The core

```
Definition ["allAllocations", any[n],
 X[n] := \left\{ x \mid_{allocation[x]} \right\}
Definition ["dominion", any[\pi, n, Y],
 D[Y, \pi, n] := \left\{ x \mid_{\text{allocation } [X]} \exists (y \in Y \land \text{dominates} [y, x, \pi, n]) \right\} \right]
Definition ["dominionOn", any[\pi, n, Y],
  DOn[Y, \pi, n, Z] := \left\{ x \mid_{\text{allocation.}[x]} x \in Z \bigwedge_{y} \exists (y \in Y \land dominates[y, x, \pi, n]) \right\} \right]
Definition["Un-dominion", any[\pi, n, Y],
  U[Y, \pi, n] := X[n] \setminus D[Y, \pi, n]
Definition ["core-paper", any[\pi, n],
  \mathbb{K}[\pi,\,n] := \left\{ \mathbf{x} \mid \underset{\text{allocation } [\mathbf{x}]}{\exists} \mathbf{x}_{\mathbf{i}} > 0 \; \Leftrightarrow \; \pi[\{\mathbf{i}\},\,\mathbf{x}] \; \geq \; \pi[\mathbb{I}[n] \setminus \; \{\mathbf{i}\},\,\mathbf{x}] \right\} \right]
Definition["core", any [\pi, n],
  K[\pi, n] := U[X[n], \pi, n]
Definition ["tyrannical", any[i, n],
  t[i,n]_{j} := \begin{cases} 1 & \leftarrow j = i \\ 0 & \leftarrow \text{ otherwise } \end{cases}
Definition ["tyrannical", any[i, n],
 t[i,n] := \left\langle \left\{ \begin{array}{ll} 1 \in j = i \\ 0 \in \text{otherwise} & \left| \begin{array}{c} 1 \\ 1 \end{array} \right| \right\rangle \right]
Lemma ["tyrannicalSum", any[j, n],
  \sum_{i \in I[n]} t[j, n]_{i} = 1
Lemma ["tyrannical allocation", any [i, n], with [n \in \mathbb{N} \land i \in I[n]],
  allocation<sub>n</sub>[t[i, n]]]
Definition ["split allocation", any[i, j, n],
  s[j, k, n]_{i} := \begin{cases} \frac{1}{2} \leftarrow ((i = j) \lor (i = k)) \\ 0 \leftarrow \text{otherwise} \end{cases}
Definition "split allocation", any[k, j, n],
 s[j,k,n] := \left\langle \left\{ \begin{array}{l} \frac{1}{2} & \leftarrow ((i=j)) \lor (i=k)) \\ 0 & \leftarrow \text{ otherwise} \end{array} \right. \right. \right|_{i \in I[n]} \right\rangle \right]
Theorem ["Jordan[2006a], 2.6a", any[\pi, n],
  \forall_{\mathtt{i} \in \mathtt{I}[\mathtt{n}]} \ (\mathtt{t}[\mathtt{i},\mathtt{n}] \ \in \ \mathtt{K}[\pi,\mathtt{n}]) \ \Leftrightarrow \ \pi[\{\mathtt{i}\},\mathtt{t}[\mathtt{i},\mathtt{n}]] \ \geq \ \pi[\mathtt{I}[\mathtt{n}] \setminus \{\mathtt{i}\},\mathtt{t}[\mathtt{i},\mathtt{n}]] \Big]
```

```
Theorem | "Jordan[2006a], 2.6a, contra", any [\pi, n],
   \forall_{\mathtt{i} \in \mathtt{I}[\mathtt{n}]} \; \left( (\mathtt{t}[\mathtt{i},\mathtt{n}] \; \notin \; \mathtt{K}[\pi,\mathtt{n}]) \; \Rightarrow \pi[\mathtt{I}[\mathtt{n}] \setminus \{\mathtt{i}\},\mathtt{t}[\mathtt{i},\mathtt{n}]] > \pi[\{\mathtt{i}\},\mathtt{t}[\mathtt{i},\mathtt{n}]] \right)
 \text{Theorem}\Big[\text{"Jordan[2006a],2.6b", any}[\pi,\,\text{n}], \text{ with}\Big[\underset{\text{if}[n]}{\forall} \pi[\{\text{i}\},\,\text{t[i,\,n]}] < \pi[\text{I[n]} \setminus \{\text{i}\},\,\text{t[i,\,n]}]\Big], 
   K[\pi, n] = \emptyset
Lemma ["Core n=3,a", any [\pi],
   (\mathtt{K}[\pi,\,3] = \emptyset) \ \Rightarrow \ \forall \\ \mathtt{i},\mathtt{j},\mathtt{k} \in \mathtt{I}[3] \ (\mathtt{distinct}[\mathtt{i},\,\mathtt{j},\,\mathtt{k}] \Rightarrow \mathtt{t}[\mathtt{i},\,3] \in \mathtt{D}[\{\mathtt{s}[\mathtt{j},\,\mathtt{k},\,3]\},\,\pi,\,3]) \ ]
