

$\sum a \text{ for } a \text{ in } L$   
 $\sum a \text{ for } a \text{ in } L$   
 $\text{if } a > 0$

## Section 6 Recursion

$\sum a \text{ if isinstance}(a, \text{int})$   
 $\text{else } \sum(a)$   
 $\text{for } a \text{ in } L$

### Recursion 三部曲：

1. Function description：根据docstring 搞清function的具体作用, input types和rtype. 在写body的时候 assume function已经写好啦
2. Base Case： 一般来说是各种数据类型的最小值, 比如integer的一般为0或1. String的一般是empty string, nested list的一般是非list或者empty list. 树可以是leaf or None.
3. Recursive step:
  - 3.1 Divide step： 搞清楚是如何把大问题化小, 是如何把当前问题分解成一个相对较小的问题.
  - 3.2 Combination step： 搞清楚是如何把子问题的结果组合成原问题的结果.

## Nested List recursion

Nested list recursion是我们在148接触的第一种recursion的考点, 每一个nested list的recursion问题简单来说可以通过使用isinstance来分为两种case:

1. 当前非list的情况.
- 2.. 当前依然是list的情况.

$\text{if not isinstance(obj, list):}$

常见 Recursion 问题中 可能存在filter condition或者来控制base case的结果. 对于这种类型, 我们基本上可以总结成只需要再base case中进行操作, 例如:

```
def count_odd(obj)
    """ Return number of odd element in obj.
```

$[\star, 1, 0]$

Assume: obj is an int or non-empty list with finite nesting depth, and obj doesn't contain any empty lists

```
"""
if not isinstance(obj, list):
    if obj % 2 == 1:
        return 1
    else:
        return 0
else:
```

```
    acc = 0
    for sub in obj:
```

$acc += \text{count\_odd}(\text{sub})$

$\text{return acc.}$

并且, 在recursion问题中, accumulation (combination)环节会根据题上不同要求存在不同的方法, 例如求depth, 每次是取sub problem 的max 加1,

def depth(obj):  
 """Return 0 if obj is a non-list, or 1 + maximum  
 depth of elements of obj, a possibly nested  
 list of objects.  
 if not isinstance(obj, list):  
 return 0.  
 else: acc = [0] # handle empty list,  
 for sub in obj:  
 sub.append(depth(sub)).  
 return max(acc) + 1  
def count\_lists(list\_):  
 """  
 Return the number of lists, including list\_ itself, contained  
 in list\_.  
 @param list list\_: list to count lists in  
 @rtype: int  
 >>> count\_lists([])  
 1  
 >>> count\_lists([5, [1, [2, 3], 4], 6])  
 3  
 """  
 if not isinstance(list\_, list):  
 return 0.  
 else:  
 acc = 1  
 for sub in obj  
 acc += count\_lists(sub)  
 return acc

在nested list recursion中, 也会出现对list进行modification的题型, 比如:

```
def add_one(obj):
    """Add one to every element in obj.
```

```
>>> L = [3, [7, 1], 5]
>>> add_one(L)
>>> L
[4, [8, 2], 6]
```

```
"""
if not isinstance(obj, list):
    for i in range(len(obj)):
        if not isinstance(obj[i], list):
            add_one(obj[i])
        else:
            obj[i] += 1
def concatenate_flat(list_):
```

1. join([a, b])  
Return the concatenation, from left to right, of strings contained in flat (depth 1) sublists contained in list\_, but no other strings.

'a-b'

Assume all non-list elements of list\_ or its nested sub-lists are strings

param list list\_: possibly nested sub-list to concatenate from  
rtype: str

```
>>> concatenate_flat(["five", [["four", "by"], "three"], ["two"]])
'fourbytwo'
```

```
"""
if helper(list_):
    return "-".join(list_)
else:
```

acc = "

for sub in list\_:

if not isinstance(sub, list):

acc += concatenate\_flat(sub)

return acc

def helper(L):
 for item in L:
 if not isinstance(item, list):
 return False

return True.

```

def width(list_, max_depth):
    """ Return the maximum number of items that occur at the same depth
    in list_ or its sub-lists combined. These could be list or non-list
    items. Elements may be lists or non-lists.

    @param list[list|object] list_: a possibly nested list
    @param max_depth int: maximum depth of list_
    @rtype: int
    """
    if not list_:
        return 0
    else:
        return max([width(item, max_depth - 1) for item in list_])

