

بسم الله الرحمن الرحيم

«سیستم عامل»

۱۸۰

جلسه ۱۹: مدیریت حافظه (۷)

Page replacement

- ❑ Assume a normal page table (e.g., BLITZ)
- ❑ User-program is executing
- ❑ A PageInvalidFault occurs!
 - ❖ The page needed is not in memory
- ❑ Select some frame and remove the page in it
 - ❖ If it has been modified, it must be written back to disk
 - the “dirty” bit in its page table entry tells us if this is necessary
- ❑ Figure out which page was needed from the faulting addr
- ❑ Read the needed page into this frame
- ❑ Restart the interrupted process by retrying the same instruction

Page replacement algorithms

- ❑ Which frame to replace?
- ❑ Algorithms:
 - ❖ The Optimal Algorithm
 - ❖ First In First Out (FIFO)
 - ❖ Not Recently Used (NRU)
 - ❖ Second Chance / Clock
 - ❖ Least Recently Used (LRU)
 - ❖ Not Frequently Used (NFU)
 - ❖ Working Set (WS)
 - ❖ WSClock

❑

The optimal page replacement algorithm

- ❑ Idea:
 - ❖ Given all the data, how to find the optimal page replacement?

Time	0	1	2	3	4	5	6	7	8	9	10
Requests		c	a	d	b	e	b	a	b	c	d

Page	0	a
Frames	1	b
	2	c
	3	d

Page faults

The optimal page replacement algorithm

- Idea:

- ❖ Given all the data, how to find the optimal page replacement?

Time	0	1	2	3	4	5	6	7	8	9	10
Requests		c	a	d	b	e	b	a	b	c	d

Page	0	a	a	a	a	a					
Frames	1	b	b	b	b	b					
	2	c	c	c	c	c					
	3	d	d	d	d	d					

Page faults

X

The optimal page replacement algorithm

- Idea:
 - ❖ Given all the data, how to find the optimal page replacement?
 - ❖ **Longest Forward Distance (LFD):** Select the page that will not be needed for the longest time

LFD

- Replace the page that will not be needed for the longest
- Example:

Time	0	1	2	3	4	5	6	7	8	9	10
Requests		c	a	d	b	e	b	a	b	c	d

Page	0	a									
Frames	1	b	a	a	a	a					
	2	c	b	b	b	b					
	3	d	c	c	c	c					
			d	d	d	d					

Page faults

X

LFD

- Select the page that will not be needed for the longest time
- Example:

Time	0	1	2	3	4	5	6	7	8	9	10
Requests		c	a	d	b	e	b	a	b	c	d
<hr/>											
Page	0	a	a	a	a	a	a	a	a	a	
Frames	1	b	b	b	b	b	b	b	b	b	
	2	c	c	c	c	c	c	c	c	c	
	3	d	d	d	d	e	e	e	e	e	
Page faults						X				X	

Is LFD optimal?

Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD

Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - First non-equal position: $i+1$

Is LFD optimal?

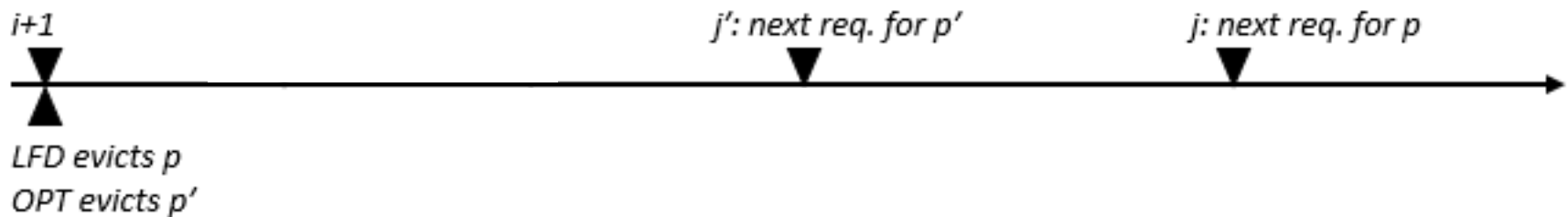
- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - First non-equal position: $i+1$
 - Case 1) $i+1$ is not a page fault.
 - Case 2) $i+1$ is page fault

Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - First non-equal position: $i+1$
 - Case 1) $i+1$ is not a page fault.
 - Case 2) $i+1$ is page fault
 - LFD evicts p
 - OPT evicts p'