Compute[W[3, t[1, 3], s[2, 3, 3]], using \rightarrow \{Definition["agents"], agents"]\}
         Definition["WinLose"], Definition["tyrannical"], Definition["split allocation"]},
   builtin → {Builtin ["Connectives"], Builtin ["Quantifiers"],
          Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}]
 {2,3}
\label{eq:compute_loss} \begin{split} &\text{Compute}[L[3,\,t[1,\,3]\,,\,s[2,\,3,\,3]]\,,\,using \rightarrow \{Definition["agents"]\,,\,\\ &\text{Compute}[L[3,\,t[1,\,3]\,,\,s[2,\,3]]\,,\,using \rightarrow \{Definition["agents"]\,,\,using \rightarrow \{Definition["agents]\,,\,using \rightarrow \{Definition["agen
         Definition["WinLose"], Definition["tyrannical"], Definition["split allocation"]},
   builtin → {Builtin ["Connectives"], Builtin ["Quantifiers"],
         Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}]
Definition["distinct", any[i, j, k],
   distinct[i, j, k]: \Leftrightarrow (i \neq j\land i \neq k\land j \neq k)]
Compute [W[3, t[1, 3], s[2, 3, 3]] = I[3] \setminus \{1\},\
   using → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
         Definition["tyrannical"], Definition["split allocation"]},
   builtin → {Builtin ["Connectives"], Builtin ["Quantifiers"],
          Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}]
True
Definition["agents"]
\operatorname{def}\left[\operatorname{agents}, \operatorname{range}\left[\operatorname{simpleRange}\left[n\right]\right], \operatorname{True}, \operatorname{flist}\left[\operatorname{lf}\left[,\operatorname{I}\left[n\right]:=\left\{i\right]\right]\right]\right]
Compute \left[ \bigvee_{i,j,k \in I[3]} \left( distinct[i,j,k] \Rightarrow (W[3,t[i,3],s[j,k,3]] = I[3] \setminus \{i\}) \right), 
   using → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
         Definition["tyrannical"], Definition["split allocation"]},
   builtin → {Builtin ["Connectives"], Builtin ["Quantifiers"]
          Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}
True
Compute \forall (distinct[i, j, k] \Rightarrow (L[3, t[i, 3], s[j, k, 3]] = {i})),
   using → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
         Definition["tyrannical"], Definition["split allocation"]},
   builtin \rightarrow \{Builtin ["Connectives"], Builtin ["Quantifiers"],
          Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}
True
ResetComputation[]
```

```
$Context
```

```
Tma`
```

```
Lemma | "s dominates t", any[i, j, k ∈ I[3]], with[distinct[i, j, k]],
 W[3, t[i, 3], s[j, k, 3]] = I[3] \setminus \{i\}
 L[3, t[i, 3], s[j, k, 3]] = \{i\}
Prove [Lemma ["s dominates t"],
 builtin → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
    Definition["tyrannical"], Definition["split allocation"],
    Builtin ["Connectives"], Builtin ["Quantifiers"], Builtin ["Tuples"],
    \label{eq:builtin} \verb|Builtin ["Sets"], Builtin ["Numbers"]| , by \rightarrow SetTheoryPCSProver,
