


A fast algorithm to compute a curve of confidence upper bounds for the False Discovery Proportion using a reference family with a forest structure

Guillermo Durand ¹ Laboratoire de Mathématiques d'Orsay, Université Paris-Saclay

Date published: 2024-07-03 Last modified: 2024-07-03

Abstract

This paper presents a new algorithm (and an additional trick) that allows to compute fastly an entire curve of post hoc bounds for the False Discovery Proportion when the underlying bound $V_{\mathfrak{R}}^*$ construction is based on a reference family \mathfrak{R} with a forest structure à la [Durand et al. \(2020\)](#). By an entire curve, we mean the values $V_{\mathfrak{R}}^*(S_1), \dots, V_{\mathfrak{R}}^*(S_m)$ computed on a path of increasing selection sets $S_1 \subsetneq \dots \subsetneq S_m, |S_t| = t$. The new algorithm leverages the fact that going from S_t to S_{t+1} is done by adding only one hypothesis.

Keywords: multiple testing, algorithmic, post hoc inference, false discovery proportion, confidence bound

1	Contents	
2	1 Introduction	2
3	2 Notation and reference family methodology	2
4	2.1 Multiple testing notation	2
5	2.2 Post hoc bounds with reference families	2
6	2.3 Deterministic regions with a forest structure	2
7	3 New algorithms	2
8	3.1 Pruning the forest	2
9	3.2 Fast algorithm to compute a curve of confidence bounds on a path of selection sets	2
10	4 Numerical experiments	2
11	5 Conclusion	2
12	6 Acknowledgments	2
13	References	3
14	Session information	3

¹Corresponding author: guillermo.durand@universite-paris-saclay.fr

1 Introduction

2 Notation and reference family methodology

2.1 Multiple testing notation

2.2 Post hoc bounds with reference families

2.3 Deterministic regions with a forest structure

3 New algorithms

3.1 Pruning the forest

Algorithm 1 Pruning of \mathfrak{R}

```
1: procedure PRUNING( $\mathfrak{R} = (R_k, \zeta_k)_{k \in \mathcal{K}}$  with  $\mathfrak{R}$  complete)
2:    $\mathcal{K}^{\text{pr}} \leftarrow \mathcal{K}$ 
3:    $H \leftarrow \max_{k \in \mathcal{K}} \phi(k)$  ▷ maximum depth
4:   for  $h = H - 1, \dots, 1$  do
5:      $\mathcal{K}^h \leftarrow \{k \in \mathcal{K} : \phi(k) = h\}$ 
6:      $\text{newVec} \leftarrow (0)_{k \in \mathcal{K}^h}$ 
7:     for  $k \in \mathcal{K}^h$  do
8:        $\text{Succ}_k \leftarrow \{k' \in \mathcal{K}^{h+1} : R_{k'} \subseteq R_k\}$ 
9:       if  $\text{Succ}_k = \emptyset$  then
10:         $\text{newVec}_k \leftarrow \zeta_k$ 
11:       else
12:        if  $\zeta_k \geq \sum_{k' \in \text{Succ}_k} \text{Vec}_{k'}$  then
13:          $\mathcal{K}^{\text{pr}} \leftarrow \mathcal{K}^{\text{pr}} \setminus \{k\}$ 
14:        end if
15:         $\text{newVec}_k \leftarrow \min(\zeta_k, \sum_{k' \in \text{Succ}_k} \text{Vec}_{k'})$ 
16:       end if
17:     end for
18:      $\text{Vec} \leftarrow \text{newVec}$ 
19:   end for return  $(\mathcal{K}^{\text{pr}}, \sum_{k \in \mathcal{K}^1} \text{Vec}_k)$ 
20: end procedure
```

3.2 Fast algorithm to compute a curve of confidence bounds on a path of selection sets

4 Numerical experiments

5 Conclusion

6 Acknowledgments

This work has been supported by ANR-20-IDEES-0002 (PIA), ANR-19-CHIA-0021 (BISCOTTE), ANR-23-CE40-0018 (BACKUP) and ANR-21-CE23-0035 (ASCAI).

