

بسم الله الرحمن الرحيم

Introduction

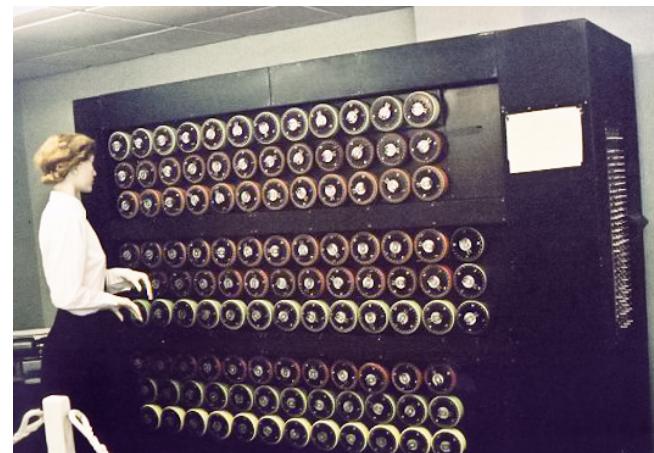


- Course Information
- A Quick History of Compilers
- The Structure of a Compiler

What is a Compiler?



History of High-level Languages



A Short History of Compilers

- First, there was nothing.
- Then, there was machine code.
- Then, there were assembly languages.
- Programming expensive; 50% of costs for machines went into programming.

First Practical Compiler



In his PhD dissertation 1951; published in 1954), Böhm describes for the first time a translation mechanism of a programming language, written in that same language.

https://en.wikipedia.org/wiki/Corrado_B%C3%B6hm

https://en.wikipedia.org/wiki/B%C3%B6hm%27s_language

High-Level Languages



Rear Admiral **Grace Hopper**, inventor of COBOL, and the term “compiler.”

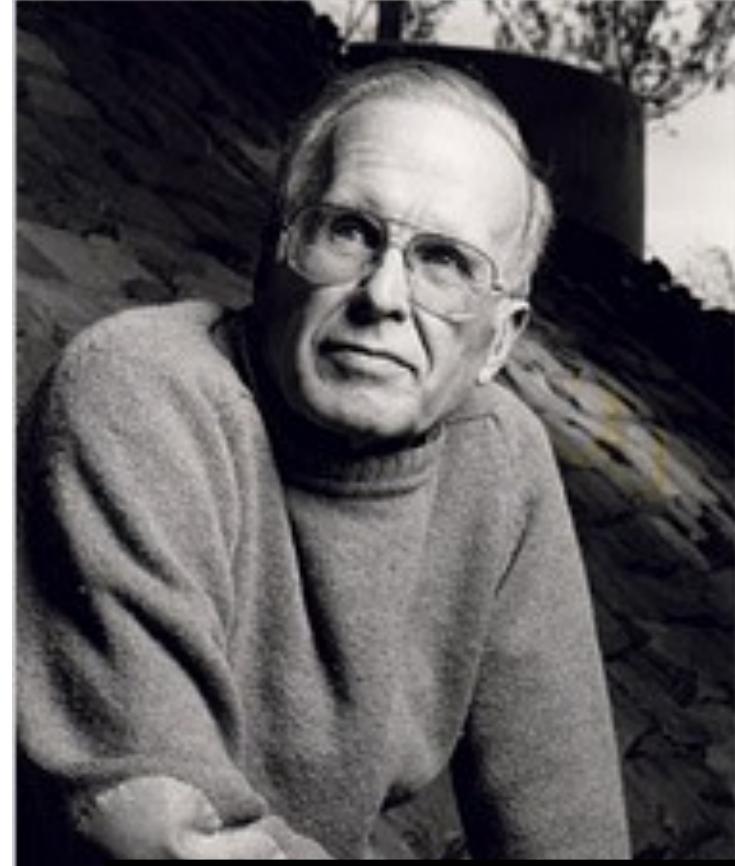
High-Level Languages



Rear Admiral **Grace Hopper**, inventor of COBOL, and the term “compiler.”

A programming language based on English

High-Level Languages



John Backus,
team lead on
FORTRAN.

Image: http://upload.wikimedia.org/wikipedia/commons/thumb/5/55/Grace_Hopper_in_Navy_uniform.jpg/800px-Grace_Hopper_in_Navy_uniform.jpg

<http://www.nytimes.com/2007/03/20/business/20backus.html>

FORTRAN I

- Translate high-level code to assembly.
- Many thought this impossible.
- Had already failed in other projects.
- Development time halved
- Performance is close to hand-written assembly!



Effect on Computer Science

- The first compiler
- Huge impact on computer science.
- Led to an enormous body of theoretical and practical work.
- Modern compilers preserve the outlines of FORTRAN I

```
INTEGER FUNCTION FCN20(NDIMS, X, NFCNS, FUNVLS)
  INTEGER NDIMS, NFCNS
  DOUBLE PRECISION X(*), FUNVLS(*)
  DOUBLE PRECISION Z
  Z = (X(1) + X(2) + X(3)) ** 2
  IF (Z .NE. 0.0) THEN
    FUNVLS(1) = 1.0 / Z
  ELSE
    FUNVLS(1) = 0.0
  ENDIF
  FCN20 = 1
  RETURN
END
```

What is a Compiler?

- Takes as input a program written in one language and **translates** it into a **functionally equivalent** program in another language.
- Source is usually high-level (e.g. Java), target is usually low-level (e.g. Assembly).
-

Hierarchy of dependency between concepts related to compiler

- computational model
- programming language
- compiler

Computational Thinking in Programming Language Design

Underlying every programming language is a **model of computation**:

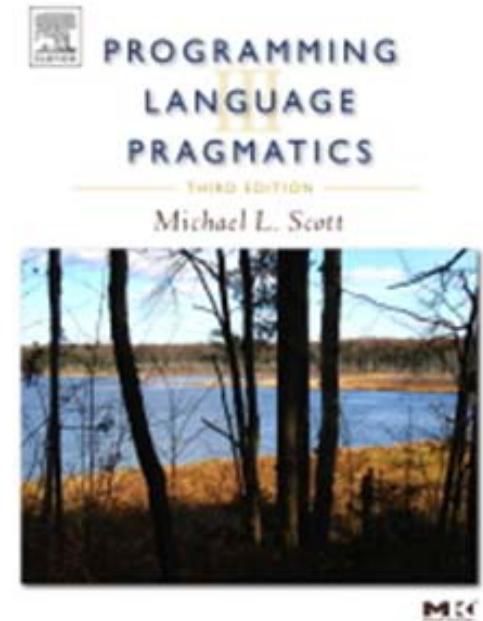
Procedural: C, C++, C#, Java

Declarative: SQL

Logic: Prolog

Functional: Haskell

Scripting: AWK, Perl, Python, Ruby



Evolutionary Forces on Languages and Compilers

More and different kinds of languages

Increasing diversity of applications

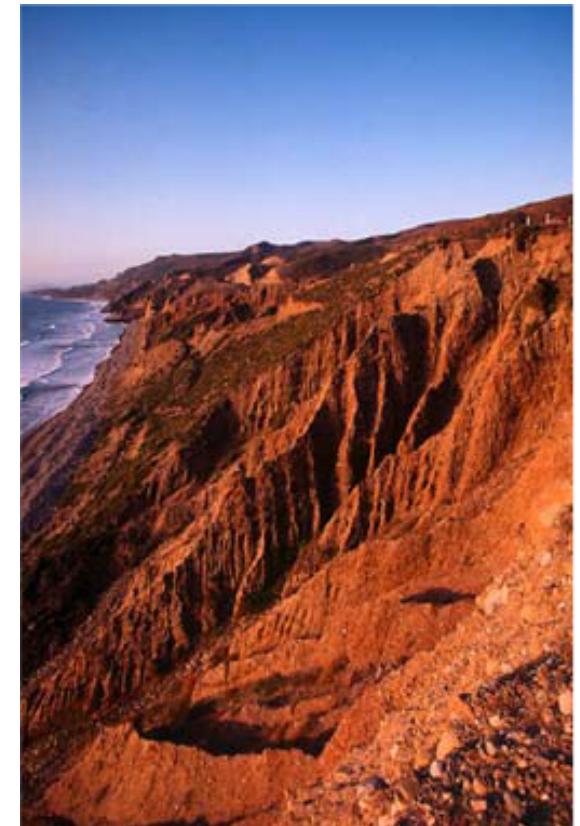
Stress on increasing productivity

Need to improve software reliability

Target machines more diverse

Parallel machine architectures

Massive compiler collections



1970

Fortran

Lisp

Cobol

Algol 60

APL

Snobol 4

Simula 67

Basic

PL/1

Pascal

2010

Java

C

PHP

C++

Visual Basic

C#

Python

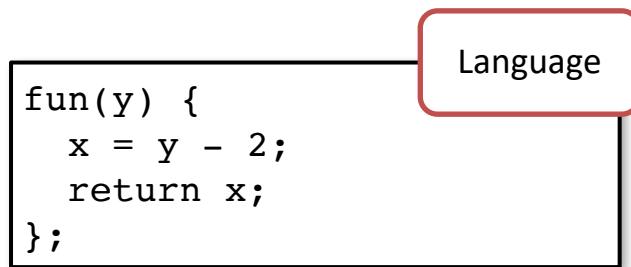
Perl

Delphi

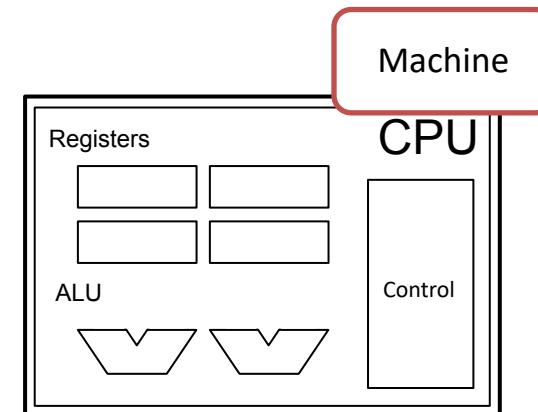
JavaScript

[<http://www.tiobe.com>]

Schema



Infinite resources
No performance
specification



- Finite resources
- Extremely performance sensitive

Language Environment Construction touches many topics in Computer Science

- Theory
 - Finite State Automata, Grammars and Parsing, data-flow
- Algorithms
 - Graph manipulation, dynamic programming
- Data structures
 - Symbol tables, abstract syntax trees
- Systems
 - Allocation and naming, multi-pass systems, compiler construction
- Computer Architecture
 - Memory hierarchy, instruction selection, interlocks and latencies, parallelism
- Security
 - Detection of and Protection against vulnerabilities
- Software Engineering
 - Software development environments, debugging
- Artificial Intelligence
 - Heuristic based search for best optimizations

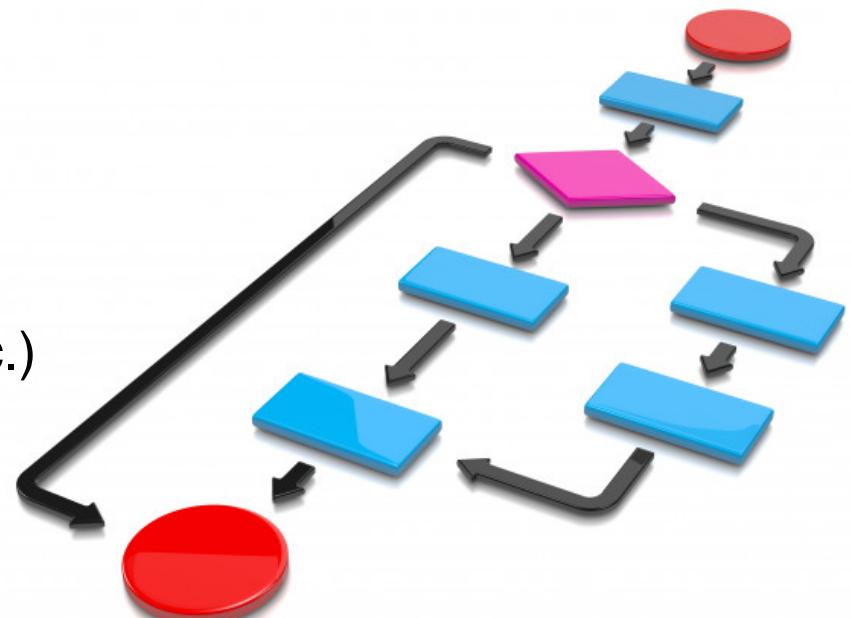
The diagram contains several mathematical equations and diagrams:

- $(x)^2 = ab$
- $dh(x) = bc$
- $e = f^2 (x + 4gh)^2 (s) \cdot (x^3 - (gh)^2 - x^2)$
- $f = gh^2 + (s)(x+2h)^2 \times 4x^2(hc)^2 + x^2 - 2x^2$
- $g = x^2 \div (x)(2x)^2 + (hc)^2 4x^2(3h)(f)^2(e)^2 + x^4 + x^2$
- $h = efg^2 - (x^2 + (3)^2(f)^3 + x(4x)^2)$
- $a = x(s^1) + (h)(c) + (d)(ef)^2 = x^2$
- $b = 2x + (h)(d) \div (s^1)(h^2)(b)^2 = 4x^2 hd$
- $c = 3x(c)d \alpha(2x)^3 \div (x)(x)^2 2x = 2s + 4x$
- $d = x^2 + 4(s^1)(s)^2 + ab \div c^2(h)$
- $j^1 = O^2 - (K\rho)^3 - (4)(-x)^6(oK)^3$
- $j^2 = K\rho(x)(-2)^2(4x)^3$
- $\rho = 40^\circ(s^1)O + \rho(-2)(-x^2)$
- $O = X^2(x.K\rho)^3 \div (4)^4 + (x)(K) = 23^\circ$
- $\frac{j^2 = K^2(\rho o)^2 x}{x + 2x}$
- $(x)^2 = ab$
- $dh(x) = bc$

designed by freePik

Inputs

- Standard language
 - State
 - Variables,
 - Structures,
 - Arrays
 - Computation
 - Expressions (arithmetic, logical, etc.)
 - Assignment statements
 - Control flow (conditionals, loops)
 - Procedures



Outputs

- State
 - Registers
 - Memory with Flat Address Space
- Machine code – load/store architecture
 - Load, store instructions
 - Arithmetic, logical operations on registers
 - Branch instructions

```
FE30- 20 B4 FC 90 F7 60 B1 3C
*F0EDL

F0ED-   6C 36 00    JMP    ($0036)
F0F0-   C0 A0        CMP    #$A0
F0F2-   00 02        BCC    $FDF6
F0F4-   00 32        AND    $32
F0F6-   04 35        STY    $35
F0F8-   40            PHA
F0F9-   20 78 FB    JSR    $FB78
F0FC-   60            PLA
F0FD-   A4 35        LDY    $35
F0FF-   00            RTS
FEG0-   00 34        DEC    $34
FEG2-   F0 9F        BEQ    $FDA3
FEG4-   CA            DEX
FEG5-   D0 16        BNE    $FE1D
FEG7-   C9 BA        CMP    #$BA
FEG9-   D0 BB        BNE    $FDC6
FEGB-   00 31        STA    $31
FEGD-   A5 3E        LDA    $3E
FEGF-   91 40        STA    ($40), Y
FE11-   E6 40        INC    $40
*
```

Translation Approaches

- Compiler Approach
- Interpreter Approach
- Dynamic Approach

Compiler Approach

Compiler Approach

- The target is not necessarily a machine code.
 - e.g. Assembly language e.g. MIPS or x86
 - e.g. VHDL: the output is C.
 - It might be intermediate code e.g. JBC

Compiler Approach

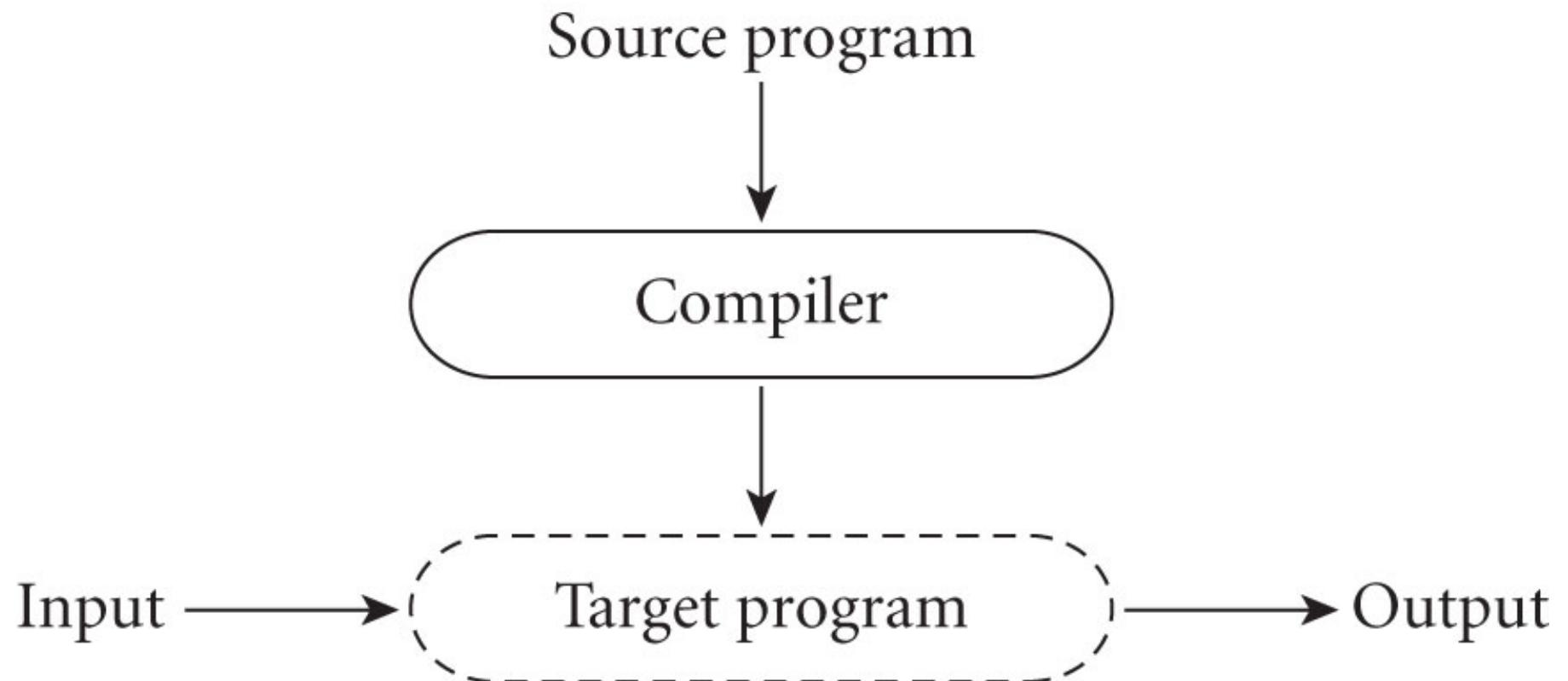
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- By following **physical structure** of program we translate it.

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- The generated code is much more faster.

Compiler Approach

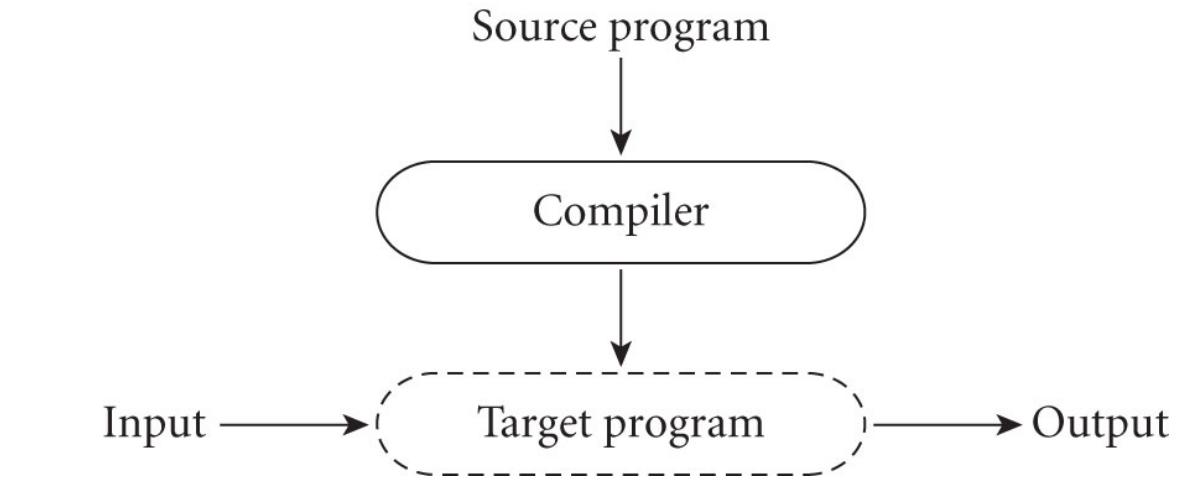
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 - e.g. Assembly language e.g. MIPS or x86
 - e.g. VHDL: the output is C.
 - It might be intermediate code e.g. JBC
- By following **physical structure** of program we translate it.
- The generated code is much more faster.
- We decide before run the code (e.g. type)



```
#include <stdio.h>

int getInt() {
    int i;
    scanf("%d", &i);
    return i;
}

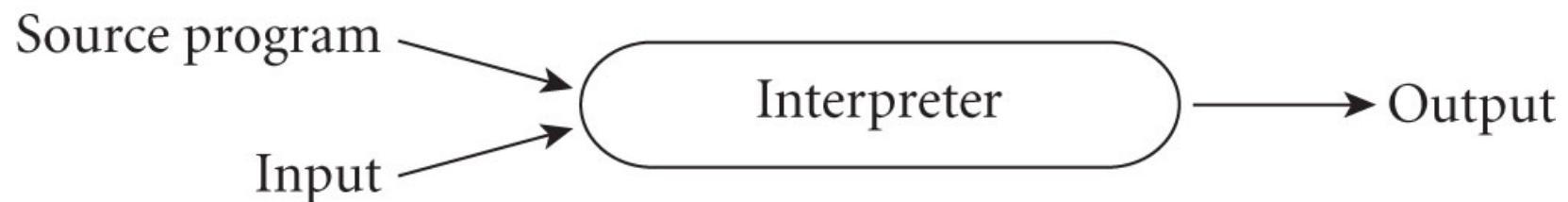
int gcd() {
    int i = getInt(), j = getInt();
    while (i != j) {
        if (i > j) i = i - j;
        else j = j - i;
    }
    return i;
}
```