 ProverOptions → {DisableProver → {STKBR, STP, QR, PND, CDP}, TransformRanges → False}]
{ - ProofObject -, - ProofObject -}
Prove[Lemma["Core n=3,a"],
 using → {Theorem["Jordan[2006a],2.6a,contra"], Lemma["s dominates t"],
    Lemma["tyrannical allocation"], Definition["dominion"], Definition["dominates"]},
 by → SetTheoryPCSProver, ProverOptions → {DisableProver → {STKBR}},
 builtin → {Builtin ["Number Domains"], Builtin ["Connectives"]}, SearchDepth → 110]
- ProofObject -
Block[{RecursionLimit = \infty}],
 {\tt Transform[\%, TransformerOptions} \rightarrow \{{\tt branches} \rightarrow {\tt Proved, steps} \rightarrow {\tt Useful}\}]]
- ProofObject -
Lemma ["Core n=3,b", any [\pi], with [AN[\pi, 3] \land powerfunction [\pi, 3]],
 \forall_{\text{i,j,keI[3]}} (\text{distinct[i,j,k]} \Rightarrow s[\text{j,k,3}] \notin D[\{\text{t[i,3]}\},\pi,3])
Lemma "AN perm, n=3", any [\pi], with [AN[\pi, 3]],
 \forall (distinct[i, j, k] \Rightarrow (\pi[\{i\}, s[j, k, 3]] = \pi[\{j\}, s[i, k, 3]]))
Lemma ["t dominates s", any[i, j, k], with[distinct[i, j, k]],
 W[3, s[j, k, 3], t[i, 3]] = \{i\}
 L[3, s[j, k, 3], t[i, 3]] = I[3] \setminus \{i\}
Lemma["distinct agents", any[i, j, k], with[distinct[i, j, k]],
 I[3] \setminus \{i\} = \{j, k\}]
Lemma | "SR, n=3", any [\pi], with [powerfunction [\pi, 3]],
 \forall_{\texttt{i},\texttt{j},\texttt{k},\alpha} \; (\texttt{distinct[i,j,k]} \; \bigwedge \pi[\{\texttt{j}\},\, \texttt{s[i,k,3]}] \; > \alpha \Rightarrow \pi[\{\texttt{j}\},\, \texttt{s[j,k,3]}] \; > \alpha) \; \Big]
Lemma "WC, n=3", any [\pi], with [powerfunction [\pi, 3]],
 \forall_{\substack{i,j,k,\alpha}} \left( \text{distinct[i,j,k]} \wedge \pi[\{j\},s[j,k,3]] > \alpha \Rightarrow \pi[\{j,k\},s[j,k,3]] > \alpha \right) \Big]
Lemma["irreflexive>", any[a],
 \neg (a > a)
```

```
Prove[Lemma["Core n=3,b"], using \rightarrow
    {Lemma["AN perm,n=3"], Lemma["t dominates s"], Lemma["distinct agents"], Lemma["SR,n=3"],
     Lemma["WC,n=3"], Definition["dominion"], Definition["dominates"], Lemma["irreflexive>"]},
  by \rightarrow SetTheoryPCSProver, ProverOptions \rightarrow {DisableProver \rightarrow {STP, STC, STS, CDP},
     RWSetOperators → True, AllowIntroduceQuantifiers → True,
     DisableInferenceRule → {¢KBInferNonEmpty, ¢KBSetEquality, ¢KBInclusion}}, SearchDepth → 90]
- ProofObject -
Block[{$RecursionLimit = Infinity}, ProofShow[]]
Definition \lceil \text{"dyadic number", any}[x], \text{ with}[x \in \mathbb{R}],
 dyadicNumber[x]: \Leftrightarrow \left((x=0) \bigvee_{k \in \mathbb{N}} \left(x=2^{-k}\right)\right)
Definition ["dyadic allocation", any [x, n],
 dyadicAllocation[x, n] : \Leftrightarrow \forall dyadicNumber[x_i]
Definition "dyadic set", any[k, n],
 dyadicSet[k,n] = \left\{ x \mid_{x} dyadicAllocation[x,n] \bigwedge_{i \in I[n]} \forall_{i} > 0 \Rightarrow x_{i} \geq 2^{-k} \right\} 
Lemma [ "D03",
  dyadicSet[0, 3] = {t[1, 3], t[2, 3], t[3, 3]}]
  dyadicSet[1, 3] = \{t[1, 3], t[2, 3], t[3, 3], s[1, 2, 3], s[1, 3, 3], s[2, 3, 3]\}
Definition ["championPowerfunction", any[C, x],
