

# Finding All $q$ -Hypergeometric Solutions of $q$ -Difference Equations

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## Abstract

We present an algorithm for finding all solutions  $y(x)$  of a linear homogeneous  $q$ -difference equation such that  $y(qx)/y(x)$  is a rational function of  $q$  and  $x$ . The algorithm can also be used to construct  $q$ -hypergeometric series solutions of  $q$ -difference equations.

## Résumé

Nous présentons un algorithme qui trouve toutes les solutions  $y(x)$  des équations linéaires homogènes aux  $q$ -différences, telles que  $y(qx)/y(x)$  est une fonction rationnelle de  $q$  et de  $x$ . On peut utiliser cet algorithme aussi pour construire les solutions des équations aux  $q$ -différences ayant la forme d'une série  $q$ -hypergéométrique.

## 1 Introduction

Let  $\mathbb{Q}$  be the rational number field,  $q$  transcendental over  $\mathbb{Q}$ ,  $K$  a computable extension of  $\mathbb{Q}(q)$ , and  $x$  transcendental over  $K$ . Denote by  $Q$  the unique automorphism of  $K(x)$  which fixes  $K$  and satisfies  $Qx = qx$ . Then  $K(x)$  together with  $Q$  is an inversive difference field.

Let  $M$  be a difference extension ring of  $K(x)$ . An element  $a \in M$  is  $q$ -polynomial if  $a \in K[x]$ , and  $q$ -rational if  $a \in K(x)$ . An element  $a \in M \setminus \{0\}$  is a  $q$ -hypergeometric term if  $Qa = ra$  for some  $r \in K(x)$ . All these concepts are relative to the field  $K$ .

We are interested in  $q$ -hypergeometric solutions  $y$  of  $Ly = 0$  where

$$L = \sum_{i=0}^{\rho} p_i Q^i$$

is a linear  $q$ -difference operator of order  $\rho$  with coefficients  $p_i \in K(x)$ , with  $p_\rho, p_0 \neq 0$ . By clearing denominators in  $Ly = 0$  we can restrict our attention to operators  $L$  with  $p_i \in K[x]$ . An algorithm for this problem is presented in Section 4. It is a  $q$ -analogue of the algorithm

for finding hypergeometric solutions of difference equations described in [6]. In preparation, we show how to find  $q$ -polynomial solutions of  $Ly = 0$  in Section 2, and give a normal form for  $q$ -rational functions in Section 3. Finally, in Section 5, we describe solution of various related problems such as solving nonhomogeneous equations, finding solutions in the form of  $q$ -hypergeometric series, and deriving  $q$ -hypergeometric identities.

We use  $\mathbb{N}$  to denote the set of nonnegative integers. By  $(a; q)_n$  we denote the expression  $(1 - a)(1 - aq) \cdots (1 - aq^{n-1})$ .

In our examples we use two algebraic settings which are special cases of the general framework described above. In one we work with sequences of elements of  $K$ , identifying sequences which agree from some point on. More precisely, we take  $M = K^{\mathbb{N}}/J$  where  $K^{\mathbb{N}}$  is the ring of sequences over  $K$ , and  $J$  is the ideal of sequences with finitely many nonzero terms. In particular, all equalities among sequences (of the form  $a_n = b_n$ ) are meant to hold for all but finitely many  $n \in \mathbb{N}$ . Further we take  $x = (q^n)_{n=0}^{\infty} + J$  and define  $Q$  as the unique automorphism of  $M$  satisfying  $Q(a + J) = Ea + J$  for all  $a \in K^{\mathbb{N}}$ . Here  $E$  denotes the shift operator acting on  $K^{\mathbb{N}}$  by  $Ea_n = a_{n+1}$ . Obviously  $K$  can be embedded in  $M$  as the subring of constant sequences. To simplify notation, we will henceforth identify  $a + J \in K^{\mathbb{N}}/J$  with its representative  $a \in K^{\mathbb{N}}$ . Note that in this context a sequence  $a_n$  is  $q$ -polynomial if  $a_n = p(q^n)$  for some  $p \in K[x]$ ,  $q$ -rational if  $a_n = r(q^n)$  for some  $r \in K(x)$ , and a  $q$ -hypergeometric term if  $a_{n+1} = r(q^n)a_n$  for some  $r \in K(x)$ .

