1 Match Constraint Language

 $\overline{\top} : \tau$ (1a)

CTUnknown

$$\overline{?:\tau}$$
 (1b)

 $\frac{\text{CTNum}}{\underline{n}: \text{num}} \tag{1c}$

CTInl $\frac{\dot{\xi}_1 : \tau_1}{\operatorname{inl}(\dot{\xi}_1) : (\tau_1 + \tau_2)}$ (1d)

CTInr $\frac{\dot{\xi}_2 : \tau_2}{\operatorname{inr}(\dot{\xi}_2) : (\tau_1 + \tau_2)} \tag{1e}$

CTPair $\frac{\dot{\xi}_1 : \tau_1 \qquad \dot{\xi}_2 : \tau_2}{(\dot{\xi}_1, \dot{\xi}_2) : (\tau_1 \times \tau_2)}$ (1f)

CTOr $\frac{\dot{\xi}_1 : \tau \qquad \dot{\xi}_2 : \tau}{\dot{\xi}_1 \lor \dot{\xi}_2 : \tau}$ (1g)

 $|\dot{\xi}|$ refutable? $|\dot{\xi}|$ is refutable

RXNum

$$\frac{}{\underline{n} \; \mathtt{refutable}_?}$$
 (2a)

RXUnknown

RXInl

$$\frac{}{\operatorname{inl}(\dot{\xi})\operatorname{refutable}_?}$$
 (2c)

RXInr

$$\frac{}{\operatorname{inr}(\dot{\xi})\operatorname{refutable}_?}$$
 (2d)

RXPairL

$$\frac{\dot{\xi}_1 \text{ refutable}?}{(\dot{\xi}_1, \dot{\xi}_2) \text{ refutable}?} \tag{2e}$$

RXPairR
$$\frac{\dot{\xi}_2 \text{ refutable}_?}{(\dot{\xi}_1, \dot{\xi}_2) \text{ refutable}_?}$$
 (2f)

 $\frac{\text{RXOr}}{\dot{\xi}_1 \text{ refutable}_?} \frac{\dot{\xi}_2 \text{ refutable}_?}{\dot{\xi}_1 \vee \dot{\xi}_2 \text{ refutable}_?} \tag{2g}$

$\mathit{refutable}_?(\dot{\xi})$

$$refutable_{?}(\underline{n}) = true$$
 (3a)

$$refutable_{?}(?) = true$$
 (3b)

$$refutable_{?}(inl(\dot{\xi})) = refutable_{?}(\dot{\xi})$$
 (3c)

$$refutable_{?}(inr(\dot{\xi})) = refutable_{?}(\dot{\xi})$$
 (3d)

$$refutable_{?}((\dot{\xi}_{1}, \dot{\xi}_{2})) = refutable_{?}(\dot{\xi}_{1}) \text{ or } refutable_{?}(\dot{\xi}_{2})$$
 (3e)

$$refutable_{?}(\dot{\xi}_{1} \vee \dot{\xi}_{2}) = refutable_{?}(\dot{\xi}_{1}) \text{ and } refutable_{?}(\dot{\xi}_{2})$$
 (3f)

Otherwise
$$refutable_{?}(\dot{\xi}) = false$$
 (3g)

 $e \models \dot{\xi}$ $e \text{ satisfies } \dot{\xi}$

$$\frac{\text{CSTruth}}{e \models \top} \tag{4a}$$

CSNum

$$\underline{\underline{n} \models \underline{n}} \tag{4b}$$

CSInl

$$\frac{e_1 \models \dot{\xi}_1}{\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)} \tag{4c}$$

CSInr

$$\frac{e_2 \models \dot{\xi}_2}{\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)} \tag{4d}$$

CSPair

$$\begin{array}{ll}
e_1 \models \dot{\xi}_1 & e_2 \models \dot{\xi}_2 \\
\hline
(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)
\end{array} \tag{4e}$$

CSNotIntroPair

$$\frac{e \text{ notintro}}{e \text{ prl}(e) \models \dot{\xi}_1 \qquad \text{prr}(e) \models \dot{\xi}_2}{e \models (\dot{\xi}_1, \dot{\xi}_2)} \tag{4f}$$

CSOrL
$$\frac{e \models \dot{\xi}_1}{e \models \dot{\xi}_1 \lor \dot{\xi}_2} \tag{4g}$$

CSOrR
$$\frac{e \models \dot{\xi}_2}{e \models \dot{\xi}_1 \lor \dot{\xi}_2}$$
 (4h)

 $satisfy(e,\dot{\xi})$

$$satisfy(e, \top) = true$$
 (5a)

$$satisfy(n_1, n_2) = (n_1 = n_2)$$
 (5b)

$$satisfy(e, \dot{\xi}_1 \lor \dot{\xi}_2) = satisfy(e, \dot{\xi}_1) \text{ or } satisfy(e, \dot{\xi}_2)$$
 (5c)

$$\mathit{satisfy}(\mathtt{inl}_{\tau_2}(e_1),\mathtt{inl}(\dot{\xi_1})) = \mathit{satisfy}(e_1,\dot{\xi_1}) \tag{5d}$$

$$satisfy(\operatorname{inr}_{\tau_1}(e_2), \operatorname{inr}(\dot{\xi}_2)) = satisfy(e_2, \dot{\xi}_2)$$
(5e)

$$satisfy((e_1, e_2), (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(e_1, \dot{\xi}_1) \text{ and } satisfy(e_2, \dot{\xi}_2)$$
 (5f)

$$\mathit{satisfy}(\lVert \rVert^u, (\dot{\xi}_1, \dot{\xi}_2)) = \mathit{satisfy}(\mathtt{prl}(\lVert \rVert^u), \dot{\xi}_1) \text{ and } \mathit{satisfy}(\mathtt{prr}(\lVert \rVert^u), \dot{\xi}_2) \tag{5g}$$

$$satisfy(\langle e \rangle^u, (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(\operatorname{prl}(\langle e \rangle^u), \dot{\xi}_1) \text{ and } satisfy(\operatorname{prr}(\langle e \rangle^u), \dot{\xi}_2)$$

$$(5h)$$

$$satisfy(e_1(e_2), (\dot{\xi_1}, \dot{\xi_2})) = satisfy(\mathtt{prl}(e_1(e_2)), \dot{\xi_1})$$

and
$$satisfy(prr(e_1(e_2)), \dot{\xi}_2)$$
 (5i)

 $\mathit{satisfy}(\texttt{match}(e)\{\hat{rs}\},(\dot{\xi_1},\dot{\xi_2})) = \mathit{satisfy}(\texttt{prl}(\texttt{match}(e)\{\hat{rs}\}),\dot{\xi_1})$

and
$$satisfy(prr(match(e)\{\hat{rs}\}), \dot{\xi}_2)$$
 (5j)

 $satisfy(\mathtt{prl}(e), (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(\mathtt{prl}(\mathtt{prl}(e)), \dot{\xi}_1)$

and
$$satisfy(prr(prl(e)), \dot{\xi}_2)$$
 (5k)

 $satisfy(\mathtt{prr}(e), (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(\mathtt{prl}(\mathtt{prr}(e)), \dot{\xi}_1)$

and
$$satisfy(prr(prr(e)), \dot{\xi}_2)$$
 (51)

Otherwise
$$satisfy(e, \dot{\xi}) = false$$
 (5m)

 $e \models_? \dot{\xi} \mid e \text{ may satisfy } \dot{\xi}$

CMSUnknown

$$\overline{e \models_?}$$
? (6a)

CMSInl

$$\frac{e_1 \models_? \dot{\xi}_1}{\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)} \tag{6b}$$

CMSInr

$$\frac{e_2 \models_? \dot{\xi}_2}{\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)} \tag{6c}$$

CMSPairL

$$\frac{e_1 \models_? \dot{\xi}_1 \qquad e_2 \models_{\dot{\xi}_2}}{(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)} \tag{6d}$$

CMSPairR
$$\underbrace{e_1 \models \dot{\xi}_1 \qquad e_2 \models_? \dot{\xi}_2}_{(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)} \tag{6e}$$

CMSPair
$$\frac{e_1 \models_? \dot{\xi}_1 \qquad e_2 \models_? \dot{\xi}_2}{(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)}$$
(6f)

CMSOrL
$$\frac{e \models_? \dot{\xi}_1 \quad e \not\models \dot{\xi}_2}{e \models_? \dot{\xi}_1 \lor \dot{\xi}_2} \tag{6g}$$

CMSOrR
$$\frac{e \not\models \dot{\xi}_{1} \qquad e \models_{?} \dot{\xi}_{2}}{e \models_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}} \tag{6h}$$

CMSNotIntro $\frac{e \text{ notintro} \qquad \dot{\xi} \text{ refutable}_?}{e \models_? \dot{\xi}} \tag{6i}$

 $e \models_{?}^{\dagger} \dot{\xi}$ e satisfies or may satisfy $\dot{\xi}$

CSMSMay $\frac{e \models_? \dot{\xi}}{e \models_?^{\dagger} \dot{\xi}} \tag{7a}$

CSMSSat
$$\frac{e \models \dot{\xi}}{e \models_{?}^{\dagger} \dot{\xi}} \tag{7b}$$

Lemma 1.0.1. $e \not\models \bot$

Proof. By rule induction over Rules (14), we notice that $e \models \bot$ is in syntactic contradiction with all rules, hence not derivable.

Lemma 1.0.2. $e \not\models_? \bot$

Proof. Assume $e \models_? \bot$. By rule induction over Rules (16) on $e \models_? \bot$, only one case applies.

Case (16b).

(1) \perp refutable? by assumption

By rule induction over Rules (10) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \bot$ is not derivable.

Lemma 1.0.3. $e \not\models_? \top$

Proof. Assume $e \models_? \top$. By rule induction over Rules (16) on $e \models_? \top$, only one case applies.

Case (16b).

(1)
$$\top$$
 refutable?

by assumption

By rule induction over Rules (10) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \top$ is not derivable.

Lemma 1.0.4. $e \not\models ?$

Proof. By rule induction over Rules (14), we notice that $e \models ?$ is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.5.
$$e \models_{?}^{\dagger} \dot{\xi} \text{ iff } e \models_{?}^{\dagger} \dot{\xi} \vee \bot$$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e \models_?^\dagger \dot{\xi}$$

by assumption

By rule induction over Rules (17) on (1).

Case (17a).

(2)
$$e \models_? \dot{\xi}$$
 by assumption
(3) $e \models_? \dot{\xi} \lor \bot$ by Rule (16c) on (2)
and Lemma 2.0.1
(4) $e \models_?^{\dagger} \dot{\xi} \lor \bot$ by Rule (17a) on (3)

by Rule (17a) on (3)

Case (17b).

(2)
$$e \models \dot{\xi}$$
 by assumption
(3) $e \models \dot{\xi} \lor \bot$ by Rule (14e) on (2)

(4) $e \models_{?}^{\dagger} \dot{\xi} \lor \bot$

by Rule (17b) on (3)

2. Necessity:

(1)
$$e \models_{2}^{\dagger} \dot{\xi} \lor \bot$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17a).

(2)
$$e \models_? \dot{\xi} \lor \bot$$
 by assumption

By rule induction over Rules (16) on (2), only two of them apply.

Case (16c).

(3)
$$e \models_? \dot{\xi}$$

by assumption

(4)
$$e \models_2^{\dagger} \dot{\xi}$$

by Rule (17a) on (3)

Case (16d).

(3)
$$e \models_? \bot$$

by assumption

$$(4) e \not\models_? \bot$$

by Lemma 2.0.2

Case (17b).

(2)
$$e \models \dot{\xi} \lor \bot$$

by assumption

By rule induction over Rules (14) on (2), only two of them apply.

Case (14e).

(3)
$$e \models \dot{\xi}$$

by assumption

(4)
$$e \models_{?}^{\dagger} \dot{\xi}$$

by Rule (17b) on (3)

Case (14f).

(3)
$$e \models \bot$$

by assumption

(4)
$$e \not\models \bot$$

by Lemma 2.0.1

Corollary 1.0.1. $\top \models_2^{\dagger} \dot{\xi} \text{ iff } \top \models_2^{\dagger} \dot{\xi} \lor \bot$

Proof. Follows directly from Definition 2.1.2 and Lemma 2.0.5.

Lemma 1.0.6. Suppose that $\dot{\xi}_1 : \tau$ and $\dot{\xi}_2 : \tau$. Then $\dot{\xi}_1 \not\models \dot{\xi}_2$ iff $\dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$ Proof.

 $(1) \ \dot{\xi}_1:\tau$

by assumption

(2) $\dot{\xi}_2 : \tau$

by assumption

(3) $\perp : \tau$

by Rule (8b)

(4) $\dot{\xi}_2 \vee \bot : \tau$

by Rule (8f) on (2) and (3)

Then we prove sufficiency and necessity separately.

- 1. Sufficiency:
 - $(5) \ \dot{\xi}_1 \not\models \dot{\xi}_2$

by assumption

To prove $\dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$, assume $\dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$.

(6)
$$\dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \dot{\xi}_{1}$ implies

(7) $e \models \dot{\xi}_2 \lor \bot$

- by Definition 2.1.1 on
- (1) and (4) and (6)

By rule induction over Rules (14) on (7).

Case (14e).

(8) $e \models \dot{\xi}_2$ (9) $\dot{\xi}_1 \models \dot{\xi}_2$ by assumption

by Definition 2.1.1 on (8)

(5) contradicts (9).

Case (14f).

(8) $e \models \bot$

by assumption

(9) $e \not\models \bot$

by Lemma 2.0.1

(8) contradicts (9).

The conclusion holds as follows:

- (a) $\dot{\xi}_1 \not\models \dot{\xi}_2 \vee \bot$
- 2. Necessity:
 - $(5) \ \dot{\xi}_1 \not\models \dot{\xi}_2 \lor \bot$

by assumption

To prove $\dot{\xi}_1 \not\models \dot{\xi}_2$, assume $\dot{\xi}_1 \models \dot{\xi}_2$.

(6) $\dot{\xi}_1 \models \dot{\xi}_2$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \dot{\xi}_{1}$ implies

(7) $e \models \dot{\xi}_2$

by Definition 2.1.1 on

(1) and (2) and (6)

(8) $e \models \dot{\xi}_2 \lor \bot$

by Rule (14e) on (7)

 $(9) \ \dot{\xi}_1 \models \dot{\xi}_2 \lor \bot$

by Definition 2.1.1 on

(8)

(9) contradicts (5).

The conclusion holds as follows:

(a) $\dot{\xi}_1 \not\models \dot{\xi}_2$

Lemma 1.0.7. $e \not\models_{?}^{\dagger} \dot{\xi}_{1}$ and $e \not\models_{?}^{\dagger} \dot{\xi}_{2}$ iff $e \not\models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency: to show $e \not\models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$, we assume $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$.
 - (1) $e \models^{\dagger}_{?} \dot{\xi}_1 \lor \dot{\xi}_2$

by assumption

(2) $e \not\models_?^\dagger \dot{\xi}_1$

by assumption

(3) $e \not\models_?^\dagger \dot{\xi}_2$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

$$(4) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$$

by assumption

By rule induction over Rules (14) on (4) and only two of them apply.

Case (14e).

(5) $e \models \dot{\xi}_1$

by assumption

(6) $e \models_{?}^{\dagger} \dot{\xi}_{1}$

by Rule (17b) on (5)

(6) contradicts (2).

Case (14f).

(5) $e \models \dot{\xi}_2$

by assumption

(6) $e \models^{\dagger}_{?} \dot{\xi}_{2}$

by Rule (17b) on (5)

(6) contradicts (3).

Case (17a).

$$(4) e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$$

by assumption

By rule induction over Rules (16) on (4) and only two of them apply.

Case (16c).

(5)
$$e \models_? \dot{\xi}_1$$

by assumption

(6)
$$e \models_{?}^{\dagger} \dot{\xi}_{1}$$

by Rule (17a) on (5)

(6) contradicts (2).

Case (16d).

(5)
$$e \models_? \dot{\xi}_2$$

by assumption

(6)
$$e \models_{?}^{\dagger} \dot{\xi}_2$$

by Rule (17a) on (5)

(6) contradicts (3).

The conclusion holds as follows:

- (a) $e \not\models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$
- 2. Necessity:

(1)
$$e \not\models_2^\dagger \dot{\xi_1} \lor \dot{\xi_2}$$

by assumption

We show $e \not\models_{2}^{\dagger} \dot{\xi}_{1}$ and $e \not\models_{2}^{\dagger} \dot{\xi}_{2}$ separately.

- (a) To show $e \not\models_?^\dagger \dot{\xi}_1$, we assume $e \models_?^\dagger \dot{\xi}_1$.
 - (2) $e \models_{?}^{\dagger} \dot{\xi}_1$

by assumption

(3) $e \models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$

by Lemma 2.0.10 on

(2)

Contradicts (1).

- (b) To show $e \not\models_{?}^{\dagger} \dot{\xi}_{2}$, we assume $e \models_{?}^{\dagger} \dot{\xi}_{2}$.
 - (2) $e \models^{\dagger}_{?} \dot{\xi}_{2}$

by assumption

by Lemma 2.0.10 on

(2)

 $(3) \ e \models^{\dagger}_{?} \dot{\xi}_1 \lor \dot{\xi}_2$

In conclusion, $e \not\models_{?}^{\dagger} \dot{\xi}_{1}$ and $e \not\models_{?}^{\dagger} \dot{\xi}_{2}$.

Contradicts (1).

Lemma 1.0.8. If $e \models^{\dagger}_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ and $e \not\models^{\dagger}_{?} \dot{\xi}_{1}$ then $e \models^{\dagger}_{?} \dot{\xi}_{2}$

Proof.

(4)
$$e \models_?^\dagger \dot{\xi}_1 \lor \dot{\xi}_2$$

by assumption

(5)
$$e \not\models_?^\dagger \dot{\xi}_1$$

by assumption

By rule induction over Rules (17) on (4).

Case (17b).

(6)
$$e \models \dot{\xi}_1 \lor \dot{\xi}_2$$

by assumption

By rule induction over Rules (14) on (6) and only two of them apply.

Case (14e).

(7)
$$e \models \dot{\xi}_1$$

by assumption

(8)
$$e \models^{\dagger}_{?} \dot{\xi}_{1}$$

by Rule (17b) on (7)

(8) contradicts (5).

Case (14f).

(7)
$$e \models \dot{\xi}_2$$

by assumption

(8)
$$e \models^{\dagger}_{?} \dot{\xi}_{2}$$

by Rule (17b) on (7)

Case (17a).

(6)
$$e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$$

by assumption

By rule induction over Rules (16) on (6) and only two of them apply.

Case (16c).

- (7) $e \models_? \dot{\xi}_1$ by assumption (8) $e \models_{?}^{\dagger} \dot{\xi}_1$ by Rule (17a) on (7)
- (8) contradicts (5).

Case (16d).

(7) $e \models_? \dot{\xi}_2$

by assumption

(8) $e \models_{?}^{\dagger} \dot{\xi}_2$

by Rule (17a) on (7)

Lemma 1.0.9. If $e \models^{\dagger}_{?} \dot{\xi}_{1}$ then $e \models^{\dagger}_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ and $e \models^{\dagger}_{?} \dot{\xi}_{2} \lor \dot{\xi}_{1}$

Proof.

(1)
$$e \models^{\dagger}_{?} \dot{\xi}_{1}$$

by assumption

By rule induction over Rules (17) on (1),

Case (17b).

(2)
$$e \models \dot{\xi}_1$$

by assumption

$$(3) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$$

by Rule (14e) on (2)

$$(4) \ e \models \dot{\xi}_2 \lor \dot{\xi}_1$$

by Rule (14f) on (2)

(5)
$$e \models^{\dagger}_{?} \dot{\xi}_1 \lor \dot{\xi}_2$$

by Rule (17b) on (3)

(6)
$$e \models_2^{\dagger} \dot{\xi}_2 \lor \dot{\xi}_1$$

by Rule (17b) on (4)

Case (17a).

(2)
$$e \models_? \dot{\xi}_1$$

by assumption

By case analysis on the result of $satisfy(e, \dot{\xi}_2)$.

Case true.

(3)
$$satisfy(e, \dot{\xi}_2) = true$$

by assumption

$$(4) e \models \dot{\xi}_2$$

by Lemma 2.0.19 on

(5)
$$e \models \dot{\xi}_1 \lor \dot{\xi}_2$$

(6) $e \models \dot{\xi}_2 \lor \dot{\xi}_1$

by Rule (14f) on (4)

(6)
$$e \models \xi_2 \vee \xi_1$$

by Rule (14e) on (4)

(7)
$$e \models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$$
 by Rule (17b) on (5)
(8) $e \models_{?}^{\dagger} \dot{\xi}_{2} \lor \dot{\xi}_{1}$ by Rule (17b) on (6)

Case false.

(3)
$$satisfy(e, \dot{\xi}_2) = false$$
 by assumption
(4) $e \not\models \dot{\xi}_2$ by Lemma 2.0.19 on
(3)
(5) $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (16c) on (2)
and (4)
(6) $e \models_?^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (17a) on (5)

Lemma 1.0.10. $e_1\models^\dagger_?\dot{\xi_1}$ iff $\mathrm{inl}_{\tau_2}(e_1)\models^\dagger_?\mathrm{inl}(\dot{\xi_1})$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e_1 \models_{?}^{\dagger} \dot{\xi}_1$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$e_1 \models \dot{\xi}_1$$
 by assumption
(3) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$ by Rule (14g) on (2)
(4) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ by Rule (17b) on (3)

Case (17a).

$$\begin{array}{lll} (2) & e_1 \models_? \dot{\xi}_1 & \text{by assumption} \\ (3) & \inf_{\tau_2}(e_1) \models_? \inf(\dot{\xi}_1) & \text{by Rule (16e) on (2)} \\ (4) & \inf_{\tau_2}(e_1) \models_?^{\dagger} \inf(\dot{\xi}_1) & \text{by Rule (17a) on (3)} \end{array}$$

2. Necessity:

(1)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$$
 by assumption
By rule induction over Rules (14) on (2), only one rule applies.

Case (14g).

(3)
$$e_1 \models \dot{\xi}_1$$
 by assumption
(4) $e_1 \models_{7}^{\dagger} \dot{\xi}_1$ by Rule (17b) on (3)

Case (17a).

(2)
$$\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi_1})$$
 by assumption

By rule induction over Rules (16) on (2), only two rules apply.

Case (16e).

(3)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption
(4) $e_1 \models_?^{\dagger} \dot{\xi}_1$ by Rule (17a) on (3)

Case (16b).

(3)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro by assumption

By rule induction over Rules (26) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.11. $e_2 \models_{?}^{\dagger} \dot{\xi}_2$ iff $\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e_2 \models^{\dagger}_{?} \dot{\xi}_2$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$e_2 \models \dot{\xi}_2$$
 by assumption
(3) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$ by Rule (14h) on (2)
(4) $\operatorname{inr}_{\tau_1}(e_2) \models_2^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (17b) on (3)

Case (17a).

(2)
$$e_2 \models_? \dot{\xi}_2$$
 by assumption
(3) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$ by Rule (16f) on (2)
(4) $\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ by Rule (17a) on (3)

2. Necessity:

(1)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\dot{\xi_2})$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi_2})$$

by assumption

By rule induction over Rules (14) on (2), only one rule applies.

Case (14h).

$$(3) e_2 \models \dot{\xi}_2$$

by assumption

(4)
$$e_2 \models_{?}^{\dagger} \dot{\xi}_2$$

by Rule (17b) on (3)

Case (17a).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi_2})$$

by assumption

By rule induction over Rules (16) on (2), only two rules apply.

Case (16f).

(3)
$$e_2 \models_? \dot{\xi}_2$$

by assumption

(4)
$$e_2 \models_?^{\dagger} \dot{\xi}_2$$

by Rule (17a) on (3)

Case (16b).

(3)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (26) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.12. $e_1 \models_{?}^{\dagger} \dot{\xi_1} \text{ and } e_2 \models_{?}^{\dagger} \dot{\xi_2} \text{ iff } (e_1, e_2) \models_{?}^{\dagger} (\dot{\xi_1}, \dot{\xi_2})$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e_1 \models^{\dagger}_{?} \dot{\xi}_1$$

by assumption

(2)
$$e_2 \models_?^{\dagger} \dot{\xi}_2$$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

$$(3) e_1 \models \dot{\xi}_1$$

by assumption

By rule induction over Rules (17) on (2).

Case (17b).

(4)
$$e_2 \models \dot{\xi}_2$$

by assumption

(5)
$$(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (14i) on (3)

and
$$(4)$$

(6)
$$(e_1, e_2) \models^{\dagger}_{?} (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (17b) on (5)

Case (17a).

(4)
$$e_2 \models_? \dot{\xi}_2$$

by assumption

(5)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$

by Rule (16h) on (3)

and (4)

(6)
$$(e_1, e_2) \models^{\dagger}_{?} (\dot{\xi}_1, \dot{\xi}_2)$$
 by Rule (17a) on (5)

Case (17a).

(4)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption

By rule induction over Rules (17) on (2).

Case (17b).

- (5) $e_2 \models \dot{\xi}_2$ by assumption (6) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (16g) on (4) and (5)
- (7) $(e_1, e_2) \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17a) on (6)

Case (17a).

- (5) $e_2 \models_? \dot{\xi}_2$ by assumption (6) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (16h) on (4) and (5) (7) $(e_1, e_2) \models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17a) on (6)
- 2. Necessity:

(1)
$$(e_1, e_2) \models_7^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption

By rule induction over Rules (14) on (2), only one rule applies.

Case (14i).

$$(3) e_1 \models \dot{\xi}_1 \qquad \qquad \text{by assumption}$$

$$(4) e_2 \models \dot{\xi}_2 \qquad \qquad \text{by assumption}$$

$$(5) e_1 \models_{?}^{\dagger} \dot{\xi}_1 \qquad \qquad \text{by Rule (17b) on (3)}$$

$$(6) e_2 \models_{?}^{\dagger} \dot{\xi}_2 \qquad \qquad \text{by Rule (17b) on (4)}$$

Case (17a).

(2)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption

By rule induction over Rules (16) on (2), only three rules apply. Case (16g).

(3)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption
(4) $e_2 \models \dot{\xi}_2$ by assumption
(5) $e_1 \models_?^{\dagger} \dot{\xi}_1$ by Rule (17a) on (3)
(6) $e_2 \models_?^{\dagger} \dot{\xi}_2$ by Rule (17b) on (4)

Case (16h).

$(3) e_1 \models \dot{\xi}_1$	by assumption
$(4) e_2 \models_? \dot{\xi}_2$	by assumption
$(5) e_1 \models^{\dagger}_{?} \dot{\xi}_1$	by Rule $(17b)$ on (3)
$(6) e_2 \models^{\dagger}_{?} \dot{\xi}_2$	by Rule (17a) on (4)
Case (16i).	
$(3) e_1 \models_? \dot{\xi}_1$	by assumption
$(4) e_2 \models_? \dot{\xi}_2$	by assumption
$(5) e_1 \models^{\dagger}_{?} \dot{\xi_1}$	by Rule $(17a)$ on (3)
$(6) e_2 \models^{\dagger}_{?} \dot{\xi}_2$	by Rule $(17a)$ on (4)
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Lemma 1.0.13 (Soundness and Completeness of Refutable Constraints). $\dot{\xi}$ refutable? iff refutable? $\dot{\xi}$ = true.

Lemma 1.0.14. There does not exist such a constraint $\dot{\xi}_1 \wedge \dot{\xi}_2$ such that $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable?

Proof. By rule induction over Rules (10), we notice that $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.15. There does not exist such a constraint $\dot{\xi}_1 \lor \dot{\xi}_2$ such that $\dot{\xi}_1 \lor \dot{\xi}_2$ refutable?

Proof. By rule induction over Rules (10), we notice that $\dot{\xi}_1 \vee \dot{\xi}_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 1.0.16. If e notintro and $e \models \dot{\xi}$ then $\dot{\xi}$ refutable?.

Proof.

 $(1)\ e\ {\tt notintro}$

by assumption

(2) $e \models \xi$

by assumption

By rule induction over Rules (14) on (2).

Case (14a).

(3)
$$\dot{\xi} = \top$$

by assumption

Assume \top refutable?. By rule induction over Rules (10), no case applies due to syntactic contradiction.

Therefore, <u>⊤refutable</u>?.

Case (14e),(14f).

 $(3) \ \dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$

by assumption

(4) $\dot{\xi_1} \vee \dot{\xi_2}$ refutable?

by Lemma 2.0.17

Case (14d).

(3)
$$\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$$
 by assumption
(4) $\dot{\xi}_1 \wedge \dot{\xi}_2$ refutable? by Lemma 2.0.16

Case (14j).

(3)
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption
(4) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption
(5) $\operatorname{prr}(e) \models \dot{\xi}_2$ by assumption
(6) $\operatorname{prl}(e)$ notintro by Rule (26e)
(7) $\operatorname{prr}(e)$ notintro by Rule (26f)
(8) $\dot{\xi}_1$ refutable? by IH on (6) and (4)
(9) $\dot{\xi}_2$ refutable? by IH on (7) and (5)

Assume $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? By rule induction over Rules (10) on it, only two cases apply.

Case (10d).

(10)
$$\dot{\xi}_1$$
 refutable? by assumption

Contradicts (8).

Case (10e).

(10)
$$\dot{\xi}_2$$
 refutable? by assumption

Contradicts (9).

Therefore, $(\dot{\xi}_1, \dot{\xi}_2)$ refutable?.

Otherwise.

(3)
$$e = \underline{n}, \text{inl}_{\tau_2}(e_1), \text{inr}_{\tau_1}(e_2), (e_1, e_2)$$
 by assumption

By rule induction over Rules (26) on (1), no case applies due to syntactic contradiction.

Lemma 1.0.17. $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi_2})$ is not derivable.

Proof. We prove by assuming $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$ and obtaining a contradiction.

(1)
$$\operatorname{inl}_{72}(e_1) \models_2^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 by assumption

By rule induction over Rules (17) on (1).

Case (17b).

$$(2) \ \operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\dot{\xi}_2)$$

by assumption

By rule induction over Rules (14) on (2), no rule applies due to syntactic contradiction.

Case (17a).

(2)
$$\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\dot{\xi}_2)$$

by assumption

By rule induction over Rules (16) on (2), only one rule applies.

Case (16b).

(3)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro

by assumption

By rule induction over Rules (26) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.18. $\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$ is not derivable.

Proof. We prove by assuming $\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$ and obtaining a contradiction.

(1)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\dot{\xi_1})$$

by assumption

By rule induction over Rules (14) on (2), no rule applies due to syntactic contradiction.

Case (17a).

(2)
$$\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\dot{\xi_1})$$

by assumption

By rule induction over Rules (16) on (2), only one rule applies.

Case (16b).

(3)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (26) on (3), no rule applies due to syntactic contradiction.

Lemma 1.0.19 (Soundness and Completeness of Satisfaction Judgment). $e \models \dot{\xi}$ iff $satisfy(e, \dot{\xi}) = true$.

Proof. We prove soundness and completeness separately.

1. Soundness:

(1)
$$e \models \dot{\xi}$$
 by assumption

By rule induction over Rules (14) on (1).

Case (14a).

(2) $\dot{\xi} = \top$ by assumption (3) $satisfy(e, \top) = true$ by Definition 15a

Case (14b).

 $\begin{array}{ll} (2) & e = \underline{n} & \text{by assumption} \\ (3) & \dot{\xi} = \underline{n} & \text{by assumption} \\ (4) & satisfy(\underline{n},\underline{n}) = (n=n) = \text{true} & \text{by Definition 15b} \end{array}$

Case (14c).

- (2) $e = \underline{n_1}$ by assumption
- (3) $\dot{\xi} = \underline{\underline{p_2}}$ by assumption
- (4) $n_1 \neq n_2$ by assumption
- (5) $satisfy(\underline{n_1}, \underline{p_2}) = (n_1 \neq n_2) = true$ by Definition 15c on (4)

Case (14d).

- (2) $\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$ by assumption
- (3) $e \models \dot{\xi}_1$ by assumption
- (4) $e \models \dot{\xi}_2$ by assumption
- (5) $satisfy(e, \dot{\xi}_1) = true$ by IH on (3)
- (6) $satisfy(e, \dot{\xi}_2) = true$ by IH on (4)
- (7) $satisfy(e, \dot{\xi}_1 \land \dot{\xi}_2) = satisfy(e, \dot{\xi}_1)$ and $satisfy(e, \dot{\xi}_2) = true$ by Definition 15d on (5) and (6)

Case (14e).

(2) $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$ by assumption (3) $e \models \dot{\xi}_1$ by assumption (4) $satisfy(e, \dot{\xi}_1) = true$ by IH on (3)

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(5) satisfy(e,\dot{\xi}_1\vee\dot{\xi}_2)=satisfy(e,\dot{\xi}_1) or satisfy(e,\dot{\xi}_2)= true by Definition 15e on (4)
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Case (14f).

- (2) $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$ by assumption (3) $e \models \dot{\xi}_2$ by assumption
- (4) $satisfy(e, \dot{\xi}_2) = true$ by IH on (3)
- (5) $satisfy(e, \dot{\xi}_1 \lor \dot{\xi}_2) = satisfy(e, \dot{\xi}_1)$ or $satisfy(e, \dot{\xi}_2) = \text{true}$ by Definition 15e on (4)

Case (14g).

- (2) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (3) $\dot{\xi} = \text{inl}(\dot{\xi}_1)$ by assumption
- (4) $e_1 \models \dot{\xi}_1$ by assumption
- (5) $satisfy(e_1, \dot{\xi}_1) = true$ by IH on (4)
- (6) $satisfy(inl_{\tau_2}(e_1), inl(\dot{\xi}_1)) = satisfy(e_1, \dot{\xi}_1) = true$ by Definition 15f on (5)

Case (14h).

- (2) $e = inr_{\tau_1}(e_2)$ by assumption
- (3) $\dot{\xi} = inl(\dot{\xi}_2)$ by assumption
- (4) $e_2 \models \dot{\xi}_2$ by assumption
- (5) $satisfy(e_2, \dot{\xi}_2) = true$ by IH on (4)
- (6) $satisfy(inr_{\tau_1}(e_2), inr(\dot{\xi}_2)) = satisfy(e_2, \dot{\xi}_2) = true$ by Definition 15g on (5)

Case (14i).

- (2) $e = (e_1, e_2)$ by assumption
- (3) $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$ by assumption
- (4) $e_1 \models \dot{\xi}_1$ by assumption
- (5) $e_2 \models \dot{\xi}_2$ by assumption
- (6) $satisfy(e_1, \dot{\xi}_1) = true$ by IH on (4)
- (7) $satisfy(e_2, \dot{\xi}_2) = true$ by IH on (5)
- (8) $satisfy((e_1,e_2),(\dot{\xi}_1,\dot{\xi}_2)) = satisfy(e_1,\dot{\xi}_1)$ and $satisfy(e_2,\dot{\xi}_2) = true$

by Definition 15h on

(6) and (7)

Case (14j).

(2)
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$
 by assumption
(3) e notintro by assumption
(4) $prl(e) \models \dot{\xi}_1$ by assumption
(5) $prr(e) \models \dot{\xi}_2$ by assumption
(6) $satisfy(prl(e), \dot{\xi}_1) = true$ by IH on (4)
(7) $satisfy(prr(e), \dot{\xi}_2) = true$ by IH on (5)

By rule induction over Rules (26) on (3).

Otherwise.

(8)
$$e = \langle \rangle^u, \langle e_0 \rangle^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \}$$
 by assumption

(9)
$$satisfy(e, (\dot{\xi}_1, \dot{\xi}_2)) = satisfy(\texttt{prl}(e), \dot{\xi}_1)$$
 and $satisfy(\texttt{prr}(e), \dot{\xi}_2) = \text{true}$ by Definition 15 on (6) and (7)

2. Completeness:

(1)
$$satisfy(e, \dot{\xi}) = true$$
 by assumption

By structural induction on $\dot{\xi}$.

Case $\dot{\xi} = \top$.

(2)
$$e \models \top$$
 by Rule (14a)

Case $\dot{\xi} = \bot$,?.

(2)
$$\operatorname{satisfy}(e,\dot{\xi})=\operatorname{false}$$
 by Definition 150

(2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = \underline{n}$.

By structural induction on e.

Case $e = \underline{n'}$.

(2)
$$n' = n$$
 by Definition 15b on (1) (3) $\underline{n'} \models \underline{n}$ by Rule (14b) on (2)

Otherwise.

(2)
$$satisfy(e, \underline{n}) = false$$
 by Definition 150

(2) contradicts (1) and thus vacuously true.

Case $\xi = \chi$.

By structural induction on e.

Case $e = \underline{n'}$.

(2) $n' \neq n$

by Definition 15c on (1)

(3) $\underline{n'} \models \underline{\mathscr{M}}$

by Rule (14c) on (2)

Otherwise.

- (2) $satisfy(e, \mathbf{x}) = false$
- by Definition 15o
- (2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = \dot{\xi}_1 \wedge \dot{\xi}_2$.

- (2) $satisfy(e, \dot{\xi}_1) = true$
- by Definition 15d on
- (3) $satisfy(e, \dot{\xi}_2) = true$
- by Definition 15d on
- (1)

(4) $e \models \dot{\xi}_1$

by IH on (2)

(5) $e \models \dot{\xi}_2$

by IH on (3)

(6) $e \models \dot{\xi}_1 \wedge \dot{\xi}_2$

by Rule (14d) on (4)

and (5)

Case $\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$.

- (2) $satisfy(e, \dot{\xi}_1)$ or $satisfy(e, \dot{\xi}_2) = true$
 - by Definition 15e on (1)

By case analysis on (2).

Case $satisfy(e, \dot{\xi}_1) = true.$

- (3) $satisfy(e, \dot{\xi}_1) = true$
- by assumption

(4) $e \models \dot{\xi}_1$

by IH on (3)

(5) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$

by Rule (14e) on (4)

Case $satisfy(e, \dot{\xi}_2) = true.$

- (3) $satisfy(e, \dot{\xi}_2) = true$
- by assumption

(4) $e \models \dot{\xi}_2$

by IH on (3)

(5) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$

by Rule (14f) on (4)

Case $\dot{\xi} = \text{inl}(\dot{\xi}_1)$.

By structural induction on e.

Case $e = \operatorname{inl}_{\tau_2}(e_1)$.

- (2) $satisfy(e_1, \dot{\xi}_1) = true$
- by Definition 15f on (1)

(3) $e_1 \models \dot{\xi}_1$

- by IH on (2)
- $(4) \ \operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$
- by Rule (14g) on (3)

Otherwise.

- (2) $satisfy(e, \mathtt{inl}(\dot{\xi_1})) = \text{false}$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = inr(\dot{\xi}_2)$.

By structural induction on e.

Case $e = \operatorname{inr}_{\tau_1}(e_2)$.

- (2) $satisfy(e_2, \dot{\xi}_2) = true$ by Definition 15g on (1)
- (3) $e_2 \models \dot{\xi}_2$ by IH on (2)
- (4) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$ by Rule (14h) on (3)

Otherwise.

- (2) $satisfy(e, inr(\dot{\xi}_2)) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$.

By structural induction on e.

Case $e = (e_1, e_2)$.

- (2) $satisfy(e_1, \dot{\xi}_1) = true$ by Definition 15h on (1)
- (3) $satisfy(e_2, \dot{\xi}_2) = true$ by Definition 15h on (1)
- (4) $e_1 \models \dot{\xi}_1$ by IH on (2)
- (5) $e_2 \models \dot{\xi}_2$ by IH on (3)
- (6) $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (14i) on (4) and (5)

Case $e = \| u, (e_0)u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{ \hat{rs} \}.$

- (2) $satisfy(prl(e), \dot{\xi}_1) = true$ by Definition 15h on (1)
- (3) $satisfy(prr(e), \dot{\xi}_2) = true$ by Definition 15h on (1)
- (4) $\operatorname{prl}(e) \models \dot{\xi}_1$ by IH on (2)
- (5) $\operatorname{prr}(e) \models \dot{\xi}_2$ by IH on (3)
- (6) e notintro by each rule in Rules (26)
- (7) $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (14j) on (6) and (4) and (5)

Otherwise.

- (2) $satisfy(e, (\dot{\xi}_1, \dot{\xi}_2)) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Lemma 1.0.20.

satormay soundness and completeness

Lemma 1.0.21. $e \not\models \dot{\xi}$ and $e \not\models_? \dot{\xi}$ iff $e \not\models_? \dot{\xi}$.

Proof. 1. Sufficiency:

(1) $e \not\models \dot{\xi}$

by assumption

(2) $e \not\models_? \dot{\xi}$

by assumption

Assume $e \models^{\dagger}_{?} \dot{\xi}$. By rule induction over Rules (17) on it.

Case (17a).

(3) $e \models \dot{\xi}$

by assumption

Contradicts (1).

Case (17b).

(3) $e \models_? \dot{\xi}$

by assumption

Contradicts (2).

Therefore, $e \models^{\dagger}_{?} \dot{\xi}$ is not derivable.

