Make FunBlocks alive

Marvin FOURASTIE

Master project

Motivations

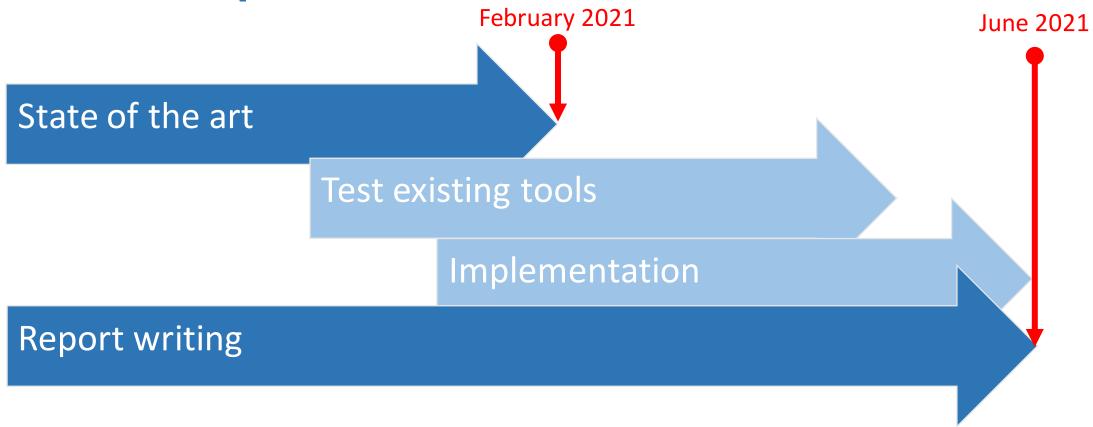






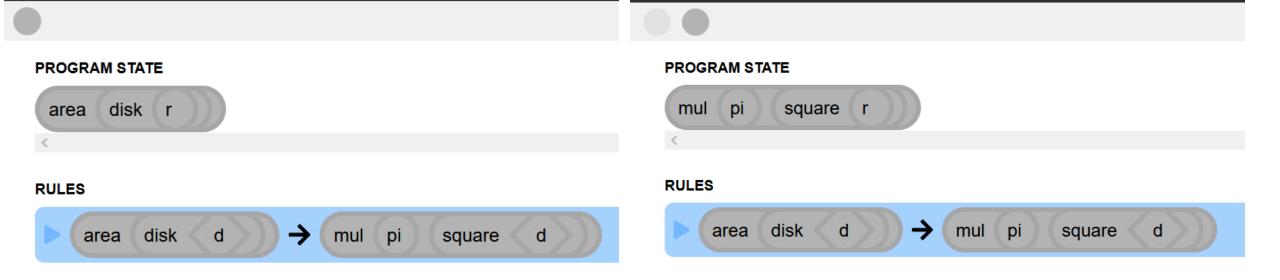
FunBlocks

Road Map



FunBlocks

```
init area(disk(r))
case area(disk($\$d$)) => mul(pi,square($\$d$))
```



FunBlocks

area(disk(r)) Declarative paradigm case area(disk(\$d)) => mul(pi,square(\$d)) Based on rewrite systems square Static typing node (Tree \$t) (Tree \$t) type Tree *\$t* :: empty | leaf *\$t* |

Goals

Provide users with valuable insights about their program

Verification of rewrite systems

Rewrite systems

Stack operators

 $Zero = \{0\}$

Nat = Zero U succ(Nat)

Empty = Λ

Stack = Empty U push(Nat, Stack)

top : Stack \rightarrow Nat

pop : Stack → Stack

alternate : Stack \times Stack \rightarrow Stack

Rewrite systems

Canonical rewrite system

$$top(push(x, y)) = x$$

$$pop(push(x, y)) = y$$

$$pop(push(x, y)) \rightarrow y$$

$$alternate(\Lambda, z) = z$$

$$alternate(push(x, y), z) = push(x, alternate(z, y))$$

$$alternate(push(x, y), z) \rightarrow push(x, alternate(z, y))$$

Rewrite systems

Termination

Confluence

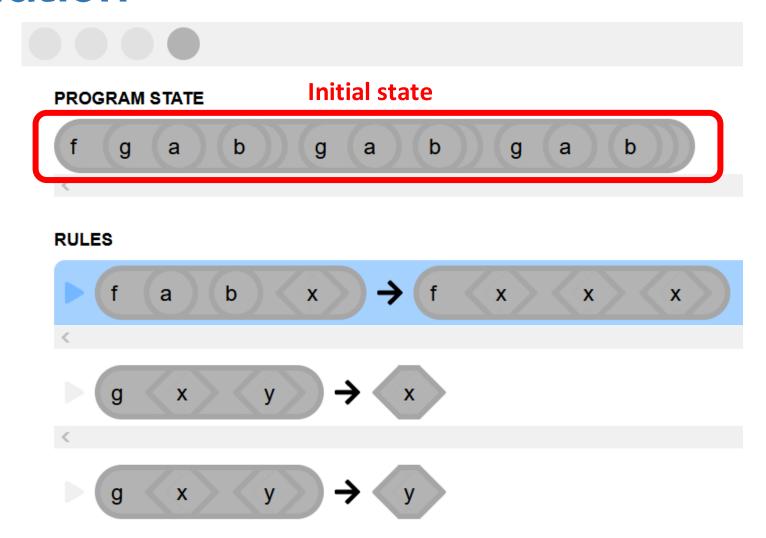
Soundness

Completeness

Correctness

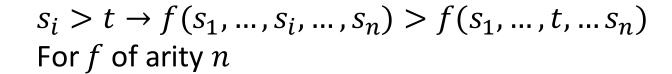
Undecidable in general

Termination



Reduction order

Monotone



Close under substitution ———

$$s>t
ightarrow \sigma s>\sigma t$$
 , for all substitution σ

Well-founded

no infinite descending chain

 $(\mathbb{N}, <)$ is well-founded

 $(\mathbb{Z}, <)$ is not well-founded

Termination

A term rewriting system is terminating

if and only if

it admits a compatible reduction order < (if l > r for every rewrite rule $l \rightarrow r$)



Verification of termination

Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(a) = 1$$

$$w(g(t)) = 1 + w(t)$$

$$w(f(t_1,t_2)) = 2w(t_1) + w(t_2)$$

Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(f(a,x)) = 2 + w(x)$$

$$w(f(g(x),y)) = 2 + 2w(x) + w(y)$$

$$w(x) = w(x)$$

$$w(g(f(x,y)) = 1 + 2w(x) + w(y)$$

$$w(f(a,x)) > w(x)$$

$$w(f(g(x),y)) > w(g(f(x,y))$$

Reduction order → Termination

Algorithms

Recursive Path Ordering

Order based on the mutisets

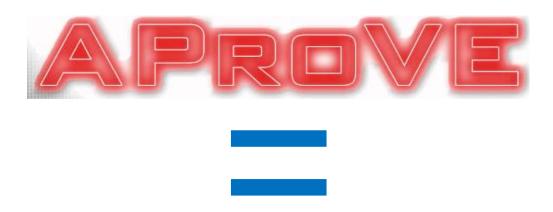
Knuth-Bendix Ordering

Based on weights assigned to operators

Dependency pairs

Prove innermost termination

Termination

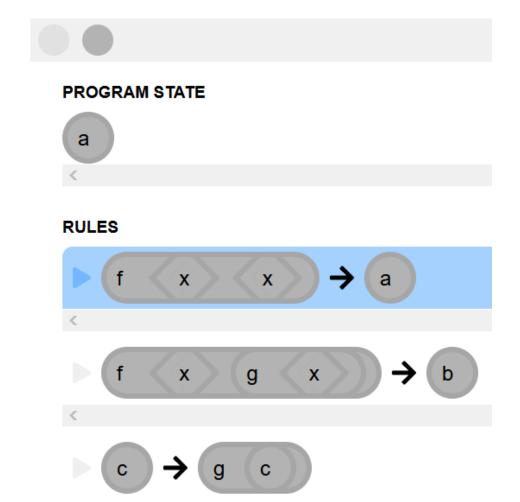


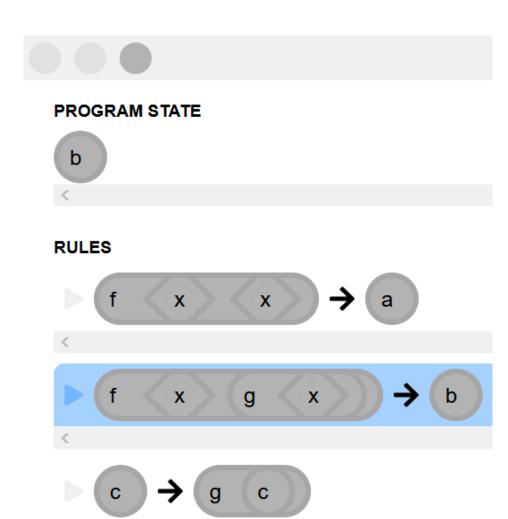
Direct proof (polynomial, LBO, KBO,...)



