Data Structures – CST 201 Module - 2

Syllabus

- Polynomial representation using Arrays
- Sparse matrix
- Stacks
 - Evaluation of Expressions
- Queues
 - Circular Queues
 - Priority Queues
 - Double Ended Queues
- Linear Search
- Binary Search

PRIORITY QUEUE

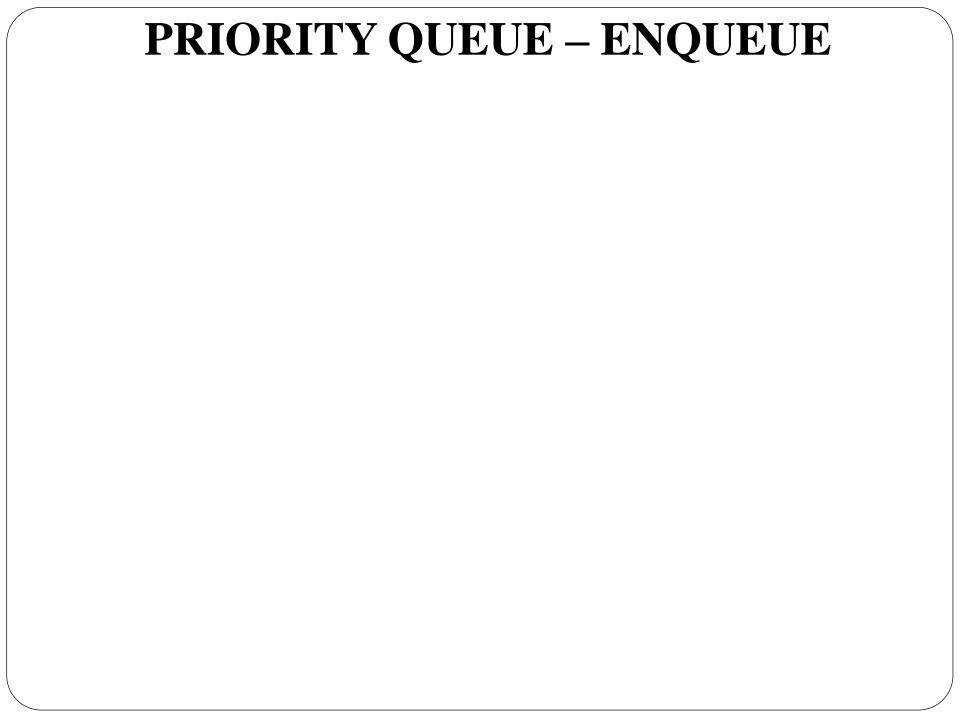
- Priority Queue is an extension of queue with following properties.
 - Every item has a priority associated with it.
 - An element with high priority is dequeued before an element with low priority.
 - If two elements have the same priority, they are served according to their order in the queue.

PRIORITY QUEUE- Operations

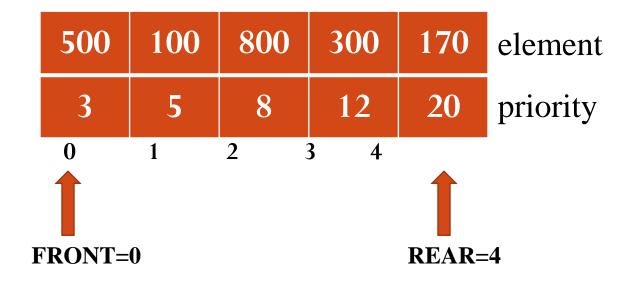
• **ENQUEUE**: Insert an element in the queue based on priority

