

SOME DETAILS ABOUT REWIRING SYSTEMS

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Strongly based on a
présentation of Roel de Vrijer
and Alexis Marechal



RECALL

- Signature $\Sigma = \langle S, F \rangle$
- From Σ we can derive the set T_Σ , *i.e.*, the set of all terms generated by Σ .
- Give a set X of S indexed variables, we can derive the set $T_{\Sigma, X}$ of terms with variables.
- We can give a semantics to all of that by adding equations (called axioms), the whole pack $\Sigma + X + \text{Axioms} = \text{Adt}$

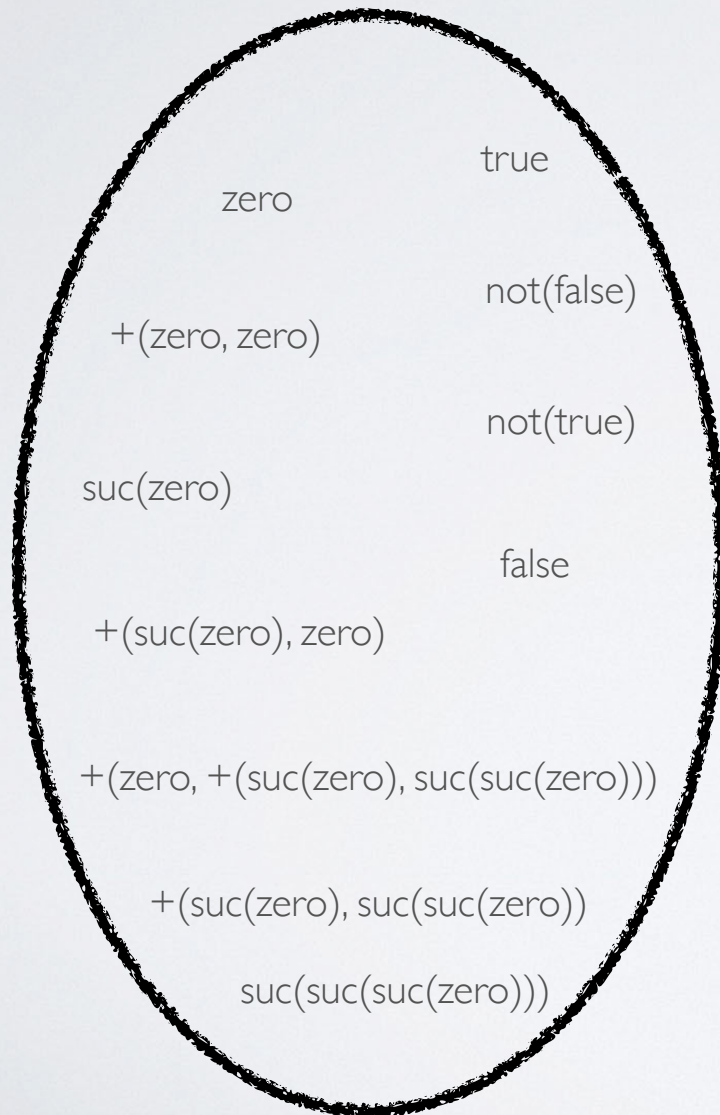
$$\Sigma = \langle S, F \rangle$$

$$S = \{\text{nat}, \text{bool}\}$$

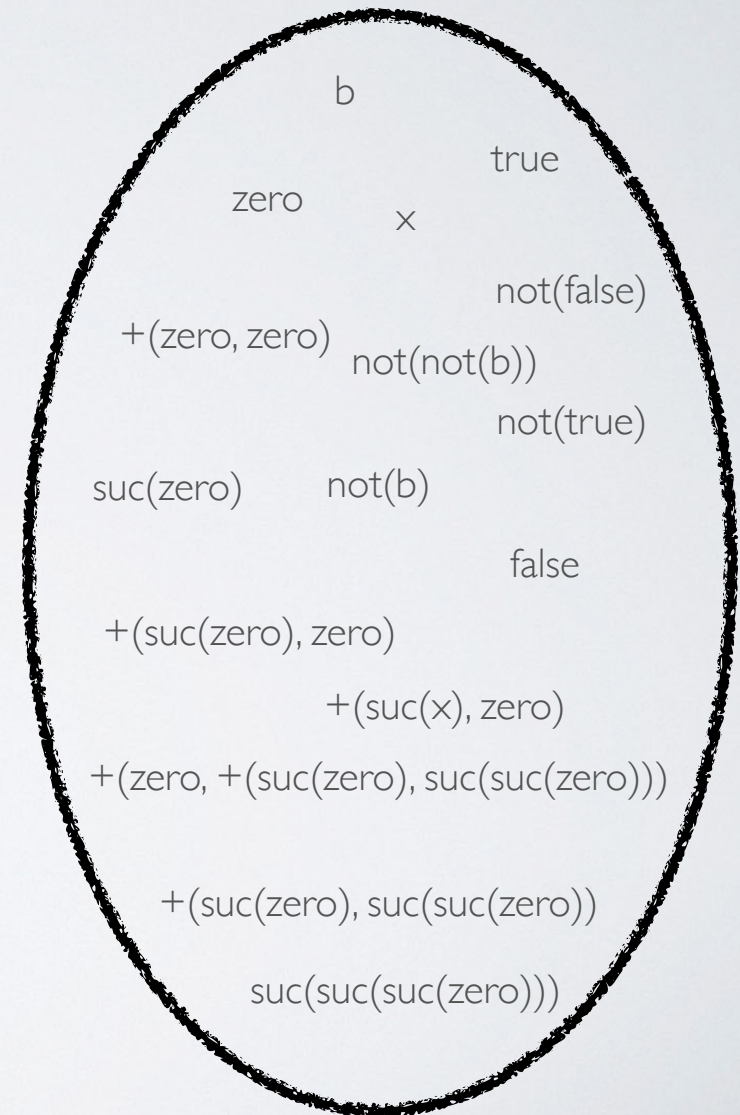
$$X_s = \{x_{\text{nat}}, b_{\text{bool}}\}$$

$$F = \{\text{true}_{\varepsilon, \text{bool}}, \text{false}_{\varepsilon, \text{bool}}, \text{zero}_{\varepsilon, \text{nat}}, \text{suc}_{\text{nat}, \text{nat}}, +_{\text{nat}, \text{nat}, \text{nat}}, \text{not}_{\text{bool}, \text{bool}}\}$$

