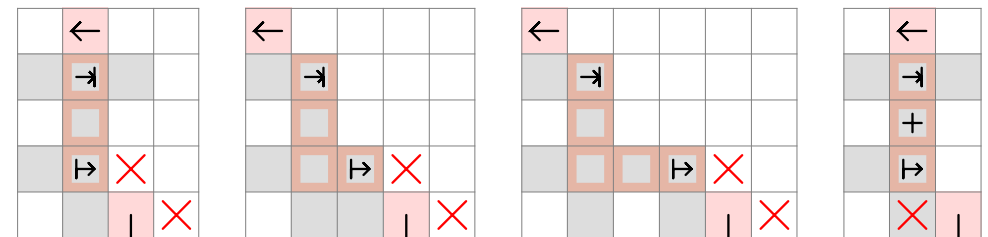
If we exit at white: Similarly to the first of the four cases, a black to black edge will remain on one side. Drawing  $(W-1) / 2$  more white lines is either not possible, or if we do so, the corner black may make up a black to black line, decreasing the difference.

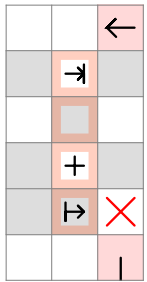


When entering later,  $B / 2$  and  $(W-1) / 2$  are the numbers. Since they match the above, no rule is applied.

Check the original 21 x 21 example. Two steps back, there will be 9 distance with the wall to the left. The number of black fields on the edge is 5, therefore there cannot be 3 more black fields in the area than white, but counting them, they are 51 and 48.

Does this procedure apply to any of the size-specific rules? Not exactly, but let's look them through. Here are all of them that deal with black and white field imbalance:

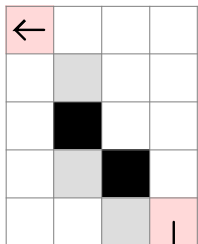




As a reminder, they indicate impair areas, and in the first two case if the count area start field (black) type is not 1 field more in the area, we cannot enter later.

In the fourth, the field marked with + (white) is the type that needs to be 1 more than black, otherwise stepping forward is forbidden.

I will take the second rule under examination as it contains both a horizontal and vertical offset.



If we enter now:

- We can exit at the end white, so 1 white line is possible.
- We can exit at the black farthest away and then make a line using the black field closest. 1 black line is possible.

I will mark it like this: 1W -> 1B

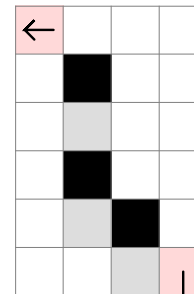
If we enter later:

- Having two black fields, 1 black line is possible.
- There is only one white field. It sits in a corner, but the two black fields will give 1 black line. 0 white line is possible.

0W -> 1B

- So if there are 1 more white fields in the area than black (1W), we cannot enter later.

- Now let's increase the vertical distance.



If we enter now, 1 white line to 2 black lines are possible. There are 2 black fields on a corner.

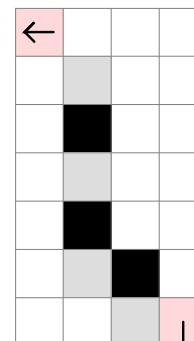
1W -> 2B

Later, 0 white line and 2 black lines can be drawn.

- 0W -> 2B

- Conclusion: in case of 1W, we cannot enter later.

5 distance:

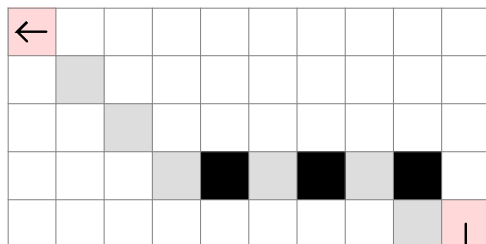


Now: 2W -> 1B

Later: 1W -> 2B

- 2W: cannot enter later
- 2B: cannot enter now

But we are not finished, we still need to examine the distance of 8.



Now: 4W -> 1B

After entry, we need to move up to fill the corner black and exit at the first white field.

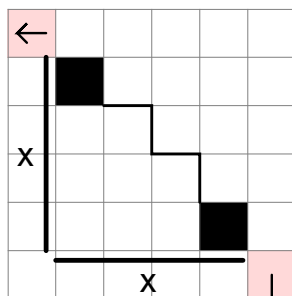
At 4 distance, only 2W was possible, but now we can exit at the first white and fill the area when entering at the second and exiting at the third.

Later: 3W -> 1B

After this practice, let's calculate how many white and black lines we can draw when we have an obstacle x and y distance away.

There are three cases to look at.

1. Equal horizontal and vertical distance



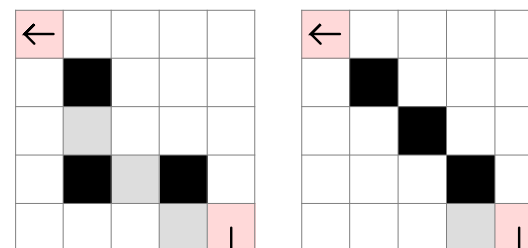
There are an x number of black fields on the area boundary.

In case of 7, there are 4 fields added to the 3-distance example. We would expect that in case of 2W, we cannot enter later, but now, 2W is possible even when entering later, because one line will be the corner white, and the other can go between the other two white fields, taking up all the blacks along the way.

From now on, increasing the numbers by 1 for every 4 distance increase will work.

The next thing to do is the horizontal increase.

3 distance:



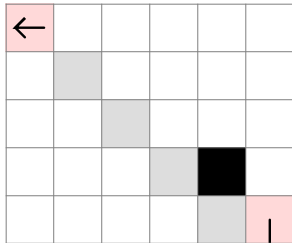
The picture on the left is the representation we have used so far. However, we cannot exit at the black in the middle, and when we exit at one of the whites, the other is only accessible for immediate entry. Therefore, I will add the extra field. If there is a taken field anywhere ahead acting as the obstacle, we can get to it by drawing a straight line and a stair.

