1 Match Constraint Language

 $\overline{\top} : \tau$ (1a)

CTUnknown

$$\overline{?:\tau}$$
 (1b)

 $\frac{\text{CTNum}}{\underline{n}: \text{num}} \tag{1c}$

CTInl $\frac{\dot{\xi}_1 : \tau_1}{\operatorname{inl}(\dot{\xi}_1) : (\tau_1 + \tau_2)}$ (1d)

CTInr $\frac{\dot{\xi}_2 : \tau_2}{\operatorname{inr}(\dot{\xi}_2) : (\tau_1 + \tau_2)} \tag{1e}$

CTPair $\frac{\dot{\xi}_1 : \tau_1 \qquad \dot{\xi}_2 : \tau_2}{(\dot{\xi}_1, \dot{\xi}_2) : (\tau_1 \times \tau_2)}$ (1f)

CTOr $\frac{\dot{\xi}_1 : \tau \qquad \dot{\xi}_2 : \tau}{\dot{\xi}_1 \lor \dot{\xi}_2 : \tau}$ (1g)

 $|\dot{\xi}|$ refutable? $|\dot{\xi}|$ is refutable

RXNum

$$\frac{}{\underline{n} \; \mathtt{refutable}_?}$$
 (2a)

RXUnknown

RXInl

$$\frac{}{\operatorname{inl}(\dot{\xi})\operatorname{refutable}_?}$$
 (2c)

RXInr

$$\frac{}{\operatorname{inr}(\dot{\xi})\operatorname{refutable}_?}$$
 (2d)

RXPairL

$$\frac{\dot{\xi}_1 \text{ refutable}?}{(\dot{\xi}_1, \dot{\xi}_2) \text{ refutable}?} \tag{2e}$$

RXPairR
$$\frac{\dot{\xi}_2 \text{ refutable}_?}{(\dot{\xi}_1, \dot{\xi}_2) \text{ refutable}_?}$$
(2f)

$$\frac{\text{RXOr}}{\dot{\xi}_1 \text{ refutable}_?} \frac{\dot{\xi}_2 \text{ refutable}_?}{\dot{\xi}_1 \vee \dot{\xi}_2 \text{ refutable}_?} \tag{2g}$$

$\mathit{refutable}_?(\dot{\xi})$

$$refutable_{?}(\underline{n}) = true$$
 (3a)

$$refutable_{?}(?) = true$$
 (3b)

$$refutable_?(\mathtt{inl}(\dot{\xi})) = refutable_?(\dot{\xi}) \tag{3c}$$

$$refutable_{?}(inr(\dot{\xi})) = refutable_{?}(\dot{\xi})$$
 (3d)

$$\mathit{refutable}_?((\dot{\xi}_1,\dot{\xi}_2)) = \mathit{refutable}_?(\dot{\xi}_1) \ \text{or} \ \mathit{refutable}_?(\dot{\xi}_2) \tag{3e}$$

$$refutable_{?}(\dot{\xi}_{1} \vee \dot{\xi}_{2}) = refutable_{?}(\dot{\xi}_{1}) \text{ and } refutable_{?}(\dot{\xi}_{2})$$
 (3f)

Otherwise
$$refutable_?(\dot{\xi}) = \text{false}$$
 (3g)

Lemma 1.0.1 (Soundness and Completeness of Refutable Constraints). ξ refutable? iff $refutable_?(\dot{\xi}) = true$.

 $e \models \dot{\xi}$ e satisfies $\dot{\xi}$

CSTruth
$$\frac{e}{\vdash \top} \tag{4a}$$

$$\underline{n \models n} \tag{4b}$$

CSInl

$$\frac{e_1 \models \dot{\xi}_1}{\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)} \tag{4c}$$

CSInr

$$\frac{e_2 \models \dot{\xi}_2}{\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)} \tag{4d}$$

CSPair
$$\frac{e_1 \models \dot{\xi}_1 \qquad e_2 \models \dot{\xi}_2}{(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)}$$
(4e)

CSNotIntroPair

$$\frac{e \text{ notintro}}{e \mid e \mid (\dot{\xi}_1, \dot{\xi}_2)} \text{prr}(e) \mid = \dot{\xi}_2$$

$$e \mid (\dot{\xi}_1, \dot{\xi}_2)$$

$$(4f)$$

CSOrL
$$\frac{e \models \dot{\xi}_1}{e \models \dot{\xi}_1 \lor \dot{\xi}_2} \tag{4g}$$

CSOrR
$$\frac{e \models \dot{\xi}_2}{e \models \dot{\xi}_1 \lor \dot{\xi}_2} \tag{4h}$$

 $\mathit{satisfy}(e,\dot{\xi})$

$$satisfy(e, \top) = true$$
 (5a)

$$satisfy(n_1, n_2) = (n_1 = n_2)$$
 (5b)

$$\mathit{satisfy}(e,\dot{\xi}_1 \lor \dot{\xi}_2) = \mathit{satisfy}(e,\dot{\xi}_1) \text{ or } \mathit{satisfy}(e,\dot{\xi}_2) \tag{5c}$$

$$\mathit{satisfy}(\mathtt{inl}_{\tau_2}(e_1),\mathtt{inl}(\dot{\xi}_1)) = \mathit{satisfy}(e_1,\dot{\xi}_1) \tag{5d}$$

$$\mathit{satisfy}(\mathsf{inr}_{\tau_1}(e_2),\mathsf{inr}(\dot{\xi}_2)) = \mathit{satisfy}(e_2,\dot{\xi}_2) \tag{5e}$$

$$\mathit{satisfy}((e_1,e_2),(\dot{\xi}_1,\dot{\xi}_2)) = \mathit{satisfy}(e_1,\dot{\xi}_1) \text{ and } \mathit{satisfy}(e_2,\dot{\xi}_2) \tag{5f}$$

$$\mathit{satisfy}(())^u, (\dot{\xi}_1, \dot{\xi}_2)) = \mathit{satisfy}(\mathtt{prl}(())^u), \dot{\xi}_1) \text{ and } \mathit{satisfy}(\mathtt{prr}(())^u), \dot{\xi}_2) \tag{5g}$$

$$\mathit{satisfy}((\!(e)\!)^u, (\dot{\xi}_1, \dot{\xi}_2)) = \mathit{satisfy}(\mathsf{prl}((\!(e)\!)^u), \dot{\xi}_1) \text{ and } \mathit{satisfy}(\mathsf{prr}((\!(e)\!)^u), \dot{\xi}_2)$$
 (5h)

$$satisfy(e_1(e_2),(\dot{\xi_1},\dot{\xi_2})) = satisfy(\mathtt{prl}(e_1(e_2)),\dot{\xi_1})$$

and
$$satisfy(prr(e_1(e_2)), \dot{\xi}_2)$$
 (5i)

 $\mathit{satisfy}(\mathtt{match}(e)\{\hat{rs}\},(\dot{\xi}_1,\dot{\xi}_2)) = \mathit{satisfy}(\mathtt{prl}(\mathtt{match}(e)\{\hat{rs}\}),\dot{\xi}_1)$

and
$$satisfy(prr(match(e)\{\hat{rs}\}), \dot{\xi}_2)$$
 (5j)

$$\mathit{satisfy}(\mathtt{prl}(e),(\dot{\xi}_1,\dot{\xi}_2)) = \mathit{satisfy}(\mathtt{prl}(\mathtt{prl}(e)),\dot{\xi}_1)$$

and
$$satisfy(prr(prl(e)), \dot{\xi}_2)$$
 (5k)

$$\mathit{satisfy}(\mathtt{prr}(e),(\dot{\xi}_1,\dot{\xi}_2)) = \mathit{satisfy}(\mathtt{prl}(\mathtt{prr}(e)),\dot{\xi}_1)$$

and
$$satisfy(prr(prr(e)), \dot{\xi}_2)$$
 (51)

Otherwise
$$satisfy(e, \dot{\xi}) = false$$
 (5m)

Lemma 1.0.2 (Soundness and Completeness of Satisfaction Judgment). $e \models \dot{\xi}$ iff $satisfy(e, \dot{\xi}) = true$.

Proof. We prove soundness and completeness separately.

1. Soundness:

(1)
$$e \models \dot{\xi}$$
 by assumption

By rule induction over Rules (15) on (1).

Case (15a).

- (2) $\dot{\xi} = \top$ by assumption (3) $satisfy(e, \top) = \text{true}$ by Definition 16a
- Case (15b).
 - (2) e = n by assumption
 - (3) $\dot{\xi} = \underline{n}$ by assumption
 - (4) $satisfy(\underline{n},\underline{n}) = (n = n) = true$ by Definition 16b

Case (15c).

- (2) $e = n_1$ by assumption
- (3) $\dot{\xi} = \underline{p_2}$ by assumption
- (4) $n_1 \neq n_2$ by assumption
- (5) $satisfy(n_1, p_2) = (n_1 \neq n_2) = true$ by Definition 16c on (4)

Case (15d).

- (2) $\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$ by assumption
- (3) $e \models \dot{\xi}_1$ by assumption
- (4) $e \models \dot{\xi}_2$ by assumption
- (5) $satisfy(e, \dot{\xi}_1) = true$ by IH on (3)
- (6) $satisfy(e, \dot{\xi}_2) = true$ by IH on (4)
- (7) $satisfy(e, \dot{\xi}_1 \land \dot{\xi}_2) = satisfy(e, \dot{\xi}_1)$ and $satisfy(e, \dot{\xi}_2) = true$ by Definition 16d on (5) and (6)

Case (15e).

- (2) $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$ by assumption
- (3) $e \models \dot{\xi}_1$ by assumption
- (4) $satisfy(e, \dot{\xi}_1) = true$ by IH on (3)
- (5) $satisfy(e, \dot{\xi}_1 \lor \dot{\xi}_2) = satisfy(e, \dot{\xi}_1)$ or $satisfy(e, \dot{\xi}_2) = \text{true}$ by Definition 16e on (4)

Case (15f).

- (2) $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$ by assumption
- (3) $e \models \dot{\xi}_2$ by assumption
- (4) $satisfy(e, \dot{\xi}_2) = true$ by IH on (3)
- (5) $satisfy(e, \dot{\xi}_1 \lor \dot{\xi}_2) = satisfy(e, \dot{\xi}_1)$ or $satisfy(e, \dot{\xi}_2) = true$ by Definition 16e on (4)

Case (15g).

 $\begin{array}{lll} (2) & e=\mathtt{inl}_{\tau_2}(e_1) & \text{by assumption} \\ (3) & \dot{\xi}=\mathtt{inl}(\dot{\xi}_1) & \text{by assumption} \\ (4) & e_1 \models \dot{\xi}_1 & \text{by assumption} \\ (5) & \textit{satisfy}(e_1,\dot{\xi}_1) = \text{true} & \text{by IH on (4)} \\ (6) & \textit{satisfy}(\mathtt{inl}_{\tau_2}(e_1),\mathtt{inl}(\dot{\xi}_1)) = \textit{satisfy}(e_1,\dot{\xi}_1) = \text{true} \\ & & \text{by Definition 16f on (5)} \end{array}$

Case (15h).

 $\begin{array}{lll} (2) & e=\inf_{\tau_1}(e_2) & \text{by assumption} \\ (3) & \dot{\xi}=\inf(\dot{\xi_2}) & \text{by assumption} \\ (4) & e_2 \models \dot{\xi_2} & \text{by assumption} \\ (5) & \textit{satisfy}(e_2,\dot{\xi_2}) = \text{true} & \text{by IH on (4)} \\ (6) & \textit{satisfy}(\inf_{\tau_1}(e_2),\inf(\dot{\xi_2})) = \textit{satisfy}(e_2,\dot{\xi_2}) = \text{true} \\ & \text{by Definition 16g on} \end{array}$

(5)

(6) and (7)

Case (15i).

 $\begin{array}{lll} (2) & e=(e_1,e_2) & \text{by assumption} \\ (3) & \dot{\xi}=(\dot{\xi}_1,\dot{\xi}_2) & \text{by assumption} \\ (4) & e_1\models\dot{\xi}_1 & \text{by assumption} \\ (5) & e_2\models\dot{\xi}_2 & \text{by assumption} \\ (6) & satisfy(e_1,\dot{\xi}_1)=\text{true} & \text{by IH on (4)} \\ (7) & satisfy(e_2,\dot{\xi}_2)=\text{true} & \text{by IH on (5)} \\ (8) & satisfy((e_1,e_2),(\dot{\xi}_1,\dot{\xi}_2))=\\ & satisfy(e_1,\dot{\xi}_1) \text{ and } satisfy(e_2,\dot{\xi}_2)=\text{true} \\ & \text{by Definition 16h on} \end{array}$

Case (15j).

(2) $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$ by assumption (3) e notintro by assumption (4) $prl(e) \models \dot{\xi}_1$ by assumption (5) $prr(e) \models \dot{\xi}_2$ by assumption (6) $satisfy(prl(e), \dot{\xi}_1) = true$ by IH on (4) (7) $satisfy(prr(e), \dot{\xi}_2) = true$ by IH on (5)

By rule induction over Rules (27) on (3).

Otherwise.

(8)
$$e = \{\|u, \|e_0\|^u, e_1(e_2), \text{prl}(e_0), \text{prr}(e_0), \text{match}(e_0) \}$$
 by assumption

(9)
$$satisfy(e, (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(\mathtt{prl}(e), \dot{\xi}_1)$$
 and $satisfy(\mathtt{prr}(e), \dot{\xi}_2) = \text{true}$ by Definition 16 on (6) and (7)

2. Completeness:

(1)
$$satisfy(e, \dot{\xi}) = true$$
 by assumption

By structural induction on $\dot{\xi}$.

Case
$$\dot{\xi} = \top$$
.

(2)
$$e \models \top$$
 by Rule (15a)

Case $\dot{\xi} = \bot, ?$.

(2)
$$satisfy(e, \dot{\xi}) = false$$
 by Definition 160

(2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = n$.

By structural induction on e.

Case $e = \underline{n'}$.

(2)
$$n' = n$$
 by Definition 16b on (1) by Puls (15b) on (2)

(3)
$$\underline{n'} \models \underline{n}$$
 by Rule (15b) on (2)

Otherwise.

(2)
$$satisfy(e, \underline{n}) = false$$
 by Definition 160

(2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = \underline{\varkappa}$.

By structural induction on e.

Case $e = \underline{n'}$.

(2)
$$n' \neq n$$
 by Definition 16c on (1)
(3) $\underline{n'} \models \underline{\varkappa}$ by Rule (15c) on (2)

Otherwise.

(2)
$$satisfy(e, \underline{\varkappa}) = false$$
 by Definition 160

(2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$.

(2)
$$satisfy(e, \dot{\xi}_1) = true$$
 by Definition 16d on (1)

(3)
$$\mathit{satisfy}(e,\dot{\xi}_2) = \mathsf{true}$$
 by Definition 16d on (1)

(4)
$$e \models \dot{\xi}_1$$
 by IH on (2)

(5)
$$e \models \dot{\xi}_2$$
 by IH on (3)

(6)
$$e \models \dot{\xi}_1 \wedge \dot{\xi}_2$$
 by Rule (15d) on (4) and (5)

Case $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$.

(2)
$$satisfy(e, \dot{\xi}_1)$$
 or $satisfy(e, \dot{\xi}_2) = true$ by Definition 16e on (1)

By case analysis on (2).

Case $satisfy(e, \dot{\xi}_1) = true.$

(3)
$$satisfy(e, \dot{\xi}_1) = true$$
 by assumption
(4) $e \models \dot{\xi}_1$ by IH on (3)

(5)
$$e \models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by Rule (15e) on (4)

Case $satisfy(e, \dot{\xi}_2) = true.$

(3)
$$satisfy(e, \dot{\xi}_2) = true$$
 by assumption
(4) $e \models \dot{\xi}_2$ by IH on (3)

(5)
$$e \models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by Rule (15f) on (4)

Case $\dot{\xi} = \text{inl}(\dot{\xi}_1)$.

By structural induction on e.

Case $e = \operatorname{inl}_{\tau_2}(e_1)$.

(2)
$$satisfy(e_1, \dot{\xi}_1) = true$$
 by Definition 16f on (1)

(3)
$$e_1 \models \dot{\xi}_1$$
 by IH on (2)

(4)
$$\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi_1})$$
 by Rule (15g) on (3)

Otherwise.

(2)
$$satisfy(e, \mathtt{inl}(\dot{\xi}_1)) = false$$
 by Definition 160

(2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = inr(\dot{\xi}_2)$.

By structural induction on e.

Case $e = \operatorname{inr}_{\tau_1}(e_2)$.

(2)
$$satisfy(e_2, \dot{\xi}_2) = true$$
 by Definition 16g on (1)

(3)
$$e_2 \models \dot{\xi}_2$$
 by IH on (2)

(4)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$$
 by Rule (15h) on (3)

Otherwise.

- (2) $satisfy(e, inr(\dot{\xi}_2)) = false$
- by Definition 16o
- (2) contradicts (1) and thus vacuously true.

Case
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$
.

By structural induction on e.

Case $e = (e_1, e_2)$.

- (2) $satisfy(e_1, \dot{\xi}_1) = true$ by Definition 16h on (1)
- (3) $satisfy(e_2, \dot{\xi}_2) = true$ by Definition 16h on (1)
- (4) $e_1 \models \dot{\xi}_1$ by IH on (2)
- (5) $e_2 \models \dot{\xi}_2$ by IH on (3)
- (6) $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (15i) on (4) and (5)

 $\mathbf{Case}\ e = \|)^u, (e_0)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{\hat{rs}\}.$

- (2) $satisfy(prl(e), \dot{\xi}_1) = true$ by Definition 16h on (1)
- (3) $satisfy(prr(e), \dot{\xi}_2) = true$ by Definition 16h on (1)
- (4) $prl(e) \models \dot{\xi}_1$ by IH on (2)
- (5) $prr(e) \models \dot{\xi}_2$ by IH on (3)
- (6) e notintro by each rule in Rules (27)
- (7) $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (15j) on (6) and (4) and (5)

Otherwise.

- (2) $satisfy(e, (\dot{\xi}_1, \dot{\xi}_2)) = false$
- by Definition 16o

(2) contradicts (1) and thus vacuously true.

 $e \models_? \dot{\xi}$ $e \text{ may satisfy } \dot{\xi}$

$$\frac{\text{CMSUnknown}}{e \models_? ?} \tag{6a}$$

CMSInl

$$\frac{e_1 \models_? \dot{\xi}_1}{\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)} \tag{6b}$$

CMSInr

$$\frac{e_2 \models_? \dot{\xi}_2}{\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)} \tag{6c}$$

CMSPairL
$$\begin{array}{ll}
e_1 \models_? \dot{\xi}_1 & e_2 \models \dot{\xi}_2 \\
\hline
(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)
\end{array} (6d)$$

$${\rm CMSPairR}$$

$$\frac{e_1 \models \dot{\xi}_1 \qquad e_2 \models_? \dot{\xi}_2}{(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)} \tag{6e}$$

$$\frac{e_1 \models_? \dot{\xi}_1 \qquad e_2 \models_? \dot{\xi}_2}{(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)} \tag{6f}$$

CMSOrL
$$\frac{e \models_? \dot{\xi}_1 \qquad e \not\models \dot{\xi}_2}{e \models_? \dot{\xi}_1 \lor \dot{\xi}_2} \tag{6g}$$

$$\frac{e \not\models \dot{\xi}_1 \quad e \models_? \dot{\xi}_2}{e \models_? \dot{\xi}_1 \lor \dot{\xi}_2} \tag{6h}$$

CMSNotIntro

$$\frac{e \text{ notintro} \qquad \dot{\xi} \text{ refutable}_?}{e \models_? \dot{\xi}} \tag{6i}$$

$\textit{maysatisfy}(e, \dot{\xi})$

$$maysatisfy(e,?) = true$$
 (7a)

$$maysatisfy(inl_{\tau_2}(e_1), inl(\dot{\xi}_1)) = maysatisfy(e_1, \dot{\xi}_1)$$
 (7b)

$$\mathit{maysatisfy}(\mathsf{inr}_{\tau_1}(e_2),\mathsf{inr}(\dot{\xi}_2)) = \mathit{maysatisfy}(e_2,\dot{\xi}_2) \tag{7c}$$

$$\begin{aligned} \mathit{maysatisfy}((e_1, e_2), (\dot{\xi}_1, \dot{\xi}_2)) &= \left(\mathit{maysatisfy}(e_1, \dot{\xi}_1) \text{ and } \mathit{satisfy}(e_2, \dot{\xi}_2) \right) \\ & \text{or } \left(\mathit{satisfy}(e_1, \dot{\xi}_1) \text{ and } \mathit{maysatisfy}(e_2, \dot{\xi}_2) \right) \\ & \text{or } \left(\mathit{maysatisfy}(e_1, \dot{\xi}_1) \text{ and } \mathit{maysatisfy}(e_2, \dot{\xi}_2) \right) \end{aligned}$$

 $\mathit{maysatisfy}(e,\dot{\xi}_1 \lor \dot{\xi}_2) = \left(\mathit{maysatisfy}(e,\dot{\xi}_1) \text{ and } \left(\mathit{not } \mathit{satisfy}(e,\dot{\xi}_2)\right)\right)$ or $\left(\left(\text{not } \mathit{satisfy}(e,\dot{\xi}_1)\right) \text{ and } \mathit{maysatisfy}(e,\dot{\xi}_2)\right)$ (7e)

$$maysatisfy(e, \dot{\xi}) = notintro(e) \text{ and } refutable_?(\dot{\xi})$$
 (7f)

Otherwise
$$maysatisfy(e, \dot{\xi}) = false$$
 (7g)

Lemma 1.0.3 (Soundness and Completeness of Maybe Satisfaction Judgment). $e \models_? \dot{\xi} iff maysatisfy(e, \dot{\xi}) = true.$

 ${\it Proof.}$ We prove soundness and completeness separately.

1. Soundness:

(1)	$e \models_? \mathcal{E}$	by	assumption
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By rule induction over Rules (17) on (1).

Case (17a).

(2) $\dot{\xi} = ?$	by assumption
$(3) \ \textit{maysatisfy}(e,?)$	by Equation 7a

Case (17e).

$(2) \ e = \mathtt{inl}_{\tau_2}(e_1)$	by assumption
$(3) \ \dot{\xi} = \mathtt{inl}(\dot{\xi_1})$	by assumption
$(4) e_1 \models_? \dot{\xi}_1$	by assumption
(5) $maysatisfy(e_1, \dot{\xi}_1)$	by IH on (4)
$(6) \ \textit{maysatisfy}(\texttt{inl}_{\tau_2}(e_1), \texttt{inl}(\dot{\xi}_1))$	by Equation 7b on (5)

Case (17f).

$(2) \ e = \operatorname{inr}_{\tau_1}(e_2)$	by assumption
$(3) \ \dot{\xi} = \mathtt{inr}(\dot{\xi}_2)$	by assumption
$(4) e_2 \models_? \dot{\xi}_2$	by assumption
(5) $maysatisfy(e_2,\dot{\xi}_2)$	by IH on (4)
$(6) \ \textit{maysatisfy}(\texttt{inr}_{\tau_1}(e_2),\texttt{inr}(\dot{\xi_2}))$	by Equation 7c on (5)

Case (17g).

(2) $e = (e_1, e_2)$	by assumption
(3) $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$	by assumption
$(4) e_1 \models_? \dot{\xi}_1$	by assumption
$(5) e_2 \models \dot{\xi}_2$	by assumption
(6) $maysatisfy(e_1,\dot{\xi}_1)$	by IH on (4)
(7) $satisfy(e_2, \dot{\xi}_2)$	by Lemma $2.0.19$ on
	(5)
(8) $maysatisfy((\dot{\xi}_1,\dot{\xi}_2),(\dot{\xi}_1,\dot{\xi}_2))$	by Equation 7d on (6)
	and (7)

Case (17h).

(2) $e = (e_1, e_2)$	by assumption
(3) $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$	by assumption
$(4) e_1 \models \dot{\xi}_1$	by assumption
$(5) e_2 \models_? \dot{\xi}_2$	by assumption
(6) $satisfy(e_1, \dot{\xi}_1)$	by Lemma $2.0.19$ on
	(4)

(7)	$\mathit{maysatisfy}(e_2,\dot{\xi}_2)$	by IH on (5)
(8)	$\mathit{maysatisfy}((\dot{\xi}_1,\dot{\xi}_2),(\dot{\xi}_1,\dot{\xi}_2))$	by Equation 7d on (6) and (7)
Case (17i).		
(2)	$e = (e_1, e_2)$	by assumption
(3)	$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$	by assumption
(4)	$e_1 \models_? \dot{\xi}_1$	by assumption
(5)	$e_2 \models_? \dot{\xi}_2$	by assumption
(6)	$\mathit{maysatisfy}(e_1,\dot{\xi}_1)$	by IH on (4)
(7)	$\mathit{maysatisfy}(e_2,\dot{\xi}_2)$	by IH on (?)
(8)	$\textit{maysatisfy}((\dot{\xi}_1,\dot{\xi}_2),(\dot{\xi}_1,\dot{\xi}_2))$	by Equation 7d on (6) and (7)
Case (17c).		
(2)	$\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$	by assumption
(3)	$e \models_? \dot{\xi}_1$	by assumption
(4)	$e \not\models \dot{\xi}_2$	by assumption
(5)	$\mathit{maysatisfy}(e, \dot{\xi}_1)$	by IH on (3)
(6)	not $satisfy(e, \dot{\xi}_2)$	by Lemma 2.0.19 on (4)
(7)	$\mathit{maysatisfy}(e, \dot{\xi}_1 \lor \dot{\xi}_2)$	by Equation 16e on (5) and (6)
Case (17d)	•	
(2)	$\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$	by assumption
(3)	$e \not\models \dot{\xi}_1$	by assumption
(4)	$e \models_? \dot{\xi}_2$	by assumption
(5)	not $satisfy(e, \dot{\xi}_1)$	by Lemma 2.0.19 on (3)
(6)	$\mathit{maysatisfy}(e,\dot{\xi}_2)$	by IH on (4)
(7)	$\textit{maysatisfy}(e, \dot{\xi}_1 \lor \dot{\xi}_2)$	by Equation 16e on (5) and (6)
Case (17b)	•	
(2)	e notintro	by assumption
(3)	$\dot{\xi}$ refutable?	by assumption
(4)	notintro(e)	by Lemma 4.0.1 on (?)
(5)	$refutable_{?}(\dot{\xi})$	by Lemma 2.0.14 on (3)
(6)	$\textit{maysatisfy}(e, \dot{\xi})$	by Equation 7f on (4) and (5)

2. Completeness:

e satisfies or may satisfy ξ

CSMSMay

$$\frac{e \models_? \dot{\xi}}{e \models_?^{\dagger} \dot{\xi}} \tag{8a}$$

CSMSSat

$$\frac{e \models \dot{\xi}}{e \models_{?}^{\dagger} \dot{\xi}} \tag{8b}$$

Lemma 1.0.4.

Lemma 1.0.5. $e \not\models \bot$

Proof. By rule induction over Rules (15), we notice that $e \models \bot$ is in syntactic contradiction with all rules, hence not derivable.

satormay soundness and completeness

Lemma 1.0.6. $e \not\models_? \bot$

Proof. Assume $e \models_? \bot$. By rule induction over Rules (17) on $e \models_? \bot$, only one case applies.

Case (17b).

(1) \perp refutable?

by assumption

By rule induction over Rules (11) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \bot$ is not derivable.

Lemma 1.0.7. $e \not\models_? \top$

Proof. Assume $e \models_? \top$. By rule induction over Rules (17) on $e \models_? \top$, only one case applies.

Case (17b).

(1) \top refutable?

by assumption

By rule induction over Rules (11) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \top$ is not derivable.

Lemma 1.0.8. $e \not\models ?$

Proof. By rule induction over Rules (15), we notice that $e \models ?$ is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.9. $e \models^{\dagger}_{?} \dot{\xi} iff e \models^{\dagger}_{?} \dot{\xi} \lor \bot$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e \models^{\dagger}_{?} \dot{\xi}$

by assumption

By rule induction over Rules (18) on (1).

Case (18a).

(2) $e \models_? \dot{\xi}$

by assumption

(3) $e \models_? \dot{\xi} \lor \bot$

by Rule (17c) on (2)

(4) $e \models_{?}^{\dagger} \dot{\xi} \lor \bot$

and Lemma 2.0.1 by Rule (18a) on (3)

Case (18b).

by assumption

(2) $e \models \dot{\xi}$ (3) $e \models \dot{\xi} \lor \bot$

by Rule (15e) on (2)

(4) $e \models_2^{\dagger} \dot{\xi} \lor \bot$

by Rule (18b) on (3)

- 2. Necessity:
 - (1) $e \models_2^{\dagger} \dot{\xi} \lor \bot$

by assumption

By rule induction over Rules (18) on (1).

Case (18a).

(2) $e \models_? \dot{\xi} \lor \bot$

by assumption

By rule induction over Rules (17) on (2), only two of them apply.

Case (17c).

(3) $e \models_? \dot{\xi}$

by assumption

(4) $e \models_{?}^{\dagger} \dot{\xi}$

by Rule (18a) on (3)

Case (17d).

(3) $e \models_? \bot$

by assumption

(4) $e \not\models_? \bot$

by Lemma 2.0.2

(3) contradicts (4).

Case (18b).

(2) $e \models \dot{\xi} \lor \bot$

by assumption

By rule induction over Rules (15) on (2), only two of them apply.

Case (15e).

(3)
$$e \models \dot{\xi}$$

by assumption

(4)
$$e \models_{?}^{\dagger} \dot{\xi}$$

by Rule (18b) on (3)

Case (15f).

(3)
$$e \models \bot$$

(4)
$$e \not\models \bot$$

by Lemma 2.0.1

(3) contradicts (4).

Corollary 1.0.1. $\top \models^{\dagger}_{?} \dot{\xi} \text{ iff } \top \models^{\dagger}_{?} \dot{\xi} \vee \bot$

Proof. Follows directly from Definition 2.1.2 and Lemma 2.0.5.

Lemma 1.0.10. Suppose that $\dot{\xi}_1 : \tau$ and $\dot{\xi}_2 : \tau$. Then $\dot{\xi}_1 \not\models \dot{\xi}_2$ iff $\dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$

Proof.

(1)
$$\dot{\xi}_1 : \tau$$

by assumption

(2)
$$\dot{\xi}_2 : \tau$$

by assumption

$$(3) \perp : \tau$$

by Rule (9b)

(4)
$$\dot{\xi}_2 \vee \bot : \tau$$

by Rule (9f) on (2) and

Then we prove sufficiency and necessity separately.

1. Sufficiency:

$$(5) \ \dot{\xi}_1 \not\models \dot{\xi}_2$$

by assumption

To prove $\dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$, assume $\dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$.

(6)
$$\dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models_{?}^{\dagger} \dot{\xi}_{1}$ implies

(7)
$$e \models \dot{\xi}_2 \lor \bot$$

by Definition 2.1.1 on

(1) and (4) and (6)

By rule induction over Rules (15) on (7).

Case (15e).

(8)
$$e \models \dot{\xi}_2$$

(9) $\dot{\xi}_1 \models \dot{\xi}_2$

by assumption

(9)
$$\xi_1 \models \xi_2$$

by Definition 2.1.1 on

(8)

(5) contradicts (9).

Case (15f).

 $(8) e \models \bot$ $(9) e \not\models \bot$

by assumption by Lemma 2.0.1

(8) contradicts (9).

The conclusion holds as follows:

- (a) $\dot{\xi}_1 \not\models \dot{\xi}_2 \vee \bot$
- 2. Necessity:
 - (5) $\dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$

by assumption

To prove $\dot{\xi}_1 \not\models \dot{\xi}_2$, assume $\dot{\xi}_1 \models \dot{\xi}_2$.

(6) $\dot{\xi}_1 \models \dot{\xi}_2$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models_{?}^{\dagger} \dot{\xi}_{1}$ implies

(7) $e \models \dot{\xi}_2$

by Definition 2.1.1 on (1) and (2) and (6)

(8) $e \models \dot{\xi}_2 \lor \bot$

by Rule (15e) on (7)

(9) $\dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$

by Definition 2.1.1 on

(8)

(9) contradicts (5).

The conclusion holds as follows:

(a) $\dot{\xi}_1 \not\models \dot{\xi}_2$

Lemma 1.0.11. $e \not\models_{?}^{\dagger} \dot{\xi}_1$ and $e \not\models_{?}^{\dagger} \dot{\xi}_2$ iff $e \not\models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency: to show $e \not\models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$, we assume $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$.
 - (1) $e \models_2^{\dagger} \dot{\xi_1} \lor \dot{\xi_2}$

by assumption

(2) $e \not\models_2^{\dagger} \dot{\xi}_1$

by assumption

(3) $e \not\models_{?}^{\dagger} \dot{\xi}_{2}$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(4) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

By rule induction over Rules (15) on (4) and only two of them apply.

Case (15e).

(5) $e \models \dot{\xi}_1$

by assumption

(6) $e \models_{?}^{\dagger} \dot{\xi}_1$

by Rule (18b) on (5)

(6) contradicts (2).

Case (15f).

(5) $e \models \dot{\xi}_2$

by assumption

(6) $e \models_2^{\dagger} \dot{\xi}_2$

by Rule (18b) on (5)

(6) contradicts (3).

Case (18a).

(4) $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

By rule induction over Rules (17) on (4) and only two of them apply.

Case (17c).

(5) $e \models_? \xi_1$

by assumption

(6) $e \models_{?}^{\dagger} \dot{\xi}_1$

by Rule (18a) on (5)

(6) contradicts (2).

Case (17d).

(5) $e \models_? \dot{\xi}_2$

by assumption

(6) $e \models_{?}^{\dagger} \dot{\xi}_2$

by Rule (18a) on (5)

(6) contradicts (3).

The conclusion holds as follows:

- (a) $e \not\models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$
- 2. Necessity:
 - (1) $e \not\models_2^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

We show $e \not\models_{?}^{\dagger} \dot{\xi_1}$ and $e \not\models_{?}^{\dagger} \dot{\xi_2}$ separately.

- (a) To show $e \not\models_?^\dagger \dot{\xi}_1$, we assume $e \models_?^\dagger \dot{\xi}_1$.
 - (2) $e \models^{\dagger}_{?} \dot{\xi}_{1}$

by assumption

(3) $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$

by Lemma 2.0.10 on

(2)

Contradicts (1).

- (b) To show $e \not\models_{?}^{\dagger} \dot{\xi}_{2}$, we assume $e \models_{?}^{\dagger} \dot{\xi}_{2}$.
 - (2) $e \models_{?}^{\dagger} \dot{\xi}_{2}$

by assumption

(3) $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$

by Lemma 2.0.10 on

(2)

Contradicts (1).

In conclusion, $e \not\models_{?}^{\dagger} \dot{\xi}_1$ and $e \not\models_{?}^{\dagger} \dot{\xi}_2$.

Lemma 1.0.12. If $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ and $e \not\models_{?}^{\dagger} \dot{\xi}_1$ then $e \models_{?}^{\dagger} \dot{\xi}_2$

Proof.

- $(4) e \models_{?}^{\dagger} \dot{\xi_1} \lor \dot{\xi_2}$
- by assumption

(5) $e \not\models_{?}^{\dagger} \dot{\xi}_1$

by assumption

By rule induction over Rules (18) on (4).

Case (18b).

(6) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

By rule induction over Rules (15) on (6) and only two of them apply.

Case (15e).

(7) $e \models \dot{\xi}_1$

by assumption

(8) $e \models^{\dagger}_{?} \dot{\xi}_{1}$

by Rule (18b) on (7)

(8) contradicts (5).

Case (15f).

(7) $e \models \dot{\xi}_2$

by assumption

(8) $e \models_{?}^{\dagger} \dot{\xi}_2$

by Rule (18b) on (7)

Case (18a).

(6) $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

By rule induction over Rules (17) on (6) and only two of them apply.

Case (17c).

(7) $e \models_? \dot{\xi}_1$

by assumption

(8) $e \models_2^{\dagger} \dot{\xi}_1$

by Rule (18a) on (7)

(8) contradicts (5).

Case (17d).

(7) $e \models_? \dot{\xi}_2$

by assumption

(8) $e \models_2^{\dagger} \dot{\xi}_2$

by Rule (18a) on (7)

Lemma 1.0.13. If
$$e \models_{?}^{\dagger} \dot{\xi}_1$$
 then $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ and $e \models_{?}^{\dagger} \dot{\xi}_2 \lor \dot{\xi}_1$

Proof.

(1) $e \models^{\dagger}_{?} \dot{\xi}_{1}$ by assumption

By rule induction over Rules (18) on (1),

Case (18b).

(2)
$$e \models \dot{\xi}_{1}$$
 by assumption
(3) $e \models \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Rule (15e) on (2)
(4) $e \models \dot{\xi}_{2} \lor \dot{\xi}_{1}$ by Rule (15f) on (2)
(5) $e \models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Rule (18b) on (3)
(6) $e \models_{?}^{\dagger} \dot{\xi}_{2} \lor \dot{\xi}_{1}$ by Rule (18b) on (4)

Case (18a).

(2)
$$e \models_? \dot{\xi}_1$$
 by assumption

By case analysis on the result of $satisfy(e, \dot{\xi}_2)$.

Case true.

(3)
$$satisfy(e, \dot{\xi}_2) = true$$
 by assumption
(4) $e \models \dot{\xi}_2$ by Lemma 2.0.19 on
(5) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (15f) on (4)
(6) $e \models \dot{\xi}_2 \lor \dot{\xi}_1$ by Rule (15e) on (4)
(7) $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (18b) on (5)
(8) $e \models_{?}^{\dagger} \dot{\xi}_2 \lor \dot{\xi}_1$ by Rule (18b) on (6)

Case false.

(3)
$$satisfy(e, \dot{\xi}_2) = false$$
 by assumption
(4) $e \not\models \dot{\xi}_2$ by Lemma 2.0.19 on
(3)
(5) $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (17c) on (2)
and (4)
(6) $e \models_?^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (18a) on (5)

Lemma 1.0.14. $e_1 \models_?^\dagger \dot{\xi_1} \ iff \ \mathrm{inl}_{\tau_2}(e_1) \models_?^\dagger \ \mathrm{inl}(\dot{\xi_1})$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e_1 \models_2^{\dagger} \dot{\xi}_1$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2) $e_1 \models \dot{\xi}_1$

- by assumption
- (3) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$
- by Rule (15g) on (2)
- $(4) \ \operatorname{inl}_{\tau_2}(e_1) \models^\dagger_? \operatorname{inl}(\dot{\xi_1})$
- by Rule (18b) on (3)

Case (18a).

(2) $e_1 \models_? \dot{\xi}_1$

- by assumption
- (3) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi_1})$
- by Rule (17e) on (2)
- (4) $\operatorname{inl}_{\tau_2}(e_1) \models_{\tau}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$
- by Rule (18a) on (3)

- 2. Necessity:
 - (1) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

- (2) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi_1})$
- by assumption

By rule induction over Rules (15) on (2), only one rule applies.

Case (15g).

(3) $e_1 \models \dot{\xi}_1$

by assumption

(4) $e_1 \models_{?}^{\dagger} \dot{\xi}_1$

by Rule (18b) on (3)

Case (18a).

- (2) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi_1})$
- by assumption

By rule induction over Rules (17) on (2), only two rules apply.

Case (17e).

(3) $e_1 \models_? \dot{\xi}_1$

by assumption

(4) $e_1 \models_2^{\dagger} \dot{\xi}_1$

by Rule (18a) on (3)

Case (17b).

- (3) $\operatorname{inl}_{\tau_2}(e_1)$ notintro
- by assumption

By rule induction over Rules (27) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.15. $e_2\models^\dagger_?\dot{\xi}_2$ iff $\operatorname{inr}_{\tau_1}(e_2)\models^\dagger_?\operatorname{inr}(\dot{\xi}_2)$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e_2 \models_{?}^{\dagger} \dot{\xi}_2$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2) $e_2 \models \dot{\xi}_2$

- by assumption
- (3) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$
- by Rule (15h) on (2)

- (4) $\operatorname{inr}_{\tau_1}(e_2) \models_{\tau_1}^{\dagger} \operatorname{inr}(\dot{\xi_2})$
- by Rule (18b) on (3)

Case (18a).

(2) $e_2 \models_? \dot{\xi}_2$

- by assumption
- (3) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$
- by Rule (17f) on (2)
- (4) $\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$
- by Rule (18a) on (3)

- 2. Necessity:
 - (1) $\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\dot{\xi_2})$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

- (2) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$
- by assumption

By rule induction over Rules (15) on (2), only one rule applies.

Case (15h).

(3) $e_2 \models \dot{\xi}_2$

by assumption

(4) $e_2 \models_?^{\dagger} \dot{\xi}_2$

by Rule (18b) on (3)

Case (18a).

 $(2) \ \operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi_2})$

by assumption

By rule induction over Rules (17) on (2), only two rules apply.

Case (17f).

(3) $e_2 \models_? \dot{\xi}_2$

by assumption

(4) $e_2 \models_{?}^{\dagger} \dot{\xi}_2$

by Rule (18a) on (3)

Case (17b).

$$(3) \ \operatorname{inr}_{\tau_1}(e_2) \ \operatorname{notintro}$$

by assumption

By rule induction over Rules (27) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.16. $e_1 \models_{?}^{\dagger} \dot{\xi_1} \text{ and } e_2 \models_{?}^{\dagger} \dot{\xi_2} \text{ iff } (e_1, e_2) \models_{?}^{\dagger} (\dot{\xi_1}, \dot{\xi_2})$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e_1 \models_{?}^{\dagger} \dot{\xi}_1$

by assumption

(2) $e_2 \models_{?}^{\dagger} \dot{\xi}_2$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(3)
$$e_1 \models \dot{\xi}_1$$

by assumption

By rule induction over Rules (18) on (2).

