

빅 데이터 혁신 공유 대학

파이썬으로 배우는 데이터 구조

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Data Structures in Python

Chapter 6

- Hash Table
- Collision Resolution
- Double Hashing & Rehashing
- HashMap Coding

Agenda & Readings

- Agenda
 - Hashing
 - Hash Table
 - Hash Function
- Reference:
 - Problem Solving with Algorithms and Data Structures
 - Chapter 5 - Hashing

Overview

- Hashing or Hash Table Data Structure:

- Data structures so far

Array of size n	unsorted list	sorted array	Trees BST – average AVL – worst	Heap, Priority Queue	Hashing
insert	find+O(1)	O(n)	O(log n)	O(log n)	
find	O(n)	O(log n)	O(log n)	O(log n)	
remove	find+O(1)	O(n)	O(log n)	O(log n)	

Overview

- Hashing or Hash Table Data Structure:
supports insertion, deletion and search in average case constant time **$O(1)$** .

- Data structures so far

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insert	find+ $O(1)$	$O(n)$	$O(\log n)$	$O(\log n)$	$O(1)$
find	$O(n)$	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(1)$
remove	find+ $O(1)$	$O(n)$	$O(\log n)$	$O(\log n)$	$O(1)$

Overview

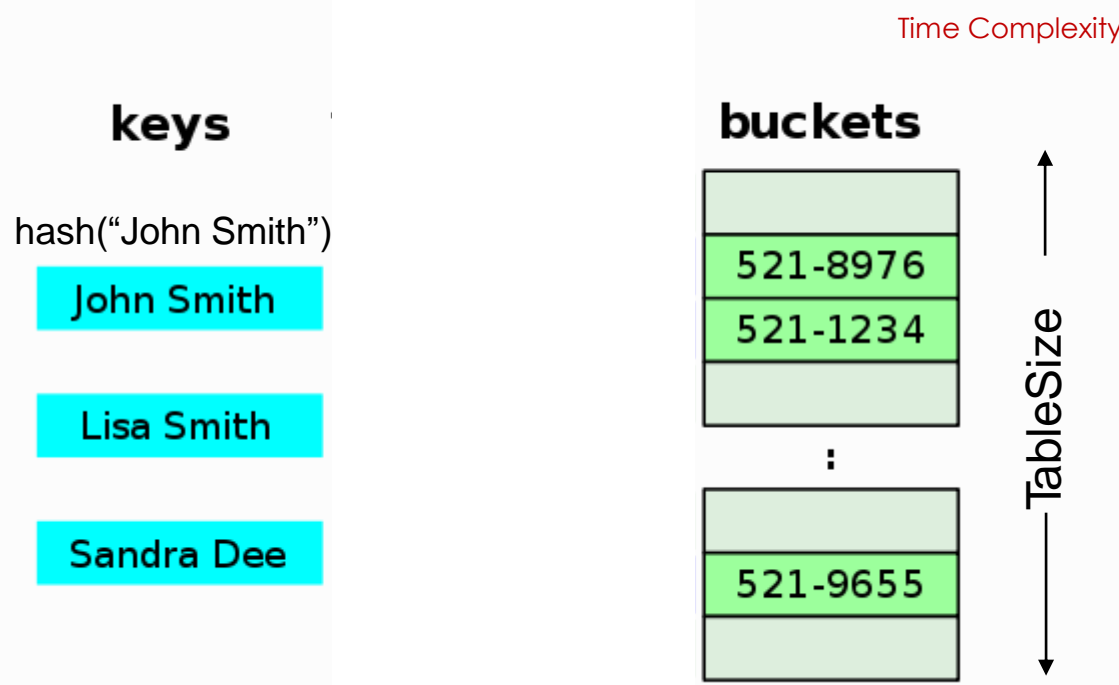
- Hashing or Hash Table Data Structure:
supports insertion, deletion and search in average case constant time **$O(1)$** .
- Hash table
 - It is data structure that stores **key-value pairs**.
 - The key is sent to a **hash function** that performs arithmetic operations on it.
 - The result is called **hash value** that is the **index of the key-value pair** in the **hash table**.
- Hash function
 - $\text{hash}(\text{key}) \rightarrow \text{integer value}$
 - $\text{hash}(\text{"string key"}) \rightarrow \text{integer value}$