$$(T_{\Sigma})_s =$$



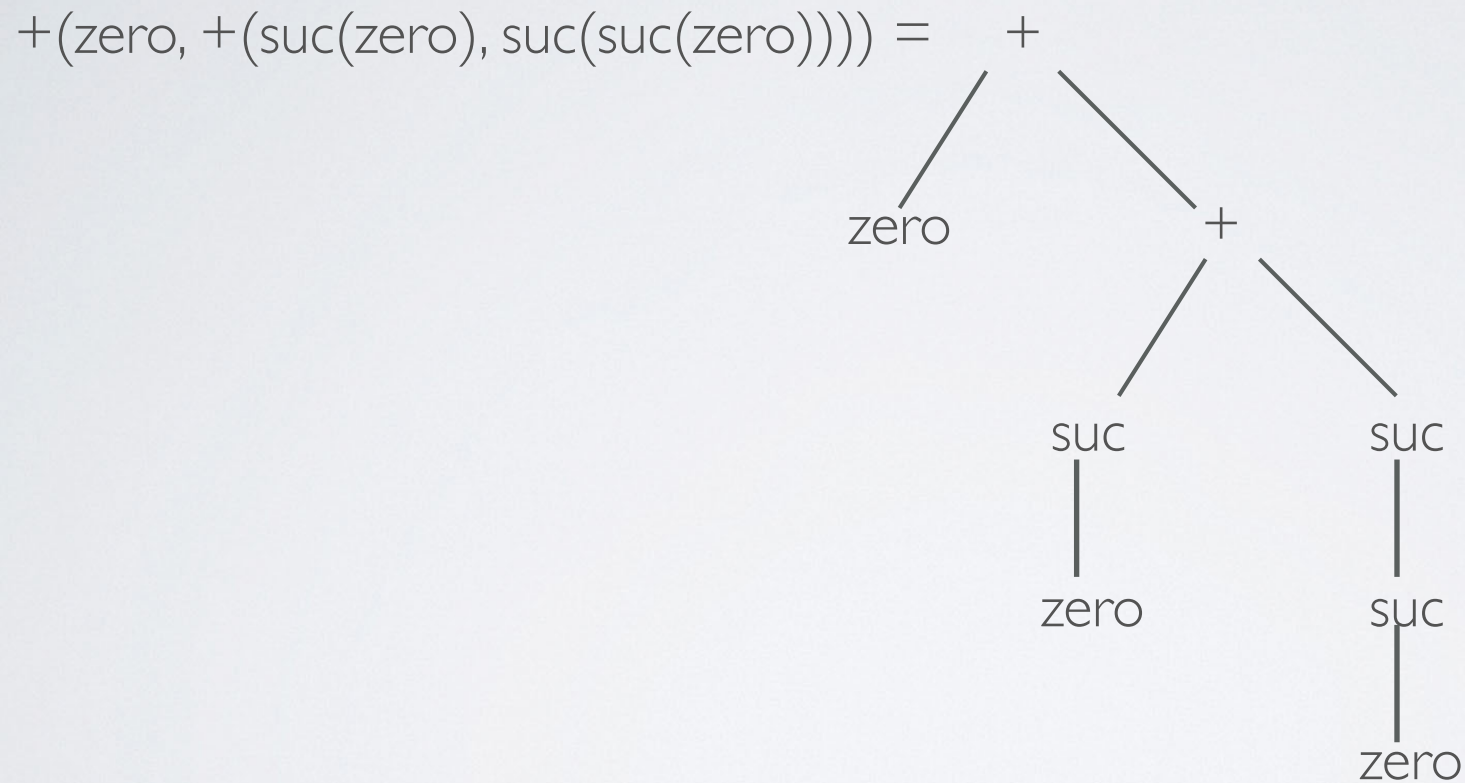
$$(T_{\Sigma, X})_s =$$



SYNTACTIC DERIVATION

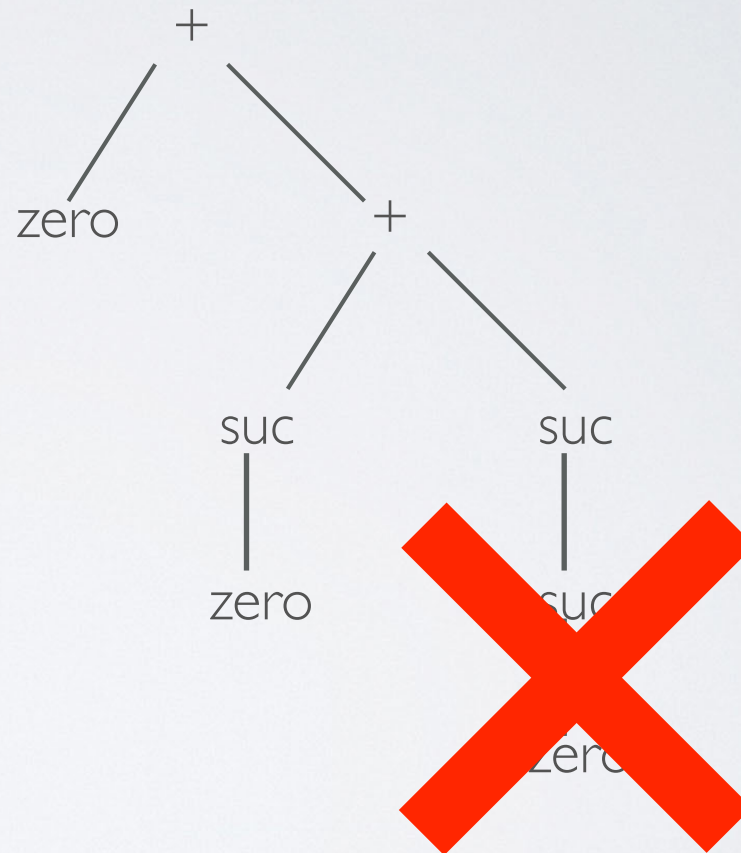
- Two types of terms: atoms and composite terms.
 - atoms: true, false, zero; composite: $\text{suc}(\text{zero})$, $+(\text{zero}, \text{zero})$, etc.)

