

# SOME STUDIES OF FACTORS OF INFINITE WORDS GENERATED BY INVERTIBLE SUBSTITUTION

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## ABSTRACT

In this paper, we introduce a class of infinite words generated by invertible substitutions and establish some elementary properties of factors of these infinite words. As a consequence, we prove these words are Sturmian words.

Keywords: Invertible substitution, singular word, sesquipower.

## 1. Introduction

Let  $S = \{a, b\}$  be the alphabet of two letters. Let  $S^*$  and  $\tilde{S}$  be respectively the free monoid with empty word  $\epsilon$  as neutral element and the free group generated by  $S$ .

A morphism  $\pi : S^* \rightarrow S^*$  is called a substitution over  $S$ . If  $\pi$  is also in  $Aut(\tilde{S})$ , then  $\pi$  is called an invertible substitution, where  $Aut(\tilde{S})$  denotes the group of automorphisms over  $\tilde{S}$ . It is known [6,7] that  $Aut(\tilde{S})$  is generated by the following three special automorphisms:  $\sigma = (ab, a)$ ,  $\phi = (b, a)$ ,  $\psi = (a, b^{-1})$ , where  $\tau = (u, v)$  means that  $\tau(a) = u, \tau(b) = v$ .

Let  $\sigma = (ab, a)$  be defined as above, then the infinite Fibonacci word is obtained by iterating  $\sigma$  starting with the letter  $a$  which is also a fixed point of  $\sigma$  [1]. The combinatorial properties of the infinite Fibonacci word are studied extensively by many authors, (see Berstel[1], Mignosi[5] and the references therein.) Notice that  $\sigma$  is an invertible substitution, we are led to consider a natural generalization of the infinite Fibonacci word: the infinite words generated by the invertible substitutions.

Brown [2], Kosa [4], Séébold [9] have considered the invertible substitutions of  $\langle \phi, \Phi_1, \Phi_2 \rangle$ , where  $\Phi_1 = (a, ab), \Phi_2 = (a, ba)$ ,  $\langle \phi, \Phi_1, \Phi_2 \rangle$  denotes the semigroup generated by  $\phi, \Phi_1, \Phi_2$ . They studied both combinatorial and arithmetic properties of the infinite words generated by these invertible substitutions.

The purpose of this note is to study the properties of factors of infinite words generated by the invertible substitutions. We recall first some preliminaries. Then in section 1, we discuss the local isomorphism of two invertible substitutions, by using the result, instead of studying the general case, we only need to study a special class of invertible substitutions. In section 2, after establishing some elementary properties of factors, we introduce singular factors and discuss their properties. As an example, by using singular factors, we determine the special factors and prove that all these infinite words are Sturmian words.

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In this paper, we shall use the following terminology.

Let  $w \in S^*$  be a word. We denote by  $|w|$  the length of  $w$ , and we denote by  $|w|_a$  (resp.  $|w|_b$ ) the number of letters  $a$  (resp.  $b$ ) in the word  $w$ , we denote by  $L(w)$  the vector  $(|w|_a, |w|_b)$ .

A word  $v$  is a factor of a word  $w$  and we write  $v \prec w$ , if there exists  $u, u' \in A^*$ , such that  $w = uvu'$ . We say that  $v$  is a left ( resp. right ) factor of a word  $w$  and we note  $v \triangleleft w$  ( resp.  $v \triangleright w$  ), if there exists  $u \in A^*$  such that  $w = vu$  ( resp.  $w = uv$  ). The notions of factor, left factor are extended in a natural way to an infinite word.

Let  $w = x_1 x_2 \dots x_n \in S^*$ , the reversed word  $\bar{w}$  of  $w$  is defined as  $\bar{w} = x_n \dots x_2 x_1$ . A word  $w$  is called palindrome if  $w = \bar{w}$  and we denote by  $\mathcal{P}$  the set of palindrome words.

Let  $w = x_1 x_2 \dots x_n$ . For  $1 \leq k \leq n$ , we define the  $k$ th conjugation of  $w$  as  $C_k(w) = x_{k+1} \dots x_n x_1 \dots x_k$ , and we denote by  $C(w) = \{C_k(w); 1 \leq k \leq |w|\}$ .

A word  $w \in S^*$  is called primitive if  $w = u^p, u \in S^*, p > 0$  implies  $w = u$ .

Let  $\tau : S^* \rightarrow S^*$  be a substitution of  $S^*$ . We denote by  $F_\tau$  the fixed point of  $\tau$  [8] (if it exists).

Let  $w$  be a word, we denote by  $\Omega_n(w)$  the set of factors of  $w$  of the length  $n$ , where  $|w| \geq n$ , and we note simply  $\Omega_n := \Omega_n(F_\tau)$ .

Let  $w = x_1 x_2 \dots x_n \in S^*$ , we denote by  $w^{-1}$  the inverse word of  $w$ , that is  $w^{-1} = x_n^{-1} \dots x_2^{-1} x_1^{-1}$ . Let  $w = uv$ , then  $wv^{-1} = u$  by convention.

## 1. Local isomorphism of invertible substitutions

**Lemma 1.1.** Let

$$\mathcal{M} = \left\{ M = \begin{pmatrix} p & r \\ q & s \end{pmatrix}, p, q, r, s \in \mathbb{N}, \det \begin{pmatrix} p & r \\ q & s \end{pmatrix} = \pm 1 \right\}$$

Then  $\mathcal{M} = \langle \alpha, \beta \rangle$ , where  $\alpha = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ ,  $\beta = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ .

**Proof.** Let

$g_1 = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ ,  $g_2 = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ ,  $g_3 = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ ,  $g_4 = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}$ . Then it is easy to check that  $g_1 = \alpha^2, g_2 = \beta\alpha, g_3 = \alpha\beta, g_4 = \alpha\beta\alpha$ .

Now let  $M = \begin{pmatrix} p & r \\ q & s \end{pmatrix}$  be an element of  $\mathcal{M}$ , such that at least one entry is larger than 1. By using  $\alpha$ , we can make  $p$  the largest entry of the matrix. Moreover,  $r \geq s$  (otherwise  $|\det M| > 1$ ).

Let  $M' = \begin{pmatrix} p-q & r-s \\ q & s \end{pmatrix}$ , then  $M = \beta\alpha M'$ . By the same argument as above, we can transform  $M'$  by  $\alpha$  such that it satisfies the same condition as  $M$ .