```

```

>>> list_ = [0, 1]
>>> width(list_, 1)
2
>>> list_ = [[0, 1], 2, [3, [[], 4]]]
>>> width(list_, 4)
4
>>> # 4 elements: 0, 1, 3, [[], 4]
"""

```

```

def helper(list_, max_depth):
    """

```

```

    if not isinstance(list_, list):
        return {0: 1}.
    """

```

```

    else:
        acc = {0: 1}.

```

for sub in list\_:

temp = helper(sub, max\_depth - 1).

for depth in temp:

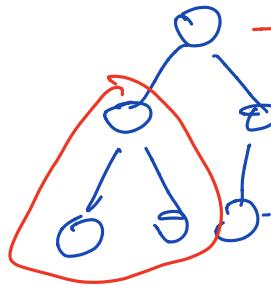
acc[depth + 1] = acc.get(depth + 1, 0) +

temp[depth].

return acc

return max(helper(list\_, max\_depth).values())

## Section 7 Tree



- root.  
 - leaf 等于 children  
 internal. 有 children  
 .height (node length). 3.  
 (Path length). 2.

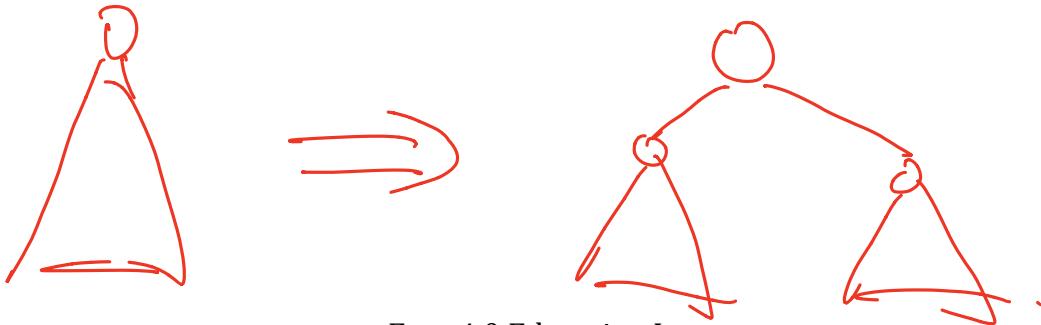
### Tree Node 的 implementation

```

class Tree:
    """
    A bare-bones Tree ADT that identifies the root with the entire tree.
    === Attributes ===
    @param object value: value of root node
    @param list[Tree|None] children: child nodes
    """
    def __init__(self, value=None, children=None):
        """
        Create Tree self with content value and 0 or more children
        @param Tree self: this tree
        @param object value: value contained in this tree
        @param list[Tree|None] children: possibly-empty list of children
        @rtype: None
        """
        self._value = value
        # copy children if not None
        # NEVER have a mutable default parameter...
        self._children = children[:] if children is not None else []
    
```

Tree, 一个倒过来的树, 每一个tree node都有 value和children, 常见的basecase就是在没有children的时候. 如果有children, 那么要对于每个children进行recursion. 这里要注意, 类似于linkedlist, 每一个node即是一个node, 同时也代表着整个subtree. 我们对一个tree object做recursion的时候, 我们可以分为两步:

1. 对node的value进行操作.
2. 对每一个child node都 recursively 操作.



Path 相关题型:

```
def get_longest_path(t: Tree) -> List[obj]:
    """Return a list contains values of the longest path from node t to the leaf.
```

```
>>> get_longest_path(Tree(5, [Tree(4, Tree(3)), Tree(2)]))
[5, 4, 3]
```

```
    """  
    if not t.children:  
        return [t.value]
```

```
    else:  
        result = []  
        for c in t.children:  
            path = get_longest_path(c).  
            if len(result) < len(path):  
                result = path[1]  
        return [t.value] + result.
```

```
def get_all_paths(t: Tree) -> List[obj]:
```