Now: 0W -> 2B

Later: 1B -> 2B

If we used all 3 black corners for a separate line, we would not enter the area.

4 distance:

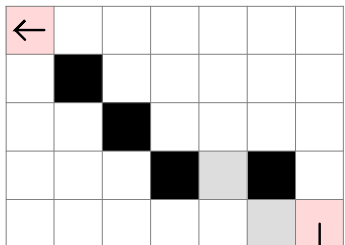


Now: 2W -> 0B

If we exited at the nearest white after entering, we have either not entered the area or not filled the black field.

Later: 2W -> 0B

5 distance:

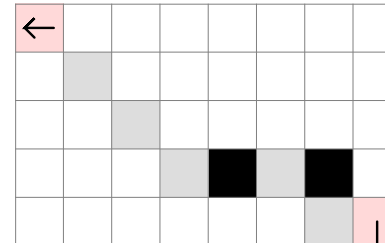


Now: 0W -> 3B

If we reserved the 3 black corners, the only way to end the first line on black is to move downwards after entry.

Later: 0W -> 3B

6 distance:

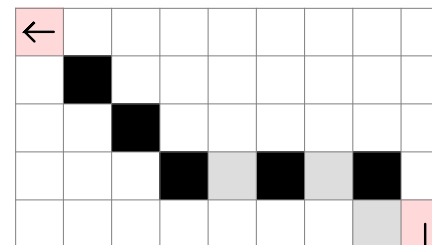


Now: 3W -> 1B

Similarly, we need to move down after entry in order to finish at the second black field, leaving the first for itself.

Later: 3W -> 1B

7 distance:



Now: 1W -> 3B

Later: 0W -> 4B

There cannot be one white line, because out of the first three black fields only 2 would be filled.

Notice that as we added 4 distance to 3, now both the white and the black end of the ranges have increased by one.



Later: Because of three inline white fields, one white line is possible. For blacks, we can use all the corners and two out of the three inline blacks.  
 $1W \rightarrow (x+1)B$

We don't need more examples.

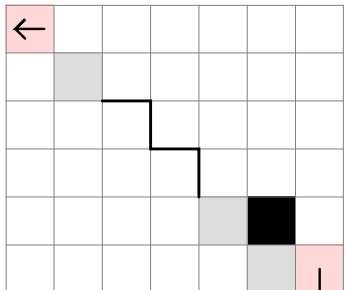
For  $2n$  added fields, there will be  $n$  amount of inline white fields.

If we enter now,  $(n+1 - (n+1) \% 2) / 2$  white lines can be drawn if  $n > 1$ .  
 The black line count is all the corners minus one for finishing the first line, plus one for each remaining pair.  
 $x + (n-1 - (n-1) \% 2) / 2$

For entering later, the white line count is  $(n - n \% 2) / 2$ , and the black line count is  $x + (n - n \% 2) / 2$ .

**If the added distance is impair:**

$2n + 1 = 1, n = 0$ :



Now: If we end the first line at the first white, we could not have filled the corner black and the area simultaneously.  $xW$  is therefore not possible, but  $(x-1)W$  is.

And since there is only one black field, the black line count will be 0.

$(x-1)W \rightarrow 0B$

$x = 1$ :

Now:  $0W \rightarrow 0B$

We cannot enter later.

$x = 2$ :

Now:  $0W \rightarrow 1B$

Later:  $1B$

$x = 3$ :

Now:  $0W \rightarrow 2B$

Later:  $1B \rightarrow 2B$

$x = 4$ :

Now:  $0W \rightarrow 3B$

Later:  $1B \rightarrow 3B$

$x = n$ :

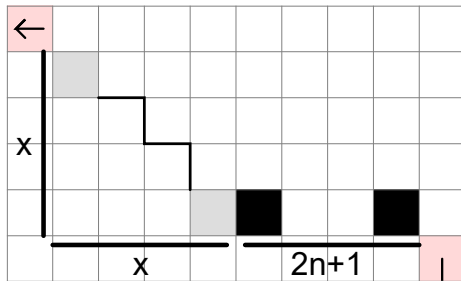
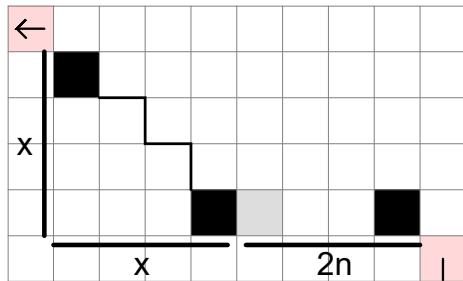
Now:  $0W \rightarrow (n-1)B$

Later:  $1B \rightarrow (n-1)B$

There is an  $x$  amount of corner blacks, but we need to enter the area as well.

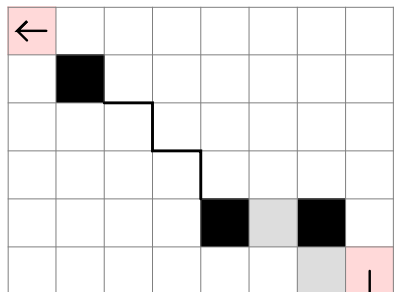
Conclusion: if the white and black fields in the area are equal, we cannot enter later.

## 2. Larger horizontal distance



**If the added distance is pair:**

$2n = 2$ :



Now: Can we make 1W? No, because that would require going through the first black field to end at the first white, while filling the other parts of the area too.

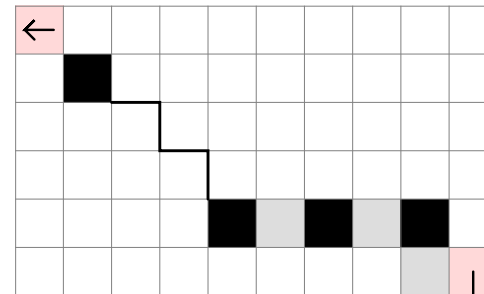
What about the black line count? Without the horizontal addition,  $(x-1)B$  was possible. Can we now make  $x$  amount? Yes, by stepping down and ending the line at the first white field and the second black, having filled everything except the corner blacks.