After a finite number of steps, the matrix  $M$  will be transformed, by the actions of  $\alpha$  and  $\beta$ , to one of the six matrices:  $\alpha, \beta, g_1, g_2, g_3, g_4$ , which may be generated by  $\alpha, \beta$ .

From the definition of the substitutive matrix, we can prove easily

**Lemma 1.2.** Let  $\tau = \tau_1 \tau_2 \dots \tau_n$ , then

$$M_\tau = M_{\tau_1} M_{\tau_2} \dots M_{\tau_n},$$

where  $M_\tau$  and  $M_{\tau_j}$  are the substitutive matrices of  $\tau$  and  $\tau_j$  respectively.

**Lemma 1.3.** For any  $M \in \mathcal{M}$ , there is a  $\tau \in \langle \sigma, \phi \rangle$ , such that  $M_\tau = M$ .

**Proof.** By lemma 1.1, for any  $M \in \mathcal{M}$ , there exist  $M_1, M_2, \dots, M_n \in \{\alpha, \beta\}$  such that  $M = M_1 M_2 \dots M_n$ .

Note that  $M_\sigma = \beta$ ,  $M_\phi = \alpha$ , so if we take  $\tau = \tau_1 \tau_2 \dots \tau_n$ , where  $\tau_j = \phi$  or  $\sigma$  according to  $M_j = \alpha$  or  $\beta$ , then by lemma 1.2, we get

$$M_\tau = M_{\tau_1} M_{\tau_2} \dots M_{\tau_n} = M_1 M_2 \dots M_n = M.$$

**Lemma 1.4.** [8] Let  $\tau_1, \tau_2$  be two invertible substitutions, then the following two assertions are equivalent:

- 1).  $M_{\tau_1} = M_{\tau_2}$ ;
- 2). there exists a  $w \in S^*$ , such that  $\tau_1 = w\tau_2 w^{-1}$ , (or  $\tau_1 = w^{-1}\tau_2 w$ ).

Let  $u = u_1 u_2 \dots u_n \dots$  and  $v = v_1 v_2 \dots v_n \dots$  be two infinite words over the alphabet  $S$ . We say that  $u$  and  $v$  are locally isomorphic if any factor ( or its mirror image ) of  $u$  is also factor of  $v$  and vice versa.

If  $u$  and  $v$  are locally isomorphic, we shall note  $u \simeq v$ .

**Lemma 1.5.** Let  $\tau_1, \tau_2$  be two invertible substitutions and let  $F_{\tau_1}, F_{\tau_2}$  be the fixed points of  $\tau_1$  and  $\tau_2$  respectively. If  $M_{\tau_1} = M_{\tau_2}$ , then  $F_{\tau_1} \simeq F_{\tau_2}$ .

**Proof.** By lemma 1.4,  $\exists w \in S^*$ , such that

$$\tau_1(a) = w\tau_2(a)w^{-1}.$$

Furthermore for any  $n$

$$\tau_1^n(a) = \tau_1^{n-1}(w)\tau_1^{n-2}(w)\dots w\tau_2^n(a)w^{-1}\dots \tau_2^{n-1}(w^{-1})\tau_2^n(w^{-1})$$

That is,  $\exists w_n \in S^*$ , such that  $\tau_1^n(a)w_n = w_n\tau_2^n(a)$ . Thus  $\tau_1^n(a)$  and  $\tau_2^n(a)$  are conjugate, (see [3]), that is, there exist  $u_n$  and  $v_n$ , such that  $\tau_1^n(a) = u_n v_n$ , and  $\tau_2^n(a) = v_n u_n$ .

This means that, both  $u_n$  and  $v_n$  are factors of  $\tau_1^n(a)$  and  $\tau_2^n(a)$ . Since  $|\tau_1^n(a)|$  tends to infinity, so does  $|u_n|$  (or  $|v_n|$ ). Let  $w$  be a factor of  $F_{\tau_1}$ , by proposition 4.12 of [8],  $w \prec u_n$  (or  $w \prec v_n$ ) for  $n$  large enough, thus  $w \prec F_{\tau_2}$ . Inversely, by the same argument, if  $w \prec F_{\tau_2}$ , then  $w \prec F_{\tau_1}$ . We thus prove that  $F_{\tau_1} \simeq F_{\tau_2}$ .

The case of  $\tau_1(a) = w^{-1}\tau_2(a)w$  can be treated in the same manner.

**Remark 1.6.** By lemma 1.3 and 1.5, we see that to study the properties of factors of an infinite word generated by an invertible substitution, it suffices to consider the elements of  $\langle \sigma, \phi \rangle$ , that is, the invertible substitutions of the following form:

$$\sigma^{n_k} \circ \phi \circ \sigma^{n_{k-1}} \circ \phi \dots \sigma^{n_2} \circ \phi \circ \sigma^{n_1},$$

where  $n_1, n_k \geq 0; n_2, \dots, n_{k-1} \geq 1$ .

Throughout this paper, we only discuss the following form:

$$\sigma^{n_k} \circ \phi \circ \sigma^{n_{k-1}} \circ \phi \dots \sigma^{n_2} \circ \phi \circ \sigma^{n_1},$$

where  $n_j \geq 1, 1 \leq j \leq k$ , and we denote by  $\mathcal{G}$  the set of these invertible substitutions.

The other three cases ( $n_k \geq 1, n_1 = 0; n_k = 0, n_1 \geq 1; n_k = n_1 = 0$ ) can be discussed in the same way.

Lemma 1.1 and lemma 1.3 give an algorithm for finding an invertible substitution of which the substitutional matrix is equal to a given matrix of  $\mathcal{M}$ . Now we give an explicit algorithm by continued fractions.

The following lemma can be checked easily.

**Lemma 1.7.** For any  $w \in S^*$ , we have

- 1).  $L(\sigma(w)) = (|w|_a + |w|_b, |w|_a), L(\phi(w)) = (|w|_b, |w|_a);$
- 2).  $L((\sigma\phi)^n(w)) = (|w|_a + n|w|_b, |w|_b), L((\phi\sigma)^n(w)) = (n|w|_a + |w|_b, |w|_a).$

**Lemma 1.8.** Let  $(p, q) = 1, q \neq 1$ . Then there exists a unique pair  $(r, s)$  (resp.  $(r', s')$ ),  $0 \leq r < p, 0 \leq s < q$  (resp.  $0 \leq r' < p, 0 \leq s' < q$ ), such that  $ps - rq = 1$  (resp.  $-1$ ).

