# Make FunBlocks alive

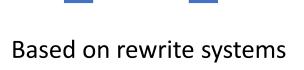
Marvin FOURASTIE

Master project

### Motivations

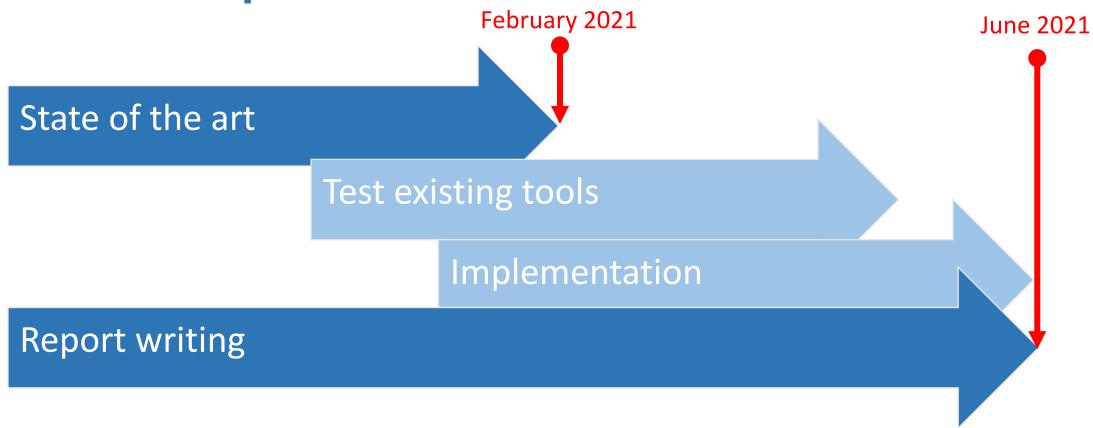






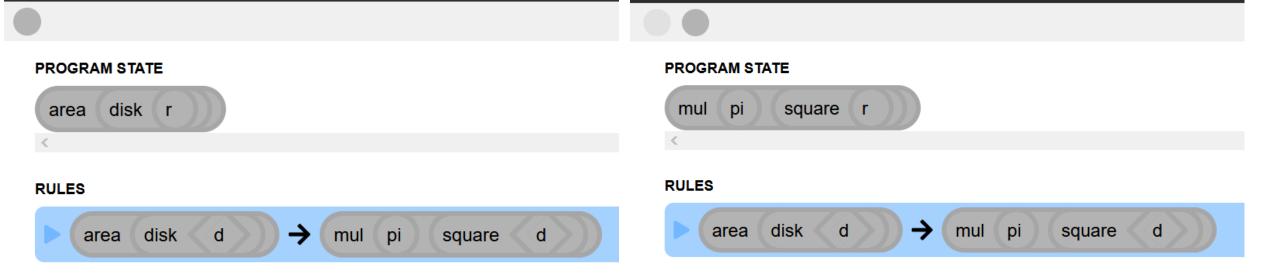
# **FunBlocks**

# Road Map



### **FunBlocks**

```
init area(disk(r))
case area(disk($d$)) => mul(pi,square($d$))
```



### **FunBlocks**

area(disk(r)) Declarative paradigm case area(disk(\$d)) => mul(pi,square(\$d)) Based on rewrite systems square Static typing node (Tree \$t) (Tree \$t) type Tree \$t :: empty | leaf \$t |

### Goals

Provide users with valuable insights about their program

Verification of rewrite systems

## Rewrite systems

### Stack operators

 $Zero = \{0\}$ 

Nat = Zero U succ(Nat)

Empty =  $\Lambda$ 

Stack = Empty U push(Nat, Stack)