Dependency pairs and size-change principle

Confluence

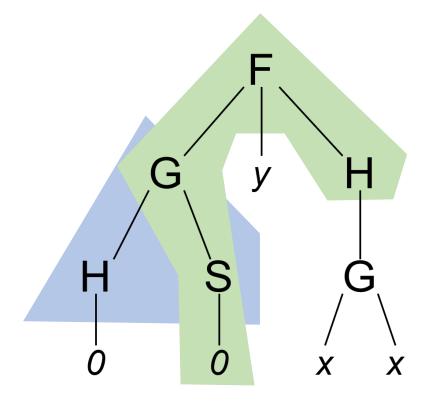




Overlap and critical pairs

$$\rho_1 : F(G(x,S(0)),y,H(z)) \to x$$

$$\rho_2 : G(H(x),S(y)) \to y$$



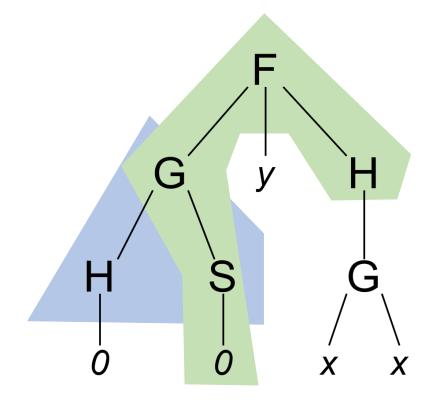
Overlap and critical pairs

Overlapping:

Term:
$$F(G(H(0),S(0)),y,H(z))$$

$$F(G(\square,S(0)),\square,H(\square))$$

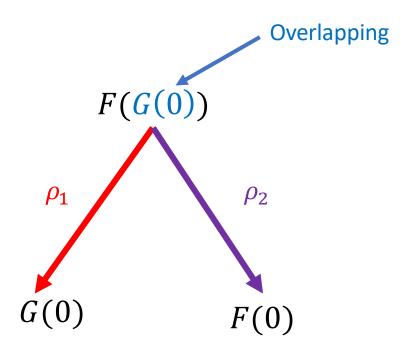
$$G(H(\square),S(\square))$$



Overlap and critical pairs

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



< G(x), F(x) > is called critical pair

Critical Pair Lemma

A terminating rewriting system is confluent

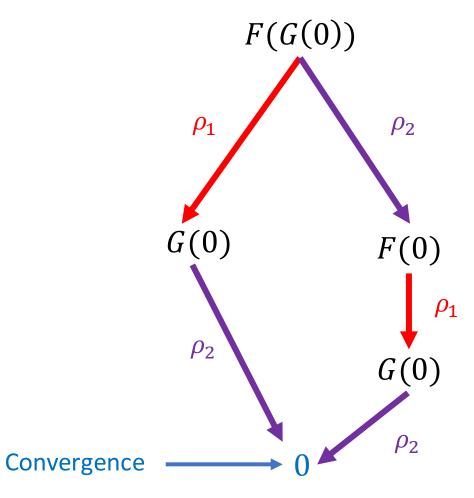
if and only if

all critical pairs are convergent

Critical Pair Lemma

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



Input:

A set of equation

A reduction ordering <

$$\begin{array}{rcl}
1 \cdot x & = x \\
x^{-1} \cdot x & = 1 \\
(x \cdot y) \cdot z & = x \cdot (y \cdot z)
\end{array}$$

Non-confluent

Output:

Terminate successfully

Terminating and confluent rewrite system

Non-terminating rewrite system

Rule which cannot be ordered (i.e. commutative operator)

Basic rules:

Orienting _____

Transform s = t to $s \rightarrow t$

Adding

 \longrightarrow

Add s = t in the set of equation

Simplifying

 \longrightarrow

Simplify s = t in s' = t'

Deleting



Delete trivial rules s = s

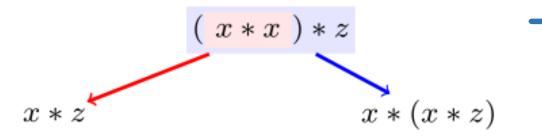
Adding



Add s = t in the set of equation

$$(x*y)*z \rightarrow x*(y*z)$$

 $x*x \rightarrow x$



No convergent

$$\begin{array}{ccc} (x*y)*z & \rightarrow x*(y*z) \\ x*x & \rightarrow x \\ \hline x*(x*z) & \rightarrow x*z \end{array}$$

New rule added

Completion process:

- 1. For each equation s = t reduce s and t to normal form s' and t'
- 2. Fill the set of rules using basic operators and reduction ordering
- 3. If the algorithm terminate successfully: terminating and confluent rewrite system

Completion for axioms of groups:

$$\begin{array}{rcl}
1 \cdot x & = x \\
x^{-1} \cdot x & = 1 \\
(x \cdot y) \cdot z & = x \cdot (y \cdot z)
\end{array}$$



$$\begin{array}{cccc}
1 \cdot x & \rightarrow x \\
x^{-1} \cdot x & \rightarrow 1 \\
(x \cdot y) \cdot z & \rightarrow x \cdot (y \cdot z) \\
x^{-1} \cdot (x \cdot y) & \rightarrow y \\
1^{-1} & \rightarrow 1 \\
x \cdot 1 & \rightarrow x \\
(x^{-1})^{-1} & \rightarrow x \\
x \cdot x^{-1} & \rightarrow 1 \\
x \cdot (x^{-1} \cdot y) & \rightarrow y \\
(x \cdot y)^{-1} & \rightarrow y^{-1}x^{-1}
\end{array}$$

Educational tools



CSI

KBCV

Termination

Confluence

Completion

TRS tool

Parcourir... Aucun fichier sélectionné. Upload $(VAR \times y)$ (RULES $f(x,y) \rightarrow x$ $f(x,y) \rightarrow f(x,g(y))$ $g(x) \rightarrow h(x)$ $F(g(x),x) \rightarrow F(x,g(x))$ $F(h(x),x) \rightarrow F(x,h(x))$ (COMMENT Example 6 of \cite{AT97}) (COMMENT %% TagRevision: 1 %%) (COMMENT %% Tags: [4ec3f85c01836]non left linear{}; [4ec3f87f0f1e0]r 50 × Rewrites Limit (Use with caution) Go!

TRS tool

	R ₀	$= f(x,y) \rightarrow x$
R ₀	is	Left-Linear
Ro	is	Right-Linear
Ro	is	Linear
R ₀	is	Collapsing
Ro	is	not Duplicating
R ₀	is	not Conservative
Ro	is	Destructive

TRS								
The	TRS	is	not Left-Linear					
The	TRS	is	not Right-Linear					
The	TRS	is	not Linear					
The	TRS	is	Collapsing					
The	TRS	is	not Duplicating					
The	TRS	is	not Conservative					
The	TRS	is	Destructive					
The	TRS	is	not Orthogonal					
The	TRS	is	not Almost Orthogonal					
The	TRS	is	not Weakly Orthogonal					
The	TRS	is	Locally Confluent					
Unkı	nown	COI	nfluence for The TRS					
1			non terminating pop: $f(x,g(y)) \rightarrow f(x,g(y))$					

TTT2

Tyrolean Termination Tool 2 (1.20)

1. Input Term Rewrite System

For input use the standard TRS format.

