

Lecture 9: A closer look at terms

- Theory
 - Introduce the == predicate
 - Take a closer look at term structure
 - Introduce strings in Prolog
 - Introduce operators
- Exercises
 - Exercises of LPN: 9.1, 9.2, 9.3, 9.4, 9.5
 - Practical session

Comparing terms: `==/2`

- Prolog contains an important predicate for comparing terms
- This is the identity predicate `==/2`
- The identity predicate `==/2` does not instantiate variables, that is, it behaves differently from `=/2`

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```
?- a==a.
```

```
yes
```

```
?- a==b.
```

```
no
```

```
?- a=='a'.
```

```
yes
```

```
?- a==X.
```

```
X = _443
```

```
no
```

Comparing variables

- Two different **uninstantiated** variables are not identical terms
- Variables **instantiated** with a term T are identical to T

Comparing variables

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```
?- X==X.
```

```
X = _443
```

```
yes
```

```
?- Y==X.
```

```
Y = _442
```

```
X = _443
```

```
no
```

```
?- a=U, a==U.
```

```
U = _443
```

```
yes
```

Comparing terms: $\neq/2$

- The predicate $\neq/2$ is defined so that it succeeds in precisely those cases where $=/2$ fails
- In other words, it succeeds whenever two terms are **not identical**, and fails otherwise

Comparing terms: $\backslash==/2$

- The predicate $\backslash==/2$ is defined so that it succeeds in precisely those cases where $==/2$ fails
- In other words, it succeeds whenever two terms are **not identical**, and fails otherwise

?- a $\backslash==$ a.

no

?- a $\backslash==$ b.

yes

?- a $\backslash==$ 'a'.

no

?- a $\backslash==$ X.

X = _443

yes

Terms with a special notation

- Sometimes terms look different, but Prolog regards them as identical
- For example: **a** and '**a**', but there are many other cases
- Why does Prolog do this?
 - Because it makes programming more pleasant
 - More natural way of coding Prolog programs

Arithmetic terms

- Recall lecture 5 where we introduced arithmetic
- $+$, $-$, $<$, $>$, etc are functors and expressions such as $2+3$ are actually ordinary complex terms
- The term $2+3$ is identical to the term $+(2,3)$

Arithmetic terms

- Recall lecture 5 where we introduced arithmetic
- $+$, $-$, $<$, $>$, etc are functors and expressions such as $2+3$ are actually ordinary complex terms
- The term $2+3$ is identical to the term $+(2,3)$

?- $2+3 == +(2,3).$

yes

?- $-(2,3) == 2-3.$

yes

?- $(4<2) == <(4,2).$

yes

Summary of comparison predicates

$=$	Unification predicate
$\backslash=$	Negation of unification predicate
$==$	Identity predicate
$\backslash==$	Negation of identity predicate
$:=$	Arithmetic equality predicate
$=\backslash=$	Negation of arithmetic equality predicate

Lists as terms

- Another example of Prolog working with one internal representation, while showing another to the user
- Using the `!` constructor, there are many ways of writing the same list

```
?- [a,b,c,d] == [a![b,c,d]].  
yes  
?- [a,b,c,d] == [a,b,c![d]].  
yes  
?- [a,b,c,d] == [a,b,c,d![]].  
yes  
?- [a,b,c,d] == [a,b![c,d]].  
yes
```

Prolog lists internally

- Internally, lists are built out of two special terms:
 - [] (which represents the empty list)
 - '.' (a functor of arity 2 used to build non-empty lists)
- These two terms are also called *list constructors*
- A recursive definition shows how they construct lists

Definition of prolog list

- The empty list is the term $[]$. It has length 0.
- A non-empty list is any term of the form $.(term, list)$, where $term$ is any Prolog term, and $list$ is any Prolog list. If $list$ has length n , then $.(term, list)$ has length $n+1$.

A few examples...

?- .(a,[]) == [a].

yes

?- .(f(d,e),[]) == [f(d,e)].

yes

?- .(a,.(b,[])) == [a,b].

yes

?- .(a,.(b,.(f(d,e),[]))) == [a,b,f(d,e)].