In another setting we take  $M = K[[x]]$  (or  $M = K((x))$ ), the ring of formal power series (resp. the field of formal Laurent series) over  $K$ . Again,  $K$ ,  $K[x]$ , and  $K(x)$  are embedded in  $M$  in a natural way. We distinguish between series that are  $q$ -hypergeometric terms, and series whose coefficients form a  $q$ -hypergeometric sequence. More precisely, a series  $f(x) = \sum_{j=0}^{\infty} \alpha_j x^j$  is a  $q$ -hypergeometric term if  $f(qx) = r(x)f(x)$  for some  $r(x) \in K(x)$ , and a  $q$ -hypergeometric series if  $\alpha_{j+1} = r(q^j)\alpha_j$  for some  $r(x) \in K(x)$  and for all large enough  $j \in \mathbb{N}$ .

Several times we will need to find the largest  $n \in \mathbb{N}$  (if any) such that  $q^n$  is a root of a given polynomial with coefficients in  $K$ . Therefore we assume that  $K$  is a *q-suitable field*, meaning that there exists an algorithm which given  $p \in K[x]$  finds all  $n \in \mathbb{N}$  such that  $p(q^n) = 0$ . For instance, if  $K = k(q)$  where  $q$  is transcendental over  $k$  we can proceed as follows: Let  $p(x) = \sum_{i=0}^d c_i x^i$  where  $c_i \in k[q]$ . Compute  $s = \min\{i; c_i \neq 0\}$  and  $t = \max\{j; q^j | c_s\}$ . Then  $p(q^n) = 0$  only if  $n \leq t$ , and the set of all such  $n$  can be found by testing the values  $n = t, t-1, \dots, 0$ .

## 2 $q$ -polynomial solutions

First we show how to find solutions  $y \in K[x]$  of  $Ly = 0$ . Let  $p_i = \sum_{k=0}^d c_{ik} x^k$  where  $c_{ik} \in k[q]$  and not all  $c_{id}$  are zero. Assume that  $y = \sum_{j=0}^N \alpha_j x^j$  where  $\alpha_N \neq 0$ . Substituting these expressions into  $Ly = 0$  and replacing  $k$  by  $l = j + k$  yields

$$\sum_{i,l,j} c_{i,l-j} \alpha_j q^{ij} x^l = 0$$

which implies that

$$\sum_{j=\max\{l-d,0\}}^{\min\{l,N\}} \sum_{i=0}^{\rho} c_{i,l-j} \alpha_j q^{ij} = 0, \quad \text{for } 0 \leq l \leq N+d. \quad (1)$$

In particular, for  $l = N + d$ ,

$$\sum_{i=0}^{\rho} c_{id} q^{iN} = 0, \quad (2)$$

and for  $l = 0$ ,

$$\alpha_0 \sum_{i=0}^{\rho} c_{i0} = 0. \quad (3)$$

From (2) it follows that  $q^N$  is a root of the polynomial  $P(x) = \sum_{i=0}^{\rho} c_{id} x^i$ . Let  $N_0$  be the largest  $n \in \mathbb{N}$  such that  $P(q^n) = 0$  (see the last paragraph of Introduction). All  $q$ -polynomial solutions  $y$  of  $Ly = 0$  can now be found by the method of undetermined coefficients. Ultimately, the problem is reduced to a system of linear algebraic equations over  $K$  with  $N_0 + 1$  unknowns. – A more efficient method leading to a system with at most  $\min\{2d, N_0 + 1\}$  unknowns is described in [2].

### 3 A normal form for $q$ -rational functions

**Theorem 1** *Let  $r \in K(x) \setminus \{0\}$ . Then there are  $z \in K$  and monic polynomials  $a, b, c \in K[x]$  such that*

$$r(x) = z \frac{a(x)}{b(x)} \frac{c(qx)}{c(x)}, \quad (4)$$

$$\gcd(a(x), b(q^n x)) = 1 \quad \text{for all } n \in \mathbb{N}, \quad (5)$$

$$\gcd(a(x), c(x)) = 1, \quad (6)$$

$$\gcd(b(x), c(qx)) = 1, \quad (7)$$

$$c(0) \neq 0. \quad (8)$$

*Proof:* Write  $r(x) = \frac{f(x)}{g(x)}$  where  $f, g$  are relatively prime polynomials. We start by finding the set  $\mathcal{S}$  of all  $n \in \mathbb{N}$  such that  $f(x)$  and  $g(q^n x)$  have a nonconstant common factor. To this end consider the polynomial  $R(h) = \text{Resultant}_x(f(x), g(hx))$ . By the well-known properties of polynomial resultants,  $\mathcal{S} = \{n \in \mathbb{N}; R(q^n) = 0\}$ .

