Project 2 - solution code

- Project 2 solution code in files for project 3:
 - Mutex solution in Synch.c
 - But this code has several flaws!
 - If you copied this, we will know!
 - Producer/Consumer and Dining Philosophers solution in Proj2Sol.pdf
- Make sure you study and understand the solutions!
 - * Try to identify errors in your own project 2 code
 - Email your thoughts on these to me and the TA over the next two days for extra credit
 - Important: if the solutions you handed in were copied let me know by email a.s.a.p.

Project 3

Part 1: The Sleeping Barber problem

- Use semaphores and mutex variables for thread synchronization
- You decide how to test your code!!
 - · We're not providing the sample output for this one!

Part 2: The Gaming Parlor problem

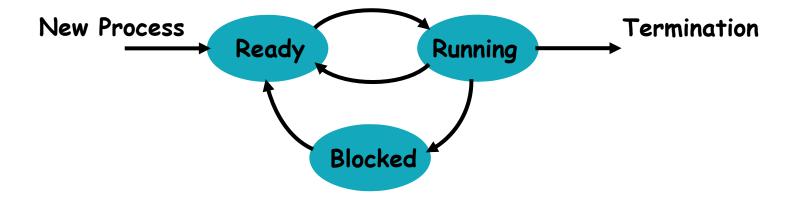
- Use monitors and condition variables for synchronization
- Avoid deadlock when handling resource requests
- Sample output given in handout

CS 333 Introduction to Operating Systems

Class 8 - Scheduling

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Process state model



CPU scheduling criteria

- CPU Utilization how busy is the CPU?
- Throughput how many jobs finished/unit time?
- Turnaround Time how long from job submission to job termination?
- Response Time how long (on average) does it take to get a "response" from a "stimulus"?
- Missed deadlines were any deadlines missed?

Scheduler options

Priorities

- * May use priorities to determine who runs next
- * Dynamic vs. Static algorithms
 - Dynamically alter the priority of the tasks while they are in the system (possibly with feedback)
 - · Static algorithms typically assign a fixed priority when the job is initially started.

Preemptive vs. Nonpreemptive

Preemptive systems allow the task to be interrupted at any time so that the O.S. can take over again.

Scheduling policies

- First-Come, First Served (FIFO)
- Shortest Job First (non-preemptive)
- Shortest Job First (with preemption)
- Round-Robin Scheduling
- Priority Scheduling
- Real-Time Scheduling

- Start jobs in the order they arrive (FIFO queue)
- Run each job until completion

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	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2

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		Arrival	Processing		Turnaround
Proc	ess	Time	Time	Delay	Time
	1	0	3		
	2	2	6		
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	4	6	5		
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3	4	4		
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- Start jobs in the order they arrive (FIFO queue)
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	2	2		6					
	3	4		4					
	4	6		5					
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2	2	6		
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Process	Time	Time	Delay	Time
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3	4	4		
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Process	Time	Time	Delay	Time
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Process	Time	Time	Delay	Time
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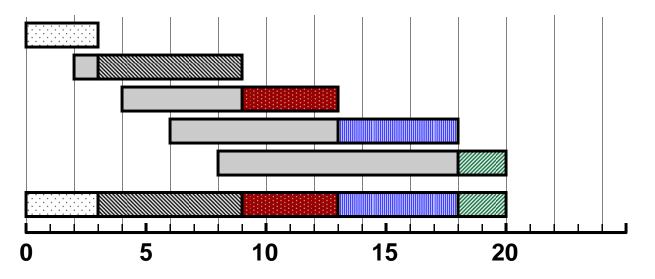
	Arrival	Processing		Turnaround
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- Run each job until completion

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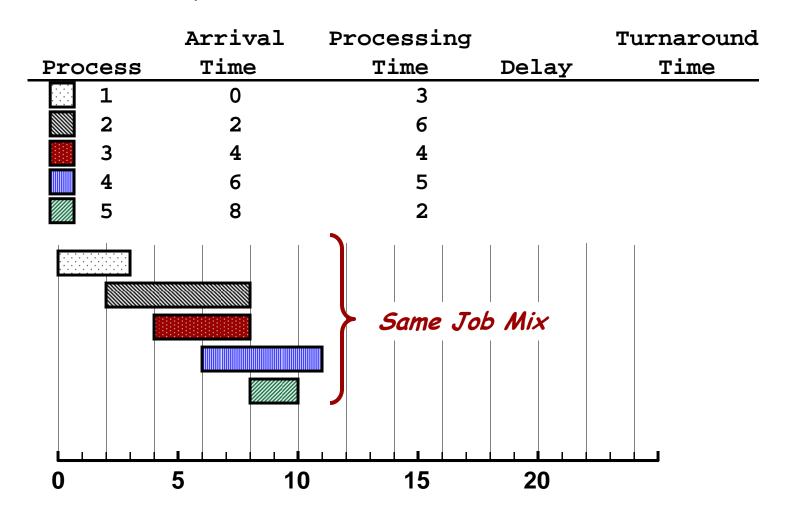
- Start jobs in the order they arrive (FIFO queue)
- Run each job until completion

		Arrival	Processing		Turnaround
Prod	cess	Time	Time	Delay	Time
	1	0	3	0	3
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- Select the job with the shortest (expected) running time
- Non-Preemptive

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Process	Arrival Time	Processing Time	Delay	Turnaround Time
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3	4	4		
4	6	5		
5	8	2		
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- Select the job with the shortest (expected) running time
- Non-Preemptive

D		Processing	Dolo	Turnaround
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Drogogg	Arrival Time	Processing Time	Delay	Turnaround Time
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- Select the job with the shortest (expected) running time
- Non-Preemptive