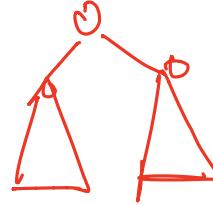
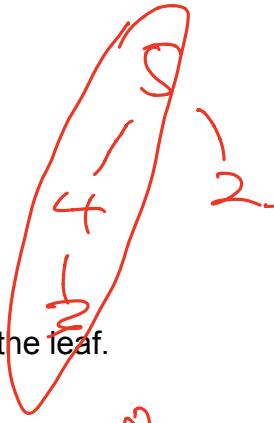
```
    """Return a list of lists contain values of the longest path from node t to all leaves.
```

```
>>> get_all_paths(Tree(5, [Tree(4, Tree(3)), Tree(2)]))
[[5, 4, 3], [5, 2]]
```

```
    """  
    if len(t.children) == 0:  
        return [[t.value]].
```

```
    else:  
        acc = []  
        for c in t.children:  
            for path in get_all_paths(c).  
                acc.append([t.value] + path)
```

```
    return acc
```



set()

∅

{0}.

add.

intersection  
difference  
union.

def pathlength\_sets(t):

"""

Replace the value of each node in Tree t by a set containing all path lengths from that node to any leaf. A path's length is the number of edges it contains.

param Tree t: tree to record path lengths in  
datatype: None

```
>>> t = Tree(5)
>>> pathlength_sets(t)
>>> print(t)
{0}
>>> t.children.append(Tree(17))
>>> t.children.append(Tree(13, [Tree(11)]))
>>> pathlength_sets(t)
>>> print(t)
{1, 2}
{0}
{1}
{0}
```

"""

if not t.children:

t.value = {0}.

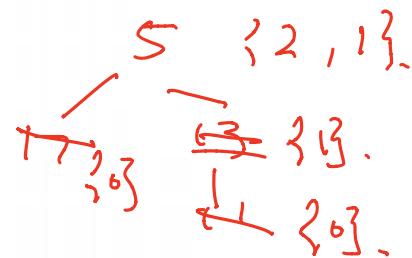
else:

acc = set()

for c in t.children:  
 pathlength\_sets(c).

for n in c.value:  
 acc.add(n+1)

t.value = acc



Read the declaration of class Tree and the docstring for function `unique_paths`. Then implement (write the body for) function `unique_paths` on the next page.

Hint: You may conclude there are unique paths in a tree if you traverse (visit) every node without finding a node that has been visited twice. A set is a convenient data structure for recording objects that have been seen.

```
class Tree:
    def __init__(self, value=None, children=None):
        ''' (Tree, object, list of Tree) -> NoneType

        Create Tree(self) with content value and 0 or more Tree children.
        '''
        self.value = value
        # copy children if not None
        self.children = children.copy() if children else []

def unique_paths(t):
    ''' (Tree) -> bool

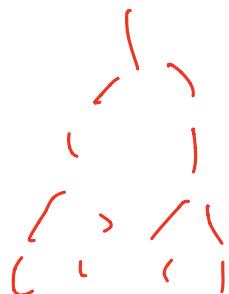
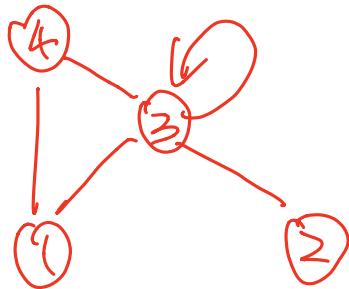
    Return whether there is a unique path from t to each
    of its descendants.

    Assume that two Trees are the same if they have the same
    memory address, that is id(t1) == id(t2)
    '''

>>> t1 = Tree(1)
>>> t2 = Tree(2)
>>> t3 = Tree(3, [t1, t2])
>>> unique_path(t3)
True
>>> t4 = Tree(4, [t3, t1])
>>> unique_path(t4)
False
>>> t3.children.append(t3)
>>> unique_path(t3)
False
'''
```

Part (a) [2 MARKS]

Draw the tree rooted at `t4`, as it is after all the doctest code has been executed.