Case (18b).

$$(4) e_2 \models \dot{\xi}_2$$

by assumption

(5)
$$(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (15i) on (3)

and
$$(4)$$

(6)
$$(e_1, e_2) \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (18b) on (5)

Case (18a).

(4)
$$e_2 \models_? \dot{\xi}_2$$

by assumption

(5)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (17h) on (3)

and
$$(4)$$

(6)
$$(e_1, e_2) \models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (18a) on (5)

Case (18a).

(4)
$$e_1 \models_? \dot{\xi}_1$$

by assumption

By rule induction over Rules (18) on (2).

Case (18b).

$$(5) e_2 \models \dot{\xi}_2$$

by assumption

(6)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (17g) on (4)

(7)
$$(e_1, e_2) \models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (18a) on (6)

Case (18a).

(5)
$$e_2 \models_? \dot{\xi}_2$$

by assumption

(6)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (17h) on (4)

and (5)

(7)
$$(e_1, e_2) \models_{?}^{+} (\dot{\xi}_1, \dot{\hat{\xi}}_2)$$
 by Rule (18a) on (6)

2. Necessity:

(1) $(e_1, e_2) \models_{?}^{+} (\dot{\xi}_1, \dot{\xi}_2)$ by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2) $(e_1, e_2) \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$ by assumption

By rule induction over Rules (15) on (2), only one rule applies.

Case (15i).

(3) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(4) $e_2 \models_{?} \dot{\xi}_2$ by assumption

(5) $e_1 \models_{?}^{+} \dot{\xi}_1$ by Rule (18b) on (3)

(6) $e_2 \models_{?}^{+} \dot{\xi}_2$ by Rule (18b) on (4)

Case (18a).

(2) $(e_1, e_2) \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$ by assumption

By rule induction over Rules (17) on (2), only three rules apply.

Case (17g).

(3) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(4) $e_2 \models_{?} \dot{\xi}_2$ by assumption

(5) $e_1 \models_{?}^{+} \dot{\xi}_1$ by Rule (18a) on (3)

(6) $e_2 \models_{?}^{+} \dot{\xi}_2$ by assumption

(7) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(8) $e_1 \models_{?} \dot{\xi}_1$ by Rule (18b) on (4)

Case (17h).

(9) $e_1 \models_{?} \dot{\xi}_1$ by Rule (18b) on (4)

Case (17i).

(10) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(21) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(32) $e_1 \models_{?} \dot{\xi}_2$ by assumption

(33) $e_1 \models_{?} \dot{\xi}_1$ by Rule (18a) on (3)

(44) $e_2 \models_{?} \dot{\xi}_2$ by Rule (18a) on (4)

Case (17i).

(5) $e_1 \models_{?} \dot{\xi}_1$ by assumption

(6) $e_2 \models_{?} \dot{\xi}_2$ by assumption

(7) $e_1 \models_{?} \dot{\xi}_1$ by Rule (18a) on (3)

(8) $e_2 \models_{?} \dot{\xi}_2$ by assumption

(9) $e_1 \models_{?} \dot{\xi}_1$ by Rule (18a) on (4)

Lemma 1.0.17. *If* e notintro and $e \models_? \xi$ then ξ refutable?

Lemma 1.0.18. There does not exist such a constraint $\dot{\xi}_1 \wedge \dot{\xi}_2$ such that $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable?

Proof. By rule induction over Rules (11), we notice that $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.19. There does not exist such a constraint $\dot{\xi}_1 \lor \dot{\xi}_2$ such that $\dot{\xi}_1 \lor \dot{\xi}_2$ refutable?.

Proof. By rule induction over Rules (11), we notice that $\dot{\xi}_1 \vee \dot{\xi}_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.20. If e notintro and $e \models \dot{\xi}$ then $\dot{\xi}$ refutable?

Proof.

(1) e notintro by assumption

(2) $e \models \dot{\xi}$ by assumption

By rule induction over Rules (15) on (2).

Case (15a).

(3)
$$\dot{\xi} = \top$$
 by assumption

Assume \top refutable?. By rule induction over Rules (11), no case applies due to syntactic contradiction.

Therefore, \top refutable?.

Case (15e), (15f).

(3) $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$ by assumption

(4) $\underline{\dot{\xi}_1 \lor \dot{\xi}_2 \text{ refutable}_?}$ by Lemma 2.0.17

Case (15d).

(3) $\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$ by assumption

(4) $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable? by Lemma 2.0.16

Case (15j).

(3) $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$ by assumption

(4) $prl(e) \models \dot{\xi}_1$ by assumption

(5) $prr(e) \models \dot{\xi}_2$ by assumption

(6) prl(e) notintro by Rule (27e)

(7) prr(e) notintro by Rule (27f)

(8) ξ_1 refutable? by IH on (6) and (4)

(9)
$$\dot{\xi}_2$$
 refutable?

by IH on (7) and (5)

Assume $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? By rule induction over Rules (11) on it, only two cases apply.

Case (11d).

(10) $\dot{\xi}_1$ refutable?

by assumption

Contradicts (8).

Case (11e).

(10) $\dot{\xi}_2$ refutable?

by assumption

Contradicts (9).

Therefore, $(\dot{\xi}_1, \dot{\xi}_2)$ refutable?.

Otherwise.

(3)
$$e = \underline{n}, \operatorname{inl}_{\tau_2}(e_1), \operatorname{inr}_{\tau_1}(e_2), (e_1, e_2)$$
 by assumption

By rule induction over Rules (27) on (1), no case applies due to syntactic contradiction.

Lemma 1.0.21. $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ is not derivable.

Proof. We prove by assuming $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ and obtaining a contradiction.

$$(1) \ \operatorname{inl}_{\tau_2}(e_1) \models^\dagger_? \operatorname{inr}(\dot{\xi_2})$$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

$$(2) \ \operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\dot{\xi}_2)$$

by assumption

By rule induction over Rules (15) on (2), no rule applies due to syntactic contradiction.

Case (18a).

$$(2) \ \operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\dot{\xi}_2)$$

by assumption

By rule induction over Rules (17) on (2), only one rule applies.

Case (17b).

(3)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro

by assumption

By rule induction over Rules (27) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.22. $\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$ is not derivable.

Proof. We prove by assuming $\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ and obtaining a contradiction.

(1)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau_1}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\dot{\xi}_1)$$
 by assumption

By rule induction over Rules (15) on (2), no rule applies due to syntactic contradiction.

Case (18a).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by assumption

By rule induction over Rules (17) on (2), only one rule applies.

Case (17b).

(3)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro by assumption

By rule induction over Rules (27) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.23. $e \not\models \dot{\xi}$ and $e \not\models_? \dot{\xi}$ iff $e \not\models_?^{\dagger} \dot{\xi}$.

Proof. 1. Sufficiency:

(1)
$$e \not\models \dot{\xi}$$
 by assumption

(2)
$$e \not\models_? \dot{\xi}$$
 by assumption

Assume $e \models_{?}^{\dagger} \dot{\xi}$. By rule induction over Rules (18) on it.

Case (18a).

(3)
$$e \models \dot{\xi}$$
 by assumption

Contradicts (1).

Case (18b).

(3)
$$e \models_? \dot{\xi}$$
 by assumption

Contradicts (2).

Therefore, $e \models^{\dagger}_{?} \dot{\xi}$ is not derivable.

- 2. Necessity:
 - (1) $e \not\models_?^\dagger \dot{\xi}$

by assumption

Assume $e \models \dot{\xi}$.

(2) $e \models^{\dagger}_{?} \dot{\xi}$

by Rule (18b) on assumption

Contradicts (1). Therefore, $e \not\models \dot{\xi}$. Assume $e \models_? \dot{\xi}$.

(3) $e \models^{\dagger}_{?} \dot{\xi}$

by Rule (18a) on assumption

Contradicts (1). Therefore, $e \not\models_? \dot{\xi}$.

Theorem 1.1 (Exclusiveness of Satisfaction Judgment). If $\dot{\xi}$: τ and \cdot ; $\Delta \vdash e$: τ and e final then exactly one of the following holds

- 1. $e \models \dot{\xi}$
- 2. $e \models_? \dot{\xi}$
- 3. $e \not\models_?^\dagger \dot{\xi}$

Proof.

 $(4) \ \dot{\xi}:\tau$

by assumption

(5) \cdot ; $\Delta \vdash e : \tau$

by assumption

(6) e final

by assumption

By rule induction over Rules (9) on (4), we would show one conclusion is derivable while the other two are not.

Case (9a).

 $(7) \ \dot{\xi} = \top$

by assumption

(8) $e \models \top$

by Rule (15a)

(9) $e \not\models_? \top$

by Lemma 2.0.3

(10) $e \models_?^\dagger \top$

by Rule (18b) on (8)

Case (9b).

$(7) \ \dot{\xi} = \bot$	by assumption
(8) $e \not\models \bot$	by Lemma 2.0.1
$(9) \ e \not\models_? \bot$	by Lemma 2.0.2
$(10) \ e \not\models^{\dagger}_{?} \bot$	by Lemma $2.0.20$ on
	(8) and (9)

Case (1b).

(7) $\dot{\xi} = ?$	by assumption
(8) $e \not\models ?$	by Lemma $2.0.4$
$(9) e \models_? ?$	by Rule (17a)
$(10) \ e \models^{\dagger}_{?} ?$	by Rule $(18a)$ on (9)

Case (9c).

(7)
$$\dot{\xi} = \underline{n_2}$$
 by assumption
(8) $\tau = \text{num}$ by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20l),(20m).

$$(9) \ \ e = ()^u, (e_0)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{ \hat{rs} \}$$
 by assumption
$$(10) \ \ e \ \mathtt{notintro} \qquad \qquad \mathsf{by Rule}$$

$$(27a), (27b), (27c), (27d), (27e), (27f)$$

Assume $e \models \underline{n_2}$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction on $\dot{\xi}$.

$(11) \ e \not\models \underline{n_2}$	by contradiction
(12) $\underline{n_2}$ refutable?	by Rule (11a)
$(13) e \models_? \underline{n_2}$	by Rule (17b) on (10)
	and (12)
$(14) \ e \models^{\dagger}_{?} \underline{n_2}$	by Rule $(18a)$ on (13)

Case (20d).

(9)
$$e = n_1$$
 by assumption

Assume $\underline{n_1} \models_? \underline{n_2}$. By rule induction over Rules (17), only one case applies.

Case (17b).

(10)
$$\underline{n_1}$$
 notintro by assumption Contradicts Lemma 4.0.6.

(11)
$$n_1 \not\models_? n_2$$
 by contradiction

By case analysis on whether n_1 is equal to n_2 .

Case $n_1 = n_2$.

$$\begin{array}{lll} (12) \;\; satisfy(\underline{n_1},\underline{n_2}) = {\rm true} & & {\rm by \; Definition \; 16} \\ (13) \;\; \underline{n_1} \models \underline{n_2} & & {\rm by \; Lemma \; 2.0.19 \; on} \\ (14) \;\; e \models_?^\dagger \underline{n_2} & & {\rm by \; Rule \; (18b) \; on \; (13)} \\ \end{array}$$

Case $n_1 \neq n_2$.

(12)
$$satisfy(\underline{n_1}, \underline{n_2}) = false$$
 by Definition 16
(13) $\underline{n_1} \not\models \underline{n_2}$ by Lemma 2.0.19 on (12)
(14) $e \not\models_{?}^{\dagger} \underline{n_2}$ by Lemma 2.0.20 on (11) and (13)

Case (9f).

(7)
$$\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$$
 by assumption

By inductive hypothesis on (5) and (6), exactly one of $e \models \dot{\xi}_1$, $e \models_? \dot{\xi}_1$, and $e \not\models_?^{\dagger} \dot{\xi}_1$ holds. The same goes for $\dot{\xi}_2$. By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $e \models \dot{\xi}_1, e \models \dot{\xi}_2$.

(8) $e \models \xi_1$	by assumption
$(9) \ e \not\models_? \dot{\xi}_1$	by assumption
$(10) \ e \models \dot{\xi}_2$	by assumption
$(11) \ e \not\models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule $(15e)$ on (8)
$(13) \ e \models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule (18b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14)
$$\dot{\xi}_1 \vee \dot{\xi}_2$$
 refutable? by assumption Contradicts Lemma 2.0.17.

Case (17c).

(14)
$$e \models_? \dot{\xi}_1$$
 by assumption Contradicts (9).

Case (17d).

(14)
$$e \models_? \dot{\xi}_2$$
 by assumption Contradicts (11).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \models \dot{\xi}_1, e \models_? \dot{\xi}_2$.

(8) $e \models \dot{\xi}_1$	by assumption
$(9) \ e \not\models_? \dot{\xi}_1$	by assumption
$(10) \ e \not\models \dot{\xi}_2$	by assumption
$(11) \ e \models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule $(15e)$ on (8)
(13) $e \models_{2}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$	by Rule (18b) on (12

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14)
$$\dot{\xi}_1 \vee \dot{\xi}_2$$
 refutable? by assumption Contradicts Lemma 2.0.17.

Case (17c).

(14)
$$e \models_? \dot{\xi}_1$$
 by assumption Contradicts (9).

Case (17d).

(14)
$$e \not\models \dot{\xi}_1$$
 by assumption Contradicts (8).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \models \dot{\xi}_1, e \not\models_{?}^{\dagger} \dot{\xi}_2$.

(8)
$$e \models \dot{\xi}_1$$
 by assumption
(9) $e \not\models_? \dot{\xi}_1$ by assumption
(10) $e \not\models \dot{\xi}_2$ by assumption
(11) $e \not\models_? \dot{\xi}_2$ by assumption
(12) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (15e) on (8)
(13) $e \models_?^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (18b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\dot{\xi}_1 \lor \dot{\xi}_2$ refutable? Contradicts Lemma 2.0.17.	by assumption
Case (17c).	
$(14) \ e \models_? \dot{\xi}_1$	by assumption
Contradicts (9).	by abbampion
Case (17d).	
$(14) \ e \not\models \dot{\xi}_1$	by assumption
Contradicts (8).	sy assamption
()	
$(15) \ e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$	by contradiction
Case $e \models_? \dot{\xi}_1, e \models \dot{\xi}_2$.	
(8) $e \not\models \dot{\xi}_1$	by assumption
$(9) e \models_? \dot{\xi}_1$	by assumption
$(10) e \models \dot{\xi}_2$	by assumption
$(11) \ e \not\models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule $(15f)$ on (10)
$(13) \ e \models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule $(18b)$ on (12)
Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction following cases apply.	over Rules (17) on it, the
Case (17b).	
(14) $\dot{\xi}_1 ee \dot{\xi}_2$ refutable $_?$	by assumption
Contradicts Lemma 2.0.17.	
Case (17c).	
$(14) \ e \not\models \dot{\xi}_2$	by assumption
Contradicts (10).	
Case (17d).	
$(14) e \models_? \dot{\xi}_2$	by assumption
Contradicts (11).	
$(15) \ e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$	by contradiction
Case $e \models_? \dot{\xi}_1, e \models_? \dot{\xi}_2$.	
(8) $e \not\models \dot{\xi}_1$	by assumption
(9) $e \models_? \dot{\xi_1}$	by assumption
$(10) e \not\models \dot{\xi}_2$	by assumption
$(11) e \models_? \dot{\xi_2}$	by assumption
	_

(12)
$$e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by Rule (17c) on (9) and (10)

(13)
$$e \models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$$
 by Rule (18a) on (12)

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14)
$$e \models \dot{\xi}_1$$
 by assumption

Contradicts (8)

Case (15f).

(14)
$$e \models \dot{\xi}_2$$
 by assumption Contradicts (10)

(15)
$$e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \models_? \dot{\xi}_1, e \not\models_?^{\dagger} \dot{\xi}_2$.

(8)
$$e \not\models \dot{\xi}_1$$
 by assumption
(9) $e \models_? \dot{\xi}_1$ by assumption
(10) $e \not\models \dot{\xi}_2$ by assumption
(11) $e \not\models_? \dot{\xi}_2$ by assumption
(12) $e \models_? \dot{\xi}_2 \lor \dot{\xi}_2$ by Rule (17c) on

(12)
$$e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by Rule (17c) on (9) and (10)

(13)
$$e \models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$$
 by Rule (18a) on (12)

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14)
$$e \models \dot{\xi}_1$$
 by assumption

Contradicts (8).

Case (15f).

(14)
$$e \models \dot{\xi}_2$$
 by assumption Contradicts (10).

(15)
$$e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \not\models_?^\dagger \dot{\xi}_1, e \models \dot{\xi}_2$.

(8)
$$e \not\models \dot{\xi}_1$$
 by assumption
(9) $e \not\models_? \dot{\xi}_1$ by assumption
(10) $e \models \dot{\xi}_2$ by assumption
(11) $e \not\models_? \dot{\xi}_2$ by assumption

(12)
$$e \models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by Rule (15f) on (10)
(13) $e \models_{7}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (18b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\dot{\xi}_1 \vee \dot{\xi}_2$ refutable? by assumption Contradicts Lemma 2.0.17.

Case (17c).

(14)
$$e \not\models \dot{\xi}_2$$
 by assumption

Contradicts (10).

Case (17d).

(14)
$$e \models_? \dot{\xi}_2$$
 by assumption Contradicts (11).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \not\models_{?}^{\dagger} \dot{\xi}_1, e \models_{?} \dot{\xi}_2$.

$$(8) \ e \not\models \dot{\xi}_1 \qquad \qquad \text{by assumption}$$

$$(9) \ e \not\models_? \dot{\xi}_1 \qquad \qquad \text{by assumption}$$

$$(10) \ e \not\models \dot{\xi}_2 \qquad \qquad \text{by assumption}$$

$$(11) \ e \models_? \dot{\xi}_2 \qquad \qquad \text{by assumption}$$

$$(12) \ e \models_? \dot{\xi}_1 \lor \dot{\xi}_2 \qquad \qquad \text{by Rule (17d) on (11)}$$

$$(13) \ e \models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2 \qquad \qquad \text{by Rule (18a) on (12)}$$

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14)
$$e \models \dot{\xi}_1$$
 by assumption

Contradicts (8)

Case (15f).

(14)
$$e \models \dot{\xi}_2$$
 by assumption Contradicts (10)

(15)
$$e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \not\models_?^\dagger \dot{\xi}_1, e \not\models_?^\dagger \dot{\xi}_2$.

(8)
$$e \not\models \dot{\xi}_1$$
 by assumption

$(9) \ e \not\models_? \dot{\xi}_1$	by assumption
$(10) \ e \not\models \dot{\xi}_2$	by assumption
(11) $e \not\models_2 \dot{\mathcal{E}}_2$	by assumption

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (15) on it, only two cases apply.

Case (15e).

(12)
$$e \models \dot{\xi}_1$$
 by assumption

Contradicts (8).

Case (15f).

(12)
$$e \models \dot{\xi}_2$$
 by assumption Contradicts (10).

(13)
$$e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14)
$$\dot{\xi}_1 \vee \dot{\xi}_2$$
 refutable? by assumption Contradicts Lemma 2.0.17.

Case (17c).

(14)
$$e \models_? \dot{\xi}_1$$
 by assumption

Contradicts (9).

Case (17d).

(14)
$$e \models_? \dot{\xi}_2$$
 by assumption Contradicts (11).

(15)
$$e \not\models_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}$$
 by contradiction
(16) $e \not\models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Lemma 2.0.20 on
(13) and (15)

Case (9g).

$$\begin{array}{ll} (7) \ \ \dot{\xi} = \mathtt{inl}(\dot{\xi}_1) & \text{by assumption} \\ (8) \ \ \tau = (\tau_1 + \tau_2) & \text{by assumption} \\ (9) \ \ \dot{\xi}_1 : \tau_1 & \text{by assumption} \end{array}$$

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20l),(20m).

$$(10) \ e = (\!()^u, (\!(e_0)\!)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{\hat{rs}\}$$

by assumption

(11)
$$e$$
 notintro by Rule

Assume $e \models \mathtt{inl}(\dot{\xi}_1)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inl}(\dot{\xi}_1)$$

by contradiction

By case analysis on the value of $refutable_{?}(inl(\dot{\xi}_{1}))$.

Case $refutable_{?}(inl(\dot{\xi}_1)) = true.$

(13)
$$refutable_{?}(inl(\dot{\xi}_1)) = true$$
 by assumption

(14)
$$\operatorname{inl}(\dot{\xi}_1)$$
 refutable? by Lemma 2.0.14 on (13)

(15)
$$e \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (17b) on (11) and (14)

(16)
$$e \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (18a) on (15)

Case $refutable_{?}(inl(\dot{\xi}_1)) = false.$

(13)
$$refutable_{?}(inl(\dot{\xi}_{1})) = false$$
 by assumption

(14)
$$\underline{\operatorname{inl}(\dot{\xi}_1)}$$
 refutable? by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(15)
$$\operatorname{inl}(\dot{\xi}_1)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inl}(\dot{\xi}_1)$$
 by contradiction
(17) $e \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ by Lemma 2.0.20 on
(12) and (16)

Case (20j).

$$\begin{array}{ll} (10) & e = \operatorname{inl}_{\tau_2}(e_1) & \text{by assumption} \\ (11) & \cdot ; \Delta \vdash e_1 : \tau_1 & \text{by assumption} \\ (12) & e_1 \text{ final} & \text{by Lemma 4.0.3 on (6)} \end{array}$$

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_1 \models \dot{\xi}_1$, $e_1 \models_? \dot{\xi}_1$, and $e_1 \not\models_?^{\dagger} \dot{\xi}_1$ holds. By case analysis on which one holds.

Case $e_1 \models \dot{\xi}_1$.

(13)
$$e_1 \models \dot{\xi}_1$$
 by assumption (14) $e_1 \not\models_? \dot{\xi}_1$ by assumption

(15)
$$\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (15g) on (13)

(16)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{\tau}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (18b) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (17) on it, only two cases apply.

Case (17b).

(17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17e).

$$(17) \ e_1 \models_? \dot{\xi}_1$$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\dot{\xi_1})$$
 by contradiction

Case $e_1 \models_? \dot{\xi}_1$.

(13)
$$e_1 \not\models \dot{\xi}_1$$
 by assumption

(14)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption

(15)
$$\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (17e) on (14)

(16)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{\tau_1}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (18a) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (15) on it, only one case applies.

Case (15g).

$$(17) e_1 \models \dot{\xi}_1$$

Contradicts (13).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\dot{\xi}_1)$$
 by contradiction

Case $e_1 \not\models_?^\dagger \dot{\xi}_1$.

(13)
$$e_1 \not\models \dot{\xi}_1$$
 by assumption

(14)
$$e_1 \not\models_? \dot{\xi}_1$$
 by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (15) on it, only one case applies.

Case (15g).

$$(15) e_1 \models \dot{\xi}_1$$

Contradicts (13).

(16)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\dot{\xi_1})$$
 by contradiction Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi_1})$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(17) inl $_{ au_2}(e_1)$ notintro

by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17e).

(17) $e_1 \models_? \dot{\xi}_1$

Contradicts (14).

(18) $\operatorname{inl}_{\tau_2}(e_1) \not\models_{?} \operatorname{inl}(\dot{\xi_1})$ by contradiction

(19) $\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ by Lemma 2.0.20 on (16) and (18)

Case (20k).

(10) $e = inr_{\tau_1}(e_2)$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(11) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inl}(\dot{\xi}_1)$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\dot{\xi_1})$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(12) $\operatorname{inr}_{\tau_1}(e_2)$ notintro

by assumption

By rule induction over Rules (27) on (12), no case applies due to syntactic contradiction.

$$(13) \ \operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inl}(\dot{\xi_1})$$

by contradiction

(14)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$$

by Lemma 2.0.20 on

(11) and (13)

Case (9h).

$$(7) \ \dot{\xi} = \operatorname{inr}(\dot{\xi}_2)$$

by assumption

(8)
$$\tau = (\tau_1 + \tau_2)$$

by assumption

(9)
$$\dot{\xi}_2 : \tau_2$$

by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

(10)
$$e = \{ \|u, \|e_0\|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \} \}$$

by assumption

(11)
$$e$$
 notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

Assume $e \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inr}(\dot{\xi}_2)$$

by contradiction

By case analysis on the value of $refutable_{?}(inr(\dot{\xi}_{2}))$.

inr is refutable

Case $refutable_7(inr(\dot{\xi}_2)) = true.$

- (13) $refutable_?(inr(\dot{\xi}_2)) = true$ by assumption
- (14) $\operatorname{inr}(\dot{\xi}_2)$ refutable? by Lemma 2.0.14 on (13)
- (15) $e \models_? \operatorname{inr}(\dot{\xi}_2)$ by Rule (17b) on (11) and (14)
- (16) $e \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (15)

Case $refutable_?(inr(\dot{\xi}_2)) = false.$

- (13) $refutable_{?}(inr(\dot{\xi}_{2})) = false$ by assumption
- (14) $\underline{\operatorname{inr}(\dot{\xi}_2)}$ refutable? by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

- (15) $\operatorname{inr}(\dot{\xi}_2)$ refutable? by assumption Contradicts (14).
- (16) $e \not\models_? \operatorname{inr}(\dot{\xi}_2)$ by contradiction (17) $e \not\models_?^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Lemma 2.0.20 on (12) and (16)

Case (20j).

(10)
$$e = inl_{\tau_2}(e_1)$$

by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inr}(\dot{\xi_2})$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

- (12) $\operatorname{inl}_{\tau_2}(e_1)$ notintro
- by assumption

By rule induction over Rules (27) on (12), no case applies due to syntactic contradiction.

- (13) $\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inr}(\dot{\xi_2})$
- by contradiction

(14)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 by Lemma 2.0.20 on (11) and (13)

Case (20k).

- (10) $e = \inf_{\tau_1}(e_2)$ by assumption (11) $\cdot : \Delta \vdash e_2 : \tau_2$ by assumption
- (12) e_2 final by Lemma 4.0.4 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_2 \models \dot{\xi}_2$, $e_2 \models_? \dot{\xi}_2$, and $e_2 \not\models_?^{\dagger} \dot{\xi}_2$ holds. By case analysis on which one holds.

Case $e_2 \models \dot{\xi}_2$.

- (13) $e_2 \models \dot{\xi}_2$ by assumption (14) $e_2 \not\models_? \dot{\xi}_2$ by assumption
- (15) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$ by Rule (15g) on (13)
- (16) $\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\dot{\xi_2})$ by Rule (18b) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (17) on it, only two cases apply.

Case (17b).

(17) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17f).

(17)
$$e_2 \models_? \dot{\xi}_2$$
 Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_{?} \operatorname{inr}(\dot{\xi_2})$$
 by contradiction

Case $e_2 \models_? \dot{\xi}_2$.

- (13) $e_2 \not\models \dot{\xi}_2$ by assumption (14) $e_2 \models_? \dot{\xi}_2$ by assumption
- (15) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi_2})$ by Rule (17f) on (14)
- (16) $\inf_{\tau_1}(e_2) \models_{?}^{\dagger} \inf(\dot{\xi_2})$ by Rule (18a) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15h).

(17)
$$e_2 \models \dot{\xi}_2$$

Contradicts (13).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\dot{\xi}_2)$$
 by contradiction Case $e_2 \not\models_7^{\dagger} \dot{\xi}_2$.

(13)
$$e_2 \not\models \dot{\xi}_2$$
 by assumption (14) $e_2 \not\models_{\dot{\gamma}} \dot{\xi}_2$ by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15h).

(15)
$$e_2 \models \dot{\xi}_2$$

Contradicts (13).

(16) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\dot{\xi}_2)$ by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(17) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17f).

(17)
$$e_2 \models_? \dot{\xi}_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\dot{\xi}_2)$$
 by contradiction
(19) $\operatorname{inl}_{\tau_2}(e_1) \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ by Lemma 2.0.20 on
(16) and (18)

Case (15i).

(7)
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption
(8) $\tau = (\tau_1 \times \tau_2)$ by assumption
(9) $\dot{\xi}_1 : \tau_1$ by assumption
(10) $\dot{\xi}_2 : \tau_2$ by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20l),(20m).

$$(11) \ e = (\!()^u, (\!(e_0)\!)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{\hat{rs}\}$$

by assumption

(12)
$$e$$
 notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

(13)
$$e$$
 indet by Lemma 4.0.10 on

(6) and (12)

$$(14)$$
 prl (e) indet
 by Rule $(25g)$ on (13)
 (15) prl (e) final
 by Rule $(26b)$ on (14)
 (16) prr (e) indet
 by Rule $(25h)$ on (13)

(17) prr(e) final by Rule (26b) on (16) (18) \cdot ; $\Delta \vdash prl(e) : \tau_1$ by Rule (20h) on (5) (19) \cdot ; $\Delta \vdash prr(e) : \tau_2$ by Rule (20i) on (5)

By inductive hypothesis on (9) and (18) and (15), exactly one of $\operatorname{prl}(e) \models \dot{\xi}_1$, $\operatorname{prl}(e) \models_? \dot{\xi}_1$, and $\operatorname{prl}(e) \not\models_?^{\dagger} \dot{\xi}_1$ holds. By inductive hypothesis on (10) and (19) and (17), exactly one of $\operatorname{prr}(e) \models_{\dot{\xi}_2}$, $\operatorname{prr}(e) \models_? \dot{\xi}_2$, and $\operatorname{prr}(e) \not\models_?^{\dagger} \dot{\xi}_2$ holds.

By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $prl(e) \models \dot{\xi}_1, prr(e) \models \dot{\xi}_2.$

(20) $prl(e) \models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption (22) $\operatorname{prr}(e) \models \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption (24) $e \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (15j) on (12) and (20) and (22) (25) $e \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18b) on (24) (26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Lemma 2.0.18 on (12) and (24)

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption Contradicts (26).

(28) $e \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$

by contradiction

Case $prl(e) \models \dot{\xi}_1, prr(e) \models_? \dot{\xi}_2$.

(20) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption (22) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

(24) $prr(e) \models \dot{\xi}_2$ by assumption Contradicts (22)

(25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

By rule induction over Rules (17) on (23), only one case applies.

assume no "or" and "and" in pair

Case (17b).

- (26) $\dot{\xi}_2$ refutable? by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Rule (11e) on (26)
- (28) $e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17b) on (12) and (27)
- (29) $e \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (28)

Case $prl(e) \models \dot{\xi}_1, prr(e) \not\models_?^{\dagger} \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption
- (21) $prl(e) \not\models_? \dot{\xi}_1$ by assumption
- (22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$ by assumption
- (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

- (24) $prr(e) \models \dot{\xi}_2$ by assumption Contradicts (22).
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption

By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

- (27) $\dot{\xi}_1$ refutable? by assumption (28) prl(e) notintro by Rule (27e)
- (29) $\operatorname{prl}(e) \models_{?} \dot{\xi}_{1}$ by Rule (17b) on (28) and (27)

Contradicts (21).

Case (11e).

- (27) $\dot{\xi}_2$ refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $\operatorname{prr}(e) \models_{?} \dot{\xi}_{2}$ by Rule (17b) on (28) and (27)

Contradicts (23).

- (30) $e \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction
- (31) $e \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \models_? \dot{\xi}_1, prr(e) \models \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \models_? \dot{\xi}_1$ by assumption
- (22) $\operatorname{prr}(e) \models \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

- (24) $prl(e) \models \dot{\xi}_1$ by assumption Contradicts (20).
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction By rule induction over Rules (17) on (21), only one case applies.

Case (17b).

- (26) $\dot{\xi}_1$ refutable? by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Rule (11e) on (26)
- (28) $e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17b) on (12) and (27)
- (29) $e \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (28)

Case $prl(e) \models_? \dot{\xi}_1, prr(e) \models_? \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption
- (21) $\operatorname{prl}(e) \models_{?} \dot{\xi}_{1}$ by assumption
- (22) $prr(e) \not\models \dot{\xi}_2$ by assumption (23) $prr(e) \models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

- (24) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption Contradicts (20).
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction By rule induction over Rules (17) on (23), only one case applies.

Case (17b).

- (26) $\dot{\xi}_2$ refutable?
- (27) $(\dot{\xi}_1,\dot{\xi}_2)$ refutable?
- (28) $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$

by assumption

by Rule (11e) on (26)

by Rule (17b) on (12) and (27)

assume no "or" and "and" in pair

assume no "or" and

"and" in

pair

(29)
$$e \models_{?}^{\uparrow} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (18a) on (28)

Case $prl(e) \models_? \dot{\xi}_1, prr(e) \not\models_?^{\dagger} \dot{\xi}_2.$

- (20) $prl(e) \not\models \dot{\xi}_1$ by assumption (21) $prl(e) \models_? \dot{\xi}_1$ by assumption
- (22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

(24)
$$prl(e) \models \dot{\xi}_1$$
 by assumption Contradicts (20)

(25)
$$e \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

By rule induction over Rules (17) on (21), only one case applies.

Case (17b).

- (26) $\dot{\xi}_1$ refutable? by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Rule (11e) on (26)

assume no "or" and

"and" in

pair

(28) $e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17b) on (12) and (27)

(29)
$$e \models^{\dagger}_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (18a) on (28)

Case $prl(e) \not\models_?^{\dagger} \dot{\xi}_1, prr(e) \models \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption
- (21) $prl(e) \not\models_? \dot{\xi}_1$ by assumption (22) $prr(e) \models \dot{\xi}_2$ by assumption
- (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

(24)
$$prl(e) \models \dot{\xi}_1$$
 by assumption Contradicts (20)

(25)
$$e \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(26)
$$(\dot{\xi}_1, \dot{\xi}_2)$$
 refutable? by assumption By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

- (27) $\dot{\xi}_1$ refutable? by assumption (28) prl(e) notintro by Rule (27e)
- (29) $\operatorname{prl}(e) \models_? \dot{\xi}_1$ by Rule (17b) on (28) and (27)

Contradicts (21).

Case (11e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $\operatorname{prr}(e) \models_? \dot{\xi}_2$ by Rule (17b) on (28) and (27)

Contradicts (23).

(30) $e \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction (31) $e \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \not\models_{?}^{\dagger} \dot{\xi}_{1}, prr(e) \models_{?} \dot{\xi}_{2}.$

(20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption (22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

- (24) $prl(e) \models \dot{\xi}_1$ by assumption Contradicts (20).
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction By rule induction over Rules (17) on (23), only one case applies.

Case (17b).

- (26) $\dot{\xi}_2$ refutable?
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable?
- (28) $e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$
- (29) $e \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$
- by assumption
- by Rule (11e) on (26)

assume no "or" and

"and" in

pair

- by Rule (17b) on (12)
- and (27)
- by Rule (18a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \dot{\xi}_{1}, prr(e) \not\models_{?}^{\dagger} \dot{\xi}_{2}.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$
- (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$

- by assumption
- by assumption

(22) $prr(e) \not\models \dot{\xi}_2$ by assumption (23) $prr(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

- (24) $prl(e) \models \dot{\xi}_1$ by assumption Contradicts (20)
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction Assume $e \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

- (27) $\dot{\xi}_1$ refutable? by assumption (28) prl(e) notintro by Rule (27e)
- (29) $\operatorname{prl}(e) \models_{?} \dot{\xi}_{1}$ by Rule (17b) on (28) and (27)

Contradicts (21).

Case (11e).

- (27) $\dot{\xi}_2$ refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $\operatorname{prr}(e) \models_{?} \dot{\xi}_{2}$ by Rule (17b) on (28) and (27)

Contradicts (23).

(30) $e \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction (31) $e \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (25) and (30)

Case (20g).

 $\begin{array}{lll} (11) & e = (e_1, e_2) & \text{by assumption} \\ (12) & \cdot; \Delta \vdash e_1 : \tau_1 & \text{by assumption} \\ (13) & \cdot; \Delta \vdash e_2 : \tau_2 & \text{by assumption} \\ (14) & e_1 \text{ final} & \text{by Lemma 4.0.5 on (6)} \\ (15) & e_2 \text{ final} & \text{by Lemma 4.0.5 on (6)} \\ \end{array}$

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 \models \dot{\xi}_1$, $e_1 \models_? \dot{\xi}_1$, and $e_1 \models \dot{\xi}_1$ holds.

By inductive hypothesis on (10) and (13) and (15), exactly one of $e_2 \models \dot{\xi}_2$, $e_2 \models_? \dot{\xi}_2$, and $e_2 \models \dot{\xi}_2$ holds.

By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $e_1 \models \dot{\xi}_1, e_2 \models \dot{\xi}_2$.

(16) $e_1 \models \dot{\xi}_1$ by assumption (17) $e_1 \not\models_? \dot{\xi}_1$ by assumption

(18) $e_2 \models \dot{\xi}_2$ by assumption

(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption

(20) $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (15i) on (16) and (18)

(21) $(e_1, e_2) \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18b) on (20)

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption Contradicts (17).

Case (17h).

(22) $e_2 \models_? \dot{\xi_2}$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \dot{\xi_1}$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Case $e_1 \models \dot{\xi}_1, e_2 \models_? \dot{\xi}_2$.

(16) $e_1 \models \dot{\xi}_1$ by assumption

(17) $e_1 \not\models_? \dot{\xi}_1$ by assumption

(18) $e_2 \not\models \dot{\xi}_2$ by assumption

(19) $e_2 \models_? \dot{\xi}_2$ by assumption

(20) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17h) on (16) and (19)

(21) $(e_1, e_2) \models^{\dagger}_{?} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (20)

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.9.

Case (15i).

(22)
$$e_2 \models \dot{\xi}_2$$
 by assumption Contradicts (18).

(23)
$$(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Case $e_1 \models \dot{\xi}_1, e_2 \not\models_?^{\dagger} \dot{\xi}_2$.

(16)
$$e_1 \models \dot{\xi}_1$$
 by assumption
(17) $e_1 \not\models_? \dot{\xi}_1$ by assumption
(18) $e_2 \not\models \dot{\xi}_2$ by assumption
(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (15i).

(20)
$$e_2 \models \dot{\xi}_2$$
 by assumption Contradicts (18).

(21)
$$(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.9.

Case (17g).

(22)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption Contradicts (17).

Case (17h).

(22)
$$e_2 \models_? \dot{\xi}_2$$
 by assumption Contradicts (19).

Case (17i).

(22)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption Contradicts (17).

(23)
$$(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Case
$$e_1 \models_? \dot{\xi}_1, e_2 \models_{?} \dot{\xi}_2$$
.

(16) $e_1 \not\models_{\dot{1}} \dot{\xi}_1$ by assumption
(17) $e_1 \models_? \dot{\xi}_1$ by assumption
(18) $e_2 \models_{\dot{2}} \dot{\xi}_2$ by assumption
(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption
(20) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17g) on (17) and (18)

(21) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (20)

Assume $(e_1, e_2) \models_{\dot{1}} \dot{\xi}_1, \dot{\xi}_2$ by assumption
it, only two cases apply.

Case (15j).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(22) $e_1 \models_{\dot{1}} \dot{\xi}_1$ by assumption

Contradicts (16).

(23) $(e_1, e_2) \not\models_{\dot{1}} \dot{\xi}_2$.

(16) $e_1 \not\models_{\dot{1}} \dot{\xi}_1$ by assumption

Contradicts (16).

Case $e_1 \models_? \dot{\xi}_1, e_2 \models_? \dot{\xi}_2$.

(16) $e_1 \not\models_{\dot{1}} \dot{\xi}_1$ by assumption

(17) $e_1 \models_? \dot{\xi}_1$ by assumption
(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption
(20) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17i) on (17) and (19)

(21) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (20)

Assume $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (20)

Assume $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (18a) on (20)

Assume $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(22) $e_1 \not\models_{\dot{\xi}} \dot{\xi}_1$ by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(22) $e_1 \not\models_{\dot{\xi}} \dot{\xi}_1$ by assumption

Contradicts Lemma 4.0.9.

Case $e_1 \models_? \dot{\xi}_1, e_2 \not\models_?^{\dagger} \dot{\xi}_2$.	
(16) $e_1 \not\models \dot{\xi}_1$	by assumption
$(17) e_1 \models_? \dot{\xi}_1$	by assumption
$(18) e_2 \not\models \dot{\xi}_2$	by assumption
$(19) e_2 \not\models_? \dot{\xi}_2$	by assumption
Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By ru	
it, only two cases apply.	` '
Case $(15j)$.	
$(20)\ (e_1,e_2)\ { t notintro}$	by assumption
Contradicts Lemma 4.0.9.	
Case $(15i)$.	
(20) $e_1 \models \dot{\xi}_1$	by assumption
Contradicts (16).	
. ,	
$(21) (e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$	by contradiction
Assume $(e_1, e_2) \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By	rule induction over Rules (17)
on it, the following cases apply.	
Case $(17b)$.	
$(22)\ (e_1,e_2)$ notintro	by assumption
Contradicts Lemma 4.0.9.	
Case $(17g)$.	
$(22) e_2 \models \dot{\xi}_2$	by assumption
Contradicts (18).	
Case $(17h)$.	
(22) $e_2 \models_? \dot{\xi}_2$	by assumption
Contradicts (19).	· -
Case (17i).	
$(22) e_2 \models_? \dot{\xi}_2$	by assumption
Contradicts (19).	T .
(1)	
(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$	by contradiction
$(24) (e_1, e_2) \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$	by Lemma 2.0.20 on
() (1) 2) ((31) 32)	(21) and (23)
Case $e_1 \not\models_?^\dagger \dot{\xi}_1, e_2 \models \dot{\xi}_2$.	
$(16) e_1 \not\models \dot{\xi}_1$	by assumption
$(10) e_1 \not\models \zeta_1$ $(17) e_1 \not\models_? \dot{\xi}_1$	by assumption
$(18) e_1 \models \dot{\xi}_2$	by assumption
$(19) e_2 \models \zeta_2 \\ (19) e_2 \not\models_? \dot{\xi}_2$	by assumption
(10) 02 V : S2	

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(20) $e_1 \models \dot{\xi}_1$ by assumption Contradicts (16).

