# ld.so(8) — Linux manual page

# NAME | SYNOPSIS | DESCRIPTION | OPTIONS | ENVIRONMENT | FILES | NOTES | SEE ALSO | COLOPHON

Search online pages

Ld.so(8) System Manager's Manual

Ld.so(8)

# NAME top

ld.so, ld-linux.so - dynamic linker/loader

#### SYNOPSIS top

The dynamic linker can be run either indirectly by running some dynamically linked program or shared object (in which case no command-line options to the dynamic linker can be passed and, in the ELF case, the dynamic linker which is stored in the .interp section of the program is executed) or directly by running:

/lib/ld-linux.so.\* [OPTIONS] [PROGRAM [ARGUMENTS]]

# **DESCRIPTION** top

The programs **ld.so** and **ld-linux.so\*** find and load the shared objects (shared libraries) needed by a program, prepare the program to run, and then run it.

Linux binaries require dynamic linking (linking at run time) unless the **-static** option was given to ld(1) during compilation.

The program <code>ld.so</code> handles a.out binaries, a binary format used long ago. The program <code>ld-linux.so\*</code> (/lib/ld-linux.so.1 for libc5, /lib/ld-linux.so.2 for glibc2) handles binaries that are in the more modern ELF format. Both programs have the same behavior, and use the same support files and programs (<code>ldd(1)</code>, <code>ldconfig(8)</code>, and /etc/ld.so.conf).

When resolving shared object dependencies, the dynamic linker first inspects each dependency string to see if it contains a slash (this can occur if a shared object pathname containing slashes was specified at link time). If a slash is found, then the dependency string is interpreted as a (relative or absolute) pathname, and the shared object is loaded using that pathname.

If a shared object dependency does not contain a slash, then it is searched for in the following order:

- (1) Using the directories specified in the DT\_RPATH dynamic section attribute of the binary if present and DT\_RUNPATH attribute does not exist.
- (2) Using the environment variable LD\_LIBRARY\_PATH, unless the executable is being run in secure-execution mode (see below), in which case this variable is ignored.
- (3) Using the directories specified in the DT\_RUNPATH dynamic section attribute of the binary if present. Such directories are searched only to find those objects required by DT\_NEEDED (direct dependencies) entries and do not apply to those objects' children, which must themselves have their own DT\_RUNPATH entries. This is unlike DT\_RPATH, which is applied to searches for all children in the dependency tree.
- (4) From the cache file /etc/ld.so.cache, which contains a compiled list of candidate shared objects previously found in the augmented library path. If, however, the binary was linked with the -z nodefaultlib linker option, shared objects in the default paths are skipped. Shared objects

installed in hardware capability directories (see below) are preferred to other shared objects.

(5) In the default path /lib, and then /usr/lib. (On some 64-bit architectures, the default paths for 64-bit shared objects are /lib64, and then /usr/lib64.) If the binary was linked with the -z nodefaultlib linker option, this step is skipped.

#### Dynamic string tokens

In several places, the dynamic linker expands dynamic string tokens:

- In the environment variables LD\_LIBRARY\_PATH, LD\_PRELOAD, and LD\_AUDIT,
- inside the values of the dynamic section tags DT\_NEEDED, DT\_RPATH, DT\_RUNPATH, DT\_AUDIT, and DT\_DEPAUDIT of ELF binaries,
- in the arguments to the ld.so command line options --audit,
  --library-path, and --preload (see below), and
- in the filename arguments to the dlopen(3) and dlmopen(3) functions.

The substituted tokens are as follows:

#### \$ORIGIN (or equivalently \${ORIGIN})

This expands to the directory containing the program or shared object. Thus, an application located in <code>somedir/app</code> could be compiled with

```
gcc -Wl,-rpath,'$ORIGIN/../lib'
```

so that it finds an associated shared object in <code>somedir/lib</code> no matter where <code>somedir</code> is located in the directory hierarchy. This facilitates the creation of "turn-key" applications that do not need to be installed into special directories, but can instead be unpacked into any directory and still find their own shared objects.

