1 Introduction

1.1 Protocol performances

G: Total load, S arrival rate of new packets.

1.1.1 Pure ALOHA

If you have data to send, send the data. If the message collides with another transmission, try resending later. On collision, sender waits random time before trying again.

$$P(k \text{ transm. in } 2Xs) = \frac{(2G)^k}{k!}e^{-2G}$$
$$S = G \cdot P(0) = Ge^{-2G}$$

1.1.2 Slotted ALOHA

Probability of k packets generated during a slot: P(k) = $\frac{G^k e^{-G}}{L!}$ Throughput: $P(1) = Ge^{-G}$

Goal: reduce the wastage of bandwidth due to packet collisions. Principle: sensing the channel before transmitting (never transmit when the channel is busy).

Non-persistent If channel is busy, directly run back off algo-

p-persistent If it is busy, they persist with sensing until the channel becomes idle. If it is idle:

- With probability p, the station transmits its packet
- With probability 1-p, the station waits for a random time and senses again

Performance of Unslotted nonpersistent CSMA : For a = t_{prop}/X , the normalized one-way propagation delay. S = G^{-aG}

$$\frac{G}{G(1+2a)+e^{-aG}}$$

Performance of Slotted nonpersistent CSMA $: S = \frac{aG^{-aG}}{C}$

renormance of Stotted holipersistent CSMA : $S = \frac{1}{1 - e^{-aG} + a}$						
Comment	Dis- advantages	Advantages	Signal separation	Terminals	ldea	Approach
used in all cellular systems	inflexible, antennas typically fixed	very simple, increases capacity per km²	cell structure, directed antennas	only one terminal can be active in one cell/one sector	segment space into cells/sectors	SDMA
standard in fixed networks, together with FDMA/SDMA used in many mobile networks	guard space needed (multipath propagation), synchronization difficult	established, fully digital, flexible	synchronization in the time domain	all terminals are active for short periods of time on the same frequency	segment sending time into disjoint time-slots, demand driven or fixed patterns	TDMA
typically combined with TDMA (frequency hopping patterns) and SDMA (frequency reuse)	inflexible, frequencies are a scarce resource	simple, established, robust	filtering in the frequency domain	every terminal has its own frequency, uninterrupted	segment the frequency band into disjoint sub-bands	FDMA
higher complexity	complex receivers, needs more complicated power control for senders	flexible, less frequency planning needed, soft handover	code plus special receivers	all terminals can be active at the same place at the same moment, uninterrupted	spread the spectrum using orthogonal codes	CDMA

1.2 Exercises

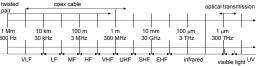
Capacity of a link vs Transmission capacity (=total capacity of all the links). Wire: $C_t = \min\{C_1, C_2\}$ Wireless: $d/C_t = d/C_1 + d/C_2 \leftrightarrow C_t = (c_1c_2/c_1 + c_2)$ ALOHA: Aloha channel with infinite number of users gives

94% of idle slots. $P(0) = e^{-G} = 0.94 \rightarrow G = 0.062$ $S = P(1) = Ge^{-G} \approx 5.8\%$

 $G < G_{peak} = 1$: channel underloaded.

Ration of busy slots occupied by collisions : $\frac{1-P(0)-P(1)}{1-P(0)} =$

2 WLAN Engineering aspects



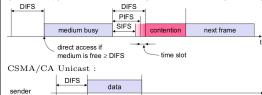
Frequency(f) and wave length(λ), $c = 3 \times 10^8 m/s$: $\lambda = c/f$

2.1 802.11

receive

other stations

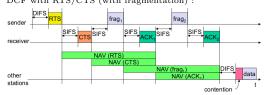
Physical layer: DSSS or FHSS, MAC Layer: best effort asynchronous data service, DCF CSMA/CA (mandatory), DCF with RTS/CTS or PCF (optional)



DIFS

data

DCF with RTS/CTS (with fragmentation)



waiting time

MAC address format :

scenario	to DS	from DS	address 1	address 2	address 3	address 4
ad-hoc network	0	0	DA	SA	BSSID	-
infrastructure network, from AP	0	1	DA	BSSID	SA	-
infrastructure network, to AP	1	0	BSSID	SA	DA	-
infrastructure network, within DS	1	1	RA	TA	DA	SA

2.2 Exercises

Wireless LAN use polling between M workstations and a central access point. Channel at 25Mbps. Stations 100 m away from AP, polling messages 64 bytes long. Packet length: 1250 bytes. No more packet indicated with 64-byte message. Maximum arrival rate $\lambda_{max} = \rho_{max} * Br/P_{length} \rho_{max} =$

Effectivetime _				
Wholetime	$M*(NT_{packet}+T_{poll}+T_{end}+2t_{prop})$			
	1050.0			
$t_{nron} = d/c T_{nacket}$	= 1200 + 0			

One station A sends a frame to another station B in a different BSS in an IEEE 802.11 infrastructure network with DCF access method without RTS/CTS.

