Taller de Programación Algoritmo de Inferencia de Tipos

Paradigmas de Lenguajes de Programación

Facultad de Ciencias Exactas y Naturales Universidad de Buenos Aires

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Algoritmo de Inferencia

Entrada una expresión U de λ sin anotaciones

Salida el juicio de tipado $\Gamma \rhd M$: σ más general para U, o bien una falla si U no es tipable

Estrategia recursión sobre la estructura de U

Expresiones de tipo:

$$\sigma$$
 ::= $s \mid Nat \mid Bool \mid \sigma \rightarrow \tau$

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En Haskell:

Expresiones anotadas:

```
M ::= x
\mid 0 \mid succ(M) \mid pred(M) \mid iszero(M)
\mid true \mid false \mid if M then P else Q
\mid \lambda x : \sigma.M
\mid M N
```

Expresiones sin anotar:

```
M ::= x
\mid 0 \mid succ(M) \mid pred(M) \mid iszero(M)
\mid true \mid false \mid if M then P else Q
\mid \lambda x.M
\mid M N
```

```
type Symbol = String
data Exp a = VarExp Symbol
            | ZeroExp
            | SuccExp (Exp a)
            | PredExp (Exp a)
            | IsZeroExp (Exp a)
            | TrueExp
            | FalseExp
            | IfExp (Exp a) (Exp a) (Exp a)
            | LamExp Symbol a (Exp a)
            | AppExp (Exp a) (Exp a)
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           | LamExp Symbol a (Exp a)
           | AppExp (Exp a) (Exp a)
type AnnotExp = Exp Type
type PlainExp = Exp ()
```

Contexto

```
emptyEnv :: Env
extendE :: Env -> Symbol -> Type -> Env
removeE :: Env -> Symbol -> Env
```

evalE :: Env -> Symbol -> Type

joinE :: [Env] -> Env

domainE :: Env -> [Symbol]

Tipos de datos y funciones auxiliares Sustitiuciones y unificación

Sustituciones

```
emptySubst :: Subst
extendS :: Int -> Type -> Subst -> Subst
```

Sustitiuciones y unificación

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class Substitutable a where
   (<.>) :: Subst -> a -> a
   instance Substitutable Type -- subst <.> t
   instance Substitutable Env -- subst <.> env
   instance Substitutable Exp -- subst <.> e
```

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```

Unificación

```
type UnifGoal = (Type, Type)
data UnifResult = UOK Subst | UError Type Type
mgu :: [UnifGoal] -> UnifResult
```

```
type TypingJudgment = (Env, AnnotExp, Type) data Result a = OK a | Error String inferType :: PlainExp \rightarrow Result TypingJudgment
```

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type TypingJudgment = (Env, AnnotExp, Type)
data Result a = OK a | Error String
inferType :: PlainExp -> Result TypingJudgment
inferType (VarExp x) = ...
...
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```

$$\mathbb{W}(x) \stackrel{\text{def}}{=} \{x:s\} \triangleright x:s, \quad s \text{ variable fresca}$$

¡A programar!

Consigna

- Completar archivo TypeInference.hs
- Definir la función inferType utilizando infer'
- Definir la función infer' para los casos
 VarExp ZeroExp LamExp AppExp
- Usar pattern matching sobre Exp

¡A programar!

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Tip

```
let x = expr1 in expr2
case expr of Pattern1 -> res1
    ...
Paternn -> resn
```

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Prueba

- Archivo Examples.hs
- expr n :: String
- ullet inferExpr :: String o Doc

```
infer'(SuccExp e) n =
```

```
infer'(SuccExp e) n =
case infer' e n of
```

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infer'(SuccExp e) n =
case infer' e n of
   OK ( n', (env', e', t') ) ->
```

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```
infer'(SuccExp e) n =
case infer' e n of
   OK ( n', (env', e', t') ) ->
        case mgu [ (t', TNat) ] of
```

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infer'(SuccExp e) n =
case infer' e n of
   OK ( n', (env', e', t') ) ->
      case mgu [ (t', TNat) ] of
      UOK subst ->
```

UError u1 u2 ->

res@(Error _) -> res

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infer'(SuccExp e) n =
case infer' e n of
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infer'(SuccExp e) n =
case infer' e n of
   OK ( n', (env', e', t') ) ->
        case mgu [ (t', TNat) ] of
            UOK subst ->
                OK ( n', (
                         ) )
            UError u1 u2 ->
                uError u1 u2
    res@(Error _) -> res
```

```
infer'(SuccExp e) n =
case infer' e n of
    OK ( n', (env', e', t') ) ->
        case mgu [ (t', TNat) ] of
            UOK subst ->
                OK ( n', (
                           subst <.> env',
                           subst <.> SuccExp e',
                          TNat
                          ) )
            UError u1 u2 ->
                uError u1 u2
    res@(Error _) -> res
```