

Math 725 Advanced Linear Algebra

HW1

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Vector Space and Subspaces.

```
In [1]: # import libraries
import numpy as np
import sympy as sym
from sympy.matrices import Matrix
import matplotlib.pyplot as plt
from IPython.display import display, Math, Latex
```

1. (1A 13)

13. Show that $(ab)x = a(bx)$ for all $x \in \mathbb{F}^n$ and $a, b \in \mathbb{F}$.

Answer:

Let $x = (x_1, \dots, x_n)$, then

$$(ab)x = ab(x_1, \dots, x_n) = ((ab)x_1, \dots, (ab)x_n)$$

$$= (a(bx_1), \dots, a(bx_n)) = a(bx_1, \dots, bx_n)$$

$$= a(bx).$$

*see also <https://linearalgebras.com/1a.html>

2. (1B 6)

6. Let ∞ and $-\infty$ denote two distinct objects, neither of which is in \mathbb{R} .

Define an addition and scalar multiplication on $\mathbb{R} \cup \infty \cup (-\infty)$ as you could guess from the notation.

Specifically, the sum and production of two real numbers is as usual, and for $t \in \mathbb{R}$ define

$$t\infty = \begin{cases} -\infty & \text{if } t < 0, \\ 0 & \text{if } t = 0, \\ \infty & \text{if } t > 0, \end{cases}, t(-\infty) = \begin{cases} -\infty & \text{if } t < 0, \\ 0 & \text{if } t = 0, \\ \infty & \text{if } t > 0, \end{cases},$$

$$t + \infty = \infty + t = \infty, t + (-\infty) = (-\infty) + t = (-\infty),$$

$$\infty + \infty = \infty, (-\infty) + (-\infty) = (-\infty), \infty + (-\infty) = 0.$$

Is $\mathbb{R} \cup \infty \cup (-\infty)$ a vector space over \mathbb{R} ?

Explain.

Answer:

Consider \mathbb{V} as $\mathbb{R} \cup \infty \cup (-\infty)$, if \mathbb{V} is a subspace of \mathbb{R} , then $\infty(1) = (\infty) \in \mathbb{R}$. This is False, thus \mathbb{V} is not a subspace of \mathbb{R} as it's not even in the vector space \mathbb{R} .

A more rigorous approach is to consider the distributive property, $\{(a+b)x = ax + bx | \forall v \in \mathbb{V}\}$ and $a, b \in \mathbb{R}$.

Then $2 \cdot \infty = \infty$ and $2 = 3 - 1$, hence,

$$\{2 \cdot \infty = (3 - 1) \cdot \infty = 3 \cdot \infty + (-1) \cdot \infty = \infty - \infty = 0\} \implies (3 - 1) \cdot \infty = 3 \cdot \infty + (-1) \cdot \infty,$$

which is False, thus \mathbb{V} is not a subspace of \mathbb{R} .

Rarely is the union of subspaces a subspace itself.

Discussed with Darcy and Felix. Felix suggested adding the distributive property argument back.

However, is the intersection of subspaces always a subspace? Think about this, seems reasonable.

```
In [2]: # scratch ideas
        sym.oo-sym.oo # contradicts the definition above
```

```
Out[2]: NaN
```

3. (1C 5)

5. Is \mathbb{R}^2 a subspace of the complex space \mathbb{C}^2 ?

Answer:

Consider if \mathbb{R}^2 is a subspace of \mathbb{C}^2 , then $i(1, 1) = (i, i) \in \mathbb{R}^2$. This is False, thus not closed under scalar multiplication. \mathbb{R}^2 is not a subspace of \mathbb{C}^2 .

Discussed with Darcy and Felix. Darcy suggested explicitly stating that it's not closed under scalar multiplication.

4. (1C 9)

9. A function $f : \mathbb{R} \rightarrow \mathbb{R}$ is called *periodic* if there exists a positive number p such that $f(x) = f(x + p)$ for all $x \in \mathbb{R}$.

Is the set of periodic functions from $\mathbb{R} \rightarrow \mathbb{R}$ a subspace of $\mathbb{R}^{\mathbb{R}}$? Explain.

Answer:

Not a subspace, as functions with different periods must have a unique p to make the statement $f(x) = f(x + p)$ for all $x \in \mathbb{R}$ True. Thus a contradiction in additive identity property, $f(x) + g(x) \neq f(x + p) + g(x + p)$.

Consider $f(x) = \cos(x)$, for $f(x) = f(x + p)$ for all $x \in \mathbb{R}$ to be True, $p = 2\pi$. Now consider $g(y) = \cos(\frac{3y}{2})$ for $g(x) = g(x + q)$ for all $x \in \mathbb{R}$ to be True, $q = \frac{4\pi}{3}$. $p \neq q$ thus contradiction.

