

A Brief Compendium of GPU-enabled Numerical Libraries

CSCS-USI Autumn School, 15.09.2013

Dr. William Sawyer, Dr. Karl Rupp, Dr. Michael Heroux

Tutorial Summary

- Introduction and library overview
- Typical numerical and semi-numerical problems
- Survey of some gpu-enabled libraries:
 - MAGMA
 - Thrust / CuSP
 - ViennaCL
 - Paralution
 - PETSc
 - Trilinos
- Case studies: Thrust, ViennaCL





Poll

- What programming languages do you work with?
- What parallel programming paradigms?
- What problems do you want to solve ?
- What algorithms do you want to use?
- What libraries, if any, do you use?



Objectives of this tutorial

- Awareness of the available libraries
- Realization that it is not necessary to "recreate the wheel"
- Bonus: some hands-on experiences



Problems you might like to solve

- Partial differential equations
- Dense systems of linear equations
- Sparse systems of linear equations
- Preconditioning of large systems
- Eigenvalue / singular value decompositions of sparse/ dense matrices
- Partitioning large graphs
- Non-linear systems and optimization
- •





Typical dense linear algebra operations

Cholesky factorization:

$$A = A^{T} = LL^{T}$$
 $i < j \Rightarrow L_{i,j} = 0$

QR factorization:

$$A = QR$$
 $Q^{T}Q = I$ $i > j \Rightarrow R_{i,j} = 0$

LU factorization:

$$A = P^T L U \qquad P^T P = I$$

Forward/back-substitution:

$$Ax = y \Rightarrow LUx = y \Rightarrow w = L^{-1}y \Rightarrow x = R^{-1}w$$

- Eigenvalue decomposition: $Ax = \lambda x \Rightarrow A = QDQ^T$
- Generalized eigen-problem: $Ax = \lambda Bx$
- Singular value decomposition: $A = U\Sigma V^T$ $U^TU = I$ $V^TV = I$



MAGMA: Matrix Algebra on GPU and Multicore Architecture

- MAGMA: Matrix Algebra on GPU and Multicore Architectures http://icl.cs.utk.edu/magma
- Soon: D-PLASMA, D-MAGMA for distributed memory platforms
- See subsequent tutorial by Stan Tomov!



Thrust: Standard Template Library for GPUs

- A library of parallel algorithms resembling the C++ STL
- Allows easy access/manipulation of vectors on both host (CPU) and device (GPU); based on data iterators
- Defines straightforward vector data operators, e.g., :
 - * initialize vectors
 - * exchange existing values
 - * copy one to another
 - * transform with an operator
 - * perform reductions (e.g., one-dimensional to scalar)
 - * sorting and other operators





Thrust: typical operations

Declare arrays on host or device

```
thrust::host_vector<int> H(4);
thrust::device_vector<int> D = H;
thrust::device_vector<int> Z(4, 1); // All ones
```

Initialize arrays, use iterators

```
thrust::sequence(H.begin(), H.end()); // H = (0,1,2,3)
thrust::fill(D.begin(), D.end(), 2); // Fill with twos
```

Transform arrays

```
thrust::transform(D.begin(), D.end(), Z.begin(), thrust::negate<int>());
thrust::replace(H.begin(), H.end(), 2, -2);
```

Perform a reduction

```
int sum = thrust::reduce(D.begin(), D.end(), (int) 0, thrust::plus<int>());
```

Sort array

```
thrust::sort(H, H + 4);
```







Thrust: simple manipulations

```
#include <thrust/host vector.h>
#include <thrust/device vector.h>
#include <thrust/copy.h>
#include <thrust/fill.h>
#include <iostream>
int main(void)
  // initialize all ten integers of a device vector to 1
  thrust::device vector<int> D(10, 1);
  // initialize a host vector with the first five elements of D
  thrust::host vector<int> H(D.begin(), D.begin() + 5);
  // set the first seven elements of a vector to 9
 thrust::fill(D.begin(), D.begin() + 7, 9);
  // copy all of H back to the beginning of D
 thrust::copy(H.begin(), H.end(), D.begin());
  // print D
  for(int i = 0; i < D.size(); i++)
    std::cout << "D[" << i << "] = " << D[i] << std::endl;
  return 0;
                      Assignment: what values are printed?
```

ETH

Sawyer / Autumn School / 15.09.2013

CSCS
Swiss National Supercomputing Centre

Thrust: reductions operations

```
#include <thrust/sort.h>
                 // initial value of the reduction
Sum array
                 int init = 0;
                 // binary operation used to reduce values
                 thrust::plus<int> binary op;
                 // compute sum on the device
                 int sum = thrust::reduce(d vec.begin(), d vec.end(), init, binary op);
#include <thrust/iterator/permutation iterator.h>
// gather locations
thrust::device vector<int> map(4);
map[0] = 3; map[1] = 1; map[2] = 0; map[3] = 5;
                                                      Perform a reduction
// array to gather from
                                                          on an indirectly
thrust::device vector<int> source(6);
source[0] = 10; source[1] = 20; source[2] = 30;
                                                         addressed array
source[3] = 40; source[4] = 50; source[5] = 60;
// fuse gather with reduction:
     sum = source[map[0]] + source[map[1]] + ...
int sum = thrust::reduce(thrust::make permutation iterator(source.begin(), map.begin()),
                        thrust::make permutation iterator(source.begin(), map.end()));
                           Sawyer / Autumn School / 15.09.2013
```