(21) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \dot{\xi_2}$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

(24) $(e_1, e_2) \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (21) and (23)

Case $e_1 \not\models_?^\dagger \dot{\xi}_1, e_2 \models_? \dot{\xi}_2$.

(16) $e_1 \not\models \dot{\xi}_1$ by assumption (17) $e_1 \not\models_? \dot{\xi}_1$ by assumption (18) $e_2 \not\models \dot{\xi}_2$ by assumption (19) $e_2 \models_? \dot{\xi}_2$ by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

Contradicts (18).
$$(21) \ (e_1,e_2) \not\models (\dot{\xi}_1,\dot{\xi}_2) \qquad \text{by contradiction}$$
 Assume $(e_1,e_2) \models_? (\dot{\xi}_1,\dot{\xi}_2)$. By rule induction over Rules (17) on it, the following cases apply.
$$\textbf{Case (17b)}.$$

$$(22) \ (e_1,e_2) \ \text{notintro} \qquad \text{by assumption}$$
 Contradicts Lemma 4.0.9.
$$\textbf{Case (17g)}.$$

$$(22) \ e_1 \models_? \dot{\xi}_1 \qquad \text{by assumption}$$
 Contradicts (17).
$$\textbf{Case (17h)}.$$

$$(22) \ e_1 \models_? \dot{\xi}_1 \qquad \text{by assumption}$$
 Contradicts (16).
$$\textbf{Case (17i)}.$$

$$(22) \ e_1 \models_? \dot{\xi}_1 \qquad \text{by assumption}$$
 Contradicts (17).
$$(23) \ (e_1,e_2) \not\models_? (\dot{\xi}_1,\dot{\xi}_2) \qquad \text{by contradiction}$$
 Contradicts (17).
$$(23) \ (e_1,e_2) \not\models_? (\dot{\xi}_1,\dot{\xi}_2) \qquad \text{by contradiction}$$

$$(24) \ (e_1,e_2) \not\models_? (\dot{\xi}_1,\dot{\xi}_2) \qquad \text{by Lemma 2.0.20 on}$$

$$(21) \ \text{and (23)}$$

$$\textbf{Case } e_1 \not\models_? \dot{\xi}_1, e_2 \not\models_? \dot{\xi}_2.$$

$$(16) \ e_1 \not\models_? \dot{\xi}_1 \qquad \text{by assumption}$$

$$(17) \ e_1 \not\models_? \dot{\xi}_1 \qquad \text{by assumption}$$

$$(18) \ e_2 \not\models_? \dot{\xi}_2 \qquad \text{by assumption}$$

$$(19) \ e_2 \not\models_? \dot{\xi}_2 \qquad \text{by assumption}$$

$$(19) \ e_2 \not\models_? \dot{\xi}_2 \qquad \text{by assumption}$$
 Assume $(e_1,e_2) \models(\dot{\xi}_1,\dot{\xi}_2)$. By rule induction over Rules (15) on it, only two cases apply. }

by assumption

Case (15j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

(20) $e_2 \models \dot{\xi}_2$

Case (15i).

(20) $e_2 \models \dot{\xi}_2$ by assumption Contradicts (18).

(21) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22)
$$(e_1, e_2)$$
 notintro

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \dot{\xi}_2$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

(24) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (21) and (23)

Definition 1.1.1 (Entailment of Constraints). Suppose that $\dot{\xi}_1 : \tau$ and $\dot{\xi}_2 : \tau$. Then $\dot{\xi}_1 \models \dot{\xi}_2$ iff for all e such that $\cdot : \Delta \vdash e : \tau$ and e val we have $e \models_{?}^{\dagger} \dot{\xi}_1$ implies $e \models \dot{\xi}_2$

Definition 1.1.2 (Potential Entailment of Constraints). Suppose that $\dot{\xi}_1: \tau$ and $\dot{\xi}_2: \tau$. Then $\dot{\xi}_1 \models_{?}^{\dagger} \dot{\xi}_2$ iff for all e such that $\cdot; \Delta \vdash e: \tau$ and e final we have $e \models_{?}^{\dagger} \dot{\xi}_1$ implies $e \models_{?}^{\dagger} \dot{\xi}_2$

Corollary 1.1.1. Suppose that $\dot{\xi} : \tau \text{ and } \cdot ; \Delta \vdash e : \tau \text{ and } e \text{ final. Then } \top \models_{?}^{\dagger} \dot{\xi} \text{ implies } e \models_{?}^{\dagger} \dot{\xi}$

Proof.

(1) $\dot{\xi}:\tau$	by assumption
$(2) \ \cdot ; \Gamma \vdash e : \tau$	by assumption
(3) e final	by assumption
$(4) \top \models^{\dagger}_{?} \dot{\xi}$	by assumption
(5) $e_1 \models \top$	by Rule (15a)
(6) $e_1 \models_?^\dagger \top$	by Rule (18b) on (5)
$(7) \ \ \top : \tau$	by Rule (9a)
$(8) e_1 \models^{\dagger}_{?} \dot{\xi_r}$	by Definition 2.1.2 of (4) on (7) and (1) and (2) and (3) and (6)

2 Match Constraint Language

 $\operatorname{CTTruth}$

$$\overline{\top : \tau}$$
 (9a)

CTFalsity

$$\frac{}{\perp : \tau} \tag{9b}$$

CTNum

$$\underline{\underline{n}:\mathtt{num}}$$
 (9c)

CTNotNum

$$\overline{\mathscr{H}:\mathtt{num}}$$
 (9d)

CTAnd

$$\frac{\xi_1 : \tau \qquad \xi_2 : \tau}{\xi_1 \wedge \xi_2 : \tau} \tag{9e}$$

CTOr

$$\frac{\xi_1 : \tau \qquad \xi_2 : \tau}{\xi_1 \vee \xi_2 : \tau} \tag{9f}$$

CTInl

$$\frac{\xi_1:\tau_1}{\operatorname{inl}(\xi_1):(\tau_1+\tau_2)}\tag{9g}$$

 CTInr

$$\frac{\xi_2 : \tau_2}{\operatorname{inr}(\xi_2) : (\tau_1 + \tau_2)} \tag{9h}$$

CTPair

$$\frac{\xi_1 : \tau_1}{(\xi_1, \xi_2) : (\tau_1 \times \tau_2)} \qquad (9i)$$

 $|\overline{\xi_1} = \xi_2|$ dual of ξ_1 is ξ_2

```
\overline{\underline{n}} = \underline{\varkappa}
                                                                                                                                 (10c)
                                             \underline{\underline{n}} = \underline{n}
                                                                                                                                (10d)
                                     \overline{\xi_1 \wedge \xi_2} = \overline{\xi_1} \vee \overline{\xi_2}
                                                                                                                                 (10e)
                                     \overline{\xi_1 \vee \xi_2} = \overline{\xi_1} \wedge \overline{\xi_2}
                                                                                                                                 (10f)
                                    \overline{\mathtt{inl}(\xi_1)} = \mathtt{inl}(\overline{\xi_1}) \vee \mathtt{inr}(\top)
                                                                                                                                (10g)
                                    \overline{\mathtt{inr}(\xi_2)} = \mathtt{inr}(\overline{\xi_2}) \vee \mathtt{inl}(\top)
                                                                                                                                (10h)
                                     \overline{(\xi_1,\xi_2)} = (\xi_1,\overline{\xi_2}) \vee (\overline{\xi_1},\xi_2) \vee (\overline{\xi_1},\overline{\xi_2})
                                                                                                                                 (10i)
\xi refutable?
                               \xi is refutable
                                                       RXNum
                                                                                                                                 (11a)
                                                        \underline{n} refutable?
                                                  RXInl
                                                                                                                                (11b)
                                                   inl(\xi) refutable?
                                                   RXInr
                                                                                                                                 (11c)
                                                   inr(\xi) refutable?
                                                  RXPairL
                                                       \xi_1 refutable?
                                                                                                                                (11d)
                                                   (\xi_1,\xi_2) refutable?
                                                  RXPairR
                                                       \xi_2 refutable?
                                                                                                                                 (11e)
                                                   (\xi_1, \xi_2) refutable?
                                      RXOr
                                      \xi_1 refutable? \xi_2 refutable?
                                                                                                                                 (11f)
                                                   \xi_1 \vee \xi_2 \text{ refutable}_?
refutable_?(\xi)
                                refutable_{?}(\underline{n}) = true
                                                                                                                                 (12a)
                                refutable_{?}(\cancel{x}) = true
                                                                                                                                (12b)
                                 refutable_{?}(?) = true
                                                                                                                                 (12c)
                       refutable_{?}(inl(\xi)) = refutable_{?}(\xi)
                                                                                                                                (12d)
                       \mathit{refutable}_?(\mathtt{inr}(\xi)) = \mathit{refutable}_?(\xi)
                                                                                                                                 (12e)
                       refutable_{?}((\xi_{1},\xi_{2})) = refutable_{?}(\xi_{1}) \text{ or } refutable_{?}(\xi_{2})
                                                                                                                                 (12f)
                       refutable_{?}(\xi_{1} \vee \xi_{2}) = refutable_{?}(\xi_{1}) and refutable_{?}(\xi_{2})
                                                                                                                                (12g)
           Otherwise refutable_{?}(\xi) = false
                                                                                                                                (12h)
```

 $\overline{\top} = \bot$

 $\overline{\perp} = \top$

(10a)

(10b)

$\dot{\top}(\xi_1) = \xi_2$

$$\dot{\top}(\top) = \top \tag{13a}$$

$$\dot{\top}(\bot) = \bot \tag{13b}$$

$$\dot{\top}(?) = \top \tag{13c}$$

$$\dot{\top}(\underline{n}) = \underline{n} \tag{13d}$$

$$\dot{\top}(\underline{n}) = \underline{n} \tag{13e}$$

$$\dot{\top}(\xi_1 \land \xi_2) = \dot{\top}(\xi_1) \land \dot{\top}(\xi_2) \tag{13f}$$

$$\dot{\top}(\xi_1 \lor \xi_2) = \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2) \tag{13g}$$

$$\dot{\top}(\sin(\xi)) = \sin(\dot{\top}(\xi)) \tag{13h}$$

$$\dot{\top}(\sin(\xi)) = \sin(\dot{\top}(\xi)) \tag{13i}$$

$$\dot{\top}((\xi_1, \xi_2)) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2)) \tag{13j}$$

(13j)

$\overline{\dot{\bot}(\xi_1) = \xi_2}$

$$\dot{\bot}(\top) = \top \tag{14a}$$

$$\dot{\bot}(\bot) = \bot \tag{14b}$$

$$\dot{\bot}(?) = \bot \tag{14c}$$

$$\dot{\bot}(\underline{n}) = \underline{n} \tag{14d}$$

$$\dot{\bot}(\underline{\varkappa}) = \underline{\varkappa} \tag{14e}$$

$$\dot{\bot}(\xi_1 \land \xi_2) = \dot{\bot}(\xi_1) \land \dot{\bot}(\xi_2) \tag{14f}$$

$$\dot{\bot}(\xi_1 \lor \xi_2) = \dot{\bot}(\xi_1) \lor \dot{\bot}(\xi_2) \tag{14g}$$

$$\dot{\bot}(\operatorname{inl}(\xi)) = \operatorname{inl}(\dot{\bot}(\xi)) \tag{14h}$$

$$\dot{\bot}(\operatorname{inr}(\xi)) = \operatorname{inr}(\dot{\bot}(\xi)) \tag{14i}$$

$$\dot{\bot}((\xi_1, \xi_2)) = (\dot{\bot}(\xi_1), \dot{\bot}(\xi_2)) \tag{14j}$$

e satisfies ξ $e \models \xi$

CSTruth(15a) $\overline{e \models \top}$

CSNum (15b) $\underline{n \models \underline{n}}$

CSNotNum $n_1 \neq n_2$ (15c) $\underline{n_1} \models \underline{n_2}$

CSAnd
$$\frac{e \models \xi_1 \quad e \models \xi_2}{e \models \xi_1 \land \xi_2} \tag{15d}$$

CSOrL
$$\frac{e \models \xi_1}{e \models \xi_1 \lor \xi_2}$$
(15e)

CSOrR
$$\frac{e \models \xi_2}{e \models \xi_1 \lor \xi_2} \tag{15f}$$

CSInl

$$\frac{e_1 \models \xi_1}{\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)} \tag{15g}$$

$$\frac{\text{CSInr}}{e_2 \models \xi_2} \\ \frac{e_2 \models \xi_2}{\text{inr}_{\tau_1}(e_2) \models \text{inr}(\xi_2)}$$
 (15h)

CSPair
$$\frac{e_1 \models \xi_1 \qquad e_2 \models \xi_2}{(e_1, e_2) \models (\xi_1, \xi_2)} \tag{15i}$$

CSNotIntroPair

$$\frac{e \text{ notintro} \qquad \text{prl}(e) \models \xi_1 \qquad \text{prr}(e) \models \xi_2}{e \models (\xi_1, \xi_2)} \tag{15j}$$

 $\mathit{satisfy}(e,\xi)$

$$satisfy(\underline{n_1},\underline{n_2}) = (n_1 = n_2) \tag{16b}$$

$$satisfy(\underline{n_1},\underline{n_2}) = (n_1 \neq n_2) \tag{16c}$$

$$satisfy(e,\xi_1 \land \xi_2) = satisfy(e,\xi_1) \text{ and } satisfy(e,\xi_2) \tag{16d}$$

$$satisfy(e,\xi_1 \land \xi_2) = satisfy(e,\xi_1) \text{ or } satisfy(e,\xi_2) \tag{16d}$$

$$satisfy(e,\xi_1 \lor \xi_2) = satisfy(e,\xi_1) \text{ or } satisfy(e,\xi_2) \tag{16e}$$

$$satisfy(inl_{\tau_2}(e_1), inl(\xi_1)) = satisfy(e_1,\xi_1) \tag{16f}$$

$$satisfy(inl_{\tau_2}(e_1), inl(\xi_2)) = satisfy(e_2,\xi_2) \tag{16h}$$

$$satisfy((e_1,e_2),(\xi_1,\xi_2)) = satisfy(e_1,\xi_1) \text{ and } satisfy(e_2,\xi_2) \tag{16h}$$

$$satisfy((e_1,e_2),(\xi_1,\xi_2)) = satisfy(prl((e_1,e_1)),\xi_1) \text{ and } satisfy(prr((e_1,e_2)),\xi_1) \tag{16i}$$

$$satisfy(e_1(e_2),(\xi_1,\xi_2)) = satisfy(prl(e_1(e_2)),\xi_1) \text{ and } satisfy(prr(e_1(e_2)),\xi_2) \tag{16h}$$

$$satisfy(match(e)\{r^s\},(\xi_1,\xi_2)) = satisfy(prl(match(e)\{r^s\}),\xi_2) \tag{16l}$$

$$satisfy(prl(e),(\xi_1,\xi_2)) = satisfy(prl(match(e)\{r^s\}),\xi_2) \tag{16l}$$

$$satisfy(prl(e),(\xi_1,\xi_2)) = satisfy(prl(match(e)\{r^s\}),\xi_2) \tag{16l}$$

$$satisfy(prl(e),(\xi_1,\xi_2)) = satisfy(prl(prl(e)),\xi_1) \text{ and } satisfy(prr(prl(e)),\xi_2) \tag{16m}$$

$$satisfy(prr(e),(\xi_1,\xi_2)) = satisfy(prl(prl(e)),\xi_1) \text{ and } satisfy(prr(prl(e)),\xi_2) \tag{16h}$$

$$Otherwise \ satisfy(e,\xi) = false \tag{16o}$$

$$e \mapsto \xi$$

$$\frac{CMSUnknown}{e \mapsto \gamma}$$

$$\frac{CMSUnknown}{e \mapsto \gamma}$$

$$\frac{e \mapsto \xi}{e}$$

$$\frac{CMSUnknown}{e \mapsto \gamma}$$

$$\frac{e \mapsto \xi}{e}$$

$$\frac{CMSOrL}{e \mapsto \gamma}$$

$$\frac{e \mapsto \xi}{e}$$

$$\frac{CMSOrL}{e \mapsto \gamma}$$

$$\frac{e \mapsto \xi}{e}$$

$$\frac{CHSOrL}{e \mapsto \gamma}$$

$$\frac{e \mapsto \xi}{e}$$

$$\frac{e \mapsto$$

(16a)

 $satisfy(e, \top) = true$

$$\frac{e_2 \models_? \xi_2}{\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)} \tag{17f}$$

 ${\rm CMSPairL}$

$$\frac{e_1 \models_? \xi_1 \qquad e_2 \models_\xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)} \tag{17g}$$

CMSPairR

$$\frac{e_1 \models \xi_1 \qquad e_2 \models_? \xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)} \tag{17h}$$

CMSPair

$$\frac{e_1 \models_? \xi_1 \qquad e_2 \models_? \xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)} \tag{17i}$$

 $e \models^{\dagger}_{?} \xi$ e satisfies or may satisfy ξ

CSMSMay

$$\frac{e \models_? \xi}{e \models_?^\dagger \xi} \tag{18a}$$

CSMSSat

$$\frac{e \models \xi}{e \models_{7}^{+} \xi} \tag{18b}$$

Lemma 2.0.1. $e \not\models \bot$

Proof. By rule induction over Rules (15), we notice that $e \models \bot$ is in syntactic contradiction with all rules, hence not derivable.

Lemma 2.0.2. $e \not\models_? \bot$

Proof. Assume $e \models_? \bot$. By rule induction over Rules (17) on $e \models_? \bot$, only one case applies.

Case (17b).

(1) \perp refutable?

by assumption

By rule induction over Rules (11) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \bot$ is not derivable.

Lemma 2.0.3. $e \not\models_? \top$

Proof. Assume $e \models_? \top$. By rule induction over Rules (17) on $e \models_? \top$, only one case applies.

Case (17b).

(1) \top refutable?

by assumption

By rule induction over Rules (11) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \top$ is not derivable.

Lemma 2.0.4. $e \not\models ?$

Proof. By rule induction over Rules (15), we notice that $e \models ?$ is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.5.
$$e \models^{\dagger}_{?} \xi \text{ iff } e \models^{\dagger}_{?} \xi \lor \bot$$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e \models_{?}^{\dagger} \xi$

by assumption

By rule induction over Rules (18) on (1).

Case (18a).

(2) $e \models_? \xi$

by assumption

(3) $e \models_? \xi \lor \bot$

by Rule (17c) on (2)

(4) $e \models_2^{\dagger} \xi \lor \bot$

and Lemma 2.0.1 by Rule (18a) on (3)

- Case (18b).

by assumption

(3) $e \models \xi \lor \bot$

(2) $e \models \xi$

by Rule (15e) on (2)

(4) $e \models_{?}^{\dagger} \xi \lor \bot$

by Rule (18b) on (3)

- 2. Necessity:
 - (1) $e \models^{\dagger}_{?} \xi \lor \bot$

by assumption

By rule induction over Rules (18) on (1).

Case (18a).

(2) $e \models_? \xi \lor \bot$

by assumption

By rule induction over Rules (17) on (2), only two of them apply.

Case (17c).

(3) $e \models_? \xi$

by assumption

(4) $e \models_2^{\dagger} \xi$

by Rule (18a) on (3)

Case (17d).

(3)
$$e \models_? \bot$$
 by assumption
(4) $e \not\models_? \bot$ by Lemma 2.0.2

(3) contradicts (4).

Case (18b).

(2)
$$e \models \xi \lor \bot$$

by assumption

By rule induction over Rules (15) on (2), only two of them apply.

Case (15e).

(3)
$$e \models \xi$$

by assumption

(4)
$$e \models_{?}^{\dagger} \xi$$

by Rule (18b) on (3)

Case (15f).

(3)
$$e \models \bot$$

by assumption

(4)
$$e \not\models \bot$$

by Lemma 2.0.1

(3) contradicts (4).

Corollary 2.0.1. $\top \models_{?}^{\dagger} \xi \text{ iff } \top \models_{?}^{\dagger} \xi \vee \bot$

Proof. By Definition 2.1.2 and Lemma 2.0.5.

Lemma 2.0.6. Suppose that $\xi_1 : \tau$ and $\xi_2 : \tau$. Then $\xi_1 \not\models \xi_2$ iff $\xi_1 \not\models \xi_2 \lor \bot$

Proof.

(1)
$$\xi_1: \tau$$
 by assumption

(2)
$$\xi_2: \tau$$
 by assumption

(3)
$$\bot:\tau$$
 by Rule (9b)

(4)
$$\xi_2 \lor \bot : \tau$$
 by Rule (9f) on (2) and (3)

Then we prove sufficiency and necessity separately.

1. Sufficiency:

(5)
$$\xi_1 \not\models \xi_2$$

by assumption

To prove $\xi_1 \not\models \xi_2 \lor \bot$, assume $\xi_1 \models \xi_2 \lor \bot$.

(6)
$$\xi_1 \models \xi_2 \lor \bot$$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \xi_{1}$ implies

(7)
$$e \models \xi_2 \lor \bot$$

by Definition 2.1.1 on

(1) and (4) and (6)

By rule induction over Rules (15) on (7).

Case (15e).

 $(8) e \models \xi_2$ $(9) \xi_1 \models \xi_2$

by assumption

by Definition 2.1.1 on

(8)

(5) contradicts (9).

Case (15f).

 $(8) e \models \bot$ $(9) e \not\models \bot$

by assumption

by Lemma 2.0.1

(8) contradicts (9).

The conclusion holds as follows:

- (a) $\xi_1 \not\models \xi_2 \lor \bot$
- 2. Necessity:
 - (5) $\xi_1 \not\models \xi_2 \vee \bot$

by assumption

To prove $\xi_1 \not\models \xi_2$, assume $\xi_1 \models \xi_2$.

(6) $\xi_1 \models \xi_2$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \xi_{1}$ implies

(7) $e \models \xi_2$

by Definition 2.1.1 on

(1) and (2) and (6)

(8) $e \models \xi_2 \lor \bot$

by Rule (15e) on (7)

(9) $\xi_1 \models \xi_2 \lor \bot$

by Definition 2.1.1 on

(8)

(9) contradicts (5).

The conclusion holds as follows:

(a) $\xi_1 \not\models \xi_2$

Lemma 2.0.7. If $e \not\models_{?}^{\dagger} \xi_1$ and $e \not\models_{?}^{\dagger} \xi_2$ then $e \not\models_{?}^{\dagger} \xi_1 \vee \xi_2$

Proof. Assume, for the sake of contradiction, that $e \models_{?}^{\dagger} \xi_1 \vee \xi_2$.

 $(1) \ e \models^{\dagger}_{?} \xi_1 \lor \xi_2$

by assumption

(2) $e \not\models_?^\dagger \xi_1$

by assumption

(3) $e \not\models_{?}^{\dagger} \xi_2$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(4)
$$e \models \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (15) on (4) and only two of them apply.

Case (15e).

(5)
$$e \models \xi_1$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{1}$$

by Rule (18b) on (5)

(6) contradicts (2).

Case (15f).

(5)
$$e \models \xi_2$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{2}$$

by Rule (18b) on (5)

(6) contradicts (3).

Case (18a).

(4)
$$e \models_? \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (17) on (4) and only two of them apply.

Case (17c).

(5)
$$e \models_{?} \xi_{1}$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{1}$$

by Rule (18a) on (5)

(6) contradicts (2).

Case (17d).

(5)
$$e \models_{?} \xi_{2}$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{2}$$

by Rule (18a) on (5)

The conclusion holds as follows:

1.
$$e \not\models_{?}^{\dagger} \xi_1 \lor \xi_2$$

Lemma 2.0.8. If $e \models^{\dagger}_{?} \xi_1 \lor \xi_2$ and $e \not\models^{\dagger}_{?} \xi_1$ then $e \models^{\dagger}_{?} \xi_2$

Proof.

$$(1) \ e \models^{\dagger}_{?} \xi_1 \lor \xi_2$$

by assumption

(2)
$$e \not\models_?^{\dagger} \xi_1$$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(3)
$$e \models \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (15) on (3) and only two of them apply.

Case (15e).

(4) $e \models \xi_1$

by assumption

(5) $e \models_{?}^{\dagger} \xi_{1}$

by Rule (18b) on (4)

(5) contradicts (2).

Case (15f).

(4) $e \models \xi_2$

by assumption

(5) $e \models_{?}^{\dagger} \xi_2$

by Rule (18b) on (4)

Case (18a).

(3)
$$e \models_{?} \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (17) on (3) and only two of them apply.

Case (17c).

(4) $e \models_? \xi_1$

by assumption

(5) $e \models_{2}^{\dagger} \xi_{1}$

by Rule (18a) on (4)

(5) contradicts (2).

Case (17d).

(4) $e \models_{?} \xi_{2}$

by assumption

(5) $e \models_{?}^{\dagger} \xi_{2}$

by Rule (18a) on (4)

Lemma 2.0.9. If $e \models_{?}^{\dagger} \xi_1$ and $e \models_{?}^{\dagger} \xi_2$ then $e \models_{?}^{\dagger} \xi_1 \wedge \xi_2$

Lemma 2.0.10. If $e \models^{\dagger}_{?} \xi_1$ then $e \models^{\dagger}_{?} \xi_1 \vee \xi_2$ and $e \models^{\dagger}_{?} \xi_2 \vee \xi_1$

Proof.

$$(1) \ e \models^{\dagger}_{?} \xi_1$$

by assumption

By rule induction over Rules (18) on (1),

Case (18b).

(2) $e \models \xi_1$

by assumption

(3) $e \models \xi_1 \lor \xi_2$

by Rule (15e) on (2)

(4) $e \models \xi_2 \vee \xi_1$

by Rule (15f) on (2)

(5)
$$e \models_{?}^{\dagger} \xi_1 \vee \xi_2$$

by Rule (18b) on (3)

(6)
$$e \models_{?}^{\dagger} \xi_2 \vee \xi_1$$

by Rule (18b) on (4)

Case (18a).

(2)
$$e \models_? \xi_1$$

by assumption

By case analysis on the result of $satisfy(e, \xi_2)$.

Case true.

(3)
$$satisfy(e, \xi_2) = true$$

by assumption

(4)
$$e \models \xi_2$$

by Lemma 2.0.19 on

(5)
$$e \models \xi_1 \lor \xi_2$$

by Rule (15f) on (4)

(6)
$$e \models \xi_2 \lor \xi_1$$

by Rule (15e) on (4)

(7)
$$e \models_{?}^{\dagger} \xi_1 \lor \xi_2$$

by Rule (18b) on (5)

(8)
$$e \models_{2}^{\dagger} \xi_{2} \vee \xi_{1}$$

by Rule (18b) on (6)

Case false.

(3)
$$satisfy(e, \xi_2) = false$$

by assumption

(4)
$$e \not\models \xi_2$$

by Lemma 2.0.19 on

(3)

(5)
$$e \models_? \xi_1 \lor \xi_2$$

by Rule (17c) on (2)

and (4)

(6)
$$e \models_2^{\dagger} \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$$

by Rule (18a) on (5)

Lemma 2.0.11. If $e_1 \models_{?}^{\dagger} \xi_1$ then $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$

Proof.

$$(1) e_1 \models^{\dagger}_{?} \xi_1$$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2)
$$e_1 \models \xi_1$$

by assumption

(3)
$$\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$$

by Rule (15g) on (2)

$$(4) \ \operatorname{inl}_{\tau_2}(e_1) \models^\dagger_? \operatorname{inl}(\xi_1)$$

by Rule (18b) on (3)

Case (18a).

(2)
$$e_1 \models_? \xi_1$$

by assumption

(3)
$$\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$$

by Rule (17e) on (2)

(4)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{\tau}^{\dagger} \operatorname{inl}(\xi_1)$$

by Rule (18a) on (3)

Lemma 2.0.12. If $e_2 \models_?^\dagger \xi_2$ then $\operatorname{inr}_{\tau_1}(e_2) \models_?^\dagger \operatorname{inr}(\xi_2)$

Proof.

(1)
$$e_2 \models_{?}^{\dagger} \xi_2$$

by assumption

By rule induction over Rules (18) on (1).

Case (18b).

(2)
$$e_2 \models \xi_2$$

by assumption

(3)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$$

by Rule (15h) on (2)

(4)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau_2}^{\dagger} \operatorname{inr}(\xi_2)$$

by Rule (18b) on (3)

Case (18a).

(2)
$$e_2 \models_? \xi_2$$

by assumption

(3)
$$\operatorname{inl}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$$

by Rule (17f) on (2)

(4)
$$\operatorname{inl}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$

by Rule (18a) on (3)

Lemma 2.0.13. If $e_1 \models_{?}^{\dagger} \xi_1$ and $e_2 \models_{?}^{\dagger} \xi_2$ then $(e_1, e_2) \models_{?}^{\dagger} (\xi_1, \xi_2)$

Lemma 2.0.14 (Soundness and Completeness of Refutable Constraints). ξ refutable? *iff* refutable?

Lemma 2.0.15. If e notintro and ξ refutable? then either $\dot{\top}(\xi)$ refutable? or $e \models \dot{\top}(\xi)$.

Proof. By structural induction on ξ .

Lemma 2.0.16. There does not exist such a constraint $\xi_1 \wedge \xi_2$ such that $\xi_1 \wedge \xi_2$ refutable?.

Proof. By rule induction over Rules (11), we notice that $\xi_1 \wedge \xi_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.17. There does not exist such a constraint $\xi_1 \vee \xi_2$ such that $\xi_1 \vee \xi_2$ refutable?

Proof. By rule induction over Rules (11), we notice that $\xi_1 \vee \xi_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.18. If e notintro and $e \models \xi$ then ξ refutable?

Proof.

- (1) e notintro by assumption
- (2) $e \models \xi$ by assumption

By rule induction over Rules (15) on (2).

Case (15a).

(3)
$$\xi = \top$$

Assume \top refutable?. By rule induction over Rules (11), no case applies due to syntactic contradiction.

by assumption

Therefore, Trefutable?.

Case (15e),(15f).

- (3) $\xi = \xi_1 \vee \xi_2$ by assumption
- (4) $\xi_1 \vee \xi_2$ refutable? by Lemma 2.0.17

Case (15d).

- (3) $\xi = \xi_1 \wedge \xi_2$ by assumption
- (4) $\xi_1 \wedge \xi_2$ refutable? by Lemma 2.0.16

Case (15j).

- (3) $\xi = (\xi_1, \xi_2)$ by assumption
- (4) $prl(e) \models \xi_1$ by assumption
- (5) $prr(e) \models \xi_2$ by assumption
- (6) prl(e) notintro by Rule (27e)
- (7) prr(e) notintro by Rule (27f)
- (8) ξ_1 refutable? by IH on (6) and (4)
- (9) ξ_2 refutable? by IH on (7) and (5)

Assume (ξ_1, ξ_2) refutable? By rule induction over Rules (11) on it, only two cases apply.

Case (11d).

(10) ξ_1 refutable? by assumption

Contradicts (8).

Case (11e).

(10)
$$\xi_2$$
 refutable?

by assumption

Contradicts (9).

Therefore, (ξ_1, ξ_2) refutable?

Otherwise.

(3)
$$e = \underline{n}, \text{inl}_{\tau_2}(e_1), \text{inr}_{\tau_1}(e_2), (e_1, e_2)$$
 by assumption

By rule induction over Rules (27) on (1), no case applies due to syntactic contradiction.

Lemma 2.0.19 (Soundness and Completeness of Satisfaction Judgment). $e \models \xi$ iff $satisfy(e, \xi) = true$.

Proof. We prove soundness and completeness separately.

1. Soundness:

(1)
$$e \models \xi$$
 by assumption

By rule induction over Rules (15) on (1).

Case (15a).

(2)
$$\xi = \top$$
 by assumption
(3) $satisfy(e, \top) = \text{true}$ by Definition 16a

Case (15b).

$$\begin{array}{ll} (2) \ e = \underline{n} & \text{by assumption} \\ (3) \ \xi = \underline{n} & \text{by assumption} \\ (4) \ satisfy(\underline{n},\underline{n}) = (n=n) = \text{true} & \text{by Definition 16b} \end{array}$$

Case (15c).

(2)
$$e = \underline{n_1}$$
 by assumption
(3) $\xi = \underline{p_2}$ by assumption
(4) $n_1 \neq n_2$ by assumption
(5) $satisfy(n_1, p_2) = (n_1 \neq n_2) = true$ by Definition 16c on (4)

Case (15d).

(2)
$$\xi = \xi_1 \wedge \xi_2$$
 by assumption

 $\begin{array}{lll} (3) & e \models \xi_1 & \text{by assumption} \\ (4) & e \models \xi_2 & \text{by assumption} \\ (5) & \textit{satisfy}(e,\xi_1) = \text{true} & \text{by IH on (3)} \\ (6) & \textit{satisfy}(e,\xi_2) = \text{true} & \text{by IH on (4)} \\ (7) & \textit{satisfy}(e,\xi_1 \land \xi_2) = \textit{satisfy}(e,\xi_1) \text{ and } \textit{satisfy}(e,\xi_2) = \text{true} \\ & \text{by Definition 16d on} \\ & (5) \text{ and (6)} \\ \end{array}$

Case (15e).

- (2) $\xi = \xi_1 \vee \xi_2$ by assumption
- (3) $e \models \xi_1$ by assumption
- (4) $satisfy(e, \xi_1) = true$ by IH on (3)
- (5) $satisfy(e, \xi_1 \lor \xi_2) = satisfy(e, \xi_1)$ or $satisfy(e, \xi_2) = true$ by Definition 16e on (4)

Case (15f).

- (2) $\xi = \xi_1 \vee \xi_2$ by assumption
- (3) $e \models \xi_2$ by assumption
- (4) $satisfy(e, \xi_2) = true$ by IH on (3)
- (5) $satisfy(e, \xi_1 \vee \xi_2) = satisfy(e, \xi_1)$ or $satisfy(e, \xi_2) = true$ by Definition 16e on (4)

Case (15g).

- (2) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (3) $\xi = inl(\xi_1)$ by assumption
- (4) $e_1 \models \xi_1$ by assumption
- (5) $satisfy(e_1, \xi_1) = true$ by IH on (4)
- (6) $satisfy(inl_{\tau_2}(e_1), inl(\xi_1)) = satisfy(e_1, \xi_1) = true$ by Definition 16f on (5)

Case (15h).

- (2) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption
- (3) $\xi = inl(\xi_2)$ by assumption
- (4) $e_2 \models \xi_2$ by assumption
- (5) $satisfy(e_2, \xi_2) = true$ by IH on (4)
- $(6) \ \ \mathit{satisfy}(\mathtt{inr}_{\tau_1}(e_2),\mathtt{inr}(\xi_2)) = \mathit{satisfy}(e_2,\xi_2) = \mathsf{true}$

by Definition 16g on (5)

Case (15i).

$$(2) \ e = (e_1, e_2) \qquad \qquad \text{by assumption}$$

$$(3) \ \xi = (\xi_1, \xi_2) \qquad \qquad \text{by assumption}$$

$$(4) \ e_1 \models \xi_1 \qquad \qquad \text{by assumption}$$

$$(5) \ e_2 \models \xi_2 \qquad \qquad \text{by assumption}$$

$$(6) \ satisfy(e_1, \xi_1) = \text{true} \qquad \qquad \text{by IH on (4)}$$

$$(7) \ satisfy(e_2, \xi_2) = \text{true} \qquad \qquad \text{by IH on (5)}$$

(8) $satisfy((e_1,e_2),(\xi_1,\xi_2))=satisfy(e_1,\xi_1)$ and $satisfy(e_2,\xi_2)=$ true by Definition 16h on (6) and (7)

Case (15j).

(2)
$$\xi = (\xi_1, \xi_2)$$
 by assumption
(3) e notintro by assumption
(4) $prl(e) \models \xi_1$ by assumption
(5) $prr(e) \models \xi_2$ by assumption
(6) $satisfy(prl(e), \xi_1) = true$ by IH on (4)
(7) $satisfy(prr(e), \xi_2) = true$ by IH on (5)

By rule induction over Rules (27) on (3).

Otherwise.

(8)
$$e = \{ \| u, \| e_0 \| u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \} \}$$
 by assumption

(9)
$$satisfy(e, (\xi_1, \xi_2)) = satisfy(\texttt{prl}(e), \xi_1)$$
 and $satisfy(\texttt{prr}(e), \xi_2) = \text{true}$ by Definition 16 on (6) and (7)

2. Completeness:

(1)
$$satisfy(e, \xi) = true$$
 by assumption

By structural induction on ξ .

Case
$$\xi = \top$$
.

(2)
$$e \models \top$$
 by Rule (15a)

Case $\xi = \bot$,?.

(2)
$$satisfy(e, \xi) = false$$
 by Definition 160

(2) contradicts (1) and thus vacuously true.

Case
$$\xi = \underline{n}$$
.

By structural induction on e.

Case $e = \underline{n'}$.

- (2) n' = n by Definition 16b on (1)
- (3) $\underline{n'} \models \underline{n}$ by Rule (15b) on (2)

Otherwise.

- (2) $satisfy(e, \underline{n}) = false$ by Definition 160
- (2) contradicts (1) and thus vacuously true.

Case $\xi = \underline{\varkappa}$.

By structural induction on e.

Case $e = \underline{n'}$.

- (2) $n' \neq n$ by Definition 16c on (1)
- (3) $\underline{n'} \models \underline{\varkappa}$ by Rule (15c) on (2)

Otherwise.

- (2) $satisfy(e, \varkappa) = false$ by Definition 160
- (2) contradicts (1) and thus vacuously true.

Case $\xi = \xi_1 \wedge \xi_2$.

- (2) $satisfy(e, \xi_1) = true$ by Definition 16d on
 - (1)
- (3) $satisfy(e, \xi_2) = true$ by Definition 16d on (1)
- (4) $e \models \xi_1$ by IH on (2)
- (5) $e \models \xi_2$ by IH on (3)
- (6) $e \models \xi_1 \land \xi_2$ by Rule (15d) on (4) and (5)

Case $\xi = \xi_1 \vee \xi_2$.

(2) $\operatorname{satisfy}(e, \xi_1)$ or $\operatorname{satisfy}(e, \xi_2) = \operatorname{true}$

by Definition 16e on (1)

By case analysis on (2).

Case $satisfy(e, \xi_1) = true.$

- (3) $satisfy(e, \xi_1) = true$ by assumption (4) $e \models \xi_1$ by IH on (3)
- (5) $e \models \xi_1 \lor \xi_2$ by Rule (15e) on (4)

Case $satisfy(e, \xi_2) = true.$

- (3) $satisfy(e, \xi_2) = true$ by assumption (4) $e \models \xi_2$ by IH on (3)
- (5) $e \models \xi_1 \vee \xi_2$ by Rule (15f) on (4)

Case $\xi = inl(\xi_1)$.

By structural induction on e.

Case $e = \operatorname{inl}_{\tau_2}(e_1)$.

- (2) $satisfy(e_1, \xi_1) = true$ by Definition 16f on (1)
- (3) $e_1 \models \xi_1$ by IH on (2)
- (4) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (15g) on (3)

Otherwise.

- (2) $satisfy(e, inl(\xi_1)) = false$ by Definition 160
- (2) contradicts (1) and thus vacuously true.

Case $\xi = inr(\xi_2)$.

By structural induction on e.

Case $e = \operatorname{inr}_{\tau_1}(e_2)$.

- (2) $satisfy(e_2, \xi_2) = true$ by Definition 16g on (1)
- (3) $e_2 \models \xi_2$ by IH on (2)
- (4) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ by Rule (15h) on (3)

Otherwise.

- (2) $satisfy(e, inr(\xi_2)) = false$ by Definition 160
- (2) contradicts (1) and thus vacuously true.

Case $\xi = (\xi_1, \xi_2)$.

By structural induction on e.

Case $e = (e_1, e_2)$.

- (2) $satisfy(e_1, \xi_1) = true$ by Definition 16h on (1)
- (3) $satisfy(e_2, \xi_2) = true$ by Definition 16h on (1)
- (4) $e_1 \models \xi_1$ by IH on (2)
- (5) $e_2 \models \xi_2$ by IH on (3)
- (6) $(e_1, e_2) \models (\xi_1, \xi_2)$ by Rule (15i) on (4) and (5)

Case $e = (||u|, ||e_0||u|, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}.$

- (2) $satisfy(prl(e), \xi_1) = true$ by Definition 16h on (1)
- (3) $satisfy(prr(e), \xi_2) = true$ by Definition 16h on (1)
- (4) $prl(e) \models \xi_1$ by IH on (2)
- (5) $prr(e) \models \xi_2$ by IH on (3)
- (6) e notintro by each rule in Rules (27)

(7)
$$(e_1, e_2) \models (\xi_1, \xi_2)$$

by Rule (15j) on (6) and (4) and (5)

Otherwise.

(2) $satisfy(e, (\xi_1, \xi_2)) = false$

by Definition 16o

(2) contradicts (1) and thus vacuously true.

Lemma 2.0.20. $e \not\models \xi$ and $e \not\models_? \xi$ iff $e \not\models_?^{\dagger} \xi$.

Proof. 1. Sufficiency:

(1) $e \not\models \xi$

by assumption

(2) $e \not\models_? \xi$

by assumption

Assume $e \models_{?}^{\dagger} \xi$. By rule induction over Rules (18) on it.

Case (18a).

(3) $e \models \xi$

by assumption

Contradicts (1).

Case (18b).

(3) $e \models_? \xi$

by assumption

Contradicts (2).

Therefore, $e \models_{?}^{\dagger} \xi$ is not derivable.

2. Necessity:

(1) $e \not\models_?^\dagger \xi$

by assumption

Assume $e \models \xi$.

(2) $e \models^{\dagger}_{?} \xi$

by Rule (18b) on assumption

Contradicts (1). Therefore, $e \not\models \xi$. Assume $e \models_? \xi$.

(3) $e \models^{\dagger}_{?} \xi$

by Rule (18a) on assumption

Contradicts (1). Therefore, $e \not\models_? \xi$.