.LC0:	main:			
.string "%d"				
getInt():				
push rbp	push	rbp		
mov rbp, rsp	mov	rbp, rsp		
sub rsp, 16	sub	rsp, 16		
lea rax, [rbp-4]	x			
mov rsi, rax			call	getInt()
mov edi, OFFSET FLAT:.LC0			mov	DWORD PTR [rbp-4], ea
mov eax, 0	x			
call __isoc99_scanf			jmp	.L5
mov eax, DWORD PTR [rbp-4]	.L7:		mov	eax, DWORD PTR [rbp-4]
leave				
ret]		
.LC1:			cmp	eax, DWORD PTR [rbp-8]
.string "%d\n"]	jle	.L6
putInt(int):			mov	eax, DWORD PTR [rbp-8]
push rbp				
mov rbp, rsp]		
sub rsp, 16			sub	DWORD PTR [rbp-4], ea
mov DWORD PTR [rbp-4], edi	x			
mov eax, DWORD PTR [rbp-4]			jmp	.L5
mov esi, eax		.L6:	mov	eax, DWORD PTR [rbp-4]
mov edi, OFFSET FLAT:.LC1				
mov eax, 0]		
call printf			sub	DWORD PTR [rbp-8], ea
leave		x		
ret		.L5:	mov	eax, DWORD PTR [rbp-4]

Interpreter Approach

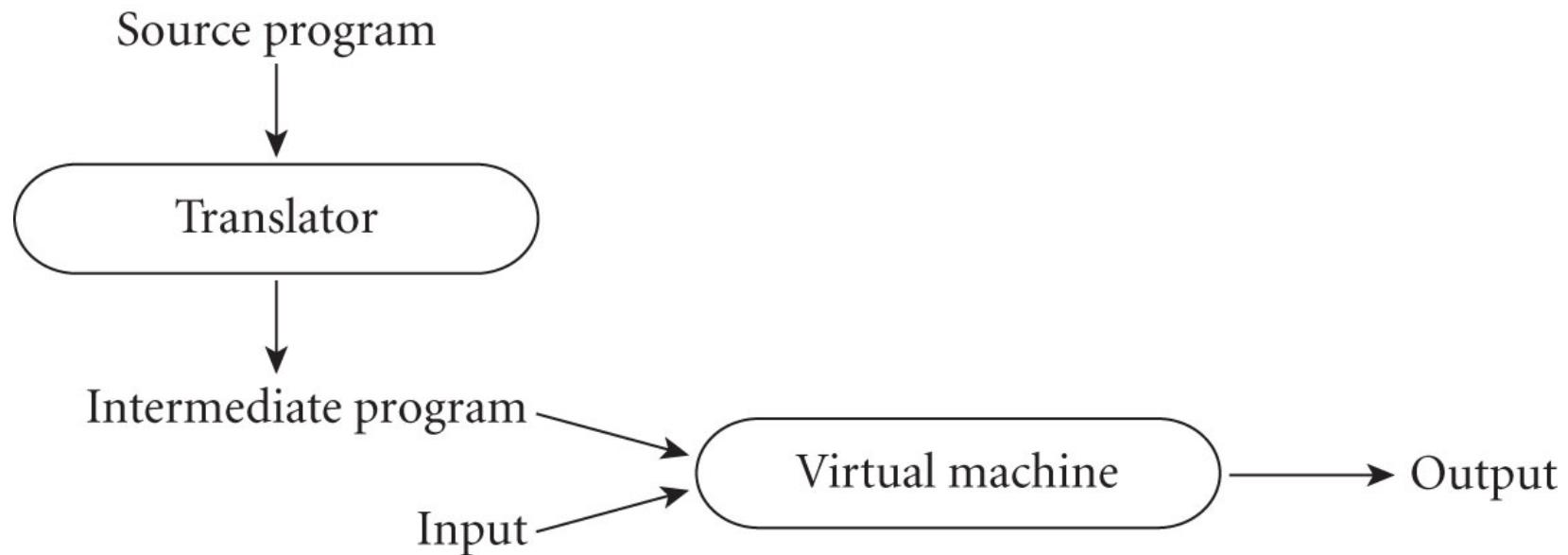
- Interpreter **translate** the source code to machine code **online**.
- The code is **translated** while it is **running**.



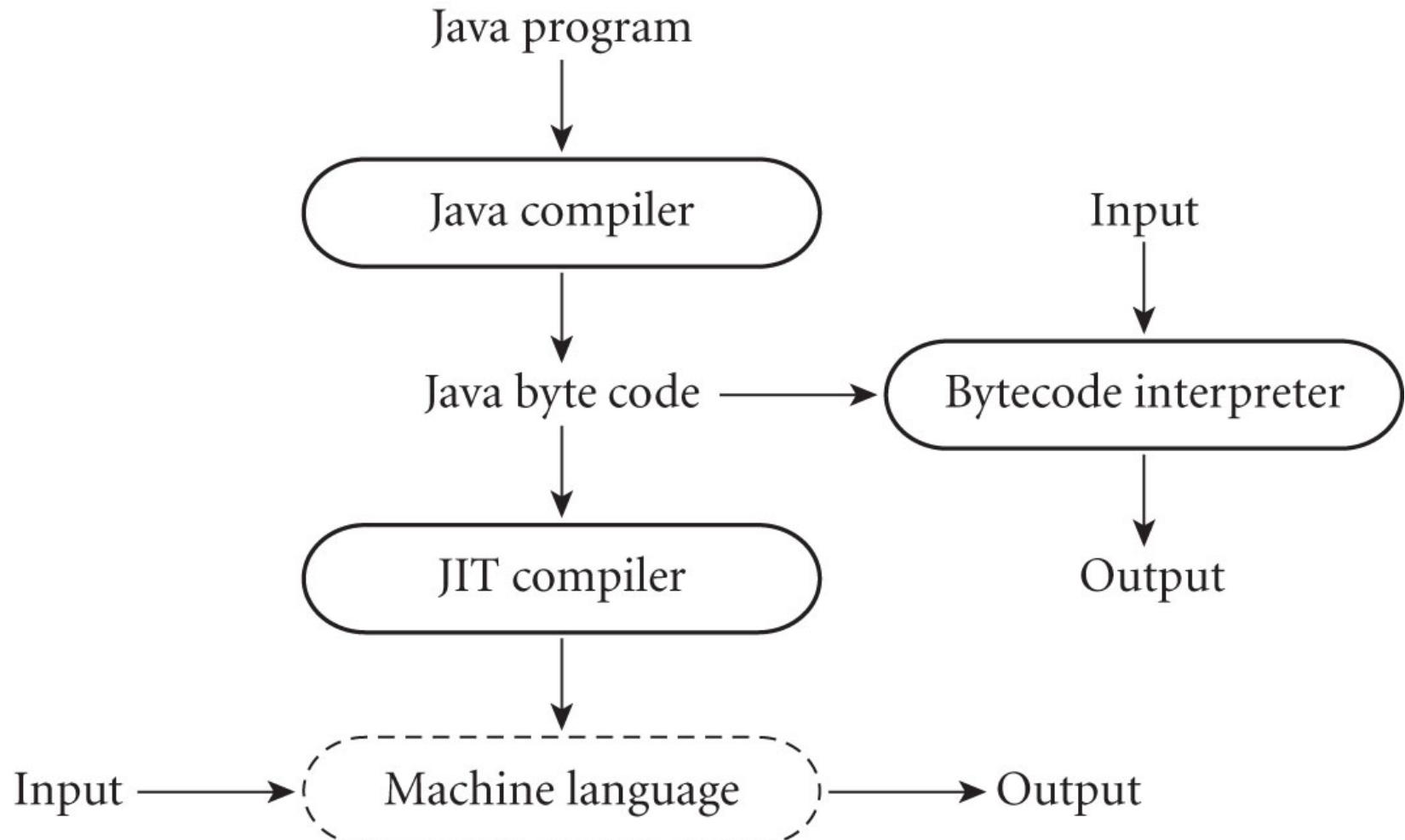
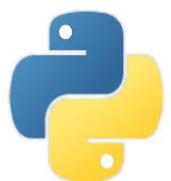
Interpreter Approach

- We follow the **execution path** to translate the code.
- leads to greater **flexibility** and better **diagnostics** (error messages) than does compilation.
- It can also cope with languages in which fundamental characteristics of the program, such as the sizes and types of variables, or even which names refer to which variables, can depend on the input data.