2. Necessity:

 $(1) \ e \not\models_{?}^{\dagger} \dot{\xi}$

by assumption

Assume $e \models \dot{\xi}$.

(2) $e \models^{\dagger}_{?} \dot{\xi}$

by Rule (17b) on assumption

Contradicts (1). Therefore, $e \not\models \dot{\xi}$. Assume $e \models_? \dot{\xi}$.

(3) $e \models^{\dagger}_{?} \dot{\xi}$

by Rule (17a) on assumption

Contradicts (1). Therefore, $e \not\models_? \dot{\xi}$.

Theorem 1.1 (Exclusiveness of Satisfaction Judgment). If $\dot{\xi}: \tau$ and $\cdot; \Delta \vdash e: \tau$ and e final then exactly one of the following holds

- 1. $e \models \dot{\xi}$
- 2. $e \models_? \dot{\xi}$
- 3. $e \not\models_?^\dagger \dot{\xi}$

Proof.

(4) $\dot{\xi}:\tau$

by assumption

(5) \cdot ; $\Delta \vdash e : \tau$

by assumption

(6) e final

by assumption

By rule induction over Rules (8) on (4), we would show one conclusion is derivable while the other two are not.

Case (8a).

(7) $\dot{\xi} = \top$

by assumption

(8) $e \models \top$

by Rule (14a)

 $(9) \ e \not\models_? \top$

by Lemma 2.0.3

(10) $e \models_{?}^{\dagger} \top$

by Rule (17b) on (8)

Case (8b).

(7) $\dot{\xi} = \bot$

by assumption

(8) $e \not\models \bot$

by Lemma 2.0.1

(9) $e \not\models_? \bot$

by Lemma 2.0.2

(10) $e \not\models_2^{\dagger} \bot$

by Lemma 2.0.20 on

(8) and (9)

Case (1b).

(7) $\dot{\xi} = ?$

by assumption

(8) $e \not\models ?$

by Lemma 2.0.4

(9) $e \models_? ?$

by Rule (16a)

(10) $e \models_{?}^{\dagger} ?$

by Rule (17a) on (9)

Case (8c).

 $(7) \ \dot{\xi} = \underline{n_2}$

by assumption

(8) $\tau = \text{num}$

by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(9) $e = (v, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$ by assumption

Assume $e \models \underline{n_2}$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction on $\dot{\xi}$.

(11)
$$e \not\models \underline{n_2}$$
 by contradiction
(12) $\underline{n_2}$ refutable? by Rule (10a)
(13) $e \models_? \underline{n_2}$ by Rule (16b) on (10)
and (12)
(14) $e \models_?^{\dagger} n_2$ by Rule (17a) on (13)

Case (19d).

(9)
$$e = n_1$$
 by assumption

Assume $\underline{n_1} \models_? \underline{n_2}$. By rule induction over Rules (16), only one case applies.

Case (16b).

(10)
$$\underline{n_1}$$
 notintro by assumption Contradicts Lemma 4.0.5.

(11)
$$\underline{n_1} \not\models_? \underline{n_2}$$
 by contradiction

By case analysis on whether n_1 is equal to n_2 .

Case $n_1 = n_2$.

$$\begin{array}{lll} (12) & \textit{satisfy}(\underline{n_1},\underline{n_2}) = \text{true} & & \text{by Definition 15} \\ (13) & \underline{n_1} \models \underline{n_2} & & \text{by Lemma 2.0.19 on} \\ (14) & e \models_{?}^{\dagger} \underline{n_2} & & \text{by Rule (17b) on (13)} \\ \end{array}$$

Case
$$n_1 \neq n_2$$
.

(12)
$$satisfy(\underline{n_1}, \underline{n_2}) = false$$
 by Definition 15
(13) $\underline{n_1} \not\models \underline{n_2}$ by Lemma 2.0.19 on (12)
(14) $e \not\models_{?}^{\dagger} \underline{n_2}$ by Lemma 2.0.20 on (11) and (13)

Case (8f).

(7)
$$\dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$$
 by assumption

By inductive hypothesis on (5) and (6), exactly one of $e \models \dot{\xi}_1$, $e \models_? \dot{\xi}_1$, and $e \not\models_?^{\dagger} \dot{\xi}_1$ holds. The same goes for $\dot{\xi}_2$. By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $e \models \dot{\xi}_1, e \models \dot{\xi}_2$.

(8) $e \models \dot{\xi}_1$	by assumption
$(9) \ e \not\models_? \dot{\xi}_1$	by assumption
$(10) \ e \models \dot{\xi}_2$	by assumption
$(11) \ e \not\models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule (14e) on (8)
$(13) \ e \models_?^{\dagger} \dot{\xi_1} \lor \dot{\xi_2}$	by Rule $(17b)$ on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\dot{\xi}_1 \vee \dot{\xi}_2$ refutable? by assumption Contradicts Lemma 2.0.17.

Case (16c).

(14)
$$e \models_{?} \dot{\xi}_{1}$$
 by assumption Contradicts (9).

Case (16d).

(14)
$$e \models_{?} \dot{\xi}_{2}$$
 by assumption Contradicts (11).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \models \dot{\xi}_1, e \models_? \dot{\xi}_2$.

(8) $e \models \dot{\xi}_1$	by assumption
$(9) \ e \not\models_? \dot{\xi}_1$	by assumption
$(10) \ e \not\models \dot{\xi}_2$	by assumption
$(11) \ e \models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule (14e) on (8)
$(13) \ e \models^{\dagger}_{?} \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule (17b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14)
$$\dot{\xi}_1 \vee \dot{\xi}_2$$
 refutable? by assumption Contradicts Lemma 2.0.17.

Case (16c).

(14)
$$e \models_? \dot{\xi}_1$$
 by assumption Contradicts (9).

Case (16d).

(14)
$$e \not\models \dot{\xi}_1$$
 by assumption Contradicts (8).

(15) $e \not\models_{?} \dot{\xi}_1 \lor \dot{\xi}_2$ by contradiction

Case $e \models \dot{\xi}_1, e \not\models_?^{\dagger} \dot{\xi}_2$.

(8) $e \models \dot{\xi}_1$ by assumption (9) $e \not\models_? \dot{\xi}_1$ by assumption (10) $e \not\models \dot{\xi}_2$ by assumption (11) $e \not\models_? \dot{\xi}_2$ by assumption (12) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (14e) on (8) (13) $e \models_?^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (17b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\dot{\xi}_1 \vee \dot{\xi}_2$ refutable? by assumption Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \models_? \dot{\xi}_1$ by assumption Contradicts (9).

Case (16d).

(14) $e \not\models \dot{\xi}_1$ by assumption Contradicts (8).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \models_? \dot{\xi}_1, e \models \dot{\xi}_2$.

(8) $e \not\models \dot{\xi}_1$	by assumption
$(9) \ e \models_? \dot{\xi}_1$	by assumption
$(10) \ e \models \dot{\xi}_2$	by assumption
$(11) \ e \not\models_? \dot{\xi}_2$	by assumption
$(12) \ e \models \dot{\xi}_1 \lor \dot{\xi}_2$	by Rule $(14f)$ on (10)
$(13) \ e \models^{\dagger}_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}$	by Rule $(17b)$ on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\dot{\xi}_1 \lor \dot{\xi}_2$ refutable? by assumption Contradicts Lemma 2.0.17.

(14) $e \not\models \dot{\xi}_2$ by assumption

Contradicts (10).

Case (16d).

(14) $e \models_? \dot{\xi}_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$ by contradiction

Case $e \models_? \dot{\xi}_1, e \models_? \dot{\xi}_2$.

(8) $e \not\models \dot{\xi}_1$ by assumption

(9) $e \models_? \dot{\xi}_1$ by assumption (10) $e \not\models \dot{\xi}_2$ by assumption

(11) $e \models_? \dot{\xi}_2$ by assumption

(12) $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (16c) on (9) and (10)

(13) $e \models_{2}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Rule (17a) on (12)

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (14), only two cases apply.

Case (14e).

(14) $e \models \dot{\xi}_1$ by assumption

Contradicts (8)

Case (14f).

(14) $e \models \dot{\xi}_2$ by assumption

Contradicts (10)

(15) $e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$ by contradiction

Case $e \models_? \dot{\xi}_1, e \not\models_?^{\dagger} \dot{\xi}_2$.

(8) $e \not\models \dot{\xi}_1$ by assumption (9) $e \models_? \dot{\xi}_1$ by assumption

 $\begin{array}{ll} (10) \ e \not\models \dot{\xi}_2 & \text{by assumption} \\ (11) \ e \not\models_? \dot{\xi}_2 & \text{by assumption} \end{array}$

(12) $e \models_{?} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Rule (16c) on (9)

 $\begin{array}{c}
\text{and } (10)
\end{array}$

(13) $e \models_{?}^{\dagger} \dot{\xi}_{1} \lor \dot{\xi}_{2}$ by Rule (17a) on (12)

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (14), only two cases apply.

(14)
$$e \models \dot{\xi}_1$$
 by assumption

Contradicts (8).

Case (14f).

(14)
$$e \models \dot{\xi}_2$$
 by assumption

Contradicts (10).

(15)
$$e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \not\models_?^\dagger \dot{\xi}_1, e \models \dot{\xi}_2$.

(8)
$$e \not\models \dot{\xi}_1$$
 by assumption
(9) $e \not\models_? \dot{\xi}_1$ by assumption
(10) $e \models \dot{\xi}_2$ by assumption
(11) $e \not\models_? \dot{\xi}_2$ by assumption
(12) $e \models \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (14f) on (10)
(13) $e \models_?^{\dot{\tau}} \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (17b) on (12)

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14)
$$\dot{\xi}_1 \vee \dot{\xi}_2$$
 refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14)
$$e \not\models \dot{\xi}_2$$
 by assumption

Contradicts (10).

Case (16d).

(14)
$$e \models_? \dot{\xi}_2$$
 by assumption Contradicts (11).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction

Case $e \not\models_?^\dagger \dot{\xi}_1, e \models_? \dot{\xi}_2$.

(8)
$$e \not\models \dot{\xi}_1$$
 by assumption
(9) $e \not\models ? \dot{\xi}_1$ by assumption
(10) $e \not\models \dot{\xi}_2$ by assumption
(11) $e \models ? \dot{\xi}_2$ by assumption
(12) $e \models ? \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (16d) on (11)
and (8)
(13) $e \models ? \dot{\xi}_1 \lor \dot{\xi}_2$ by Rule (17a) on (12)

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (14), only two cases apply.

Case (14e).

 $(14) \ e \models \dot{\xi}_1$

by assumption

Contradicts (8)

Case (14f).

(14) $e \models \dot{\xi}_2$

by assumption

Contradicts (10)

 $(15) \ e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$

by contradiction

Case $e \not\models_?^\dagger \dot{\xi}_1, e \not\models_?^\dagger \dot{\xi}_2$.

(8) $e \not\models \dot{\xi}_1$

by assumption

 $(9) \ e \not\models_? \dot{\xi}_1$

by assumption by assumption

(10) $e \not\models \dot{\xi}_2$ (11) $e \not\models_? \dot{\xi}_2$

by assumption

Assume $e \models \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (14) on it, only two cases apply.

Case (14e).

 $(12) \ e \models \dot{\xi}_1$

by assumption

Contradicts (8).

Case (14f).

(12) $e \models \dot{\xi}_2$

by assumption

Contradicts (10).

 $(13) \ e \not\models \dot{\xi}_1 \lor \dot{\xi}_2$

by contradiction

Assume $e \models_? \dot{\xi}_1 \lor \dot{\xi}_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\dot{\xi}_1 \lor \dot{\xi}_2$ refutable?

by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \models_? \dot{\xi}_1$

by assumption

Contradicts (9).

Case (16d).

 $(14) \ e \models_? \dot{\xi}_2$

by assumption

Contradicts (11).

(15)
$$e \not\models_? \dot{\xi}_1 \lor \dot{\xi}_2$$
 by contradiction
(16) $e \not\models_?^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$ by Lemma 2.0.20 on
(13) and (15)

Case (8g).

(7)
$$\dot{\xi} = \text{inl}(\dot{\xi}_1)$$
 by assumption
(8) $\tau = (\tau_1 + \tau_2)$ by assumption
(9) $\dot{\xi}_1 : \tau_1$ by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(10)
$$e = \langle | \rangle^u, \langle | e_0 \rangle^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{ \hat{rs} \}$$
 by assumption
(11) e notintro by Rule
$$(26a), (26b), (26c), (26d), (26e), (26f)$$

Assume $e \models \mathtt{inl}(\dot{\xi}_1)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inl}(\dot{\xi}_1)$$
 by contradiction

By case analysis on the value of $refutable_7(inl(\dot{\xi}_1))$.

Case $refutable_7(inl(\dot{\xi}_1)) = true.$

(13)
$$refutable_{?}(inl(\dot{\xi}_{1})) = true$$
 by assumption (14) $inl(\dot{\xi}_{1})$ refutable? by Lemma 2.0.14 on (13)

(15)
$$e \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (16b) on (11) and (14)

(16)
$$e \models^{\dagger}_{?} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (17a) on (15)

Case $refutable_?(\mathtt{inl}(\dot{\xi}_1)) = \mathrm{false.}$

(13)
$$refutable_{?}(inl(\dot{\xi}_{1})) = false$$
 by assumption

(14)
$$\underline{\operatorname{inl}(\dot{\xi}_1)}$$
 refutable? by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(15)
$$\operatorname{inl}(\dot{\xi}_1)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inl}(\dot{\xi}_1)$$
 by contradiction

(17)
$$e \not\models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_{1})$$
 by Lemma 2.0.20 on (12) and (16)

Case (19j).

- (10) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption (11) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (12) e_1 final by Lemma 4.0.2 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_1 \models \dot{\xi}_1$, $e_1 \models_? \dot{\xi}_1$, and $e_1 \not\models_?^{\dagger} \dot{\xi}_1$ holds. By case analysis on which one holds.

Case $e_1 \models \dot{\xi}_1$.

- (13) $e_1 \models \dot{\xi}_1$ by assumption
- (14) $e_1 \not\models_? \dot{\xi}_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi_1})$ by Rule (14g) on (13)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$ by Rule (17b) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (16) on it, only two cases apply.

Case (16b).

(17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16e).

(17)
$$e_1 \models_? \dot{\xi}_1$$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\dot{\xi_1})$$
 by contradiction

Case $e_1 \models_? \dot{\xi}_1$.

- (13) $e_1 \not\models \dot{\xi}_1$ by assumption
- (14) $e_1 \models_? \dot{\xi}_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by Rule (16e) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$ by Rule (17a) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (14) on it, only one case applies.

Case (14g).

$$(17) e_1 \models \dot{\xi}_1$$

Contradicts (13).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\dot{\xi_1})$$
 by contradiction

Case $e_1 \not\models_?^\dagger \dot{\xi}_1$.

 $(13) e_1 \not\models \dot{\xi}_1$

by assumption

(14) $e_1 \not\models_? \dot{\xi}_1$

by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (14) on it, only one case applies.

Case (14g).

 $(15) e_1 \models \dot{\xi}_1$

Contradicts (13).

(16) $\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\dot{\xi_1})$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro

by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16e).

(17) $e_1 \models_? \dot{\xi}_1$

Contradicts (14).

(18) $\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\dot{\xi_1})$

by contradiction

(19) $\operatorname{inl}_{\tau_2}(e_1) \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$

by Lemma 2.0.20 on

(16) and (18)

Case (19k).

(10)
$$e = inr_{\tau_1}(e_2)$$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\dot{\xi}_1)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inl}(\dot{\xi}_1)$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\dot{\xi_1})$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(12) $\operatorname{inr}_{\tau_1}(e_2)$ notintro

by assumption

By rule induction over Rules (26) on (12), no case applies due to syntactic contradiction.

(13) $\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inl}(\dot{\xi}_1)$

by contradiction

(14) $\operatorname{inr}_{\tau_1}(e_2) \not\models_{?}^{\dagger} \operatorname{inl}(\dot{\xi}_1)$

by Lemma 2.0.20 on

(11) and (13)

Case (8h).

$$(7) \ \dot{\xi} = \mathtt{inr}(\dot{\xi}_2)$$

by assumption

(8)
$$\tau = (\tau_1 + \tau_2)$$

by assumption

(9)
$$\dot{\xi}_2 : \tau_2$$

by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(10)
$$e = \langle || u, || e_0 \rangle u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{ \hat{rs} \}$$

by assumption

$$(11)$$
 e notintro

by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

Assume $e \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inr}(\dot{\xi}_2)$$

by contradiction

By case analysis on the value of $refutable_2(inr(\dot{\xi}_2))$.

inr is refutable

Case $refutable_{?}(inr(\dot{\xi}_{2})) = true.$

(13)
$$refutable_?(inr(\dot{\xi}_2)) = true$$

by assumption

(14)
$$\operatorname{inr}(\dot{\xi}_2)$$
 refutable?

by Lemma 2.0.14 on

(13)

(15)
$$e \models_? \operatorname{inr}(\dot{\xi}_2)$$

by Rule (16b) on (11)

and (14)

(16)
$$e \models^{\dagger}_{?} \operatorname{inr}(\dot{\xi}_2)$$

by Rule (17a) on (15)

Case $refutable_7(inr(\dot{\xi}_2)) = false.$

(13)
$$refutable_?(inr(\dot{\xi}_2)) = false$$

by assumption

(14)
$$\operatorname{inr}(\dot{\xi}_2)$$
 refutable?

by Lemma 2.0.14 on (13)

Assume $e \models_? inr(\dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(15)
$$\operatorname{inr}(\dot{\xi}_2)$$
 refutable? Contradicts (14).

by assumption

(16)
$$e \not\models_? \operatorname{inr}(\dot{\xi}_2)$$

by contradiction

(17)
$$e \not\models^{\dagger}_{?} \operatorname{inr}(\dot{\xi}_2)$$

by Lemma 2.0.20 on

(12) and (16)

Case (19j).

(10)
$$e = inl_{\tau_2}(e_1)$$

by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inr}(\dot{\xi}_2)$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\dot{\xi_2})$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(12)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro

by assumption

By rule induction over Rules (26) on (12), no case applies due to syntactic contradiction.

(13)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inr}(\dot{\xi}_2)$$

by contradiction

(14)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_?^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$

by Lemma 2.0.20 on

(11) and (13)

Case (19k).

(10)
$$e = inr_{\tau_1}(e_2)$$

by assumption

(11)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$

by assumption

(12)
$$e_2$$
 final

by Lemma 4.0.3 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_2 \models \dot{\xi}_2$, $e_2 \models_? \dot{\xi}_2$, and $e_2 \not\models_?^{\dagger} \dot{\xi}_2$ holds. By case analysis on which one holds.

Case $e_2 \models \dot{\xi}_2$.

$$(13) \ e_2 \models \dot{\xi}_2$$

by assumption

(14)
$$e_2 \not\models_? \dot{\xi}_2$$

by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$$

by Rule (14g) on (13)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inr}(\dot{\xi_2})$$

by Rule (17b) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (16) on it, only two cases apply.

Case (16b).

(17)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16f).

$$(17) \ e_2 \models_? \dot{\xi}_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\dot{\xi_2})$$

by contradiction

Case $e_2 \models_? \dot{\xi}_2$.

(13)
$$e_2 \not\models \dot{\xi}_2$$

by assumption

$$(14) e_2 \models_? \dot{\xi}_2$$

by assumption

$$(15) \ \operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi_2})$$

by Rule (16f) on (14)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_?^{\dagger} \operatorname{inr}(\dot{\xi_2})$$

by Rule (17a) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14h).

(17)
$$e_2 \models \dot{\xi}_2$$

Contradicts (13).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\dot{\xi_2})$$

by contradiction

Case $e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$.

$$(13) \ e_2 \not\models \dot{\xi}_2$$

by assumption

(14)
$$e_2 \not\models_? \dot{\xi}_2$$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14h).

$$(15) \ e_2 \models \dot{\xi}_2$$

Contradicts (13).

$$(16) \ \operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\dot{\xi_2})$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(17)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16f).

(17)
$$e_2 \models_? \dot{\xi}_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\dot{\xi_2})$$

by contradiction

(19)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_?^{\dagger} \operatorname{inl}(\dot{\xi_1})$$

by Lemma 2.0.20 on

(16) and (18)

Case (14i).

(7)
$$\dot{\xi} = (\dot{\xi}_1, \dot{\xi}_2)$$

by assumption

(8)
$$\tau = (\tau_1 \times \tau_2)$$

by assumption

(9)
$$\dot{\xi}_1 : \tau_1$$

by assumption

(10)
$$\dot{\xi}_2 : \tau_2$$

by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(11)
$$e = ()^u, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$
 by assumption

(12) e notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

(13) e indet by Lemma 4.0.9 on (6)

and (12)

(14) prl(e) indet by Rule (24g) on (13)

(15) prl(e) final by Rule (25b) on (14)

(16) prr(e) indet by Rule (24h) on (13)

(17) prr(e) final by Rule (25b) on (16) (18) \cdot ; $\Delta \vdash prl(e) : \tau_1$ by Rule (19h) on (5)

(19) \cdot ; $\Delta \vdash \mathsf{prr}(e) : \tau_2$ by Rule (19i) on (5)

By inductive hypothesis on (9) and (18) and (15), exactly one of $\operatorname{prl}(e) \models \dot{\xi}_1$, $\operatorname{prl}(e) \models_? \dot{\xi}_1$, and $\operatorname{prl}(e) \not\models_?^{\dagger} \dot{\xi}_1$ holds. By inductive hypothesis on (10) and (19) and (17), exactly one of

 $\operatorname{prr}(e) \models \dot{\xi}_2, \operatorname{prr}(e) \models_? \dot{\xi}_2, \text{ and } \operatorname{prr}(e) \not\models_?^{\dagger} \dot{\xi}_2 \text{ holds.}$

By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $prl(e) \models \dot{\xi}_1, prr(e) \models \dot{\xi}_2.$

(20) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption

(22) $prr(e) \models \dot{\xi}_2$ by assumption

(23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$ by assumption

(24) $e \models (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (14j) on (12) and (20) and (22)

(25) $e \models^{\dagger}_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17b) on (24)

(26) $(\dot{\xi}_1,\dot{\xi}_2)$ refutable? by Lemma 2.0.18 on (12) and (24)

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption Contradicts (26).

(28) $e \not\models_{?} (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Case $prl(e) \models \dot{\xi}_1, prr(e) \models_? \dot{\xi}_2.$

(20) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption

- (21) $prl(e) \not\models_? \dot{\xi}_1$ by assumption
- (22) $prr(e) \not\models \dot{\xi}_2$ by assumption (23) $prr(e) \models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24)
$$prr(e) \models \dot{\xi}_2$$
 by assumption Contradicts (22)

(25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

By rule induction over Rules (16) on (23), only one case applies.

Case (16b).

- (26) $\dot{\xi}_2$ refutable? by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Rule (10e) on (26)

assume no "or" and

"and" in

pair

- (28) $e \models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (16b) on (12) and (27)
- (29) $e \models^{\dagger}_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17a) on (28)

Case $prl(e) \models \dot{\xi}_1, prr(e) \not\models_?^{\dagger} \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \models \dot{\xi}_1$ by assumption
- (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption (22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$ by assumption
- (23) $prr(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

- (24) $prr(e) \models \dot{\xi}_2$ by assumption Contradicts (22).
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption

By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

- (27) $\dot{\xi}_1$ refutable? by assumption
- (28) prl(e) notintro by Rule (26e)

$$(29) \ \operatorname{prl}(e) \models_? \dot{\xi}_1 \qquad \text{by Rule (16b) on (28)} \\ \operatorname{and (27)} \\ \operatorname{Contradicts (21)}. \\ \operatorname{Case (10e)}. \\ (27) \ \dot{\xi}_2 \ \operatorname{refutable}_? \qquad \text{by assumption} \\ (28) \ \operatorname{prr}(e) \ \operatorname{notintro} \qquad \operatorname{by Rule (26f)} \\ (29) \ \operatorname{prr}(e) \models_? \dot{\xi}_2 \qquad \operatorname{by Rule (16b) on (28)} \\ \operatorname{and (27)} \\ \operatorname{Contradicts (23)}. \\ \\ (30) \ e \not\models_? (\dot{\xi}_1, \dot{\xi}_2) \qquad \operatorname{by contradiction} \\ (31) \ e \not\models_? (\dot{\xi}_1, \dot{\xi}_2) \qquad \operatorname{by Lemma 2.0.20 on} \\ (25) \ \operatorname{and (30)} \\ \\ \operatorname{Case prl}(e) \models_? \dot{\xi}_1, \operatorname{prr}(e) \models_{\dot{\xi}_2}. \\ (20) \ \operatorname{prl}(e) \models_{\dot{\xi}_1} & \operatorname{by assumption} \\ (21) \ \operatorname{prl}(e) \models_{\dot{\gamma}} \dot{\xi}_2 & \operatorname{by assumption} \\ (22) \ \operatorname{prr}(e) \models_{\dot{\gamma}} \dot{\xi}_2 & \operatorname{by assumption} \\ (23) \ \operatorname{prr}(e) \models_{\dot{\gamma}} \dot{\xi}_2 & \operatorname{by assumption} \\ \operatorname{Assume } e \models (\dot{\xi}_1, \dot{\xi}_2). & \operatorname{By rule induction over Rules (14), only one case applies.} \\ \\ \operatorname{Case (14i)}. & (24) \ \operatorname{prl}(e) \models_{\dot{\gamma}} \dot{\xi}_1 & \operatorname{by assumption} \\ \operatorname{Contradicts (20)}. & (25) \ e \not\models (\dot{\xi}_1, \dot{\xi}_2) & \operatorname{by contradiction} \\ \operatorname{By rule induction over Rules (16) on (21), only one case applies.} \\ \\ \hline{Case (16b)}. & (26) \ \dot{\xi}_1 \ \operatorname{refutable}_? & \operatorname{by assumption} \\ (27) \ (\dot{\xi}_1, \dot{\xi}_2) \ \operatorname{refutable}_? & \operatorname{by Rule (10e) on (26)} \\ (28) \ e \models_? (\dot{\xi}_1, \dot{\xi}_2) & \operatorname{by Rule (16b) on (12)} \\ \operatorname{and (27)} & \operatorname{and (27)} \\ (29) \ e \models_?^2 (\dot{\xi}_1, \dot{\xi}_2) & \operatorname{by Rule (17a) on (28)} \\ \\ \hline{Case } \operatorname{prl}(e) \models_? \dot{\xi}_1, \operatorname{prr}(e) \models_? \dot{\xi}_2. \\ (20) \ \operatorname{prl}(e) \models_? \dot{\xi}_1 & \operatorname{by assumption} \\ (21) \ \operatorname{prl}(e) \models_? \dot{\xi}_1 & \operatorname{by assumption} \\ \operatorname{by$$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only

(23) $prr(e) \models_? \dot{\xi}_2$

one case applies.

Case (14j).

(24) $prl(e) \models \dot{\xi}_1$ Contradicts (20).

by assumption

(25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$

by contradiction

By rule induction over Rules (16) on (23), only one case applies.

assume no
"or" and
"and" in
pair

assume no "or" and

"and" in

pair

Case (16b).

- (26) $\dot{\xi}_2$ refutable?
- by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable?
- by Rule (10e) on (26) by Rule (16b) on (12)
- (28) $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$
- and (27)
- (29) $e \models^{\dagger}_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$
- by Rule (17a) on (28)

Case $prl(e) \models_? \dot{\xi}_1, prr(e) \not\models_?^{\dagger} \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$
- by assumption

(21) $prl(e) \models_? \dot{\xi}_1$

by assumption

(22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$

by assumption

- (23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$
- by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24) $\operatorname{prl}(e) \models \dot{\xi}_1$

by assumption

Contradicts (20)

(25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$

by contradiction

By rule induction over Rules (16) on (21), only one case applies.

Case (16b).

(26) $\dot{\xi}_1$ refutable?

by assumption

(20) ζ_1 refutable?

by Rule (10e) on (26)

(27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable?

by Rule (16b) on (12)

(28) $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$

and (27)

(29) $e \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$

by Rule (17a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \dot{\xi}_{1}, prr(e) \models \dot{\xi}_{2}.$

(20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$

by assumption

(21) $\operatorname{prl}(e) \not\models_? \xi_1$

by assumption

(22) $prr(e) \models \dot{\xi}_2$

by assumption

(23) $\operatorname{prr}(e) \not\models_? \dot{\xi}_2$

by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

(24)
$$prl(e) \models \dot{\xi}_1$$
 by assumption Contradicts (20)

(25)
$$e \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Assume $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

- (27) $\dot{\xi}_1$ refutable? by assumption (28) prl(e) notintro by Rule (26e)
- (29) $prl(e) \models_{?} \dot{\xi}_{1}$ by Rule (16b) on (28) and (27)

Contradicts (21).

Case (10e).

- (27) $\dot{\xi}_2$ refutable? by assumption (28) prr(e) notintro by Rule (26f)
- (29) $prr(e) \models_{?} \dot{\xi}_{2}$ by Rule (16b) on (28) and (27)

Contradicts (23).

(30)
$$e \not\models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by contradiction
(31) $e \not\models_{?}^{\dagger} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Lemma 2.0.20 on
(25) and (30)

Case $prl(e) \not\models_?^{\dagger} \dot{\xi}_1, prr(e) \models_? \dot{\xi}_2.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption (22) $\operatorname{prr}(e) \not\models \dot{\xi}_2$ by assumption (23) $\operatorname{prr}(e) \models_? \dot{\xi}_2$ by assumption
- Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24)
$$prl(e) \models \dot{\xi}_1$$
 by assumption Contradicts (20).

(25)
$$e \not\models (\dot{\xi}_1, \dot{\xi}_2)$$

by contradiction

assume no "or" and

"and" in

pair

By rule induction over Rules (16) on (23), only one case applies.

Case (16b).

- (26) $\dot{\xi}_2$ refutable? by assumption
- (27) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by Rule (10e) on (26)
- (28) $e \models_{?} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (16b) on (12) and (27)
- (29) $e \models_{?}^{\dagger} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Rule (17a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \dot{\xi}_{1}, prr(e) \not\models_{?}^{\dagger} \dot{\xi}_{2}.$

- (20) $\operatorname{prl}(e) \not\models \dot{\xi}_1$ by assumption
- (21) $\operatorname{prl}(e) \not\models_? \dot{\xi}_1$ by assumption
- (22) $prr(e) \not\models \dot{\xi}_2$ by assumption
- (23) $prr(e) \not\models_? \dot{\xi}_2$ by assumption

Assume $e \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

- (24) $prl(e) \models \dot{\xi}_1$ by assumption Contradicts (20)
- (25) $e \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $e \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(26) $(\dot{\xi}_1, \dot{\xi}_2)$ refutable? by assumption

By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

- (27) $\dot{\xi}_1$ refutable? by assumption
- (28) prl(e) notintro by Rule (26e)
- (29) $prl(e) \models_{?} \dot{\xi}_{1}$ by Rule (16b) on (28) and (27)

Contradicts (21).

Case (10e).

- (27) $\dot{\xi}_2$ refutable? by assumption (28) prr(e) notintro by Rule (26f)
- (29) $\operatorname{prr}(e) \models_{?} \dot{\xi}_{2}$ by Rule (16b) on (28) and (27)

Contradicts (23).

(30)
$$e \not\models_{?} (\dot{\xi}_{1}, \dot{\xi}_{2})$$
 by contradiction
(31) $e \not\models_{?}^{+} (\dot{\xi}_{1}, \dot{\xi}_{2})$ by Lemma 2.0.20 on
(25) and (30)

Case (19g).

$(11) \ e = (e_1, e_2)$	by assumption
$(12) \cdot ; \Delta \vdash e_1 : \tau_1$	by assumption
$(13) \cdot ; \Delta \vdash e_2 : \tau_2$	by assumption
(14) e_1 final	by Lemma $4.0.4$ on (6)
(15) e_2 final	by Lemma $4.0.4$ on (6)

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 \models \dot{\xi}_1, e_1 \models_? \dot{\xi}_1$, and $e_1 \models \overline{\dot{\xi}_1}$ holds. By inductive hypothesis on (10) and (13) and (15), exactly one of

 $e_2 \models \dot{\xi}_2, e_2 \models_? \dot{\xi}_2$, and $e_2 \models \overline{\dot{\xi}_2}$ holds. By case analysis on which conclusion holds for $\dot{\xi}_1$ and $\dot{\xi}_2$.

Case $e_1 \models \dot{\xi}_1, e_2 \models \dot{\xi}_2$.

$(16) e_1 \models \dot{\xi}_1$	by assumption
(17) $e_1 \not\models_? \dot{\xi}_1$	by assumption
$(18) e_2 \models \dot{\xi}_2$	by assumption
(19) $e_2 \not\models_? \dot{\xi}_2$	by assumption
$(20) (e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$	by Rule (14i) on (16)
	and (18)
$(21) (e_1, e_2) \models_2^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$	by Rule (17b) on (20)

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1,e_2) notintro by assumption Contradicts Lemma 4.0.8.

Case (16g).

(22)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption Contradicts (17).

Case (16h).

(22)
$$e_2 \models_? \dot{\xi}_2$$
 by assumption Contradicts (19).

Case (16i).

(22)
$$e_1 \models_? \dot{\xi}_1$$
 by assumption Contradicts (17).

(23)
$$(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Case $e_1 \models \dot{\xi}_1, e_2 \models_? \dot{\xi}_2$.

(16) $e_1 \models \dot{\xi}_1$ by assumption (17) $e_1 \not\models_? \dot{\xi}_1$ by assumption

(18) $e_2 \not\models \dot{\xi}_2$ by assumption

(19) $e_2 \models_? \dot{\xi}_2$ by assumption

(20) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (16h) on (16) and (19)

(21) $(e_1, e_2) \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(22) $e_2 \models \dot{\xi}_2$ by assumption Contradicts (18).

(23) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Case $e_1 \models \dot{\xi}_1, e_2 \not\models_?^{\dagger} \dot{\xi}_2$.

(16) $e_1 \models \dot{\xi}_1$ by assumption (17) $e_1 \not\models_? \dot{\xi}_1$ by assumption (18) $e_2 \not\models \dot{\xi}_2$ by assumption (19) $e_2 \not\models_? \dot{\xi}_2$ by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(20) $e_2 \models \dot{\xi}_2$ by assumption Contradicts (18).

(21) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

Case (16h).

(22) $e_2 \models_? \dot{\xi}_2$ by assumption

Contradicts (19).

Case (16i).

(22) $e_1 \models_? \dot{\xi_1}$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

(24) $(e_1, e_2) \not\models_?^\dagger (\dot{\xi_1}, \dot{\xi_2})$ by Lemma 2.0.20 on

(21) and (23)

Case $e_1 \models_? \dot{\xi}_1, e_2 \models \dot{\xi}_2$.

(16) $e_1 \not\models \dot{\xi}_1$ by assumption

(17) $e_1 \models_? \dot{\xi}_1$ by assumption

(18) $e_2 \models \dot{\xi}_2$ by assumption

(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption

(20) $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (16g) on (17) and (18)

(21) $(e_1, e_2) \models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(22) $e_1 \models \dot{\xi}_1$ by assumption

Contradicts (16).

(23) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Case $e_1 \models_? \dot{\xi}_1, e_2 \models_? \dot{\xi}_2$.

(16) $e_1 \not\models \dot{\xi}_1$ by assumption (17) $e_1 \models_? \dot{\xi}_1$ by assumption

(18) $e_2 \not\models \dot{\xi}_2$ by assumption

(19) $e_2 \models_? \dot{\xi}_2$ by assumption

(20)
$$(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$$
 by Rule (16i) on (17) and (19)

(21)
$$(e_1, e_2) \models_7^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(22)
$$e_1 \models \dot{\xi}_1$$
 by assumption Contradicts (16).

(23)
$$(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Case $e_1 \models_? \dot{\xi}_1, e_2 \not\models_?^{\dagger} \dot{\xi}_2$.

(16)
$$e_1 \not\models \dot{\xi}_1$$
 by assumption
(17) $e_1 \models_? \dot{\xi}_1$ by assumption
(18) $e_2 \not\models \dot{\xi}_2$ by assumption
(19) $e_2 \not\models_? \dot{\xi}_2$ by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(20)
$$e_1 \models \dot{\xi}_1$$
 by assumption Contradicts (16).

(21)
$$(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$$
 by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (16g).

(22)
$$e_2 \models \dot{\xi}_2$$
 by assumption Contradicts (18).

Case (16h).

(22)
$$e_2 \models_? \dot{\xi}_2$$
 by assumption

Contradicts (19).

Case (16i).

(22) $e_2 \models_? \dot{\xi}_2$ by assumption Contradicts (19).

(23) $(e_1, e_2) \not\models_{\stackrel{\cdot}{\underline{}}} (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

(24) $(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$ by Lemma 2.0.20 on (21) and (23)

Case $e_1 \not\models_?^\dagger \dot{\xi}_1, e_2 \models \dot{\xi}_2$.

(16) $e_1 \not\models \dot{\xi}_1$ by assumption (17) $e_1 \not\models \dot{\xi}_1$ by assumption (18) $e_2 \models \dot{\xi}_2$ by assumption (19) $e_2 \not\models \dot{\xi}_2$ by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(20) $e_1 \models \dot{\xi}_1$ by assumption Contradicts (16).

(21) $(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1, e_2) notintro by assumption Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption Contradicts (17).

Case (16h).

(22) $e_2 \models_? \dot{\xi}_2$ by assumption

Contradicts (19).

Case (16i).

(22) $e_1 \models_? \dot{\xi}_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$ by contradiction

(19)
$$e_2 \not\models_? \dot{\xi}_2$$

by assumption

Assume $(e_1, e_2) \models (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (14i).

$$(20) e_2 \models \dot{\xi}_2$$

by assumption

Contradicts (18).

(21)
$$(e_1, e_2) \not\models (\dot{\xi}_1, \dot{\xi}_2)$$

by contradiction

Assume $(e_1, e_2) \models_? (\dot{\xi}_1, \dot{\xi}_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22)
$$e_1 \models_? \dot{\xi}_1$$

by assumption

Contradicts (17).

Case (16h).

(22)
$$e_2 \models_? \dot{\xi}_2$$

by assumption

Contradicts (19).

Case (16i).

(22)
$$e_1 \models_? \dot{\xi}_1$$

by assumption

Contradicts (17).

(23)
$$(e_1, e_2) \not\models_? (\dot{\xi}_1, \dot{\xi}_2)$$

by contradiction

(24)
$$(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$

by Lemma 2.0.20 on

(21) and (23)

Definition 1.1.1 (Entailment of Constraints). Suppose that $\dot{\xi}_1 : \tau$ and $\dot{\xi}_2 : \tau$. Then $\dot{\xi}_1 \models \dot{\xi}_2$ iff for all e such that $\cdot ; \Delta \vdash e : \tau$ and e val we have $e \models_?^{\dagger} \dot{\xi}_1$ implies $e \models \dot{\xi}_2$

Definition 1.1.2 (Potential Entailment of Constraints). Suppose that $\dot{\xi}_1: \tau$ and $\dot{\xi}_2: \tau$. Then $\dot{\xi}_1 \models_?^{\dagger} \dot{\xi}_2$ iff for all e such that $\cdot; \Delta \vdash e: \tau$ and e final we have $e \models_?^{\dagger} \dot{\xi}_1$ implies $e \models_?^{\dagger} \dot{\xi}_2$

Corollary 1.1.1. Suppose that $\dot{\xi} : \tau \ and \cdot ; \Delta \vdash e : \tau \ and \ e \ final.$ Then $\top \models_{?}^{\dagger} \dot{\xi}$ implies $e \models_{?}^{\dagger} \dot{\xi}$

Proof.