COMPOSITE TERM AS TREE



CONTEXT

$+(zero, +(suc(zero), suc(_))) =$

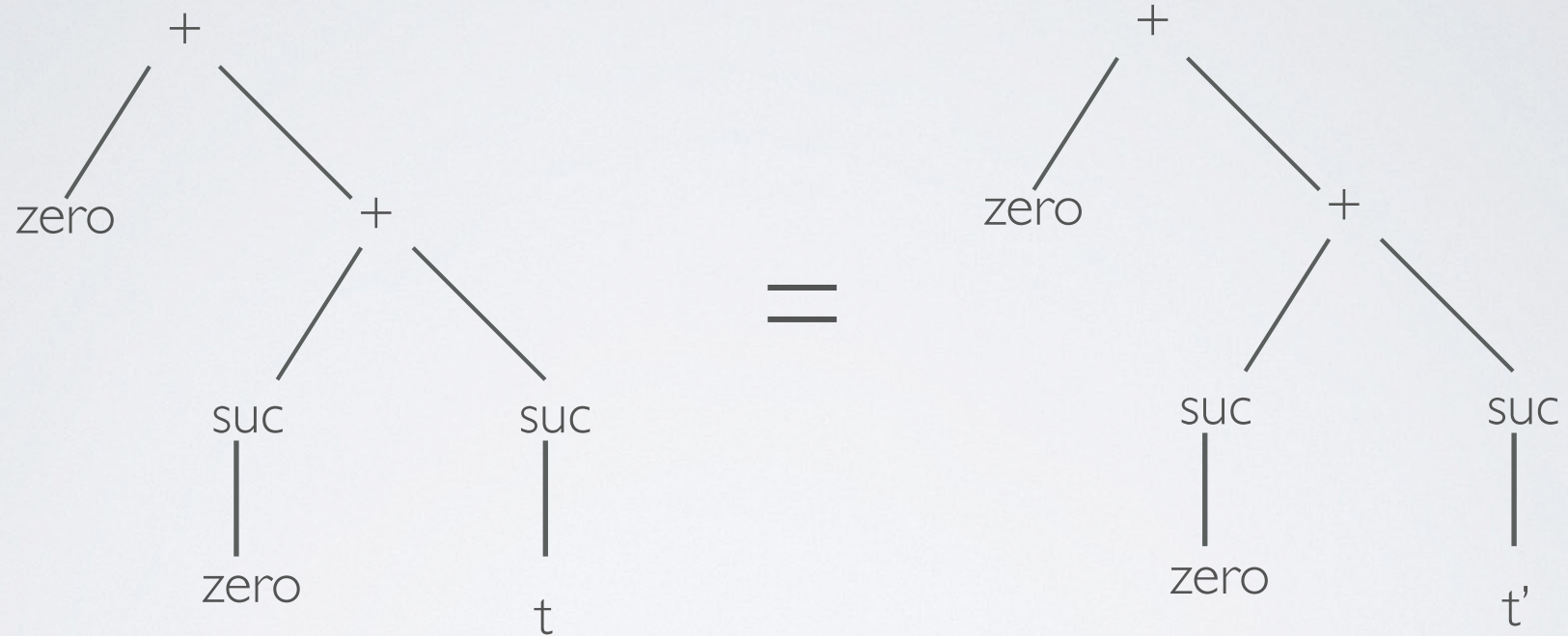


$c[_] = +(zero, +(suc(zero), suc(_)))$

SYNTACTIC DERIVATION PRINCIPLE

- $c[t] = c[t']$ iff $t = t'$, which is the same as
 - $c[t] = c[t']$ implies $t = t'$, and
 - $t = t'$ implies $c[t] = c[t']$
- Example: if $\text{not}(\text{true}) = \text{false}$ is the equation, and $c[_] = \text{not}(_)$ the context, we can **deduce by syntactic derivation that**:
 - $\text{not}(\text{true}) = \text{false} \vdash \text{not}(\text{not}(\text{true})) = \text{not}(\text{false})$

GRAPHICALLY



if and only if, $t = t'$

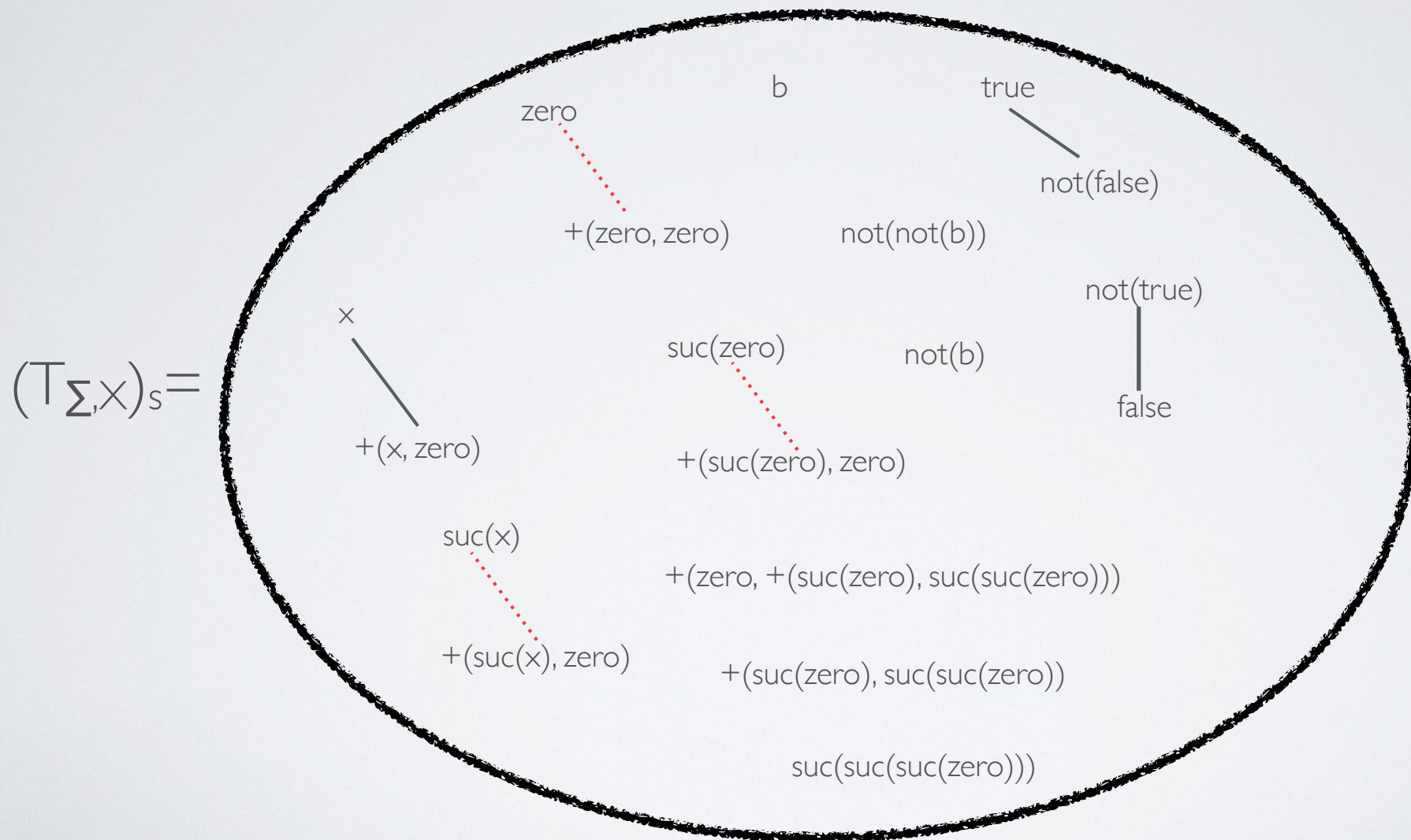
REPLACEMENT

- The operation that takes a term $c[t]$ and a new term t' and returns a $c[t']$ is called a **replacement**.

$$\Sigma = \langle S, F \rangle \quad S = \{\text{nat}, \text{bool}\} \quad X_s = \{x_{\text{nat}}, b_{\text{bool}}\}$$

$$F = \{\text{true}_{\varepsilon, \text{bool}}, \text{false}_{\varepsilon, \text{bool}}, \text{zero}_{\varepsilon, \text{nat}}, \text{suc}_{\text{nat}, \text{nat}}, +_{\text{nat}, \text{nat}, \text{nat}}, \text{not}_{\text{bool}, \text{bool}}\}$$

$$\text{Axioms} = \{\text{not}(\text{true}) = \text{false}, \text{not}(\text{false}) = \text{true}, +(x, \text{zero}) = x\}$$



Systèmes axiomatiques:

$$\textcircled{1} f(\$a) = g(\$a)$$

$$\textcircled{2} g(\$a) = 0$$

Systèmes de réécriture:

$$\textcircled{1} f(\$a) \rightarrow g(\$a)$$

$$\textcircled{2} g(\$a) \rightarrow 0$$

Orientation!