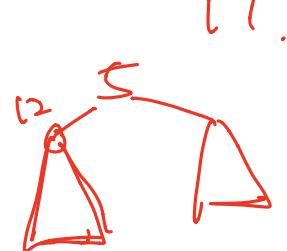
```

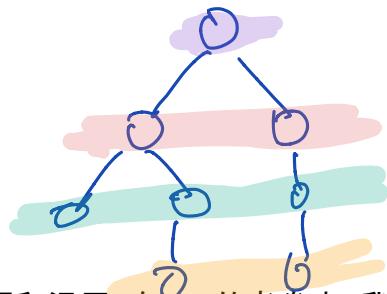
def unique_path(t):
    Seen = set()
    return helper(t)

def helper(t):
    if t in seen:
        return False
    if not t.children:
        Seen.add(t)
        return True
    else:
        Seen.add(t)
        for c in t.children:
            if not helper(c):
                return False
    return True

def has_path_sum(t, total):
    """Given binary tree, t return a boolean value
    indicating whether there is a path in t to any
    leaf whose values add up to int total. Precondition: t
    is not None, total is integer."""
    if not t.children:
        return t.value == total
    else:
        for c in t.children:
            if has_path_sum(c, total - t.value):
                return True
    return False

```





层层叠叠相关题:

在树的问题中, 我们经常会考察关于层数 depth / level 的相关理解和运用. 在148的考卷中, 我们经常默认root所在的是depth 0. 会出现两种常见题型:

1. 针对某一特定depth之上, 之下, 或者这一层进行操作.

```
def count_above_depth(t: Tree) > number:
```

"""Return the number of nodes above the depth in t.

```
>>> count_above_depth(Tree(5, [Tree(4, Tree(3)), Tree(2)]), 2)
```

3

.....

```
if d <= 0:
    return 0
```

```
else: acc = 1
```

```
for c in t.children:
    acc += count_above_depth(c, d-1).
```

```
return acc.
```

```
def count_below_depth(t: Tree) > number:
```

"""Return the number of nodes above the below in t.

```
>>> count_above_depth(Tree(5, [Tree(4, Tree(3)), Tree(2)]), 1)
```

1

.....

```
acc = 0.
```

```
if d < 0:
```

```
    acc += 1
```

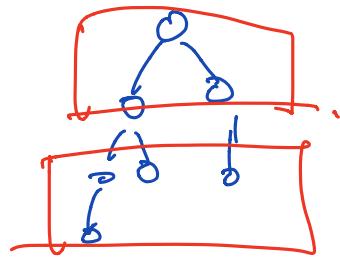
```
for c in t.children:
```

```
    acc += count_below_depth(c, d-1)
```

```
return acc
```

2. 针对奇数或者偶数层进行操作.

```
def count_on_odd_depth(t: Tree) -> number:
    """Return the number of nodes on odd depth.
```



```
>>> count_on_odd_depth(Tree(5, [Tree(4, Tree(3)), Tree(2)]), 1)
```

1

acc = len(t.children)

for c in t.children:

    for g in c.children

        acc += count\_on\_odd\_depth(g).

return acc

```
def duplicate_even_level(t: Tree) -> None:
```

"""Duplicate all nodes on even level by adding a new one after each of them.

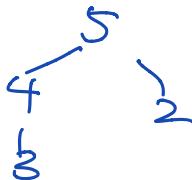
```
>>> t = Tree(5, [Tree(4, Tree(3)), Tree(2)])
```

```
>>> duplicate_even_level(t)
```

```
>>> t
```

```
Tree(5, [Tree(4, Tree(3)), Tree(4), Tree(2), Tree(2)])
```

.....



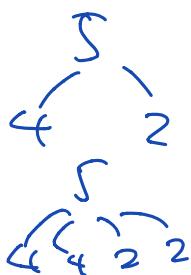
def helper(t, do):

for c in t.children:

helper(c, not do)

if do:

t.children = [Tree(t.value, t.children)]



helper(t, False)

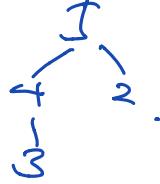
## 3. 针对每一层进行的操作

```
def count_by_level(t: Tree) -> List[int]:
    """Return a list that counts number of nodes on each level, index 0 is the root.
```

```
>>> t = Tree(5, [Tree(4, Tree(3)), Tree(2)])
>>> count_by_level(t)
[1, 2, 1]
```

```
"""
if not t.children:
    return [1]
```

```
else: acc = [1]
    for c in t.children:
        temp = count_by_level(c)
        for i in range(len(temp)):
            if i+1 >= len(acc):
                acc.append(temp[i])
            else:
                acc[i+1] += temp[i]
```