Algorithm 2 Formal computation of $(V_{\mathfrak{R}}^*(S_t))_{0 \leq t \leq m}$

```
1: procedure CURVE( $\mathfrak{R} = (R_k, \zeta_k)_{k \in \mathcal{K}}$  with  $\mathfrak{R}$  complete, path  $(S_t)_{1 \leq t \leq m}$  with  $S_t = \{i_1, \dots, i_t\}$ )
2:    $\mathcal{P}^0 \leftarrow \{(i, i) : 1 \leq i \leq n\}$  ▷ the set of all atoms indices
3:    $\mathcal{K}_0^- \leftarrow \{k \in \mathcal{K} : \zeta_k = 0\}$ 
4:    $\eta_k^0 \leftarrow 0$  for all  $k \in \mathcal{K}$ 
5:   for  $t = 1, \dots, m$  do
6:     if  $i_t \in \bigcup_{k \in \mathcal{K}_{t-1}^-} R_k$  then
7:        $\mathcal{P}^t \leftarrow \mathcal{P}^{t-1}$ 
8:        $\mathcal{K}_t^- \leftarrow \mathcal{K}_{t-1}^-$ 
9:        $\eta_k^t \leftarrow \eta_k^{t-1}$  for all  $k \in \mathcal{K}$ 
10:    else
11:      for  $h = 1, \dots, h_{\max}(t)$  do
12:         $\eta_{k^{(t,h)}}^t \leftarrow \eta_{k^{(t,h)}}^{t-1} + 1$ 
13:        if  $\eta_{k^{(t,h)}}^t < \zeta_k$  then
14:          Pass
15:        else
16:           $h_t^f \leftarrow h.$ 
17:           $\mathcal{P}^t \leftarrow \left( \mathcal{P}^{t-1} \setminus \{k \in \mathcal{P}^{t-1} : R_k \subseteq R_{k^{(t,h_t^f)}}\} \right) \cup \{k^{(t,h_t^f)}\}$ 
18:           $\mathcal{K}_t^- \leftarrow \mathcal{K}_{t-1}^- \cup \{k^{(t,h_t^f)}\}$ 
19:          Break the loop
20:        end if
21:      end for
22:      if the loop has been broken then
23:         $\eta_k^t \leftarrow \eta_k^{t-1}$  for all  $k \in \mathcal{K}$  not visited during the loop, that is all  $k \notin \{k^{(t,h)}, 1 \leq h \leq h_t^f\}$ 
24:      else
25:         $\mathcal{P}^t \leftarrow \mathcal{P}^{t-1}$ 
26:         $\mathcal{K}_t^- \leftarrow \mathcal{K}_{t-1}^-$ 
27:         $\eta_k^t \leftarrow \eta_k^{t-1}$  for all  $k \in \mathcal{K}$  not visited during the loop, that is all  $k \notin \{k^{(t,h)}, 1 \leq h \leq$ 
28:           $h_{\max}(t)\}$ 
29:        end if
30:      end if
31:    end for return  $\mathcal{P}^t, \eta_k^t$  for all  $t = 1, \dots, m$  and  $k \in \mathcal{K}$ 
32: end procedure
```

References

Guillermo Durand, Gilles Blanchard, Pierre Neuvial, and Etienne Roquain. Post hoc false positive control for structured hypotheses. *Scand. J. Stat.*, 47(4):1114–1148, 2020. ISSN 0303-6898. doi: 10.1111/sjos.12453. URL <https://doi.org/10.1111/sjos.12453>.

Session information

R version 4.4.0 (2024-04-24)
Platform: x86_64-pc-linux-gnu
Running under: Ubuntu 22.04.4 LTS
Matrix products: default

Algorithm 3 Implementation of $(V_{\mathfrak{R}}^*(S_t))_{0 \leq t \leq m}$

```
1: procedure CURVE( $\mathfrak{R} = (R_k, \zeta_k)_{k \in \mathcal{K}}$  with  $\mathfrak{R}$  complete, path  $(S_t)_{1 \leq t \leq m}$  with  $S_t = \{i_1, \dots, i_t\}$ )
2:    $V_0 \leftarrow 0$ 
3:    $\mathcal{K}^- \leftarrow \{k \in \mathcal{K} : \zeta_k = 0\}$ 
4:    $\eta_k \leftarrow 0$  for all  $k \in \mathcal{K}$ 
5:   for  $t = 1, \dots, m$  do
6:     if  $i_t \in \bigcup_{k \in \mathcal{K}^-} R_k$  then
7:        $V_t \leftarrow V_{t-1}$ 
8:     else
9:       for  $h = 1, \dots, h_{\max}(t)$  do
10:        find  $k^{(t,h)} \in \mathcal{K}^h$  such that  $i_t \in R_{k^{(t,h)}}$ 
11:         $\eta_{k^{(t,h)}} \leftarrow \eta_{k^{(t,h)}} + 1$ 
12:        if  $\eta_{k^{(t,h)}} < \zeta_k$  then
13:          pass
14:        else
15:           $\mathcal{K}^- \leftarrow \mathcal{K}^- \cup \{k^{(t,h)}\}$ 
16:          break the loop
17:        end if
18:      end for
19:       $V_t \leftarrow V_{t-1} + 1$ 
20:    end if
21:  end for return  $(V_t)_{1 \leq t \leq m}$ 
22: end procedure
```

```
39 BLAS:    /usr/lib/x86_64-linux-gnu/openblas-pthread/libblas.so.3
40 LAPACK:  /usr/lib/x86_64-linux-gnu/openblas-pthread/libopenblas-p-r0.3.20.so;  LAPACK version 3.10.0
41
42 locale:
43   [1] LC_CTYPE=C.UTF-8      LC_NUMERIC=C          LC_TIME=C.UTF-8
44   [4] LC_COLLATE=C.UTF-8    LC_MONETARY=C.UTF-8   LC_MESSAGES=C.UTF-8
45   [7] LC_PAPER=C.UTF-8      LC_NAME=C             LC_ADDRESS=C
46   [10] LC_TELEPHONE=C        LC_MEASUREMENT=C.UTF-8 LC_IDENTIFICATION=C
47
48 time zone: UTC
49 tzcode source: system (glibc)
50
51 attached base packages:
52 [1] stats      graphics  grDevices datasets  utils      methods    base
53
54 loaded via a namespace (and not attached):
55 [1] compiler_4.4.0    fastmap_1.1.1      cli_3.6.2          htmltools_0.5.8.1
56 [5] tools_4.4.0       yaml_2.3.8         rmarkdown_2.26     knitr_1.46
57 [9] jsonlite_1.8.8    xfun_0.43          digest_0.6.35      rlang_1.1.3
58 [13] renv_1.0.7        evaluate_0.23
```