Assume that  $\mathcal{S} = \{n_1, n_2, \dots, n_t\}$  where  $t \geq 0$  and  $n_1 < n_2 < \dots < n_t$ . In addition, let  $n_{t+1} = +\infty$ . Define polynomials  $f_i$  and  $g_i$  inductively by setting

$$f_0(x) = f(x), \quad g_0(x) = g(x),$$

and for  $i = 1, 2, \dots, t$ ,

$$\begin{aligned} s_i(x) &= \gcd(f_{i-1}(x), g_{i-1}(q^{n_i} x)), \\ f_i(x) &= f_{i-1}(x)/s_i(x), \\ g_i(x) &= g_{i-1}(x)/s_i(q^{-n_i} x). \end{aligned}$$

Now take

$$\begin{aligned} z &= \alpha/\beta, \\ a(x) &= f_t(x)/\alpha, \\ b(x) &= g_t(x)/\beta, \\ c(x) &= \prod_{i=1}^t \prod_{j=1}^{n_i} s_i(q^{-j}x), \end{aligned}$$

where  $\alpha$  and  $\beta$  denote the leading coefficients of  $f_t(x)$  and  $g_t(x)$ , respectively. Before proving (4) – (8) we state a lemma.

**Lemma 1** *Let  $n \in \mathbb{N}$ . If  $0 \leq l \leq i, j \leq t$  and  $n < n_{l+1}$ , then  $\gcd(f_i(x), g_j(q^n x)) = 1$ .*

*Proof:* Assume first that  $n \notin \mathcal{S}$ . Then  $R(q^n) \neq 0$ , hence  $\gcd(f(x), g(q^n x)) = 1$ . Since  $f_i(x) | f(x)$  and  $g_j(x) | g(x)$  it follows that  $\gcd(f_i(x), g_j(q^n x)) = 1$ , too.

To prove the lemma for  $n \in \mathcal{S}$  we use induction on  $l$ .

$l = 0$ : In this case there is nothing to prove since there is no  $n \in \mathcal{S}$  such that  $n < n_1$ .

$l > 0$ : Assume that the lemma holds for all  $n < n_l$ . It remains to show that it also holds for  $n = n_l$ . Since  $f_i(x) | f_l(x)$  and  $g_j(x) | g_l(x)$  it follows that  $\gcd(f_i(x), g_j(q^{n_l} x))$  divides  $\gcd(f_l(x), g_l(q^{n_l} x)) = \gcd(f_{l-1}(x)/s_l(x), g_{l-1}(q^{n_l} x)/s_l(x))$ . By the definition of  $s_l(x)$ , the latter gcd is 1, completing the proof.  $\square$

Now we proceed to verify properties (4) – (8).

(4):

$$\begin{aligned} z \frac{a(x)}{b(x)} \frac{c(qx)}{c(x)} &= \frac{f_t(x)}{g_t(x)} \prod_{i=1}^t \prod_{j=1}^{n_i} \frac{s_i(q^{1-j}x)}{s_i(q^{-j}x)} \\ &= \frac{f_0(x)}{\prod_{i=1}^t s_i(x)} \frac{\prod_{i=1}^t s_i(q^{-n_i}x)}{g_0(x)} \prod_{i=1}^t \frac{s_i(x)}{s_i(q^{-n_i}x)} = \frac{f(x)}{g(x)} = r(x). \end{aligned}$$

(5): Let  $i = j = l = t$  in Lemma 1. Then  $\gcd(f_t(x), g_t(q^n x)) = 1$  for all  $n < n_{t+1} = +\infty$ . In other words,  $\gcd(a(x), b(q^n x)) = 1$  for all  $n \in \mathbb{N}$ .

(6): If  $a(x)$  and  $c(x)$  have a non-constant common factor then so do  $f_t(x)$  and  $s_i(q^{-j}x)$ , for some  $i$  and  $j$  such that  $1 \leq i \leq t$  and  $1 \leq j \leq n_i$ . Since  $g_{i-1}(q^{n_i-j}x) = g_i(q^{n_i-j}x)s_i(q^{-j}x)$ , it follows that  $g_{i-1}(q^{n_i-j}x)$  contains this factor as well. As  $n_i - j < n_i$ , this contradicts Lemma 1. Hence  $a(x)$  and  $c(x)$  are relatively prime.