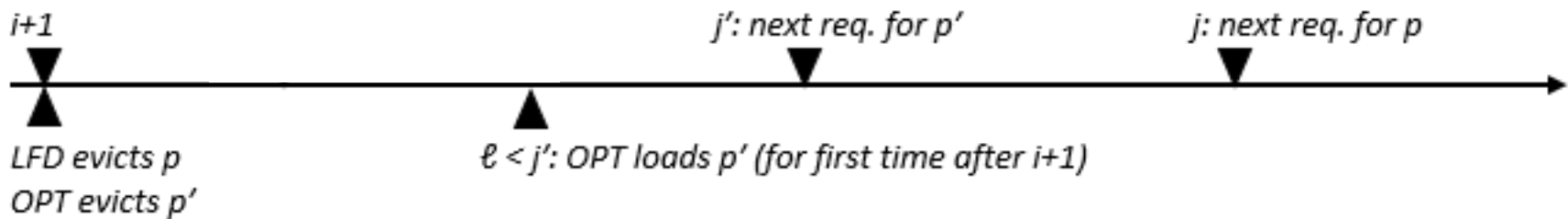
Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - First non-equal position: $i+1$
 - Case 1) $i+1$ is not a page fault.
 - Case 2) $i+1$ is page fault
 - LFD evicts p
 - OPT evicts p'



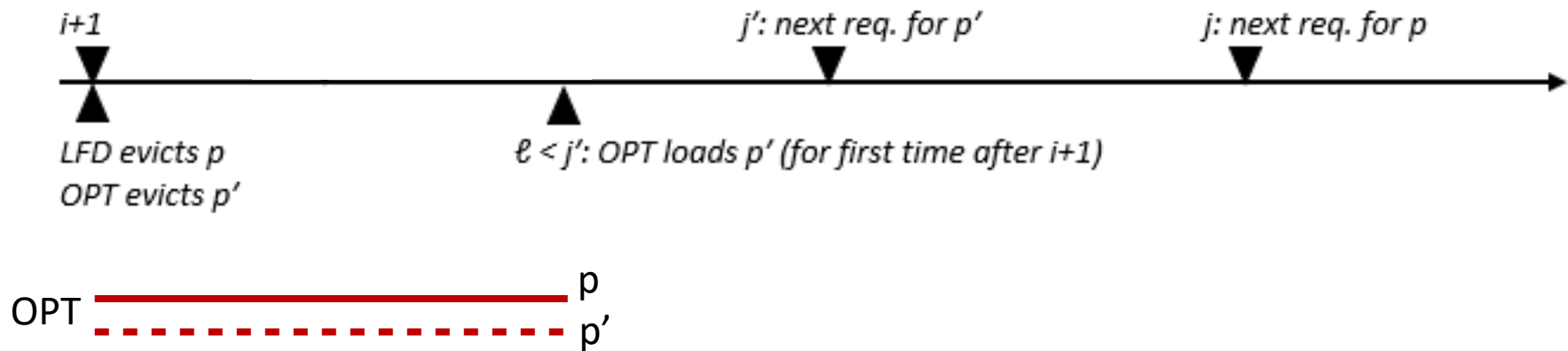
Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



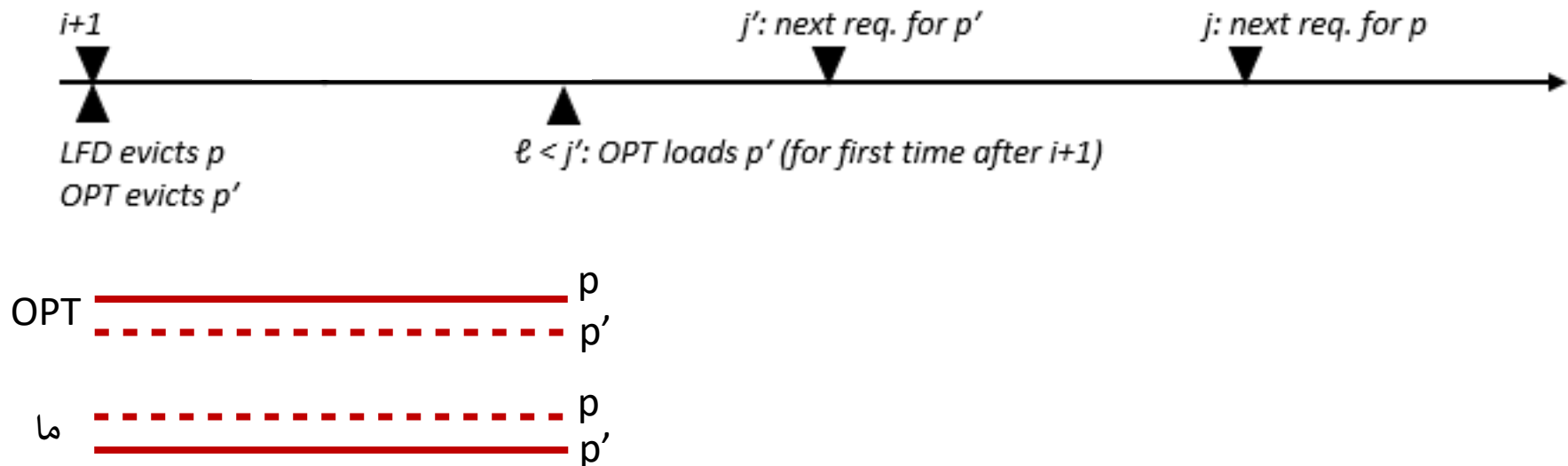
Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