(1) $\dot{\xi}:\tau$	by assumption
$(2) \ \cdot ; \Gamma \vdash e : \tau$	by assumption
(3) e final	by assumption
$(4) \top \models^{\dagger}_{?} \dot{\xi}$	by assumption
$(5) e_1 \models \top$	by Rule (14a)
$(6) e_1 \models^{\dagger}_{?} \top$	by Rule $(17b)$ on (5)
$(7) \ \top : \tau$	by Rule (8a)
$(8) e_1 \models_?^{\dagger} \dot{\xi}_r$	by Definition 2.1.2 of
	(4) on (7) and (1) and (2)
	(2) and (3) and (6)

2 Match Constraint Language

 $\begin{array}{ll} \xi & ::= & \top \mid \bot \mid \underline{n} \mid \underline{\varkappa} \mid \xi_1 \wedge \xi_2 \mid \xi_1 \vee \xi_2 \mid \mathtt{inl}(\xi) \mid \mathtt{inr}(\xi) \mid (\xi_1, \xi_2) \\ \hline \xi : \tau & \xi \text{ constrains final expressions of type } \tau \end{array}$

 $\frac{\text{CTTruth}}{\top : \tau}$

CTFalsity
——— (8b)

(8a)

 $\frac{}{\perp : \tau} \tag{8b}$

 $\frac{\text{CTNum}}{\underline{n}: \text{num}} \tag{8c}$

CTNotNum

 $\underline{\mathscr{H}}: \mathtt{num}$ (8d)

CTAnd

 $\frac{\xi_1:\tau\quad \xi_2:\tau}{\xi_1\wedge\xi_2:\tau} \tag{8e}$

CTOr $\frac{\xi_1 : \tau \qquad \xi_2 : \tau}{\xi_1 \vee \xi_2 : \tau}$ (8f)

CTInl

 $\frac{\xi_1:\tau_1}{\mathtt{inl}(\xi_1):(\tau_1+\tau_2)} \tag{8g}$

CTInr

$$\frac{\xi_2:\tau_2}{\operatorname{inr}(\xi_2):(\tau_1+\tau_2)}\tag{8h}$$

CTPair
$$\frac{\xi_1:\tau_1\qquad \xi_2:\tau_2}{(\xi_1,\xi_2):(\tau_1\times\tau_2)} \tag{8i}$$

 $|\overline{\xi_1} = \xi_2|$ dual of ξ_1 is ξ_2

$$\overline{\top} = \bot$$
 (9a)

$$\overline{\perp} = \top$$
 (9b)

$$\underline{\overline{n}} = \underline{\varkappa} \tag{9c}$$

$$\underline{\varkappa} = \underline{n}$$
 (9d)

$$\overline{\xi_1 \wedge \xi_2} = \overline{\xi_1} \vee \overline{\xi_2} \tag{9e}$$

$$\overline{\xi_1 \vee \xi_2} = \overline{\xi_1} \wedge \overline{\xi_2} \tag{9f}$$

$$\overline{\operatorname{inl}(\xi_1)} = \operatorname{inl}(\overline{\xi_1}) \vee \operatorname{inr}(\top) \tag{9g}$$

$$\overline{\operatorname{inr}(\xi_2)} = \operatorname{inr}(\overline{\xi_2}) \vee \operatorname{inl}(\top) \tag{9h}$$

$$\overline{(\xi_1, \xi_2)} = (\xi_1, \overline{\xi_2}) \vee (\overline{\xi_1}, \xi_2) \vee (\overline{\xi_1}, \overline{\xi_2}) \tag{9i}$$

 ξ refutable? ξ is refutable

RXNum

$$\frac{}{\underline{n} \text{ refutable}_?}$$
 (10a)

RXInl

$$\frac{}{\mathsf{inl}(\xi)\,\mathtt{refutable}_?}\tag{10b}$$

RXInr

$$\frac{}{\mathsf{inr}(\xi)\;\mathsf{refutable}_?}\tag{10c}$$

RXPairL

$$\frac{\xi_1 \text{ refutable}?}{(\xi_1, \xi_2) \text{ refutable}?} \tag{10d}$$

RXPairR

$$\frac{\xi_2 \text{ refutable}?}{(\xi_1, \xi_2) \text{ refutable}?} \tag{10e}$$

RXO

$$\frac{\xi_1 \text{ refutable}?}{\xi_1 \vee \xi_2 \text{ refutable}?} \tag{10f}$$

 $refutable_{?}(\xi)$

$$refutable_{?}(\underline{n}) = \text{true} \qquad (11a)$$

$$refutable_{?}(\underline{\mathscr{P}}) = \text{true} \qquad (11b)$$

$$refutable_{?}(\underline{\mathscr{P}}) = \text{true} \qquad (11c)$$

$$refutable_{?}(\operatorname{inl}(\xi)) = refutable_{?}(\xi) \qquad (11d)$$

$$refutable_{?}(\operatorname{inn}(\xi)) = refutable_{?}(\xi) \qquad (11d)$$

$$refutable_{?}(\operatorname{inn}(\xi)) = refutable_{?}(\xi) \qquad (11e)$$

$$refutable_{?}((\xi_{1}, \xi_{2})) = refutable_{?}(\xi_{1}) \text{ or } refutable_{?}(\xi_{2}) \qquad (11f)$$

$$refutable_{?}(\xi_{1} \vee \xi_{2}) = refutable_{?}(\xi_{1}) \text{ and } refutable_{?}(\xi_{2}) \qquad (11g)$$

$$\operatorname{Otherwise} \quad refutable_{?}(\xi) = \text{false} \qquad (12h)$$

$$\dot{\uparrow}(\xi_{1}) = \xi_{2}$$

$$\dot{\uparrow}(T) = T \qquad (12a)$$

$$\dot{\uparrow}(x) = T \qquad (12b)$$

$$\dot{\uparrow}(x) = \frac{\pi}{2} \qquad (12e)$$

$$\dot{\uparrow}(x_{1}) = \frac{\pi}{2} \qquad (12e)$$

$$\dot{\uparrow}(\xi_{1}) + \xi_{2} = \frac{\pi}{2} \qquad (12e)$$

$$\dot{\uparrow}((\xi_{1}, \xi_{2})) = \frac{\pi}{2} \qquad (12e)$$

$$\dot{\uparrow}((\xi_{1}, \xi_{2})) = \frac{\pi}{2} \qquad (12e)$$

$$\dot{\downarrow}(\xi_{1}) = \sin(\dot{\uparrow}(\xi)) \qquad (12e)$$

$$\dot{\downarrow}(\xi_{1}) = \xi_{2}$$

$$\dot{\downarrow}(T) = T \qquad (13a)$$

$$\dot{\downarrow}(\xi_{1}) = \xi_{2} \qquad (13e)$$

$$\dot{\downarrow}(\xi_{1}) = \xi_{2} \qquad (13e$$

$$e \models \xi$$
 e satisfies ξ

CSTruth
$$\overline{e \models \top} \tag{14a}$$

CSNum

$$\underline{n \models n} \tag{14b}$$

CSNotNum

$$\frac{n_1 \neq n_2}{\underline{n_1} \models \underline{p_2}} \tag{14c}$$

CSAnd

$$\frac{e \models \xi_1 \qquad e \models \xi_2}{e \models \xi_1 \land \xi_2} \tag{14d}$$

 $\frac{\text{CSOrL}}{e \models \xi_1}$ $e \models \xi_1 \lor \xi_2$ (14e)

$$\frac{e \models \xi_2}{e \models \xi_1 \lor \xi_2} \tag{14f}$$

CSInl

$$\frac{e_1 \models \xi_1}{\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)} \tag{14g}$$

CSInr
$$\frac{e_2 \models \xi_2}{\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)}$$
(14h)

CSPair
$$\begin{array}{ll}
e_1 \models \xi_1 & e_2 \models \xi_2 \\
\hline
(e_1, e_2) \models (\xi_1, \xi_2)
\end{array} (14i)$$

CSNotIntroPair

$$\frac{e \text{ notintro}}{e \mid e \mid (\xi_1, \xi_2)} \text{prr}(e) \models \xi_2 \qquad (14j)$$

 $\mathit{satisfy}(e,\xi)$

$$satisfy(\underline{n_1},\underline{n_2}) = (n_1 = n_2) \tag{15b}$$

$$satisfy(\underline{n_1},\underline{n_2}) = (n_1 \neq n_2) \tag{15c}$$

$$satisfy(\underline{n_1},\underline{n_2}) = (n_1 \neq n_2) \tag{15c}$$

$$satisfy(e, \xi_1 \land \xi_2) = satisfy(e, \xi_1) \text{ and } satisfy(e, \xi_2) \tag{15d}$$

$$satisfy(e, \xi_1 \land \xi_2) = satisfy(e, \xi_1) \text{ or } satisfy(e, \xi_2) \tag{15e}$$

$$satisfy(inl_{\tau_2}(e_1), inl(\xi_1)) = satisfy(e_1, \xi_1) \text{ or } satisfy(e, \xi_2) \tag{15f}$$

$$satisfy(inl_{\tau_2}(e_2), inr(\xi_2)) = satisfy(e_2, \xi_2) \tag{15h}$$

$$satisfy((e_1, e_2), (\xi_1, \xi_2)) = satisfy(e_1, \xi_1) \text{ and } satisfy(e_2, \xi_2) \tag{15h}$$

$$satisfy((e_1)^u, (\xi_1, \xi_2)) = satisfy(prl(((e_1)^u), \xi_1)) \text{ and } satisfy(prr(((e_1)^u), \xi_2))$$

$$(15i)$$

$$satisfy((e_1)^u, (\xi_1, \xi_2)) = satisfy(prl(e_1(e_2)), \xi_1) \text{ and } satisfy(prr((e_1(e_2)), \xi_2))$$

$$satisfy(e_1(e_2), (\xi_1, \xi_2)) = satisfy(prl(e_1(e_2)), \xi_2) \text{ (15h)}$$

$$satisfy(prr(e_1(e_2), (\xi_1, \xi_2)) = satisfy(prr(match(e)\{r^is\}), \xi_2) \text{ (15h)}$$

$$satisfy(prr(e), (\xi_1, \xi_2)) = satisfy(prr(match(e)\{r^is\}), \xi_2) \text{ (15h)}$$

$$satisfy(prr(e), (\xi_1, \xi_2)) = satisfy(prr(prr(e)), \xi_1) \text{ and } satisfy(prr(prr(e)), \xi_2) \text{ (15h)}$$

$$satisfy(prr(e), (\xi_1, \xi_2)) = satisfy(prr(prr(e)), \xi_2) \text{ (15h)}$$

$$condoting(prr(e), \xi_1, \xi_2) = satisfy(prr(prr(e), \xi_2), \xi_2) \text{ (15h)}$$

$$condoting(prr(e), \xi_1, \xi_2) = satisfy(prr(prr(e), \xi_2), \xi_2) \text{ (15h)}$$

$$condoting(prr(e),$$

(15a)

 $satisfy(e, \top) = true$

$$\frac{e_2 \models_? \xi_2}{\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)} \tag{16f}$$

${\rm CMSPairL}$

$$\frac{e_1 \models_? \xi_1 \qquad e_2 \models_\xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)} \tag{16g}$$

CMSPairR

$$\frac{e_1 \models \xi_1 \qquad e_2 \models_? \xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)}$$

$$(16h)$$

$\operatorname{CMSPair}$

$$\frac{e_1 \models_? \xi_1 \qquad e_2 \models_? \xi_2}{(e_1, e_2) \models_? (\xi_1, \xi_2)} \tag{16i}$$

 $e \models_{?}^{\dagger} \xi$ e satisfies or may satisfy ξ

CSMSMay

$$\frac{e \models_? \xi}{e \models_?^\dagger \xi} \tag{17a}$$

CSMSSat

$$\frac{e \models \xi}{e \models_{2}^{+} \xi} \tag{17b}$$

Lemma 2.0.1. $e \not\models \bot$

Proof. By rule induction over Rules (14), we notice that $e \models \bot$ is in syntactic contradiction with all rules, hence not derivable.

Lemma 2.0.2. $e \not\models_? \bot$

Proof. Assume $e \models_? \bot$. By rule induction over Rules (16) on $e \models_? \bot$, only one case applies.

Case (16b).

(1) \perp refutable?

by assumption

By rule induction over Rules (10) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \bot$ is not derivable.

Lemma 2.0.3. $e \not\models_? \top$

Proof. Assume $e \models_? \top$. By rule induction over Rules (16) on $e \models_? \top$, only one case applies.

Case (16b).

(1) \top refutable?

by assumption

By rule induction over Rules (10) on (1), no case applies due to syntactic contradiction.

Therefore, $e \models_? \top$ is not derivable.

Lemma 2.0.4. $e \not\models ?$

Proof. By rule induction over Rules (14), we notice that $e \models ?$ is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.5.
$$e \models^{\dagger}_{?} \xi \text{ iff } e \models^{\dagger}_{?} \xi \lor \bot$$

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (1) $e \models_{?}^{\dagger} \xi$

by assumption

By rule induction over Rules (17) on (1).

Case (17a).

(2) $e \models_? \xi$

by assumption

(3) $e \models_? \xi \lor \bot$

by Rule (16c) on (2)

and Lemma 2.0.1

(4) $e \models_2^{\dagger} \xi \lor \bot$

by Rule (17a) on (3)

Case (17b).

(2) $e \models \xi$

by assumption

(3) $e \models \xi \lor \bot$

by Rule (14e) on (2)

(4) $e \models_{?}^{\dagger} \xi \lor \bot$

by Rule (17b) on (3)

- 2. Necessity:
 - (1) $e \models_{?}^{\dagger} \xi \lor \bot$

by assumption

By rule induction over Rules (17) on (1).

Case (17a).

(2) $e \models_? \xi \lor \bot$

by assumption

By rule induction over Rules (16) on (2), only two of them apply.

Case (16c).

(3) $e \models_? \xi$

by assumption

(4) $e \models_2^{\dagger} \xi$

by Rule (17a) on (3)

Case (16d).

(3)
$$e \models_? \bot$$
 by assumption
(4) $e \not\models_? \bot$ by Lemma 2.0.2

(3) contradicts (4).

Case (17b).

(2)
$$e \models \xi \lor \bot$$

by assumption

By rule induction over Rules (14) on (2), only two of them apply.

Case (14e).

(3)
$$e \models \xi$$

by assumption

(4)
$$e \models_{?}^{\dagger} \xi$$

by Rule (17b) on (3)

Case (14f).

(3)
$$e \models \bot$$

by assumption

(4)
$$e \not\models \bot$$

by Lemma 2.0.1

(3) contradicts (4).

Corollary 2.0.1. $\top \models^{\dagger}_{?} \xi \text{ iff } \top \models^{\dagger}_{?} \xi \vee \bot$

Proof. By Definition 2.1.2 and Lemma 2.0.5.

Lemma 2.0.6. Suppose that $\xi_1 : \tau$ and $\xi_2 : \tau$. Then $\xi_1 \not\models \xi_2$ iff $\xi_1 \not\models \xi_2 \lor \bot$

Proof.

(1)
$$\xi_1: \tau$$
 by assumption

(2)
$$\xi_2: \tau$$
 by assumption

(3)
$$\perp : \tau$$
 by Rule (8b)

(4)
$$\xi_2 \lor \bot : \tau$$
 by Rule (8f) on (2) and (3)

Then we prove sufficiency and necessity separately.

1. Sufficiency:

(5)
$$\xi_1 \not\models \xi_2$$

by assumption

To prove $\xi_1 \not\models \xi_2 \lor \bot$, assume $\xi_1 \models \xi_2 \lor \bot$.

(6)
$$\xi_1 \models \xi_2 \lor \bot$$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \xi_{1}$ implies

(7)
$$e \models \xi_2 \lor \bot$$

by Definition 2.1.1 on

(1) and (4) and (6)

By rule induction over Rules (14) on (7).

Case (14e).

 $(8) e \models \xi_2$ $(9) \xi_1 \models \xi_2$

by assumption

by Definition 2.1.1 on

(8)

(5) contradicts (9).

Case (14f).

(8) $e \models \bot$

by assumption

(9) $e \not\models \bot$

by Lemma 2.0.1

(8) contradicts (9).

The conclusion holds as follows:

- (a) $\xi_1 \not\models \xi_2 \vee \bot$
- 2. Necessity:
 - (5) $\xi_1 \not\models \xi_2 \vee \bot$

by assumption

To prove $\xi_1 \not\models \xi_2$, assume $\xi_1 \models \xi_2$.

(6) $\xi_1 \models \xi_2$

by assumption

For all e such that \cdot ; $\Delta \vdash e : \tau$ and e final we have $e \models^{\dagger}_{?} \xi_{1}$ implies

(7) $e \models \xi_2$

by Definition 2.1.1 on

(1) and (2) and (6)

(8) $e \models \xi_2 \lor \bot$

by Rule (14e) on (7)

(9) $\xi_1 \models \xi_2 \lor \bot$

by Definition 2.1.1 on

(8)

(9) contradicts (5).

The conclusion holds as follows:

(a) $\xi_1 \not\models \xi_2$

Lemma 2.0.7. If $e \not\models_{?}^{\dagger} \xi_1$ and $e \not\models_{?}^{\dagger} \xi_2$ then $e \not\models_{?}^{\dagger} \xi_1 \vee \xi_2$

Proof. Assume, for the sake of contradiction, that $e \models_{7}^{\dagger} \xi_1 \vee \xi_2$.

 $(1) \ e \models^{\dagger}_{?} \xi_1 \lor \xi_2$

by assumption

(2) $e \not\models_2^{\dagger} \xi_1$

by assumption

(3) $e \not\models_{?}^{\dagger} \xi_2$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(4)
$$e \models \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (14) on (4) and only two of them apply.

Case (14e).

(5)
$$e \models \xi_1$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{1}$$

by Rule (17b) on (5)

(6) contradicts (2).

Case (14f).

(5)
$$e \models \xi_2$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{2}$$

by Rule (17b) on (5)

(6) contradicts (3).

Case (17a).

(4)
$$e \models_? \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (16) on (4) and only two of them apply.

Case (16c).

(5)
$$e \models_{?} \xi_{1}$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{1}$$

by Rule (17a) on (5)

(6) contradicts (2).

Case (16d).

(5)
$$e \models_{?} \xi_{2}$$

by assumption

(6)
$$e \models_{?}^{\dagger} \xi_{2}$$

by Rule (17a) on (5)

(6) contradicts (3).

The conclusion holds as follows:

1.
$$e \not\models_{?}^{\dagger} \xi_1 \vee \xi_2$$

Lemma 2.0.8. If $e \models^{\dagger}_{?} \xi_1 \lor \xi_2$ and $e \not\models^{\dagger}_{?} \xi_1$ then $e \models^{\dagger}_{?} \xi_2$

Proof.

$$(1) \ e \models^{\dagger}_{?} \xi_1 \lor \xi_2$$

by assumption

(2)
$$e \not\models_{?}^{\dagger} \xi_{1}$$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(3)
$$e \models \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (14) on (3) and only two of them apply.

Case (14e).

(4)
$$e \models \xi_1$$

by assumption

(5)
$$e \models_{?}^{\dagger} \xi_{1}$$

by Rule (17b) on (4)

(5) contradicts (2).

Case (14f).

(4)
$$e \models \xi_2$$

by assumption

(5)
$$e \models^{\dagger}_{?} \xi_{2}$$

by Rule (17b) on (4)

Case (17a).

(3)
$$e \models_{?} \xi_1 \lor \xi_2$$

by assumption

By rule induction over Rules (16) on (3) and only two of them apply.

Case (16c).

(4)
$$e \models_? \xi_1$$

by assumption

(5)
$$e \models^{\dagger}_{?} \xi_1$$

by Rule (17a) on (4)

(5) contradicts (2).

Case (16d).

(4)
$$e \models_{?} \xi_{2}$$

by assumption

(5)
$$e \models_{2}^{\dagger} \xi_{2}$$

by Rule (17a) on (4)

Lemma 2.0.9. If $e \models_{?}^{\dagger} \xi_1$ and $e \models_{?}^{\dagger} \xi_2$ then $e \models_{?}^{\dagger} \xi_1 \wedge \xi_2$

Lemma 2.0.10. If $e \models^{\dagger}_{?} \xi_1$ then $e \models^{\dagger}_{?} \xi_1 \lor \xi_2$ and $e \models^{\dagger}_{?} \xi_2 \lor \xi_1$

Proof.

$$(1) \ e \models^{\dagger}_{?} \xi_1$$

by assumption

By rule induction over Rules (17) on (1),

Case (17b).

(2)
$$e \models \xi_1$$

by assumption

(3)
$$e \models \xi_1 \lor \xi_2$$

by Rule (14e) on (2)

(4)
$$e \models \xi_2 \vee \xi_1$$

by Rule (14f) on (2)

(5)
$$e \models_{?}^{\dagger} \xi_1 \vee \xi_2$$

by Rule (17b) on (3)

(6)
$$e \models_{?}^{\dagger} \xi_2 \vee \xi_1$$

by Rule (17b) on (4)

Case (17a).

(2)
$$e \models_? \xi_1$$

by assumption

By case analysis on the result of $satisfy(e, \xi_2)$.

Case true.

(3)
$$satisfy(e, \xi_2) = true$$

by assumption

(4)
$$e \models \xi_2$$

by Lemma 2.0.19 on

(3)

(5)
$$e \models \xi_1 \lor \xi_2$$

by Rule (14f) on (4)

(6)
$$e \models \xi_2 \lor \xi_1$$

by Rule (14e) on (4)

(7)
$$e \models_{?}^{\dagger} \xi_1 \vee \xi_2$$

by Rule (17b) on (5)

(8)
$$e \models_{2}^{\dagger} \xi_{2} \vee \xi_{1}$$

by Rule (17b) on (6)

Case false.

(3)
$$satisfy(e, \xi_2) = false$$

by assumption

(4)
$$e \not\models \xi_2$$

by Lemma 2.0.19 on

(3)

(5)
$$e \models_? \xi_1 \lor \xi_2$$

by Rule (16c) on (2)

and (4)

(6)
$$e \models_2^{\dagger} \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$$

by Rule (17a) on (5)

Lemma 2.0.11. If $e_1 \models_{?}^{\dagger} \xi_1$ then $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$

Proof.

$$(1) e_1 \models^{\dagger}_? \xi_1$$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$e_1 \models \xi_1$$

by assumption

(3)
$$\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$$

by Rule (14g) on (2)

$$(4) \ \operatorname{inl}_{\tau_2}(e_1) \models^\dagger_? \operatorname{inl}(\xi_1)$$

by Rule (17b) on (3)

Case (17a).

(2)
$$e_1 \models_? \xi_1$$

by assumption

(3)
$$\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$$

by Rule (16e) on (2)

(4)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{\tau}^{\dagger} \operatorname{inl}(\xi_1)$$

by Rule (17a) on (3)

Lemma 2.0.12. If $e_2 \models_?^\dagger \xi_2$ then $\operatorname{inr}_{\tau_1}(e_2) \models_?^\dagger \operatorname{inr}(\xi_2)$

Proof.

(1)
$$e_2 \models_{?}^{\dagger} \xi_2$$

by assumption

By rule induction over Rules (17) on (1).

Case (17b).

(2)
$$e_2 \models \xi_2$$

by assumption

$$(3) \ \operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$$

by Rule (14h) on (2)

(4)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau_2}^{\dagger} \operatorname{inr}(\xi_2)$$

by Rule (17b) on (3)

Case (17a).

(2)
$$e_2 \models_? \xi_2$$

by assumption

(3)
$$\operatorname{inl}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$$

by Rule (16f) on (2)

(4)
$$\operatorname{inl}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$

by Rule (17a) on (3)

Lemma 2.0.13. If $e_1 \models_{?}^{\dagger} \xi_1$ and $e_2 \models_{?}^{\dagger} \xi_2$ then $(e_1, e_2) \models_{?}^{\dagger} (\xi_1, \xi_2)$

Lemma 2.0.14 (Soundness and Completeness of Refutable Constraints). ξ refutable? *iff* refutable?

Lemma 2.0.15. If e notintro and ξ refutable? then either $\dot{\top}(\xi)$ refutable? or $e \models \dot{\top}(\xi)$.

Proof. By structural induction on ξ .

Lemma 2.0.16. There does not exist such a constraint $\xi_1 \wedge \xi_2$ such that $\xi_1 \wedge \xi_2$ refutable?

Proof. By rule induction over Rules (10), we notice that $\xi_1 \wedge \xi_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.17. There does not exist such a constraint $\xi_1 \vee \xi_2$ such that $\xi_1 \vee \xi_2$ refutable?

Proof. By rule induction over Rules (10), we notice that $\xi_1 \vee \xi_2$ refutable? is in syntactic contradiction with all the cases, hence not derivable.

Lemma 2.0.18. If e notintro and $e \models \xi$ then ξ refutable?.

Proof.

- (1) e notintro by assumption
- (2) $e \models \xi$ by assumption

By rule induction over Rules (14) on (2).

Case (14a).

(3)
$$\xi = \top$$
 by assumption

Assume \top refutable?. By rule induction over Rules (10), no case applies due to syntactic contradiction.

Therefore, Trefutable?.

Case (14e),(14f).

- (3) $\xi = \xi_1 \vee \xi_2$ by assumption (4) $\xi_1 \vee \xi_2$ refutable? by Lemma 2.0.17
- Case (14d).
 - (3) $\xi = \xi_1 \wedge \xi_2$ by assumption (4) $\xi_1 \wedge \xi_2$ refutable? by Lemma 2.0.16
- Case (14j).
 - (3) $\xi = (\xi_1, \xi_2)$ by assumption (4) $\operatorname{prl}(e) \models \xi_1$ by assumption (5) $\operatorname{prr}(e) \models \xi_2$ by assumption (6) $\operatorname{prl}(e)$ notintro by Rule (26e) (7) $\operatorname{prr}(e)$ notintro by Rule (26f) (8) $\underline{\xi_1}$ refutable? by IH on (6) and (4) (9) $\underline{\xi_2}$ refutable? by IH on (7) and (5)

Assume (ξ_1, ξ_2) refutable? By rule induction over Rules (10) on it, only two cases apply.

Case (10d).

(10) ξ_1 refutable? by assumption Contradicts (8).

Case (10e).

(10)
$$\xi_2$$
 refutable?

by assumption

Contradicts (9).

Therefore, (ξ_1, ξ_2) refutable?

Otherwise.

(3)
$$e = \underline{n}, \text{inl}_{\tau_2}(e_1), \text{inr}_{\tau_1}(e_2), (e_1, e_2)$$
 by assumption

By rule induction over Rules (26) on (1), no case applies due to syntactic contradiction.

Lemma 2.0.19 (Soundness and Completeness of Satisfaction Judgment). $e \models \xi$ iff $satisfy(e, \xi) = true$.

Proof. We prove soundness and completeness separately.

1. Soundness:

(1)
$$e \models \xi$$
 by assumption

By rule induction over Rules (14) on (1).

Case (14a).

$$\begin{array}{ll} (2) \ \ \xi = \top & \text{by assumption} \\ (3) \ \ satisfy(e,\top) = \text{true} & \text{by Definition 15a} \end{array}$$

Case (14b).

$$\begin{array}{ll} (2) \ \ e = \underline{n} & \text{by assumption} \\ (3) \ \ \xi = \underline{n} & \text{by assumption} \\ (4) \ \ satisfy(\underline{n},\underline{n}) = (n=n) = \text{true} & \text{by Definition 15b} \end{array}$$

Case (14c).

(2)
$$e = \underline{n_1}$$
 by assumption
(3) $\xi = \underline{p_2}$ by assumption
(4) $n_1 \neq n_2$ by assumption
(5) $satisfy(n_1, p_2) = (n_1 \neq n_2) = true$ by Definition 15c on (4)

Case (14d).

(2)
$$\xi = \xi_1 \wedge \xi_2$$
 by assumption

 $\begin{array}{lll} (3) & e \models \xi_1 & \text{by assumption} \\ (4) & e \models \xi_2 & \text{by assumption} \\ (5) & \textit{satisfy}(e,\xi_1) = \text{true} & \text{by IH on (3)} \\ (6) & \textit{satisfy}(e,\xi_2) = \text{true} & \text{by IH on (4)} \\ (7) & \textit{satisfy}(e,\xi_1 \land \xi_2) = \textit{satisfy}(e,\xi_1) \text{ and } \textit{satisfy}(e,\xi_2) = \text{true} \\ & \text{by Definition 15d on} \\ & & (5) \text{ and (6)} \\ \end{array}$

Case (14e).

- (2) $\xi = \xi_1 \vee \xi_2$ by assumption
- (3) $e \models \xi_1$ by assumption
- (4) $satisfy(e, \xi_1) = true$ by IH on (3)
- (5) $satisfy(e, \xi_1 \lor \xi_2) = satisfy(e, \xi_1)$ or $satisfy(e, \xi_2) = true$ by Definition 15e on (4)

Case (14f).

- (2) $\xi = \xi_1 \vee \xi_2$ by assumption
- (3) $e \models \xi_2$ by assumption
- (4) $satisfy(e, \xi_2) = true$ by IH on (3)
- (5) $satisfy(e, \xi_1 \lor \xi_2) = satisfy(e, \xi_1)$ or $satisfy(e, \xi_2) = true$ by Definition 15e on (4)

Case (14g).

- (2) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (3) $\xi = inl(\xi_1)$ by assumption
- (4) $e_1 \models \xi_1$ by assumption
- (5) $satisfy(e_1, \xi_1) = true$ by IH on (4)
- (6) $satisfy(inl_{\tau_2}(e_1), inl(\xi_1)) = satisfy(e_1, \xi_1) = true$ by Definition 15f on (5)

Case (14h).

- (2) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption
- (3) $\xi = inl(\xi_2)$ by assumption
- (4) $e_2 \models \xi_2$ by assumption
- (5) $satisfy(e_2, \xi_2) = true$ by IH on (4)
- (6) $satisfy(inr_{\tau_1}(e_2), inr(\xi_2)) = satisfy(e_2, \xi_2) = true$

by Definition 15g on (5)

Case (14i).

$$\begin{array}{lll} (2) & e=(e_1,e_2) & & \text{by assumption} \\ (3) & \xi=(\xi_1,\xi_2) & & \text{by assumption} \\ (4) & e_1 \models \xi_1 & & \text{by assumption} \\ (5) & e_2 \models \xi_2 & & \text{by assumption} \\ (6) & \textit{satisfy}(e_1,\xi_1) = \text{true} & & \text{by IH on (4)} \\ (7) & \textit{satisfy}(e_2,\xi_2) = \text{true} & & \text{by IH on (5)} \\ \end{array}$$

(8)
$$satisfy((e_1,e_2),(\xi_1,\xi_2))=satisfy(e_1,\xi_1)$$
 and $satisfy(e_2,\xi_2)=$ true by Definition 15h on (6) and (7)

Case (14j).

(2)
$$\xi = (\xi_1, \xi_2)$$
 by assumption
(3) e notintro by assumption
(4) $prl(e) \models \xi_1$ by assumption
(5) $prr(e) \models \xi_2$ by assumption
(6) $satisfy(prl(e), \xi_1) = true$ by IH on (4)
(7) $satisfy(prr(e), \xi_2) = true$ by IH on (5)

By rule induction over Rules (26) on (3).

Otherwise.

(8)
$$e = (||)^u, (|e_0|)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$
 by assumption

(9)
$$satisfy(e, (\xi_1, \xi_2)) = satisfy(\texttt{prl}(e), \xi_1)$$
 and $satisfy(\texttt{prr}(e), \xi_2) = \text{true}$ by Definition 15 on (6) and (7)

2. Completeness:

(1)
$$satisfy(e, \xi) = true$$
 by assumption

By structural induction on ξ .

Case
$$\xi = \top$$
.

(2)
$$e \models \top$$
 by Rule (14a)

Case
$$\xi = \bot$$
,?.

(2)
$$satisfy(e, \xi) = false$$
 by Definition 150

(2) contradicts (1) and thus vacuously true.

Case
$$\xi = \underline{n}$$
.

By structural induction on e.

Case $e = \underline{n'}$.

- (2) n' = n by Definition 15b on (1)
- (3) $\underline{n'} \models \underline{n}$ by Rule (14b) on (2)

Otherwise.

- (2) $satisfy(e, \underline{n}) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\xi = \underline{\varkappa}$.

By structural induction on e.

Case $e = \underline{n'}$.

- (2) $n' \neq n$ by Definition 15c on (1)
- (3) $\underline{n'} \models \underline{\varkappa}$ by Rule (14c) on (2)

Otherwise.

- (2) $satisfy(e, \mathbf{x}) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\xi = \xi_1 \wedge \xi_2$.

- (2) $satisfy(e, \xi_1) = true$ by Definition 15d on
- (3) $satisfy(e, \xi_2) = true$ by Definition 15d on (1)
- (4) $e \models \xi_1$ by IH on (2)
- (5) $e \models \xi_2$ by IH on (3)
- (6) $e \models \xi_1 \land \xi_2$ by Rule (14d) on (4) and (5)

Case $\xi = \xi_1 \vee \xi_2$.

(2) $\operatorname{satisfy}(e,\xi_1)$ or $\operatorname{satisfy}(e,\xi_2)=\operatorname{true}$

by Definition 15e on (1)

By case analysis on (2).

Case $satisfy(e, \xi_1) = true.$

- (3) $satisfy(e, \xi_1) = true$ by assumption (4) $e \models \xi_1$ by IH on (3)
- (5) $e \models \xi_1 \lor \xi_2$ by Rule (14e) on (4)

Case $satisfy(e, \xi_2) = true.$

- (3) $satisfy(e, \xi_2) = true$ by assumption (4) $e \models \xi_2$ by IH on (3)
- (5) $e \models \xi_1 \lor \xi_2$ by Rule (14f) on (4)

Case $\xi = inl(\xi_1)$.

By structural induction on e.

Case $e = \operatorname{inl}_{\tau_2}(e_1)$.

- (2) $satisfy(e_1, \xi_1) = true$ by Definition 15f on (1)
- (3) $e_1 \models \xi_1$ by IH on (2)
- (4) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (14g) on (3)

Otherwise.

- (2) $satisfy(e, inl(\xi_1)) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\xi = inr(\xi_2)$.

By structural induction on e.

Case $e = \operatorname{inr}_{\tau_1}(e_2)$.

- (2) $satisfy(e_2, \xi_2) = true$ by Definition 15g on (1)
- (3) $e_2 \models \xi_2$ by IH on (2)
- (4) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ by Rule (14h) on (3)

Otherwise.

- (2) $satisfy(e, inr(\xi_2)) = false$ by Definition 150
- (2) contradicts (1) and thus vacuously true.

Case $\xi = (\xi_1, \xi_2)$.

By structural induction on e.

Case $e = (e_1, e_2)$.

- (2) $satisfy(e_1, \xi_1) = true$ by Definition 15h on (1)
- (3) $satisfy(e_2, \xi_2) = true$ by Definition 15h on (1)
- (4) $e_1 \models \xi_1$ by IH on (2)
- (5) $e_2 \models \xi_2$ by IH on (3)
- (6) $(e_1, e_2) \models (\xi_1, \xi_2)$ by Rule (14i) on (4) and (5)

Case $e = (||u|, ||e_0||u|, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}.$

- (2) $satisfy(prl(e), \xi_1) = true$ by Definition 15h on (1)
- (3) $satisfy(prr(e), \xi_2) = true$ by Definition 15h on (1)
- (4) $\operatorname{prl}(e) \models \xi_1$ by IH on (2)
- (5) $prr(e) \models \xi_2$ by IH on (3)
- (6) e notintro by each rule in Rules (26)

(7)
$$(e_1, e_2) \models (\xi_1, \xi_2)$$

by Rule (14j) on (6) and (4) and (5)

Otherwise.

(2) $satisfy(e, (\xi_1, \xi_2)) = false$

by Definition 15o

(2) contradicts (1) and thus vacuously true.

Lemma 2.0.20. $e \not\models \xi$ and $e \not\models_? \xi$ iff $e \not\models_?^{\dagger} \xi$.

Proof. 1. Sufficiency:

(1) $e \not\models \xi$

by assumption

(2) $e \not\models_? \xi$

by assumption

Assume $e \models_{?}^{\dagger} \xi$. By rule induction over Rules (17) on it.

Case (17a).

(3) $e \models \xi$

by assumption

Contradicts (1).

Case (17b).

(3) $e \models_? \xi$

by assumption

Contradicts (2).

Therefore, $e \models^{\dagger}_{?} \xi$ is not derivable.

2. Necessity:

(1) $e \not\models_?^\dagger \xi$

by assumption

Assume $e \models \xi$.

(2) $e \models^{\dagger}_{?} \xi$

by Rule (17b) on assumption

Contradicts (1). Therefore, $e \not\models \xi$. Assume $e \models_? \xi$.

 $(3) \ e \models^\dagger_? \xi$

by Rule (17a) on assumption

Contradicts (1). Therefore, $e \not\models_? \xi$.

Theorem 2.1 (Exclusiveness of Satisfaction Judgment). If $\xi : \tau$ and $\cdot ; \Delta \vdash e : \tau$ and e final then exactly one of the following holds

1. $e \models \xi$

$$2. e \models_? \xi$$

3.
$$e \not\models_?^\dagger \xi$$

Proof.

- (4) $\xi:\tau$ by assumption
- (5) \cdot ; $\Delta \vdash e : \tau$ by assumption
- (6) e final by assumption

By rule induction over Rules (8) on (4), we would show one conclusion is derivable while the other two are not.

Case (8a).

- (7) $\xi = \top$ by assumption
- (8) $e \models \top$ by Rule (14a)
- (9) $e \not\models_? \top$ by Lemma 2.0.3
- (10) $e \models^{\dagger}_{?} \top$ by Rule (17b) on (8)

Case (8b).

- (7) $\xi = \bot$ by assumption
- (8) $e \not\models \bot$ by Lemma 2.0.1 (9) $e \not\models_? \bot$ by Lemma 2.0.2
- (10) $e \not\models_2^{\dagger} \bot$ by Lemma 2.0.20 on
 - (8) and (9)

Case (1b).

- (7) $\xi = ?$ by assumption
- (8) $e \not\models$? by Lemma 2.0.4
- (9) $e \models_?$? by Rule (16a)
- (10) $e \models^{\dagger}_{?}$? by Rule (17a) on (9)

Case (8c).

- (7) $\xi = \underline{n_2}$ by assumption
- (8) $\tau = \text{num}$ by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(9)
$$e = (v_0)^u, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$

by assumption

(10)
$$e$$
 notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

Assume $e \models \underline{n_2}$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction on ξ .

(11)
$$e \not\models \underline{n_2}$$
 by contradiction (12) n_2 refutable? by Rule (10a)

(13)
$$\overline{e} \models_{?} \underline{n_2}$$
 by Rule (16b) on (10) and (12)

(14)
$$e \models_{?}^{\dagger} n_2$$
 by Rule (17a) on (13)

Case (19d).

(9)
$$e = n_1$$
 by assumption

Assume $\underline{n_1} \models_? \underline{n_2}$. By rule induction over Rules (16), only one case applies.

Case (16b).

(10)
$$\underline{n_1}$$
 notintro by assumption

Contradicts Lemma 4.0.5.

(11)
$$\underline{n_1} \not\models_? \underline{n_2}$$
 by contradiction

By case analysis on whether n_1 is equal to n_2 .

Case $n_1 = n_2$.

(12)
$$satisfy(\underline{n_1},\underline{n_2}) = true$$
 by Definition 15
(13) $\underline{n_1} \models \underline{n_2}$ by Lemma 2.0.19 on (12)

(14)
$$e \models_{7}^{\dagger} n_2$$
 by Rule (17b) on (13)

Case $n_1 \neq n_2$.

(12)
$$satisfy(\underline{n_1},\underline{n_2}) = false$$
 by Definition 15
(13) $\underline{n_1} \not\models \underline{n_2}$ by Lemma 2.0.19 on
(12)

(14)
$$e \not\models_{?}^{\dagger} \underline{n_2}$$
 by Lemma 2.0.20 on (11) and (13)

Case (8f).

(7)
$$\xi = \xi_1 \vee \xi_2$$
 by assumption

By inductive hypothesis on (5) and (6), exactly one of $e \models \xi_1$, $e \models_? \xi_1$, and $e \not\models_? \xi_1$ holds. The same goes for ξ_2 . By case analysis on which conclusion holds for ξ_1 and ξ_2 .

Case $e \models \xi_1, e \models \xi_2$.

(8) $e \models \xi_1$	by assumption
$(9) \ e \not\models_? \xi_1$	by assumption
$(10) \ e \models \xi_2$	by assumption
(11) $e \not\models_? \xi_2$	by assumption
$(12) \ e \models \xi_1 \vee \xi_2$	by Rule $(14e)$ on (8)
$(13) \ e \models_{?}^{\dagger} \xi_1 \vee \xi_2$	by Rule (17b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9).

Case (16d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \models \xi_1, e \models_? \xi_2$.

(8) $e \models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \models_? \xi_2$ by assumption (12) $e \models \xi_1 \lor \xi_2$ by Rule (14e) on (8) (13) $e \models_?^\dagger \xi_1 \lor \xi_2$ by Rule (17b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9). Case (16d). (14) $e \not\models \xi_1$ by assumption Contradicts (8). (15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction Case $e \models \xi_1, e \not\models_{?}^{\dagger} \xi_2$. (8) $e \models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption (12) $e \models \xi_1 \lor \xi_2$ by Rule (14e) on (8) (13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (17b) on (12) Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (16) on it, the following cases apply. Case (16b). (14) $\xi_1 \vee \xi_2$ refutable? by assumption Contradicts Lemma 2.0.17. Case (16c). (14) $e \models_{?} \xi_{1}$ by assumption Contradicts (9). Case (16d). (14) $e \not\models \xi_1$ by assumption Contradicts (8). (15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction Case $e \models_? \xi_1, e \models \xi_2$. (8) $e \not\models \xi_1$ by assumption (9) $e \models_{?} \xi_{1}$ by assumption (10) $e \models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption (12) $e \models \xi_1 \vee \xi_2$ by Rule (14f) on (10)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (16) on it, the following cases apply.

by Rule (17b) on (12)

Case (16b).

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \not\models \xi_2$ by assumption

Contradicts (10).