# \$LIB (or equivalently \${LIB})

This expands to lib or lib64 depending on the architecture (e.g., on x86-64, it expands to lib64 and on x86-32, it expands to lib).

#### \$PLATFORM (or equivalently \${PLATFORM})

This expands to a string corresponding to the processor type of the host system (e.g., "x86\_64"). On some architectures, the Linux kernel doesn't provide a platform string to the dynamic linker. The value of this string is taken from the AT\_PLATFORM value in the auxiliary vector (see getauxval(3)).

Note that the dynamic string tokens have to be quoted properly when set from a shell, to prevent their expansion as shell or environment variables.

#### OPTIONS top

#### --argv0 string (since glibc 2.33)

Set  $argv[\theta]$  to the value string before running the program.

#### --audit list

Use objects named in *list* as auditors. The objects in *list* are delimited by colons.

# --glibc-hwcaps-mask list

only search built-in subdirectories if in *list*.

# --glibc-hwcaps-prepend list

Search glibc-hwcaps subdirectories in *list*.

# --inhibit-cache

Do not use /etc/ld.so.cache.

#### --library-path path

Use path instead of LD\_LIBRARY\_PATH environment variable setting (see below). The names ORIGIN, LIB, and PLATFORM are interpreted as for the LD\_LIBRARY\_PATH environment variable.

#### --inhibit-rpath list

Ignore RPATH and RUNPATH information in object names in *list*. This option is ignored when running in secure-execution mode (see below). The objects in *list* are delimited by colons or spaces.

--list List all dependencies and how they are resolved.

#### --list-diagnostics (since glibc 2.33)

Print system diagnostic information in a machine-readable format, such as some internal loader variables, the auxiliary vector (see getauxval(3)), and the environment variables. On some architectures, the command might print additional information (like the cpu features used in GNU indirect function selection on x86). --list-tunables (since glibc 2.33) Print the names and values of all tunables, along with the minimum and maximum allowed values.

# --preload list (since glibc 2.30)

Preload the objects specified in *list*. The objects in *list* are delimited by colons or spaces. The objects are preloaded as explained in the description of the **LD PRELOAD** environment variable below.

By contrast with LD\_PRELOAD, the --preload option provides a way to perform preloading for a single executable without affecting preloading performed in any child process that executes a new program.

#### --verify

Verify that program is dynamically linked and this dynamic linker can handle it.

# ENVIRONMENT top

Various environment variables influence the operation of the dynamic linker.

#### Secure-execution mode

For security reasons, if the dynamic linker determines that a binary should be run in secure-execution mode, the effects of some environment variables are voided or modified, and furthermore those environment variables are stripped from the environment, so that the program does not even see the definitions. Some of these environment variables affect the operation of the dynamic linker itself, and are described below. Other environment variables treated in this way include:

GCONV\_PATH, GETCONF\_DIR, HOSTALIASES, LOCALDOMAIN, LD\_AUDIT, LD\_DEBUG, LD\_DEBUG\_OUTPUT, LD\_DYNAMIC\_WEAK, LD\_HWCAP\_MASK, LD\_LIBRARY\_PATH, LD\_ORIGIN\_PATH, LD\_PRELOAD, LD\_PROFILE, LD\_SHOW\_AUXV, LOCALDOMAIN, LOCPATH, MALLOC\_TRACE, NIS\_PATH, NLSPATH, RESOLV\_HOST\_CONF, RES\_OPTIONS, TMPDIR, and TZDIR.

A binary is executed in secure-execution mode if the AT\_SECURE entry in the auxiliary vector (see getauxval(3)) has a nonzero value. This entry may have a nonzero value for various reasons, including:

- The process's real and effective user IDs differ, or the real and effective group IDs differ. This typically occurs as a result of executing a set-user-ID or set-group-ID program.
- A process with a non-root user ID executed a binary that conferred capabilities to the process.
- A nonzero value may have been set by a Linux Security Module.

