

بسم الله الرحمن الرحيم

جلسه هفدهم

خلاصه سازی برای مدداده

احساس فردگی

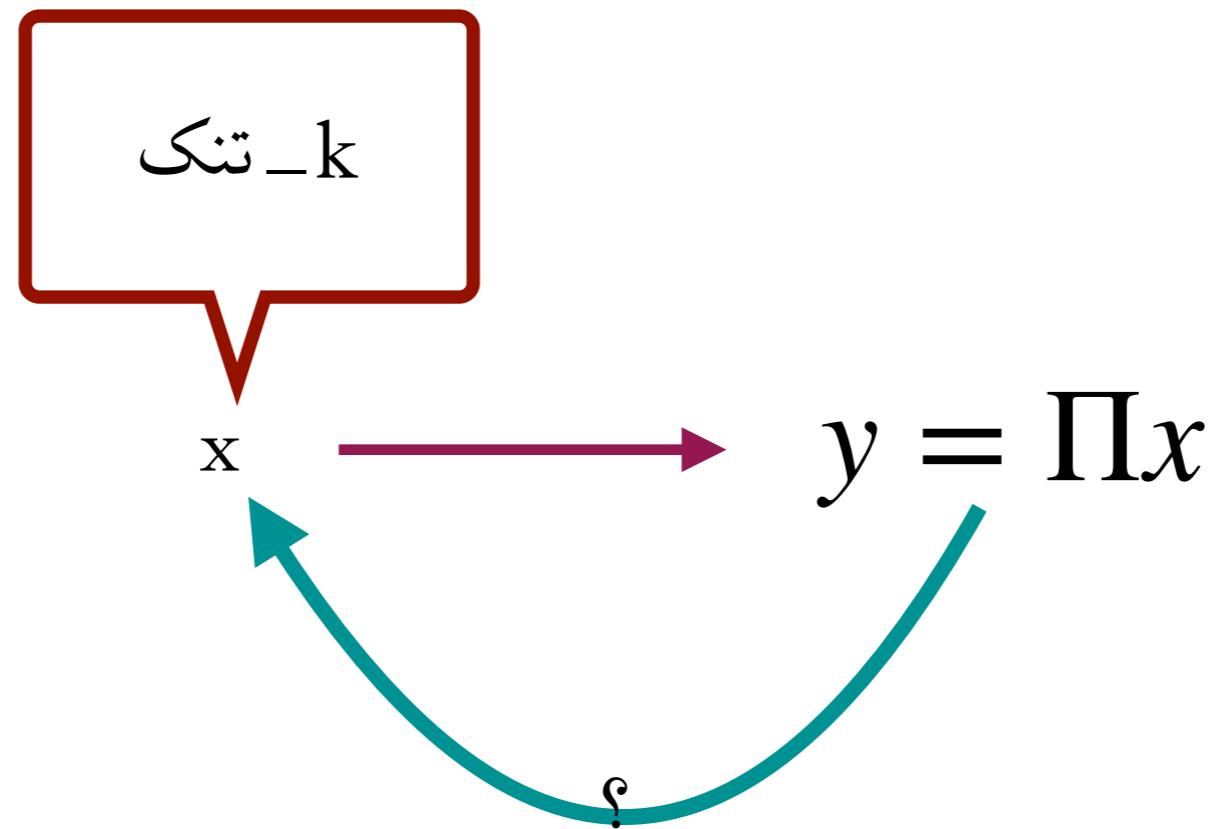


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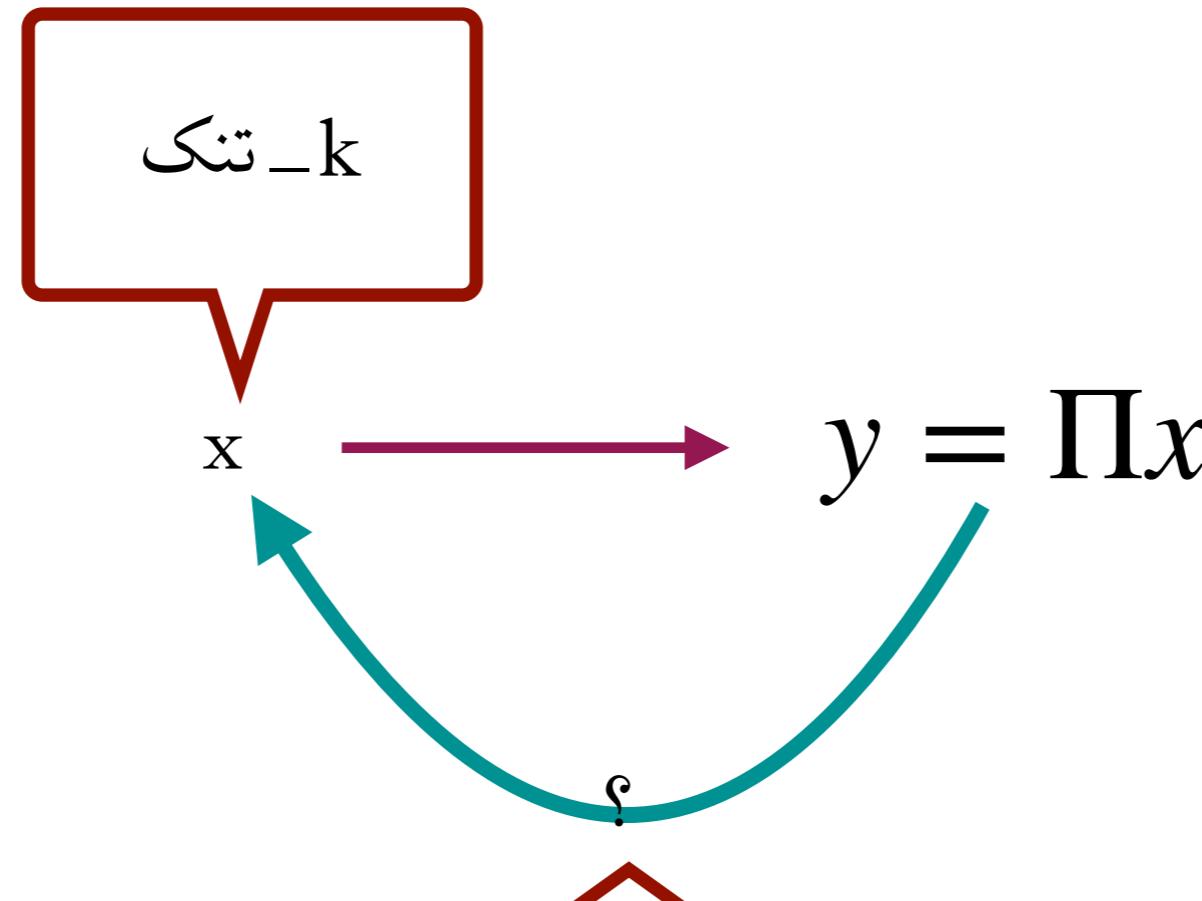
الگوریتم کمینه نرم ۱



