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برنامه ریزی ریاضی (SDP)

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Question 1 :

G is a complete graph K_n , where n is even. What is the size of the maximum cut of G ?

Question 2 :

We consider a communication channel, with input $X(t) \in \{1, \dots, n\}$, and output $Y(t) \in \{1, \dots, m\}$, for $t = 1, 2, \dots$ (in seconds, say). The relation between the input and the output is given statistically:

$$p_{ij} = \text{prob}(Y(t) = i | X(t) = j), \quad i = 1, 2, \dots, m, \quad j = 1, 2, \dots, n$$

The matrix $P \in R^{m \times n}$ is called the channel transition matrix, and the channel is called a discrete memoryless channel.

A famous result of Shannon states that information can be sent over the communication channel, with arbitrarily small probability of error, at any rate less than a number C , called the channel capacity, in bits per second. Shannon also showed that the capacity of a discrete memoryless channel can be found by solving an optimization problem. Assume that X has a probability distribution denoted $x \in R^n$, i.e.,

$$x_j = \text{prob}(X = j), \quad j = 1, \dots, n$$

The mutual information between X and Y is given by

$$I(X; Y) = \sum_{i=1}^m \sum_{j=1}^n x_j p_{ij} \log_2 \frac{p_{ij}}{\sum_{k=1}^n x_k p_{ik}}$$

Then the channel capacity C is given by

$$C = \sup_x I(X; Y),$$

where the supremum is over all possible probability distributions for the input X , i.e., over $x \succeq 0$, $1^T x = 1$.

Show how the channel capacity can be computed with concave objective function.

Hint. Introduce the variable $y = Px$, which gives the probability distribution of the output Y , and show that the mutual information can be expressed as

$$I(X; Y) = c^T x - \sum_{i=1}^m y_i \log_2 y_i,$$

(The entropy of y is concave), where $c_j = \sum_{i=1}^m p_{ij}$, $j = 1, 2, \dots, n$.

Question 3 :

Suppose that C and D are disjoint subsets of R^n .

Consider the set of $(a, b) \in R^{n+1}$ for which $a^T x \leq b$ for all $x \in C$, and $a^T x \geq b$ for all $x \in D$. Show that this set is a convex cone (which is the singleton $\{0\}$ if there is no hyperplane that separates C and D).

Question 4 :

Suppose $K \subseteq R^2$ is a closed convex cone.

(a) Give a simple description of K in terms of the polar coordinates of its elements ($x = r(\cos\theta, \sin\theta)$ with $r \geq 0$).

(b) Give a simple description of K^* , and draw a plot illustrating the relation between K and K^* .

Question 5 :

Let $G = (V, E)$ be a graph. First, show that $\vartheta(G)$ can be expressed as the value of the following optimization problem:

$$\begin{aligned} \vartheta(G) = \min \quad & \lambda_{\max}(11^T + X) \\ \text{s.t.} \quad & X_{ij} = 0, \text{ if } \{i, j\} \in \bar{E}, \text{ or } i = j \\ & X \in SYM_n. \end{aligned}$$

Second, what is the dual of this SDP formulation?

Question 6 :

Formulate the following optimization problems as semidefinite programs. The variable is $x \in R^n$; $F(x)$ is defined as

$$F(x) = F_0 + x_1 F_1 + x_2 F_2 + \dots + x_n F_n$$

with $F_i \in S^m$. The domain of f in each subproblem is $\text{dom } f = \{x \in R^n | F(x) \succ 0\}$.

(a) Minimize $f(x) = \max_{i=1, \dots, k} c_i^T F(x)^{-1} c_i$ where $c_i \in R^m$, $i = 1, \dots, k$.

(b) Minimize $f(x) = \sup_{\|c\|_2 \leq 1} c^T F(x)^{-1} c$

Hint. Schur Complements

Question 7 :

Suppose $A : R^n \rightarrow S^m$ is affine, i.e.,

$$A(x) = A_0 + x_1 A_1 + x_2 A_2 + \dots x_n A_n$$

where $A_i \in S^m$. Let $\lambda_1(x) \geq \lambda_2(x) \geq \dots \geq \lambda_m(x)$ denote the eigenvalues of $A(x)$. Show how to pose the following problems as SDPs.

(a) Minimize the maximum eigenvalue $\lambda_1(x)$

(b) Minimize the spread of the eigenvalues, $\lambda_1(x) - \lambda_m(x)$