

## Lecture 2. Reconstruction Attacks.

- De-identified data  $\times$ ; releasing "Aggregate" Statistics?
- Warmup : Difference Attacks
- Reconstruction examples
- Reconstruction Formulation  
Linear Attacks [Dinur & Nissim 03]

## 2. Targeting

**Location**

Country:

Everywhere  
 By State/Province  
 By City  
San Francisco, CA  
 Include cities within 50 miles.

**Demographics**

Age: 24 - 30

Sex:  All  Men  Women

Birthday:  Target people on their birthdays

Interested In:  All  Men  Women

Relationship:  All  Single  Engaged  
 In a Relationship  Married

Languages:   
 Fewer Demographic Options

**Likes & Interests**

Enter an interest

**Education & Work**

Education:  All  College Grad  
Harvard  
Enter a major  
 In College  
 In High School

Workplaces:   
 Hide Education & Work Options

Facebook ad campaign targeting interface.

Ref: Korolova,

"Privacy violation Using  
Microtargeted Ads: A Case Study"

## Warmup : Difference Attacks

Q: How many people were born in 1992 and live in Zipcode 15206 and have a heart disease?

A: ~~1~~ less than 5.

Q: How many faculty members @ CMU joined before 9/1/2020 and have had a heart disease?

A: 37

Q: How many faculty members @ CMU joined before 9/2/2020 and have had a heart disease?

A: 38

## Reconstruction in the US Census.

- 3 Males
- Ages  $A \leq B \leq C$
- $1 \leq A \leq B \leq C \leq 125$
- Median = 30.  
 $B = 30$   
 $(A \leq 30), C \geq 30.$
- Mean = 44.  
 $\frac{A+B+C}{3} = 44.$   
 $\Rightarrow A+C = 102.$

$(A, C)$  has 30 possibilities.

Before:  $(125)^3$  possibilities

TABLE 1: FICTIONAL STATISTICAL DATA FOR A FICTIONAL BLOCK

STATISTIC	GROUP	AGE		
		COUNT	MEDIAN	MEAN
1A	total population	7	30	38
2A	female	4	30	33.5
2B	male	3	30	44
2C	black or African American	4	51	48.5
2D	white	3	24	24
3A	single adults	[D]	[D]	[D]
3B	married adults	4	51	54
4A	black or African American female	3	36	36.7
4B	black or African American male	[D]	[D]	[D]
4C	white male	[D]	[D]	[D]
4D	white female	[D]	[D]	[D]
5A	persons under 5 years	[D]	[D]	[D]
5B	persons under 18 years	[D]	[D]	[D]
5C	persons 64 years or over	[D]	[D]	[D]

Note: Married persons must be 15 or over

Garfinkel, Aboud, Martindale 2018.

TABLE 2: POSSIBLE AGES FOR A MEDIAN OF 30 AND MEAN OF 44

A	B	C	A	B	C	A	B	C
1	30	101	11	30	91	21	30	81
2	30	100	12	30	90	22	30	80
3	30	99	13	30	89	23	30	79
4	30	98	14	30	88	24	30	78
5	30	97	15	30	87	25	30	77
6	30	96	16	30	86	26	30	76
7	30	95	17	30	85	27	30	75
8	30	94	18	30	84	28	30	74
9	30	93	19	30	83	29	30	73
10	30	92	20	30	82	30	30	72

# Reconstruction in the US Census. 2010.

Variable	Range
Block	6,207,027 inhabited blocks
Sex	2 (Female/Male)
Age	103 (0-99 single age year categories, 100-104, 105-109, 110+)
Race	63 allowable race combinations
Ethnicity	2 (Hispanic/Not)
Relationship	17 values

Publication	Released counts
PL94-171 Redistricting	2,771,998,263
Balance of Summary File 1	2,806,899,669
Total Statistics in PL94-171 and Balance of SF1:	5,578,897,932
Published Statistics/person	18
Recall: Collected variables/person:	6
<b>Published Statistics/collected variable</b>	<b>18 ÷ 6 ffi 3</b>

Survey ↑

5.5 billion simultaneous equations

on

1.8 billion unknown integers

## Reconstruction Formulation

Dataset  $X$

Statistics  $f_1, \dots, f_k$

answers

$$\begin{aligned} a_1 &\approx f_1(X) \\ a_2 &\approx f_2(X) \\ &\vdots \\ a_k &\approx f_k(X) \end{aligned} \quad \approx \text{"approx"}$$

Reconstruction Problem: Given "constraints"  $\{f_i(X) \approx a_i\}$ ,  
find a dataset  $\tilde{X}$  that is consistent w/ the constants.

