Make FunBlocks alive

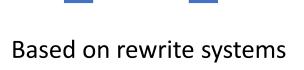
Marvin FOURASTIE

Master project

Motivations

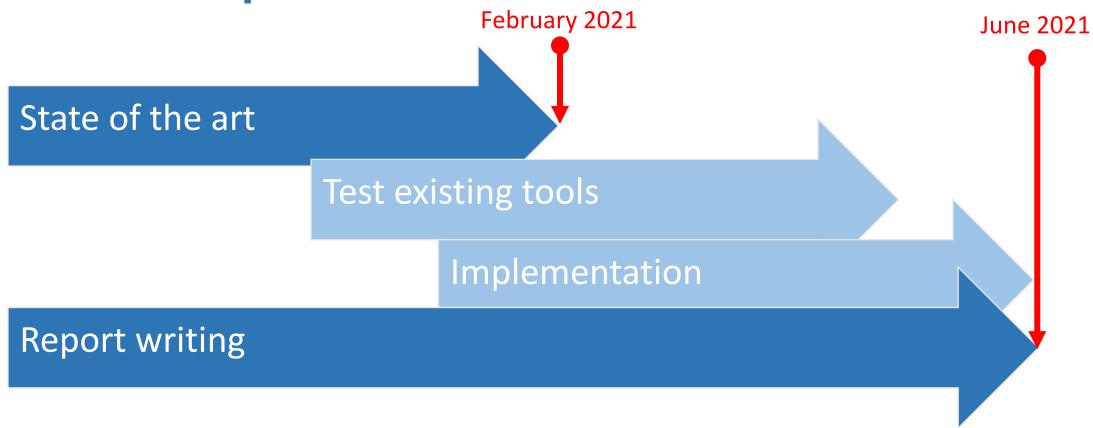






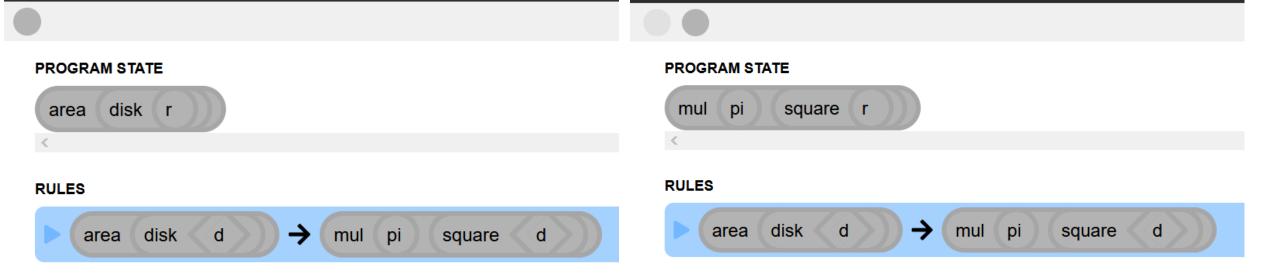
FunBlocks

Road Map



FunBlocks

```
init area(disk(r))
case area(disk($d$)) => mul(pi,square($d$))
```



FunBlocks

area(disk(r)) Declarative paradigm case area(disk(\$d)) => mul(pi,square(\$d)) Based on rewrite systems square Static typing node (Tree \$t) (Tree \$t) type Tree \$t :: empty | leaf \$t |

Goals

Provide users with valuable insights about their program

Verification of rewrite systems

Rewrite systems

Stack operators

 $Zero = \{0\}$

Nat = Zero U succ(Nat)