  champion\pi[C, x] = \max_{i \in C} x_i
Lemma \left[ \text{"max", } \exists_{k \in C} \left( \max_{i \in C} \mathbf{x}_i = \mathbf{x}_k \right) \bigwedge_{k \in C} \max_{i \in C} \mathbf{x}_i \ge \mathbf{x}_k \right]
Definition ["leaderPowerfunction", any [C, x, \nu], with [0 < \nu \wedge \nu < 1],
 leader \pi_{v}[C, x] = \begin{cases} \max_{i \in C} x_{i} & \Leftarrow \max_{i \in C} x_{i} \leq \frac{1}{3} \\ \sum_{i \in C} ((1 - v) * x_{i} + v) & \Leftarrow \text{ otherwise} \end{cases}
Definition | "centre allocation",
 \overline{p} = \left\langle \frac{1}{3}, \frac{1}{3}, \frac{1}{3} \right\rangle
Lemma ["championCore",
 K[\operatorname{champion}\pi, 3] = \operatorname{dyadicSet}[1, 3] \bigcup \{\overline{p}\}]
Lemma ["leaderCore", any [\nu], with [0 < \nu \land \nu < 1],
 K[leader\pi_v, 3] = dyadicSet[0, 3] \cup \{\overline{p}\}]
Theorem["coreCharacterization", any[\pi], with[powerfunction[\pi, 3] \wedge AN[\pi, 3]],
  (K[\pi, 3] = \emptyset) \lor (K[\pi, 3] = dyadicSet[0, 3]) \lor
    (K[\pi, 3] = dyadicSet[0, 3] \cup \{\overline{p}\}) \vee (K[\pi, 3] = dyadicSet[1, 3]) \vee
    (K[\pi, 3] = dyadicSet[1, 3] \bigcup \{\overline{p}\})]
```

```
Theorem ["coreCharacterization: \subseteq",  \left\{ \mathbb{K}[\pi, 3] \mid \text{powerfunction}[\pi, 3] \land \text{AN}[\pi, 3] \right\} \subseteq \\  \left\{ \emptyset, \, \text{dyadicSet}[0, 3], \, \text{dyadicSet}[0, 3] \cup \left\{ \overline{p} \right\}, \, \text{dyadicSet}[1, 3], \, \text{dyadicSet}[1, 3] \cup \left\{ \overline{p} \right\} \right]  Theorem ["coreCharacterization: \supseteq",  \left\{ \mathbb{K}[\pi, 3] \mid \text{powerfunction}[\pi, 3] \land \text{AN}[\pi, 3] \right\} \supseteq \\  \left\{ \emptyset, \, \text{dyadicSet}[0, 3], \, \text{dyadicSet}[0, 3] \cup \left\{ \overline{p} \right\}, \, \text{dyadicSet}[1, 3], \, \text{dyadicSet}[1, 3] \cup \left\{ \overline{p} \right\} \right]  Prove [Theorem ["coreCharacterization: \subseteq"], using \rightarrow {}, by \rightarrow SetTheoryPCSProver, ProverOptions \rightarrow {AllowIntroduceQuantifiers \rightarrow False}, SearchDepth \rightarrow 40] 
ES
```

In the rest of the document is the variant of the core characterisation, in which the lemma is not proved for arbitrary disjoint i,j,k, but 1,2,3

```
Lemma["Core n=3,a,123", any[\pi], with[distinct[1, 2, 3]],
    (K[\pi, 3] = \emptyset) \Rightarrow t[1, 3] \in D[\{s[2, 3, 3]\}, \pi, 3]]
Lemma s dominates t,123",
   W[3, t[1, 3], s[2, 3, 3]] = I[3] \setminus \{1\}
   L[3, t[1, 3], s[2, 3, 3]] = \{1\}
Prove[Lemma["s dominates t,123"],
   builtin → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
        Definition["tyrannical"], Definition["split allocation"],
        Builtin ["Connectives"], Builtin ["Quantifiers"], Builtin ["Tuples"],
        Builtin ["Sets"], Builtin ["Numbers"]}, by \rightarrow SetTheoryPCSProver,
   ProverOptions → {DisableProver → {STKBR, STP, QR, PND, CDP}, TransformRanges → False}]
Theorema::noEnvironment: No environment of type def named "distinct" available.