2. Select Strategy

• FAST O	FBI C	HYDRA	O LPO	○ кво	OPOLY	\bigcirc MA	T(2)	MAT(3)	\bigcirc	COMP	\bigcirc	COMPLEXITY
O EXPERT												

3. Encode State into URL (optional)

encode URL clear URL

4. Start TTT2

check use HTML output if available (*experimental feature*)

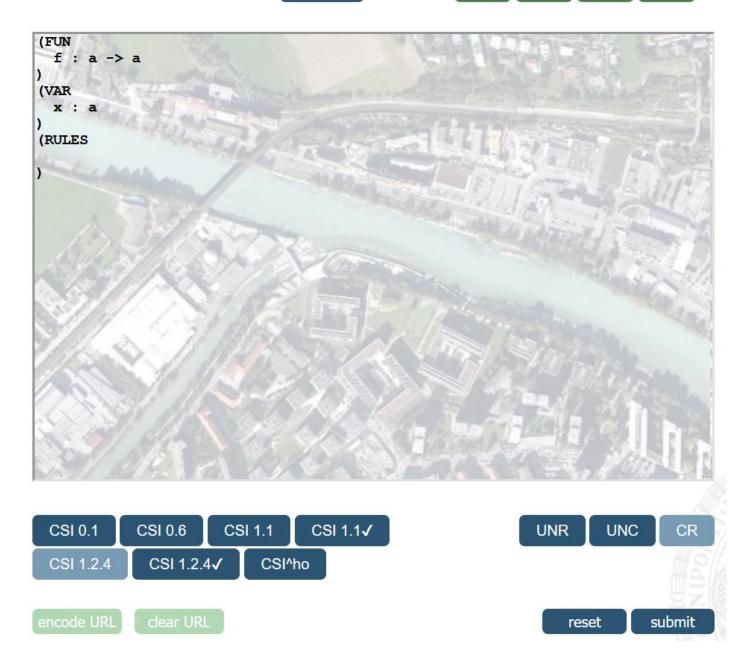
Enter a TRS or HRS or upload a file browse... Examples: trs1

trs2

hrs1

hrs2

CSI



KBCV

Equations	from[1	to[500	Rules	from 1	500
Orient → Orient ← Simplify Delete Compose Collapse Deduce Completion		to[500	Rules	from[1 to	500
Undo					
LPO Precedence		111			
		III			
Undo / Redo Stack 1. start : Welcome to the 'Knuth-Bendix Completic	an Mauglizor I				
2. add : equations f(f(x,y),z)=f(x,f(y,z)), f(x,c)=x,					
Status Messages		III			
Welcome to the 'Knuth-Bendix Completion Visualizer'					
equations $f(f(x,y),z)=f(x,f(y,z))$, $f(x,c)=x$, $f(x,g(x))=c$ we	ere added				
file '.kbcvinit' loaded!					

Performance tools

Termination

Confluence

Completion

Proof verification

MU-TERM

ACP

Maxcomp

CoLoR

NaTT

T | Saigawa





Hybrid tools







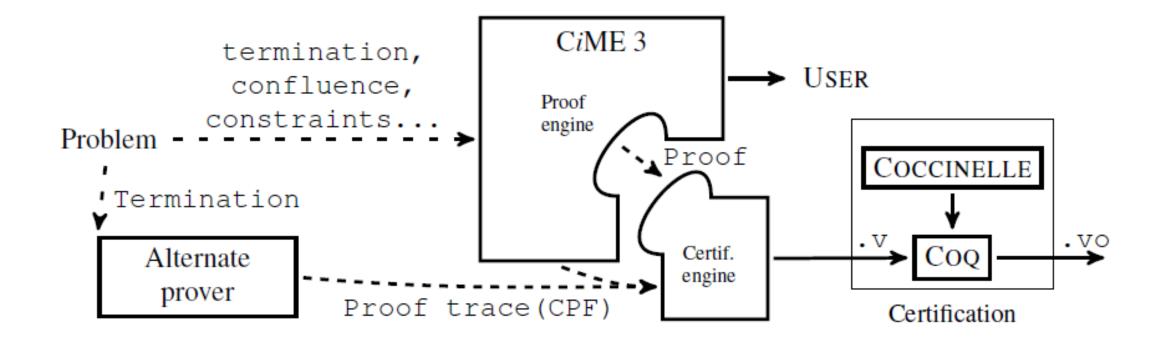


Rewriting toolkit





Proof certification



Examples of declarations

```
let X = variables "x,y";
let F = signature "plus : binary; 0:constant; S:unary;";
let T = algebra F;
let t1 = term T "S(0)";
let R = trs T "plus(0,x) -> x; plus(S x, y) -> S(plus(x,y));";
let c = order\_constraints T "0 < S(0) / S(plus(x,y)) < plus(S(x),y)";
```

Definition of signatures

```
CiME>let X = variables "x,y,z";
X : variable_set = variables "z,x,y"
```

Definition of algebra and terms

```
CiME> let A_peano = algebra F_peano ;
A_peano : F_peano algebra = algebra F_peano
```

```
CiME> let t = term A_peano "s(s(s(0)))*(s(0)+s(s(0)))";
t : F_peano term = s(s(s(0))) *(s(0)+s(s(0)))
```

Term rewriting system

```
CiME> let R_peano = trs A_peano "
      x+0 -> x;
      X+S(Y) \longrightarrow S(X+Y);
      x*0 -> 0;
      x*s(y) -> (x*y) +x;
       ";
 R_peano : F_peano trs = trs A_peano "
            x+0 -> x;
            x+s(y) \rightarrow s(x+y);
            x * 0 -> 0;
             x * s (y) -> (x * y) + x "
```

```
CiME> termination R_peano;
CiME> coq_certify_proof R_peano;
CiME> convergence R_peano ;
...
```





Simplicity



Expressiveness



Performance

```
fmod BASIC-NAT is
        sort Nat .
       op 0 : -> Nat .
       op s : Nat -> Nat .
       op + : Nat Nat -> Nat .
       vars N M : Nat .
       eq 0 + N = N.
        eq s(M) + N = s(M + N).
endfm
```

```
fmod FACTORIAL is
protecting NAT .

op _! : Nat -> NzNat .
var N : Nat .
eq 0 ! = 1 .
eq (s N) ! = (s N) * N ! .

endfm

fmod FACTORIAL is
protecting NAT .

op _! : Nat -> NzNat .
var N : Nat .
eq 0 ! = 1 .
eq (s N) ! = (s N) * N ! .
```

```
> load factorial.maude
> red 100 ! .
Reduce in FACTORIAL : 100 ! .
rewrites: 201 in Oms cpu (Oms real) (~ rewrites/second)
result NzNAT:
9332621544394415268169923885626670049071596826438162146
8592963895217599993229915608941463976156518286253697920
8272237582511852109168640000000000000000000000
```

```
mod VENDING-MACHINE is
 including VENDING-MACHINE-SIGNATURE .
 var M : Marking .
  rl [add-q] : M => M q.
  rl [add-\$] : M => M \$ .
  rl [buy-c] : $ => c .
  rl [buy-a] : $ => a q .
  rl [change] : q q q => $.
endm
```

Inductive Theorem Prover (ITP)

Sufficient Completeness Checker (SCC)

Church-Rosser Checker (CRC)

Coherence Checker (ChC)

Maude Termination Tool (MTT)

Maude Formal Environment (MFE)

Tools overview

	Maude	CiME
Extensibility	+	≈
Still active	≈	
I/O files	+	-
Syntax	+	-
Documentation		



Standard input/output



File handling



Sockets

```
mod VENDING-MACHINE is
 including VENDING-MACHINE-SIGNATURE .
 var M : Marking .
  rl [add-q] : M => M q.
  rl [add-\$] : M => M \$ .
  rl [buy-c] : $ => c .
  rl [buy-a] : $ => a q .
  rl [change] : q q q => $.
endm
```

```
fmod VENDING-MACHINE-SIGNATURE is
sorts Coin Item Marking .
subsorts Coin Item < Marking .
op __ : Marking Marking -> Marking [assoc comm id: null] .
op null : -> Marking .
op $ : -> Coin [format (r! o)] .
op q : -> Coin [format (r! o)] .
op a : -> Item [format (b! o)] .
op c : -> Item [format (b! o)] .
endfm
```

```
load vending-machine-signature.maude
fmod VENDING-MACHINE-GRAMMAR is
  protecting VENDING-MACHINE-SIGNATURE .
  protecting NAT .
  sort Action .
  op insert $ : -> Action .
 op insert q : -> Action .
  op show basket : -> Action .
  op show credit : -> Action .
  op buy__(s) : Nat Item -> Action .
endfm
```

```
load vending-machine-grammar.maude
load buying-strats.maude
load file.maude
mod VENDING-MACHINE-IO is
  rl < O : X | action : insert $, marking : M, Atts >
  => < 0 : X | action : idle,
       marking : downTerm(insertCoin('add-$, upTerm(M)), null), Atts >
    write(stdout, O, "one dollar introduced\n") .
```

Parse input?