top : Stack  $\rightarrow$  Nat

pop : Stack → Stack

alternate : Stack × Stack → Stack

## Rewrite systems

### Canonical rewrite system

$$top(push(x, y)) = x$$
 
$$top(push(x, y)) \rightarrow x$$
 
$$pop(push(x, y)) = y$$
 
$$pop(push(x, y)) \rightarrow y$$
 
$$alternate(\Lambda, z) = z$$
 
$$alternate(push(x, y), z) = push(x, alternate(z, y))$$
 
$$alternate(push(x, y), z) \rightarrow push(x, alternate(z, y))$$

# Rewrite systems

**Termination** 

Confluence

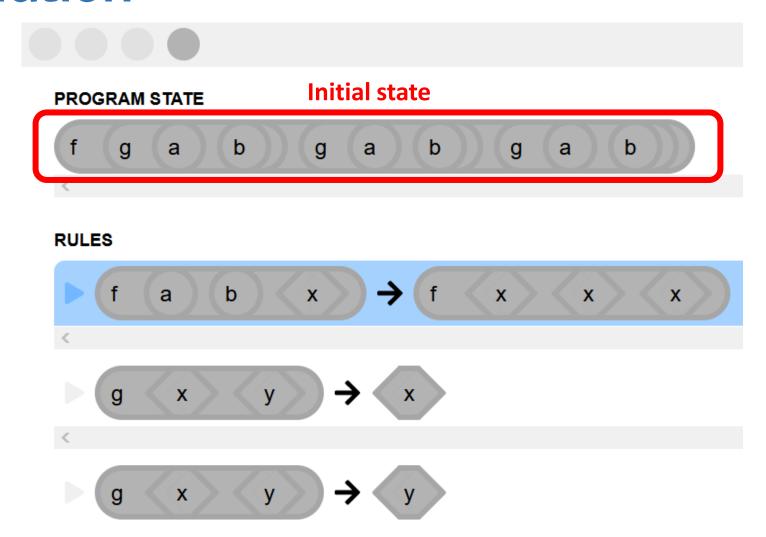
Soundness

Completeness

Correctness

Undecidable in general

### **Termination**



### Reduction order

Monotone

$$s_i > t \rightarrow f(s_1, ..., s_i, ..., s_n) > f(s_1, ..., t, ..., s_n)$$
  
For  $f$  of arity  $n$ 

Close under substitution ————

$$s>t\to\sigma s>\sigma t$$
 , for all substitution  $\sigma$ 

Well-founded

no infinite descending chain

 $(\mathbb{N}, <)$  is well-founded

 $(\mathbb{Z}, <)$  is not well-founded

### **Termination**

A term rewriting system is terminating

if and only if

it admits a compatible reduction order < (if l > r for every rewrite rule  $l \rightarrow r$ )