Theorem 2.1 (Exclusiveness of Satisfaction Judgment). If $\xi : \tau$ and $\cdot ; \Delta \vdash e : \tau$ and e final then exactly one of the following holds

1. $e \models \xi$

$$2. e \models_? \xi$$

3.
$$e \not\models_?^\dagger \xi$$

Proof.

- (4) $\xi : \tau$ by assumption
- (5) \cdot ; $\Delta \vdash e : \tau$ by assumption
- (6) e final by assumption

By rule induction over Rules (9) on (4), we would show one conclusion is derivable while the other two are not.

Case (9a).

- (7) $\xi = \top$ by assumption
- (8) $e \models \top$ by Rule (15a)
- (9) $e \not\models_? \top$ by Lemma 2.0.3
- (10) $e \models^{\dagger}_{?} \top$ by Rule (18b) on (8)

Case (9b).

- (7) $\xi = \bot$ by assumption
- (8) $e \not\models \bot$ by Lemma 2.0.1 (9) $e \not\models_? \bot$ by Lemma 2.0.2
- (10) $e \not\models_?^{\dagger} \bot$ by Lemma 2.0.20 on

(8) and (9)

Case (1b).

- (7) $\xi = ?$ by assumption
- (8) $e \not\models$? by Lemma 2.0.4
- (9) $e \models_?$? by Rule (17a)

(10) $e \models^{\dagger}_{?}$? by Rule (18a) on (9)

Case (9c).

- (7) $\xi = \underline{n_2}$ by assumption
- (8) $\tau = \text{num}$ by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20l),(20m).

(9)
$$e = (v_0)^u, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$

by assumption

(10)
$$e$$
 notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

Assume $e \models n_2$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction on ξ .

(11)
$$e \not\models \underline{n_2}$$
 by contradiction (12) n_2 refutables by Rule (11a)

(12)
$$\underline{n_2}$$
 refutable? by Rule (11a)

(13)
$$e \models_{?} \underline{n_2}$$
 by Rule (17b) on (10) and (12)

(14)
$$e \models_{7}^{\dagger} n_2$$
 by Rule (18a) on (13)

Case (20d).

(9)
$$e = n_1$$
 by assumption

Assume $\underline{n_1} \models_? \underline{n_2}$. By rule induction over Rules (17), only one case applies.

Case (17b).

(10)
$$n_1$$
 notintro by assumption

Contradicts Lemma 4.0.6.

(11)
$$\underline{n_1} \not\models_? \underline{n_2}$$
 by contradiction

By case analysis on whether n_1 is equal to n_2 .

Case $n_1 = n_2$.

$$\begin{array}{ll} (12) \;\; \mathit{satisfy}(\underline{n_1},\underline{n_2}) = \mathsf{true} & \qquad \mathsf{by Definition 16} \\ (13) \;\; \underline{n_1} \models \underline{n_2} & \qquad \mathsf{by Lemma 2.0.19 on} \\ (12) \end{array}$$

(14)
$$e \models_{?}^{\dagger} \underline{n_2}$$
 by Rule (18b) on (13)

Case $n_1 \neq n_2$.

(12)
$$satisfy(\underline{n_1},\underline{n_2}) = false$$
 by Definition 16
(13) $\underline{n_1} \not\models \underline{n_2}$ by Lemma 2.0.19 on
(12)

(14)
$$e \not\models_{?}^{\dagger} \underline{n_2}$$
 by Lemma 2.0.20 on (11) and (13)

Case (9f).

(7)
$$\xi = \xi_1 \vee \xi_2$$
 by assumption

By inductive hypothesis on (5) and (6), exactly one of $e \models \xi_1$, $e \models_? \xi_1$, and $e \not\models_?^{\dagger} \xi_1$ holds. The same goes for ξ_2 . By case analysis on which conclusion holds for ξ_1 and ξ_2 .

Case $e \models \xi_1, e \models \xi_2$.

(8)	$e \models \xi_1$	by assumption
(9)	$e \not\models_? \xi_1$	by assumption
(10)	$e \models \xi_2$	by assumption
(11)	$e \not\models_? \xi_2$	by assumption
(12)	$e \models \xi_1 \lor \xi_2$	by Rule $(15e)$ on (8)
(13)	$e \models_{2}^{\dagger} \xi_{1} \lor \xi_{2}$	by Rule (18b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (17c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9).

Case (17d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \models \xi_1, e \models_? \xi_2$.

(8) $e \models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \models_? \xi_2$ by assumption (12) $e \models \xi_1 \lor \xi_2$ by Rule (15e) on (8) (13) $e \models_?^1 \xi_1 \lor \xi_2$ by Rule (18b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (17c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9). Case (17d). (14) $e \not\models \xi_1$ by assumption Contradicts (8). (15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction Case $e \models \xi_1, e \not\models_{?}^{\dagger} \xi_2$. (8) $e \models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption (12) $e \models \xi_1 \lor \xi_2$ by Rule (15e) on (8) (13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (18b) on (12) Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (17) on it, the following cases apply. Case (17b). (14) $\xi_1 \vee \xi_2$ refutable? by assumption Contradicts Lemma 2.0.17. Case (17c). (14) $e \models_? \xi_1$ by assumption Contradicts (9). Case (17d). (14) $e \not\models \xi_1$ by assumption Contradicts (8). (15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction Case $e \models_? \xi_1, e \models \xi_2$. (8) $e \not\models \xi_1$ by assumption (9) $e \models_{?} \xi_{1}$ by assumption (10) $e \models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption (12) $e \models \xi_1 \lor \xi_2$ by Rule (15f) on (10)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (17) on it, the following cases apply.

by Rule (18b) on (12)

Case (17b).

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (17c).

(14) $e \not\models \xi_2$ by assumption

Contradicts (10).

Case (17d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \models_? \xi_1, e \models_? \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \models_{?} \xi_1$ by assumption

(10) $e \not\models \xi_2$ by assumption (11) $e \models_? \xi_2$ by assumption

(12) $e \models_{?} \xi_1 \lor \xi_2$ by Rule (17c) on (9)

and (10)

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (18a) on (12)

Assume $e \models \xi_1 \lor \xi_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14) $e \models \xi_1$ by assumption

Contradicts (8)

Case (15f).

(14) $e \models \xi_2$ by assumption

Contradicts (10)

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \models_? \xi_1, e \not\models_?^{\dagger} \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \models_? \xi_1$ by assumption

(10) $e \not\models \xi_2$ by assumption

(11) $e \not\models_? \xi_2$ by assumption

(12) $e \models_? \xi_1 \lor \xi_2$ by Rule (17c) on (9)

and (10)

(13) $e \models_{2}^{\dagger} \xi_{1} \lor \xi_{2}$ by Rule (18a) on (12)

Assume $e \models \xi_1 \vee \xi_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14) $e \models \xi_1$ by assumption

Contradicts (8).

Case (15f).

(14) $e \models \xi_2$ by assumption

Contradicts (10).

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \models \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \not\models_? \xi_1$ by assumption

(10) $e \models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption

(12) $e \models \xi_1 \lor \xi_2$ by Rule (15f) on (10)

(13) $e \models_{2}^{\dagger} \xi_{1} \lor \xi_{2}$ by Rule (18b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (17c).

(14) $e \not\models \xi_2$ by assumption

Contradicts (10).

Case (17d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \models_? \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption

(11) $e \models_? \xi_2$ by assumption

(12) $e \models_? \xi_1 \lor \xi_2$ by Rule (17d) on (11)

and (8)

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (18a) on (12)

Assume $e \models \xi_1 \vee \xi_2$. By rule induction over Rules (15), only two cases apply.

Case (15e).

(14) $e \models \xi_1$ by assumption

Contradicts (8)

Case (15f).

(14) $e \models \xi_2$ by assumption

Contradicts (10)

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \not\models_?^\dagger \xi_2$.

(8) $e \not\models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption

Assume $e \models \xi_1 \lor \xi_2$. By rule induction over Rules (15) on it, only two cases apply.

Case (15e).

(12) $e \models \xi_1$ by assumption

Contradicts (8).

Case (15f).

(12) $e \models \xi_2$ by assumption

Contradicts (10).

(13) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Assume $e \models_{?} \xi_1 \vee \xi_2$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (17c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9).

Case (17d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

(16)
$$e \not\models_{?}^{\dagger} \xi_1 \vee \xi_2$$
 by Lemma 2.0.20 on (13) and (15)

Case (9g).

(7)
$$\xi = \text{inl}(\xi_1)$$
 by assumption
(8) $\tau = (\tau_1 + \tau_2)$ by assumption
(9) $\xi_1 : \tau_1$ by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

(10)
$$e = \langle || u, || e_0 || u, e_1 (e_2), prl(e_0), prr(e_0), match(e_0) \{ \hat{rs} \}$$
 by assumption

(11) e notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

Assume $e \models \text{inl}(\xi_1)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inl}(\xi_1)$$
 by contradiction

By case analysis on the value of $refutable_{?}(inl(\xi_1))$.

Case $refutable_{?}(inl(\xi_1)) = true.$

- (13) $refutable_{?}(inl(\xi_{1})) = true$ by assumption (14) $inl(\xi_{1})$ refutable_? by Lemma 2.0.14 on (13) (15) $e \models_{?} inl(\xi_{1})$ by Rule (17b) on (11) and (14)
- (16) $e \models^{\dagger}_{?} \operatorname{inl}(\xi_1)$ by Rule (18a) on (15)

Case $refutable_{?}(inl(\xi_1)) = false.$

(13) $refutable_?(inl(\xi_1)) = false$ by assumption (14) $inl(\xi_1)$ $refutable_?$ by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(15)
$$\operatorname{inl}(\xi_1)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inl}(\xi_1)$$
 by contradiction
(17) $e \not\models_?^{\dagger} \operatorname{inl}(\xi_1)$ by Lemma 2.0.20 on
(12) and (16)

Case (20j).

- (10) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (11) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (12) e_1 final by Lemma 4.0.3 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_1 \models \xi_1$, $e_1 \models_? \xi_1$, and $e_1 \not\models_?^{\dagger} \xi_1$ holds. By case analysis on which one holds.

Case $e_1 \models \xi_1$.

- (13) $e_1 \models \xi_1$ by assumption
- (14) $e_1 \not\models_? \xi_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (15g) on (13)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Rule (18b) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (17) on it, only two cases apply.

Case (17b).

- (17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption
- By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17e).

(17) $e_1 \models_? \xi_1$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\xi_1)$$
 by contradiction

Case $e_1 \models_? \xi_1$.

- (13) $e_1 \not\models \xi_1$ by assumption
- (14) $e_1 \models_? \xi_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by Rule (17e) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Rule (18a) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (15) on it, only one case applies.

Case (15g).

(17) $e_1 \models \xi_1$

Contradicts (13).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\xi_1)$$
 by contradiction

Case $e_1 \not\models_{?}^{\dagger} \xi_1$.

(13)
$$e_1 \not\models \xi_1$$
 by assumption

(14)
$$e_1 \not\models_? \xi_1$$

by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (15) on it, only one case applies.

Case (15g).

(15)
$$e_1 \models \xi_1$$

Contradicts (13).

(16)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\xi_1)$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro

by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17e).

(17)
$$e_1 \models_? \xi_1$$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\xi_1)$$
 by contradiction

(19)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inl}(\xi_1)$$
 by Lemma 2.0.20 on (16) and (18)

Case (20k).

(10)
$$e = inr_{\tau_1}(e_2)$$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inl}(\xi_1)$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(12)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (27) on (12), no case applies due to syntactic contradiction.

(13)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inl}(\xi_1)$$

by contradiction

$$(14) \ \operatorname{inr}_{\tau_1}(e_2) \not\models_?^\dagger \operatorname{inl}(\xi_1)$$

by Lemma 2.0.20 on

(11) and (13)

Case (9h).

(7)
$$\xi = \operatorname{inr}(\xi_2)$$

by assumption

(8)
$$\tau = (\tau_1 + \tau_2)$$

by assumption

(9) $\xi_2 : \tau_2$

by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20m).

$$(10) \ e = ()^u, (e_0)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{ \hat{rs} \}$$

by assumption

(11)
$$e$$
 notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

Assume $e \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inr}(\xi_2)$$

by contradiction

By case analysis on the value of $refutable_{?}(inr(\xi_{2}))$.

inr is refutable

Case $refutable_2(inr(\xi_2)) = true.$

(13)
$$refutable_2(inr(\xi_2)) = true$$
 by assumption

(14)
$$inr(\xi_2)$$
 refutable? by Lemma 2.0.14 on (13)

(15)
$$e \models_? \operatorname{inr}(\xi_2)$$
 by Rule (17b) on (11) and (14)

(16)
$$e \models_{7}^{\dagger} inr(\xi_2)$$
 by Rule (18a) on (15)

Case $refutable_{?}(inr(\xi_2)) = false.$

(13)
$$refutable_{?}(inr(\xi_2)) = false$$
 by assumption

(14)
$$\operatorname{inr}(\xi_2)$$
 refutable? by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(15)
$$inr(\xi_2)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction
(17) $e \not\models_?^{\dagger} \operatorname{inr}(\xi_2)$ by Lemma 2.0.20 on
(12) and (16)

Case (20j).

(10)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (15) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inr}(\xi_2)$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(12) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

By rule induction over Rules (27) on (12), no case applies due to syntactic contradiction.

(13)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction

(14)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Lemma 2.0.20 on (11) and (13)

Case (20k).

(10)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption

(11)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$ by assumption

(12)
$$e_2$$
 final by Lemma 4.0.4 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_2 \models \xi_2$, $e_2 \models_? \xi_2$, and $e_2 \not\models^{\dagger}_? \xi_2$ holds. By case analysis on which one holds.

Case $e_2 \models \xi_2$.

(13)
$$e_2 \models \xi_2$$
 by assumption

(14)
$$e_2 \not\models_? \xi_2$$
 by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$$
 by Rule (15g) on (13)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Rule (18b) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (17) on it, only two cases apply.

Case (17b).

(17)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17f).

(17)
$$e_2 \models_? \xi_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction

Case $e_2 \models_? \xi_2$.

(13)
$$e_2 \not\models \xi_2$$
 by assumption

(14)
$$e_2 \models_? \xi_2$$
 by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{?} \operatorname{inr}(\xi_2)$$
 by Rule (17f) on (14)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Rule (18a) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15h).

$$(17) \ e_2 \models \xi_2$$

Contradicts (13).

(18) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\xi_2)$ by contradiction

Case $e_2 \not\models_?^\dagger \xi_2$.

 $(13) e_2 \not\models \xi_2$

by assumption

(14) $e_2 \not\models_? \xi_2$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15h).

(15)
$$e_2 \models \xi_2$$

Contradicts (13).

(16) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\xi_2)$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(17) $\operatorname{inr}_{\tau_1}(e_2)$ notintro

by assumption

By rule induction over Rules (27) on (17), no case applies due to syntactic contradiction.

Case (17f).

(17)
$$e_2 \models_? \xi_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction

(19)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inl}(\xi_1)$$
 by Lemma 2.0.20 on (16) and (18)

Case (15i).

(7)
$$\xi = (\xi_1, \xi_2)$$

by assumption

(8) $\tau = (\tau_1 \times \tau_2)$

by assumption

(9) $\xi_1 : \tau_1$

by assumption

(10) $\xi_2 : \tau_2$

by assumption

By rule induction over Rules (20) on (5), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

```
(11) e = \{ \|u, \|e_0\|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \} \}
                                                             by assumption
    (12) e notintro
                                                             by Rule
                                                             (27a),(27b),(27c),(27d),(27e),(27f)
    (13) e indet
                                                             by Lemma 4.0.10 on
                                                             (6) and (12)
   (14) prl(e) indet
                                                             by Rule (25g) on (13)
    (15) prl(e) final
                                                             by Rule (26b) on (14)
    (16) prr(e) indet
                                                             by Rule (25h) on (13)
    (17) prr(e) final
                                                             by Rule (26b) on (16)
    (18) \cdot; \Delta \vdash \mathtt{prl}(e) : \tau_1
                                                             by Rule (20h) on (5)
                                                            by Rule (20i) on (5)
    (19) \cdot; \Delta \vdash \mathsf{prr}(e) : \tau_2
 By inductive hypothesis on (9) and (18) and (15), exactly one of
\operatorname{prl}(e) \models \xi_1, \operatorname{prl}(e) \models_? \xi_1, \text{ and } \operatorname{prl}(e) \not\models_?^{\dagger} \xi_1 \text{ holds.}
By inductive hypothesis on (10) and (19) and (17), exactly one of
\operatorname{prr}(e) \models \xi_2, \operatorname{prr}(e) \models_? \xi_2, \text{ and } \operatorname{prr}(e) \not\models_?^{\dagger} \xi_2 \text{ holds.}
By case analysis on which conclusion holds for \xi_1 and \xi_2.
 Case prl(e) \models \xi_1, prr(e) \models \xi_2.
         (20) prl(e) \models \xi_1
                                                             by assumption
         (21) prl(e) \not\models_? \xi_1
                                                             by assumption
          (22) prr(e) \models \xi_2
                                                             by assumption
         (23) prr(e) \not\models_? \xi_2
                                                             by assumption
         (24) e \models (\xi_1, \xi_2)
                                                             by Rule (15j) on (12)
                                                             and (20) and (22)
         (25) e \models^{\dagger}_{?} (\xi_1, \xi_2)
                                                             by Rule (18b) on (24)
         (26) (\xi_1, \xi_2) refutable?
                                                            by Lemma 2.0.18 on
                                                            (12) and (24)
       Assume e \models_{?} (\xi_1, \xi_2). By rule induction over Rules (17) on it,
      only one case applies.
       Case (17b).
               (27) (\xi_1, \xi_2) refutable?
                                                             by assumption
            Contradicts (26).
         (28) e \not\models_{?} (\xi_1, \xi_2)
                                                             by contradiction
 Case prl(e) \models \xi_1, prr(e) \models_? \xi_2.
         (20) prl(e) \models \xi_1
                                                             by assumption
          (21) prl(e) \not\models_? \xi_1
                                                             by assumption
          (22) prr(e) \not\models \xi_2
                                                             by assumption
          (23) prr(e) \models_? \xi_2
                                                            by assumption
```

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

(24)
$$prr(e) \models \xi_2$$

by assumption

Contradicts (22)

(25)
$$e \not\models (\xi_1, \xi_2)$$

by contradiction

By rule induction over Rules (17) on (23), only one case applies.

Case (17b).

(26) ξ_2 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (11e) on (26)

assume no or and

"and" in

pair

(28) $e \models_{?} (\xi_1, \xi_2)$

by Rule (17b) on (12) and (27)

(29) $e \models_{?}^{\dagger} (\xi_1, \xi_2)$

by Rule (18a) on (28)

Case $prl(e) \models \xi_1, prr(e) \not\models_{?}^{\dagger} \xi_2.$

(20) $prl(e) \models \xi_1$

by assumption

(21) $prl(e) \not\models_? \xi_1$

by assumption

(22) $prr(e) \not\models \xi_2$

by assumption

(23) $prr(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

(24)
$$prr(e) \models \xi_2$$

by assumption

Contradicts (22).

(25)
$$e \not\models (\xi_1, \xi_2)$$

by contradiction

Assume $e \models_{?} (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(26) (ξ_1, ξ_2) refutable?

by assumption

By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

(27) ξ_1 refutable?

by assumption

(28) prl(e) notintro

by Rule (27e)

(29) $prl(e) \models_? \xi_1$

by Rule (17b) on (28)

and (27)

Contradicts (21).

Case (11e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $prr(e) \models_{?} \xi_{2}$ by Rule (17b) on (28) and (27)

Contradicts (23).

(30) $e \not\models_{?} (\xi_{1}, \xi_{2})$ by contradiction (31) $e \not\models_{?}^{\dagger} (\xi_{1}, \xi_{2})$ by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \models_? \xi_1, prr(e) \models \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$ by assumption (21) $\operatorname{prl}(e) \models_? \xi_1$ by assumption (22) $\operatorname{prr}(e) \models \xi_2$ by assumption (23) $\operatorname{prr}(e) \not\models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

- (24) $prl(e) \models \xi_1$ by assumption Contradicts (20).
- (25) $e \not\models (\xi_1, \xi_2)$ by contradiction By rule induction over Rules (17) on (21), only one case applies.

Case (17b).

- (26) ξ_1 refutable? by assumption
- (27) (ξ_1, ξ_2) refutable? by Rule (11e) on (26)

assume no "or" and

"and" in

pair

- (28) $e \models_{?} (\xi_1, \xi_2)$ by Rule (17b) on (12) and (27)
- (29) $e \models_{?}^{\dagger} (\xi_1, \xi_2)$ by Rule (18a) on (28)

Case $prl(e) \models_? \xi_1, prr(e) \models_? \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$ by assumption (21) $\operatorname{prl}(e) \models_? \xi_1$ by assumption (22) $\operatorname{prr}(e) \not\models \xi_2$ by assumption (23) $\operatorname{prr}(e) \models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

(24) $prl(e) \models \xi_1$ by assumption Contradicts (20).

(25) $e \not\models (\xi_1, \xi_2)$

by contradiction

By rule induction over Rules (17) on (23), only one case applies.

Case (17b).

(26) ξ_2 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (11e) on (26)

assume no "or" and

"and" in

assume no "or" and

"and" in

pair

pair

(28) $e \models_{?} (\xi_1, \xi_2)$

by Rule (17b) on (12)

and (27)

(29) $e \models_{?}^{\dagger} (\xi_1, \xi_2)$

by Rule (18a) on (28)

Case $prl(e) \models_? \xi_1, prr(e) \not\models_?^{\dagger} \xi_2.$

(20) $prl(e) \not\models \xi_1$

by assumption

(21) $prl(e) \models_? \xi_1$

by assumption

(22) $\operatorname{prr}(e) \not\models \xi_2$

by assumption

(23) $\operatorname{prr}(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

(24) $prl(e) \models \xi_1$

by assumption

Contradicts (20)

(25) $e \not\models (\xi_1, \xi_2)$

by contradiction

By rule induction over Rules (17) on (21), only one case applies.

Case (17b).

(26) ξ_1 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (11e) on (26)

(28) $e \models_{?} (\xi_1, \xi_2)$

by Rule (17b) on (12)

and (27)

(29) $e \models^{\dagger}_{?} (\xi_1, \xi_2)$

by Rule (18a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \xi_1, prr(e) \models \xi_2.$

(20) $prl(e) \not\models \xi_1$

by assumption

(21) $prl(e) \not\models_? \xi_1$

by assumption

(22) $prr(e) \models \xi_2$

by assumption

(23) $prr(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

(24) $prl(e) \models \xi_1$

Contradicts (20)

by assumption

89

(25)
$$e \not\models (\xi_1, \xi_2)$$
 by contradiction
Assume $e \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it,

only one case applies.

Case (17b).

(26) (ξ_1, ξ_2) refutable? by assumption

By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

- (27) ξ_1 refutable? by assumption (28) prl(e) notintro by Rule (27e)
- (29) $prl(e) \models_{?} \xi_{1}$ by Rule (17b) on (28) and (27)

Contradicts (21).

Case (11e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $\operatorname{prr}(e) \models_{?} \xi_{2}$ by Rule (17b) on (28) and (27)

Contradicts (23).

(30) $e \not\models_? (\xi_1, \xi_2)$ by contradiction (31) $e \not\models_?^{\dagger} (\xi_1, \xi_2)$ by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \not\models_{?}^{\dagger} \xi_1, prr(e) \models_{?} \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \xi_1$ by assumption (22) $\operatorname{prr}(e) \not\models \xi_2$ by assumption (23) $\operatorname{prr}(e) \models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15), only one case applies.

Case (15j).

- (24) $prl(e) \models \xi_1$ by assumption Contradicts (20).
- (25) $e \not\models (\xi_1, \xi_2)$ by contradiction

By rule induction over Rules (17) on (23), only one case applies.

Case (17b).

- (26) ξ_2 refutable?
- by assumption
- (27) (ξ_1, ξ_2) refutable? by Rule (11e) on (26)

assume no
"or" and
"and" in
pair

(28)
$$e \models_{?} (\xi_{1}, \xi_{2})$$
 by Rule (17b) on (12) and (27)

(29) $e \models_{?}^{\dagger} (\xi_{1}, \xi_{2})$ by Rule (18a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \xi_{1}, prr(e) \not\models_{?}^{\dagger} \xi_{2}$.

(20) $prl(e) \not\models \xi_{1}$ by assumption

(20) $prl(e) \not\models \xi_1$

(21) $prl(e) \not\models_? \xi_1$ by assumption (22) $prr(e) \not\models \xi_2$ by assumption (23) $prr(e) \not\models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only one case applies.

Case (15j).

(24) $prl(e) \models \xi_1$ by assumption Contradicts (20)

(25) $e \not\models (\xi_1, \xi_2)$ by contradiction

Assume $e \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, only one case applies.

Case (17b).

(26) (ξ_1, ξ_2) refutable? by assumption

By rule induction over Rules (11) on (26), only two cases apply.

Case (11d).

- (27) ξ_1 refutable? by assumption (28) prl(e) notintro by Rule (27e)
- (29) $prl(e) \models_? \xi_1$ by Rule (17b) on (28) and (27)

Contradicts (21).

Case (11e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (27f)
- (29) $prr(e) \models_? \xi_2$ by Rule (17b) on (28) and (27)

Contradicts (23).

(30)
$$e \not\models_{?} (\xi_{1}, \xi_{2})$$
 by contradiction
(31) $e \not\models_{?}^{\dagger} (\xi_{1}, \xi_{2})$ by Lemma 2.0.20 on
(25) and (30)

Case (20g).

(11)
$$e = (e_1, e_2)$$
 by assumption
(12) $\cdot : \Delta \vdash e_1 : \tau_1$ by assumption

(13) \cdot ; $\Delta \vdash e_2 : \tau_2$ by assumption (14) e_1 final by Lemma 4.0.5 on (6)

(15) e_2 final by Lemma 4.0.5 on (6)

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 \models \xi_1$, $e_1 \models_? \xi_1$, and $e_1 \models_{\overline{\xi_1}}$ holds.

By inductive hypothesis on (10) and (13) and (15), exactly one of $e_2 \models \xi_2$, $e_2 \models_? \xi_2$, and $e_2 \models \overline{\xi_2}$ holds.

By case analysis on which conclusion holds for ξ_1 and ξ_2 .

Case $e_1 \models \xi_1, e_2 \models \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \models \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

(20) $(e_1, e_2) \models (\xi_1, \xi_2)$ by Rule (15i) on (16)

and (18) $(21) (e_1, e_2) \models_{?}^{\dagger} (\xi_1, \xi_2)$ by Rule (18b) on (20)

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Case $e_1 \models \xi_1, e_2 \models_? \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \not\models \xi_2$ by assumption

(19) $e_2 \models_? \xi_2$ by assumption

(20) $(e_1, e_2) \models_? (\xi_1, \xi_2)$ by Rule (17h) on (16)

and (19)

(21)
$$(e_1, e_2) \models_{?}^{\uparrow} (\xi_1, \xi_2)$$
 by Rule (18a) on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(22) $e_2 \models \xi_2$ by assumption

Contradicts (18).

(23) $(e_1, e_2) \not\models (\xi_1, \xi_2)$ by contradiction

Case $e_1 \models \xi_1, e_2 \not\models_?^{\dagger} \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption (18) $e_2 \not\models \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(20) $e_2 \models \xi_2$ by assumption

Contradicts (18).

(21) $(e_1, e_2) \not\models (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

emma 2.0.20 on and (23)

Case $e_1 \models_? \xi_1, e_2 \models \xi_2$.

$$\begin{array}{lll} (16) & e_1 \not\models \xi_1 & & \text{by assumption} \\ (17) & e_1 \models_? \xi_1 & & \text{by assumption} \\ (18) & e_2 \models \xi_2 & & \text{by assumption} \\ (19) & e_2 \not\models_? \xi_2 & & \text{by assumption} \\ \end{array}$$

(20)
$$(e_1, e_2) \models_? (\xi_1, \xi_2)$$
 by Rule (17g) on (17) and (18)

(21)
$$(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$$
 by Rule (18a) on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.9.

Case (15i).

(22)
$$e_1 \models \xi_1$$
 by assumption Contradicts (16).

(23)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction

Case $e_1 \models_? \xi_1, e_2 \models_? \xi_2$.

$(16) e_1 \not\models \xi_1$	by assumption
(17) $e_1 \models_? \xi_1$	by assumption
$(18) \ e_2 \not\models \xi_2$	by assumption
(19) $e_2 \models_? \xi_2$	by assumption
$(20) (e_1, e_2) \models_? (\xi_1, \xi_2)$	by Rule (17i) on (17) and (19)
(21) $(e_1, e_2) \models_?^{\dagger} (\xi_1, \xi_2)$	by Rule $(18a)$ on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.9.

Case (15i).

(22)
$$e_1 \models \xi_1$$
 by assumption Contradicts (16).

Case
$$e_1 \models_? \xi_1, e_2 \not\models^? \xi_2$$
.

 $(16) \ e_1 \not\models \xi_1$ by assumption
 $(17) \ e_1 \models_? \xi_1$ by assumption
 $(18) \ e_2 \not\models \xi_2$ by assumption
 $(19) \ e_2 \not\models_? \xi_2$ by assumption
Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case $(15j)$.

 $(20) \ (e_1, e_2)$ not intro by assumption
Contradicts Lemma 4.0.9.

Case $(15i)$.

 $(20) \ e_1 \models \xi_1$ by assumption
Contradicts (16).

$$(21) \ (e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction
Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case $(17b)$.

 $(22) \ (e_1, e_2)$ not intro by assumption
Contradicts Lemma 4.0.9.

Case $(17g)$.

 $(22) \ (e_1, e_2)$ not intro
Contradicts (18).

Case $(17h)$.

 $(22) \ e_2 \models_? \xi_2$ by assumption
Contradicts (19).

Case $(17i)$.

 $(22) \ e_2 \models_? \xi_2$ by assumption
Contradicts (19) .

Case $(17i)$.

 $(22) \ e_2 \models_? \xi_2$ by assumption
Contradicts (19) .

Case $(17i)$.

 $(22) \ e_2 \models_? \xi_2$ by assumption
Contradicts (19) .

Case $(17i)$.

 $(22) \ e_2 \models_? \xi_2$ by assumption
Contradicts (19) .

Case $(17i)$.

 $(22) \ e_2 \models_? \xi_2$ by contradiction
Contradicts (19) .

Case $(17i)$.

 $(23) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction
 $(24) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on
 (21) and (23)

Case $e_1 \not\models_? \xi_1$ by assumption

Case $e_1 \not\models_? \xi_1, e_2 \models_\xi 2$.

 $(16) \ e_1 \not\models_? \xi_1$ by assumption
by assumption

by assumption

(18) $e_2 \models \xi_2$

(19)
$$e_2 \not\models_? \xi_2$$
 by assumption ssume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) $e_2 \not\models_? \xi_2$

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (15i).

(20) $e_1 \models \xi_1$ by assumption Contradicts (16).

(21) $(e_1, e_2) \not\models (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction (24) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Case $e_1 \not\models_?^{\dagger} \xi_1, e_2 \models_? \xi_2$.

 $\begin{array}{lll} (16) & e_1 \not\models \xi_1 & & \text{by assumption} \\ (17) & e_1 \not\models_? \xi_1 & & \text{by assumption} \\ (18) & e_2 \not\models \xi_2 & & \text{by assumption} \\ (19) & e_2 \models_? \xi_2 & & \text{by assumption} \\ \end{array}$

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

Contradicts (18).

(21)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (16).

Case (17i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

(24) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Case $e_1 \not\models_? \xi_1$ by assumption

Case $e_1 \not\models_? \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \not\models_? \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (15) on it, only two cases apply.

Case (15i).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (15i).

(20) $e_2 \models \xi_2$ by assumption

Contradicts (18).

(21) $(e_1, e_2) \not\models (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by rule induction over Rules (17) on it, the following cases apply.

Case (17b).

(22)
$$(e_1, e_2)$$
 notintro by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (17h).

(22) $e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (17i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

(24) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Definition 2.1.1 (Entailment of Constraints). Suppose that $\xi_1 : \tau$ and $\xi_2 : \tau$. Then $\xi_1 \models \xi_2$ iff for all e such that $\cdot ; \Delta \vdash e : \tau$ and e val we have $e \models_{?}^{\dagger} \xi_1$ implies $e \models \xi_2$

Definition 2.1.2 (Potential Entailment of Constraints). Suppose that $\xi_1: \tau$ and $\xi_2: \tau$. Then $\xi_1 \models_{?}^{\dagger} \xi_2$ iff for all e such that $\cdot; \Delta \vdash e: \tau$ and e final we have $e \models_{?}^{\dagger} \xi_1$ implies $e \models_{?}^{\dagger} \xi_2$

Corollary 2.1.1. Suppose that $\xi : \tau \text{ and } \cdot ; \Delta \vdash e : \tau \text{ and } e \text{ final. Then } \top \models_{?}^{\uparrow} \xi \text{ implies } e \models_{?}^{\dagger} \xi$

Proof.

(1) $\xi:\tau$	by assumption
$(2) \ \cdot ; \Gamma \vdash e : \tau$	by assumption
(3) e final	by assumption
$(4) \top \models^{\dagger}_{?} \xi$	by assumption
(5) $e_1 \models \top$	by Rule (15a)
(6) $e_1 \models_?^\dagger \top$	by Rule (18b) on (5)
$(7) \ \ \top : \tau$	by Rule (9a)
$(8) e_1 \models^{\dagger}_{?} \xi_r$	by Definition 2.1.2 of (4) on (7) and (1) and (2) and (3) and (6)

3 Static Semantics

$$\begin{array}{lll} \tau & ::= & \operatorname{num} \mid (\tau_1 \to \tau_2) \mid (\tau_1 \times \tau_2) \mid (\tau_1 + \tau_2) \\ e & ::= & x \mid \underline{n} \\ & \mid & (\lambda x : \tau.e) \mid e_1(e_2) \\ & \mid & (e_1, e_2) \\ & \mid & \inf_{\tau}(e) \mid \operatorname{inr}_{\tau}(e) \mid \operatorname{match}(e) \{ \hat{rs} \} \\ & \mid & \emptyset^u \mid (e)^u \\ \hat{rs} & ::= & (rs \mid r \mid rs) \\ rs & ::= & \cdot \mid (r \mid rs') \\ r & ::= & p \Rightarrow e \\ \underline{p} & ::= & x \mid \underline{n} \mid _ \mid (p_1, p_2) \mid \operatorname{inl}(p) \mid \operatorname{inr}(p) \mid (\emptyset^w \mid (p)^w) \\ \hline (\hat{rs})^\diamond = rs & rs \text{ can be obtained by erasing pointer from } \hat{rs} \end{array}$$

$$(\cdot \mid r \mid rs)^{\diamond} = r \mid rs \tag{19a}$$

$$((r' \mid rs') \mid r \mid rs)^{\diamond} = r' \mid (rs' \mid r \mid rs)^{\diamond} \tag{19b}$$

 $\Gamma ; \Delta \vdash e : \tau \mid e \text{ is of type } \tau$

$$\frac{\text{TVar}}{\Gamma, x : \tau \; ; \Delta \vdash x : \tau} \tag{20a}$$

TEHole

$$\frac{1}{\Gamma; \Delta, u :: \tau \vdash ()^u : \tau} \tag{20b}$$

THole

$$\frac{\Gamma; \Delta, u :: \tau \vdash e : \tau'}{\Gamma; \Delta, u :: \tau \vdash (e)^u : \tau}$$
(20c)

TNum

$$\frac{}{\Gamma ; \Delta \vdash n : \mathtt{num}} \tag{20d}$$

TLam

$$\frac{\Gamma, x : \tau_1 ; \Delta \vdash e : \tau_2}{\Gamma; \Delta \vdash (\lambda x : \tau_1.e) : (\tau_1 \to \tau_2)}$$
 (20e)

TAp

$$\frac{\Gamma ; \Delta \vdash e_1 : (\tau_2 \to \tau) \qquad \Gamma ; \Delta \vdash e_2 : \tau_2}{\Gamma ; \Delta \vdash e_1(e_2) : \tau}$$
 (20f)

TPair

$$\frac{\Gamma; \Delta \vdash e_1 : \tau_1 \qquad \Gamma; \Delta \vdash e_2 : \tau_2}{\Gamma; \Delta \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$
(20g)

TPrl
$$\frac{\Gamma ; \Delta \vdash e : (\tau_1 \times \tau_2)}{\Gamma : \Delta \vdash \mathtt{prl}(e) : \tau_1}$$
(20h)

$$\frac{\text{TPrr}}{\Gamma; \Delta \vdash e : (\tau_1 \times \tau_2)}$$

$$\frac{\Gamma; \Delta \vdash \text{prr}(e) : \tau_2}{\Gamma; \Delta \vdash \text{prr}(e) : \tau_2}$$

$$(20i)$$

TInl

$$\frac{\Gamma ; \Delta \vdash e : \tau_1}{\Gamma ; \Delta \vdash \operatorname{inl}_{\tau_2}(e) : (\tau_1 + \tau_2)}$$
(20j)

TInr

$$\frac{\Gamma \; ; \; \Delta \vdash e : \tau_2}{\Gamma \; ; \; \Delta \vdash \mathsf{inr}_{\tau_1}(e) : (\tau_1 + \tau_2)} \tag{20k}$$

TMatchZPre

$$\frac{\Gamma ; \Delta \vdash e : \tau \qquad \Gamma ; \Delta \vdash [\bot]r \mid rs : \tau[\xi] \Rightarrow \tau' \qquad \top \models_{?}^{\dagger} \xi}{\Gamma ; \Delta \vdash \mathsf{match}(e)\{\cdot \mid r \mid rs\} : \tau'} \tag{201}$$

TMatchNZPre

$$\frac{\Gamma ; \Delta \vdash e : \tau \quad e \text{ final} \quad \Gamma ; \Delta \vdash [\bot] r s_{pre} : \tau[\xi_{pre}] \Rightarrow \tau'}{\Gamma ; \Delta \vdash [\bot \lor \xi_{pre}] r \mid r s_{post} : \tau[\xi_{rest}] \Rightarrow \tau' \quad e \not\models_{?}^{\dagger} \xi_{pre} \quad \top \models_{?}^{\dagger} \xi_{pre} \lor \xi_{rest}}{\Gamma ; \Delta \vdash \mathsf{match}(e) \{r s_{pre} \mid r \mid r s_{post}\} : \tau'}$$

$$(20m)$$

 $p:\tau[\xi]\dashv \Gamma;\Delta$ p is assigned type τ and emits constraint ξ

PTVai

$$\frac{}{x:\tau[\top] \dashv \cdot ; x:\tau} \tag{21a}$$

PTWild

PTEHole

PTHole

$$\frac{p : \tau[\xi] \dashv \Gamma; \Delta}{(p)^w : \tau'[?] \dashv \Gamma; \Delta, w :: \tau'}$$
(21d)

PTNum

$$\frac{\underline{n}: \operatorname{num}[\underline{n}] \dashv \cdot;}{\underline{n}} \tag{21e}$$

PTInl

$$\frac{p:\tau_1[\xi]\dashv \Gamma;\Delta}{\mathtt{inl}(p):(\tau_1+\tau_2)[\mathtt{inl}(\xi)]\dashv \Gamma;\Delta} \tag{21f}$$

PTInr

$$\frac{p : \tau_2[\xi] \dashv \Gamma ; \Delta}{\operatorname{inr}(p) : (\tau_1 + \tau_2)[\operatorname{inr}(\xi)] \dashv \Gamma ; \Delta}$$
(21g)

PTPair

$$\frac{p_1:\tau_1[\xi_1]\dashv \Gamma_1\;;\; \Delta_1 \qquad p_2:\tau_2[\xi_2]\dashv \Gamma_2\;;\; \Delta_2}{(p_1,p_2):(\tau_1\times\tau_2)[(\xi_1,\xi_2)]\dashv \Gamma_1\uplus\Gamma_2\;;\; \Delta_1\uplus\Delta_2} \tag{21h}$$

$$\frac{\Gamma; \Delta \vdash (p \Rightarrow e) : \tau[\xi] \Rightarrow \tau'}{\text{CTRule}} \qquad \begin{array}{c} r \text{ transforms a final expression of type } \tau \\ \text{to a final expression of type } \tau' \\ \hline \frac{p : \tau[\xi] \dashv \Gamma_p; \Delta_p \qquad \Gamma \uplus \Gamma_p; \Delta \uplus \Delta_p \vdash e : \tau'}{\Gamma; \Delta \vdash p \Rightarrow e : \tau[\xi] \Rightarrow \tau'} \end{aligned} \tag{22a}$$

 $\frac{\Gamma ; \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau'}{\text{CTOneRules}} \quad \begin{array}{c} rs \text{ transforms a final expression of type } \tau \\ \text{to a final expression of type } \tau' \end{array}$

$$\frac{\Gamma ; \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \qquad \xi_r \not\models \xi_{pre}}{\Gamma ; \Delta \vdash [\xi_{pre}](r \mid \cdot) : \tau[\xi_r] \Rightarrow \tau'}$$
(23a)

CTRules

$$\frac{\Gamma ; \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \qquad \Gamma ; \Delta \vdash [\xi_{pre} \lor \xi_r] rs : \tau[\xi_{rs}] \Rightarrow \tau' \qquad \xi_r \not\models \xi_{pre}}{\Gamma ; \Delta \vdash [\xi_{pre}] r \mid rs : \tau[\xi_r \lor \xi_{rs}] \Rightarrow \tau'}$$
(23b)

Lemma 3.0.1. If $p : \tau[\xi] \dashv \Gamma$; Δ then $\xi : \tau$.

Proof. By rule induction over Rules
$$(21)$$
.

Lemma 3.0.2. If \cdot ; $\Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau \text{ then } \xi_r : \tau_1$.

Proof. By rule induction over Rules
$$(22)$$
.

Lemma 3.0.3. If
$$\cdot$$
; $\Delta \vdash [\xi_{pre}]rs : \tau_1[\xi_{rs}] \Rightarrow \tau \ then \ \xi_{rs} : \tau_1$.