Empty = Λ

Stack = Empty U push(Nat, Stack)

top : Stack \rightarrow Nat

pop : Stack → Stack

alternate : Stack \times Stack \rightarrow Stack

Rewrite systems

Canonical rewrite system

$$top(push(x, y)) = x$$

$$top(push(x, y)) \rightarrow x$$

$$pop(push(x, y)) \Rightarrow y$$

$$alternate(\Lambda, z) = z$$

$$alternate(push(x, y), z) = push(x, alternate(z, y))$$

$$alternate(push(x, y), z) \rightarrow push(x, alternate(z, y))$$

Rewrite systems

Termination

Confluence

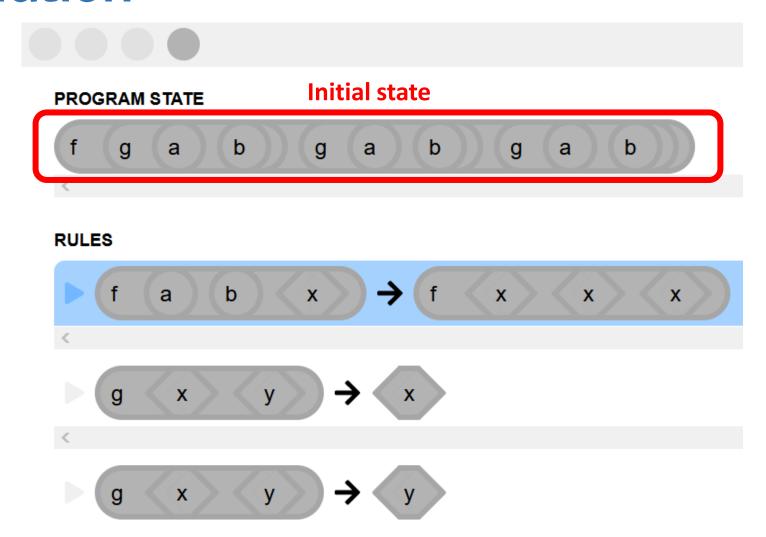
Soundness

Completeness

Correctness

Undecidable in general

Termination



Reduction order

Monotone

$$s_i > t \rightarrow f(s_1, ..., s_i, ..., s_n) > f(s_1, ..., t, ..., s_n)$$

For f of arity n

Close under substitution ————

$$s>t\to\sigma s>\sigma t$$
 , for all substitution σ

Well-founded

no infinite descending chain

 $(\mathbb{N}, <)$ is well-founded

 $(\mathbb{Z}, <)$ is not well-founded

Termination

A term rewriting system is terminating

if and only if

it admits a compatible reduction order < (if l > r for every rewrite rule $l \rightarrow r$)



