

Lecture 8. DP and Mechanism Design. (Reverse game theory)

- Revisit Digital Goods Example
- DP as a tool for truthfulness.
- (If have time): continual Release ↴

Logistics: HW Posted on Canvas
Due Sunday.

Example 2: Pricing a digital good.

- Selling an app; what price?
- n buyers report their valuations: "How much are they willing
bidders to pay?"
 x_i

Revenue:

$$g(p; x) = p \cdot \#\{i : x_i \geq p\}$$

Toy Example: \$1, \$1, \$2.01, \$4.
 $\hookrightarrow \text{Rev} = 4.02$

Utility Theoretic View.

Each buyer may have a utility function $f: Y \rightarrow \mathbb{R}_{\geq 0}$

Notation:
 (from game theory) $(x_{-i}, x'_i) = (x_1, x_2, \dots, x_{i-1}, x'_i, x_{i+1}, \dots, x_n)$
 ↑
 Unilateral deviation.
 ≈ "Neighboring datasets"

Proposition (Utility Theoretic View of DP)

An algorithm $A: X^n \rightarrow Y$ is ϵ -DP.
 if and only if for every utility function $f: Y \rightarrow \mathbb{R}_+$
 and every $x \in X^n$, $x'_i \in X$:

$$\mathbb{E}_{Y \sim A(x)} [f(Y)] \geq \exp(\epsilon) \mathbb{E}_{A' \sim A(x_{-i}, x'_i)} [f(Y')]$$

↓
 utility for
 bidder i

↓
 Expected utility
 for truthtelling

↓
 Expected utility
 for misreporting

Proposition (Utility Theoretic View of DP)

An algorithm $A: X^n \rightarrow Y$ is ε -DP.

if and only if for every utility function $f: Y \rightarrow \mathbb{R}_+$ and every $x \in X^n$, $x_i' \in X$:

$$\mathbb{E}_{Y \sim A(x)} [f(Y)] \geq \exp(\varepsilon) \mathbb{E}_{A \sim A(x_i, x_i')} [f(Y)] \quad (*)$$

Proof. $DP \Rightarrow (*)$

$$\begin{aligned} \mathbb{E}_{Y \sim A(x)} [f(Y)] &= \int_y f(y) \cdot \underbrace{\mathbb{P}[A(x) = y]}_{\text{def}} dy \\ &\geq \int_y f(y) \exp(-\varepsilon) \cdot \mathbb{P}[A(x) = y] dy \\ &= e^{-\varepsilon} \int_y f(y) \mathbb{P}[A(x) = y] dy \\ &= \mathbb{E}_{Y \sim A(x_i, x_i')} [f(Y)] \cdot e^{-\varepsilon}. \end{aligned}$$

$(*) \Rightarrow DP$.

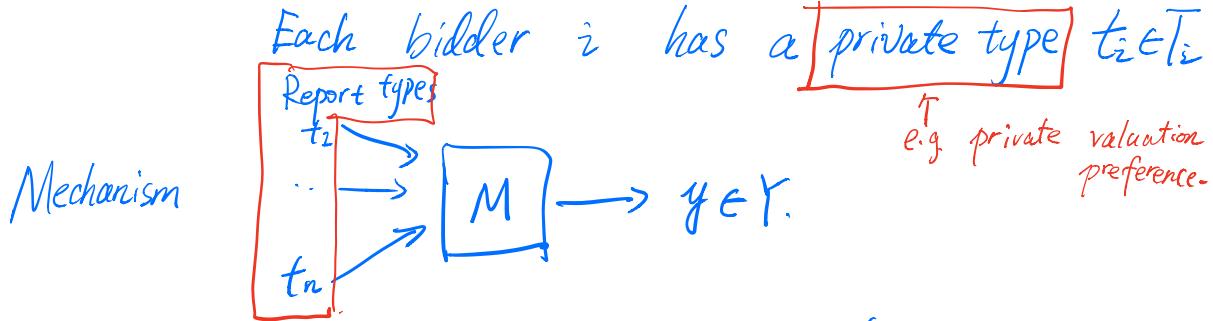
$$\text{Fix } E \subseteq Y, \quad f(y) = \mathbb{1}[y \in E]$$

$$\mathbb{P}[A(x) \in E] \geq \exp(-\varepsilon) \mathbb{P}[A(x') \in E].$$

□

Mechanism Design Basics.

n bidders indexed by i
 (agents in general)



Utility function: $U_i(y) \equiv u(t_i, y)$

for agent i

$u: T \times Y \rightarrow [0,1]$

Definition (Dominant Strategy Truthfulness)

A mechanism $M: T^n \rightarrow Y$ is ϵ -approximately

Dominant Strategy truthful if for all

$t \in T^n$ and for all i , $t'_i \in T$

$$U_i(M(t)) \geq U_i(M(t_{-i}, t'_i)) - \epsilon$$

t_i depends on t_{-i}

ε -DP $\implies \varepsilon$ -approx Dominant Strategy Truthful.