 $A \rightarrow AP1$

·o	From	Type	Dur	A1	A2	A431 Free Space Loss
1	0	Data	$T_d + SIFS + T_A$	BSS1	A	Pree space loss, ideal isotropic antenna:
	_					Tree space loss, ideal isotropic antenna.

To DS	From DS	Type	Duration	Addr. 1	
0	0	ACK	0	A	
$\overline{\text{AP1} \rightarrow \text{AP1} : 1 , 1 , \text{Data}, T_d + S + T_A , \text{AP1}, B , A}$					

 $AP2 \rightarrow AP1 : 0, 0, ACK, 0, AP1$

 $AP2 \rightarrow B: 0, 1, Data, T_d + S + T_A, B, BSS2, A$ $B \rightarrow AP2 : 0, 0, ACK, 0, BSS2$

3 Bianchi model

 π , probability of transmission, p, probability of collision, $b_{i,k}$ stationary probability of state i, k: $p = 1 - (1 - \pi)^{N-1}$ $\pi = \sum_{i=0}^m b_{i,0} = \frac{b_{0,0}}{1-p} = \frac{2(1-2p)}{(1-2p)(W_{min}+1) + pW_{min}(1-(2p)^m)}$
$$\begin{split} &= \frac{i=0}{1+W_{\min}+pW_{\min}\sum_{k=0}^{m-1}(2p)^k} \\ &b_{i,k} = \frac{CW_i-k}{CW_i} \cdot \left\{ \begin{array}{ll} (1-p)\sum_{j=0}^{m}b_{j,0} & i=0 \\ p \cdot b_{i-1,0} & 0 < i < m \\ p \cdot (b_{m-1,0}+b_{m,0}) & i=m \end{array} \right. \end{split}$$

3.1 Saturation throughput

$$\begin{split} \tau &= \frac{E[\text{Payload Transmitted by user i in a slot time}]}{E[\text{Duration of slot time}]} \\ &= \frac{P_{\text{s}}P_{\text{tr}}L}{P_{\text{s}}P_{\text{tr}}T_{\text{s}} + P_{\text{tr}}(1-P_{\text{s}})T_{\text{c}} + (1-P_{\text{tr}})T_{\text{id}}}, \\ P_{\text{s}} &= \frac{N\pi(1-\pi)^{N-1}}{1-(1-\pi)^{N}}, \\ P_{\text{tr}} &= 1-(1-\pi)^{N}, \end{split}$$

$$T_{
m s} = t_{
m header} + t_{
m payload} + {
m SIFS} + t_{
m ACK} + {
m DIFS} + 2\sigma,$$

 $T_{
m c} = t_{
m header} + t_{
m payload} + {
m SIFS} + \sigma$

3.2 DOMINO Cheating detection

3			
Cheating Method	Detection Test		
Frame scrambling	Number of retransmissions		
Oversized NAV1	Comparison of the declared and actual NAV values		
Transmission before	Comparison of the idle time after		
DIFS	the last ACK with DIFS		
Backoff manipulation	Actual Backoff/ Consecutive		
	Backoff		
Frame scrambling with MAC forging	Periodic dummy frame injection		

4 Antennas & Propagation

Free space propagation, received power: $P_{\rm R} = P_{\rm T} \frac{A_{\rm R}}{4\pi d^2} \eta_{\rm R}$ with $\eta_{\rm B}$ an efficiency parameter, $A_{\rm B}$ the receiving antenna

Focusing capability, depends on size in wavelength λ : $G_{\mathrm{T}} = 4\pi \eta_{\mathrm{T}} A_{\mathrm{T}} / \lambda^2$

Directional emitter, received power:
$$P_{\rm R} = P_{\rm T}G_{\rm T}\frac{A_{\rm R}}{4\pi d^2}\eta_{\rm R}$$