**Proof.** The existence of a such pair is well known. Now suppose that there is another pair  $(r^*, s^*)$  satisfying the condition, then we have

$$p(s - s^*) = q(r - r^*).$$

Since  $s - s^* < q, r - r^* < p$ , we get a contradiction to the hypothesis for  $(p, q) = 1$ .

**Remark 1.9.** If  $q = 1$ , the conclusion holds trivially.

Now let  $M = \begin{pmatrix} p & r \\ q & s \end{pmatrix} \in \mathcal{M}$  with  $p$  the largest entry of  $M$ . Let  $t_0, t_1, \dots, t_n$  be the continued fraction expansion of  $p/q$ . Then by the lemmas 1.7 and 1.8, it is not difficult to prove that

**Proposition 1.10.** With the notations above. Let

$$\begin{aligned} \tau_1 &= (\sigma\phi)^{t_0-1}\sigma(\sigma\phi)^{t_1-1}\sigma(\sigma\phi)^{t_{n-1}-1}\sigma(\sigma\phi)^{t_n-1}\sigma, \\ \tau_2 &= (\sigma\phi)^{t_0-1}\sigma(\sigma\phi)^{t_1-1}\sigma(\sigma\phi)^{t_{n-1}-1}\sigma(\sigma\phi)^{t_n-1}\phi\sigma. \end{aligned}$$

If  $\det M = 1$ , then  $M_{\tau_1} = M$  if  $n$  is odd and  $M_{\tau_2} = M$  if  $n$  is even.

If  $\det M = -1$ , then  $M_{\tau_1} = M$  if  $n$  is even and  $M_{\tau_2} = M$  if  $n$  is odd.

## 2. Studies of factors

In this section, we first establish some elementary properties of the infinite words generated by invertible substitution, then we introduce singular words and study their properties. As we shall see, the singular words play an important role in the study of the properties of the factors.

If  $\alpha\beta$  appears in the context, we mean that  $\alpha, \beta \in S$  with  $\alpha \neq \beta$ . Moreover if we write  $\alpha\beta \triangleright \tau(b)$ , we assume  $|\tau(b)| \geq 2$ .

The following lemma can easily be proved:

**Lemma 2.1** Let  $\tau \in \mathcal{G}$ , then

- 1).  $(\sigma\phi)^n(a) = a, (\sigma\phi)^n(b) = a^n b;$
- 2). If  $\sum_{j=1}^k n_j + k \in 2N+1$ , then  $ab \triangleright \tau(a)$ ; if  $\sum_{j=1}^k n_j + k \in 2N$ , then  $ba \triangleright \tau(a)$ ;
- 3).  $\tau(b) \triangleleft \tau(a), \tau(a) \triangleleft \tau^n(a), \tau(b) \triangleleft \tau^n(b), n \geq 2;$
- 4). If  $|\tau(b)| = 1$ , then  $\tau = (\sigma\phi)^n\sigma$  for some  $n$ ; if  $|\tau(b)| = 2$ , then  $\tau = \sigma(\sigma\phi)^n\sigma$  for some  $n$ ; if  $|\tau(a)| = 2$ , then  $\tau = \sigma$ ;
- 5). If  $|\tau(b)| > 1$  and  $\alpha\beta \triangleright \tau(a)$ , then  $\beta\alpha \triangleright \tau(b)$ .

**Lemma 2.2** let  $\pi = (\sigma\phi)^n\sigma^2$  and let  $w \in S^*$ . Then

- 1).  $a^{n+1}b \triangleleft \pi(w);$
- 2).  $a \triangleright w \Rightarrow ba \triangleright \pi(w), b \triangleright w \Rightarrow a^{n+1}b \triangleright \pi(w), ab \triangleright w \Rightarrow a^{n+2}b \triangleright \pi(w);$
- 3). The highest power of  $a$  appearing in  $\pi(w)$  is  $a^{n+2}$  if  $aab \prec \sigma^2(w)$  and  $a^{n+1}$  otherwise; The factor of  $\pi(w)$  situated between two adjacent  $b$  is either  $a^{n+1}$  or  $a^{n+2}$ ;
- 4). The factor of  $\pi(w)$  situated between two adjacent  $b$  is either  $a^{n+1}$  or  $a^{n+2}$ .

**Proof.** Let  $\sigma^2(w) = x_1x_2\dots x_{m-1}x_m$ . Since  $\sigma^2(a) = aba, \sigma^2(b) = ab$ , it is easy to see that: 1.  $x_1x_2 = ab$ ; 2.  $x_{m-1}x_m = ab$  or  $ba$  according to  $b \triangleright w$  or  $a \triangleright w$ ; 3.  $a^3 \not\prec \sigma^2(w), b^2 \not\prec \sigma^2(w)$ ; 4. the factor of  $\sigma^2(w)$  situated between two adjacent  $b$  is either  $a^2$  or  $a$ . Since

$$\pi(w) = (\sigma\phi)^n\sigma^2(w) = (\sigma\phi)^n(x_1)(\sigma\phi)^n(x_2)\dots(\sigma\phi)^n(x_{m-1})(\sigma\phi)^n(x_m),$$

the lemma follows from the discussions above and lemma 2.1.1.

**Lemma 2.3** Let  $k(\tau(a))$  (resp.  $k(\tau(b))$ ) be the integer such that  $a^{k(\tau(a))}$  (resp.  $b^{k(\tau(b))}$ ) is the highest power of  $a$  (resp.  $b$ ) appearing in  $\tau(a)$ . (resp.  $\tau(b)$ ) Then

- 1).  $k(\tau(b)) \leq k(\tau(a));$
- 2).  $a^{p(\tau(a))}b \triangleleft \tau(a)$ , where  $p(\tau(a))$  is  $k(\tau(a))$  or  $k(\tau(a)) - 1$ ;
- 3). Either  $a^{p(\tau(a))}b \triangleright \tau(a)$ , or  $ba \triangleright \tau(a)$ ;
- 4).  $a^{p(\tau(a))} \prec F_\tau, a^{p(\tau(a))+2} \not\prec F_\tau, b^2 \not\prec F_\tau;$
- 5). Any factor of  $F_\tau$  situated between two adjacent  $b$  is either  $a^{p(\tau(a))+1}$  or  $a^{p(\tau(a))}$ .

**Proof.** Let  $\tau = \sigma^{n_k}\phi\sigma^{n_{k-1}}\phi\dots\sigma^{n_2}\phi\sigma^{n_1}$ .

If  $n_j = 1$  for all  $j, 1 \leq j \leq k$ , then  $\tau = (\sigma\phi)^k\sigma$ . In this case, the lemma can be proved easily.