META-LEVEL module



Parse input

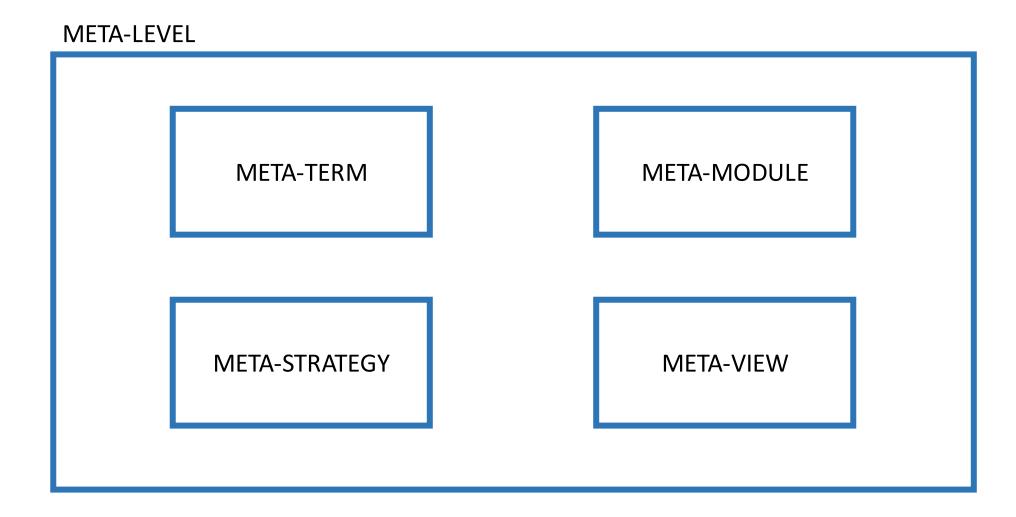


Execute



Print

META-LEVEL module



META-LEVEL operations

Metalanguage

Funblocks

```
init area(disk(r))

case area(disk($d)) => mul(pi, square($d))
```

```
sorts Const Var Func State Rule Term Expr .
```

Metalanguage

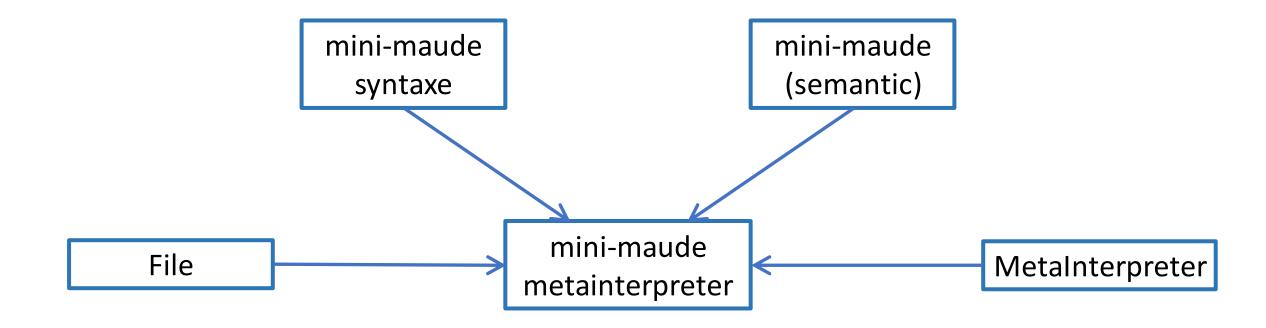
case
$$f(\$x) => a$$

Tokens are represented by quoted identifiers (QID)

Metalanguage

```
> result: {'case_=>_[bubble[...]],bubble['a.Qid]],'Rule}
```

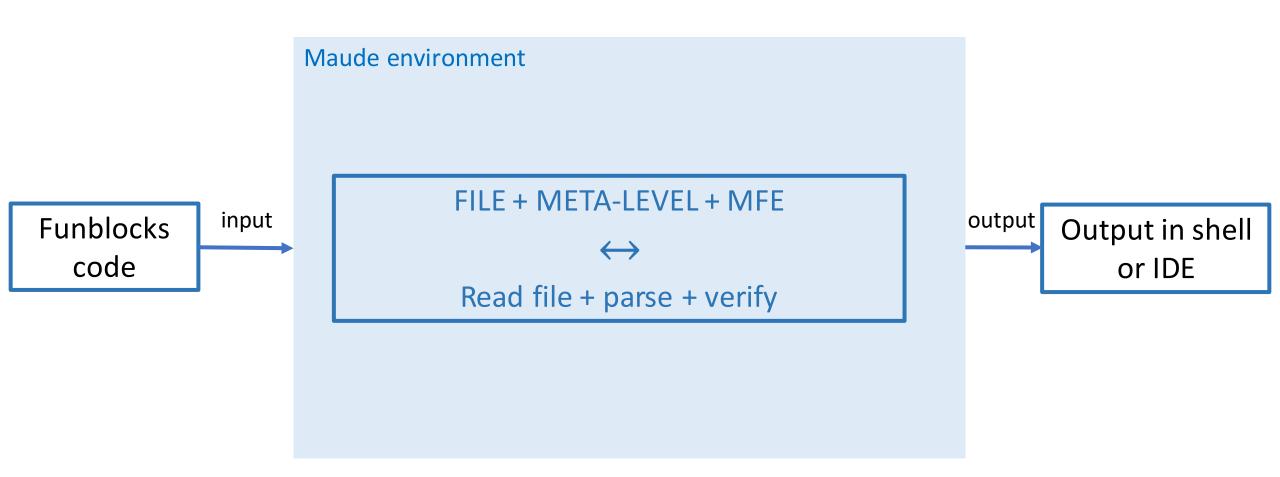
MINI-MAUDE

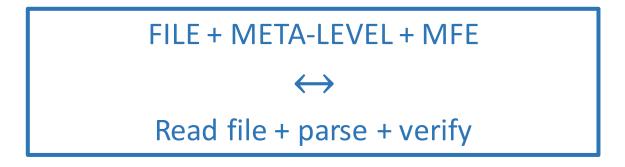


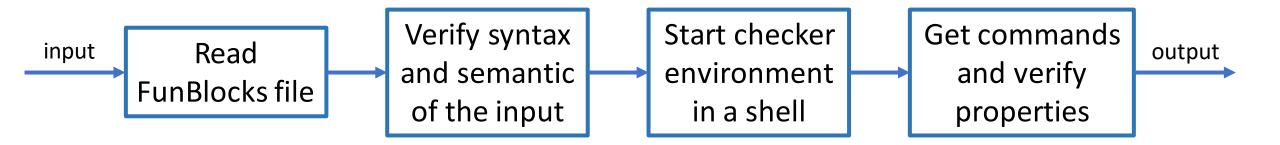
MINI-MAUDE

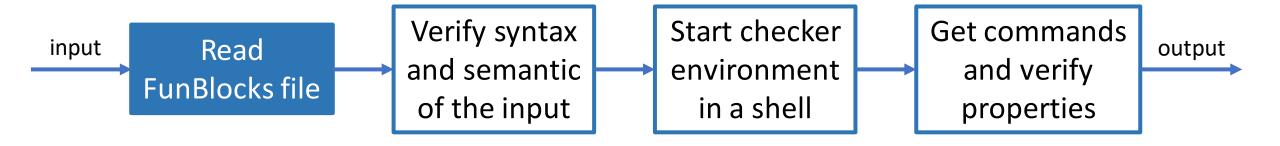
```
Maude> erew minimaude .
erewrite in MINI-MAUDE-META-INTERPRETER: minimaude.
MiniMaude Execution Environment
minimaude> fmod NAT3 is
      > sort Nat3.
      > op s_ : Nat3 -> Nat3 .
      > op 0 : -> Nat3 .
      > eq s s s 0 = 0.
      > endfm
Module loaded successfully
minimaude> reduce s s s s 0 .
result Nat3: s 0
```