Free space received power: $P_{\rm R} = P_{\rm T}G_{\rm T}G_{\rm R}(\frac{\lambda}{4\pi d})^2$
Loss: $L = \frac{P_T}{P_R} = \frac{(4\pi d)^2}{G_RG_T\lambda^2}$
 $ERP = P_TG_T$
Waves: $\lambda \cdot f = c$; $c = 3 \cdot 10^8$
Parabola: $G = \frac{fA}{\lambda^2}$

Mobnet Decibels : $B = 10 \log(\frac{P}{P_0})$

Propagation modes Ground Wave: $f \leq 2 \text{ Mhz}$, Sky Wave, Line of Sight: f > 30 Mhz

4.0.1 Line of sight equations

Horizon distance d[km] in **kilometers**, antenna height h[m]and refraction adjustment factor K = 4/3:

Optical LOS : $d = 3.57\sqrt{h}$

Effective LOS : $d = 3.57\sqrt{Kh}$

Max LOS distance for two antennas

$$3.57(\sqrt{Kh_1} + \sqrt{Kh_2})$$

ree Space Loss

 $\frac{P_t}{P_t} = \frac{(4\pi d)^2}{\sqrt{2}} = \frac{(4\pi f d)^2}{2}$

Free space loss equation can be recast:

$$L_{DB} = 10\log\frac{P_t}{P_r} = 20\log(f) + 20\log(d) - 147.56dB$$

Free space loss accounting for gain of other antennas

$$\frac{P_t}{P_r} = \frac{(4\pi d)^2}{G_r G_t \lambda^2} = \frac{(cd)^2}{f^2 A_r A_t}$$

 $G_t = gain of transmitting antenna$

 A_r = effective area of receiving antenna

Categories of noise : Thermal Noise, Intermodulation Noise, Cross-talk, Impulse Noise.

Thermal Noise $N_0 = kT \quad (W/Hz)$

For signal power S, bitrate R, $k = 1.3806 \cdot 10^{-23} JK^{-1}$ the Boltzmann constant and T the temperature: $\frac{E_b}{N_0} = \frac{S/R}{N_0} =$ $\frac{S}{kTR}$

4.2 Forward Error Correction (FEC)

Redundancy in packets to allow limited error correction at the receiver: used in 802.11a (Convolutional), HSDPA (Turbo Codes) and 802.11n (LDPC).

5 Cellular Networks

For a trunk of N channels, an offered load $A = \lambda E[X]$, X the call duration, Y the call arrival per sec $\sim \text{Poisson}(\lambda)$ and ρ the traffic carried by each channel:

> $P_{\text{Blocking}} = P(\text{Drop a call because busy line})$ $= \frac{A^N}{N! \sum_{i=0}^N (\frac{A^i}{i!})}$ $\rho = \frac{(1 - P_{\text{blocking}})A}{N}$

Cellular efficiency $E = \frac{Conversations}{collection}$

Area: $A = 1.5R^2\sqrt{3}$

Distance btw. adjacent cells: $d = \sqrt{3}R$

5.1 Co-channel interference

Co-channel reuse ratio : $Q = \frac{D}{R} = \sqrt{3N}$ with D the distance to the nearest co-channel cell, R the radius of a cell and N the cluster size.

Signal to Interference ratio (SIR) : $\mathrm{SIR} = \frac{S}{T} = \frac{S}{\sum_{i=1}^{i_0} I_i}.$ With

S the desired signal power, I_i the interference power from the ith interfering co-channel base-station, i₀ the number of co-channel interfering cells

Signal to Interference plus Noise ratio (SINR) : SINR = $\frac{S}{I+No}$

Average received power $P_r: P_r = P_0(\frac{d}{d_0})^{-\alpha}$ or

 $P_r(dBm) = P_0(dBm) - 10\alpha \log(\frac{d}{d_0})$ with P_0 the power received from a small distance d_0 from the transmitter and α the path loss exponent.

SIR in the corner of a cell : $\frac{S}{I} = \frac{R^{-\alpha}}{\sum_{i=1}^{i_0} D_i^{-\alpha}}$

First interfering layer approximation : $\frac{S}{I} = \frac{(\frac{D}{R})^{\alpha}}{i_0} = \frac{(\sqrt{3N})^{\alpha}}{i_0}$ eg. $= (\frac{D}{R})^2 \frac{1}{2} \text{ for two first layer interferers (cell divided into 3 sectors with directional antennas.)}$

5.2 Capacity of a cellular network

For B_{t} the total allocated spectrum and B_{c} the channel

$$m = \frac{B_t}{B_c \frac{Q^2}{3}} = \frac{B_t}{B_c \left(\frac{6}{3\frac{\alpha}{2}} \left(\frac{S}{I}\right)_{\min}\right)^{\frac{2}{\alpha}}} = \lfloor \frac{C}{N} \rfloor$$

For a cluster size N, $N = (i + j)^2 - ij$ for i, j = 0, 1, 2, ...and number of channels C.