Hash Table

- Hash table is an array of fixed size elements

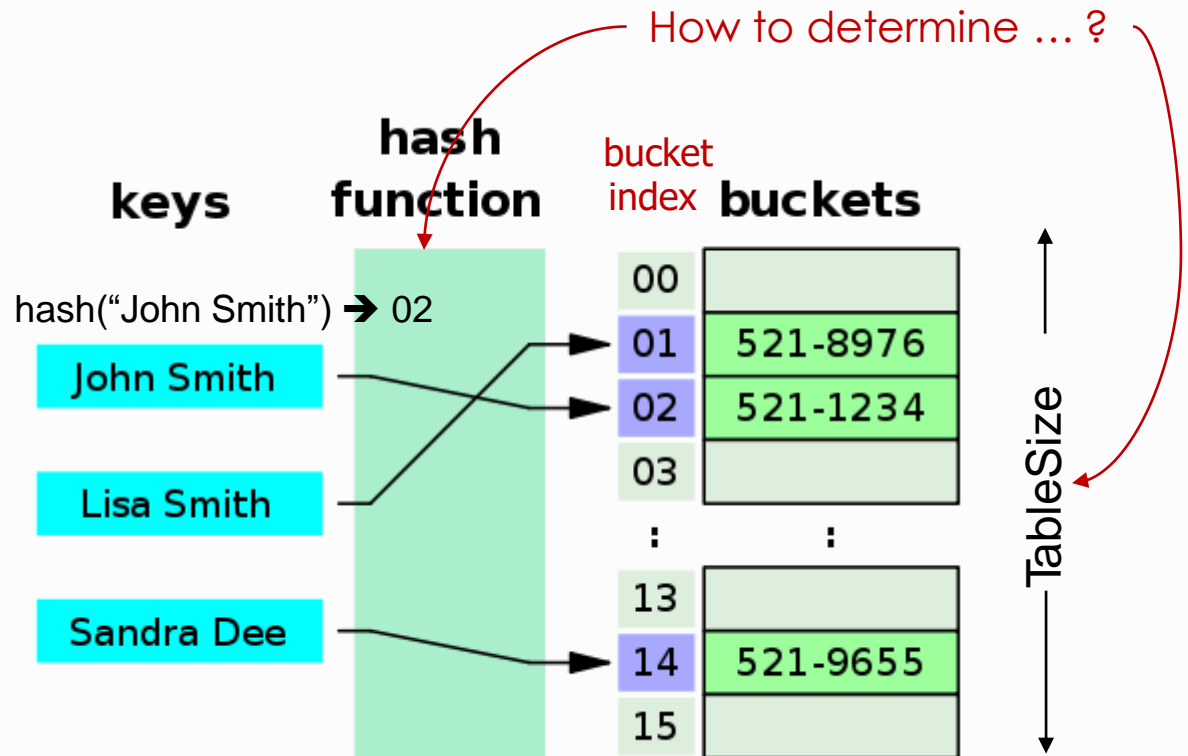
Let us suppose that there are one billion of names and numbers.

- Find, insert, and remove a number by a given name in **$O(1)$** .