Verification of termination

## Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(a) = 1$$

$$w(g(t)) = 1 + w(t)$$

$$w(f(t_1,t_2)) = 2w(t_1) + w(t_2)$$

## Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(f(a,x)) = 2 + w(x)$$

$$w(f(g(x),y)) = 2 + 2w(x) + w(y)$$

$$w(g(f(x,y)) = 1 + 2w(x) + w(y)$$

$$w(f(a,x)) > w(x)$$

$$w(f(g(x),y)) > w(g(f(x,y))$$

Reduction order → Termination

## Algorithms

Recursive Path Ordering

Order based on the mutisets

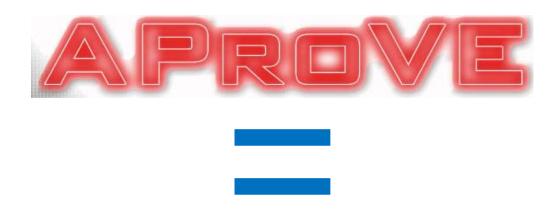
Knuth-Bendix Ordering

Based on weights assigned to operators

Dependency pairs

Prove innermost termination

### **Termination**

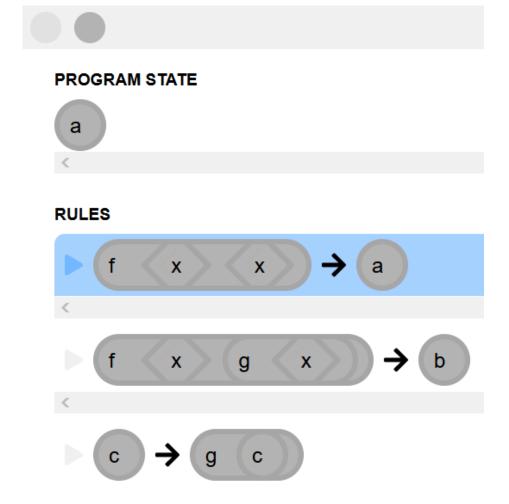


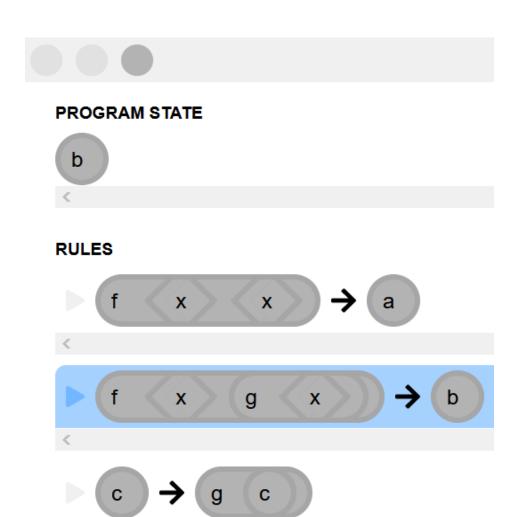
Direct proof (polynomial, LBO, KBO,...)



Dependency pairs and size-change principle

### Confluence

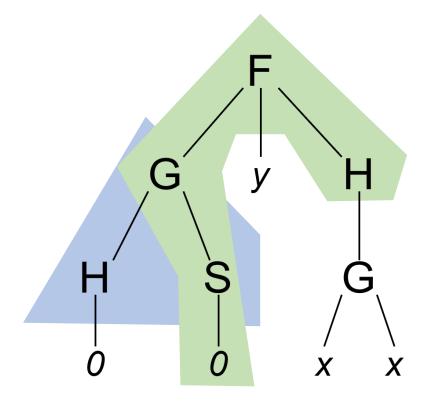




## Overlap and critical pairs

$$\rho_1 : F\left(G(x, S(0)), y, H(z)\right) \to x$$

$$\rho_2 : G(H(x), S(y)) \to y$$



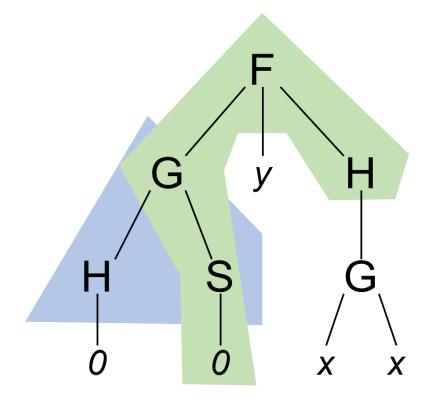
## Overlap and critical pairs

### Overlapping:

Term: 
$$F(G(H(0),S(0)),y,H(z))$$

$$F(G(\square,S(0)),\square,H(\square))$$

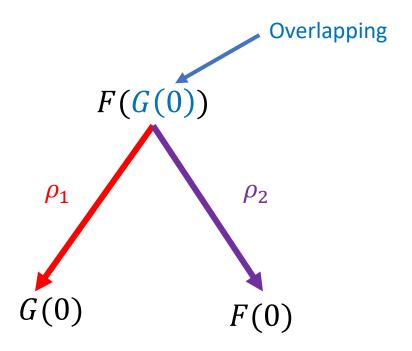
$$G(H(\square),S(\square))$$



## Overlap and critical pairs

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



< G(x), F(x) > is called critical pair

### Critical Pair Lemma

A terminating rewriting system is confluent

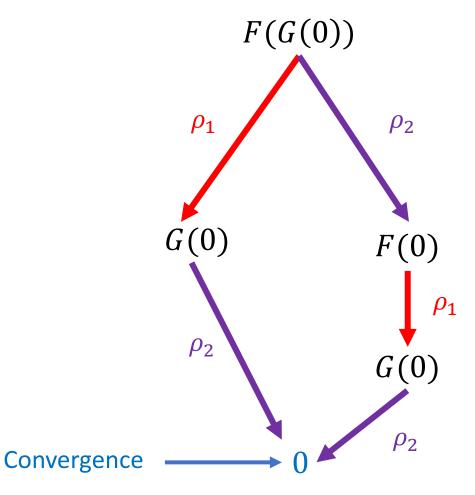
if and only if

all critical pairs are convergent

### Critical Pair Lemma

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



### Input:

A set of equation

A reduction ordering <

$$\begin{array}{ll} 1 \cdot x & = x \\ x^{-1} \cdot x & = 1 \\ (x \cdot y) \cdot z & = x \cdot (y \cdot z) \end{array}$$