(7): If  $b(x)$  and  $c(qx)$  have a non-constant common factor then so do  $g_t(x)$  and  $s_i(q^{-j}x)$ , for some  $i$  and  $j$  such that  $1 \leq i \leq t$  and  $1 \leq j+1 \leq n_i$ . Since  $f_{i-1}(q^{-j}x) = f_i(q^{-j}x)s_i(q^{-j}x)$ , it follows that  $f_{i-1}(x)$  and  $g_t(q^j x)$  contain this factor as well. As  $j < n_i$ , this contradicts Lemma 1. Hence  $b(x)$  and  $c(qx)$  are relatively prime.

(8): It is easy to see that  $s_i(x)$  divides both  $f(x)$  and  $g(q^{n_i} x)$ . Hence  $s_i(0) = 0$  would imply that  $f(0) = g(0) = 0$ , contrary to the assumption that  $f$  and  $g$  are relatively prime. It follows that  $s_i(0) \neq 0$  for all  $i$ , and consequently  $c(0) \neq 0$ .  $\square$

**Theorem 2** Let  $a, b, c, A, B, C \in K[x]$  be polynomials such that  $c(0) \neq 0$  and  $\gcd(a(x), c(x)) = \gcd(b(x), c(qx)) = \gcd(A(x), B(q^n x)) = 1$ , for all  $n \in \mathbb{N}$ . If

$$\frac{a(x)}{b(x)} \frac{c(qx)}{c(x)} = \frac{A(x)}{B(x)} \frac{C(qx)}{C(x)}, \quad (9)$$

then  $c(x)$  divides  $C(x)$ .

*Proof:* Let

$$\begin{aligned} g(x) &= \gcd(c(x), C(x)), \\ d(x) &= c(x)/g(x), \\ D(x) &= C(x)/g(x). \end{aligned}$$

Then  $\gcd(d(x), D(x)) = \gcd(a(x), d(x)) = \gcd(b(x), d(qx)) = 1$  and  $d(0) \neq 0$ . Clear denominators in (9) and cancel  $g(x)g(qx)$  on both sides. The result  $A(x)b(x)d(x)D(qx) = a(x)B(x)D(x)d(qx)$  shows that

$$\begin{aligned} d(x) &\mid B(x)d(qx), \\ d(qx) &\mid A(x)d(x). \end{aligned}$$

Using these two relations repeatedly we find that

$$\begin{aligned} d(x) &\mid B(x)B(qx)\cdots B(q^{n-1}x)d(q^n x), \\ d(x) &\mid A(q^{-1}x)A(q^{-2}x)\cdots A(q^{-n}x)d(q^{-n} x), \end{aligned}$$

for all  $n \in \mathbb{N}$ . It is easy to see that since  $d(0) \neq 0$  and  $q$  is not a root of unity,  $d(x)$  and  $d(q^n x)$  are relatively prime for all large enough  $n$ . It follows that  $d(x)$  divides both  $B(x)B(qx)\cdots B(q^{n-1}x)$  and  $A(q^{-1}x)A(q^{-2}x)\cdots A(q^{-n}x)$  for all large enough  $n$ . But these polynomials are relatively prime by assumption, so  $d(x)$  is constant. Hence  $c(x) \mid C(x)$ .  $\square$

**Corollary 1** The factorization of  $r(x)$  described in Theorem 1 is unique.

*Proof:* If

$$z \frac{a(x)}{b(x)} \frac{c(qx)}{c(x)} = Z \frac{A(x)}{B(x)} \frac{C(qx)}{C(x)}$$

are two such factorizations then  $c(x) \mid C(x)$  and  $C(x) \mid c(x)$ , by Theorem 2. Since these polynomials are monic,  $c = C$ . It follows that  $z = Z$  and  $aB = Ab$ . Hence  $a \mid A$  and  $A \mid a$ , so  $a = A$  and  $b = B$ .  $\square$

**Corollary 2** Among all factorizations of  $r(x)$  satisfying (4) and (5) of Theorem 1, the one satisfying (4) – (8) has  $c(x)$  of least degree.