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich Swiss National Supercomputing Cen

Thrust: Exercise

- Compare host and device sorting of integer vectors
- Create device version from host version (given)
- https://github.com/fomics/GPU_Libraries_2013/wiki/Thrust-Exercise
- Performance comparison



Thrust: advanced operators

Scan operations (parse and alter vector)

```
#include <thrust/scan.h>
int data[6] = {1, 0, 2, 2, 1, 3};
thrust::inclusive_scan(data, data + 6, data); // data now {1, 1, 3, 5, 6, 9}
thrust::exclusive_scan(data, data + 6, data); // data now {0, 1, 2, 5, 10, 16}
```

Iterator transformation -- bind an operator to an iterator.

```
#include <thrust/transform_iterator.h>
thrust::device_vector<int> vec(3);
vec[0] = 10; vec[1] = 20; vec[2] = 30;
thrust::device_vector<int>::iterator first = thrust::make_transform_iterator(vec.begin(), negate<int>());
thrust::device_vector<int>::iterator last = thrust::make_transform_iterator(vec.end(), negate<int>());
// first[0] returns -10, first[1] returns -20, first[2] returns -30
thrust::reduce(first, last); // returns -60 (i.e. -10 + -20 + -30)
```

• Zip operator: turns multiple input arguments into tuples

```
#include <thrust/iterator/zip_iterator.h>
thrust::device_vector<int> A(3); A[0] = 10; A[1] = 20; A[2] = 30;
thrust::device_vector<char> B(3); rust::device_vector<char> B(3);
first = thrust::make_zip_iterator(thrust::make_tuple(A.begin(), B.begin()));
last = thrust::make_zip_iterator(thrust::make_tuple(A.end(), B.end()));
thrust::maximum< tuple<int,char> > binary_op;
thrust::tuple<int,char> init = first[0];
thrust::reduce(first, last, init, binary_op); // returns tuple(30, 'z')
```





Thrust: conclusions

- Attempt to extend C++ STL functionality for host/device
- Based on CUDA, thus bound to NVIDIA GPUs
- Uses template meta-programming to find correct implementation at compile time
- Is an community, open-source project, but appears to have long-term approval from NVIDIA (bundled in SDK releases)
- Development is demand-driven by community; let your opinions / needs / usages / suggestions be known!





Linear solvers

Goal: Support the solution of linear systems, Ax=b,

particularly for *sparse*, parallel problems arising from PDE-based models.

User provides:

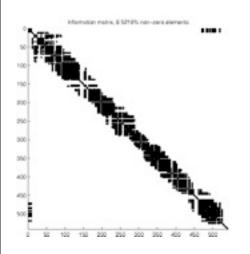
- A (matrix or operator)
- b (right-hand side)
- u (initial guess)

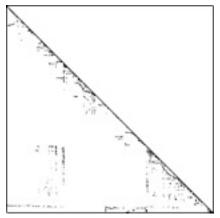


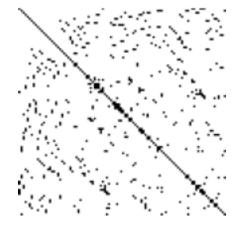


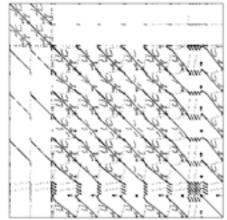
Libraries for Sparse Linear Algebra

- MAGMA limited to $m \times n$ matrices with $m,n = O(10^4)$
- Sparse matrices typically contain at least 90% zeros
- Number of non-zero (nz) elements, large: $nz = O(10^7)$
- Matrix market http://math.nist.gov/MatrixMarket/











Sawyer / Autumn School / 15.09.2013



Linear System Solution: Ax = b

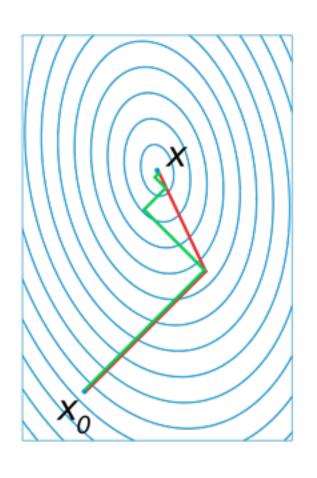
Two basic techniques:

- Direct methods, i.e. factorize matrix
 - good for multiple right hand sides
 - tend to be more robust
- Iterative methods
 - good if matrix known via operators
 - possibilities for approximate solutions





Hestenes/Stiefel, 1952: Conjugate Gradient



$$k = 0; \quad x_{0} = 0; \quad r_{0} = 0$$

$$while \quad r_{k} \neq 0 \quad \{$$

$$k = k + 1$$

$$if(k = 0) \Rightarrow p_{1} = r_{0}$$

$$if(k > 0) \Rightarrow \beta_{k} = r_{k-1}^{T} r_{k-1} / r_{k-2}^{T} r_{k-2}; \quad p_{k} = r_{k-1} + \beta_{k} p_{k-1}$$

$$\alpha_{k} = r_{k-1}^{T} r_{k-1} / p_{k}^{T} A p_{k}$$

$$x_{k} = x_{k-1} + \alpha_{k} p_{k}$$

$$r_{k} = r_{k-1} - \alpha_{k} A p_{k}$$

$$\}$$

1980's: led to a wide class of iterative Krylov subspace methods



CSCS
Swiss National Supercomputing Centre

Preconditioners: KSM alone insufficient!

- CG method initially ignored due to slow convergence
 - Theoretical convergence after 2*n steps, but n is huge
 - Convergence rate related to ratio largest/smallest eigenvalue
- Easier problem: preconditioner $M \approx A$ $Ax = b \Rightarrow M^{-1}Ax = M^{-1}b$
 - Find an approximation for A where $M^{-1}x$ is 'easily' calculated
 - Possibilities:
 - Approximate inverse known through physical description
 - Incomplete LU decomposition
 - Sparse approximative inverse (assume inverse also sparse)
 - Multilevel (multigrid) preconditioners
 - More...





CUSP: Sparse Lin. Alg. for GPUs

- CUda SParse: a templated library for GPUs and CPUs, providing a high-level interface that hides GPU complexities (NVIDIA, Apache license)
- Built on top of Thrust (NVIDIA)

```
#include <cusp/hyb_matrix.h>
#include <cusp/io/matrix_market.h>
#include <cusp/krylov/cq.h>
int main(void)
    // create an empty sparse matrix structure (HYB format)
    cusp::hyb_matrix<int, float, cusp::device_memory> A;
    // load a matrix stored in MatrixMarket format
    cusp::io::read_matrix_market_file(A, "5pt_10x10.mtx");
    // allocate storage for solution (x) and right hand side (b)
    cusp::array1d<float, cusp::device_memory> x(A.num_rows, 0);
    cusp::array1d<float, cusp::device_memory> b(A.num_rows, 1);
    // solve the linear system A * x = b with the Conjugate Gradient
method
    cusp::krylov::cq(A, x, b);
    return 0;
```



CSCS
Swiss National Supercomputing Centre

CUSP: some extensions

- CUSP implementation of sparse approximate inverse preconditioner at CSCS
- Requires least-squares minimization (QR factorization)
- GMRES solver (like CG for non-symmetric matrices)





CUSP: Some conclusions

- Template metaprogramming: conceptually easy to specify one template for different data types, e.g., single/double precision
- Only supports single node execution (multi-node implementation should be at a higher level, anyway)
- Has only CUDA backends: only for NVIDIA GPUs
- Not supported by NVIDIA! Future: uncertain
- Community effort, driven by user demand
- Location: https://github.com/cusplibrary





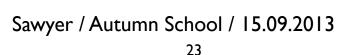
ViennaCL: Sparse Lin. Algebra on multiple platforms

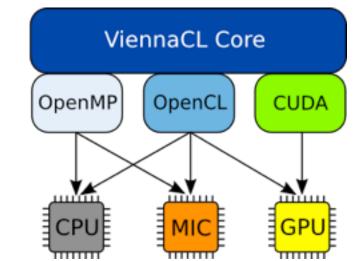
API

Backend

Hardware

- Linear algebra library for many core architectures (GPUs, Intel Xeon Phi)
- Supports BLAS I-3
- Iterative solvers
- Sparse row matrix-vector multiplication
- Goals:
 - ⇒ Simplicity, minimal dependencies
 - → Compatible with Boost.uBLAS
 - → Open source, header-only library









Swiss Federal Institute of Technology Zurich

Boost: Solve linear system

```
using namespace boost::numeric::ublas;
matrix<double> A(1000, 1000);
vector<double> x(1000), rhs(1000);
/* Fill A, x, rhs here */
// Some operations
rhs += 2.0 * x;
double val = inner prod(x, rhs);
matrix += val * outer prod(x, rhs);
x = solve(A, rhs, upper tag()); // Upper triangular solver
std::cout << " 2-norm: " << norm 2(x) << std::endl;
std::cout << "sup-norm: " << norm inf(x) << std::endl;</pre>
```