Proof. By rule induction over Rules
$$(23)$$
.

Lemma 3.0.4. If $\Gamma : \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau' \text{ and } \Gamma : \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \text{ and } \xi_r \not\models \xi_{pre} \lor \xi_{rs} \text{ then } \Gamma : \Delta \vdash [\xi_{pre}](rs \mid r \mid \cdot)^{\diamond} : \tau[\xi_{rs} \lor \xi_r] \Rightarrow \tau'$

Proof.

- (1) $\Gamma : \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau'$ by assumption
- (2) $\Gamma : \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau'$ by assumption
- (3) $\xi_r \not\models \xi_{pre} \vee \xi_{rs}$ by assumption

By rule induction over Rules (23) on (1).

Case (23a).

(4)
$$rs = r' \mid \cdot$$
 by assumption

(5)
$$\xi_{rs} = \xi'_r$$
 by assumption

(6)
$$\Gamma : \Delta \vdash r' : \tau[\xi_r'] \Rightarrow \tau'$$
 by assumption

(7)
$$\xi_r' \not\models \xi_{pre}$$
 by assumption

(8)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r]r \mid \cdot : \tau[\xi_r] \Rightarrow \tau'$$
 by Rule (23a) on (2) and (3)

(9)
$$\Gamma ; \Delta \vdash [\xi_{pre}](r' \mid r \mid \cdot) : \tau[\xi'_r \lor \xi_r] \Rightarrow \tau'$$
 by Rule (23b) on (6) and (8) and (7)

$$\begin{array}{ll} (10) \ \Gamma \, ; \Delta \vdash [\xi_{pre}]((r' \mid \cdot) \mid r \mid \cdot)^{\diamond} : \tau[\xi'_r \lor \xi_r] \Rightarrow \tau' \\ & \text{by Definition 19 on (9)} \end{array}$$

Case (23b).

(4)
$$rs = r' \mid rs'$$
 by assumption

(5)
$$\xi_{rs} = \xi'_r \vee \xi'_{rs}$$
 by assumption

(6)
$$\Gamma : \Delta \vdash r' : \tau[\xi'_r] \Rightarrow \tau'$$
 by assumption

(7)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r]rs' : \tau[\xi'_{rs}] \Rightarrow \tau'$$
 by assumption

(8)
$$\xi_r' \not\models \xi_{pre}$$
 by assumption

(9)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r](rs' \mid r \mid \cdot)^{\diamond} : \tau[\xi'_{rs} \lor \xi_r] \Rightarrow \tau'$$
 by IH on (7) and (2) and (3)

(10)
$$\Gamma$$
; $\Delta \vdash [\xi_{pre}](r' \mid (rs' \mid r \mid \cdot)^{\diamond}) : \tau[\xi'_r \lor \xi'_{rs} \lor \xi_r] \Rightarrow \tau'$ by Rule (23b) on (6) and (9) and (8)

(11)
$$\Gamma ; \Delta \vdash [\xi_{pre}]((r' \mid rs') \mid r \mid \cdot)^{\diamond} : \tau[\xi'_r \lor \xi'_{rs} \lor \xi_r] \Rightarrow \tau'$$
 by Definition 19 on (10)

Lemma 3.0.5 (Substitution). If $\Gamma, x : \tau ; \Delta \vdash e_0 : \tau_0 \text{ and } \Gamma ; \Delta \vdash e : \tau \text{ then } \Gamma ; \Delta \vdash [e/x]e_0 : \tau_0$

Lemma 3.0.6 (Simultaneous Substitution). *If* $\Gamma \uplus \Gamma'$; $\Delta \vdash e : \tau$ *and* $\theta : \Gamma'$ *then* $\Gamma ; \Delta \vdash [\theta]e : \tau$

Lemma 3.0.7 (Substitution Typing). If $e \rhd p \dashv \theta$ and \cdot ; $\Delta_e \vdash e : \tau$ and $p : \tau[\xi] \dashv \Gamma$; Δ then $\theta : \Gamma$

Proof by induction on the derivation of $e \triangleright p \dashv \theta$.

Theorem 3.1 (Determinism). If \cdot ; $\Delta \vdash e : \tau$ then exactly one of the following holds

- 1. e val
- 2. e indet
- 3. $e \mapsto e'$ for some unique e'

102

4 Dynamic Semantics

e val e is a value

$$\frac{\text{VNum}}{n \text{ val}} \tag{24a}$$

VLam

$$\frac{}{(\lambda x:\tau.e)\;\mathrm{val}} \tag{24b}$$

VPair

$$\frac{e_1 \text{ val} \qquad e_2 \text{ val}}{(e_1, e_2) \text{ val}} \tag{24c}$$

VInl

$$\frac{e \text{ val}}{\text{inl}_{\tau}(e) \text{ val}} \tag{24d}$$

VInr

$$\frac{e \; \mathtt{val}}{\mathtt{inr}_\tau(e) \; \mathtt{val}} \tag{24e}$$

e indet e is indeterminate

IEHole

$$\sqrt{\|)^u \text{ indet}}$$

IHole

$$\frac{e \; \text{final}}{(e)^u \; \text{indet}} \tag{25b}$$

IAp

$$\frac{e_1 \; \mathtt{indet} \quad e_2 \; \mathtt{final}}{e_1(e_2) \; \mathtt{indet}} \tag{25c}$$

IPairL

$$\frac{e_1 \; \mathtt{indet} \qquad e_2 \; \mathtt{val}}{(e_1, e_2) \; \mathtt{indet}} \tag{25d}$$

IPairR

$$\frac{e_1 \text{ val} \qquad e_2 \text{ indet}}{(e_1, e_2) \text{ indet}} \tag{25e}$$

IPair

$$\frac{e_1 \; \mathtt{indet} \quad e_2 \; \mathtt{indet}}{(e_1,e_2) \; \mathtt{indet}} \tag{25f}$$

IPrl

$$\frac{e \; \mathtt{indet}}{\mathtt{prl}(e) \; \mathtt{indet}} \tag{25g}$$

$$\frac{\operatorname{IPrr}}{e \text{ indet}} \qquad (25h)$$

$$\operatorname{IInL} \qquad e \text{ indet} \qquad (25i)$$

$$\operatorname{IInR} \qquad e \text{ indet} \qquad (25i)$$

$$\operatorname{IInR} \qquad e \text{ indet} \qquad (25i)$$

$$\operatorname{IIInR} \qquad e \text{ indet} \qquad (25j)$$

$$\operatorname{IIInthe} \qquad e \text{ final} \qquad e \text{ operator} \qquad (25j)$$

$$\operatorname{IIInthe} \qquad e \text{ final} \qquad e \text{ operator} \qquad (25k)$$

$$e \text{ final} \qquad e \text{ is final}$$

$$\operatorname{FVal} \qquad e \text{ val} \qquad e \text{ operator} \qquad (26a)$$

$$\operatorname{FIndet} \qquad e \text{ indet} \qquad e \text{ in$$

$$notintro(\emptyset)^u) = true$$
 (28a)

$$notintro(\langle e \rangle^u) = true$$
 (28b)

$$notintro(e_1(e_2)) = true$$
 (28c)

$$notintro(match(e)\{\hat{rs}\}) = true$$
 (28d)

$$notintro(prl(e)) = true$$
 (28e)

$$notintro(prr(e)) = true$$
 (28f)

Otherwise
$$notintro(e) = false$$
 (28g)

Lemma 4.0.1 (Soundness and Completeness of NotIntro Judgment). e notintro $\it iff\ not intro(e)$.

 $e' \in \mathtt{values}(e)$ e' is one of the possible values of e

$$\frac{e \text{ val} \qquad \cdot ; \Delta \vdash e : \tau}{e \in \text{values}(e)}$$
 (29a)

IVIndet

$$\frac{e \text{ notintro} \qquad \cdot ; \Delta \vdash e : \tau \qquad e' \text{ val} \qquad \cdot ; \Delta \vdash e' : \tau}{e' \in \text{values}(e)}$$
 (29b)

IVInl

$$\frac{\mathtt{inl}_{\tau_2}(e_1)\ \mathtt{indet} \quad \cdot \ ; \ \Delta \vdash \mathtt{inl}_{\tau_2}(e_1) : \tau \quad e_1' \in \mathtt{values}(e_1)}{\mathtt{inl}_{\tau_2}(e_1') \in \mathtt{values}(\mathtt{inl}_{\tau_2}(e_1))} \tag{29c}$$

 ${\rm IVInr}$

$$\frac{\operatorname{inr}_{\tau_1}(e_2) \operatorname{indet} \quad \cdot ; \Delta \vdash \operatorname{inr}_{\tau_1}(e_2) : \tau \quad e_2' \in \operatorname{values}(e_2)}{\operatorname{inr}_{\tau_1}(e_2') \in \operatorname{values}(\operatorname{inr}_{\tau_1}(e_2))} \tag{29d}$$

IVPair

$$\frac{(e_1,e_2) \text{ indet } \cdot ; \Delta \vdash (e_1,e_2) : \tau \qquad e_1' \in \mathtt{values}(e_1) \qquad e_2' \in \mathtt{values}(e_2)}{(e_1',e_2') \in \mathtt{values}((e_1,e_2))} \tag{29e}$$

Lemma 4.0.2. If e indet and \cdot ; $\Delta \vdash e : \tau$ and $\dot{\xi} : \tau$ and $e \not\models_{2}^{\dagger} \dot{\xi}$ then $e' \not\models_{2}^{\dagger} \dot{\xi}$ whenever $e' \in \mathtt{values}(e)$.

Proof.

(1)
$$e$$
 indet by assumption
(2) \cdot ; $\Delta \vdash e : \tau$ by assumption
(3) $\dot{\xi} : \tau$ by assumption
(4) $e \not\models_2^{\dagger} \dot{\xi}$ by assumption

By rule induction over Rules (9) on (3).

Case (9a).

 $(5) \ \dot{\xi} = \top$

by assumption

(6) $e \models \top$

by Rule (15a)

(7) $e \models_?^\dagger \top$

by Rule (18b) on (6)

Contradicts (4).

Case (1b).

(5) $\dot{\xi} = ?$

by assumption

(6) $e \models_? ?$

by Rule (17a)

(7) $e \models_{?}^{\dagger} ?$

by Rule (18a) on (6)

Contradicts (4).

Case (9c).

(5) $\dot{\xi} = \underline{n}$

by assumption

(6) $\tau = \text{num}$

by assumption

(7) \underline{n} refutable?

by Rule (11a)

By rule induction over Rules (25) on (1).

Case (25a).

(8) $e = (1)^u$

by assumption

(9) $()^u$ notintro

by Rule (27a)

 $(10) \ (\!)^u \models_? \underline{n}$

- by Rule (17b) on (9)
- and (7)

 $(11) \ (\!)^u \models^{\dagger}_? \underline{n}$

by Rule (18a) on (10)

Contradicts (4).

Case (25b).

(8) $e = (e_1)^u$

by assumption

(9) $(e_1)^u$ notintro

by Rule (27b)

 $(10) (e_1)^u \models_? \underline{n}$

- by Rule (17b) on (9)
- and (7)
- by Rule (18a) on (10)

 $(11) (|e_1|)^u \models^{\dagger}_{?} \underline{n}$

Contradicts (4).

Case (25c).

(8) $e = e_1(e_2)$

by assumption

(9) $e_1(e_2)$ notintro

by Rule (27c)

(10)
$$e_1(e_2) \models_? \underline{n}$$
 by Rule (17b) on (9) and (7)

(11)
$$e_1(e_2) \models_{?}^{\dagger} \underline{n}$$
 by Rule (18a) on (10)

Contradicts (4).

Case (25g).

(8)
$$e = prl(e_1)$$
 by assumption
(9) $prl(e_1)$ notintro by Rule (27e)
(10) $prl(e_1) \models_2 n$ by Rule (17b) on

(10)
$$\operatorname{prl}(e_1) \models_? \underline{n}$$
 by Rule (17b) on (9) and (7)

(11)
$$\operatorname{prl}(e_1) \models_{?}^{\dagger} \underline{n}$$
 by Rule (18a) on (10)

Contradicts (4).

Case (25h).

(8)
$$e = prr(e_1)$$
 by assumption
(9) $prr(e_1)$ notintro by Rule (27f)
(10) $prr(e_1) \models_? \underline{n}$ by Rule (17b) on (9)
and (7)

(11)
$$\operatorname{prr}(e_1) \models_{?}^{\dagger} \underline{n}$$
 by Rule (18a) on (10)

Contradicts (4).

Case (25k).

(8)
$$e = \text{match}(e_1)\{\hat{rs}\}\$$
 by assumption
(9) $\text{match}(e_1)\{\hat{rs}\}\$ notintro by Rule (27d)
(10) $\text{match}(e_1)\{\hat{rs}\}\models_? \underline{n}$ by Rule (17b) on (9)
and (7)
(11) $\text{match}(e_1)\{\hat{rs}\}\models_? n$ by Rule (18a) on (10)

Contradicts (4).

Case (25d), (25e), (25f).

(8)
$$e = (e_1, e_2)$$
 by assumption

By rule induction over Rules (20) on (2), no rule applies due to syntactic contradiction.

Case (25i).

(8)
$$e = inl_{\tau_2}(e_1)$$
 by assumption

By rule induction over Rules (20) on (2), no rule applies due to syntactic contradiction.

Case (25j).

(8)
$$e = inr_{\tau_1}(e_2)$$
 by assumption

By rule induction over Rules (20) on (2), no rule applies due to syntactic contradiction.

Case (9g).

$(5) \ \dot{\xi} = \mathtt{inl}(\dot{\xi}_1)$	by assumption
(6) $\tau = (\tau_1 + \tau_2)$	by assumption
(7) $\dot{\xi}_1 : \tau_1$	by assumption
(8) inl $(\dot{\xi}_1)$ refutable?	by Rule (11b)

By rule induction over Rules (25) on (1).

Case (25a).

(9)
$$e = \emptyset^u$$
 by assumption
(10) \emptyset^u notintro by Rule (27a)
(11) $\emptyset^u \models_? \operatorname{inl}(\dot{\xi_1})$ by Rule (17b) on (10)
and (8)
(12) $\emptyset^u \models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$ by Rule (18a) on (11)

Contradicts (4).

Case (25b).

(9)
$$e = (e_1)^u$$
 by assumption
(10) $(e_1)^u$ notintro by Rule (27b)
(11) $(e_1)^u \models_? \text{inl}(\dot{\xi_1})$ by Rule (17b) on (10)
and (8)
(12) $(e_1)^u \models_?^{\dagger} \text{inl}(\dot{\xi_1})$ by Rule (18a) on (11)

Contradicts (4).

Case (25c).

(9)
$$e = e_1(e_2)$$
 by assumption
(10) $e_1(e_2)$ notintro by Rule (27c)
(11) $e_1(e_2) \models_? inl(\dot{\xi_1})$ by Rule (17b) on (10)
and (8)
(12) $e_1(e_2) \models_?^{\dagger} inl(\dot{\xi_1})$ by Rule (18a) on (11)

Contradicts (4).

Case (25g).

(9)
$$e = prl(e_1)$$
 by assumption
(10) $prl(e_1)$ notintro by Rule (27e)
(11) $prl(e_1) \models_? inl(\dot{\xi_1})$ by Rule (17b) on (10)
and (8)
(12) $prl(e_1) \models_?^{\dagger} inl(\dot{\xi_1})$ by Rule (18a) on (11)

Contradicts (4).

Case (25h).

- (9) $e = prr(e_1)$ by assumption (10) $prr(e_1)$ notintro by Rule (27f)
- (11) $\operatorname{prr}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$ by Rule (17b) on (10) and (8)
- (12) $\operatorname{prr}(e_1) \models^{\dagger}_{?} \operatorname{inl}(\dot{\xi}_1)$ by Rule (18a) on (11)

Contradicts (4).

Case (25k).

- (9) $e = \text{match}(e_1)\{\hat{rs}\}$ by assumption (10) $\text{match}(e_1)\{\hat{rs}\}$ notintro by Rule (27d)
- (11) $\operatorname{match}(e_1)\{\hat{rs}\} \models_? \operatorname{inl}(\dot{\xi_1})$ by Rule (17b) on (10) and (8)
- (12) $\mathrm{match}(e_1)\{\hat{rs}\}\models^{\dagger}_{?}\mathrm{inl}(\dot{\xi_1})$ by Rule (18a) on (11)

Contradicts (4).

Case (25d), (25e), (25f).

(9)
$$e = (e_1, e_2)$$
 by assumption

By rule induction over Rules (25) on (1), no rule applies due to syntactic contradiction.

Case (25i).

- (9) $e = \operatorname{inl}_{\tau_2'}(e_1)$ by assumption
- (10) e_1 indet by assumption

By rule induction over Rules (20) on (2), only one rule applies.

Case (20j).

- (11) $\tau_2' = \tau_2$ by assumption (12) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (13) $e_1 \not\models_?^{\dagger} \dot{\xi}_1$ by Lemma 2.0.11 on (4)
- (14) if $e_1' \in \mathtt{values}(e_1)$ then $e_1' \not\models_?^\dagger \dot{\xi}_1$

by IH on (10) and (12) and (7) and (13)

To show if $e' \in \mathtt{values}(\mathtt{inl}_{\tau_2}(e_1))$ then $e' \not\models_?^\dagger \mathtt{inl}(\dot{\xi}_1)$, we assume $e' \in \mathtt{values}(\mathtt{inl}_{\tau_2}(e_1))$.

(15) $e' \in \text{values}(\text{inl}_{\tau_2}(e_1))$ by assumption

By rule induction over Rules (29) on (15).

Case (29a).

(16) $\operatorname{inl}_{\tau_2}(e_1)$ val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(16) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

Contradicts Lemma 4.0.7

Case (29c).

- (16) $e' = \operatorname{inl}_{\tau_2}(e'_1)$ by assumption (17) $e'_1 \in \operatorname{values}(e_1)$ by assumption
- (18) $e'_1 \not\models_2^{\dagger} \dot{\xi}_1$ by (14) on (17)
- (19) $\operatorname{inl}_{\tau_2}(e'_1) \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$ by Lemma 2.0.11 on (18)

Case (25j).

(9)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption

To show if $e' \in \mathtt{values}(\mathtt{inr}_{\tau_1}(e_2))$ then $e' \not\models_?^{\dagger} \mathtt{inl}(\dot{\xi}_1)$, we assume $e' \in \mathtt{values}(\mathtt{inr}_{\tau_1}(e_2))$.

(10)
$$e' \in values(inr_{\tau_1}(e_2))$$
 by assumption

By rule induction over Rules (29) on (10).

Case (29a).

(11) $\operatorname{inr}_{\tau_1}(e_2)$ val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(11) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption

Contradicts Lemma 4.0.8

Case (29d).

(11) $e' = \operatorname{inr}_{\tau_1}(e'_2)$ by assumption (12) $\operatorname{inr}_{\tau_1}(e'_2) \not\models_{\tau}^{\dagger} \operatorname{inl}(\dot{\xi_1})$ by Lemma 1.0.22

Case (9h).

(5) $\dot{\xi} = \operatorname{inr}(\dot{\xi}_2)$ by assumption (6) $\tau = (\tau_1 + \tau_2)$ by assumption (7) $\dot{\xi}_2 : \tau_2$ by assumption (8) $\operatorname{inr}(\dot{\xi}_2)$ refutable? by Rule (11c)

By rule induction over Rules (25) on (1).

Case (25a).

- (9) $e = ()^u$ by assumption (10) $()^u$ notintro by Rule (27a)
- (11) $\emptyset^u \models_? \operatorname{inr}(\dot{\xi}_2)$ by Rule (17b) on (10) and (8)
- (12) $\emptyset^u \models_?^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11)

Contradicts (4).

Case (25b). (9) $e = (e_1)^u$ by assumption (10) $(e_1)^u$ notintro by Rule (27b) (11) $(e_1)^u \models_? inr(\dot{\xi}_2)$ by Rule (17b) on (10) and (8)(12) $(e_1)^u \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11) Contradicts (4). Case (25c). (9) $e = e_1(e_2)$ by assumption (10) $e_1(e_2)$ notintro by Rule (27c) (11) $e_1(e_2) \models_? inr(\dot{\xi}_2)$ by Rule (17b) on (10) and (8)(12) $e_1(e_2) \models_2^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11) Contradicts (4). Case (25g). (9) $e = prl(e_1)$ by assumption (10) $prl(e_1)$ notintro by Rule (27e) (11) $\operatorname{prl}(e_1) \models_? \operatorname{inr}(\xi_2)$ by Rule (17b) on (10) and (8)(12) $\operatorname{prl}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11) Contradicts (4). Case (25h). (9) $e = prr(e_1)$ by assumption (10) $prr(e_1)$ notintro by Rule (27f) (11) $\operatorname{prr}(e_1) \models_? \operatorname{inr}(\xi_2)$ by Rule (17b) on (10) and (8)(12) $\operatorname{prr}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11) Contradicts (4). Case (25k). (9) $e = \operatorname{match}(e_1)\{\hat{rs}\}$ by assumption (10) $match(e_1)\{\hat{rs}\}\ notintro$ by Rule (27d) (11) $\operatorname{match}(e_1)\{\hat{rs}\} \models_? \operatorname{inr}(\dot{\xi}_2)$ by Rule (17b) on (10) and (8)(12) $\operatorname{match}(e_1)\{\hat{rs}\} \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (18a) on (11) Contradicts (4).

by assumption

Case (25d), (25e), (25f).

(9) $e = (e_1, e_2)$

By rule induction over Rules (25) on (1), no rule applies due to syntactic contradiction.

Case (25i).

(9)
$$e = inl_{\tau_2}(e_1)$$

by assumption

To show if $e' \in \mathtt{values}(\mathtt{inl}_{\tau_2}(e_1))$ then $e' \not\models_?^\dagger \mathtt{inr}(\dot{\xi}_2)$, we assume $e' \in \mathtt{values}(\mathtt{inl}_{\tau_2}(e_1))$.

(10)
$$e' \in \text{values}(\text{inl}_{\tau_2}(e_1))$$

by assumption

By rule induction over Rules (29) on (10).

Case (29a).

(11)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 val

by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(11)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro

by assumption

Contradicts Lemma 4.0.7

Case (29c).

$$(11) \ e'=\operatorname{inl}_{\tau_2}(e_1')$$

by assumption

(12)
$$\operatorname{inl}_{\tau_2}(e_1') \not\models_?^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$

by Lemma 1.0.21

Case (25j).

$$(9) \ e = \operatorname{inr}_{\tau_1'}(e_2)$$

by assumption

(10) e_2 indet

by assumption

By rule induction over Rules (20) on (2), only one rule applies.

Case (20k).

(11)
$$\tau_1' = \tau_1$$

by assumption

(12)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$

by assumption

(13)
$$e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$$

by Lemma 2.0.11 on

(4)

(14) if
$$e_2' \in \mathtt{values}(e_2)$$
 then $e_2' \not\models_?^\dagger \dot{\xi_2}$

by IH on (10) and (12) and (7) and (13)

To show if $e' \in \mathtt{values}(\mathtt{inr}_{\tau_1}(e_2))$ then $e' \not\models_?^\dagger \mathtt{inr}(\dot{\xi}_2)$, we assume $e' \in \mathtt{values}(\mathtt{inr}_{\tau_1}(e_2))$.

(15) $e' \in values(inr_{\tau_1}(e_2))$ by assumption

By rule induction over Rules (29) on (15).

Case (29a).

(16)
$$inr_{\tau_1}(e_2) val$$

by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

 $(16) \ \operatorname{inr}_{\tau_1}(e_2) \ \operatorname{notintro}$

by assumption

Contradicts Lemma 4.0.8

Case (29d).

(16)
$$e' = \inf_{\tau_1}(e'_2)$$
 by assumption
(17) $e'_2 \in \text{values}(e_2)$ by assumption
(18) $e'_2 \not\models_?^{\dagger} \dot{\xi}_2$ by (14) on (17)
(19) $\inf_{\tau_1}(e'_2) \not\models_?^{\dagger} \inf(\dot{\xi}_2)$ by Lemma 2.0.12 on (18)

Case (9i).

(5)
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption
(6) $\tau = (\tau_1 \times \tau_2)$ by assumption
(7) $\dot{\xi}_1 : \tau_1$ by assumption
(8) $\dot{\xi}_2 : \tau_2$ by assumption

By rule induction over Rules (25) on (1).

Case (25a), (25b), (25c), (25g), (25h), (25k).

$$(9) \ \ e = \{ \}^u, \{ e_1 \}^u, e_1(e_2), \operatorname{prl}(e_1), \operatorname{prr}(e_1), \operatorname{match}(e_1) \{ \hat{rs} \} \}$$
 by assumption
$$(10) \ \ e \ \operatorname{notintro} \qquad \qquad \operatorname{by} \ \operatorname{Rules} \ (27)$$

$$(11) \ \ \operatorname{prl}(e) \ \operatorname{notintro} \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (27e)$$

$$(12) \ \ \operatorname{prr}(e) \ \operatorname{notintro} \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (27f)$$

$$(13) \ \ \operatorname{prl}(e) \ \operatorname{indet} \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (25g) \ \operatorname{on} \ (1)$$

$$(14) \ \ \operatorname{prr}(e) \ \operatorname{indet} \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (25h) \ \operatorname{on} \ (1)$$

$$(15) \ \cdot \ ; \Delta \vdash \operatorname{prl}(e) : \tau_1 \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (20h) \ \operatorname{on} \ (2)$$

$$(16) \ \cdot \ ; \Delta \vdash \operatorname{prr}(e) : \tau_2 \qquad \qquad \operatorname{by} \ \operatorname{Rule} \ (20i) \ \operatorname{on} \ (2)$$

By case analysis on the result of $\mathit{satisfyormay}(\mathtt{prl}(e), \dot{\xi}_1).$

Case true.

(17)
$$satisfyormay(\mathtt{prl}(e),\dot{\xi_1}) = \mathrm{true}$$

(18) $\operatorname{prl}(e) \models_{?}^{\dagger} \dot{\xi}_{1}$ by assumption by Lemma 1.0.4 on (17)

By case analysis on the result of $satisfyormay(prr(e), \dot{\xi}_2)$. Case true.

(19)
$$satisfyormay(prr(e), \dot{\xi}_2) = true$$

by assumption

(20)
$$prr(e) \models_{?}^{\dagger} \dot{\xi}_{2}$$
 by Lemma 1.0.4 on (19)

By rule induction over Rules (18) on (18).

Case (18b).

(21)
$$\operatorname{prl}(e) \models \dot{\xi}_1$$
 by assumption

By rule induction over Rules (18) on (20).

Case (18b).

(22)
$$\operatorname{prr}(e) \models \dot{\xi}_2$$
 by assumption
(23) $e \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (15j) on (10)
and (21) and (22)

(24)
$$e \models_{?}^{\dagger} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (18b) on (23)

Contradicts (4).

Case (18a).

(22)
$$prr(e) \models_? \dot{\xi}_2$$
 by assumption
(23) $\dot{\xi}_2$ refutables by Lemma 1.0.17

(23)
$$\dot{\xi}_2$$
 refutable? by Lemma 1.0.17 on (12) and (22)

(24)
$$(\dot{\xi}_1, \dot{\xi}_2)$$
 refutable? by Rule (11e) on (23)
(25) $e \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17b) on (10)

(26)
$$e \models_{7}^{\dagger} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (18a) on (25)

and (24)

Case (18a).

(21)
$$prl(e) \models_? \dot{\xi}_1$$
 by assumption

(22)
$$\dot{\xi}_1$$
 refutable? by Lemma 1.0.17 on (11) and (21)

(23)
$$(\dot{\xi}_1, \dot{\xi}_2)$$
 refutable? by Rule (11d) on (22)

(24)
$$e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (17b) on (10) and (23)

(25)
$$e \models^{\dagger}_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by Rule (18a) on (24)

Case false.

(19)
$$satisfyormay(prr(e), \dot{\xi}_2) = false$$

by assumption

(20)
$$\operatorname{prr}(e) \models_{?}^{\dagger} \dot{\xi}_{2}$$
 by Lemma 1.0.4 on (19)

(21) if
$$e_2' \in \mathtt{values}(\mathtt{prr}(e))$$
 then $e_2' \not\models_?^\dagger \dot{\xi}_2$

by IH on (14) and (16) and (8) and (20)

To show if $e' \in \mathtt{values}(e)$ then $e' \not\models_?^\dagger (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}(e)$.

(22)
$$e' \in \mathtt{values}(e)$$
 by assumption

By rule induction over Rules (29) on (22), only two rules apply.

Case (29a).

(23)
$$e$$
 val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

$$(23)$$
 e' val

by assumption

$$(24) \cdot ; \Delta \vdash e' : (\tau_1 \times \tau_2)$$

by assumption

By rule induction over Rules (24) on (23).

Case (24a).

(25)
$$e' = \underline{n}$$

by assumption

By rule induction over Rules (20) on (24), no rule applies due to syntactic contradiction.

Case (24b).

(25)
$$e' = (\lambda x : \tau'.e'_1)$$

by assumption

By rule induction over Rules (20) on (24), no rule applies due to syntactic contradiction.

Case (24c).

$$(25) \ e' = (e'_1, e'_2)$$

by assumption

$$(26)$$
 e_2' val

by assumption

By rule induction over Rules (20) on (24), only one rule applies.

Case (20g).

(27)
$$\cdot$$
; $\Delta \vdash e_2' : \tau_2$

by assumption

$$(28)\ e_2' \in \mathtt{values}(\mathtt{prr}(e))$$

by Rule (29b) on (12)and (16) and (26) and

(27)

(29)
$$e_2' \not\models_?^\dagger \dot{\xi}_2$$

by (21) on (28)

(30)
$$(e'_1, e'_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Lemma 2.0.13 on (27)

Case (24d).

(25)
$$e' = \operatorname{inl}_{\tau_2}(e'_1)$$

by assumption

By rule induction over Rules (20) on (24), no rule applies due to syntactic contradiction.

Case (24e).

(25)
$$e' = inr_{\tau_1}(e'_2)$$

by assumption

By rule induction over Rules (20) on (24), no rule applies due to syntactic contradiction.

Case false.

(17) $satisfyormay(prl(e), \dot{\xi}_1) = false$

by assumption

(18)
$$\operatorname{prl}(e) \not\models_?^\dagger \dot{\xi_1}$$

by Lemma 1.0.4 on

 $\begin{array}{ll} \text{(19)} \ \text{if} \ e_1' \in \mathtt{values}(\mathtt{prl}(e)) \ \text{then} \ e_1' \not\models_?^\dagger \dot{\xi}_1 \\ \qquad \qquad \qquad \text{by IH on (13) and (15)} \end{array}$

and (7) and (18)

To show if $e' \in \mathtt{values}(e)$ then $e' \not\models^\dagger_? (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}(e)$.

(20)
$$e' \in \mathtt{values}(e)$$

by assumption

By rule induction over Rules (29) on (20), only two rules apply. Case (29a).

(21) e val

by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(21) e' val

by assumption

(22) \cdot ; $\Delta \vdash e' : (\tau_1 \times \tau_2)$

by assumption

By rule induction over Rules (24) on (21).

Case (24a).

(23)
$$e' = \underline{n}$$

by assumption

By rule induction over Rules (20) on (22), no rule applies due to syntactic contradiction.

Case (24b).

(23)
$$e' = (\lambda x : \tau'.e'_1)$$

by assumption

By rule induction over Rules (20) on (22), no rule applies due to syntactic contradiction.

Case (24c).

(23)
$$e' = (e'_1, e'_2)$$

by assumption

$$(24)$$
 e_1' val

by assumption

By rule induction over Rules (20) on (22), only one rule applies.

Case (20g).

$$(25) \cdot ; \Delta \vdash e_1' : \tau_1$$

by assumption

$$(26) \ e_1' \in \mathtt{values}(\mathtt{prl}(e))$$

by Rule (29b) on (11)

and (15) and (24) and

(25)

(27)
$$e'_1 \not\models_{?}^{\dagger} \dot{\xi}_1$$

by (19) on (26)

(28)
$$(e'_1, e'_2) \not\models_2^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Lemma 2.0.13 on

(27)

Case (24d).

(23)
$$e' = \operatorname{inl}_{\tau_2}(e'_1)$$

by assumption

By rule induction over Rules (20) on (22), no rule applies due to syntactic contradiction.

Case (24e).

(23)
$$e' = \operatorname{inr}_{\tau_1}(e'_2)$$

by assumption

By rule induction over Rules (20) on (22), no rule applies due to syntactic contradiction.

Case (25d).

(9)
$$e = (e_1, e_2)$$
 by assumption
(10) e_1 indet by assumption
(11) e_2 val by assumption
(12) $e_1 \not\models_?^{\dagger} \dot{\xi}_1$ or $e_2 \not\models_?^{\dagger} \dot{\xi}_2$ by Lemma 2.0.13 on
(4)

By case analysis on the disjunction in (12).

Case $e_1 \not\models_?^\dagger \dot{\xi}_1$.

(13)
$$e_1 \not\models_{\hat{1}}^{\dagger} \dot{\xi_1}$$
 by assumption

By rule induction over Rules (20) on (2), only one rule applies. Case (20g).

(14)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$ by assumption

(15) if
$$e_1' \in \mathtt{values}(e_1)$$
 then $e_1' \not\models_?^\dagger \dot{\xi}_1$ by IH on (10) and (14) and (7) and (13)

To show that if $e' \in \mathtt{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}((e_1, e_2))$.

(16)
$$e' \in values((e_1, e_2))$$
 by assumption

By rule induction over Rules (29) on (16).

Case (29a).

(17)
$$(e_1, e_2)$$
 val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(17)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.9.

Case (29e).

(17)
$$e' = (e'_1, e'_2)$$
 by assumption
(18) $e'_1 \in \text{values}(e_1)$ by assumption
(19) $e'_1 \not\models_?^{\dagger} \dot{\xi}_1$ by (15) on (18)
(20) $(e'_1, e'_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.13 on (19)

Case $e_2 \not\models_?^\dagger \dot{\xi}_2$.

(13)
$$e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$$
 by assumption

To show that if $e' \in \mathtt{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_?^\dagger (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}((e_1, e_2))$.

(14)
$$e' \in values((e_1, e_2))$$
 by assumption

By rule induction over Rules (29) on (14).

Case (29a).

(15)
$$(e_1, e_2)$$
 val by assumption

Contradicts (1) by Lemma 4.0.11. Case (29b). (15) (e_1,e_2) notintro by assumption Contradicts Lemma 4.0.9. Case (29e). (15) $e' = (e'_1, e'_2)$ by assumption (16) $e_2' \in \mathtt{values}(e_2)$ by assumption By rule induction over Rules (29) on (16). Case (29a). (17) $e_2' = e_2$ by assumption (18) $e_2' \not\models_2^{\dagger} \dot{\xi}_2$ by (17) and (13)(19) $(e'_1, e'_2) \not\models_2^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.13 on (18)Case (29b). (17) e_2 notintro by assumption Contradicts (11) by Lemma 4.0.12. Case (29c), (29d), (29e). (17) e_2 indet by assumption Contradicts (11) by Lemma 4.0.11. Case (25e). (9) $e = (e_1, e_2)$ by assumption (10) e_1 val by assumption (11) e_2 indet by assumption (12) $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1 \text{ or } e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$ by Lemma 2.0.13 on (4)By case analysis on the disjunction in (12). Case $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1$. (13) $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1$ by assumption To show that if $e' \in \text{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_7^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in values((e_1, e_2))$. (14) $e' \in \text{values}((e_1, e_2))$ by assumption By rule induction over Rules (29) on (14). Case (29a). (15) (e_1, e_2) val by assumption Contradicts (1) by Lemma 4.0.11. Case (29b). by assumption (15) (e_1,e_2) notintro Contradicts Lemma 4.0.9.

Case (29e).

```
(15) e' = (e'_1, e'_2)
                                        by assumption
(16) e_1' \in \text{values}(e_1)
                                       by assumption
```

By rule induction over Rules (29) on (16).

Case (29a).

(17)
$$e_1' = e_1$$
 by assumption

(18)
$$e'_1 \not\models_2^{\dagger} \dot{\xi}_1$$
 by (17) and (13)

(19)
$$(e'_1, e'_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 by Lemma 2.0.13 on (18)

Case (29b).

(17)
$$e_1$$
 notintro by assumption

Contradicts (10) by Lemma 4.0.12.

Case (29c), (29d), (29e).

(17)
$$e_1$$
 indet by assumption

Contradicts (10) by Lemma 4.0.11.

Case $e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$.

(13)
$$e_2 \not\models_7^{\dagger} \dot{\xi}_2$$
 by assumption

By rule induction over Rules (20) on (2), only one rule applies. Case (20g).

(14)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$ by assumption

(15) if
$$e_2' \in \mathtt{values}(e_2)$$
 then $e_2' \not\models_?^\dagger \dot{\xi}_2$ by IH on (11) and (14) and (8) and (13)

To show that if $e' \in \mathtt{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in values((e_1, e_2))$.

$$(16) \ e' \in \mathtt{values}((e_1, e_2)) \qquad \quad \text{by assumption}$$

By rule induction over Rules (29) on (16).

Case (29a).

(17)
$$(e_1, e_2)$$
 val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(17)
$$(e_1, e_2)$$
 notintro by assumption

Contradicts Lemma 4.0.9.

Case (29e).

(17)
$$e' = (e'_1, e'_2)$$
 by assumption
(18) $e'_2 \in \text{values}(e_2)$ by assumption
(19) $e'_2 \not\models_?^{\dagger} \dot{\xi}_2$ by (15) on (18)
(20) $(e'_1, e'_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.13 on

(19)

Case (25f).

(9) $e = (e_1, e_2)$ by assumption (10) e_1 indet by assumption (11) e_2 indet by assumption (12) $e_1 \not\models_?^\dagger \dot{\xi}_1$ or $e_2 \not\models_?^\dagger \dot{\xi}_2$ by Lemma 2.0.13 on (4)

By rule induction over Rules (20) on (2), only one rule applies.

Case (20g).

- (13) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (14) \cdot ; $\Delta \vdash e_2 : \tau_2$ by assumption

By case analysis on the disjunction in (12).

Case $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1$.

- (15) $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1$ by assumption
- (16) if $e'_1 \in \text{values}(e_1)$ then $e'_1 \not\models_{?}^{\dagger} \dot{\xi}_1$ by IH on (10) and (13) and (7) and (15)

To show that if $e' \in \mathtt{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}((e_1, e_2))$.

(17) $e' \in values((e_1, e_2))$ by assumption

By rule induction over Rules (29) on (17).

Case (29a).

(18) (e_1, e_2) val by assumption Contradicts (1) by Lemma 4.0.11.

Case (29b).

(18) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (29e).

- (18) $e' = (e'_1, e'_2)$ by assumption
- (19) $e_1' \in values(e_1)$ by assumption
- (20) $e'_1 \not\models_{?}^{\dagger} \dot{\xi}_1$ by (16) on (19)
- (21) $(e_1', e_2') \not\models_?^\dagger (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.13 on (20)

Case $e_2 \not\models_?^\dagger \dot{\xi}_2$.

- (15) $e_2 \not\models_?^{\dagger} \dot{\xi}_2$ by assumption
- (16) if $e'_2 \in \text{values}(e_2)$ then $e'_2 \not\models_{?}^{\dagger} \dot{\xi}_2$ by IH on (11) and (14) and (8) and (15)

To show that if $e' \in \mathtt{values}((e_1, e_2))$ then $(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$, we assume $e' \in \mathtt{values}((e_1, e_2))$.

(17) $e' \in values((e_1, e_2))$ by assumption

By rule induction over Rules (29) on (17).

Case (29a).

(18) (e_1, e_2) val by assumption

Contradicts (1) by Lemma 4.0.11.

Case (29b).

(18) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (29e).

(18) $e' = (e'_1, e'_2)$ by assumption (19) $e'_2 \in \text{values}(e_2)$ by assumption (20) $e'_2 \not\models_?^{\dagger} \dot{\xi}_2$ by (16) on (19) (21) $(e'_1, e'_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.13 on (20)

Case (25i).

(9)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

By rule induction over Rules (20) on (2), no rule applies due to syntactic contradiction.

Case (25j).

(9)
$$e = \operatorname{inr}_{\tau_1'}(e_2)$$
 by assumption

By rule induction over Rules (20) on (2), no rule applies due to syntactic contradiction.

Case (9f).