$0W \rightarrow xB$

Later: We can use all corner blacks, and the line starting at the first white and ending at the second black will fill the area.

$0W \rightarrow xB$

$2n = 4$ :



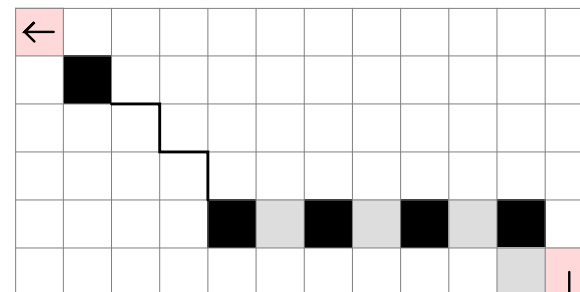
Now: 1W is now possible by ending at the second white, and everything can be filled on the way, while  $(x+1)B$  is not possible. Two inline blacks remain after taking off all the corner blacks, and we have to use one of them to exit the first line.

$1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

The line connecting the two inline blacks can fill the rest of the area.

$2n = 6$ :

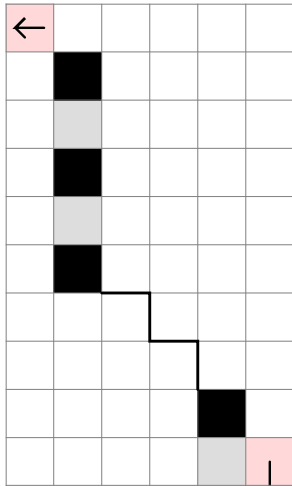


Now: If the first line goes through the first black and exits at the first white, the second line can go between the remaining two whites.

As far as the blacks concerned, if we exit at the second black, we can use the third and fourth for a line, plus the  $x$  amount of corners.

$2W \rightarrow (x+1)B$

$2n = 4$ :



Now: There cannot be more than 1W. And the black line count is unchanged too. One of the two inline blacks will be used for completing the first line, and one remains plus x amount of corner.

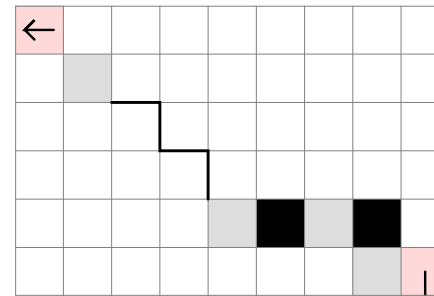
$1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

Later: All the corner whites plus the neutral line makes  $(x-1)W$ . The black line count is still 0. The black field is a corner, but it will be counterbalanced by at least one white to white line.

$(x-1)W \rightarrow 0B$

$2n + 1 = 3, n = 1$ :



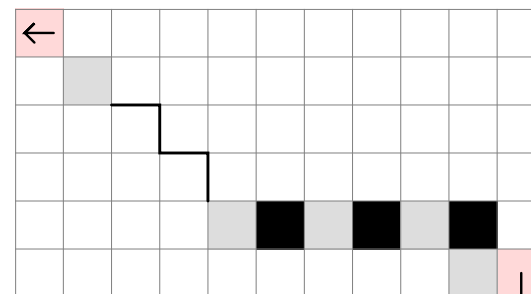
Now: Aside the corner whites, 1W is possible by ending at the second white, just like in the  $2n = 4$  case.

The black line count is now 1, but we need to step downwards and finish at the first white and second black in order to have the corner black available.

$xW \rightarrow 1B$

Later:  $xW \rightarrow 1B$

$2n + 1 = 5, n = 2$ :



Now: We can step up and finish at the first white, because line connecting the remaining two whites can fill the area.

$(x+1)W \rightarrow 1B$

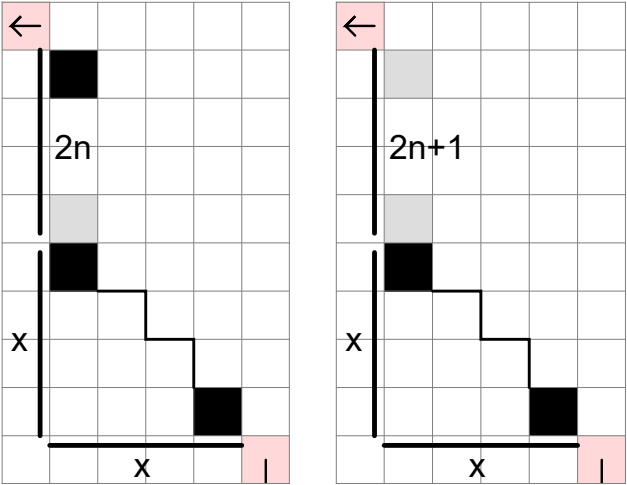
Later:  $xW \rightarrow 2B$

From now on, the calculations are the following:

If we enter now,  $x + (n - n \% 2) / 2$  white lines can be drawn if  $n > 0$ .  
The number of black fields is  $n+1$ , and when we use up one to finish the first line,  $n$  amount remains, one of which is a corner. Add one to make pairs, and the formula will be  $(n+1 - (n+1) \% 2) / 2$ .

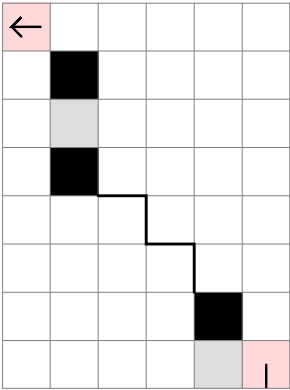
When entering later, the white line count is  $x-1 + (n+1 - (n+1) \% 2) / 2$ , and the black line count is  $(n+2 - (n+2) \% 2) / 2$  if  $n > 0$ .