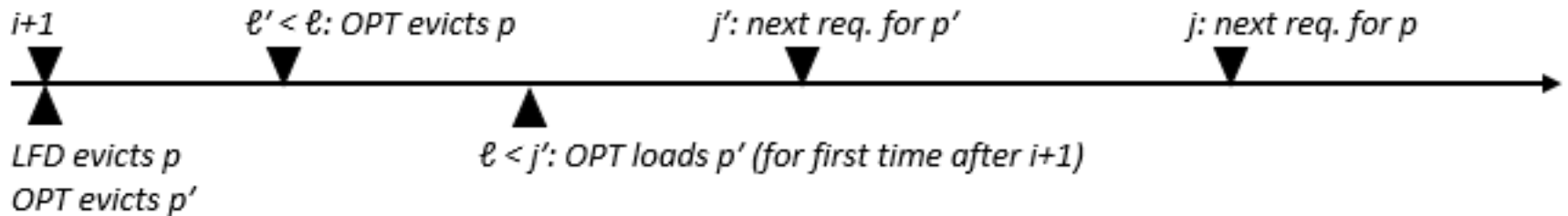
Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



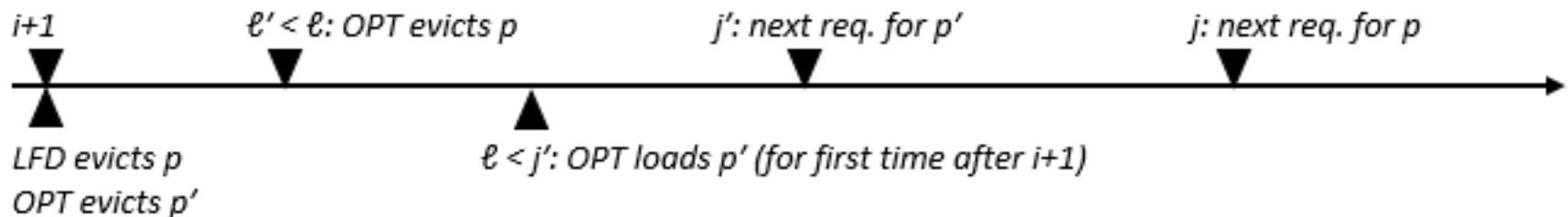
Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



Is LFD optimal?

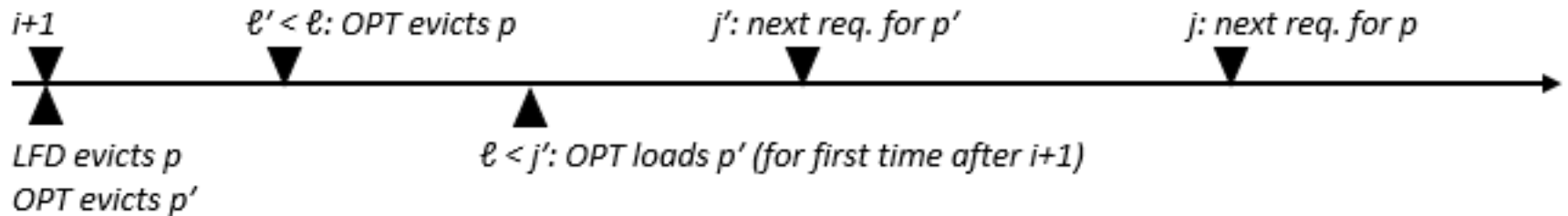
- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



OPT ——— p
- - - - - p'

Is LFD optimal?

- Proof (by contradiction):
 - **OPT**: Optimum with longest prefix equal to LFD
 - Case 2) $i+1$ is page fault
 - Case 2-A) OPT keeps p until l
 - Case 2-B) OPT evicts p at $l' < l$



OPT ——— p
- - - - - p'

ω - - - - - p
——— p'

$$\text{LFD} = \text{OPT}$$