## Linear Reconstruction Attack

- Introduced by Dinur & Nissim in 2003 development of Differential Privacy. 06

Data Set  
X

Name	Postal Code	Age	Sex	Has Disease?
Alice	02445	36	F	1
Bob	02446	18	M	0
Charlie	02118	66	M	1
:	:	:	:	:
Zora	02120	40	F	1

$Z = \text{identifiers}$        $S = \text{Secret bit}$

Release count statistics: # people satisfy some property

- How many people are older than 40 & have secret bit = 1?

$$f(X) = \sum_{j=1}^n \varphi(z_j) s_j \quad \text{for some } \varphi: Z \mapsto \{0,1\}$$

$$f(X) = \underbrace{(\varphi(z_1), \varphi(z_2), \dots, \varphi(z_n))}_{\text{"property" bit vector}} \cdot \begin{matrix} \uparrow \\ \text{dot product} \\ \text{inner product} \end{matrix} \underbrace{(s_1, \dots, s_n)}_{\substack{\text{secret bit} \\ \text{vector.}}}$$

Identifiers	Secret
$z_1$	$s_1$
$z_2$	$s_2$
$z_3$	$s_3$
:	:
$z_n$	$s_n$

← Format.

Releasing  $k$  linear Statistics

Released Statistics  $\begin{bmatrix} f_1(X) \\ \vdots \\ f_k(X) \end{bmatrix} = \begin{bmatrix} \varphi_1(z_1) & \cdots & \varphi_1(z_n) \\ \vdots & F_i & \vdots \\ \varphi_k(z_1) & \cdots & \varphi_k(z_n) \end{bmatrix} \begin{bmatrix} s_1 \\ \vdots \\ s_n \end{bmatrix}$  ← Secret bits

$$f_i(X) = F_i \cdot s$$

Examples :

$\varphi_1(z_j) = 1$  :  $z_j$  is older than 40

$\varphi_2(z_j) = 1$  :  $z_j$  is older than 40 and male

$\varphi_3(z_j) = 1$  :  $z_j$  is older than 20 and male

## First Reconstruction Attack

" You can't release all count statistics with non-trivial accuracy."

Queries :  $k=2^n$

$n = \text{number of people}$

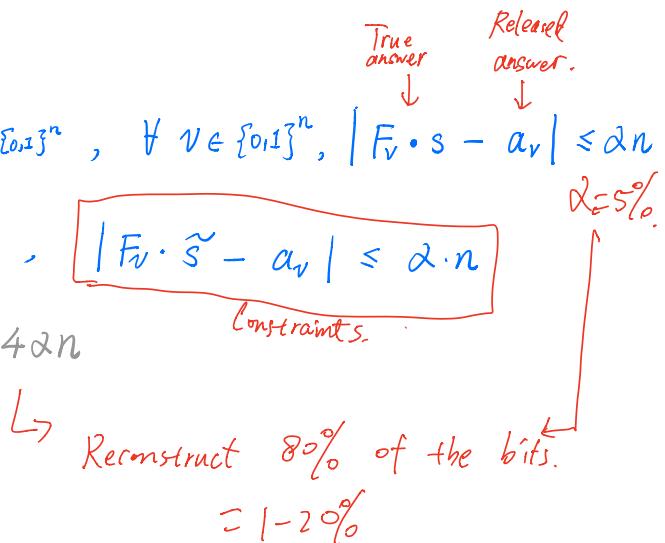
For every  $v \in \{0,1\}^n$ ,  $F_v = v$

Reconstruction :

Suppose the answers  $(a_v)_{v \in \{0,1\}^n}$ ,  $\forall v \in \{0,1\}^n$ ,  $|F_v \cdot s - a_v| \leq \alpha n$

Choose  $\tilde{s} \in \{0,1\}^n$ ,  $\forall v$ ,  $|F_v \cdot \tilde{s} - a_v| \leq \alpha \cdot n$

Theorem.  $\|s - \tilde{s}\|_1 \leq 4\alpha n$



**Theorem.** If all  $2^n$  counts are within  $\alpha n$  error, then  $s, \tilde{s}$  disagree on  $\leq 4\alpha n$  bits.

