



East West University

Department of CSE

Project Report

Course Code: CSE207

Course Name: Data Structure

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Introduction

XOR linked list is a data structure that uses the bitwise XOR operation to store the memory addresses of the next and previous nodes like a doubly linked list.

Algorithm

1. X-OR Function:

Algorithm XOR(p, q):

```
return (node*)((uintptr_t)(p) ^ (uintptr_t)(q))
```

2. Create Node Function:

Algorithm CreateNode(d):

```
node = malloc(sizeof(node))
```

```
if node is NULL:
```

```
    print "Memory allocation failed."
```

```
    exit(EXIT_FAILURE)
```

```
node->data = d
```

```
node->link = NULL
```

```
return node
```

3.Create List Function:

Algorithm CreateList():

Input: None

Output: Head of XOR Linked List

h = NULL

t = NULL

print "Enter the size of the list: "

read x

for i = 1 to x:

 print "Enter the element ", i, ": "

 read d

 m = CreateNode(d)

 if h is NULL:

 h = m

 t = m

 else:

 t->link = XOR(m, t->link)

 m->link = t

 t = m

return h

4.Display Function:

Algorithm Display(h):

Input: Head of XOR Linked List

Output: None

c = h

prev = NULL

next = NULL

while c is not NULL:

 print c->data, "--> "

 next = XOR(prev, c->link)

 prev = c

 c = next

print "NULL"

5.Insert at Beginning Function:

Algorithm InsertAtBeginning(h, d):

Input: Head of XOR Linked List, Data to be inserted

Output: Updated Head of XOR Linked List

m = CreateNode(d)

m->link = h

if h is not NULL:

 h->link = XOR(m, h->link)

```
h = m  
return h
```

6.Insert at position Function:

Algorithm InsertAtPosition(h, d, p):

Input: Head of XOR Linked List, Data to be inserted, Position

Output: None

```
if p <= 0:  
    call InsertAtBeginning(h, d)  
    return
```

```
m = CreateNode(d)
```

```
if h is NULL:  
    m->link = NULL  
    h = m  
    return
```

```
c = h  
prev = NULL  
next = NULL  
a = 1
```

```
while c is not NULL and a < p:
```

next = XOR(prev, c->link)

prev = c

c = next

a++

if c is NULL:

prev->link = XOR(m, prev->link)

m->link = prev

else:

m->link = XOR(prev, c)

prev->link = XOR(XOR(prev->link, c), m)

if c->link is not NULL:

c->link = XOR(m, XOR(prev, c->link))

7.Delete at Beginning Function:

Algorithm DeleteAtBeginning(h):

Input: Head of XOR Linked List

Output: None

if h is NULL:

print "List is empty. Cannot delete."

return

next = XOR(NULL, h->link)

if next is not NULL:

next->link = XOR(NULL, XOR(h, next->link))

free(h)

h = next

8.Delete at position Function:

Algorithm DeleteAtPosition(h, p):

Input: Head of XOR Linked List, Position

Output: None

if $p \leq 0$:

 call DeleteAtBeginning(h)

 return

if h is NULL:

 print "List is empty. Cannot delete."

 return

c = h

prev = NULL

next = NULL

a = 1

while c is not NULL and $a < p$:

 next = XOR(prev, c->link)

prev = c

c = next

a++

if c is NULL:

print "Position not found. Cannot delete."

else:

next = XOR(prev, c->link)

if next is not NULL:

next->link = XOR(prev, XOR(c, next->link))

if prev is not NULL:

prev->link = XOR(XOR(prev->link, c), next)

free(c)

9. Search By Key Function:

Algorithm SearchByKey(h, k):

Input: Head of XOR Linked List, Key to search

Output: Node with matching key or NULL

c = h

prev = NULL

next = NULL

while c is not NULL and c->data != k:

next = XOR(prev, c->link)

prev = c

c = next

if c is not NULL and c->data == k:

return c

else:

return NULL

10. **Reversed list Function:**

Algorithm ReverseList(h):

Input: Head of XOR Linked List

Output: None

c = h

prev = NULL

next = NULL

while c is not NULL:

next = XOR(prev, c->link)

prev = c

c = next

h = prev

Conclusion

In conclusion, the XOR linked list approach offers a means of representing doubly linked list nodes using a single pointer instead of the traditional two, resulting in a notable reduction in memory usage. This method allows for optimization in various aspects, including memory efficiency, traversal capabilities, and addressing complexities and challenges associated with implementing XOR linked lists in code.