CSCS
Swiss National Supercomputing Centre

ViennaCL: Solve linear system

```
using namespace viennacl;
using namespace viennacl::linalg;
matrix<double> A(1000, 1000);
vector<double> x(1000), rhs(1000);
/* Fill A, x, rhs here */
// Some operations
rhs += 2.0 * x;
double val = inner prod(x, rhs);
matrix += val * outer prod(x, rhs);
x = solve(A, rhs, upper_tag()); // Upper triangular solver
std::cout << " 2-norm: " << norm 2(x) << std::endl;
std::cout << "sup-norm: " << norm inf(x) << std::endl;</pre>
```





ViennaCL: Memory Model

Memory buffers need to be managed differently for each of the compute backends (OpenMP, CUDA, OpenCL)

- Memory domain abstraction in class viennacl::backend::mem_handle
- Raw handles from cuda_handle(), opencl_handle() and ram_handle()
- backend is required to support:
 - memory_create(): Create a memory buffer
 - memory copy(): Copy the (partial) contents of one buffer to another
 - → memory write(): Write from a memory location in CPU RAM to the buffer
 - → memory_read(): Read from the buffer to a memory location in CPU RAM





ViennaCL: Interoperability

Standard C++ vectors and Boost uBLAS vectors can be passed to/from ViennaCL vectors:

```
std::vector<double> std_x(100)
ublas::vector<double> ublas_x(100);
viennacl::vector<double> vcl_x1, vcl_x2;

/* setup of std_x and ublas_x omitted */
viennacl::copy(std_x.begin(), std_x.end(), vcl_x1.begin());
viennacl::copy(ublas_x.begin(), ublas_x.end(), vcl_x2.begin());
```





ViennaCL: Solve sparse system

```
using namespace viennacl;
using namespace viennacl::linalg;
compressed_matrix<double> A(1000, 1000); // sparse matrix format
vector<double> x(1000), rhs(1000);
/* Fill A, x, rhs here */
x = solve(A, rhs, cg_tag()); // Conjugate Gradient solver
x = solve(A, rhs, bicgstab_tag()); // BiCGStab solve
x = solve(A, rhs, gmres tag()); // GMRES solver
```

uBLAS has no iterative solvers, but thanks to compatibility

```
using namespace boost::numeric::ublas;
using namespace viennacl::linalg;
compressed_matrix<double> A(1000, 1000);
vector<double> x(1000), rhs(1000);
/* Fill A, x, rhs here */
x = solve(A, rhs, cg_tag()); // Conjugate Gradient solver
x = solve(A, rhs, bicgstab_tag()); // BiCGStab solver
x = solve(A, rhs, gmres_tag()); // GMRES solver
```





ViennaCL: temporaries

Consider the expression

```
vec1 = vec2 + alpha * vec3 - beta * vec4;
```

With naive C++ this could be equivalent to

```
tmp1 <- alpha * vec3
tmp2 <- beta * vec4;
tmp3 <- tmp1 - tmp2;
tmp4 <- vec2 + tmp3;
vec1 <- tmp4;</pre>
```

Temporaries are costly on CPUs, even more so on GPUs

Expression templates reduce usage of temporaries





ViennaCL: expression templates

Example expression

```
vec1 += alpha * vec3;
```

Typical operator overload signatures

```
vector & vector::operator+=(vector const & T);
vector operator*(double value, vector const & v);
```

Avoid the temporary returned by operator* as follows

```
vector &
vector::operator+=(vector_expression<double, op_mult,vector> const & expr);
vector_expression<double, op_mult, vector> operator*(double value, vector const & v);
```

Implementation of operator+= calls in-place mult-add kernel

```
vector &
vector::operator+=(vector_expression<double, op_mult,vector> const & expr)
{
    inplace_mult_add(*this, expr.lhs(), expr.rhs());
}
```



Sawyer / Autumn School / 15.09.2013



ViennaCL: Exercise

- Compile and run uBLAS version of sparse matrix-vector multiply version on CPU
- Create device version from host version (given)
- https://github.com/fomics/GPU_Libraries_2013/wiki/Thrust-Exercise
- Performance comparison



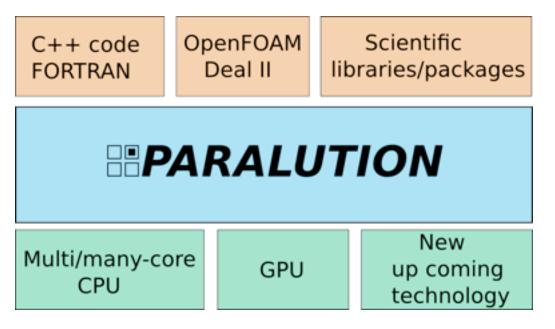
ViennaCL Conclusions

- A logical extension of Boost uBLAS template library
- Backends for OpenMP, OpenCL and CUDA
- Runs on CPUs, Intel Xeon Phi, NVIDIA + AMD GPUs
- Actively supported by ViennaCL team
- Krylov-subspace solvers, minimal spectrum of preconditioners
- Performance results are favorable
- Interoperates with other libraries:
 - Eigen
 - ▶ PETSc
 - others...
- → Promising library supporting high-level linear algebra