انگیزش



انگیزش

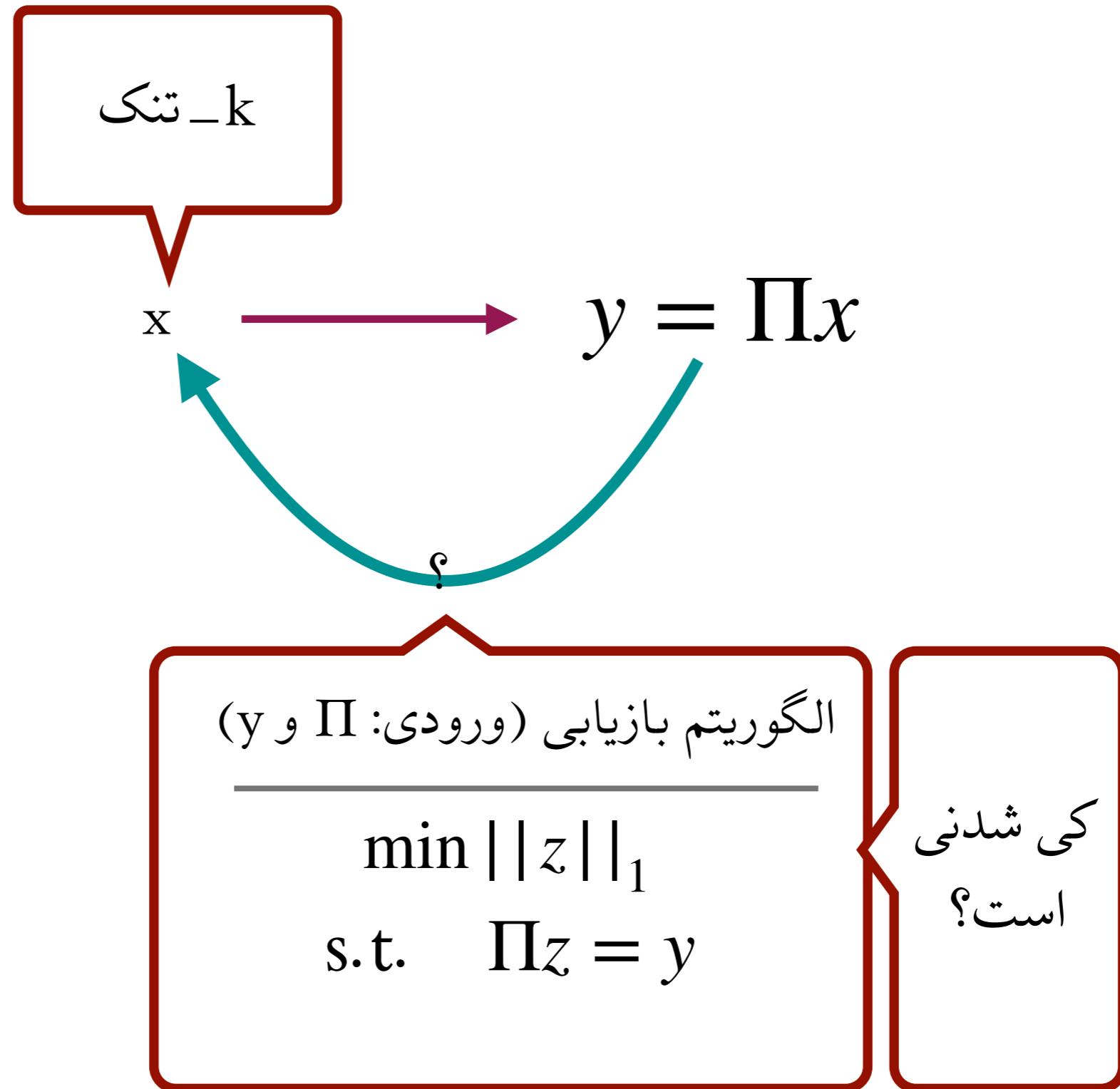


الگوریتم بازیابی (ورودی: y و Π)

