Web Structure Mining

Community detection and link prediction

— Session 3 —

Community Detection

Community

Community. It is formed by individuals such that those within a group interact with each other more frequently than with those outside the group, a.k.a. group, cluster, cohesive subgroup, module in different contexts

Community detection: discovering groups in a network where individuals' group memberships are not explicitly given

Why communities in social media?

- Human beings are social
- Easy-to-use social media allows people to extend their social life in unprecedented ways
- Difficult to meet friends in the physical world, but much easier to find friend online with similar interests
- Interactions between nodes can help determine communities

Community in Social Media

Two types of groups in social media

- Explicit Groups: formed by user subscriptions
- Implicit Groups: implicitly formed by social interactions

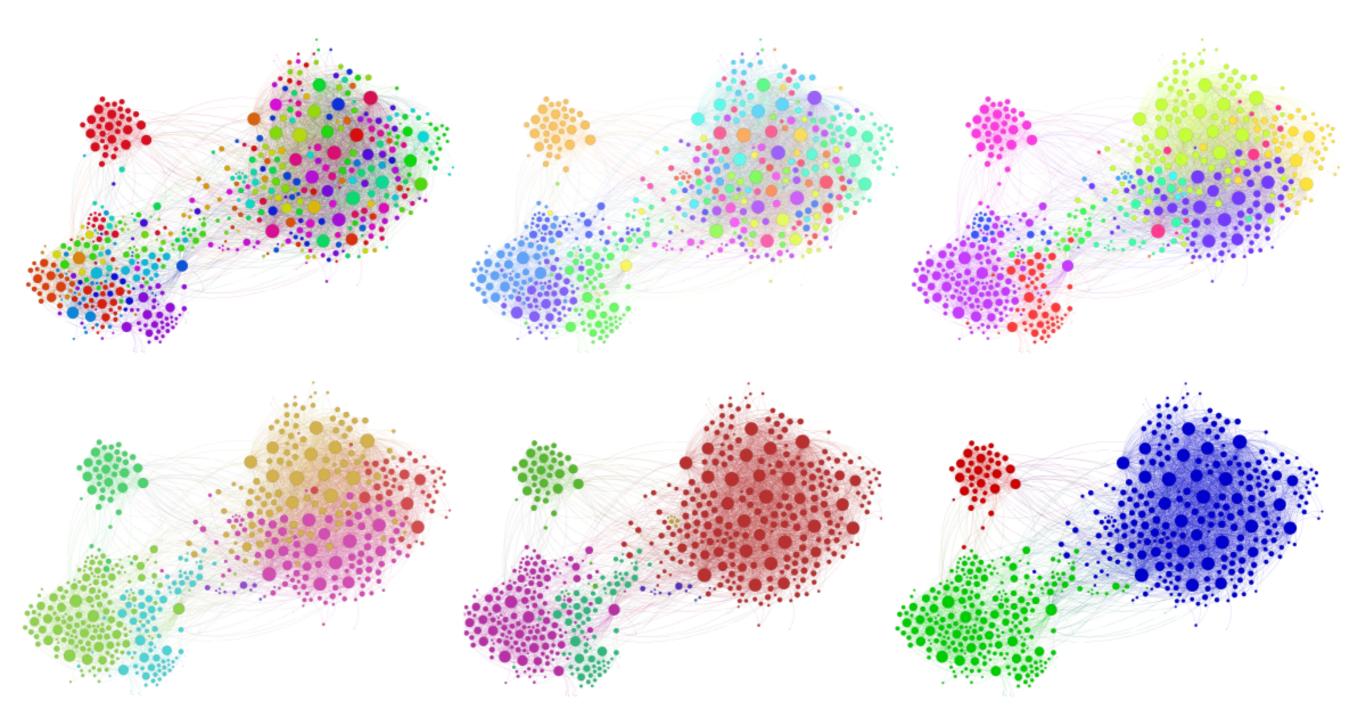
Some social media sites allow people to join groups, is it necessary to extract groups based on network topology?

- Not all sites provide community platform
- Not all people want to make effort to join groups
- Groups can change dynamically

Network interaction provides rich information about the relationship between users

- Can complement other kinds of information
- Help network visualization and navigation
- Provide basic information for other tasks

Subjectivity in Community Detection



Taxonomy of Community Criteria

Criteria vary depending on the tasks

Community detection methods can be divided into 4 categories (not exclusive):

1 - Node-Centric Community

Each node in a group satisfies certain properties

2 - Group-Centric Community

Consider the connections within a group as a whole. The group has to satisfy certain properties without zooming into node-level

3 - Network-Centric Community

Partition the whole network into several disjoint sets

4 - Hierarchy-Centric Community

Construct a hierarchical structure of communities

Node-Centric Community Detection

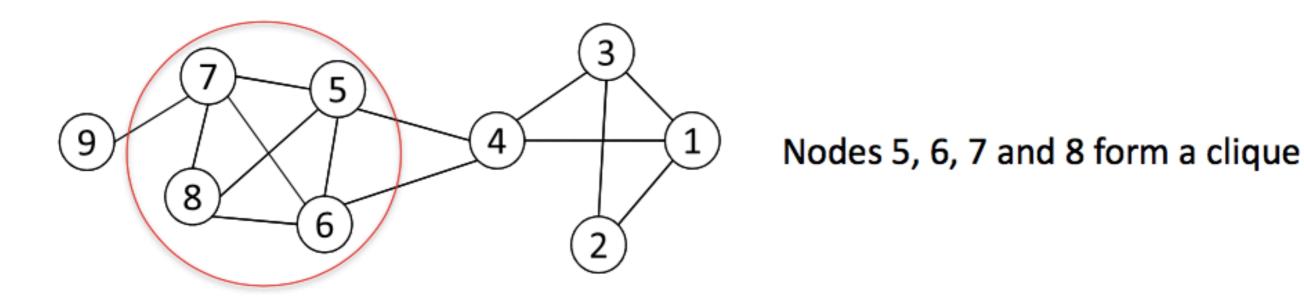
Nodes satisfy different properties

- Complete Mutuality
 - Cliques
- Reachability of members
 - k-clique, k-clan, k-club
- Nodal degrees
 - k-plex, k-core
- Relative frequency of Within-Outside Ties
 - LS sets, Lambda sets