5.2.1 CDMA Capacity: single cell case

For the bitrate R, available bandwidth W, noise spectral density N_0 , thermal noise η , received user signal (at base station) S, we have a possible number N of users:

$$N = 1 + \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

With a duty cycle δ (Discontinuous transmission mode: takes advantage of intermittent nature of speech):

$$N = 1 + \frac{1}{\delta} \frac{W/R}{E_b/N_0} - (\frac{\eta}{S})$$

And if we have m sectors, the effective capacity becomes mN. 5.2.2 CDMA multiple cells

Frequency reuse factor on the uplink $f=\frac{N_0}{N_0+\sum_i U_i N_{ai}}$ where $N_0=$ total interference power received from N-1 in-cell users, U_i = number of users in the i^{th} adjacent cell and N_{ai} = average interference power from a user located in the i^{th}

Average received power from users in adjacent cell $N_{ai}\,=\,$ $\sum_{i} N_{ij}/U_{i}$ where N_{ij} = power received at the base station of interest from the ith user in the ith cell

5.3 Ad-hoc Netowrks

Upper Bound for the Throughput If we have n identical randomly located nodes each capable of transmitting W bits/s. Then the throughput $\lambda(n)$ obtainable by each node for a randomly chosen destination is $\lambda(n) = \Theta\left(\frac{W}{\sqrt{n \log n}}\right)$

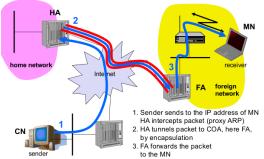
Routing proactive: DSDV, OLSR. reactive: AODV, DSR DSR: Route discovery only when source S attempts to send a packet to destination D, by flooding Route Requests (RREQ). Route maintenance by allowing S to detect when a link is broken with a Route Error message RERR. S try other route in its cache, otherwise route disc. AODV: Similar to DSR but maintains routing tables at the nodes (smaller header). AODV ages the routes and maintains a hop count.

6 Mobile Network Laver

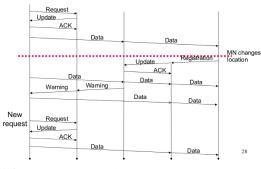
Mobile Network Laver : Transparency, Compatibility, Security, Efficiency, Scalability.

Network Layers Top-down: Application, Transport, (HIP layer), Network, Data-link, Physical.

Mobile IP :Issues : Security(Authentication to FA is problematic), Firewalls, QoS. IPSec can provide CIA by adding layer btwn IP and TCP/UDP. Mobile IPv6 : no FA, COA alwys co-loc, IPsec, optim. bidirectional tunnel HA<->COA.



Agent Discovery: -MN discovers its location, HA and FA periodically send advertisement messages or via DHCP -MN learns a COA. Registration : -MN securely signals the COA to the HA (via the FA) Tunneling: -HA encapsulates IP packets from CN and sends them to the COA -FA (or MN) decapsulates these packets and sends them to the MN (Route optimization: HA provides the CN with the current location of MN (FA).CN sends tunneled traffic directly to FA. Picture of route and FA handover optim :



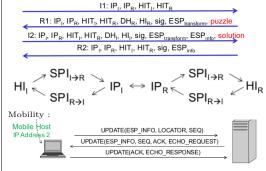
HIP New layer btw IP and transport, integrate security, mobility and multi-homing, decouple name and locator role of IP HI = public key. HIT=h(HI), DH : Diffie-Hellman key

FDD Frequency Division Duplex

FDMA Frequency Division Multiple Access

FEC Forward Error Correction FHSS Frequency Hopping Spread Spec-

material, sig signature generated with private key of $HI_{I/R}$ Initiator (I) Responder (R)



7 TCP

7.1 Standard

Tahoe Basic TCP. Three duplicate ACK's provoke fast retransmit (resend 1st missing packet), set ssthresh to cwnd/2, cwnd to 1 and provoke slow start.