(5)	$\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$	by assumption
(6)	$\dot{\xi}_1: au_1$	by assumption
(7)	$\dot{\xi}_2: au_2$	by assumption
(8)	$e \not\models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$	by assumption
(9)	$e \not\models_?^\dagger \dot{\xi}_1$	by Lemma 2.0.10 on (8)
(10)	$e \not\models_?^\dagger \dot{\xi}_2$	by Lemma 2.0.10 on (8)
(11)	if $e' \in \mathtt{values}(e)$ then $e' \not\models_?^\dagger \dot{\xi}_1$	by IH on (1) and (2) and (6) and (9)
(12)	if $e' \in \mathtt{values}(e)$ then $e' \not\models_?^\dagger \dot{\xi}_2$	by IH on (1) and (2) and (7) and (10)

To show that if $e' \in \mathtt{values}(e)$ then $e' \not\models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$, we assume $e' \in \mathtt{values}(e)$.

```
(14) e' \not\models_?^\dagger \dot{\xi}_1
                                                                             by (11) on (13)
          (15) \ e' \not\models_{?}^{\dagger} \dot{\xi}_2
                                                                             by (12) on (13)
          (16) \ e' \not\models_?^\dagger \dot{\xi_1} \lor \dot{\xi_2}
                                                                             by Lemma 2.0.10 on
                                                                             (14) and (15)
                                                                                                            \theta:\Gamma
            \theta is of type \Gamma
                                                STEmpty
                                                                                                        (30a)
                                                 \overline{\emptyset:\cdot}
                                       STExtend
                                                  \Gamma ; \Delta \vdash e : \tau
                                       \theta : \Gamma_{\theta}
                                                                                                       (30b)
                                           \theta, x/e : \Gamma_{\theta}, x : \tau
p refutable?
                          p is refutable
                                            RNum
                                                                                                        (31a)
                                            \underline{n} refutable?
                                           REHole
                                                                                                        (31b)
                                           \overline{(\!(\!)^w \text{ refutable}_?}
                                          RHole
                                                                                                        (31c)
                                          \overline{(p)^w \text{ refutable}_?}
                                         RInl
                                                                                                       (31d)
                                         inl(p) refutable?
                                         RInr
                                                                                                        (31e)
                                         inr(p) refutable?
                                        RPairL
                                            p_1 refutable?
                                                                                                        (31f)
                                         (p_1, p_2) refutable?
                                        RPairR
                                            p_2 refutable?
                                                                                                        (31g)
                                         (p_1, p_2) refutable?
e matches p, emitting \theta
                                              MVar
                                                                                                        (32a)
```

by assumption

(13) $e' \in \mathtt{values}(e)$

$$\frac{\text{MWild}}{e \rhd _ \dashv \cdot} \tag{32b}$$

$$\frac{1}{n > n \dashv \cdot} \tag{32c}$$

MPair

$$\frac{e_1 \rhd p_1 \dashv \theta_1}{(e_1, e_2) \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2}$$

$$(32d)$$

MInl

$$\frac{e \rhd p \dashv \theta}{\operatorname{inl}_{\tau}(e) \rhd \operatorname{inl}(p) \dashv \theta} \tag{32e}$$

MInr

$$\frac{e \rhd p \dashv \theta}{\operatorname{inr}_{\tau}(e) \rhd \operatorname{inr}(p) \dashv \theta}$$
 (32f)

 ${\bf MNotIntroPair}$

$$\frac{e \text{ notintro}}{e \text{ notintro}} \frac{\text{prl}(e) \rhd p_1 \dashv \theta_1 \qquad \text{prr}(e) \rhd p_2 \dashv \theta_2}{e \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2}$$
(32g)

e?p e may match p

$$\frac{1}{e? \left(\right)^{w}} \tag{33a}$$

$$\frac{}{e?(p)^w} \tag{33b}$$

MMNotIntro

$$\frac{e \; \mathtt{notintro} \quad p \; \mathtt{refutable}_?}{e \; ? \; p} \tag{33c}$$

MMPairL

$$\frac{e_1? p_1}{(e_1, e_2)? (p_1, p_2)}$$
 (33d)

MMPairR

$$\frac{e_1 \triangleright p_1 \dashv \theta_1 \qquad e_2 ? p_2}{(e_1, e_2) ? (p_1, p_2)}$$
(33e)

MMPair

$$\frac{e_1?p_1 - e_2?p_2}{(e_1, e_2)?(p_1, p_2)}$$
 (33f)

MMInl

$$\frac{e?p}{\operatorname{inl}_{\tau}(e)?\operatorname{inl}(p)} \tag{33g}$$

$$\frac{\text{MMInr}}{e?p} \frac{e?p}{\text{inr}_{\tau}(e)? \text{inr}(p)}$$
 (33h)

 $e \perp p$ e does not match p

NMNum
$$\frac{n_1 \neq n_2}{n_1 \perp n_2} \tag{34a}$$

NMPairL

$$\frac{e_1 \perp p_1}{(e_1, e_2) \perp (p_1, p_2)} \tag{34b}$$

 ${\bf NMPairR}$

$$\frac{e_2 \perp p_2}{(e_1, e_2) \perp (p_1, p_2)} \tag{34c}$$

 ${\rm NMConfL}$

$$\frac{-}{\operatorname{inr}_{\tau}(e) \perp \operatorname{inl}(p)} \tag{34d}$$

 ${\rm NMConfR}$

$$\frac{}{\operatorname{inl}_{\tau}(e) \perp \operatorname{inr}(p)} \tag{34e}$$

NMInl

$$\frac{e \perp p}{\mathsf{inr}_{\tau}(e) \perp \mathsf{inl}(p)} \tag{34f}$$

NMInr

$$\frac{e \perp p}{\mathtt{inl}_{\tau}(e) \perp \mathtt{inr}(p)} \tag{34g}$$

 $e \mapsto e'$ e takes a step to e'

ITApFun

$$\frac{e_1 \mapsto e_1'}{e_1(e_2) \mapsto e_1'(e_2)} \tag{35b}$$

ITApArg

$$\frac{e_1 \text{ val} \qquad e_2 \mapsto e'_2}{e_1(e_2) \mapsto e_1(e'_2)} \tag{35c}$$

ITAP

$$\frac{e_2 \text{ val}}{(\lambda x: \tau.e_1)(e_2) \mapsto [e_2/x]e_1} \tag{35d}$$

ITPairL
$$\frac{e_1 \mapsto e_1'}{(e_1, e_2) \mapsto (e_1', e_2)}$$
(35e)

ITPairR
$$\frac{e_1 \text{ val} \qquad e_2 \mapsto e_2'}{(e_1, e_2) \mapsto (e_1, e_2')} \tag{35f}$$

$$\frac{(e_1, e_2) \text{ final}}{\text{prl}((e_1, e_2)) \mapsto e_1} \tag{35g}$$

$$\frac{(e_1, e_2) \text{ final}}{\text{prr}((e_1, e_2)) \mapsto e_2} \tag{35h}$$

ITInl

$$\frac{e \mapsto e'}{\operatorname{inl}_{\tau}(e) \mapsto \operatorname{inl}_{\tau}(e')} \tag{35i}$$

$$\frac{e \mapsto e'}{\operatorname{inr}_{\tau}(e) \mapsto \operatorname{inr}_{\tau}(e')} \tag{35j}$$

ITExpMatch

$$\frac{e \mapsto e'}{\operatorname{match}(e)\{\hat{rs}\} \mapsto \operatorname{match}(e')\{\hat{rs}\}}$$
(35k)

$$\begin{split} & \underset{\mathsf{match}(e)\{rs_{pre} \mid (p_r \Rightarrow e_r) \mid rs_{post}\} \mapsto [\theta](e_r)}{e \; \mathsf{final}} & e \rhd p_r \dashv \theta \\ & \underbrace{\mathsf{match}(e)\{rs_{pre} \mid (p_r \Rightarrow e_r) \mid rs_{post}\} \mapsto [\theta](e_r)} \end{split} \tag{35l}$$

ITFailMatch

$$\frac{e \; \text{final} \qquad e \perp p_r}{ \text{match}(e) \{ rs \mid (p_r \Rightarrow e_r) \mid (r' \mid rs') \} \mapsto \text{match}(e) \{ (rs \mid (p_r \Rightarrow e_r) \mid \cdot)^{\diamond} \mid r' \mid rs' \} }{(35\text{m})}$$

Lemma 4.0.3. If $\operatorname{inl}_{\tau_2}(e_1)$ final $\operatorname{then} e_1$ final.

Proof. By rule induction over Rules (26) on $\operatorname{inl}_{\tau_2}(e_1)$ final.

Case (26a).

(17)
$$inl_{\tau_2}(e_1) val$$

by assumption

By rule induction over Rules (24) on (17), only one case applies.

Case (24d).

(18)
$$e_1$$
 val by assumption
(19) e_1 final by Rule (26a) on (18)

Case (26b).

(17)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 indet by assumption

By rule induction over Rules (25) on (17), only one case applies.

Case (25i).

- (18) e_1 indet
- (19) e_1 final

by assumption

by Rule (26b) on (18)

Lemma 4.0.4. If $\operatorname{inr}_{\tau_1}(e_2)$ final $\operatorname{then} e_2$ final.

Proof. By rule induction over Rules (26) on $\operatorname{inr}_{\tau_1}(e_2)$ final.

Case (26a).

(1) $\operatorname{inr}_{\tau_1}(e_2)$ val

by assumption

By rule induction over Rules (24) on (1), only one case applies.

Case (24d).

(2) e_2 val

by assumption

(3) e_2 final

by Rule (26a) on (2)

Case (26b).

 $(1) \ \operatorname{inr}_{\tau_1}(e_2) \ \operatorname{indet}$

by assumption

By rule induction over Rules (25) on (1), only one case applies.

Case (25i).

(2) e_2 indet

by assumption

(3) e_2 final

by Rule (26b) on (2)

Lemma 4.0.5. If (e_1,e_2) final $then\ e_1$ final $and\ e_2$ final.

Proof. By rule induction over Rules (26) on (e_1, e_2) final.

Case (26a).

 $(1) \ (e_1,e_2) \ {\tt val}$

by assumption

By rule induction over Rules (24) on (1), only one case applies.

Case (24c).

(2) e_1 val

by assumption

(3) e_2 val

by assumption

(4) e_1 final

by Rule (26a) on (2)

(5) e_2 final

by Rule (26a) on (3)

Case (26b).	
(1) (e_1,e_2) indet	by assumption
By rule induction over Rules (25) on ((1), only three cases apply.
Case (25d).	
(2) e_1 indet	by assumption
(3) e_2 val	by assumption
(4) e_1 final	by Rule (26b) on (2)
(5) e_1 final	by Rule (26a) on (3)
Case (25e).	
(2) e_1 val	by assumption
(3) e_2 indet	by assumption
(4) e_1 final	by Rule $(26a)$ on (2)
(5) e_1 final	by Rule $(26b)$ on (3)
Case (25f).	
(2) e_1 indet	by assumption
(3) e_2 indet	by assumption
(4) e_1 final	by Rule $(26b)$ on (2)
(5) e_1 final	by Rule $(26b)$ on (3)
Lemma 4.0.6. There doesn't exist \underline{n} such the	$hat \; \underline{n} \; ext{notintro}.$
<i>Proof.</i> By rule induction over Rules (27) on a syntactic contradiction.	\underline{n} notintro, no case applies due to
Lemma 4.0.7. There doesn't exist $\operatorname{inl}_{\tau}(e)$	$such \ that \ {\tt inl}_{\tau}(e) \ {\tt notintro}.$
${\it Proof.}$ By rule induction over Rules (27) on due to syntactic contradiction.	$\operatorname{inl}_{\tau}(e)$ notintro, no case applies \Box
Lemma 4.0.8. There doesn't exist $\operatorname{inr}_{\tau}(e)$	$such\ that\ {\tt inr}_{\tau}(e)\ {\tt notintro}.$
${\it Proof.}$ By rule induction over Rules (27) on due to syntactic contradiction.	$\operatorname{inr}_{\tau}(e)$ notintro, no case applies \Box
Lemma 4.0.9. There doesn't exist (e_1, e_2) s	such that (e_1,e_2) notintro.
<i>Proof.</i> By rule induction over Rules (27) on	(e_1, e_2) notintro, no case applies

Lemma 4.0.10. If e final and e notintro then e indet.

due to syntactic contradiction.

Proof Sketch. By rule induction over Rules (27) on e notintro, for each case, by rule induction over Rules (24) on e val and we notice that e val is not derivable. By rule induction over Rules (26) on e final, Rule (26a) result in a contradiction with the fact that e val is not derivable while Rule (26b) tells us e indet.

Lemma 4.0.11. There doesn't exist such an expression e such that both e val and e indet.

Lemma 4.0.12. There doesn't exist such an expression e such that both e val and e notintro.

Lemma 4.0.13 (Finality). There doesn't exist such an expression e such that both e final and $e \mapsto e'$ for some e'

Proof. Assume there exists such an e such that both e final and $e \mapsto e'$ for some e' then proof by contradiction.

By rule induction over Rules (26) and Rules (35), *i.e.*, over Rules (24) and Rules (35) and over Rules (25) and Rules (35) respectively. The proof can be done by straightforward observation of syntactic contradictions. \Box

Lemma 4.0.14 (Matching Determinism). If e final $and \cdot ; \Delta_e \vdash e : \tau$ and $p : \tau[\xi] \dashv \Gamma ; \Delta$ then exactly one of the following holds

- 1. $e > p \dashv \theta$ for some θ
- 2. e?p
- 3. $e \perp p$

Proof.

- (1) *e* final by assumption
- (2) $\cdot : \Delta_e \vdash e : \tau$ by assumption
- (3) $p:\tau[\xi]\dashv \Gamma;\Delta$ by assumption

By rule induction over Rules (21) on (3), we would show one conclusion is derivable while the other two are not.

Case (21a).

- (4) p = x by assumption
- (5) $e \triangleright x \dashv e/x$ by Rule (32a)

Assume e? x. By rule induction over Rules (33) on it, only one case applies.

Case (33c).

(6) x refutable?

by assumption

By rule induction over Rules (31) on (6), no case applies due to syntactic contradiction.

(7) e^{2x} by contradiction

Assume $e \perp x$. By rule induction over Rules (34) on it, no case applies due to syntactic contradiction.

(8) $e \pm x$ by contradiction

Case (21b).

(4) p =_ by assumption

(5) $e \rhd \dashv \cdot$ by Rule (32b)

Assume e? _. By rule induction over Rules (33) on it, only one case applies.

Case (33c).

(6) refutable? by assumption

By rule induction over Rules (31) on (6), no case applies due to syntactic contradiction.

(7) e^{2} by contradiction

Assume $e \perp$ _. By rule induction over Rules (34) on it, no case applies due to syntactic contradiction.

(8) e by contradiction

Case (21c).

(4) $p = \emptyset^w$ by assumption

(5) e? \emptyset^w by Rule (33a)

Assume $e \rhd \bigoplus^w \dashv \theta$ for some θ . By rule induction over Rules (33) on it, no case applies due to syntactic contradiction.

(6) $e \rightarrow \theta = \theta$ by contradiction

Assume $e \perp \emptyset^w$. By rule induction over Rules (34) on it, no case applies due to syntactic contradiction.

(7) e by contradiction

Case (21d).

(4)
$$p = (p_0)^w$$
 by assumption
(5) $e ? (p_0)^w$ by Rule (33b)

Assume $e \rhd (p_0)^w \dashv \theta$ for some θ . By rule induction over Rules (33) on it, no case applies due to syntactic contradiction.

(6)
$$e \triangleright p_0 \stackrel{w}{=} \overline{\theta}$$
 by contradiction

Assume $e \perp (p_0)^w$. By rule induction over Rules (34) on it, no case applies due to syntactic contradiction.

(7)
$$e \perp p_0 \uparrow^{\sigma}$$
 by contradiction

Case (21e).

(4)
$$p = \underline{n_2}$$
 by assumption
(5) $\tau = \text{num}$ by assumption
(6) $\xi = \underline{n_2}$ by assumption
(7) n_2 refutable? by Rule (31a)

By rule induction over Rules (20) on (2), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

$$(8) \ e = ()^{u}, (e_{0})^{u}, e_{1}(e_{2}), \operatorname{prl}(e_{0}), \operatorname{prr}(e_{0}), \operatorname{match}(e_{0}) \{\hat{rs}\}$$
 by assumption
$$(9) \ e \ \operatorname{notintro} \qquad \qquad \operatorname{by \ Rule} \qquad \qquad (27a), (27b), (27c), (27d), (27e), (27f)$$

$$(10) \ e ? \underline{n_{2}} \qquad \qquad \operatorname{by \ Rule} \ (17b) \ \operatorname{on} \ (7)$$
 and
$$(9)$$

Assume $e
ightharpoonup \underline{n_2} \dashv \theta$ for some θ . By rule induction over it, no case applies due to syntactic contradiction.

(11)
$$e \triangleright n_2 + \theta$$
 by contradiction

Assume $e \perp \underline{n_2}$. By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \perp n_2$$
 by contradiction

Case (20d).

(8)
$$e = n_1$$

Assume $\underline{n_1}$? $\underline{n_2}$. By rule induction over Rules (33) on it, only two cases apply.

Case (33c).

(9)
$$\underline{n_1}$$
 notintro by assumption Contradicts Lemma 4.0.6.

(10)
$$n_1 2 n_2$$
 by contradiction

By case analysis on whether $n_1 = n_2$.

Case $n_1 = n_2$.

(11)
$$n_1 = n_2$$
 by assumption

(12)
$$\underline{n_1} \rhd \underline{n_2} \dashv \cdots$$
 by Rule (32c)

Assume $\underline{n_1} \perp \underline{n_2}$. By rule induction over Rules (34) on it, only one case applies.

Case (34a).

(13)
$$n_1 \neq n_2$$
 by assumption Contradicts (11).

(14)
$$\underline{n_1} + \underline{n_2}$$
 by contradiction

Case $n_1 \neq n_2$.

(11)
$$n_1 \neq n_2$$
 by assumption

(12)
$$n_1 \perp n_2$$
 by Rule (34a) on (11)

Assume $\underline{n_1} \rhd \underline{n_2} \dashv \theta$ for some θ . By rule induction over Rules (32) on it, no case applies due to syntactic contradiction.

(13)
$$\underline{n_1} \triangleright \underline{n_2} \dashv \theta$$
 by contradiction

Case (21f).

$$\begin{array}{ll} (4) \ \ p = \mathtt{inl}(p_1) & \text{by assumption} \\ (5) \ \ \tau = (\tau_1 + \tau_2) & \text{by assumption} \\ (6) \ \ \xi = \mathtt{inl}(\xi_1) & \text{by assumption} \\ (7) \ \ p_1 : \tau_1[\xi_1] \dashv \Gamma \,; \, \Delta & \text{by assumption} \\ (8) \ \ \mathtt{inl}(p_1) \ \ \mathtt{refutable}_? & \text{by Rule (31d)} \\ \end{array}$$

By rule induction over Rules (20) on (2), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

(9)
$$e = (v, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$
 by assumption

(10) e notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

(11)
$$e$$
? $inl(p_1)$ by Rule (17b) on (8) and (10)

Assume $e \rhd \operatorname{inl}(p_1) \dashv \theta_1$ for some θ_1 . By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \triangleright \operatorname{inl}(p_1) \dashv \theta_1$$

by contradiction

Assume $e \perp inl(p_1)$. By rule induction over it, no case applies due to syntactic contradiction.

(13)
$$e \perp \operatorname{inl}(p_1)$$

by contradiction

by assumption

by assumption

Case (20j).

- (9) $e = inl_{\tau_2}(e_1)$ (10) \cdot ; $\Delta_e \vdash e_1 : \tau_1$
- (11) e_1 final by Lemma 4.0.3 on (1)

By inductive hypothesis on (10) and (11) and (7), exactly one of $e_1 \triangleright p_1 \dashv \theta_1$ for some $\theta_1, e_1 ? p_1$, and $e_1 \perp p_1$ holds. By case analysis on which one holds.

Case $e_1 \triangleright p_1 \dashv \theta_1$.

- (12) $e_1 \triangleright p_1 \dashv \theta_1$ by assumption
- (13) $e_1 ? p_1$ by assumption by assumption
- $(14) e_1 + p_1$ (15) $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta_1$ by Rule (32e) on (12)

Assume $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$. By rule induction over Rules (33) on it, only two cases apply.

Case (33c).

(16) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

Contradicts Lemma 4.0.7.

Case (33g).

 $(16) e_1? p_1$ by assumption

Contradicts (13).

(17) $\operatorname{inl}_{\tau_2}(e_1) ? \operatorname{inl}(p_1)$ by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1)$. By rule induction over Rules (34) on it, only one case applies.

Case (34f).

- (18) $e_1 \perp p_1$ by assumption Contradicts (14).
- (19) $\operatorname{inl}_{\tau_2}(e_1) \pm \operatorname{inl}(p_1)$ by contradiction

Case $e_1 ? p_1$.

- (12) $e_1 \triangleright p_1 + \theta_1$ by assumption (13) $e_1 ? p_1$ by assumption $(14) e_1 + p_1$ by assumption
- (15) $inl_{\tau_2}(e_1)$? $inl(p_1)$ by Rule (33g) on (13)

Assume $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (32) on it, only one case applies.

Case (32e).

(16)
$$e_1 \triangleright p_1 \dashv \theta$$
 by assumption Contradicts (12).

(17)
$$\operatorname{inl}_{\tau_2}(e_1) \triangleright \operatorname{inl}(p_1) \dashv \theta$$
 by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1)$. By rule induction over Rules (34) on it, only one case applies.

Case (34f).

(18)
$$e_1 \perp p_1$$
 by assumption Contradicts (14).

(19)
$$\operatorname{inl}_{T_2}(e_1) \pm \operatorname{inl}(p_1)$$
 by contradiction

Case $e_1 \perp p_1$.

$$\begin{array}{lll} (12) & \underline{e_1 \triangleright p_1 \dashv \theta_1} & \text{by assumption} \\ (13) & \underline{e_1 \cdot p_1} & \text{by assumption} \\ (14) & \underline{e_1 \perp p_1} & \text{by assumption} \\ (15) & \operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1) & \text{by Rule (34f) on (14)} \\ \end{array}$$

Assume $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (32) on it, only one case applies.

Case (32e).

(16)
$$e_1 \triangleright p_1 \dashv \theta$$
 by assumption Contradicts (12).

(17)
$$\operatorname{inl}_{\tau_2}(e_1) \triangleright \operatorname{inl}(p_1) \dashv \theta$$
 by contradiction Assume $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$. By rule induction over Rules (33) on it, only two cases apply.

Case (33c).

(18)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro by assumption Contradicts Lemma 4.0.7.

Case (33g).

(18)
$$e_1$$
? p_1 by assumption Contradicts (13).

(19)
$$\underline{\operatorname{inl}_{\tau_2}(e_1) ? \operatorname{inl}(p_1)}$$
 by contradiction

Case (21g).

(4)
$$p = inr(p_2)$$
 by assumption
(5) $\tau = (\tau_1 + \tau_2)$ by assumption

(6) $\xi = \operatorname{inr}(\xi_2)$ by assumption (7) $p_2 : \tau_2[\xi_2] \dashv \Gamma; \Delta$ by assumption (8) $\operatorname{inr}(p_2)$ refutable? by Rule (31e)

By rule induction over Rules (20) on (2), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20m).

(9)
$$e = \{ \| u, \| e_0 \|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \} \{ \hat{rs} \}$$

by assumption

(10) e notintro by Rule

(27a),(27b),(27c),(27d),(27e),(27f)

(11) e? $inr(p_2)$ by Rule (17b) on (8) and (10)

Assume $e \triangleright \operatorname{inr}(p_2) \dashv \theta_2$ for some θ_2 . By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \triangleright inr(p_2) \dashv \theta_2$$
 by contradiction

Assume $e \perp inr(p_2)$. By rule induction over it, no case applies due to syntactic contradiction.

(13)
$$e \perp inr(p_2)$$
 by contradiction

Case (20k).

 $\begin{array}{ll} (9) & e = \operatorname{inr}_{\tau_1}(e_2) & \text{by assumption} \\ (10) & \cdot \; ; \Delta_e \vdash e_2 : \tau_2 & \text{by assumption} \\ (11) & e_2 \; \text{final} & \text{by Lemma 4.0.4 on (1)} \end{array}$

By inductive hypothesis on (10) and (11) and (7), exactly one of $e_2 \triangleright p_2 \dashv \theta_2$ for some θ_2 , $e_2 ? p_2$, and $e_2 \perp p_2$ holds.

By case analysis on which one holds.

Case $e_2 \triangleright p_2 \dashv \theta_2$.

(12) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption (13) $e_2 \not p_2$ by assumption (14) $e_2 \not p_2$ by assumption

(15) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta_2$ by Rule (32f) on (12)

Assume $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$. By rule induction over Rules (33) on it, only two cases apply.

Case (33c).

(16) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption Contradicts Lemma 4.0.8.

Case (33h).

(16) $e_2 ? p_2$ by assumption

Contradicts (13).

(17)
$$inr_{\tau_1}(e_2) ? inr(p_2)$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \perp \operatorname{inr}(p_2)$. By rule induction over Rules (34) on it, only one case applies.

Case (34g).

(18)
$$e_2 \perp p_2$$

by assumption

Contradicts (14).

(19)
$$\operatorname{inr}_{\tau_1}(e_2) \pm \operatorname{inr}(p_2)$$

by contradiction

Case $e_2 ? p_2$.

(12)
$$\underline{e_2} \triangleright \underline{p_2} + \underline{\theta}$$

by assumption

$$(13)$$
 $e_2 ? p_2$

by assumption

$$(14) \ e_2 \perp p_2$$

by assumption

(15)
$$inr_{\tau_1}(e_2)$$
? $inr(p_2)$

by Rule (33h) on (13)

Assume $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ . By rule induction over Rules (32) on it, only one case applies.

Case (32f).

$$(16) e_2 \rhd p_2 \dashv \theta$$

by assumption

Contradicts (12).

(17)
$$\operatorname{inr}_{\tau_1}(e_2) \supset \operatorname{inr}(p_2) \dashv \theta$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \perp \operatorname{inr}(p_2)$. By rule induction over Rules (34) on it, only one case applies.

Case (34g).

(18)
$$e_2 \perp p_2$$

by assumption

Contradicts (14).

(19)
$$\underline{\operatorname{inr}_{\tau_1}(e_2) \pm \operatorname{inr}(p_2)}$$

by contradiction

Case $e_2 \perp p_2$.

(12) $\underline{e_2} \triangleright \underline{p_2} + \theta$

by assumption

(13)
$$e_2 ? p_2$$

by assumption

(14)
$$e_2 \perp p_2$$

by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \perp \operatorname{inr}(p_2)$$

by Rule (34g) on (14)

Assume $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ . By rule induction over Rules (32) on it, only one case applies.

Case (32f).

(16)
$$e_2 \triangleright p_2 \dashv \theta$$

by assumption

Contradicts (12).

(17)
$$\operatorname{inr}_{\tau_1}(e_2) \supset \operatorname{inr}(p_2) \dashv \theta$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$. By rule induction over Rules (33) on it, only two cases apply.

Case (33c).

(18) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption Contradicts Lemma 4.0.8.

Case (33h).

(18) e_2 ? p_2 by assumption Contradicts (13).

(19) $\operatorname{inr}_{\mathcal{I}_1}(e_2)$? $\operatorname{inr}(p_2)$ by contradiction

Case (21h).

$$(4) \ \ p = (p_1, p_2) \qquad \qquad \text{by assumption}$$

$$(5) \ \ \tau = (\tau_1 \times \tau_2) \qquad \qquad \text{by assumption}$$

$$(6) \ \ \xi = (\xi_1, \xi_2) \qquad \qquad \text{by assumption}$$

$$(7) \ \ \Gamma = \Gamma_1 \uplus \Gamma_2 \qquad \qquad \text{by assumption}$$

$$(8) \ \ \Delta = \Delta_1 \uplus \Delta_2 \qquad \qquad \text{by assumption}$$

$$(9) \ \ p_1 : \tau_1[\xi_1] \dashv \Gamma_1 ; \Delta_1 \qquad \qquad \text{by assumption}$$

$$(10) \ \ p_2 : \tau_2[\xi_2] \dashv \Gamma_2 ; \Delta_2 \qquad \qquad \text{by assumption}$$

By rule induction over Rules (20) on (2), the following cases apply.

(11) $e = \{ \|u, \|e_0\|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \} \}$

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

$$\begin{array}{c} \text{by assumption} \\ \text{(12)} \ e \ \text{notintro} \\ \text{(27a),(27b),(27c),(27d),(27e),(27f)} \\ \text{(13)} \ e \ \text{indet} \\ \text{by Lemma 4.0.10 on} \\ \text{(1) and (12)} \\ \end{array}$$

Assume $e \perp (p_1, p_2)$. By rule induction on it, no case applies due to syntactic contradiction.

(20) $e \perp (p_1, p_2)$ by contradiction

By inductive hypothesis on (15) and (18) and (9), exactly one of $\mathtt{prl}(e) \rhd p_1 \dashv\! \theta_1$, $\mathtt{prl}(e) ? p_1$, and $\mathtt{prl}(e) \perp p_1$ holds.

By inductive hypothesis on (17) and (19) and (10), exactly one of $prr(e) \triangleright p_2 \dashv \theta_2$, prr(e)? p_2 , and $prr(e) \perp p_2$ holds.

By case analysis on which conclusion holds for p_1 and p_2 . Note that we have already shown $e \perp (p_1, p_2)$.

Case $prl(e) \triangleright p_1 \dashv \theta_1, prr(e) \triangleright p_2 \dashv \theta_2$.

(21) $\operatorname{prl}(e) \rhd p_1 \dashv \theta_1$	by assumption
(22) $prl(e) ? p_1$	by assumption
(23) $\underline{\operatorname{prl}(e) \perp p_1}$	by assumption
$(24) \ \operatorname{prr}(e) \rhd p_2 \dashv \mid \theta_2$	by assumption
$(25) \ \underline{\operatorname{prr}(e)?p_2}$	by assumption
(26) $\operatorname{prr}(e) \perp p_2$	by assumption
$(27) e \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2$	by Rule (32g) on (12

(27) $e \triangleright (p_1, p_2) \dashv \theta_1 \uplus \theta_2$ by Rule (32g) on (12) and (21) and (24) Assume $e?(p_1, p_2)$. By rule induction over Rules (33) on it, only

one case applies. Case (33c).

(28) (p_1, p_2) refutable? by assumption By rule induction over Rules (31), only two cases apply.

Case (31f).

(29) p_1 refutable?	by assumption
(30) $\mathtt{prl}(e)$ notintro	by Rule (27e)
(31) $prl(e) ? p_1$	by Rule $(33c)$ on (29)
	and (30)

Contradicts (22).

Case (31g).

(29) p_2 refutable?	by assumption
(30) $\mathtt{prr}(e)$ notintro	by Rule (27f)
(31) $prl(e) ? p_1$	by Rule (33c) on (29)
	and (30)

Contradicts (22).

(32)
$$e?(p_1, p_2)$$
 by contradiction

Case $prl(e) \triangleright p_1 \dashv \theta_1, prr(e) ? p_2.$

$(21) \ \mathtt{prl}(e) \rhd p_1 \dashv\! \mid \theta_1$	by assumption
(22) $prl(e) ? p_1$	by assumption
(23) $\operatorname{prl}(e) \pm p_1$	by assumption
(24) $prr(e) \Rightarrow p_2 \dashv \overline{\theta_2}$	by assumption
$(25) \ prr(e) ? p_2$	by assumption
(26) $prr(e) \perp p_2$	by assumption

Assume $e \rhd (p_1, p_2) \dashv \theta$. By rule induction over Rules (32), only one case applies.

Case (32g).

- (27) $\theta = \theta_1 \uplus \theta_2$ by assumption (28) $prr(e) \rhd p_2 \dashv \theta_2$ by assumption
- Contradicts (24).

(29)
$$e \triangleright (p_1, p_2) \dashv \theta$$
 by contradiction

By rule induction over Rules (33) on (25), the following cases apply.

Case (33a),(33b).

- (30) $p_2 = \langle | \rangle^w, \langle | p \rangle^w$ by assumption
- (31) p_2 refutable? by Rule (31b) and Rule (31c)
- (32) (p_1, p_2) refutable? by Rule (31g) on (31)
- (33) e? (p_1, p_2) by Rule (33c) on (12) and (32)

Case (33c).

- (30) p_2 refutable? by assumption
- (31) (p_1, p_2) refutable? by Rule (31g) on (30)
- (32) $e ? (p_1, p_2)$ by Rule (33c) on (12) and (31)

Case $prl(e) \triangleright p_1 \dashv \theta_1, prr(e) \perp p_2$.

- (21) $\operatorname{prl}(e) \rhd p_1 \dashv \theta_1$ by assumption
- (22) prl(e)? p_1 by assumption
- (23) $\underline{\operatorname{prl}(e) \perp p_1}$ by assumption
- (24) $prr(e) \Rightarrow p_2 \dashv \theta_2$ by assumption
- (25) prr(e)? p_2 by assumption (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (34) on (26), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case prl(e) ? $p_1, prr(e) \triangleright p_2 \dashv \theta_2$.

(21) $\underline{\text{prl}(e)} \Rightarrow p_1 \dashv \theta_1$ by assumption (22) $\underline{\text{prl}(e)} ? p_1$ by assumption (23) $\underline{\text{prl}(e)} \perp p_1$ by assumption (24) $\underline{\text{prr}(e)} \Rightarrow p_2 \dashv \theta_2$ by assumption (25) $\underline{\text{prr}(e)} ? p_2$ by assumption (26) $\underline{\text{prr}(e)} \perp p_2$ by assumption

Assume $e \rhd (p_1, p_2) \dashv \theta$. By rule induction over Rules (32), only one case applies.

Case (32g).

(27) $\theta = \theta_1 \uplus \theta_2$

by assumption

(28) $prl(e) \triangleright p_1 \dashv \theta_1$

by assumption

Contradicts (21).

(29) $e \triangleright (p_1, p_2) \dashv \overline{\theta}$

by contradiction

By rule induction over Rules (33) on (22), the following cases apply.

Case (33a),(33b).

(30) $p_1 = (||w|, ||p||^w)$

by assumption

(31) p_1 refutable?

by Rule (31b) and Rule

(31c)

(32) (p_1, p_2) refutable?

by Rule (31g) on (31)

(33) $e?(p_1,p_2)$

by Rule (33c) on (12)

and (32)

Case (33c).

(30) p_1 refutable?

by assumption

(31) (p_1, p_2) refutable?

by Rule (31g) on (30)

(32) e? (p_1, p_2)

by Rule (33c) on (12)

and (31)

Case prl(e) ? p_1 , prr(e) ? p_2 .

(21) $\operatorname{prl}(e) \Rightarrow p_1 \dashv \theta_1$

by assumption

(22) $prl(e) ? p_1$

by assumption

(23) $prl(e) \pm p_1$ (24) $prr(e) \rightarrow p_2 \dashv \theta_2$ by assumption

 $(25) prr(e) ? p_2$

by assumption

by assumption

(26) $prr(e) \pm p_2$

by assumption

Assume $e > (p_1, p_2) \dashv \theta$. By rule induction over Rules (32), only one case applies.

Case (32g).

(27) $\theta = \theta_1 \uplus \theta_2$

by assumption

(28) $prl(e) \triangleright p_1 \dashv \theta_1$

by assumption

Contradicts (21).

(29) $e \triangleright (p_1, p_2) \dashv \overline{\theta}$

by contradiction

By rule induction over Rules (33) on (22), the following cases apply.

Case (33a),(33b).

(30) $p_1 = (1)^w, (1p)^w$

by assumption

(31) p_1 refutable?

by Rule (31b) and Rule

(31c)

- (32) (p_1, p_2) refutable? by Rule (31g) on (31)
- (33) e? (p_1, p_2) by Rule (33c) on (12) and (32)

Case (33c).

- (30) p_1 refutable? by assumption
- (31) (p_1, p_2) refutable? by Rule (31g) on (30)
- (32) $e ? (p_1, p_2)$ by Rule (33c) on (12) and (31)

Case prl(e)? $p_1, prr(e) \perp p_2$.

- (21) $\underline{\mathsf{prl}(e)} \triangleright p_1 \dashv \theta_1$ by assumption
- (22) $\operatorname{prl}(e)$? p_1 by assumption
- (23) $prl(e) \perp p_1$ by assumption
- (24) $\underline{\operatorname{prr}(e)} \triangleright p_2 \dashv \theta_2$ by assumption
- (25) $\underline{prr(e)? p_2}$ by assumption
- (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (34) on (26), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) \rhd p_2 \dashv \theta_2$.

- (21) $prl(e) \rightarrow p_1 \dashv \theta_1$ by assumption
- (22) prl(e)? p_1 by assumption
- (23) $prl(e) \perp p_1$ by assumption
- (24) $prr(e) > p_2 \dashv \theta_2$ by assumption (25) $prr(e) ? p_2$ by assumption
- (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (34) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) ? p_2$.

- (21) $\operatorname{pr}_{\underline{1}}(e) \rightarrow p_{\underline{1}} \dashv \overline{\theta_{1}}$ by assumption
- (22) $prl(e) ? p_1$ by assumption
- (23) $prl(e) \perp p_1$ by assumption
- (24) $prr(e) \rightarrow p_2 + \theta_2$ by assumption
- (25) prr(e)? p_2 by assumption
- (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (34) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) \perp p_2$.

(21) $prl(e) \rightarrow p_1 \dashv \theta_1$ by assumption

(22) $\underline{\operatorname{prl}(e)}$? $\overline{p_1}$	by assumption
(23) $\operatorname{prl}(e) \perp p_1$	by assumption
(24) $\underline{\operatorname{prr}(e)} \Rightarrow p_2 \dashv \overline{\theta_2}$	by assumption
$(25) \ prr(e) ? p_2$	by assumption
(26) $\operatorname{prr}(e) \perp p_2$	by assumption

By rule induction over Rules (34) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case (20g).

$(11) \ e = (e_1, e_2)$	by assumption
$(12) \cdot ; \Delta \vdash e_1 : \tau_1$	by assumption
$(13) \cdot ; \Delta \vdash e_2 : \tau_2$	by assumption
(14) e_1 final	by Lemma $4.0.5$ on (1)
(15) e_2 final	by Lemma $4.0.5$ on (1)

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 > p_1 \dashv \theta_1$, $e_1 ? p_1$, and $e_1 \perp p_1$ holds.

By inductive hypothesis on (10) and (13) and (15), exactly one of $e_2
ightharpoonup p_2 \dashv \theta_2, \ e_2 ? \ p_2, \ \text{and} \ e_2 \perp p_2 \ \text{holds}.$

By case analysis on which conclusion holds for p_1 and p_2 .

Case $e_1 \rhd p_1 \dashv \mid \theta_1, e_2 \rhd p_2 \dashv \mid \theta_2$.

$(16) e_1 \rhd p_1 \dashv \mid \theta_1$	by assumption
$(17) e_1 ? p_1$	by assumption
(18) $e_1 + p_1$	by assumption
$(19) \ e_2 \rhd p_2 \dashv \theta_2$	by assumption
$(20) e_2 ? p_2$	by assumption
$(21) \underline{e_2 + p_2}$	by assumption
$(22) (e_1, e_2) \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2$	by Rule (32d) on (16)
	and (19)

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (33) on it, only four cases apply.

Case (33c).

(23) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (33d).

(23) $e_1 ? p_1$ by assumption

Contradicts (17).

Case (33e).

(23) e_2 ? p_2 by assumption

Contradicts (20).

Case (33f).

Contradicts (17).

(24)
$$(e_1, e_2) ? (p_1, p_2)$$
 by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (34) on it, only two cases apply.

Case (34b).

(25) $e_1 \perp p_1$ by assumption

Contradicts (18).

Case (34c).

(25) $e_2 \perp p_2$ by assumption

Contradicts (21).

(26) $(e_1, e_2) \perp (p_1, p_2)$ by contradiction

Case $e_1 \triangleright p_1 \dashv \theta_1, e_2 ? p_2$.

(16) $e_1 \triangleright p_1 \dashv \theta_1$ by assumption

(17) $e_1 + p_1$ by assumption

(18) $e_1 \perp p_1$ by assumption

(19) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption

(20) $e_2 ? p_2$ by assumption

(21) $e_2 \perp p_2$ by assumption

(22) $(e_1, e_2) \vdash (p_1, p_2)$ by Rule (33e) on (16) and (20)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies.

Case (32d).

(23) $\theta = \theta_1 \uplus \theta_2$

(24) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption

Contradicts (19).

(25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2) \dashv \theta$ by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2) \dashv \theta$ by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2) \dashv \theta$ by assumption

Contradicts (19).

(26) $e_1 \perp p_1$ by assumption

Contradicts (18).

Case (34c).

(26) $e_1 \perp p_1$ by assumption

Contradicts (18).

Case (34c).

Contradicts (21).

$(27) \ \underline{(e_1,e_2) \pm (p_1,p_2)}$	by contradiction
Case $e_1 \rhd p_1 \dashv \theta_1, e_2 \perp p_2$.	
$(16) e_1 \rhd p_1 \dashv \theta_1$	by assumption
$(17) e_1 ? p_1$	by assumption
$(18) \underbrace{e_1 + p_1}$	by assumption
$(19) e_2 \triangleright p_2 + \theta_2$	by assumption
$(20) e_2 ? p_2$	by assumption
$(21) e_2 \perp p_2$	by assumption
$(22) (e_1, e_2) \perp (p_1, p_2)$	by Rule (34c) on (21)
, , , , , , , , , , , , , , , , , , , ,	By rule induction over Rules (32)
on it, only one case applies.	•
Case (32d).	
$(23) \ \theta = \theta_1 \uplus \theta_2$	
$(24) \ e_2 \rhd p_2 \dashv \mid \theta_2$	by assumption
Contradicts (19).	
$(25) \ \underline{(e_1,e_2) \triangleright (p_1,p_2) \dashv \theta}$	by contradiction
	rule induction over Rules (33) on
it, only four cases apply.	
Case (33c).	
(26) (e_1,e_2) notintro	by assumption
Contradicts Lemma 4.0.9.	
Case $(33d)$.	
$(26) e_1? p_1$	by assumption
Contradicts (17) .	
Case $(33e)$.	
$(26) e_2 ? p_2$	by assumption
Contradicts (20).	
Case $(33f)$.	
$(26) e_1 ? p_1$	by assumption
Contradicts (17).	
$(27) \ \underline{(e_1,e_2)?(p_1,p_2)}$	by contradiction
Case $e_1 ? p_1, e_2 > p_2 \dashv ! \theta_2$.	
(16) $\underline{e_1 \triangleright p_1 \dashv \theta_1}$	by assumption
(17) $e_1 ? p_1$	by assumption
(18) $e_1 \pm p_1$	by assumption
$(19) \ e_2 \rhd p_2 \dashv \vdash \theta_2$	by assumption
$(20) \ \underline{e_2 ? p_2}$	by assumption
(21) $e_2 \perp p_2$	by assumption

(22)
$$(e_1, e_2)$$
? (p_1, p_2) by Rule (33d) on (17) and (19)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies.

Case (32d).

(23)
$$\theta = \theta_1 \uplus \theta_2$$

(24)
$$e_1 > p_1 \dashv \theta_1$$
 by assumption

Contradicts (16).