PARALUTION: Sparse Linear Algebra on multiple platforms



- Sparse Iterative solvers
- Preconditioners
- Will be treated in tutorial of Dimitar Lukarski





Beyond linear algebra

So perhaps single-node sparse/dense linear algebra is covered

- What about distributed memory parallelism?
- What about problems beyond linear algebra?
- Fact: there are solid MPI-GPU development efforts ongoing:
 - → D-MAGMA (see Stan's tutorial)
 - → PETSc (next slides)
 - → Trilinos (next slides)
- Fact: current multi-node GPU support for non-linear problems, optimizations, eigenvalue problems and others is provisional at best!





What is PETSc?

- Supports non-linear PDE problems
- A freely available (and supported!) research code
- Available via http://www.mcs.anl.gov/petsc
- Free for everyone, including industrial users
- Hyperlinked documentation and manual pages for all routines
- Many tutorial-style examples
- Support via email: petsc-maint@mcs.anl.gov
- Current version: 3.3 (released Jun. 5, 2012)
- Portable to any parallel system supporting MPI
- Tightly coupled systems, e.g., Cray XK6, XE6





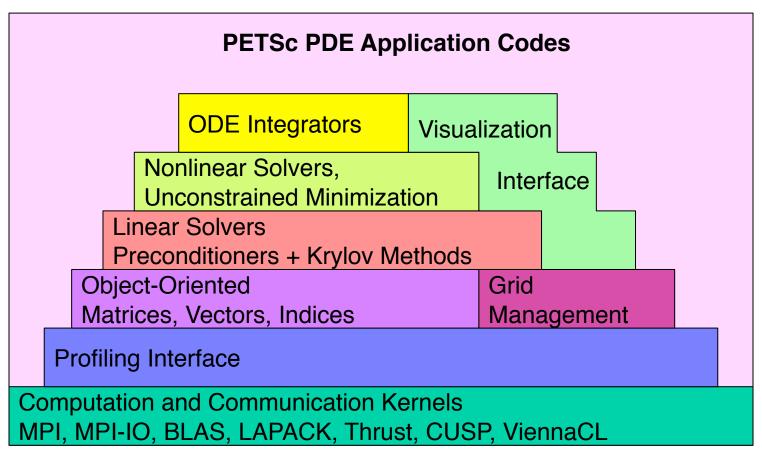
PETSc Concepts

- To specify the mathematics of the problem:
- Programmer manipulates mathematical objects (sparse matrices, nonlinear equations), algorithmic objects (solvers) and discrete geometry (meshes)
- To solve the problem:
- Solvers: linear, nonlinear, and time stepping (ODE)
- Parallel computing considerations:
- Parallel data layout, e.g., structured and unstructured meshes





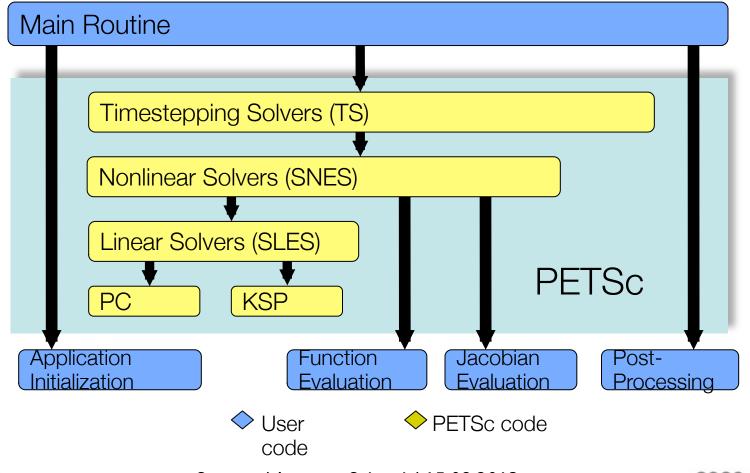
Structure of PETSc







Flow control for PDE solution





Sawyer / Autumn School / 15.09.2013



PETSc Programming Model

Goals

- Portable, runs everywhere, including heterogeneous multi-core
- Performance
- Scalable parallelism

Approach

- Distributed memory, "shared-nothing"
 - Access to data on remote machines or nodes through MPI
- Can still exploit node parallelism on each node (e.g., SMP), with limitations (see PETSc home page)
- Hide within parallel objects the details of the communication
- User orchestrates communication at a higher abstract level than message passing
- Additional classes added for GPU support



Swiss Federal Institute of Technology Zurich



PETSc Data Objects

- Vectors (Vec)
 - focus: field data arising in nonlinear PDEs
- Matrices (Mat)
 - focus: linear operators arising in nonlinear PDEs (i.e., Jacobians)

beginner

Object creation

beginner

Object assembly

intermediate

Setting options

intermediate

Viewing

advanced

User-defined customizations





PETSc Data Objects

What are PETSc vectors?