• **DEQUEUE:** Delete highest priority element from the queue

• **DISPLAY**: Display the contents of the Queue



Case 1: FRONT=0 and REAR=SIZE-1
Priority Queue is FULL

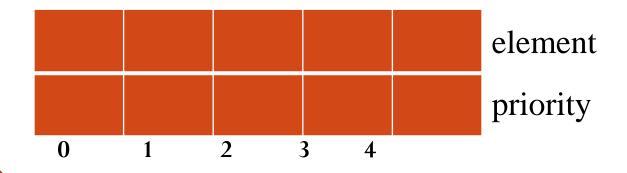


Case 2: FRONT=-1 and REAR=-1

FRONT=REAR=0

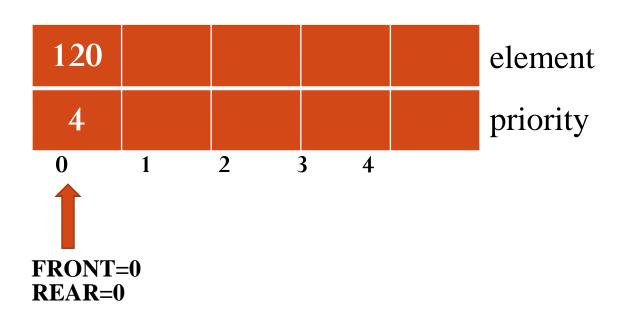
A[REAR].item=120

A[REAR].priority=4



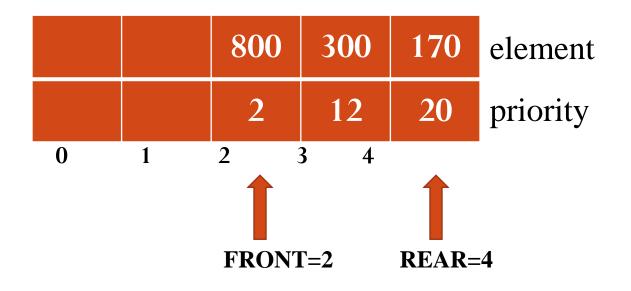
FRONT=-1 REAR=-1

Case 2: FRONT=-1 and REAR=-1
FRONT=REAR=0
A[REAR].item=120
A[REAR].priority=4



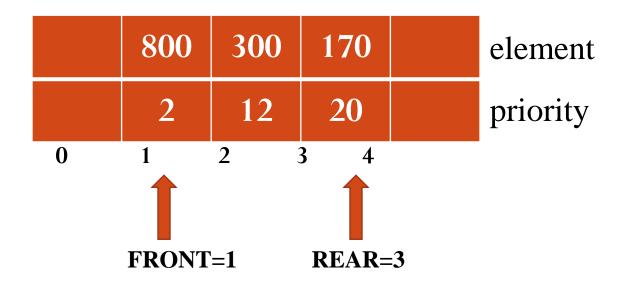
Case 3: if REAR=SIZE-1

Shift all elements one position to left



Case 3: if REAR=SIZE-1

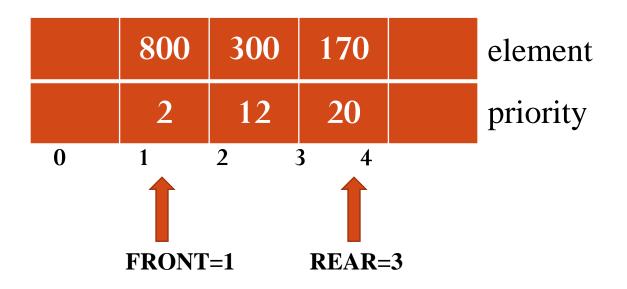
Shift all elements one position to left



Case 3: if REAR=SIZE-1

Shift all elements one position to left

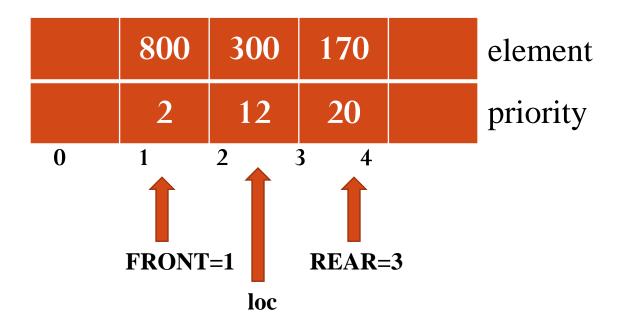
Find the location where the new elmt is to be inserted