Now suppose that there exists at least one index  $j$ , such that  $n_j > 1$ . Let  $j_0$  be the largest index of these indices. Then

$$\tau = (\sigma\phi)^{k-j_0}\sigma^2(\sigma^{n_{j_0}-2}\dots\sigma^{n_2}\phi\sigma^{n_1}). \quad (1)$$

Hence the proof of this lemma follows from the lemma 2.2 by treating the different cases of (1) (the discussion is tedious, but not difficult, and we omit the details.)

We now determine the length of  $\tau(a)$  and  $\tau(b)$ . Let  $\mu, \nu \in \mathbb{N}$ , we denote by  $F_n(\mu, \nu)$  the Fibonacci number with initial condition  $F_{-1}(\mu, \nu) = \nu, F_0(\mu, \nu) = \mu$ . Thus for  $w \in S^*$ , it is readily seen that

$$L(\sigma^n(w)) = (F_n(|w|_a, |w|_b), F_{n-1}(|w|_a, |w|_b)), L(\phi(w)) = (|w|_b, |w|_a).$$

Thus it is not difficult to get

**Lemma 2.4.** Let  $\tau = \sigma^{n_k} \phi \dots \sigma^{n_2} \phi \sigma^{n_1}$ , then

$$\begin{aligned} |\tau(a)| &= F_{n_k+1}(\mu_{k-1}, \nu_{k-1}), |\tau(b)| = F_{n_k}(\mu_{k-1}, \nu_{k-1}); \\ L(\tau(a)) &= (F_{n_k}(\mu_{k-1}, \nu_{k-1}), F_{n_k-1}(\mu_{k-1}, \nu_{k-1}), \\ L(\tau(b)) &= (F_{n_k-1}(\mu_{k-1}, \nu_{k-1}), F_{n_k-2}(\mu_{k-1}, \nu_{k-1})); \end{aligned}$$

$$M_\tau = \begin{pmatrix} F_{n_k}(\mu_{k-1}, \nu_{k-1}) & F_{n_k-1}(\mu_{k-1}, \nu_{k-1}) \\ F_{n_k-1}(\mu_{k-1}, \nu_{k-1}) & F_{n_k-2}(\mu_{k-1}, \nu_{k-1}) \end{pmatrix},$$

where  $\mu_0 = 1, \nu_0 = 0, \mu_{k-1} = F_{n_{k-1}}(\mu_{k-2}, \nu_{k-2}), \nu_{k-1} = F_{n_k}(\mu_{k-2}, \nu_{k-2})$ .

**Lemma 2.5.**  $\forall \gamma \in S$ , the word  $\tau(\gamma)$  is primitive.

**Proof.** Suppose that  $\tau(\gamma)$  is not primitive, then there will exist a word  $w \in S^*$  and an integer  $p \geq 2$ , such that  $\tau(\gamma) = w^p$ . Hence  $L(\tau(\gamma)) = (p|w|_a, p|w|_b)$ , thus  $\det(M_\tau)$  will have factor  $p$ , which is in contradiction with the fact  $\det(M_\tau) = \pm 1$ .

Now we are going to establish an important factorization of  $\tau(a)$  which will be used in the following.

Since  $\tau(b) \triangleleft \tau(a)$ , we can write  $\tau(a) = \tau(b)u_n^*$ . Let  $u \triangleleft \tau(a)$  with  $|u| = |u_n^*|$ , we can write also  $\tau(a) = uv$  with  $|v| = |\tau(b)|$ .

**Lemma 2.6.** Let  $\tau(a) = \tau(b)u^* = uv$  as above. Let  $\alpha\beta \triangleright \tau(a), |\tau(b)| \geq 2$ .

Then

- 1).  $u = u^*$ ;
- 2).  $\tau(b)\alpha^{-1}\beta^{-1} = v\beta^{-1}\alpha^{-1}$ ;
- 3).  $\tau(a) = \tau(b)u = \tau(b)\alpha^{-1}\beta^{-1}\alpha\beta$ .

**Proof.** Since  $\alpha\beta \triangleright \tau(a), \beta\alpha \triangleright \tau(b)$ .

We prove the lemma by induction with respect to the length of  $\tau$  in  $\Sigma := \{\sigma, \phi\}$ .

- 1). If  $\tau = \sigma$ , then  $\sigma(a) = ab = \tau(b)b$ , and the statement is true.

Suppose that for  $\tau = \tau_n \dots \tau_1$ , we have

$$\tau_n \dots \tau_1(a) = \tau_n \dots \tau_1(b)u_n^* = u_nv_n \quad (*)$$

with  $u_n^* = u_n$ , (so  $|v_n| = |\tau_n \dots \tau_1(b)|$ .)

Let  $\tau = \tau_{n+1} \tau_n \dots \tau_1$  and let

$$\tau(a) = \tau_{n+1} \tau_n \dots \tau_1(a) = \tau_{n+1} \tau_n \dots \tau_1(b)u_{n+1}^* = u_{n+1} v_{n+1}$$

with  $|u_{n+1}^*| = |u_{n+1}|$ . We have to prove  $u_{n+1}^* = u_{n+1}$ .

If  $\tau_{n+1} = \phi$ , the proof is trivial. Now let  $\tau_{n+1} = \sigma$ , we have by (\*)

$$\begin{aligned} \sigma(\tau_n \dots \tau_1(a)) &= \sigma(u_nv_n) \\ &= \sigma(u_n)\sigma(v_n) = \sigma(\tau_n \dots \tau_1(b))\sigma(u_n^*) \end{aligned}$$

Thus  $u_{n+1}^* = \sigma(u_n^*, u_{n+1}) = \sigma(u_n)$ .

Since  $u_n^* = u_n$  by hypothesis of induction, we get  $u_{n+1}^* = u_{n+1}$ .