		Processing		Turnaround
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- Select the job with the shortest (expected) running time
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		Processing		Turnaround
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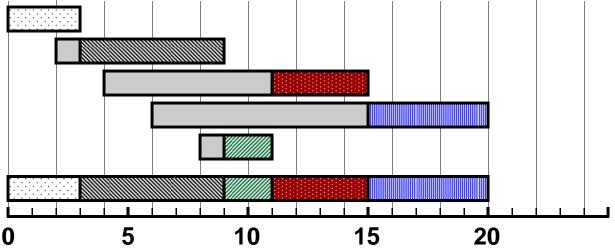
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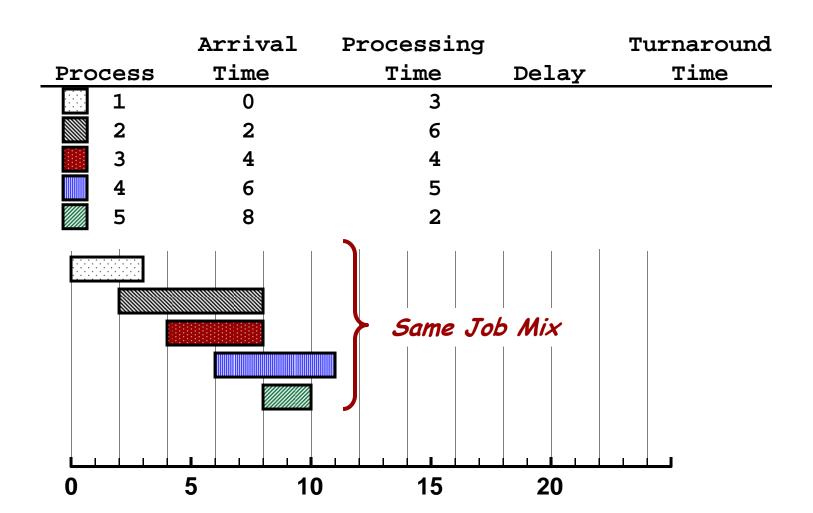
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- Select the job with the shortest (expected) running time
- Non-Preemptive

	Arrival	Processing		Turnaround
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Process	Arrival Time	Processing Time	Delay	Turnaround Time
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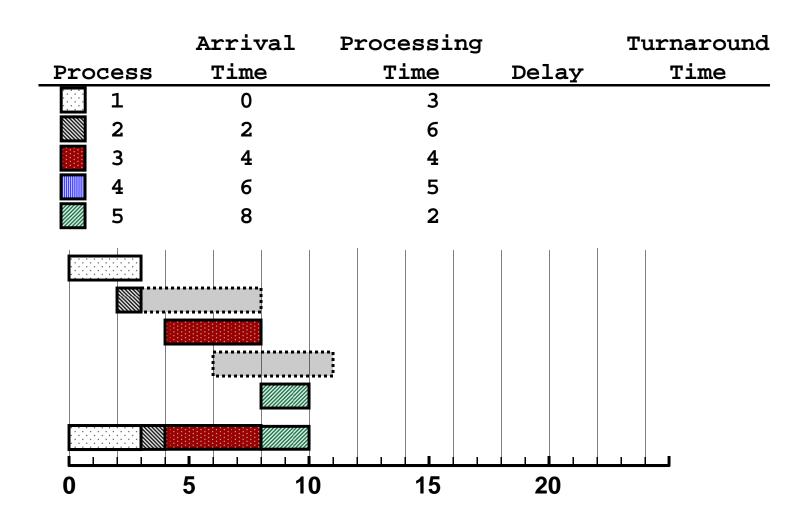
_	Arrival	_		Turnaround
Process	Time	Time	Delay	Time
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3	4	4		
4	6	5		
5	8	2		
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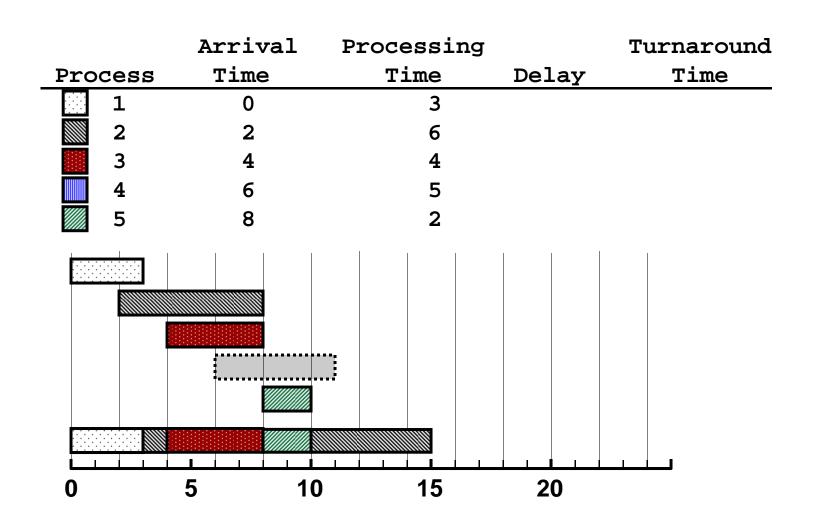
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3	4	4		
4	6	5		
5	8	2		
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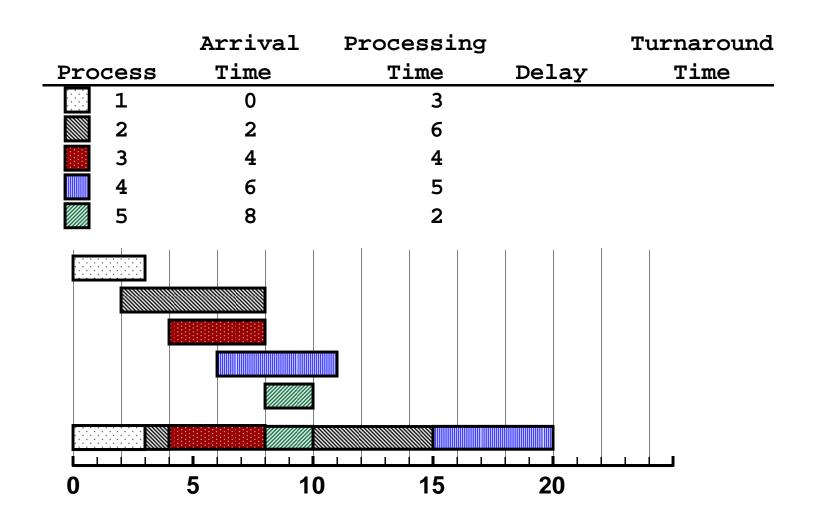
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0	5	10 15	20	