Commonly used in traditional social network analysis Here, we discuss some representative ones

Complete Mutuality: Cliques

A clique is a complete maximal subgraph



NP-hard to find the maximum clique in a network

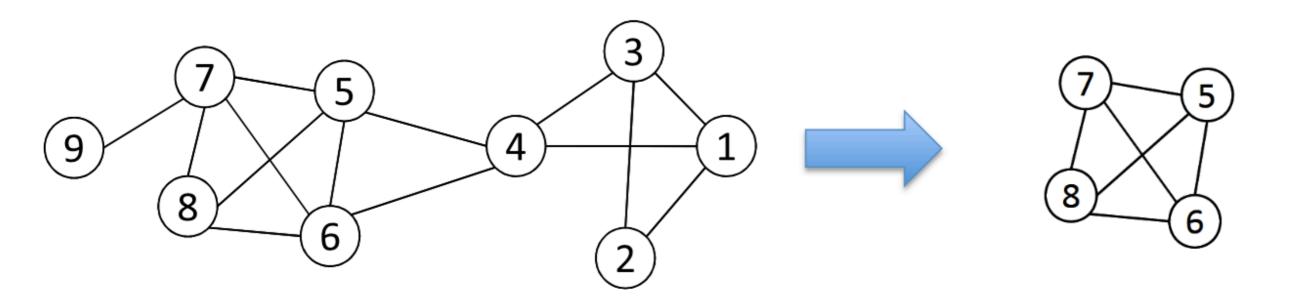
Straightforward implementation to find cliques is very expensive in time complexity

Finding the Maximum Clique

- In a clique of size k, each node has a degree >= k-1
- Nodes with degree < k-1 will not be included in the maximum clique
- Recursive pruning procedure :
 - Sample a sub-network from the given network, and find a clique in the subnetwork, say, by a greedy approach
 - Suppose the clique above is size k, in order to find out a larger clique, all nodes with degree <= k-1 should be removed
- Repeat until the network is small enough
- In social media, many nodes are removed as social networks follow a power law distribution for node degrees

Maximum Clique

Example



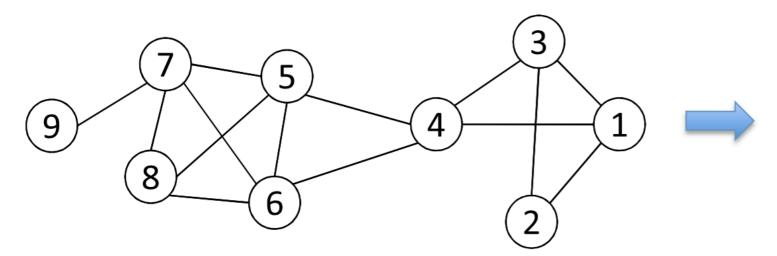
- Suppose we sample a sub-network with nodes 1 to 5 and find a 3-clique {1,2,3}
- In order to find a clique > 3, remove nodes with degree <= 2

Clique Percolation Method (CPM)

- Clique is a very strict definition, unstable
- Normally use cliques as a core or a seed to find larger communities
- CPM is such a method to find overlapping communities
 - Input
 - A parameter k, and a network
 - Procedure
 - Find out all cliques of size k in a given network
 - Construct a clique graph. Two cliques are adjacent if they share k-1 nodes
 - Each connected components in the clique graph form a community

CPM

Example



Cliques of size 3:

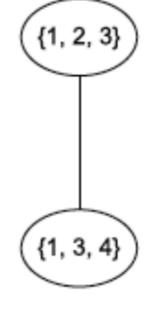
{1, 2, 3}, {1, 3, 4}, {4, 5, 6}, {5, 6, 7}, {5, 6, 8}, {5, 7, 8}, {6, 7, 8}

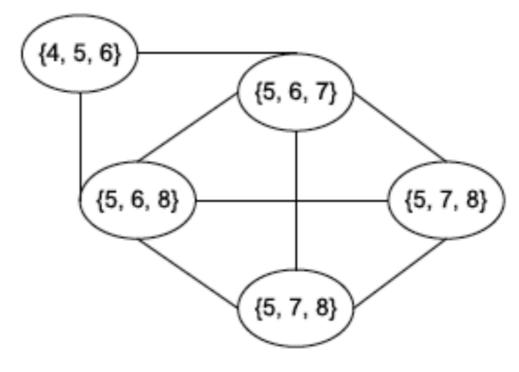


Communities:

{1, 2, 3, 4} {4, 5, 6, 7, 8}







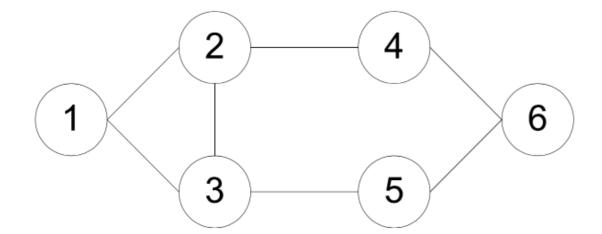
Reachability

k-clique and k-club

Any node in a group should be reachable in k hops

k-clique: a maximal subgraph in which the largest geodesic distance between any nodes <= k

k-club: a substructure of diameter <= k



Cliques: {1, 2, 3}

2-cliques: {1, 2, 3, 4, 5}, {2, 3, 4, 5, 6}

2-clubs: {1,2,3,4}, {1, 2, 3, 5}, {2, 3, 4, 5, 6}

A k-clique might have diameter larger than k in the subgraph

Commonly used in traditional SNA

Often involves combinatorial optimization

Group-Centric Community Detection

Density-Based Groups

The group-centric criterion requires the whole group to satisfy a certain condition E.g., the group density >= a given threshold