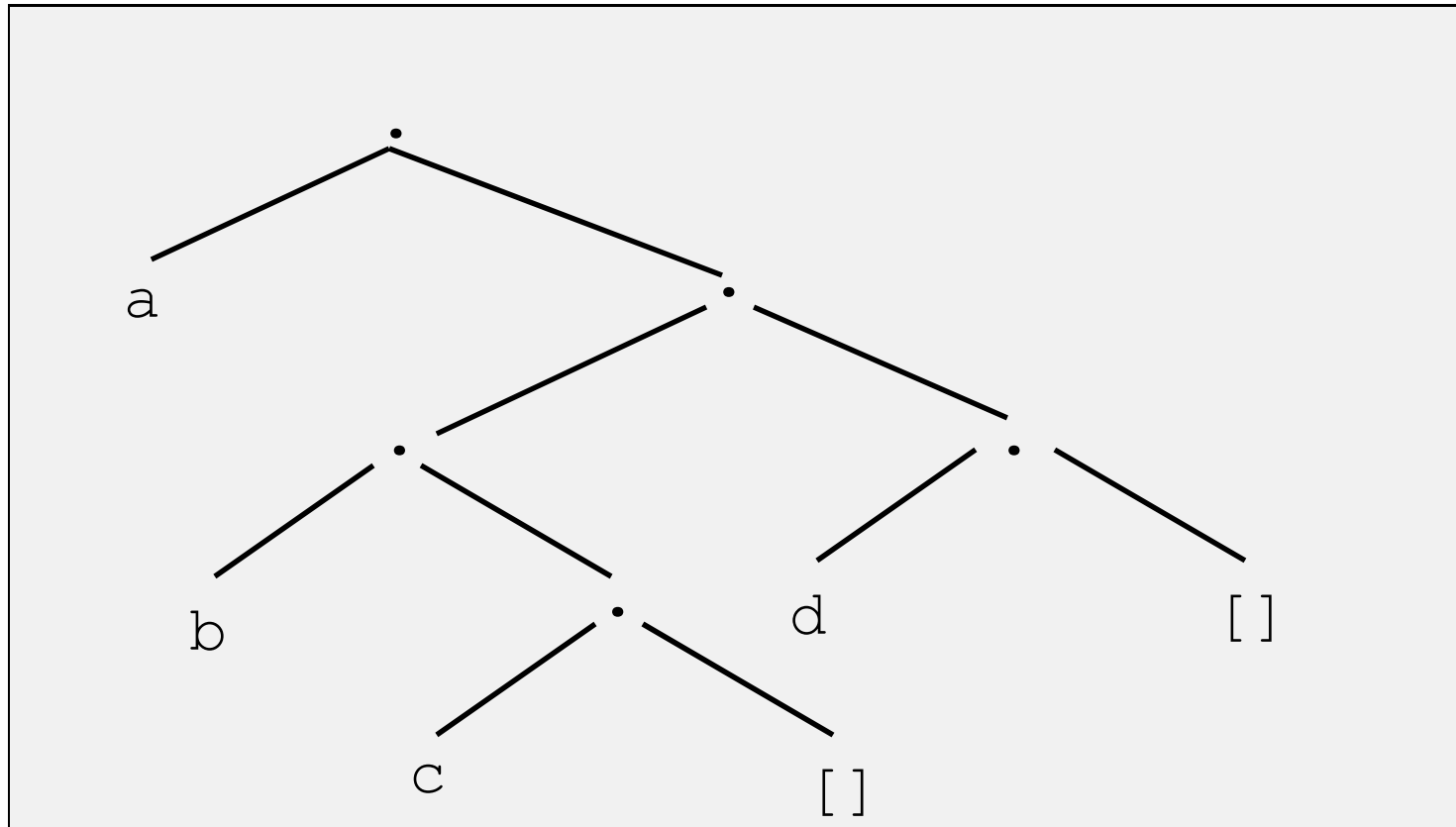
yes

Internal list representation

- Works similar to the l notation:
- It represents a list in two parts
 - Its first element, the *head*
 - the rest of the list, the *tail*
- The trick is to read these terms as trees
 - Internal nodes are labeled with .
 - All nodes have two daughter nodes
 - Subtree under left daughter is the head
 - Subtree under right daughter is the tail

