

**Notice**

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.

[Learn more and customize](#)[Reject](#)[Accept](#)

1 MARCH 2023 • JOE DESIMONE

# Stopping Vulnerable Driver Attacks

Using vulnerable drivers to gain kernel mode execution.

⌚ 7 min read   ⚡ Security operations, Detection science



## Key takeaways

- Ransomware actors are leveraging vulnerable drivers to tamper with endpoint security products.
- Elastic Security [released](#) 65 YARA rules to detect vulnerable driver abuse.
- Elastic Endpoint (8.3+) protects users from this threat.

## Jump to section

[Key takeaways](#)[Background](#)[Attack flow](#)[Blocklisting](#)[Allowlisting](#)[Behavior control](#)[First seen](#)[Conclusion](#)

## Background

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.

Fast forward to 2022, and attacks leveraging vulnerable drivers are a growing concern due to a **proliferation** of open source **tools** to perform these **attacks**. Vulnerable drivers have now been used by ransomware to terminate security software before encrypting the system. Organizations can reduce their risk by limiting administrative user permissions. However, it is also imperative for security vendors to protect the user-to-kernel boundary because once an attacker can execute code in the kernel, security tools can no longer effectively protect the host. Kernel access gives attackers free rein to tamper or terminate endpoint security products or inject code into protected processes.

This post includes a primer on kernel mode attacks, along with Elastic’s recommendations for securing users from kernel attacks leveraging vulnerable drivers.

## Attack flow

There are a number of flaws in drivers that can allow attackers to gain kernel mode access to fully compromise the system and remain undetected. Some of the **most common** flaws include granting user mode processes write access to virtual memory, physical memory, or **model-specific registers** (MSR). Classic buffer overflows and missing bounds checks are also common.

A less common driver flaw is unrestricted **handle duplication**. While this may seem like innocuous functionality at first glance, handle duplication can be leveraged to gain full kernel code execution by user mode processes. For example, the latest **Process Explorer** driver by Microsoft exposes such a function.

An attacker can leverage this vulnerability to duplicate a **sensitive handle** to raw physical memory present in the System (PID 4) process.