A subgraph $G_s(V_s, E_s)$ is a γ -dense quasi-clique if:

$$\frac{|E_s|}{|V_s|(|V_s|-1)/2} \ge \gamma$$

A similar strategy to that of cliques can be used

- Sample a subgraph, and find a maximal γ -dense quasi-clique (say, of size k)
- Remove nodes with degree $< k\gamma$

Network-Centric Community Detection

Network-centric criterion needs to consider the connections within a network globally

Goal: partition nodes of a network into disjoint sets

Approaches:

- Clustering based on vertex similarity
- Latent space models
- Block model approximation
- Spectral clustering
- Modularity maximization

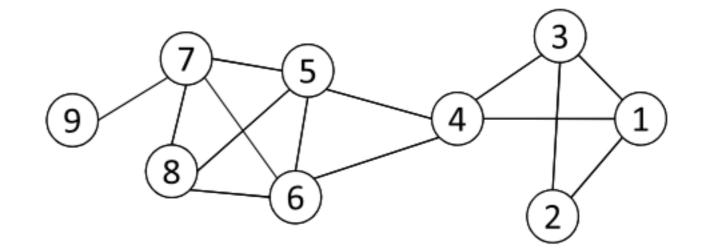
Clustering Based on Vertex Similarity

Apply k-means or similarity-based clustering to nodes

Vertex similarity is defined in terms of the similarity of their neighborhood

Structural equivalence: two nodes are structurally equivalent iff they are connecting to the same set of actors

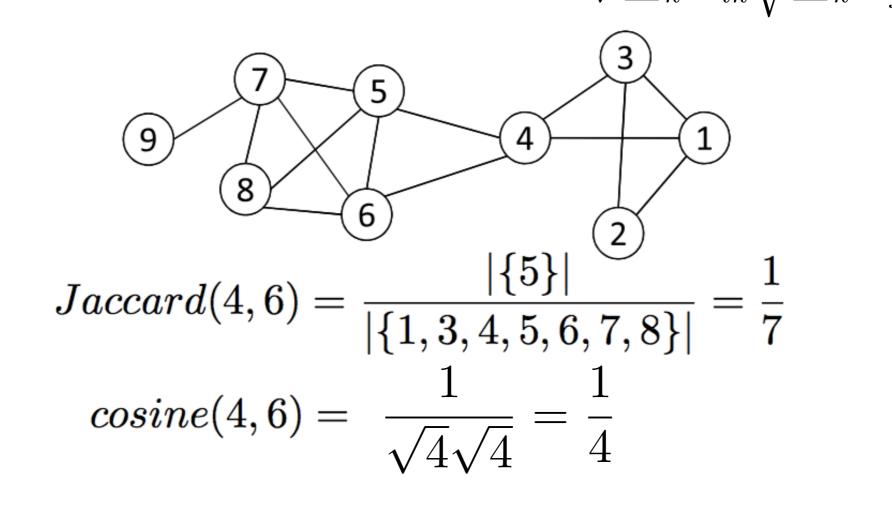
Nodes 1 and 3 are structurally equivalent; So are nodes 5 and 7.



Structural equivalence is too restrict for practical use

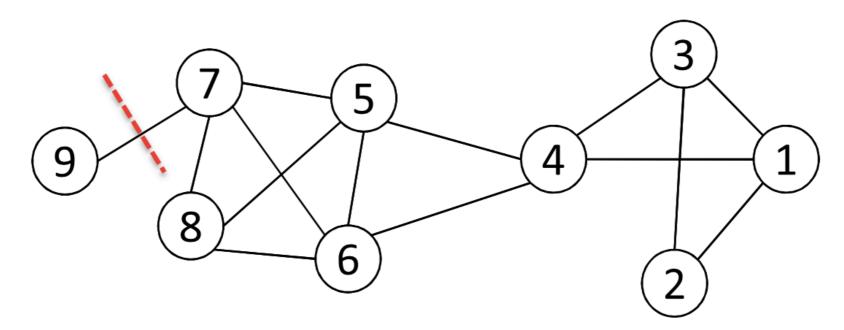
Vertex Similarity

Jaccard Similarity
$$Jaccard(v_i,v_j)=rac{|N_i\cup N_j|}{|N_i\cap N_j|}$$
 Cosine similarity $cosine(v_i,v_j)=rac{\sum_k A_{ik}A_{jk}}{\sqrt{\sum_k A_{ik}^2}\sqrt{\sum_k A_{jk}^2}}$

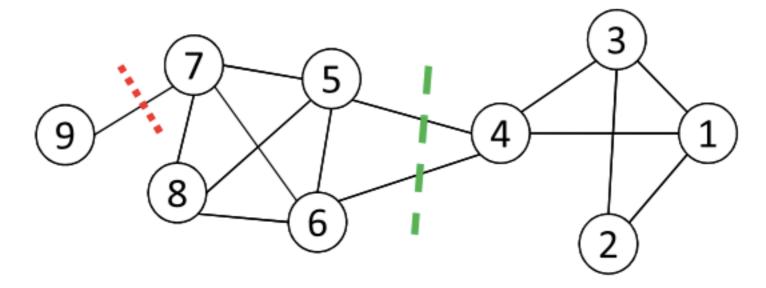


Cut

- Most interactions are within group whereas interactions between groups are few
- Community detection: minimum cut problem
- Cut: A partition of ver-ces of a graph into two disjoint sets
- Minimum cut problem: find a graph partition such that the number of edges between the two sets is minimized