2.) The case of  $|\tau(b)| = 2$  can be checked directly. Let  $\tau = \tau_n \dots \tau_1$ , then by 1) we have

$$\tau_{n+1} \tau_n \dots \tau_1(a) = \tau_{n+1} \tau_n \dots \tau_1(b)u_{n+1} = u_{n+1} v_{n+1}.$$

Let  $\alpha\beta \triangleright \tau_n \dots \tau_1(a)$  and suppose that  $v_n \beta^{-1} \alpha^{-1} = \tau_n \dots \tau_1(b) \alpha^{-1} \beta^{-1}$ . We shall prove

$$v_{n+1} \alpha^{-1} \beta^{-1} = \tau_{n+1} \tau_n \dots \tau_1(b) \beta^{-1} \alpha^{-1}.$$

As in 1), it suffices to consider  $\tau_{n+1} = \sigma$ . Since  $\alpha\beta \triangleright \tau_{n+1} \tau_n \dots \tau_1(b), \beta\alpha \triangleright \sigma(v_{n+1})$ , we have

$$\begin{aligned} v_{n+1} = \sigma(v_n) &= \sigma(\tau_n \dots \tau_1(b) \alpha^{-1} \beta^{-1} \alpha \beta) \\ &= \sigma(\tau_n \dots \tau_1(b))(\sigma(\beta \alpha))^{-1} \sigma(\alpha \beta). \end{aligned}$$

Notice that  $\sigma(\alpha \beta) = a \beta \alpha$  for any  $\alpha \neq \beta$ , and therefore we obtain

$$v_{n+1} \alpha^{-1} \beta^{-1} = \tau_{n+1} \tau_n \dots \tau_1(b) \beta^{-1} \alpha^{-1}.$$

3). The conclusion follows immediately from 1) and 2).

**Lemma 2.7.** Let  $w \in \mathcal{P}$ , then  $\sigma(w)a, a^{-1}\sigma(w) \in \mathcal{P}$ .

**Proof.** Since  $w \in \mathcal{P}$ , we can write  $w$  in the following form:

$$a^{k_1} b^{l_1} a^{k_2} b^{l_2} \dots a^{k_m} \gamma a^{k_m} \dots b^{l_2} a^{k_2} b^{l_1} a^{k_1}$$

where  $\gamma \in \{a, b, \epsilon\}, k_j, l_j \geq 1$ , if  $1 \leq j \leq m-1$ , and  $k_m \geq 0$ .  
Thus we have

$$\begin{aligned} \sigma(w)a &= (ab)^{k_1} a^{l_1} \dots (ab)^{k_m} \sigma(\gamma)(ab)^{k_m} \dots a^{l_2} (ab)^{k_2} a^{l_1} (ab)^{k_1} a \\ &= (ab)^{k_1} a^{l_1} \dots (ab)^{k_m} \sigma(\gamma) a (ba)^{k_m} \dots a^{l_2} (ba)^{k_2} a^{l_1} (ba)^{k_1}, \end{aligned}$$

Since  $\sigma(\gamma)a$  is equal to  $a, aba, aa$  respectively according to  $\gamma$  being  $\epsilon, a, b$ ,  $\sigma(w)a$  is a palindrome with center  $a, b, \epsilon$  respectively by the formula above.

**Corollary 2.8.** Let  $\beta \triangleright \tau(\gamma), \gamma \in S$ . Then  $\alpha\tau(\gamma)\beta^{-1} \in \mathcal{P}$ . In particular, if  $|\tau(\gamma)| \geq 3$ , and  $\alpha\beta \triangleright \tau(\gamma)$ , then  $\tau(\gamma)\beta^{-1}\alpha^{-1} \in \mathcal{P}$ .

**Proof.** Let  $\tau = \tau_n \tau_{n-1} \dots \tau_1$ . (Notice that  $\tau_1 = \sigma$  by the definition of  $\tau$ .)

We prove this lemma by induction as in lemma 2.6. We only prove the case  $\gamma = a$ , the case of  $\gamma = b$  can be proved in a similar way.

The case of  $n = 1$  is evident. Now suppose that  $\beta \triangleright \tau_n \tau_{n-1} \dots \tau_1(a)$  and  $\alpha\tau_n \tau_{n-1} \dots \tau_1(a)\beta^{-1} \in \mathcal{P}$ .

If  $\tau_{n+1} = \phi$ , notice that the word  $\phi(\tau_n \tau_{n-1} \dots \tau_1(a))$  is obtained by exchanging the letters  $a$  and  $b$  in the word  $\tau_n \tau_{n-1} \dots \tau_1(a)$ . So by the induction hypothesis, we get  $\beta\phi(\tau_n \tau_{n-1} \dots \tau_1(a))\alpha^{-1} \in \mathcal{P}$ .

Now suppose that  $\tau_{n+1} = \sigma$ . Notice that  $\alpha \triangleright \sigma\tau_n \tau_{n-1} \dots \tau_1(a)$ . Thus

$$\begin{aligned} \beta\sigma\tau_n \tau_{n-1} \dots \tau_1(a)\alpha^{-1} &= \beta\sigma(\alpha^{-1}\alpha\tau_n \tau_{n-1} \dots \tau_1(a)\beta^{-1}\beta)\alpha^{-1} \\ &= \beta\sigma(\alpha^{-1})\sigma(\alpha\tau_n \tau_{n-1} \dots \tau_1(a)\beta^{-1})\sigma(\beta)\alpha^{-1}. \end{aligned}$$

By the induction hypothesis  $\alpha\tau_n \tau_{n-1} \dots \tau_1(a)\beta^{-1} \in \mathcal{P}$ . On the other hand, by a simple calculation, we have either  $\beta\sigma(\alpha^{-1}) = a^{-1}, \sigma(\beta)\alpha^{-1} = \epsilon$ , or  $\beta\sigma(\alpha^{-1}) = \epsilon, \sigma(\beta)\alpha^{-1} = a$ . Hence by lemma 2.7, we get  $\beta\sigma\tau_n \tau_{n-1} \dots \tau_1(a)\alpha^{-1} \in \mathcal{P}$ .

Let  $\tau(a) = \tau(b)u$  be the factorization of lemma 2.6. Let  $\tau = \sigma^{n_k} \phi \dots \sigma^{n_2} \phi \sigma^{n_1}$ , then  $u = \sigma^{n_k} \phi \dots \sigma^{n_2} \phi \sigma^{n_1-1}(b)$ . Thus by lemmas 2.1.4, 2.6 and corollary 2.8, we get

**Corollary 2.9.** Let  $\tau \in \mathcal{G}$  and let  $\alpha\beta \triangleright \tau(a)$ .

- 1). If  $|\tau(b)| = 1$ , then  $\tau(a) = a^n b$  for some  $n$ , thus  $\tau(a)$  is the product of palindromes  $a^n$  and  $b$ ;
- 2). If  $|\tau(b)| = 2$ , then  $\tau(a)$  is a palindrome;
- 3). If  $|\tau(b)| > 2$ , then  $\tau(a)$  is the product of palindromes  $\tau(b)\alpha^{-1}\beta^{-1}$  and  $\beta\alpha u$ .