# **Environment variables**

Among the more important environment variables are the following:

#### LD\_ASSUME\_KERNEL (from glibc 2.2.3 to glibc 2.36)

Each shared object can inform the dynamic linker of the minimum kernel ABI version that it requires. (This requirement is encoded in an ELF note section that is viewable via readelf -n as a section labeled NT\_GNU\_ABI\_TAG.) At run time, the dynamic linker determines the ABI version of the running kernel and will reject loading shared objects that specify minimum ABI versions that exceed that ABI version.

LD\_ASSUME\_KERNEL can be used to cause the dynamic linker to assume that it is running on a system with a different kernel ABI version. For example, the following command line causes the dynamic linker to assume it is running on Linux 2.2.5 when loading the shared objects required by myprog:

#### \$ LD\_ASSUME\_KERNEL=2.2.5 ./myprog

On systems that provide multiple versions of a shared object (in different directories in the search path) that have different minimum kernel ABI version requirements, LD\_ASSUME\_KERNEL can be used to select the version of the object that is used (dependent on the directory search order).

Historically, the most common use of the LD\_ASSUME\_KERNEL feature was to manually select the older LinuxThreads POSIX threads implementation on systems that provided both LinuxThreads and NPTL (which latter was typically the default on such systems); see pthreads(7).

#### LD\_BIND\_NOW (since glibc 2.1.1)

If set to a nonempty string, causes the dynamic linker to resolve all symbols at program startup instead of deferring function call resolution to the point when they are first referenced. This is useful when using a debugger.

# LD\_LIBRARY\_PATH

A list of directories in which to search for ELF libraries at execution time. The items in the list are separated by either colons or semicolons, and there is no support for escaping either separator. A zero-length directory name indicates the current working directory.

This variable is ignored in secure-execution mode.

Within the pathnames specified in LD\_LIBRARY\_PATH, the dynamic linker expands the tokens \$ORIGIN, \$LIB, and \$PLATFORM (or the versions using curly braces around the names) as described above in Dynamic string tokens. Thus, for example, the following would cause a library to be searched for in either the lib or lib64 subdirectory below the directory containing the program to be executed:

#### \$ LD\_LIBRARY\_PATH='\$ORIGIN/\$LIB' prog

(Note the use of single quotes, which prevent expansion of *\$ORIGIN* and *\$LIB* as shell variables!)

#### LD\_PRELOAD

A list of additional, user-specified, ELF shared objects to be loaded before all others. This feature can be used to selectively override functions in other shared objects.

The items of the list can be separated by spaces or colons, and there is no support for escaping either separator. The objects are searched for using the rules given under DESCRIPTION. Objects are searched for and added to the link map in the left-to-right order specified in the list.

In secure-execution mode, preload pathnames containing slashes are ignored. Furthermore, shared objects are preloaded only from the standard search directories and only if they have set-user-ID mode bit enabled (which is

not typical).

Within the names specified in the LD\_PRELOAD list, the dynamic linker understands the tokens \$ORIGIN, \$LIB, and \$PLATFORM (or the versions using curly braces around the names) as described above in Dynamic string tokens. (See also the discussion of quoting under the description of LD\_LIBRARY\_PATH.)

There are various methods of specifying libraries to be preloaded, and these are handled in the following order:

- (1) The LD\_PRELOAD environment variable.
- (2) The **--preload** command-line option when invoking the dynamic linker directly.
- (3) The /etc/ld.so.preload file (described below).

#### LD\_TRACE\_LOADED\_OBJECTS

If set (to any value), causes the program to list its dynamic dependencies, as if run by ldd(1), instead of running normally.

Then there are lots of more or less obscure variables, many obsolete or only for internal use.

#### LD\_AUDIT (since glibc 2.4)

A list of user-specified, ELF shared objects to be loaded before all others in a separate linker namespace (i.e., one that does not intrude upon the normal symbol bindings that would occur in the process) These objects can be used to audit the operation of the dynamic linker. The items in the list are colon-separated, and there is no support for escaping the separator.