## Proof Intuition.

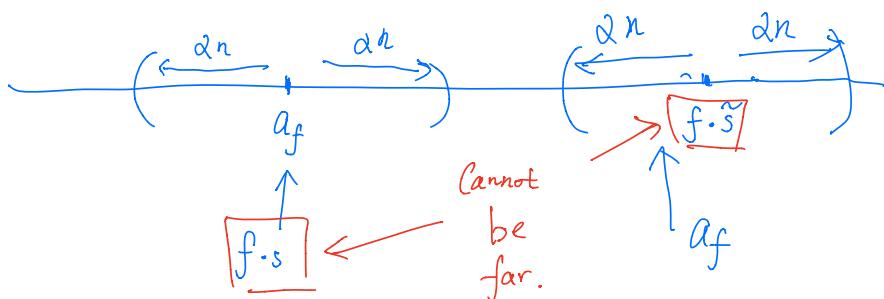
$$S = [1011 \quad \underline{\hspace{2cm}}]$$

$$\mathfrak{Z} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ \underbrace{\phantom{0100}}_{\text{4 columns}} & \xrightarrow{\text{4 rows}} \end{bmatrix}$$

Property  $\varphi_j$  that captures the diff.

Statistic  $f$

- | Assumption
- |  $Af = \text{Releaved answer.}$
- |  $f \cdot s = \text{true answer.}$
- |  $|Af - f \cdot s| \leq \alpha n$
- | Reconstruction:
- | Find  $\tilde{s}$  such that
- |  $|Af - f \cdot \tilde{s}| \leq \alpha n$



Theorem. If all  $2^n$  counts are within  $\alpha n$  error,  
then  $s, \tilde{s}$  disagree on  $\leq 4\alpha n$  bits.

Proof Sketch.

Two sets:  
 $S_{01} = \{j : s_j=0 \text{ & } \tilde{s}_j=1\}$   
 $S_{10} = \{j : s_j=1 \text{ & } \tilde{s}_j=0\}$

Proof by Contradiction

If  $\|s - \tilde{s}\|_1 > 4\alpha n$   $\xrightarrow{\ell_1 \text{ norm}} \sum_j |s_j - \tilde{s}_j|$

 $\Rightarrow |S_{01}| > 2\alpha n \text{ or } |S_{10}| > 2\alpha n = \sum_{j \in S_{01}} \underbrace{|s_j - \tilde{s}_j|}_{1} + \sum_{j \in S_{10}} \underbrace{|s_j - \tilde{s}_j|}_{1}$ 
 $\Rightarrow \text{Then there exists } v \in \{0,1\}^n \text{ such that } |v \cdot (s - \tilde{s})| > 2\alpha n$ 
 $\Rightarrow |v \cdot \tilde{s} - v \cdot v| > 2\alpha n - |v \cdot s - v \cdot v| > \alpha n$ 

Triangle Inequality

 $\Rightarrow \text{Contradiction} \quad \left( \text{Since } |v \cdot \tilde{s} - v \cdot v| \leq \alpha n \text{ in our reconstruction} \right)$

No Class 9/6.

No Recitation this Friday

Reading for next Weds.

Reconstruction Using Fewer Queries  
 # Released Statistics <<  $2^n$  ?

Attack : Choose  $k = 20n$  random  $\varphi_i : \mathbb{Z} \mapsto \{0,1\}$ ,  $\forall i \in [k]$ .  
 $\Rightarrow k$  random vectors/queries  $F_i \in \{0,1\}^n$   
 Suppose that answers :  $\forall i \in [k]$ ,  $|F_i \cdot s - a_i| \leq \alpha n$   
 Find  $\tilde{s} \in \{0,1\}^n$  such that:  $\forall i \in [k]$ ,  $|F_i \cdot \tilde{s} - a_i| \leq \alpha n$

Theorem.  $\|s - \tilde{s}\|_1 \leq 256 \alpha^2 n^2$