Pillage Games

```
In[1]:= Needs["Theorema`"]
```

Basic Definitions

```
Definition["1", f[] := "Hello World!"]
         Compute[f[], using → Definition["1"]]
         Hello World!
TS_In[2]:= Definition | "agents", any[n],
            Compute[I[10], using → Definition["agents"],
          builtin → {Builtin ["Sets"], Builtin ["Quantifiers"], Builtin ["Numbers"]}]
         {1, 2, 3, 4, 5, 6, 7, 8, 9, 10}
        Definition["appL", any[n],
          appL[n] : \Leftrightarrow (n \in \mathbb{N} \land n \ge 2)]
TS_In[3]:= Lemma["linI", any[n], with[appL[n]],
            \{1\} \subseteq I[n]
TS_ln[4]:= Lemma["2inI", any[n], with[appL[n]],
            \{2\} \subseteq I[n]
         Prove [Lemma["linI"], using → {Definition["appL"], Definition["agents"]},
          builtin → {Builtin ["Numbers"]}, by → SetTheoryPCSProver,
          ProverOptions → {SimplifyFormula → True}]
         - ProofObject -
TS_ln[5]:= Definition ["allocation", any[x, n],
            \text{allocation}_n[\mathbf{x}] : \Leftrightarrow \left( n \in \mathbb{N} \bigwedge \text{istuple } [\mathbf{x}] \bigwedge \left( |\mathbf{x}| = n \right) \bigwedge \bigvee_{i \in I[n]} \mathbf{x}_i \ge 0 \bigwedge \left( \sum_{i \in I[n]} \mathbf{x}_i = 1 \right) \right) \right]
TS_In[6]:=
          Definition ["WC", any [\pi, n], bound [allocation<sub>n</sub>[x]],
            WC[\pi, n] : \Leftrightarrow \bigvee_{\substack{C1,C2 \\ C1 \in C2 \land C2 \subseteq I[n]}} \forall \pi[C2, x] \ge \pi[C1, x] \Big]
TS In[7]:=
          Definition ["WR", any [\pi, n], bound [allocation, [x], allocation, [y]],
           WR[\pi, n] : \Leftrightarrow \left[ n \in \mathbb{N} \bigwedge \bigvee_{\substack{C \\ C \in L[n]}} \bigvee_{x,y} \left( \left( \bigvee_{i \in C} y_i \geq x_i \right) \Rightarrow \pi[C, y] \geq \pi[C, x] \right) \right]
```

```
 \text{TS}_{\text{In}[8]:=} \text{ Definition} \left[ \text{"WR contrapositive", any}[\pi, n], \text{ bound}[\text{allocation}_n[x], \text{ allocation}_n[y]], \\ \text{WR}[\pi, n] : \Leftrightarrow \left( n \in \mathbb{N} \bigwedge_{\substack{C \\ CSI[n]}} \bigvee_{x,y} \left( \pi[C, x] > \pi[C, y] \Rightarrow \left( \frac{3}{3} x_i > y_i \right) \right) \right]   \text{TS}_{\text{In}[9]:=}   \text{Definition} \left[ \text{"SR", any}[\pi, n], \text{ bound}[\text{allocation}_n[x], \text{ allocation}_n[y]], \\ \text{SR}[\pi, n] : \Leftrightarrow \left( n \in \mathbb{N} \bigwedge_{\substack{C \\ CSI[n] \land C \neq \emptyset}} \bigvee_{x,y} \left( \left( \bigvee_{i \in C} y_i > x_i \right) \Rightarrow \pi[C, y] > \pi[C, x] \right) \right) \right]   \text{TS}_{\text{In}[10]:=}   \text{Definition} \left[ \text{"powerfunction", any}[\pi, n], \\ \text{powerfunction}[\pi, n] : \Leftrightarrow \bigwedge_{\substack{C \\ CSI[n] \land C \neq \emptyset}} \bigvee_{x,y} \left( \bigvee_{i \in C} y_i > x_i \right) \Rightarrow \pi[C, y] > \pi[C, x] \right) \right]   \text{TS}_{\text{In}[11]:=} \text{ Theory} \left[ \text{"powerfunction", } \\ \text{Definition} \left[ \text{"agents"} \right] \\ \text{Definition} \left[ \text{"agents"} \right] \\ \text{Definition} \left[ \text{"WC"} \right] \\ \text{Definition} \left[ \text{"WC"} \right] \\ \text{Definition} \left[ \text{"gowerfunction"} \right]   \text{Definition} \left[ \text{"gowerfunction"} \right]   \text{Definition} \left[ \text{"gowerfunction"} \right]   \text{Definition} \left[ \text{"gowerfunction"} \right]   \text{Definition} \left[ \text{"gowerfunction"} \right]
```