Case 3: if REAR=SIZE-1

Shift all elements one position to left

Find the location where the new elmt is to be inserted

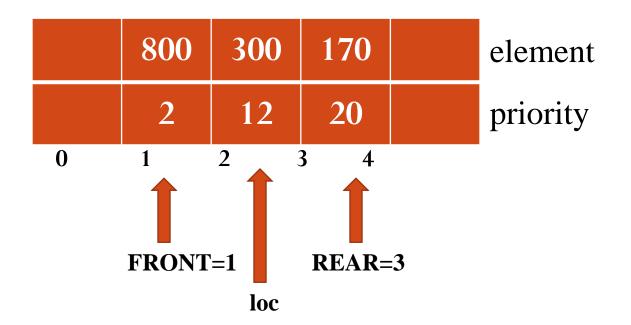


Case 3: if REAR=SIZE-1

Shift all elements one position to left

Find the location where the new elmt is to be inserted

Shift loc to REAR elements one position to right

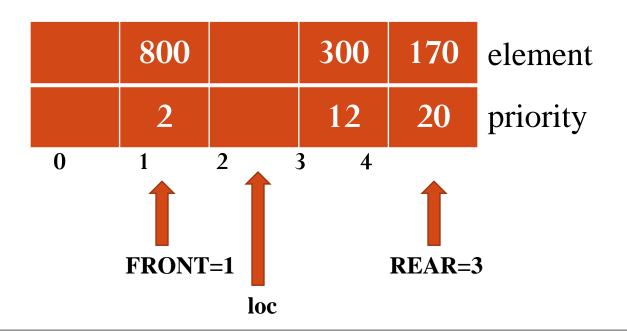


Case 3: if REAR=SIZE-1

Shift all elements one position to left

Find the location where the new elmt is to be inserted

Shift loc to REAR elements one position to right



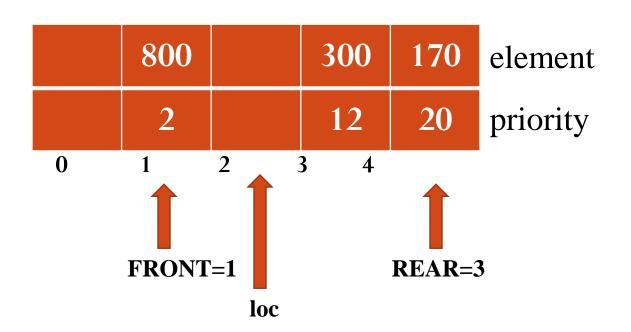
Case 3: if REAR=SIZE-1

Shift all elements one position to left

Find the location where the new elmt is to be inserted

Shift loc to REAR elements one position to right

Insert the data at the index loc



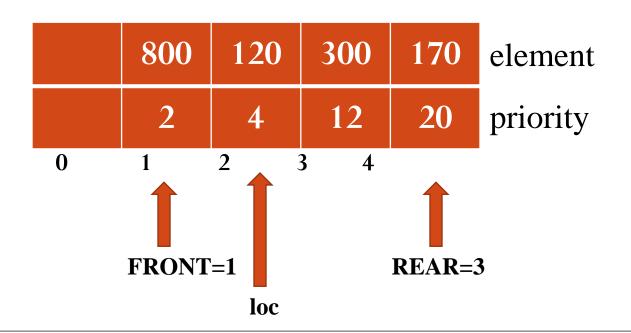
Case 3: if REAR=SIZE-1

Shift all elements one position to left

Find the location where the new elmt is to be inserted

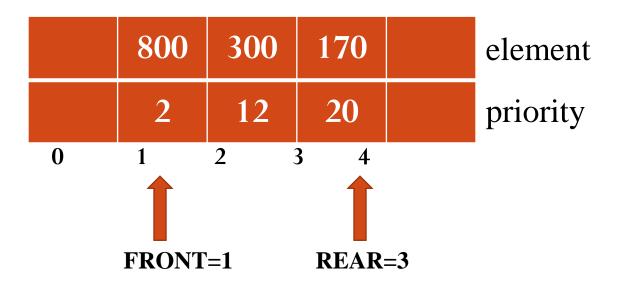
Shift loc to REAR elements one position to right