Case (16d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \models_? \xi_1, e \models_? \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \models_{?} \xi_1$ by assumption

(10) $e \not\models \xi_2$ by assumption (11) $e \models_? \xi_2$ by assumption

(12) $e \models_{?} \xi_1 \lor \xi_2$ by Rule (16c) on (9)

and (10)

(13) $e \models^{\dagger}_{?} \xi_1 \vee \xi_2$ by Rule (17a) on (12)

Assume $e \models \xi_1 \lor \xi_2$. By rule induction over Rules (14), only two cases apply.

Case (14e).

(14) $e \models \xi_1$ by assumption

Contradicts (8)

Case (14f).

(14) $e \models \xi_2$ by assumption

Contradicts (10)

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \models_? \xi_1, e \not\models_?^{\dagger} \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \models_? \xi_1$ by assumption

(10) $e \not\models \xi_2$ by assumption

(11) $e \not\models_? \xi_2$ by assumption

(12) $e \models_{?} \xi_1 \lor \xi_2$ by Rule (16c) on (9)

and (10)

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (17a) on (12)

Assume $e \models \xi_1 \vee \xi_2$. By rule induction over Rules (14), only two cases apply.

Case (14e).

(14) $e \models \xi_1$ by assumption

Contradicts (8).

Case (14f).

(14) $e \models \xi_2$ by assumption

Contradicts (10).

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \models \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \not\models_? \xi_1$ by assumption

(10) $e \models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption

(12) $e \models \xi_1 \lor \xi_2$ by Rule (14f) on (10)

(13) $e \models_{7}^{\dagger} \xi_{1} \lor \xi_{2}$ by Rule (17b) on (12)

Assume $e \models_? \xi_1 \lor \xi_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \not\models \xi_2$ by assumption

Contradicts (10).

Case (16d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \models_? \xi_2$.

(8) $e \not\models \xi_1$ by assumption

(9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption

(11) $e \models_? \xi_2$ by assumption

(12) $e \models_{?} \xi_1 \lor \xi_2$ by Rule (16d) on (11)

and (8)

(13) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$ by Rule (17a) on (12)

Assume $e \models \xi_1 \vee \xi_2$. By rule induction over Rules (14), only two cases apply.

Case (14e).

(14) $e \models \xi_1$ by assumption

Contradicts (8)

Case (14f).

(14) $e \models \xi_2$ by assumption

Contradicts (10)

(15) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Case $e \not\models_?^\dagger \xi_1, e \not\models_?^\dagger \xi_2$.

(8) $e \not\models \xi_1$ by assumption (9) $e \not\models_? \xi_1$ by assumption (10) $e \not\models \xi_2$ by assumption (11) $e \not\models_? \xi_2$ by assumption

Assume $e \models \xi_1 \lor \xi_2$. By rule induction over Rules (14) on it, only two cases apply.

Case (14e).

(12) $e \models \xi_1$ by assumption

Contradicts (8).

Case (14f).

(12) $e \models \xi_2$ by assumption

Contradicts (10).

(13) $e \not\models \xi_1 \vee \xi_2$ by contradiction

Assume $e \models_{?} \xi_1 \vee \xi_2$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(14) $\xi_1 \vee \xi_2$ refutable? by assumption

Contradicts Lemma 2.0.17.

Case (16c).

(14) $e \models_? \xi_1$ by assumption

Contradicts (9).

Case (16d).

(14) $e \models_? \xi_2$ by assumption

Contradicts (11).

(15) $e \not\models_? \xi_1 \lor \xi_2$ by contradiction

(16)
$$e \not\models_?^{\dagger} \xi_1 \vee \xi_2$$
 by Lemma 2.0.20 on (13) and (15)

Case (8g).

(7)
$$\xi = \text{inl}(\xi_1)$$
 by assumption
(8) $\tau = (\tau_1 + \tau_2)$ by assumption
(9) $\xi_1 : \tau_1$ by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(10)
$$e = \langle | \rangle^u, \langle | e_0 \rangle^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{ \hat{rs} \}$$
 by assumption

(11) e notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

Assume $e \models \mathtt{inl}(\xi_1)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inl}(\xi_1)$$
 by contradiction

By case analysis on the value of $refutable_{?}(inl(\xi_1))$.

Case $refutable_{?}(inl(\xi_1)) = true.$

(13) $refutable_{?}(inl(\xi_{1})) = true$ by assumption (14) $inl(\xi_{1})$ refutable_? by Lemma 2.0.14 on (13) (15) $e \models_{?} inl(\xi_{1})$ by Rule (16b) on (11) and (14) (16) $e \models_{?}^{\dagger} inl(\xi_{1})$ by Rule (17a) on (15)

Case $refutable_{?}(inl(\xi_1)) = false.$

(13) $refutable_?(inl(\xi_1)) = false$ by assumption (14) $inl(\xi_1)$ $refutable_?$ by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(15)
$$\operatorname{inl}(\xi_1)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inl}(\xi_1)$$
 by contradiction
(17) $e \not\models_?^{\dagger} \operatorname{inl}(\xi_1)$ by Lemma 2.0.20 on
(12) and (16)

Case (19j).

- (10) $e = \operatorname{inl}_{\tau_2}(e_1)$ by assumption
- (11) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption
- (12) e_1 final by Lemma 4.0.2 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_1 \models \xi_1$, $e_1 \models_? \xi_1$, and $e_1 \not\models_?^{\dagger} \xi_1$ holds. By case analysis on which one holds.

Case $e_1 \models \xi_1$.

- (13) $e_1 \models \xi_1$ by assumption
- (14) $e_1 \not\models_? \xi_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (14g) on (13)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Rule (17b) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (16) on it, only two cases apply.

Case (16b).

- (17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption
- By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16e).

(17) $e_1 \models_? \xi_1$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\xi_1)$$
 by contradiction

Case $e_1 \models_? \xi_1$.

- (13) $e_1 \not\models \xi_1$ by assumption
- (14) $e_1 \models_? \xi_1$ by assumption
- (15) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by Rule (16e) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Rule (17a) on (15)

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (14) on it, only one case applies.

Case (14g).

(17)
$$e_1 \models \xi_1$$

Contradicts (13).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\xi_1)$$
 by contradiction

Case $e_1 \not\models_{?}^{\dagger} \xi_1$.

(13)
$$e_1 \not\models \xi_1$$
 by assumption

(14)
$$e_1 \not\models_? \xi_1$$

by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (14) on it, only one case applies.

Case (14g).

(15)
$$e_1 \models \xi_1$$

Contradicts (13).

(16)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inl}(\xi_1)$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(17) $\operatorname{inl}_{\tau_2}(e_1)$ notintro

by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16e).

(17)
$$e_1 \models_? \xi_1$$

Contradicts (14).

(18)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_? \operatorname{inl}(\xi_1)$$

by contradiction

(19)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_?^{\dagger} \operatorname{inl}(\xi_1)$$

by Lemma $2.0.20~\rm on$

(16) and (18)

Case (19k).

(10)
$$e = inr_{\tau_1}(e_2)$$

by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inl}(\xi_1)$$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inl}(\xi_1)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(12)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro

by assumption

By rule induction over Rules (26) on (12), no case applies due to syntactic contradiction.

(13)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inl}(\xi_1)$$

by contradiction

$$(14) \ \operatorname{inr}_{\tau_1}(e_2) \not\models_?^\dagger \operatorname{inl}(\xi_1)$$

by Lemma 2.0.20 on

(11) and (13)

Case (8h).

(7)
$$\xi = \operatorname{inr}(\xi_2)$$

by assumption

(8)
$$\tau = (\tau_1 + \tau_2)$$

by assumption

(9) $\xi_2 : \tau_2$

by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

$$(10) \ e = ()^u, (e_0)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{ \hat{rs} \}$$

by assumption

(11)
$$e$$
 notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

Assume $e \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(12)
$$e \not\models \operatorname{inr}(\xi_2)$$

by contradiction

By case analysis on the value of $refutable_{?}(inr(\xi_2))$.

inr is refutable

Case $refutable_{?}(inr(\xi_2)) = true.$

(13)
$$refutable_2(inr(\xi_2)) = true$$
 by assumption

(14)
$$inr(\xi_2)$$
 refutable? by Lemma 2.0.14 on

$$(13)$$

(15)
$$e \models_? \operatorname{inr}(\xi_2)$$
 by Rule (16b) on (11) and (14)

Case
$$refutable_7(inr(\xi_2)) = false$$
.

(16) $e \models_2^{\dagger} \operatorname{inr}(\xi_2)$

(13)
$$refutable_?(inr(\xi_2)) = false$$
 by assumption

(14)
$$\underline{\operatorname{inr}(\xi_2)}$$
 refutable? by Lemma 2.0.14 on (13)

Assume $e \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(15)
$$\operatorname{inr}(\xi_2)$$
 refutable? by assumption Contradicts (14).

(16)
$$e \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction
(17) $e \not\models_?^{\dagger} \operatorname{inr}(\xi_2)$ by Lemma 2.0.20 on
(12) and (16)

Case (19j).

(10)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

Assume $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (14) on it, no case applies due to syntactic contradiction.

(11)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models \operatorname{inr}(\xi_2)$$

by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(12)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro by assumption

By rule induction over Rules (26) on (12), no case applies due to syntactic contradiction.

(13)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{\uparrow} \operatorname{inr}(\xi_2)$$
 by contradiction

(14)
$$\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Lemma 2.0.20 on (11) and (13)

Case (19k).

(10)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption

(11)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$ by assumption

(12)
$$e_2$$
 final by Lemma 4.0.3 on (6)

By inductive hypothesis on (9) and (11) and (12), exactly one of $e_2 \models \xi_2$, $e_2 \models_? \xi_2$, and $e_2 \not\models_?^{\dagger} \xi_2$ holds. By case analysis on which one holds.

Case $e_2 \models \xi_2$.

(13)
$$e_2 \models \xi_2$$
 by assumption

(14)
$$e_2 \not\models_? \xi_2$$
 by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$$
 by Rule (14g) on (13)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Rule (17b) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (16) on it, only two cases apply.

Case (16b).

(17)
$$\operatorname{inr}_{\tau_1}(e_2)$$
 notintro by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16f).

(17)
$$e_2 \models_? \xi_2$$

Contradicts (14).

(18)
$$\operatorname{inr}_{\tau_1}(e_2) \not\models_? \operatorname{inr}(\xi_2)$$
 by contradiction

Case $e_2 \models_? \xi_2$.

(13)
$$e_2 \not\models \xi_2$$
 by assumption

(14)
$$e_2 \models_? \xi_2$$
 by assumption

(15)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{?} \operatorname{inr}(\xi_2)$$
 by Rule (16f) on (14)

(16)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Rule (17a) on (15)

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14h).

$$(17) \ e_2 \models \xi_2$$

Contradicts (13).

(18) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\xi_2)$ by contradiction

Case $e_2 \not\models_?^\dagger \xi_2$.

(13) $e_2 \not\models \xi_2$ by assumption

(14) $e_2 \not\models_? \xi_2$ by assumption

Assume $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14h).

(15)
$$e_2 \models \xi_2$$

Contradicts (13).

(16) $\operatorname{inr}_{\tau_1}(e_2) \not\models \operatorname{inr}(\xi_2)$ by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(17) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption

By rule induction over Rules (26) on (17), no case applies due to syntactic contradiction.

Case (16f).

(17)
$$e_2 \models_? \xi_2$$

Contradicts (14).

(18) $\operatorname{inr}_{\tau_1}(e_2) \not\models_{?} \operatorname{inr}(\xi_2)$ by contradiction

(19) $\operatorname{inl}_{\tau_2}(e_1) \not\models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Lemma 2.0.20 on (16) and (18)

Case (14i).

(7) $\xi = (\xi_1, \xi_2)$ by assumption

(8) $\tau = (\tau_1 \times \tau_2)$ by assumption

(9) $\xi_1: \tau_1$ by assumption

(10) $\xi_2 : \tau_2$ by assumption

By rule induction over Rules (19) on (5), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

```
(11) e = \{ \|u, \|e_0\|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \} \}
                                                             by assumption
   (12) e notintro
                                                             by Rule
                                                             (26a),(26b),(26c),(26d),(26e),(26f)
    (13) e indet
                                                             by Lemma 4.0.9 on (6)
                                                             and (12)
   (14) prl(e) indet
                                                             by Rule (24g) on (13)
    (15) prl(e) final
                                                             by Rule (25b) on (14)
    (16) prr(e) indet
                                                             by Rule (24h) on (13)
    (17) prr(e) final
                                                             by Rule (25b) on (16)
    (18) \cdot; \Delta \vdash \mathtt{prl}(e) : \tau_1
                                                             by Rule (19h) on (5)
    (19) \cdot; \Delta \vdash \mathsf{prr}(e) : \tau_2
                                                            by Rule (19i) on (5)
 By inductive hypothesis on (9) and (18) and (15), exactly one of
\operatorname{prl}(e) \models \xi_1, \operatorname{prl}(e) \models_? \xi_1, \text{ and } \operatorname{prl}(e) \not\models_?^{\dagger} \xi_1 \text{ holds.}
By inductive hypothesis on (10) and (19) and (17), exactly one of
\operatorname{prr}(e) \models \xi_2, \operatorname{prr}(e) \models_? \xi_2, \text{ and } \operatorname{prr}(e) \not\models_?^{\dagger} \xi_2 \text{ holds.}
By case analysis on which conclusion holds for \xi_1 and \xi_2.
 Case prl(e) \models \xi_1, prr(e) \models \xi_2.
         (20) prl(e) \models \xi_1
                                                             by assumption
         (21) prl(e) \not\models_? \xi_1
                                                             by assumption
                                                             by assumption
          (22) prr(e) \models \xi_2
         (23) prr(e) \not\models_? \xi_2
                                                             by assumption
         (24) e \models (\xi_1, \xi_2)
                                                             by Rule (14j) on (12)
                                                             and (20) and (22)
         (25) e \models^{\dagger}_{?} (\xi_1, \xi_2)
                                                             by Rule (17b) on (24)
         (26) (\xi_1, \xi_2) refutable?
                                                            by Lemma 2.0.18 on
                                                            (12) and (24)
       Assume e \models_? (\xi_1, \xi_2). By rule induction over Rules (16) on it,
      only one case applies.
       Case (16b).
               (27) (\xi_1, \xi_2) refutable?
                                                             by assumption
            Contradicts (26).
         (28) e \not\models_{?} (\xi_1, \xi_2)
                                                             by contradiction
 Case prl(e) \models \xi_1, prr(e) \models_? \xi_2.
         (20) prl(e) \models \xi_1
                                                             by assumption
          (21) prl(e) \not\models_? \xi_1
                                                             by assumption
          (22) prr(e) \not\models \xi_2
                                                             by assumption
```

by assumption

(23) $prr(e) \models_? \xi_2$

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24)
$$prr(e) \models \xi_2$$

by assumption

Contradicts (22)

(25)
$$e \not\models (\xi_1, \xi_2)$$

by contradiction

By rule induction over Rules (16) on (23), only one case applies.

Case (16b).

(26) ξ_2 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (10e) on (26)

assume no "or" and

"and" in

pair

(28) $e \models_{?} (\xi_1, \xi_2)$

by Rule (16b) on (12) and (27)

an

(29)
$$e \models^{\dagger}_{?} (\xi_1, \xi_2)$$

by Rule (17a) on (28)

Case $prl(e) \models \xi_1, prr(e) \not\models_{?}^{\dagger} \xi_2.$

(20) $prl(e) \models \xi_1$

by assumption

(21) $\operatorname{prl}(e) \not\models_? \xi_1$

by assumption

 $(22) \ \operatorname{prr}(e) \not\models \xi_2$

by assumption

(23) $prr(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

(24)
$$prr(e) \models \xi_2$$

by assumption

Contradicts (22).

(25)
$$e \not\models (\xi_1, \xi_2)$$

by contradiction

Assume $e \models_{?} (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(26) (ξ_1, ξ_2) refutable?

by assumption

By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

(27) ξ_1 refutable?

by assumption

(28) prl(e) notintro

by Rule (26e)

(29) $prl(e) \models_? \xi_1$

by Rule (16b) on (28)

and (27)

Contradicts (21).

Case (10e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (26f)
- (29) $prr(e) \models_? \xi_2$ by Rule (16b) on (28) and (27)

Contradicts (23).

(30) $e \not\models_{?} (\xi_{1}, \xi_{2})$ by contradiction (31) $e \not\models_{?}^{\dagger} (\xi_{1}, \xi_{2})$ by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \models_? \xi_1, prr(e) \models \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$ by assumption (21) $\operatorname{prl}(e) \models_? \xi_1$ by assumption (22) $\operatorname{prr}(e) \models \xi_2$ by assumption (23) $\operatorname{prr}(e) \not\models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

- (24) $prl(e) \models \xi_1$ by assumption Contradicts (20).
- (25) $e \not\models (\xi_1, \xi_2)$ by contradiction By rule induction over Rules (16) on (21), only one case applies.

Case (16b).

- (26) ξ_1 refutable? by assumption
- (27) (ξ_1, ξ_2) refutable? by Rule (10e) on (26)

assume no "or" and

"and" in

pair

- (28) $e \models_{?} (\xi_1, \xi_2)$ by Rule (16b) on (12) and (27)
- (29) $e \models_{?}^{\dagger} (\xi_1, \xi_2)$ by Rule (17a) on (28)

Case $prl(e) \models_? \xi_1, prr(e) \models_? \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$ by assumption (21) $\operatorname{prl}(e) \models_? \xi_1$ by assumption (22) $\operatorname{prr}(e) \not\models \xi_2$ by assumption (23) $\operatorname{prr}(e) \models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24) $prl(e) \models \xi_1$ by assumption Contradicts (20).

(25) $e \not\models (\xi_1, \xi_2)$

by contradiction

By rule induction over Rules (16) on (23), only one case applies.

Case (16b).

(26) ξ_2 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (10e) on (26)

(28) $e \models_{?} (\xi_1, \xi_2)$

by Rule (16b) on (12)

and (27)

(29) $e \models^{\dagger}_{?} (\xi_1, \xi_2)$

by Rule (17a) on (28)

Case $prl(e) \models_? \xi_1, prr(e) \not\models_?^{\dagger} \xi_2.$

(20) $\operatorname{prl}(e) \not\models \xi_1$

by assumption

(21) $prl(e) \models_? \xi_1$

by assumption by assumption

(22) $\operatorname{prr}(e) \not\models \xi_2$ (23) $\operatorname{prr}(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24) $prl(e) \models \xi_1$

by assumption

Contradicts (20)

(25) $e \not\models (\xi_1, \xi_2)$

by contradiction

By rule induction over Rules (16) on (21), only one case applies.

Case (16b).

(26) ξ_1 refutable?

by assumption by Rule (10e) on (26)

(27) (ξ_1, ξ_2) refutable? (28) $e \models_? (\xi_1, \xi_2)$

by Rule (16b) on (12)

and (27)

(29) $e \models_{?}^{\dagger} (\xi_1, \xi_2)$

by Rule (17a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \xi_1, prr(e) \models \xi_2.$

(20) $prl(e) \not\models \xi_1$

by assumption

(21) $prl(e) \not\models_? \xi_1$

by assumption

(22) $prr(e) \models \xi_2$

by assumption

(23) $prr(e) \not\models_? \xi_2$

by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

(24) $prl(e) \models \xi_1$

Contradicts (20)

by assumption

assume no
"or" and
"and" in
pair

assume no "or" and

"and" in

pair

$$(25) \quad e \not\models (\xi_1, \xi_2)$$

by contradiction

Assume $e \models_{?} (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

(26) (ξ_1, ξ_2) refutable?

by assumption

By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

- (27) ξ_1 refutable?
- by assumption
- $(28) \ \operatorname{prl}(e) \ \operatorname{notintro}$
- by Rule (26e) by Rule (16b) on (28)
- (29) $\operatorname{prl}(e) \models_? \xi_1$

and (27)

Contradicts (21).

Case (10e).

(27) ξ_2 refutable?

(29) $prr(e) \models_{?} \xi_{2}$

- by assumption
- (28) $\operatorname{prr}(e)$ notintro
- by Rule (26f) by Rule (16b) on (28)
- and (27)

Contradicts (23).

(30) $e \not\models_? (\xi_1, \xi_2)$

by contradiction

(31) $e \not\models_{?}^{\dagger} (\xi_1, \xi_2)$

by Lemma 2.0.20 on (25) and (30)

Case $prl(e) \not\models_{?}^{\dagger} \xi_1, prr(e) \models_{?} \xi_2.$

(20) $prl(e) \not\models \xi_1$

by assumption

(21) $\operatorname{prl}(e) \not\models_? \xi_1$

by assumption

(22) $prr(e) \not\models \xi_2$ (23) $prr(e) \models_? \xi_2$ by assumption by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14), only one case applies.

Case (14j).

(24) $prl(e) \models \xi_1$

by assumption

Contradicts (20).

(25) $e \not\models (\xi_1, \xi_2)$

by contradiction

By rule induction over Rules (16) on (23), only one case applies.

Case (16b).

(26) ξ_2 refutable?

by assumption

(27) (ξ_1, ξ_2) refutable?

by Rule (10e) on (26)

assume no "or" and "and" in pair

(28)
$$e \models_{?} (\xi_{1}, \xi_{2})$$
 by Rule (16b) on (12) and (27)

(29)
$$e \models^{\dagger}_{?} (\xi_1, \xi_2)$$
 by Rule (17a) on (28)

Case $prl(e) \not\models_{?}^{\dagger} \xi_1, prr(e) \not\models_{?}^{\dagger} \xi_2.$

(20) $prl(e) \not\models \xi_1$ by assumption (21) $prl(e) \not\models_? \xi_1$ by assumption (22) $prr(e) \not\models \xi_2$ by assumption (23) $prr(e) \not\models_? \xi_2$ by assumption

Assume $e \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only one case applies.

Case (14j).

(24)
$$prl(e) \models \xi_1$$
 by assumption Contradicts (20)

(25)
$$e \not\models (\xi_1, \xi_2)$$
 by contradiction

Assume $e \models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, only one case applies.

Case (16b).

- (26) (ξ_1, ξ_2) refutable? by assumption
- By rule induction over Rules (10) on (26), only two cases apply.

Case (10d).

- (27) ξ_1 refutable? by assumption (28) prl(e) notintro by Rule (26e)
- (29) $prl(e) \models_? \xi_1$ by Rule (16b) on (28) and (27)

Contradicts (21).

Case (10e).

- (27) ξ_2 refutable? by assumption (28) prr(e) notintro by Rule (26f)
- (29) $prr(e) \models_? \xi_2$ by Rule (16b) on (28) and (27)

Contradicts (23).

(30)
$$e \not\models_? (\xi_1, \xi_2)$$
 by contradiction
(31) $e \not\models_?^{\dagger} (\xi_1, \xi_2)$ by Lemma 2.0.20 on
(25) and (30)

Case (19g).

(11)
$$e = (e_1, e_2)$$
 by assumption
(12) \cdot ; $\Delta \vdash e_1 : \tau_1$ by assumption

(13) \cdot ; $\Delta \vdash e_2 : \tau_2$ by assumption (14) e_1 final by Lemma 4.0.4 on (6)

(15) e_2 final by Lemma 4.0.4 on (6)

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 \models \xi_1$, $e_1 \models_? \xi_1$, and $e_1 \models \overline{\xi_1}$ holds.

By inductive hypothesis on (10) and (13) and (15), exactly one of $e_2 \models \xi_2$, $e_2 \models_? \xi_2$, and $e_2 \models \overline{\xi_2}$ holds.

By case analysis on which conclusion holds for ξ_1 and ξ_2 .

Case $e_1 \models \xi_1, e_2 \models \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \models \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

(20) $(e_1, e_2) \models (\xi_1, \xi_2)$ by Rule (14i) on (16)

and (18)

(21) $(e_1, e_2) \models_?^{\dagger} (\xi_1, \xi_2)$ by Rule (17b) on (20)

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (16h).

(22) $e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (16i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Case $e_1 \models \xi_1, e_2 \models_? \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \not\models \xi_2$ by assumption

(19) $e_2 \models_? \xi_2$ by assumption

(20) $(e_1, e_2) \models_? (\xi_1, \xi_2)$ by Rule (16h) on (16)

and (19)

(21)
$$(e_1, e_2) \models_{?}^{\uparrow} (\xi_1, \xi_2)$$
 by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(22) $e_2 \models \xi_2$

by assumption

Contradicts (18).

(23)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$

by contradiction

Case $e_1 \models \xi_1, e_2 \not\models_{?}^{\dagger} \xi_2$.

(16) $e_1 \models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption (18) $e_2 \not\models \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(20) $e_2 \models \xi_2$

by assumption

Contradicts (18).

(21)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$

by contradiction

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \xi_1$

by assumption

Contradicts (17).

Case (16h).

(22) $e_2 \models_? \xi_2$

by assumption

Contradicts (19).

Case (16i).

(22) $e_1 \models_? \xi_1$

by assumption

Contradicts (17).

ontradiction
emma 2.0.20 on and (23)

Case $e_1 \models_? \xi_1, e_2 \models \xi_2$.

$$\begin{array}{lll} (16) & e_1 \not\models \xi_1 & & \text{by assumption} \\ (17) & e_1 \models_? \xi_1 & & \text{by assumption} \\ (18) & e_2 \models \xi_2 & & \text{by assumption} \\ (19) & e_2 \not\models_? \xi_2 & & \text{by assumption} \\ \end{array}$$

(20)
$$(e_1, e_2) \models_? (\xi_1, \xi_2)$$
 by Rule (16g) on (17) and (18)

(21)
$$(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$$
 by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(22)
$$e_1 \models \xi_1$$
 by assumption Contradicts (16).

(23)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction

Case $e_1 \models_? \xi_1, e_2 \models_? \xi_2$.

(16) $e_1 \not\models \xi_1$	by assumption
(17) $e_1 \models_? \xi_1$	by assumption
$(18) \ e_2 \not\models \xi_2$	by assumption
(19) $e_2 \models_? \xi_2$	by assumption
$(20) (e_1, e_2) \models_{?} (\xi_1, \xi_2)$	by Rule (16i) on (17) and (19)
(21) $(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$	by Rule (17a) on (20)

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (14i).

(22)
$$e_1 \models \xi_1$$
 by assumption Contradicts (16).

Case
$$e_1 \models_? \xi_1, e_2 \not\models^? \xi_2$$
.

 $(16) \ e_1 \not\models \xi_1$ by assumption

 $(17) \ e_1 \models_? \xi_1$ by assumption

 $(18) \ e_2 \not\models \xi_2$ by assumption

 $(19) \ e_2 \not\models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

 $(20) \ (e_1, e_2)$ not intro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

 $(20) \ e_1 \models \xi_1$ by assumption

Contradicts (16).

$$(21) \ (e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

 $(22) \ (e_1, e_2)$ not intro by assumption

Contradicts Lemma 4.0.8.

Case (16g).

 $(22) \ e_2 \models_? \xi_2$ by assumption

Contradicts (18).

Case (16h).

 $(22) \ e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (16i).

 $(22) \ e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (16i).

 $(22) \ e_2 \models_? \xi_2$ by assumption

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

Case (16) .

Case (16) .

 $(21) \ (e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Contradicts (19).

by assumption

(18) $e_2 \models \xi_2$

(19)
$$e_2 \not\models_? \xi_2$$
 by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1,e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(20) $e_1 \models \xi_1$ by assumption Contradicts (16).

 $(21) (e_1, e_2) \not\models (\xi_1, \xi_2)$

by contradiction

Assume $(e_1, e_2) \models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1,e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (16h).

by assumption (22) $e_2 \models_? \xi_2$

Contradicts (19).

Case (16i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

 $(24) (e_1, e_2) \not\models_{?}^{\dagger} (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Case $e_1 \not\models_?^{\dagger} \xi_1, e_2 \models_? \xi_2$.

(16) $e_1 \not\models \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \not\models \xi_2$ by assumption

(19) $e_2 \models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14j).

(20) (e_1,e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

Contradicts (18).

(21)
$$(e_1, e_2) \not\models (\xi_1, \xi_2)$$
 by contradiction

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

Case (16h).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (16).

Case (16i).

(22) $e_1 \models_? \xi_1$ by assumption

Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

(24) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Case $e_1 \not\models_? \xi_1$ by assumption

Case $e_1 \not\models_? \xi_1$ by assumption

(17) $e_1 \not\models_? \xi_1$ by assumption

(18) $e_2 \not\models_? \xi_2$ by assumption

(19) $e_2 \not\models_? \xi_2$ by assumption

Assume $(e_1, e_2) \models (\xi_1, \xi_2)$. By rule induction over Rules (14) on it, only two cases apply.

Case (14i).

(20) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (14i).

(20) $e_2 \models \xi_2$ by assumption

Contradicts (18).

(21) $(e_1, e_2) \not\models (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction

Assume $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$. By rule induction over Rules (16) on it, the following cases apply.

Case (16b).

(22)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (16g).

(22) $e_1 \models_? \xi_1$ by assumption Contradicts (17).

Case (16h).

(22) $e_2 \models_? \xi_2$ by assumption Contradicts (19).

Case (16i).

(22) $e_1 \models_? \xi_1$ by assumption Contradicts (17).

(23) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by contradiction (24) $(e_1, e_2) \not\models_? (\xi_1, \xi_2)$ by Lemma 2.0.20 on

(21) and (23)

Definition 2.1.1 (Entailment of Constraints). Suppose that $\xi_1 : \tau$ and $\xi_2 : \tau$. Then $\xi_1 \models \xi_2$ iff for all e such that $\cdot ; \Delta \vdash e : \tau$ and e val we have $e \models_{?}^{\dagger} \xi_1$ implies $e \models \xi_2$

Definition 2.1.2 (Potential Entailment of Constraints). Suppose that $\xi_1: \tau$ and $\xi_2: \tau$. Then $\xi_1 \models_{?}^{\dagger} \xi_2$ iff for all e such that $\cdot; \Delta \vdash e: \tau$ and e final we have $e \models_{?}^{\dagger} \xi_1$ implies $e \models_{?}^{\dagger} \xi_2$

Corollary 2.1.1. Suppose that $\xi : \tau \text{ and } \cdot ; \Delta \vdash e : \tau \text{ and } e \text{ final. Then } \top \models_{?}^{\dagger} \xi \text{ implies } e \models_{?}^{\dagger} \xi$

Proof.

(1) $\xi:\tau$	by assumption
$(2) \ \cdot ; \Gamma \vdash e : \tau$	by assumption
(3) e final	by assumption
$(4) \top \models^{\dagger}_{?} \xi$	by assumption
(5) $e_1 \models \top$	by Rule (14a)
(6) $e_1 \models_?^\dagger \top$	by Rule (17b) on (5)
$(7) \ \top : \tau$	by Rule (8a)
$(8) e_1 \models^{\dagger}_{?} \xi_r$	by Definition 2.1.2 of (4) on (7) and (1) and (2) and (3) and (6)

3 Static Semantics

$$\begin{array}{lll} \tau & ::= & \operatorname{num} \mid (\tau_1 \to \tau_2) \mid (\tau_1 \times \tau_2) \mid (\tau_1 + \tau_2) \\ e & ::= & x \mid \underline{n} \\ & \mid & (\lambda x \colon \tau.e) \mid e_1(e_2) \\ & \mid & (e_1, e_2) \\ & \mid & \inf_{\tau}(e) \mid \operatorname{inr}_{\tau}(e) \mid \operatorname{match}(e) \{ \hat{rs} \} \\ & \mid & \|^u \mid (e)^u \\ \hat{rs} & ::= & (rs \mid r \mid rs) \\ rs & ::= & \cdot \mid (r \mid rs') \\ r & ::= & p \Rightarrow e \\ \underline{p} & ::= & x \mid \underline{n} \mid _ \mid (p_1, p_2) \mid \operatorname{inl}(p) \mid \operatorname{inr}(p) \mid (\|^w \mid (p)^w) \\ \hline (\hat{rs})^\diamond = rs & rs \text{ can be obtained by erasing pointer from } \hat{rs} \end{array}$$

$$(\cdot \mid r \mid rs)^{\diamond} = r \mid rs \tag{18a}$$

$$((r'\mid rs')\mid r\mid rs)^{\diamond} = r'\mid (rs'\mid r\mid rs)^{\diamond} \tag{18b}$$

 $|\Gamma; \Delta \vdash e : \tau|$ e is of type τ

$$\frac{\text{TVar}}{\Gamma, x : \tau \; ; \Delta \vdash x : \tau} \tag{19a}$$

TEHole

$$\frac{\Gamma ; \Delta, u :: \tau \vdash ()^u : \tau}{\Gamma ; \Delta, u :: \tau \vdash ()^u : \tau}$$
 (19b)

THole

$$\frac{\Gamma; \Delta, u :: \tau \vdash e : \tau'}{\Gamma; \Delta, u :: \tau \vdash (e)^u : \tau}$$
(19c)

TNum

$$\frac{}{\Gamma \; ; \Delta \vdash n : \mathtt{num}} \tag{19d}$$

TLam

$$\frac{\Gamma, x : \tau_1 ; \Delta \vdash e : \tau_2}{\Gamma; \Delta \vdash (\lambda x : \tau_1.e) : (\tau_1 \to \tau_2)}$$
(19e)

TAp

$$\frac{\Gamma; \Delta \vdash e_1 : (\tau_2 \to \tau) \qquad \Gamma; \Delta \vdash e_2 : \tau_2}{\Gamma; \Delta \vdash e_1(e_2) : \tau}$$
(19f)

TPair

$$\frac{\Gamma; \Delta \vdash e_1 : \tau_1 \qquad \Gamma; \Delta \vdash e_2 : \tau_2}{\Gamma; \Delta \vdash (e_1, e_2) : (\tau_1 \times \tau_2)}$$
(19g)

$$\frac{\text{TPrl}}{\Gamma : \Delta \vdash e : (\tau_1 \times \tau_2)} \\
\frac{\Gamma : \Delta \vdash \text{prl}(e) : \tau_1}{\Gamma : \Delta \vdash \text{prl}(e) : \tau_1}$$
(19h)

$$\frac{\text{TPrr}}{\Gamma; \Delta \vdash e : (\tau_1 \times \tau_2)} \frac{\Gamma; \Delta \vdash e : (\tau_1 \times \tau_2)}{\Gamma; \Delta \vdash \text{prr}(e) : \tau_2}$$
(19i)

TInl

$$\frac{\Gamma ; \Delta \vdash e : \tau_1}{\Gamma ; \Delta \vdash \operatorname{inl}_{\tau_2}(e) : (\tau_1 + \tau_2)}$$
(19j)

TInr

$$\frac{\Gamma \; ; \; \Delta \vdash e : \tau_2}{\Gamma \; ; \; \Delta \vdash \mathsf{inr}_{\tau_1}(e) : (\tau_1 + \tau_2)} \tag{19k}$$

TMatchZPre

$$\frac{\Gamma; \Delta \vdash e : \tau \qquad \Gamma; \Delta \vdash [\bot]r \mid rs : \tau[\xi] \Rightarrow \tau' \qquad \top \models_?^{\dagger} \xi}{\Gamma; \Delta \vdash \mathsf{match}(e)\{\cdot \mid r \mid rs\} : \tau'} \tag{191}$$

 ${\bf TMatchNZPre}$

$$\frac{\Gamma ; \Delta \vdash e : \tau \quad e \text{ final} \quad \Gamma ; \Delta \vdash [\bot] r s_{pre} : \tau[\xi_{pre}] \Rightarrow \tau'}{\Gamma ; \Delta \vdash [\bot \lor \xi_{pre}] r \mid r s_{post} : \tau[\xi_{rest}] \Rightarrow \tau' \quad e \not\models_{?}^{\dagger} \xi_{pre} \quad \top \models_{?}^{\dagger} \xi_{pre} \lor \xi_{rest}} \Gamma ; \Delta \vdash \text{match}(e) \{r s_{pre} \mid r \mid r s_{post}\} : \tau'}$$

$$(19m)$$

 $p:\tau[\xi]\dashv \Gamma;\Delta$ p is assigned type τ and emits constraint ξ

PTVai

$$\frac{}{x:\tau[\top] \dashv \cdot ; x:\tau} \tag{20a}$$

PTWild

PTEHole

$$\overline{(\!()^w : \tau[?] \dashv \cdot; w :: \tau}$$

PTHole

$$\frac{p : \tau[\xi] \dashv \Gamma; \Delta}{(p)^w : \tau'[?] \dashv \Gamma; \Delta, w :: \tau'}$$
(20d)

PTNum

$$\frac{\underline{n}: \operatorname{num}[\underline{n}] \dashv \cdot; \cdot}{\underline{n}} \tag{20e}$$

PTInl

$$\frac{p:\tau_1[\xi]\dashv \Gamma;\Delta}{\mathtt{inl}(p):(\tau_1+\tau_2)[\mathtt{inl}(\xi)]\dashv \Gamma;\Delta} \tag{20f}$$

PTInr

$$\frac{p : \tau_2[\xi] \dashv \Gamma ; \Delta}{\operatorname{inr}(p) : (\tau_1 + \tau_2)[\operatorname{inr}(\xi)] \dashv \Gamma ; \Delta}$$
(20g)

PTPair

$$\frac{p_1:\tau_1[\xi_1]\dashv \Gamma_1\;;\Delta_1 \qquad p_2:\tau_2[\xi_2]\dashv \Gamma_2\;;\Delta_2}{(p_1,p_2):(\tau_1\times\tau_2)[(\xi_1,\xi_2)]\dashv \Gamma_1\uplus\Gamma_2\;;\Delta_1\uplus\Delta_2} \tag{20h}$$

$$\frac{\Gamma; \Delta \vdash (p \Rightarrow e) : \tau[\xi] \Rightarrow \tau'}{\text{CTRule}} \qquad \begin{array}{c} r \text{ transforms a final expression of type } \tau \\ \text{to a final expression of type } \tau' \\ \frac{p : \tau[\xi] \dashv \Gamma_p; \Delta_p \qquad \Gamma \uplus \Gamma_p; \Delta \uplus \Delta_p \vdash e : \tau'}{\Gamma; \Delta \vdash p \Rightarrow e : \tau[\xi] \Rightarrow \tau'} \end{aligned} \tag{21a}$$

 $\frac{\Gamma ; \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau'}{\text{CTOneRules}} \quad \begin{array}{c} rs \text{ transforms a final expression of type } \tau \\ \text{to a final expression of type } \tau' \end{array}$

$$\frac{\Gamma; \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \qquad \xi_r \not\models \xi_{pre}}{\Gamma; \Delta \vdash [\xi_{pre}](r \mid \cdot) : \tau[\xi_r] \Rightarrow \tau'}$$
(22a)

CTRules

$$\frac{\Gamma ; \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \qquad \Gamma ; \Delta \vdash [\xi_{pre} \lor \xi_r] rs : \tau[\xi_{rs}] \Rightarrow \tau' \qquad \xi_r \not\models \xi_{pre}}{\Gamma ; \Delta \vdash [\xi_{pre}] r \mid rs : \tau[\xi_r \lor \xi_{rs}] \Rightarrow \tau'}$$
(22b)

Lemma 3.0.1. If $p : \tau[\xi] \dashv \Gamma ; \Delta \text{ then } \xi : \tau$.

Proof. By rule induction over Rules
$$(20)$$
.

Lemma 3.0.2. If \cdot ; $\Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau \ then \ \xi_r : \tau_1$.

Proof. By rule induction over Rules
$$(21)$$
.

Lemma 3.0.3. If
$$\cdot$$
; $\Delta \vdash [\xi_{pre}]rs : \tau_1[\xi_{rs}] \Rightarrow \tau \ then \ \xi_{rs} : \tau_1$.

Proof. By rule induction over Rules
$$(22)$$
.

Lemma 3.0.4. If $\Gamma : \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau' \text{ and } \Gamma : \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau' \text{ and } \xi_r \not\models \xi_{pre} \lor \xi_{rs} \text{ then } \Gamma : \Delta \vdash [\xi_{pre}](rs \mid r \mid \cdot)^{\diamond} : \tau[\xi_{rs} \lor \xi_r] \Rightarrow \tau'$

Proof.

- (1) $\Gamma : \Delta \vdash [\xi_{pre}]rs : \tau[\xi_{rs}] \Rightarrow \tau'$ by assumption
- (2) $\Gamma : \Delta \vdash r : \tau[\xi_r] \Rightarrow \tau'$ by assumption
- (3) $\xi_r \not\models \xi_{pre} \vee \xi_{rs}$ by assumption

By rule induction over Rules (22) on (1).

Case (22a).

(4)
$$rs = r' \mid \cdot$$
 by assumption

(5)
$$\xi_{rs} = \xi'_r$$
 by assumption

(6)
$$\Gamma : \Delta \vdash r' : \tau[\xi_r'] \Rightarrow \tau'$$
 by assumption

(7)
$$\xi_r' \not\models \xi_{pre}$$
 by assumption

(8)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r]r \mid \cdot : \tau[\xi_r] \Rightarrow \tau'$$
 by Rule (22a) on (2) and (3)

(9)
$$\Gamma ; \Delta \vdash [\xi_{pre}](r' \mid r \mid \cdot) : \tau[\xi'_r \lor \xi_r] \Rightarrow \tau'$$
 by Rule (22b) on (6) and (8) and (7)

$$\begin{array}{ll} (10) \ \Gamma \, ; \Delta \vdash [\xi_{pre}]((r' \mid \cdot) \mid r \mid \cdot)^{\diamond} : \tau[\xi'_r \lor \xi_r] \Rightarrow \tau' \\ & \text{by Definition 18 on (9)} \end{array}$$

Case (22b).