Lemma 1

```
 \text{TS\_In[12]:= Lemma} \left[ \text{"powerfunction-independent", any}[\pi, n, C, x, y], \right. \\ \text{with[allocation}_n[x] \land \text{allocation}_n[y] \land C \subseteq I[n] \land \text{powerfunction}[\pi, n]], \\ \text{\forall } \left( x_i = y_i \right) \Rightarrow \left( \pi[C, x] = \pi[C, y] \right) \right]
```

There is a potential problem in the next axiom, since it should be stated for real numbers only. Philosophy, trust in a proof is to read the proof.

? branches

```
Option of ProofSimplifier with possible values: Proved, Pending, Failed,
                             Disproved and list combinations of these. All (default) means list of all.
                                               PLemmalSimp = Block[{$RecursionLimit = Infinity}, Transform[PLemmal,
                                                                          by → ProofSimplifier, TransformerOptions → {branches → Proved, steps → Useful}]]
                                                 - ProofObject -
                                               Save["ecosimp1.m", {PLemma1Simp}]
                                               PLemma2Simp = Block[{$RecursionLimit = Infinity}, Transform[PLemma2,
                                                                          by → ProofSimplifier, TransformerOptions → {branches → Proved, steps → Useful}]]
Not used: {Definition (allocation), Definition (WC), Definition (SR)}
Not used: {Definition (allocation), Definition (WC), Definition (SR)}
                                                Save["ecosimp1full.m", {PLemma2Simp}]
                                                Options[SetTheoryPCSProver]
                                                   \{\texttt{DisableProver} \rightarrow \{\texttt{PND}\}\,,\,\, \texttt{EarlyRewriting} \rightarrow \texttt{False}\,,\,\, \texttt{TransformRanges} \rightarrow \texttt{True}\,,\,\,
                                                         {\tt AllowIntroduceQuantifiers} \rightarrow {\tt False, ApplyBuiltIns} \rightarrow {\tt True, BackChaining} \rightarrow {\tt False, ApplyBuiltIns} \rightarrow {\tt True, 
                                                         {\tt BackChainingDisjunction} \rightarrow {\tt False}, \ {\tt BackChainingEquivalence} \rightarrow {\tt False}, \\ {\tt BackChainingEquivalence} \rightarrow {\tt False}, \\ {\tt BackChainingDisjunction} \rightarrow {\tt False}, \\ {\tt BackChainingEquivalence} \rightarrow {\tt False}, \\ {\tt False}, \\ {\tt BackChainingEquivalence} \rightarrow {\tt False}, \\ {\tt BackCha
                                                         BackChainingImplication \rightarrow False, ChooseFromFiniteSet \rightarrow False,
                                                         DisableInferenceRule → {KBComposeIntersection, KBFiniteChoice, KBInferNonEmpty},
                                                         \texttt{EarlyCaseDistinction} \rightarrow \texttt{True}, \ \texttt{GRWTarget} \rightarrow \{\texttt{kb}, \ \texttt{goal}\}, \ \texttt{InferMembershipIntersection} \rightarrow \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \ \texttt{False}, \ \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \ \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \ \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \\ \texttt{False}, \ \texttt{False}, \\ \texttt{Fal
                                                         \texttt{KBRWOnlySimplification} \rightarrow \texttt{False} \,, \, \texttt{MatchExistential} \, \rightarrow \texttt{TryAllAtOnce} \,, \, \texttt{ModusPonensException} \, \rightarrow \, \texttt{None} \,, \, \\ \texttt{ModusPonensException
                                                       \texttt{ModusPonensKB} \rightarrow \texttt{True, PNDLevel} \rightarrow \texttt{1, RWBooleanLiteralCombinations} \rightarrow \texttt{True,}
                                                         \texttt{RWCombine} \rightarrow \texttt{False}, \ \texttt{RWExistentialGoal} \rightarrow \texttt{False}, \ \texttt{RWHigherOrder} \rightarrow \texttt{False}, \ \texttt{RWInnermost} \rightarrow \texttt{True},
                                                         {\tt RWInsideQuantifiers} \rightarrow {\tt False}, \ {\tt RWSetOperators} \rightarrow {\tt False}, \ {\tt RWTuples} \rightarrow {\tt True}, \ {\tt SemanticMatch} \rightarrow {\tt True},
                                                         SimplifyFormula \rightarrow False, STPFunctionProperties \rightarrow True, STPLevel \rightarrow 100,
                                                         {\tt STPMembershipByInclusion} \rightarrow {\tt False, TryAlternatives} \rightarrow {\tt False, TrySubgoal} \rightarrow {\tt False, TryAlternatives} \rightarrow {\tt False, TrySubgoal} \rightarrow {\tt False, TryAlternatives} \rightarrow {\tt False, TryAlternati
                                                         TrySubgoalDisjunction → False, TrySubgoalEquivalence → False, TrySubgoalImplication → False,
                                                         {\tt UseCyclicRules} \rightarrow {\tt False}, \ {\tt UseEqualitiesFirst} \rightarrow {\tt True}, \ {\tt UseNonMembership} \rightarrow {\tt True} \}
                                                Options[PredicateProver]
                                                 \{ \texttt{TryAlternatives} \rightarrow \texttt{False}, \, \texttt{BackChaining} \rightarrow \texttt{False}, \, \texttt{BackChainingDisjunction} \rightarrow \texttt{False}, \, \\
                                                         BackChainingImplication → False, BackChainingEquivalence → False,
                                                         TrySubgoalDisjunction \rightarrow False, TrySubgoalImplication \rightarrow False,
                                                         TrySubgoalEquivalence \rightarrow False, TrySubgoal \rightarrow False, PNDLevel \rightarrow 1\}
                                                ProofShow[]
                                               Depth[PLemma1]
                                                188
                                               Depth[PLemma2]
                                                372
                                               Depth[PLemma1Simp]
                                                70
                                               Depth[PLemma2Simp]
                                                74
```

Auxiliary Concepts

Permuted Allocation

```
 \begin{split} & \text{TS\_In[14]:= Definition["perm", any[x, \sigma, k],} \\ & \text{perm}[x, \sigma]_k := x_{\sigma[k]}] \\ & \text{TS\_In[15]:= Proposition["alloc perm", any[n, \sigma, x], with[allocation_n[x] \land permutation[\sigma, I[n]]],} \\ & \text{allocation}_n[perm[x, \sigma]]] \\ & \text{TS\_In[16]:= Proposition} \Big[ \text{"sum perm", any}[\sigma, A, x], \text{ with}[permutation[\sigma, A]],} \\ & \sum_{i \in A} x_{\sigma[i]} = \sum_{i \in A} x_i \Big] \\ & \text{Prove}[Proposition["alloc perm"],} \\ & \text{using} \rightarrow \{\text{Proposition}["alloc perm"],} \\ & \text{by} \rightarrow \text{SetTheoryPCSProver, ProverOptions} \rightarrow \{\text{AllowIntroduceQuantifiers} \rightarrow \text{True}\}] \\ & \text{- ProofObject -} \\ \end{split}
```