Ratio Cut and Normalized Cut



- Minimum cut often returns an unbalanced partition, with one set being a singleton
- Change the objective function to take the size of the communities into account

Ratio
$$\operatorname{Cut}(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{\operatorname{cut}(C_i, \bar{C}_i)}{|C_i|},$$

Normalized
$$\operatorname{Cut}(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{\operatorname{cut}(C_i, \bar{C}_i)}{\operatorname{vol}(C_i)}$$

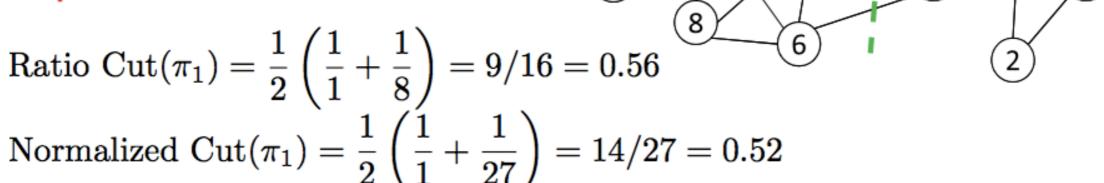
C_{i,}: a community

|C_i|: number of nodes in C_i vol(C_i): sum of degrees in C_i

Ratio Cut and Normalized Cut

Example

For partition in red: π_1



For partition in green: π_2

Ratio
$$Cut(\pi_2) = \frac{1}{2} \left(\frac{2}{4} + \frac{2}{5} \right) = 9/20 = 0.45 < Ratio $Cut(\pi_1)$
Normalized $Cut(\pi_2) = \frac{1}{2} \left(\frac{2}{12} + \frac{2}{16} \right) = 7/48 = 0.15 < Normalized $Cut(\pi_1)$$$$

Both ratio cut and normalized cut prefer a balanced partition

Hierarchy-Centric Community Detection

Goal: build a hierarchical structure of communities based on network topology

Allow the analysis of a network at different resolutions

Representative approaches:

- Divisive Hierarchical Clustering
- Agglomerative Hierarchical clustering

Divisive Hierarchical Clustering

Divisive clustering

- Partition nodes into several sets
- Each set is further divided into smaller ones
- Network-centric partition can be applied for the partition

One particular example: recursively remove the "weakest" tie

- Find the edge with the least strength
- Remove the edge and update the corresponding strength of each edge

Recursively apply the above two steps until a network is discomposed into desired number of connected components.

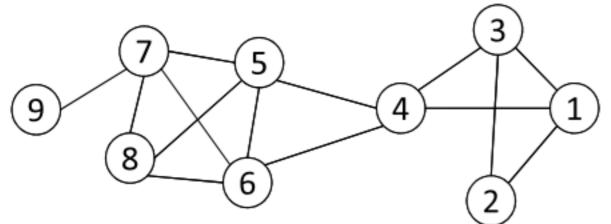
Each component forms a community

Edge Betweenness

The strength of a tie can be measured by edge betweenness

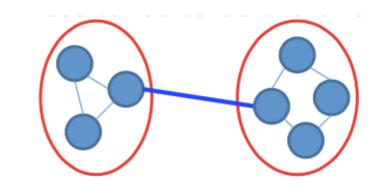
Edge betweenness: the number of shortest paths that pass along with the edge

edge-betweenness(e) =
$$\Sigma_{s < t} \frac{\sigma_{st}(e)}{\sigma_{s,t}}$$

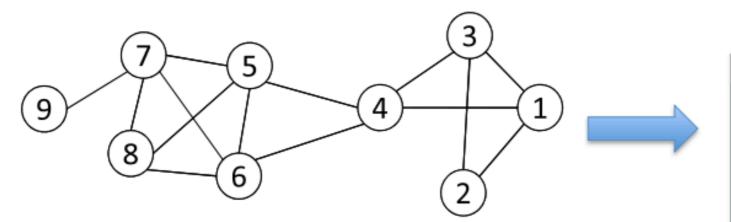


The edge betweenness of e(1, 2) is 4, as all the shortest paths from 2 to {4, 5, 6, 7, 8, 9} have to either pass e(1, 2) or e(2, 3), and e(1,2) is the shortest path between 1 and 2

The edge with the highest betweenness tends to be a bridge between 2 communities

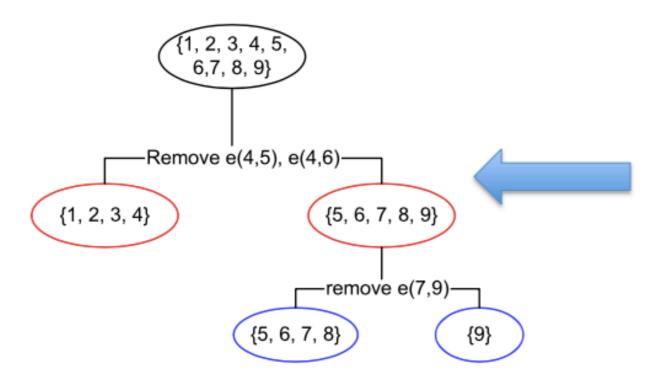


Divisive Clustering based on Edge Betweenness



Initial betweenness value

Table 3.3: Edge Betweenness									
	1	2	3	4	5	6	7	8	9
1	0	4	1	9	0	0	0	0	0
2	4	0	4	0	0	0	0	0	0
3	1	4	0	9	0	0	0	0	0
4	9	0	9	0	10	10	0	0	0
5	0	0	0	10	0	1	6	3	0
6	0	0	0	10	1	0	6	3	0
7	0	0	0	0	6	6	0	2	8
8	0	0	0	0	3	3	2	0	0
9	0	0	0	0	0	0	8	0	0



After remove e(4,5), the betweenness of e(4, 6) becomes 20, which is the highest;

After remove e(4,6), the edge e(7,9) has the highest betweenness value 4, and should be removed.