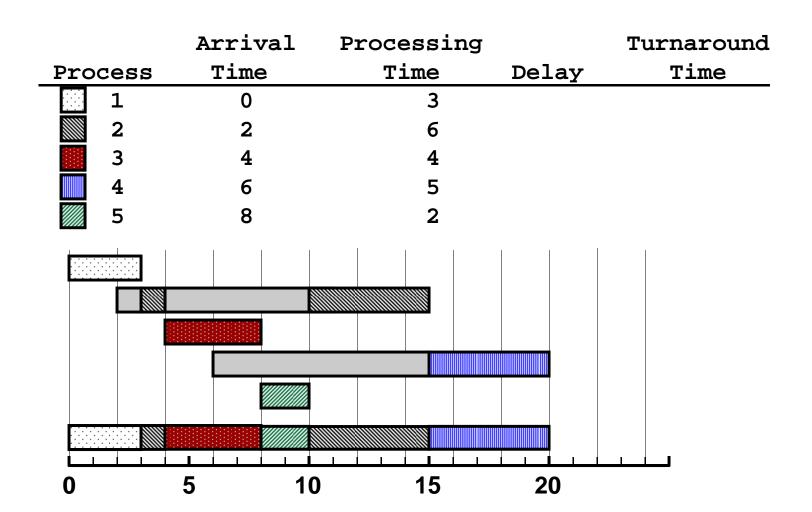
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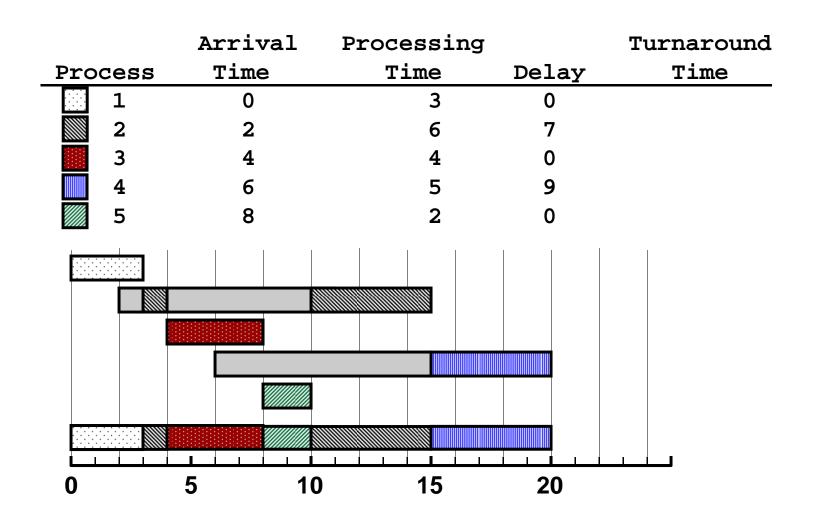
	Arrival	Processing		Turnaround
Process	Time	Time	Delay	Time
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2	2	6		
3	4	4		
4	6	5		
5	8	2		
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	Arrival	Processin	g	Turnaround
Process	Time	Time	Delay	Time
3 1	0	3	0	3
2	2	6	7	13
3	4	4	0	4
4	6	5	9	14
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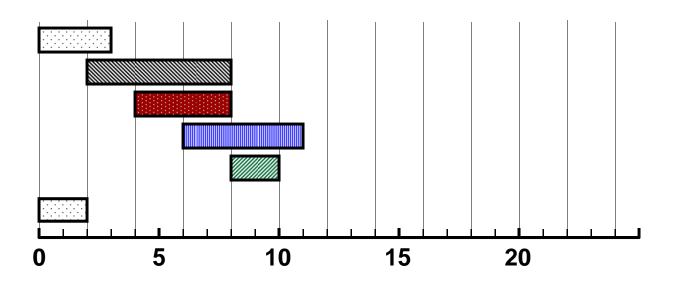
Goal: Enable interactivity

Limit the amount of CPU that a process can have at one time.

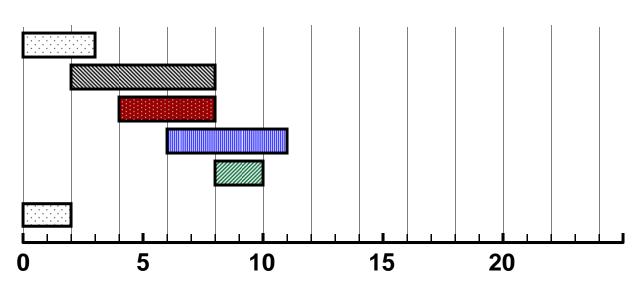
Time quantum

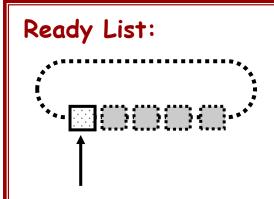
- Amount of time the OS gives a process before intervention
- * The "time slice"
- * Typically: 1 to 100ms

•	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2

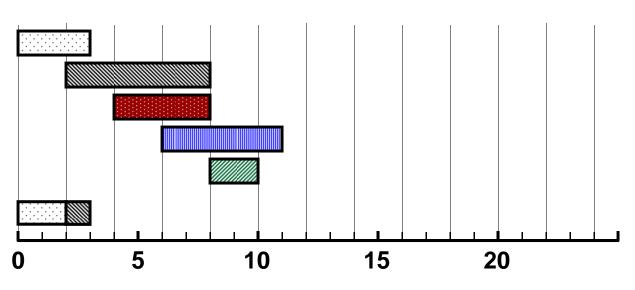


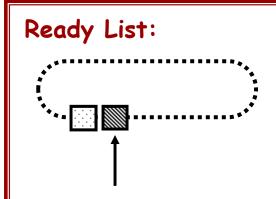
•		Arrival	Processing
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	1	0	3
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	3	4	4
	4	6	5
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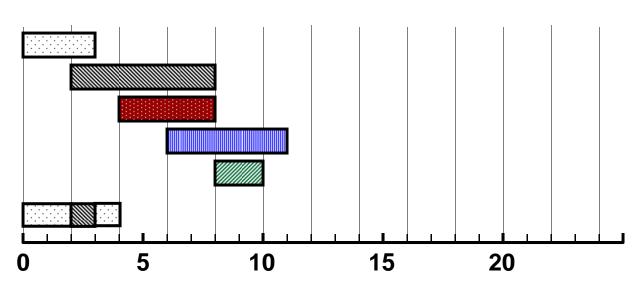