Non-confluent

### Output:

Terminate successfully

Terminating and confluent rewrite system

Non-terminating rewrite system

Rule which cannot be ordered (i.e. commutative operator)

#### Basic rules:

Orienting \_\_\_\_\_

Transform s = t to  $s \rightarrow t$ 

Adding

 $\longrightarrow$ 

Add s = t in the set of equation

Simplifying

 $\longrightarrow$ 

Simplify s = t in s' = t'

Deleting



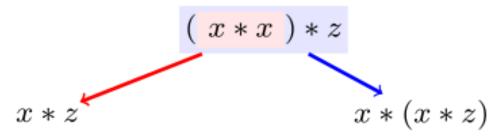
Delete trivial rules s = s

#### Adding

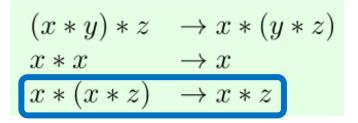


Add s = t in the set of equation

$$\begin{array}{ccc} (x*y)*z & \rightarrow x*(y*z) \\ x*x & \rightarrow x \end{array}$$







New rule added

#### Completion process:

- 1. For each equation s = t reduce s and t to normal form s' and t'
- 2. Fill the set of rules using basic operators and reduction ordering
- 3. If the algorithm terminate successfully: terminating and confluent rewrite system

Completion for axioms of groups:

$$\begin{array}{rcl}
1 \cdot x & = x \\
x^{-1} \cdot x & = 1 \\
(x \cdot y) \cdot z & = x \cdot (y \cdot z)
\end{array}$$



$$\begin{array}{cccc}
1 \cdot x & \rightarrow x \\
x^{-1} \cdot x & \rightarrow 1 \\
(x \cdot y) \cdot z & \rightarrow x \cdot (y \cdot z) \\
x^{-1} \cdot (x \cdot y) & \rightarrow y \\
1^{-1} & \rightarrow 1 \\
x \cdot 1 & \rightarrow x \\
(x^{-1})^{-1} & \rightarrow x \\
x \cdot x^{-1} & \rightarrow 1 \\
x \cdot (x^{-1} \cdot y) & \rightarrow y \\
(x \cdot y)^{-1} & \rightarrow y^{-1}x^{-1}
\end{array}$$

### **Educational tools**



**CSI** 

**KBCV** 

**Termination** 

Confluence

Completion

### TRS tool

Parcourir... Aucun fichier sélectionné. Upload  $(VAR \times y)$ (RULES  $f(x,y) \rightarrow x$  $f(x,y) \rightarrow f(x,g(y))$  $g(x) \rightarrow h(x)$  $F(g(x),x) \rightarrow F(x,g(x))$  $F(h(x),x) \rightarrow F(x,h(x))$ (COMMENT Example 6 of \cite{AT97}) (COMMENT %% TagRevision: 1 %%) (COMMENT %% Tags: [4ec3f85c01836]non left linear{}; [4ec3f87f0f1e0]r 50 × Rewrites Limit (Use with caution) Go!

## TRS tool

	R <sub>0</sub>	$= f(x,y) \rightarrow x$									
R <sub>0</sub>	is	Left-Linear									
R <sub>0</sub>	is	Right-Linear									
Ro	is	Linear									
R <sub>0</sub>	is	Collapsing									
Ro	is	not Duplicating									
R <sub>0</sub>	is	not Conservative									
R <sub>0</sub>	is	Destructive									

TRS										
The	TRS	is	not Left-Linear							
The	TRS	is	not Right-Linear							
The	TRS	is	not Linear							
The	TRS	is	Collapsing							
The	TRS	is	not Duplicating							
The	TRS	is	not Conservative							
The	TRS	is	Destructive							
The	TRS	is	not Orthogonal							
The	TRS	is	not Almost Orthogonal							
The	TRS	is	not Weakly Orthogonal							
The	TRS	is	Locally Confluent							
Unkı	nown	COI	nfluence for The TRS							
			non terminating pop: $f(x,g(y)) \rightarrow f(x,g(y))$							

### TTT2

#### **Tyrolean Termination Tool 2 (1.20)**

#### 1. Input Term Rewrite System

For input use the standard TRS format.

