

## Finite-State Machines

A finite-state recognizing machine is described by:

- A finite set of states
- A finite set of input symbols (input alphabet)
- A transition function  $\delta$  which assigns a new state to every combination of state and input
- A subset of states designated as accepting states
- A state designated as the starting state

The transition function  $\delta$  defines a new state  $S_{new}$  in terms of a current state  $S_{old}$  and a current input symbol  $x$ .

A finite-state processing machine is a finite-state recognizing machine that exits to a specified routine on reaching the end of its input. Extend the input alphabet of the machine that processes the language by the end-marker symbol  $\dagger$ , (note: the input alphabet of the machine that recognizes the language remains unchanged).

## State Equivalence

State  $S$  in finite-state recognizer  $M$  is equivalent to state  $T$  in finite-state recognizer  $N$  if and only if machine  $M$  starting in state  $S$  will accept exactly the same sequences as machine  $N$  starting in state  $T$ .

Finite-state recognizers  $M$  and  $N$  are said to be equivalent if and only if their starting states are equivalent.

If two states are not equivalent then any sequence which causes one state to make a transition into an accepting state and the other state to go into a rejecting state is called a distinguishing sequence.

Two states are equivalent if and only if they have no distinguishing sequence.

State equivalence is:

- Reflexive — each state is equivalent to itself
- Symmetric — state  $S$  equivalent to state  $T$  implies state  $T$  equivalent to state  $S$
- Transitive — if states  $S$  and  $T$  are equivalent and states  $T$  and  $U$  are equivalent then states  $S$  and  $U$  are equivalent

## Extraneous states

States in a finite-state recognizer that can never be reached by any possible input sequence when the machine is initially in its starting state are called extraneous states.

To prepare a list of non-extraneous states for any given finite-state machine:

- i. Initialize the list with the starting state.
- ii. Add to the list all states which can be reached from the starting state under single inputs.
- iii. For every new state on the list add any unlisted state which can be reached from this state.

## Reduced (Minimal) Machines

A finite-state machine is reduced if it has no extraneous states and if no two states are equivalent to each other.