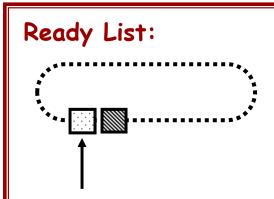
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Process	Time	Time
1	0	3
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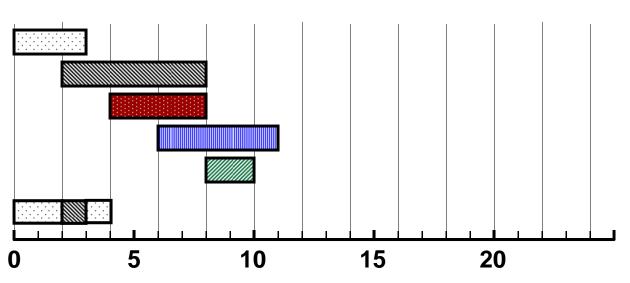


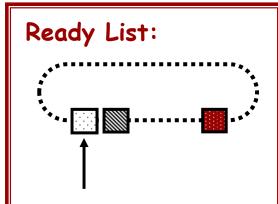
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	1	0	3
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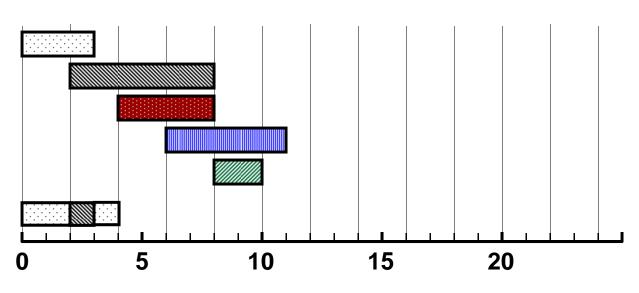


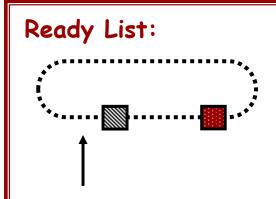
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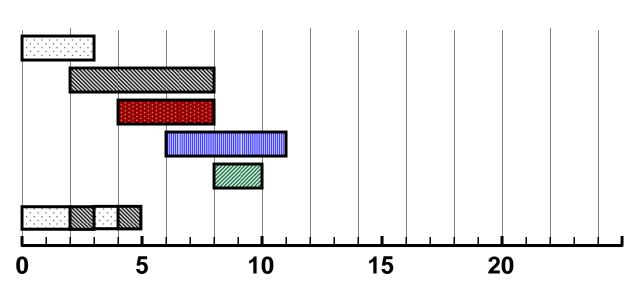


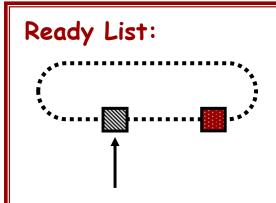
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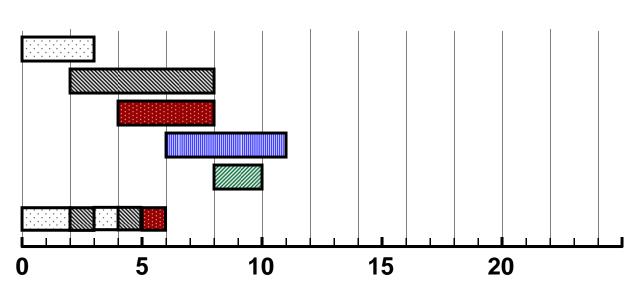


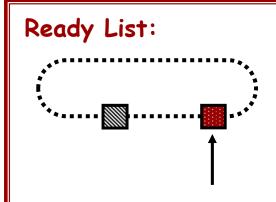
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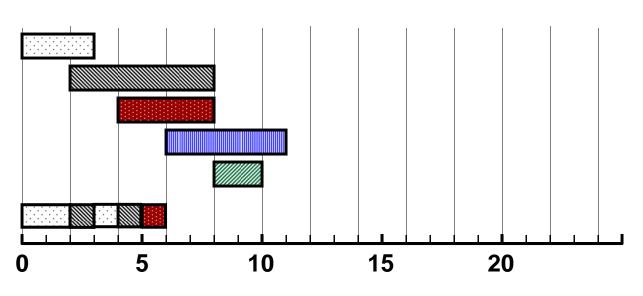


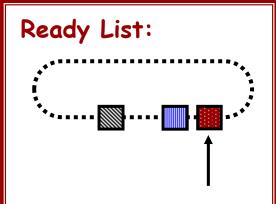
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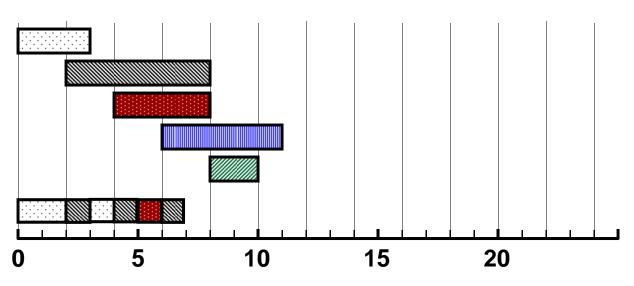


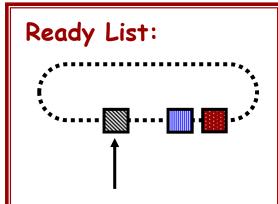
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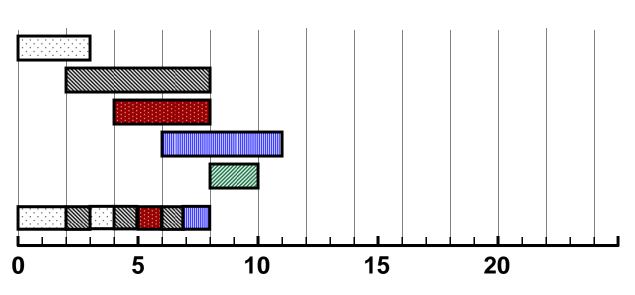