Allocation Swap

```
TS_In[17]:= Definition | "swap", any[i, j, k],
           \sigma_{\text{i,j}}[k] := \begin{cases} j \in k = i \\ i \in k = j \\ k \in \text{ otherwise} \end{cases}
TS_ln[18]:= Definition["permutation", any[\sigma, A],
           permutation[\sigma, A]: \Leftrightarrow \sigma :: A \xrightarrow{\text{bij}} A]
TS_ln[19]:= Proposition["swapperm", any[n], with[appL[n]],
           permutation[\sigma_{1,2}, I[n]]]
       Prove[Proposition["swapperm"],
        using -> {Definition["swap"], Definition["permutation"]}, by -> SetTheoryPCSProver]
       InverseFunction::ifun:
         Inverse functions are being used. Values may be lost for multivalued inverses. >>
       $Aborted
       ProofShow[]
TS_In[20]:= Proposition["swap idempotent", any[i], \sigma_{1,2}[\sigma_{1,2}[i]] = i]
       Prove[Proposition["swap idempotent"], using → Definition["swap"], by → SetTheoryPCSProver]
       - ProofObject -
TS_In[21]:= Lemma["perm swap", any[x],
           \mathtt{perm}\left[\mathtt{x,}\ \sigma_{1,2}\right]_2=\mathtt{x}_1]
       Prove[Lemma["perm swap"],
         using → {Definition["perm"], Definition["swap"]}, by → SetTheoryPCSProver]
       - ProofObject -
```