3. Larger vertical distance



If the added distance is pair:

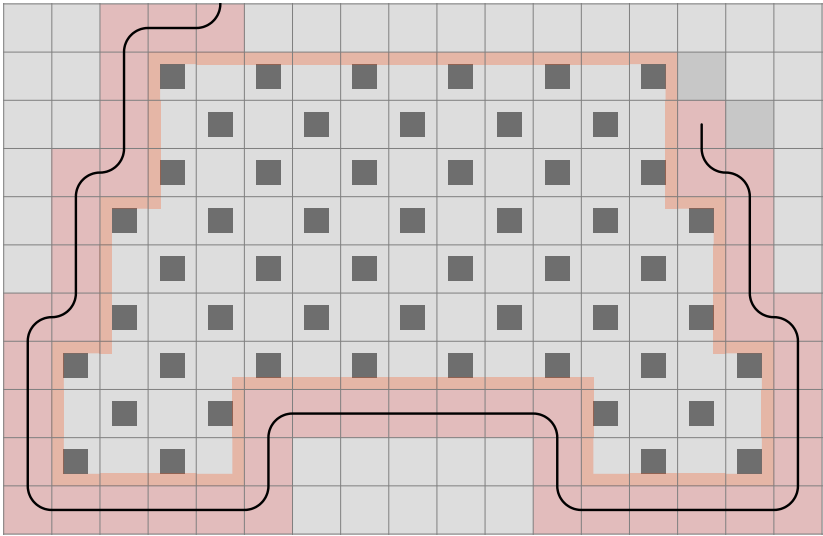
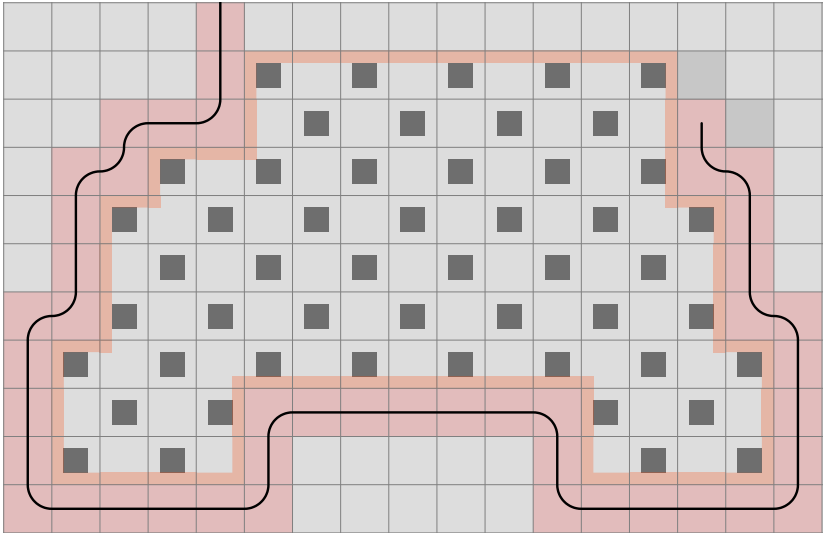
$2n = 2$ :



Now: It is possible to exit at the white field, having filled everything. The previous step must have been the last black.  
When counting the blacks, we exit the first line at the first black field above the stair, and  $x$  amount of corner blacks will remain.  
 $1W \rightarrow xB$

Later: Same as the  $2n = 2$  case previously.  
 $0W \rightarrow xB$

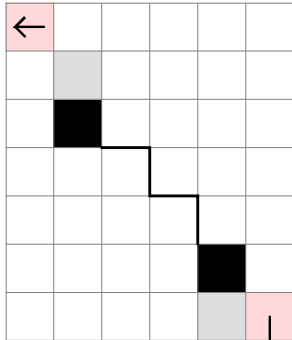
Compare these two cases:



When entering now, the general formula will be  $(n+1 - (n+1) \% 2) / 2$  white and  $x + (n-1 - (n-1) \% 2) / 2$  black.  
The later case is the same as previously,  $(n - n \% 2) / 2$  for white and  $x + (n - n \% 2) / 2$  for black.

### If the added distance is impair:

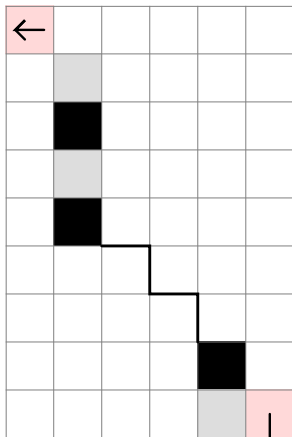
$$2n + 1 = 1, n = 0:$$



Now: 1W -> (x-1)B

Later: (see horizontal case) 0W -> (x-1)B

$$2n + 1 = 3, n = 1:$$



Now: 2W -> (x-1)B

Later: 1W -> xB

Now: The corner white always gives 1. Then make pairs with the remaining inline whites plus the white field we are stepping first.

$$1 + (n+1 - (n+1) \% 2) / 2$$

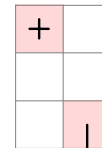
For the blacks, one of the inline black fields will be taken by the first line. We can then make pairs with the remaining inline blacks and add the corners.

$$x - 1 + (n - n \% 2) / 2.$$

Later: The white line count is  $1 + (n - n \% 2) / 2$  if  $n > 0$ , and the black line count is  $x - 1 + (n+1 - (n+1) \% 2) / 2$ .

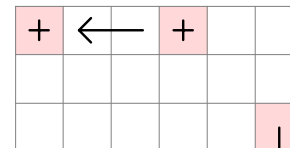
Next, we will look at the corner discovery algorithm.

Starting with 1 horizontal and 2 vertical distance, we check if that field is taken.