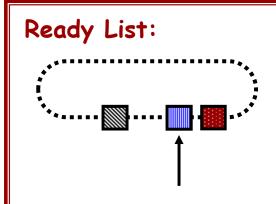
•	Arrival	Processing
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<u> </u>	0	3
2	2	6
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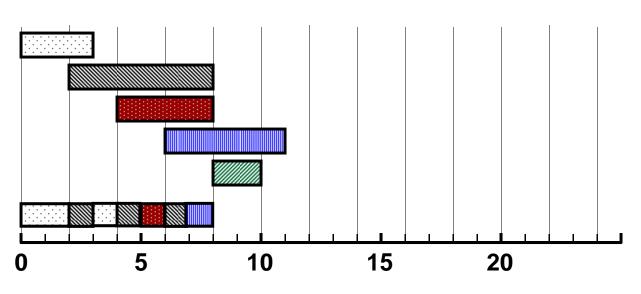


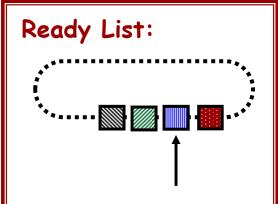
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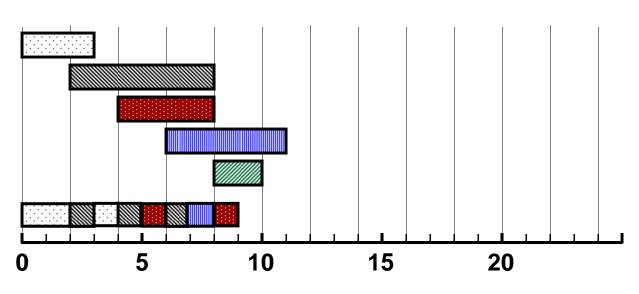


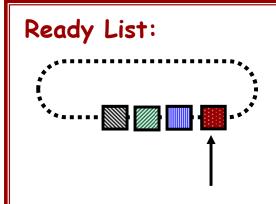
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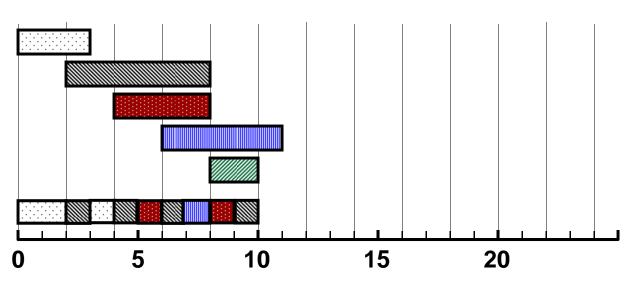


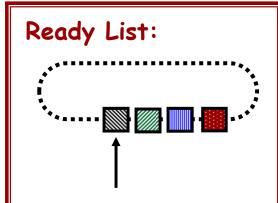
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Process	Time	Time
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2	2	6
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5	8	2



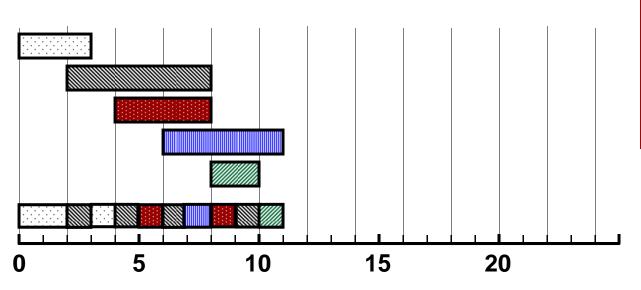


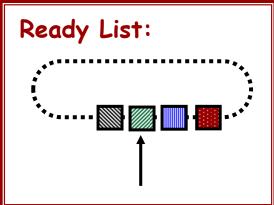
•		Arrival	Processing
_	Process	Time	Time
_	1	0	3
	2	2	6
	3	4	4
	4	6	5
	5	8	2



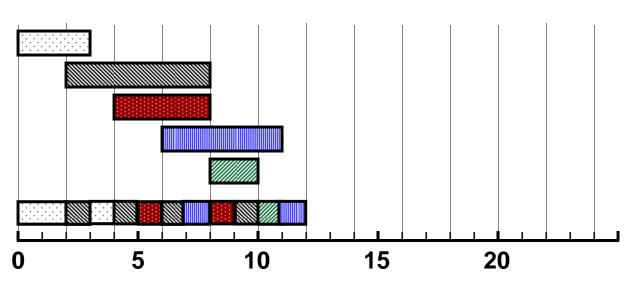


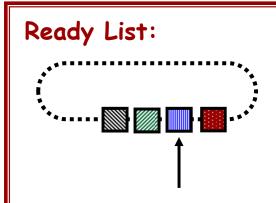
•		Arrival	Processing
Pro	ocess	Time	Time
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	2	2	6
	3	4	4
	4	6	5
	5	8	2



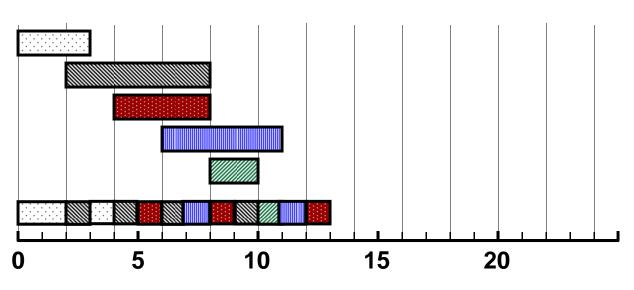


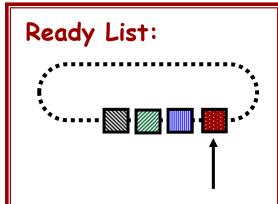
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_	Process	Time	Time
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	3	4	4
	4	6	5
	5	8	2



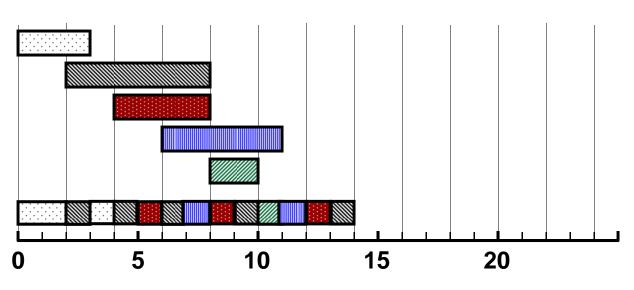


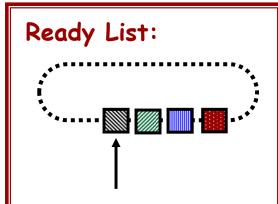
•	Arrival	Processing
Process	Time	Time
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2	2	6
3	4	4
4	6	5
5	8	2