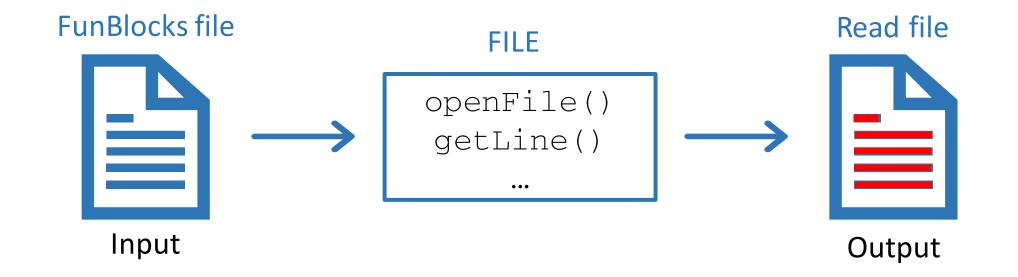
System diagram

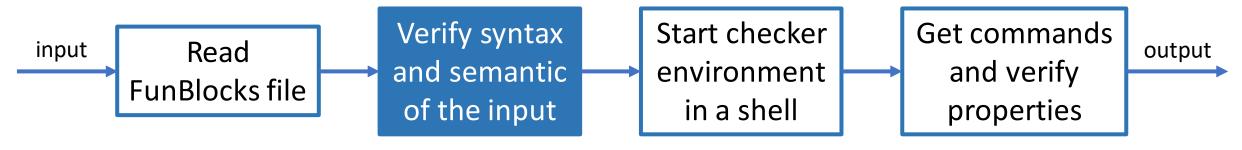


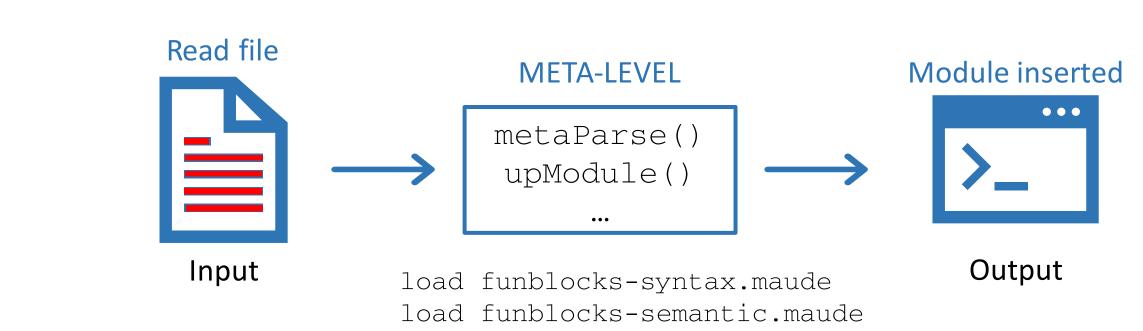


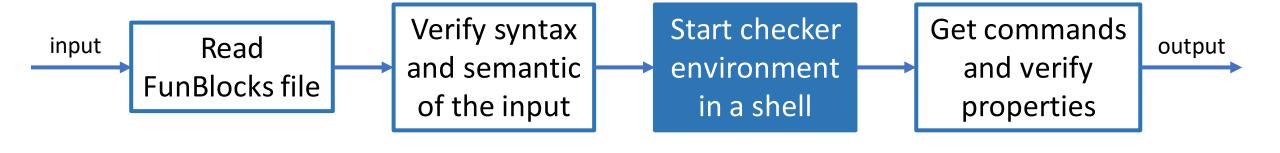


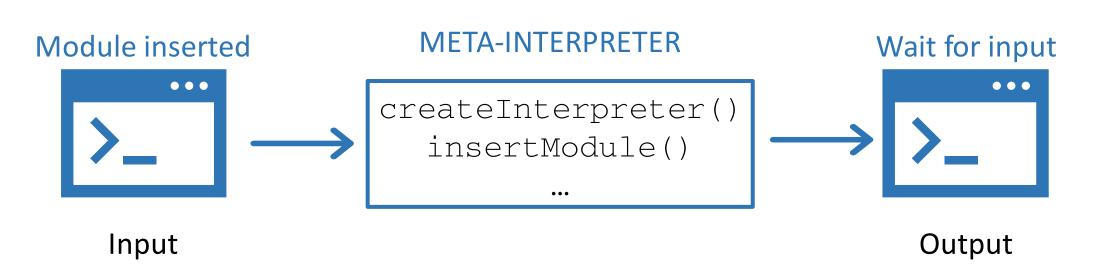


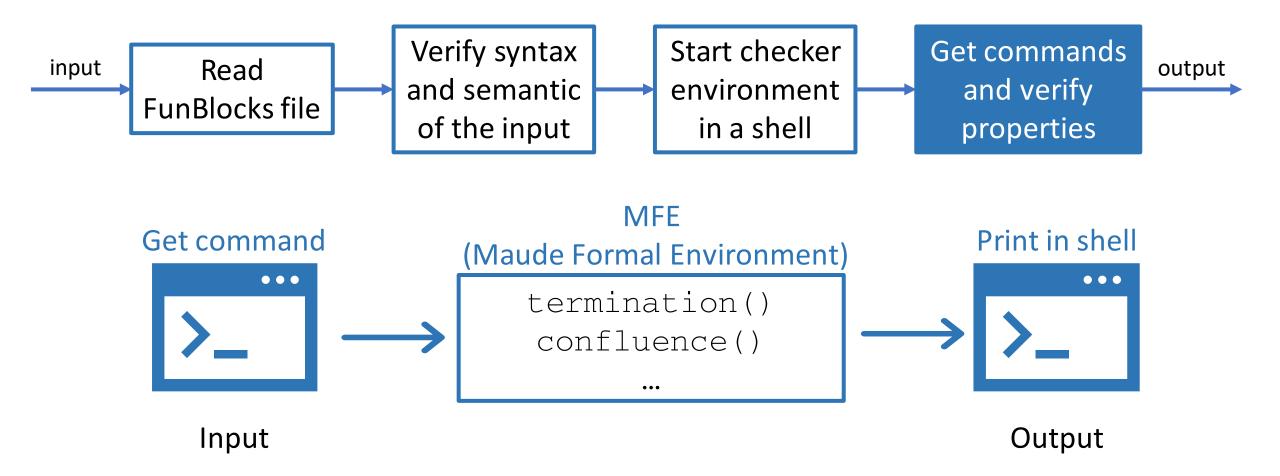




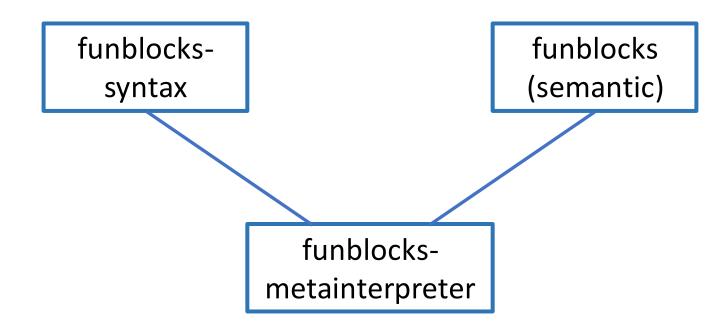








FunBlocks checker system



funblocksmetainterpreter

```
rl < 0 : FunBlocks | mi: null, st: 0, Atts >
    wrote(0, 0')
     createdInterpreter(0, Y, MI)
 => < 0 : FunBlocks | mi: MI, st: 1, Atts >
     insertModule(MI, 0, upModule('FUNBLOCKS-SYNTAX, true)) .
 rl < 0 : FunBlocks | mi: MI, st: 1, Atts >
     insertedModule(0, 0')
 => < 0 : FunBlocks | mi: MI, st: 2, Atts >
     getLine(stdin, 0, "funblocks> ") .
```

funblocksmetainterpreter

```
rl < 0 : FunBlocks | mi: null, st: 0, Atts >
     wrote(0, 0')
     createdInterpreter(0, Y, MI)
 => < 0 : FunBlocks | mi: MI, st: 1, Atts >
    insertModule(MI, 0, upModule('FUNBLOCKS-SYNTAX, true)) .
 rl < 0 : FunBlocks | mi: MI, st: 1, Atts >
    insertedModule(0, 0')
  => < 0 : FunBlocks | mi: MI, st: 2, Atts >
     getLine(stdin, 0, "funblocks> ") .
```

Mini-Maude: insert module

```
marvinf@marvinf-X405UA:~/Desktop/master_project/funblocks-checker/maude-3.1 🦲 🕜 😵
File Edit View Search Terminal Help
minimaude> fmod NAT3 is
     sort Nat3 .
     op s_ : Nat3 -> Nat3 .
     op 0 : -> Nat3 .
     eq s s s 0 = 0.
> endfm
Advisory: redefining module NAT3.
Module loaded successfully
minimaude>
```

Mini-Maude: use « reduce » command

```
marvinf@marvinf-X405UA:~/Desktop/master_project/funblocks-checker/maude-3.1 🦲 🕜 😵
File Edit View Search Terminal Help
minimaude> reduce s s s s 0 .
result Nat3: s 0
minimaude>
```