Additional Axioms

```
TS_In[22]:=
           Definition "AN", any [\pi, n], bound [allocation<sub>n</sub>[x], allocation<sub>n</sub>[y]],
            (\pi[Cx, x] = \pi[Cy, y])
TS_[n][23]:= Lemma["AN perm", any [n, \pi, x], with [appL[n] \land AN[\pi, n] \land allocation_n[x]],
            \pi[\{1\}, x] = \pi[\{2\}, perm[x, \sigma_{1,2}]]]
        PLemmaANperm =
         Prove[Lemma["AN perm"], using → {Proposition["swapperm"], Proposition["swap idempotent"],
              Proposition["alloc perm"], Lemma["linI"], Lemma["2inI"], Definition["AN"],
              \texttt{Definition["perm"], Definition["swap"]}, \ by \rightarrow \texttt{SetTheoryPCSProver, SearchDepth} \rightarrow \texttt{80,}
           ProverOptions → {AllowIntroduceQuantifiers → True, UseCyclicRules → False,
              EarlyRewriting → False, GRWTarget -> {"goal", "kb"}, DisableProver → {STC}}]
        $Aborted
        Transform[PLemmaANperm, TransformerOptions → {branches → Proved, steps → Useful}]
        ProofShow[]
Continuous 2 means that the function is continuous in its 2nd component.
TS_ln[24]:= Definition["CX", any[\pi, n],
            CX[\pi, n] : \Leftrightarrow continuous2[\pi]]
TS_ln[25]:= Definition TRE", any [\pi, n], bound [allocation [x]],
            RE[\pi, n] : \Leftrightarrow \bigvee_{C,i,x} (i \notin C \land \pi[\{i\}, x] > 0 \Rightarrow \pi[C \cup \{i\}, x] > \pi[C, x])
TS_In[26]:=
           Definition ["WinLose", any[n, x, y],
            W[n, x, y] := \left\{ i \mid_{i \in I[n]} y_i > x_i \right\} "W"
            L[n, x, y] := \left\{ i \bigcup_{i \in I[n]} x_i > y_i \right\} "L"
TS_{n[27]} = Definition["dominates", any[\pi, n, x, y],
            dominates[y, x, \pi, n] : \Leftrightarrow \pi[W[n, x, y], x] > \pi[L[n, x, y], x]]
        Lemma "ANdominates-aux", any[n, x, y],
          \text{with} [\text{allocation}_n[\textbf{x}] \ \land \text{allocation}_n[\textbf{y}] \ \land \ (\textbf{y}_i > \textbf{x}_i \ \Leftrightarrow \ (\textbf{i} = \textbf{1})) \ \land \ (\textbf{y}_i < \textbf{x}_i \ \Leftrightarrow \ (\textbf{i} = \textbf{2}))] \ , 
                                     dominates[y, x, \pi, n] \Leftrightarrow x_1 > x_2
         powerfunction[\pi,n]\bigwedgeAN[\pi,n]
TS_In[28]:= Lemma "ANdominates", any[n, x, y],
             \text{with} \left[ n \in \mathbb{N} \bigwedge n \geq 2 \bigwedge \text{allocation}_n[x] \ \bigwedge \text{allocation}_n[y] \ \bigwedge \ (\mathbb{W}[n,\,x,\,y] = \{1\}) \ \bigwedge \ (\mathbb{L}[n,\,x,\,y] = \{2\}) \right], 
                                         (dominates[y, x, \pi, n] \Rightarrow x_1 > x_2)
            AN[\pi,n] \(\rangle\) powerfunction [\pi,n]
```

```
PLemmaANdominates =