Modification相关题: `return acc`

```
def extend_every_third_level(t) -> None:
    """ Add a new node after the node on every third level.
```

```
>>> t = Tree(5, [Tree(4, Tree(3, Tree(7))), Tree(2)])
>>> extend_every_third_level(t)
>>> t
Tree(5, [Tree(5, Tree(4, Tree(3, Tree(7, [Tree(7)])))), Tree(2))])
```

```
"""
def helper(t, level):
    for c in t.children:
        helper(c, level+1)
    if level % 3 == 0:
        t.children = [Tree(t.value, t.children)]
```

`helper(t, 1).`

```

def prune(t, predicate):
    """ Return a new tree with the same values as t, except it prunes (omits) all paths of t that start
    with nodes where predicate(node.value) == False. If predicate(t.value) == False, then prune
    returns None.

    Assume that all values in t are ints, and that predicate always takes an int and returns a bool.

    @param Tree t: tree to prune
    @param function[int, bool] predicate: predicate on t's values.
    @rtype: Tree|None

    >>> t1 = Tree(6, [Tree(8), Tree(9)])
    >>> t2 = Tree(4, [Tree(11), Tree(10)])
    >>> t3 = Tree(7, [Tree(3), Tree(12)])
    >>> t = Tree(5, [t1, t2, t3])
    >>> t3_pruned = Tree(7, [Tree(12)])
    >>> def predicate(v): return v > 4
    >>> prune(t, predicate) == Tree(5, [t1, t3_pruned])
    True
    """
    if not predicate(t.value):
        return None.

    else:
        acc = []
        for c in t.children:
            new_c = prune(c)
            if new_c:
                acc.append(new_c)

        return Tree(t.value, acc)

```

## Section 8 Iterative instead of Recursion

我们可以通过使用我们在148中学过的Stack 和 Queue 两种不同的ADT 来对一个Tree进行操作, 在这里要注意两件事情, Stack 会得到一个按照path 顺序进行 recursion的结果, 而 Queue会得到一个按照level 顺序 得到的结果.

def mark\_tree\_by\_depth(t) -> None:

""" Mark the nodes in the tree using numbers by depth

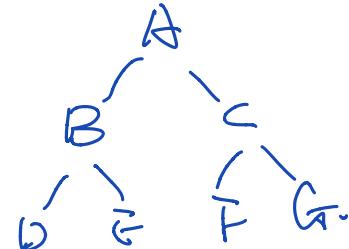
```
>>> t = Tree('A', [Tree('B', Tree('C')), Tree('D')])  
>>> mark_tree_by_depth(t)  
>>> t  
Tree(1, [Tree(3, Tree(4)), Tree(2)])
```

"""  
s=Stack()

num = 1  
s.add(t)

while not s.is\_empty():

temp = s.remove()  
temp.value = num  
num = num + 1



for c in temp.children[::-1]  
s.add(c).

def mark\_tree\_by\_level(t) -> None:

""" Mark the nodes in the tree using numbers by level

```
>>> t = Tree('A', [Tree('B', Tree('C')), Tree('D')])  
>>> mark_tree_by_depth(t)  
>>> t  
Tree(1, [Tree(2, Tree(4)), Tree(3)])
```

"""  
q=Queue()

q.add(t)

num =

while not q.is\_empty():

temp = q.remove()

temp.value = num

num = num + 1

for c in temp.children:

q.add(c)

## Section 4. Binary Tree

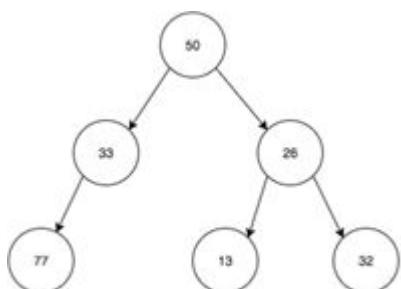
Binary Tree, 顾名思义, arity最大只有2的Tree, 但是, 我们的implementation有改变, 不再是node, 变成了 left和right. ( 虽然我们还是可以 for child in [t.left, t.right]).