LD\_AUDIT is ignored in secure-execution mode.

The dynamic linker will notify the audit shared objects at so-called auditing checkpoints—for example, loading a new shared object, resolving a symbol, or calling a symbol from another shared object—by calling an appropriate function within the audit shared object. For details, see rtld-audit(7). The auditing interface is largely compatible with that provided on Solaris, as described in its Linker and Libraries Guide, in the chapter Runtime Linker Auditing Interface.

Within the names specified in the LD\_AUDIT list, the dynamic linker understands the tokens \$ORIGIN, \$LIB, and \$PLATFORM (or the versions using curly braces around the names) as described above in *Dynamic string tokens*. (See also the discussion of quoting under the description of LD\_LIBRARY\_PATH.)

Since glibc 2.13, in secure-execution mode, names in the audit list that contain slashes are ignored, and only shared objects in the standard search directories that have the set-user-ID mode bit enabled are loaded.

#### LD\_BIND\_NOT (since glibc 2.1.95)

If this environment variable is set to a nonempty string, do not update the GOT (global offset table) and PLT (procedure linkage table) after resolving a function symbol. By combining the use of this variable with **LD\_DEBUG** (with the categories *bindings* and *symbols*), one can observe all run-time function bindings.

#### LD\_DEBUG (since glibc 2.1)

Output verbose debugging information about operation of the dynamic linker. The content of this variable is one of more of the following categories, separated by colons, commas, or (if the value is quoted) spaces:

help Specifying help in the value of this variable does not run the specified program, and displays a help message about which categories can be specified in this environment variable.

all Print all debugging information (except statistics and unused; see below).

#### bindings

Display information about which definition each symbol is bound to.

files Display progress for input file.

libs Display library search paths.

reloc Display relocation processing.

scopes Display scope information.

statistics

Display relocation statistics.

symbols

Display search paths for each symbol look-up.

unused Determine unused DSOs.

versions

Display version dependencies.

Since glibc 2.3.4, **LD\_DEBUG** is ignored in secure-execution mode, unless the file /etc/suid-debug exists (the content of the file is irrelevant).

#### LD\_DEBUG\_OUTPUT (since glibc 2.1)

By default, LD\_DEBUG output is written to standard error. If LD\_DEBUG\_OUTPUT is defined, then output is written to the pathname specified by its value, with the suffix "." (dot) followed by the process ID appended to the pathname.

LD\_DEBUG\_OUTPUT is ignored in secure-execution mode.

# LD\_DYNAMIC\_WEAK (since glibc 2.1.91)

By default, when searching shared libraries to resolve a symbol reference, the dynamic linker will resolve to the first definition it finds.

Old glibc versions (before glibc 2.2), provided a different behavior: if the linker found a symbol that was weak, it would remember that symbol and keep searching in the remaining shared libraries. If it subsequently found a strong definition of the same symbol, then it would instead use that definition. (If no further symbol was found, then the dynamic linker would use the weak symbol that it initially found.)

The old glibc behavior was nonstandard. (Standard practice is that the distinction between weak and strong symbols should have effect only at static link time.) In glibc 2.2, the dynamic linker was modified to provide the current behavior (which was the behavior that was provided by most other implementations at that time).

Defining the LD\_DYNAMIC\_WEAK environment variable (with any value) provides the old (nonstandard) glibc behavior, whereby a weak symbol in one shared library may be overridden by a strong symbol subsequently discovered in another shared library. (Note that even when this variable is set, a strong symbol in a shared library will not override a weak definition of the same symbol in the main program.)

Since glibc 2.3.4, LD\_DYNAMIC\_WEAK is ignored in secure-execution mode.

#### LD\_HWCAP\_MASK (from glibc 2.1 to glibc 2.38)

Mask for hardware capabilities. Since glibc 2.26, the option might be ignored if glibc does not support tunables.

