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Top Down Processing

- $\bullet\,$ Considering the Grammar
- 1. $\langle s \rangle \rightarrow D \langle s \rangle \langle A \rangle$
- 2. $\langle s \rangle \rightarrow B \langle A \rangle C$
- $3. \langle A \rangle \rightarrow D \langle A \rangle$
- $4. \langle A \rangle \rightarrow C$
- This is an S-Grammar
 - S stand for simple

DBBCCDC

- $\langle s \rangle \stackrel{1}{\Rightarrow} D \langle s \rangle \langle A \rangle$
- D<s><A> $\stackrel{2}{\Rightarrow}$ DB<A>C<A>
- DB<A>C<A> $\stackrel{4}{\Rightarrow}$ DBCC<A> DBCC<A> $\stackrel{3}{\Rightarrow}$ DBCCD<A>
- DBCCD<A>→ DBCCDC

The parse is deterministic. Can only be parsed correctly.

- Inputs = $\{B, C, D, \exists\}$
- Stack Symbols = $\{\langle s \rangle, \langle A \rangle, C, \nabla\}$
- Starting Stack = ∇S

	В	С	D	4
<s></s>	#2		#1	
<a>		#4	#3	
\mathbf{C}		POP,ADVANCE		
∇				ACCEPT

All blank non table entries represent reject

- <A> \rightarrow B α
- 1. REPLACE(<A><s>), ADVANCE
- 2. REPLACE(C<A>), ADVANCE
- 3. REPLACE(A), ADVANCE
- 4. POP, ADVANCE

At each step during the parse there is an assertion that the input is correct if and only if the string of remaining terminals can be derived from the sequence of symbols on the stack

Stack	Inputs
∇ <s></s>	DBCCBD-
⊽ <a><s></s>	$\mathrm{BCCDC}\dashv$
∇ <a>C<a>	$\mathrm{CCDC}\dashv$
∇ <a>C	$\mathrm{CDC}\dashv$
▽<a>	$\mathrm{DC}\dashv$
▽<a>	$\mathrm{C}\dashv$
∇	\dashv

ACCEPT

Past Inputs	Stack	String	
	<s>∇</s>	<s></s>	
D	<s><a>▽</s>	D <s><a></s>	
DB	<a>C<a>▽	DB <a>C<a>	
DBC	$\mathrm{C} extsf{<}\mathrm{A} extsf{>} riangle$	DBCC <a>	
DBCC	< A> ▽	DBCC <a>	
DBCCDC	∇	DBCCDC	

Stack from right to left

Intermediate string concatenation of the past inputs and the symbols on the stack