(4)
$$rs = r' \mid rs'$$
 by assumption

(5)
$$\xi_{rs} = \xi'_r \vee \xi'_{rs}$$
 by assumption

(6)
$$\Gamma : \Delta \vdash r' : \tau[\xi'_r] \Rightarrow \tau'$$
 by assumption

(7)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r]rs' : \tau[\xi'_{rs}] \Rightarrow \tau'$$
 by assumption

(8)
$$\xi_r' \not\models \xi_{pre}$$
 by assumption

(9)
$$\Gamma : \Delta \vdash [\xi_{pre} \lor \xi'_r](rs' \mid r \mid \cdot)^{\diamond} : \tau[\xi'_{rs} \lor \xi_r] \Rightarrow \tau'$$
 by IH on (7) and (2) and (3)

(10)
$$\Gamma$$
; $\Delta \vdash [\xi_{pre}](r' \mid (rs' \mid r \mid \cdot)^{\diamond}) : \tau[\xi'_r \lor \xi'_{rs} \lor \xi_r] \Rightarrow \tau'$ by Rule (22b) on (6) and (9) and (8)

$$\begin{array}{ll} (11) & \Gamma \; ; \Delta \vdash [\xi_{pre}]((r' \mid rs') \mid r \mid \cdot)^{\diamond} : \tau[\xi'_r \lor \xi'_{rs} \lor \xi_r] \Rightarrow \tau' \\ & \text{by Definition 18 on} \\ & (10) \end{array}$$

Lemma 3.0.5 (Substitution). If $\Gamma, x : \tau ; \Delta \vdash e_0 : \tau_0 \text{ and } \Gamma ; \Delta \vdash e : \tau \text{ then } \Gamma ; \Delta \vdash [e/x]e_0 : \tau_0$

Lemma 3.0.6 (Simultaneous Substitution). *If* $\Gamma \uplus \Gamma'$; $\Delta \vdash e : \tau$ *and* $\theta : \Gamma'$ *then* $\Gamma ; \Delta \vdash [\theta]e : \tau$

Lemma 3.0.7 (Substitution Typing). *If* $e \rhd p \dashv \theta$ and \cdot ; $\Delta_e \vdash e : \tau$ and $p : \tau[\xi] \dashv \Gamma$; Δ then $\theta : \Gamma$

Proof by induction on the derivation of $e \triangleright p \dashv \theta$.

Theorem 3.1 (Determinism). If \cdot ; $\Delta \vdash e : \tau$ then exactly one of the following holds

- 1. e val
- 2. e indet
- 3. $e \mapsto e'$ for some unique e'

4 Dynamic Semantics

e val e is a value

$$\frac{\text{VNum}}{n \text{ val}} \tag{23a}$$

VLam

$$\frac{}{(\lambda x:\tau.e)\;\mathrm{val}} \tag{23b}$$

VPair

$$\frac{e_1 \text{ val} \qquad e_2 \text{ val}}{(e_1, e_2) \text{ val}} \tag{23c}$$

VInl

$$\frac{e \text{ val}}{\text{inl}_{\tau}(e) \text{ val}} \tag{23d}$$

VInr

$$\frac{e \; \mathrm{val}}{\mathrm{inr}_{\tau}(e) \; \mathrm{val}} \tag{23e}$$

e indet e is indeterminate

IEHole

$$\sqrt{\|)^u \text{ indet}}$$

IHole

$$\frac{e \; \mathtt{final}}{(e)^u \; \mathtt{indet}} \tag{24b}$$

IAp

$$\frac{e_1 \; \mathtt{indet} \qquad e_2 \; \mathtt{final}}{e_1(e_2) \; \mathtt{indet}} \tag{24c}$$

IPairL

$$\frac{e_1 \text{ indet} \qquad e_2 \text{ val}}{(e_1, e_2) \text{ indet}} \tag{24d}$$

IPairR

$$\frac{e_1 \text{ val} \qquad e_2 \text{ indet}}{(e_1, e_2) \text{ indet}} \tag{24e}$$

IPair

$$\frac{e_1 \text{ indet} \qquad e_2 \text{ indet}}{(e_1, e_2) \text{ indet}} \tag{24f}$$

IPrl

$$\frac{e \; \mathtt{indet}}{\mathtt{prl}(e) \; \mathtt{indet}} \tag{24g}$$

```
IPrr
                                                      e \; \mathtt{indet}
                                                                                                                 (24h)
                                                 prr(e) indet
                                                IInL
                                                      e \; \mathtt{indet}
                                                                                                                  (24i)
                                                 {\tt inl}_{	au}(e) indet
                                                IInR
                                                      e \; \mathtt{indet}
                                                                                                                  (24j)
                                                 \mathtt{inr}_{	au}(e) indet
                           IMatch
                                              e \; \mathtt{final} \qquad e \; ? \; p_r
                                                                                                                 (24k)
                           \overline{\mathrm{match}(e)\{rs_{pre}\mid (p_r\Rightarrow e_r)\mid rs_{post}\} \text{ indet}}
                   e is final
e \; \mathtt{final}
                                                     FVal
                                                       e \; \mathtt{val}
                                                                                                                 (25a)
                                                      \overline{e} final
                                                     FIndet
                                                      e \; {\tt indet}
                                                                                                                 (25b)
                                                      e \; \mathtt{final}
                        e cannot be a value syntactically
e \; \mathtt{notintro}
                                                 NVEHole
                                                                                                                 (26a)

\frac{}{(\!(\!)^u \text{ notintro}\!)}

                                                NVHole
                                                                                                                 (26b)
                                                (e)^u notintro
                                               NVAp
                                                                                                                 (26c)
                                               e_1(e_2) notintro
                                         NVMatch
                                                                                                                 (26d)
                                         \mathrm{match}(e)\{\hat{rs}\} notintro
                                              NVPrl
                                                                                                                 (26e)
                                              \mathtt{prl}(e) notintro
                                              NVPrr
                                                                                                                 (26f)
                                              \overline{\text{prr}(e) \text{ notintro}}
                          for e final and \cdot; \Delta \vdash e : \tau
{\tt complete}(e)
```

$$\begin{split} \operatorname{\texttt{complete}}(e) &= \{e\} & \text{if } e \text{ val} \\ \operatorname{\texttt{complete}}(e) &= \{e' \mid e : \tau \text{ and } e \text{ val}\} & \text{if } e \text{ notintro} \end{split} \tag{27a} \end{split}$$

$$\mathtt{complete}(\mathtt{inl}_{\tau_2}(e_1)) = \{\mathtt{inl}_{\tau_2}(e_1') \mid e_1' \in \mathtt{complete}(e_1)\} \tag{27c}$$

$$\mathtt{complete}(\mathtt{inr}_{\tau_1}(e_2)) = \{\mathtt{inr}_{\tau_1}(e_2') \mid e_2' \in \mathtt{complete}(e_2)\} \tag{27d}$$

$$\texttt{complete}((e_1,e_2)) = \{(e_1',e_2') \mid e_1' \in \texttt{complete}(e_1) \text{ and } e_2' \in \texttt{complete}(e_2)\}$$
 (27e)

Lemma 4.0.1. If e indet $and \cdot ; \Delta \vdash e : \tau \ and \ \dot{\xi} : \tau \ and \ e \not\models_?^\dagger \dot{\xi} \ then \ e' \not\models_?^\dagger \dot{\xi}$ for all $e' \in \mathtt{complete}(e)$.

Proof.

- (1) e indet by assumption
- (2) \cdot ; $\Delta \vdash e : \tau$ by assumption
- (3) $\dot{\xi}:\tau$ by assumption
- (4) $e \not\models_{?}^{\dagger} \dot{\xi}$ by assumption

By rule induction over Rules (8) on (3).

Case (8a).

- (5) $\dot{\xi} = \top$ by assumption
- (6) $e \models \top$ by Rule (14a)
- (7) $e \models_2^{\dagger} \top$ by Rule (17b) on (6)

Contradicts (4).

Case (1b).

- (5) $\dot{\xi} = ?$ by assumption
- (6) $e \models_?$? by Rule (16a)
- (7) $e \models^{\dagger}_{?}$? by Rule (17a) on (6)

Contradicts (4).

Case (8c).

(5) $\dot{\xi} = \underline{n}$ by assumption (6) $\tau = \text{num}$ by assumption (7) \underline{n} refutable? by Rule (10a)

By rule induction over Rules (24) on (1).

Case (24a).		
(8)	$e = ()^u$	by assumption
(9)	$(\!(\!)^u$ notintro	by Rule (26a)
(10)	$()^u \models_? \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$\emptyset^u \models^{\dagger}_{?} \underline{n}$	by Rule $(17a)$ on (10)
Case $(24b)$		
(8)	$e = (e_1)^u$	by assumption
(9)	$(\![e_1]\!]^u$ notintro	by Rule (26b)
(10)	$(e_1)^u \models_? \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$(e_1)^u \models_?^\dagger \underline{n}$	by Rule $(17a)$ on (10)
Case (24c).		
(8)	$e = e_1(e_2)$	by assumption
(9)	$e_1(e_2)$ notintro	by Rule (26c)
(10)	$e_1(e_2) \models_? \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$e_1(e_2) \models^{\dagger}_{?} \underline{n}$	by Rule $(17a)$ on (10)
Case (24g).		
(8)	$e = \mathtt{prl}(e_1)$	by assumption
(9)	$\mathtt{prl}(e_1)$ notintro	by Rule (26e)
(10)	$\mathtt{prl}(e_1) \models_? \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$\mathtt{prl}(e_1) \models^\dagger_? \underline{n}$	by Rule $(17a)$ on (10)
Case (24h)		
(8)	$e = \mathtt{prr}(e_1)$	by assumption
(9)	$\mathtt{prr}(e_1)$ notintro	by Rule (26f)
(10)	$prr(e_1) \models_? \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$\mathtt{prr}(e_1) \models^\dagger_? \underline{n}$	by Rule (17a) on (10)
Case (24k)	•	
(8)	$e = \mathtt{match}(e_1)\{\hat{rs}\}$	by assumption
, ,	$\mathtt{match}(e_1)\{\hat{rs}\}$ notintro	by Rule (26d)
	$\mathtt{match}(e_1)\{\hat{rs}\} \models_{?} \underline{n}$	by Rule (16b) on (9) and (7)
(11)	$\mathtt{match}(e_1)\{\hat{rs}\}\models^\dagger_?\underline{n}$	by Rule (17a) on (10)

Case (24d), (24e), (24f).

(8)
$$e = (e_1, e_2)$$
 by assumption

By rule induction over Rules (19) on (2), no rule applies due to syntactic contradiction.

Case (24i).

(8)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

By rule induction over Rules (19) on (2), no rule applies due to syntactic contradiction.

Case (24j).

(8)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption

By rule induction over Rules (19) on (2), no rule applies due to syntactic contradiction.

Case (8g).

(5)
$$\dot{\xi} = \text{inl}(\dot{\xi}_1)$$
 by assumption
(6) $\tau = (\tau_1 + \tau_2)$ by assumption
(7) $\dot{\xi}_1 : \tau_1$ by assumption
(8) $\text{inl}(\dot{\xi}_1)$ refutable? by Rule (10b)

By rule induction over Rules (24) on (1).

Case (24a).

$(9) e = \emptyset^u$	by assumption
(10) ()) u notintro	by Rule (26a)
$(11) \ (\!)^u \models_? \mathtt{inl}(\dot{\xi}_1)$	by Rule (16b) on (10)
	and (8)
(12) $\mathbb{Q}^u \models_2^{\dagger} \operatorname{inl}(\dot{\xi_1})$	by Rule (17a) on (11)

Case (24b).

(9) $e = (e_1)^u$	by assumption
(10) $(e_1)^u$ notintro	by Rule (26b)
$(11) \ (e_1)^u \models_? \operatorname{inl}(\dot{\xi}_1)$	by Rule (16b) on (10)
	and (8)
$(12) \ (e_1)^u \models^{\dagger}_? \mathtt{inl}(\dot{\xi}_1)$	by Rule (17a) on (11)

Case (24c).

(9) $e = e_1(e_2)$	by assumption
(10) $e_1(e_2)$ notintro	by Rule (26c)
$(11) e_1(e_2) \models_? \operatorname{inl}(\dot{\xi}_1)$	by Rule (16b) on (10)
	and (8)
$(12) e_1(e_2) \models^{\dagger}_{?} \operatorname{inl}(\dot{\xi}_1)$	by Rule (17a) on (11)

Case (24g).

(9)
$$e = prl(e_1)$$
 by assumption
(10) $prl(e_1)$ notintro by Rule (26e)

(11)
$$\operatorname{prl}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (16b) on (10) and (8)

(12)
$$\operatorname{prl}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (17a) on (11)

Case (24h).

(9)
$$e = prr(e_1)$$
 by assumption
(10) $prr(e_1)$ notintro by Rule (26f)

(11)
$$\operatorname{prr}(e_1) \models_? \operatorname{inl}(\dot{\xi}_1)$$
 by Rule (16b) on (10) and (8)

(12)
$$\operatorname{prr}(e_1) \models_2^{\dagger} \operatorname{inl}(\dot{\xi_1})$$
 by Rule (17a) on (11)

Case (24k).

(9)
$$e = \text{match}(e_1)\{\hat{rs}\}$$
 by assumption
(10) $\text{match}(e_1)\{\hat{rs}\}$ notintro by Rule (26d)

(11)
$$\operatorname{match}(e_1)\{\hat{rs}\} \models_? \operatorname{inl}(\dot{\xi_1})$$
 by Rule (16b) on (10) and (8)

(12)
$$\operatorname{match}(e_1)\{\hat{rs}\}\models_{?}^{\dagger}\operatorname{inl}(\dot{\xi_1})$$
 by Rule (17a) on (11)

Case (24d), (24e), (24f).

(9)
$$e = (e_1, e_2)$$
 by assumption

By rule induction over Rules (24) on (1), no rule applies due to syntactic contradiction.

Case (24i).

(9)
$$e = \operatorname{inl}_{\tau_2'}(e_1)$$
 by assumption
(10) e_1 indet by assumption

By rule induction over Rules (19) on (2), only one rule applies.

Case (19j).

$$\begin{array}{ll} (11) & \tau_2' = \tau_2 & \text{by assumption} \\ (12) & \cdot ; \Delta \vdash e_1 : \tau_1 & \text{by assumption} \\ (13) & e_1 \not\models_?^\dagger \dot{\xi_1} & \text{by Lemma 2.0.11 on} \end{array}$$

(14)
$$e_1' \not\models_?^{\dagger} \dot{\xi}_1$$
 for any $e_1' \in \mathsf{complete}(e_1)$ by IH on (10) and (12) and (7) and (13)

(15)
$$e' \not\models_?^\dagger \operatorname{inl}(\dot{\xi}_1)$$
 for any $e' \in \{ \operatorname{inl}_{\tau_2}(e'_1) \mid e'_1 \in \operatorname{complete}(e_1) \}$ by Lemma 2.0.11 on (14)

$$(16) \ e' \not\models_{?}^{\dagger} \ \operatorname{inl}(\dot{\xi}_1) \ \operatorname{for \ any} \ e' \in \operatorname{complete}(\operatorname{inl}_{\tau_2}(e_1)) \\ \operatorname{by \ Equation \ } 27 \ \operatorname{on} \ (15)$$

$$\mathbf{Case} \ (24j).$$

$$(9) \ e = \operatorname{inr}_{\tau_1}(e_2) \qquad \operatorname{by \ assumption} \\ (10) \ e' \not\models_{?}^{\dagger} \ \operatorname{inl}(\dot{\xi}_1) \ \operatorname{for \ any} \ e' \in \{\operatorname{inr}_{\tau_1}(e'_2)\} \\ \operatorname{by \ Lemma \ } 1.0.18 \\ (11) \ e' \not\models_{?}^{\dagger} \ \operatorname{inl}(\dot{\xi}_1) \ \operatorname{for \ any} \ e' \in \{\operatorname{inr}_{\tau_1}(e'_2)\} \\ \operatorname{by \ Weakening \ on} \ (10) \\ (12) \ e' \not\models_{?}^{\dagger} \ \operatorname{inl}(\dot{\xi}_1) \ \operatorname{for \ any} \ e' \in \operatorname{complete}(\operatorname{inr}_{\tau_1}(e'_2)) \\ \operatorname{by \ Equation} \ 27 \ \operatorname{on} \ (11) \\ \mathbf{Case} \ (8h).$$

$$(5) \ \dot{\xi} = \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ assumption} \\ (6) \ \tau = (\tau_1 + \tau_2) \qquad \operatorname{by \ assumption} \\ (6) \ \tau = (\tau_1 + \tau_2) \qquad \operatorname{by \ assumption} \\ (8) \ \operatorname{inr}(\dot{\xi}_2) \ \operatorname{refutable}_? \qquad \operatorname{by \ Rule} \ (10c) \\ \mathbf{By \ rule \ induction \ over \ Rules} \ (24) \ \operatorname{on} \ (1).}$$

$$\mathbf{Case} \ (24a).$$

$$(9) \ e = \emptyset^u \qquad \operatorname{by \ Rule} \ (26a) \\ (11) \ \emptyset^u \models_{?} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (16b) \ \operatorname{on} \ (10) \\ \operatorname{and} \ (8) \\ (12) \ \emptyset^u \models_{?}^{\dagger} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (16b) \ \operatorname{on} \ (10) \\ \operatorname{and} \ (8) \\ (12) \ \emptyset^e_1)^u \models_{?}^{\dagger} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (16b) \ \operatorname{on} \ (10) \\ \operatorname{and} \ (8) \\ (12) \ \emptyset^e_1)^u \models_{?}^{\dagger} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (17a) \ \operatorname{on} \ (11) \\ \mathbf{Case} \ (24c).$$

$$(9) \ e = e_1(e_2) \qquad \operatorname{by \ Rule} \ (16b) \ \operatorname{on} \ (10) \\ \operatorname{and} \ (8) \\ (12) \ \theta_1)^u \models_{?}^{\dagger} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (26c) \\ (10) \ e_1(e_2) \mapsto_{?} \ \operatorname{inr}(\dot{\xi}_2) \qquad \operatorname{by \ Rule} \ (26c) \\ \operatorname{ond} \ (8) \ \operatorname{ond} \ (8) \\ \operatorname{ond} \ (8)$$

by Rule (17a) on (11)

(12) $e_1(e_2) \models_2^{\dagger} \operatorname{inr}(\dot{\xi}_2)$

Case (24g).

(9)
$$e = \operatorname{prl}(e_1)$$
 by assumption
(10) $\operatorname{prl}(e_1)$ notintro by Rule (26e)
(11) $\operatorname{prl}(e_1) \models_? \operatorname{inr}(\dot{\xi}_2)$ by Rule (16b) on (10)
and (8)

(12)
$$\operatorname{prl}(e_1) \models_{2}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 by Rule (17a) on (11)

Case (24h).

(9)
$$e = \operatorname{prr}(e_1)$$
 by assumption
(10) $\operatorname{prr}(e_1)$ notintro by Rule (26f)
(11) $\operatorname{prr}(e_1) \models_? \operatorname{inr}(\dot{\xi_2})$ by Rule (16b) on (10)
and (8)

(12)
$$\operatorname{prr}(e_1) \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 by Rule (17a) on (11)

Case (24k).

$$(9) \ e = \mathtt{match}(e_1)\{\hat{rs}\} \qquad \qquad \text{by assumption}$$

$$(10) \ \mathtt{match}(e_1)\{\hat{rs}\} \ \mathtt{notintro} \qquad \qquad \text{by Rule (26d)}$$

$$(11) \ \mathtt{match}(e_1)\{\hat{rs}\} \models_? \mathtt{inr}(\dot{\xi}_2) \qquad \qquad \text{by Rule (16b) on (10)}$$
 and (8)

(12)
$$\operatorname{match}(e_1)\{\hat{rs}\} \models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 by Rule (17a) on (11)

Case (24d), (24e), (24f).

(9)
$$e = (e_1, e_2)$$
 by assumption

By rule induction over Rules (24) on (1), no rule applies due to syntactic contradiction.

Case (24i).

(9)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

(10)
$$e' \not\models^{\dagger}_{?} \operatorname{inr}(\dot{\xi}_2)$$
 for any $e' \in \{\operatorname{inl}_{\tau_2}(e'_1)\}$

by Lemma 1.0.17

$$(11) \ e' \not\models_?^\dagger \mathtt{inr}(\dot{\xi}_2) \ \mathrm{for \ any} \ e' \in \{\mathtt{inl}_{\tau_2}(e_1') \ | \ e_1' \in \mathtt{complete}(e_1)\} \\ \mathrm{by \ Weakening \ on} \ (10)$$

(12)
$$e' \not\models_{?}^{\dagger} \operatorname{inr}(\dot{\xi}_2)$$
 for any $e' \in \operatorname{complete}(\operatorname{inl}_{\tau_2}(e_1))$ by Equation 27 on (11)

Case (24j).

(9)
$$e = \operatorname{inr}_{\tau'_1}(e_2)$$
 by assumption (10) e_2 indet by assumption

By rule induction over Rules (19) on (2), only one rule applies.

Case (19k).

(11)
$$\tau_1' = \tau_1$$
 by assumption
(12) \cdot ; $\Delta \vdash e_2 : \tau_2$ by assumption

(13)
$$e_{2} \not\models_{?}^{\dagger} \dot{\xi}_{2}$$
 by Lemma 2.0.11 on (4)

(14) $e'_{2} \not\models_{?}^{\dagger} \dot{\xi}_{2}$ for any $e'_{2} \in \text{complete}(e_{2})$ by IH on (10) and (12) and (7) and (13)

(15) $e' \not\models_{?}^{\dagger} \text{inr}(\dot{\xi}_{2})$ for any $e' \in \{\text{inr}_{\tau_{1}}(e'_{2}) \mid e'_{2} \in \text{complete}(e_{2})\}$ by Lemma 2.0.12 on (14)

(16) $e' \not\models_{?}^{\dagger} \text{inr}(\dot{\xi}_{2})$ for any $e' \in \text{complete}(\text{inr}_{\tau_{1}}(e_{2}))$ by Equation 27 on (15)

Case (8i).

$$\begin{array}{ll} (5) \ \dot{\xi} = (\dot{\xi}_1,\dot{\xi}_2) & \text{by assumption} \\ (6) \ \tau = (\tau_1 \times \tau_2) & \text{by assumption} \\ (7) \ \dot{\xi}_1 : \tau_1 & \text{by assumption} \\ (8) \ \dot{\xi}_2 : \tau_2 & \text{by assumption} \end{array}$$

By rule induction over Rules (24) on (1).

Case
$$(24a)$$
, $(24b)$, $(24c)$, $(24g)$, $(24h)$, $(24k)$.

$$(9) \ e = ()^u, (e_1)^u, e_1(e_2), \mathtt{prl}(e_1), \mathtt{prr}(e_1), \mathtt{match}(e_1) \{ \hat{rs} \}$$

by assumption

(10)
$$e$$
 notintro by Rules (26)

(11)
$$complete(e) = \{e' \mid e' \text{ val } and \ e' : (\tau_1 \times \tau_2)\}$$

by Equation 27 on (10)

(12)
$$prl(e)$$
 indet by Rule (24g) on (1)

(13)
$$prr(e)$$
 indet by Rule (24h) on (1)

(14)
$$\cdot$$
; $\Delta \vdash \operatorname{prl}(e) : \tau_1$ by Rule (19h) on (2)

(15)
$$\cdot$$
; $\Delta \vdash prr(e) : \tau_2$ by Rule (19i) on (2)

By case analysis on the result of $satisfyormay(prl(e), \dot{\xi}_1)$.

Case true.

(16)
$$satisfyormay(prl(e), \dot{\xi}_1) = true$$

(17) $\operatorname{prl}(e) \models_{?}^{\dagger} \dot{\xi}_{1}$ by assumption by Lemma 1.0.20 on (16)

Case false.

(16)
$$satisfyormay(prl(e), \dot{\xi}_1) = false$$

(17)
$$\operatorname{prl}(e) \not\models_{?}^{\dagger} \dot{\xi}_{1}$$
 by assumption by Lemma 1.0.20 on (16)

$$(18) \ e_1' \not\models_?^\dagger \dot{\xi}_1 \ \text{for any} \ e_1' \in \texttt{complete}(\texttt{prl}(e))$$

by IH on (12) and (14) and (7) and (17)

(19)
$$e'_1 \not\models_{?}^{\dagger} \dot{\xi}_1$$
 for any $e'_1 \in \{e'_1 \mid e'_1 \text{ val and } e'_1 : \tau_1\}$ by Equation 27 on (18)

Then for any $e' \in \{e' \mid e' \text{ val and } e' : (\tau_1 \times \tau_2)\},\$

(20)
$$e'$$
 val by assumption

(21)
$$\cdot$$
; $\Delta \vdash e' : (\tau_1 \times \tau_2)$ by assumption

By rule induction over Rules (23) on (20).

Case (23a).

(22)
$$e' = n$$
 by assumption

By rule induction over Rules (19) on (21), no rule applies due to syntactic contradiction.

Case (23b).

(22)
$$e' = (\lambda x : \tau'.e'_1)$$
 by assumption

By rule induction over Rules (19) on (21), no rule applies due to syntactic contradiction.

Case (23c).

(22)
$$e' = (e'_1, e'_2)$$
 by assumption

(23)
$$e'_1$$
 val by assumption

By rule induction over Rules (19) on (21), only one rule applies.

Case (19g).

(24)
$$\cdot$$
; $\Delta \vdash e'_1 : \tau_1$ by assumption

(25)
$$e'_1 \not\models_?^{\dagger} \dot{\xi}_1$$
 by (19) on (23) and (24)

(26)
$$(e'_1, e'_2) \not\models^{\dagger}_{?} (\dot{\xi}_1, \dot{\xi}_2)$$
 by Lemma 2.0.13 on (25)

Case (23d).

(22)
$$e' = \operatorname{inl}_{\tau_2}(e'_1)$$
 by assumption

By rule induction over Rules (19) on (21), no rule applies due to syntactic contradiction.

Case (23e).

(22)
$$e' = \operatorname{inr}_{\tau_1}(e'_2)$$
 by assumption

By rule induction over Rules (19) on (21), no rule applies due to syntactic contradiction.

To conclude, $(e'_1, e'_2) \not\models^{\dagger}_? (\dot{\xi}_1, \dot{\xi}_2)$ for any $e' \in \mathtt{complete}(e)$.

Case (24d).

(9)
$$e = (e_1, e_2)$$
 by assumption

(10)
$$e_1$$
 indet by assumption

```
(11) e_2 val by assumption
(12) e_1 \not\models_?^\dagger \dot{\xi}_1 or e_2 \not\models_?^\dagger \dot{\xi}_2 by Lemma 2.0.13 on (4)
```

By case analysis on the disjunction in (12).

Case $e_1 \not\models_?^\dagger \dot{\xi}_1$.

(13)
$$e_1 \not\models_2^{\dagger} \dot{\xi}_1$$
 by assumption

By rule induction over Rules (19) on (2), only one rule applies. Case (19g).

(14)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$ by assumption

(15)
$$e_1' \not\models_?^{\dagger} \dot{\xi}_1$$
 for any $e_1' \in \text{complete}(e_1)$

by IH on (10) and (14) and (7) and (13)

$$(16) \ (e_1,e_2) \not\models_?^\dagger (\dot{\xi}_1,\dot{\xi}_2) \ \text{for any} \ e' \in \{(e'_1,e'_2) \mid e'_1 \in \texttt{complete}(e_1) \ \text{and} \ e'_2 \in \texttt{complete}(e_2)\}$$

by Lemma 2.0.13 on (15)

Case $e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$.

(13)
$$e_2' \not\models_7^{\dagger} \dot{\xi}_2$$
 by assumption

(14)
$$e_2' \not\models_?^\dagger \dot{\xi}_2$$
 for any $e_2' \in \mathtt{complete}(e_2)$

by Equation 27 and (13)

(15)
$$(e_1, e_2) \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 for any $e' \in \{(e'_1, e'_2) \mid e'_1 \in \text{complete}(e_1) \text{ and } e'_2 \in \text{complete}(e_2)\}$

by Lemma 2.0.13 on (14)

Case (24e).

$$\begin{array}{lll} (9) & e=(e_1,e_2) & \text{by assumption} \\ (10) & e_1 \text{ val} & \text{by assumption} \\ (11) & e_2 \text{ indet} & \text{by assumption} \\ (12) & e_1 \not\models_?^\dagger \dot{\xi}_1 \text{ or } e_2 \not\models_?^\dagger \dot{\xi}_2 & \text{by Lemma 2.0.13 on} \end{array}$$

By case analysis on the disjunction in (12).

Case $e_1 \not\models_{?}^{\dagger} \dot{\xi}_1$.

(13)
$$e_1' \not\models_{?}^{\dagger} \dot{\xi}_1$$
 by assumption

(14)
$$e_1' \not\models_?^{\dagger} \dot{\xi}_1$$
 for any $e_1' \in \mathtt{complete}(e_1)$

by Equation 27 and

(4)

(15)
$$(e_1, e_2) \not\models_{?}^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 for any $e' \in \{(e'_1, e'_2) \mid e'_1 \in \text{complete}(e_1) \text{ and } e'_2 \in \text{complete}(e_2)\}$ by Lemma 2.0.13 on (14)

Case $e_2 \not\models_?^\dagger \dot{\xi}_2$.

(13)
$$e_2 \not\models_{?}^{\dagger} \dot{\xi}_2$$

by assumption

By rule induction over Rules (19) on (2), only one rule applies. Case (19g).

$$(14) \cdot ; \Delta \vdash e_2 : \tau_2$$

by assumption

(15)
$$e_2' \not\models_7^{\dagger} \dot{\xi}_2 \text{ for any } e_2' \in \text{complete}(e_2)$$

by IH on (11) and (14)

and (8) and (13)

$$(16) \ (e_1,e_2) \not\models_?^\dagger (\dot{\xi}_1,\dot{\xi}_2) \ \text{for any} \ e' \in \{(e'_1,e'_2) \mid e'_1 \in \texttt{complete}(e_1) \ \text{and} \ e'_2 \in \texttt{complete}(e_2)\}$$

by Lemma 2.0.13 on (15)

Case (24f).

(9)
$$e = (e_1, e_2)$$

by assumption

$$(10)$$
 e_1 indet

by assumption

$$(11)$$
 e_2 indet

by assumption

(12)
$$e_1 \not\models_?^{\dagger} \dot{\xi}_1 \text{ or } e_2 \not\models_?^{\dagger} \dot{\xi}_2$$

by Lemma 2.0.13 on

(4)

By case analysis on the disjunction in (12).

Case $e_1 \not\models_?^\dagger \dot{\xi}_1$.

(13)
$$e_1 \not\models_?^{\dagger} \dot{\xi}_1$$

by assumption

By rule induction over Rules (19) on (2), only one rule applies. Case (19g).

(14)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$

by assumption

(15)
$$e_1' \not\models_{?}^{\dagger} \dot{\xi}_1 \text{ for any } e_1' \in \text{complete}(e_1)$$

by IH on (10) and (14)

and (7) and (13)

$$(16) \ (e_1,e_2) \not\models_?^\dagger (\dot{\xi}_1,\dot{\xi}_2) \ \text{for any} \ e' \in \{(e'_1,e'_2) \mid e'_1 \in \texttt{complete}(e_1) \ \text{and} \ e'_2 \in \texttt{complete}(e_2)\}$$

by Lemma 2.0.13 on (15)

Case $e_2 \not\models_?^\dagger \dot{\xi}_2$.

(13)
$$e_2 \not\models_{?}^{\dagger} \xi$$

by assumption

By rule induction over Rules (19) on (2), only one rule applies. Case (19g).

(14)
$$\cdot$$
; $\Delta \vdash e_2 : \tau_2$

by assumption

(15)
$$e_2' \not\models_7^{\dagger} \dot{\xi}_2$$
 for any $e_2' \in \mathsf{complete}(e_2)$

by IH on (11) and (14)

(16)
$$(e_1, e_2) \not\models_?^{\dagger} (\dot{\xi}_1, \dot{\xi}_2)$$
 for any $e' \in \{(e'_1, e'_2) \mid e'_1 \in \text{complete}(e_1) \text{ and } e'_2 \in \text{complete}(e_2)\}$ by Lemma 2.0.13 on (15)

Case (24i).

(9)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption

By rule induction over Rules (19) on (2), no rule applies due to syntactic contradiction.

Case (24j).

(9)
$$e = \operatorname{inr}_{\tau_1'}(e_2)$$
 by assumption

By rule induction over Rules (19) on (2), no rule applies due to syntactic contradiction.

Case (8f).

$(5) \ \dot{\xi} = \dot{\xi}_1 \vee \dot{\xi}_2$	by assumption
(6) $\dot{\xi}_1: \tau_1$	by assumption
(7) $\dot{\xi}_2 : \tau_2$	by assumption
$(8) \ e \not\models_{?}^{\dagger} \dot{\xi}_1 \lor \dot{\xi}_2$	by assumption
$(9) \ e \not\models_?^\dagger \dot{\xi}_1$	by Lemma 2.0.10 on (8)
$(10) \ e \not\models_?^\dagger \dot{\xi}_2$	by Lemma 2.0.10 on (8)
(11) $e' \not\models_?^{\dagger} \dot{\xi}_1$ for any $e' \in \mathtt{complete}(e)$	by IH on (1) and (2) and (6) and (9)
(12) $e' \not\models_?^\dagger \dot{\xi}_2$ for any $e' \in \mathtt{complete}(e)$	by IH on (1) and (2) and (7) and (10)
(13) $e' \not\models_?^{\dagger} \dot{\xi_1} \lor \dot{\xi_2} \text{ for any } e' \in \text{complete}(e)$	by Lemma $2.0.10$ on

 $\theta:\Gamma$ θ is of type Γ

STEmpty
$$\overline{\emptyset : \cdot}$$
 (28a)

(11) and (12)

STExtend
$$\frac{\theta: \Gamma_{\theta} \qquad \Gamma; \Delta \vdash e: \tau}{\theta, x/e: \Gamma_{\theta}, x: \tau}$$
(28b)

p refutable? p is refutable

$$(p)^{w} \text{ refutable?}$$

$$RInl$$

$$in1(p) \text{ refutable?}$$

$$RInr$$

$$inr(p) \text{ refutable?}$$

$$RPairL$$

$$p_{1} \text{ refutable?}$$

$$(p_{1}, p_{2}) \text{ refutable?}$$

$$(p_{1}, p_{2}) \text{ refutable?}$$

$$RPairR$$

$$p_{2} \text{ refutable?}$$

$$(p_{1}, p_{2}) \text{ refutable?}$$

$$(p_{2}, p_{2}) \text{ refutable?}$$

$$(p_{2},$$

RNum

REHole

RHole

 \underline{n} refutable?

 $\frac{}{(\!(\!)^w \text{ refutable}?}$

(29a)

(29b)

(29c)

$$e?p$$
 e may match p

$$\frac{\text{MMEHole}}{e? \, (\!\!\!)^w} \tag{31a}$$

$$\frac{e?(p)^w}{e?(p)^w}$$

MMNotIntro

$$\frac{e \; \mathtt{notintro} \qquad p \; \mathtt{refutable}_?}{e \; ? \; p} \tag{31c}$$

$$\frac{e_1?p_1 \qquad e_2 \rhd p_2 \dashv \theta_2}{(e_1, e_2)?(p_1, p_2)}$$
(31d)

$\operatorname{MMPairR}$

$$\frac{e_1 \rhd p_1 \dashv \theta_1 \qquad e_2 ? p_2}{(e_1, e_2) ? (p_1, p_2)}$$
(31e)

$$\frac{e_1? p_1 - e_2? p_2}{(e_1, e_2)? (p_1, p_2)}$$
(31f)

MMInl

$$\frac{e?p}{\operatorname{inl}_{\tau}(e)?\operatorname{inl}(p)} \tag{31g}$$

MMInr

$$\frac{e?p}{\operatorname{inr}_{\tau}(e)?\operatorname{inr}(p)} \tag{31h}$$

$e \perp p$ e does not match p

$$\frac{n_1 \neq n_2}{n_1 \perp n_2} \tag{32a}$$

${\rm NMPairL}$

$$\frac{e_1 \perp p_1}{(e_1, e_2) \perp (p_1, p_2)} \tag{32b}$$

${\bf NMPairR}$

$$\frac{e_2 \perp p_2}{(e_1, e_2) \perp (p_1, p_2)} \tag{32c}$$

${\rm NMConfL}$

$$\frac{-}{\operatorname{inr}_{\tau}(e) \perp \operatorname{inl}(p)} \tag{32d}$$

NMConfR

$$\frac{}{\operatorname{inl}_{\tau}(e) \perp \operatorname{inr}(p)} \tag{32e}$$

$$\frac{\text{NMInl}}{e \perp p} \\ \frac{e \perp p}{\text{inr}_{\tau}(e) \perp \text{inl}(p)}$$
 (32f)

NMInr

$$\frac{e \perp p}{\operatorname{inl}_{\tau}(e) \perp \operatorname{inr}(p)} \tag{32g}$$

 $e \mapsto e'$ e takes a step to e'

$$\frac{e \mapsto e'}{ (\!\! | e \!\!)^u \mapsto (\!\! | e' \!\!)^u }$$
 (33a)

ITApFun

$$\frac{e_1 \mapsto e_1'}{e_1(e_2) \mapsto e_1'(e_2)} \tag{33b}$$

$$\begin{array}{ll} \text{ITApArg} \\ \underline{e_1 \text{ val}} & \underline{e_2 \mapsto e_2'} \\ \underline{e_1(e_2) \mapsto e_1(e_2')} \end{array} \tag{33c}$$

ITAP

$$\frac{e_2 \text{ val}}{(\lambda x : \tau . e_1)(e_2) \mapsto [e_2/x]e_1} \tag{33d}$$

ITPairL
$$\frac{e_1 \mapsto e_1'}{(e_1, e_2) \mapsto (e_1', e_2)}$$
(33e)

$$\begin{array}{ll} \text{ITPairR} \\ \underline{e_1 \text{ val}} & \underline{e_2 \mapsto e_2'} \\ \underline{(e_1, e_2) \mapsto (e_1, e_2')} \end{array}$$
 (33f)

$$\frac{(e_1, e_2) \text{ final}}{\text{prl}((e_1, e_2)) \mapsto e_1}$$
(33g)

$$\frac{(e_1, e_2) \text{ final}}{\text{prr}((e_1, e_2)) \mapsto e_2}$$
(33h)

$$\frac{e \mapsto e'}{\operatorname{inl}_{\tau}(e) \mapsto \operatorname{inl}_{\tau}(e')} \tag{33i}$$

ITInr

$$\frac{e \mapsto e'}{\operatorname{inr}_{\tau}(e) \mapsto \operatorname{inr}_{\tau}(e')}$$
(33j)

ITExpMatch

$$\frac{e \mapsto e'}{\mathtt{match}(e)\{\hat{rs}\} \mapsto \mathtt{match}(e')\{\hat{rs}\}} \tag{33k}$$

$$ITSuccMatch\\$$

$$\frac{e \; \text{final} \qquad e \rhd p_r \dashv \theta}{\text{match}(e) \{ rs_{pre} \mid (p_r \Rightarrow e_r) \mid rs_{post} \} \mapsto [\theta](e_r)} \tag{33l}$$

ITFailMatch

$$\frac{e \; \mathtt{final} \qquad e \perp p_r}{\mathtt{match}(e) \{ rs \mid (p_r \Rightarrow e_r) \mid (r' \mid rs') \} \mapsto \mathtt{match}(e) \{ (rs \mid (p_r \Rightarrow e_r) \mid \cdot)^{\diamond} \mid r' \mid rs' \}} \tag{33m}$$

Lemma 4.0.2. If $\operatorname{inl}_{\tau_2}(e_1)$ final then e_1 final.

Proof. By rule induction over Rules (25) on $\operatorname{inl}_{\tau_2}(e_1)$ final.

Case (25a).

(14)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 val

by assumption

By rule induction over Rules (23) on (14), only one case applies.

Case (23d).

$$(15)$$
 e_1 val

by assumption

$$(16)$$
 e_1 final

by Rule (25a) on (15)

Case (25b).

$$(14)$$
 $\operatorname{inl}_{ au_2}(e_1)$ indet

by assumption

By rule induction over Rules (24) on (14), only one case applies.

Case (24i).

$$(15)$$
 e_1 indet

by assumption

$$(16)$$
 e_1 final

by Rule (25b) on (15)

Lemma 4.0.3. If $\operatorname{inr}_{\tau_1}(e_2)$ final then e_2 final.

Proof. By rule induction over Rules (25) on $\operatorname{inr}_{\tau_1}(e_2)$ final.

Case (25a).

$$(1)$$
 $\operatorname{inr}_{ au_1}(e_2)$ val

by assumption

By rule induction over Rules (23) on (1), only one case applies.