- Fundamental objects for storing field solutions, right-hand sides, etc.
- Each process locally owns a subvector of contiguously numbered global indices
- Create vectors via
 - VecCreate(...,Vec *)
 - MPI_Comm processors that share the vector
 - number of elements local to this processor
 - or total number of elements
 - VecSetType(Vec,VecType)
 - Where VecType is
 - VEC_SEQ, VEC_MPI, or VEC_SHARED

proc 0

proc 1

proc 2

proc 3

proc 4



Sawyer / Autumn School / 15.09.2013

Parallel Vector (and Matrix) Assembly

- Processors may generate any entries in vectors and matrices
- Entries need not be generated on the processor on which they ultimately will be stored
- PETSc automatically moves data during the assembly process if necessary



PETSc Communication

- MPI communicators (MPI_Comm) specify collectivity (processors involved in a computation)
- All PETSc creation routines for solver and data objects are collective with respect to a communicator, e.g.,
- VecCreate(MPI_Comm comm, int m, int M, Vec *x)
- Some operations are collective, while others are not, e.g.,
- collective: VecNorm()
- not collective: VecGetLocalSize()
- If a sequence of collective routines is used, they must be called in the same order on each processor





PETSc Vector Assembly

- VecSetValues(Vec,...)
- number of entries to insert/add
- indices of entries
- values to add
- mode: [INSERT_VALUES, ADD_VALUES]
- VecAssemblyBegin(Vec)
- VecAssemblyEnd(Vec)





PETSc Matrices

- PETSc matrices are fundamental objects for storing linear operators (e.g., Jacobians)
- Create matrices via
- MatCreate(...,Mat *)
 - MPI_Comm processors that share the matrix
 - number of local/global rows and columns
- MatSetType(Mat, MatType), where MatType is one of
 - default sparse AIJ: MPIAIJ, SEQAIJ
 - block sparse AIJ (for multi-component PDEs): MPIAIJ, SEQAIJ
 - symmetric block sparse AIJ: MPISBAIJ, SAEQSBAIJ
 - block diagonal: MPIBDIAG, SEQBDIAG
 - dense: MPIDENSE, SEQDENSE
 - matrix-free





PETSc Matrix Assembly

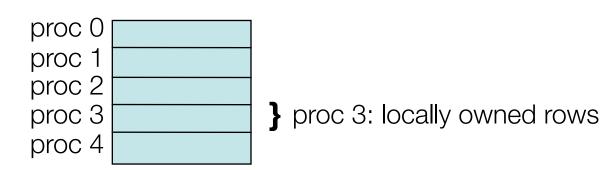
- MatSetValues(Mat,...)
- number of rows to insert/add
- indices of rows and columns
- number of columns to insert/add
- values to add
- mode: [INSERT_VALUES, ADD_VALUES]
- MatAssemblyBegin(Mat, MAT_FINAL_ASSEMBLY)
- MatAssemblyEnd(Mat, MAT_FINAL_ASSEMBLY)





PETSc Parallel Matrix Distribution

Each process locally owns a submatrix of contiguously numbered global rows.



MatGetOwnershipRange(Mat A, int *rstart, int *rend)

- rstart: first locally owned row of global matrix
- − rend −1: last locally owned row of global matrix





PETSc Linear Solvers (subset)

Krylov Methods (KSP) Preconditioners (PC)

- Conjugate Gradient
- GMRES
- CG-Squared
- Bi-CG-stab
- Transpose-free QMR
- etc.

- Block Jacobi
- Overlapping Additive Schwarz
- ICC, ILU via BlockSolve95
- ILU(k), LU (sequential only)
- etc.





PETSc Krylov Subspace Solvers

Only KSP solvers discussed here

Create a KSP solver:

```
ierr = KSPCreate(PETSC_COMM_WORLD,&ksp)
```

Use matrix to define an operator

```
ierr = KSPSetOperators(ksp,A,A,XXXX)
```

Set the solver from the command-line options

```
ierr = KSPSetFromOptions(ksp);
```

Set the solver from the command-line options

```
ierr = KSPSolve(ksp,b,u);
```





PETSc: GPU support

PETSc GPU Model

- Each MPI process has access to a single GPU, which has its own memory
- Backends for CUSP, CUSparse and ViennaCL available
- New implementations of Vec and Mat (type at run-time)
 - → Vectors with types vecsegouse, vecmpiouse, or vecouse
 - → Matrices with types matseqaijcusp, matmpiaijcusp, or mataijcusp
 - → Matrices with types matseqaijcusparse, matmpiaijcusparse, or mataijcusparse