#### 2. Select Strategy

$\odot$	FAST	$\bigcirc$	FBI	$\bigcirc$	HYDRA	$\bigcirc$	LPO	$\bigcirc$	KBO	$\circ$	POLY	$\bigcirc$	MAT(2)	$\bigcirc$	MAT(3)	$\bigcirc$	COMP	$\bigcirc$	COMPLEXITY
$\bigcirc$	EXPE	RT																	

#### 3. Encode State into URL (optional)

encode URL clear URL

#### 4. Start TTT2

**check** use HTML output if available (*experimental feature*)

Enter a TRS or HRS or upload a file browse...

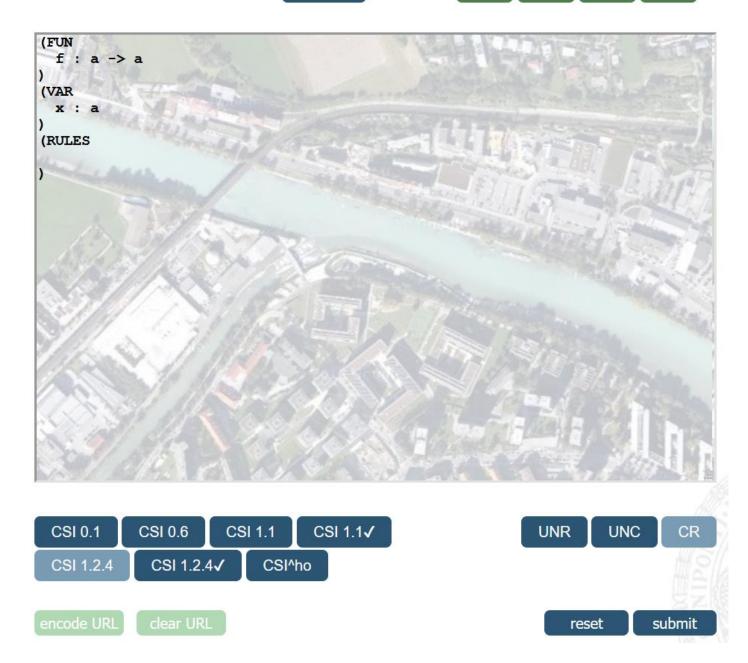
Examples: trs1

trs2

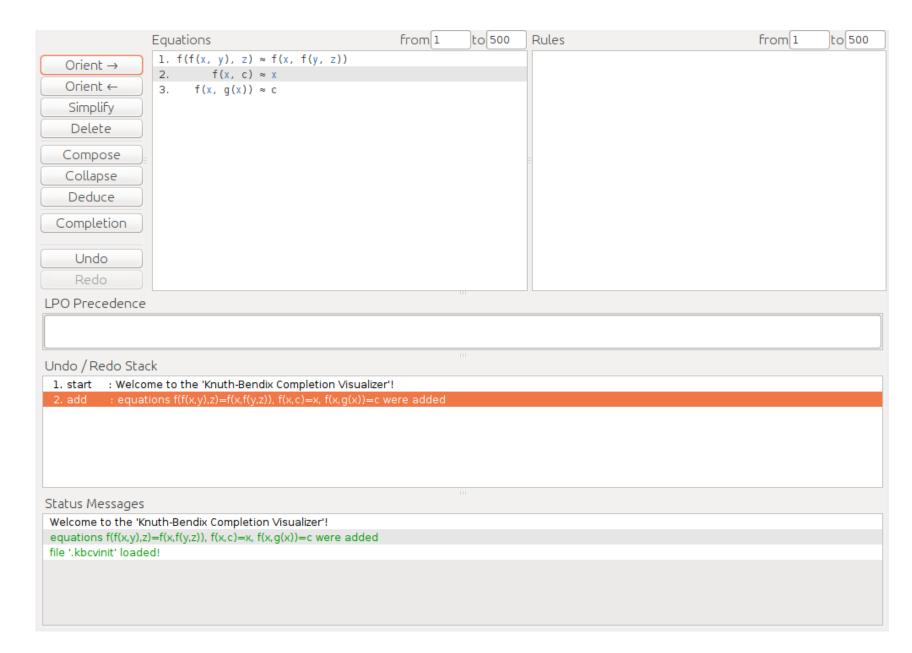
hrs1

hrs2

## **CSI**



### **KBCV**



### Performance tools

**Termination** 

Confluence

Completion

**Proof verification** 

**MU-TERM** 

**ACP** 

Maxcomp

CoLoR

**NaTT** 

Saigawa

**CeTA** 



# Hybrid tools







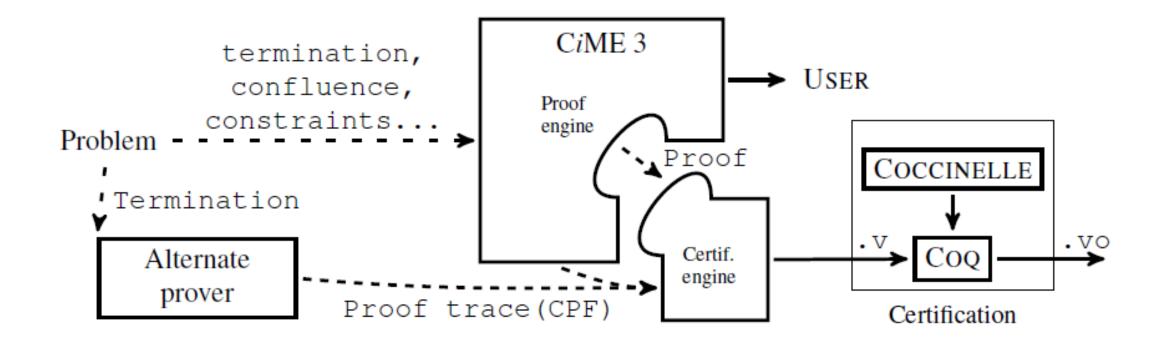


Rewriting toolkit





**Proof certification** 



#### Examples of declarations

```
let X = variables "x,y";
let F = signature "plus : binary; 0:constant; S:unary;";
let T = algebra F;
let t1 = term T "S(0)";
let R = trs T "plus(0,x) -> x; plus(S x, y) -> S(plus(x,y));";
let c = order\_constraints T "0 < S(0) / S(plus(x,y)) < plus(S(x),y)";
```

#### Definition of signatures

```
CiME>let X = variables "x,y,z";
X : variable_set = variables "z,x,y"
```

#### Definition of algebra and terms

```
CiME> let A_peano = algebra F_peano ;
A_peano : F_peano algebra = algebra F_peano
```

```
CiME> let t = term A_peano "s(s(s(0)))*(s(0)+s(s(0)))";
t : F_peano term = s(s(s(0))) *(s(0)+s(s(0)))
```

#### Term rewriting system

```