## 4 $q$ -hypergeometric solutions

After this preparation we turn to the algorithm for finding  $q$ -hypergeometric solutions  $y$  of  $Ly = 0$ . Let  $Qy = ry$  where  $r \in K(x)$ , then  $Q^i y = \prod_{j=0}^{i-1} r(q^j x) y$ . We look for  $r(x)$  in the normal form described in Theorem 1. After inserting (4) into  $Ly = 0$ , clearing denominators and cancelling  $y$  we obtain

$$\sum_{i=0}^{\rho} z^i f_i(x) c(q^i x) = 0 \quad (10)$$

where

$$f_i(x) = p_i(x) \prod_{j=0}^{i-1} a(q^j x) \prod_{j=i}^{\rho-1} b(q^j x).$$

Since all terms in (10) except for  $i = 0$  are divisible by  $a(x)$  it follows that  $a(x)$  divides  $p_0(x) \prod_{j=0}^{\rho-1} b(q^j x) c(x)$ . Because of (5) and (6),  $a(x)$  divides  $p_0(x)$ . Similarly, all terms in (10) except for  $i = \rho$  are divisible by  $b(q^{\rho-1} x)$ , therefore  $b(q^{\rho-1} x)$  divides  $z^\rho p_\rho(x) \prod_{j=0}^{\rho-1} a(q^j x) c(q^\rho x)$ . Because of (5) and (7),  $b(q^{\rho-1} x)$  divides  $p_\rho(x)$ . Thus we have a finite choice for  $a(x)$  and  $b(x)$ .

For each choice of  $a(x)$  and  $b(x)$ , equation (10) is a  $q$ -difference equation for the unknown polynomial  $c(x)$ . However,  $z \in K$  is also not known yet. Let  $u_{ik}$  denote the coefficient of  $x^k$  in  $f_i$ . Since  $c(0) \neq 0$ , we have  $\alpha_0 \neq 0$  in (3), hence applying (3) to (10) we obtain

$$\sum_{i=0}^{\rho} u_{i0} z^i = 0. \quad (11)$$

We may assume that not all  $u_{i0}$  are zero, or else we start by first cancelling a power of  $x$  from the coefficients of (10). Thus  $z$  is a nonzero root of  $f(z) = \sum_{i=0}^{\rho} u_{i0} z^i$ , and is algebraic over  $K$ .

If  $N = \deg c(x)$  then by (2),

$$\sum_{i=0}^{\rho} u_{id} z^i q^{iN} = 0, \quad (12)$$

hence  $w = zq^N$  is a nonzero root of  $g(w) = \sum_{i=0}^{\rho} u_{id} w^i$ . It follows that  $q^N$  is a root of  $p(x) = \text{Resultant}_w(f(w), g(wx))$ , thus to obtain an upper bound on  $N$  computation in algebraic extensions of  $K$  is not necessary.

In summary, we find the factors of  $r(x)$  as follows:

1.  $a(x)$  is a monic factor of  $p_0(x)$ ,
2.  $b(x)$  is a monic factor of  $p_\rho(q^{1-\rho} x)$ ,
3.  $z$  is a root of Eqn. (11),
4.  $c(x)$  is a nonzero  $q$ -polynomial solution of (10).

Then  $r = z(a/b)(Qc/c)$  and  $Qy = ry$ .

**Example 1** Let us find a  $q$ -hypergeometric solution  $y$  of  $Ly = 0$  where

$$L = xQ^3 - q^3x^2Q^2 - (x^2 + q)Q + qx(x^2 + q).$$

The candidates for  $a(x)$  are

$$1, x, x^2 + q, x(x^2 + q),$$

and the candidates for  $b(x)$  are

$$1, x.$$

Here we explore only the choice  $a(x) = x$  and  $b(x) = 1$ . The corresponding equation (10) is, after cancelling one  $x$ ,

$$z^3q^3x^3c(q^3x) - z^2q^4x^3c(q^2x) - z(x^2 + q)c(qx) + q(x^2 + q)c(x) = 0, \quad (13)$$

whence  $f(z) = -qz + q^2$  with unique root  $z = q$ , and  $g(w) = q^3w^3 - q^4w^2$  with unique nonzero root  $w = q = zq^N = q^{N+1}$ . It follows that  $N = 0$  is the only possible degree for  $c$ . Equation (13) is satisfied by  $c = 1$ . Thus we have found  $r = z(a/b)(Qc/c) = qx$ , and the corresponding  $q$ -hypergeometric solution of  $Ly = 0$  satisfies  $Qy = qxy$ . We can take, for instance,  $y_n = x(x/q)(x/q^2) \cdots (x/q^n) = q^{\binom{n+1}{2}}$ .