       Prove [Lemma ["ANdominates"], using → {Lemma ["AN perm"], Definition ["dominates"],
           {\tt Definition["powerfunction"], Definition["SR"]}, \ by \rightarrow {\tt SetTheoryPCSProver,}
         SearchDepth \rightarrow 50, ProverOptions \rightarrow {AllowIntroduceQuantifiers \rightarrow True,
           UseCyclicRules → True, EarlyRewriting → False, DisableProver → {STKBR, STC}}]
      $Aborted
      ProofShow[]
TS In[29]:= Definition["appAlloc", any[n, x],
          appAlloc[n, x] : \Leftrightarrow (appL[n] \land allocation_n[x])]
TS_ln[30]:= Lemma "ANdominates 1", any[n, x, y],
          (dominates[y, x, \pi, n] \Rightarrow x_1 > x_2)
         \forall \pi
AN[\pi,n]\powerfunction[\pi,n]
TS_ln[31]:= Lemma "ANdominates 2", any[n, x, y],
          (\mathbf{x}_1 > \mathbf{x}_2 \Rightarrow \text{dominates}[\mathbf{y}, \mathbf{x}, \pi, n])
          AN[\pi,n] \land powerfunction[\pi,n]
TS_ln[32]:= Definition["appPowAlloc", any[n, \pi, x],
          appPowAlloc[n, \pi, x]: \Leftrightarrow (appAlloc[n, x] \land AN[\pi, n])]
TS_ln[33]:= Lemma "AN all", any[n, \pi, x], with[appAlloc[n, x] \wedge AN[\pi, n]],
          Prove [Lemma ["AN all"], using → {Lemma ["AN perm"], Proposition ["alloc perm"],
          Proposition["swapperm"], Definition["appAlloc"]}, by → PredicateProver]
      - ProofObject -
      PLemmaANdominates = Prove[Lemma["ANdominates 1"],
         using → {Definition["dominates"], Lemma["AN all"], Definition["WR contrapositive"],
           Lemma["2inI"], Definition["powerfunction"], Lemma["perm swap"], Definition["appAlloc"]},
         by \rightarrow SetTheoryPCSProver, SearchDepth \rightarrow 80, ProverOptions \rightarrow
          {AllowIntroduceQuantifiers → False, RWCombine → True, DisableProver → {STC},
           DisableInferenceRule → {¢KBSetEquality, ¢KBInclusion}}]
      - ProofObject -
      Block[{RecursionLimit = \infty}, ProofShow[]]
      Block[{RecursionLimit = \infty}],
       Transform[PLemmaANdominates, TransformerOptions → {branches → Proved, steps → Useful}]]
      - ProofObject -
      Block[{RecursionLimit = \infty}, ProofShow[PLemmaANdominates]]
      Prove[Lemma["ANdominates 2"],
       using → {Definition["dominates"], Lemma["AN all"], Definition["SR"], Lemma["2inI"],
          Definition["powerfunction"], Lemma["perm swap"], Definition["appAlloc"]},
       builtin → {Builtin ["Operators"]["¬", "="]}, by → SetTheoryPCSProver,
       \texttt{SearchDepth} \rightarrow \texttt{80, ProverOptions} \rightarrow \{\texttt{AllowIntroduceQuantifiers} \rightarrow \texttt{False, RWCombine} \rightarrow \texttt{True,}
          \label{eq:distance} \mbox{DisableInferenceRule} \rightarrow \{\mbox{$\varsigma$KBSetEquality, $\varsigma$KBInclusion}\}\}]
      - ProofObject -
```