Insert the data at the index loc



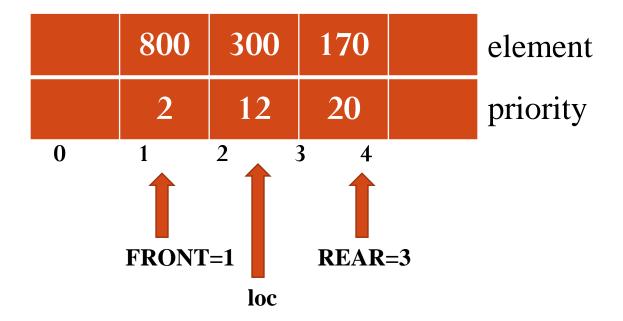
Case 4: All other cases

Find the location where the new elmt is to be inserted



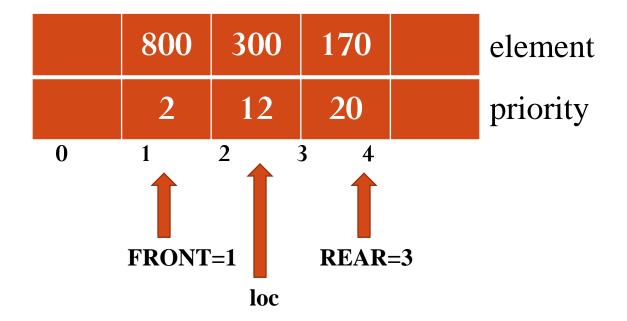
Case 4: All other cases

Find the location where the new elmt is to be inserted



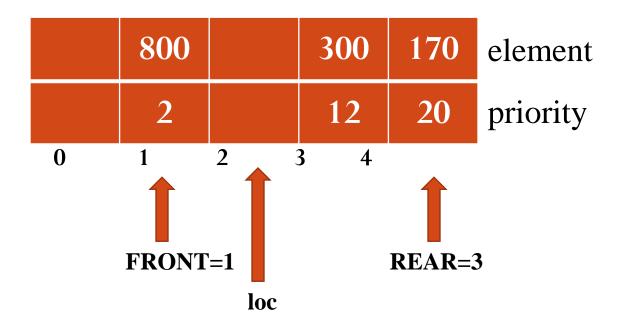
Case 4: All other cases

Find the location where the new elmt is to be inserted Shift elements from loc to REAR one position right



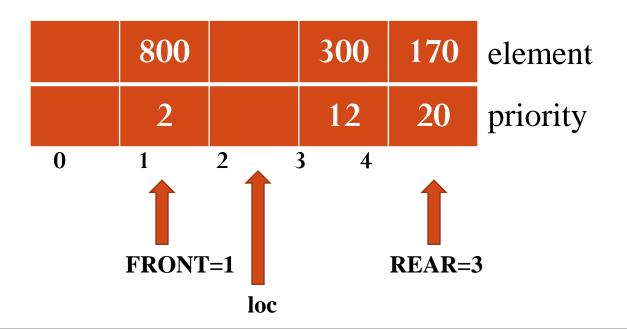
Case 4: All other cases

Find the location where the new elmt is to be inserted Shift elements from loc to REAR one position right