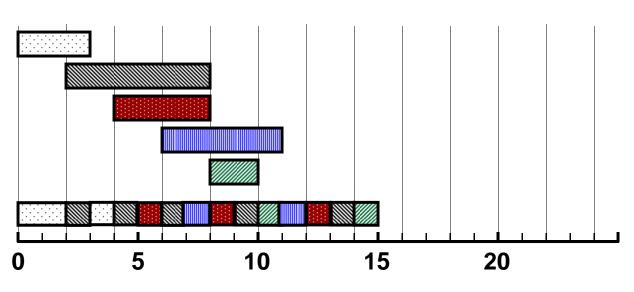


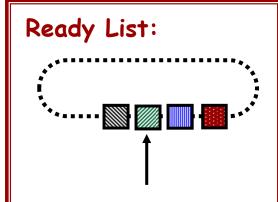
•	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2



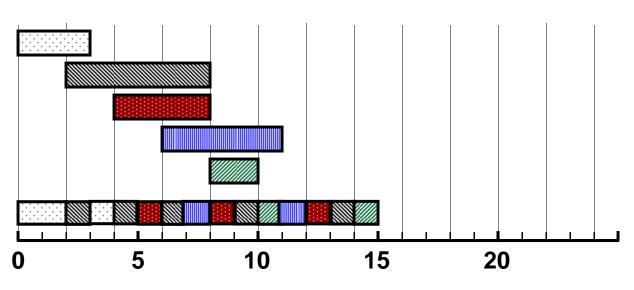


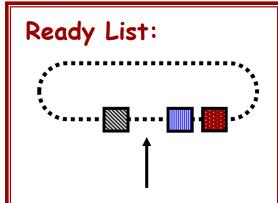
•	Arrival	Processing
Process	Time	Time
<u> </u>	0	3
2	2	6
3	4	4
4	6	5
5	8	2



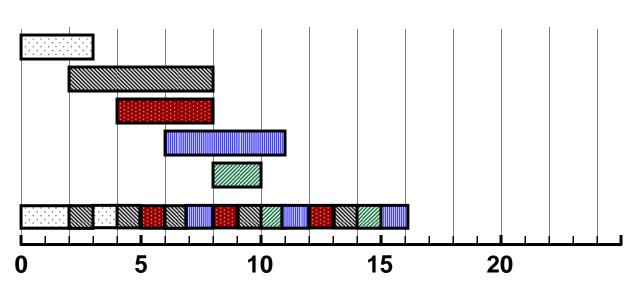


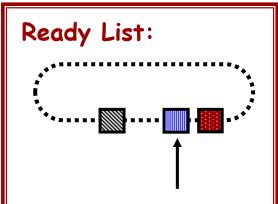
•		Arrival	Processing
_	Process	Time	Time
_	1	0	3
	2	2	6
	3	4	4
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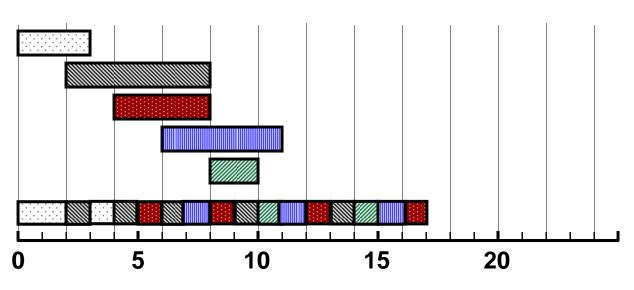


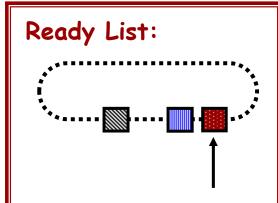
•	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2



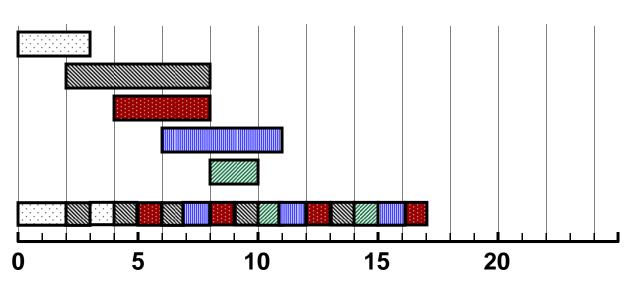


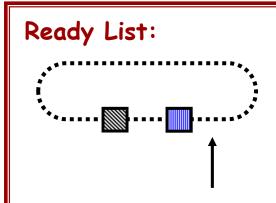
•	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2



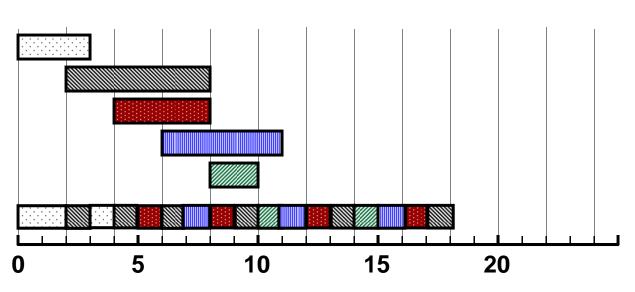


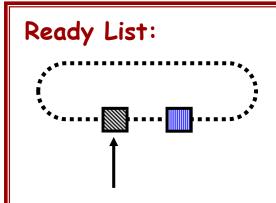
•	Arrival	Processing
Process	Time	Time
1	0	3
2	2	6
3	4	4
4	6	5
5	8	2