$$\min ||z||_1$$

$$\text{s.t. } \Pi z = y$$

انگیزش

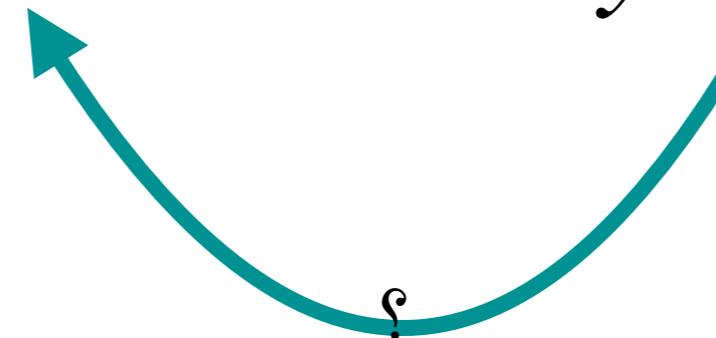


انگیزش

تقریبا k -تنک؟

k -تنک

$$x \longrightarrow y = \Pi x$$



الگوریتم بازیابی (ورودی: y و Π)

$$\begin{aligned} & \min ||z||_1 \\ \text{s.t. } & \Pi z = y \end{aligned}$$

کی شدنی
است؟

Definition 28. We say a matrix $\Pi \in \mathbb{R}^{m \times n}$ satisfies the (ε, k) -restricted isometry property (or RIP for short) if for all k -sparse vectors x of unit Euclidean norm,

$$1 - \delta \leq \|\Pi x\|_2^2 \leq 1 + \delta.$$

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$$\sup_{T \subset [n] | T|=k} \|I_k - (\Pi^{(T)})^* \Pi^{(T)}\| < \delta,$$

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مثال: RIP با احتمال مثبت: برای

$$\sqrt{\frac{d}{m}} S H$$

ماتریس نمونه‌گیری:
 m سطر، هر سطر یک $1 \pm$

$\times \frac{1}{\sqrt{d}}$ ماتریس هادامار

Definition 1 An $m \times n$ matrix A satisfies a null-space property of order k with constant C if for any $\eta \in \mathbb{R}^n$ such that $A\eta = 0$, and any set $T \subset \{1 \dots n\}$ of size k , we have

$$\|\eta\|_1 \leq C\|\eta_{-T}\|_1$$

where $-T$ denotes a complement of T .

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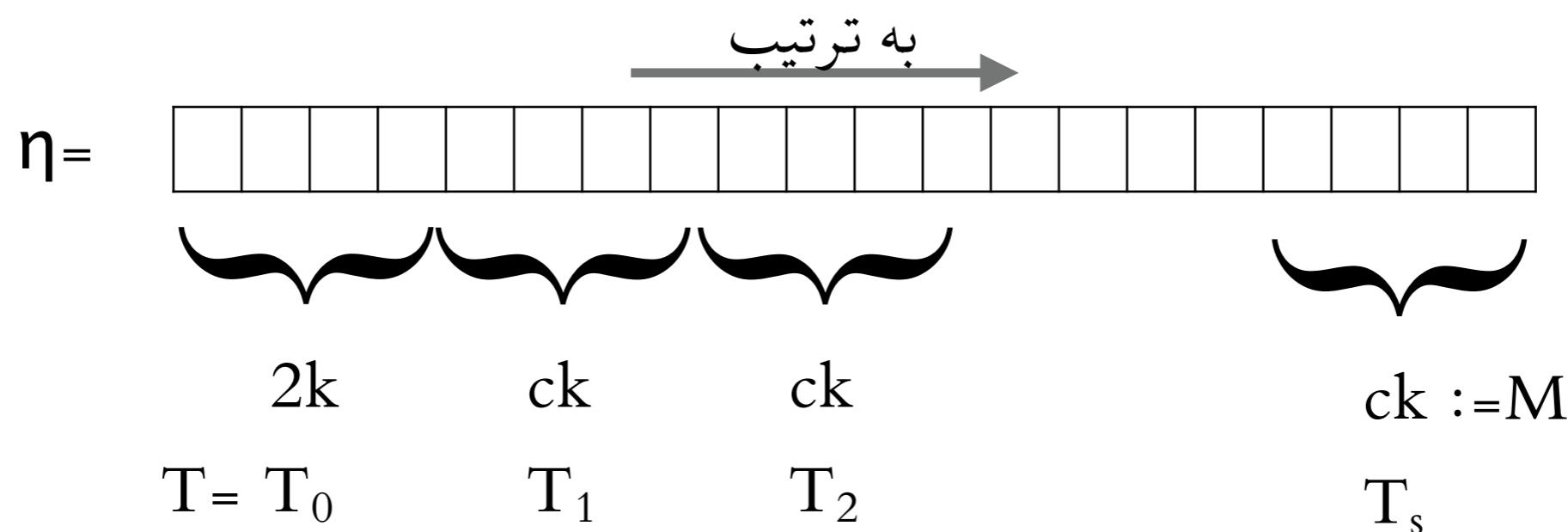
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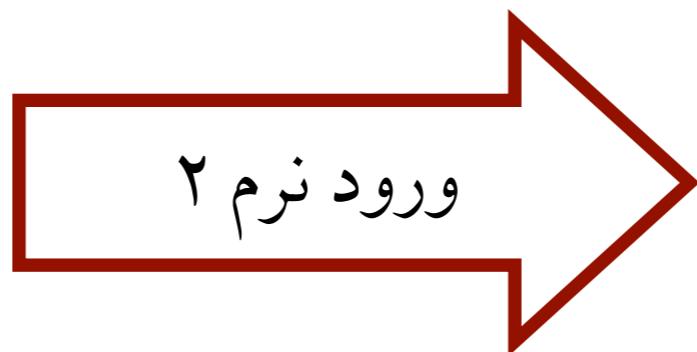
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$$\| \eta_T \|_1 \leq (C - 1) \| \eta_{-T} \|_1$$