Claim. If M is ε -DP, then
 M is also ε -approx dominant strategy truthful.

Proof.

$$\mathbb{E}[u_i(Y)] \geq \exp(-\varepsilon) \cdot \mathbb{E}[u_i(Y)] \geq (1-\varepsilon) \cdot \mathbb{E}[u_i(Y)]$$

$\uparrow \quad \uparrow \quad \uparrow$
from claim
before $e^{-\varepsilon} \geq (1-\varepsilon)$

$Y \sim M(t) \quad Y \sim M(t_i, t_{i'}) \quad Y \sim M(t, t_i, t_{i'})$

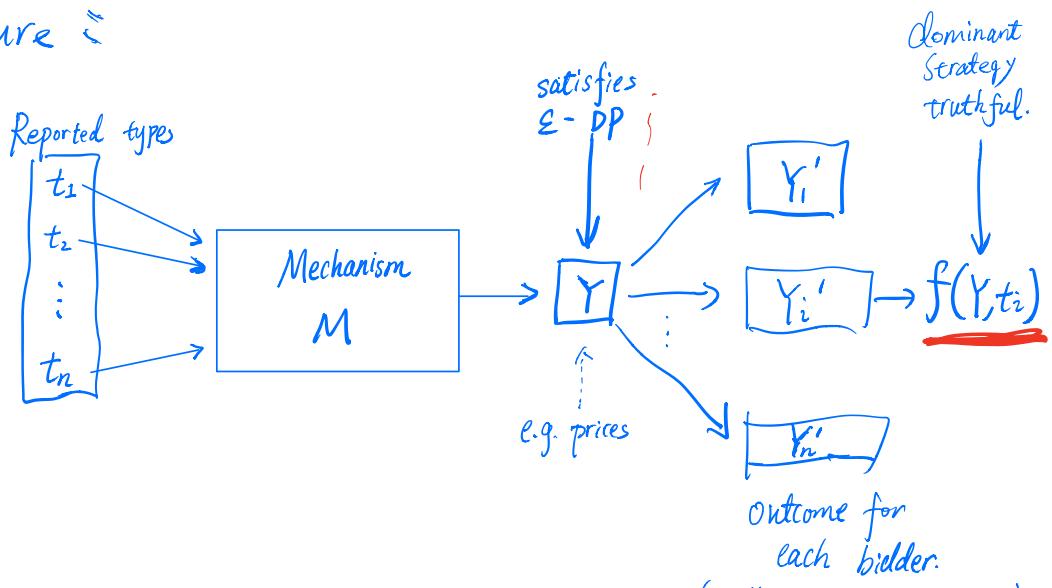
$$u_i \in [0,1] \longrightarrow \geq \mathbb{E}[u_i(Y)] - \varepsilon$$

$Y \sim M(t, t_i, t_{i'}) \quad \square$

Observations :

- ① Group privacy \Rightarrow collusion proof.
(k people collude $\rightarrow k\varepsilon$ -DST)
- ② Every strategy is approx dominant strategy.
(not just being truthful)

Structure :



Claim. If $M: T^n \rightarrow Y$ is ϵ -DP
and $f: Y \times T \rightarrow Y'_i$ is dominant strategy truthful. $\forall y \in Y$

Then truthful reporting is ϵ -approx dominant strategy truthful.
for the mechanism $M': T^n \rightarrow Y'$ that returns
 $f(M(t), t_i)$ for agent i .

Pricing a digital good.

private type $t_i \in [0,1]$ (re-scaled)

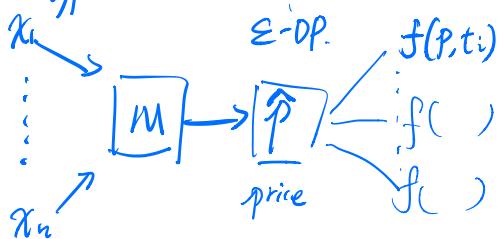
Outcome space $\mathcal{Y} = 2^{[n]} \times [0,1]$
 "who gets the item" \uparrow price. \downarrow

For any (S, p)

$$U_i(S, p) = \mathbb{1}[i \in S] \cdot (t_i - p)$$

$$\equiv \begin{cases} t_i - p & \text{if } i \in S \\ 0 & \text{if } i \notin S. \end{cases} \quad \leftarrow$$

Reported types



① M: Run Exp mech. (Report noisy max)
 w/ revenue as the quality score.

$$Y_2 = \{0, 2\alpha, \dots, 1\}$$

$$f = \#\{x_i \geq p\} \cdot p.$$

$$② f(\hat{p}, x_i) \equiv \mathbb{1}[x_i \geq \hat{p}]$$

$i \in S$ if and only if
 reported $x_i \geq \hat{p}$.

$\Rightarrow \varepsilon$ -approx DS truthful.