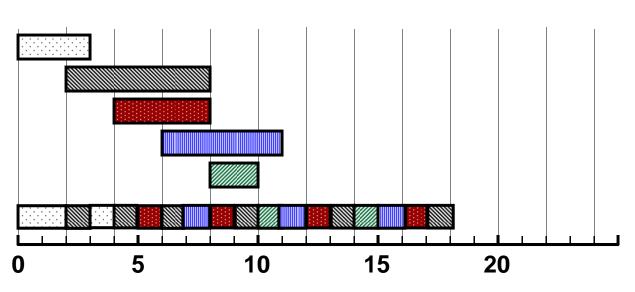


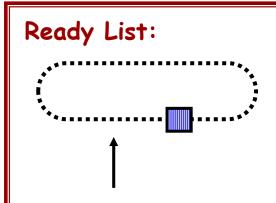
•		Arrival	Processing
_	Process	Time	Time
_	1	0	3
	2	2	6
	3	4	4
	4	6	5
	5	8	2



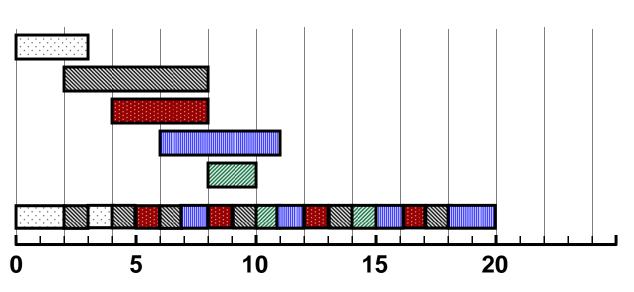


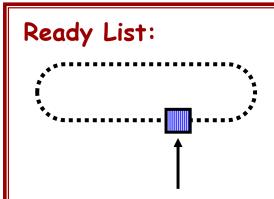
•		Arrival	Processing
_	Process	Time	Time
_	1	0	3
	2	2	6
	3	4	4
	4	6	5
	5	8	2



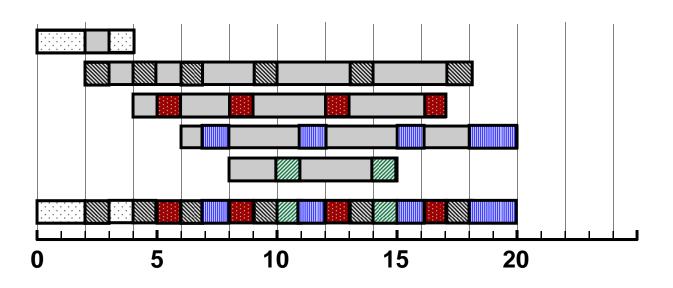


•	Arrival	Processing
Process	Time	Time
<u> </u>	0	3
2	2	6
3	4	4
4	6	5
5	8	2

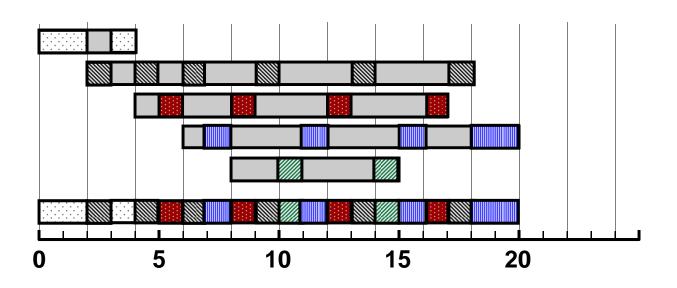




•		Arrival	Processing	
_	Process	Time	Time	
	1	0	3	
	2	2	6	
	3	4	4	
	4	6	5	
	5	8	2	



•		Arrival	Processing		Turnaround
_	Process	Time	Time	Delay	Time
_	1	0	3	1	4
	2	2	6	10	16
	3	4	4	9	13
	4	6	5	9	14
	5	8	2	5	7



- Effectiveness of round-robin depends on
 - The number of jobs, and
 - * The size of the time quantum.
- Large # of jobs means that the time between scheduling of a single job increases
 - Slow responses
- Larger time quantum means that the time between the scheduling of a single job also increases
 - Slow responses
- Smaller time quantum means higher processing rates but also more overhead!

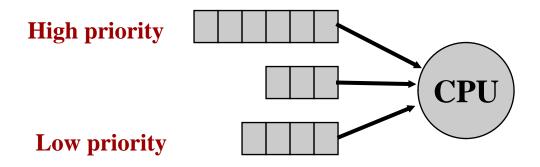
Scheduling in general purpose systems

Priority scheduling

- Assign a priority (number) to each process
- Schedule processes based on their priority
- Higher priority processes get more CPU time
- Managing priorities
 - * Can use "nice" to reduce your priority
 - Can periodically adjust a process' priority
 - · Prevents starvation of a lower priority process
 - Can improve performance of I/O-bound processes by basing priority on fraction of last quantum used

Multi-Level Queue Scheduling

- Multiple queues, each with its own priority.
 - * Equivalently: each priority level has its own ready queue
- Within each queue...Round-robin scheduling.
- Simplist Approach:
 - * A Process's priority is fixed & unchanging

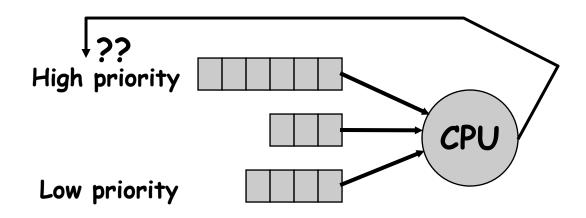


- Problem: Fixed priorities are too restrictive
 - Processes exhibit varying ratios of CPU to I/O times.
- Dynamic Priorities
 - Priorities are altered over time, as process behavior changes!

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- Problem: Fixed priorities are too restrictive
 - Processes exhibit varying ratios of CPU to I/O times.
- Dynamic Priorities
 - Priorities are altered over time, as process behavior changes!
- Issue: When do you change the priority of a process and how often?
- Solution: Let the amount of CPU used be an indication of how a process is to be handled
 - Expired time quantum → more processing needed
 - Unexpired time quantum → less processing needed
- Adjusting quantum and frequency vs. adjusting priority?

- n priority levels, round-robin scheduling within a level
- Quanta increase as priority decreases
- Jobs are demoted to lower priorities if they do not complete within the current quantum



- 🗅 Details, details, details...
 - Starting priority?
 - Frequency of moving between priorities?
 - * How long should the time quantum be?