Community Evaluation

Evaluating Community Detection

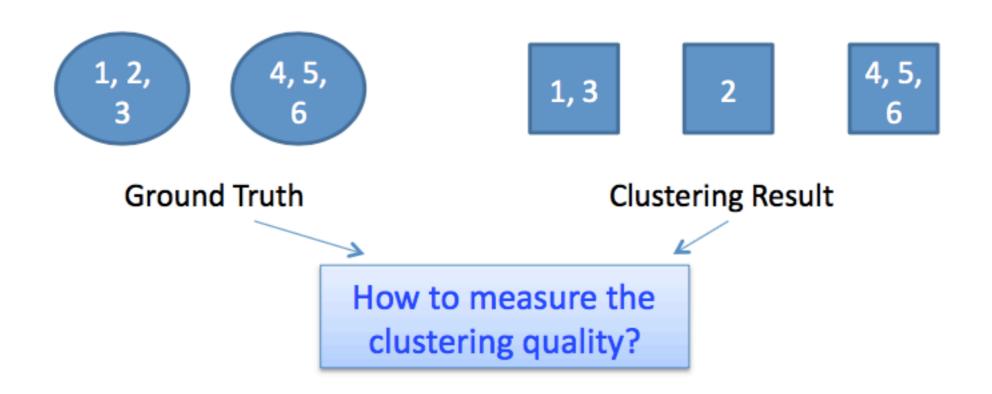
For groups with a clear and formal definition

- E.g., cliques, k-cliques, k-clubs, ...
- Verify if the extracted communities satisfy the definition

For networks with ground truth information

- Normalized Mutual Information
- Accuracy of pairwise community memberships

Measuring a Clustering Result



- The number of communities after grouping can be different from the ground truth
- No clear community correspondance between clustering result and the ground truth
- Normalized Mutual Information can be used

Normalized Mutual Information

Entropy

The information contained in a distribution

$$H(X) = \sum_{x \in X} p(x) \log p(x)$$

Mutual Information

The shared information between two distributions

$$I(X;Y) = \sum_{y \in Y} \sum_{x \in X} p(x,y) \log \left(\frac{p(x,y)}{p_1(x)p_2(y)} \right)$$

Normalized Mutual Information (between 0 and 1)

$$NMI(X;Y) = \frac{I(X;Y)}{\sqrt{H(X)H(Y)}}$$

Consider a partition as a distribution (probability of one node falling into one community), we can compute the matching between two clusterings

Normalized Mutual Information

$$H(X) = \sum_{x \in X} p(x) \log p(x)$$

$$H(\pi^b) = \sum_{\ell=0}^{k^{(a)}} \frac{n_h^a}{n} \log(\frac{n_h^a}{n})$$

$$H(\pi^b) = \sum_{\ell=0}^{k^{(b)}} \frac{n_\ell^b}{n} \log(\frac{n_\ell^b}{n})$$

$$I(X;Y) = \sum_{y \in Y} \sum_{x \in X} p(x,y) \log \left(\frac{p(x,y)}{p_1(x)p_2(y)} \right) \Longrightarrow I(\pi^a,\pi^b) = \sum_h \sum_\ell \frac{n_{h,\ell}}{n} \log \left(\frac{\frac{n_{h,\ell}}{n}}{\frac{n_h^a}{n} \frac{n_\ell^b}{n}} \right)$$

$$NMI(X;Y) = \frac{I(X;Y)}{\sqrt{H(X)H(Y)}}$$

$$NMI(\pi^{a}, \pi^{b}) = \frac{\sum_{h=1}^{k^{(a)}} \sum_{\ell=1}^{k^{(b)}} n_{h,\ell} \log \left(\frac{n \cdot n_{h,l}}{n_{h}^{(a)} \cdot n_{\ell}^{(b)}}\right)}{\sqrt{\left(\sum_{h=1}^{k^{(a)}} n_{h}^{(a)} \log \frac{n_{h}^{a}}{n}\right) \left(\sum_{\ell=1}^{k^{(b)}} n_{\ell}^{(b)} \log \frac{n_{\ell}^{b}}{n}\right)}}$$

Normalized Mutual Information

Example

Partition a: [1, 1, 1, 2, 2, 2]

Partition b: [1, 2, 1, 3, 3, 3]

1, 2, 3	4, 5, 6
1,3 2	4, 5,6

$$n = 6$$

$$k^{(a)} = 2$$

$$k^{(b)} = 3$$

	n_h^a
h=1	3
h=2	3

	n_l^b
l=1	2
l=2	1
l=3	3

$$n_{h,l}$$
 $l=1$
 $l=2$
 $l=3$

 h=1
 2
 1
 0

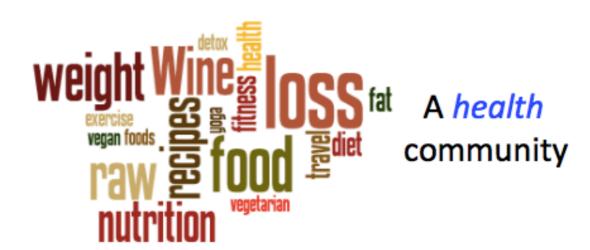
 h=2
 0
 0
 3

$$NMI(\pi^{a}, \pi^{b}) = \frac{\sum_{h=1}^{k^{(a)}} \sum_{\ell=1}^{k^{(b)}} n_{h,\ell} \log \left(\frac{n \cdot n_{h,l}}{n_{h}^{(a)} \cdot n_{\ell}^{(b)}} \right)}{\sqrt{\left(\sum_{h=1}^{k^{(a)}} n_{h}^{(a)} \log \frac{n_{h}^{a}}{n}\right) \left(\sum_{\ell=1}^{k^{(b)}} n_{\ell}^{(b)} \log \frac{n_{\ell}^{b}}{n}\right)}} \quad = 0.8278$$