CiME> let R_peano = trs A_peano "
      x+0 -> x;
      x+s(y) \rightarrow s(x+y);
      x*0 -> 0;
      x*s(y) -> (x*y) +x;
       ";
 R_peano : F_peano trs = trs A_peano "
            x+0 -> x;
            x+s(y) \rightarrow s(x+y);
            x * 0 -> 0;
            x * s (y) -> (x * y) + x "
```

```
CiME> termination R_peano;
CiME> coq_certify_proof R_peano;
CiME> convergence R_peano ;
...
```





**Simplicity** 



Expressiveness



Performance

```
fmod BASIC-NAT is
        sort Nat .
       op 0 : -> Nat .
       op s : Nat -> Nat .
       op + : Nat Nat -> Nat .
        vars N M : Nat .
        eq 0 + N = N.
       eq s(M) + N = s(M + N).
endfm
```

```
fmod FACTORIAL is
protecting NAT .

op _! : Nat -> NzNat .
var N : Nat .

eq 0 ! = 1 .
eq (s N) ! = (s N) * N ! .

endfm
```

```
> load factorial.maude
> red 100 ! .
Reduce in FACTORIAL : 100 ! .
rewrites: 201 in 0ms cpu (0ms real) (~ rewrites/second)
result NzNAT:
9332621544394415268169923885626670049071596826438162146
8592963895217599993229915608941463976156518286253697920
82722375825118521091686400000000000000000000000
```

```
mod VENDING-MACHINE is
  including VENDING-MACHINE-SIGNATURE .
  var M : Marking .
  rl [add-q] : M \Rightarrow M q.
  rl [add-$] : M => M $ .
  rl [buy-c] : $ => c .
  rl [buy-a] : $ => a q .
  rl [change] : q q q => $.
endm
```

Inductive Theorem Prover (ITP)

Sufficient Completeness Checker (SCC)

Church-Rosser Checker (CRC)

Coherence Checker (ChC)

Maude Termination Tool (MTT)