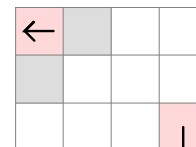
If so, we mirror sides and start the algorithm on the right side.

If not, we increase the horizontal distance by one until we find a taken field or run into the border.

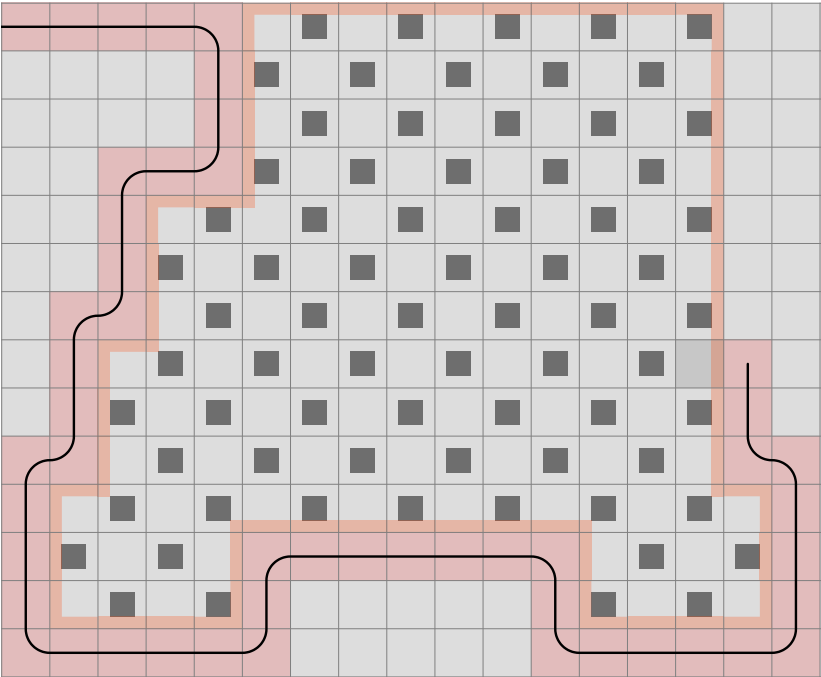


If it is a border field, we increase the vertical distance and start with 1 horizontal distance again.

Otherwise, we check if the bottom field is free, and by comparing the index of the corner field with the field above, we can determine if the line is going down and left, so the enclosed area is on the side we want.

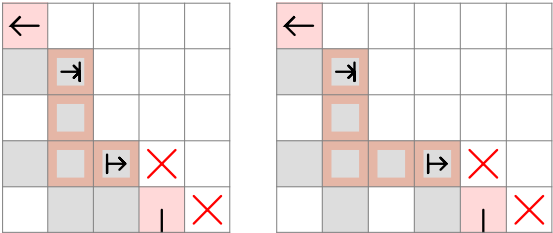


And as we step back, we find the point where the line should move in another direction.

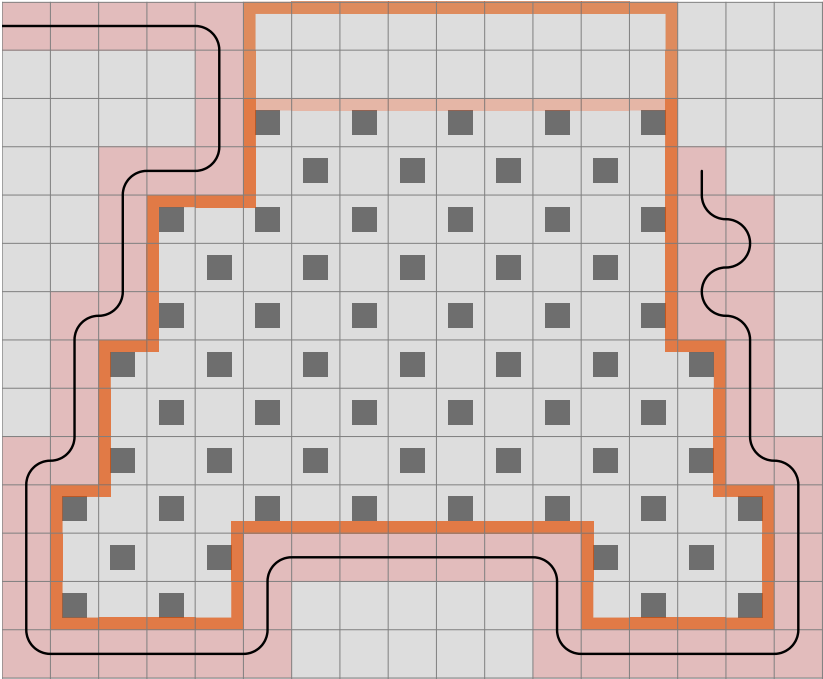


The only difference is the added 2x2 area. To the first, we apply the straight-to-side algorithm, while at the second, we have a corner that defines the area, but essentially, the procedure is the same.

Having the universal algorithm, these two size-specific rules can be deleted:

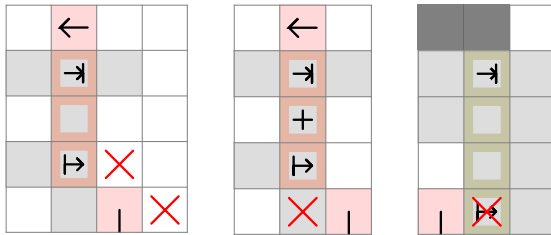


As we continue the case above, soon we will discover a deficiency which actually has been visible all along.



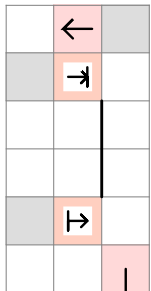
With the border on top, now we have no option to move.