Definition (Rewrite step)

Let $\Sigma = \langle S, F \rangle$ be a signature and X be a S -sorted set of variables and $l \rightsquigarrow r, l, r \in T_\Sigma(X)$ a rewrite rule.

- $\text{filter}(t, l) = \langle \sigma, c \rangle \Leftrightarrow \exists \sigma \in S \exists c, c[\sigma l] = t^a$
- $t' = c[\sigma r]$
- $\langle t, t' \rangle \in \text{Rew}_{l \rightsquigarrow r}$ a rewrite step

a. $c[-]$ denotes the context of a term, i.e. a term with a place holder

Rew. rule: $f(\$a) \rightsquigarrow \a

Rew. step: $g(f(0)) \rightsquigarrow g(0)$

$l \Leftrightarrow f(\$a)$ $t \Leftrightarrow g(f(0))$ $\sigma \Leftrightarrow [0/a]$

$r \Leftrightarrow \$a$ $t' \Leftrightarrow g(0)$ $c[t] \Leftrightarrow g(t)$

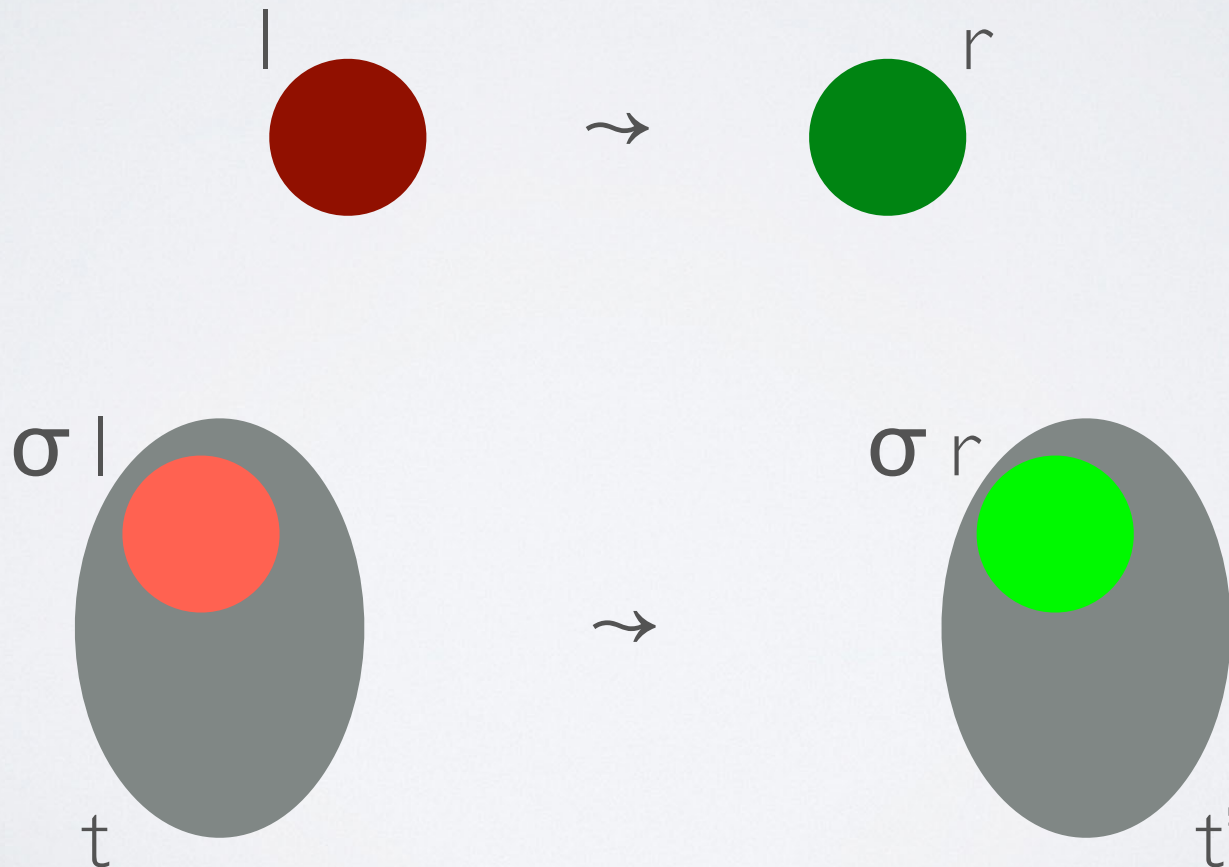
Notations non-officielles!

Definition (Rewrite step)

Let $\Sigma = \langle S, F \rangle$ be a signature and X be a S -sorted set of variables and $l \rightsquigarrow r, l, r \in T_\Sigma(X)$ a rewrite rule.

- $\text{filter}(t, l) = \langle \sigma, c \rangle \Leftrightarrow \exists \sigma \in S \exists c, c[\sigma l] = t^a$
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- $\langle t, t' \rangle \in \text{Rew}_{l \rightsquigarrow r}$ a rewrite step