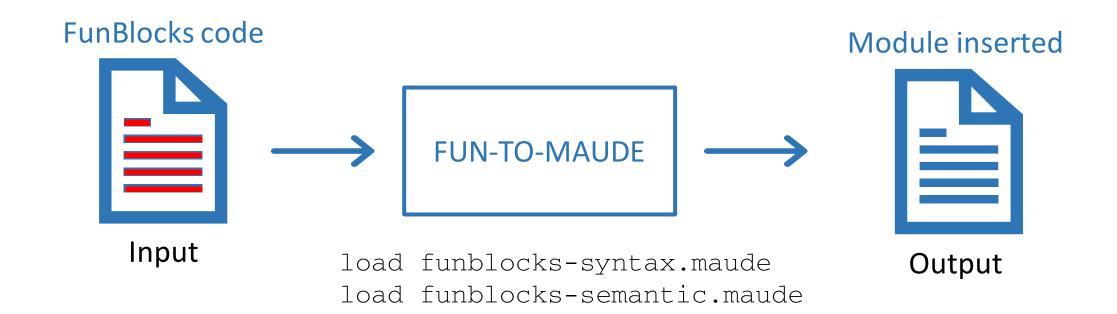
Mini-Maude: wrong command

```
marvinf@marvinf-X405UA:~/Desktop/master_project/funblocks-checker/maude-3.1 🦲 🕜 😵
File Edit View Search Terminal Help
minimaude> wrong input
Parse error
minimaude>
```

Mini-Maude: quit environment

```
marvinf@marvinf-X405UA:~/Desktop/master_project/funblocks-checker/maude-3.1 🕒 🕝 🛭
File Edit View Search Terminal Help
minimaude> q
goodbye
rewrites: 300 in 20ms cpu (185349ms real) (15000 rewrites/second)
result Portal: <>
     reduce in NAT3 : s s s s 0 .
rewrites: 1 in Oms cpu (Oms real) (~ rewrites/second)
result Nat3: s 0
Maude>
```

Funblocks checker



Funblocks checker

```
1 case f => g

Input
```

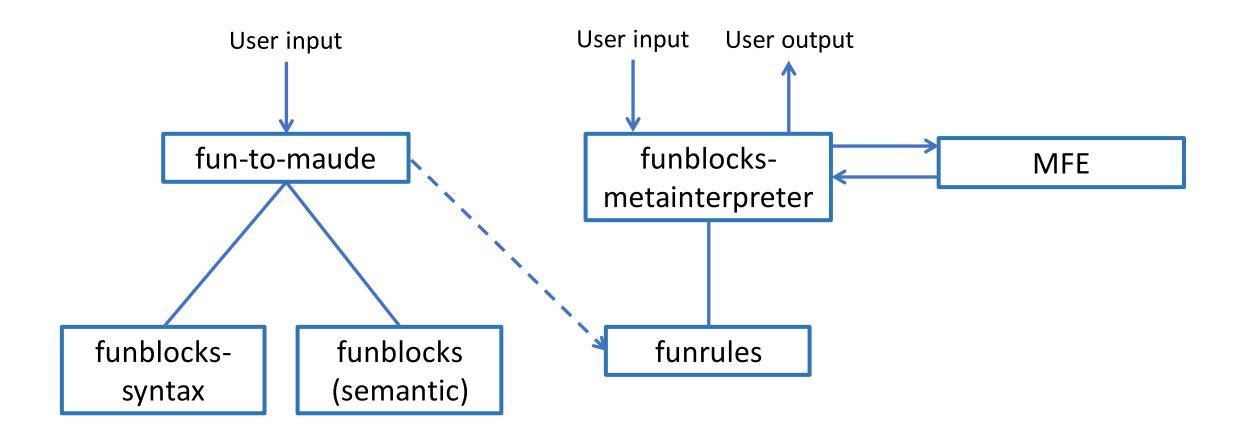
```
1  mod FUN-RULES is
2  sort Term .
3  op f : -> Term .
4  op g : -> Term .
5  rl f => g .
6  endm
```

Output

What's next?

```
1 mod FUN-RULES is
2 sort Term .
3 op f : -> Term . Typing?
4 op g : -> Term .
5 rl f => g . Type checking
6 endm
```

FunBlocks checker system



FUNBLOCK-SYNTAX

```
sorts Term Var Const Func Decl .

subsort Var Func Const < Term .

op case_=>_ : Term Term -> Decl .

op $_ : Token -> Var .

op _(_) : Token Token -> Func .

op _(_) : Token Term -> Func .
```

FUNBLOCKS

```
'case '$ 'x '=> 'id '( 'a ')
              funblocks
             (semantic)
'var ''$'x ': 'Untyped '. '\n
'op 'id ': 'Untyped '-> 'Untyped '. '\n
'op 'a ': '-> 'Untyped '. '\n
'rl '$'x '=> ''id '( ''a ') '. '\n
```

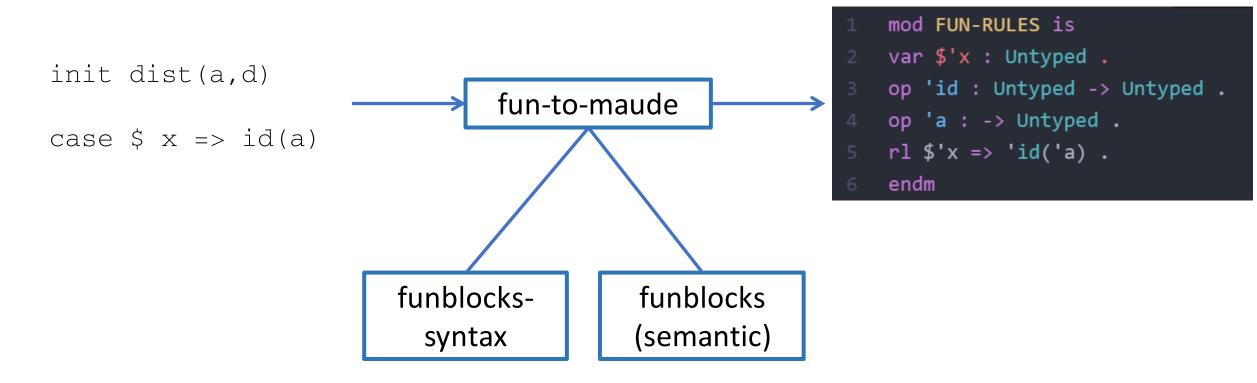
FUNBLOCKS

'case 'compose '('f ', 'g ', '\$ 'x ') '=> 'f '('g '('\$ 'x ') ')

funblocks
(semantic)

noParse()

FUN-TO-MAUDE



FUNRULES

```
1 mod FUN-RULES is
2 var $'x : Untyped .
3 op 'id : Untyped -> Untyped .
4 op 'a : -> Untyped .
5 rl $'x => 'id('a) .
6 endm
```

```
1  mod FUN-RULES is
2  inc FUNTYPES .
3  sorts Numbers String .
4  var $'x : Untyped .
5  op 'id : Untyped -> Untyped .
6  op 'a : -> Untyped .
7  rl $'x => 'id('a) .
8  endm
```

Pre-processing (script)

Post-processing (script)

case
$$id(\$'x) => \$'x$$

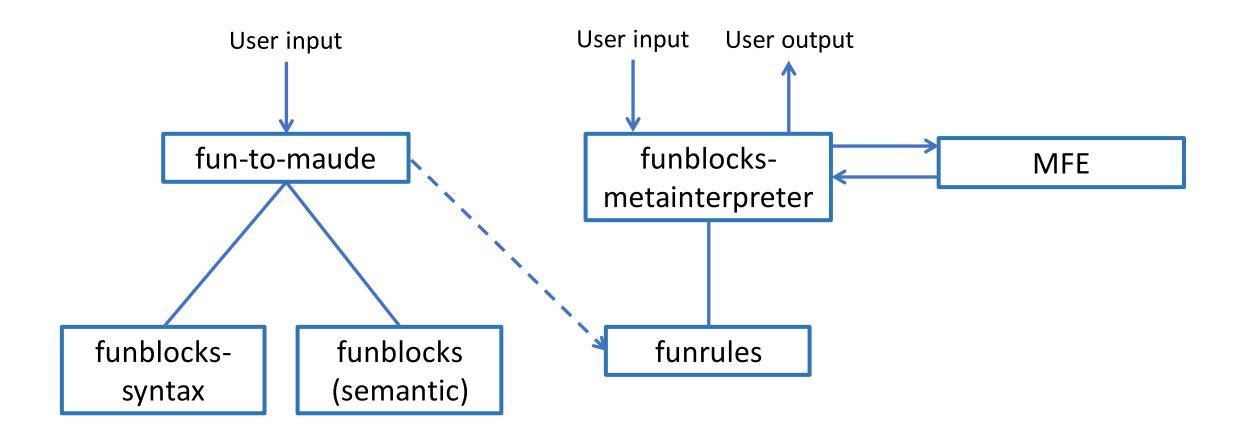
op a : -> Nat

op a : -> Nat

FUN-TO-MAUDE

file_process (script) mod FUN-RULES is var \$x : Untyped . init dist(a,d) op id : Untyped -> Untyped . fun-to-maude op a : -> Untyped . case \$x => id(a) $rl $x \Rightarrow id(a)$. endm funblocksfunblocks (semantic) syntax