Case (23d).

(2) e_2 val

by assumption

(3) e_2 final

by Rule (25a) on (2)

Case (25b).

(1) $\operatorname{inr}_{\tau_1}(e_2)$ indet

by assumption

By rule induction over Rules (24) on (1), only one case applies.

Case (24i).

(2) e_2 indet

by assumption

(3) e_2 final

by Rule (25b) on (2)

Lemma 4.0.4. If (e_1, e_2) final $then e_1$ final $and e_2$ final.

Proof. By rule induction over Rules (25) on (e_1, e_2) final.

Case (25a).

(1) (e_1, e_2) val

by assumption

By rule induction over Rules (23) on (1), only one case applies.

Case (23c).

(2) e_1 val

by assumption

(3) e_2 val

by assumption

(4) e_1 final

by Rule (25a) on (2)

(5) e_2 final

by Rule (25a) on (3)

Case (25b).

(1) (e_1, e_2) indet

by assumption

By rule induction over Rules (24) on (1), only three cases apply.

Case (24d).

(2) e_1 indet

by assumption

(3) e_2 val

by assumption

(4) e_1 final

by Rule (25b) on (2)

(5) e_1 final

by Rule (25a) on (3)

Case (24e).

(2) e_1 val

by assumption

(3) e_2 indet

by assumption

(4) e_1 final

by Rule (25a) on (2)

(5) e_1 final

by Rule (25b) on (3)

Case (24f).

(2) e_1 indet

by assumption

 $\begin{array}{lll} (3) & e_2 & \text{indet} & & \text{by assumption} \\ (4) & e_1 & \text{final} & & \text{by Rule (25b) on (2)} \\ (5) & e_1 & \text{final} & & \text{by Rule (25b) on (3)} \\ \end{array}$

Lemma 4.0.5. There doesn't exist \underline{n} such that \underline{n} notintro.

Proof. By rule induction over Rules (26) on \underline{n} notintro, no case applies due to syntactic contradiction.

Lemma 4.0.6. There doesn't exist $\operatorname{inl}_{\tau}(e)$ such that $\operatorname{inl}_{\tau}(e)$ notintro.

Proof. By rule induction over Rules (26) on $\operatorname{inl}_{\tau}(e)$ notintro, no case applies due to syntactic contradiction.

Lemma 4.0.7. There doesn't exist $\operatorname{inr}_{\tau}(e)$ such that $\operatorname{inr}_{\tau}(e)$ notintro.

Proof. By rule induction over Rules (26) on $\operatorname{inr}_{\tau}(e)$ notintro, no case applies due to syntactic contradiction.

Lemma 4.0.8. There doesn't exist (e_1, e_2) such that (e_1, e_2) notintro.

Proof. By rule induction over Rules (26) on (e_1, e_2) notintro, no case applies due to syntactic contradiction.

Lemma 4.0.9. If e final and e notintro then e indet.

Proof Sketch. By rule induction over Rules (26) on e notintro, for each case, by rule induction over Rules (23) on e val and we notice that e val is not derivable. By rule induction over Rules (25) on e final, Rule (25a) result in a contradiction with the fact that e val is not derivable while Rule (25b) tells us e indet.

Lemma 4.0.10 (Finality). There doesn't exist such an expression e such that both e final and $e \mapsto e'$ for some e'

Proof. Assume there exists such an e such that both e final and $e \mapsto e'$ for some e' then proof by contradiction.

By rule induction over Rules (25) and Rules (33), *i.e.*, over Rules (23) and Rules (33) and over Rules (24) and Rules (33) respectively. The proof can be done by straightforward observation of syntactic contradictions. \Box

Lemma 4.0.11 (Matching Determinism). *If* e **final** $and \cdot ; \Delta_e \vdash e : \tau$ and $p : \tau[\xi] \dashv \Gamma ; \Delta$ then exactly one of the following holds

- 1. $e > p \dashv \theta$ for some θ
- 2. e ? p
- 3. $e \perp p$

Proof.

(1) e final by assumption

(2) \cdot ; $\Delta_e \vdash e : \tau$ by assumption

(3) $p:\tau[\xi]\dashv \Gamma;\Delta$ by assumption

By rule induction over Rules (20) on (3), we would show one conclusion is derivable while the other two are not.

Case (20a).

(4) p = x by assumption

(5) $e \triangleright x \dashv e/x$ by Rule (30a)

Assume e? x. By rule induction over Rules (31) on it, only one case applies.

Case (31c).

(6) x refutable? by assumption

By rule induction over Rules (29) on (6), no case applies due to syntactic contradiction.

(7) $e^{2}\overline{x}$ by contradiction

Assume $e \perp x$. By rule induction over Rules (32) on it, no case applies due to syntactic contradiction.

(8) $e \pm x$ by contradiction

Case (20b).

(4) p =_ by assumption

(5) $e > \exists \exists$ by Rule (30b)

Assume e? _. By rule induction over Rules (31) on it, only one case applies.

Case (31c).

(6) refutable? by assumption

By rule induction over Rules (29) on (6), no case applies due to syntactic contradiction.

(7) e^{2} by contradiction

Assume $e \perp$ _. By rule induction over Rules (32) on it, no case applies due to syntactic contradiction.

Case (20c).

(4)
$$p = \emptyset^w$$
 by assumption

(5)
$$e ? \emptyset^w$$
 by Rule (31a)

Assume $e \rhd \oplus^w \dashv \theta$ for some θ . By rule induction over Rules (31) on it, no case applies due to syntactic contradiction.

(6)
$$e \rightarrow \theta \theta$$
 by contradiction

Assume $e \perp \emptyset^w$. By rule induction over Rules (32) on it, no case applies due to syntactic contradiction.

Case (20d).

(4)
$$p = (p_0)^w$$
 by assumption

(5)
$$e ? (p_0)^w$$
 by Rule (31b)

Assume $e \rhd (p_0)^w \dashv \theta$ for some θ . By rule induction over Rules (31) on it, no case applies due to syntactic contradiction.

(6)
$$e \triangleright p_{\theta} = \theta$$
 by contradiction

Assume $e \perp (p_0)^w$. By rule induction over Rules (32) on it, no case applies due to syntactic contradiction.

Case (20e).

(4)
$$p = n_2$$
 by assumption

(5)
$$\tau = \text{num}$$
 by assumption

(6)
$$\xi = n_2$$
 by assumption

(7)
$$n_2$$
 refutable? by Rule (29a)

By rule induction over Rules (19) on (2), the following cases apply.

(8)
$$e = \{ \| u, \| e_0 \|^u, e_1(e_2), \text{prl}(e_0), \text{prr}(e_0), \text{match}(e_0) \{ \hat{rs} \} \}$$
 by assumption

(9)
$$e$$
 notintro by Rule

(10)
$$e ? \underline{n_2}$$
 by Rule (16b) on (7) and (9)

Assume $e > \underline{n_2} \dashv \theta$ for some θ . By rule induction over it, no case applies due to syntactic contradiction.

(11)
$$e \triangleright n_2 + \theta$$
 by contradiction

Assume $e \perp \underline{n_2}$. By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \perp n_2$$
 by contradiction

Case (19d).

(8)
$$e = n_1$$

Assume $\underline{n_1}$? $\underline{n_2}$. By rule induction over Rules (31) on it, only two cases apply.

Case (31c).

(9)
$$\underline{n_1}$$
 notintro by assumption Contradicts Lemma 4.0.5.

(10)
$$n_1 ? n_2$$
 by contradiction

By case analysis on whether $n_1 = n_2$.

Case $n_1 = n_2$.

(11)
$$n_1 = n_2$$
 by assumption
(12) $n_1 \triangleright n_2 \dashv \cdot$ by Rule (30c)

Assume $\underline{n_1} \perp \underline{n_2}$. By rule induction over Rules (32) on it, only one case applies.

Case (32a).

(13)
$$n_1 \neq n_2$$
 by assumption Contradicts (11).

(14)
$$n_1 + n_2$$
 by contradiction

Case $n_1 \neq n_2$.

(11)
$$n_1 \neq n_2$$
 by assumption
(12) $\underline{n_1} \perp \underline{n_2}$ by Rule (32a) on (11)

Assume
$$\underline{n_1} \rhd \underline{n_2} \dashv \theta$$
 for some θ . By rule induction over Rules (30) on it, no case applies due to syntactic contradiction.

(13)
$$n_1 \triangleright n_2 \dashv \theta$$
 by contradiction

Case (20f).

$(4) \ p = \mathtt{inl}(p_1)$	by assumption
(5) $\tau = (\tau_1 + \tau_2)$	by assumption
$(6) \ \xi = \mathtt{inl}(\xi_1)$	by assumption
$(7) \ p_1: \tau_1[\xi_1] \dashv \Gamma; \Delta$	by assumption
(8) $inl(p_1)$ refutable?	by Rule (29d)

By rule induction over Rules (19) on (2), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(9)
$$e = \langle || ||^u, \langle || e_0 \rangle||^u, e_1(e_2), \text{prl}(e_0), \text{prr}(e_0), \text{match}(e_0) \{ \hat{rs} \}$$
 by assumption

(10) e notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

(11) e? $inl(p_1)$ by Rule (16b) on (8) and (10)

Assume $e \triangleright \text{inl}(p_1) \dashv \theta_1$ for some θ_1 . By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \triangleright \operatorname{inl}(p_1) \dashv \theta_1$$
 by contradiction

Assume $e \perp \text{inl}(p_1)$. By rule induction over it, no case applies due to syntactic contradiction.

(13)
$$e \perp int(p_1)$$
 by contradiction

Case (19j).

(9)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption
(10) \cdot ; $\Delta_e \vdash e_1 : \tau_1$ by assumption
(11) e_1 final by Lemma 4.0.2 on (1)

By inductive hypothesis on (10) and (11) and (7), exactly one of $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 , $e_1 ? p_1$, and $e_1 \perp p_1$ holds. By case analysis on which one holds.

Case $e_1 > p_1 \dashv \theta_1$.

$$\begin{array}{lll} (12) & e_1 \rhd p_1 \dashv \theta_1 & \text{by assumption} \\ (13) & e_1 \nearrow p_1 & \text{by assumption} \\ (14) & e_1 \not \perp p_1 & \text{by assumption} \\ (15) & & & \text{inl}_{\tau_2}(e_1) \rhd & & \text{inl}(p_1) \dashv \theta_1 & \text{by Rule (30e) on (12)} \\ \end{array}$$

Assume $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$. By rule induction over Rules (31) on it, only two cases apply.

Case (31c).

(16)
$$\operatorname{inl}_{\tau_2}(e_1)$$
 notintro by assumption Contradicts Lemma 4.0.6.

Case (31g).

(16)
$$e_1$$
? p_1 by assumption Contradicts (13).

(17)
$$\underline{\operatorname{inl}_{\tau_2}(e_1) ? \operatorname{inl}(p_1)}$$
 by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1)$. By rule induction over Rules (32) on it, only one case applies.

Case (32f).

(18)
$$e_1 \perp p_1$$
 by assumption Contradicts (14).

(19)
$$\operatorname{inl}_{\tau_2}(e_1) \pm \operatorname{inl}(p_1)$$
 by contradiction

Case $e_1 ? p_1$.

$$\begin{array}{lll} (12) & \underline{e_1} \triangleright p_1 + | \theta_1 \\ (13) & e_1 ? p_1 \\ (14) & \underline{e_1} + | p_1 \\ (15) & \operatorname{inl}_{\tau_2}(e_1) ? \operatorname{inl}(p_1) \end{array} \qquad \begin{array}{ll} \text{by assumption} \\ \text{by assumption} \\ \text{by Rule (31g) on (13)} \end{array}$$

Assume $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (30) on it, only one case applies.

Case (30e).

(16)
$$e_1 \triangleright p_1 \dashv \theta$$
 by assumption Contradicts (12).

(17)
$$\operatorname{inl}_{\tau_2}(e_1)
ightharpoonup \operatorname{inl}(p_1) \dashv \theta$$
 by contradiction Assume $\operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1)$. By rule induction over Rules (32) on it, only one case applies.

Case (32f).

(18)
$$e_1 \perp p_1$$
 by assumption Contradicts (14).

(19)
$$\operatorname{inl}_{\tau_2}(e_1) + \operatorname{inl}(p_1)$$
 by contradiction

Case $e_1 \perp p_1$.

$$\begin{array}{lll} (12) & \underline{e_1} \triangleright p_1 \# f_1 & \text{by assumption} \\ (13) & \underline{e_1} \cdot ? p_1 & \text{by assumption} \\ (14) & \underline{e_1} \perp p_1 & \text{by assumption} \\ (15) & \operatorname{inl}_{\tau_2}(e_1) \perp \operatorname{inl}(p_1) & \text{by Rule (32f) on (14)} \end{array}$$

Assume $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (30) on it, only one case applies.

Case (30e).

(16)
$$e_1 \triangleright p_1 \dashv \theta$$
 by assumption Contradicts (12).

(17)
$$\operatorname{inl}_{\tau_2}(e_1) \triangleright \operatorname{inl}(p_1) \dashv \theta$$
 by contradiction

Assume $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$. By rule induction over Rules (31) on it, only two cases apply.

Case (31c).

(18) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption Contradicts Lemma 4.0.6.

Case (31g).

(18) $e_1 ? p_1$ by assumption Contradicts (13).

(19) $\operatorname{inl}_{\tau_2}(e_1) ? \operatorname{inI}(p_1)$ by contradiction

Case (20g).

(4)
$$p = inr(p_2)$$
 by assumption
(5) $\tau = (\tau_1 + \tau_2)$ by assumption
(6) $\xi = inr(\xi_2)$ by assumption
(7) $p_2 : \tau_2[\xi_2] \dashv \Gamma; \Delta$ by assumption
(8) $inr(p_2)$ refutable? by Rule (29e)

By rule induction over Rules (19) on (2), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(9)
$$e = (||u|, ||e_0||^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$$

by assumption

(10)
$$e$$
 notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

(11)
$$e$$
? $inr(p_2)$ by Rule (16b) on (8)

and (10)

Assume $e
ightharpoonup \operatorname{inr}(p_2) \dashv \theta_2$ for some θ_2 . By rule induction over it, no case applies due to syntactic contradiction.

(12)
$$e \triangleright inr(p_2) \dashv \theta_2$$
 by contradiction

Assume $e \perp inr(p_2)$. By rule induction over it, no case applies due to syntactic contradiction.

(13)
$$e \perp inr(p_2)$$
 by contradiction

Case (19k).

(9)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption
(10) \cdot ; $\Delta_e \vdash e_2 : \tau_2$ by assumption

(11)
$$e_2$$
 final by Lemma 4.0.3 on (1)

By inductive hypothesis on (10) and (11) and (7), exactly one of $e_2 \triangleright p_2 \dashv \theta_2$ for some θ_2 , $e_2 ? p_2$, and $e_2 \perp p_2$ holds. By case analysis on which one holds.

Case $e_2 > p_2 \dashv \theta_2$.

 $(12) e_2 \rhd p_2 \dashv \theta_2$

by assumption

(13) $e_2 ? p_2$

by assumption

 $(14) \ e_2 \perp p_2$

by assumption

(15) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta_2$

by Rule (30f) on (12)

Assume $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$. By rule induction over Rules (31) on it, only two cases apply.

Case (31c).

(16) $\operatorname{inr}_{ au_1}(e_2)$ notintro

by assumption

Contradicts Lemma 4.0.7.

Case (31h).

 $(16) e_2 ? p_2$

by assumption

Contradicts (13).

(17) $\inf_{\tau_1}(e_2)$? $\inf(p_2)$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \perp \operatorname{inr}(p_2)$. By rule induction over Rules (32) on it, only one case applies.

Case (32g).

(18) $e_2 \perp p_2$

by assumption

Contradicts (14).

(19) $\operatorname{inr}_{\tau_1}(e_2) \pm \operatorname{inr}(p_2)$

by contradiction

Case e_2 ? p_2 .

 $(12) \ \underline{e_2 \triangleright p_2 \# \theta}$

by assumption

(13) $e_2 ? p_2$

by assumption

 $(14) \ \underline{e_2 \perp p_2}$

by assumption

(15) $inr_{\tau_1}(e_2)$? $inr(p_2)$

by Rule (31h) on (13)

Assume $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ . By rule induction over Rules (30) on it, only one case applies.

Case (30f).

 $(16) e_2 \rhd p_2 \dashv\!\!\dashv \theta$

by assumption

Contradicts (12).

(17) $\operatorname{inr}_{\tau_1}(e_2) \triangleright \operatorname{inr}(p_2) \dashv \theta$

by contradiction

Assume $\operatorname{inr}_{\tau_1}(e_2) \perp \operatorname{inr}(p_2)$. By rule induction over Rules (32) on it, only one case applies.

Case (32g).

Contradicts (14).

(19)
$$\inf_{T_1}(e_2) + \inf(p_2)$$
 by contradiction

Case $e_2 \perp p_2$.

(12) $\underbrace{e_2 \triangleright p_2 + \theta}$ by assumption

(13) $\underbrace{e_2 \cdot p_2}$ by assumption

(14) $e_2 \perp p_2$ by assumption

(15) $\inf_{T_1}(e_2) \perp \inf(p_2)$ by Rule (32g) on (14)

Assume $\inf_{T_1}(e_2) \triangleright \inf(p_2) + \theta$ for some θ . By rule induction over Rules (30) on it, only one case applies.

Case (30f).

(16) $e_2 \triangleright p_2 + \theta$ by assumption

Contradicts (12).

(17) $\inf_{T_1}(e_2) \triangleright \inf(p_2) + \theta$ by contradiction

Assume $\inf_{T_1}(e_2) \triangleright \inf(p_2) + \theta$ by contradiction

Assume $\inf_{T_1}(e_2) \triangleright \inf(p_2) + \theta$ by contradiction

Contradicts (12).

(18) $\inf_{T_1}(e_2) \triangleright \inf(p_2) + \theta$ by assumption

Contradicts Lemma 4.0.7.

Case (31c).

(18) $e_2 \cdot p_2$ by assumption

Contradicts (13).

(19) $\inf_{T_1}(e_2) \cdot \inf(p_2)$ by contradiction

Case (20h).

(4) $p = (p_1, p_2)$ by assumption

Contradicts (13).

(5) $\tau = (\tau_1 \times \tau_2)$ by assumption

(6) $\xi = (\xi_1, \xi_2)$ by assumption

(7) $\Gamma = \Gamma_1 \uplus \Gamma_2$ by assumption

(8) $\Delta = \Delta_1 \uplus \Delta_2$ by assumption

by assumption

By rule induction over Rules (19) on (2), the following cases apply.

by assumption

by assumption

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(9) $p_1 : \tau_1[\xi_1] \dashv \Gamma_1 ; \Delta_1$

(10) $p_2 : \tau_2[\xi_2] \dashv \Gamma_2 ; \Delta_2$

```
(11) e = \{ \|u, \|e_0\|^u, e_1(e_2), \operatorname{prl}(e_0), \operatorname{prr}(e_0), \operatorname{match}(e_0) \{ \hat{rs} \} \}
                                                     by assumption
(12) e notintro
                                                     by Rule
                                                     (26a),(26b),(26c),(26d),(26e),(26f)
(13) e indet
                                                     by Lemma 4.0.9 on (1)
                                                     and (12)
(14) prl(e) indet
                                                     by Rule (24g) on (13)
(15) prl(e) final
                                                     by Rule (25b) on (14)
(16) prr(e) indet
                                                     by Rule (24h) on (13)
(17) prr(e) final
                                                     by Rule (25b) on (16)
(18) \cdot; \Delta \vdash \mathtt{prl}(e) : \tau_1
                                                     by Rule (19h) on (2)
(19) \cdot; \Delta \vdash \mathsf{prr}(e) : \tau_2
                                                    by Rule (19i) on (2)
```

Assume $e \perp (p_1, p_2)$. By rule induction on it, no case applies due to syntactic contradiction.

(20)
$$e \perp (p_1, p_2)$$
 by contradiction

By inductive hypothesis on (15) and (18) and (9), exactly one of $\operatorname{prl}(e) \rhd p_1 \dashv \theta_1$, $\operatorname{prl}(e) ? p_1$, and $\operatorname{prl}(e) \perp p_1$ holds. By inductive hypothesis on (17) and (19) and (10), exactly one of $\operatorname{prr}(e) \rhd p_2 \dashv \theta_2$, $\operatorname{prr}(e) ? p_2$, and $\operatorname{prr}(e) \perp p_2$ holds. By case analysis on which conclusion holds for p_1 and p_2 . Note that we have already shown $e \perp (p_1, p_2)$.

Case $\operatorname{prl}(e) \rhd p_1 \dashv \theta_1, \operatorname{prr}(e) \rhd p_2 \dashv \theta_2.$

$(21) \ \operatorname{prl}(e) \rhd p_1 \dashv \theta_1$	by assumption
(22) $prl(e) ? p_1$	by assumption
(23) $\underline{\operatorname{prl}(e) \perp p_1}$	by assumption
$(24) \ \operatorname{prr}(e) \rhd p_2 \dashv \theta_2$	by assumption
$(25) \ \underline{\operatorname{prr}(e)?p_2}$	by assumption
(26) $\operatorname{prr}(e) \perp p_2$	by assumption
$(27) \ e \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2$	by Rule $(30g)$ on (12)
	and (21) and (24)

Assume $e?(p_1, p_2)$. By rule induction over Rules (31) on it, only one case applies.

Case (31c).

(28) (p_1, p_2) refutable? by assumption By rule induction over Rules (29), only two cases apply. Case (29f).

(29) p_1 refutable?by assumption(30)prl(e) notintroby Rule (26e)(31)prl(e)? p_1 by Rule (31c) on (29)and (30)

Contradicts (22).

Case (29g).

- (29) p_2 refutable? by assumption (30) prr(e) notintro by Rule (26f)
- (31) prl(e) ? p_1 by Rule (31c) on (29) and (30)

Contradicts (22).

(32) $e_{1}(p_{1}, p_{2})$

by contradiction

Case $prl(e) > p_1 \dashv \theta_1, prr(e) ? p_2.$

(21) $\operatorname{prl}(e) \triangleright p_1 \dashv \theta_1$ by assumption (22) $\operatorname{prl}(e) ? p_1$ by assumption (23) $\operatorname{prl}(e) \not p_1$ by assumption (24) $\operatorname{prr}(e) \triangleright p_2 \dashv \theta_2$ by assumption (25) $\operatorname{prr}(e) ? p_2$ by assumption (26) $\operatorname{prr}(e) \not p_2$ by assumption

Assume $e \rhd (p_1, p_2) \dashv \theta$. By rule induction over Rules (30), only one case applies.

Case (30g).

(27) $\theta = \theta_1 \uplus \theta_2$ by assumption (28) $prr(e) \rhd p_2 \dashv \theta_2$ by assumption

Contradicts (24).

(29) $e \triangleright (p_1, p_2) \dashv \theta$

by contradiction

By rule induction over Rules (31) on (25), the following cases apply.

Case (31a),(31b).

- (30) $p_2 = \langle | \rangle^w, \langle | p \rangle^w$ by assumption
- (31) p_2 refutable? by Rule (29b) and Rule (29c)
- $(32) \ (p_1,p_2) \ \mathtt{refutable}_? \qquad \qquad \mathrm{by} \ \mathrm{Rule} \ (29\mathrm{g}) \ \mathrm{on} \ (31)$
- (33) e? (p_1, p_2) by Rule (31c) on (12) and (32)

Case (31c).

- (30) p_2 refutable? by assumption
- (31) (p_1, p_2) refutable? by Rule (29g) on (30)
- (32) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (31)

Case $prl(e) \triangleright p_1 \dashv \theta_1, prr(e) \perp p_2$.

(21) $\operatorname{prl}(e) \triangleright p_1 \dashv \theta_1$ by assumption (22) $\operatorname{prl}(e) ? p_1$ by assumption

(23) $prl(e) \perp p_1$ by assumption (24) $prr(e) \triangleright p_2 \dashv \theta_2$ by assumption (25) $prr(e) ? p_2$ by assumption (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (32) on (26), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case prl(e) ? $p_1, prr(e) \triangleright p_2 \dashv \theta_2$.

(21) $\underline{\operatorname{prl}(e)} \triangleright p_1 \dashv \overline{\theta_1}$	by assumption
$(22) \ \operatorname{prl}(e) ? p_1$	by assumption
(23) $\operatorname{prl}(e) \pm p_1$	by assumption
$(24) \ prr(e) \rhd p_2 \dashv \theta_2$	by assumption
(25) $prr(e)$? p_2	by assumption
(26) $prr(e) \perp p_2$	by assumption

Assume $e \rhd (p_1, p_2) \dashv \theta$. By rule induction over Rules (30), only one case applies.

Case (30g).

$$\begin{array}{ll} (27) \ \theta = \theta_1 \uplus \theta_2 & \text{by assumption} \\ (28) \ \operatorname{prl}(e) \rhd p_1 \dashv\! \mid \! \theta_1 & \text{by assumption} \\ \text{Contradicts (21)}. & \end{array}$$

(29)
$$e \triangleright (p_1, p_2) \dashv \theta$$
 by contradiction

By rule induction over Rules (31) on (22), the following cases apply.

Case (31a),(31b).

$(30) p_1 = ()^w, (p)^w$	by assumption
(31) p_1 refutable?	by Rule (29b) and Rule
	(29c)
(32) (p_1,p_2) refutable?	by Rule $(29g)$ on (31)
$(33) e? (p_1, p_2)$	by Rule (31c) on (12)
	and (32)

Case (31c).

(30) p_1 refutable?	by assumption
(31) (p_1,p_2) refutable?	by Rule $(29g)$ on (30)
$(32) e? (p_1, p_2)$	by Rule (31c) on (12)
	and (31)

Case prl(e) ? p_1 , prr(e) ? p_2 .

(21) $\underline{\operatorname{prl}(e)} \Rightarrow p_1 \dashv \overline{\theta_1}$	by assumption
$(22) \ \mathtt{prl}(e) \mathbin{?} p_1$	by assumption
(23) $prl(e) \pm p_1$	by assumption

(24) $\underline{\operatorname{prr}(e)} \triangleright p_2 \dashv \overline{\theta_2}$ by assumption

(25) prr(e)? p_2 by assumption

(26) $prr(e) \perp p_2$ by assumption

Assume $e \rhd (p_1, p_2) \dashv \theta$. By rule induction over Rules (30), only one case applies.

Case (30g).

(27) $\theta = \theta_1 \uplus \theta_2$ by assumption

(28) $\operatorname{prl}(e) \rhd p_1 \dashv \theta_1$ by assumption

Contradicts (21).

(29) $e \triangleright (p_1, p_2) \dashv \theta$ by contradiction

By rule induction over Rules (31) on (22), the following cases apply.

Case (31a),(31b).

(30) $p_1 = \langle \rangle^w, \langle p \rangle^w$ by assumption

(31) p_1 refutable? by Rule (29b) and Rule (29c)

(32) (p_1, p_2) refutable? by Rule (29g) on (31)

(33) e? (p_1, p_2) by Rule (31c) on (12) and (32)

Case (31c).

(30) p_1 refutable? by assumption

(31) (p_1, p_2) refutable? by Rule (29g) on (30)

(32) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (31)

Case prl(e)? $p_1, prr(e) \perp p_2$.

(21) $prl(e) \rightarrow p_1 \dashv \theta_1$ by assumption

(22) prl(e)? p_1 by assumption

(23) $prl(e) \perp p_1$ by assumption

(24) $prr(e) \rightarrow p_2 \dashv \overline{\theta_2}$ by assumption

(25) prr(e)? p_2 by assumption

(26) $\operatorname{prr}(e) \perp p_2$ by assumption

By rule induction over Rules (32) on (26), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) \triangleright p_2 \dashv \theta_2$.

(21) $\underline{\text{prl}(e)} \Rightarrow p_1 \dashv \overline{\theta_1}$ by assumption (22) $\underline{\text{prl}(e)}, p_1$ by assumption

(23) $prl(e) \perp p_1$ by assumption

(24) $prr(e) > p_2 \dashv \theta_2$ by assumption

(25) prr(e)? p_2 by assumption (26) $prr(e) \perp p_2$ by assumption

By rule induction over Rules (32) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) ? p_2$.

(21) $\underline{\operatorname{prl}(e)} \rightarrow p_1 \dashv \theta_1$	by assumption
$(22) \ \underline{\mathtt{prl}(e)?p_1}$	by assumption
(23) $\operatorname{prl}(e) \perp p_1$	by assumption
(24) $\underline{\operatorname{prr}(e)} \Rightarrow p_2 \dashv \overline{\theta_2}$	by assumption
$(25) \ prr(e) ? p_2$	by assumption
(26) $\operatorname{prr}(e) \perp p_2$	by assumption

By rule induction over Rules (32) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case $prl(e) \perp p_1, prr(e) \perp p_2$.

(21) $\underline{\operatorname{prl}(e)} \Rightarrow p_1 \dashv \overline{\theta_1}$	by assumption
(22) $prl(e) ? p_1$	by assumption
(23) $\operatorname{prl}(e) \perp p_1$	by assumption
(24) $\underline{\operatorname{prr}(e) \triangleright p_2 \dashv \theta_2}$	by assumption
$(25) \ prr(e) ? p_2$	by assumption
(26) $prr(e) \pm p_2$	by assumption

By rule induction over Rules (32) on (23), no case applies due to syntactic contradiction.

Therefore, vacuously true.

Case (19g).

$(11) \ e = (e_1, e_2)$	by assumption
$(12) \cdot ; \Delta \vdash e_1 : \tau_1$	by assumption
$(13) \cdot ; \Delta \vdash e_2 : \tau_2$	by assumption
(14) e_1 final	by Lemma $4.0.4$ on (1)
(15) e_2 final	by Lemma 4.0.4 on (1)

By inductive hypothesis on (9) and (12) and (14), exactly one of $e_1 \triangleright p_1 \dashv \theta_1$, $e_1 ? p_1$, and $e_1 \perp p_1$ holds.

By inductive hypothesis on (10) and (13) and (15), exactly one of $e_2 \rhd p_2 \dashv \theta_2, \ e_2 ? p_2$, and $e_2 \perp p_2$ holds.

By case analysis on which conclusion holds for p_1 and p_2 .

Case $e_1 \rhd p_1 \dashv \mid \theta_1, e_2 \rhd p_2 \dashv \mid \theta_2$.

$(16) e_1 \rhd p_1 \dashv \theta_1$	by assumption
$(17) e_1 ? p_1$	by assumption
$(18) e_1 + p_1$	by assumption

 $\begin{array}{lll} (19) & e_{2} \rhd p_{2} \dashv \theta_{2} & \text{by assumption} \\ (20) & \underline{e_{2} ? p_{2}} & \text{by assumption} \\ (21) & \underline{e_{2} \bot p_{2}} & \text{by assumption} \\ (22) & (e_{1}, e_{2}) \rhd (p_{1}, p_{2}) \dashv \theta_{1} \uplus \theta_{2} & \text{by Rule (30d) on (16)} \\ & & \text{and (19)} \end{array}$

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (31) on it, only four cases apply.

Case (31c).

(23) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (31d).

(23) $e_1 ? p_1$ by assumption

Contradicts (17). Case (31e).

(23) e_2 ? p_2 by assumption

Contradicts (20).

Case (31f).

(23) $e_1 ? p_1$ by assumption

Contradicts (17).

(24) (e_1, e_2) ? (p_1, p_2) by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (32) on it, only two cases apply.

Case (32b).

(25) $e_1 \perp p_1$ by assumption

Contradicts (18).

Case (32c).

(25) $e_2 \perp p_2$ by assumption

Contradicts (21).

(26) $(e_1, e_2) \pm (p_1, p_2)$ by contradiction

Case $e_1 > p_1 \dashv \theta_1, e_2 ? p_2$.

 $\begin{array}{lll} (16) & e_1 \rhd p_1 \dashv \theta_1 & \text{by assumption} \\ (17) & e_1 \not \vdash p_1 & \text{by assumption} \\ (18) & e_1 \not \vdash p_1 & \text{by assumption} \\ (19) & e_2 \rhd p_2 \dashv \theta_2 & \text{by assumption} \\ (20) & e_2 \not \vdash p_2 & \text{by assumption} \\ (21) & e_2 \not \vdash p_2 & \text{by assumption} \\ \end{array}$

(22) (e_1, e_2) ? (p_1, p_2) by Rule (31e) on (16)

and (20)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies.

Case (30d).

(23)
$$\theta = \theta_1 \uplus \theta_2$$

$$(24) \ e_2 \rhd p_2 \dashv \theta_2$$

by assumption

Contradicts (19).

$$(25) \quad (e_1, e_2) \mathrel{\triangleright} (p_1, p_2) \dashv \theta$$

by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (32) on it, only two cases apply.

Case (32b).

(26)
$$e_1 \perp p_1$$

by assumption

Contradicts (18).

Case (32c).

(26)
$$e_2 \perp p_2$$

by assumption

Contradicts (21).

$$(27) \ \underline{(e_1,e_2) \pm (p_1,p_2)}$$

by contradiction

Case $e_1 \triangleright p_1 \dashv \theta_1, e_2 \perp p_2$.

$$(16) e_1 \rhd p_1 \dashv \theta_1$$

by assumption

$$(17) e_1 ? p_1$$

by assumption

(18)
$$e_1 + p_1$$

by assumption

(19)
$$e_2 \triangleright p_2 \# \theta_2$$

by assumption

$$(20) e_2 ? p_2$$

by assumption

(21)
$$e_2 \perp p_2$$

by assumption

(22)
$$(e_1, e_2) \perp (p_1, p_2)$$

by Rule (32c) on (21)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies.

Case (30d).

$$(23) \ \theta = \theta_1 \uplus \theta_2$$

$$(24) \quad e_2 \rhd p_2 \dashv \mid \theta_2$$

by assumption

Contradicts (19).

(25)
$$(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$

by contradiction

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (31) on it, only four cases apply.

Case (31c).

(26) (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (31d).

(26) e_1 ? p_1	by assumption
Contradicts (17) . Case $(31e)$.	
(26) $e_2 ? p_2$	by assumption
Contradicts (20) .	by assumption
Case (31f).	
$(26) e_1? p_1$	by assumption
Contradicts (17).	
,	
$(27) \ \underline{(e_1,e_2) ? (p_1,p_2)}$	by contradiction
Case $e_1 ? p_1, e_2 \rhd p_2 \dashv \theta_2$.	
$(16) \ \underline{e_1} \triangleright \underline{p_1} + \underline{\theta_1}$	by assumption
$(17) e_1 ? p_1$	by assumption
$(18) \ \underline{e_1} + \underline{p_1}$	by assumption
$(19) e_2 \rhd p_2 \dashv \mid \theta_2$	by assumption
$(20) e_2 ? p_2$	by assumption
$(21) \underline{e_2} \perp \overline{p_2}$	by assumption
$(22) (e_1, e_2) ? (p_1, p_2)$	by Rule (31d) on (17) and (19)
Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$.	By rule induction over Rules (30)
on it, only one case applies.	(00)
Case (30d).	
$(23) \ \theta = \theta_1 \uplus \theta_2$	
$(24) e_1 \rhd p_1 \dashv \theta_1$	by assumption
Contradicts (16).	
(25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$	by contradiction
	rule induction over Rules (32) on
it, only two cases apply.	
Case (32b).	
$(26) e_1 \perp p_1$	by assumption
Contradicts (18).	
Case (32c).	l
(26) $e_2 \perp p_2$	by assumption
Contradicts (21).	
(27) $(e_1, e_2) + (p_1, p_2)$	by contradiction
Case $e_1 ? p_1, e_2 ? p_2$.	
$(16) \underline{e_1 \triangleright p_1 \dashv \theta_1}$	by assumption
$(17) e_1? p_1$	by assumption

(18) $e_1 + p_1$ by assumption (19) $e_2 + p_2 + p_2$ by assumption (20) $e_2 ? p_2$ by assumption (21) $e_2 + p_2$ by assumption (22) $(e_1, e_2) ? (p_1, p_2)$ by Rule (31f) on (17) and (20)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies.

Case (30d).

- (23) $\theta = \theta_1 \uplus \theta_2$
- (24) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption

Contradicts (19).

(25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ by contradiction

Assume $(e_1, e_2) \perp (p_1, p_2)$. By rule induction over Rules (32) on it, only two cases apply.

Case (32b).

(26) $e_1 \perp p_1$ by assumption

Contradicts (18).

Case (32c).

- (26) $e_2 \perp p_2$ by assumption Contradicts (21).
- (27) $(e_1, e_2) \pm (p_1, p_2)$ by contradiction

Case $e_1 ? p_1, e_2 \perp p_2$.

 $\begin{array}{lll} (16) & \underline{e_1} \triangleright p_1 \# \theta_1 & \text{by assumption} \\ (17) & \underline{e_1} ? p_1 & \text{by assumption} \\ (18) & \underline{e_1} \# p_1 & \text{by assumption} \\ (19) & \underline{e_2} \triangleright p_2 \# \theta_2 & \text{by assumption} \\ (20) & \underline{e_2} ? p_2 & \text{by assumption} \\ (21) & \underline{e_2} \perp p_2 & \text{by assumption} \\ (22) & (\underline{e_1}, \underline{e_2}) \perp (p_1, p_2) & \text{by Rule (32c) on (21)} \end{array}$

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies.

Case (30d).

- (23) $\theta = \theta_1 \uplus \theta_2$
- (24) $e_2 \triangleright p_2 \dashv \theta_2$ by assumption Contradicts (19).
- (25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ by contradiction

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (31) on it, only four cases apply.

Case (31c).

(26) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (31d).

(26) $e_2 > p_2 \dashv \theta_2$ by assumption

Contradicts (19).

Case (31e).

(26) e_2 ? p_2 by assumption

Contradicts (20).

Case (31f).

(26) e_2 ? p_2 by assumption

Contradicts (20).

(27)
$$(e_1, e_2)$$
? (p_1, p_2) by contradiction

Case $e_1 \perp p_1, e_2 \rhd p_2 \dashv \theta_2$.

 $\begin{array}{ll} (16) \ \underline{e_1} \triangleright p_1 \# \theta_1 & \text{by assumption} \\ (17) \ \underline{e_1} \nearrow p_1 & \text{by assumption} \\ (18) \ e_1 \perp p_1 & \text{by assumption} \end{array}$

(19) $e_2 > p_2 \dashv \theta_2$ by assumption (20) $e_2 ? p_2$ by assumption

(21) $e_2 \perp p_2$ by assumption

(22) $(e_1, e_2) \perp (p_1, p_2)$ by Rule (32b) on (18)

Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies.

Case (30d).

(23) $\theta = \theta_1 \uplus \theta_2$

(24) $e_1 \rhd p_1 \dashv \theta_1$ by assumption

Contradicts (16).

(25)
$$(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$$
 by contradiction

Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (31) on it, only four cases apply.

Case (31c).

(26) (e_1, e_2) notintro by assumption

Contradicts Lemma 4.0.8.

Case (31d).

(26) $e_1 ? p_1$ by assumption

Contradicts (17).

Case (31e). $(26) e_2? p_2$ by assumption Contradicts (20). Case (31f). (26) $e_1 ? p_1$ by assumption Contradicts (17). (27) (e_1,e_2) ? (p_1,p_2) by contradiction Case $e_1 \perp p_1, e_2 ? p_2$. (16) $\underline{e_1} \triangleright p_1 + \theta_1$ by assumption $(17) e_1 ? p_1$ by assumption (18) $e_1 \perp p_1$ by assumption (19) $e_2 \triangleright p_2 + \theta_2$ by assumption (20) $e_2 ? p_2$ by assumption $(21) e_2 + p_2$ by assumption (22) $(e_1, e_2) \perp (p_1, p_2)$ by Rule (32b) on (18) Assume $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$. By rule induction over Rules (30) on it, only one case applies. Case (30d). (23) $\theta = \theta_1 \uplus \theta_2$ (24) $e_2 \triangleright p_2 \dashv \mid \theta_2$ by assumption Contradicts (19). $(25) (e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ by contradiction Assume (e_1, e_2) ? (p_1, p_2) . By rule induction over Rules (31) on it, only four cases apply. Case (31c). (26) (e_1,e_2) notintro by assumption Contradicts Lemma 4.0.8. Case (31d). (26) $e_2 \triangleright p_2 \dashv \mid \theta_2$ by assumption Contradicts (19). Case (31e). (26) $e_1 \triangleright p_1 \dashv \theta_1$ by assumption Contradicts (16). Case (31f). (26) $e_1 ? p_1$ by assumption Contradicts (17).

by contradiction

 $(27) (e_1,e_2)?(p_1,p_2)$

```
(16) e_1 \triangleright p_1 + \theta_1
                                                  by assumption
   (17) e_1 ? p_1
                                                  by assumption
   (18) e_1 \perp p_1
                                                  by assumption
   (19) e_2 > p_2 + \theta_2
                                                  by assumption
   (20) e_2 ? p_2
                                                  by assumption
   (21) e_2 \pm p_2
                                                  by assumption
   (22) (e_1, e_2) \perp (p_1, p_2)
                                                  by Rule (32b) on (18)
 Assume (e_1, e_2) \triangleright (p_1, p_2) \dashv \theta. By rule induction over Rules (30)
on it, only one case applies.
 Case (30d).
        (23) \theta = \theta_1 \uplus \theta_2
        (24) e_2 \triangleright p_2 \dashv \mid \theta_2
                                                  by assumption
      Contradicts (19).
   (25) (e_1, e_2) \supset (p_1, p_2) \dashv \theta
                                                  by contradiction
 Assume (e_1, e_2)? (p_1, p_2). By rule induction over Rules (31) on
it, only four cases apply.
 Case (31c).
        (26) (e_1,e_2) notintro
                                                  by assumption
      Contradicts Lemma 4.0.8.
 Case (31d).
        (26) e_2 \triangleright p_2 \dashv \mid \theta_2
                                                  by assumption
      Contradicts (19).
 Case (31e).
        (26) e_1 \triangleright p_1 \dashv \theta_1
                                                  by assumption
      Contradicts (16).
 Case (31f).
        (26) e_1 ? p_1
                                                  by assumption
      Contradicts (17).
   (27) (e_1,e_2)? (p_1,p_2)
                                                  by contradiction
```

Lemma 4.0.12 (Matching Coherence of Constraint). Suppose that $\cdot; \Delta_e \vdash e : \tau$ and e final and $p : \tau[\xi] \dashv \Gamma; \Delta$. Then we have

1.
$$e \models \xi \text{ iff } e \rhd p \dashv \theta$$

Case $e_1 \perp p_1, e_2 \perp p_2$.