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$$\|x\|_1 = \langle x, 1 \rangle \leq \|x\|_2 \sqrt{2k}$$

$$\|\eta_T\|_1 \leq (2k)^{1/2} \|\eta_T\|_2$$

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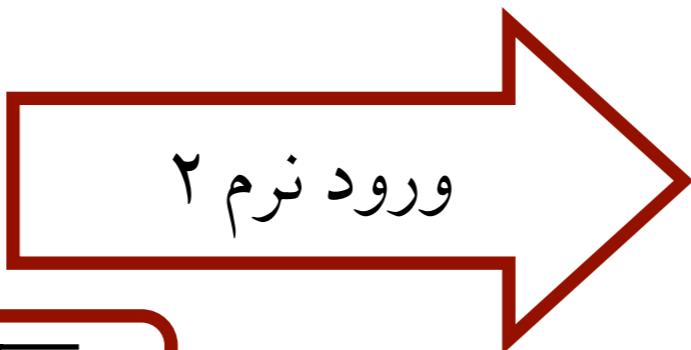
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$$\leq (1-\delta)^{-1}(1+\delta) \sum_{j=2}^s \|\eta_{T_j}\|_2$$

$$1 - \delta \leq \|\Pi x\|_2^2 \leq 1 + \delta.$$

$$A \sum_j \eta_j = 0$$

$$\text{مثلثی}$$

$$\|\Pi x\|_2^2 \leq 1 + \delta.$$

$$\|\eta_T\|_2 ~\leq~ (1-\delta)^{-1}(1+\delta)\sum_{j=2}^s \|\eta_{T_j}\|_2$$

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بازگشت به نرم ا

$$\|\eta_T\|_2 \leq (1 - \delta)^{-1}(1 + \delta) \sum_{j=2}^s \|\eta_{T_j}\|_2$$

بازگشت به نرم ا

$$\forall i \in T_{j+1} \quad |\eta_i| \leq \|\eta_{T_j}\|_1/M$$

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$$\forall i \in T_{j+1} \quad |\eta_i| \leq \|\eta_{T_j}\|_1/M$$

$$\|\eta_{T_{j+1}}\|_2 \leq (M(\|\eta_{T_j}\|_1/M)^2)^{1/2}$$

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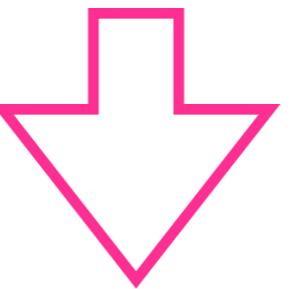
$$\begin{aligned} &\leq (1 - \delta)^{-1}(1 + \delta)/M^{1/2} \sum_{j=1}^s \|\eta_{T_j}\|_1 \\ &= (1 - \delta)^{-1}(1 + \delta)/M^{1/2} \|\eta_{-T}\|_1 \end{aligned}$$

$$\|\eta_T\|_1 \leq (2k)^{1/2}~\|\eta_T\|_2$$

$$\|\eta_T\|_2 ~~~\leq~~~ (1-\delta)^{-1}(1+\delta)/M^{1/2}\|\eta_{-T}\|_1$$

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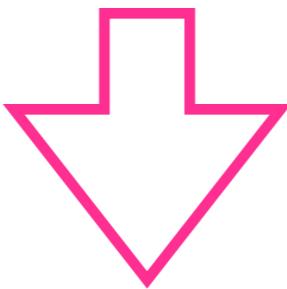
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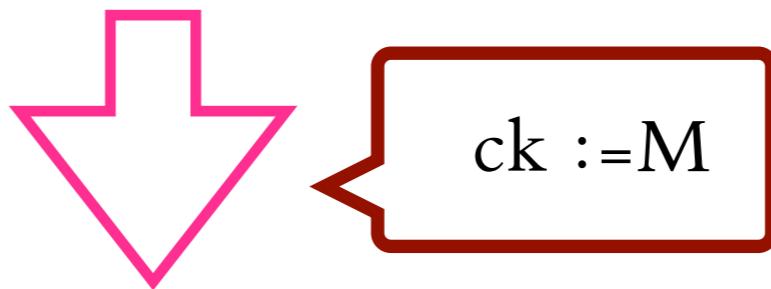
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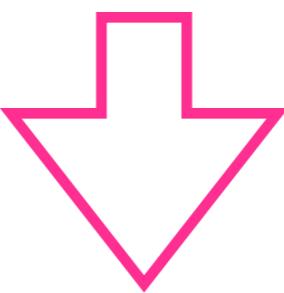
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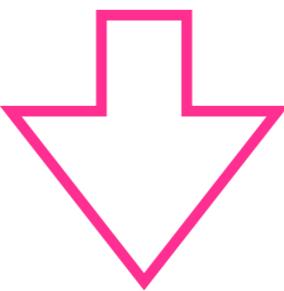
$$\|\eta\|_1 \leq [1 + (1 - \delta)^{-1}(1 + \delta)(2/c)^{1/2}] \|\eta_{-T}\|_1$$