Verification of termination

Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(a) = 1$$

$$w(g(t)) = 1 + w(t)$$

$$w(f(t_1,t_2)) = 2w(t_1) + w(t_2)$$

Polynomial interpretation

$$f(a,x) \to x$$

$$f(g(x),y) \to g(f(x,y))$$

$$w(f(a,x)) = 2 + w(x)$$

$$w(f(g(x),y)) = 2 + 2w(x) + w(y)$$

$$w(g(f(x,y)) = 1 + 2w(x) + w(y)$$

$$w(f(a,x)) > w(x)$$

$$w(f(g(x),y)) > w(g(f(x,y))$$

Reduction order → Termination

Algorithms

Recursive Path Ordering

Order based on the mutisets

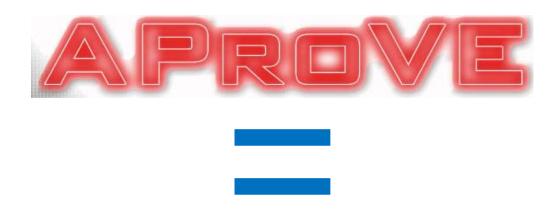
Knuth-Bendix Ordering

Based on weights assigned to operators

Dependency pairs

Prove innermost termination

Termination

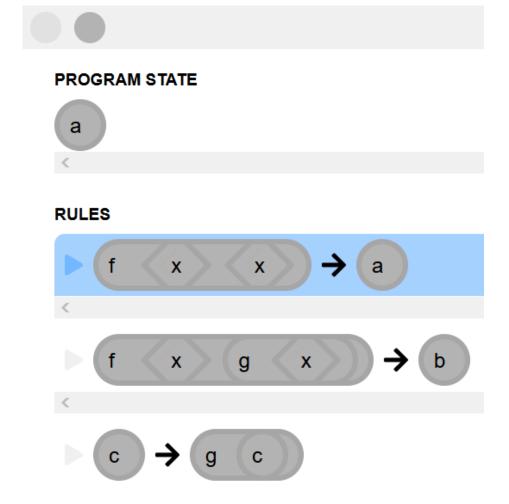


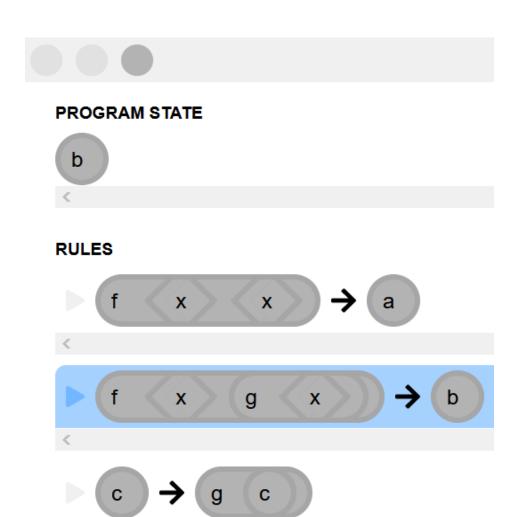
Direct proof (polynomial, LBO, KBO,...)



Dependency pairs and size-change principle

Confluence

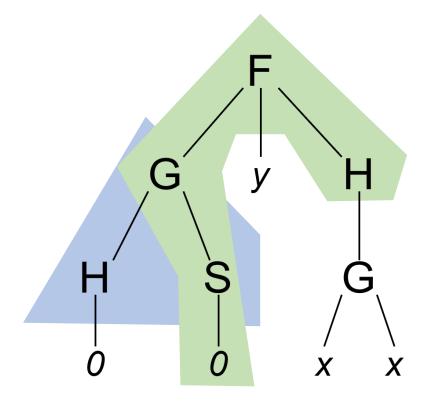




Overlap and critical pairs

$$\rho_1 : F\left(G(x, S(0)), y, H(z)\right) \to x$$

$$\rho_2 : G(H(x), S(y)) \to y$$



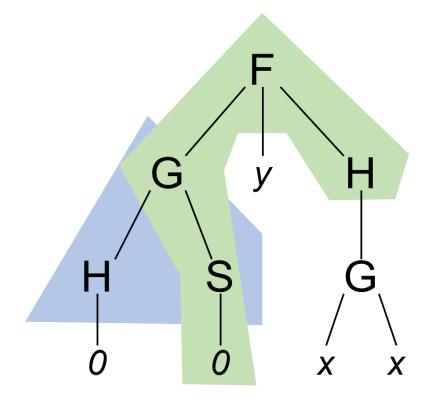
Overlap and critical pairs

Overlapping:

Term:
$$F(G(H(0),S(0)),y,H(z))$$

$$F(G(\square,S(0)),\square,H(\square))$$

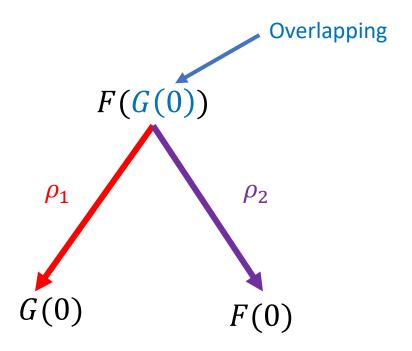
$$G(H(\square),S(\square))$$



Overlap and critical pairs

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



< G(x), F(x) > is called critical pair

Critical Pair Lemma

A terminating rewriting system is confluent

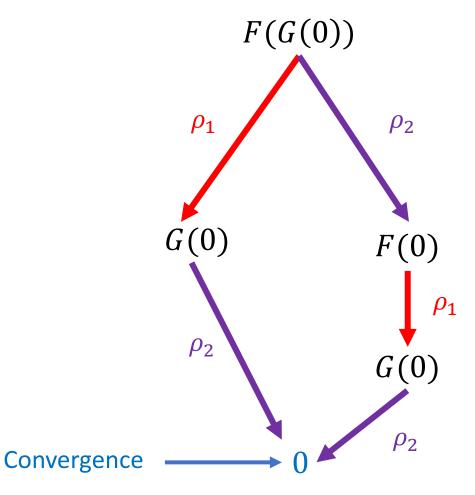
if and only if

all critical pairs are convergent

Critical Pair Lemma

$$\rho_1: F(x) \to G(0)$$

$$\rho_2:G(x)\to 0$$



Input:

A set of equation

A reduction ordering <

$$\begin{array}{ll} 1 \cdot x & = x \\ x^{-1} \cdot x & = 1 \\ (x \cdot y) \cdot z & = x \cdot (y \cdot z) \end{array}$$

Non-confluent

Output:

Terminate successfully

Terminating and confluent rewrite system

Non-terminating rewrite system

Rule which cannot be ordered (i.e. commutative operator)

Basic rules:

Orienting _____

Transform s = t to $s \rightarrow t$

Adding

 \longrightarrow

Add s = t in the set of equation

Simplifying

 \longrightarrow

Simplify s = t in s' = t'

Deleting



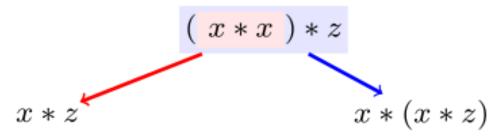
Delete trivial rules s = s

Adding

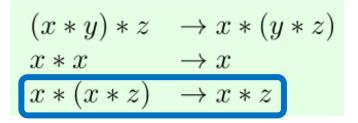


Add s = t in the set of equation

$$\begin{array}{ccc} (x*y)*z & \rightarrow x*(y*z) \\ x*x & \rightarrow x \end{array}$$







New rule added

Completion process:

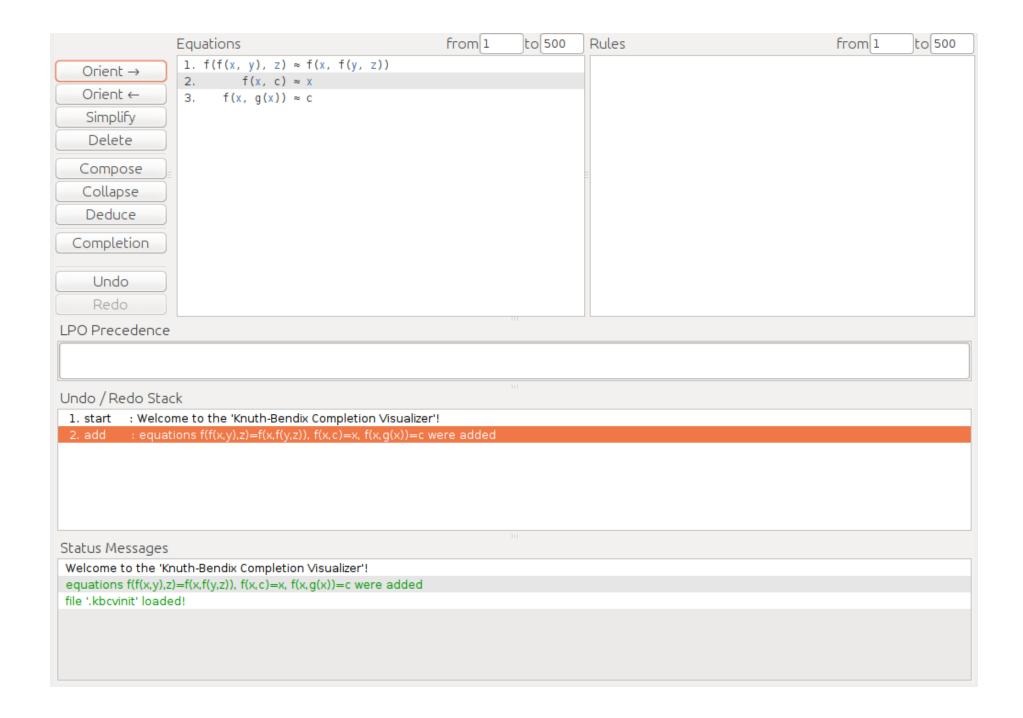
- 1. For each equation s = t reduce s and t to normal form s' and t'
- 2. Fill the set of rules using basic operators and reduction ordering
- 3. If the algorithm terminate successfully: terminating and confluent rewrite system