The following rules are active, in addition to the universal one that disables the left field:

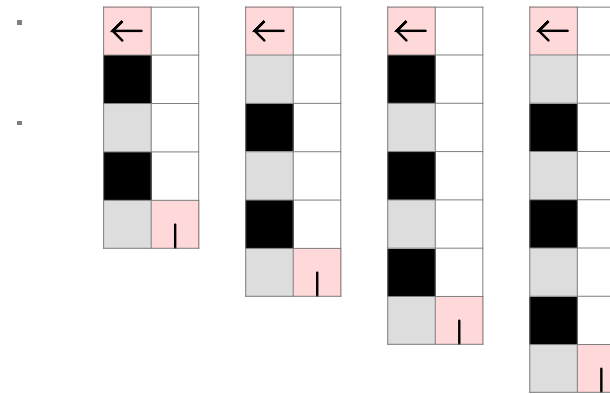


There are three more black fields in the area than white, so there is not enough vertical distance for entering and exiting that many times.

The straight-to-side algorithm has to be rotated upwards, so we get this:



Let's look at the distances from 3 to 6 to have an example of each case of modulo 4.



$D \pmod 4 = 3$ :

Now:  $1W \rightarrow 0B$ ,  $2W \rightarrow 1B$  etc. =  $(D+1)/4$   $W \rightarrow (D-3)/4$   $B$   
 Later:  $0W \rightarrow 1B$ ,  $1W \rightarrow 2B$  etc. =  $(D-3)/4$   $W \rightarrow (D+1)/4$   $B$   
 (double rule)

$D \pmod 4 = 0$ :

Now:  $1W \rightarrow 0B$ ,  $2W \rightarrow 1B$  etc. =  $D/4$   $W \rightarrow D/4 - 1$   $B$   
 Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $D/4$   $W \rightarrow D/4$   $B$   
 (single rule)

$D \pmod 4 = 1$ :

Now:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-1)/4$   $W \rightarrow (D-1)/4$   $B$   
 Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-1)/4$   $W \rightarrow (D-1)/4$   $B$   
 (no rule)

$D \pmod 4 = 2$ :

Now:  $2W \rightarrow 1B$ ,  $3W \rightarrow 2B$  etc. =  $(D+2)/4$   $W \rightarrow (D-2)/4$   $B$   
 Later:  $1W \rightarrow 1B$ ,  $2W \rightarrow 2B$  etc. =  $(D-2)/4$   $W \rightarrow (D-2)/4$   $B$   
 (single rule)

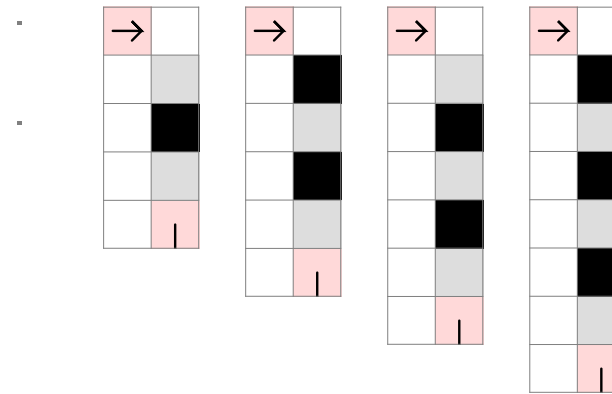
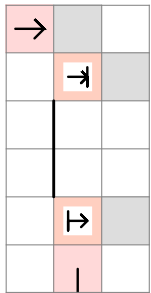


Figure 1 consists of two side-by-side diagrams illustrating a method for handling missing data in a 10x10 grid. Both diagrams show a 10x10 grid with a black boundary. The left diagram shows a grid with missing data (gray squares) and a black boundary. The right diagram shows the same grid with a red boundary and a red square indicating the location of the missing data.

- 0 vertical distance small area
- x horizontal and y vertical distance small area
- 0 horizontal distance small area
- 0 horizontal distance big area

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An extension of the universal rule is necessary to include cases of a "big" area where the obstacle is on the other side of the live end.



Now:  $1W \rightarrow 0B = (D+1)/4$   $W \rightarrow (D-3)/4$  B  
 Later:  $1W \rightarrow 0B = (D+1)/4$   $W \rightarrow (D-3)/4$  B  
 (no rule)

Now:  $1W \rightarrow 0B = D/4$   $W \rightarrow D/4 - 1 B$   
 Later:  $1W \rightarrow 1B = D/4$   $W \rightarrow D/4 B$   
 (single rule)

Now:  $2W \rightarrow 0B = (D+3)/4$   $W \rightarrow (D-5)/4$  B  
 Later:  $1W \rightarrow 1B = (D-1)/4$   $W \rightarrow (D-1)/4$  B  
 (double rule)

Now:  $2W \rightarrow 1B = (D+2)/4$   $W \rightarrow (D-2)/4 B$   
 Later:  $1W \rightarrow 1B = (D-2)/4$   $W \rightarrow (D-2)/4 B$   
 (single rule)

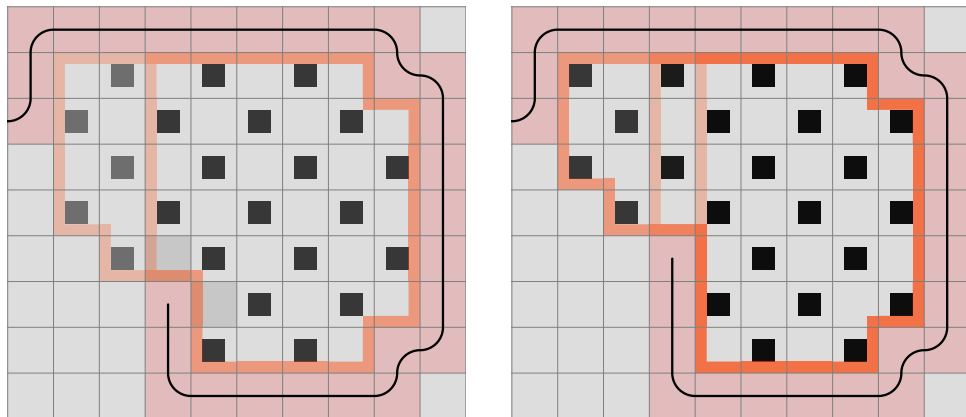
General:

Now: If  $n > 0$ , we can use all corner blacks after exiting the first line.