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ck := M

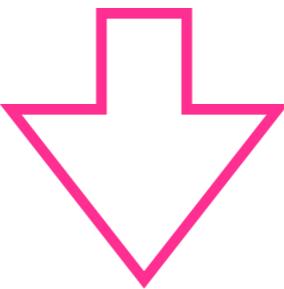
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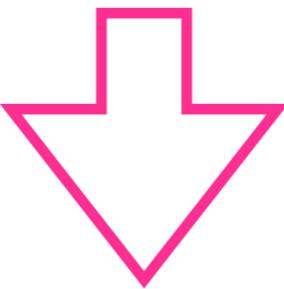
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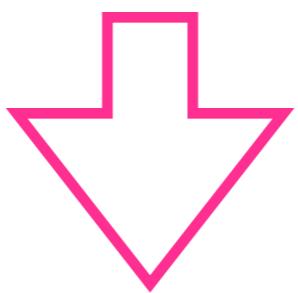


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$\|\eta\|_1 \leq C \|\eta_{-T}\|_1$

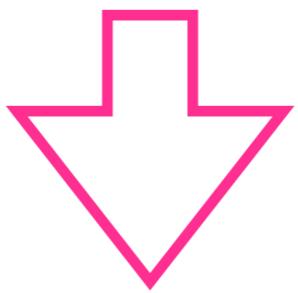
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$$\|\eta_T\|_2 \leq O\left(\frac{1}{\sqrt{k}}\right) \|\eta_{-T}\|_1$$

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Lemma 2 Assume A satisfies the nullspace property of order $2k$ with constant $C < 2$. Then for x^* that minimizes $\|x^*\|_1$ subject to $Ax^* = Ax$ we have

$$\|x - x^*\|_1 \leq \frac{2C}{2-C} Err_1^k(x)$$

$$Err_1^k(x) := \|x_{tail(k)}\|_1$$

: $A\eta = 0$ که η

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$$x^* = \arg \min_{z: \Pi z = \Pi x} \|z\|_1$$

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حکم:

: حکم

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: ایم

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و

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و

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$$\wedge \\ (C - 1)\|\eta_{-T}\|_1$$

:) ن

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$$\|\hat{x} - x\|_2 \leq O\left(\frac{1}{\sqrt{k}}\right) \cdot \|x_{\text{tail}(k)}\|_1$$

احساس فشدگی

روش تکراری



الگوریتم بازیابی (ورودی: y و Π)

$$\min | | z | |_1$$

$$\text{s.t. } \Pi z = y$$

$$\|x - x^*\|_1 \leq \frac{2C}{2-C} Err_1^k(x)$$

$$||\hat{x} - x||_2 \leq O(\frac{1}{\sqrt{k}}).||x_{\text{tail}(k)}||_1$$

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الگوریتم بازیابی با خطأ (ورودی: Π و y و α)

$$\min ||z||_1$$

$$\text{s.t. } ||\Pi z - y||_2 \leq \alpha$$

الگوریتم بازیابی (ورودی: Π و y)

$$\min ||z||_1$$

$$\text{s.t. } \Pi z = y$$

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$$x^{[t]} \rightarrow x^{[t+1]}$$

الگوریتم بازیابی (وروودی: Π و y)

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Theorem 55 ([BD09]). If Π satisfies $(\varepsilon, 3k)$ -RIP for $\varepsilon < \frac{1}{4\sqrt{2}}$, then
 $\forall T \geq 1$

$$||x^{[T+1]} - x||_2 \lesssim 2^{-T} ||x||_2 + ||x_{\text{tail}(k)}||_2 + \frac{1}{\sqrt{k}} ||x_{\text{tail}(k)}||_1 + ||e||_2$$

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فاصله تا جواب

الگوریتم بازیابی (ورودی: Π و y)

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فاصله تا جواب

کران قدیمی

الگوریتم بازیابی (وروودی: Π و y)

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فاصله تا جواب

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قدیمی

کران قدیمی

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فاصله تا جواب

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قدیمی

کران قدیمی

خطای اندازه‌گیری

الگوریتم بازیابی (ورودی: Π و y)