**Lemma 2.10.** Suppose that  $w \in S^*$  is a primitive word and suppose that  $w = u_1 u_2$  is a product of two palindromes  $u_1$  and  $u_2$ .

- 1). If  $|u_1|, |u_2| \in 2N + 1$ , then for any  $k$ ,  $C_k(w) \notin \mathcal{P}$ ;
- 2). If  $|u_1|, |u_2| \in 2N$ , then  $C_{|u_1|/2}(w) \in \mathcal{P}$ ,  $C_{|u_1|+|u_2|/2}(w) \in \mathcal{P}$ . For other  $k$ ,  $C_k(w) \notin \mathcal{P}$ ;
- 3). If  $|u_1| \in 2N$ ,  $|u_2| \in 2N + 1$ , (or  $|u_1| \in 2N + 1$ ,  $|u_2| \in 2N$ ), then  $C_{|u_1|/2}(w) \in \mathcal{P}$ , (or  $C_{|u_1|+|u_2|/2}(w) \in \mathcal{P}$ .) For other  $k$ ,  $C_k(w) \notin \mathcal{P}$ .

**Proof.** We only prove 1), the other cases can be discussed in the same way.

Notice that at least one of  $|u_1|, |u_2|$  is larger than 2 (otherwise  $u_1 u_2$  will be  $a^2$  or  $b^2$ .) Suppose that without loosing generality  $|u_1| \geq 3$ . Then we can write  $u_1 = xu'_1x$ , where  $x \in S$ ,  $u'_1 \in \mathcal{P}$ . Thus  $C_1(w) = C_1(u_1 u_2) = (u'_1)(xu_2 x)$  is a product of two palindromes, by the proposition 9 of [3],  $C_1(u_1 u_2) \notin \mathcal{P}$ . Since  $|u'_1|, |xu_2 x| \in 2N + 1$  and  $u'_1, xu_2 x \in \mathcal{P}$ , we can repeat the discussion above which follows the proof .

**Lemma 2.11.** Suppose that  $|\tau(a)| \in 2N$ . If  $\tau(a) = u_1 u_2$  is a product of two palindromes, then  $|u_1|, |u_2| \in 2N + 1$ .

**Proof.** If  $|u_1|, |u_2| \in 2N$ , then  $|u_1|_a, |u_1|_b, |u_2|_a, |u_2|_b \in 2N$ . Thus both components of  $L(\tau(a))$  are even which follows that  $|\det(M_\tau)| \in 2N$ . But we know that  $\det(M_\tau)$  must be  $\pm 1$ .

From corollary 2.9, lemma 2.10, lemma 2.11 and the propositions 9,10,11 of [3], we obtain

**Proposition 2.12.** Let  $\tau \in \mathcal{G}$ .

- 1). Any conjugate of  $\tau(a)$  is primitive;
- 2).  $|C(\tau(a))| = |\tau(a)|$ . That is, all conjugates of  $\tau(a)$  are distinct;
- 3). Any conjugation of  $\tau(a)$  (containing  $\tau(a)$  itself) is either a palindrome, or a product of two palindromes. In the later case, the factorization is unique;
- 4).  $C(\tau(a)) = \overline{C(\tau(a))}$ , where  $\overline{C(\tau(a))} = \{\bar{w}; w \in C(\tau(a))\}$ .
- 5). If  $|\tau(a)| \in 2N$ , there is no palindrome in  $C(\tau(a))$ ; if  $|\tau(a)| \in 2N + 1$ , there is only one palindrome in  $C(\tau(a))$ .

Let  $F_\tau = \tau(F_\tau) = x_1 x_2 \dots x_k \dots$  be the fixed point of  $\tau$ . Then for any  $n \in N$ , we have

$$F_\tau = \tau^n(F_\tau) = \tau^n(x_1) \tau^n(x_2) \dots \tau^n(x_k) \dots \quad (**)$$

We now discuss the properties of the factors of the infinite word  $F_\tau$ .

We denote by  $A_n = \tau^n(a)$ ,  $B_n = \tau^n(b)$ , and  $l_{n,a} = |A_n|$ ,  $l_{n,b} = |B_n|$ .

Notice that if  $\tau \in \mathcal{G}$ , then for any  $n \geq 2$ ,  $\tau^n \in \mathcal{G}$ . Therefore the conclusions about  $\tau$  which we have obtained above hold also for  $\tau^n$ .

Let  $w$  be a factor of  $F_\tau$  of length  $|A_n|$ . By (\*\*),  $w$  will be contained in one of the following four words:  $A_n A_n, A_n B_n, B_n A_n, A_n B_n A_n$ . By lemma 2.13,  $B_n \triangleleft A_n$ , so  $A_n B_n \prec A_n A_n$ .

Let  $w \prec A_n B_n A_n$ , write  $w = w_1 B_n w_2$ . Then  $w_1 \triangleright A_n, w_2 \triangleleft A_n$ . Let  $A_n = B_n u_n$  be the factorization as in lemma 2.6. Then  $u_n \triangleleft A_n$  by lemma 2.6.3. Since

$|w| = |A_n| = |B_n| + |u_n|$ ,  $|w_2| < |u_n|$ . Consequently  $w_2 \triangleleft u_n$ , and hence  $w = w_1 B_n w_2 \prec A_n B_n u_n = A_n A_n$ .

By the discussions above, any factor of length  $|A_n|$  of  $F_\tau$  will be contained either in  $A_n A_n$  or  $B_n A_n$  which we are going to study respectively.

**Lemma 2.13.** Let  $w \prec A_n A_n$  with  $|w| = |A_n|$ , then  $w \in C(A_n)$ .

**Proof.** Since  $w \prec A_n A_n$ , and  $|w| = |A_n|$ , we can write  $w = uv$  with  $u \triangleright A_n$ , and  $|u| + |v| = |A_n|$ , which follows that  $A_n = vu$ . That is  $w = C_{|v|}(A_n)$ .