PETSc: GPU support

Objects support both CPU and GPU copy of data, and carry flags indicated the validity of the data

| PETSC_CUDA_UNALLOCATED | MEMORY NOT ALLOCATED ON GPU |
|------------------------|--------------------------------|
| PETSC_CUDA_GPU | VALUES ON GPU ARE CURRENT |
| PETSC_CUDA_CPU | VALUES ON CPU ARE CURRENT |
| PETSC_CUDA_BOTH | VALUES ON BOTH DEVICES CURRENT |

Implementations for GPU-CPU data movement

- VecCUDACopyToGPU
- VecCUDACopyFromGPU
- . . .

these are generally used internally in solvers





PETSc: GPU support

How it works, at least conceptually

- User needs to specify types of vectors and matrices at run time
- Functionality from backend (CUSP, CUSparse, ViennaCL invoked, if available
- Implementation can be transparent to user
- Alternatively, user can program in CUDA and access device objects directly, directly call thrust:: and cusp:: operators, etc.



PETSc Summary

- PETSc library of PDE/ODE solvers
- Extensive selection of solvers, high quality, good support, free
- However: much more effective for new code
- Saddled by design choices, i.e., not thread-safe
- Monolithic: one package tries to solve all, though there are adaptors to other libraries
- MPI-GPU support in development version, in principle can be invoked at run-time. Backends for CUSP, CUSparse, ViennaCL





Trilinos: a 'pearl necklace' of packages

Object-oriented software framework for the solution of large-scale, complex multi-physics engineering and scientific problems

- http://trilinos.sandia.gov
- http://code.google.com/p/trilinos/wiki/TrilinosHandsOnTutorial





Trilinos: parallel packages (with GPU support)

- Basic linear algebra: Epetra/EpetraExt (C++), Tpetra (C++ templates)
- Preconditioners: AztecOO, Ifpack2, ML, Meros
- Iterative linear solvers: AztecOO, Belos
- Direct linear solvers: Amesos (SuperLU, UMFPACK, MUMPS, ScaLAPACK, ...)
- Non-linear / optimization solvers: NOX, MOOCHO
- Eigensolvers: Anasazi
- Mesh generation / adaptivity: Mesquite, PAMGEN
- Domain decomposition: Claps
- Partitioning / load balance: Isorropia, Zoltan2





Trilinos: Kokkos Compute Model

- How to make shared-memory programming generic:
 - →Parallel reduction is the intersection of dot() and norm1()
 - → Parallel for loop is the intersection of axpy() and mat-vec
 - →We need a way of fusing kernels with these basic constructs.
- Template meta-programming is the answer
 - → This is the same approach that Intel TBB and Thrust take
 - → Has the effect of requiring that Tpetra objects be templated on Node type.





Trilinos: Generic Parallel Constructs

Node provides generic parallel constructs, user fills in the rest **Parallel FOR:**

```
template <class WDP> void Node::parallel_for(int beg, int end, WDP workdata);
```

Work-Data Pair (WDP) struct provides:

loop body via wdp::execute(i)

Parallel REDUCE:

```
template <class WDP> WDP::ReductionType
Node::parallel_reduce(int beg, int end, WDP workdata);
```

Work-Data Pair (WDP) struct provides:

- Reduction type wdp::ReductionType
- Element generation via WDP::generate(i)
- Reduction via wdp::reduce(x,y)







Kokkos: axpy implementation

```
template <class WDP> void
Node::parallel_for(int beg, int end, WDP workdata);

template <class T> struct AxpyOp
{const T* x;
    T* y;
    T alpha, beta;
    void execute(int i) { y[i] = alpha*x[i] + beta*y[i]; }
};

AxpyOp<double> op;
op.x = ...; op.alpha = ...; op.y = ...; op.beta = ...;
node.parallel_for< AxpyOp<double> > (0, length, op);
```





Kokkos: dot product

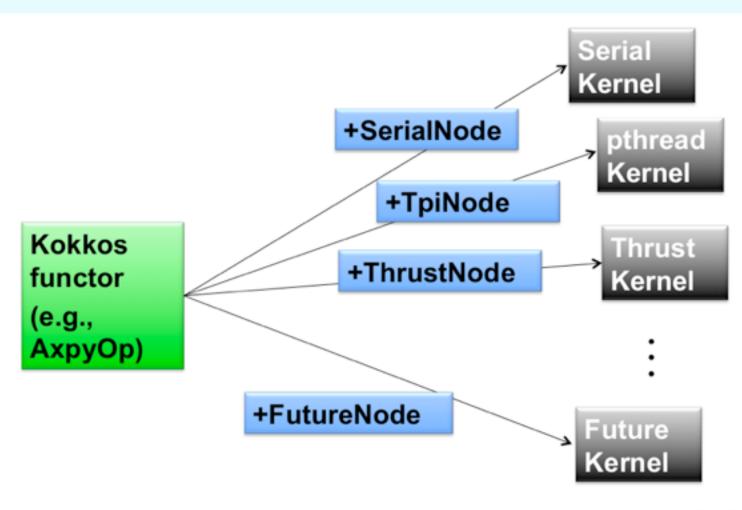
```
template <class WDP> WDP::ReductionType
Node::parallel_reduce(int beg, int end, WDP workdata);

template <class T> struct DotOp
{typedef T ReductionType;
    const T* x, y;
    T identity() { return (T)0;}T alpha, beta;
    T generate(int i) { return x[i]*y[i]; }
    T reduce(T x, T y) { return x + y; }
};

DotOp <float> op;
    op.x = ...; op.y = ...;
float dot;
dot = node.parallel_reduce < DotOp<float> > (0, length, op);
```