(25)
$$(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$
 by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (34) on it, only two cases apply.

Case (34b).

(26)
$$e_1 \perp p_1$$
 by assumption

Contradicts (18).

Case (34c).

(26)
$$e_2 \perp p_2$$
 by assumption

Contradicts (21).

(27)
$$(e_1, e_2) \perp (p_1, p_2)$$
 by contradiction

Case $e_1 ? p_1, e_2 ? p_2$.

(16)
$$e_1 \triangleright p_1 \dashv \theta_1$$
 by assumption
(17) $e_1 ? p_1$ by assumption
(18) $e_1 \not p_1 \not p_1$ by assumption
(19) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption

(20)
$$e_2 ? p_2$$
 by assumption
(21) $e_2 + p_2$ by assumption

(22)
$$(e_1, e_2)$$
? (p_1, p_2) by Rule (33f) on (17) and (20)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies.

Case (32d).

$$(23) \ \theta = \theta_1 \uplus \theta_2$$

(24)
$$e_2 \triangleright p_2 \dashv \theta_2$$
 by assumption

Contradicts (19).

(25)
$$(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$
 by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (34) on it, only two cases apply.

Case (34b).

(26)
$$e_1 \perp p_1$$
 by assumption

Contradicts (18). Case (34c). (26) $e_2 \perp p_2$ by assumption Contradicts (21). (27) $(e_1,e_2) \pm (p_1,p_2)$ by contradiction Case $e_1 ? p_1, e_2 \perp p_2$. (16) $e_1 \triangleright p_1 + \theta_1$ by assumption (17) $e_1 ? p_1$ by assumption (18) $e_1 + p_1$ by assumption (19) $e_2 > p_2 + \theta_2$ by assumption $(20) e_2 ? p_2$ by assumption (21) $e_2 \perp p_2$ by assumption (22) $(e_1, e_2) \perp (p_1, p_2)$ by Rule (34c) on (21) Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies. Case (32d). (23) $\theta = \theta_1 \uplus \theta_2$ (24) $e_2 \triangleright p_2 \dashv \mid \theta_2$ by assumption Contradicts (19). $(25) (e_1, e_2) \supset (p_1, p_2) \dashv \theta$ by contradiction Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (33) on it, only four cases apply. **Case** (33c). (26) (e_1,e_2) notintro by assumption Contradicts Lemma 4.0.9. Case (33d). $(26) e_2 \rhd p_2 \dashv \theta_2$ by assumption Contradicts (19). Case (33e). (26) $e_2 ? p_2$ by assumption Contradicts (20). Case (33f). (26) $e_2 ? p_2$ by assumption Contradicts (20). $(27) (e_1,e_2)?(\overline{p_1,p_2})$ by contradiction Case $e_1 \perp p_1, e_2 \rhd p_2 \dashv \theta_2$.

by assumption

(16) $e_1 \triangleright p_1 + \theta_1$

- $\begin{array}{lll} (17) & \underbrace{e_1 \cdot p_1} & & \text{by assumption} \\ (18) & e_1 \perp p_1 & & \text{by assumption} \\ (19) & e_2 \rhd p_2 \dashv \theta_2 & & \text{by assumption} \\ (20) & \underbrace{e_2 \cdot p_2} & & \text{by assumption} \\ (21) & \underbrace{e_2 \perp p_2} & & \text{by assumption} \\ (22) & (e_1, e_2) \perp (p_1, p_2) & & \text{by Rule (34b) on (18)} \\ \end{array}$
- Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies.

Case (32d).

- (23) $\theta = \theta_1 \uplus \theta_2$
- (24) $e_1 \triangleright p_1 \dashv \theta_1$ by assumption Contradicts (16).
- (25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ by contradiction Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (33) on it, only four cases apply.

Case (33c).

(26) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.9.

Case (33d).

(26) $e_1 ? p_1$ by assumption

Contradicts (17).

Case (33e).

(26) $e_2 ? p_2$ by assumption

Contradicts (20).

Case (33f).

(26) $e_1 ? p_1$ by assumption

Contradicts (17).

(27) (e_1, e_2) ? (p_1, p_2) by contradiction

Case $e_1 \perp p_1, e_2 ? p_2$.

- $\begin{array}{lll} (16) & \underline{e_1} \triangleright p_1 \# \theta_1 & \text{by assumption} \\ (17) & \underline{e_1} \mathbin{?} p_1 & \text{by assumption} \\ (18) & \underline{e_1} \perp p_1 & \text{by assumption} \\ (19) & \underline{e_2} \triangleright p_2 \# \theta_2 & \text{by assumption} \\ (20) & \underline{e_2} \mathbin{?} p_2 & \text{by assumption} \\ (21) & \underline{e_2} \perp p_2 & \text{by assumption} \\ (22) & (\underline{e_1}, \underline{e_2}) \perp (p_1, p_2) & \text{by Rule (34b) on (18)} \\ \end{array}$
- Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32) on it, only one case applies.

Case (32d).

(23)
$$\theta = \theta_1 \uplus \theta_2$$

(24) $e_2 \triangleright p_2 \dashv \mid \theta_2$

by assumption

Contradicts (19).

$$(25) (e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$

by contradiction

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (33) on it, only four cases apply.

Case (33c).

(26) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.9.

Case (33d).

(26) $e_2 \rhd p_2 \dashv \mid \theta_2$

by assumption

Contradicts (19).

Case (33e).

(26) $e_1 \triangleright p_1 \dashv \theta_1$

by assumption

Contradicts (16).

Case (33f).

(26) $e_1 ? p_1$

by assumption

Contradicts (17).

$$(27) \ \underline{(e_1,e_2)?(p_1,p_2)}$$

by contradiction

Case $e_1 \perp p_1, e_2 \perp p_2$.

(16)
$$e_1 \triangleright p_1 \dashv \theta_1$$
 by assumption
(17) $e_1 \stackrel{?}{?} p_1$ by assumption
(18) $e_1 \perp p_1$ by assumption
(19) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption
(20) $e_2 \stackrel{?}{?} p_2$ by assumption
(21) $e_2 \perp p_2$ by assumption

(22) $(e_1, e_2) \perp (p_1, p_2)$

by Rule (34b) on (18) Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (32)

on it, only one case applies.

Case (32d).

(23)
$$\theta = \theta_1 \uplus \theta_2$$

$$(24)$$
 $e_2 \triangleright p_2 \dashv \mid \theta_2$

by assumption

Contradicts (19).

$$(25) \quad (e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$

by contradiction

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (33) on it, only four cases apply.

(26)
$$(e_1, e_2)$$
 notintro by assumption

Contradicts Lemma 4.0.9.

Case (33d).

(26)
$$e_2 > p_2 \dashv \theta_2$$
 by assumption

Contradicts (19).

Case (33e).

(26)
$$e_1 > p_1 \dashv \theta_1$$
 by assumption

Contradicts (16).

Case (33f).

(26)
$$e_1 ? p_1$$
 by assumption

Contradicts (17).

(27)
$$(e_1, e_2)$$
? (p_1, p_2) by contradiction

Lemma 4.0.15 (Matching Coherence of Constraint). Suppose that $\cdot; \Delta_e \vdash e : \tau$ and e final and $p : \tau[\xi] \dashv \Gamma; \Delta$. Then we have

1.
$$e \models \xi \text{ iff } e \triangleright p \dashv \theta$$

2.
$$e \models_? \xi \text{ iff } e ? p$$

3.
$$e \not\models_{?}^{\dagger} \xi \text{ iff } e \perp p$$

Proof.

(1)
$$\cdot$$
; $\Delta_e \vdash e : \tau$ by assumption

(2)
$$e$$
 final by assumption

(3)
$$p:\tau[\xi]\dashv \Gamma;\Delta$$
 by assumption

Given Lemma 3.0.1, Theorem 2.1, and Lemma 4.0.14, it is sufficient to prove

1.
$$e \models \xi \text{ iff } e \triangleright p \dashv \theta$$

2.
$$e \models_? \xi$$
 iff $e ? p$

By rule induction over Rules (21) on (3).

Case (21a).

(4)
$$p = x$$
 by assumption

(5)
$$\xi = \top$$
 by assumption

1. Prove $e \models \top$ implies $e \triangleright x \dashv \theta$ for some θ .

(6)
$$e \triangleright x \dashv e/x$$
 by Rule (32a)

- 2. Prove $e > x \dashv \theta$ implies $e \models \top$.
 - (6) $e \models \top$

by Rule (15a)

- 3. Prove $e \models_? \top$ implies e ? x.
 - (6) $e \not\models_? \top$

by Lemma 2.0.3

Vacuously true.

4. Prove e ? x implies $e \models_? \top$.

By rule induction over Rules (33), we notice that either, e?x is in syntactic contradiction with all the cases, or the premise x refutable? is not derivable. Hence, e?x are not derivable. And thus vacuously true.

Case (21b).

(4) $p = _{-}$

by assumption

(5) $\xi = \top$

by assumption

- 1. Prove $e \models \top$ implies $e \triangleright _ \dashv \theta$ for some θ .
 - (6) $e \rhd _ \dashv \cdot$

by Rule (32a)

- 2. Prove $e > \exists \theta \text{ implies } e \models \top$.
 - (6) $e \models \top$

by Rule (15a)

- 3. Prove $e \models_? \top$ implies e? .
 - (6) $e \not\models_? \top$

by Lemma 2.0.3

Vacuously true.

4. Prove e? _ implies $e \models_? \xi$.

By rule induction over Rules (33), we notice that either, e? _ is in syntactic contradiction with all the cases, or the premise _ refutable? is not derivable. Hence, e? _ are not derivable. And thus vacuously true

Case (21c).

(4) $p = ()^w$

by assumption

(5) $\xi = ?$

by assumption

(6) $\bar{\xi} = ?$

by Definition 10

- 1. Prove $e \models ?$ implies $e \rhd ()^w \dashv \theta$ for some θ .
 - $(7) e \not\models ?$

by Rule (32a)

Vacuously true.

- 2. Prove $e \rhd ()^w \dashv \theta$ implies $e \models ?$. By rule induction over Rules (32), we notice that $e \rhd ()^w \dashv \theta$ is in syntactic contradiction with all the cases, hence not derivable. And thus vacuously true.
- 3. Prove $e \models_?$? implies e? \emptyset^w .

(7)
$$e$$
? \mathbb{D}^w by Rule (33a)

4. Prove e? $()^w$ implies $e \models_?$?.

(7)
$$e \models_?$$
? by Rule (17a)

Case (21d).

(4)
$$p = (p_0)^w$$
 by assumption

(5)
$$\xi = ?$$
 by assumption

1. Prove $e \models ?$ implies $e \rhd (p_0)^w \dashv \theta$ for some θ .

(6)
$$e \not\models ?$$
 by Rule (32a)

Vacuously true.

thus vacuously true.

2. Prove $e \rhd (p_0)^w \dashv \theta$ implies $e \models ?$. By rule induction over Rules (32), we notice that $e \rhd (p_0)^w \dashv \theta$ is in syntactic contradiction with all the cases, hence not derivable. And

3. Prove $e \models_?$? implies e? $(p_0)^w$.

(6)
$$e ? (p_0)^w$$
 by Rule (33b)

4. Prove $e ? (p_0)^w$ implies $e \models_? ?$.

(6)
$$e \models_?$$
? by Rule (17a)

Case (21e).

(4)
$$p = \underline{n}$$
 by assumption

(5)
$$\xi = \underline{n}$$
 by assumption

1. Prove $e \models \underline{n}$ implies $e \triangleright \underline{n} \dashv \theta$ for some θ .

(6)
$$e \models \underline{n}$$
 by assumption

By rule induction over Rules (15) on (6), only one case applies.

Case (15b).

(7)
$$e = \underline{n}$$
 by assumption
(8) $n > n \dashv \cdot$ by Rule (32c)

2. Prove $e \triangleright \underline{n} \dashv \theta$ implies $e \models \underline{n}$.

(6)
$$e \triangleright \underline{n} \dashv \theta$$
 by assumption

By rule induction over Rules (32) on (6), only one case applies.

Case (32c).

- (7) $e = \underline{n}$ by assumption (8) $\theta = \cdot$ by assumption (9) $\underline{n} \models \underline{n}$ by Rule (15b)
- 3. Prove $e \models_{?} \underline{n}$ implies $e ? \underline{n}$.
 - (6) $e \models_? \underline{n}$

by assumption

By rule induction over Rules (17) on (6), only one case applies.

Case (17b).

- $\begin{array}{lll} (7) & e \text{ notintro} & & \text{by assumption} \\ (8) & \underline{n} \text{ refutable}_? & & \text{by Rule (31a)} \\ (9) & e ? \underline{n} & & \text{by Rule (33c) on (7)} \\ & & \text{and (8)} \end{array}$
- 4. Prove $e ? \underline{n}$ implies $e \models_{?} \underline{n}$.
 - (6) $e ? \underline{n}$

by assumption

By rule induction over Rules (33) on (6), only one case applies.

Case (33c).

 $\begin{array}{ll} (7) \ e \ {\rm notintro} & {\rm by \ assumption} \\ (8) \ \underline{n} \ {\rm refutable?} & {\rm by \ Rule} \ (11a) \\ (9) \ e \models_{?} \underline{n} & {\rm by \ Rule} \ (17) \ {\rm on} \ (7) \\ {\rm and} \ (8) & \end{array}$

Case (21f).

 $\begin{array}{ll} (4) \ \ p = \mathtt{inl}(p_1) & \text{by assumption} \\ (5) \ \ \xi = \mathtt{inl}(\xi_1) & \text{by assumption} \\ (6) \ \ \tau = (\tau_1 + \tau_2) & \text{by assumption} \\ (7) \ \ p_1 : \tau_1[\xi_1] \dashv \Gamma \; ; \Delta & \text{by assumption} \end{array}$

By rule induction over Rules (20) on (1), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20l),(20m).

(8)
$$e = \emptyset^u, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0)\{\hat{rs}\}\$$
 by assumption
(9) e notintro by Rule $(27a), (27b), (27c), (27d), (27e), (27f)$

1. Prove $e \models \operatorname{inl}(\xi_1)$ implies $e \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (15) on $e \models \operatorname{inl}(\xi_1)$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

2. Prove $e
ightharpoonup \operatorname{inl}(p_1) \dashv \theta$ implies $e \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (32) on $e
ightharpoonup \operatorname{inl}(p_1) \dashv \theta$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

- 3. Prove $e \models_? inl(\xi_1)$ implies e? $inl(p_1)$.
 - (10) $inl(p_1)$ refutable? by Rule (31d)
 - (11) e? $inl(p_1)$ by Rule (33c) on (9) and (10)
- 4. Prove e? $inl(p_1)$ implies $e \models_? inl(\xi_1)$.
 - (10) $inl(\xi_1)$ refutable? by Rule (11b)
 - (11) $e \models_? \operatorname{inl}(\xi_1)$ by Rule (17b) on (9) and (10)

Case (20j).

- (8) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (9) \cdot ; $\Delta_e \vdash e_1 : \tau_1$ by assumption
- (10) e_1 final by Lemma 4.0.3 on (2)

By inductive hypothesis on (10) and (9) and (7).

- (11) $e_1 \models \xi_1 \text{ iff } e_1 \triangleright p_1 \dashv \theta \text{ for some } \theta$
- (12) $e_1 \models_? \xi_1 \text{ iff } e_1 ? p_1$
- 1. Prove $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ implies $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ .
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by assumption

By rule induction over Rules (15) on (13), only one case applies. Case (15g).

- (14) $e_1 \models \xi_1$ by assumption
- (15) $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 by (11) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta_1$ by Rule (32e) on (15)
- 2. Prove $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ implies $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ by assumption

By rule induction over Rules (32) on (13), only one case applies.

Case (32e).

- (14) $e_1 \triangleright p_1 \dashv \theta$ by assumption
- (15) $e_1 \models \xi_1$ by (11) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (15g) on (15)
- 3. Prove $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ implies $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by assumption

By rule induction over Rules (17) on (13), only two cases apply. Case (17b).

(14) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

Contradicts Lemma 4.0.7.

Case (17e).

- (14) $e_1 \models_? \xi_1$ by assumption (15) $e_1 ? p_1$ by (12) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ by Rule (33g) on (15)
- 4. Prove $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ implies $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ by assumption

By rule induction over Rules (33) on (13), only two cases apply.

Case (33c).

(14) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption Contradicts Lemma 4.0.7.

Case (33g).

- $\begin{array}{lll} (14) & e_1 ? \, p_1 & & \text{by assumption} \\ (15) & e_1 \models_? \xi_1 & & \text{by (12) on (14)} \end{array}$
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by Rule (17e) on (15)

Case (21g).

(4) $p = inr(p_2)$ by assumption (5) $\xi = inr(\xi_2)$ by assumption (6) $\tau = (\tau_1 + \tau_2)$ by assumption (7) $p_2 : \tau_2[\xi_2] \dashv \Gamma ; \Delta$ by assumption

By rule induction over Rules (20) on (1), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20l),(20m).

- (8) $e = ()^u, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$ by assumption (9) e notintro by Rule (27a), (27b), (27c), (27d), (27e), (27f)
- 1. Prove $e \models \operatorname{inr}(\xi_2)$ implies $e \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ . By rule induction over Rules (15) on $e \models \operatorname{inr}(\xi_2)$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

2. Prove $e
ightharpoonup \operatorname{inr}(p_2) \dashv \theta$ implies $e \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (32) on $e \rhd \operatorname{inr}(p_2) \dashv \theta$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

- 3. Prove $e \models_? \operatorname{inr}(\xi_2)$ implies e? $\operatorname{inr}(p_2)$.
 - (10) $inr(p_2)$ refutable? by Rule (31e)
 - (11) e? $inr(p_2)$ by Rule (33c) on (9) and (10)

- 4. Prove e? $inr(p_2)$ implies $e \models_? inr(\xi_2)$.
 - (10) $inr(\xi_2)$ refutable?

by Rule (11c)

(11) $e \models_? inr(\xi_2)$

by Rule (17b) on (9)

and (10)

Case (20k).

 $(8) \ e = \operatorname{inr}_{\tau_1}(e_2)$

by assumption

 $(9) \cdot ; \Delta_e \vdash e_2 : \tau_2$

by assumption

(10) e_2 final

by Lemma 4.0.3 on (2)

By inductive hypothesis on (10) and (9) and (7).

- (11) $e_2 \models \xi_2 \text{ iff } e_2 \triangleright p_2 \dashv \theta \text{ for some } \theta$
- (12) $e_2 \models_? \xi_2 \text{ iff } e_2 ? p_2$
- 1. Prove $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ implies $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ .
 - $(13) \ \operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$

by assumption

By rule induction over Rules (15) on (13), only one case applies.

Case (15g).

 $(14) e_2 \models \xi_2$

by assumption

- (15) $e_2 \triangleright p_2 \dashv \theta_1$ for some θ_1
- by (11) on (14)
- (16) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta_1$ by Rule (32e) on (15)
- 2. Prove $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ implies $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$.
 - (13) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$

by assumption

By rule induction over Rules (32) on (13), only one case applies.

Case (32e).

(14) $e_2 \triangleright p_2 \dashv \theta$

by assumption

(15) $e_2 \models \xi_2$

- by (11) on (14)
- (16) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$
- by Rule (15g) on (15)
- 3. Prove $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$ implies $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$.
 - (13) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$

by assumption

By rule induction over Rules (17) on (13), only two cases apply.

Case (17b).

(14) $\operatorname{inr}_{\tau_1}(e_2)$ notintro

by assumption

Contradicts Lemma 4.0.7.

Case (17e).

 $(14) e_2 \models_? \xi_2$

by assumption

 $(15) e_2? p_2$

- by (12) on (14)
- (16) $inr_{\tau_1}(e_2)$? $inr(p_2)$
- by Rule (33g) on (15)
- 4. Prove $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$ implies $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$.
 - (13) $inr_{\tau_1}(e_2)$? $inr(p_2)$

by assumption

By rule induction over Rules (33) on (13), only two cases apply.

Case (33c).

(14) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption Contradicts Lemma 4.0.7.

Case (33g).

(14) e_2 ? p_2 by assumption (15) $e_2 \models_? \xi_2$ by (12) on (14) (16) $\inf_{\tau_1}(e_2) \models_? \inf(\xi_2)$ by Rule (17e) on (15)

Case (21h).

$$(4) \ \ p = (p_1, p_2) \qquad \qquad \text{by assumption}$$

$$(5) \ \ \xi = (\xi_1, \xi_2) \qquad \qquad \text{by assumption}$$

$$(6) \ \ \tau = (\tau_1 \times \tau_2) \qquad \qquad \text{by assumption}$$

$$(7) \ \ \Gamma = \Gamma_1 \uplus \Gamma_2 \qquad \qquad \text{by assumption}$$

$$(8) \ \ \Delta = \Delta_1 \uplus \Delta_2 \qquad \qquad \text{by assumption}$$

$$(9) \ \ p_1 : \tau_1[\xi_1] \dashv \Gamma_1 ; \Delta_1 \qquad \qquad \text{by assumption}$$

$$(10) \ \ p_2 : \tau_2[\xi_2] \dashv \Gamma_2 ; \Delta_2 \qquad \qquad \text{by assumption}$$

By rule induction over Rules (20) on (1), the following cases apply.

Case (20b),(20c),(20f),(20h),(20i),(20i),(20m).

```
(11) e = \{ \| u, \| e_0 \| u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \} \}
                                                    by assumption
(12) e notintro
                                                    by Rule
                                                    (27a),(27b),(27c),(27d),(27e),(27f)
(13) e indet
                                                    by Lemma 4.0.10 on
                                                    (2) and (12)
(14) prl(e) indet
                                                    by Rule (25g) on (13)
(15) prl(e) final
                                                    by Rule (26b) on (14)
(16) prr(e) indet
                                                    by Rule (25h) on (13)
(17) prr(e) final
                                                    by Rule (26b) on (16)
                                                    by Rule (20h) on (1)
(18) \cdot; \Delta \vdash \mathtt{prl}(e) : \tau_1
(19) \cdot; \Delta \vdash \mathsf{prr}(e) : \tau_2
                                                   by Rule (20i) on (1)
```

By inductive hypothesis on (9) and (18) and (15) and by inductive hypothesis on (10) and (19) and (17).

- (20) $\operatorname{prl}(e) \models \xi_1 \text{ iff } \operatorname{prl}(e) \triangleright p_1 \dashv \theta_1 \text{ for some } \theta_1$
- (21) $prl(e) \models_? \xi_1 \text{ iff } prl(e) ? p_1$
- (22) $\operatorname{prr}(e) \models \xi_2 \text{ iff } \operatorname{prr}(e) \triangleright p_2 \dashv \theta_2 \text{ for some } \theta_2$
- (23) $prr(e) \models_? \xi_2 \text{ iff } prr(e) ? p_2$

```
1. Prove e \models (\xi_1, \xi_2) implies e \triangleright (p_1, p_2) \dashv \theta for some \theta.
```

(24)
$$e \models (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (15) on (24), only one case applies. Case (15j).

$$(25) \operatorname{prl}(e) \models \xi_1$$

by assumption

(26)
$$prr(e) \models \xi_2$$

by assumption

$$(27) \ \operatorname{prl}(e) \rhd p_1 \dashv \theta_1$$

by (20) on (25)

(28)
$$\operatorname{prr}(e) \rhd p_2 \dashv \theta_2$$

by (22) on (26)

$$(29) \quad e \rhd (p_1, p_2) \dashv \mid \theta_1 \uplus \theta_2$$

by Rule (32g) on (12) $\,$

and
$$(27)$$
 and (28)

2. Prove $e \rhd (p_1, p_2) \dashv \theta$ implies $e \models (\xi_1, \xi_2)$.

$$(24) e \rhd (p_1, p_2) \dashv \theta$$

by assumption

By rule induction over Rules (32) on (24), only one case applies.

Case (32g).

$$(25) \ \theta = \theta_1 \uplus \theta_2$$

by assumption

(26)
$$\operatorname{prl}(e) \rhd \xi_1 \dashv \theta_1$$

by assumption

(27)
$$\operatorname{prr}(e) \rhd \xi_2 \dashv \theta_2$$

(28) $\operatorname{prl}(e) \models \xi_1$

by assumption

(29)
$$\operatorname{prr}(e) \models \xi_2$$

by (20) on (26)

(30)
$$e \models (\xi_1, \xi_2)$$

by (22) on (27) by Rule (15j) on (12)

and (28) and (29)

3. Prove $e \models_{?} (\xi_1, \xi_2)$ implies $e ? (p_1, p_2)$.

(24)
$$e \models_{?} (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (17) on (24), only one case applies. Case (17b).

(25)
$$(\xi_1, \xi_2)$$
 refutable?

by assumption

By rule induction over Rules (11) on (25), only two cases apply.

Case (11d).

(26)
$$\xi_1$$
 refutable?

by assumption

$$(27)$$
 prl (e) notintro

by Rule (27e)

(28)
$$prl(e) \models_? \xi_1$$

by Rule (17b) on (26)

and
$$(27)$$

(29)
$$prl(e) ? p_1$$

by (21) on (28)

By rule induction over Rules (33) on (29), only three cases apply.

Case (33a),(33b).

(30)
$$p_1 = (1)^w, (p_0)^w$$

by assumption

$$(31)$$
 p_1 refutable?

by Rule (31b) and Rule

(31c)

$$(32)$$
 (p_1, p_2) refutable?

by Rule (31f) on (31)

(33) $e ? (p_1, p_2)$	by Rule (33c) on (12) and (32)	
Case (33c).		
(30) p_1 refutable?	by assumption	
(31) (p_1,p_2) refutable?	by Rule (31f) on (30)	
(32) $e ? (p_1, p_2)$	by Rule $(33c)$ on (12)	
	and (31)	
Case (11e).		
(26) ξ_2 refutable?	by assumption	
(27) $prr(e)$ notintro	by Rule (27e)	
$(28) \ prr(e) \models_? \xi_2$	by Rule (17b) on (26) and (27)	
$(29) \ \mathtt{prr}(e) \mathbin{?} p_2$	by (23) on (28)	
By rule induction over Rules (33) on (29), only three cases apply.		
Case (33a),(33b).		
$(30) \ p_2 = ()^w, (p_0)^w$	by assumption	
(31) p_2 refutable?	by Rule (31b) and Rule (31c)	
$(32) \ (p_1,p_2) \ { t refutable}_?$	by Rule (31g) on (31)	
(33) $e ? (p_1, p_2)$	by Rule (33c) on (12) and (32)	
Case (33c).		
(30) p_2 refutable?	by assumption	
$(31) \;\; (p_1,p_2) \; exttt{refutable}_?$	by Rule $(31g)$ on (30)	
(32) $e ? (p_1, p_2)$	by Rule (33c) on (12) and (31)	
4. Prove $e ? (p_1, p_2)$ implies $e \models_? (\xi_1, \xi_2)$.		
$(24) e? (p_1, p_2)$	by assumption	
By rule induction over Rules (33) on (24), only one case applies.	
Case (33c).		
$(25) \ (p_1,p_2) \ { t refutable}_?$	by assumption	
By rule induction over Rules (31) on (25), only two cases	
apply.		
Case (31f).		
(26) p_1 refutable?	by assumption	
(27) prl (e) notintro	by Rule (27e)	
$(28) \ prl(e) ? p_1$	by Rule (33c) on (26) and (27)	
(29) $prl(e) \models_? \xi_1$	by (21) on (28)	

By rule induction over Rules (17) on (29), only three cases apply.

Case (17a).

- (30) $\xi_1 = ?$ by assumption (31) ξ_1 refutable? by Rule (2b)
- (32) (ξ_1, ξ_2) refutable? by Rule (11d) on (31)
- (33) $e \models_{?} (\xi_1, \xi_2)$ by Rule (17b) on (12) and (32)

Case (17b).

- (30) ξ_1 refutable? by assumption
- (31) (ξ_1, ξ_2) refutable? by Rule (11d) on (30)
- (32) $e \models_{?} (\xi_1, \xi_2)$ by Rule (17b) on (12) and (31)

Case (31g).

- (26) p_2 refutable? by assumption
- (27) prr(e) notintro by Rule (27e) (28) prr(e)? p_2 by Rule (33c) on (26)
- and (27)
- (29) $prr(e) \models_{?} \xi_{2}$ by (23) on (28)

By rule induction over Rules (17) on (29), only three cases apply.

Case (17a).

- (30) $\xi_2 = ?$ by assumption (31) ξ_2 refutable? by Rule (2b)
- (32) (ξ_1, ξ_2) refutable? by Rule (11e) on (31)
- (33) $e \models_{?} (\xi_1, \xi_2)$ by Rule (17b) on (12) and (32)

Case (17b).

- (30) ξ_2 refutable? by assumption
- (31) (ξ_1, ξ_2) refutable? by Rule (11e) on (30)
- (32) $e \models_{?} (\xi_1, \xi_2)$ by Rule (17b) on (12) and (31)

Case (20g).

- (11) $e = (e_1, e_2)$ by assumption (12) \cdot ; $\Delta_e \vdash e_1 : \tau_1$ by assumption
- (13) $\cdot : \Delta_e \vdash e_2 : \tau_2$ by assumption
- (14) e_1 final by Lemma 4.0.5 on (2)
- (15) e_2 final by Lemma 4.0.5 on (2)

By inductive hypothesis on (14) and (12) and (9) and by inductive hypothesis on (15) and (13) and (10).

```
(16) e_1 \models \xi_1 \text{ iff } e_1 \triangleright p_1 \dashv \theta_1 \text{ for some } \theta_1
```

- (17) $e_1 \models_? \xi_1 \text{ iff } e_1 ? p_1$
- (18) $e_2 \models \xi_2 \text{ iff } e_2 \triangleright p_2 \dashv \theta_2 \text{ for some } \theta_2$
- (19) $e_2 \models_? \xi_2 \text{ iff } e_2 ? p_2$
- 1. Prove $(e_1, e_2) \models (\xi_1, \xi_2)$ implies $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ for some θ .

$$(20) (e_1, e_2) \models (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (15) on (20), only two cases apply.

Case (15i).

- (21) $e_1 \models \xi_1$ by assumption
- (22) $e_2 \models \xi_2$ by assumption
- (23) $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 by (16) on (21)
- (24) $e_2 \triangleright p_2 \dashv \theta_2$ for some θ_2 by (18) on (22)
- (25) $(e_1, e_2) \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2$ by Rule (32d) on (23) and (24)

Case (15j).

(21) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.9.

2. Prove $(e_1, e_2) > (p_1, p_2) \dashv \theta$ implies $(e_1, e_2) \models (\xi_1, \xi_2)$.

(20)
$$(e_1, e_2) \rhd (p_1, p_2) \dashv \theta$$

by assumption

By rule induction over Rules (32) on (20), only two cases apply. Case (32d).

- (21) $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 by assumption
- (22) $e_2 \triangleright p_2 \dashv \theta_2$ for some θ_2 by assumption
- (23) $e_1 \models \xi_1$ by (16) on (21)
- (24) $e_2 \models \xi_2$ by (18) on (22)
- (25) $(e_1, e_2) \models (\xi_1, \xi_2)$ by Rule (15i) on (23) and (24)

Case (32g).

(21) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.9.

- 3. Prove $(e_1, e_2) \models_? (\xi_1, \xi_2)$ implies $(e_1, e_2) ? (p_1, p_2)$.
 - $(20) (e_1, e_2) \models_? (\xi_1, \xi_2)$

by assumption

By rule induction over Rules (17) on (20), only four cases apply. Case (17b).

(21) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.9.

Case (17g).

(21) $e_1 \models_? \xi_1$

by assumption

```
(22) e_2 \models \xi_2
                                           by assumption
(23) e_1 ? p_1
                                           by (17) on (21)
(24) \ e_2 \rhd p_2 \dashv \mid \theta_2
                                           by (18) on (22)
(25) (e_1, e_2)? (p_1, p_2)
                                           by Rule (33d) on (23)
                                           and (24)
```

Case (17h).

- (21) $e_1 \models \xi_1$ by assumption (22) $e_2 \models_? \xi_2$ by assumption (23) $e_1 \triangleright p_1 \dashv \mid \theta_1$ by (16) on (21) (24) $e_2 ? p_2$ by (19) on (22) (25) (e_1, e_2) ? (p_1, p_2) by Rule (33e) on (23) and (24)

Case (17i).

- (21) $e_1 \models_? \xi_1$ by assumption (22) $e_2 \models_? \xi_2$ by assumption (23) $e_1 ? p_1$ by (17) on (21) (24) $e_2 ? p_2$ by (19) on (22) (25) (e_1, e_2) ? (p_1, p_2) by Rule (33f) on (23)and (24)
- 4. Prove (e_1, e_2) ? (p_1, p_2) implies $(e_1, e_2) \models_? (\xi_1, \xi_2)$.
 - (20) (e_1, e_2) ? (p_1, p_2)

by assumption

By rule induction over Rules (33) on (20), only four cases apply.

Case (33c).

(21) (e_1,e_2) notintro by assumption Contradicts Lemma 4.0.9.

Case (33d).

(21) $e_1 ? p_1$ by assumption $(22) e_2 \rhd p_2 \dashv \theta_2$ by assumption (23) $e_1 \models_? \xi_1$ by (17) on (21) $(24) e_2 \models \xi_2$ by (18) on (22) (25) (e_1, e_2) ? (p_1, p_2) by Rule (17g) on (23) and (24)

Case (33e).

(21) $e_1 \triangleright p_1 \dashv \theta_1$ by assumption (22) $e_2 ? p_2$ by assumption (23) $e_1 \models \xi_1$ by (16) on (21) (24) $e_2 \models_? \xi_2$ by (19) on (22) (25) (e_1, e_2) ? (p_1, p_2) by Rule (17h) on (23) and (24)

Case (33f).

(21)	$e_1 ? p_1$	by assumption
(22)	$e_2 ? p_2$	by assumption
(23)	$e_1 \models_? \xi_1$	by (17) on (21)
(24)	$e_2 \models_? \xi_2$	by (19) on (22)
(25)	$(e_1, e_2) ? (p_1, p_2)$	by Rule (17i) on (23)
		and (24)

5 Preservation and Progress

Theorem 5.1 (Preservation). If \cdot ; $\Delta \vdash e : \tau$ and $e \mapsto e'$ then \cdot ; $\Delta \vdash e' : \tau$

Proof. By rule induction over Rules (20) on typing judgment of e. For simplicity, we only consider two cases for match expressions here.

Case (201).

$(1) \ \cdot ; \Delta \vdash \mathtt{match}(e_1) \{ \cdot \mid r \mid rs \} : \tau$	by assumption
$(2) \ \mathtt{match}(e_1)\{\cdot \mid r \mid rs\} \mapsto e'$	by assumption
$(3) \cdot ; \Delta \vdash e_1 : \tau_1$	by assumption
$(4) \ \cdot ; \Delta \vdash [\bot](r \mid rs) : \tau_1[\xi] \Rightarrow \tau$	by assumption
(5) $\top \models_2^{\dagger} \xi$	by assumption

By rule induction over Rules (35) on (2).

Case (35k).

$(6) e' = \mathtt{match}(e_1')\{\cdot \mid r \mid rs\}$	by assumption
$(7) e_1 \mapsto e_1'$	by assumption
(8) \cdot ; $\Delta \vdash e'_1 : \tau_1$	by IH on (3) and (7)
$(9) \cdot ; \Delta \vdash \mathtt{match}(e_1') \{ \cdot \mid r \mid rs \} : \tau$	by Rule (20l) on (8)
	and (4) and (5)

Case (351).

(6)
$$r = p_r \Rightarrow e_r$$
 by assumption
(7) $e' = [\theta](e_r)$ by assumption
(8) $e_1 \triangleright p_r \dashv \theta$ by assumption

By rule induction over Rules (23) on (4).

Case (23a).

$$\begin{array}{ll} (9) & \xi = \xi_r & \text{by assumption} \\ (10) & \cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau & \text{by assumption} \\ (11) & p_r : \tau_1[\xi_r] \dashv \Gamma_r \; ; \Delta_r & \text{by Inversion of Rule} \\ & (22a) \text{ on } (10) \end{array}$$

(12) $\Gamma_r ; \Delta \uplus \Delta_r \vdash e_r : \tau$	by Inversion of Rule (22a) on (10)
(13) $\theta:\Gamma_r$	by Lemma 3.0.7 on (3) and (11) and (8)
$(14) \cdot ; \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau$	by Lemma 3.0.6 on (12) and (13)
Case $(23b)$.	
$(9) \ \xi = \xi_r \vee \xi_{rs}$	by assumption
$(10) \cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$	by assumption
$(11) p_r : \tau_1[\xi_r] \dashv \mid \Gamma_r ; \Delta_r$	by Inversion of Rule (22a) on (10)
$(12) \ \Gamma_r ; \Delta \uplus \Delta_r \vdash e_r : \tau$	by Inversion of Rule (22a) on (10)
(13) $\theta:\Gamma_r$	by Lemma 3.0.7 on (3) and (11) and (8)
$(14) \cdot ; \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau$	by Lemma 3.0.6 on (12) and (13)
Case (35m).	
$(6) rs = r' \mid rs'$	by assumption
(7) $e' = \mathtt{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid r\}$	
(1) (Pr , Sr) (Pr	by assumption
(8) e_1 final	by assumption
(9) $e_1 \perp p_r$	by assumption
By rule induction over Rules (23) on (4).	
Case (23a). Syntactic contradiction of rs	
Case $(23b)$.	
$(10) \ \xi = \xi_r \vee \xi_{rs}$	by assumption
$(11) \cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$	by assumption
$(12) \cdot ; \Delta \vdash [\bot \lor \xi_r](r' \mid rs') : \tau_1[\xi_{rs}]$	
(40) + 1/ 1	by assumption
$(13) \ \xi_r \not\models \bot$	by assumption
$(14) \ p_r : \tau_1[\xi_r] \dashv \mid \Gamma_r ; \Delta_r$	by Inversion of Rule (22a) on (11)
$(15) \ \Gamma_r ; \Delta \uplus \Delta_r \vdash e_r : \tau$	by Inversion of Rule (22a) on (11)
$(16) \cdot ; \Delta \vdash [\bot](p_r \Rightarrow e_r \mid \cdot) : \tau_1[\xi_r] =$	by Rule (23a) on (11)
$(17) e_1 \not\models_?^\dagger \xi_r$	and (13) by Lemma 4.0.15 on (3) and (8) and (14) and (9)

(18)
$$\cdot$$
; $\Delta \vdash \mathtt{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'\} : \tau$ by Rule (20m) on (3) and (8) and (16) and (12) and (17) and (5)

Case (20m).

(1)
$$rs_{pre} = r_{pre} \mid rs'_{pre}$$
 by assumption

(2)
$$\cdot$$
; $\Delta \vdash \mathtt{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\} : \tau$ by assumption

(3)
$$\mathsf{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\} \mapsto e'$$
 by assumption

(4)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$ by assumption

(5)
$$e_1$$
 final by assumption

(6)
$$\cdot ; \Delta \vdash [\bot] rs_{pre} : \tau_1[\xi_{pre}] \Rightarrow \tau$$
 by assumption

$$(7) \quad \cdot \; ; \Delta \vdash [\bot \lor \xi_{pre}](r \mid rs_{post}) : \tau_1[\xi_{rest}] \Rightarrow \tau$$
 by assumption

(8)
$$e_1 \not\models_2^{\dagger} \xi_{pre}$$
 by assumption

(9)
$$\top \models_{?}^{\dagger} \xi_{pre} \vee \xi_{rest}$$
 by assumption

By rule induction over Rules (35) on (3).

Case (35k).

(10)
$$e' = \text{match}(e'_1)\{rs_{pre} \mid r \mid rs_{post}\}$$
 by assumption (11) $e_1 \mapsto e'_1$ by assumption

By Lemma 4.0.13, (11) contradicts (5).

Case (351).

$$\begin{array}{ll} (10) & r = p_r \Rightarrow e_r & \text{by assumption} \\ (11) & e' = [\theta](e_r) & \text{by assumption} \\ (12) & e_1 \rhd p_r \dashv \theta & \text{by assumption} \\ \end{array}$$

By rule induction over Rules (23) on (7).