$Block[{RecursionLimit = \infty}, ProofShow[]]$

Options[SetTheoryPCSProver]

 $\{ \mbox{DisableProver} \rightarrow \{\mbox{PND}\}, \mbox{ EarlyRewriting} \rightarrow \mbox{False}, \mbox{ TransformRanges} \rightarrow \mbox{True}, \\ \mbox{AllowIntroduceQuantifiers} \rightarrow \mbox{False}, \mbox{ ApplyBuiltIns} \rightarrow \mbox{True}, \mbox{ BackChaining} \rightarrow \mbox{False}, \\ \mbox{BackChainingImplication} \rightarrow \mbox{False}, \mbox{ BackChainingEquivalence} \rightarrow \mbox{False}, \\ \mbox{BackChainingImplication} \rightarrow \mbox{False}, \mbox{ ChooseFromFiniteSet} \rightarrow \mbox{False}, \\ \mbox{DisableInferenceRule} \rightarrow \{\mbox{KBComposeIntersection}, \mbox{KBFiniteChoice}, \mbox{KBInferNonEmpty}\}, \\ \mbox{EarlyCaseDistinction} \rightarrow \mbox{True}, \mbox{ GRWTarget} \rightarrow \{\mbox{kb}, \mbox{goal}\}, \mbox{InferMembershipIntersection} \rightarrow \mbox{False}, \\ \mbox{KBRWOnlySimplification} \rightarrow \mbox{False}, \mbox{ MatchExistential} \rightarrow \mbox{TryAllAtOnce}, \mbox{ ModusPonensException} \rightarrow \mbox{None}, \\ \mbox{ModusPonensKB} \rightarrow \mbox{True}, \mbox{ PNDLevel} \rightarrow \mbox{1}, \mbox{ RWBooleanLiteralCombinations} \rightarrow \mbox{True}, \\ \mbox{RWCombine} \rightarrow \mbox{False}, \mbox{ RWExistentialGoal} \rightarrow \mbox{False}, \mbox{ RWHigherOrder} \rightarrow \mbox{False}, \mbox{ RWInnermost} \rightarrow \mbox{True}, \\ \mbox{ RWInsideQuantifiers} \rightarrow \mbox{False}, \mbox{ RWSetOperators} \rightarrow \mbox{ False}, \mbox{ RWTuples} \rightarrow \mbox{ True}, \mbox{ SemanticMatch} \rightarrow \mbox{ True}, \\ \mbox{ SimplifyFormula} \rightarrow \mbox{ False}, \mbox{ STPFunctionProperties} \rightarrow \mbox{ True}, \mbox{ STPLevel} \rightarrow \mbox{ 100}, \\ \mbox{ STPMembershipByInclusion} \rightarrow \mbox{ False}, \mbox{ TrySubgoalEquivalence} \rightarrow \mbox{ False}, \mbox{ TrySubgoalImplication} \rightarrow \mbox{ False}, \\ \mbox{ UseCyclicRules} \rightarrow \mbox{ False}, \mbox{ UseEqualitiesFirst} \rightarrow \mbox{ True}, \mbox{ UseNonMembership} \rightarrow \mbox{ True} \}$

ResetComputation[]