Kokkos: compile-time kernel specialization





Sawyer / Autumn School / 15.09.2013



Eigen- and Singular values/vectors

- Formulations:
- Non-symmetric, real:
- Symmetric, real:

$$Ax = \lambda x \Rightarrow Q^T A Q = diag(\lambda_1, ..., \lambda_n)$$



– Non-symmetric, non-defective:

$$Ax = \lambda x \Rightarrow X^{-1}AX = diag(\lambda_1, ..., \lambda_n)$$

– Generalized, symm:

$$Ax = \lambda Bx \Rightarrow Q^T AQ = diag(a_1,...,a_n)$$
 $Q^T BQ = diag(b_1,...,b_n)$

– Singular values:

$$U^T A V = diag(\sigma_1,...,\sigma_n)$$

- Interpretation: eigenvectors of A when multiplied by A are parallel to themselves
- Applications in numerous fields





Eigenvalues/vectors of large, sparse matrices

Techniques based on Lanczos iteration (symm. A)

$$r_{0} = q_{1}; \quad \beta_{0} = 1; \quad q_{0} = 0; \quad j = 0$$

$$while \quad \beta_{j} \neq 0 \quad \{$$

$$q_{j+1} = r_{j} / \beta_{j}; \quad j = j+1; \quad \alpha_{j} = q_{j}^{T} A q_{j}$$

$$r_{j} = (A - \alpha_{j} I) q_{j} - \beta_{j-1} q_{j-1}; \quad \beta_{j} = \|r_{j}\|_{2}$$

$$\}$$

- Lanczos vectors: q
- Form tridiagonal matrix T: diagonal α_j subdiagonal β_j
- Diagonalization of T is stable iterative procedure





Anasazi: design objectives

- Opaque objects: hide low level complexity
- Flexibility to allow for various linear algebra primitives;
 ease of incorporation with other frameworks
- Provide a small set of 'turn-key' eigen-solvers for large (sparse) matrices
- Provide a 'workbench' framework in which new methods can be implemented



Anasazi: classes for $Ax = \lambda Bx$

- Anasazi::Eigenproblem
- Contains components of eigen-problem
- setOperator, SetA, SetB, setPrec, setInitVec
- Anasazi::Eigensolution
- Manages the solution of the eigen-problem
- Anasazi::Eigensolver
- Defines interface which must be met by any solver
- Currently implemented solvers: BlockDavidson, BlockKrylovSchur, LOBPCG
- Anasazi::SolverManager
- 'Turn-key' class to use existing eigen-solvers





NOX: non-linear equations

Solve
$$F(x) = 0$$
 with $F(x) = \begin{pmatrix} f_1(x_1, \dots, x_n) \\ \vdots \\ f_n(x_1, \dots, x_n) \end{pmatrix}$ and $J_{i,j} = \frac{\partial F_i}{\partial x_j}(x)$

- User supplies:
- Function F(x) evaluation
- Optional: Jacobian evaluation, preconditioner
- With good guess, convergence quadratic
- Heuristics used to improve first guess
- PETSc interface available
- http://trilinos.sandia.gov/packages/nox/





Trilinos Summary

- Non-monolithic set of packages, some interoperating tightly, some loosely, some not at all
- Large development team, free software, technically advanced, latest solvers, following emerging technologies (e.g. GPUs)
- Solvers are opaque, hard to see internals, bugs can be hard to deal with
- GPU implementation through Kokkos abstraction, only NOX, Tpetra, Belos, Anasazi, Ifpack2, Zoltan2





GPU-enabled Libraries Summary

The take home message is:

Don't recreate the wheel

- But: a limited number of GPU-libraries available,
 e.g. CUxxxx (vendor), Thrust, ViennaCL, ...
- Parallel message-passing libraries in development, e.g., D-MAGMA, PETSc, Trilinos



Acknowledgments

- The PETSc team
- Trilinos team
- Roberto Croce: Autumn School organization
- FoMICS and CSCS: sponsoring Autumn School
- ... and thanks to you for attending!