Lottery Scheduling

- Scheduler gives each thread some lottery tickets
- To select the next process to run...
 - The scheduler randomly selects a lottery number
 - The winning process gets to run
- Thread A gets 50 tickets
 Thread B gets 15 tickets
 Thread C gets 35 tickets
 There are 100 tickets outstanding.

Lottery Scheduling

- Scheduler gives each thread some lottery tickets.
- To select the next process to run...
 - The scheduler randomly selects a lottery number
 - The winning process gets to run
- □ Example Thread A gets 50 tickets → 50% of CPU Thread B gets 15 tickets → 15% of CPU Thread C gets 35 tickets → 35% of CPU There are 100 tickets outstanding.

Lottery Scheduling

- Scheduler gives each thread some lottery tickets.
- To select the next process to run...
- The scheduler randomly selects a lottery number
- The winning process gets to run
- □ <u>Example</u> Thread A gets 50 tickets → 50% of CPU

 Thread B gets 15 tickets → 15% of CPU

 Thread C gets 35 tickets → 35% of CPU

There are 100 tickets outstanding.

- Flexible
- · Fair
- Responsive

- Assume processes are relatively periodic
 - * Fixed amount of work per period (e.g. sensor systems or multimedia data)

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- Two "main" types of schedulers...
- Rate-Monotonic Schedulers

Earliest-Deadline-First Schedulers

- Assume processes are relatively periodic
 - Fixed amount of work per period (e.g. sensor systems or multimedia data)
- Two "main" types of schedulers...
- Rate-Monotonic Schedulers
 - Assign a fixed, unchanging priority to each process
 - No dynamic adjustment of priorities
 - Less aggressive allocation of processor
- Earliest-Deadline-First Schedulers
 - Assign dynamic priorities based upon deadlines

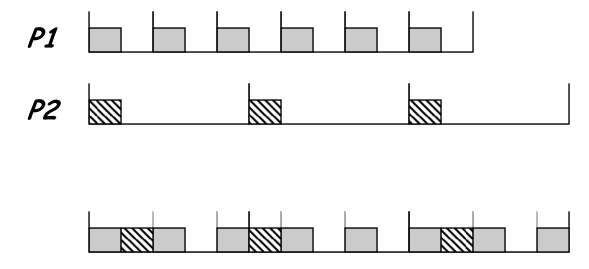
- Typically real-time systems involve several steps (that aren't in traditional systems)
- Admission control
 - All processes must ask for resources ahead of time.
 - * If sufficient resources exist, the job is "admitted" into the system.
- Resource allocation
 - Upon admission...
 - the appropriate resources need to be reserved for the task.
- Resource enforcement
 - Carry out the resource allocations properly

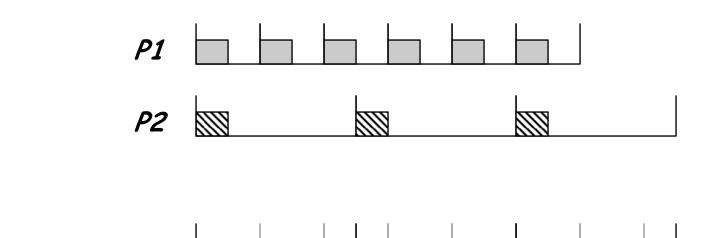
- For preemptable, periodic processes (tasks)
- Assigns a fixed priority to each task
 - * T = The period of the task
 - C = The amount of processing per task period

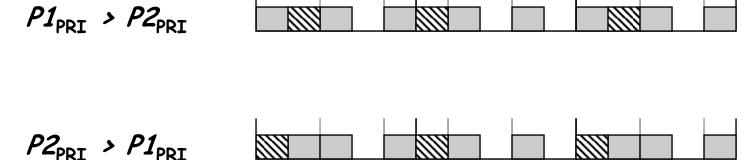
Process P1

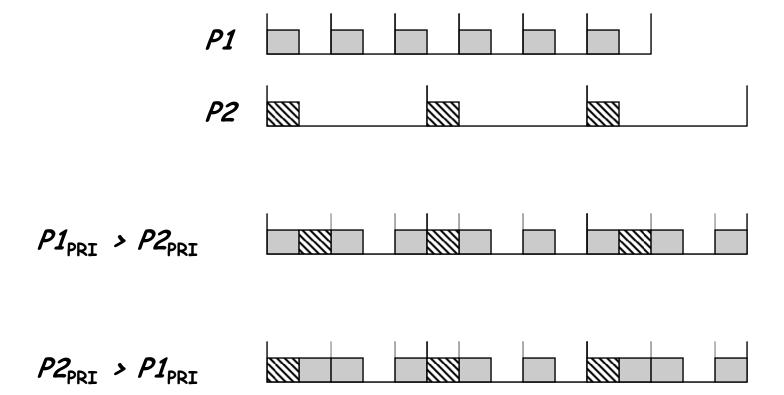


- In RMS scheduling, the question to answer is...
 - What priority should be assigned to a given task?









Which is best?

- Assign shortest period tasks to the highest priorities
- Admission control is difficult...
 - < 69% There is a standard formula
 - > 69% It may be possible, but need to do more complex analysis

¬ Assumptions:

- * Processes complete (yield) within their period
- * Independent processes
- Same CPU requirements per burst
- Other non-periodic processes have no deadlines
- Instantaneous preemption with no overhead

Earliest Deadline First

- When processes do not have periodic execution or constant CPU requirements...
- When processes have deadline specifications...
- Unlike RMS, EDF uses dynamic priorities (based upon earliest deadline first)
 - (+) 100% processor utilization ...?
 - (-) Need to keep track of deadlines
- Admission Control
 - Just check to see if 100% processor utilization.
 - * Sum the C_i/T_i 's and see if less than or equal to 1
 - * What about overhead?

Quiz

- What are the main tasks of the scheduler?
- What is the difference between preemptive and nonpreemptive scheduling?
- What is the advantage of a shorter scheduling quantum?
- What is the advantage of a longer scheduling quantum?
- Why is feedback scheduling useful for interactive jobs?
- Are these scheduling policies subject to starvation?
 - Shortest Job First scheduling?
 - Round Robin scheduling?
 - First Come First Served scheduling?
 - Priority scheduling?
 - Earliest Deadline First scheduling?