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فاصله تا جواب

هزینه سرعت

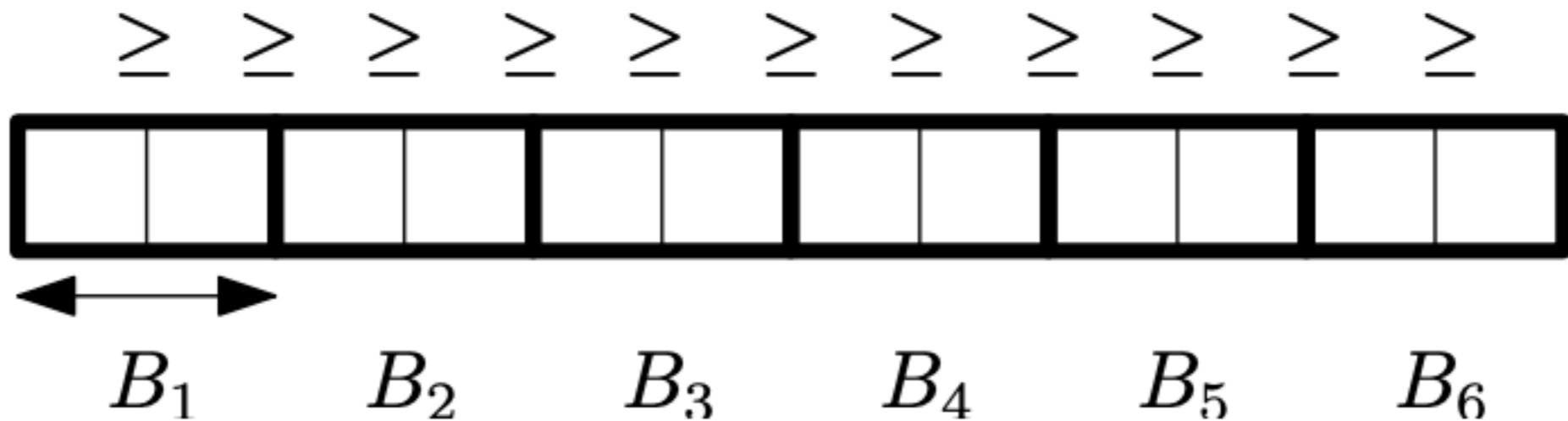
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کران قدیمی

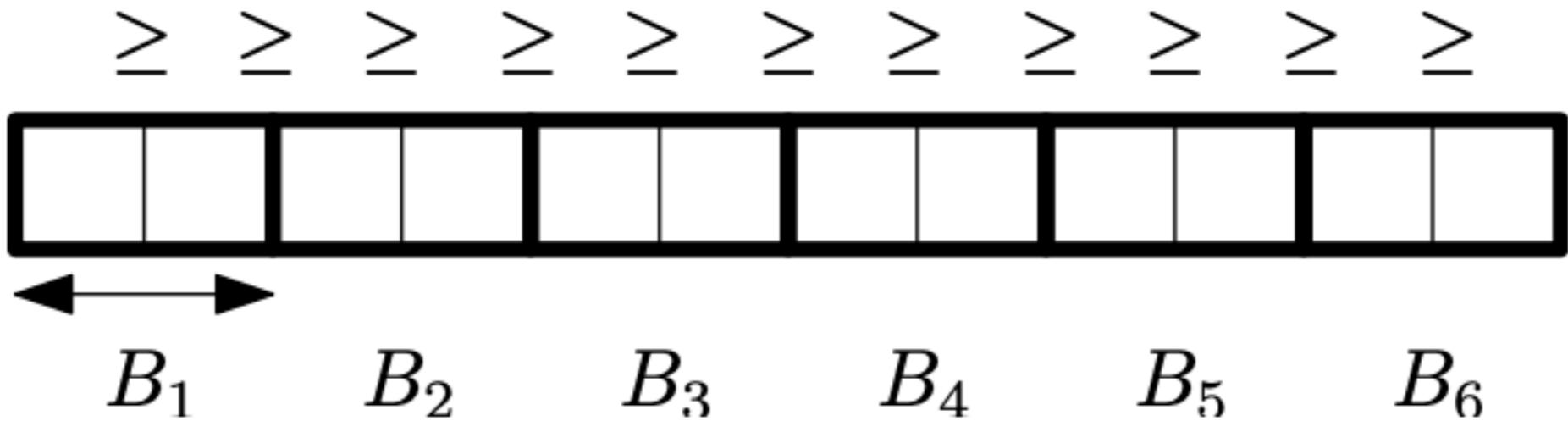
خطای اندازه‌گیری

Claim 56. $\|x_{\text{tail}(2k)}\|_2 \leq \frac{1}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$

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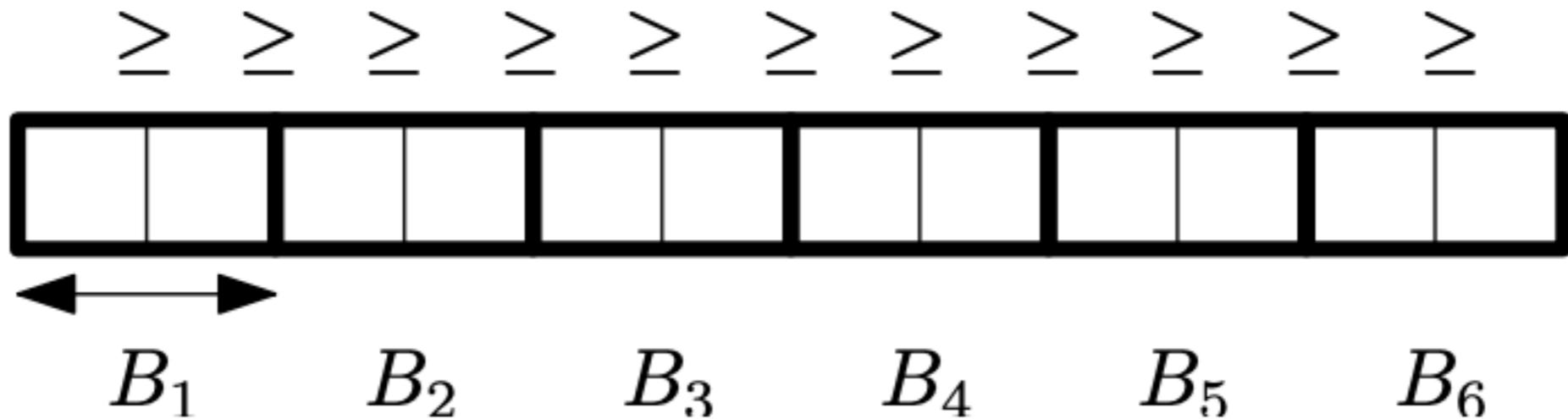