Evaluation with Semantics

- For networks with semantics
 - Networks come with semantic or attribute information of nodes or connections
 - Human subjects can check whether the extracted communities are coherent and homogeneous
- Evaluation is qualitative
- It is intuitive and helps in understanding a community





Evaluation without Ground Truth

- For networks without ground truth or semantic information
- This is the most common situation
- A option is to resort cross-validation
 - Extract communities from a (training) network
 - Evaluate the quality of the community detection on a network constructed from a different date or based on a related type of interaction

Quantitative evaluation

- Modularity
- Block model approximation error

Link Prediction

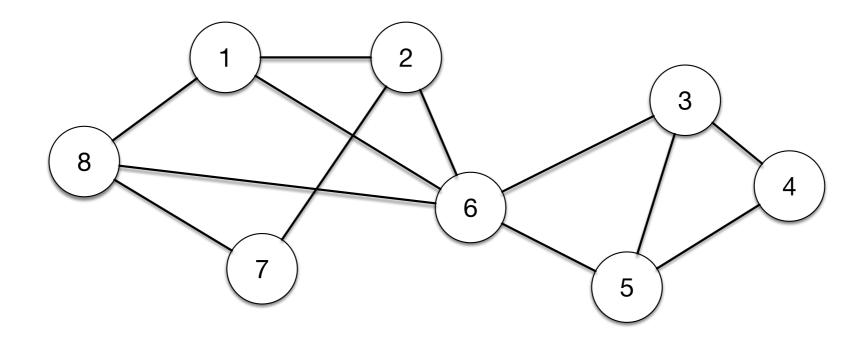
Outline

- 1. Link prediction problem
- 2. Proximity measures
- 3. Prediction by supervised learning

Link prediction

- Link prediction. Given a snapshot of a dynamic network at time t, predict edges added in the interval (t,t')
- Link completion. Given a network, infer links that are consistent with the structure, but missing
- Link reliability. Estimate the reliability of given links in the network
- What to predict?
 - Link existence
 - Link weight
 - Link type

Link prediction



- Number of missing edges = |V| (|V| 1)/2 |E|
- In sparse graphs, |E| << |V|²</p>
- Probability of correct random guess O(1/|V|²)

Scoring algorithm

- Link prediction by proximity scoring
 - 1. For each pair of nodes compute proximity score c(v,v')
 - 2. Sort all pairs by the decreasing score
 - 3. Select top n pairs (or above some threshold) as new links

Many metrics have been summarised in :

David Liben-Nowell and Jon Kleinberg. 2007. The link-prediction problem for social networks. J. Am. Soc. Inf. Sci. Technol. 58, 7 (May 2007), 1019-1031.

- Based on the local neighbourhood of v_i and v_j
 - Number of common neighbours

$$|\mathcal{N}(v_i) \cap \mathcal{N}(v_j)|$$

Jaccard's coefficient

$$\frac{|\mathcal{N}(v_i) \cap \mathcal{N}(v_j)|}{|\mathcal{N}(v_i) \cup \mathcal{N}(v_j)|}$$

Adamic / Adar

$$\sum_{v \in \mathcal{N}(v_i) \cap \mathcal{N}(v_i)} \frac{1}{\log |\mathcal{N}(v)|}$$

- Based on paths and ensemble of paths between v_i and v_j
 - Shortest path

$$-min\{path_{ij} > 0\}$$

Katz score

$$\sum_{l=1}^{\infty} \beta^{(l)} |paths_{ij}^{(l)}|$$

Personalized (rooted) PageRank

$$PR = \alpha (D^{-1}A)^T PR + (1 - \alpha)$$

- Expected number of random walk steps:
 - ightharpoonup Hitting time: $-H_{ij}$
 - Commute time: $-(H_{ij}+H_{ji})$
 - Normalized hitting / commute time: $-(H_{ij}\pi_j + H_{ji}\pi_i)$ with π_i (resp. π_j) be the stationary probability of v_i (resp. v_j)

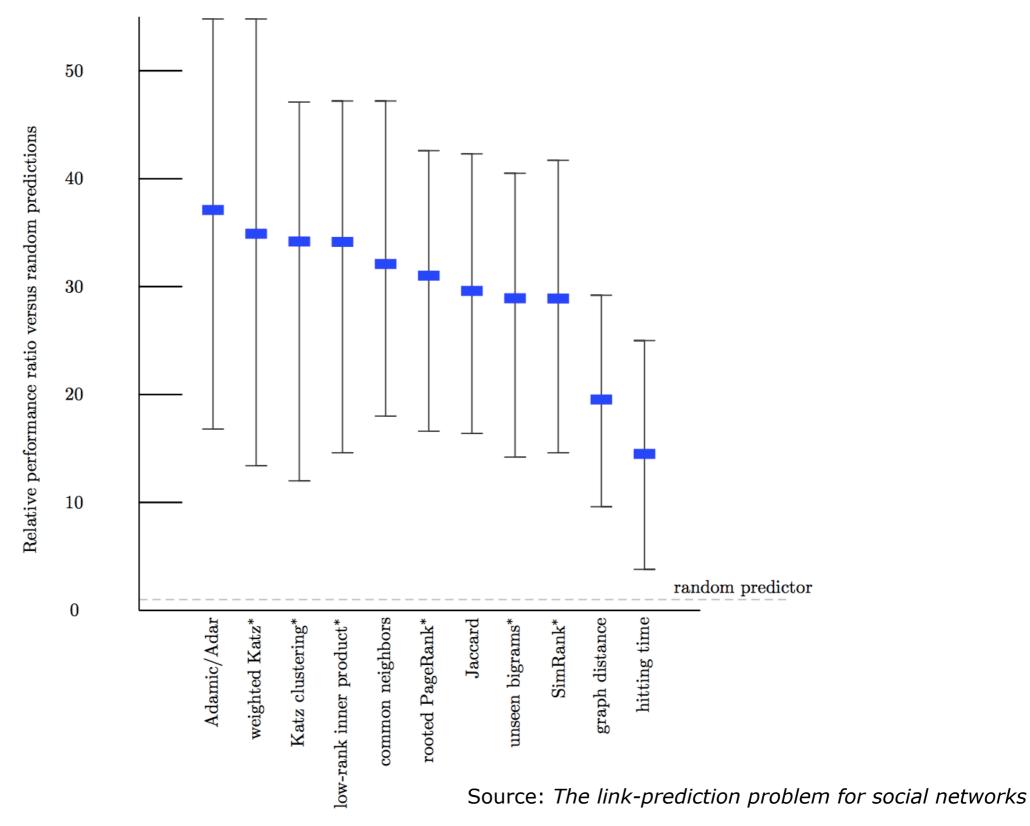
SimRank:

$$SimRank(v_i, v_j) = \gamma \cdot \frac{\sum_{a \in \mathcal{N}_i} \sum_{b \in \mathcal{N}_j} SimRank(a, b)}{|\mathcal{N}_i| \cdot |\mathcal{N}_j|}$$

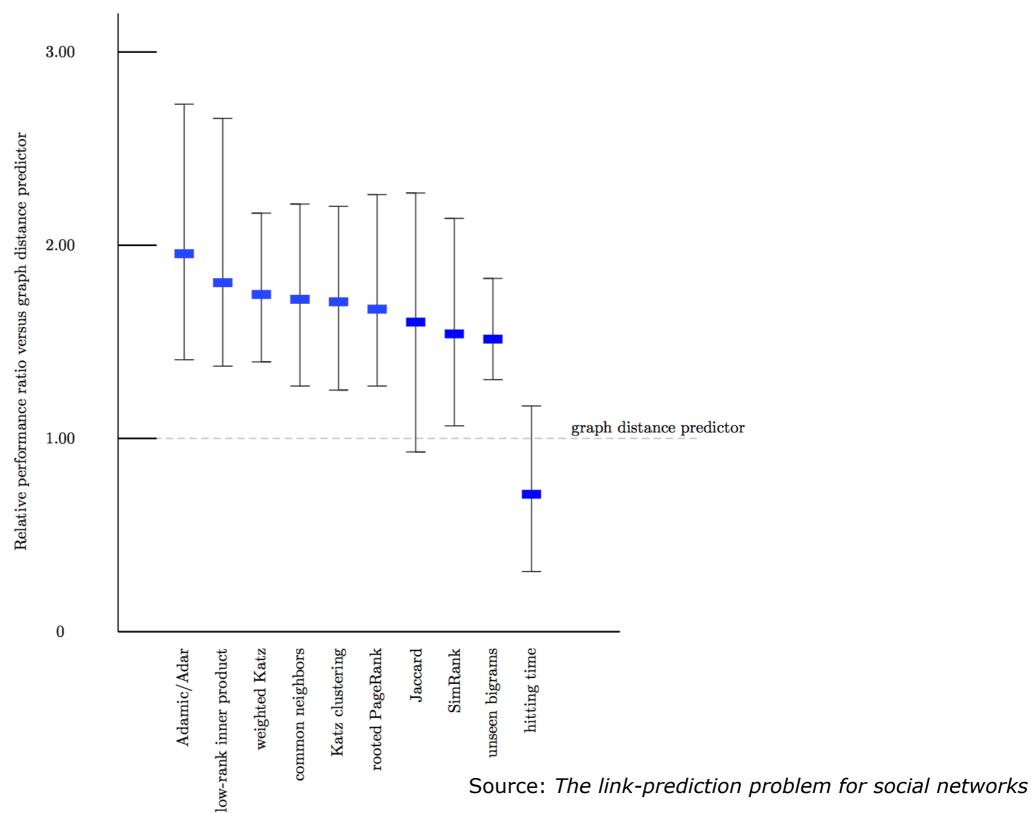
- Preferential attachment (2 alternative versions)
 - $k_i \cdot k_j = |\mathcal{N}_i| \cdot |\mathcal{N}_j|$
 - $k_i + k_j = |\mathcal{N}_i| + |\mathcal{N}_j|$