# LD\_ORIGIN\_PATH (since glibc 2.1)

Path where the binary is found.

Since glibc 2.4, LD\_ORIGIN\_PATH is ignored in secure-execution mode.

#### LD\_POINTER\_GUARD (from glibc 2.4 to glibc 2.22)

Set to 0 to disable pointer guarding. Any other value enables pointer guarding, which is also the default. Pointer guarding is a security mechanism whereby some pointers to code stored in writable program memory (return addresses saved by setjmp(3) or function pointers used by various glibc internals) are mangled semi-randomly to make it more difficult for an attacker to hijack the pointers for use in the event of a buffer overrun or stack-smashing attack. Since glibc 2.23, LD\_POINTER\_GUARD can no longer be used to disable pointer guarding, which is now always enabled.

#### LD\_PROFILE (since glibc 2.1)

The name of a (single) shared object to be profiled, specified either as a pathname or a soname. Profiling output is appended to the file whose name is: \$LD\_PROFILE\_OUTPUT/\$LD\_PROFILE.profile.

Since glibc 2.2.5, **LD\_PROFILE** uses a different default path in secure-execution mode.

#### LD\_PROFILE\_OUTPUT (since glibc 2.1)

Directory where **LD\_PROFILE** output should be written. If this variable is not defined, or is defined as an empty string, then the default is /var/tmp.

LD\_PROFILE\_OUTPUT is ignored in secure-execution mode; instead /var/profile is always used.

#### LD\_SHOW\_AUXV (since glibc 2.1)

If this environment variable is defined (with any value), show the auxiliary array passed up from the kernel (see also getauxval(3)).

Since glibc 2.3.4, LD\_SHOW\_AUXV is ignored in secure-execution mode.

# LD\_TRACE\_PRELINKING (from glibc 2.4 to glibc 2.35)

If this environment variable is defined, trace prelinking of the object whose name is assigned to this environment variable. (Use ldd(1) to get a list of the objects that might be traced.) If the object name is not recognized, then all prelinking activity is traced.

#### LD\_USE\_LOAD\_BIAS (from glibc 2.3.3 to glibc 2.35)

By default (i.e., if this variable is not defined), executables and prelinked shared objects will honor base addresses of their dependent shared objects and (nonprelinked) position-independent executables (PIEs) and other shared objects will not honor them. If LD\_USE\_LOAD\_BIAS is defined with the value 1, both executables and PIEs will honor the base addresses. If LD\_USE\_LOAD\_BIAS is defined with the value 0, neither executables nor PIEs will honor the base addresses.

Since glibc 2.3.3, this variable is ignored in secureexecution mode.

#### LD\_VERBOSE (since glibc 2.1)

If set to a nonempty string, output symbol versioning information about the program if the LD\_TRACE\_LOADED\_OBJECTS environment variable has been set.

# LD\_WARN (since glibc 2.1.3)

If set to a nonempty string, warn about unresolved symbols.

# LD\_PREFER\_MAP\_32BIT\_EXEC (x86-64 only; since glibc 2.23) According to the Intel Silvermont software optimization guide, for 64-bit applications, branch prediction performance can be negatively impacted when the target of a branch is more than 4 GB away from the branch. If this

environment variable is set (to any value), the dynamic linker will first try to map executable pages using the mmap(2) MAP\_32BIT flag, and fall back to mapping without that flag if that attempt fails. NB: MAP\_32BIT will map to the low 2 GB (not 4 GB) of the address space.

Because MAP\_32BIT reduces the address range available for address space layout randomization (ASLR), LD\_PREFER\_MAP\_32BIT\_EXEC is always disabled in secure-execution mode.

#### FILES top

/lib/ld.so

a.out dynamic linker/loader

/lib/ld-linux.so.{1,2}

ELF dynamic linker/loader

/etc/ld.so.cache

File containing a compiled list of directories in which to search for shared objects and an ordered list of candidate shared objects. See ldconfig(8).