Claim 56. $\|x_{\text{tail}(2k)}\|_2 \leq \frac{1}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$



(از مجموعه قبلی) $|x_j| \leq \frac{1}{k} \sum_{i \in B_{t-1}} |x_i| = \frac{1}{k} \|x_{B_{t-1}}\|_1$

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$$|x_j| \leq \frac{1}{k} \sum_{i \in B_{t-1}} |x_i| = \frac{1}{k} \|x_{B_{t-1}}\|_1 \quad (\text{از مجموعه قبلی})$$

$$\|x_{\text{tail}(2k)}\|_2^2 = \sum_{t=3}^{n/k} \|x_{B_t}\|_2^2 \leq \sum_{t=3}^{n/k} k \cdot \left(\frac{\|x_{B_{t-1}}\|_1}{k} \right)^2 = \frac{1}{k} \sum_{t=2}^{n/k} \|x_{B_t}\|_1^2$$

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$$\|x_{\text{tail}(2k)}\|_2 \leq \frac{1}{k} \cdot \sqrt{\sum_{t=2}^{n/k} \|x_{B_t}\|_1^2} \leq \frac{1}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$$

$$x^{[t]} \rightarrow x^{[t+1]}$$

Theorem 55 ([BD09]). If Π satisfies $(\varepsilon, 3k)$ -RIP for $\varepsilon < \frac{1}{4\sqrt{2}}$, then
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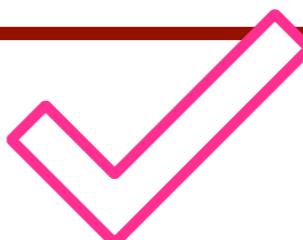
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فرض می کنیم $x : k$ -تنک

فرض می‌کنیم x_{-k} -تنک

$$\Pi x + e = \Pi(x_{\text{head}(k)} + x_{\text{tail}(k)}) + e = \Pi x_{\text{head}(k)} + \underbrace{(\Pi x_{\text{tail}(k)} + e)}_{\tilde{e}}$$

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$$\|\tilde{e}\|_2 \underset{\Delta\text{-inequality}}{\leq} \|e\|_2 + \|\Pi x_{\text{tail}(k)}\|_2$$

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$$\leq \|e\|_2 + \sum_{t=2} \|\Pi x_{B_t}\|_2$$

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$$\leq \|e\|_2 + \sum_{t=2} \|\Pi x_{B_t}\|_2 \stackrel{\text{RIP}}{\leq} \|e\|_2 + (1 + \varepsilon) \sum_{t=2} \|x_{B_t}\|$$

فرض می کنیم $x_{\text{tail}(k)}$ -تنک

$$\Pi x + e = \Pi(x_{\text{head}(k)} + x_{\text{tail}(k)}) + e = \Pi x_{\text{head}(k)} + \underbrace{(\Pi x_{\text{tail}(k)} + e)}_{\tilde{e}}$$

$$\|\tilde{e}\|_2 \stackrel{\Delta\text{-inequality}}{\leq} \|e\|_2 + \|\Pi x_{\text{tail}(k)}\|_2 = \|e\|_2 + \left\| \sum_{t=2} \Pi x_{B_t} \right\|_2$$

$$\leq \|e\|_2 + \sum_{t=2} \|\Pi x_{B_t}\|_2 \stackrel{\text{RIP}}{\leq} \|e\|_2 + (1 + \varepsilon) \sum_{t=2} \|x_{B_t}\|_1$$

$$\leq \|e\|_2 + \frac{1 + \varepsilon}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$$

$$x^{[t]} \rightarrow x^{[t+1]}$$

Algorithm 1 Iterative Hard Thresholding (IHT).

```
1: function IHT( $\Pi, y (= \Pi x + e), k, T$ )
2:    $x^{[1]} \leftarrow 0$ 
3:   for  $t = 1 \dots T$  do
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5:   end for
6:   return  $x^{T+1}$ 
7: end function
```

تکساز $\rightarrow H_k$

تحليل الگوریتم : HIT

$$r^{[t]} := x - x^{[t]}$$

فاصله تا جواب

تحليل الگوریتم : HIT

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فاصله تا جواب

$$\|r^{[t+1]}\|_2 = \|x - x^{[t+1]}\|_2$$

تحليل الگوریتم : HIT

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فاصله تا جواب

$$\|r^{[t+1]}\|_2 = \|x - x^{[t+1]}\|_2$$

•
•
•

$$\leq \frac{1}{2} \|r^{[t]}\|_2 + 3\|e\|_2$$

تحليل الگوریتم : HIT

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فاصله تا جواب

$$\|r^{[t+1]}\|_2 = \|x - x^{[t+1]}\|_2$$

⋮

$$\leq \frac{1}{2} \|r^{[t]}\|_2 + 3\|e\|_2$$

$$\leq 2^{-t} \|x\|_2 + 3(\sum_{i=0}^{t-1} 2^{-i}) \|e\|_2$$

تحليل الگوریتم : HIT

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فاصله تا جواب

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$\|r^{[0]}\|_2$

< 6

$$x^{[t]} \rightarrow x^{[t+1]}$$

Theorem 55 ([BD09]). If Π satisfies $(\varepsilon, 3k)$ -RIP for $\varepsilon < \frac{1}{4\sqrt{2}}$, then
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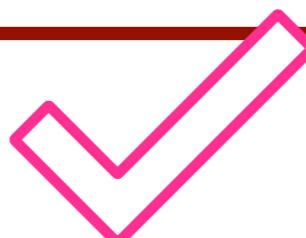
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$$2^{-t} \|x\|_2 + 6 \|e\|_2$$