- Clustering coefficient
 - $CC(v_i) \cdot CC(v_j)$
 - $CC(v_i) + CC(v_j)$

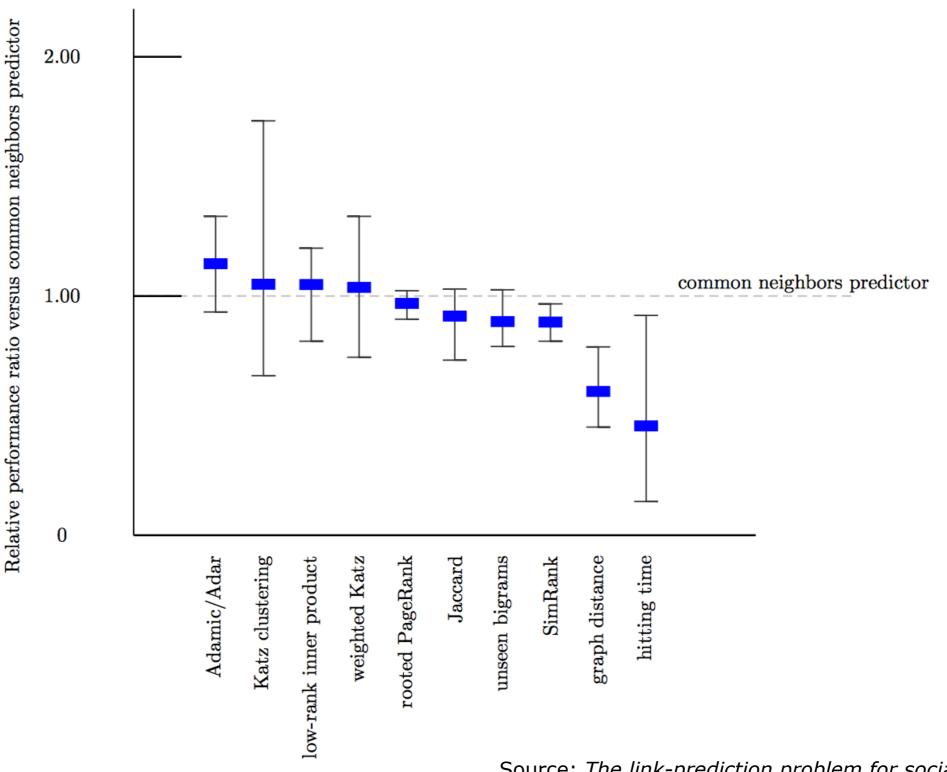




Some results



Some results



Take away message

 Node-based topological similarity measures (common neighbours, Jaccard, Adamic/Adar, preferential attachment) perform the best but does not scale well

 Path-based topological similarity measure (Katz, Hitting time, rooted PageRank) have to be preferred when dealing with relatively big networks (>10K vertices)

Binary classification

- A challenging classification problem:
 - A very large number of possible edges (quadratic in number of nodes)
 - Highly unbalanced class distribution
 - Positive examples: linear growth with number of nodes
 - Negative example : quadratic growth with number of nodes

A very challenging problem

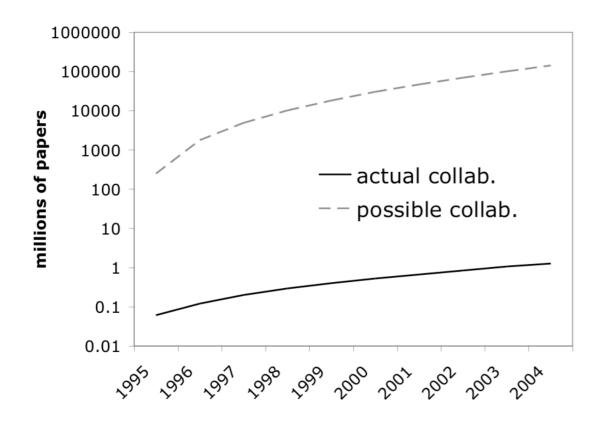


Figure 1. Logarithmic plot of actual and possible collaborations between DBLP authors, 1995-2004.

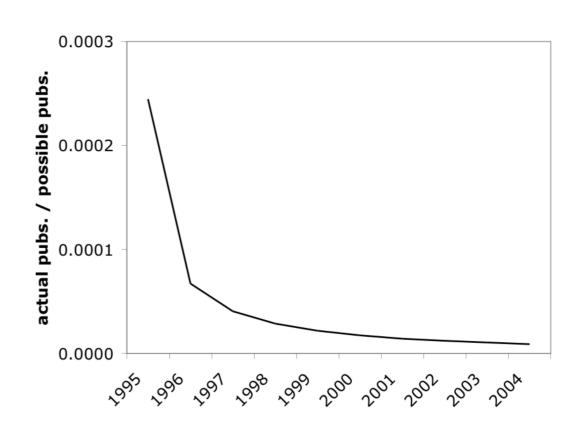


Figure 2. Publications of DBLP authors as a proportion of possible collaborations, 1995-2004.

Source: M. Rattigan, D. Jensen. The case for anomalous link discovery. ACM SIGKDD Explorations Newsletter. v 7, n 2, pp 41-47, 2005

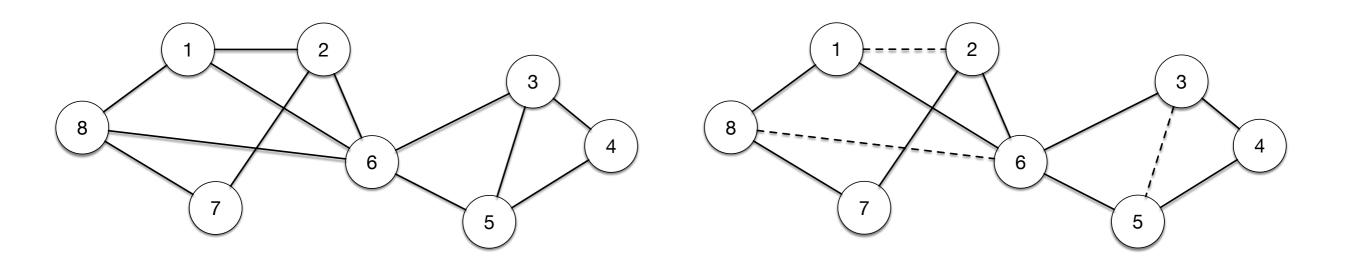
Link prediction by supervised learning

- Supervised learning process
 - 1. Feature generation
 - 2. Model training
 - 3. Testing

- Features
 - Topological proximity features
 - Aggregated features
 - Content based node proximity features

Evaluation

Simple « hold out set » evaluation



Whole graph

Training graph

More sophisticated evaluation method is preferable (cross-validation)

Evaluation metrics

Precision, recall, F-measure

$$Precision = \frac{TP}{TP + FP}, \ Recall = \frac{TP}{TP + FN}$$

$$F = \frac{2.Precision.Recall}{Precision + Recall}$$

True rate positive (TPR), False positive rate (FPR), ROC curve, AUC

$$TPR = \frac{TP}{TP + FN}, \ FPR = \frac{FP}{FP + TN}$$