To find other  $q$ -hypergeometric solutions (if any), the remaining combinations for  $a(x)$  and  $b(x)$  could be tried; or even better, the order of the equation could be reduced using the obtained solution, and the algorithm used recursively on the reduced equation. Our *Mathematica* implementation of this algorithm (which we call `qHyper`) shows that up to a constant factor, there are in fact no other  $q$ -hypergeometric solutions:

```
In[1]:= qHyper[x y[q^3 x] - q^3 x^2 y[q^2 x] -
               (x^2 + q) y[q x] + q x (x^2 + q) y[x]] == 0, y[x]]
Out[1]= {q x}
```

Note that `qHyper` returns a list of quotients  $Qy/y$  rather than solutions  $y$  themselves.  $\square$

**Example 2** Consider the equation  $Ly = 0$  where  $L = Q^2 - (1 + q)Q + q(1 - qx^2)$ . As shown by `qHyper`,

```
In[2]:= qHyper[y[q^2 x] - (1 + q) y[q x] + q (1 - q x^2) y[x], y[x]]
Out[2]= {1 - Sqrt[q] x, 1 + Sqrt[q] x}
```

this equation has two linearly independent  $q$ -hypergeometric solutions,  $(\sqrt{q}; q)_n$  and  $(-\sqrt{q}; q)_n$ . Here  $K$  is the splitting field of  $1 - qx^2$ .  $\square$

## 5 Some related problems

### 5.1 Nonhomogeneous equations

Consider the problem of finding  $q$ -hypergeometric solutions  $y$  of the nonhomogeneous equation  $Ly = b$  where  $b \neq 0$ . Let  $Qy = ry$  where  $r \in K(x)$ . Then  $Ly = fy$  where  $f = \sum_{i=0}^p p_i \prod_{j=0}^{i-1} Q^j r \in K(x)$ . This simple fact has two important consequences:

1.  $b = fy$  is  $q$ -hypergeometric,
2.  $y = b/f$  is a  $q$ -rational multiple of  $b$ .

Let  $Qb = sb$  where  $s \in K(x)$  is given. We look for  $y$  in the form  $y = fb$  where  $f \in K(x)$  is an unknown  $q$ -rational function. Substituting this into  $Ly = b$  gives

$$\sum_{i=0}^{\rho} p_i \left( \prod_{j=0}^{i-1} Q^j s \right) Q^i f = 1.$$

Now  $q$ -rational solutions of this equation can be found using the algorithm given in [1].

In particular, this gives an algorithm for the problem of *indefinite  $q$ -hypergeometric summation*: Given a  $q$ -hypergeometric sequence  $b_n$ , decide if  $y_n = \sum_{j=0}^{n-1} b_j$  is  $q$ -hypergeometric, and if so, express it in closed form. Obviously  $y_n$  satisfies  $y_{n+1} - y_n = b_n$ . Since we are interested in  $q$ -hypergeometric solutions, we can rewrite this as  $Qy - y = b$  and use the technique described above.

**Example 3** Let  $y_n = \sum_{j=0}^{n-1} b_j$  where  $b_n = q^n (q; q)_n$ . Then  $y$  satisfies the equation

$$Qy - y = b \tag{14}$$

where  $s = Qb/b = q(1 - qx)$ . The equation for  $f$  is

$$q(1 - qx)Qf - f = 1,$$

with unique  $q$ -rational solution  $f = -1/(qx)$ . Hence  $y_n = C - (q; q)_n/q$  where  $C$  is a constant. Since  $y_0 = 0$  it follows that  $C = 1/q$  and  $y_n = (1 - (q; q)_n)/q$ .  $\square$

The same technique for solving nonhomogeneous equations also works when we look for  $q$ -hypergeometric term solutions in  $M = K[[x]]$ .