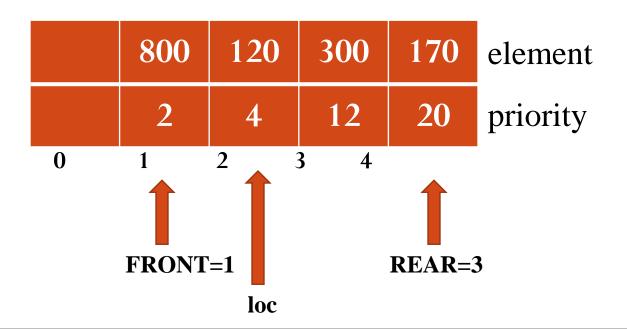
Case 4: All other cases

Find the location where the new elmt is to be inserted
Shift elements from loc to REAR one position right
Insert the data at the index loc



Case 4: All other cases

Find the location where the new elmt is to be inserted Shift elements from loc to REAR one position right Insert the data at the index loc



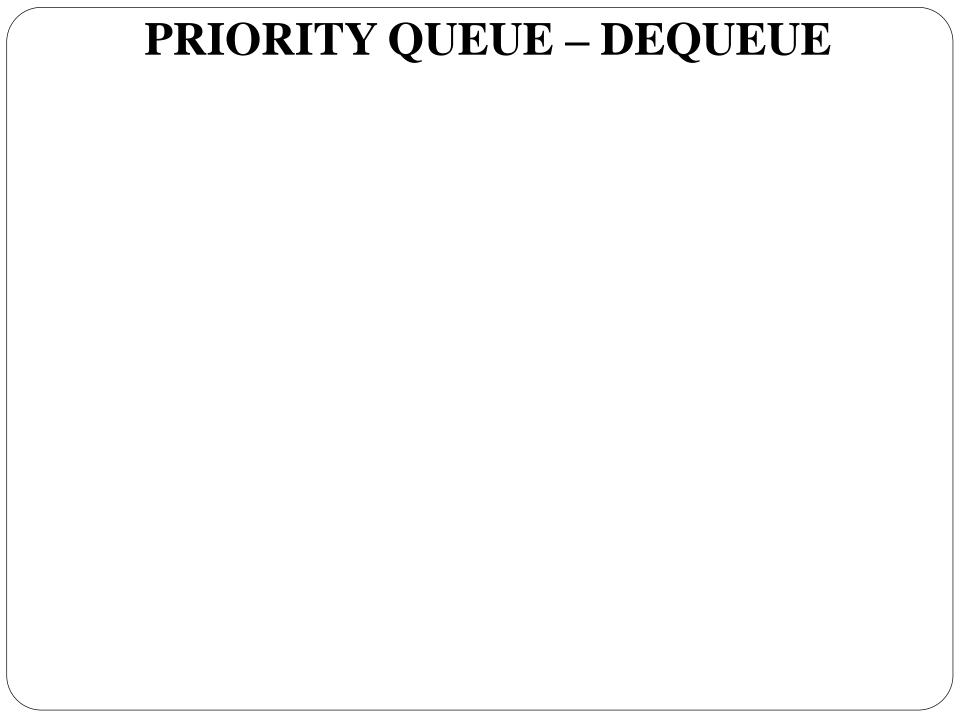
PRIORITY QUEUE – ENQUEUE Algorithm

```
Algorithm ENQUEUE_PQ(ITEM,PRIORITY)
```

```
if FRONT=0 and REAR=SIZE-1 then
      Print "Priority Queue is FULL"
else if FRONT=-1 then
  FRONT=REAR=0
  A[FRONT].item=ITEM
  A[FRONT].priority=PRIORITY
```