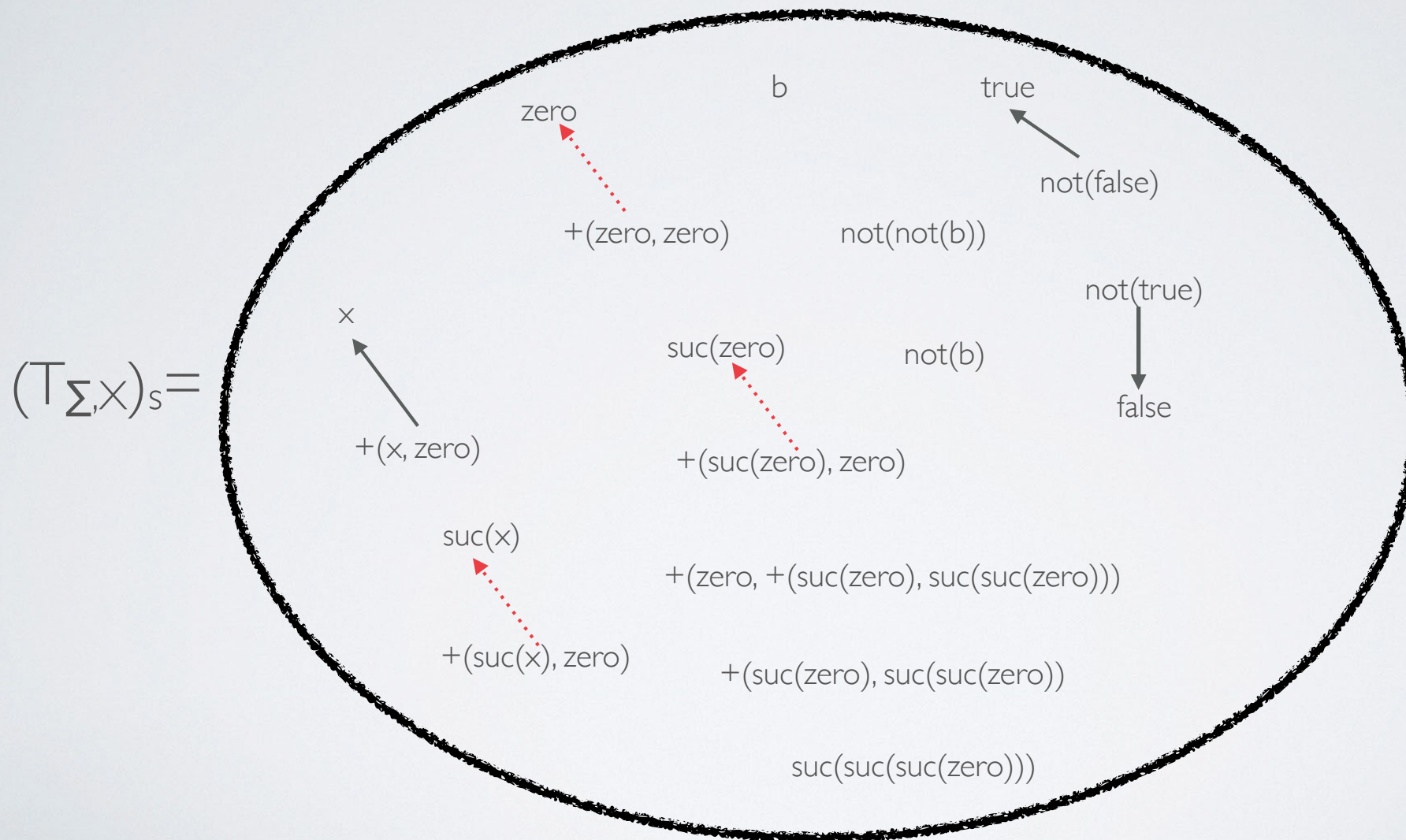
a. $c[-]$ denotes the context of a term, i.e. a term with a place holder



$$\Sigma = \langle S, F \rangle \quad S = \{\text{nat}, \text{bool}\} \quad X_s = \{x_{\text{nat}}, b_{\text{bool}}\}$$

$$F = \{\text{true}_{\varepsilon, \text{bool}}, \text{false}_{\varepsilon, \text{bool}}, \text{zero}_{\varepsilon, \text{nat}}, \text{suc}_{\text{nat}, \text{nat}}, +_{\text{nat}, \text{nat}, \text{nat}}, \text{not}_{\text{bool}, \text{bool}}\}$$

$$\text{Axioms} = \{\text{not}(\text{true}) \leadsto \text{false}, \text{not}(\text{false}) \leadsto \text{true}, +(x, \text{zero}) \leadsto x\}$$



$$\textcircled{1} \quad f(0) \rightsquigarrow s(0)$$

$$\textcircled{2} \quad g(f(\$x)) \rightsquigarrow g(\$x)$$

$$\textcircled{3} \quad h(\$x) \rightsquigarrow h(s(\$x))$$

$$g(f(0)) \overset{\textcircled{1}}{\rightsquigarrow} g(s(0))$$

$$g(f(0)) \overset{\textcircled{2}}{\rightsquigarrow} g(0)$$

Problème : confluence

$$\textcircled{1} \quad f(0) \leadsto s(0)$$

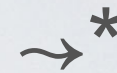
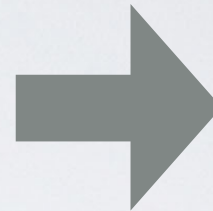
$$\textcircled{2} \quad g(f(\$x)) \leadsto g(\$x)$$

$$\textcircled{3} \quad h(\$x) \leadsto h(s(\$x))$$

$$h(0) \xrightarrow{\textcircled{3}} h(s(0)) \xrightarrow{\textcircled{3}} h(s(s(0))) \xrightarrow{\textcircled{3}} \dots$$

Problème: terminaison

Confluence / terminaison



Clôture transitive

Confluence :

Knuth-Bendix

Terminaison :



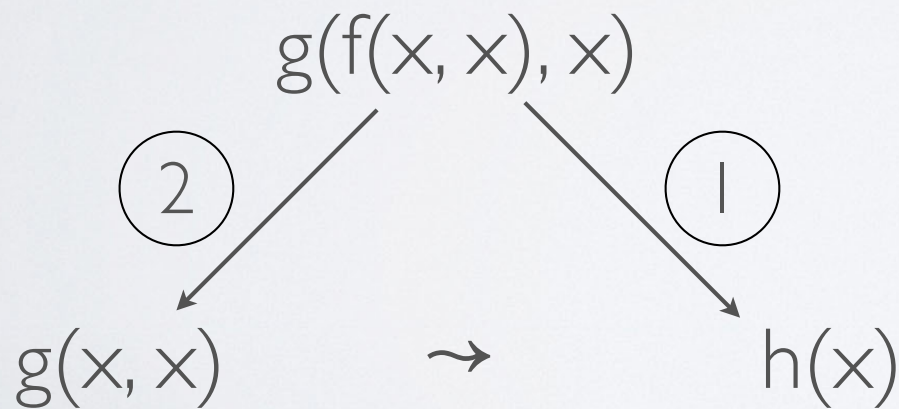
Critical Pairs - Knuth-Bendix theorem

Let $l_1 \rightsquigarrow r_1$ and $l_2 \rightsquigarrow r_2$ be two rules of a term rewriting system.
we suppose that these rules have no variables in common.

If l_1^{sub} is a subterm (and **not a variable**) of l_1 (or the term itself)
with $l_1^{context}[l_1^{sub}] = l_1$ and there exist a most general unifier σ such
that $l_1^{sub}\sigma = l_2\sigma$, then $r_1\sigma$ and $l_1^{context}[r_2\sigma]$ are called a critical pair.

$$\textcircled{1} \quad g(f(x, y), x) \rightsquigarrow h(x)$$

$$\textcircled{2} \quad f(x, x) \rightsquigarrow x$$



$$l_1^{sub} \Leftrightarrow f(x, y)$$

$$l_1^{context}[t] \Leftrightarrow g(t, x)$$

$$l_2 \Leftrightarrow f(x, x)$$

$$\sigma \Leftrightarrow [x/y]$$

$$r_1 \sigma \Leftrightarrow h(x)$$

$$l_1^{context}[r_2 \sigma] \Leftrightarrow g(x, x)$$

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$

Objectif:

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $x.e \leadsto x$
- ⑤ $x.l(x) \leadsto e$
- ⑥ $x.(y.z) \leadsto (x.y).z$
- ⑦ $l(e) \leadsto e$
- ⑧ $l(l(x)) \leadsto x$
- ⑨ $l(x.y) \leadsto l(y).l(x)$