Maude Formal Environment (MFE)

## Tools overview

	Maude	CiME
Extensibility	-	<b>~</b>
Still active		
I/O files	+	+
Syntax	-	+
Documentation		

## Maude MSOS Tool (MMT)





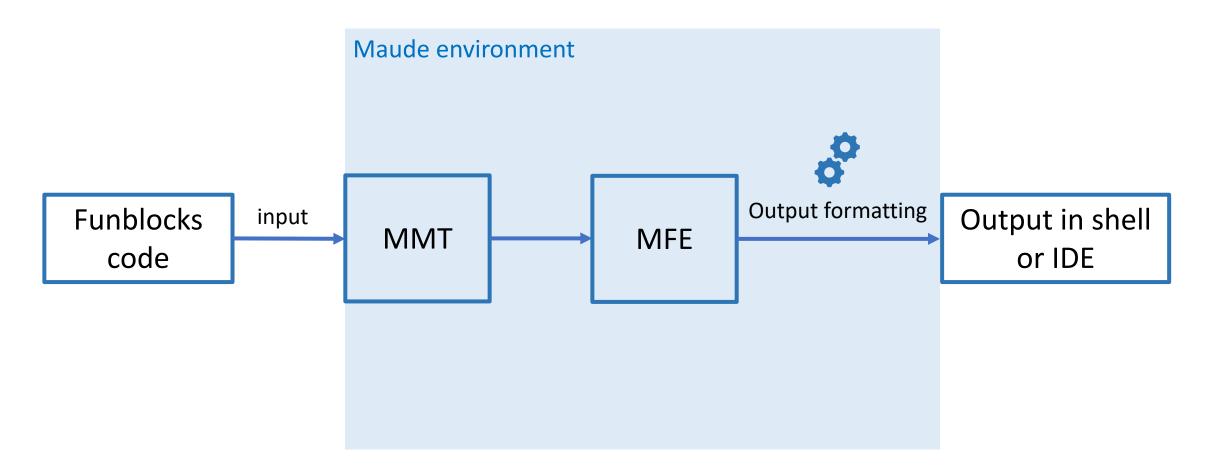


Based on transition rules



Modularity

## System diagram





Standard input/output



File handling



Sockets

```
mod VENDING-MACHINE is
  including VENDING-MACHINE-SIGNATURE .
  var M : Marking .
  rl [add-q] : M \Rightarrow M q.
  rl [add-\$] : M => M \$ .
  rl [buy-c] : $ => c .
  rl [buy-a] : $ => a q .
  rl [change] : q q q => $.
endm
```

```
fmod VENDING-MACHINE-SIGNATURE is
sorts Coin Item Marking .
subsorts Coin Item < Marking .
op __ : Marking Marking -> Marking [assoc comm id: null] .
op null : -> Marking .
op $ : -> Coin [format (r! o)] .
op q : -> Coin [format (r! o)] .
op a : -> Item [format (b! o)] .
op c : -> Item [format (b! o)] .
endfm
```

```
load vending-machine-signature.maude
fmod VENDING-MACHINE-GRAMMAR is
  protecting VENDING-MACHINE-SIGNATURE .
  protecting NAT .
  sort Action .
  op insert $ : -> Action .
  op insert q : -> Action .
  op show basket : -> Action .
  op show credit : -> Action .
  op buy__(s) : Nat Item -> Action .
endfm
```

```
load vending-machine-grammar.maude
load buying-strats.maude
load file.maude
mod VENDING-MACHINE-IO is
  rl < 0 : X | action : insert $, marking : M, Atts >
  => < 0 : X | action : idle,
       marking : downTerm(insertCoin('add-$, upTerm(M)), null), Atts >
    write(stdout, O, "one dollar introduced\n") .
```