$(n+2 - (n+2) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$  if  $n > 0$ .

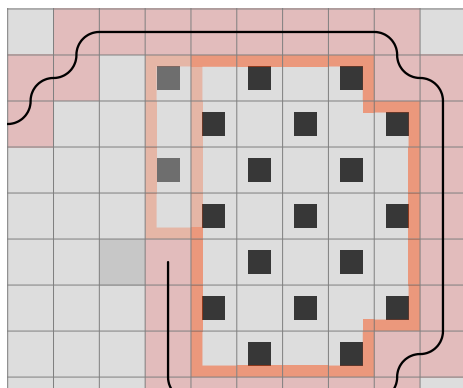
Later:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$  if  $n > 0$ .

While creating a case to verify the newly created rule set, I have encountered this:



It is not possible to continue after stepping upwards.  
But where is the missing part?

On the left, the largest area contains 1 more white field than black, and on the right it is 2 more black. It would be possible if the upper left corner was filled, like this:



I will now take the case of  $x$  horizontal and  $y$  vertical distance big area.

All representations are the same as in the small area case, only the area is on the other side.

So I will just summarize the black and white limits here.

### 1. Equal horizontal and vertical distance

$x = 2$ :

Now:  $0W \rightarrow 1B$

Later:  $1B$

$x = 3$ :

Now:  $0W \rightarrow 2B$

Later:  $1B \rightarrow 2B$

$x = n$ :

Now:  $0W \rightarrow (n-1)B$

Later:  $1B \rightarrow (n-1)B$

Exactly the same as with the small area.

### 2. Larger horizontal distance

#### If the added distance is pair:

$2n = 2$ :

Now:  $1W$  is possible. The entry field and the first white is at least 2 distance from each other, the whole area can be filled between them.

$1W \rightarrow xB$

Later: All corner blacks can be used. The line between the first black and first white will fill the area.

$0W \rightarrow xB$

General:

The only difference is the  $n = 1$  case. Otherwise, the number of inline and corner fields are the same.

Now:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

**If the added distance is impair:**

$2n + 1 = 1, n = 0:$

Now:  $xW$  is possible.

$xW \rightarrow 0B$

Later: Same values as previously.

$(x-1)W \rightarrow 0B$

$2n + 1 = 3, n = 1:$

Note that in the small area case, there were  $x-1$  corner whites and 1 corner black. Now there are  $x$  amount of corner whites and 2 inline blacks instead of 1.

Now:  $xW \rightarrow 0B$

Later:  $xW \rightarrow 1B$

$2n + 1 = 5, n = 2:$

Now:  $(x+1)W \rightarrow 1B$

Later:  $(x+1)W \rightarrow 1B$

General:

Now:  $x + (n+1 - (n+1) \% 2) / 2 W \rightarrow (n - n \% 2) / 2 B$

Later:  $x + (n - n \% 2) / 2 W$  if  $n > 0 \rightarrow (n+1 - (n+1) \% 2) / 2 B$

3. Larger vertical distance

**If the added distance is pair:**

We will find it is the same as the small area.

$2n = 2:$

Now:  $1W \rightarrow xB$

Later:  $0W \rightarrow xB$

$2n = 4:$

Now:  $1W \rightarrow xB$

Later:  $1W \rightarrow (x+1)B$

General:

Now:  $(n+1 - (n+1) \% 2) / 2 W \rightarrow x + (n-1 - (n-1) \% 2) / 2 B$

Later:  $(n - n \% 2) / 2 W \rightarrow x + (n - n \% 2) / 2 B$

**If the added distance is impair:**

$2n + 1 = 1, n = 0:$

Now:  $1W \rightarrow (x-1)B$

Later:  $0W \rightarrow (x-1)B$

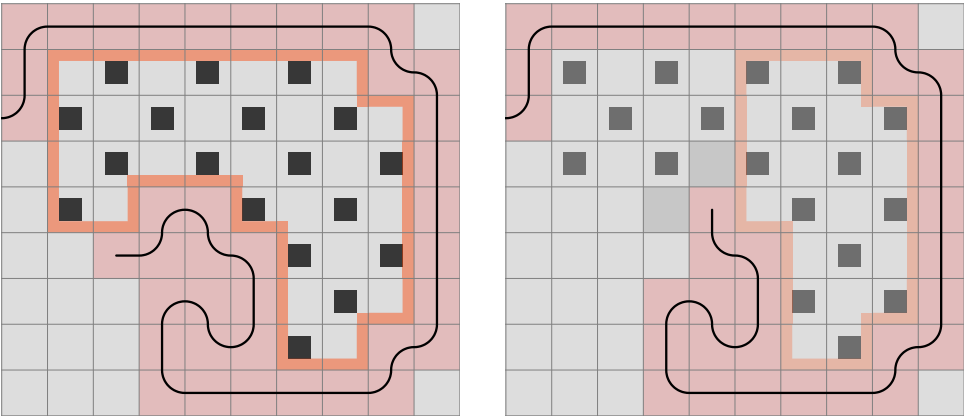
$2n + 1 = 3, n = 1:$

Here comes the change again, due to having one more corner black field and one less corner white field than in the small area case.