Completion for axioms of groups:

$$\begin{array}{rcl}
1 \cdot x & = x \\
x^{-1} \cdot x & = 1 \\
(x \cdot y) \cdot z & = x \cdot (y \cdot z)
\end{array}$$



$$\begin{array}{cccc}
1 \cdot x & \rightarrow x \\
x^{-1} \cdot x & \rightarrow 1 \\
(x \cdot y) \cdot z & \rightarrow x \cdot (y \cdot z) \\
x^{-1} \cdot (x \cdot y) & \rightarrow y \\
1^{-1} & \rightarrow 1 \\
x \cdot 1 & \rightarrow x \\
(x^{-1})^{-1} & \rightarrow x \\
x \cdot x^{-1} & \rightarrow 1 \\
x \cdot (x^{-1} \cdot y) & \rightarrow y \\
(x \cdot y)^{-1} & \rightarrow y^{-1}x^{-1}
\end{array}$$



Existing tools







Simplicity

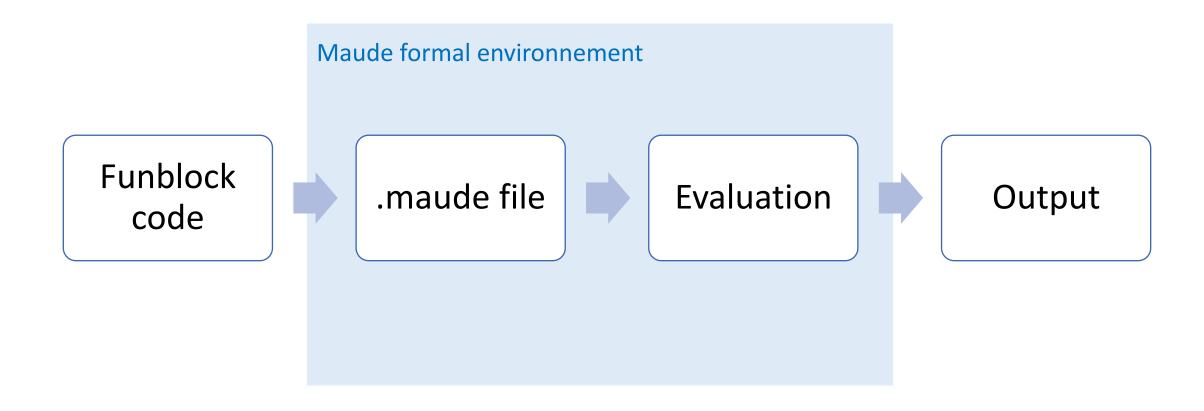


Expressiveness



Performance

Procedure



```
fmod BASIC-NAT is
        sort Nat .
       op 0 : -> Nat .
       op s : Nat -> Nat .
       op + : Nat Nat -> Nat .
        vars N M : Nat .
        eq 0 + N = N.
       eq s(M) + N = s(M + N).
endfm
```

```
fmod FACTORIAL is
protecting NAT .

op _! : Nat -> NzNat .
var N : Nat .

eq 0 ! = 1 .
eq (s N) ! = (s N) * N ! .

endfm
```

```
> load factorial.maude
> red 100 ! .
Reduce in FACTORIAL : 100 ! .
rewrites: 201 in Oms cpu (Oms real) (~ rewrites/second)
result NzNAT:
9332621544394415268169923885626670049071596826438162146
8592963895217599993229915608941463976156518286253697920
82722375825118521091686400000000000000000000000
```

```
mod VENDING-MACHINE is
  including VENDING-MACHINE-SIGNATURE .
  var M : Marking .
  rl [add-q] : M \Rightarrow M q.
  rl [add-\$] : M => M \$.
  rl [buy-c] : $ => c .
  rl [buy-a] : $ => a q .
  rl [change] : q q q => $.
endm
```