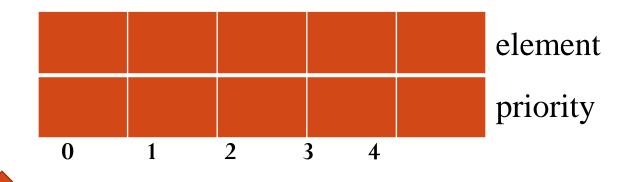
```
else if REAR=SIZE-1 then
        for i=FRONT to REAR do
                A[i-1]=A[i]
        FRONT=FRONT-1
        REAR=REAR-1
        for i=REAR to FRONT do
                if A[i].priority<PRIORITY then
                        break;
        loc=i+1
        for i=REAR to loc do
                A[i+1]=A[i]
        A[loc].item=ITEM
        A[loc].priority=PRIORITY
        REAR=REAR+1
```

```
else
       for i=REAR to FRONT do
               if A[i].priority<PRIORITY then
                        break;
        loc=i+1
       for i=REAR to loc do
               A[i+1]=A[i]
        A[loc].item=ITEM
       A[loc].priority=PRIORITY
        REAR=REAR+1
```



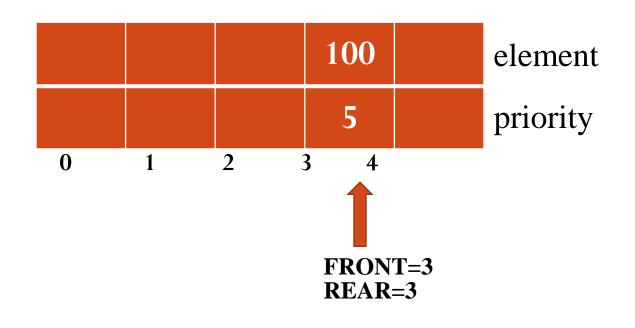
Case 1: FRONT=-1 and REAR=-1

Print "Priority Queue is EMPTY"

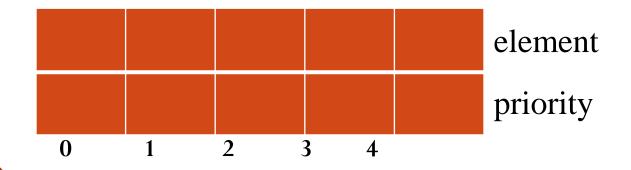


FRONT=-1 REAR=-1

Case 2: FRONT=REAR
FRONT=REAR=-1

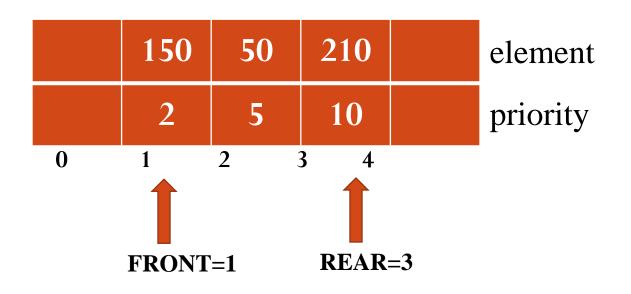


Case 2: FRONT=REAR
FRONT=REAR=-1

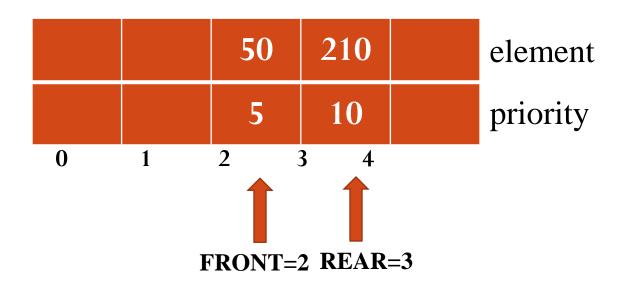


FRONT=-1 REAR=-1

Case 3: Queue contains more than one elements FRONT=FRONT+1



Case 3: Queue contains more than one elements FRONT=FRONT+1



PRIORITY QUEUE - DEQUEUE Algorithm

```
Algorithm DEQUEUE_PQ()
```

```
if FRONT=-1then
      Print "Priority Queue is EMPTY"
else if FRONT=REAR then
      Print "Dequeued item is "A[FRONT].item
      FRONT=REAR=-1
else
      Print "Dequeued item is "A[FRONT].item
      FRONT=FRONT+1
```

