

Applied Algorithms, Summer 2017
Homework 1,
Due date: 9/7/17

1. **(16 pts.)** Use the Traditional Marriage Algorithm to find a stable pairing of the following instance with $n=4$. Boys = {A,B,C,D,E}, Girls = {V,W,X,Y,Z} and preferences list as follows. Show your work in a table.

A: V,W,X,Y,Z	V: C,D,E,A,B
B: Z,Y,X,W,V	W: A,B,C,D,E
C: Z,Y,W,X,V	X: E,D,A,C,B
D: X,Z,Y,V,W	Y: B,D,E,C,A
E: X,Z,W,Y,V	Z: A,B,C,D,E

2. **(24 pts)** Consider an input to the stable pairing problem with n men and n women. Out of the n men, there are k smart men and $n-k$ stupid men. Also, there are k smart women and $n-k$ stupid women (for some k between 1 and $n-1$). Everyone would rather marry any smart person than any stupid person (so the first k entries in any preference list are of smart people of the opposite gender in some order). Prove that in every stable pairing, every smart man is matched with a smart woman.

3.a. (12 pts.) The problem $1 \mid \sum_j C_j$ is solved optimally on an instance of n jobs. The solutions' value is z . The processing time of each of the n jobs is increased by 1. The problem $1 \mid \sum_j C_j$ is then solved on the resulting instance. What is the solution's value, as a function of n and z . Explain.

3.b. (12 pts.) For some instance of n jobs, and $m=2$ machines, the value of an optimal solution to $P \mid C_{\max}$ is z . The processing time of each of the n jobs is increased by 1. Prove or give a counter example: The value of an optimal solution for $P \mid C_{\max}$ on the resulting instance is at least $z + \lfloor n/2 \rfloor$.

4.a. (16 pts.) Find an optimal solution for $P3 \mid \sum_j C_j$ and the set of jobs of lengths {3,5,31,4,8,11,6,8,1,7}. What is C_{\max} in your schedule? What is the makespan of the schedule produced for this set of jobs and machines by list-scheduling (i) when the jobs are processed from left to right (ii) from right to left (ii) in LPT order?

4.b (20 pts.) Let $n=zm$ (n = number of jobs, m = number of machines). Prove that in the optimal solution to $P \mid \sum_j C_j$ there are exactly z jobs on each machine.

Note: The algorithm SPT uses this property, so your proof cannot simply say 'because this is what we get from SPT'.

Directions: Assume the claim is false. Therefore, there must be an optimal solution in which there is a machine $M1$ with $z+x$ jobs ($x \geq 1$) and a machine $M2$ with $z-y$ jobs ($y \geq 1$). Proceed using an exchange argument.