Inductive Theorem Prover (ITP)

Sufficient Completeness Checker (SCC)

Church-Rosser Checker (CRC)

Coherence Checker (ChC)

Maude Termination Tool (MTT)

Maude Formal Environment (MFE)



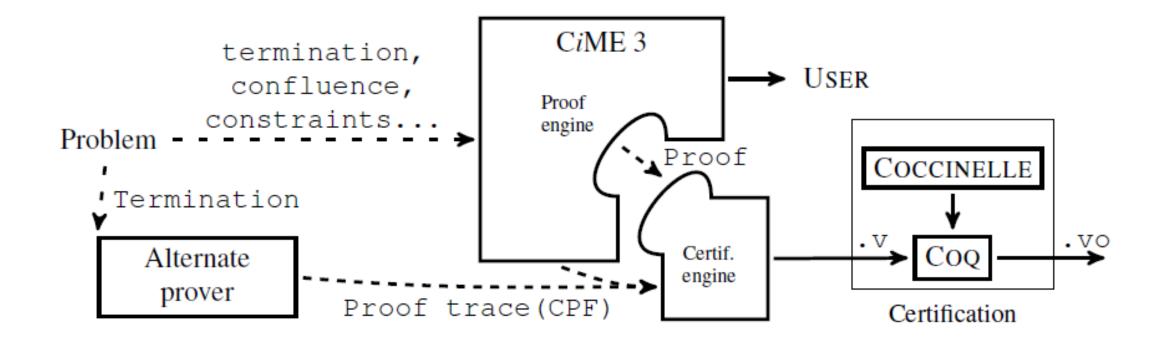


Rewriting toolkit





Proof certification



Definition of signatures

```
CiME>let X = variables "x,y,z";
X : variable_set = variables "z,x,y"
```

Definition of algebra and terms

```
CiME> let A_peano = algebra F_peano ;
A_peano : F_peano algebra = algebra F_peano
```

```
CiME> let t = term A_peano "s(s(s(0)))*(s(0)+s(s(0)))";
t : F_peano term = s(s(s(0))) *(s(0)+s(s(0)))
```

Term rewriting system

```
CiME> let R_peano = trs A_peano "
      x+0 -> x;
      x+s(y) \rightarrow s(x+y);
      x*0 -> 0;
      x*s(y) -> (x*y) +x;
       ";
 R_peano : F_peano trs = trs A_peano "
            x+0 -> x;
            x+s(y) \rightarrow s(x+y);
            x * 0 -> 0;
            x * s (y) -> (x * y) + x "
```

```
CiME> termination R_peano;
CiME> coq_certify_proof R_peano;
CiME> convergence R_peano ;
...
```

What's next?



Integration of the tools



Test the existing rewrite system checker

References

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http://rewriting.loria.fr/systems.html

Knuth-Bendix Completion Visualizer

http://cl-informatik.uibk.ac.at/software/kbcv/

Knuth-Bendix Completion subject based thesis

• https://homepage.divms.uiowa.edu/~astump/papers/thesis-wehrman.pdf

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Prolog implementation of the Knuth-Bendix completion procedure

https://www.metalevel.at/trs/

Maude tools

- http://maude.lcc.uma.es/CRChC/
- http://www.lcc.uma.es/%7Eduran/MTT/
- http://maude.sip.ucm.es/debugging/

Wikipedia

- https://fr.wikipedia.org/wiki/Compl%C3%A9tion_de_Knuth-Bendix
- https://fr.wikipedia.org/wiki/Paire_critique

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