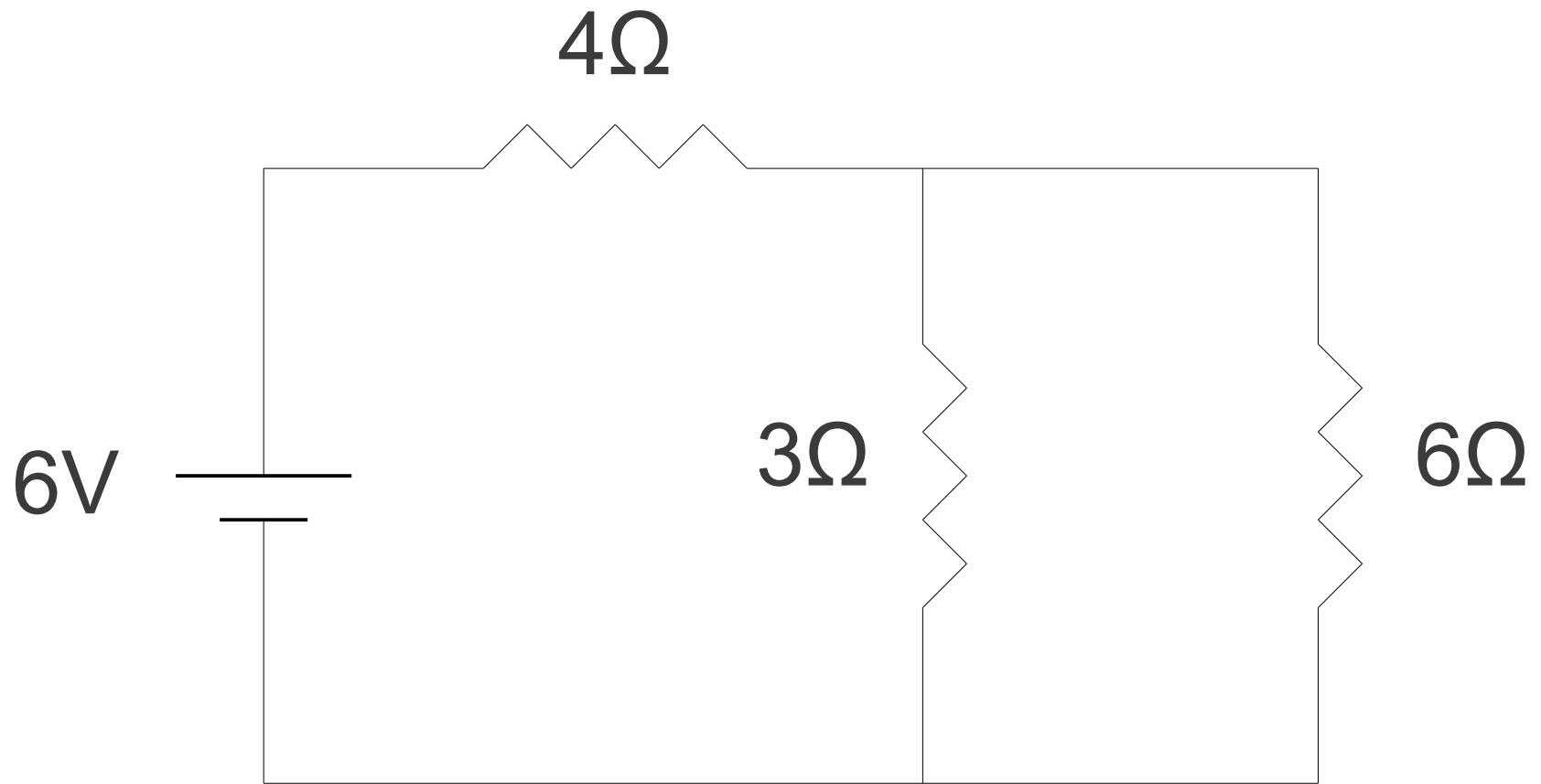
A Hybrid Approach

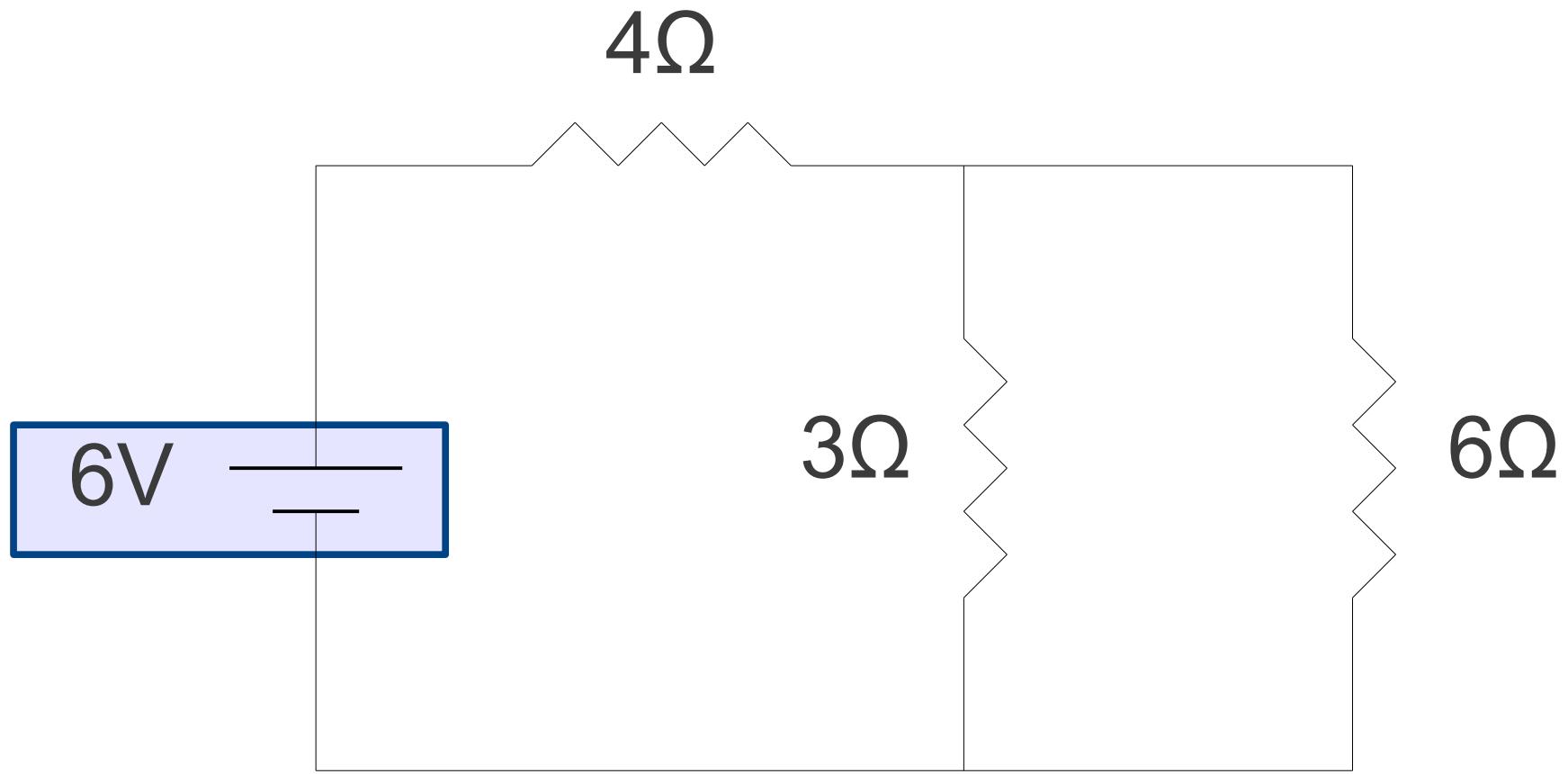


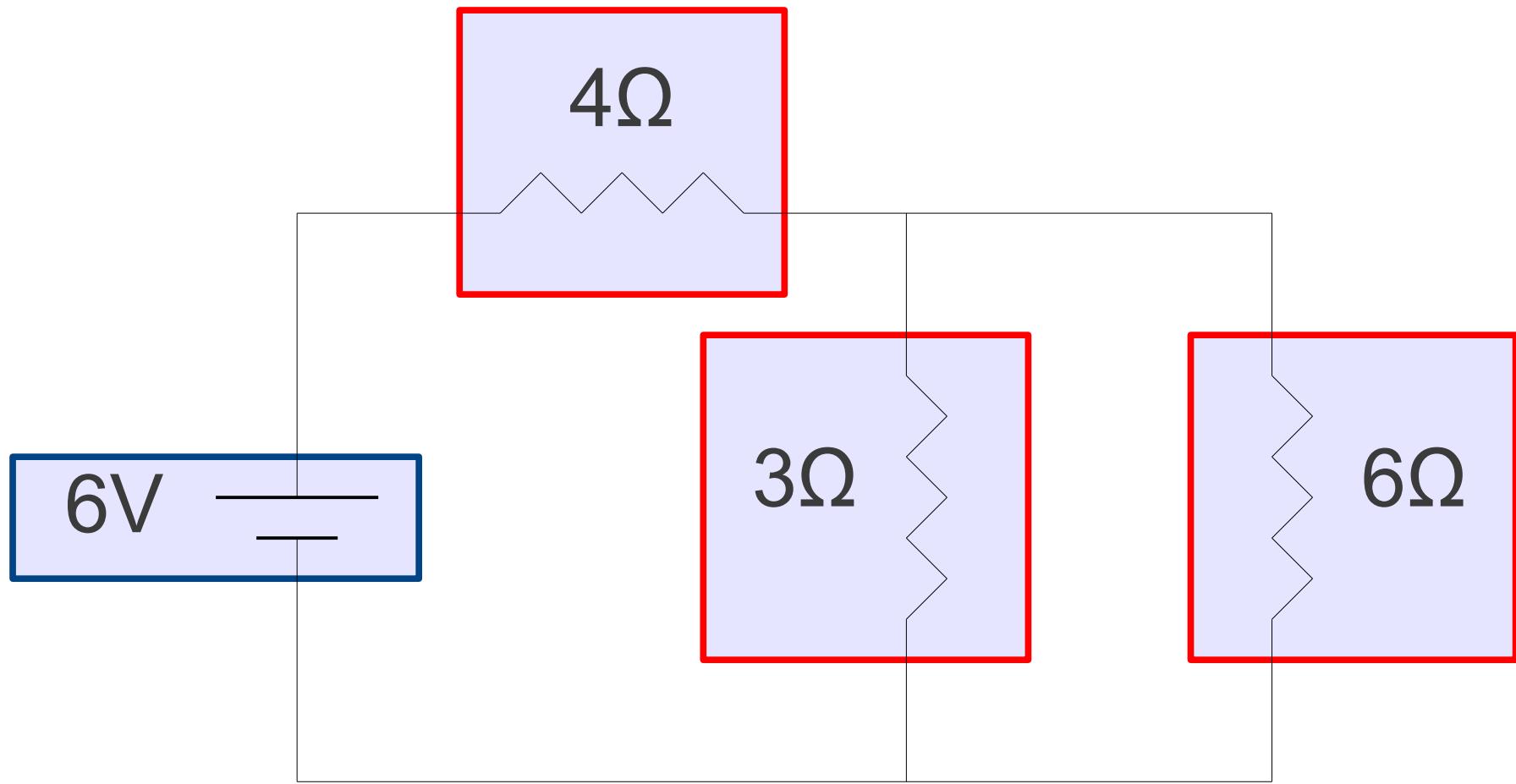
Hybrid Approach

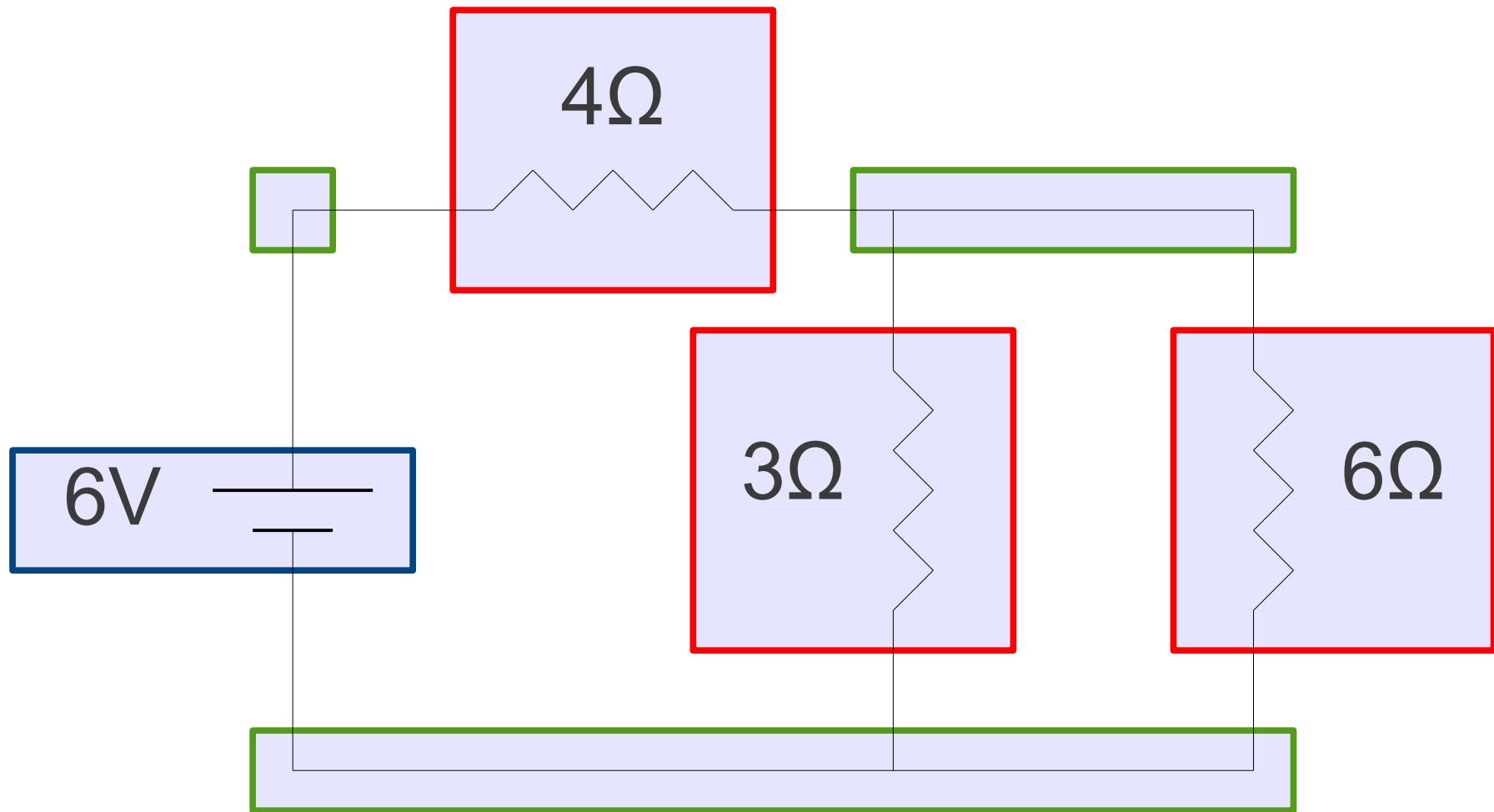


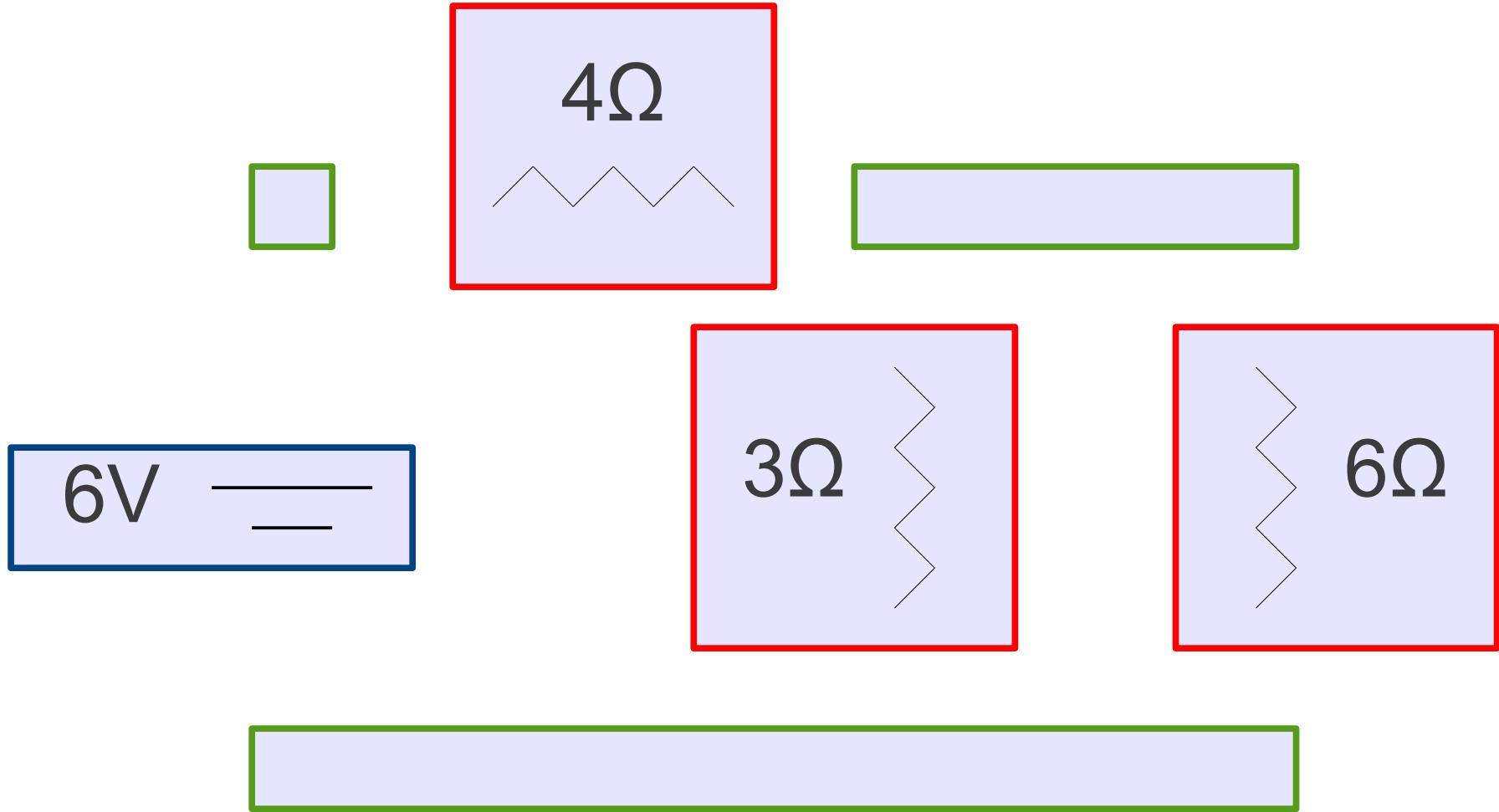
How does a compiler work?

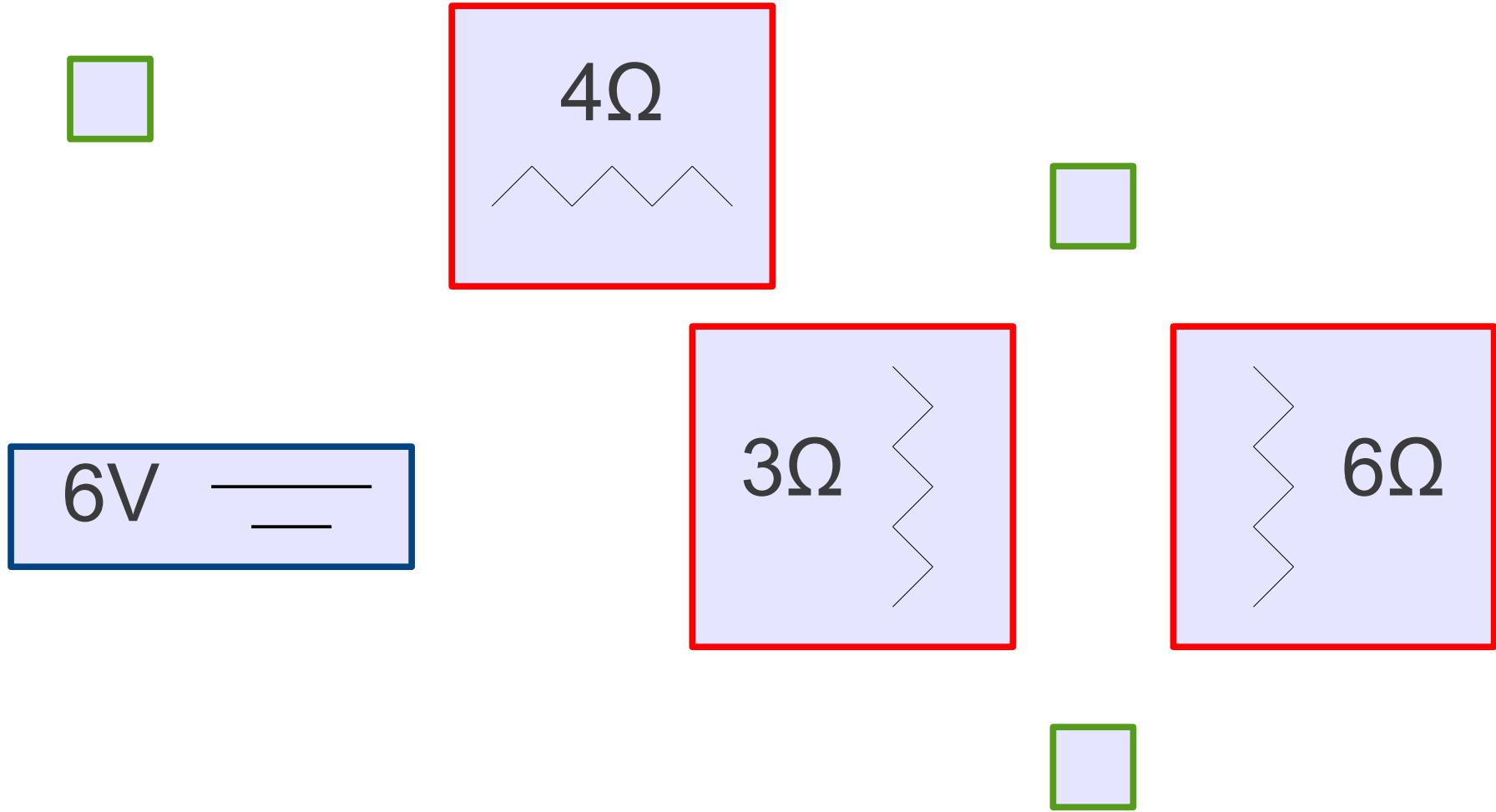


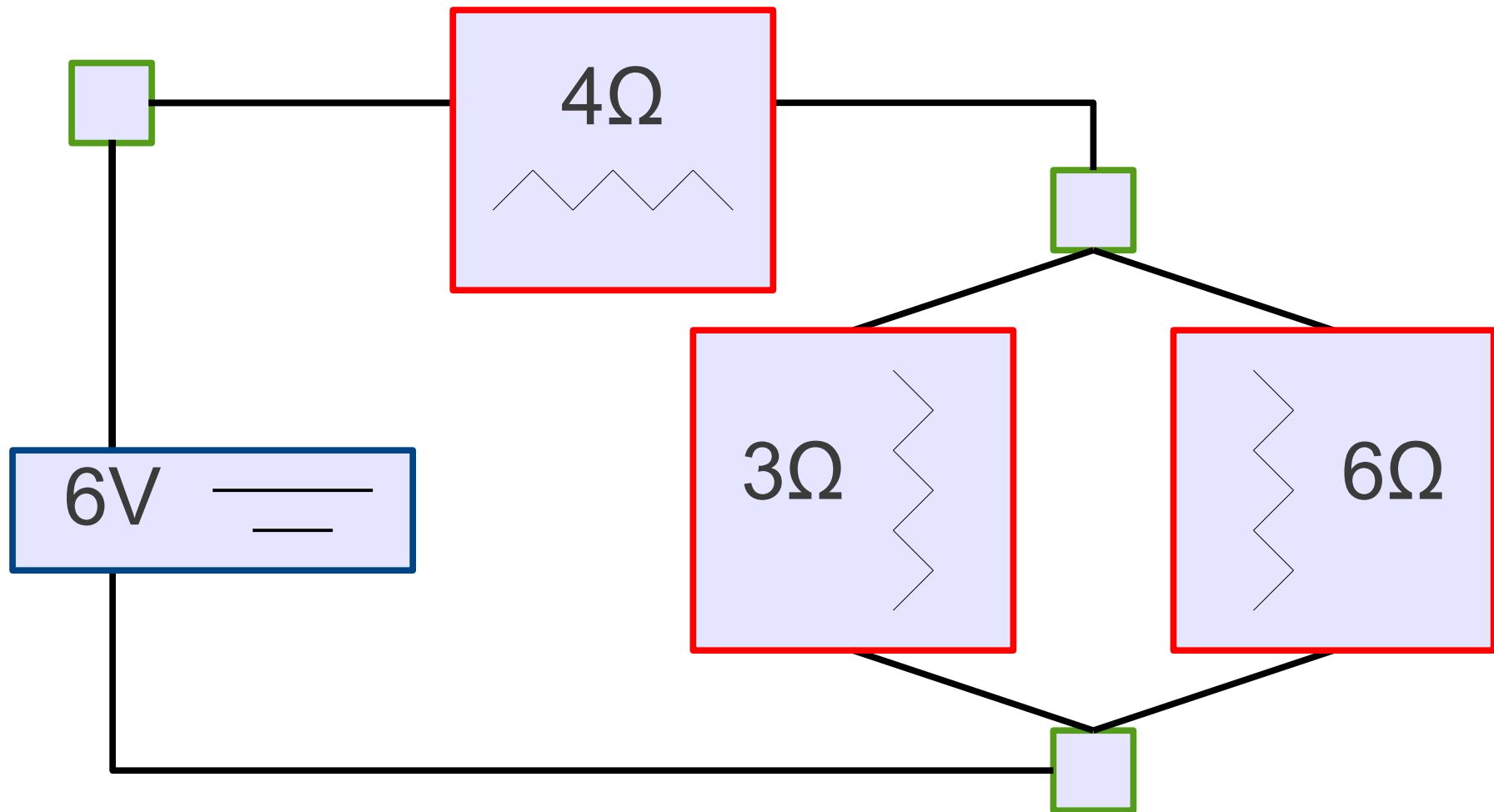


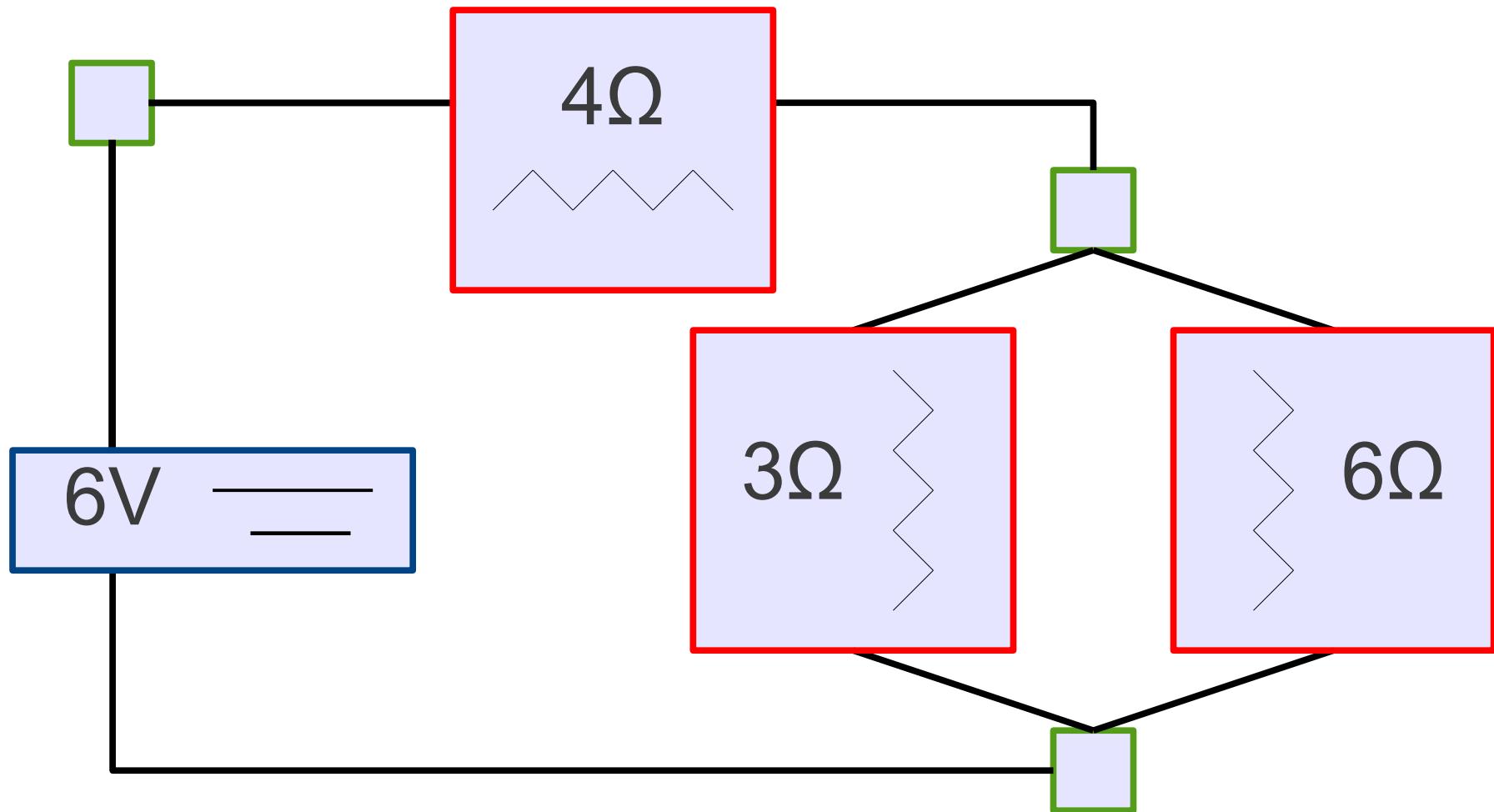




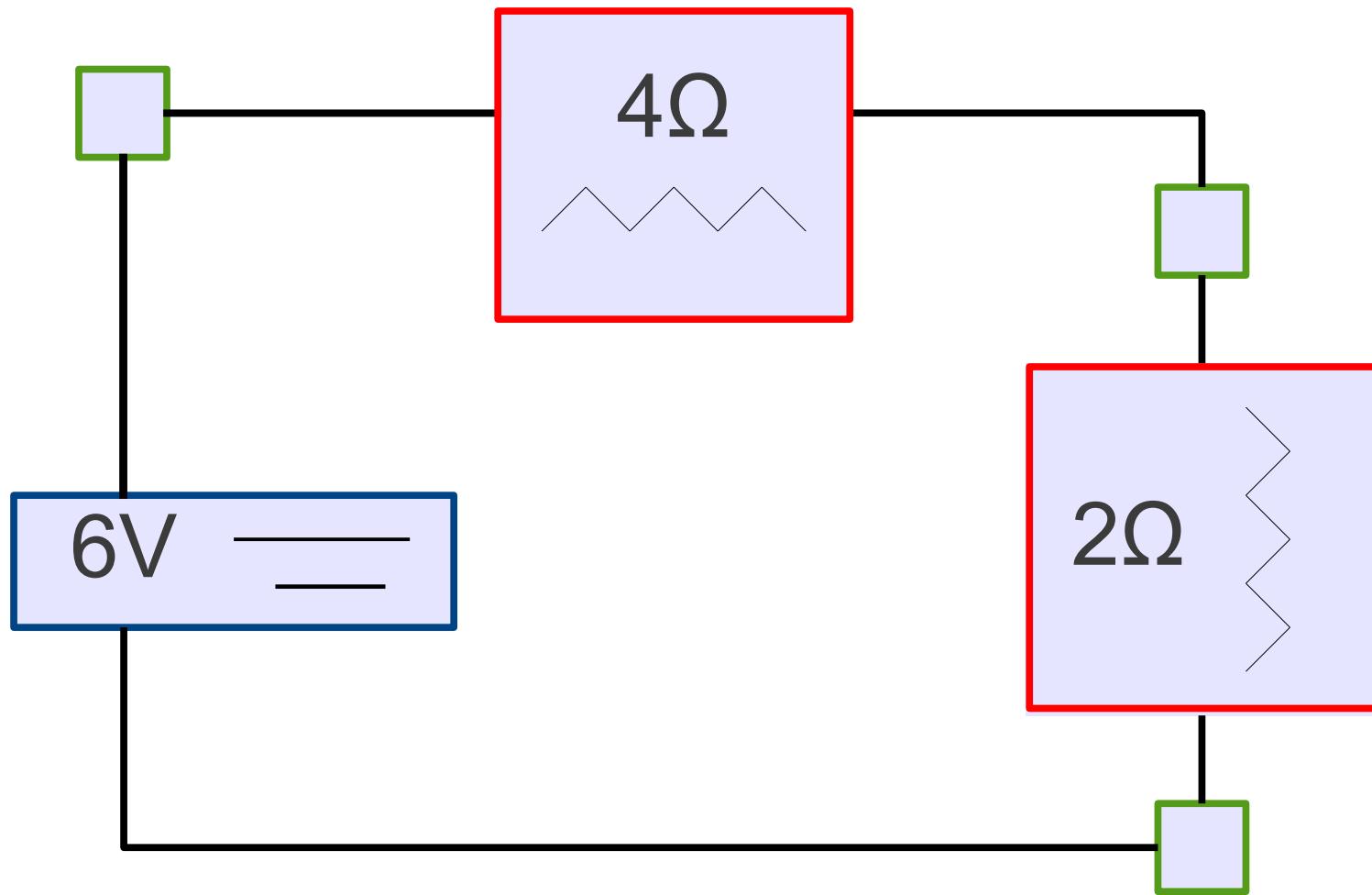




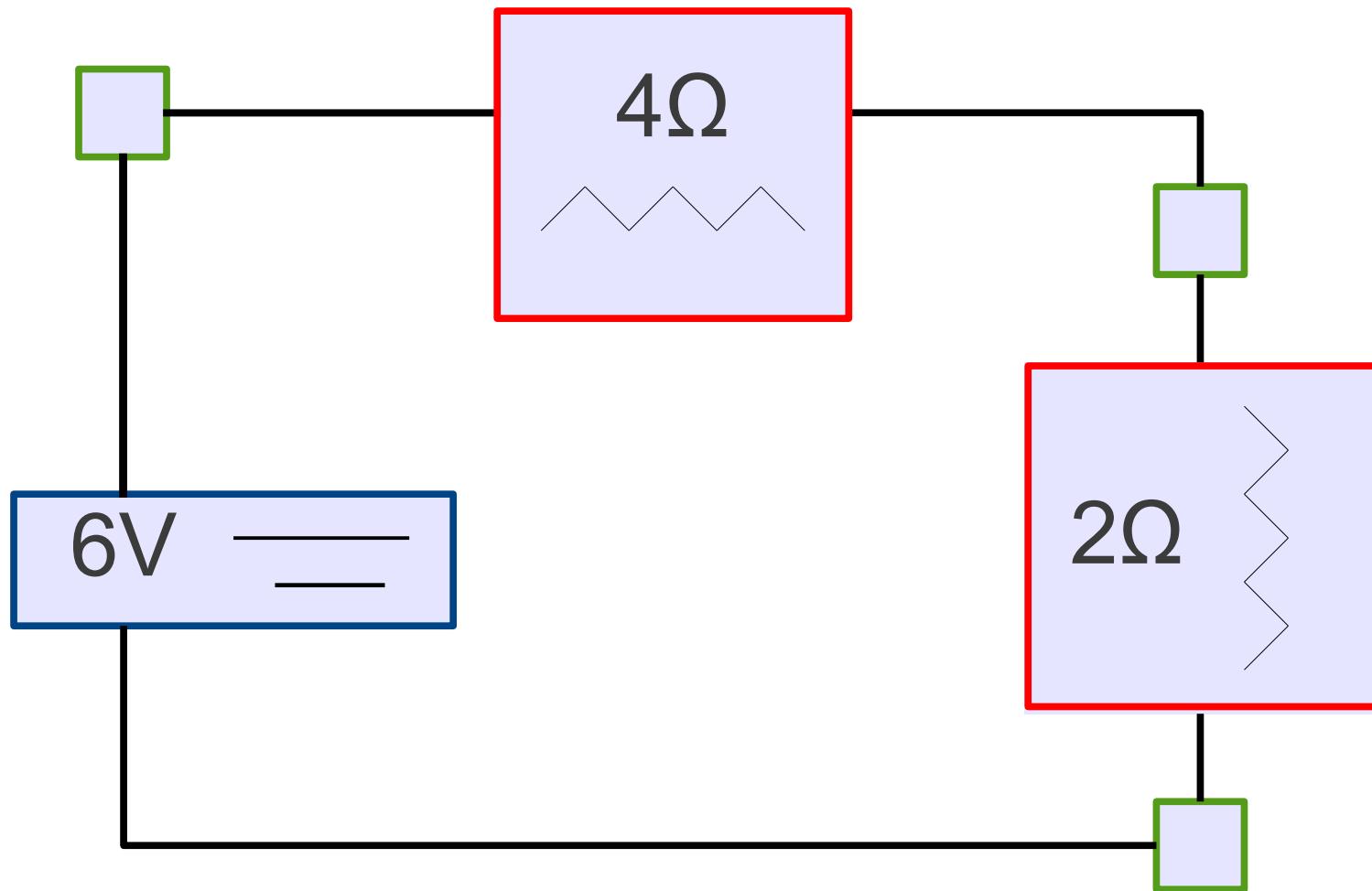


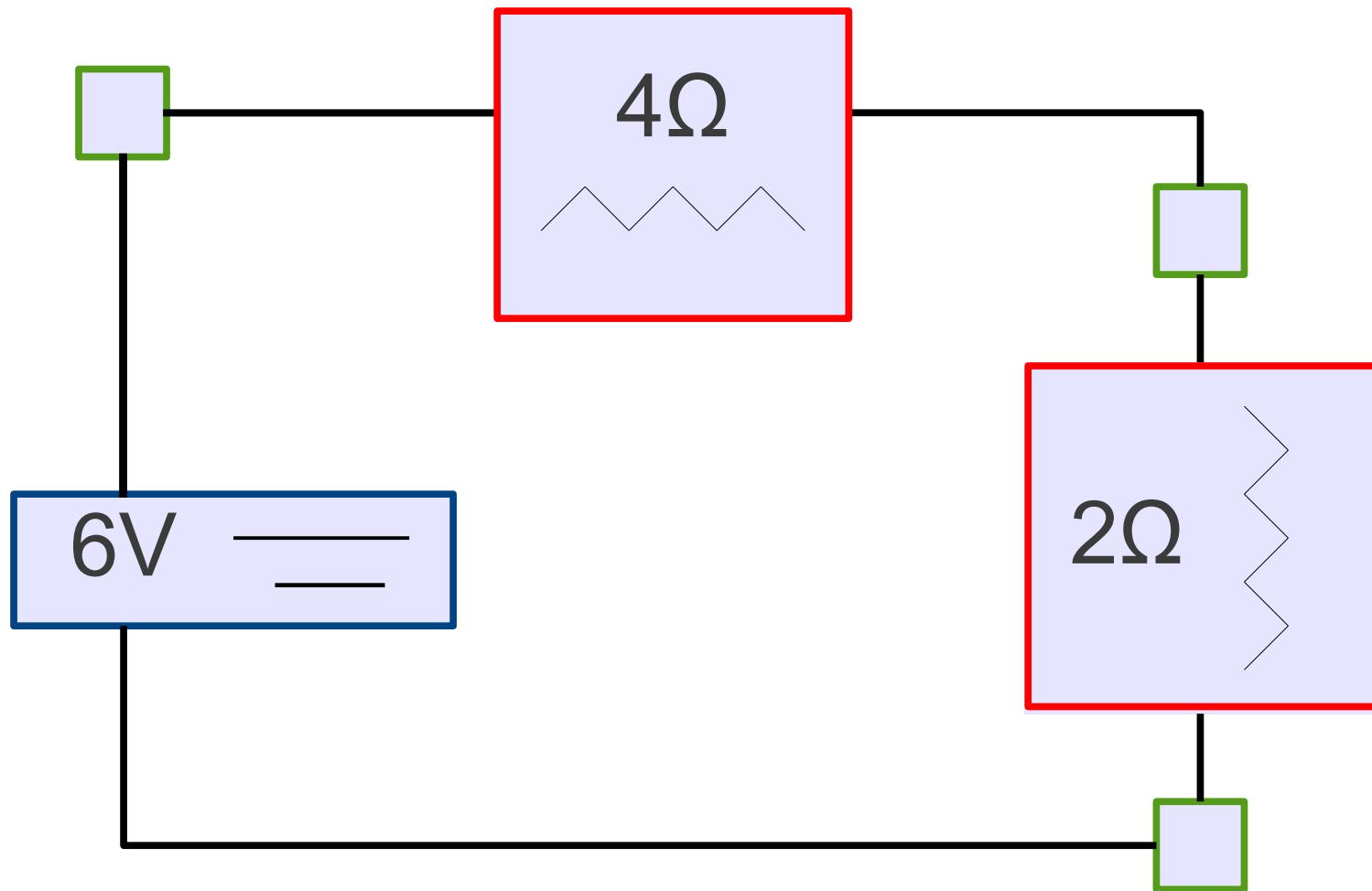


$$\frac{1}{3\Omega} + \frac{1}{6\Omega} = 2 \Omega$$

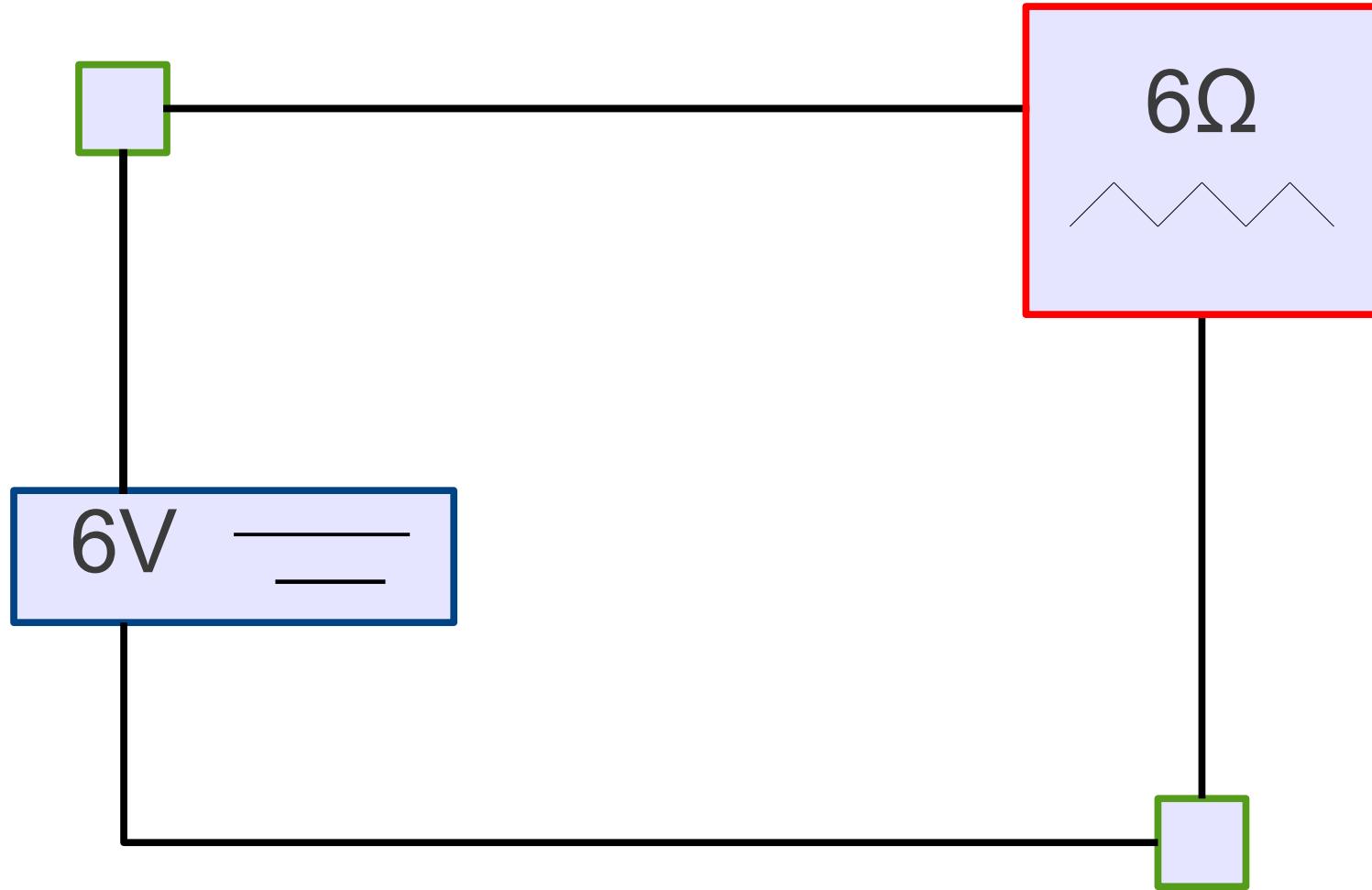


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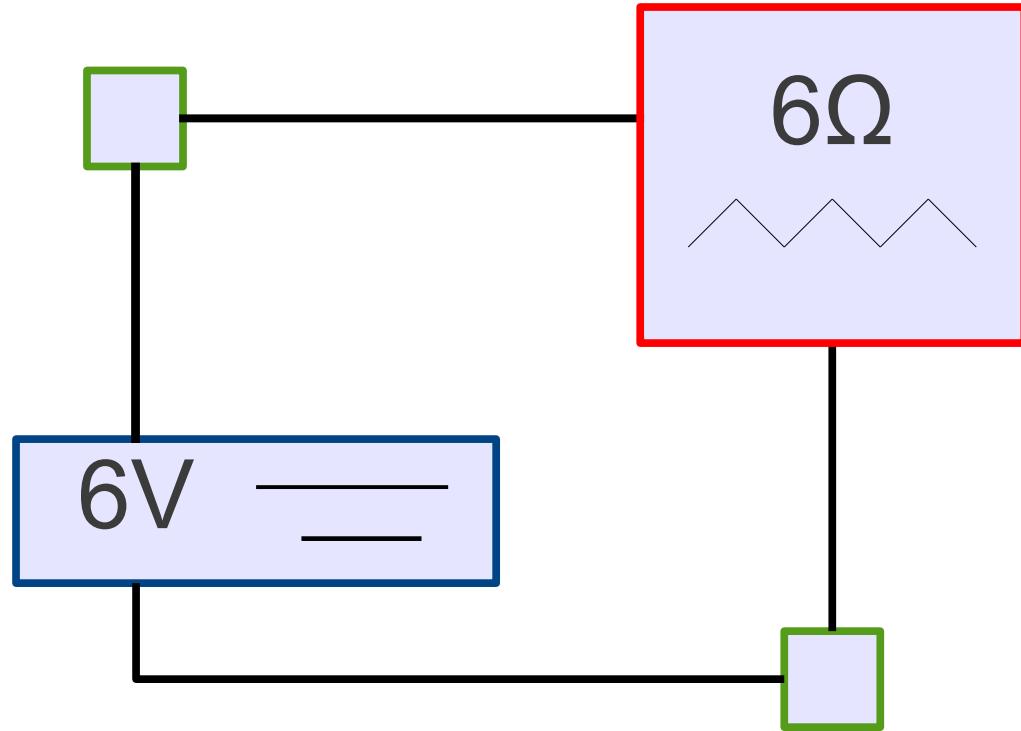


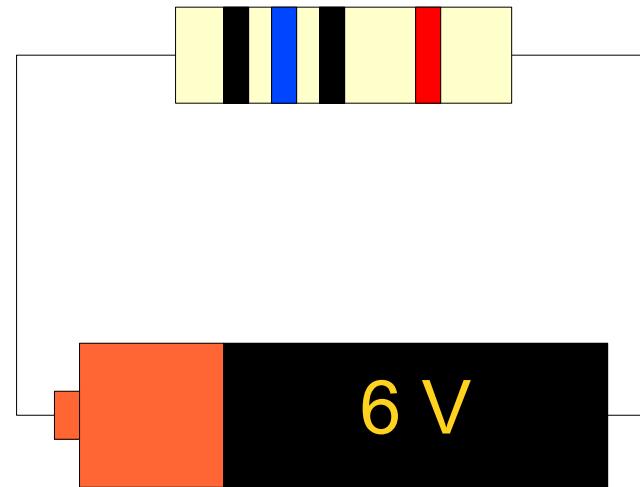
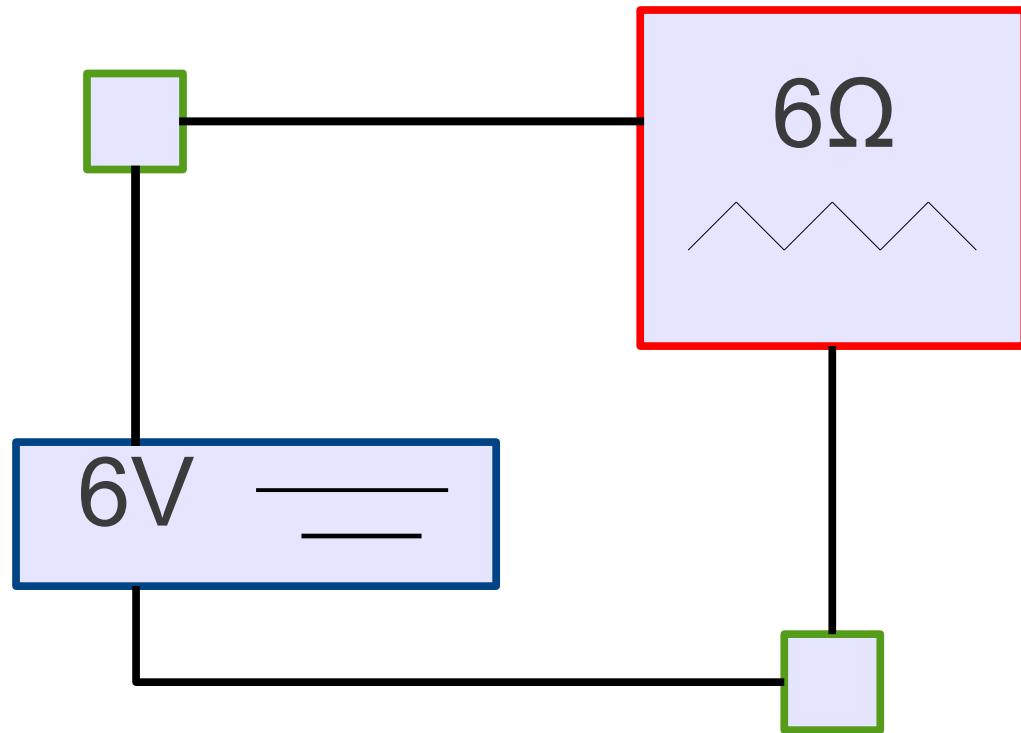


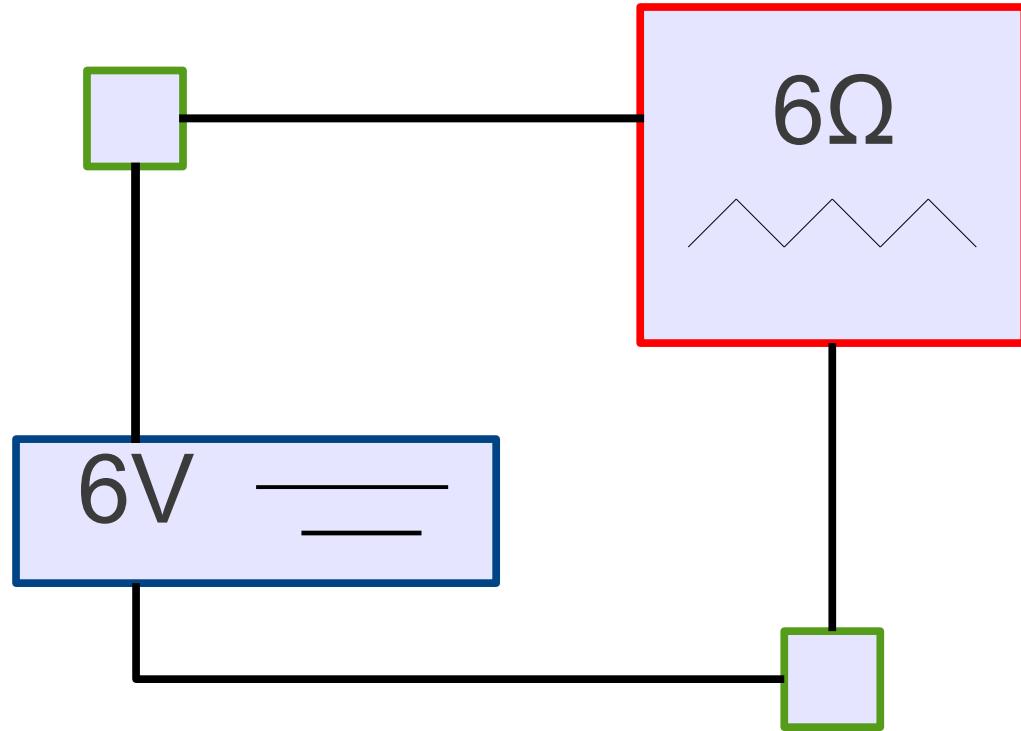
$$4\Omega + 2\Omega = 6 \Omega$$



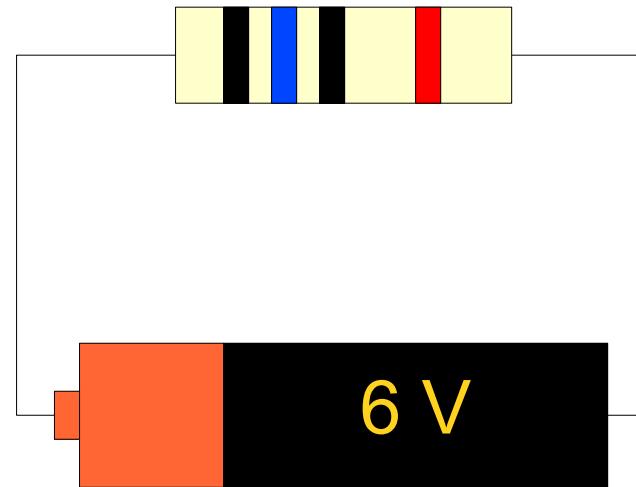
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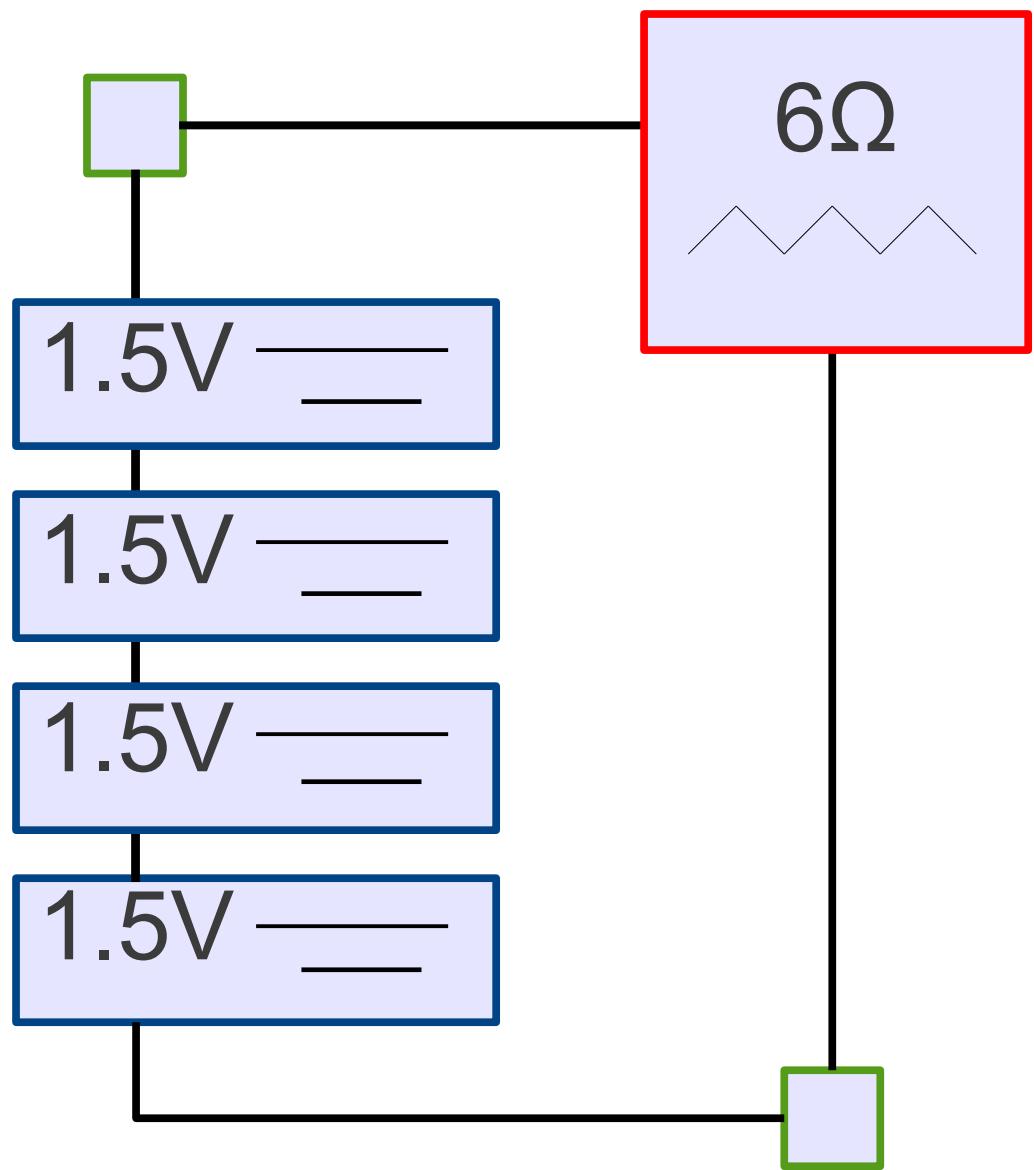


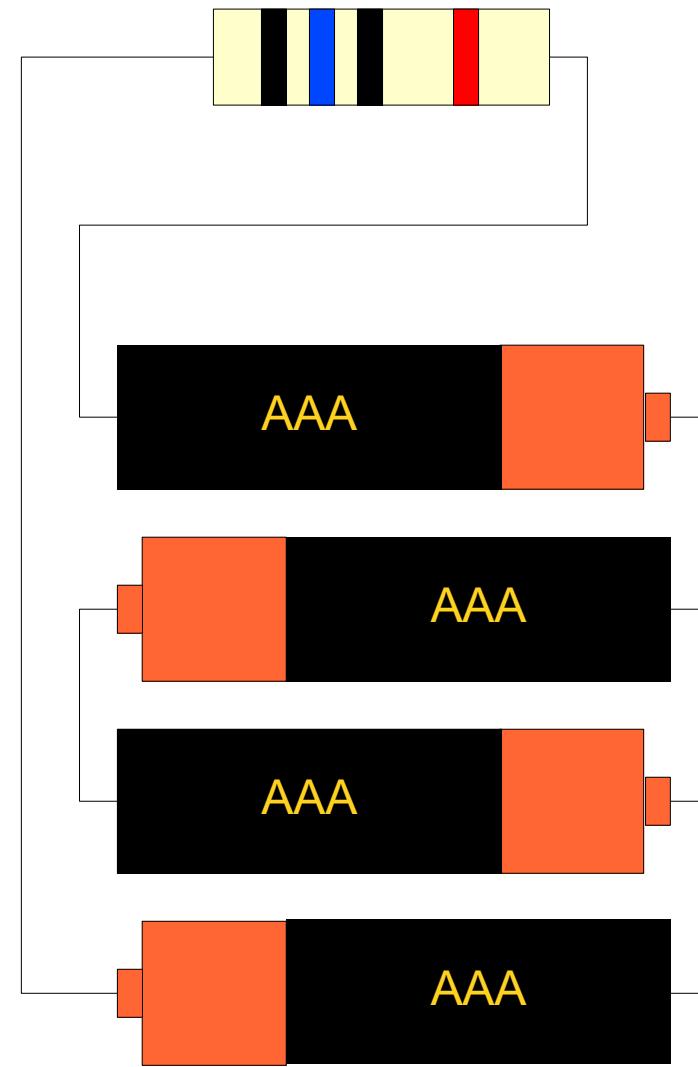
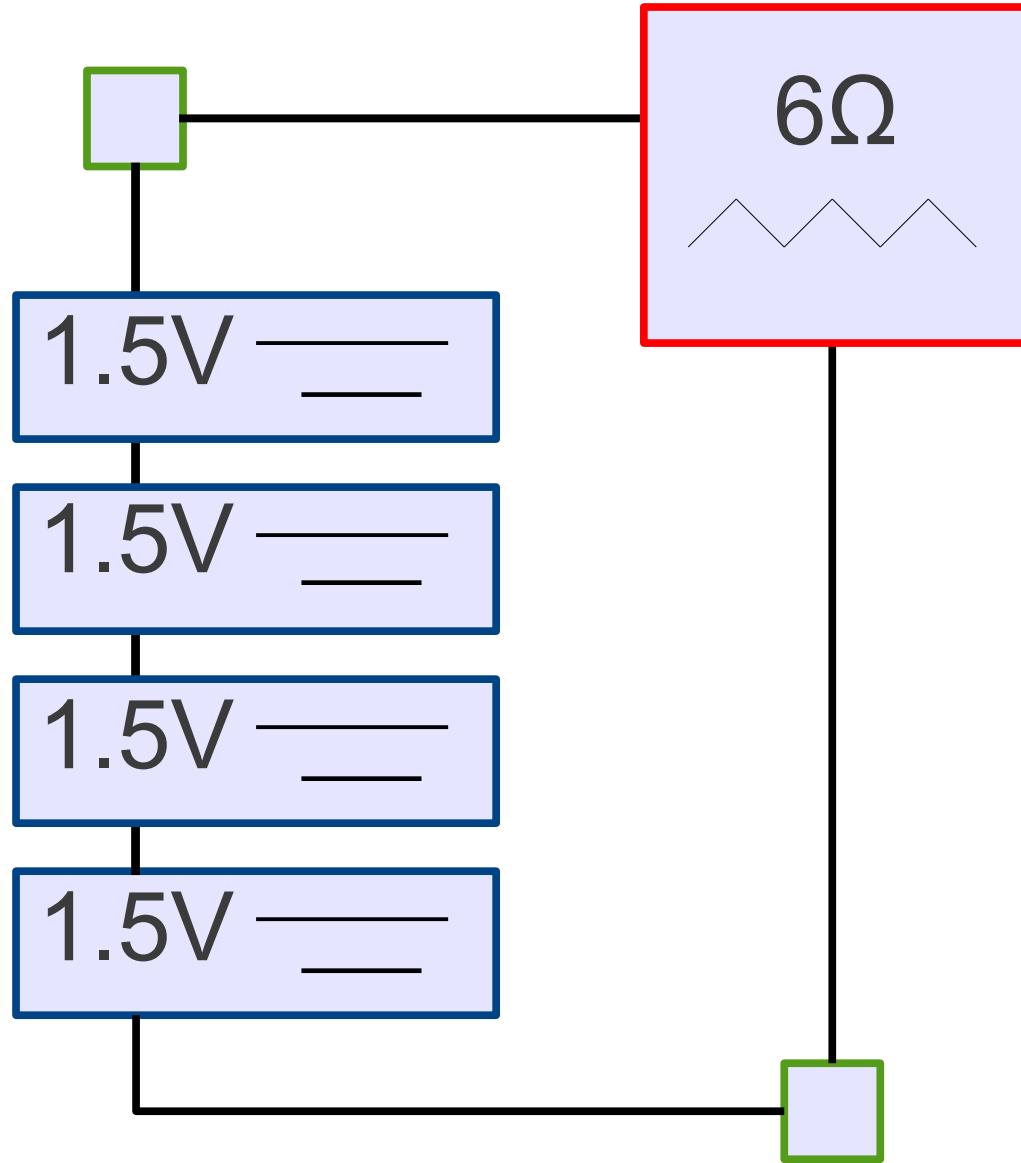


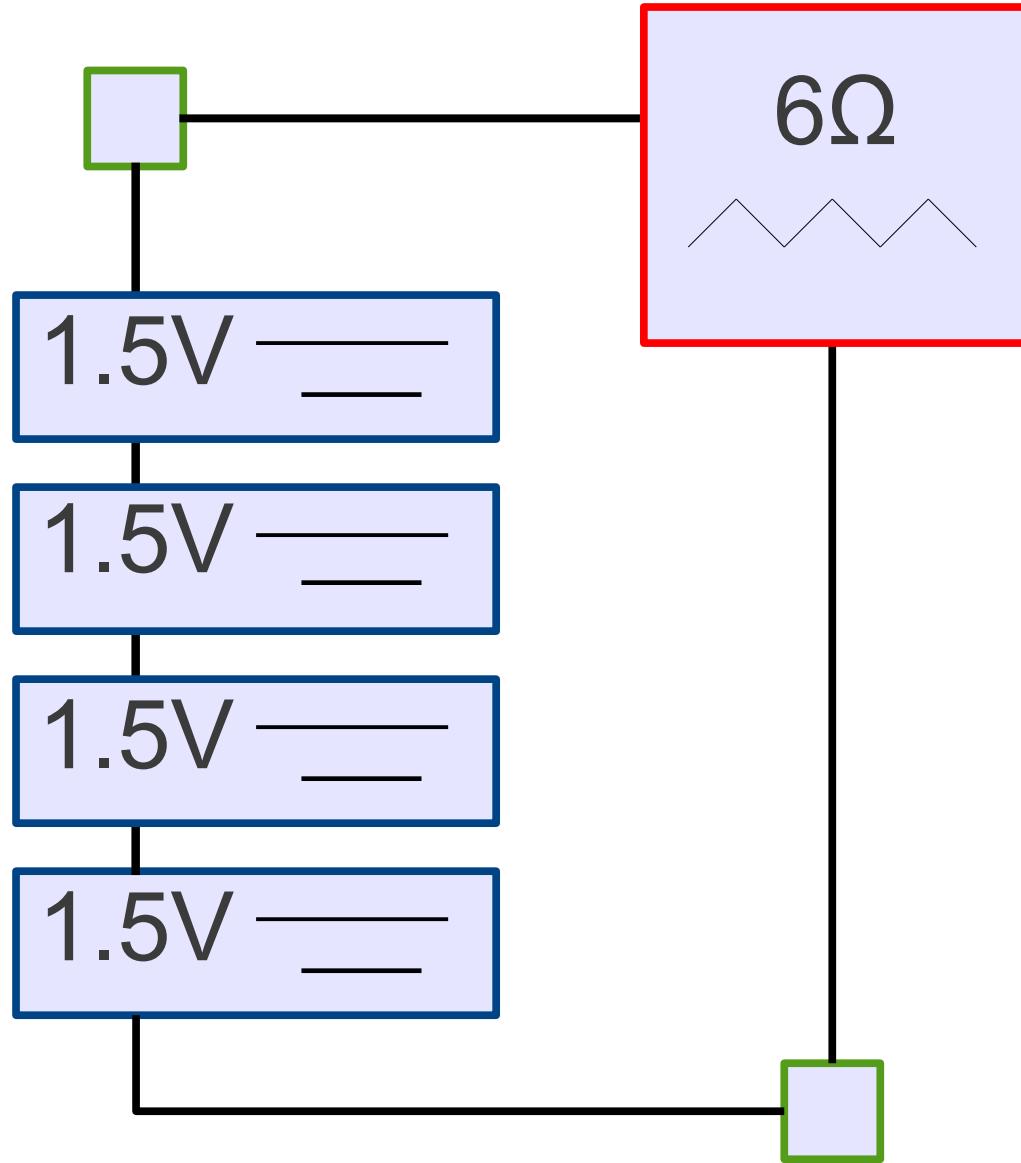


Total Cost: \$4.75

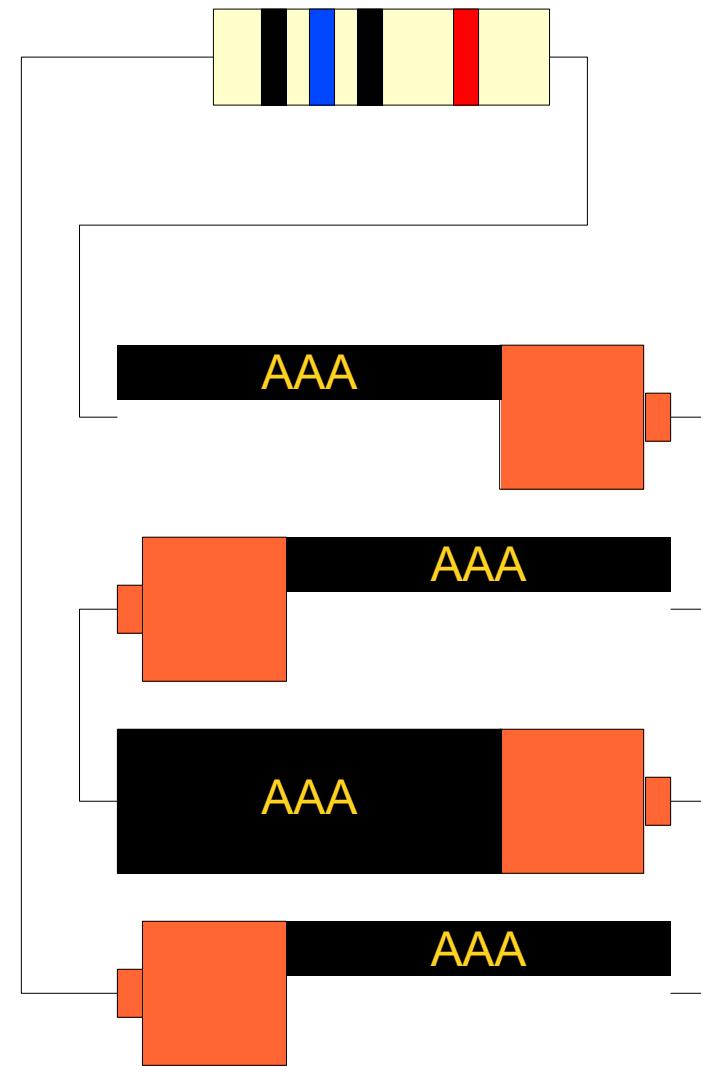








Total Cost: \$1.00



From Description to Implementation

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- **Lexical analysis (Scanning):** Identify logical pieces of the description.

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- **IR Generation:** Design one possible structure.
- **IR Optimization:** Simplify the intended structure.
- **Generation:** Fabricate the structure.
- **Optimization:** Improve the resulting structure.

Lexical Analysis

- First step: recognize words.
 - Smallest unit above letters

This is a sentence.

- Lexical analysis is not trivial. Consider:
ist his ase nte nce

Real world Example: [Watch it On YouTube!](#)



- Lexical analyzer divides program text into “words” or “tokens”

If $x == y$ then $z = 1$; else $z = 2$;

- Units:

- Once words are understood, the next step is to understand sentence structure
- Parsing = Diagramming Sentences
 - The diagram is a tree

This line is a longer sentence

article noun verb article adjective noun

subject

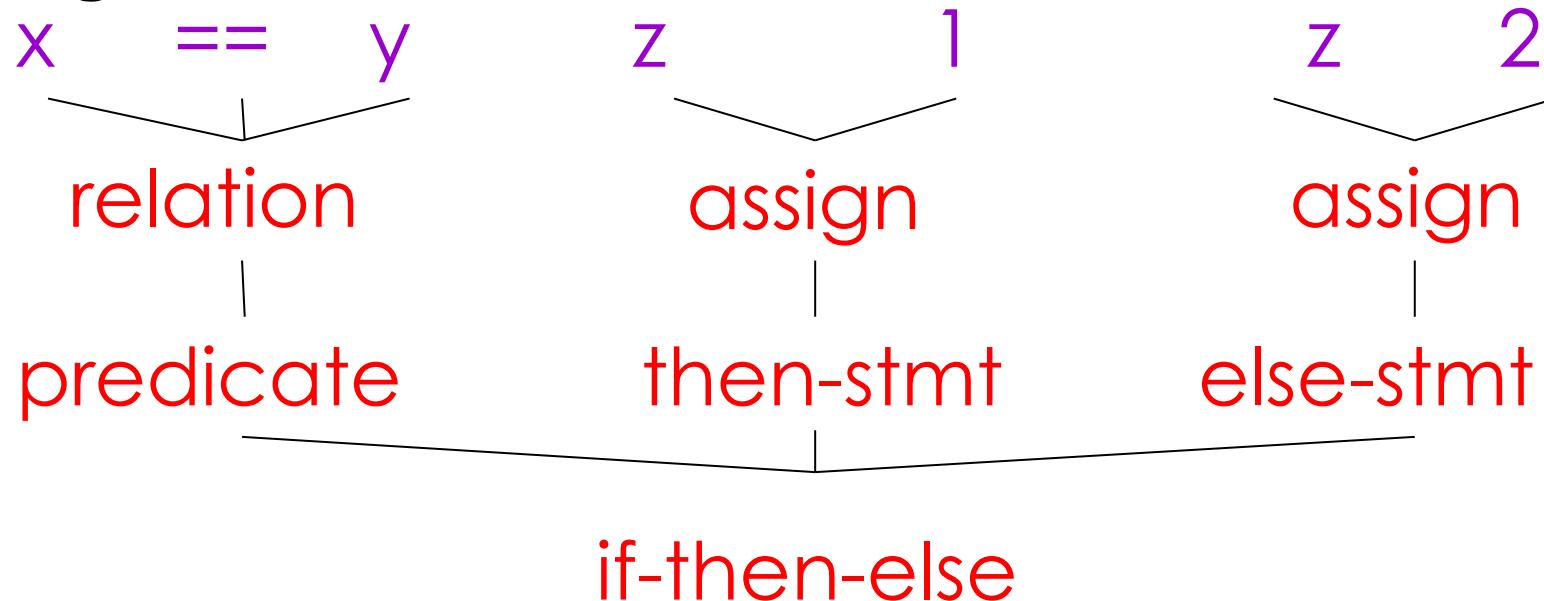
object

sentence

- Parsing program expressions is the same
- Consider:

If $x == y$ then $z = 1$; else $z = 2$;

- Diagrammed:



- Once sentence structure is understood, we can try to understand “meaning”
 - But meaning is too hard for compilers
- Compilers perform limited semantic analysis to catch inconsistencies

Semantic Analysis in English

- Example:

Jack said Jerry left his assignment at home.

What does “his” refer to? Jack or Jerry?

- Even worse:

Jack said Jack left his assignment at home?

How many Jacks are there?

Which one left the assignment?

Semantic Analysis in Programming

- Programming languages define strict rules to avoid such ambiguities
 - This C++ code prints “4”; the inner definition is used
- ```
{
 int Jack = 3;
 {
 int Jack = 4;
 cout << Jack;
 }
}
```

## More Semantic Analysis

- Compilers perform many semantic checks besides variable bindings
- Example:

Jack left her homework at home.
- Possible type mismatch between **her** and **Jack**
  - If Jack is male

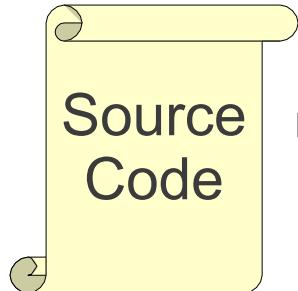
# Optimization

- No strong counterpart in English, but akin to editing
- Automatically modify programs so that they
  - Run faster
  - Use less memory
  - In general, to use or conserve some resource

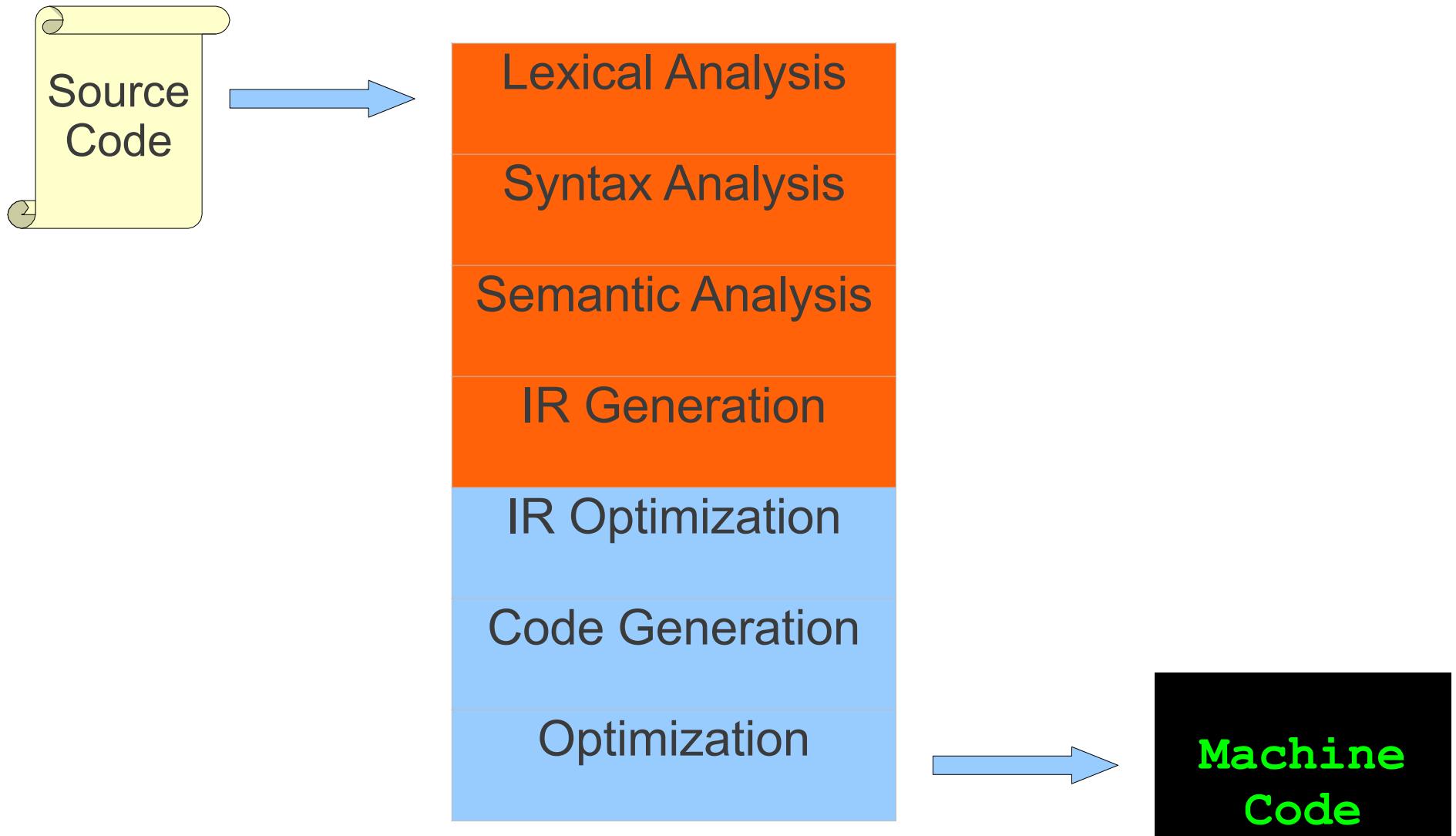
# Optimization Example

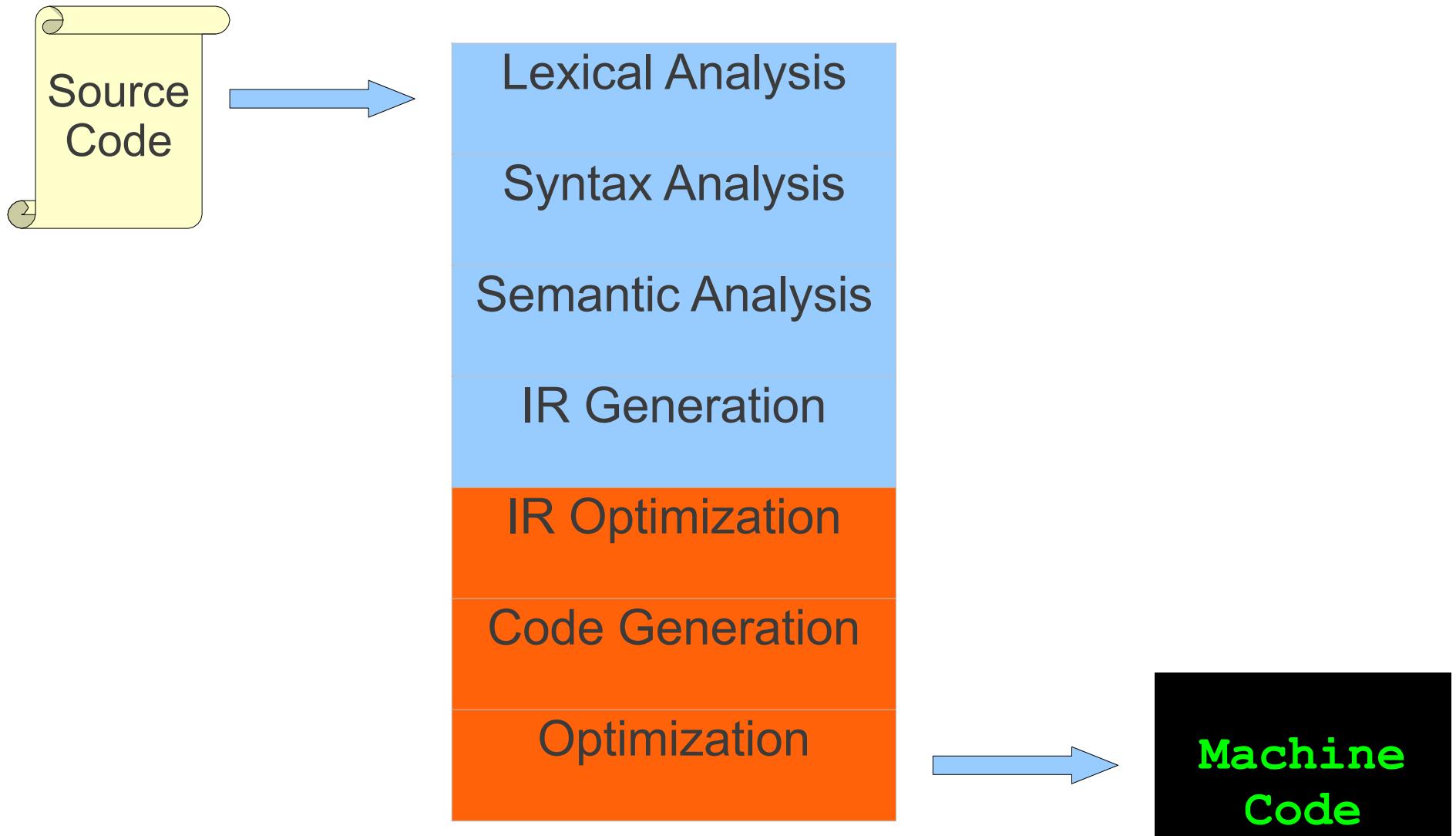
$X = Y * 0$  is the same as  $X = 0$

(the \* operator is annihilated by zero)



**Machine  
Code**





```
while (y < z) {
 int x = a + b;
 y += x;
}
```

Lexical Analysis

Syntax Analysis

Semantic Analysis

IR Generation

IR Optimization

Code Generation

Optimization

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Lexical Analysis

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Optimization

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```

```
T_While
T_LeftParen
T_Identifier y
T_Less
T_Identifier z
T_RightParen
T_OpenBrace
T_Int
T_Identifier x
T_Assign
T_Identifier a
T_Plus
T_Identifier b
T_Semicolon
T_Identifier y
T_PlusAssign
T_Identifier x
T_Semicolon
T_CloseBrace
```

Lexical Analysis

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IR Generation

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Lexical Analysis

Syntax Analysis

Semantic Analysis

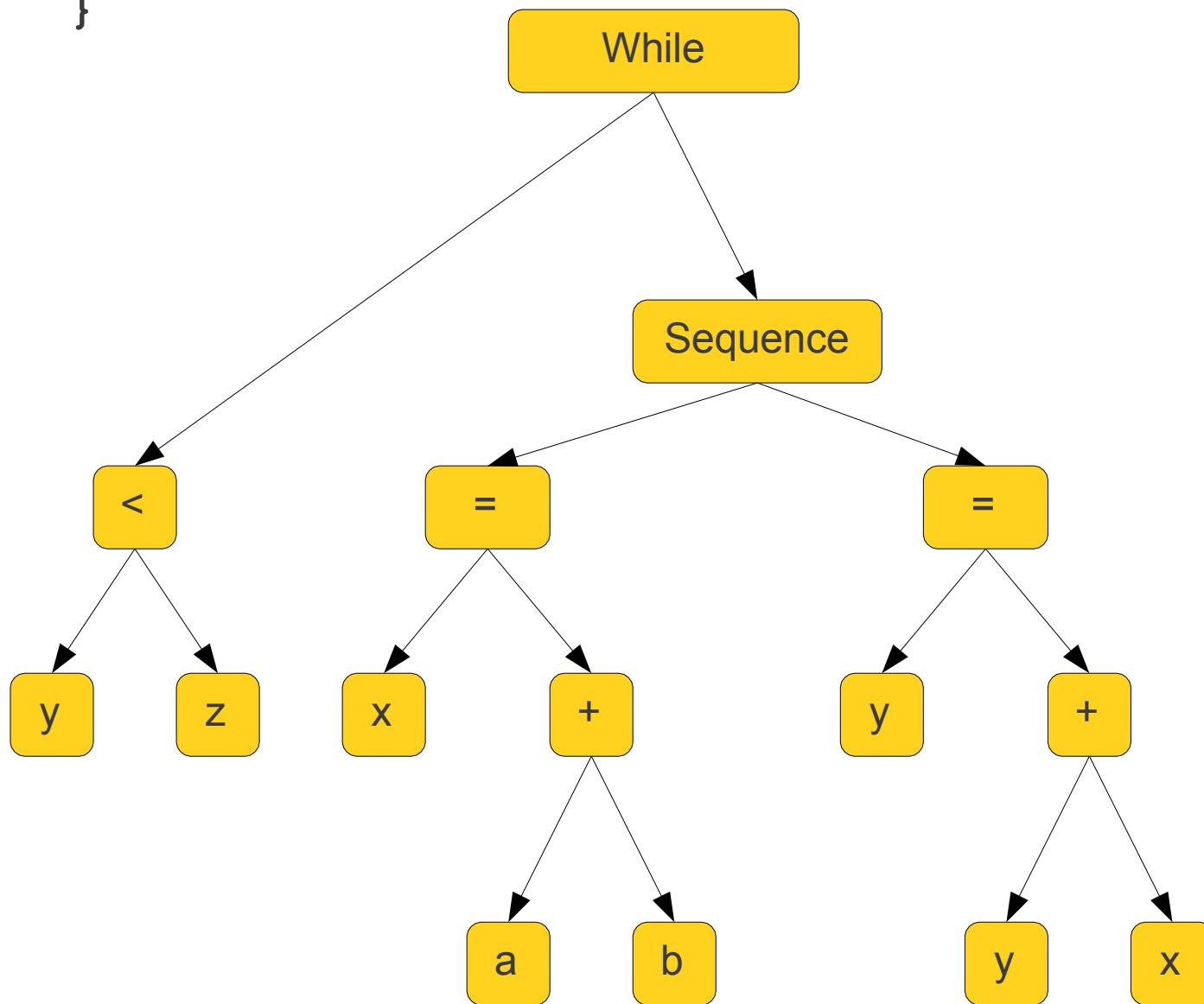
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while (y < z) {
 int x = a + b;
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```



Lexical Analysis

Syntax Analysis

Semantic Analysis

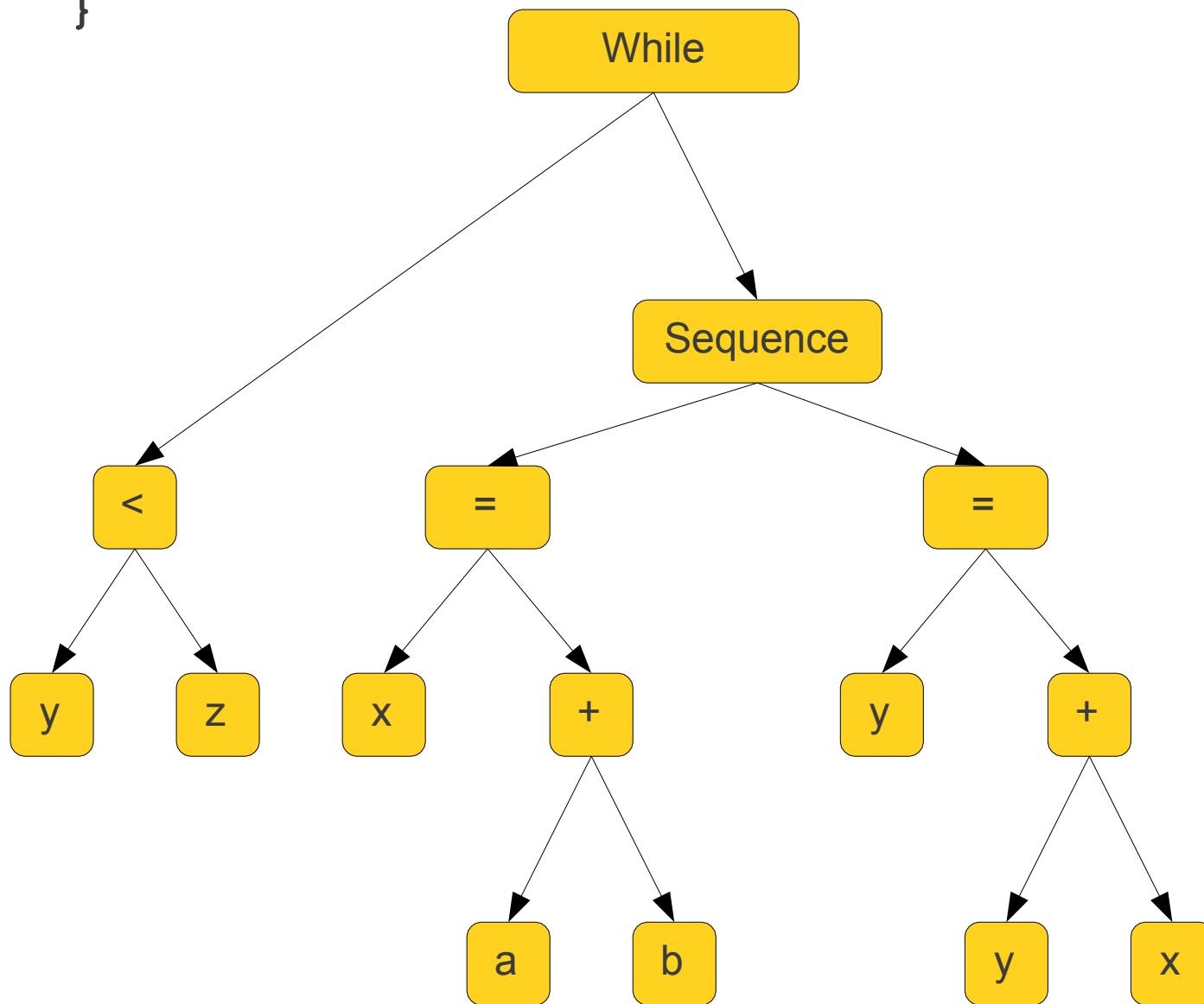
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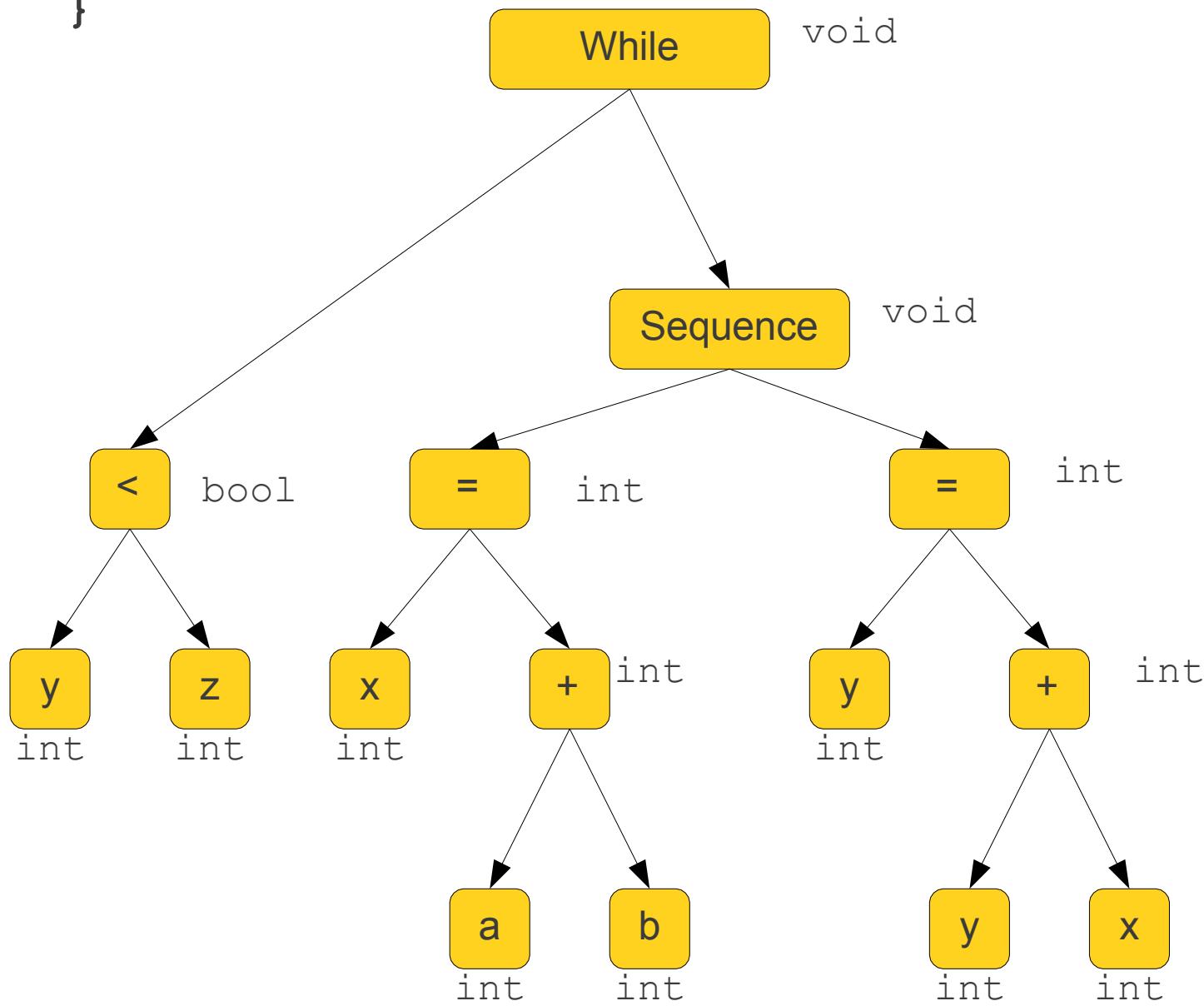
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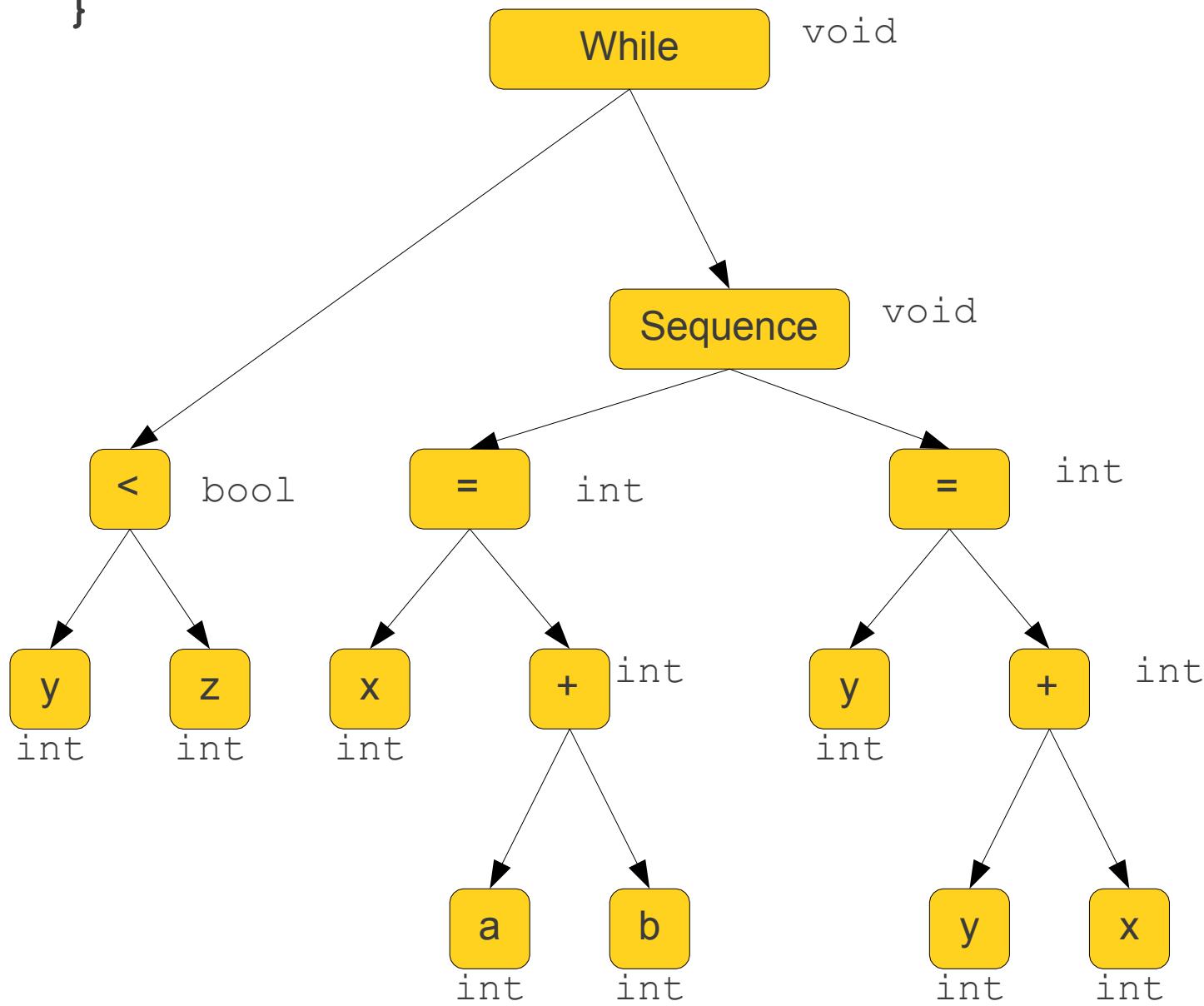
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Loop: add $1, $2, $3
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 slt $6, $4, $5
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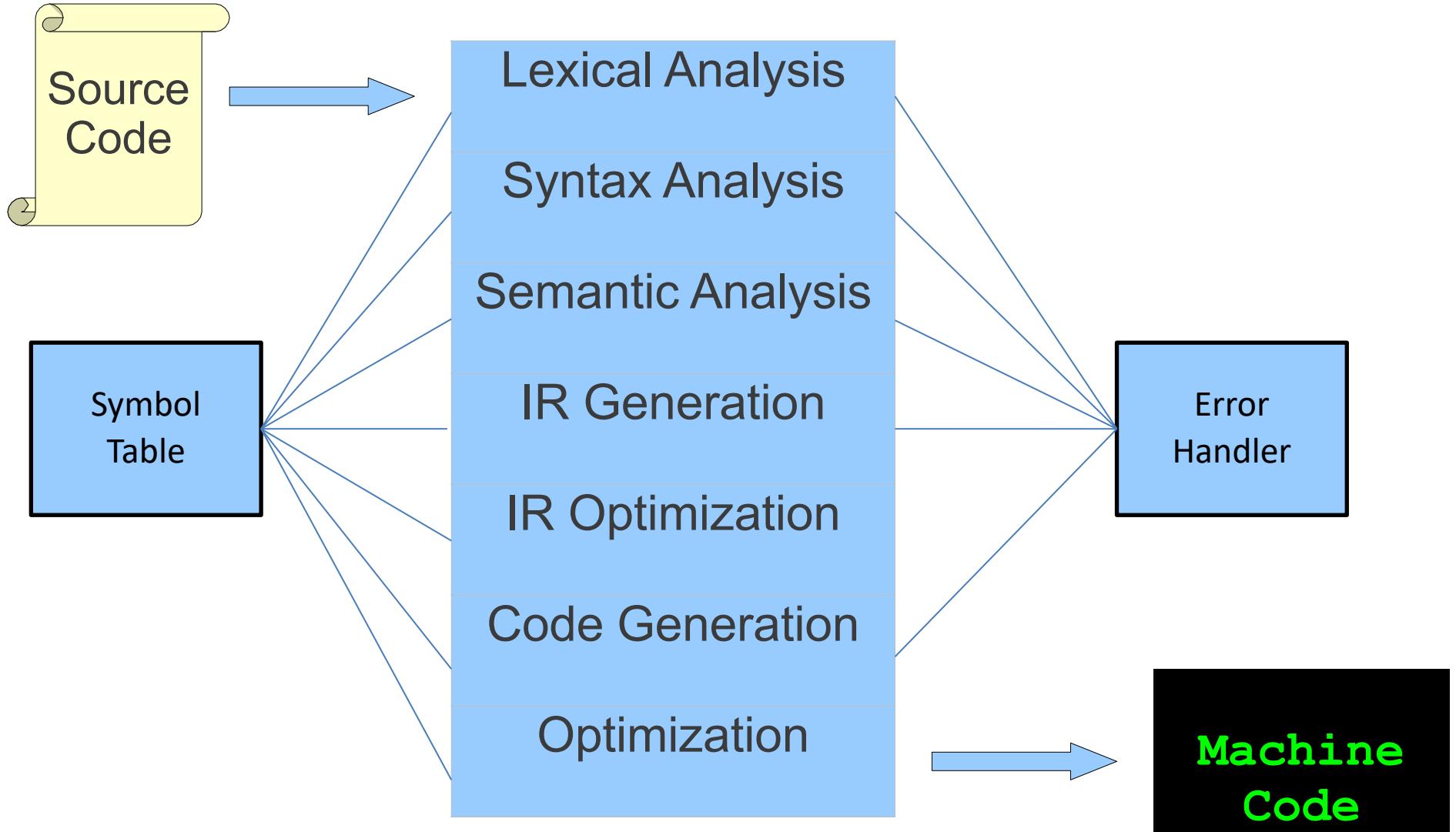
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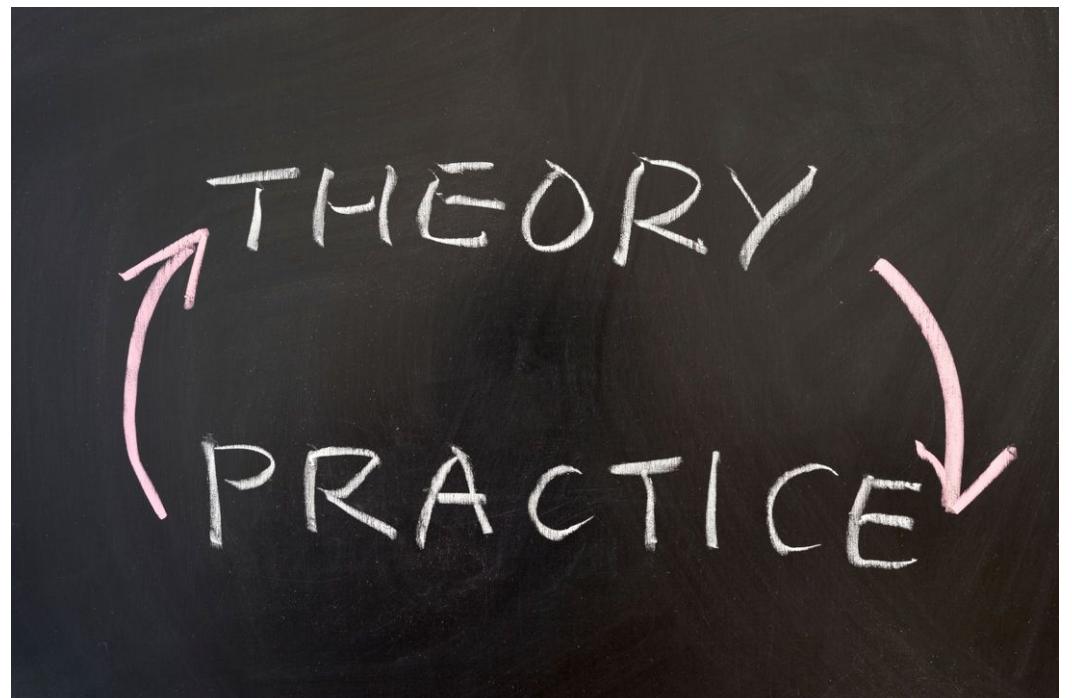
# Course Theme

- Studying design and implementation a very complicated software.
- We break it apart!
- We are going to meet weird-useful language features.



# Course Theme

- We use theory as much as possible to survive.
- Excellency at programming is needed.



# Why Take Programming Languages and Compilers?

To appreciate the **marriage** of theory and practice



**“Theory and practice are not mutually exclusive;  
they are intimately connected. They live together  
and support each other.”**

[D.E. Knuth, 1989]

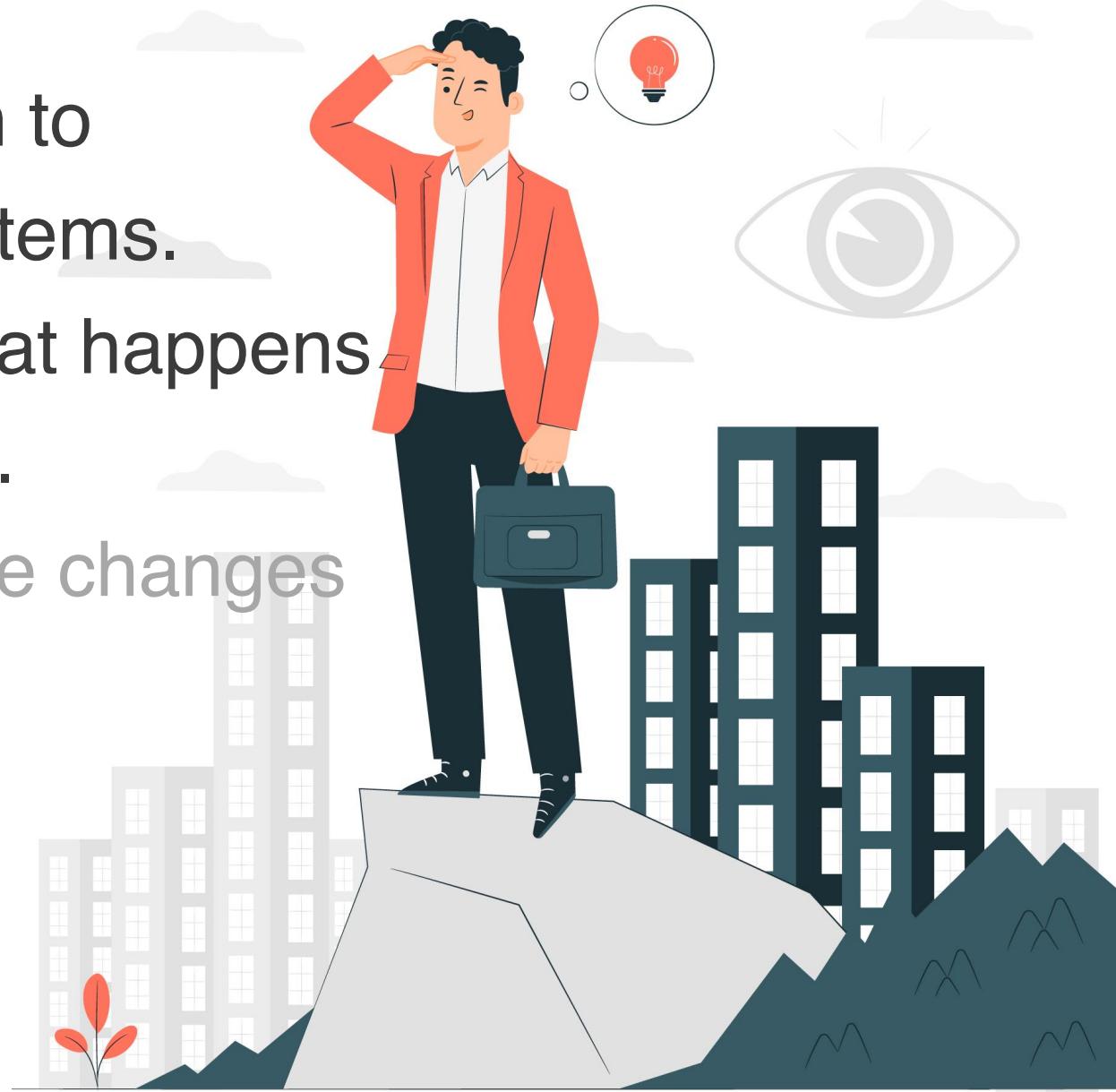
# Why Study Compilers?



# Why Study Compilers?

We Seek for a vision to

- build better systems.
- Understand what happens under the hood.
- Cope with future changes and era.

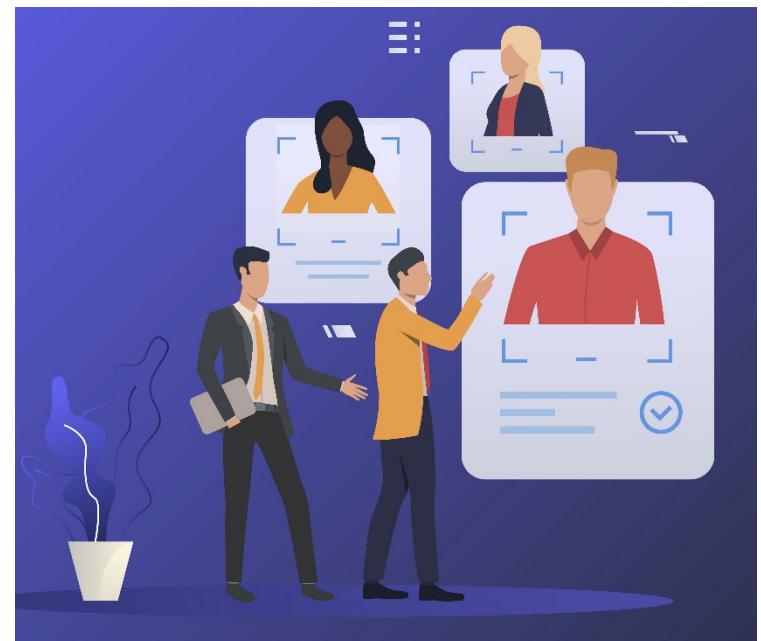


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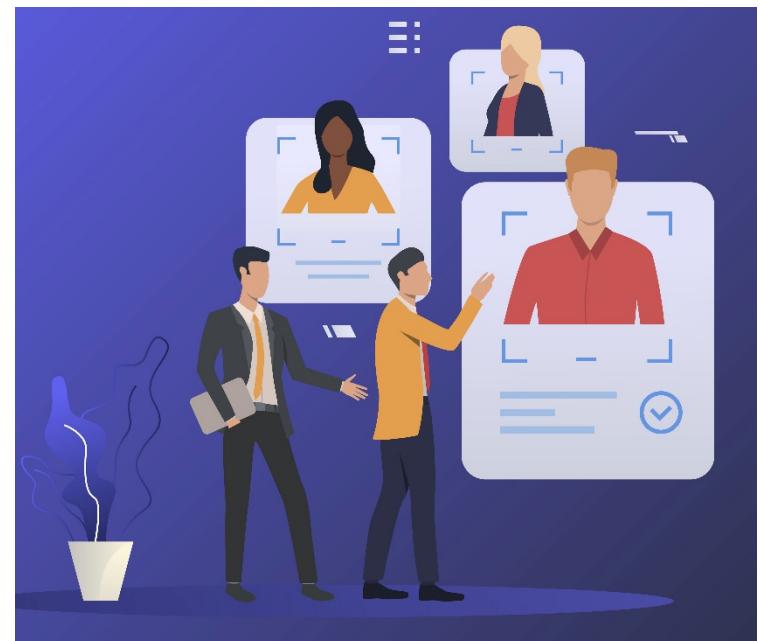
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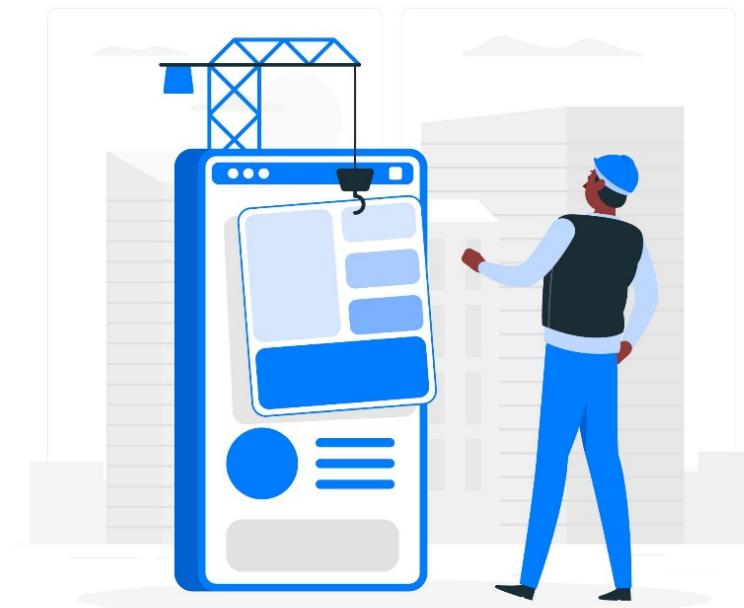


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- See theory **come to life.**

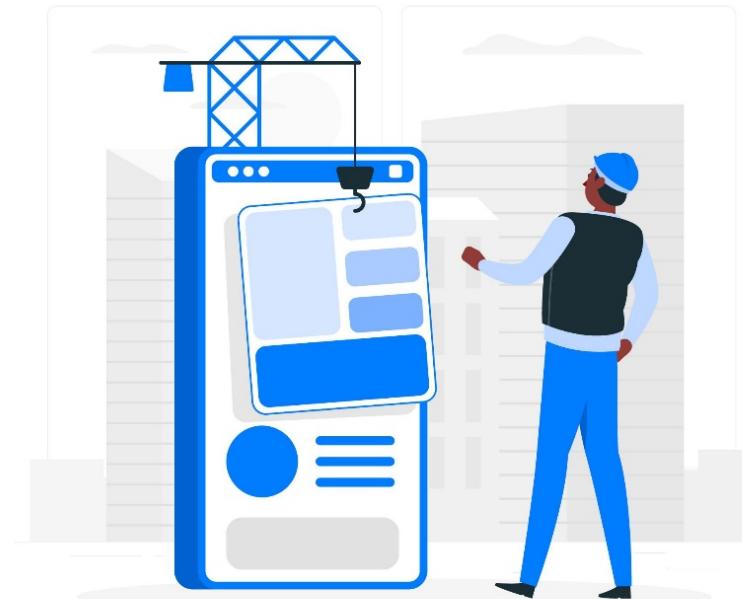


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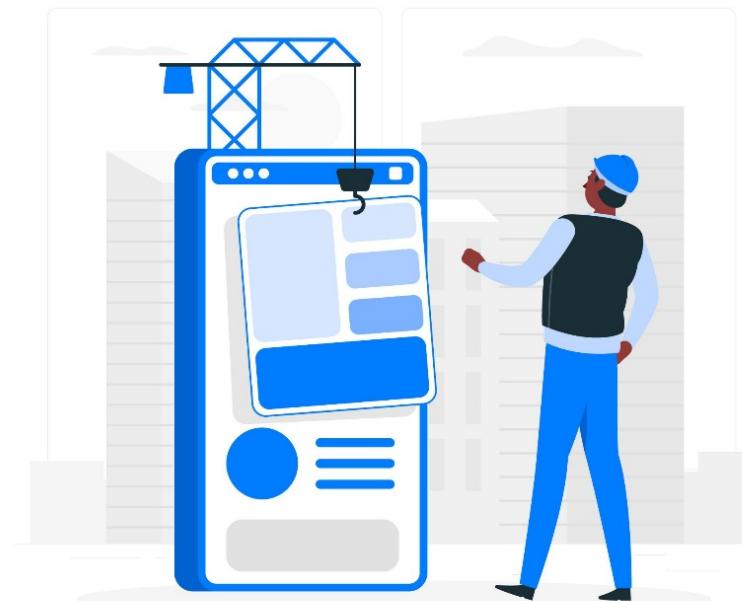
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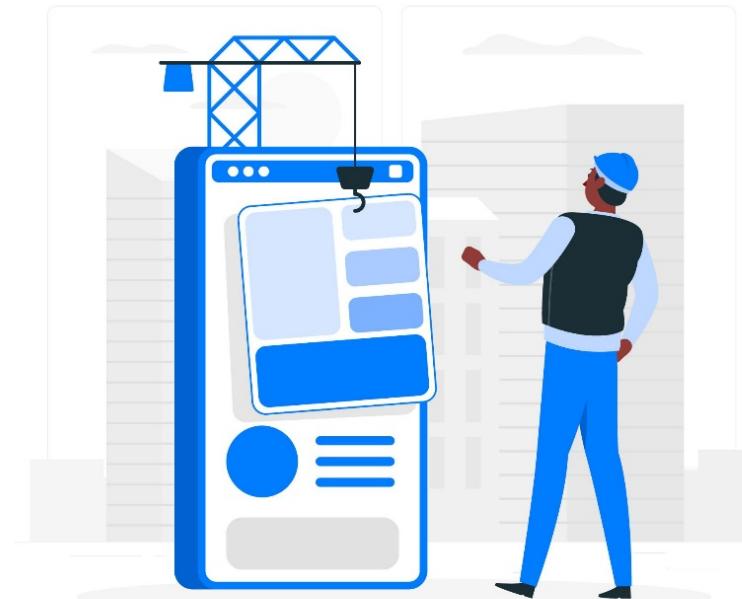
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- Learn how to **build programming languages**.
- Learn **how programming languages work**.
- Learn **tradeoffs in language design**.



# Why Study Compilers?



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**This Course Adapted from  
Stanford CS 143  
And  
MIT CS 6.s081  
(but with changes!)**

**TA: Soroosh Zare**  
**(ihaveint [] gmail)**



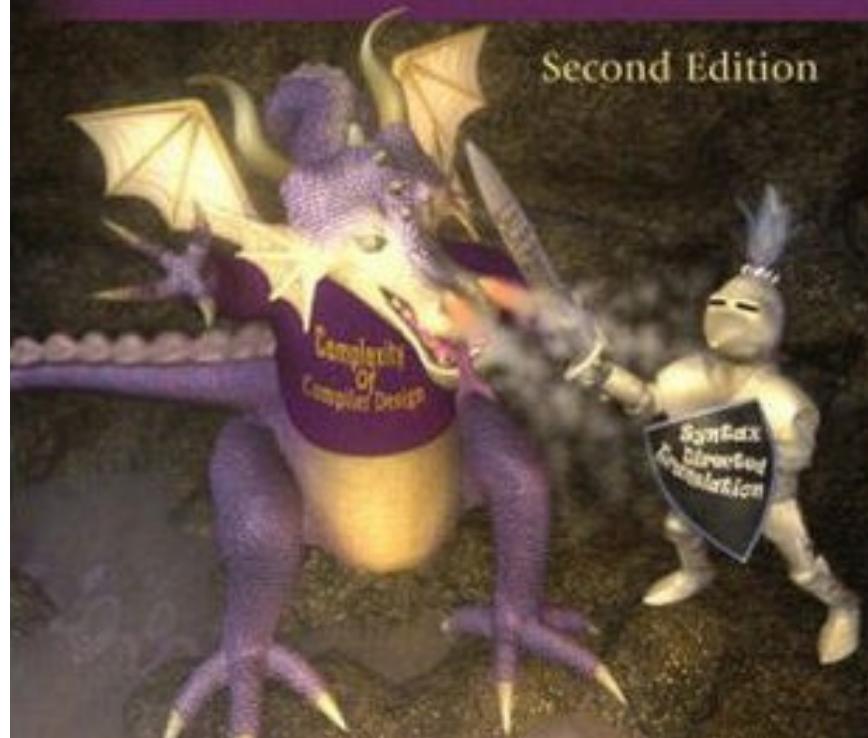
**Course Page:** [http://foroughmand.ir/?page\\_id=819](http://foroughmand.ir/?page_id=819)

- Course info
- Important Announcements
- Calendar
- Some resources

# Compilers

*Principles, Techniques, & Tools*

Second Edition



Alfred V. Aho  
Monica S. Lam  
Ravi Sethi  
Jeffrey D. Ullman



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- Quera for project, CW for homeworks and exams.

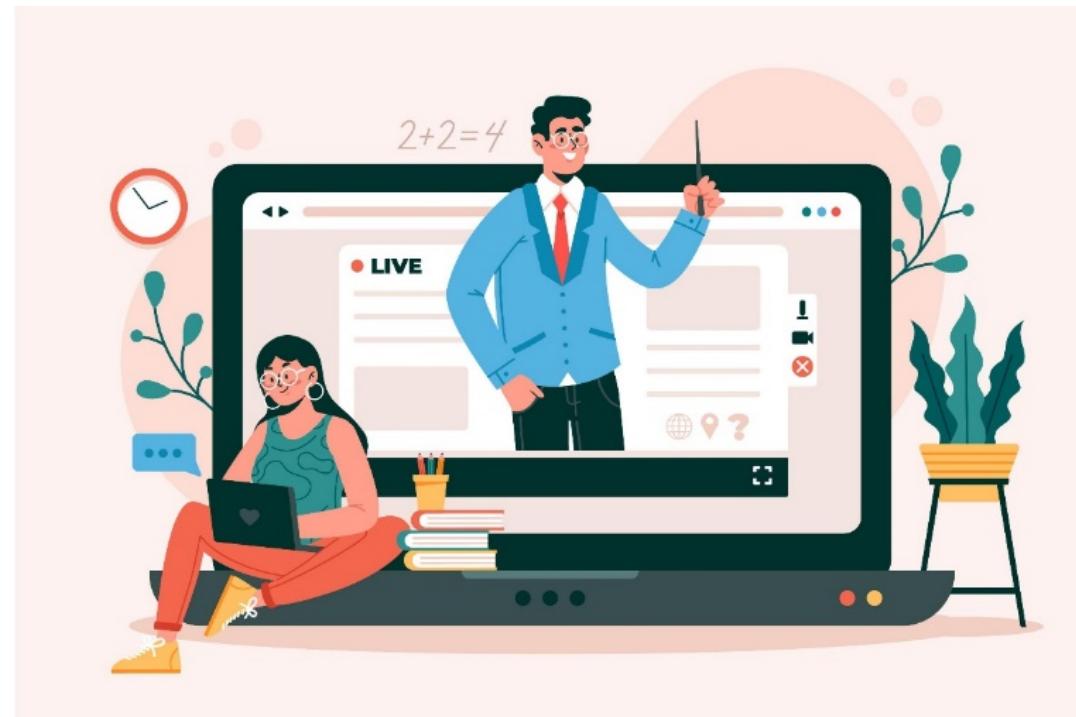


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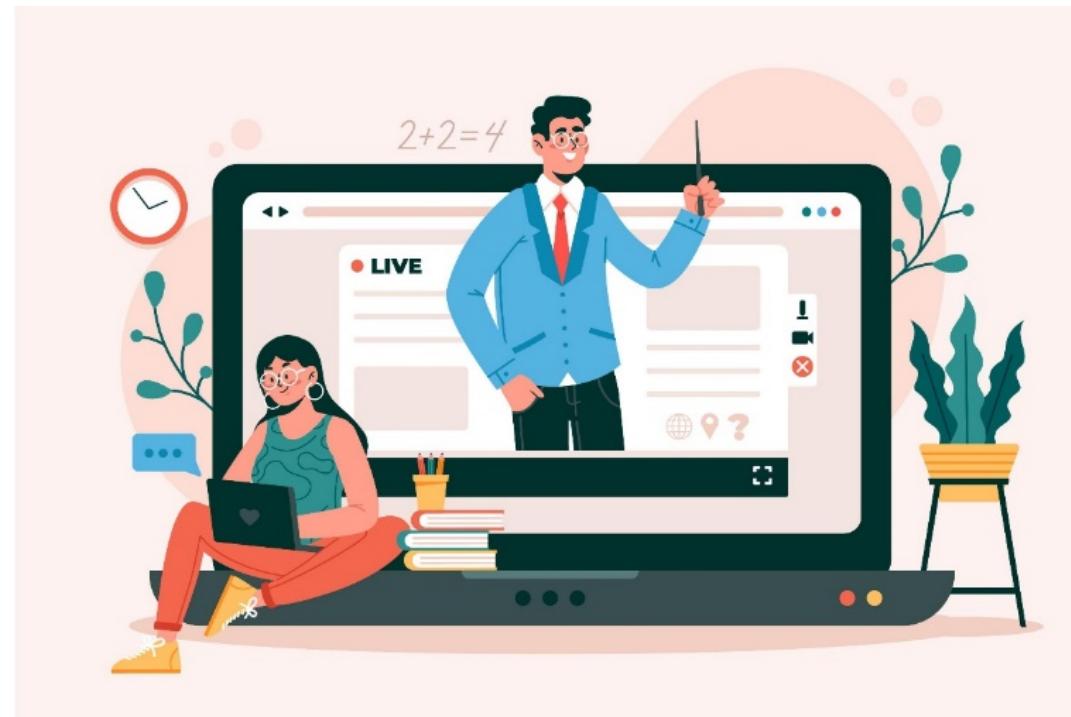




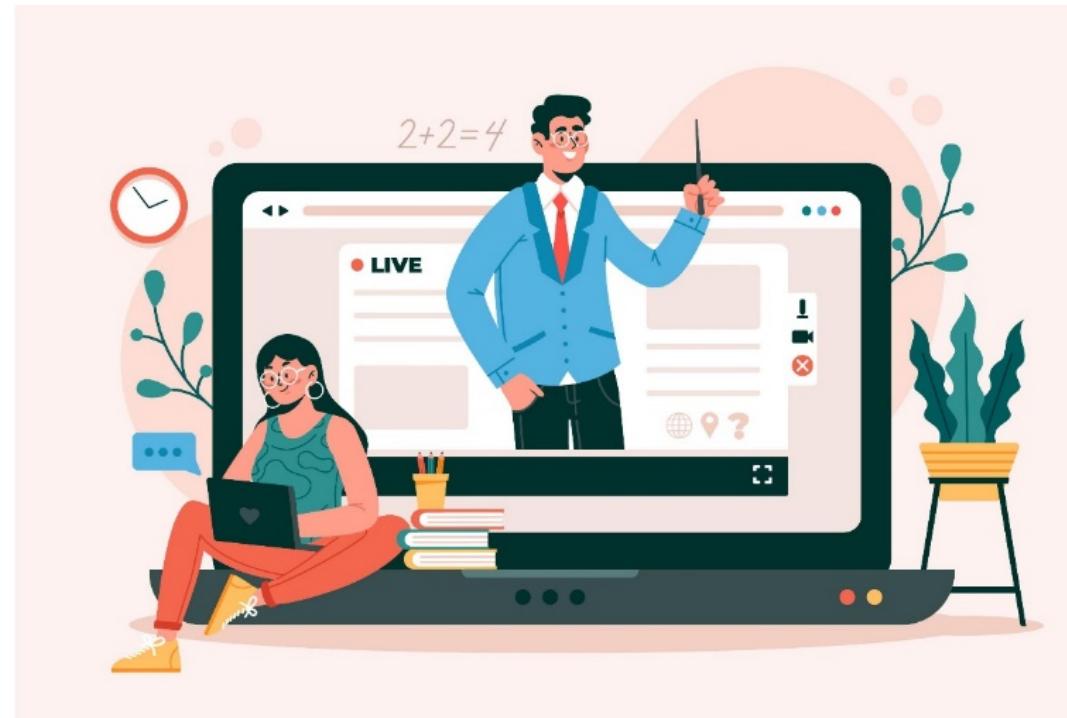
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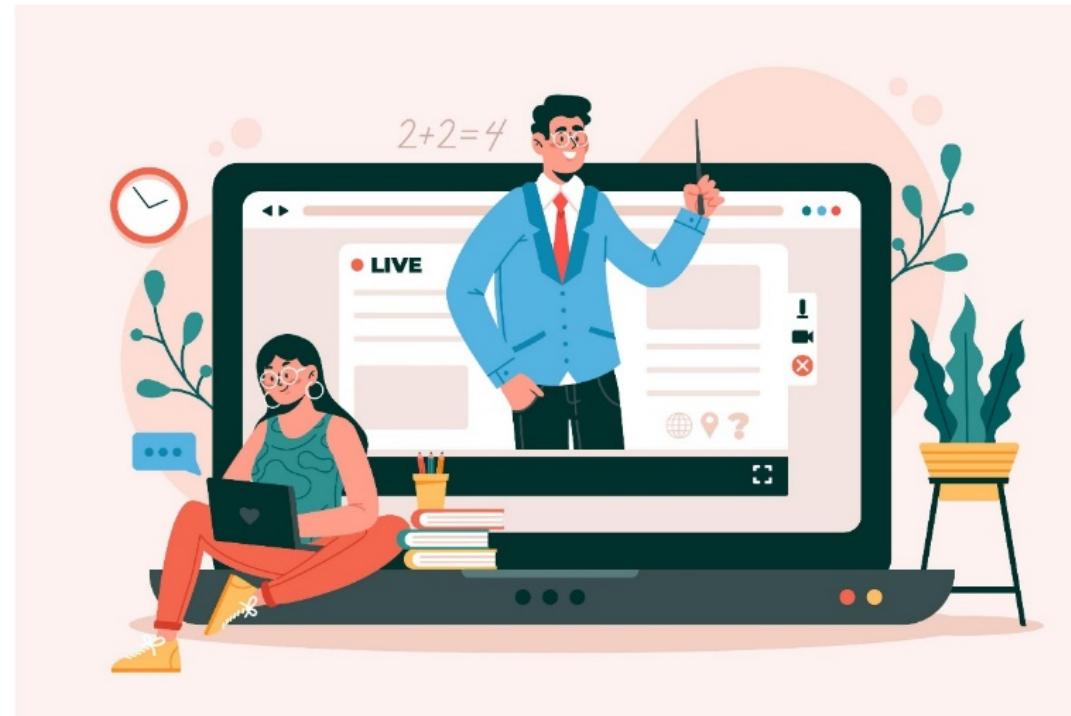
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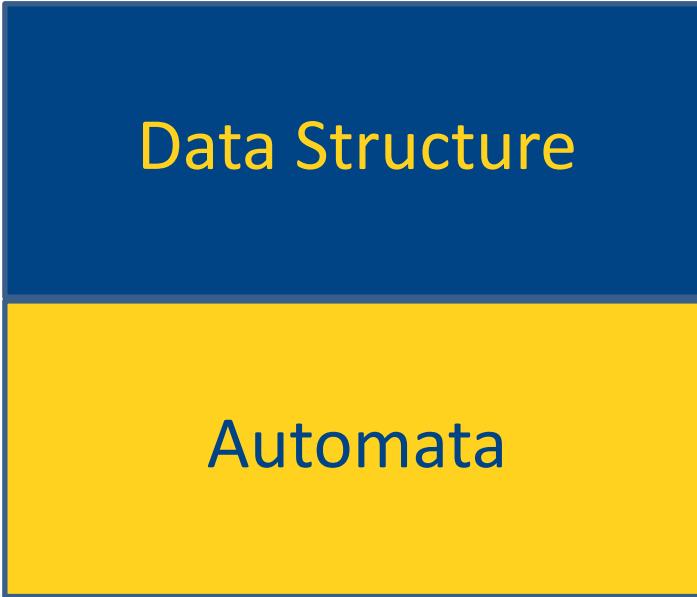


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- I'll try to record.





Data Structure

Automata

Maturity in programming and patience is also required.

# بارم‌بندی



- تمرین: ۴ نمره
- آزمونک: ۲ نمره
- پروژه: ۶ نمره
- میان‌ترم: ۲ نمره
- پایان‌ترم: ۶ نمره

# Homework

- About 7 assignments.
  - About 2 weeks each
  - About 3 problems.
- 
- Do not postpone until the deadline!
  - You have 10 days for late submission allowance.



LA<sup>T</sup>E<sub>X</sub>

# Exams (mid-term, final, and Quiz)



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- Reasonable exams, if you study you can get a good mark.



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# Exams



# Exams

- Open book and lectures in online exam.



# Exams

- Open book and lectures in online exam.
- On regular, you can bring on paper.



# Exams (mid-term, final, and Quiz)



# Exams (mid-term, final, and Quiz)

- Mid-term: ..../۲۲/۱۰..



# Exams (mid-term, final, and Quiz)

- Mid-term: ..../۲۲/۱۰..



# Mini Quiz

- 2 Quizzes
- With announcement.
- Takes 30 min.
- Not in group. Open slides.
- From previous or current lectures.
- Rules of mid-term and final applies.



# Project



# Project

- It is a complete compiler in 3-4 phases.
  - Lexical Analyzer (Scanner)
  - Parser
  - Code Generation



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# Project



**LINUX**

# Project

- The generated code **must** run on a MIPS machine.



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- The generated code **must** run on a MIPS machine.
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- It must **WORK** as the instructions.



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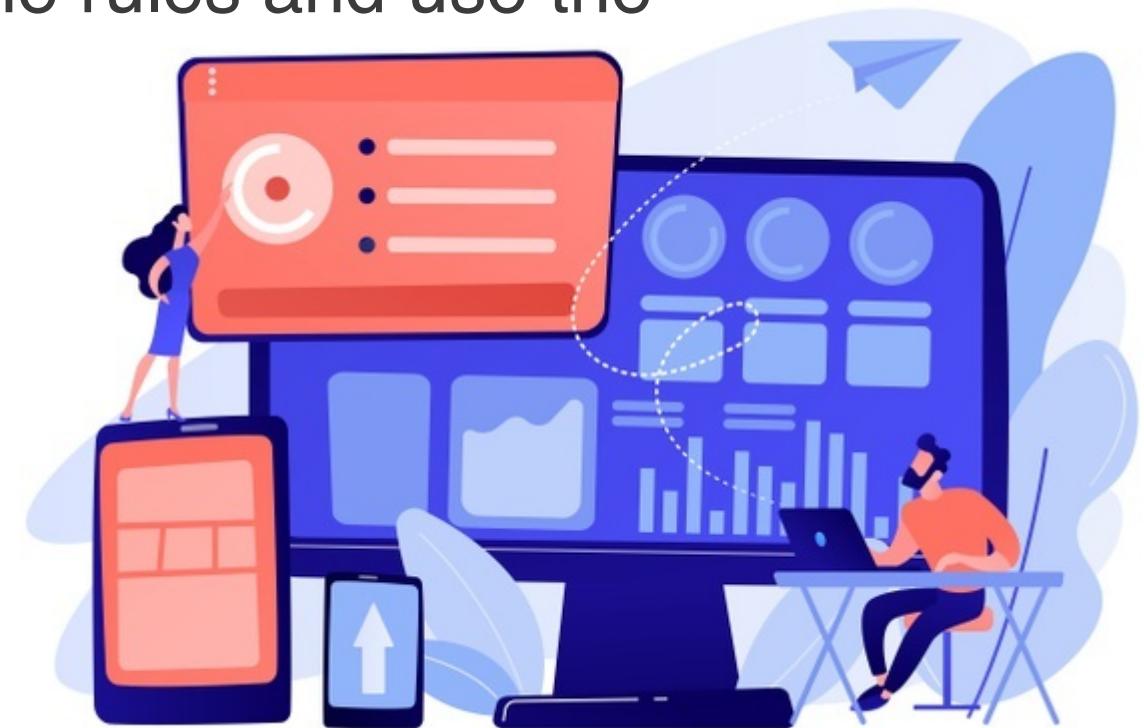
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- Project must follow the rules and use the framework.



# Warning



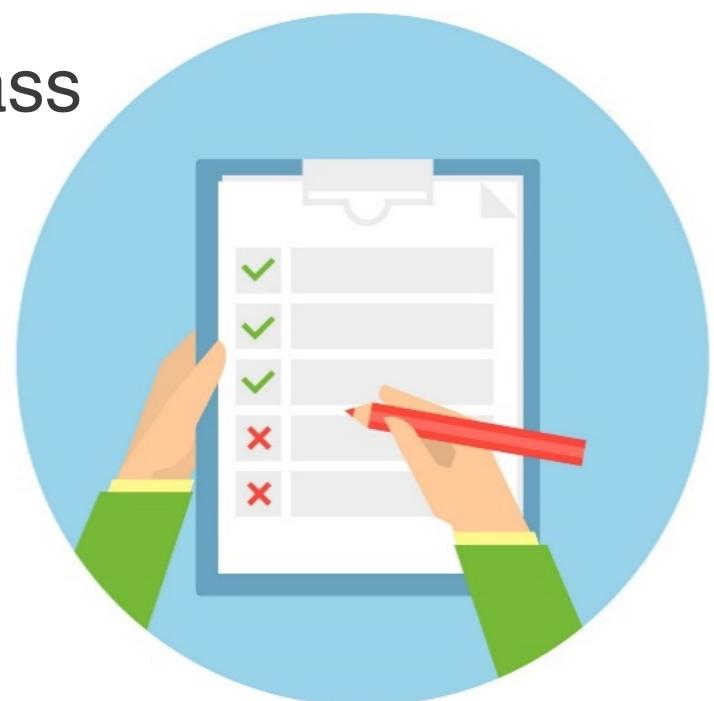
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- We have many tests either you pass or fail each one.



# Submission Guide line



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- Just submit once unless you contact with the corresponding assistant.



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- Just upload it to **Quera** or **CW**. E-mails lost usually.



# Project team

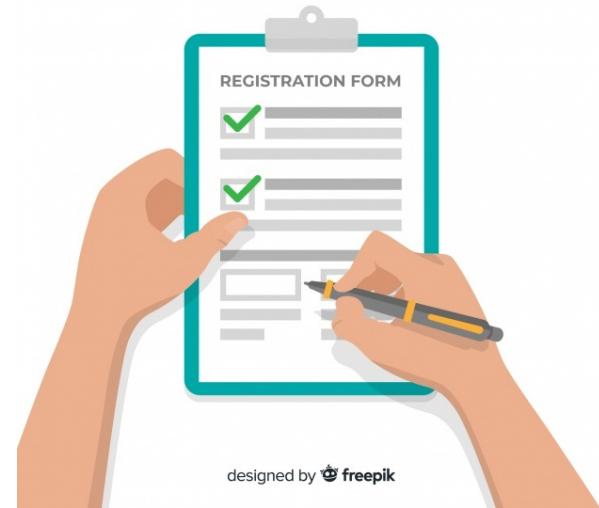


# Project team

- Team of two or one person.



# Group Rules



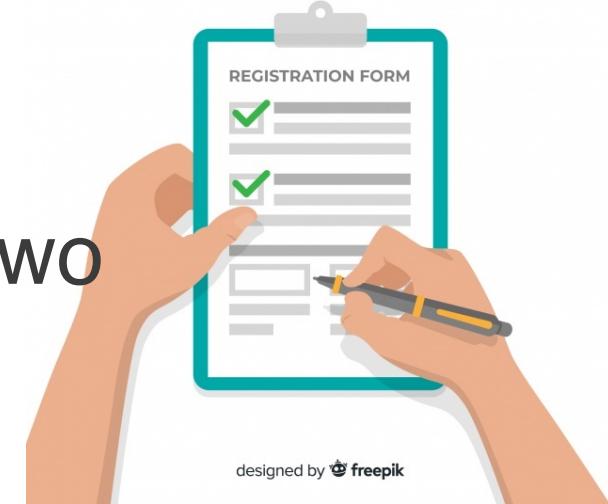
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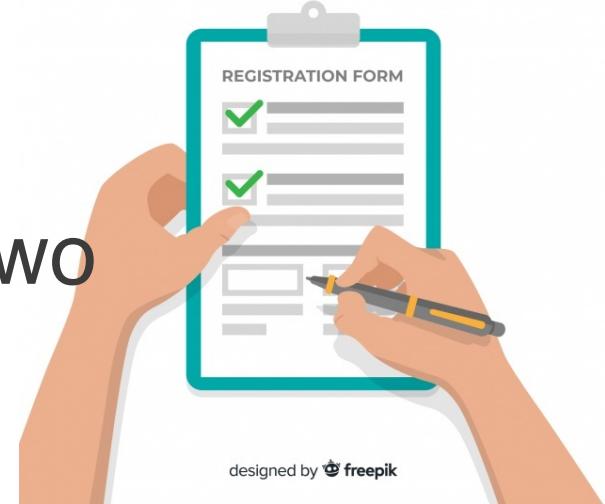
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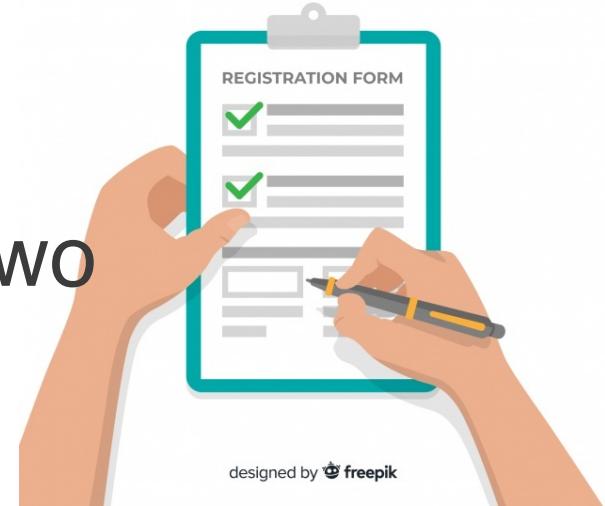
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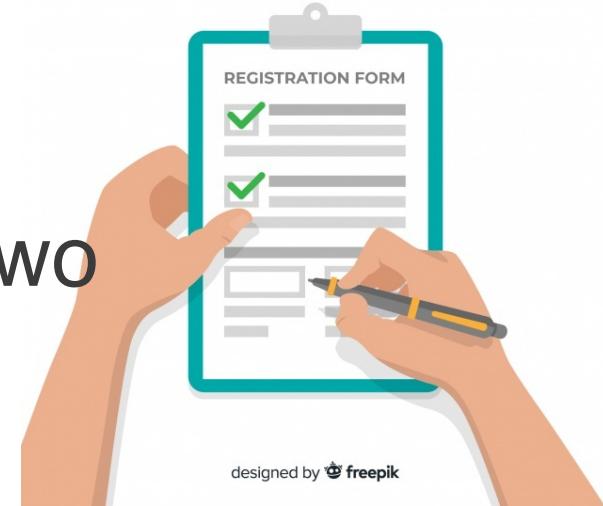
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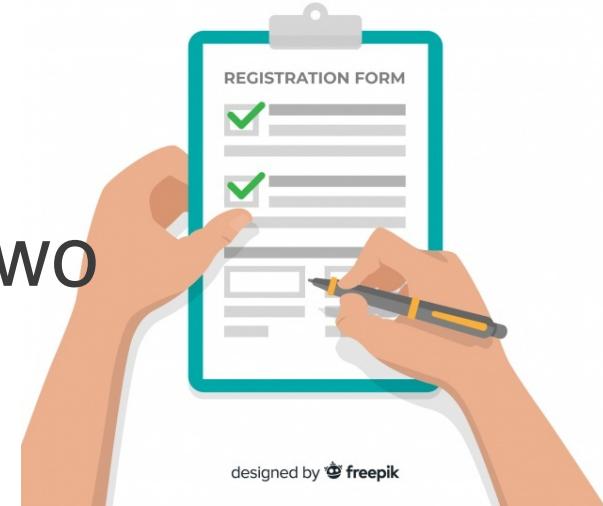
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# Group Rules

- Projects are for groups of one or two students.
- Pick an innovative name.
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- A piece of advice: create a group with friends with whom you can work easily.



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# Some Issues

# Some Issues

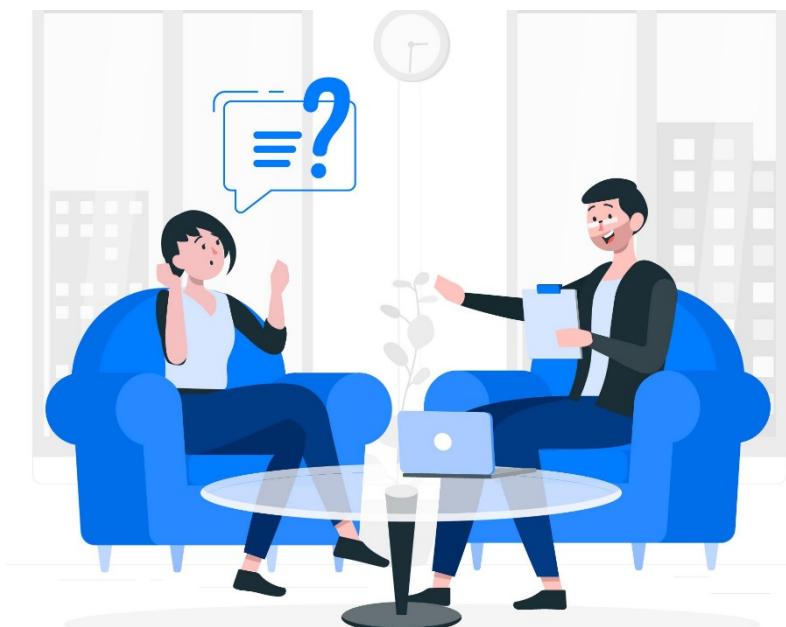
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- You must do a fair portion and have a good vision about everything else.
- There is a delivery and we check it there.



تقلب



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# Use your own knowledge

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- You can use your book and notes during online exams. But you must obtain them before the exam.



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- If we see any violation our reaction will be hard.



فیلم دکتر خسروی

# What happens if...

- In projects the penalty is a -100 grade.
- In homework, plagiarism have a -100 grade.
- In the exams/quiz (One of these can occurs):
  - Send it to education affairs (you may receive 0.25 as the final grade).
  - Fail the course.
  - 0/-100 for the whole exam.
  - -100 for the question
    - in exceptional cases
- And no extra grade, if we decide.





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## PHD Comics: Extension