#### Parse input?

## **META-LEVEL** module



Parse input

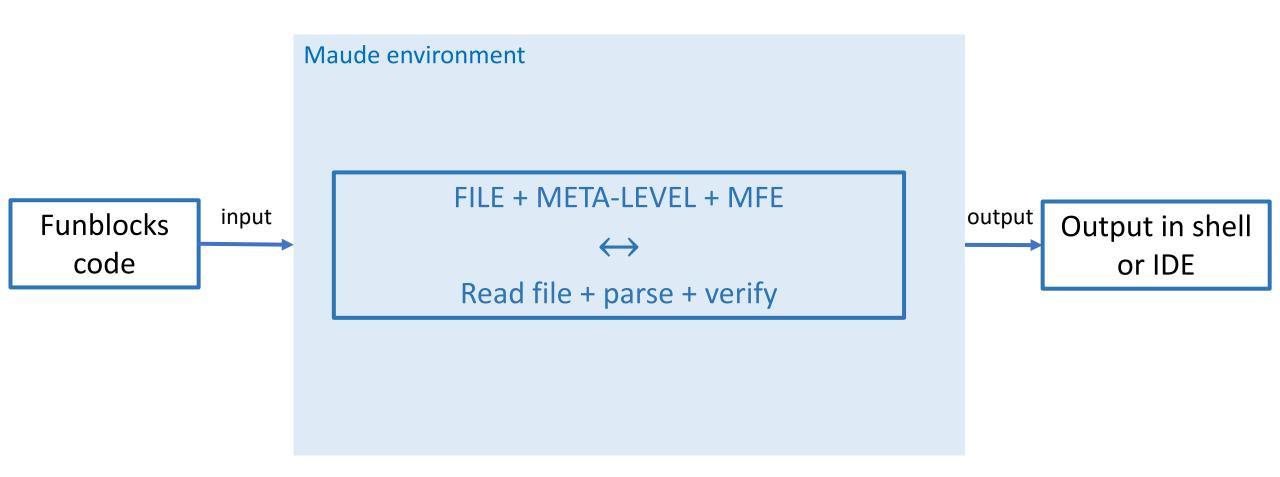


**Execute** 



Print

## System diagram



### What's next?



META-LEVEL module



Handle FunBlock file input

### References

- 1. Didier Buchs, modelisation verification course material, 2018
- Dimitri Racordon, Emmanouela Stachtiri, Damien Morard, Didier Buchs, Functional Block Programming and Debugging, 2020
- 3. Nachum Dershowitz, Jean-Pierre Jouannaud, Rewrite Systems, 1990
- 4. Nachum Dershowitz, Computing with Rewrite Systems, 1985
- 5. Terese, Term Rewriting Systems, 2003
- Thomas Sternagel and Harald Zankl, KBCV-Knuth-Bendix Completion Visualizer, 2012
- 7. Thomas Artsa, Jürgen Giesl, Termination of term rewriting using dependency pairs, 2000

### References

- 8. Jürgen Giesl, René Thiemann, Peter Schneider-Kamp, Stephan Falke, Automated Termination Proofs with AProVE, 2004
- 9. D. Kapur, P. Narendran, Path ordering for proving termination of term rewriting systems, 1985
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- 11. E. Contejean, P. Courtieu, J. Forest, O. Pons, X. Urbain, Automated Certified Proofs with CiME3, 2011
- 12. F. Chalub, C. Braga, Maude MSOS Tool, 2005
- 13. S. Winkler, A. Middeldorp, Tools in Term Rewriting for Education, 2020
- 14. A. Salvador, L. Salvador, Term Rewriting Systems .Net Framework, 2013

## References (links)

#### Database of Rewriting Systems

- http://rewriting.loria.fr/systems.html
- http://www.jaist.ac.jp/~hirokawa/tool/

#### **Knuth-Bendix Completion Visualizer**

http://cl-informatik.uibk.ac.at/software/kbcv/

#### Knuth-Bendix Completion subject-based thesis

• <a href="https://homepage.divms.uiowa.edu/~astump/papers/thesis-wehrman.pdf">https://homepage.divms.uiowa.edu/~astump/papers/thesis-wehrman.pdf</a>

## References (links)

Prolog implementation of the Knuth-Bendix completion procedure

https://www.metalevel.at/trs/

#### Maude tools

- http://maude.lcc.uma.es/CRChC/
- http://www.lcc.uma.es/%7Eduran/MTT/
- http://maude.sip.ucm.es/debugging/

#### Wikipedia

- https://fr.wikipedia.org/wiki/Compl%C3%A9tion\_de\_Knuth-Bendix
- https://fr.wikipedia.org/wiki/Paire critique

## References (links)

#### TRS tool:

• <a href="http://tfmserver.dsic.upv.es:8080/Home.html">http://tfmserver.dsic.upv.es:8080/Home.html</a>

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Marvin FOURASTIE

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