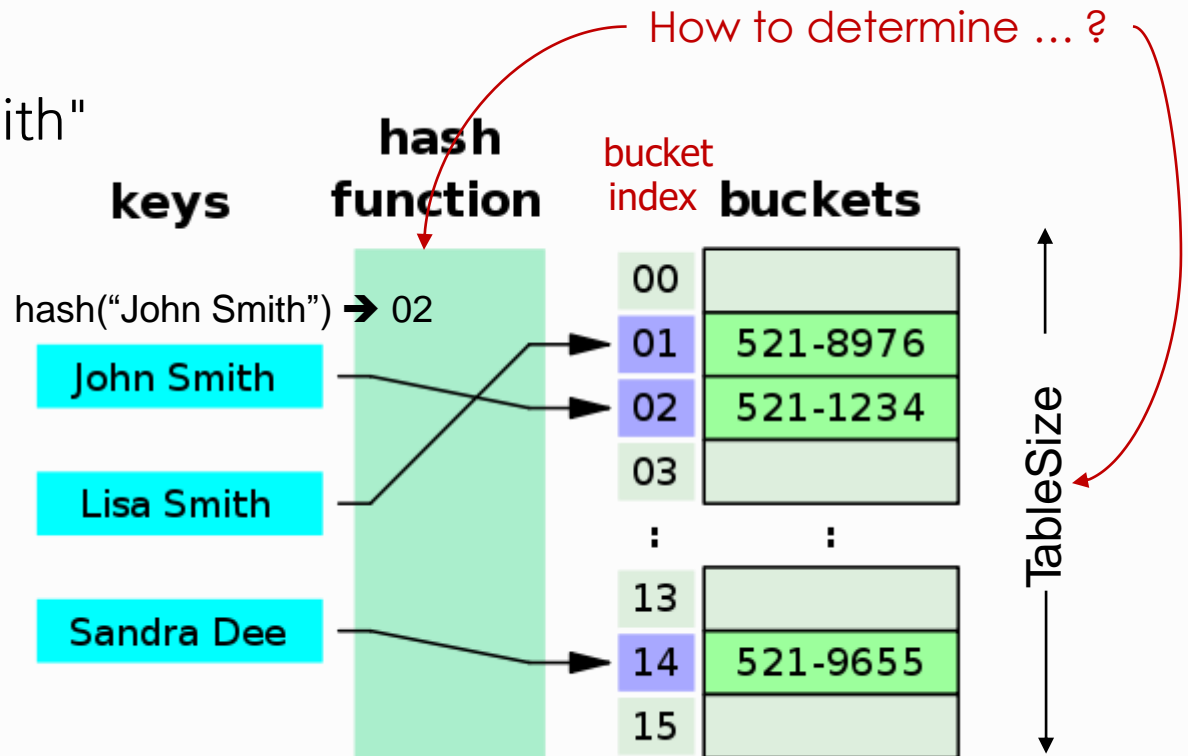
Hash Table

- Hash table is an array of fixed size elements
- Array elements indexed by a key mapped to a bucket index[0 .. TableSize-1]
- Mapping from key to index using hash(), hash function
 - e.g., hash("John Smith") → 02



Hash Table

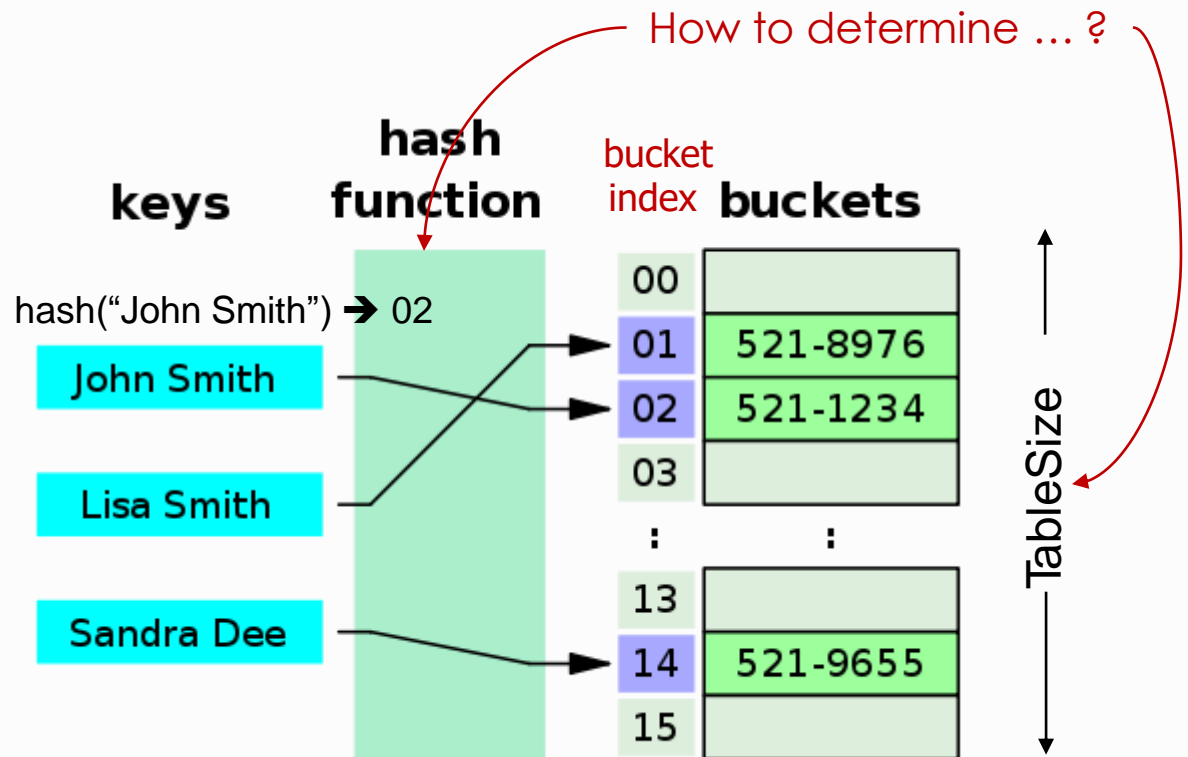
- insert
 - `HashTable[hash("John Smith")] = <"John Smith", 521-1234>`
- remove
 - `HashTable[hash("John Smith")] = None`
- find
 - `HashTable[hash("John Smith")]` returns the element hashed for "John Smith"