 { - ProofObject -, - ProofObject -}
Prove [Lemma ["Core n=3,a,123"],
   using → {Theorem["Jordan[2006a],2.6a,contra"], Lemma["s dominates t,123"],
        Lemma["tyrannical allocation"], Definition["dominion"], Definition["dominates"]},
   by \rightarrow SetTheoryPCSProver, ProverOptions \rightarrow {DisableProver \rightarrow {STKBR}},
   builtin → {Definition["agents"], Builtin ["Numbers"], Builtin ["Number Domains"],
        Builtin ["Quantifiers"], Builtin ["Connectives"]}, SearchDepth → 110]
Block[{$RecursionLimit = Infinity}, ProofShow[]]
Compute[L[3, t[1, 3], s[2, 3, 3]], using \rightarrow \{Definition["agents"], agents"], agents = \{Definition["agents"], agents = \{Defini
        Definition["WinLose"], Definition["tyrannical"], Definition["split allocation"]},
   \texttt{built} \pm \texttt{n} \rightarrow \{\texttt{Built} \pm \texttt{n} \ [\texttt{"Connectives"}], \texttt{Built} \pm \texttt{n} \ [\texttt{"Quantifiers"}],
        Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}]
 {1}
```

```
Compute [W[3, t[1, 3], s[2, 3, 3]] = I[3] \setminus \{1\},
 using → {Definition["distinct"], Definition["agents"], Definition["WinLose"],
    Definition["tyrannical"], Definition["split allocation"]},
 \texttt{built} \pm \texttt{n} \rightarrow \{\texttt{Built} \pm \texttt{n} \ [\texttt{"Connectives"}], \texttt{Built} \pm \texttt{n} \ [\texttt{"Quantifiers"}],
    Builtin ["Tuples"], Builtin ["Sets"], Builtin ["Numbers"]}]
True
Lemma "t dominates s,123",
 W[3, s[2, 3, 3], t[1, 3]] = \{1\}
 L[3, s[2, 3, 3], t[1, 3]] = I[3] \setminus \{1\}
\texttt{Lemma["SR, n=3,123", any}[\pi, \alpha], \texttt{with[powerfunction}[\pi, 3]],\\
   \pi[\{2\},\,\mathtt{s}[1,\,3,\,3]] \,>\, \alpha \,\Rightarrow\, \pi[\{2\},\,\mathtt{s}[2,\,3,\,3]] \,>\, \alpha]
Lemma ["WC n=3,b,123", any [\pi, \alpha], with [powerfunction [\pi, 3]],
  \pi[\{2\}, s[2, 3, 3]] > \alpha \Rightarrow \pi[\{2, 3\}, s[2, 3, 3]] > \alpha]
Lemma ["Core n=3,b,123", any [\pi], with [powerfunction [\pi, 3]],
 AN[\pi, 3] \Rightarrow s[2, 3, 3] \notin D[\{t[1, 3]\}, \pi, 3]]
Lemma["I[3]\{1}",
 I[3] \setminus \{1\} = \{2, 3\}]
Prove [Lemma ["Core n=3,b,123"],
 using → {Definition["dominion"], Definition["dominates"], Lemma["t dominates s,123"],
    Lemma["AN perm,n=3,123"], Lemma["SR, n=3,123"], Lemma["WC n=3,b,123"],
    \label{lemma} $$ Lemma["I[3] \setminus \{1\}"], Lemma["irreflexive>"]}, \ by \to SetTheoryPCSProver, $$
```

 $ProverOptions \rightarrow \{DisableProver \rightarrow \{STP, STC, STS, CDP\}, RWSetOperators \rightarrow True, \\$

{¢KBInferNonEmpty, ¢KBSetEquality, ¢KBInclusion}}, SearchDepth → 60]

AllowIntroduceQuantifiers → True, DisableInferenceRule →