FunBlocks checker system



Static Typing

```
type List $T :: empty | cons $T (List $T)
    rule ajoute $T :: $T -> List $T => List $T
    case ajoute($x, cons($y, $tail)) => cons($y, $tail)

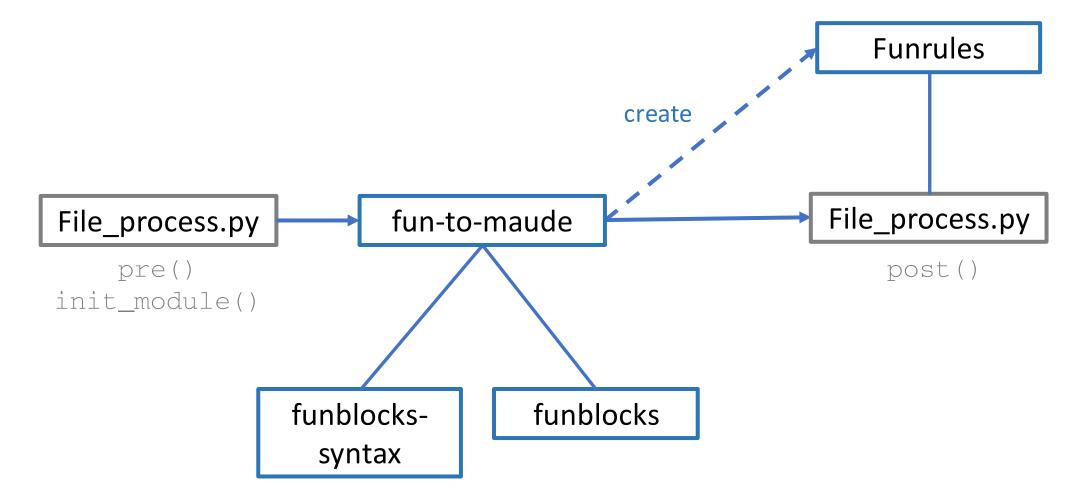
{ ajoute:
    [TypeVarRef { range: [Object], label: 'T' },
    TypeDeclRef { range: [Object], name: 'List', arguments: [Array] },
    TypeDeclRef { range: [Object], name: 'List', arguments: [Array] } ] }
```

Static Typing

```
case ajoute($x, cons($y, $tail)) => cons($y, $tail)
{ajoute:[TypeVarRef{range: null, label: 'TO'}, <2 empty items>]}
```

```
fmod MY-LIST{T :: TRIV} is
  protecting LIST{T} .
  protecting NAT .
  protecting QID .
  subsort Qid < T$Elt < List{T} .
  vars $x $y : T$Elt .
  vars $tail : List{T} .
  op empty : -> List{T} .
 op ajoute : T$Elt List{T} -> List{T} .
  op cons : T$Elt List{T} -> List{T} .
  eq ajoute($x, cons($y, $tail)) = cons($y, ajoute($x, $tail)) .
  eq ajoute($x, empty) = cons($x, empty) .
endfm
```

Execution



Maude tools

Inductive Theorem Prover (ITP)

Sufficient Completeness Checker (SCC)

Church-Rosser Checker (CRC)

Coherence Checker (ChC)

Maude Termination Tool (MTT)

Maude Formal Environment (MFE)

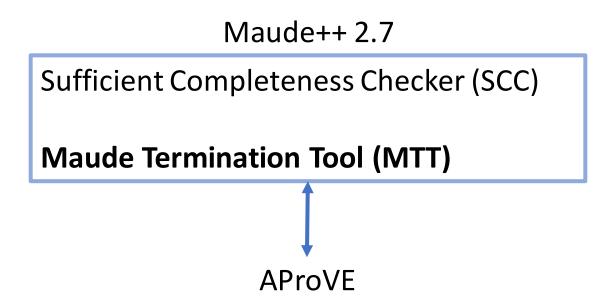
Maude Formal Environment

Maude 3.1 (last version)

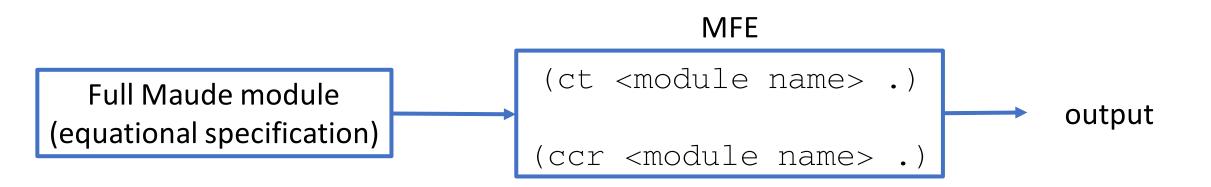
Inductive Theorem Prover (ITP)

Church-Rosser Checker (CRC)

Coherence Checker (ChC)



Maude Formal Environment





Works only with equational specifications

Church-Rosser Checker

```
1  mod TEST2 is
2  sort Generic .
3  op id : Nat -> Nat .
4  var x : Nat .
5  op a : -> Nat .
6  eq id(x) = x .
7  eq id(x) = a .
8  endm
```

```
Maude> (ccr TEST2 .)
rewrites: 36429 in 4ms cpu (3ms real) (9107250 rewrites/second)
Church-Rosser check for TEST2
The following critical pairs must be proved joinable:
    cp TEST21
    a
    = x:Nat .
        The module is sort-decreasing.
```

Church-Rosser Checker

```
1  mod TEST3 is
2  sort Generic .
3  op id : Nat -> Nat .
4  var x : Nat .
5  op a : -> Nat .
6  eq id(x) = x .
7  endm
```

```
Maude> (ccr TEST3 .)
rewrites: 36299 in 4ms cpu (3ms real) (9074750 rewrites/second)
Church-Rosser check for TEST3
All critical pairs have been joined.
The specification is locally-confluent.
The module is sort-decreasing.
```

Church-Rosser Checker

```
1  mod TEST3 is
2  sort Generic .
3  op id : Nat -> Nat .
4  var x : Nat .
5  op a : -> Nat .
6  eq id(x) = x .
7  endm
```

```
Maude> (ccr TEST3 .)
rewrites: 36299 in 4ms cpu (3ms real) (9074750 rewrites/second)
Church-Rosser check for TEST3
All critical pairs have been joined.
The specification is locally-confluent.
The module is sort-decreasing.
```

Termination tool

```
(fmod TEST2 is
 sort Nat .
 vars x y : Nat .
 op a : -> Nat .
 op b : -> Nat .
 op f : Nat Nat Nat -> Nat .
 op g : Nat Nat -> Nat .
 eq f(a,b,x) = f(x,x,x).
 eq g(x,y) = x.
 eq g(x,y) = y.
endfm)
```

```
rewrites: 66 in 4ms cpu (4ms real) (16500 rewrites/second)
The MTT has been set as current tool.
rewrites: 29 in 4ms cpu (4ms real) (7250 rewrites/second)
aprove is now the current external tool.
rewrites: 75 in Oms cpu (1ms real) (~ rewrites/second)
Success: The module TEST2 is non-terminating.
rewrites: 66 in 4ms cpu (0ms real) (16500 rewrites/second)
The CRC has been set as current tool.
rewrites: 126 in Oms cpu (3ms real) (~ rewrites/second)
Church-Rosser check for TEST2
The following critical pairs must be proved joinable:
  CD TEST22
    x:Nat
   = y:Nat .
        The module is sort-decreasing.
```

Equation vs rules (in Maude)



Equations



Rules

Represent theory

Terminating and confluent

Represent states changes

No constraints on termination and confluence

Reductions using rewrite logic

Parsing in Maude

reduce

red 1 + (1 + 1).

rewrite

rew [6] \$ \$ q q .