What happens
if `hash("John Smith") == hash("Joe Blow")`?
"Collision"

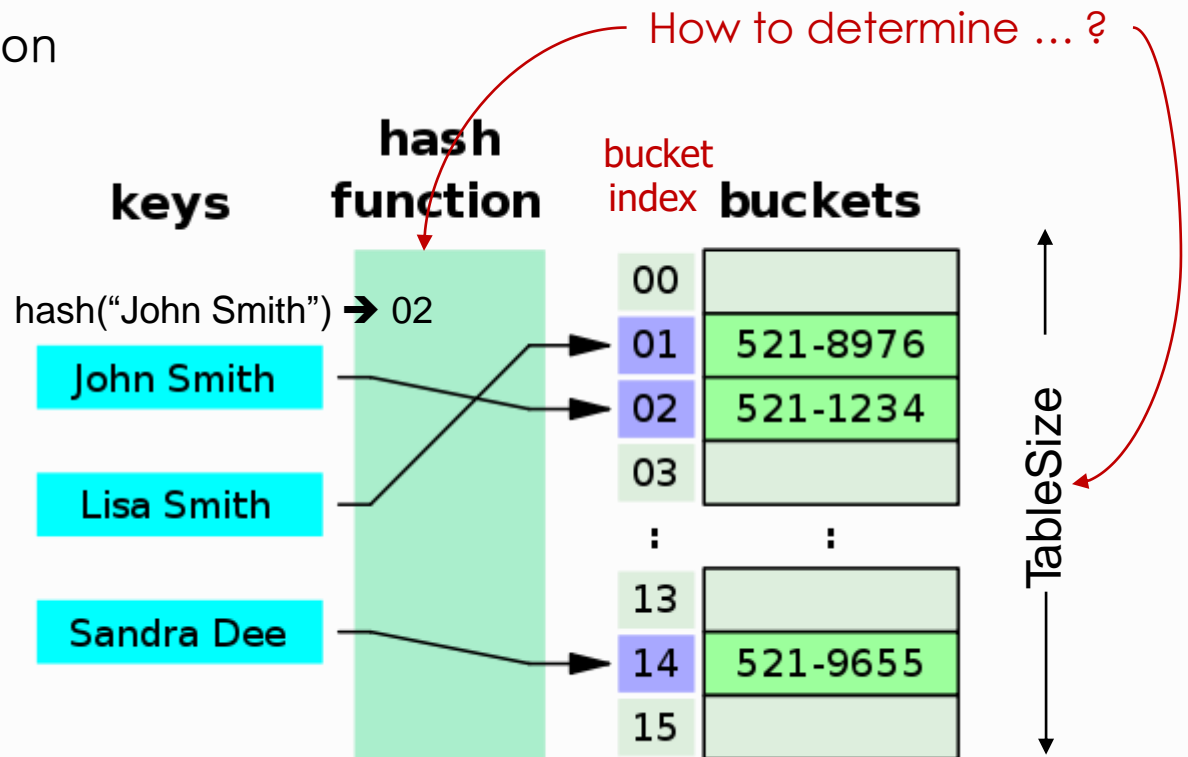
Hash Table

- Factors affecting Hash Table Design
 - Hash function
 - Table size - Usually fixed at the start
 - Collision handling schemes - Array or Linked List



Hash Function

- It maps an element's key into a valid hash table index
 - $\text{hash}(\text{key}) \rightarrow \text{hash table index}$
- Note that this is (slightly) different from saying:
 - $\text{hash}(\text{string}) \rightarrow \text{int}$
 - Because the key can be of any type
 - e.g., " $\text{hash}(\text{int}) \rightarrow \text{int}$ " is also a hash function



Hash Function Properties

- It maps an element's key into a valid hash table index
 - $\text{hash}(\text{key}) \rightarrow \text{hash table index}$
- It maps key to integer
 - Constraint: Integer should be between **[0, TableSize-1]**
- A hash function can result in a many-to-one mapping (causing collision)
 - Collision occurs when hash function maps two or more keys to same array index
- Collisions **cannot** be avoided but its chances can be reduced using a "good" hash function


Hash Function - Effective use of table size

- Simple hash function (assume integer keys)
 - $\text{hash}(\text{Key}) = \text{Key} \% \text{TableSize}$
- For random keys, $\text{hash}()$ distributes keys evenly over table
 - What if $\text{TableSize} = 100$ and keys are ALL multiples of 10?
 - Better if TableSize is a **prime number**

Hash Function Example: String Keys

- Using a very simple function to map strings to integers:
 - Add up character ASCII values (0-255) to produce integer keys
 - e.g., "abcd" = 97 + 98 + 99 + 100 = 394
 - $\text{hash}(\text{"abcd"}) = 394 \% \text{TableSize}$
- Potential problems:
 - Anagrams will map to the same index