2.
$$e \models_? \xi \text{ iff } e ? p$$

3.
$$e \not\models_{?}^{\dagger} \xi \text{ iff } e \perp p$$

Proof.

$(1) \cdot ; \Delta_e \vdash e : \tau$	by assumption
--	---------------

(2)
$$e$$
 final by assumption

(3)
$$p:\tau[\xi]\dashv \Gamma;\Delta$$
 by assumption

Given Lemma 3.0.1, Theorem 2.1, and Lemma 4.0.11, it is sufficient to prove

1.
$$e \models \xi \text{ iff } e \triangleright p \dashv \theta$$

2.
$$e \models_? \xi \text{ iff } e ? p$$

By rule induction over Rules (20) on (3).

Case (20a).

(4)
$$p = x$$
 by assumption

(5)
$$\xi = \top$$
 by assumption

1. Prove $e \models \top$ implies $e \triangleright x \dashv \theta$ for some θ .

(6)
$$e > x \dashv e/x$$
 by Rule (30a)

2. Prove $e > x \dashv \theta$ implies $e \models \top$.

(6)
$$e \models \top$$
 by Rule (14a)

3. Prove $e \models_? \top$ implies e ? x.

(6)
$$e \not\models_? \top$$
 by Lemma 2.0.3

Vacuously true.

4. Prove e ? x implies $e \models_? \top$.

By rule induction over Rules (31), we notice that either, e?x is in syntactic contradiction with all the cases, or the premise x refutable? is not derivable. Hence, e?x are not derivable. And thus vacuously true.

Case (20b).

(4)
$$p =$$
_ by assumption

(5)
$$\xi = \top$$
 by assumption

1. Prove $e \models \top$ implies $e \rhd _ \dashv \theta$ for some θ .

(6)
$$e \rhd _ \dashv \cdot$$
 by Rule (30a)

2. Prove $e \rhd _ \dashv \theta$ implies $e \models \top$.

(6)
$$e \models \top$$
 by Rule (14a)

3. Prove $e \models_? \top$ implies $e ? _$.

(6)	P	H-2	\top
(0)	e	₩?	- 1

by Lemma 2.0.3

Vacuously true.

4. Prove e? _ implies $e \models_? \xi$.

By rule induction over Rules (31), we notice that either, e?_ is in syntactic contradiction with all the cases, or the premise _ refutable? is not derivable. Hence, e?_ are not derivable. And thus vacuously true.

Case (20c).

- (4) $p = ()^w$ by assumption
- (5) $\xi = ?$ by assumption
- (6) $\overline{\xi} = ?$ by Definition 9
- 1. Prove $e \models ?$ implies $e \rhd ()^w \dashv \theta$ for some θ .
 - (7) $e \not\models ?$ by Rule (30a)

Vacuously true.

2. Prove $e \rhd ()^w \dashv \theta$ implies $e \models ?$.

By rule induction over Rules (30), we notice that $e \rhd ()^w \dashv \theta$ is in syntactic contradiction with all the cases, hence not derivable. And thus vacuously true.

- 3. Prove $e \models_?$? implies e? \emptyset^w .
 - (7) e? \emptyset^w by Rule (31a)
- 4. Prove e? $()^w$ implies $e \models_?$?.
 - (7) $e \models_?$? by Rule (16a)

Case (20d).

- (4) $p = (p_0)^w$ by assumption
- (5) $\xi = ?$ by assumption
- 1. Prove $e \models ?$ implies $e \rhd (p_0)^w \dashv \theta$ for some θ .
 - (6) $e \not\models$? by Rule (30a)

Vacuously true.

2. Prove $e \rhd (p_0)^w \dashv \theta$ implies $e \models ?$.

By rule induction over Rules (30), we notice that $e \rhd (p_0)^w \dashv \theta$ is in syntactic contradiction with all the cases, hence not derivable. And thus vacuously true.

- 3. Prove $e \models_?$? implies $e ? (p_0)^w$.
 - (6) $e ? (p_0)^w$ by Rule (31b)

4. Prove $e ? (p_0)^w$ implies $e \models_? ?$.

(6)
$$e \models_? ?$$

by Rule (16a)

Case (20e).

- (4) $p = \underline{n}$ by assumption
- (5) $\xi = \underline{n}$ by assumption
- 1. Prove $e \models \underline{n}$ implies $e \triangleright \underline{n} \dashv \theta$ for some θ .

(6)
$$e \models \underline{n}$$

by assumption

By rule induction over Rules (14) on (6), only one case applies.

Case (14b).

- (7) $e = \underline{n}$ by assumption
- (8) $\underline{n} \rhd \underline{n} \dashv l$ by Rule (30c)
- 2. Prove $e \triangleright \underline{n} \dashv \theta$ implies $e \models \underline{n}$.

(6)
$$e \rhd \underline{n} \dashv \theta$$

by assumption

By rule induction over Rules (30) on (6), only one case applies.

Case (30c).

(7) $e = \underline{n}$

by assumption

(8) $\theta = \cdot$

by assumption

(9) $\underline{n} \models \underline{n}$

- by Rule (14b)
- 3. Prove $e \models_{?} \underline{n}$ implies $e ? \underline{n}$.

(6)
$$e \models_? n$$

by assumption

By rule induction over Rules (16) on (6), only one case applies.

Case (16b).

(7) e notintro

by assumption

(8) \underline{n} refutable?

by Rule (29a)

(9) $e ? \underline{n}$

by Rule (31c) on (7)

and (8)

- 4. Prove e? \underline{n} implies $e \models_? \underline{n}$.
 - (6) e? \underline{n}

by assumption

By rule induction over Rules (31) on (6), only one case applies. Case (31c).

(7) e notintro

by assumption

(8) \underline{n} refutable?

by Rule (10a)

(9) $e \models_{?} \underline{n}$

by Rule (16) on (7)

and (8)

Case (20f).

(4) $p = \operatorname{inl}(p_1)$ by assumption (5) $\xi = \operatorname{inl}(\xi_1)$ by assumption (6) $\tau = (\tau_1 + \tau_2)$ by assumption (7) $p_1 : \tau_1[\xi_1] \dashv \Gamma; \Delta$ by assumption

By rule induction over Rules (19) on (1), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

$$(8) \ e = (\!|\!|)^u, (\!|\!|e_0|\!|)^u, e_1(e_2), \texttt{prl}(e_0), \texttt{prr}(e_0), \texttt{match}(e_0) \{\hat{rs}\}$$
 by assumption

- (9) e notintro by Rule (26a),(26b),(26c),(26d),(26e),(26f)
- 1. Prove $e \models \mathtt{inl}(\xi_1)$ implies $e \rhd \mathtt{inl}(p_1) \dashv \theta$ for some θ . By rule induction over Rules (14) on $e \models \mathtt{inl}(\xi_1)$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

2. Prove $e
ightharpoonup \operatorname{inl}(p_1) \dashv \theta$ implies $e \models \operatorname{inl}(\xi_1)$. By rule induction over Rules (30) on $e
ightharpoonup \operatorname{inl}(p_1) \dashv \theta$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

- 3. Prove $e \models_? inl(\xi_1)$ implies e? $inl(p_1)$.
 - (10) $\operatorname{inl}(p_1)$ refutable? by Rule (29d) (11) e? $\operatorname{inl}(p_1)$ by Rule (31c) on (9)

and (10)

- 4. Prove e? $inl(p_1)$ implies $e \models_? inl(\xi_1)$.
 - (10) $\operatorname{inl}(\xi_1)$ refutable? by Rule (10b) (11) $e \models_? \operatorname{inl}(\xi_1)$ by Rule (16b) on (9) and (10)

Case (19j).

 $\begin{array}{ll} (8) \ e = \operatorname{inl}_{\tau_2}(e_1) & \text{by assumption} \\ (9) \ \cdot \ ; \Delta_e \vdash e_1 : \tau_1 & \text{by assumption} \\ (10) \ e_1 \ \text{final} & \text{by Lemma 4.0.2 on (2)} \end{array}$

By inductive hypothesis on (10) and (9) and (7).

- (11) $e_1 \models \xi_1 \text{ iff } e_1 \rhd p_1 \dashv \theta \text{ for some } \theta$
- (12) $e_1 \models_? \xi_1 \text{ iff } e_1 ? p_1$
- 1. Prove $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ implies $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ for some θ .
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by assumption

By rule induction over Rules (14) on (13), only one case applies. Case (14g).

- (14) $e_1 \models \xi_1$ by assumption
- (15) $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 by (11) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta_1$ by Rule (30e) on (15)
- 2. Prove $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ implies $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \rhd \operatorname{inl}(p_1) \dashv \theta$ by assumption

By rule induction over Rules (30) on (13), only one case applies. Case (30e).

- $(14) e_1 \rhd p_1 \dashv \theta$
- (15) $e_1 \models \xi_1$ by (11) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\xi_1)$ by Rule (14g) on (15)

by assumption

- 3. Prove $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ implies $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by assumption

By rule induction over Rules (16) on (13), only two cases apply. Case (16b).

(14) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption Contradicts Lemma 4.0.6.

Case (16e).

- (14) $e_1 \models_? \xi_1$ by assumption
- (15) $e_1 ? p_1$ by (12) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ by Rule (31g) on (15)
- 4. Prove $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ implies $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$.
 - (13) $\operatorname{inl}_{\tau_2}(e_1)$? $\operatorname{inl}(p_1)$ by assumption

By rule induction over Rules (31) on (13), only two cases apply. Case (31c).

(14) $\operatorname{inl}_{\tau_2}(e_1)$ notintro by assumption Contradicts Lemma 4.0.6.

Case (31g).

- (14) $e_1 ? p_1$ by assumption
- (15) $e_1 \models_? \xi_1$ by (12) on (14)
- (16) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\xi_1)$ by Rule (16e) on (15)

Case (20g).

- (4) $p = inr(p_2)$ by assumption
- (5) $\xi = \operatorname{inr}(\xi_2)$ by assumption
- (6) $\tau = (\tau_1 + \tau_2)$ by assumption
- (7) $p_2: \tau_2[\xi_2] \dashv \Gamma; \Delta$ by assumption

By rule induction over Rules (19) on (1), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

 $(8) \ e = (\!()^u, (\!(e_0)\!)^u, e_1(e_2), \mathtt{prl}(e_0), \mathtt{prr}(e_0), \mathtt{match}(e_0) \{\hat{rs}\}$

by assumption

(9) e notintro by Rule

(26a),(26b),(26c),(26d),(26e),(26f)

1. Prove $e \models \operatorname{inr}(\xi_2)$ implies $e \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ . By rule induction over Rules (14) on $e \models \operatorname{inr}(\xi_2)$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

2. Prove $e
ightharpoonup \operatorname{inr}(p_2) \dashv \theta$ implies $e \models \operatorname{inr}(\xi_2)$. By rule induction over Rules (30) on $e \rhd \operatorname{inr}(p_2) \dashv \theta$, no case applies due to syntactic contradiction.

Therefore, vacuously true.

- 3. Prove $e \models_? \operatorname{inr}(\xi_2)$ implies $e ? \operatorname{inr}(p_2)$.
 - (10) $inr(p_2)$ refutable? by Rule (29e)
 - (11) e? $inr(p_2)$ by Rule (31c) on (9) and (10)
- 4. Prove e? $inr(p_2)$ implies $e \models_? inr(\xi_2)$.
 - (10) $inr(\xi_2)$ refutable? by Rule (10c)
 - (11) $e \models_? \operatorname{inr}(\xi_2)$ by Rule (16b) on (9) and (10)

Case (19k).

- (8) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption
- (9) \cdot ; $\Delta_e \vdash e_2 : \tau_2$ by assumption
- (10) e_2 final by Lemma 4.0.2 on (2)

By inductive hypothesis on (10) and (9) and (7).

- (11) $e_2 \models \xi_2 \text{ iff } e_2 \triangleright p_2 \dashv \theta \text{ for some } \theta$
- (12) $e_2 \models_? \xi_2 \text{ iff } e_2 ? p_2$
- 1. Prove $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ implies $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ for some θ .
 - (13) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ by assumption

By rule induction over Rules (14) on (13), only one case applies.

Case (14g).

- (14) $e_2 \models \xi_2$ by assumption
- (15) $e_2 \triangleright p_2 \dashv \theta_1$ for some θ_1 by (11) on (14)
- (16) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta_1$ by Rule (30e) on (15)
- 2. Prove $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ implies $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$.
 - (13) $\operatorname{inr}_{\tau_1}(e_2) \rhd \operatorname{inr}(p_2) \dashv \theta$ by assumption

By rule induction over Rules (30) on (13), only one case applies. Case (30e).

- (14) $e_2 > p_2 \dashv \theta$ by assumption (15) $e_2 \models \xi_2$ by (11) on (14)
- (16) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\xi_2)$ by Rule (14g) on (15)
- 3. Prove $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$ implies $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$.
 - (13) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$ by assumption

By rule induction over Rules (16) on (13), only two cases apply. Case (16b).

(14) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption Contradicts Lemma 4.0.6.

Case (16e).

- (14) $e_2 \models_? \xi_2$ by assumption
- (15) $e_2 ? p_2$ by (12) on (14)
- (16) $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$ by Rule (31g) on (15)
- 4. Prove $\operatorname{inr}_{\tau_1}(e_2)$? $\operatorname{inr}(p_2)$ implies $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\xi_2)$.
 - (13) $inr_{\tau_1}(e_2)$? $inr(p_2)$

by assumption

By rule induction over Rules (31) on (13), only two cases apply. Case (31c).

(14) $\operatorname{inr}_{\tau_1}(e_2)$ notintro by assumption Contradicts Lemma 4.0.6.

Case (31g).

- (14) e_2 ? p_2 by assumption
- (15) $e_2 \models_? \xi_2$ by (12) on (14) (16) $\inf_{\tau_1}(e_2) \models_? \inf(\xi_2)$ by Rule (16e) on (15)

Case (20h).

- (4) $p = (p_1, p_2)$ by assumption
- (5) $\xi = (\xi_1, \xi_2)$ by assumption
- (6) $\tau = (\tau_1 \times \tau_2)$ by assumption
- (7) $\Gamma = \Gamma_1 \uplus \Gamma_2$ by assumption
- (8) $\Delta = \Delta_1 \uplus \Delta_2$ by assumption
- (9) $p_1: \tau_1[\xi_1] \dashv \Gamma_1; \Delta_1$ by assumption
- (10) $p_2: \tau_2[\xi_2] \dashv \Gamma_2; \Delta_2$ by assumption

By rule induction over Rules (19) on (1), the following cases apply.

Case (19b),(19c),(19f),(19h),(19i),(19l),(19m).

(11) $e = (v, (e_0)^u, e_1(e_2), prl(e_0), prr(e_0), match(e_0) \{\hat{rs}\}\$ by assumption

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by Rule
  (12) e notintro
                                                (26a),(26b),(26c),(26d),(26e),(26f)
  (13) e indet
                                                by Lemma 4.0.9 on (2)
                                                and (12)
  (14) prl(e) indet
                                                by Rule (24g) on (13)
  (15) prl(e) final
                                                by Rule (25b) on (14)
 (16) prr(e) indet
                                                by Rule (24h) on (13)
  (17) prr(e) final
                                                by Rule (25b) on (16)
  (18) \cdot; \Delta \vdash \mathtt{prl}(e) : \tau_1
                                                by Rule (19h) on (1)
  (19) \cdot; \Delta \vdash \mathsf{prr}(e) : \tau_2
                                               by Rule (19i) on (1)
By inductive hypothesis on (9) and (18) and (15) and by inductive
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hypothesis on (10) and (19) and (17).

- (20) $\operatorname{prl}(e) \models \xi_1 \text{ iff } \operatorname{prl}(e) \triangleright p_1 \dashv \theta_1 \text{ for some } \theta_1$
- (21) $prl(e) \models_? \xi_1 \text{ iff } prl(e) ? p_1$
- (22) $\operatorname{prr}(e) \models \xi_2 \text{ iff } \operatorname{prr}(e) \triangleright p_2 \dashv \theta_2 \text{ for some } \theta_2$
- (23) $prr(e) \models_? \xi_2 \text{ iff } prr(e) ? p_2$
- 1. Prove $e \models (\xi_1, \xi_2)$ implies $e \triangleright (p_1, p_2) \dashv \theta$ for some θ .

(24)
$$e \models (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (14) on (24), only one case applies.

Case (14j).

- (25) $prl(e) \models \xi_1$ by assumption (26) $prr(e) \models \xi_2$ by assumption (27) $prl(e) \triangleright p_1 \dashv \theta_1$ by (20) on (25) (28) $prr(e) \triangleright p_2 \dashv \theta_2$ by (22) on (26) by Rule (30g) on (12)(29) $e \rhd (p_1, p_2) \dashv \theta_1 \uplus \theta_2$
- 2. Prove $e \triangleright (p_1, p_2) \dashv \theta$ implies $e \models (\xi_1, \xi_2)$.

$$(24) e \rhd (p_1, p_2) \dashv \theta$$

by assumption

and (27) and (28)

By rule induction over Rules (30) on (24), only one case applies.

Case (30g).

$$\begin{array}{lll} (25) & \theta = \theta_1 \uplus \theta_2 & \text{by assumption} \\ (26) & \mathtt{prl}(e) \rhd \xi_1 \dashv \theta_1 & \text{by assumption} \\ (27) & \mathtt{prr}(e) \rhd \xi_2 \dashv \theta_2 & \text{by assumption} \\ (28) & \mathtt{prl}(e) \vDash \xi_1 & \text{by } (20) \text{ on } (26) \\ (29) & \mathtt{prr}(e) \vDash \xi_2 & \text{by } (22) \text{ on } (27) \\ (30) & e \vDash (\xi_1, \xi_2) & \text{by Rule } (14\text{j}) \text{ on } (12) \\ & \text{and } (28) \text{ and } (29) \\ \end{array}$$

3. Prove $e \models_{?} (\xi_1, \xi_2)$ implies $e ? (p_1, p_2)$.

(24)
$$e \models_{?} (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (16) on (24), only one case applies. **Case** (16b).

- (25) (ξ_1, ξ_2) refutable? by assu
 - by assumption

By rule induction over Rules (10) on (25), only two cases apply.

Case (10d).

- (26) ξ_1 refutable? by assumption (27) prl(e) notintro by Rule (26e)
- (28) $prl(e) \models_? \xi_1$ by Rule (16b) on (26) and (27)
- (29) prl(e)? p_1 by (21) on (28)

By rule induction over Rules (31) on (29), only three cases apply.

Case (31a),(31b).

- (30) $p_1 = \{ \}^w, \{ p_0 \}^w$ by assumption
- (31) p_1 refutable? by Rule (29b) and Rule (29c)
- (32) (p_1, p_2) refutable? by Rule (29f) on (31)
- (33) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (32)

Case (31c).

- (30) p_1 refutable? by assumption
- (31) (p_1, p_2) refutable? by Rule (29f) on (30)
- (32) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (31)

Case (10e).

- (26) ξ_2 refutable? by assumption (27) prr(e) notintro by Rule (26e)
- (28) $\operatorname{prr}(e) \models_{?} \xi_{2}$ by Rule (16b) on (26)
- (29) prr(e)? p_2 and (27) by (23) on (28)

By rule induction over Rules (31) on (29), only three cases apply.

Case (31a),(31b).

- (30) $p_2 = \{ \}^w, \{ p_0 \}^w$ by assumption
- (31) p_2 refutable? by Rule (29b) and Rule (29c)
- (32) (p_1, p_2) refutable? by Rule (29g) on (31)
- (33) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (32)

Case (31c).

- (30) p_2 refutable? by assumption (31) (p_1, p_2) refutable? by Rule (29g) on (30)
- (32) $e ? (p_1, p_2)$ by Rule (31c) on (12) and (31)
- 4. Prove $e ? (p_1, p_2)$ implies $e \models_? (\xi_1, \xi_2)$.

$$(24)$$
 $e?(p_1,p_2)$

by assumption

By rule induction over Rules (31) on (24), only one case applies. Case (31c).

(25) (p_1, p_2) refutable? by assumption

By rule induction over Rules (29) on (25), only two cases apply.

Case (29f).

- (26) p_1 refutable? by assumption
- (27) prl(e) notintro by Rule (26e)
- (28) prl(e) ? p_1 by Rule (31c) on (26) and (27)
- (29) $prl(e) \models_{?} \xi_1$ by (21) on (28)

By rule induction over Rules (16) on (29), only three cases apply.

Case (16a).

- (30) $\xi_1 = ?$ by assumption
- (31) ξ_1 refutable? by Rule (2b)
- (32) (ξ_1, ξ_2) refutable? by Rule (10d) on (31)
- (33) $e \models_{?} (\xi_1, \xi_2)$ by Rule (16b) on (12) and (32)

Case (16b).

- (30) ξ_1 refutable? by assumption
- (31) (ξ_1, ξ_2) refutable? by Rule (10d) on (30)
- (32) $e \models_{?} (\xi_1, \xi_2)$ by Rule (16b) on (12) and (31)

Case (29g).

- (26) p_2 refutable? by assumption (27) prr(e) notintro by Rule (26e)
- (28) prr(e) ? p_2 by Rule (31c) on (26) and (27)
- (29) $prr(e) \models_? \xi_2$ by (23) on (28)

By rule induction over Rules (16) on (29), only three cases apply.

Case (16a).

(30) $\xi_2 = ?$ by assumption (31) ξ_2 refutable? by Rule (2b)

(32)
$$(\xi_1, \xi_2)$$
 refutable? by Rule (10e) on (31)

(33)
$$e \models_{?} (\xi_1, \xi_2)$$
 by Rule (16b) on (12) and (32)

Case (16b).

- (30) ξ_2 refutable? by assumption
- (31) (ξ_1, ξ_2) refutable? by Rule (10e) on (30)
- (32) $e \models_{?} (\xi_1, \xi_2)$ by Rule (16b) on (12) and (31)

Case (19g).

(11)
$$e = (e_1, e_2)$$
 by assumption

(12)
$$\cdot$$
; $\Delta_e \vdash e_1 : \tau_1$ by assumption

(13)
$$\cdot$$
; $\Delta_e \vdash e_2 : \tau_2$ by assumption

(14)
$$e_1$$
 final by Lemma 4.0.4 on (2)

(15)
$$e_2$$
 final by Lemma 4.0.4 on (2)

By inductive hypothesis on (14) and (12) and (9) and by inductive hypothesis on (15) and (13) and (10).

(16)
$$e_1 \models \xi_1 \text{ iff } e_1 \triangleright p_1 \dashv \theta_1 \text{ for some } \theta_1$$

(17)
$$e_1 \models_? \xi_1 \text{ iff } e_1 ? p_1$$

(18)
$$e_2 \models \xi_2 \text{ iff } e_2 \triangleright p_2 \dashv \theta_2 \text{ for some } \theta_2$$

(19)
$$e_2 \models_? \xi_2 \text{ iff } e_2 ? p_2$$

1. Prove $(e_1, e_2) \models (\xi_1, \xi_2)$ implies $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ for some θ .

$$(20) (e_1, e_2) \models (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (14) on (20), only two cases apply.

Case (14i).

- (21) $e_1 \models \xi_1$ by assumption
- (22) $e_2 \models \xi_2$ by assumption
- (23) $e_1 \triangleright p_1 \dashv \theta_1$ for some θ_1 by (16) on (21)
- (24) $e_2 \triangleright p_2 \dashv \theta_2$ for some θ_2 by (18) on (22)
- (25) $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta_1 \uplus \theta_2$ by Rule (30d) on (23) and (24)

Case (14j).

(21) (e_1, e_2) notintro by ass

by assumption

Contradicts Lemma 4.0.8.

- 2. Prove $(e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$ implies $(e_1, e_2) \models (\xi_1, \xi_2)$.
 - $(20) (e_1, e_2) \triangleright (p_1, p_2) \dashv \theta$

by assumption

By rule induction over Rules (30) on (20), only two cases apply. **Case** (30d).

(21) $e_1 \triangleright p_1 \dashv \mid \theta_1 \text{ for some } \theta_1 \quad \text{by assumption}$

```
(22) e_2 \triangleright p_2 \dashv \theta_2 for some \theta_2 by assumption
```

(23)
$$e_1 \models \xi_1$$
 by (16) on (21)

(24)
$$e_2 \models \xi_2$$
 by (18) on (22)

(25)
$$(e_1, e_2) \models (\xi_1, \xi_2)$$
 by Rule (14i) on (23) and (24)

Case (30g).

(21)
$$(e_1, e_2)$$
 notintro by assumption

Contradicts Lemma 4.0.8.

3. Prove
$$(e_1, e_2) \models_? (\xi_1, \xi_2)$$
 implies $(e_1, e_2) ? (p_1, p_2)$.

$$(20) (e_1, e_2) \models_? (\xi_1, \xi_2)$$

by assumption

By rule induction over Rules (16) on (20), only four cases apply. Case (16b).

(21)
$$(e_1, e_2)$$
 notintro by assumption Contradicts Lemma 4.0.8.

Case (16g).

(21)
$$e_1 \models_? \xi_1$$
 by assumption

(22)
$$e_2 \models \xi_2$$
 by assumption

(23)
$$e_1 ? p_1$$
 by (17) on (21)

(24)
$$e_2 > p_2 \dashv \theta_2$$
 by (18) on (22)

(25)
$$(e_1, e_2)$$
? (p_1, p_2) by Rule (31d) on (23) and (24)

Case (16h).

(21)
$$e_1 \models \xi_1$$
 by assumption

(22)
$$e_2 \models_? \xi_2$$
 by assumption

(23)
$$e_1 > p_1 \dashv \theta_1$$
 by (16) on (21)

(24)
$$e_2$$
? p_2 by (19) on (22)

(25)
$$(e_1, e_2)$$
? (p_1, p_2) by Rule (31e) on (23) and (24)

Case (16i).

(21)
$$e_1 \models_? \xi_1$$
 by assumption

(22)
$$e_2 \models_? \xi_2$$
 by assumption

(23)
$$e_1 ? p_1$$
 by (17) on (21)

(24)
$$e_2$$
? p_2 by (19) on (22)

(25)
$$(e_1, e_2)$$
? (p_1, p_2) by Rule (31f) on (23) and (24)

4. Prove (e_1, e_2) ? (p_1, p_2) implies $(e_1, e_2) \models_? (\xi_1, \xi_2)$.

(20)
$$(e_1, e_2)$$
? (p_1, p_2) by assumption

By rule induction over Rules (31) on (20), only four cases apply. Case (31c).

$$(21)$$
 (e_1,e_2) notintro

by assumption

Contradicts Lemma 4.0.8.

Case (31d).

(21)	$e_1 ? p_1$	by assumption
(22)	$e_2\rhd p_2\dashv \mid \theta_2$	by assumption
(23)	$e_1 \models_? \xi_1$	by (17) on (21)
(24)	$e_2 \models \xi_2$	by (18) on (22)
(25)	$(e_1, e_2) ? (p_1, p_2)$	by Rule (16g) on (23)
		and (24)

Case (31e).

$(21) e_1 \rhd p_1 \dashv \theta_1$	by assumption
$(22) e_2? p_2$	by assumption
(23) $e_1 \models \xi_1$	by (16) on (21)
$(24) e_2 \models_? \xi_2$	by (19) on (22)
$(25) (e_1, e_2)? (p_1, p_2)$	by Rule (16h) on (23)
	and (24)

Case (31f).

$(21) e_1? p_1$	by assumption
$(22) e_2? p_2$	by assumption
(23) $e_1 \models_? \xi_1$	by (17) on (21)
$(24) e_2 \models_? \xi_2$	by (19) on (22)
$(25) (e_1, e_2) ? (p_1, p_2)$	by Rule (16i) on (23)
	and (24)

5 Preservation and Progress

Theorem 5.1 (Preservation). If \cdot ; $\Delta \vdash e : \tau$ and $e \mapsto e'$ then \cdot ; $\Delta \vdash e' : \tau$

Proof. By rule induction over Rules (19) on typing judgment of e. For simplicity, we only consider two cases for match expressions here.

Case (191).

$$\begin{array}{lll} (1) & \cdot ; \Delta \vdash \mathtt{match}(e_1)\{\cdot \mid r \mid rs\} : \tau & \text{by assumption} \\ (2) & \mathtt{match}(e_1)\{\cdot \mid r \mid rs\} \mapsto e' & \text{by assumption} \\ (3) & \cdot ; \Delta \vdash e_1 : \tau_1 & \text{by assumption} \\ (4) & \cdot ; \Delta \vdash [\bot](r \mid rs) : \tau_1[\xi] \Rightarrow \tau & \text{by assumption} \\ (5) & \top \models_{7}^{+} \xi & \text{by assumption} \end{array}$$

By rule induction over Rules (33) on (2).

Case (33k).

$$(6) \ e' = \mathtt{match}(e'_1) \{\cdot \mid r \mid rs\} \qquad \text{by assumption}$$

$$(7) \ e_1 \mapsto e'_1 \qquad \text{by assumption}$$

$$(8) \ \cdot ; \Delta \vdash e'_1 : \tau_1 \qquad \text{by Ho } (3) \text{ and } (7)$$

$$(9) \ \cdot ; \Delta \vdash \mathtt{match}(e'_1) \{\cdot \mid r \mid rs\} : \tau \qquad \text{by Rule } (191) \text{ on } (8)$$

$$\text{and } (4) \text{ and } (5)$$

$$\text{Case } (331).$$

$$(6) \ r = p_r \Rightarrow e_r \qquad \text{by assumption}$$

$$(7) \ e' = [\theta](e_r) \qquad \text{by assumption}$$

$$(8) \ e_1 \rhd p_r \dashv \theta \qquad \text{by assumption}$$

$$(8) \ e_1 \rhd p_r \dashv \theta \qquad \text{by assumption}$$

$$\text{By rule induction over Rules } (22) \text{ on } (4).$$

$$\text{Case } (22a).$$

$$(9) \ \xi = \xi_r \qquad \text{by assumption}$$

$$(10) \ \cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau \qquad \text{by assumption}$$

$$(11) \ p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r \qquad \text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$(13) \ \theta : \Gamma_r \qquad \text{by Lemma } 3.0.7 \text{ on } (3)$$

$$\text{and } (11) \text{ and } (8)$$

$$\text{by Lemma } 3.0.6 \text{ on }$$

$$(12) \ \Gamma_r : \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau \qquad \text{by assumption}$$

$$\text{by assumption}$$

$$\text{by Lemma } 3.0.6 \text{ on }$$

$$(12) \text{ and } (13)$$

$$\text{Case } (22b). \qquad (9) \ \xi = \xi_r \lor \xi_{rs} \qquad \text{by assumption}$$

$$\text{by assumption}$$

$$\text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$\text{(12)} \ \Gamma_r : \Delta \uplus \Delta_r \vdash e_r : \tau \qquad \text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$\text{(21)} \ \Gamma_r : \Delta \uplus \Delta_r \vdash e_r : \tau \qquad \text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$\text{(21)} \ \text{on } (10)$$

$$\text{(3)} \ \theta : \Gamma_r \qquad \text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$\text{(21)} \ \text{on } (10)$$

$$\text{(3)} \ \theta : \Gamma_r \qquad \text{by Inversion of Rule}$$

$$(21a) \text{ on } (10)$$

$$\text{(21)} \ \text{on } (10)$$

$$\text{(21)} \ \text{on } (10)$$

$$\text{(3)} \ \text{and } (11) \text{ and } (8)$$

$$\text{(41)} \ \cdot ; \Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau \qquad \text{by Lemma } 3.0.7 \text{ on } (3)$$

$$\text{and } (11) \text{ and } (8)$$

$$\text{(21a)} \text{ on } (10)$$

$$\text{(21a)} \text{ on }$$

By rule induction over Rules (22) on (4).

(8) e_1 final

(9) $e_1 \perp p_r$

by assumption

by assumption

by assumption

Case (22a). Syntactic contradiction of rs.

Case (22b).

(10)
$$\xi = \xi_r \vee \xi_{rs}$$
 by assumption

(11)
$$\cdot ; \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$$
 by assumption

$$(12) \ \cdot ; \Delta \vdash [\bot \lor \xi_r](r' \mid rs') : \tau_1[\xi_{rs}] \Rightarrow \tau$$

by assumption

(13)
$$\xi_r \not\models \bot$$
 by assumption

(14)
$$p_r: \tau_1[\xi_r] \dashv \Gamma_r; \Delta_r$$
 by Inversion of Rule (21a) on (11)

(15)
$$\Gamma_r$$
; $\Delta \uplus \Delta_r \vdash e_r : \tau$ by Inversion of Rule (21a) on (11)

$$\begin{array}{cc} (16) & \cdot\; ; \Delta \vdash [\bot](p_r \Rightarrow e_r \mid \cdot) : \tau_1[\xi_r] \Rightarrow \tau \\ & \text{by Rule (22a) on (11)} \\ & \text{and (13)} \end{array}$$

(17)
$$e_1 \not\models_{?}^{\dagger} \xi_r$$
 by Lemma 4.0.12 on (3) and (8) and (14) and (9)

(18)
$$\cdot$$
; $\Delta \vdash \mathtt{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'\} : \tau$ by Rule (19m) on (3) and (8) and (16) and (12) and (17) and (5)

Case (19m).

(1)
$$rs_{pre} = r_{pre} \mid rs'_{pre}$$
 by assumption

(2)
$$\cdot$$
; $\Delta \vdash \mathtt{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\} : \tau$ by assumption

(3)
$$\operatorname{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\} \mapsto e'$$
 by assumption

(4)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$ by assumption

(5)
$$e_1$$
 final by assumption

(6)
$$\cdot ; \Delta \vdash [\bot] rs_{pre} : \tau_1[\xi_{pre}] \Rightarrow \tau$$
 by assumption

$$(7) \quad \cdot \; ; \Delta \vdash [\bot \lor \xi_{pre}](r \mid rs_{post}) : \tau_1[\xi_{rest}] \Rightarrow \tau$$
 by assumption

(8)
$$e_1 \not\models_?^{\dagger} \xi_{pre}$$
 by assumption

(9)
$$\top \models_{?}^{\dagger} \xi_{pre} \vee \xi_{rest}$$
 by assumption

By rule induction over Rules (33) on (3).

Case (33k).

(10)
$$e' = \text{match}(e'_1)\{rs_{pre} \mid r \mid rs_{post}\}$$
 by assumption (11) $e_1 \mapsto e'_1$ by assumption

By Lemma 4.0.10, (11) contradicts (5).

Case (331).

(10)
$$r = p_r \Rightarrow e_r$$
 by assumption

(11)
$$e' = [\theta](e_r)$$
 by assumption

(12)
$$e_1 \triangleright p_r \dashv \theta$$
 by assumption

By rule induction over Rules (22) on (7).

Case (22a).

(13)
$$\xi_{rest} = \xi_r$$
 by assumption
(14) \cdot ; $\Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$ by assumption

(15)
$$p_r : \tau_1[\xi_r] \dashv \Gamma_r : \Delta_r$$
 by Inversion of Rule

(16)
$$\Gamma_r$$
; $\Delta \uplus \Delta_r \vdash e_r : \tau$ (21a) on (14) by Inversion of Rule (21a) on (14)

(17)
$$\theta: \Gamma_r$$
 by Lemma 3.0.7 on (4) and (15) and (12)

(18)
$$\cdot$$
; $\Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau$ by Lemma 3.0.6 on (16) and (17)

Case (22b).

(13)
$$\xi_{rest} = \xi_r \vee \xi_{rs}$$
 by assumption

(14)
$$\cdot : \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$$
 by assumption

(15)
$$p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r$$
 by assumption

(16)
$$\Gamma_r : \Delta \uplus \Delta_r \vdash e_r : \tau$$
 by assumption

(17)
$$\theta: \Gamma_r$$
 by Lemma 3.0.7 on (4) and (15) and (12)

(18)
$$\cdot$$
; $\Delta \uplus \Delta_r \vdash [\theta](e_r) : \tau$ by Lemma 3.0.6 on (16) and (17)

Case (33m).

(10)
$$r = p_r \Rightarrow e_r$$
 by assumption

(11)
$$rs_{post} = r' \mid rs'$$
 by assumption

(12)
$$e' = \operatorname{match}(e_1)\{(rs_{pre} \mid p_r \Rightarrow e_r \mid \cdot)^{\diamond} \mid r' \mid rs'\}$$
 by assumption

(13)
$$e_1 \perp p_r$$
 by assumption

By rule induction over Rules (22) on (7).

Case (22a). Syntactic contradiction of rs_{post} .

Case (22b).

(14)
$$\xi_{rest} = \xi_r \vee \xi_{post}$$
 by assumption

(15)
$$\cdot : \Delta \vdash (p_r \Rightarrow e_r) : \tau_1[\xi_r] \Rightarrow \tau$$
 by assumption

(16)
$$\cdot ; \Delta \vdash [\bot \lor \xi_{pre} \lor \xi_r](r' \mid rs') : \tau_1[\xi_{post}] \Rightarrow \tau$$
 by assumption

(17)
$$\xi_r \not\models \xi_{pre}$$
 by assumption

Theorem 5.2 (Progress). If \cdot ; $\Delta \vdash e : \tau$ then either e final or $e \mapsto e'$ for some e'.

Proof. By rule induction over Rules (19) on typing judgment of e. For simplicity, we only consider two cases for match expressions here.

Case (191).

$$\begin{array}{ll} (1) \quad \cdot \; ; \; \Delta \vdash \mathsf{match}(e_1)\{\cdot \mid r \mid rs\} : \tau & \text{by assumption} \\ (2) \quad \cdot \; ; \; \Delta \vdash e_1 : \tau_1 & \text{by assumption} \\ (3) \quad \cdot \; ; \; \Delta \vdash [\bot](r \mid rs) : \tau_1[\xi] \Rightarrow \tau & \text{by assumption} \\ (4) \quad \top \models_{?}^{\dagger} \; \xi & \text{by assumption} \end{array}$$

By IH on (2).

Case Scrutinee takes a step.

(5)
$$e_1 \mapsto e_1'$$
 by assumption
(6) $\operatorname{match}(e_1)\{\cdot \mid r \mid rs\} \mapsto \operatorname{match}(e_1')\{\cdot \mid r \mid rs\}$ by Rule (33k) on (5)

Case Scrutinee is final.

$$(5)$$
 e_1 final

by assumption

By rule induction over Rules (22) on (3).

Case (22a).

- (6) $rs = \cdot$ by assumption
- (7) $\xi = \xi_r$ by assumption
- (8) \cdot ; $\Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau$ by assumption
- (9) $r = p_r \Rightarrow e_r$ by Inversion of Rule (21a) on (8)
- (10) $p_r: \tau_1[\xi_r] \dashv \Gamma_r; \Delta_r$ by Inversion of Rule (21a) on (8)
- (11) $e_1 \models^{\dagger}_{?} \xi_r$ by Corollary 2.1.1 on (5) and (4)

By rule induction over Rules (17) on (11).

Case (17a).

- (12) $e_1 \models_? \xi_r$ by assumption
- (13) $e_1 ? p_r$ by Lemma 4.0.12 on (2) and (5) and (10) and (12)
- (15) $\mathrm{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid\cdot\}$ final by Rule (25b) on (14)

Case (17b).

- (12) $e_1 \models \xi_r$ by assumption
- (13) $e_1 \triangleright p_r \dashv \theta$ by Lemma 4.0.12 on (2) and (5) and (10) and (12)
- (14) $\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\} \mapsto [\theta](e_r)$ by Rule (331) on (5) and (13)

Case (22b).