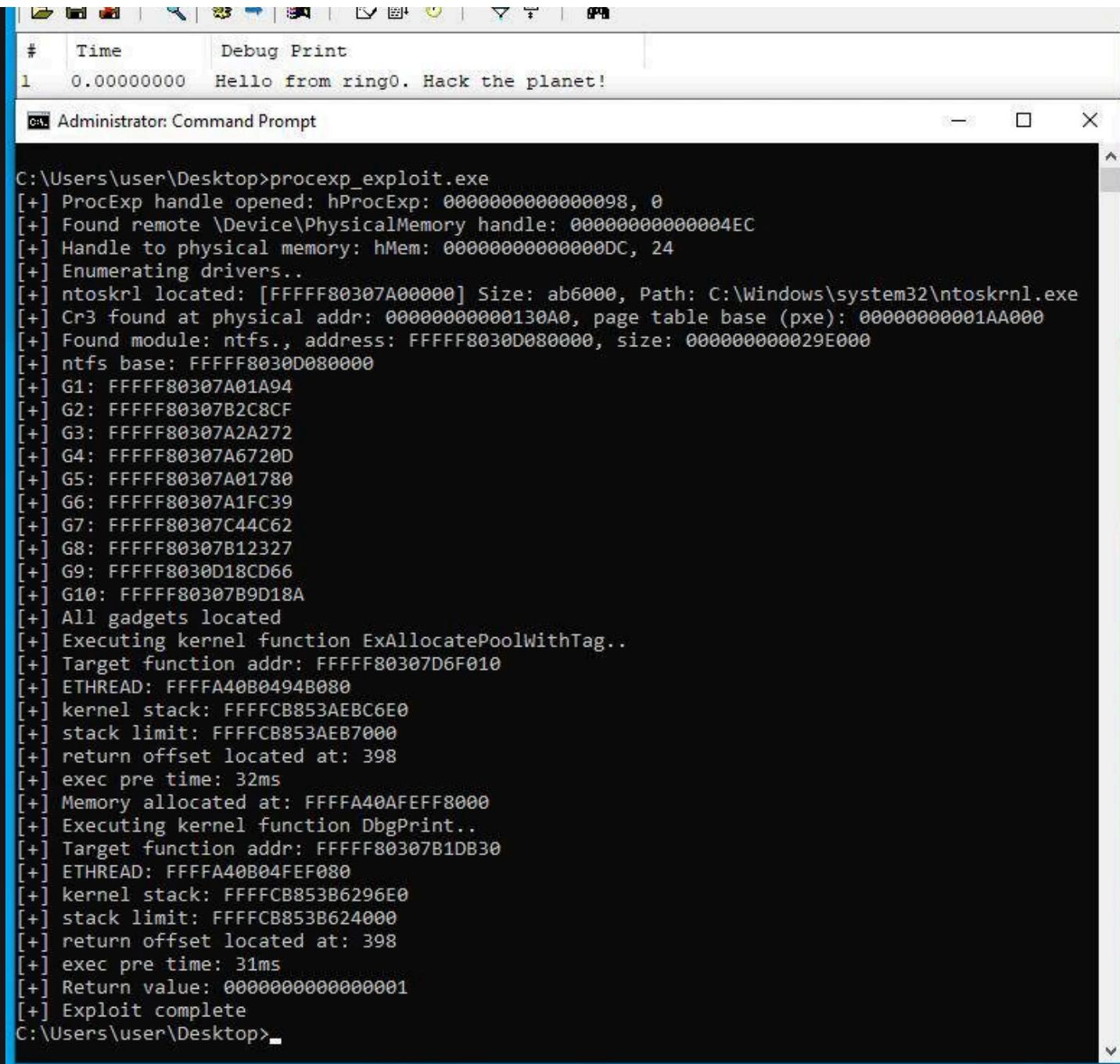
Type	Name	Handle	Granted acces...
File	\Device\NetBT_Tcpip_{35...	0x47b0	
Section	\Device\PhysicalMemory	0x494	Full control
Section	\Device\PhysicalMemory	0x4a8	Full control
File	\Device\Tcp	0x414	Write, Read
File	\Device\Tcp	0x92c	Write, Read

Handle to Physical Memory in the System process

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.



```
# Time Debug Print
1 0.00000000 Hello from ring0. Hack the planet!
Administrator: Command Prompt

C:\Users\user\Desktop>procexp_exploit.exe
[+] ProcExp handle opened: hProcExp: 0000000000000098, 0
[+] Found remote \Device\PhysicalMemory handle: 000000000000004EC
[+] Handle to physical memory: hMem: 00000000000000DC, 24
[+] Enumerating drivers..
[+] ntoskrnl located: [FFFFF80307A0000] Size: ab6000, Path: C:\Windows\system32\ntoskrnl.exe
[+] Cr3 found at physical addr: 00000000000130A0, page table base (pxe): 00000000001AA000
[+] Found module: ntfs., address: FFFF8030D080000, size: 00000000029E000
[+] ntfs base: FFFF8030D080000
[+] G1: FFFF80307A01A94
[+] G2: FFFF80307B2C8CF
[+] G3: FFFF80307A2A272
[+] G4: FFFF80307A6720D
[+] G5: FFFF80307A01780
[+] G6: FFFF80307A1FC39
[+] G7: FFFF80307C44C62
[+] G8: FFFF80307B12327
[+] G9: FFFF8030D18CD66
[+] G10: FFFF80307B9D18A
[+] All gadgets located
[+] Executing kernel function ExAllocatePoolWithTag..
[+] Target function addr: FFFF80307D6F010
[+] ETHREAD: FFFF40B0494B080
[+] kernel stack: FFFF853AEBC6E0
[+] stack limit: FFFF853AE87000
[+] return offset located at: 398
[+] exec pre time: 32ms
[+] Memory allocated at: FFFF40AFFE8000
[+] Executing kernel function DbgPrint..
[+] Target function addr: FFFF80307B1DB30
[+] ETHREAD: FFFF40B04FEF080
[+] kernel stack: FFFF853B6296E0
[+] stack limit: FFFF853B624000
[+] return offset located at: 398
[+] exec pre time: 31ms
[+] Return value: 0000000000000001
[+] Exploit complete
C:\Users\user\Desktop>
```