**Example 4** Let

$$Q^2y(x) - (1 - qx)Qy(x) + qy(x) = b(x) \tag{15}$$

where

$$b(x) = \sum_{i=0}^{\infty} \frac{x^i}{(q; q)_i}.$$

Here  $b(qx) = (1 - x)b(x)$ , as can be easily verified. Thus  $s = 1 - x$  and the equation for  $f$  is

$$(1 - qx)(1 - x)Q^2f - (1 - qx)(1 - x)Qf + qf = 1$$

with  $q$ -rational solution  $f = 1/q$ . Hence  $y(x) = b(x)/q$  solves (15).  $\square$

## 5.2 q-hypergeometric series solutions

Assume that  $y = \sum_{j=0}^{\infty} \alpha_j x^j$  and  $Ly = b$  where  $b = \sum_{j=0}^{\infty} \beta_j x^j$ . As in (1), we obtain

$$\sum_{j=\max\{l-d,0\}}^l \sum_{i=0}^{\rho} c_{i,l-j} \alpha_j q^{ij} = \beta_l, \quad \text{for } l \geq 0. \quad (16)$$

We separate the cases  $0 \leq l < d$  and  $l \geq d$ . In the former case, (16) yields initial conditions

$$\sum_{j=0}^l \alpha_j \sum_{i=0}^{\rho} c_{i,l-j} q^{ij} = \beta_l, \quad \text{for } 0 \leq l < d, \quad (17)$$

while in the latter, substitutions  $m = l - d$ ,  $s = j - m$ , and  $X = q^m$  transform (16) into the *associated q-difference equation*

$$\sum_{s=0}^d \alpha_{m+s} \sum_{i=0}^{\rho} c_{i,d-s} q^{is} X^i = \beta_{m+d}, \quad \text{for } m \geq 0, \quad (18)$$

for the unknown sequence  $(\alpha_m)_{m=0}^{\infty}$ . We use the algorithms of Sections 4 and 5.1 to find all solutions of (18) which are linear combinations of  $q$ -hypergeometric terms, then select the constants in these combinations so that conditions (17) are satisfied (if possible).

**Example 5** Let us find  $q$ -hypergeometric series solutions  $y$  of

$$q^2 x^2 Q^3 y + (1+q)xQ^2 y + (1-x)Qy - y = 0. \quad (19)$$

The associated equation (18) in this case is

$$(q^2 X - 1)\alpha_{m+2} + (q^2(q+1)X^2 - qX)\alpha_{m+1} + q^2 X^3 \alpha_m = 0 \quad (20)$$

and qHyper finds two solutions:

```
In[3]:= qHyper[(q^2 X - 1) y[q^2 X] + (q^2 (1 + q) X^2 - q X) y[q X] +
q^2 X^3 y[X] == 0, y[X]]
```

```
Out[3]= {-x, -q X^2 / (1 - q X)}
```

Thus the general solution of (20) is  $\alpha_m = Cq^{m^2}/(q;q)_m + D(-1)^m q^{\binom{m}{2}}$  where  $C$  and  $D$  are arbitrary constants. Equations (17) imply that  $D = 0$ . Hence  $y^{(1)} = \sum_{m=0}^{\infty} q^{m^2} x^m/(q;q)_m$  is a  $q$ -hypergeometric series solution of (19).

Note that running qHyper on equation (19) itself we obtain another solution  $y^{(2)} = (-1)^n/q^{\binom{n}{2}}$ .  $\square$

**Example 6** The right-hand side of the equation (15) is both a  $q$ -hypergeometric term and a  $q$ -hypergeometric series. The associated nonhomogeneous equation

$$(qX^2 - X + 1)\alpha_{m+1} + X\alpha_m = \frac{1}{q(q; q)_{m+1}}$$

can be solved as described in Section 5.1. Here  $s = 1/(1 - q^2X)$  and the equation for  $f$

$$\frac{1 - X + qX^2}{1 - q^2X} Qf + Xf = 1$$

is satisfied by the  $q$ -rational function  $f = 1 - qX$ . Thus  $\alpha_m = (1 - qX)/(q(q; q)_{m+1}) = 1/(q(q; q)_m)$ , and we find the same solution  $y(x) = b(x)/q$  as in Example 4.  $\square$

### 5.3 Deriving $q$ -hypergeometric identities

Another important application is definite  $q$ -hypergeometric summation. The corresponding algorithm of [7] will produce a  $q$ -difference equation for the sum, but in general it will not be of minimal order. Thus it can happen that the equation will be of order 2 or more while the sum can actually be expressed in closed form. In this case one can use our algorithm to find the  $q$ -hypergeometric solutions of the equation, and then test them to see which linear combination – if any – gives the initial sum.

In analogy with the ordinary hypergeometric case [4], we also expect our algorithm to play an important role in the factorization algorithm for linear  $q$ -difference operators.

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