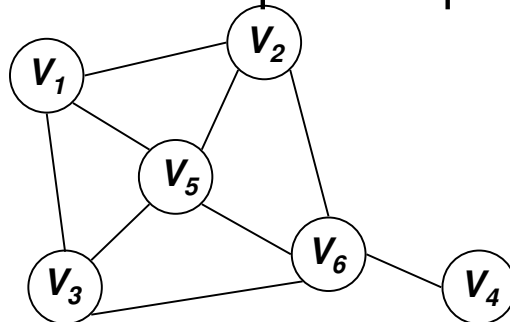
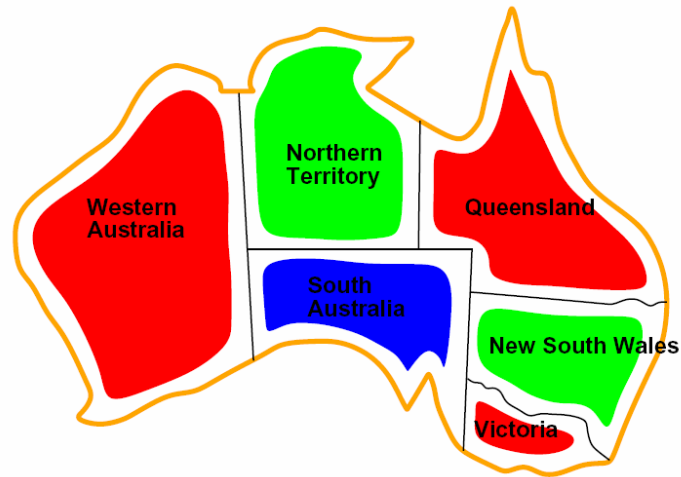


Canonical Example: Graph Coloring



- Consider N nodes in a graph
- Assign values V_1, \dots, V_N to each of the N nodes
- The values are taken in $\{R, G, B\}$
- Constraints: If there is an edge between i and j , then V_i must be different of V_j

Canonical Example: Graph Coloring



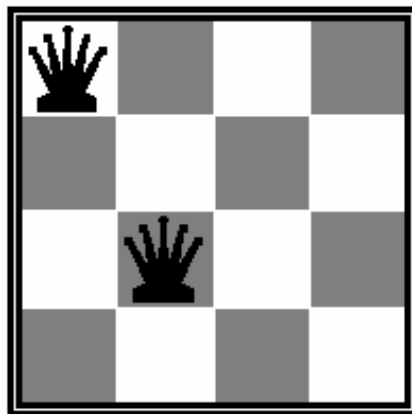
CSP Definition

- $CSP = \{V, D, C\}$
- *Variables:* $V = \{V_1, \dots, V_N\}$
 - Example: The values of the nodes in the graph
- *Domain:* The set of d values that each variable can take
 - Example: $D = \{R, G, B\}$
- *Constraints:* $C = \{C_1, \dots, C_K\}$
- Each constraint consists of a tuple of variables and a list of values that the tuple is allowed to take for this problem
 - Example: $[(V_2, V_3), \{(R, B), (R, G), (B, R), (B, G), (G, R), (G, B)\}]$
- Constraints are usually defined implicitly \rightarrow A function is defined to test if a tuple of variables satisfies the constraint
 - Example: $V_i \neq V_j$ for every edge (i, j)

Binary CSP

- Variable V and V' are connected if they appear in a constraint
- Neighbors of V = variables that are connected to V
- The domain of V , $D(V)$, is the set of candidate values for variable V
- $D_i = D(V_i)$
- Constraint graph for binary CSP problem:
 - Nodes are variables
 - Links represent the constraints
 - Same as our canonical graph-coloring problem