①

$$e.x \leadsto x$$

②

$$l(x).x \leadsto e$$

③

$$(x.y).z \leadsto x.(y.z)$$

	1	2	3
1			
2			
3			

①

$$e.x \leadsto x$$

②

$$l(x).x \leadsto e$$

③

$$(x.y).z \leadsto x.(y.z)$$

	1	2	3
1			
2			
3			

①

$$e.x \leadsto x$$

②

$$l(x).x \leadsto e$$

③

$$(x.y).z \leadsto x.(y.z)$$

		1	2	3
1				
2				
3				

①

$$e.x \leadsto x$$

②

$$l(x).x \leadsto e$$

③

$$(x.y).z \leadsto x.(y.z)$$

		1	2	3
	1			
2				
3				

①

$$e.x \rightsquigarrow x$$

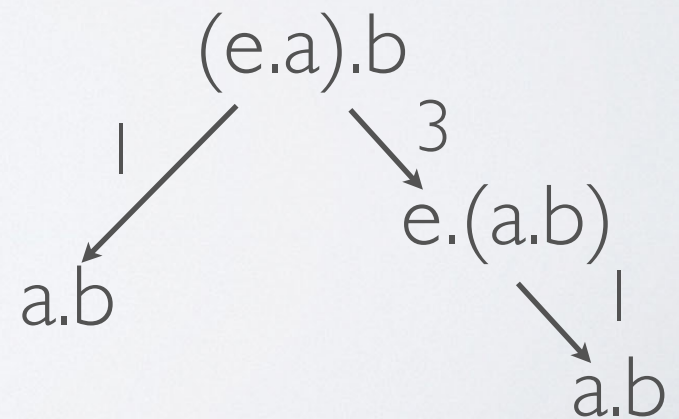
②

$$l(x).x \rightsquigarrow e$$

③

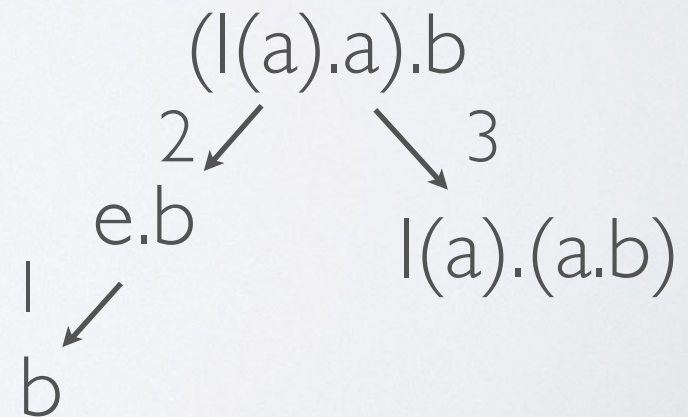
$$(x.y).z \rightsquigarrow x.(y.z)$$

		1	2	3
1				
2				
3				



- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$

		1	2	3
	1			
2				
3				



- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$

	1	2	3	4
1				
2				
3				
4				

$((a.b).c).d$

3 ↙

↘ 3

$(a.(b.c)).d$

$(a.b).(c.d)$

3 ↙

↘ 3

$a.((b.c).d)$

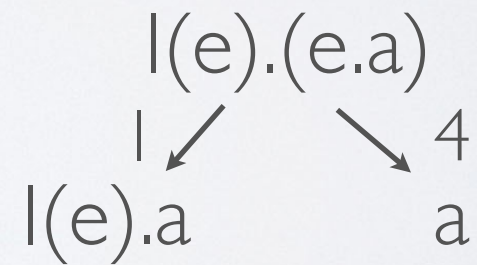
$a.(b.(c.d))$

3 ↙

$a.(b.(c.d))$

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$

		1	2	3	4
1					
2					
3					
4					



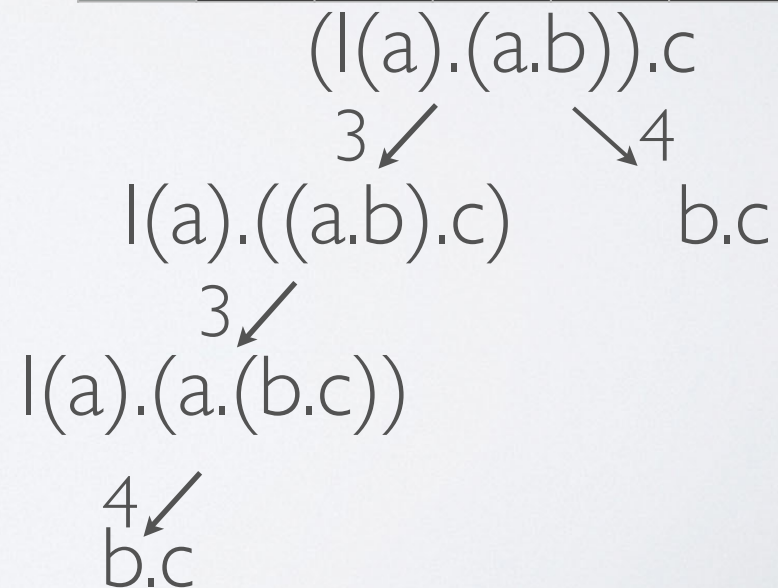
- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 6 $l(l(x)).e \leadsto x$

	1	2	3	4	5
1					
2					
3					
4					
5					

$l(l(a)).(l(a).a)$
 $\begin{array}{cc} 2 \swarrow & \searrow 4 \\ l(l(a)).e & a \end{array}$

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 6 $l(l(x)).e \leadsto x$

	1	2	3	4	5	6
1						
2						
3						
4						
5						
6						



Mêmes règles, autre substitution!

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$

	1	2	3	4	5	6
1						
2						
3						
4						
5						
6						

$l(a.b).((a.b).c)$

$\begin{matrix} 3 \swarrow & \searrow 4 \\ l(a.b).(a.(b.c)) & c \end{matrix}$

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 6 $l(l(x)).e \leadsto x$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 8 $l(l(x)).y \leadsto x.y$