Discussed with Darcy and Felix. Fixed up some symbol names with their feedback. Thanks guys!

```
In [3]: # scratch it out
T1 = 1
T2 = sym.Rational(3,2)
def f(x): return sym.cos(T1*x)
def g(x): return sym.sin(T2*x)
x=3
p=sym.Rational(4,3)*sym.pi

# show contradiction
f(x)==f(x+p), g(x)==g(x+p)
```

```
Out[3]: (False, True)
```

Note to grader, the following approach was done as an exercise in proofs and maybe ignored during grading, however your feedback and insight is welcome.

Euler approach

Consider, periodic functions as infinite series,

$$\sin x = \sum_{n=0}^{\infty} \frac{(-1)^n x^{2n+1}}{(2n+1)!} = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, \text{ convergent for all } x;$$

$$\cos x = \sum_{n=0}^{\infty} \frac{(-1)^n x^{2n}}{(2n)!} = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, \text{ convergent for all } x;$$

It is obvious that for $n = 0$, $\frac{(-1)^n x^{2n+1}}{(2n+1)!} = \frac{(-1)^n x^{2n}}{(2n)!}$, is false. Hence by induction, $f_n(x) + g_n(x) = f_n(x+p) + g_n(x+p)$ is also False, thus not closed under addition.

I wonder if exhaustive negative induction is actually a thing? * See Appendix 1.

5. (1C 24)

24. A function $f : \mathbb{R} \rightarrow \mathbb{R}$ is called *even* if $f(-x) = f(x)$ for all $x \in \mathbb{R}$.

A function $f : \mathbb{R} \rightarrow \mathbb{R}$ is called *odd* if $f(-x) = -f(x)$ for all $x \in \mathbb{R}$.

Let U_e denote the set of real-valued even functions on \mathbb{R} and let U_o denote the set of real-valued odd functions on \mathbb{R} .

Show that $\mathbb{R}^{\mathbb{R}} = U_e \oplus U_o$.

Answer:

Let $f_e = \frac{f(x) + f(-x)}{2} \in U_e$ and $f_o = \frac{f(x) - f(-x)}{2} \in U_o$. Suppose $f(x) = f_e(x) + f_o(x)$, where $f(x) \in \mathbb{R}^{\mathbb{R}}$. This is the direction sum by definition, thus $\mathbb{R}^{\mathbb{R}} = U_e \oplus U_o$. *See also Axler 1.45

```
In [4]: # scratch the itch
def f(x): return x
def f_e(x): return (f(x)+f(-x))/2
def f_o(x): return (f(x)-f(-x))/2
def ff(x): return f_e(x)+f_o(x)
[f(x)==ff(x) for x in [-1,0,1,sym.sqrt(2)]] # not exhaustive but validates proof

Out[4]: [True, True, True, True]
```

6. (Midpoint Map)

X. Let T be the midpoint map on \mathbb{C}^6 , $T(p_1, \dots, p_6) = (\frac{p_1+p_2}{2}, \dots, \frac{p_6+p_1}{2})$.

Find six vectors v_1, \dots, v_6 , such that for every $i \in \{1, \dots, 6\}$, there is $\lambda_i \in \mathbb{C}$ with $T(v_i) = \lambda_i v_i$ and that no two vectors are complex rescaling of each other.

No need to include graphics but please describe in words the geometric pictures of the vectors.

Answer:

In Euler's Form, $v_i = \lambda_i v_0$ where $\lambda_i = \left(\frac{1}{2} + \frac{e^{\frac{j\pi}{3}}}{2}\right)^{i+1}$, $v_0 = [\zeta^0, \dots, \zeta^{n-1}]$, $\zeta = e^{\frac{i2\pi}{n}}$, $i = \{0, 1, \dots\}$, $j = \sqrt{-1}$ and $n = 6$.

The vector $v_0 = [\zeta^0, \dots, \zeta^{n-1}]$ was formed by rotating a point on the unit circle. The scalar λ_i , transforms v_i by rotating and scaling the vector. The graphical result for the above transformation may be seen below.

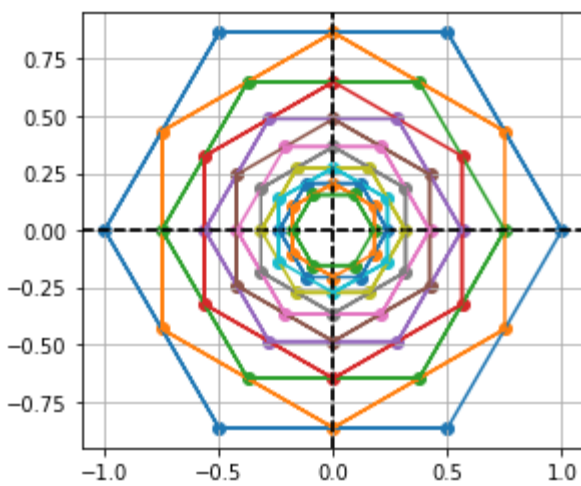
Discussed with Lekha, Darcy and Felix. Improved clarity of notation and fixed v0 of the Matrix v in the code. Really helpful talking it thru with the team! Thanks again guys!

```
In [5]: def z_plot(z, c=None):
plt.scatter(z.real, z.imag, c=c)
z = np.append(z, z[0]) # close the shape
plt.plot(z.real, z.imag, c=c)
plt.grid(visible=True);
plt.gca().set_aspect("equal") # square grids are pretty
plt.axhline(0, color='black', alpha = .2, linestyle='--')
plt.axvline(0, color='black', alpha = .2, linestyle='--')
```

```
In [8]: n=6
zeta = sym.exp(2*sym.I*sym.pi/n)
lambda_ = (1+zeta)/2
ss = 13 # how deep

v0 = Matrix([zeta**k for k in range(n)]).T # our vector in C^6
v = v0
v = v.col_join(Matrix([(lambda_**(i+1)*v0) for i in range(ss-1)])) # this one line does

# display our results
#display(v) # IPython.display is so pretty
for i in range(v.shape[0]):
    z_plot(np.array((v.row(i)).tolist()).astype(np.complex64)) # Aww Yeah!!! This is fas
```