### Hijacking Threat Flow Control

We reported this issue to Microsoft in the vulnerable driver [submission portal](#) on July 26, but as of this writing have not received a response. We hope Microsoft will consider this a serious security issue worth addressing. Ideally, they will release a fixed version without the vulnerable [IOCTLs](#) and include it in the default HVCI blocklist. This would be consistent with the [blocking](#) of the ProcessHacker (now known as [System Informer](#)) driver for the [same flaw](#).

## Blocklisting

Blocklisting prevents known vulnerable drivers from loading on a system, and is a great first step to the vulnerable driver problem. Blocklisting can raise the cost of kernel attacks to levels out of reach for some criminal groups, while maintaining low false positive rates. The downside is it does not stop more [advanced groups](#), which can identify new, previously-unknown, vulnerable drivers.

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.

To expand on the existing list of known vulnerable drivers, we pivoted through VirusTotal data with known vulnerable import hashes and other metadata. We also combed through public attack tooling to identify additional vulnerable drivers. As common practice for Elastic Security, we made our [blocklist](#) available to the community. In Elastic [Endpoint Security](#) version 8.3 and newer, all drivers are validated against the blocklist in-line before they are allowed to load onto the system (shown below).

*enter image description here*

## Allowlisting

One of the most robust defenses against this driver threat is to only allow the combination of driver signer, internal file name, version, and/or hashes, which are known to be in use. We recommend organizations be as strict as feasible. For example, do not blanket trust all [WHQL](#) signed drivers. This is the classic application control method, albeit focusing on drivers. An organization’s diversity of drivers should be more manageable than the entirety of user mode applications. Windows Defender Application Control ([WDAC](#)) is a powerful built-in feature that can be configured this way. However, the learning curve and maintenance costs may still be too high for organizations without well-staffed security teams. To reap most of the benefits of the allowlisting approach, but reduce the cost of implementation to the users (ideally to blocklisting levels), we recommend two approaches in tandem: behavior control and alert on first seen.

## Behavior control

The concept behind behavior control is to produce a more manageable set of allowlistable behavior choke points that can be tuned for high confidence. For example, we can create a

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.

### *Example EQL Query*

From there, we can allowlist the benign applications that are known to exhibit this behavior. Then we receive and triage hits, tune the rule until it becomes high confidence, and then ship as part of our [malicious behavior protection](#). Elastic SIEM users can use the same technique to [create custom](#) Detection Engine [rules](#) tuned specifically for their environment.

## First seen

Elastic Security in 8.4 adds another powerful tool that can be used to identify suspicious drivers. This is the [“New Terms” rule type](#), which can be used to create an alert when a term (driver hash, signer, version, internal file name, etc) is observed for the first time.

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.

### *First Seen*

This empowers security teams to quickly surface unusual drivers the first time they’re seen in their environment. This supports a detection opportunity for even previously unknown vulnerable drivers or other driver-based adversary tradecraft.

### *Visualizing It*

## Conclusion

## Notice

We and selected third parties use cookies or similar technologies for technical purposes and, with your consent, for other purposes as specified in the [cookie policy](#).

Use the “Accept” button to consent. Use the “Reject” button to continue without accepting.



Twitter



Facebook



LinkedIn



Reddit