/etc/ld.so.preload

File containing a whitespace-separated list of ELF shared objects to be loaded before the program. See the discussion of LD\_PRELOAD above. If both LD\_PRELOAD and /etc/ld.so.preload are employed, the libraries specified by LD\_PRELOAD are preloaded first. /etc/ld.so.preload has a system-wide effect, causing the specified libraries to be preloaded for all programs that are executed on the system. (This is usually undesirable, and is typically employed only as an emergency remedy, for example, as a temporary workaround to a library misconfiguration issue.)

lib\*.so\*

shared objects

# NOTES to

# Legacy Hardware capabilities (from glibc 2.5 to glibc 2.37)

Some shared objects are compiled using hardware-specific instructions which do not exist on every CPU. Such objects should be installed in directories whose names define the required hardware capabilities, such as /usr/lib/sse2/. The dynamic linker checks these directories against the hardware of the machine and selects the most suitable version of a given shared object. Hardware capability directories can be cascaded to combine CPU features. The list of supported hardware capability names depends on the CPU. The following names are currently recognized:

Alpha ev4, ev5, ev56, ev6, ev67

MIPS loongson2e, loongson2f, octeon, octeon2

#### PowerPC

4xxmac, altivec, arch\_2\_05, arch\_2\_06, booke, cellbe, dfp, efpdouble, efpsingle, fpu, ic\_snoop, mmu, notb, pa6t, power4, power5, power5+, power6x, ppc32, ppc601, ppc64, smt, spe, ucache, vsx

SPARC flush, muldiv, stbar, swap, ultra3, v9, v9v, v9v2

**s390** dfp, eimm, esan3, etf3enh, g5, highgprs, hpage, ldisp, msa, stfle, z900, z990, z9-109, z10, zarch

# x86 (32-bit only)

acpi, apic, clflush, cmov, cx8, dts, fxsr, ht, i386, i486, i586, i686, mca, mmx, mtrr, pat, pbe, pge, pn, pse36, sep, ss, sse, sse2, tm

The legacy hardware capabilities support has the drawback that each new feature added grows the search path exponentially, because it has to be added to every combination of the other

```
existing features.
       For instance, on x86 32-bit, if the hardware supports i686 and
       sse2, the resulting search path will be i686/sse2:i686:sse2:.. A
       new capability newcap will set the search path to
       newcap/i686/sse2:newcap/i686:newcap/sse2:newcap:i686/sse2:i686:sse2:.
   glibc Hardware capabilities (from glibc 2.33)
       glibc 2.33 added a new hardware capability scheme,
              where under each CPU architecture, certain levels can be
              defined, grouping support for certain features or special
              instructions. Each architecture level has a fixed set of
              paths that it adds to the dynamic linker search list,
              depending on the hardware of the machine. Since each new
              architecture level is not combined with previously
              existing ones, the new scheme does not have the drawback
              of growing the dynamic linker search list uncontrollably.
       For instance, on x86 64-bit, if the hardware supports x86_64-v3
       (for instance Intel Haswell or AMD Excavator), the resulting
       search path will be glibc-hwcaps/x86-64-v3:glibc-
       hwcaps/x86-64-v2:. The following paths are currently supported,
       in priority order.
       PowerPC (64-bit little-endian only)
              power10, power9
       s390 (64-bit only)
              z16, z15, z14, z13
       x86 (64-bit only)
              x86-64-v4, x86-64-v3, x86-64-v2
       glibc 2.37 removed support for the legacy hardware capabilities.
SEE ALSO
       ld(1), ldd(1), pldd(1), sprof(1), dlopen(3), getauxval(3),
       elf(5), capabilities(7), rtld-audit(7), ldconfig(8), sln(8)
COLOPHON
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PR SET FP MODE(2const), PR SET MM(2const), uselib(2), dladdr(3), dlinfo(3), dl iterate phdr(3),
dlopen(3), dlsym(3), getauxval(3), lttng-ust(3), lttng-ust-cyg-profile(3), elf(5), capabilities(7), environ(7),
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