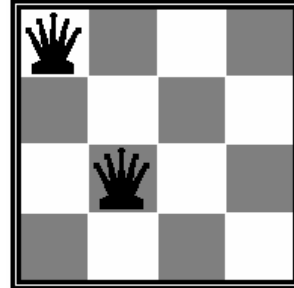
N-Queens



$$Q_1 = 1 \quad Q_2 = 3$$

Example: N-Queens

- Variables: Q_i
- Domains: $D_i = \{1, 2, 3, 4\}$
- Constraints
 - $Q_i \neq Q_j$ (cannot be in same row)
 - $|Q_i - Q_j| \neq |i - j|$ (or same diagonal)



$Q_1 = 1 \quad Q_2 = 3$

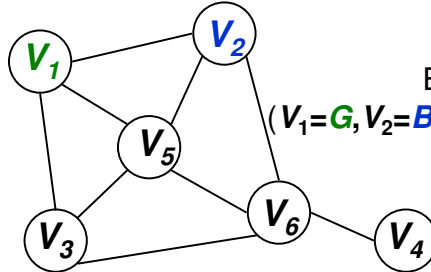
- Valid values for (Q_1, Q_2) are
 (1,3) (1,4) (2,4) (3,1) (4,1)
 (4,2)

Cryptarithmic

S E N D
 + M O R E

 M O N E Y

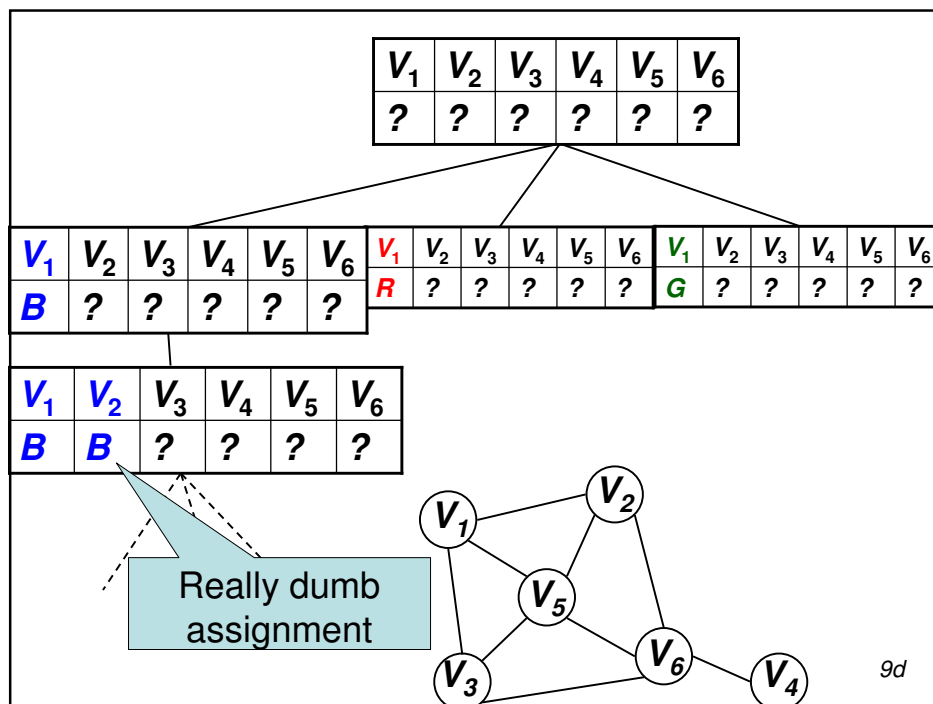
Search Space

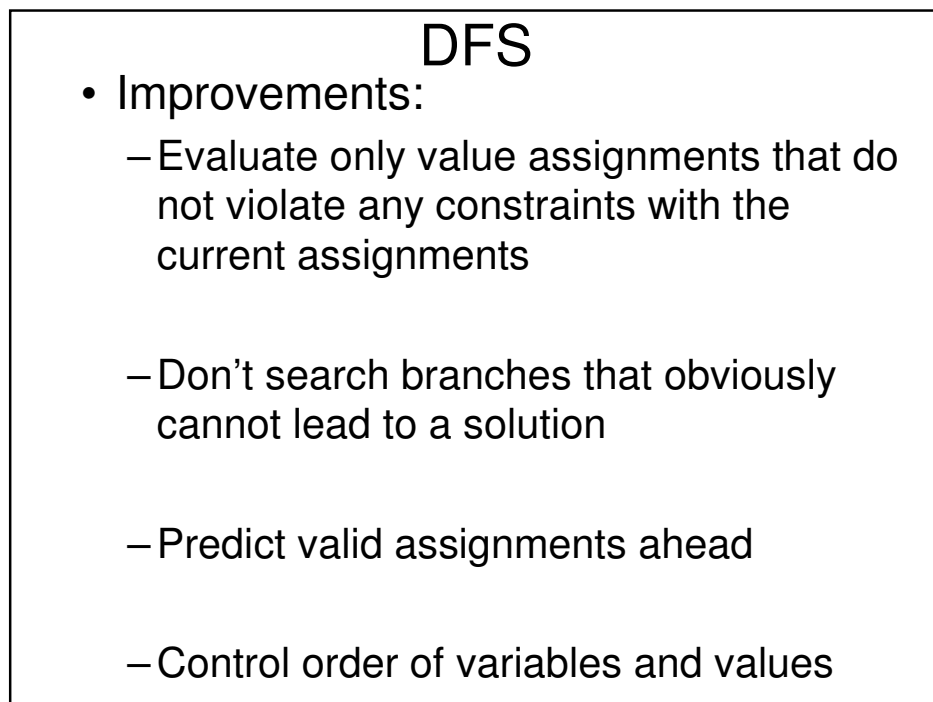
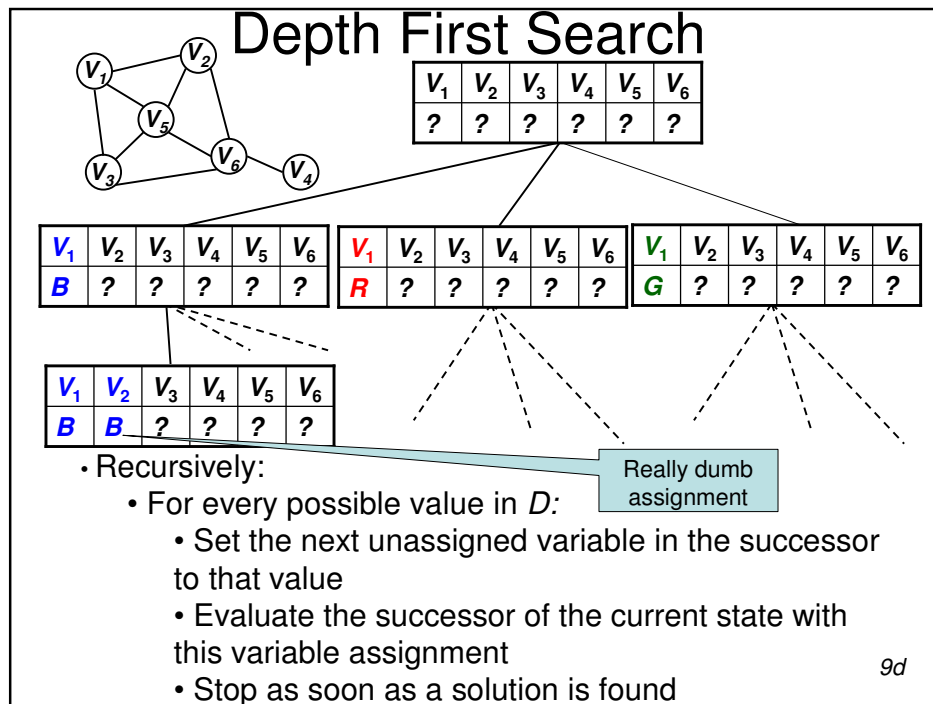


Example state:

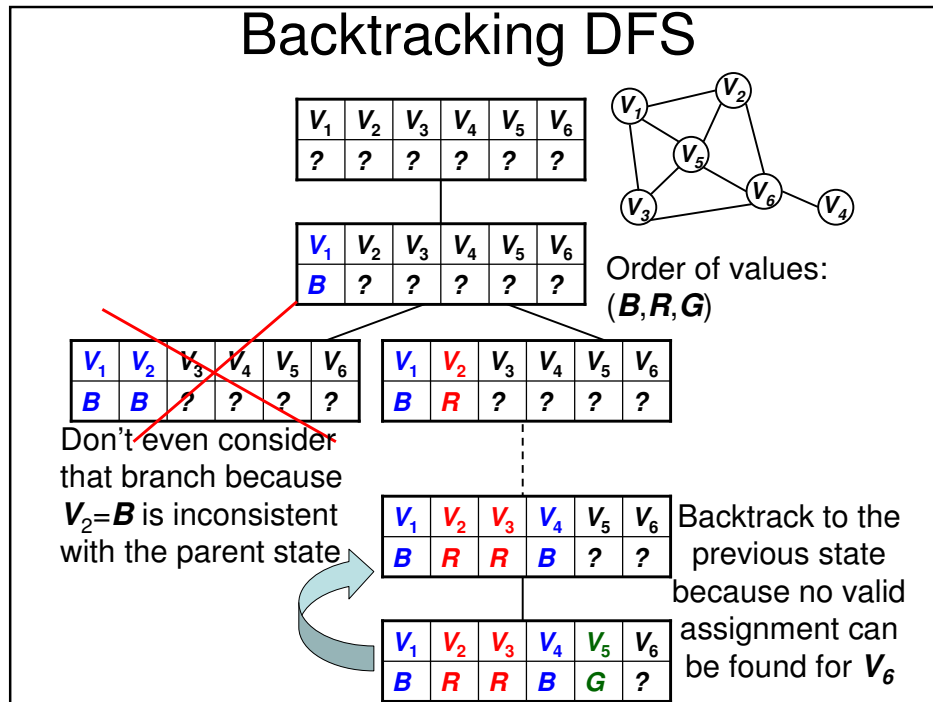
$(V_1=G, V_2=B, V_3=?, V_4=?, V_5=?, V_6=?)$

- *State*: assignment to k variables with $k+1, \dots, N$ unassigned
- *Successor*: The successor of a state is obtained by assigning a value to variable $k+1$, keeping the others unchanged
- *Start state*: $(V_1=?, V_2=?, V_3=?, V_4=?, V_5=?, V_6=?)$
- *Goal state*: All variables assigned with constraints satisfied
- No concept of cost on transition \rightarrow We just want to find a solution, we don't worry how we get there





Backtracking DFS



Backtracking DFS

- For every possible value x in D :
 - If assigning x to the next unassigned variable V_{k+1} does not violate any constraint with the k already assigned variables:
 - Set the variable V_{k+1} to x
 - Evaluate the successors of the current state with this variable assignment
- If no valid assignment is found:
Backtrack to previous state
- Stop as soon as a solution is found

9b, 27b

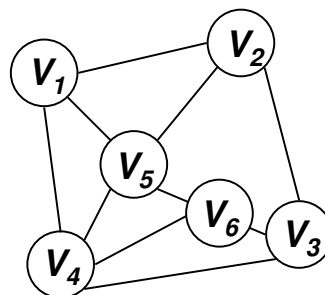
Backtracking DFS Comments

- Additional computation: At each step, we need to evaluate the constraints associated with the current candidate assignment (variable, value).
- Uninformed search, we can improve by predicting:
 - What is the effect of assigning a variable on all of the other variables?
 - Which variable should be assigned next and in which order should the values be evaluated?
 - When a branch fails, how can we avoid repeating the same mistake?

Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

	V_1	V_2	V_3	V_4	V_5	V_6
R	?	?	?	?	?	?
B	?	?	?	?	?	?
G	?	?	?	?	?	?