Example of a list as tree

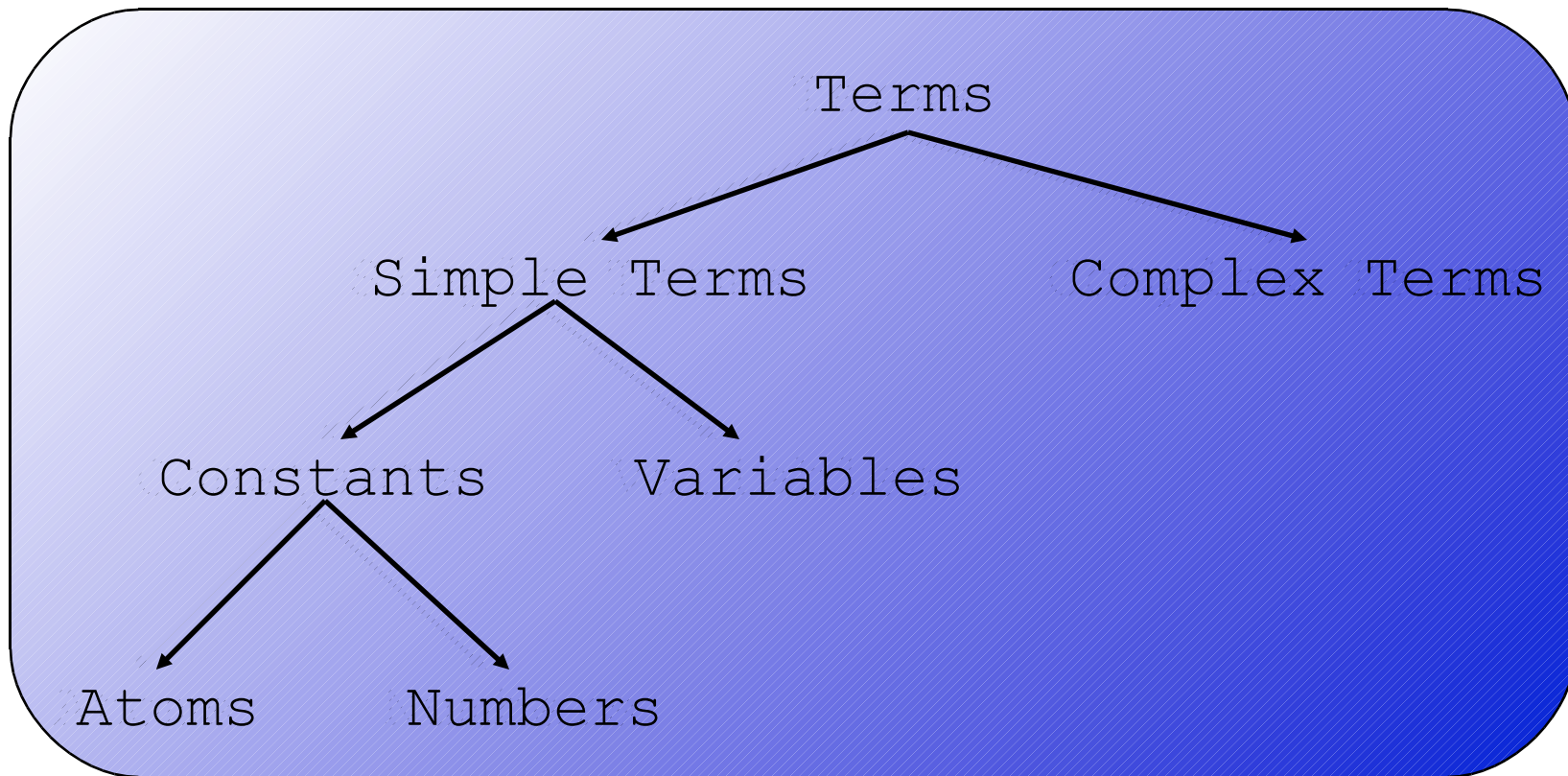
- Example: $[a, [b, c], d]$



Examining terms

- We will now look at built-in predicates that let us examine Prolog terms more closely
 - Predicates that determine the type of terms
 - Predicates that tell us something about the internal structure of terms

Type of terms



Checking the type of a term

atom/1	<i>Is the argument an atom?</i>
integer/1	<i>... an interger?</i>
float/1	<i>... a floating point number?</i>
number/1	<i>... an integer or float?</i>
atomic/1	<i>... a constant?</i>
var/1	<i>... an uninstantiated variable?</i>
nonvar/1	<i>... an instantiated variable or another term that is not an uninstantiated variable</i>

Type checking: atom/1

?- atom(a).

yes

?- atom(7).

no

?- atom(X).

no

Type checking: atom/1

?- X=a, atom(X).

X = a

yes

?- atom(X), X=a.

no

Type checking: atomic/1

?- atomic(mia).

yes

?- atomic(5).

yes

?- atomic(loves(vincent,mia)).

no

Type checking: var/1

?- var(mia).

no

?- var(X).

yes

?- X=5, var(X).

no

Type checking: nonvar/1

?- nonvar(X).

no

?- nonvar(mia).

yes

?- nonvar(23).

yes

The structure of terms

- Given a complex term of unknown structure, what kind of information might we want to extract from it?
- Obviously:
 - The functor
 - The arity
 - The argument
- Prolog provides built-in predicates to produce this information

The functor/3 predicate

- The functor/3 predicate gives the functor and arity of a complex predicate

The functor/3 predicate

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 - ?- functor(friends(lou,andy),F,A).
 - F = friends
 - A = 2
 - yes

The functor/3 predicate

- The functor/3 predicate gives the functor and arity of a complex predicate

□?- functor(friends(lou,andy),F,A).