Appendix 1. Contradiction

Ladies and gentleman skinny and scout
I'll tell you a tale I know nothing about
The admission is free so pay at the door
Now pull out a chair and sit on the floor

On one bright day in the middle of the night
Two dead boys got up to fight
Back to back they faced each other
Drew their swords and shot each other

The blind man came to see fair play
The mute man came to shout hooray
The deaf policeman heard the noise
And came to stop those two dead boys

He lived on the corner in the middle of the block
In a two story house on a vacant lot
A man with no legs came walking by
And kicked the lawman in his thigh

He crashed through a wall without making a sound
Into a dry creek bed and suddenly drowned
A long black hearse came to cart him away
But he ran for his life and is still gone today

I watched from the corner of the table
The only eyewitness to facts of my fable
If you doubt my lies are true
Just ask the blind man, he saw it too

* Two dead boys, author unknown

Appendix 2. Proof Techniques

To prove goal of the form:

- $\neg P$:
 - Reexpress as a positive statement.
 - use proof by contradiction; that is, assume that P is true and try to reach a contradiction.
- $P \implies Q$:
 - Assume P is true and prove Q .
 - Prove the contrapositive; that is, assume that Q is false and prove that P is false.
- $P \wedge Q$:
 - Prove P and Q separately. In other words, treat this as two separate goals: P and Q .
- $P \vee Q$:
 - Assume P is false and prove Q , or assume Q is false and prove P .
 - Use proof by cases. In each case, either prove P or prove Q .
- $P \iff Q$:
 - Prove $P \implies Q$ and $Q \implies P$, see method above for $P \implies Q$
- $\forall x P(x)$:
 - Let x stand for an arbitrary object, and prove $P(x)$. (If the letter x already stands for something in the proof, you will have to use a different letter for the arbitrary object.)
- $\exists x P(x)$:
 - Find a value of x that make $P(x)$ true. Prove $P(x)$ for this value of x .
- $\exists! x P(x)$:
 - Prove $\exists x P(x)$ (existence) and $\forall y \forall z ((P(y) \wedge P(z)) \implies y = z)$ (uniqueness).
 - Prove the equivalent statement $\exists x (P(x) \wedge P(y) \implies y = x)$.
- $\forall n \in \mathbb{N} P(n)$:
 - Mathematical induction: Prove $P(0)$ (base case) and $\forall n \in \mathbb{N} P(n) \implies P(n+1)$ (induction step).
 - Strong induction: Prove $\forall n \in \mathbb{N} [\forall k < n P(k) \implies P(n)]$.

To use a given form:

- $\neg P$:
 - Reexpress as a positive statement.
 - In a proof by contradiction, you can reach a contradiction by proving P .
- $P \rightarrow Q$:
 - If you are also given P , or you can prove that P is true, then you can conclude that Q is true.
 - Use the contrapositive: If you are given or can prove that Q is false, then you can conclude that P is false.
- $P \wedge Q$:
 - Treat this as two givens: P and Q .
- $P \vee Q$:
 - Use proof by cases. In the first case assume that P is true, then in the second case assume the Q is true.
 - If you are also given that P is false, or you can prove that P is false, then you can conclude that Q

if you are also given that P is false, or you can prove that P is false, then you can conclude that Q is true. Similarly, if you know that Q is false then you can conclude that P is true.

- $P \iff Q$:
 - Treat this as two givens: $P \implies Q$ and $Q \implies P$.
- $\forall x P(x)$:
 - You can plug in any value, say a , for x , and conclude that $P(a)$ is true.
- $\exists x P(x)$:
 - Introduce a new variable, say x_0 , into the proof, to stand for a particular object for which $P(x_0)$ is true.
- $\exists! x P(x)$:
 - Introduce a new variable, say x_0 , into the proof, to stand for a particular object for which $P(x_0)$ is true. You may assume that $\forall y (P(y) \implies y = x_0)$.

Techniques that can be used in any proof:

- Proof by contradiction: Assume the goal is false and derive a contradiction.
- Proof by cases: Consider several cases that are *exhaustive*, that is, that include all possibilities. Prove the goal in each case.

* See also, How to Prove It, Velleman