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Conditionals

- For expressions, we can write a conditional if `if...then...else`
 - `exp -> if exp then exp else exp`
- The else-part is compulsory and cannot be left out
- The boolean expressed after `if` is evaluated:
 - If `true`, the value is of the expression after `then`
 - If `false`, the value is of the expression after `else`

Local Declerations

- A let-expression has the form
 - `let{d1;...dn}in e`
 - d_i are declarations, e is an expression
- Scope of each d_i is e and the righthand side of all the d_j s (mutual recursion)
- Example: $ax^2 + bx + c = 0$ means $x = \frac{-b \pm (\sqrt{b^2 - 4ac})}{2a}$

```
solve a b c
= let twoa = 2 * a
    discr = b*b - 2 * twoa * c
```

```

    droot = sqrt discr
  in ((droot-b)/twoa, negate ((droot+b)/twoa))

```

- A **where** expression has the form
 - $where\{d_1; \dots; d_n\}$
 - d_i are declarations
- Scope of each d_i is the expression that precedes **where** and righthand side of all the d_i s (mutual recursion)

```

solve a b c
= ((droot-b)/twoa, negate ((droot+b)/twoa))
where
  twoa = 2 * a
  discr = b*b - 2 * twoa * c
  droot = sqrt discr

```

Let vs Where

- What is the difference between **let** and **where**?
- The **let...in...** is a full expression and can occur anywhere an expression is expected
- There **where** keyword occurs at certain places in declarations of
 - case expressions
 - modules
 - classes
 - instances
 - function and pattern righthand sides
- Both allow mutual recursion among the declarations

Case Expression

- A case expression has the form
 - $case\ of\{p_1 \rightarrow e_1; \dots; p_n \rightarrow e_n\}$
 - p_i are patterns, e_i are expressions

```

odd x =
  case (x `mod` 2) of
    True  -> False
    False -> True

```

```

vowel x =
  case x of
    'a' -> True
    'e' -> True
    'i' -> True
    'o' -> True
    'u' -> True
    _   -> False

empty x =
  case x of
    [] -> True
    _  -> False

```

Prefix vs Infix

- Functions with identifier names are prefix
 - `myfun x y = 2*x + y`
 - However, 2-argument identifiers can be used infix-style
 - * `1 `myfun` 2`
- Functions with symbol names are infix
 - `x <+> y = 2*x - y`
 - However can be used prefix-style (`<+>`) `5 7`

Using Other Functions

- Function `even` returns true if its integer argument is even
 - `even n = n `mod` 2 == 0`
 - We use the modulo function `mod` from the Prelude
- Function `recip` calculates the reciprocal of its arguments
 - `recip n = 1/n`
 - We use the division function `/` from the Prelude
 - Function call `splitAt n xs` returns two lists, the first with the first `n` elements of `xs`, the second with the rest of the elements
 - `splitAt n xs = (take n xs, drop n xs)`
 - We use the list functions `take` and `drop` from the Prelude

Decomposing Problems

- In a very real sense, programming *is* problem decomposition
- We break a big problem down into small problems
- Solve all the small problems
- Connect the solutions to the small problems together into a solution to the bigger problem
- In a lot of languages, you can get away with a certain bad habit
 1. Start writing a solution to the big problem
 2. Keep programming - when two parts need to share data, make a piece of shared data
 3. Keep Programming - eventually end up with a solution with lots of sections that depend on the value of a variable shared with other parts
- What's wrong with this?
 - No way to track how the different parts talk to each other
 - No defined interfaces between parts
- So where someone tries to modify the code, they need to keep *the entire structure of the application in their head* (spaghetti code)
- In Haskell, this is impossible
 - No mutation - shared variables can't ever change
 - It's possible to *really* program yourself into a corner and be unable to fix the code
 - The keep-going-till-it-works approach is a recipe for pain and frustration

Doing it 'right'

1. What do I have? - this is the initial type **a**
 2. What do I want? - this is the final type **b**
 3. How do I get there? - this is a function **a -> b**
 4. Implement the first piece
 5. Go to 1
- At each step, there is a defined interface that the compiler will enforce - the *type* of the function
 - If a function changes, then the program will not compile until you have fixed *every* place where you call it