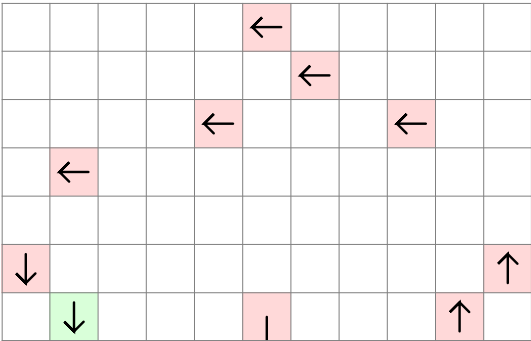
Now:  $1W \rightarrow xB$

Later:  $1W \rightarrow xB$

Now we can continue the case, but we will find that we cannot go past this point:

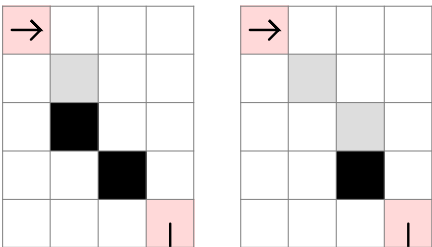


The picture on the right is the crossroad. If we step upwards, the area can be completed.  
 The area between the live end and the corner on the left contains one more black field than white, so the number could be made by stepping left and up, only we cannot step left because of the C-shape. In other words, when stepping left, the line will not be able to continue, but we need to change the programming algorithm in order to see it.  
 Now when there is a C-Shape, it is possible to step there, and no other rules will be checked. But C-Shape is just one of the nine cases of area checking.



So it is not any of the small area rules and neither the 0 horizontal distance big area rule that has something to do with it.

One thing is sure, we have been using the small area representations when defining the rule set, which does not give us the minimal area in this case. See the difference:

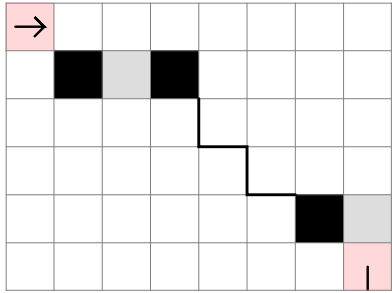


Can it be a problem?  
 We will see, but let's look at one detail: If we step upwards, we have to step left to fill the corner white, otherwise it is only good for an exit, which we do not want if the area contains 1 more white fields than black (for the left representation).  
 After this, we step upwards and then left again. We did not enter the area.  
 The 1 added distance case is therefore  $0W \rightarrow (x-1)B$  when entering now by stepping upwards and  $1W \rightarrow (x-1)B$  when stepping right.

To simplify things, I will specify the cases again using the minimal area representation.

Notice that these are the same as the small area patterns, things are just mirrored, so that the previous horizontal expansion is now vertical.

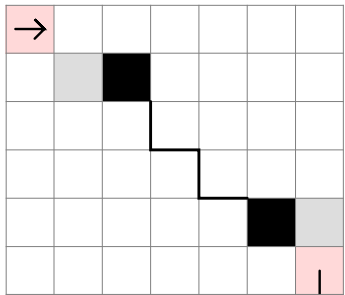
n = 1



Now:  $(n+1 - (n+1) \% 2) / 2$  W  $\rightarrow x + (n-1 - (n-1) \% 2) / 2$  B

Later:  $(n - n \% 2) / 2$  W  $\rightarrow x + (n - n \% 2) / 2$  B

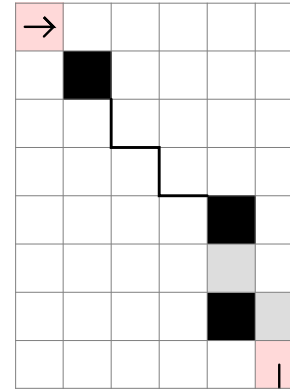
n = 0



Now:  $1 + (n+1 - (n+1) \% 2) / 2$  W  $\rightarrow x - 1 + (n - n \% 2) / 2$  B

Later:  $1 + (n - n \% 2) / 2$  W if  $n > 0$  (0 if  $n = 0$ )  $\rightarrow x - 1 + (n+1 - (n+1) \% 2) / 2$  B

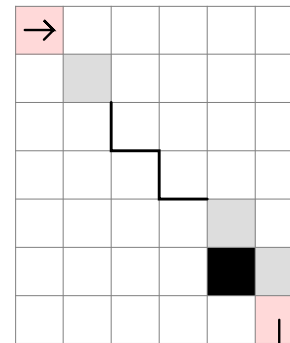
n = 1



Now:  $(n+1 - (n+1) \% 2) / 2$  W if  $n > 1$  (0 if  $n = 1$ )  $\rightarrow x + (n-1 - (n-1) \% 2) / 2$  B

Later:  $(n - n \% 2) / 2$  W  $\rightarrow x + (n - n \% 2) / 2$  B

n = 0



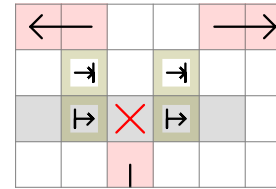
Now:  $x - 1 + (n+2 - (n+2) \% 2) / 2$  W if  $n > 0$  ( $x - 1$  if  $n = 0$ )  $\rightarrow (n+1 - (n+1) \% 2) / 2$  B

Later:  $x - 1 + (n+1 - (n+1) \% 2) / 2$  W  $\rightarrow (n+2 - (n+2) \% 2) / 2$  B if  $n > 0$  (0 if  $n = 0$ )

The difference between stepping up and right still remains in these vertical expansion cases. We didn't have to deal with it at the small area, because the line could not step backwards.

It just means we have to remove the Now W conditions for stepping right.

Remember we have made rules previously that take both sides into account, like this:



To combine all cases where after the next step, on the left side only the left field will be possible, and on the right, only the right, is not feasible.

Let's say instead that if one of the possible fields will lead to an impossible situation, we cannot step there. It still results in a usable algorithm.

(An algorithm is usable if the number of operations are on the order of  $n^2$  where  $n$  indicates the size of the table.

An unusable algorithm would require a number of operations on the order of  $2^n$  or more: that is the random path where you just guess the next step until you get an impossible situation, but the crossroad might have been as far back as almost  $n \times n$  steps.)

We are still missing the straight obstacle case:

