

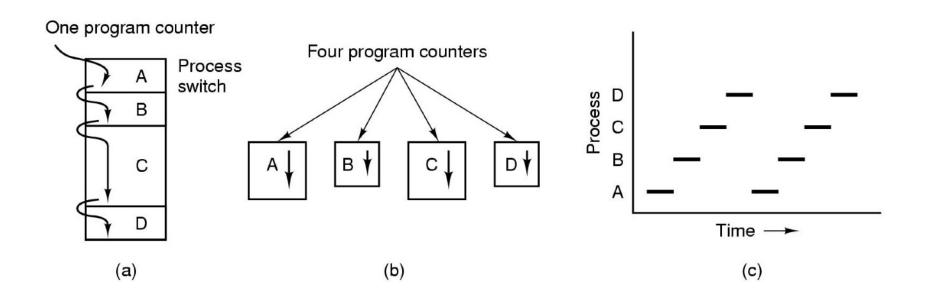
جلسه ۴: پردازه

The Process Concept

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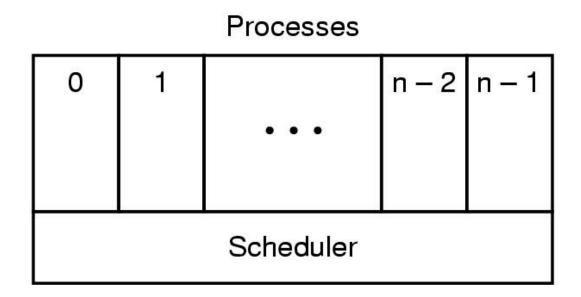
- Process a program in execution
 - * Program
 - description of how to perform an activity
 - instructions and static data values
 - * Process
 - a snapshot of a program in execution
 - memory (program instructions, static and dynamic data values)
 - CPU state (registers, PC, SP, etc)
 - operating system state (open files, accounting statistics etc)

Why use the process abstraction?



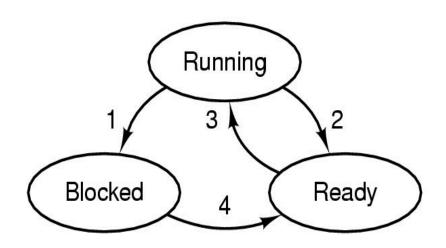
- Multiprogramming of four programs in the same address space
- Conceptual model of 4 independent, sequential processes
- Only one program active at any instant

The role of the scheduler



- Lowest layer of process-structured OS
 - * handles interrupts & scheduling of processes
- Sequential processes only exist above that layer

Process states



- 1. Process blocks for input
- 2. Scheduler picks another process
- 3. Scheduler picks this process
- 4. Input becomes available

Possible process states

- * running
- * blocked
- * ready

How do processes get created?

Principal events that cause process creation

- System initialization
- Initiation of a batch job
- User request to create a new process
- Execution of a process creation system call from another process

Process hierarchies

- Parent creates a child process,
 - * special system calls for communicating with and waiting for child processes
 - * each process is assigned a unique identifying number or process ID (PID)
- Child processes can create their own child processes
 - * Forms a hierarchy
 - UNIX calls this a "process group"

Process creation in UNIX

- All processes have a unique process id
 - getpid(), getppid() system calls allow processes to get their information
- Process creation
 - fork() system call creates a copy of a process and returns in both processes (parent and child), but with a different return value
 - exec() replaces an address space with a new program
- Process termination, signaling
 - * signal(), kill() system calls allow a process to be terminated or have specific signals sent to it

Example: process creation in UNIX

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

csh (pid = 24)

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 24)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 24)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
csh (pid = 22)
```

```
pid = fork()
if (pid == 0) {
  // child...
  exec();
else {
  // parent
  wait();
```

```
ls (pid = 24)
```

```
//ls program
main(){
  //look up dir
```

Process creation (fork)

- Fork creates a new process by copying the calling process
- The new process has its own
 - * memory address space (copied from parent)
 - Instructions
 - · Data
 - Stack
 - * Register set (copied from parent)
 - * Process table entry in the OS

Killing a process

- Sending kill signal to kernel
- Killing a process does not kill its descendants

wait()

- Waits until:
 - A child is killed, or
 - A signal is received from OS