Since all conjugates of  $A_n$  are distinct by proposition 2.12.2, we get

**Corollay 2.14.** The set of the factors of length  $|A_n|$  of  $A_n^k, k \geq 2$  is exactly  $C(A_n)$ .

Now we study the factors of the word  $B_n A_n$ . For this aim, we introduce singular words of  $F_\tau$  as follows:

Let  $\alpha \triangleright A_n$  (so  $\beta \triangleright B_n$ ). Then the  $n$ th singular words with respect  $a$  or  $b$  of  $F_\tau$  are defined respectively as:

$$w_{n,a} = \beta A_n \alpha^{-1}, w_{n,b} = \alpha B_n \beta^{-1}.$$

From the definition of the singular words, we get immediately

**Lemma 2.15.** With the notations above, we have

$$L(w_{n,a}) = \begin{cases} (|A_n|_a + 1, |A_n|_b - 1) & \text{if } \alpha = b \\ (|A_n|_a - 1, |A_n|_b + 1) & \text{if } \alpha = a \end{cases}$$

$$L(w_{n,b}) = \begin{cases} (|B_n|_a + 1, |B_n|_b - 1) & \text{if } \alpha = a \\ (|B_n|_a - 1, |B_n|_b + 1) & \text{if } \alpha = b \end{cases}$$

Since  $L(w) = L(A_n)$  for any  $w \in C(A_n)$ , we get from lemma 2.15,

**Corollay 2.16.**  $\forall n \in N, w_{n,a} \notin C(A_n), w_{n,b} \notin C(B_n)$ .

Let  $\alpha\beta \triangleright A_n$ , (so  $\beta\alpha \triangleright B_n$ ) By lemma 2.6.3,

$$B_n A_n = B_n u_n B_n \alpha^{-1} \beta^{-1} \alpha \beta = A_n B_n \alpha^{-1} \beta^{-1} \alpha \beta,$$

Since  $A_n B_n \alpha^{-1} \beta^{-1} \prec A_n B_n \prec A_n A_n$ , we see that the first  $|B_n| - 2$  factors of length  $|B_n|$  of  $B_n A_n$  are those of  $A_n A_n$  which are distinct from each other by corollary 2.14, and the  $(|B_n| - 1)$ th factor is exactly  $w_{n,a}$  which is not in  $C(A_n)$  by corollary 2.16. We get from the analysis above

**Lemma 2.17.** Let  $w \prec B_n A_n$  with  $|w| = |A_n|$ . Then  $w$  is either in  $C(A_n)$  or equal to  $w_{n,a}$ . In particular, as a factor,  $w_{n,a}$  appears only once in  $B_n A_n$ .

From proposition 2.12.4, lemma 2.13, corollary 2.14, and lemma 2.17, we obtain

**Proposition 2.18.** Let  $\tau \in \mathcal{G}, n \in N$ . Then

- 1).  $\Omega_{l_{n,a}} = C(A_n(a)) \cup w_{n,a}, \Omega_{l_{n,b}} = C(B_n(b)) \cup w_{n,b}$ ;

- 2).  $|\Omega_{l_{n,a}}| = l_{n,a} + 1$ ,  $|\Omega_{l_{n,b}}| = l_{n,b} + 1$ ;  
3).  $\Omega_{l_{n,a}} = \overline{\Omega_{l_{n,a}}}$ ,  $\Omega_{l_{n,b}} = \overline{\Omega_{l_{n,b}}}$ .

We now discuss the properties of the singular words.

**Proposition 2.19.**  $\forall n \in \mathbb{N}$ ,  $w_{n,a}, w_{n,b} \in \mathcal{P}$ .

**Proof.** This follows immediately from the definition of the singular words and lemma 2.8.

**Lemma 2.20.**  $w_{n,a} \triangleright w_{n+2,a}$ ,  $w_{n,b} \triangleright w_{n+2,b}$ .

**Proof.** Let  $\alpha\beta \triangleright \tau^n(a)$ , then  $\alpha\beta \triangleright \tau^{n+2}(a)$  by lemma 2.1.4, thus

$$w_{n+2,a} = \alpha\tau^{n+2}(a)\beta^{-1} = \alpha\tau^n(\tau^2(a))\beta^{-1}.$$

By lemma 2.1.2,  $a \triangleleft \tau^2(a)$ , so  $\tau^2(a) = ua$  for some  $u \in S$ . Thus  $w_{n+2,a} = \alpha\tau^n(u)\tau^n(a)\beta^{-1}$  which yields the conclusion of the lemma.

**Lemma 2.21.** Either  $a^{p(\tau(a))+1} \triangleright w_{n,a}$ ,  $a^{p(\tau(a))+1} \triangleleft w_{n,a}$ , or  $b \triangleright w_{n,a}$ ,  $b \triangleleft w_{n,a}$ , where  $p(\tau(a))$  is defined as in lemma 2.3.2.

This follows from lemma 2.3.2, proposition 2.18 and the definition of  $w_{n,a}$ .

**Proposition 2.22.**  $\forall n \in \mathbb{N}$ , if  $w_{n,a}$  is not a power of  $a$ , then  $w_{n,a}$  is primitive.

**Proof.** Suppose that  $w_{n,a}$  is not a power of  $a$ . If  $w_{n,a} = w^p$  for some  $w \in S^*$  and  $p \geq 2$ , then  $w_{n,a}$  will be equal to one of its conjugates. But by lemma 2.21, we will then have either  $b^2 \prec F_\tau$  or  $a^{2p(\tau(a))+2} \prec F_\tau$ . This is impossible because of lemma 2.2.4.

Let  $F_\infty = x_1 \dots x_n \dots$  be an infinite word over  $S$ . Let  $w \prec F_\infty$  be a word which appears infinitely many times in  $F_\infty$ . If any two adjacent  $w$ 's appearing in  $F_\infty$  are separated by a factor of  $F_\infty$ , (this assertion is equivalent to  $\exists u \prec F_\infty$  such that  $u = xyz$  with  $w = xy = yz$ ) we say that the word  $w$  possesses the positive separation property, and the factor is called a separate factor (with respect to  $w$ ). If any separate factor is not equal to  $\epsilon$ , we say that  $w$  is strongly separated, (otherwise, we say that  $w$  is weakly separated.)

A word  $w \in S^*$  of the form  $w = (uv)^k u$ , where  $k > 0$ ,  $uv$  is a primitive word, is called a sesquipower [3]. The positive integer  $k$  is said to be the order of the sesquipower. A sesquipower of order larger than 1 is called a strong sesquipower.

Let  $n \geq 1$  and let  $\alpha\beta \triangleright A_n$ , from lemma 2.3, lemma 2.13, lemma 2.17, the formula (\*\*) and the definition of  $w_{n,a}$ , we can write  $F_\tau$  as

$$F_\tau = (A_n^{p(\tau(a))} B_n \alpha^{-1}) w_{n,a} z_1 w_{n,a} z_2 \dots z_k w_{n,a} z_{k+1} \dots$$

where  $z_j$  is either  $\beta A_n^{p(\tau(a))} B_n \alpha^{-1}$  or  $\beta A_n^{p(\tau(a))+1} B_n \alpha^{-1}$  and  $p(\tau(a))$  is defined as in lemma 2.2.2.

