Implementation of the TIL-Script Language

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Project Goals

- ▶ Define the TIL-Script programming language
 - Build upon the current grammar
 - Define the semantics of the language
- Create a working TIL-Script interpreter
- Document the language properly

Transparent Intensional Logic

- Logical analysis of natural language
- Based on typed Lambda calculi
- Rigorously defined type hierarchy
- Procedural and hyperintensional
 - Constructions can mention other constructions
- Sentence meaning is carried by a procedure
 - We mostly care for the procedure itself, seldom do we care for the value it produces

TIL-Script

- Grammar closely resembles TIL grammar and conventions
- Adds upon TIL to form a useable programming language
 - Lists
 - Tuples
 - Structures
 - Strings
- Also adds restrictions imposed by computers with finite resources

Original Features

- ► Grammar
- ► All TIL constructions
- ► Type aliases
- Lists, Tuples
- Semantics not fully defined

New Features

- Comments
- ► Imports
- ► Nil value
- Distinction between definitions and declarations
- String type
- ► Tuples are now heterogenous
- Structs (user defined types)
- Types as language objects

Nil Value

- Used to represent a state when no real value was produced
- Internally carries a position in the source code where the value was produced
- ▶ Also carries the reason as to why a construction was improper

Declarations And Definitions

- Declaration only specifies a name and a type
 - ▶ I.e. if we only know, or only care for, the name
 - ► Halting problem we may want to refer to it, but we can't solve it
 - ► Function Halts/(Bool Program) cannot be implemented
- Definition also specifies semantics or value

Tuples And Lists

- Originally
 - Tuples were homogenous and finite
 - Lists were homogenous and infinite
- Now
 - Tuples are heterogenous and fixed in size
 - Lists are homogenous, size is unlimited

Structs

- User defined compound types
- Cannot be coerced into other types with same layout
- ▶ Allows for better use of the TIL-Script type system
 - ► E.g. cartesian vs. polar coordinates

Code Example

```
numbers \rightarrow List(Int) := ['ListOf '1 '6 '2 '5 '3 '4].
Defn LessThanFive(num: Int) \rightarrow Bool := ['< num '5].
Defn Filter(list: List(Int), pred: (Bool Int))
        -> List(Int) :=
    ['If ['IsEmpty list]
        list
        ['If [pred ['Head list]]
             [ 'Cons
                 ['Head list]
                 ['Filter ['Tail list] pred]]
             ['Filter ['Tail list] pred]]].
['Println ['Filter numbers 'LessThanFive]].
```

Implementation

- ► Kotlin
 - Algebraic data types
 - Null safety
 - Immutable interfaces for collections
- Gradle
- ► Antlr

Project Structure

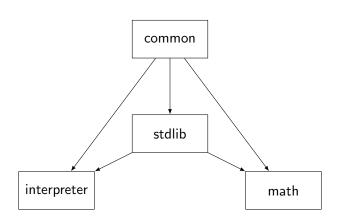


Figure: Project Structure

Interpreter Implementation

- Written in Kotlin
- Parser is autogenerated
- Functions can be written in TIL-Script or JVM languages
 - ► Allows for access to the entire JVM ecosystem
 - I.e. database libraries don't have to be implemented from scratch
 - ▶ JDBC can be used
 - No need to handle syscalls
 - No need to reimplement math functions

Interpreter Implementation

- Object oriented design
 - Not the best choice for an interpreter
 - Simplifies dealing with functions of different implementations
 - ► The interpreter can be replaced with a better, drop-in replacement
 - Limited choice on the JVM
- Functional approach where possible
- Immutability is preferred

Current State

- Working prototype
- All features work as intended
- Lists, tuples, user-defined structures, mutually recursive functions, variables, imports, etc.
- Type coherency checking

Limitations

- Working prototype
- No bytecode
 - Negative performance implications
 - Callstack is tied to the JVM callstack
 - Bytecode conversion must be bijective
- No TCO
 - Unfortunate, but cannot be done without bytecode

Possible Improvements

- ► REPL
- ► Editor support and an LSP
- Compilation to bytecode

The End

Questions

- ▶ What is a degenerate function?
 - ▶ A function which is undefined on the entirety of its domain
 - $\lambda x \cdot x \div 0$
 - $\triangleright \lambda x . log_x 0$