$$x^{[t]} \rightarrow x^{[t+1]}$$

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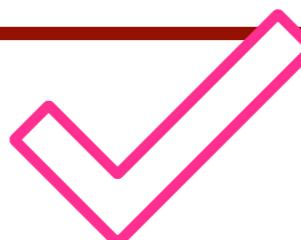
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$$2^{-t} \|x\|_2 + 6 \|e\|_2 + \frac{1 + \varepsilon}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$$

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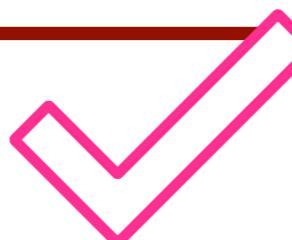
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زمان:

$$2^{-t} \|x\|_2 + 6 \|e\|_2 + \frac{1 + \varepsilon}{\sqrt{k}} \|x_{\text{tail}(k)}\|_1$$

احساس فُشَرْدَگی

مدلی



حفظ فاصله برای بردارهای k -تنک

$$x \rightarrow y = \Pi x$$

k -تنک، عدد x

$$\Omega_{n,k} = \binom{[n]}{k}$$

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حفظ فاصله برای بردارهای k -تنک

$$x \rightarrow y = \Pi x$$

M : عضو x

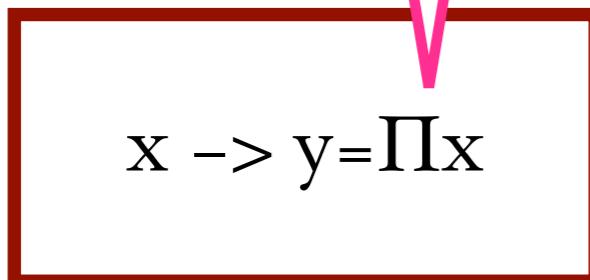
k -تنک، k عدد x

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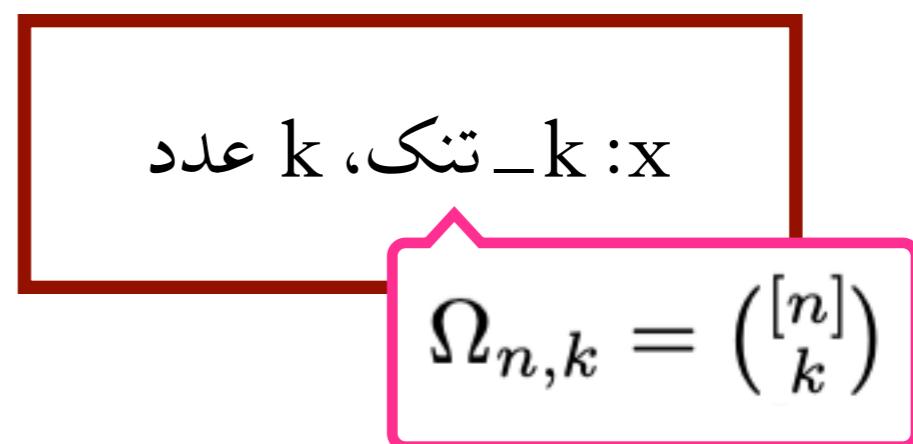
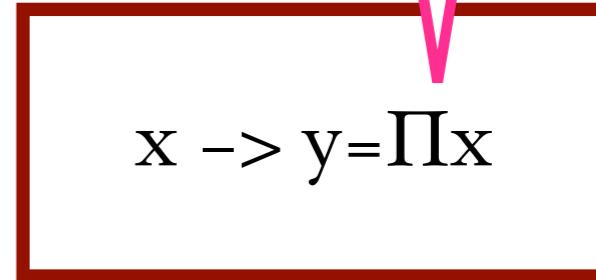
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```

حفظ فاصله برای بردارهای سه تا M



حفظ فاصله برای بردارهای ۳k_تنک



Algorithm 1 Iterative Hard Thresholding (IHT).

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```

حفظ فاصله برای بردارهای سه تا M

$$x \rightarrow y = \Pi x$$

M : عضو x

حفظ فاصله برای بردارهای k -تنک

$$x \rightarrow y = \Pi x$$

k -تنک، k عدد x

$$\Omega_{n,k} = \binom{[n]}{k}$$

Algorithm 2 Model Based Iterative Hard Thresholding (MB-IHT).

```
1: function IHT( $\Pi, y (= \Pi x + e), k, T$ )
2:    $x^{[1]} \leftarrow 0$ 
3:   for  $t = 1 \cdots T$  do
4:      $x^{[t+1]} \leftarrow P_{\mathcal{M}}(x^{[t]} + \Pi^{\top}(y - \Pi x^{[t]}))$      $\triangleright$  Hard thresholding
      operator (project  $a^{[t+1]}$  on  $\mathcal{M}$ )
5:   end for
6:   return  $x^{T+1}$ 
7: end function
```