F = friends

A = 2

yes

□?- functor([lou,andy,vicky],F,A).

F = .

A = 2

yes

functor/3 and constants

- What happens when we use functor/3 with constants?

functor/3 and constants

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 - ?- functor(mia,F,A).
F = mia
A = 0
yes

functor/3 and constants

- What happens when we use functor/3 with constants?

☐?- functor(mia,F,A).

F = mia

A = 0

yes

☐?- functor(14,F,A).

F = 14

A = 0

yes

functor/3 for constructing terms

- You can also use functor/3 to construct terms:
 - ?- functor(Term,friends,2).
Term = friends(__,__)
yes

Checking for complex terms

```
complexTerm(X):-  
    nonvar(X),  
    functor(X,_,A),  
    A > 0.
```

Arguments: arg/3

- Prolog also provides us with the predicate `arg/3`
- This predicate tells us about the arguments of complex terms
- It takes three arguments:
 - A number N
 - A complex term T
 - The N th argument of T

Arguments: arg/3

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```
?- arg(2,likes(lou,andy),A).  
A = andy  
yes
```

Strings

- Strings are represented in Prolog by a list of character codes
- Prolog offers double quotes for an easy notation for strings

```
?- S = "Vicky".
```

```
S = [86,105,99,107,121]
```

```
yes
```

Working with strings

- There are several standard predicates for working with strings
- A particular useful one is atom_codes/2

```
?- atom_codes(vicky,S).  
S = [118,105,99,107,121]  
yes
```

Operators

- As we have seen, in certain cases, Prolog allows us to use operator notations that are more user friendly
- Recall, for instance, the arithmetic expressions such as $2+2$ which internally means $+(2,2)$
- Prolog also has a mechanism to add your own operators

Properties of operators

- Infix operators
 - Functors written between their arguments
 - Examples: + - = == , ; . -->
- Prefix operators
 - Functors written before their argument
 - Example: - (to represent negative numbers)
- Postfix operators
 - Functors written after their argument
 - Example: ++ in the C programming language

Precedence

- Every operator has a certain precedence to work out ambiguous expressions
- For instance, does $2+3*3$ mean $2+(3*3)$, or $(2+3)*3$?
- Because the precedence of $+$ is greater than that of $*$, Prolog chooses $+$ to be the main functor of $2+3*3$

Associativity

- Prolog uses associativity to disambiguate operators with the same precedence value
- Example: $2+3+4$
Does this mean $(2+3)+4$ or $2+(3+4)$?
 - Left associative
 - Right associative
- Operators can also be defined as non-associative, in which case you are forced to use bracketing in ambiguous cases
 - Examples in Prolog: `:-` `-->`

Defining operators

- Prolog lets you define your own operators
- Operator definitions look like this:

```
:- op(Precedence, Type, Name.
```

- Precedence:
number between 0 and 1200
- Type: the type of operator

Types of operators in Prolog

- yfx left-associative, infix
- xfy right-associative, infix
- xfx non-associative, infix
- fx non-associative, prefix
- fy right-associative, prefix
- xf non-associative, postfix
- yf left-associative, postfix

Operators in SWI Prolog

1200	<i>xfx</i>	-->, :-
1200	<i>fx</i>	:-, ?-
1150	<i>fx</i>	dynamic, discontinuous, initialization, module-transparent, multifile, thread_local, volatile
1100	<i>xfy</i>	!,
1050	<i>xfy</i>	->, op*->
1000	<i>xfy</i>	,
954	<i>xfy</i>	\
900	<i>fy</i>	\+
900	<i>fx</i>	~
700	<i>xfx</i>	<, =, =.., =@=, =:=, =<, ==, =\=, >, >=, @<, @=<, @>, @>=, \=, \==, is
600	<i>xfy</i>	:
500	<i>yfx</i>	+, -, /\, \/, xor
500	<i>fx</i>	+, -, ?, \
400	<i>yfx</i>	*, /, //, rdiv, <<, >>, mod, rem
200	<i>xfx</i>	**
200	<i>xfy</i>	^

Next lecture

- Cuts and negation
 - How to control Prolog`s backtracking behaviour with the help of the cut predicate
 - Explain how the cut can be packaged into a more structured form, namely negation as failure