 - $\text{hash}(\text{"abcd"}) = \text{hash}(\text{"dbac"})$
 - Small strings may not use all of table
 - $\text{strlen}(s) * 255 < \text{TableSize}$
 - Time proportional to length of the string

Hash Function Example: String Keys

- Another approach:
 - Treat first 3 characters of string as base-27 integer (26 letters plus space)
 - e.g., $\text{Key} = s[0] + (27^1 * s[1]) + (27^2 * s[2])$
 - Better than previous approach because ...
- But, potential problems:
- - Apple
 - Apply
 - Appointment
 - Apricot collision

Hash Function Example: String Keys

- Last approach:
 - Use all N characters of string as an N-digit and base-K number
 - Choose K to be prime number larger than number of different digits (characters)
 - i.e., $K = 29, 31, 37$
 - If $L = \text{Length of string } S$, then

$$\text{hash}(S) = \sum_{i=0}^{L-1} S[L-1-i] * 37^i \% \text{TableSize} \quad (1)$$

- Use Horner's rule to compute $\text{hash}(S)$.
 - Limit L for long strings
- Potential problems
 - Overflow
 - Larger runtime

```
# a hash function for strings
hash(key, tablesize)
    code = 0
    for x in key:
        code = code * 37 + x
    code %= tablesize
    if code < 0: code += tablesize
    return code
```


Summary

- Using a hash table we can, on average (if table large enough and hash function suitable), insert, delete and search for items in constant time - **$O(1)$** .
- The **hash function** is the mapping between an item and the slot where the item is stored.
- A **collision** occurs when an item is mapped to an occupied slot.
- A **perfect hash function** is able to map m items into a table of size m with no collisions. Perfect hash functions are hard to come by.
- Handling collisions systematically is required - **collision resolution**.

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