Precedence and gathering

```
1 + (2 * 3)
sort Nat .
ops 1 2 3 : -> Nat .
                                              1 + 2 + 3 - \begin{cases} 1 + (2 + 3) \\ \\ (1 + 2) + 3 \end{cases}
ops _+_ _*_ : Nat Nat -> Nat .
```

Precedence and gathering

gather

E: precedence value ≤ precedence value of the operator

e: precedence value < precedence value of the operator

&: any precedence value

Precedence and gathering

```
op _+_ : Nat Nat -> Nat [prec 33 gather (E e)] .
op _*_ : Nat Nat -> Nat [prec 31 gather (E e)] .
```

Default precedence values

```
0 for the constants (zero)
0 for mixfix operator which not begin nor end with _ ( (__))
41 for mixfix operator which begin or and with _ ( to_:_ )
-> 15 for unary operator ( not_ )
41 : for mixfix operator which begin and end with _ ( __+_)
```

Default gathering values

```
& for default prefix operators
& for operator which not begin nor end with _ ( (_))

E for adjacent _ or _ at the edges (_?_:_ is (E & E))
```

Special cases of binary operators which start and end with _ and have precedence value > 0 :

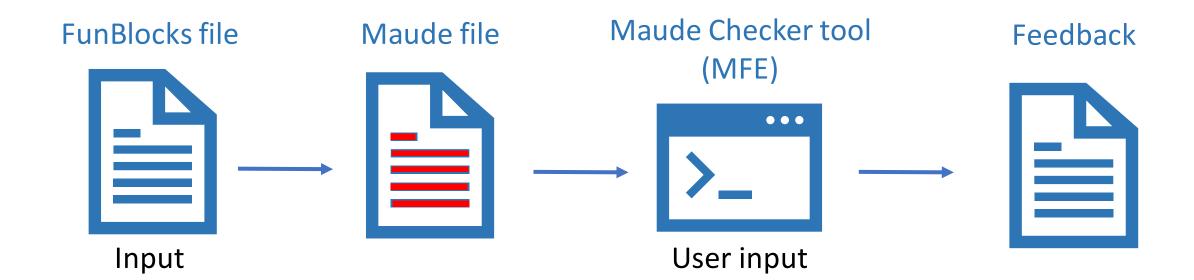
(e E) if it has the assoc attribute

If not, we have to look at the subsorts relationship...

Default gathering values

```
If we have Int < IntList</pre>
```

Pipeline



Pipeline

FunBlocks file Maude file Input Parser (Swift)

rule add :: Nat -> Nat => Nat op add : Nat Nat -> Nat . case add(
$$\$x$$
, zero) => $\$x$ eq add($\$x$, zero) = $\$x$.

```
1 type Nat :: a | b | succ Nat
2 rule f :: Nat -> Nat -> Nat => Nat
3 rule g :: Nat -> Nat => Nat
4 case f(a,b,$x) => f($x,$x,$x)
5 case g($x,$y) => $x
6 case g($x,$y) => $y
```

```
(fmod Nat is
 sort Nat .
op a : -> Nat .
op b : -> Nat .
op succ : Nat -> Nat .
endfm)
(fmod FUNRULES is
including Nat .
op f : Nat Nat Nat -> Nat .
op g : Nat Nat -> Nat .
var $x : Nat .
eq f(a, b, $x) = f($x, $x, $x).
var $x : Nat .
var $y : Nat .
eq g(x, y) = x.
var $x : Nat .
var $y : Nat .
eq g(x, y) = y.
 endfm)
```

```
rewrites: 66 in 4ms cpu (0ms real) (16500 rewrites/second)
The CRC has been set as current tool.

rewrites: 128 in 4ms cpu (4ms real) (32000 rewrites/second)
Church-Rosser check for FUNRULES
The following critical pairs must be proved joinable:
    cp FUNRULES2
    $x:Nat
    = $y:Nat .
        The module is sort-decreasing.
```

Term Rewriting System R:

[\$x, \$y]
f-`[Nat`]-`[Nat`]-`[Nat`](a, b,
\$x) -> f-`[Nat`]-`[Nat`]`[Nat`](\$x, \$x, \$x)
g-`[Nat`]-`[Nat`](\$x, \$y) -> \$x
g-`[Nat`]-`[Nat`](\$x, \$y) -> \$y

Termination of R to be shown.

R ->Dependency Pair Analysis

Found an infinite P-chain over R:

Thus, s starts an infinite chain.

Non-Termination of R could be shown.

```
op zero : -> Nat .
                                                                         op succ : Nat -> Nat .
                                                                         endfm)
                                                                         (view $T from TRIV to Nat is sort Elt to Nat
                                                                         (fmod Tree{T :: TRIV} is
                                                                          sort Tree{T} .
                                                                         op empty : -> Tree{T} .
   type Nat :: zero | succ Nat
                                                                         op leaf : T$Elt -> Tree{T} .
   type Tree \$T:: empty | leaf \$T | node (Tree \$T) (Tree \$T)
                                                                         op node : Tree{T} Tree{T} -> Tree{T} .
   rule depth \$T :: Tree \$T => Nat
                                                                         endfm)
   rule add :: Nat -> Nat => Nat
                                                                         (fmod FUNRULES is
   case add(\$x, add(\$y, zero)) => \$x
                                                                         including Nat .
   case depth(empty) => zero
6
                                                                         including Tree{$T} .
   case succ(\$z) \Rightarrow zero
                                                                         op depth : Tree{$T} -> Nat .
                                                                         op add : Nat Nat -> Nat .
                                                                         var $x : Nat .
                                                                         var $y : Nat .
                                                                         eq add($x, add($y, zero)) = $x.
                                                                         eq depth(empty) = zero .
                                                                         var $z : Nat .
                                                                         eq succ(\$z) = zero.
                                                                          endfm)
```

(fmod Nat is

sort Nat .

```
Success: The module FUNRULES is terminating.

rewrites: 66 in 0ms cpu (1ms real) (~ rewrites/second)

The CRC has been set as current tool.

rewrites: 114 in 4ms cpu (4ms real) (28500 rewrites/second)

Church-Rosser check for FUNRULES

All critical pairs have been joined.

The specification is locally-confluent.

The module is sort-decreasing.
```

```
1 type Nat :: a | b | c | zero | succ Nat
2 rule f :: Nat -> Nat => Nat
3 rule g :: Nat => Nat
4 case f($x, $x$) => a
5 case f($x, g($x$)) => b
6 case c => g(c)
```

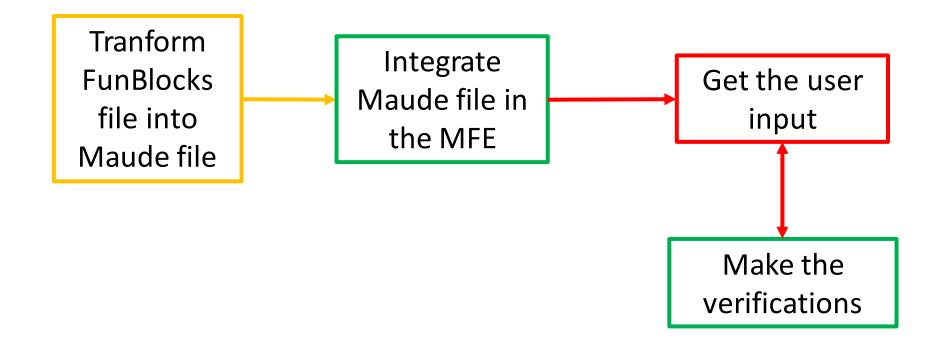
No critical pairs to show. One of the checking commands must be successfully executed previously.

Fatal error: stack overflow.

This can happen because you have an infinite computation, say a runaway recursion, or model checking an infinite model. It can also happen because the stacksize limit in your environment is set too low for the computation you are trying to do. You can find the value of your stacksize with the tcsh command 'limit stacksize' or the bash command 'ulimit -s'.

Depending on your operating system configuration you may be able to increase your stacksize with the tcsh command 'unlimit stacksize' or the bash command 'ulimit -s unlimited'.

« TODO » list



```
(fmod Bool is ... endfm)
(fmod Rel is ... endfm)
(view $T from TRIV to Rel is sort Elt to Rel . endv)
(fmod List{T :: TRIV} is
  op cons : List{T} T$Elt -> List{T} .
endfm)
(fmod FUNRULES is
 including Bool .
 including Rel .
 including List{$T} .
 op size : List{$T} -> Rel .
 op isEmpty : List{$T} -> Bool .
  eq size(empty) = zero .
 var $x : Rel .
  eq size(cons(empty, $x)) = succ(zero).
endfm)
```

```
(fmod Bool is ... endfm)
(fmod Rel is ... endfm)
(view $T1 from TRIV to Bool is sort Elt to Bool . endv)
(view $T2 from TRIV to Rel is sort Elt to Rel . endv)
(fmod List{T :: TRIV} is
 op cons : List{T} T$Elt -> List{T} .
endfm)
(fmod FUNRULES is
 including Bool .
  including Rel .
  including List($T1) .
 including List($T2) .
 op size : List{$T2} -> Rel .
 op isEmpty : List{$T2} -> Bool .
  eq size(empty) = zero .
 var $x : Rel .
  eq size(cons(empty, $x)) = succ(zero).
endfm)
```

« TODO » list



Type inference of variables



Generic types



Check FunBlocks code before verifications

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Make FunBlocks alive

Marvin FOURASTIE

Master project