- (6) $rs = r' \mid rs'$ by assumption
- (7) $\xi = \xi_r \vee \xi_{rs}$ by assumption
- (8) $\cdot : \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau$ by assumption
- (9) $r = p_r \Rightarrow e_r$ by Inversion of Rule (21a) on (8)
- (10) $p_r : \tau_1[\xi_r] \dashv \Gamma_r ; \Delta_r$ by Inversion of Rule (21a) on (8)

By Lemma 4.0.11 on (2) and (5) and (10).

Case Scrutinee matches pattern.

(11)
$$e_1 \triangleright p_r \dashv \theta$$
 by assumption

(12)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs\} \mapsto [\theta](e_r)$$
 by Rule (331) on (5) and (11)

Case Scrutinee may matches pattern.

(11)
$$e_1 ? p_r$$
 by assumption

(12)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs\}$$
 indet by Rule (24k) on (5) and (11)

(13)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs\}$$
 final by Rule (25b) on (12)

Case Scrutinee doesn't matche pattern.

(11)
$$e_1 \perp p_r$$
 by assumption

$$\begin{array}{ll} (12) \ \operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid (r' \mid rs')\} \\ \mapsto \operatorname{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'\} \\ & \text{by Rule (33m) on (5)} \\ & \text{and (11)} \end{array}$$

Case (19m).

(1)
$$rs_{pre} = r_{pre} \mid rs'_{pre}$$
 by assumption

(2)
$$\cdot$$
; $\Delta \vdash \text{match}(e_1)\{rs_{pre} \mid r \mid rs_{post}\}: \tau$ by assumption

(3)
$$\cdot$$
; $\Delta \vdash e_1 : \tau_1$ by assumption

(4)
$$e_1$$
 final by assumption

(5)
$$\cdot ; \Delta \vdash [\xi_{pre}](r \mid rs_{post}) : \tau_1[\xi_{rest}] \Rightarrow \tau$$
 by assumption

(6)
$$e_1 \not\models_{?}^{\dagger} \xi_{pre}$$
 by assumption

(7)
$$\top \models_{2}^{\dagger} \xi_{pre} \vee \xi_{rest}$$
 by assumption

By rule induction over Rules (22) on (5).

Case (22a).

$$\begin{array}{ll} \text{(5)} \ \ rs_{post} = \cdot & \text{by assumption} \\ \text{(6)} \ \ \xi_{rest} = \xi_r & \text{by assumption} \\ \text{(7)} \ \ \cdot \ ; \ \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau & \text{by assumption} \\ \text{(8)} \ \ r = p_r \Rightarrow e_r & \text{by Inversion of Rule} \\ \text{(9)} \ \ p_r : \tau_1[\xi_r] \dashv \Gamma_r \ ; \ \Delta_r & \text{by Inversion of Rule} \\ \text{(9)} \ \ p_r : \tau_1[\xi_r] \dashv \Gamma_r \ ; \ \Delta_r & \text{by Inversion of Rule} \\ \text{(21a) on (7)} \end{array}$$

(10)
$$e_1 \models^{\dagger}_{?} \xi_{pre} \lor \xi_r$$
 by Corollary 2.1.1 on (4) and (7)

(11)
$$e_1 \models^{\dagger}_{?} \xi_r$$
 by Lemma 2.0.8 on (10) and (6)

By rule induction over Rules (17) on (11).

Case (17a).

(12)
$$e_1 \models_? \xi_r$$
 by assumption

(13)
$$e_1$$
 ? p_r by Lemma 4.0.12 on (3) and (4) and (9) and (12)

(14)
$$\mathrm{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid\cdot\}$$
 indet

(15)
$$\mathrm{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid\cdot\}$$
 final

by Rule (25b) on (14)

Case (17b).

(12)
$$e_1 \models \xi_r$$
 by assumption

(13)
$$e_1 \rhd p_r \dashv \theta$$
 by Lemma 4.0.12 on (3) and (4) and (9) and (12)

(14)
$$\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid \cdot\} \mapsto [\theta](e_r)$$
 by Rule (331) on (4) and (13)

Case (22b).

$$\begin{array}{ll} (5) \ \ rs_{post} = r' \mid rs'_{post} & \text{by assumption} \\ (6) \ \ \cdot \ ; \Delta \vdash r : \tau_1[\xi_r] \Rightarrow \tau & \text{by assumption} \\ (7) \ \ r = p_r \Rightarrow e_r & \text{by Inversion of Rule} \end{array}$$

7)
$$r = p_r \Rightarrow e_r$$
 by Inversion of Rule (21a) on (6)

(8)
$$p_r: \tau_1[\xi_r] \dashv \Gamma_r; \Delta_r$$
 by Inversion of Rule (21a) on (6)

By Lemma 4.0.11 on (3) and (4) and (8).

Case Scrutinee matches pattern.

$$\begin{array}{ll} (9) & e_1 \rhd p_r \dashv \theta & \text{by assumption} \\ (10) & \mathtt{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs_{post}\} \mapsto [\theta](e_r) \\ & \text{by Rule (33l) on (4)} \\ & \text{and (9)} \end{array}$$

Case Scrutinee may matches pattern.

(9)
$$e_1 ? p_r$$
 by assumption
(10) $\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid rs_{post}\}$ indet
by Rule (24k) on (4)
and (9)

(11)
$$\operatorname{match}(e_1)\{\cdot\mid p_r\Rightarrow e_r\mid rs_{post}\}$$
 final by Rule (25b) on (10)

Case Scrutinee doesn't matche pattern.

(9)
$$e_1 \perp p_r$$
 by assumption
(10) $\operatorname{match}(e_1)\{\cdot \mid p_r \Rightarrow e_r \mid (r' \mid rs'_{post})\}$
 $\mapsto \operatorname{match}(e_1)\{(p_r \Rightarrow e_r \mid \cdot) \mid r' \mid rs'_{post}\}$ by Rule (33m) on (4)
and (9)

6 Decidability

 Ξ incon A finite set of constraints, Ξ , is inconsistent

$$\begin{array}{c} \text{CINCTruth} \\ \underline{\Xi \; \text{incon}} \\ \overline{\Xi, \top \; \text{incon}} \end{array} \tag{34a}$$

$$\Xi, \perp \mathtt{incon}$$
 (34b)

CINCNum

$$\frac{n_1 \neq n_2}{\Xi, \underline{n_1}, \underline{n_2} \text{ incon}} \tag{34c}$$

CINCNotNum

$$\frac{}{\Xi,\underline{n},\underline{\mathscr{K}}\mathtt{incon}}\tag{34d}$$

CINCAnd

$$\frac{\Xi, \xi_1, \xi_2 \text{ incon}}{\Xi, \xi_1 \wedge \xi_2 \text{ incon}}$$
 (34e)

CINCOr

$$\frac{\Xi, \xi_1 \text{ incon} \qquad \Xi, \xi_2 \text{ incon}}{\Xi, \xi_1 \vee \xi_2 \text{ incon}}$$
(34f)

CINCInj

$$\frac{\Xi, \operatorname{inl}(\xi_1), \operatorname{inr}(\xi_2) \operatorname{incon}}{\Xi, \operatorname{inl}(\xi_1), \operatorname{inr}(\xi_2) \operatorname{incon}}$$

CINCInl

$$\frac{\Xi \; {\tt incon}}{{\tt inl}(\Xi) \; {\tt incon}} \tag{34h}$$

CINCInr

$$\frac{\Xi \ {\tt incon}}{{\tt inr}(\Xi) \ {\tt incon}} \tag{34i}$$

CINCPairL
$$\frac{\Xi_1 \text{ incon}}{(\Xi_1, \Xi_2) \text{ incon}}$$
(34j)

CINCPairR
$$\frac{\Xi_2 \text{ incon}}{(\Xi_1, \Xi_2) \text{ incon}}$$
(34k)

Lemma 6.0.1 (Decidability of Inconsistency). Suppose $\dot{\top}(\xi) = \xi$. It is decidable whether ξ incon.

Lemma 6.0.2 (Inconsistency and Entailment of Constraint). Suppose that $\dot{\top}(\xi) = \xi$. Then $\bar{\xi}$ incon iff $\top \models \xi$

Lemma 6.0.3. If $e \models \xi$ then $e \models \dot{\top}(\xi)$

Proof. By rule induction over Rules (14), it is obvious to see that $\dot{T}(\xi) = \xi$. \Box

Lemma 6.0.4. *If* $e \models_? \xi$ *then* $e \models_?^{\dagger} \dot{\top}(\xi)$.

Proof.

(11)
$$e \models_? \xi$$

by assumption

By Rule Induction over Rules (16) on (11).

Case (16a).

(12) $\xi = ?$

by assumption

(13) $e \models \top$

by Rule (14a)

(14) $e \models_{?}^{\dagger} \top$

by Rule (17b) on (13)

Case (16b).

(12) e notintro

by assumption

(13) ξ refutable?

by assumption

By Lemma 2.0.15 on (12) and (13) and case analysis on its conclusion. By rule induction over Rules (10).

Case $\dot{\top}(\xi)$ refutable?.

(14) $\dot{\top}(\xi)$ refutable?

by assumption

(15) $e \models_? \dot{\top}(\xi)$

by Rule (16b) on (12)

and (14)

(16) $e \models_2^{\dagger} \dot{\top}(\xi)$

by Rule (17b) on (15)

Case $e \models \dot{\top}(\xi)$.

(14)
$$e \models \dot{\top}(\xi)$$

(15)
$$e \models_2^{\dagger} \top$$

by assumption

by Rule (17b) on (14)

Case (16c).

(12)
$$\xi = \xi_1 \vee \xi_2$$

(13) $e \models_? \xi_1$

(14) $e \models_?^\dagger \dot{\top}(\xi_1)$

(15) $e \models^{\dagger}_{?} \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$

by assumption

by assumption

by IH on (13)

by Lemma 2.0.10 on

(14)

Case (16d).

(12)
$$\xi = \xi_1 \vee \xi_2$$

(13) $e \models_? \xi_2$

(14) $e \models^{\dagger}_{?} \dot{\top}(\xi_2)$

(15) $e \models^{\dagger}_{?} \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$

by assumption

by assumption

by IH on (13)

by Lemma 2.0.10 on

(14)

Case (16e).

(12)
$$e = inl_{\tau_2}(e_1)$$

(13) $\xi = \operatorname{inl}(\xi_1)$

(14) $e_1 \models_? \xi_1$

(15) $e_1 \models_?^\dagger \dot{\top}(\xi_1)$

by assumption

by assumption

by assumption

by IH on (14)

by Lemma 2.0.11 on

(15)

Case (16f).

(12)
$$e = inr_{\tau_1}(e_2)$$

(13) $\xi = \operatorname{inr}(\xi_2)$

(14) $e_2 \models_? \xi_2$

(15) $e_2 \models^{\dagger}_{?} \dot{\top}(\xi_2)$

(16) $\operatorname{inr}_{\tau_1}(e_2) \models_{\tau}^{\dagger} \operatorname{inr}(\dot{\tau}(\xi_2))$

(16) $\operatorname{inl}_{\tau_2}(e_1) \models_{\tau}^{\dagger} \operatorname{inl}(\dot{\top}(\xi_1))$

by assumption

by assumption

by assumption

by IH on (14)

by Lemma 2.0.12 on

(15)

Case (16g).

$(12) \ e = (e_1, e_2)$	by assumption
$(13) \ \xi = (\xi_1, \xi_2)$	by assumption
$(14) e_1 \models_? \xi_1$	by assumption
$(15) e_2 \models \xi_2$	by assumption
$(16) e_1 \models^{\dagger}_? \dot{\top}(\xi_1)$	by IH on (14)
$(17) \ e_2 \models \dot{\top}(\xi_2)$	by Lemma 6.0.3 on (15)
$(18) \ e_2 \models^{\dagger}_? \dot{\top}(\xi_2)$	by Rule (17b) on (17)
(19) $(e_1, e_2) \models^{\dagger}_{?} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$	by Lemma 2.0.13 on (16) and (18)

Case (16h).

$(12) \ e = (e_1, e_2)$	by assumption
(13) $\xi = (\xi_1, \xi_2)$	by assumption
$(14) e_1 \models \xi_1$	by assumption
$(15) e_2 \models_? \xi_2$	by assumption
$(16) \ e_1 \models \dot{\top}(\xi_1)$	by Lemma $6.0.3$ on (14)
$(17) e_1 \models^{\dagger}_{?} \dot{\top}(\xi_1)$	by Rule (17b) on (16)
$(18) \ e_2 \models^{\dagger}_{?} \dot{\top}(\xi_2)$	by IH on (15)
(19) $(e_1, e_2) \models^{\dagger}_{?} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$	by Lemma 2.0.13 on (17) and (18)

Case (16i).

`	
$(12) \ e = (e_1, e_2)$	by assumption
$(13) \ \xi = (\xi_1, \xi_2)$	by assumption
$(14) \ e_1 \models_? \xi_1$	by assumption
(15) $e_2 \models_? \xi_2$	by assumption
$(16) \ e_1 \models^{\dagger}_{?} \dot{\top}(\xi_1)$	by IH on (14)
$(17) \ e_2 \models^{\dagger}_? \dot{\top}(\xi_2)$	by IH on (15)
(18) $(e_1, e_2) \models_{?}^{\dagger} (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$	by Lemma 2.0.13 on (16) and (17)

Lemma 6.0.5. $e \models^{\dagger}_{?} \xi \text{ iff } e \models^{\dagger}_{?} \dot{\top}(\xi)$

1. Sufficiency: Proof.

(1)
$$e \models_{?}^{\dagger} \xi$$

by assumption

By rule induction over Rules (17) on (1)

Case (17b).

 $\begin{array}{ll} (2) & e \models \xi \\ (3) & e \models \dot{\top}(\xi) \end{array}$

by assumption

by Lemma 6.0.3 on (2)

(4) $e \models_{?}^{\dagger} \dot{\top}(\xi)$

by Rule (17b) on (3)

Case (17a).

(2) $e \models_? \xi$

by assumption

(3) $e \models_{?}^{\dagger} \dot{\top}(\xi)$

by Lemma 6.0.4 on (2)

2. Necessity:

(1) $e \models_2^{\dagger} \dot{\top}(\xi)$

by assumption

By structural induction on ξ ,

Case $\xi = \top, \bot, \underline{n}, \varkappa$.

(2) $e \models_{?}^{\dagger} \xi$

by (1) and Definition 12

Case $\xi = ?$.

(2) $e \models_? ?$

by Rule (16a)

(3) $e \models_{2}^{\dagger} ?$

by Rule (17a) on (2)

Case $\xi = \xi_1 \vee \xi_2$.

(2) $\dot{\top}(\xi_1 \vee \xi_2) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$

by Definition 12

By rule induction over Rules (17) on (1),

Case (17b).

(3) $e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$

by assumption

By rule induction over Rules (14) on (3) and two cases apply, Case (14e).

(4) $e \models \dot{\top}(\xi_1)$

by assumption

(5) $e \models_2^{\dagger} \dot{\top}(\xi_1)$

by Rule (17b) on (4)

(6) $e \models_{?}^{\dagger} \xi_{1}$

- by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \vee \xi_2$
- by Lemma 2.0.10 on

Case (14f).

- (4) $e \models \dot{\top}(\xi_2)$
- by assumption
- (5) $e \models^{\dagger}_{?} \dot{\top}(\xi_2)$
- by Rule (17b) on (4)

(6) $e \models_{?}^{\dagger} \xi_{2}$

- by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \vee \xi_2$
- by Lemma 2.0.10 on

Case (17a).

- (3) $e \models_? \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$
- by assumption

By rule induction over Rules (16) on (3) and two cases apply, Case (16c).

- (4) $e \models_? \dot{\top}(\xi_1)$
- by assumption
- (5) $e \models_?^\dagger \dot{\top}(\xi_1)$
- by Rule (17a) on (4)

(6) $e \models_{?}^{\dagger} \xi_{1}$

- by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \vee \xi_2$
- by Lemma 2.0.10 on

Case (16d).

- (4) $e \models_? \dot{\top}(\xi_2)$
- by assumption
- (5) $e \models_?^\dagger \dot{\top}(\xi_2)$
- by Rule (17a) on (4)

(6) $e \models_{?}^{\dagger} \xi_{2}$

- by IH on (5)
- (7) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$
- by Lemma 2.0.10 on (6)

Case $\xi = inl(\xi_1)$.

(2) $e = inl_{\tau_2}(e_1)$

- by assumption
- (3) $\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$
- by assumption

By rule induction over Rules (17) on (1),

Case (17b).

- (4) $\operatorname{inl}_{\tau_2}(e_1) \models \operatorname{inl}(\dot{\top}(\xi_1))$
- by assumption

By rule induction over Rules (14) and only one case applies, Case (14g).

- (5) $e_1 \models \dot{\top}(\xi_1)$
- by assumption
- (6) $e_1 \models_?^\dagger \dot{\top}(\xi_1)$
- by Rule (17b) on (5)

(7) $e_1 \models_{?}^{\dagger} \xi_1$

by IH on (6)

(8)
$$\operatorname{inl}_{\tau_2}(e_1) \models_?^{\dagger} \operatorname{inl}(\xi_1)$$
 by Lemma 2.0.11 on (7)

Case (17a).

(4) $\operatorname{inl}_{\tau_2}(e_1) \models_? \operatorname{inl}(\dot{\top}(\xi_1))$ by assumption By rule induction over Rules (16) and only one case applies, Case (16e).

- (5) $e_1 \models_? \dot{\top}(\xi_1)$ by assumption (6) $e_1 \models_?^{\dagger} \dot{\top}(\xi_1)$ by Rule (17a) on (5) (7) $e_1 \models_?^{\dagger} \xi_1$ by IH on (6)
- (8) $\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$ by Lemma 2.0.11 on (7)

Case $\xi = inr(\xi_2)$.

(2) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption (3) $\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$ by assumption

By rule induction over Rules (17) on (1),

Case (17b).

(4) $\operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\top}(\xi_2))$ by assumption By rule induction over Rules (14) and only one case applies, Case (14h).

(5) $e_2 \models \dot{\top}(\xi_2)$ by assumption (6) $e_2 \models_{?}^{\dagger} \dot{\top}(\xi_2)$ by Rule (17b) on (5) (7) $e_2 \models_{?}^{\dagger} \xi_2$ by IH on (6) (8) $\inf_{\tau_1}(e_2) \models_{?}^{\dagger} \inf(\xi_2)$ by Lemma 2.0.12 on (7)

Case (17a).

(4) $\operatorname{inr}_{\tau_1}(e_2) \models_? \operatorname{inr}(\dot{\top}(\xi_2))$ by assumption By rule induction over Rules (16) and only one case applies, Case (16f).

(5) $e_{2} \models_{?} \dot{\top}(\xi_{2})$ by assumption (6) $e_{2} \models_{?}^{\dagger} \dot{\top}(\xi_{2})$ by Rule (17a) on (5) (7) $e_{2} \models_{?}^{\dagger} \xi_{2}$ by IH on (6) (8) $inr_{\tau_{1}}(e_{2}) \models_{?}^{\dagger} inr(\xi_{2})$ by Lemma 2.0.12 on (7)

Case $\xi = (\xi_1, \xi_2)$.

(2) $e = (e_1, e_2)$ by assumption (3) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \wedge \dot{\top}(\xi_2)$ by Definition 12

By rule induction over Rules (17) on (1),

Case (17b).

(4) $(e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by assumption

By rule induction over Rules (14) on (4) and only one case applies,

Case (14i).

- (5) $e_1 \models \dot{\top}(\xi_1)$ by assumption (6) $e_2 \models \dot{\top}(\xi_2)$ by assumption
- (7) $e_1 \models_7^{\uparrow} \dot{\uparrow}(\xi_1)$ by Rule (17b) on (5)
- (8) $e_2 \models_7^{\frac{1}{7}} \dot{\top}(\xi_2)$ by Rule (17b) on (6)
- (9) $e_1 \models_{?}^{\dagger} \xi_1$ by IH on (7)
- (10) $e_2 \models_7^{\dagger} \xi_2$ by IH on (8)
- (11) $(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$ by Lemma 2.0.13 on (9) and (10)

Case (17a).

(4) $(e_1, e_2) \models_? (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$ by assumption

By rule induction over Rules (16) on (4) and three cases apply, Case (16g).

- (5) $e_1 \models_? \dot{\top}(\xi_1)$ by assumption (6) $e_2 \models \dot{\top}(\xi_2)$ by assumption
- (6) $e_2 \models \top(\xi_2)$ by assumption (7) $e_1 \models_2^{\dagger} \dot{\top}(\xi_1)$ by Rule (17a) on (5)
- (8) $e_2 \models_7^{\dagger} \dot{\top}(\xi_2)$ by Rule (17b) on (6)
- (9) $e_1 \models_{?}^{\dagger} \xi_1$ by IH on (7)
- (10) $e_2 \models_7^+ \xi_2$ by IH on (8)
- (11) $(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$ by Lemma 2.0.13 on (9) and (10)

Case (16h).

- (5) $e_1 \models \dot{\top}(\xi_1)$ by assumption (6) $e_2 \models_? \dot{\top}(\xi_2)$ by assumption
- (7) $e_1 \models_2^{\dagger} \dot{\top}(\xi_1)$ by Rule (17b) on (5)
- (8) $e_2 \models^{\dagger}_? \dot{\top}(\xi_2)$ by Rule (17a) on (6)
- (9) $e_1 \models^{\dagger}_{?} \xi_1$ by IH on (7)
- (10) $e_2 \models_{?}^{\dagger} \xi_2$ by IH on (8)
- (11) $(e_1, e_2) \models^{\dagger}_{?} (\xi_1, \xi_2)$ by Lemma 2.0.13 on (9) and (10)

Case (16i).

(5) $e_1 \models_? \dot{\top}(\xi_1)$ by assumption (6) $e_2 \models_? \dot{\top}(\xi_2)$ by assumption

$$\begin{array}{lll} (7) & e_1 \models_?^\dagger \dot{\top}(\xi_1) & \text{by Rule (17a) on (5)} \\ (8) & e_2 \models_?^\dagger \dot{\top}(\xi_2) & \text{by Rule (17a) on (6)} \\ (9) & e_1 \models_?^\dagger \xi_1 & \text{by IH on (7)} \\ (10) & e_2 \models_?^\dagger \xi_2 & \text{by IH on (8)} \\ (11) & (e_1, e_2) \models_?^\dagger (\xi_1, \xi_2) & \text{by Lemma 2.0.13 on (9) and (10)} \\ \end{array}$$

Lemma 6.0.6. Assume $\dot{\top}(\xi) = \xi$. Then $\top \models_{?}^{\dagger} \xi$ iff $\top \models \xi$.

Proof. We prove sufficiency and necessity separately.

- 1. Sufficiency:
- 2. Necessity:

Theorem 6.1. $\top \models_{?}^{\dagger} \xi \text{ iff } \top \models \dot{\top}(\xi).$

Lemma 6.1.1. Assume that e val. Then $e \models^{\dagger}_? \xi$ iff $e \models \dot{\top}(\xi)$

Proof.

(1) e val by assumption

We prove sufficiency and necessity separately.

- 1. Sufficiency:
 - (2) $e \models^{\dagger}_{?} \xi$ by assumption

By rule induction over Rules (17) on (2).

Case (17b).

(3)
$$e \models \xi$$
 by assumption
 (4) $e \models \dot{\top}(\xi)$ by Lemma 6.0.3 on (3)

Case (17a).

(3)
$$e \models_? \xi$$
 by assumption

By rule induction over Rules (16) on (3).

Case (16a).

(4)
$$\xi = ?$$
 by assumption
(5) $e \models \dot{\top}(\xi)$ by Rule (14a) and Definition 12

Case (16b).

(4) e notintro

by assumption

By rule induction over Rules (26) on (4), for each case, by rule induction over Rules (23) on (1), no case applies due to syntactic contradiction.

Case (16c).

 $(4) \quad \xi = \xi_1 \vee \xi_2$

by assumption

(5) $e \models_{?} \xi_{1}$

- by assumption
- (6) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$
- by Equation 12

(7) $e \models_{?}^{\dagger} \xi_1$

by Rule (17a) on (5)

(8) $e \models \dot{\top}(\xi_1)$

- by IH on (7)
- (9) $e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$
- by Rule (14e) on (8)

Case (16d).

 $(4) \quad \xi = \xi_1 \vee \xi_2$

by assumption

(5) $e \models_{?} \xi_{2}$

- by assumption
- (6) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$
- by Equation 12

(7) $e \models^{\dagger}_{?} \xi_2$

by Rule (17a) on (5)

(8) $e \models \dot{\top}(\xi_2)$

- by IH on (7)
- (9) $e \models \dot{\top}(\xi_1) \lor \dot{\top}(\xi_2)$
- by Rule (14f) on (8)

Case (16e).

(4) $\xi = \operatorname{inl}(\xi_1)$

by assumption

(5) $e \models_{?} \xi_{1}$

- by assumption
- (6) $\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$
- by Equation 12

(7) $e \models_{?}^{\dagger} \xi_{1}$

by Rule (17a) on (5)

(8) $e \models \dot{\top}(\xi_1)$

- by IH on (7)
- (9) $e \models \operatorname{inl}(\dot{\top}(\xi_1))$
- by Rule (14g) on (8)

Case (16f).

(4) $\xi = \operatorname{inr}(\xi_2)$

by assumption

(5) $e \models_{?} \xi_{2}$

- by assumption
- (6) $\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$
- by Equation 12

(7) $e \models_{?}^{\dagger} \xi_2$

by Rule (17a) on (5)

(8) $e \models \dot{\top}(\xi_2)$

- by IH on (7)
- (9) $e \models \operatorname{inr}(\dot{\top}(\xi_2))$
- by Rule (14h) on (8)

Case (16g).

(4) $e = (e_1, e_2)$

by assumption

(5) $\xi = (\xi_1, \xi_2)$

by assumption

$$\begin{array}{lll} (6) & e_1 \models_? \xi_1 & \text{by assumption} \\ (7) & e_2 \models \xi_2 & \text{by assumption} \\ (8) & \dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2)) & \text{by Equation 12} \\ (9) & e_1 \models_?^{\dagger} \xi_1 & \text{by Rule (17a) on (6)} \\ (10) & e_2 \models_?^{\dagger} \xi_2 & \text{by Rule (17b) on (7)} \\ (11) & e_1 \models \dot{\top}(\xi_1) & \text{by IH on (9)} \\ (12) & e_2 \models \dot{\top}(\xi_2) & \text{by IH on (10)} \\ (13) & (e_1, e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2)) & \text{by Rule (14i) on (11)} \\ & & \text{and (12)} \\ \end{array}$$

Case (16h).

(4)	$e = (e_1, e_2)$	by assumption
(5)	$\xi = (\xi_1, \xi_2)$	by assumption
(6)	$e_1 \models \xi_1$	by assumption
(7)	$e_2 \models_? \xi_2$	by assumption
(8)	$\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$	by Equation 12
(9)	$e_1 \models^{\dagger}_{?} \xi_1$	by Rule (17b) on (6)
(10)	$e_2 \models^{\dagger}_{?} \xi_2$	by Rule $(17a)$ on (7)
(11)	$e_1 \models \dot{\top}(\xi_1)$	by IH on (9)
(12)	$e_2 \models \dot{\top}(\xi_2)$	by IH on (10)
(13)	$(e_1,e_2)\models(\dot{\top}(\xi_1),\dot{\top}(\xi_2))$	by Rule (14i) on (11)
		and (12)

Case (16i).

(101)•		
(4) e =	(e_1,e_2)	by assumption
(5) $\xi =$	(ξ_1,ξ_2)	by assumption
(6) $e_1 \models$	$=$? ξ_1	by assumption
$(7) e_2 \models$	$=$? ξ_2	by assumption
(8) $\dot{\top}(\xi)$	$=(\dot{\top}(\xi_1),\dot{\top}(\xi_2))$	by Equation 12
(9) $e_1 \models$	$= \stackrel{\dagger}{?} \xi_1$	by Rule (17a) on (6)
$(10) e_2 \models$	$=\stackrel{\dagger}{?}\xi_2$	by Rule (17a) on (7)
(11) $e_1 \models$	$=\dot{\top}(\xi_1)$	by IH on (9)
$(12) e_2 \models$	$=\dot{ op}(\xi_2)$	by IH on (10)
(13) $(e_1,$	$e_2) \models (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$	by Rule (14i) on (11)
		and (12)

2. Necessity:

(2)
$$e \models \dot{\top}(\xi)$$
 by assumption

By structural induction on ξ .

Case $\xi = \top, \bot, \underline{n}, \underline{\mathscr{n}}$.

(3) $\xi = \dot{\top}(\xi)$

by Equation 12

(4) $e \models_{?}^{\dagger} \xi$

by Rule (17b) on (2)

Case $\xi = ?$.

(3) $e \models_? ?$

by Rule (16a)

(4) $e \models_{?}^{\dagger} ?$

by Rule (17a) on (3)

Case $\xi = \xi_1 \wedge \xi_2$.

- (3) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \wedge \dot{\top}(\xi_2)$
- by Equation 12

By rule induction over Rules (14) on (2), only one case applies.

Case (14d).

(4) $e \models \dot{\top}(\xi_1)$

by assumption

(5) $e \models \dot{\top}(\xi_2)$

by assumption

(6) $e \models^{\dagger}_{?} \xi_1$

by IH on (4)

(7) $e \models^{\dagger}_{?} \xi_{2}$ (8) $e \models \xi_{1} \land \xi_{2}$

- by IH on (5) by Lemma 2.0.9 on (6)
- and (7)

Case $\xi = \xi_1 \vee \xi_2$.

- (3) $\dot{\top}(\xi) = \dot{\top}(\xi_1) \vee \dot{\top}(\xi_2)$
- by Equation 12

By rule induction over Rules (14) on (2) and only two cases apply. Case (14e).

(4) $e \models \dot{\top}(\xi_1)$

by assumption

(5) $e \models_{?}^{\dagger} \xi_1$

by IH on (4)

 $(6) \ e \models^{\dagger}_{?} \xi_1 \lor \xi_2$

by Lemma 2.0.10 on (5)

Case (14f).

(4) $e \models \dot{\top}(\xi_2)$

by assumption

(5) $e \models^{\dagger}_{?} \xi_{2}$

by IH on (4)

(6) $e \models_{?}^{\dagger} \xi_1 \lor \xi_2$

by Lemma 2.0.10 on

(5)

Case $\xi = inl(\xi_1)$.

- (3) $\dot{\top}(\xi) = \operatorname{inl}(\dot{\top}(\xi_1))$
- by Equation 12

By rule induction over Rules (14) on (2) and only one case applies. Case (14g).

(4)
$$e = \operatorname{inl}_{\tau_2}(e_1)$$
 by assumption
(5) $e_1 \models \dot{\top}(\xi_1)$ by assumption

(6)
$$e_1 \models_{?}^{\dagger} \xi_1$$
 by IH on (5)

(7)
$$\operatorname{inl}_{\tau_2}(e_1) \models_{?}^{\dagger} \operatorname{inl}(\xi_1)$$
 by Lemma 2.0.11 on (6)

Case $\xi = inr(\xi_2)$.

(3)
$$\dot{\top}(\xi) = \operatorname{inr}(\dot{\top}(\xi_2))$$
 by Equation 12

By rule induction over Rules (14) on (2) and only one case applies. Case (14h).

(4)
$$e = \operatorname{inr}_{\tau_1}(e_2)$$
 by assumption
(5) $e_2 \models \dot{\top}(\xi_2)$ by assumption
(6) $e_2 \models_{7}^{\dagger} \xi_2$ by IH on (5)

(7)
$$\operatorname{inr}_{\tau_1}(e_2) \models_{?}^{\dagger} \operatorname{inr}(\xi_2)$$
 by Lemma 2.0.12 on (6)

Case $\xi = (\xi_1, \xi_2)$.

(3)
$$\dot{\top}(\xi) = (\dot{\top}(\xi_1), \dot{\top}(\xi_2))$$
 by Equation 12

By rule induction over Rules (14) on (2) and only one case applies.

$$(4) \ e = (e_1, e_2)$$
 by assumption
$$(5) \ e_1 \models \dot{\bot}(\xi_1)$$
 by assumption
$$(6) \ e_2 \models \dot{\bot}(\xi_2)$$
 by assumption
$$(7) \ e_1 \models_{?}^{\dagger} \xi_1$$
 by IH on (5)

(8)
$$e_2 \models^{\dagger}_{?} \xi_2$$
 by IH on (6)

(9)
$$(e_1, e_2) \models_?^{\dagger} (\xi_1, \xi_2)$$
 by Lemma 2.0.13 on (7) and (8)

Lemma 6.1.2. $e \models \xi \text{ iff } e \models \dot{\bot}(\xi)$

Proof. We prove sufficiency and necessity separately.

1. Sufficiency:

(1)
$$e \models \xi$$
 by assumption

By rule induction over Rules (14) on (1).

Case (14a).

(2)
$$\xi = \top$$
 by assumption
(3) $e \models \dot{\bot}(\top)$ by (1) and Definition
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Case (14b).

(2) $\xi = \underline{n}$ (3) $e \models \dot{\perp}(\underline{n})$ by assumption by (1) and Definition 13

Case (14c).

(2) $\xi = \underline{\varkappa}$ (3) $e \models \dot{\bot}(\varkappa)$ by assumption by (1) and Definition 13

Case (14d).

- $(2) \ \xi = \xi_1 \wedge \xi_2$
- (3) $e \models \xi_1$
- (4) $e \models \xi_2$
- (5) $e \models \dot{\perp}(\xi_1)$
- (6) $e \models \dot{\perp}(\xi_2)$
- (7) $e \models \dot{\perp}(\xi_1) \land \dot{\perp}(\xi_2)$
- (8) $e \models \dot{\perp}(\xi_1 \land \xi_2)$

- by assumption
- by assumption
- by assumption
- by IH on (3)
- by IH on (4)
- by Rule (14d) on (5)
- and (6)
- by (7) and Definition 13

Case (14e).

- (2) $\xi = \xi_1 \vee \xi_2$
- (3) $e \models \xi_1$
- (4) $e \models \dot{\perp}(\xi_1)$
- (5) $e \models \dot{\perp}(\xi_1) \lor \dot{\perp}(\xi_2)$
- (6) $e \models \dot{\perp}(\xi_1 \vee \xi_2)$

- by assumption
- by assumption
- by IH on (3)
- by Rule (14e) on (4)
- by (5) and Definition 13

Case (14f).

- (2) $\xi = \xi_1 \vee \xi_2$
- (3) $e \models \xi_2$
- (4) $e \models \dot{\perp}(\xi_2)$
- (5) $e \models \dot{\perp}(\xi_1) \lor \dot{\perp}(\xi_2)$
- (6) $e \models \dot{\perp}(\xi_1 \vee \xi_2)$

- by assumption
- by assumption
- by IH on (3)
- by Rule (14f) on (4)
- by (5) and Definition 13

Case (14g).

- $(2) e = \operatorname{inl}_{\tau_2}(e_1)$
- (3) $\xi = \operatorname{inl}(\xi_1)$
- (4) $e_1 \models \xi_1$

- by assumption
- by assumption
- by assumption

 $\begin{array}{lll} (5) & e_1 \models \dot{\bot}(\xi_1) & & \text{by IH on (4)} \\ (6) & \inf_{\tau_2}(e_1) \models \inf(\dot{\bot}(\xi_1)) & & \text{by Rule (14g) on (5)} \\ (7) & \inf_{\tau_2}(e_1) \models \dot{\bot}(\inf(\xi_1)) & & \text{by (6) and Definition} \\ \end{array}$

13

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Case (14h).

 $\begin{array}{lll} (2) & e = \operatorname{inr}_{\tau_1}(e_2) & \text{by assumption} \\ (3) & \xi = \operatorname{inr}(\xi_2) & \text{by assumption} \\ (4) & e_2 \models \xi_2 & \text{by assumption} \\ (5) & e_2 \models \dot{\bot}(\xi_2) & \text{by IH on (4)} \\ (6) & \operatorname{inr}_{\tau_1}(e_2) \models \operatorname{inr}(\dot{\bot}(\xi_2)) & \text{by Rule (14h) on (5)} \\ (7) & \operatorname{inr}_{\tau_1}(e_2) \models \dot{\bot}(\operatorname{inr}(\xi_2)) & \text{by (6) and Definition} \\ & & 13 \end{array}$

Case (14i).

(2) $e = (e_1, e_2)$ by assumption (3) $\xi = (\xi_1, \xi_2)$ by assumption (4) $e_1 \models \xi_1$ by assumption (5) $e_2 \models \xi_2$ by assumption (6) $e_1 \models \dot{\bot}(\xi_1)$ by IH on (4) (7) $e_2 \models \dot{\perp}(\xi_2)$ by IH on (5)(8) $(e_1, e_2) \models (\dot{\bot}(\xi_1), \dot{\bot}(\xi_2))$ by Rule (14i) on (6) and (7)(9) $(e_1, e_2) \models \dot{\bot}((\xi_1, \xi_2))$ by (8) and Definition

2. Necessity:

(1) $e \models \dot{\perp}(\xi)$ by assumption

By structural induction on ξ .

Case
$$\xi = \top, \bot, \underline{n}, \underline{\mathscr{N}}$$
.

(2) $e \models \xi$ by (1) and Definition 13

Case
$$\xi = ?$$
.

- (2) $e \models \bot$ by (1) and Definition 13 by Lemma 2.0.1
- (3) contradicts (2).

Case
$$\xi = \xi_1 \wedge \xi_2$$
.

(2)
$$e \models \dot{\bot}(\xi_1) \land \dot{\bot}(\xi_2)$$
 by (1) and Definition 13

By rule induction over Rules (14) on (2) and only case applies.

Case (14d).

- (3) $e \models \dot{\perp}(\xi_1)$ by assumption (4) $e \models \dot{\perp}(\xi_2)$ by assumption
- (5) $e \models \xi_1$ by IH on (3)
- (6) $e \models \xi_2$ by IH on (4)
- (7) $e \models \xi_1 \land \xi_2$ by Rule (14d) on (5) and (6)

Case $\xi = \xi_1 \vee \xi_2$.

(2)
$$e \models \dot{\bot}(\xi_1) \lor \dot{\bot}(\xi_2)$$
 by (1) and Definition 13

By rule induction over Rules (14) on (2) and only two cases apply. Case (14e).

- (3) $e \models \dot{\bot}(\xi_1)$ by assumption (4) $e \models \xi_1$ by IH on (3)
- (5) $e \models \xi_1 \lor \xi_2$ by Rule (14e) on (4)

Case (14f).

- (3) $e \models \dot{\bot}(\xi_2)$ by assumption (4) $e \models \xi_2$ by IH on (3)
- (5) $e \models \xi_1 \vee \xi_2$ by Rule (14f) on (4)

Case $\xi = inl(\xi_1)$.

(2)
$$e \models \mathtt{inl}(\dot{\bot}(\xi_1))$$
 by (1) and Definition 13

By rule induction over Rules (14) on (2) and only one case applies.

Case (14g).

- $\begin{array}{ll} (3) & e=\mathtt{inl}_{\tau_2}(e_1) & \text{by assumption} \\ (4) & e_1 \models \dot{\bot}(\xi_1) & \text{by assumption} \end{array}$
- (5) $e_1 \models \xi_1$ by IH on (4)
- (6) $e \models \operatorname{inl}(\xi_1)$ by Rule (14g) on (5)

Case $\xi = inr(\xi_2)$.

(2)
$$e \models \operatorname{inr}(\dot{\bot}(\xi_2))$$
 by (1) and Definition

By rule induction over Rules (14) on (2) and only one case applies. Case (14h).

(3) $e = \operatorname{inr}_{\tau_1}(e_2)$ by assumption (4) $e_2 \models \dot{\bot}(\xi_2)$ by assumption

$$(5) \ e_2 \models \xi_2 \qquad \qquad \text{by IH on (4)}$$

$$(6) \ e \models \mathsf{inr}(\xi_2) \qquad \qquad \text{by Rule (14h) on (5)}$$

$$\mathbf{Case} \ \xi = (\xi_1, \xi_2).$$

$$(2) \ e \models (\dot{\bot}(\xi_1), \dot{\bot}(\xi_2)) \qquad \qquad \text{by (1) and Definition}$$

13

By rule induction over Rules (14) on (2) and only case applies. Case (14i).

$$\begin{array}{lll} (3) & e = (e_1, e_2) & & \text{by assumption} \\ (4) & e_1 \models \dot{\bot}(\xi_1) & & \text{by assumption} \\ (5) & e_2 \models \dot{\bot}(\xi_2) & & \text{by assumption} \\ (6) & e_1 \models \xi_1 & & \text{by IH on (4)} \\ (7) & e_2 \models \xi_2 & & \text{by IH on (5)} \\ (8) & e \models (\xi_1, \xi_2) & & \text{by Rule (14i) on (6)} \\ \end{array}$$

Lemma 6.1.3. Assume e val and $\dot{\top}(\xi) = \xi$. Then $e \not\models \xi$ iff $e \models \overline{\xi}$.

Theorem 6.2. $\xi_r \models \xi_{rs} \text{ iff } \top \models \overline{\dot{\top}(\xi_r)} \lor \dot{\bot}(\xi_{rs}).$