Let  $A_n = B_n u_n$  be the factorization as in lemma 2.6, then for any  $m \geq 1$ ,

$$\beta A_n^m B_n \alpha^{-1} = ((\beta B_n \alpha^{-1})(\alpha u_n \beta^{-1}))^m (\beta B_n \alpha^{-1})$$

by corollary 2.9,  $\beta B_n \alpha^{-1}, \alpha u_n \beta^{-1} \in \mathcal{P}$ . Thus  $\beta A_n^m B_n \alpha^{-1} \in \mathcal{P}$ .

From the discussions above, we obtain

**Theorem 2.23.** Let  $\tau \in \mathcal{G}$ , if  $\alpha\beta \triangleright \tau(a)$ , then for any  $n$ ,

- 1).  $w_{n,a}$  is strongly separated;
- 2). The separate factor is either  $\beta A_n^{p(\tau(a))+1} B_n \alpha^{-1}$ , or  $\beta A_n^{p(\tau(a))} B_n \alpha^{-1}$ ;
- 3). The separate factors are sesquipowers palindromes of order  $p(\tau(a))$  and  $p(\tau(a)) + 1$  respectively.

Berstel [2] introduced the special words of an infinite word  $F_\infty$  as follows : if  $ua, ub \prec F_\infty$ , then the word  $u$  is called a special word of  $F_\infty$ . Since  $\{w, w \prec F_\tau\} = \{\bar{w}, w \prec F_\tau\}$ , the special words of  $F_\tau$  are equivalent to  $au, bu \prec F_\tau$ .

From corollary 2.14, lemma 2.17, theorem 2.23, by an analogous argument with that the theorem 5 of [10], we have the following result which has been found by Séébold [9]. Proposition 4.7],(but he does not determine the special words).

**Theorem 2.24** Let  $\tau \in \mathcal{G}$ . Then for any  $n \geq 1$ , there is a unique special word of length  $n$ . Moreover  $w$  is a special word if and only if  $w \triangleleft F_\tau$ .

**Final remarks.** As said in the introduction, the motivation of this paper is to generalize the properties of the factors of the infinite Fibonacci word to the case of all invertible substitutions. By comparing with the results of [10], we see that one main difference is the structure of the singular words ( and the singular decomposition of the infinite word according to the singular words). In fact, the structure of the singular words of a general invertible substitution is much more complicated than that of the Fibonacci substitution (for example, all singular words and separate words of the Fibonacci word are sesquipowers of order 1, i.e, weak sesquipowers. On the other hand, those of a general invertible substitution are strong sesquipowers.) But as we shall see in another paper, the singular words still play an important role in the studies of the factors, such as power of factors, overlap of factors and local isomorphism.

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**SUR LES SYMÉTRIES TERNAIRES LIÉES  
AUX NOMBRES DE GENOCCHI**

Par

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**Abstract.** — In [Du2] DUMONT stated several conjectures about some symmetric polynomial sequences which are the refinements of the Genocchi numbers. In this paper we shall prove all of his conjectures. We first show that some special cases of his conjecture can be readily derived from a result of Wall and then prove this conjecture by computing Hankel determinants. Finally we present a new symmetric model for the Dumont-Foata polynomials.

**1. Introduction.** — Un escalier  $F$  de taille  $n \geq 1$  est le graphe d'une application surjective  $f$  de  $[2n] := \{1, 2, \dots, 2n\}$  sur  $\{2, 4, 6, \dots, 2n\}$  telle que pour tout  $k$ ,  $f(k) \geq k$ . Un point  $(k, f(k))$ ,  $1 \leq k \leq 2n - 2$ , de  $F$  est dit *maximal* si  $f(k) = 2n$ , est dit *point fixe* s'il n'est pas maximal et si  $f(k) = k$  (ce qui implique que  $k$  est pair), est dit *suffixe* s'il n'est pas maximal et si  $f(k) = k + 1$  (ce qui implique que  $k$  est impair). On note  $m(F)$  le nombre de ses points maximaux,  $f(F)$  le nombre de ses points fixes, et  $s(F)$  le nombre de ses points suffixes. Si l'on note  $E_n$  l'ensemble des escaliers de taille  $n$ , il est alors bien connu (voir [Du1, Du-Fo, Ha, Vi]) que le cardinal de  $E_n$  est le nombre de Genocchi  $G_{2n+2}$ , qui peuvent être définis par leur fonction génératrice :

$$\frac{2t}{e^t + 1} = t - \frac{t^2}{2!} + \frac{t^4}{4!} - 3\frac{t^6}{6!} + 17\frac{t^8}{8!} + \cdots + (-1)^n G_{2n} \frac{t^{2n}}{(2n)!} + \cdots$$

DUMONT et FOATA [Du-Fo] ont introduit une suite de polynômes  $F_n(x, y, z)$  qui sont définis par  $F_1(x, y, z) = 1$  et

$$F_n(x, y, z) = (x + y)(x + z)F_{n-1}(x + 1, y, z) - x^2 F_{n-1}(x, y, z).$$

Ils ont montré que ces polynômes sont symétriques dans les variables  $x, y, z$  et raffinent les nombres de Genocchi en ce sens que  $F_n(1, 1, 1) = G_{2n+2}$ . On connaît plusieurs interprétations combinatoires de ces polynômes (voir [Du1, Du-Fo, Ha, Vi]), mais il semble que celle la plus simple, en ce sens que la vérification de la récurrence est facile, est le résultat de HAN [Ha] suivant :

$$(1) \quad F_n(x, y, z) = \sum_{F \in E_n} x^{m(F)} y^{f(F)} z^{s(F)}.$$

De là, DUMONT [Du2] a récemment proposé une extension des polynômes  $F_n(x, y, z)$  de la manière suivante. Pour un escalier  $F$  de taille  $n$ , un point  $(k, f(k))$  de  $F$  est dit *doublé* s'il n'est pas seul sur sa ligne, c'est-à-dire s'il existe  $j \neq k$  tel que  $f(j) = f(k)$ . On note  $fdF$  (resp.  $fndF$ ) le nombre de ses points fixes doublés (resp. fixes non doublés),  $sdF$  (resp.  $sndF$ ) le nombre de ses points suffixes doublés (resp. suffixes