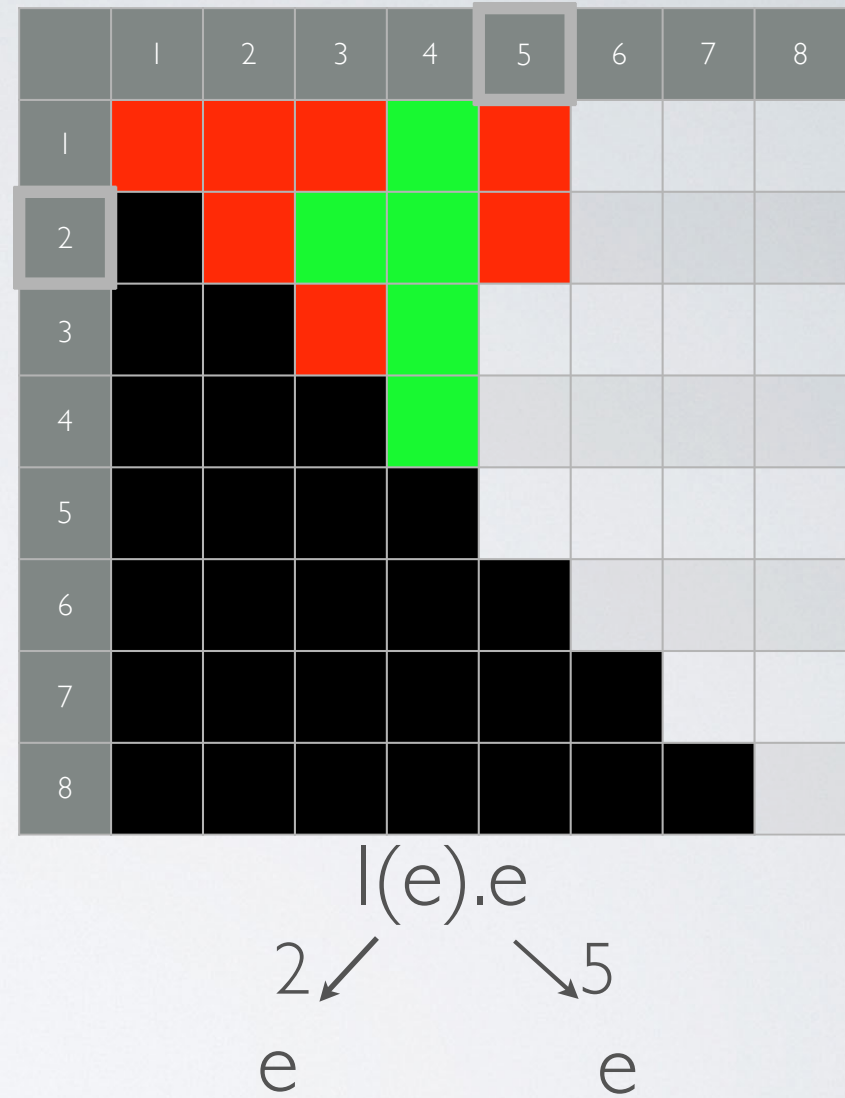
	1	2	3	4	5	6	7
1							
2							
3							
4							
5							
6							
7							

$l(l(a)).(l(a).(a.b))$
 $\swarrow 4 \quad \searrow 4$
 $l(l(a)).b \quad (a.b)$

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$
- ⑧ $l(l(x)).y \leadsto x.y$

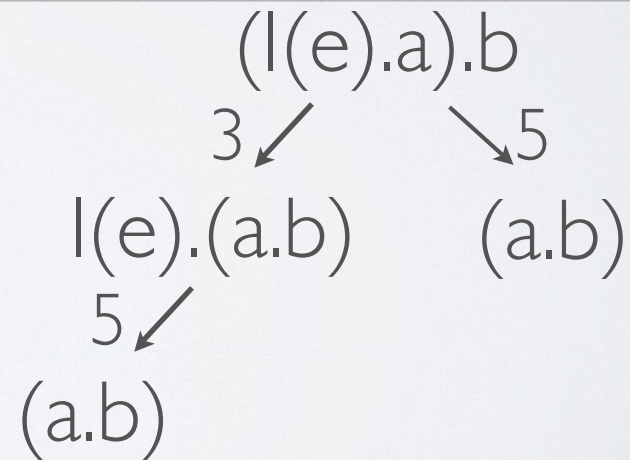
	1	2	3	4	5	6	7	8
1								
2								
3								
4								
5								
6								
7								
8								

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$
- ⑧ $l(l(x)).y \leadsto x.y$



- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$
- ⑧ $l(l(x)).y \leadsto x.y$

	1	2	3	4	5	6	7	8
1	red	red	red	green	red			
2	black	red	green	green	red			
3	black	black	red	green	red			
4	black	black	black	green				
5	black	black	black	black				
6	black	black	black	black	black			
7	black	black	black	black	black	black		
8	black	black	black	black	black	black	black	



- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 6 $l(l(x)).e \leadsto x$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 8 $l(l(x)).y \leadsto x.y$

	1	2	3	4	5	6	7	8
1	red	red	red	green	red			
2	black	red	green	green	red			
3	black	black	red	green	red			
4	black	black	black	green	red			
5	black	black	black	black				
6	black	black	black	black	black			
7	black	black	black	black	black	black		
8	black	black	black	black	black	black	black	

$l(l(e)).(l(e).a)$

4 ↙

a

↘ 5

$l(l(e)).a$

↘ 8

$e.a$

↘ 1

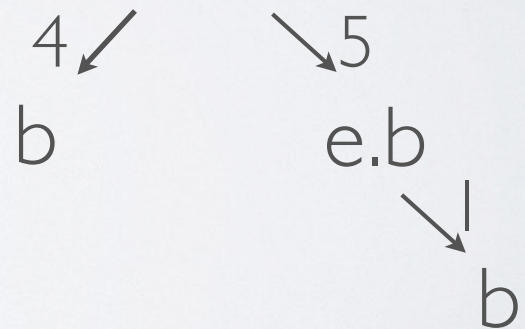
a

Mêmes règles, autre substitution!

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$
- ⑧ $l(l(x)).y \leadsto x.y$

	1	2	3	4	5	6	7	8
1								
2								
3								
4								
5								
6								
7								
8								

$l(e).(e.b)$



Trichons un peu pour accélérer le processus...

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$
- ④ $l(x).(x.z) \leadsto z$
- ⑤ $l(e).z \leadsto z$
- ⑥ $l(l(x)).e \leadsto x$
- ⑦ $l(x.y).(x.(y.z)) \leadsto z$
- ⑧ $l(l(x)).y \leadsto x.y$
- ⑨ $x.l(x) \leadsto e$

	1	2	3	4	5	6	7	8
1	red	red	red	green	red			
2	black	red	green	green	red			green
3	black	black	red	green	red			
4	black	black	black	green	red			
5	black	black	black	black	red			
6	black	black	black	black	black			
7	black	black	black	black	black	black		
8	black	black	black	black	black	black	black	

$l(l(a)).l(a)$
 $\begin{matrix} 2 \swarrow & \searrow 8 \\ e & a.l(a) \end{matrix}$

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 6 $l(l(x)).e \leadsto x$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 8 $l(l(x)).y \leadsto x.y$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$

	1	2	3	4	5	6	7	8	9
1									
2									
3									
4									
5									
6									
7									
8									
9									

$l(l(a)).e$
 $\swarrow \quad \searrow$
 6 8
 a a.e

(1)

$$e.x \leadsto x$$

(2)

$$l(x).x \leadsto e$$

(3)

$$(x.y).z \leadsto x.(y.z)$$

(4)

$$l(x).(x.z) \leadsto z$$

(5)

$$l(e).z \leadsto z$$

~~(6)~~

~~$$l(l(x)).e \leadsto x$$~~

(8) + (10)

(7)

$$l(x.y).(x.(y.z)) \leadsto z$$

(8)

$$l(l(x)).y \leadsto x.y$$

(9)

$$x.l(x) \leadsto e$$

(10)

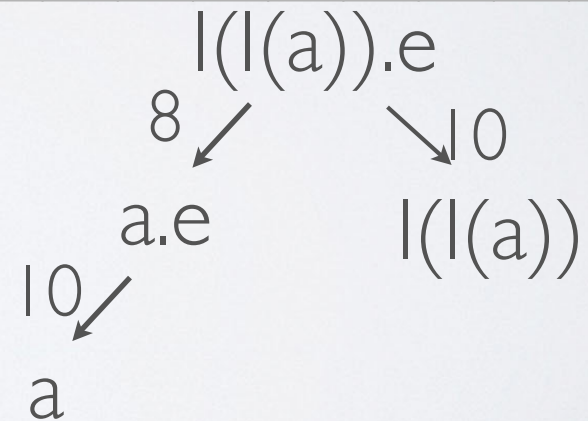
$$x.e \leadsto x$$

	1	2	3	4	5	6	7	8	9
1	red	red	red	green	red				
2	black	red	green	green	red			green	
3	black	black	red	green	red				
4	black	black	black	green	red				
5	black	black	black	black	red				
6	black	black	black	black	black			green	
7	black	black	black	black	black	black			
8	black	black	black	black	black	black	black		
9	black	black	black	black	black	black	black	black	

Useless rule!