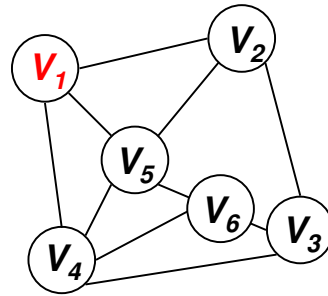


Warning: Different example with order (R,B,G)

Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

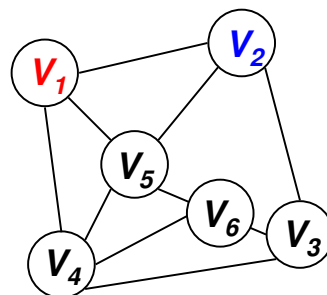
	V₁	V ₂	V ₃	V ₄	V ₅	V ₆
R	O	X	?	X	X	?
B		?	?	?	?	?
G		?	?	?	?	?



Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

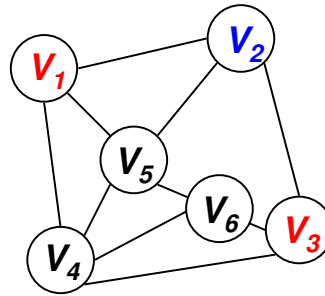
	V₁	V₂	V ₃	V ₄	V ₅	V ₆
R	O		?	X	X	?
B		O	X	?	X	?
G			?	?	?	?



Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when no variable has a legal value

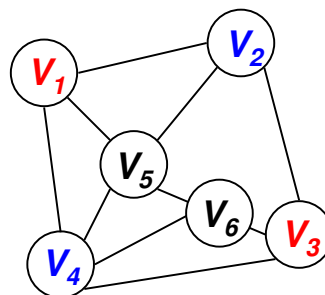
	V_1	V_2	V_3	V_4	V_5	V_6
R	O		O	X	X	X
B		O		$?$	X	$?$
G				$?$	$?$	$?$



Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

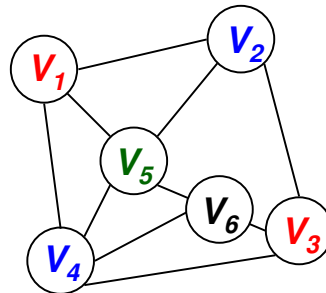
	V_1	V_2	V_3	V_4	V_5	V_6
R	O		O		X	X
B		O		O	X	X
G					$?$	$?$



Forward Checking

- Keep track of remaining legal values for unassigned variables
- Backtrack when any variable has no legal values

	V_1	V_2	V_3	V_4	V_5	V_6
R	O		O			X
B		O		O		X
G					O	X



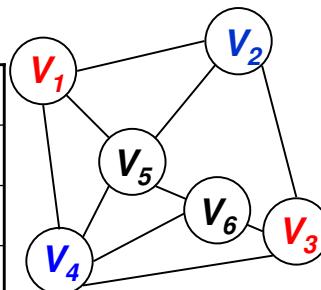
There are no valid assignments left for V_6 we need to backtrack

27f

Constraint Propagation

- Forward checking does not detect all the inconsistencies, only those that can be detected by looking at the constraints which contain the current variable.
- Can we look ahead further?

	V_1	V_2	V_3	V_4	V_5	V_6
R	O		O		X	X
B		O		O	X	X
G					$?$	$?$



At this point, it is already obvious that this branch will not lead to a solution because there are no consistent values in the remaining domain for V_5 and V_6 .

Constraint Propagation

- V = variable being assigned at the current level of the search
- Set variable V to a value in $D(V)$
- For every variable V' connected to V :
 - Remove the values in $D(V')$ that are inconsistent with the assigned variables
 - For every variable V'' connected to V' :
 - Remove the values in $D(V'')$ that are no longer possible candidates
 - And do this again with the variables connected to V''
 -until no more values can be discarded

Constraint Propagation

New: Constraint Propagation

Forward Checking as before

- V = variable being assigned at the current level of the search
- Set variable V to a value in $D(V)$
- For every variable V' connected to V :
 - Remove the values in $D(V')$ that are inconsistent with the assigned variables
 - For every variable V'' connected to V' :
 - Remove the values in $D(V'')$ that are no longer possible candidates
 - And do this again with the variables connected to V''
 -until no more values can be discarded