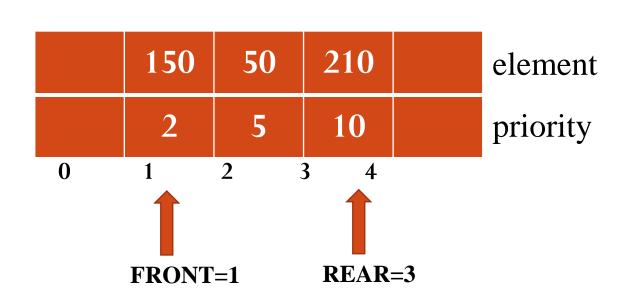
PRIORITY QUEUE – DEQUEUE Algorithm

Algorithm DISPLAY_PQ()

```
{ if FRONT=-1then
Print "Priority Queue is EMPTY" else
```

for i=FRONT to REAR do

Print A[i].item



PRIORITY QUEUE- Various States

- 1. Priority Queue is Empty: FRONT=-1 & REAR=-1
- 2. Priority Queue is Full: FRONT = 0 and REAR = SIZE 1
- 3. Priority Queue contains only one element: FRONT=REAR
- 4. Total elements in the Priority Queue : REAR-FRONT+1

PRIORITY QUEUE IMPLEMENTATION

Using ordered Array:

- Elements are inserted in the sorted order of their priority. The time complexity = O(n)
- Deletion operation is performed from the front end.
 The time complexity= O(1)

Using unordered Array:

- Elements are inserted at any end. The time complexity = O(1)
- For deletion, search an element in the Queue with highest priority. The time complexity = O(n)

DIFFERENT PRIORITY QUEUES

• Max-Priority Queue: Element with highest priority is served first

• Min-Priority Queue: Element with lowest priority is served first.

APPLICATIONS OF PRIORITY QUEUE

- CPU Scheduling
- Graph algorithms like Dijkstra's shortest path algorithm, Prim's Minimum Spanning Tree, etc
- All queue applications where priority is involved

```
PRIORITY QUEUE USING ARRAY-PROGRAM
#include<stdio.h>
int size, front, rear;
struct PQ
      int item, priority;
}A[20];
void display()
      int i;
      if(front==-1)
             printf("queue is EMPTY");
      else
             for(i=front;i<=rear;i++)
                    printf("%d\t",A[i].item);
```

```
void enqueue(int ITEM,int PRIORITY)
      int i,loc;
      if(front==0 && rear==size-1)
            printf("queue is FULL");
      else if(front==-1)
            front=0;
            rear=0;
             A[rear].item=ITEM;
            A[rear].priority=PRIORITY;
```

```
else
   if(rear==size-1)
       for(i=front;i<=rear;i++)
               A[i-1]=A[i];
       front--;
       rear--;
                                          for(i=rear;i>=loc;i--)
   for(i=rear;i>=front;i--)
                                                  A[i+1]=A[i];
       if(A[i].priority<PRIORITY)
               break;
                                          A[loc].item=ITEM;
                                          A[loc].priority=PRIORITY;
                                          rear++;
   loc=i+1;
```

```
void dequeue()
       if(front==-1)
               printf("Queue is empty");
       else if(front==rear)
               printf("deleted item is %d",A[front].item);
               front=-1;
               rear=-1;
       else
               printf("deleted item is %d",A[front].item);
               front++;
```

```
void main()
 int opt, item, prio;
 front=-1;
 rear=-1;
 printf("Enter the size of the queue");
  scanf("%d",&size);
  do
      printf("\nEnter the option\n");
      printf("1.Enqueue\n2.Dequeue\n3.Display\n4.Exit\n");
      scanf("%d",&opt);
```

```
switch(opt)
           case 1:printf("Enter the item and priority");
                  scanf("%d%d",&item,&prio);
                  enqueue(item,prio);
                  break;
           case 2:dequeue();
                  break;
           case 3:display();
                  break;
           case 4:break;
           default:printf("Enter a valid option");
}while(opt!=4);
```