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 8 $l(l(x)).y \leadsto x.y$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$
- 11 $l(l(x)) \leadsto x$

	1	2	3	4	5	6	7	8	9	10
1										
2										
3										
4										
5										
6										
7										
8										
9										
10										



- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 5 $l(e).z \leadsto z$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$
- 11 $l(l(x)) \leadsto x$
- 12 $l(e) \leadsto e$

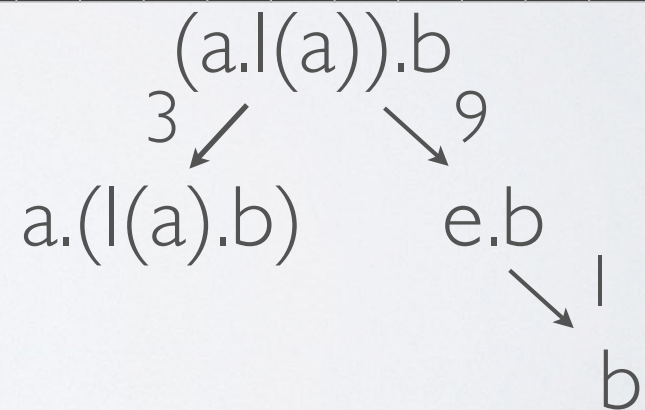
	1	2	3	4	5	6	7	8	9	10	11
1	Red	Red	Red	Green	Red						
2	Black	Red	Green	Green	Red			Green			
3	Black	Black	Red	Green	Red						
4	Black	Black	Black	Green	Red						
5	Black	Black	Black	Black	Red					Green	
6	Black	Black	Black	Black	Black			Green			
7	Black	Black	Black	Black	Black	Black					
8	Black	Black	Black	Black	Black	Black	Black			Green	
9	Black	Black	Black	Black	Black	Black	Black	Black			
10	Black	Black	Black	Black	Black	Black	Black	Black	Black		
11	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	

$l(e).e$

\swarrow 5 \searrow 10
 e $l(e)$

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$
- 11 $l(l(x)) \leadsto x$
- 12 $l(e) \leadsto e$
- 13 $y.(l(y).x) \leadsto x$

	1	2	3	4	5	6	7	8	9	10	11	12
1												
2												
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7												
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11												
12												



Mêmes règles, autre substitution!

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$
- 11 $l(l(x)) \leadsto x$
- 12 $l(e) \leadsto e$
- 13 $y.(l(y).x) \leadsto x$
- 14 $x.(y.l(x.y)) \leadsto e$

	1	2	3	4	5	6	7	8	9	10	11	12	13
1													
2													
3													
4													
5													
6													
7													
8													
9													
10													
11													
12													
13													

$(a.b).l(a.b)$
 $\begin{matrix} 3 & 9 \\ \swarrow & \searrow \end{matrix}$
 $a.(b.(l(a.b)))$ e

- 1 $e.x \leadsto x$
- 2 $l(x).x \leadsto e$
- 3 $(x.y).z \leadsto x.(y.z)$
- 4 $l(x).(x.z) \leadsto z$
- 7 $l(x.y).(x.(y.z)) \leadsto z$
- 9 $x.l(x) \leadsto e$
- 10 $x.e \leadsto x$
- 11 $l(l(x)) \leadsto x$
- 12 $l(e) \leadsto e$
- 13 $y.(l(y).x) \leadsto x$
- ~~14 $x.(y.l(x.y)) \leadsto e$~~
- 15 $x.l(y.x) \leadsto l(y)$

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	Red	Red	Red	Green	Red										
2	Black	Red	Green	Green	Red			Green							
3	Black	Black	Red	Green	Red				Blue						
4	Black	Black	Black	Green	Red									Green	
5	Black	Black	Black	Black	Red					Green					
6	Black	Black	Black	Black	Black			Green							
7	Black	Black	Black	Black	Black	Black									
8	Black	Black	Black	Black	Black	Black	Black			Green					
9	Black	Black	Black	Black	Black	Black	Black	Black							
10	Black	Black	Black	Black	Black	Black	Black	Black	Black						
11	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black					
12	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black				
13	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black			
14	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black		
15	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	Black	

$$(15) + (9)$$

$$e.x \leadsto x$$

$$l(x).x \leadsto e$$

$$(x.y).z \leadsto x.(y.z)$$

$$l(x).(x.z) \leadsto z$$

$$x.l(x) \leadsto e$$

$$x.e \leadsto x$$

$$l(l(x)) \leadsto x$$

$$l(e) \leadsto e$$

$$y.(l(y).x) \leadsto x$$

$$l(x.y) \leadsto l(y).l(x)$$

Objectif:

$$① \quad e.x \leadsto x$$

$$② \quad l(x).x \leadsto e$$

$$③ \quad (x.y).z \leadsto x.(y.z)$$

$$④ \quad x.e \leadsto x$$

$$⑤ \quad x.l(x) \leadsto e$$

$$⑥ \quad x.(y.z) \leadsto (x.y).z$$

$$⑦ \quad l(e) \leadsto e$$

$$⑧ \quad l(l(x)) \leadsto x$$

$$⑨ \quad l(x.y) \leadsto l(y).l(x)$$

Objectif:

- ① $e.x \leadsto x$
- ② $l(x).x \leadsto e$
- ③ $(x.y).z \leadsto x.(y.z)$

$l(x).(x.z) \leadsto z$ ❌

⑤ $x.l(x) \leadsto e$

④ $x.e \leadsto x$

⑧ $l(l(x)) \leadsto x$

⑦ $l(e) \leadsto e$

$y.(l(y).x) \leadsto x$ ❌

⑨ $l(x.y) \leadsto l(y).l(x)$

① $e.x \leadsto x$

② $l(x).x \leadsto e$

③ $(x.y).z \leadsto x.(y.z)$

④ $x.e \leadsto x$

⑤ $x.l(x) \leadsto e$

⑥ $x.(y.z) \leadsto (x.y).z$ ❌

⑦ $l(e) \leadsto e$

⑧ $l(l(x)) \leadsto x$

⑨ $l(x.y) \leadsto l(y).l(x)$