### 需要Check None

```
class BinaryTree:
    """
    A Binary Tree, i.e. arity 2.
    """
    def __init__(self, value, left=None, right=None):
        """
        Create BinaryTree self with value and children left and right.
        @param BinaryTree self: this binary tree
        @param object value: value of this node
        @param BinaryTree|None left: left child
        @param BinaryTree|None right: right child
        @rtype: None
        """
        self.value, self.left, self.right = value, left, right
```

### Traversals:

我们一共接触了三种traversal, pre-order, in-order和post-order. 这里pre, in和post可以理解为 root放的位置. Pre-order也就是先root再children, in-order为root在左右children中间, post-order为root在最后. 注意,



Pre-order: 50 33 77 26 13 32.

Post-order: 77 33 13 32 26 50.

In-order: 77 33 50 13 26 32.

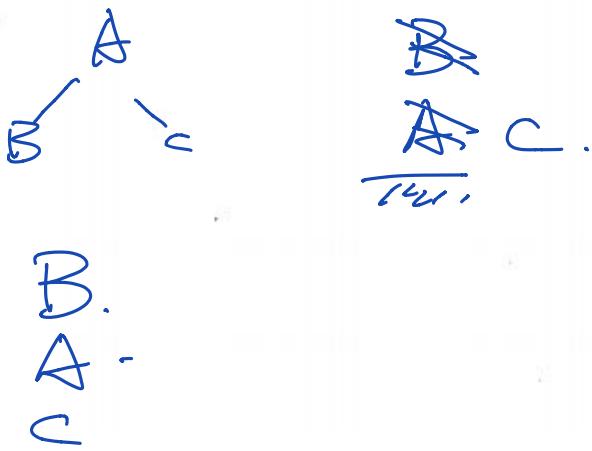
```

def inorder(node):
    if not node:
        return []
    else:
        return inorder(node.left) +
               [node.value] +
               inorder(node.right),
  
```

Here is some stack-based code that implements one of the tree traversals. Which traversal is this? Trace the code on a sample tree to support your choice.

```

def mystery_order(bt: BTNode) -> None:
    s = Stack()
    current = bt
    while current or not s.is_empty():
        if current:
            s.push(current)
            current = current.left
        else:
            current = s.pop()
            print(current.item)
            current = current.right
  
```



I wonder.

```

def occurs(root, s):
    ''' (BTNode or None, str) -> bool

    Return whether or not s equals a sequence of characters
    along some path from the root to a leaf, inclusive, and in
    that order. The empty str ("") is considered to occur in
    the empty tree, denoted None.
    
```

Assume each node in the tree rooted at root contains a str of length 1.

```

>>> left = BTNode('b', None, BTNode('d', BTNode('e'), None))
>>> right = BTNode('c', BTNode('e'), BTNode('f', BTNode('h'), BTNode('i'))))
>>> whole = BTNode('a', left, right)
>>> occurs(whole, 'acfhi')
True
>>> occurs(whole, 'ace')
True
>>> occurs(whole, 'bde')
False
...
if not root:
    return s == ""
else:
    if root.value == s[0]:
        return occurs(root.left, s[1:]) or
               occurs(root.right, s[1:])
    return False.

```

A perfect binary tree is one in which (a) every non-leaf node has exactly two children, and (b) all leaves occur at the same level. Notice that a tree consisting of a single node is a perfect binary tree of height 1. An empty tree, represented by `None`, has is a perfect binary tree of height 0.

**Part (a) [3 MARKS]**

Recall that we defined height of a tree in such a way that a tree containing just a root node has height one. Draw a perfect binary tree of height 3.

Draw a binary tree of height 4 this time (not 3) that is not perfect because it satisfies condition (a) but not condition (b).

Draw a binary tree of height 4 that is not perfect because it satisfies condition (b) but not condition (a).

Write the body of function `TPBT`.

```
def TPBT(root):
    ''' (BTNode or None) -> (int, bool)

    Return a tuple containing: (1) the height of the tallest perfect binary
    tree within the tree rooted at root, and (2) whether or not that tallest
    perfect binary tree occurs at the root itself.
    '''
```