Case (23a).

$$\begin{array}{lll} (13) & \xi_{rest} = \xi_r & \text{by assumption} \\ (14) & \cdot \; ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau & \text{by assumption} \\ (15) & p_r : \tau_1[\xi_r] \dashv \Gamma_r \; ; \Delta_r & \text{by Inversion of Rule} \\ (16) & \Gamma_r \; ; \Delta \uplus \Delta_r \vdash e_r : \tau & \text{by Inversion of Rule} \\ (17) & \theta : \Gamma_r & \text{by Lemma 3.0.7 on (4)} \\ (18) & \cdot \; ; \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau & \text{by Lemma 3.0.6 on} \end{array}$$

(16) and (17)

Case (23b). (13) $\xi_{rest} = \xi_r \vee \xi_{rs}$ by assumption $(14) \cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$ by assumption (15) $p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r$ by assumption (16) $\Gamma_r : \Delta \uplus \Delta_r \vdash e_r : \tau$ by assumption (17) $\theta : \Gamma_r$ by Lemma 3.0.7 on (4)and (15) and (12) (18) $\cdot ; \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau$ by Lemma 3.0.6 on (16) and (17)Case (35m). (10) $r = p_r \Rightarrow e_r$ by assumption (11) $rs_{post} = r' \mid rs'$ by assumption $(12) e' = \mathtt{match}(e_1) \{ (rs_{pre} \mid p_r \Rightarrow e_r \mid \cdot)^{\diamond} \mid r' \mid rs' \}$ by assumption (13) $e_1 \perp p_r$ by assumption By rule induction over Rules (23) on (7). Case (23a). Syntactic contradiction of rs_{post} . Case (23b). (14) $\xi_{rest} = \xi_r \vee \xi_{post}$ by assumption (15) $\cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$ by assumption $(16) \cdot ; \Delta \vdash [\bot \lor \xi_{pre} \lor \xi_r](r' \mid rs') : \tau_1[\xi_{post}] \Rightarrow \tau$ by assumption (17) $\xi_r \not\models \xi_{pre}$ by assumption (18) $p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r$ by Inversion of Rule (22a) on (15)(19) $\Gamma_r : \Delta \uplus \Delta_r \vdash e_r : \tau$ by Inversion of Rule (22a) on (15)(20) $\xi_r : \tau_1$ by Lemma 3.0.2 on (21) $\xi_{pre} : \tau_1$ by Lemma 3.0.3 on (6)(22) $\xi_r \not\models \bot \lor \xi_{pre}$ by Lemma 2.0.6 on (20) and (21) and (17) $(23) \ \cdot ; \Delta \vdash [\bot] (rs_{pre} \mid p_r \Rightarrow e_r \mid \cdot)^{\diamond} : \tau_1[\xi_{pre} \vee \xi_r] \Rightarrow \tau$ by Lemma 3.0.4 on (6)

(24) $e_1 \not\models_{?}^{\dagger} \xi_r$

(25) $e_1 \not\models_2^{\dagger} \xi_{pre} \vee \xi_r$

and (15) and (22)

and (13)

and (24)

by Lemma 4.0.15 on (4) and (5) and (18)

by Lemma 2.0.7 on (8)

(26)
$$\cdot$$
; $\Delta \vdash \mathtt{match}(e_1)\{(rs_{pre} \mid p_r \Rightarrow e_r \mid \cdot)^{\diamond} \mid r' \mid rs'\} : \tau$ by Rule (20m) on (4) and (5) and (23) and (16) and (25) and (9)

Theorem 5.2 (Progress). If \cdot ; $\Delta \vdash e : \tau$ then either e final or $e \mapsto e'$ for some e'.

Proof. By rule induction over Rules (20) on typing judgment of e. For simplicity, we only consider two cases for match expressions here.

Case (201).

- (1) \cdot ; $\Delta \vdash \mathsf{match}(e_1)\{\cdot \mid r \mid rs\} : \tau$ by assumption
- (2) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (3) $\cdot : \Delta \vdash [\bot](r \mid rs) : \tau_1[\xi] \Rightarrow \tau$ by assumption
- (4) $\top \models^{\dagger}_{?} \xi$ by assumption

By IH on (2).

Case Scrutinee takes a step.

- (5) $e_1 \mapsto e_1'$ by assumption
- $\begin{array}{c} (6) \ \mathtt{match}(e_1)\{\cdot \mid r \mid rs\} \mapsto \mathtt{match}(e_1')\{\cdot \mid r \mid rs\} \\ \qquad \qquad \qquad \text{by Rule (35k) on (5)} \end{array}$

Case Scrutinee is final.

(5) e_1 final by assumption

By rule induction over Rules (23) on (3).

Case (23a).

- (6) $rs = \cdot$ by assumption
- (7) $\xi = \xi_r$ by assumption
- (8) $\cdot : \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau$ by assumption
- (9) $r = p_r \Rightarrow e_r$ by Inversion of Rule (22a) on (8)
- (10) $p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r$ by Inversion of Rule (22a) on (8)
- (11) $e_1 \models^{\dagger}_{?} \xi_r$ by Corollary 2.1.1 on (5) and (4)

By rule induction over Rules (18) on (11).

Case (18a).

(12) $e_1 \models_? \xi_r$ by assumption

(13)
$$e_1$$
 ? p_r by Lemma 4.0.15 on (2) and (5) and (10) and (12)

(14)
$$\mathrm{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid\cdot\}$$
 indet

by Rule (25k) on (5) and (13)

(15)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\} \text{ final}$$
 by Rule (26b) on (14)

Case (18b).

(12)
$$e_1 \models \xi_r$$
 by assumption

(13)
$$e_1 \triangleright p_r \dashv \theta$$
 by Lemma 4.0.15 on (2) and (5) and (10) and (12)

(14)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\} \mapsto [\theta](e_r)$$
 by Rule (351) on (5) and (13)

Case (23b).

(6)
$$rs = r' \mid rs'$$
 by assumption

(7)
$$\xi = \xi_r \vee \xi_{rs}$$
 by assumption

(8)
$$\cdot ; \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau$$
 by assumption

(9)
$$r = p_r \Rightarrow e_r$$
 by Inversion of Rule (22a) on (8)

(10)
$$p_r: \tau_1[\xi_r] \dashv \Gamma_r; \Delta_r$$
 by Inversion of Rule (22a) on (8)

By Lemma 4.0.14 on (2) and (5) and (10).

Case Scrutinee matches pattern.

(11)
$$e_1 \triangleright p_r \dashv \theta$$
 by assumption

(12)
$$\operatorname{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid rs\}\mapsto [\theta](e_r)$$
 by Rule (351) on (5) and (11)

Case Scrutinee may matches pattern.

(11)
$$e_1 ? p_r$$
 by assumption

(13)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs\}$$
 final by Rule (26b) on (12)

Case Scrutinee doesn't matche pattern.

(11)
$$e_1 \perp p_r$$
 by assumption

$$\begin{array}{ll} (12) \ \operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid (r' \mid rs')\} \\ \mapsto \operatorname{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'\} \\ & \text{by Rule (35m) on (5)} \\ & \text{and (11)} \end{array}$$

Case (20m).

(1) $rs_{pre} = r_{pre} \mid rs'_{pre}$ by assumption (2) $\cdot ; \Delta \vdash \mathsf{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\} : \tau$ by assumption (3) $\cdot ; \Delta \vdash e_1 : \tau_1$ by assumption (4) e_1 final by assumption (5) $\cdot ; \Delta \vdash [\xi_{pre}](r \mid rs_{post}) : \tau_1[\xi_{rest}] \Rightarrow \tau$ by assumption (6) $e_1 \not\models_7^\dagger \xi_{pre}$ by assumption (7) $\top \models_7^\dagger \xi_{pre} \lor \xi_{rest}$ by assumption

By rule induction over Rules (23) on (5).

Case (23a).

$$\begin{array}{lll} (5) & rs_{post} = \cdot & & \text{by assumption} \\ (6) & \xi_{rest} = \xi_r & & \text{by assumption} \\ (7) & \cdot ; \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau & & \text{by assumption} \\ (8) & r = p_r \Rightarrow e_r & & \text{by Inversion of Rule} \\ (9) & p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r & & \text{by Inversion of Rule} \\ (10) & e_1 \models_{?}^{\dagger} \xi_{pre} \lor \xi_r & & \text{by Corollary 2.1.1 on} \\ (11) & e_1 \models_{?}^{\dagger} \xi_r & & \text{by Lemma 2.0.8 on} \\ (11) & e_1 \models_{?}^{\dagger} \xi_r & & \text{by Lemma 2.0.8 on} \\ (10) & \text{and } (6) & & \\ \end{array}$$

By rule induction over Rules (18) on (11).

Case (18a).

(12)
$$e_1 \models_? \xi_r$$
 by assumption
(13) $e_1 ? p_r$ by Lemma 4.0.15 on
(3) and (4) and (9) and
(12)
(14) $\mathsf{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\}$ indet
by Rule (25k) on (4)
and (13)
(15) $\mathsf{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\}$ final
by Rule (26b) on (14)

Case (18b).

Case (23b).

$$\begin{array}{lll} \text{(5)} & rs_{post} = r' \mid rs'_{post} & \text{by assumption} \\ \text{(6)} & \cdot ; \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau & \text{by assumption} \\ \text{(7)} & r = p_r \Rightarrow e_r & \text{by Inversion of Rule} \\ \text{(8)} & p_r : \tau_1[\xi_r] \dashv \mid \Gamma_r \; ; \Delta_r & \text{by Inversion of Rule} \\ \text{(22a) on (6)} & \text{(22a) on (6)} \\ \end{array}$$

By Lemma 4.0.14 on (3) and (4) and (8).

Case Scrutinee matches pattern.

$$\begin{array}{ll} (9) & e_1 \rhd p_r \dashv \theta & \text{by assumption} \\ (10) & \mathtt{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs_{post}\} \mapsto [\theta](e_r) \\ & \text{by Rule (35l) on (4)} \\ & \text{and (9)} \end{array}$$

Case Scrutinee may matches pattern.

(9)
$$e_1 ? p_r$$
 by assumption
(10) $\mathsf{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs_{post}\}$ indet by Rule (25k) on (4) and (9)
(11) $\mathsf{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs_{post}\}$ final by Rule (26b) on (10)

Case Scrutinee doesn't matche pattern.

$$\begin{array}{ll} (9) & e_1 \perp p_r & \text{by assumption} \\ (10) & \mathtt{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid (r' \mid rs'_{post})\} \\ & \mapsto \mathtt{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'_{post}\} \\ & \text{by Rule (35m) on (4)} \\ & \text{and (9)} \end{array}$$

6 Decidability

 Ξ incon A finite set of constraints, Ξ , is inconsistent

CINCTruth
$$\underline{\Xi}$$
 incon $\underline{\Xi}, \top$ incon (36a)

CINCFalsity $\underline{\Xi}, \perp$ incon (36b)

CINCNum $n_1 \neq n_2$ $\underline{\Xi}, \underline{n_1}, \underline{n_2}$ incon (36c)

CINCNotNum $\underline{\Xi}, \underline{n}, \underline{\mathscr{M}}$ incon (36d)

CINCAnd $\underline{\Xi}, \underline{\xi_1}, \underline{\xi_2}$ incon (36e)

CINCOr $\underline{\Xi}, \underline{\xi_1}$ incon $\underline{\Xi}, \underline{\xi_1}$ incon (36f)

CINCOr $\underline{\Xi}, \underline{\xi_1}$ incon $\underline{\Xi}, \underline{\xi_2}$ incon (36f)

CINCInj $\underline{\Xi}$ incon (36f)

CINCInl $\underline{\Xi}$ incon (36h)

CINCInr $\underline{\Xi}$ incon (36h)

CINCInr $\underline{\Xi}$ incon (36i)

CINCPairL $\underline{\Xi}_1$ incon (36j)

CINCPairR

(36k)

Lemma 6.0.1 (Decidability of Inconsistency). Suppose $\dot{\top}(\xi) = \xi$. It is decidable whether ξ incon.

 $\Xi_2 \ {\tt incon}$

 $\overline{(\Xi_1,\Xi_2)}$ incon

Lemma 6.0.2 (Inconsistency and Entailment of Constraint). Suppose that $\dot{\top}(\xi) = \xi$. Then $\bar{\xi}$ incon iff $\top \models \xi$

Lemma 6.0.3. If $e \models \xi$ then $e \models \dot{\top}(\xi)$

Proof. By rule induction over Rules (15), it is obvious to see that $\dot{\top}(\xi) = \xi$. \Box

Lemma 6.0.4. If $e \models_? \xi$ then $e \models_?^{\dagger} \dot{\top}(\xi)$.

Proof.

(11)
$$e \models_? \xi$$

By Rule Induction over Rules (17) on (11).

Case (17a).

(12)
$$\xi = ?$$
 by assumption (13) $e \models \top$ by Rule (15a)

(14)
$$e \models_2^{\dagger} \top$$
 by Rule (18b) on (13)

by assumption

Case (17b).

(13) ξ refutable? by assumption

By Lemma 2.0.15 on (12) and (13) and case analysis on its conclusion. By rule induction over Rules (11).

Case $\dot{\top}(\xi)$ refutable?.

(14)
$$\dot{\top}(\xi)$$
 refutable? by assumption

(15)
$$e \models_? \dot{\top}(\xi)$$
 by Rule (17b) on (12) and (14)

(16)
$$e \models_{?}^{\dagger} \dot{\top}(\xi)$$
 by Rule (18b) on (15)

Case $e \models \dot{\top}(\xi)$.

(14)
$$e \models \dot{\top}(\xi)$$
 by assumption

(15)
$$e \models^{\dagger}_{?} \top$$
 by Rule (18b) on (14)

Case (17c).

(12)
$$\xi = \xi_1 \vee \xi_2$$
 by assumption
(13) $e \models_? \xi_1$ by assumption
(14) $e \models_?^{\dagger} \dot{\top}(\xi_1)$ by IH on (13)

(15)
$$e \models^{\dagger}_{?} \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$$
 by Lemma 2.0.10 on (14)

Case (17d).

(12)
$$\xi = \xi_1 \vee \xi_2$$
 by assumption

- (13) $e \models_? \xi_2$ by assumption (14) $e \models_?^\dagger \dot{\top}(\xi_2)$ by IH on (13)
- (15) $e \models_{?}^{\uparrow} \dot{\top}(\xi_{1}) \lor \dot{\top}(\xi_{2})$ by Lemma 2.0.10 on (14)

Case (17e).

- $(12) \ e = \operatorname{inl}_{\tau_2}(e_1) \qquad \qquad \text{by assumption}$ $(13) \ \xi = \operatorname{inl}(\xi_1) \qquad \qquad \text{by assumption}$ $(14) \ e_1 \models_? \xi_1 \qquad \qquad \text{by assumption}$
- (15) $e_1 \models^{\dagger}_{?} \dot{\top}(\xi_1)$ by IH on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\top}(\xi_1))$ by Lemma 2.0.11 on (15)

Case (17f).

(12) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption (13) $\xi = \operatorname{inr}(\xi_2)$ by assumption (14) $e_2 \models_? \xi_2$ by assumption (15) $e_2 \models_?^{\dagger} \dot{\top}(\xi_2)$ by IH on (14) (16) $\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inr}(\dot{\top}(\xi_2))$ by Lemma 2.0.12 on

(15)

(10) $\operatorname{Im}_{\tau_1}(e_2) \models_? \operatorname{Im}(+(\xi_2))$

Case (17g).

- $(12) \ e = (e_1, e_2)$ by assumption $(13) \ \xi = (\xi_1, \xi_2)$ by assumption $(14) \ e_1 \models_? \xi_1$ by assumption $(15) \ e_2 \models \xi_2$ by assumption $(16) \ e_1 \models_?^\dagger \dot{\top}(\xi_1)$ by IH on (14)
- (17) $e_2 \models \dot{\top}(\xi_2)$ by Lemma 6.0.3 on (15)
- (18) $e_2 \models_?^{\dagger} \dot{\top}(\xi_2)$ by Rule (18b) on (17) (19) $(e_1, e_2) \models_?^{\dagger} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Lemma 2.0.13 on (16) and (18)

Case (17h).

(12)
$$e = (e_1, e_2)$$

by assumption

(13)
$$\xi = (\xi_1, \xi_2)$$

by assumption

(14)
$$e_1 \models \xi_1$$

by assumption

(15)
$$e_2 \models_? \xi_2$$

by assumption

(16)
$$e_1 \models \dot{\top}(\xi_1)$$

by Lemma 6.0.3 on (14)

(17)
$$e_1 \models^{\dagger}_{?} \dot{\top}(\xi_1)$$

by Rule (18b) on (16)

(18)
$$e_2 \models^{\dagger}_{?} \dot{\top}(\xi_2)$$

by IH on (15)

(19)
$$(e_1, e_2) \models_?^{\dagger} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$$

by Lemma 2.0.13 on

(17) and (18)

Case (17i).

(12)
$$e = (e_1, e_2)$$

by assumption

(13)
$$\xi = (\xi_1, \xi_2)$$

by assumption

(14)
$$e_1 \models_? \xi_1$$

by assumption

$$(15) e_2 \models_? \xi_2$$

by assumption

$$(16) \ e_1 \models^{\dagger}_? \dot{\top}(\xi_1)$$

by IH on (14)

(17)
$$e_2 \models_{?}^{\dagger} \dot{\top}(\xi_2)$$

by IH on (15)

(18) $(e_1, e_2) \models^{\dagger}_{?} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$

by Lemma 2.0.13 on (16) and (17)

Lemma 6.0.5. $e \models^{\dagger}_{?} \xi \text{ iff } e \models^{\dagger}_{?} \dot{\top}(\xi)$

Proof. 1. Sufficiency:

(1)
$$e \models_{?}^{\dagger} \xi$$

by assumption

By rule induction over Rules (18) on (1)

Case (18b).

(2)
$$e \models \xi$$

by assumption

(3)
$$e \models \dot{\top}(\xi)$$

by Lemma 6.0.3 on (2)

(4)
$$e \models_2^{\dagger} \dot{\top}(\xi)$$

by Rule (18b) on (3)

Case (18a).

(2)
$$e \models_? \xi$$

by assumption

(3)
$$e \models_?^\dagger \dot{\top}(\xi)$$

by Lemma 6.0.4 on (2)

2. Necessity:

(1)
$$e \models_2^{\dagger} \dot{\top}(\xi)$$

by assumption

By structural induction on ξ ,

Case $\xi = \top, \bot, \underline{n}, \underline{\varkappa}$.

(2)
$$e \models^{\dagger}_{?} \xi$$

by (1) and Definition 13

Case $\xi = ?$.

(2)
$$e \models_? ?$$

by Rule (17a)

(3)
$$e \models_{?}^{\dagger} ?$$

by Rule (18a) on (2)

Case $\xi = \xi_1 \vee \xi_2$.

(2)
$$\dot{\top}(\xi_1 \vee \xi_2) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$$

by Definition 13

By rule induction over Rules (18) on (1),

Case (18b).

(3)
$$e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$$

by assumption

By rule induction over Rules (15) on (3) and two cases apply, Case (15e).

(4)
$$e \models \dot{\top}(\xi_1)$$

by assumption

(5)
$$e \models_?^\dagger \dot{\top}(\xi_1)$$

by Rule (18b) on (4)

(6)
$$e \models^{\dagger}_{?} \xi_{1}$$

by IH on (5)

(7)
$$e \models^{\dagger}_{?} \xi_1 \lor \xi_2$$

by Lemma 2.0.10 on (6)

Case (15f).

(4)
$$e \models \dot{\top}(\xi_2)$$

by assumption

(5)
$$e \models^{\dagger}_{?} \dot{\top}(\xi_2)$$

by Rule (18b) on (4)

(6)
$$e \models_{?}^{\dagger} \xi_{2}$$

by IH on (5)

$$(7) e \models^{\dagger}_{?} \xi_1 \vee \xi_2$$

by Lemma 2.0.10 on

(6)

Case (18a).

(3)
$$e \models_? \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$$

by assumption

By rule induction over Rules (17) on (3) and two cases apply, Case (17c).

- (4) $e \models_? \dot{\top}(\xi_1)$ by assumption (5) $e \models_?^{\dagger} \dot{\top}(\xi_1)$ by Rule (18a) on (4)
- (6) $e \models_{2}^{\dagger} \xi_{1}$ by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Lemma 2.0.10 on (6)

Case (17d).

- (4) $e \models_? \dot{\top}(\xi_2)$ by assumption
- (5) $e \models_{2}^{\dagger} \dot{\top}(\xi_{2})$ by Rule (18a) on (4)
- (6) $e \models_{?}^{\dagger} \xi_2$ by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Lemma 2.0.10 on (6)

Case $\xi = inl(\xi_1)$.

(2) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption (3) $\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$ by assumption

By rule induction over Rules (18) on (1),

Case (18b).

- (4) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\top}(\xi_1))$ by assumption By rule induction over Rules (15) and only one case applies, Case (15g).
 - (5) $e_1 \models \dot{\top}(\xi_1)$ by assumption (6) $e_1 \models_2^{\dagger} \dot{\top}(\xi_1)$ by Rule (18b) on (5)
 - (7) $e_1 \models_{7}^{+} \xi_1$ by IH on (6)
 - (8) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Lemma 2.0.11 on (7)

Case (18a).

- (4) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\top}(\xi_1))$ by assumption By rule induction over Rules (17) and only one case applies, Case (17e).
 - (5) $e_1 \models_? \dot{\top}(\xi_1)$
- by assumption
- (6) $e_1 \models^{\dagger}_? \dot{\top}(\xi_1)$
- by Rule (18a) on (5)

(7) $e_1 \models_{?}^{\dagger} \xi_1$

- by IH on (6)
- $(8) \ \operatorname{inl}_{\tau_2}(e_1) \models^\dagger_? \operatorname{inl}(\xi_1)$
- by Lemma 2.0.11 on (7)

Case $\xi = inr(\xi_2)$.

 $(2) e = \operatorname{inr}_{\tau_1}(e_2)$

- by assumption
- (3) $\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$
- by assumption

By rule induction over Rules (18) on (1),

Case (18b).

- (4) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\top}(\xi_2))$ by assumption By rule induction over Rules (15) and only one case applies, Case (15h).
 - (5) $e_2 \models \dot{\top}(\xi_2)$

by assumption

(6) $e_2 \models_{?}^{\dagger} \dot{\top}(\xi_2)$

by Rule (18b) on (5)

(7) $e_2 \models_?^{\dagger} \xi_2$

by IH on (6)

 $(8) \ \operatorname{inr}_{\tau_1}(e_2) \models^\dagger_? \operatorname{inr}(\xi_2)$

by Lemma 2.0.12 on (7)

Case (18a).

- (4) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\top}(\xi_2))$ by assumption By rule induction over Rules (17) and only one case applies, Case (17f).
 - (5) $e_2 \models_? \dot{\top}(\xi_2)$

by assumption

(6) $e_2 \models^{\dagger}_? \dot{\top}(\xi_2)$

by Rule (18a) on (5)

 $(7) e_2 \models^{\dagger}_? \xi_2$

by IH on (6)

 $(8) \ \operatorname{inr}_{\tau_1}(e_2) \models_?^\dagger \operatorname{inr}(\xi_2)$

by Lemma 2.0.12 on (7)

Case $\xi = (\xi_1, \xi_2)$.

(2)
$$e = (e_1, e_2)$$

by assumption

(3)
$$\dot{\top}(\xi) = \dot{\top}(\xi_1) \wedge \dot{\top}(\xi_2)$$

by Definition 13

By rule induction over Rules (18) on (1),

Case (18b).

$$(4) (e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$$

by assumption

By rule induction over Rules (15) on (4) and only one case applies,

Case (15i).

(5) $e_1 \models \dot{\top}(\xi_1)$

by assumption

(6) $e_2 \models \dot{\top}(\xi_2)$

by assumption

(7) $e_1 \models_?^{\dagger} \dot{\top}(\xi_1)$

by Rule (18b) on (5)

(8) $e_2 \models_{?}^{\dagger} \dot{\top}(\xi_2)$

by Rule (18b) on (6)

(9) $e_1 \models_{?}^{\dagger} \xi_1$

by IH on (7)

 $(10) e_2 \models_{?}^{\dagger} \xi_2$

by IH on (8)

- (11) $(e_1, e_2) \models_{?}^{\dagger} (\xi_1, \xi_2)$
- by Lemma 2.0.13 on
- (9) and (10)

Case (18a).

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(4) (e_1, e_2) \models_? (\dot{\top}(\xi_1), \dot{\top}(\xi_2)) by assumption
By rule induction over Rules (17) on (4) and three cases apply,
Case (17g).
```

- $\begin{array}{lll} (5) & e_1 \models_? \dot{\top}(\xi_1) & \text{by assumption} \\ (6) & e_2 \models \dot{\top}(\xi_2) & \text{by assumption} \\ (7) & e_1 \models_?^\dagger \dot{\top}(\xi_1) & \text{by Rule (18a) on (5)} \\ (8) & e_2 \models_?^\dagger \dot{\top}(\xi_2) & \text{by Rule (18b) on (6)} \\ (9) & e_1 \models_?^\dagger \xi_1 & \text{by IH on (7)} \\ (10) & e_2 \models_?^\dagger \xi_2 & \text{by IH on (8)} \\ (11) & (e_1, e_2) \models_?^\dagger (\xi_1, \xi_2) & \text{by Lemma 2.0.13 on (9) and (10)} \\ \end{array}$
- Case (17h).
 - $\begin{array}{lll} (5) & e_1 \models \dot{\top}(\xi_1) & \text{by assumption} \\ (6) & e_2 \models_? \dot{\top}(\xi_2) & \text{by assumption} \\ (7) & e_1 \models_?^{\dagger} \dot{\top}(\xi_1) & \text{by Rule (18b) on (5)} \\ (8) & e_2 \models_?^{\dagger} \dot{\top}(\xi_2) & \text{by Rule (18a) on (6)} \\ (9) & e_1 \models_?^{\dagger} \xi_1 & \text{by IH on (7)} \\ (10) & e_2 \models_?^{\dagger} \xi_2 & \text{by IH on (8)} \\ (11) & (e_1, e_2) \models_?^{\dagger} (\xi_1, \xi_2) & \text{by Lemma 2.0.13 on (9) and (10)} \end{array}$

Case (17i).

(5) $e_1 \models_? \dot{\top}(\xi_1)$ by assumption (6) $e_2 \models_? \dot{\top}(\xi_2)$ by assumption (7) $e_1 \models_?^{\dagger} \dot{\top}(\xi_1)$ by Rule (18a) on (5) (8) $e_2 \models_?^{\dagger} \dot{\top}(\xi_2)$ by Rule (18a) on (6) (9) $e_1 \models_?^{\dagger} \xi_1$ by IH on (7) (10) $e_2 \models_?^{\dagger} \xi_2$ by IH on (8) (11) $(e_1, e_2) \models_?^{\dagger} (\xi_1, \xi_2)$ by Lemma 2.0.13 on (9) and (10)

Lemma 6.0.6. Assume $\dot{\top}(\xi) = \xi$. Then $\top \models^{\dagger}_{?} \xi$ iff $\top \models \xi$.

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
- 2. Necessity:

Theorem 6.1. $\top \models^{\dagger}_{?} \xi \text{ iff } \top \models \dot{\top}(\xi).$

Lemma 6.1.1. Assume that e val. Then $e \models_{?}^{\dagger} \xi$ iff $e \models \dot{\top}(\xi)$

Proof.

(1) e val

by assumption

We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (2) $e \models_{?}^{\dagger} \xi$

by assumption

By rule induction over Rules (18) on (2).

Case (18b).

(3) $e \models \xi$

by assumption

(4) $e \models \dot{\top}(\xi)$

by Lemma 6.0.3 on (3)

Case (18a).

(3) $e \models_? \xi$

by assumption

By rule induction over Rules (17) on (3).

Case (17a).

(4) $\xi = ?$

by assumption

(5) $e \models \dot{\top}(\xi)$

by Rule (15a) and

Definition 13

Case (17b).

(4) e notintro

by assumption

By rule induction over Rules (27) on (4), for each case, by rule induction over Rules (24) on (1), no case applies due to syntactic contradiction.

Case (17c).

(4) $\xi = \xi_1 \vee \xi_2$

by assumption

(5) $e \models_{?} \xi_{1}$

by assumption

(6) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$

by Equation 13

 $(7) \ e \models_{?}^{\dagger} \xi_{1}$

by Rule (18a) on (5)

(8) $e \models \dot{\top}(\xi_1)$

by IH on (7)

(9) $e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$

by Rule (15e) on (8)

Case (17d).

(4) $\xi = \xi_1 \vee \xi_2$

by assumption

(5) $e \models_? \xi_2$

by assumption

- (6) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$ by Equation 13
- (7) $e \models_{?}^{\dagger} \xi_{2}$ by Rule (18a) on (5)
- (8) $e \models \dot{\top}(\xi_2)$ by IH on (7)
- (9) $e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$ by Rule (15f) on (8)

Case (17e).

- (4) $\xi = \operatorname{inl}(\xi_1)$ by assumption
- (5) $e \models_? \xi_1$ by assumption (6) $\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$ by Equation 13
- (7) $e \models_{?}^{\dagger} \xi_{1}$ by Rule (18a) on (5)
- (8) $e \models \dot{\top}(\xi_1)$ by IH on (7)
- (9) $e \models \operatorname{inl}(\dot{\top}(\xi_1))$ by Rule (15g) on (8)

Case (17f).

- (4) $\xi = inr(\xi_2)$ by assumption
- (5) $e \models_{?} \xi_{2}$ by assumption (6) $\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$ by Equation 13
- (7) $e \models_{?}^{\dagger} \xi_{2}$ by Rule (18a) on (5)
- (8) $e \models \dot{\top}(\xi_2)$ by IH on (7)
- (9) $e \models \operatorname{inr}(\dot{\top}(\xi_2))$ by Rule (15h) on (8)

Case (17g).

- (4) $e = (e_1, e_2)$ by assumption
- (5) $\xi = (\xi_1, \xi_2)$ by assumption
- (6) $e_1 \models_? \xi_1$ by assumption (7) $e_2 \models \xi_2$ by assumption
- (8) $\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Equation 13
- (9) $e_1 \models_{?}^{\dagger} \xi_1$ by Rule (18a) on (6)
- (10) $e_2 \models_{?}^{\dagger} \xi_2$ by Rule (18b) on (7)
- (11) $e_1 \models \dot{\top}(\xi_1)$ by IH on (9)(12) $e_2 \models \dot{\top}(\xi_2)$
- by IH on (10)
- (13) $(e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Rule (15i) on (11) and (12)

Case (17h).

- (4) $e = (e_1, e_2)$ by assumption
- (5) $\xi = (\xi_1, \xi_2)$ by assumption
- (6) $e_1 \models \xi_1$ by assumption
- (7) $e_2 \models_? \xi_2$ by assumption (8) $\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Equation 13

$$\begin{array}{lll} (9) & e_1 \models_{?}^{\dagger} \xi_1 & \text{by Rule (18b) on (6)} \\ (10) & e_2 \models_{?}^{\dagger} \xi_2 & \text{by Rule (18a) on (7)} \\ (11) & e_1 \models \dot{\top}(\xi_1) & \text{by IH on (9)} \\ (12) & e_2 \models \dot{\top}(\xi_2) & \text{by IH on (10)} \\ (13) & (e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2)) & \text{by Rule (15i) on (11)} \\ & & \text{and (12)} \\ \end{array}$$

Case (17i).

(4)
$$e = (e_1, e_2)$$
 by assumption
(5) $\xi = (\xi_1, \xi_2)$ by assumption
(6) $e_1 \models_? \xi_1$ by assumption
(7) $e_2 \models_? \xi_2$ by assumption
(8) $\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Equation 13
(9) $e_1 \models_?^{\dagger} \xi_1$ by Rule (18a) on (6)
(10) $e_2 \models_?^{\dagger} \xi_2$ by Rule (18a) on (7)
(11) $e_1 \models \dot{\top}(\xi_1)$ by IH on (9)
(12) $e_2 \models \dot{\top}(\xi_2)$ by IH on (10)
(13) $(e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by Rule (15i) on (11) and (12)

2. Necessity:

(2)
$$e \models \dot{\top}(\xi)$$
 by assumption

By structural induction on ξ .

Case
$$\xi = \top, \bot, \underline{n}, \underline{\mathscr{M}}$$
.

(3)
$$\xi = \dot{\top}(\xi)$$
 by Equation 13
(4) $e \models_{?}^{\dagger} \xi$ by Rule (18b) on (2)

Case $\xi = ?$.

(3)
$$e \models_? ?$$
 by Rule (17a)
(4) $e \models_? † ?$ by Rule (18a) on (3)

Case $\xi = \xi_1 \wedge \xi_2$.

(3)
$$\dot{\top}(\xi) = \dot{\top}(\xi_1) \wedge \dot{\top}(\xi_2)$$
 by Equation 13

By rule induction over Rules (15) on (2), only one case applies.

Case (15d). (4) $e \models \dot{\top}(\xi_1)$

by assumption

(5)
$$e \models \dot{\top}(\xi_2)$$
 by assumption
(6) $e \models^{\dagger}_{?} \xi_1$ by IH on (4)

(7)
$$e \models_{?}^{\dagger} \xi_2$$
 by IH on (5)

(8)
$$e \models \xi_1 \land \xi_2$$
 by Lemma 2.0.9 on (6) and (7)

Case $\xi = \xi_1 \vee \xi_2$.

(3)
$$\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$$
 by Equation 13

By rule induction over Rules (15) on (2) and only two cases apply. Case (15e).

(4)
$$e \models \dot{\top}(\xi_1)$$
 by assumption

(5)
$$e \models_{?}^{\dagger} \xi_1$$
 by IH on (4)

(6)
$$e \models^{\dagger}_{?} \xi_1 \lor \xi_2$$
 by Lemma 2.0.10 on (5)

Case (15f).

(4)
$$e \models \dot{\top}(\xi_2)$$
 by assumption
(5) $e \models_2^{\dagger} \xi_2$ by IH on (4)

(5)
$$e \models_{?}^{\uparrow} \xi_{2}$$
 by IH on (4)
(6) $e \models_{?}^{\uparrow} \xi_{1} \lor \xi_{2}$ by Lemma 2.0.10 on (5)

Case $\xi = \text{inl}(\xi_1)$.

(3)
$$\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$$
 by Equation 13

By rule induction over Rules (15) on (2) and only one case applies. Case (15g).

(4)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption
(5) $e_1 \models \dot{\top}(\xi_1)$ by assumption

(6)
$$e_1 \models_{?}^{\dagger} \xi_1$$
 by IH on (5)

(7)
$$\operatorname{inl}_{\tau_2}(e_1) \models_?^{\dagger} \operatorname{inl}(\xi_1)$$
 by Lemma 2.0.11 on (6)

Case $\xi = inr(\xi_2)$.

(3)
$$\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$$
 by Equation 13

By rule induction over Rules (15) on (2) and only one case applies. Case (15h).

(4)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption
(5) $e_2 \models \dot{\top}(\xi_2)$ by assumption

(6)
$$e_2 \models_7^{\dagger} \xi_2$$
 by IH on (5)

(7)
$$\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inr}(\xi_2)$$
 by Lemma 2.0.12 on (6)

Case $\xi = (\xi_1, \xi_2)$.

(3)
$$\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$$

by Equation 13

By rule induction over Rules (15) on (2) and only one case applies.

Case (15i).

(4)
$$e = (e_1, e_2)$$
 by assumption
(5) $e_1 \models \dot{\bot}(\xi_1)$ by assumption
(6) $e_2 \models \dot{\bot}(\xi_2)$ by assumption
(7) $e_1 \models_7^{\dagger} \xi_1$ by IH on (5)
(8) $e_2 \models_7^{\dagger} \xi_2$ by IH on (6)

(9)
$$(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$$

by Lemma 2.0.13 on

(7) and (8)

Lemma 6.1.2. $e \models \xi \text{ iff } e \models \dot{\bot}(\xi)$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e \models \xi$$

by assumption

By rule induction over Rules (15) on (1).

Case (15a).

(2)
$$\xi = \top$$
 by assumption
(3) $e \models \dot{\bot}(\top)$ by (1) and Definition

Case (15b).

(2)
$$\xi = \underline{n}$$
 by assumption
(3) $e \models \dot{\perp}(\underline{n})$ by (1) and Definition

Case (15c).

(2)
$$\xi = \underline{\varkappa}$$
 by assumption
(3) $e \models \dot{\bot}(\underline{\varkappa})$ by (1) and Definition

Case (15d).

(2)
$$\xi = \xi_1 \wedge \xi_2$$
 by assumption
(3) $e \models \xi_1$ by assumption
(4) $e \models \xi_2$ by assumption
(5) $e \models \dot{\perp}(\xi_1)$ by IH on (3)

- (6) $e \models \dot{\perp}(\xi_2)$
- (7) $e \models \dot{\perp}(\xi_1) \land \dot{\perp}(\xi_2)$
- (8) $e \models \dot{\perp}(\xi_1 \wedge \xi_2)$

- by IH on (4)
- by Rule (15d) on (5)
- and (6)
- by (7) and Definition 14

Case (15e).

- (2) $\xi = \xi_1 \vee \xi_2$
- (3) $e \models \xi_1$
- (4) $e \models \dot{\perp}(\xi_1)$
- (5) $e \models \dot{\perp}(\xi_1) \lor \dot{\perp}(\xi_2)$
- (6) $e \models \dot{\perp}(\xi_1 \vee \xi_2)$

- by assumption
- by assumption
- by IH on (3)
- by Rule (15e) on (4)
- by (5) and Definition 14

Case (15f).

- (2) $\xi = \xi_1 \vee \xi_2$
- (3) $e \models \xi_2$
- (4) $e \models \dot{\perp}(\xi_2)$
- (5) $e \models \dot{\perp}(\xi_1) \lor \dot{\perp}(\xi_2)$
- (6) $e \models \dot{\perp}(\xi_1 \lor \xi_2)$

- by assumption
- by assumption
- by IH on (3)
- by Rule (15f) on (4)
- by (5) and Definition 14

Case (15g).

- $(2) e = \operatorname{inl}_{\tau_2}(e_1)$
- (3) $\xi = \operatorname{inl}(\xi_1)$
- (4) $e_1 \models \xi_1$
- (5) $e_1 \models \dot{\bot}(\xi_1)$
- (6) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\bot}(\xi_1))$
- (7) $\operatorname{inl}_{\tau_2}(e_1) \models \dot{\bot}(\operatorname{inl}(\xi_1))$
- by assumption
- by assumption
- by assumption
- by IH on (4)
- by Rule (15g) on (5)
- by (6) and Definition 14

Case (15h).

- (2) $e = inr_{\tau_1}(e_2)$
- (3) $\xi = inr(\xi_2)$
- (4) $e_2 \models \xi_2$
- (5) $e_2 \models \dot{\perp}(\xi_2)$
- (6) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\perp}(\xi_2))$
- (7) $\operatorname{inr}_{\tau_1}(e_2) \models \dot{\perp}(\operatorname{inr}(\xi_2))$
- by assumption
- by assumption
- by assumption
- by IH on (4)
- by Rule (15h) on (5)
- by (6) and Definition 14

Case (15i).

(2) $e = (e_1, e_2)$

by assumption

(3) $\xi = (\xi_1, \xi_2)$	by assumption
$(4) e_1 \models \xi_1$	by assumption
$(5) e_2 \models \xi_2$	by assumption
(6) $e_1 \models \dot{\perp}(\xi_1)$	by IH on (4)
$(7) \ e_2 \models \dot{\bot}(\xi_2)$	by IH on (5)
(8) $(e_1, e_2) \models (\dot{\bot}(\xi_1), \dot{\bot}(\xi_2))$	by Rule $(15i)$ on (6)
	and (7)
(9) $(e_1, e_2) \models \dot{\bot}((\xi_1, \xi_2))$	by (8) and Definition
	14

2. Necessity:

(1)
$$e \models \dot{\perp}(\xi)$$
 by assumption

By structural induction on ξ .

Case $\xi = ?$.

(2)
$$e \models \bot$$
 by (1) and Definition 14 by Lemma 2.0.1

(3) contradicts (2).

Case $\xi = \xi_1 \wedge \xi_2$.

(2)
$$e \models \dot{\perp}(\xi_1) \land \dot{\perp}(\xi_2)$$
 by (1) and Definition 14

By rule induction over Rules (15) on (2) and only case applies. Case (15d).

(3)
$$e \models \dot{\bot}(\xi_1)$$
 by assumption
(4) $e \models \dot{\bot}(\xi_2)$ by assumption
(5) $e \models \xi_1$ by IH on (3)
(6) $e \models \xi_2$ by IH on (4)
(7) $e \models \xi_1 \land \xi_2$ by Rule (15d) on (5)
and (6)

Case $\xi = \xi_1 \vee \xi_2$.

(2)
$$e \models \dot{\bot}(\xi_1) \lor \dot{\bot}(\xi_2)$$
 by (1) and Definition 14

By rule induction over Rules (15) on (2) and only two cases apply. Case (15e).

(3)
$$e \models \dot{\bot}(\xi_1)$$
 by assumption
(4) $e \models \xi_1$ by IH on (3)
(5) $e \models \xi_1 \lor \xi_2$ by Rule (15e) on (4)

Case (15f).

(3)
$$e \models \dot{\bot}(\xi_2)$$
 by assumption
(4) $e \models \xi_2$ by IH on (3)
(5) $e \models \xi_1 \lor \xi_2$ by Rule (15f) on (4)

Case $\xi = inl(\xi_1)$.

(2)
$$e \models \operatorname{inl}(\dot{\perp}(\xi_1))$$
 by (1) and Definition 14

By rule induction over Rules (15) on (2) and only one case applies. Case (15g).

(3)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption
(4) $e_1 \models \dot{\bot}(\xi_1)$ by assumption
(5) $e_1 \models \xi_1$ by IH on (4)
(6) $e \models \operatorname{inl}(\xi_1)$ by Rule (15g) on (5)

Case $\xi = inr(\xi_2)$.

(2)
$$e \models \operatorname{inr}(\dot{\perp}(\xi_2))$$
 by (1) and Definition

By rule induction over Rules (15) on (2) and only one case applies. Case (15h).

(3)
$$e = \inf_{\tau_1}(e_2)$$
 by assumption
(4) $e_2 \models \dot{\bot}(\xi_2)$ by assumption
(5) $e_2 \models \xi_2$ by IH on (4)
(6) $e \models \inf(\xi_2)$ by Rule (15h) on (5)

Case $\xi = (\xi_1, \xi_2)$.

(2)
$$e \models (\dot{\perp}(\xi_1), \dot{\perp}(\xi_2))$$
 by (1) and Definition

By rule induction over Rules (15) on (2) and only case applies. Case (15i).

(3) $e = (e_1, e_2)$	by assumption
$(4) e_1 \models \dot{\bot}(\xi_1)$	by assumption
$(5) e_2 \models \dot{\bot}(\xi_2)$	by assumption
(6) $e_1 \models \xi_1$	by IH on (4)
$(7) e_2 \models \xi_2$	by IH on (5)
(8) $e \models (\xi_1, \xi_2)$	by Rule (15i) on (6)
	and (7)

 $\textbf{Lemma 6.1.3.} \ \textit{Assume e val and $\dot{\top}(\xi)=\xi$. Then $e\not\models\xi$ iff $e\models\overline{\xi}$.}$

Theorem 6.2. $\xi_r \models \xi_{rs} \ \textit{iff} \ \top \models \overline{\dot{\top}(\xi_r)} \lor \dot{\bot}(\xi_{rs}).$