Reno Three duplicate ACK's provoke fast retransmit, ssthresh to cwnd/2, cwnd to ssthresh + 3 and enter fast recov-

Fast Recovery Increase cwnd by 1 segment for every received duplicate ACK. (Warning, unlogical: When new ACK is received, cwnd = ssthresh and enter congestion avoidance). If a timeout occurs, set cwnd to 1 and enter slow start.

New Reno Fast Recovery More intelligent fast recovery where you remember the last received ACK.

7.2 Mobile

Indirect TCP (I-TCP) Connection split at FA. Standard TCP on the wire line, wireless optimized TCP on the wifi side: shorter timeout, faster retransmission. Loss of end-to-end semantics, security issues. AP acks all seg.

Mobile TCP (M-TCP) Split connection at FA. Monitor packets, if a disconnect is detected, report receiver window = 0: sender will go into persist mode and doesn't timeout or modify his congestion window. Preserves end-to-end semantics. Disady.: wifi losses propagate to the wire network, link-errors pkt loss must be resent by sender, security issues. Summary: only handles mobility errors, no transmission errors.

Snooping-TCP TCP-aware link layer: Split connections, FA buffers and retransmits segments, does not ACK buffered packets (preserves end-to-end semantics). PEP (content caching, compression): breaks security end-to-

end semantics.

Link layer retransmission only if error detected (FEC) NACK mech. for missing seg.

Transaction oriented TCP (T-TCP) TCP phases: connection setup, data transmission, connection release. T-TCP combines these steps and only 2-3 packets are needed for short messages. Efficient for single packet transactions, but requires TCP modifications on all hosts.

SIP Application layer protocol

SIP invite: Port number, IP, Preferred encoding. 200 OK: Port number, IP, Preferred encoding. (over TCP or UDP). How to find IP? User register to a SIP registrar server. Then use SIP proxies to find IP (similar to DNS). RTP is an alternative to SIP. VoIP over IEEE 802.11 DCF : Best effort

8 Security

Security Requirements : Confidentiality, Authenticity, Replay Detection, Integrity, Access Control, Jamming Protection.

GSM Shared secret and challenge responses, one-way auth.

3GPP (Improvements from GSM) Two-way auth., avoid fake base station, cipher keys and auth data is now encrypted, integrity. Privacy/Anonymity not completely protected.

9 Privacy

Privacy Related Notions Anonymity, untraceability, unlinkability, unobservability, pseudonymity

Best to worst against information leakage: GPS: no third-party, determined 'alone'. Cell-ID: requires the operator database that is relatively protected (they won't easily mine you). Wireless: requires one or several third-party owned databases that can track you, and it is relatively precise due to short radio range.

9.1 Privacy Metrics

Entropy-Based Anonymity A the anonymity set (set of subjects that might have performed the action), p_x the probability for an external observer that the action was performed by x: $H = -\sum_{\forall x \in A} p_x \log_2(p_x).$

Entropy-Based Unlinkability I_1, I_2 , sets of elements to be related, p_r , the probability two elements are related for an external observator: $H = -\sum_{\forall R \subset I_1 \times I_2} p_r \log_2(p_r)$

Expected error-based measure for correctness $\sum_{v \in V} p_v \times d(v, v_0)$

9.2 Mix-zone

 $p_{i,j} = \mathbb{P} \{ \text{exiting at } i | \text{entering at } i \}$

 $D_{i,j}$: random variable (delay) representing time that elapses between entering at i and exiting at j.

 $d_{i,j}(t) = \mathbb{P}\{D_{i,j} = t\}$

 $\mathbb{P}\left\{\text{exiting at } j \text{ at } t | \text{entering at } i \text{ at } \tau\right\} = p_{i,j} d_{i,j} (t - \tau)$

Correctness : $C = \sum_{\sigma} p_{\sigma} d(\sigma(i), j)$

LEAP Light EAP

LF Low Frequency

LTE Long Term Evolution

ACO Authenticated Cipher Offset AIFS Arbitrary Inter-Frame Space AMF Authentication and Key management Field AODV Ad Hoc On-demand Distance-Vector AP Access Point AP Access Point ATIM Ad-hoc Traffic Indication Map AUTN Ad-noc trainic indication Map AUTN Authentication Token AV Authentication Vector BO BackOff BSSID Basic Service Set Identifier BSS Basic Service Set CARMA Collision Avoidance and Resolu-tion Multiple Access CA Collision Avoidance CCA Clear Channel Assessment CDMA Code Division Multiple Access CH Correspondant Host

CN Correspondant Node COA Care-Of Address CRC packet received CoRreCtly CSMA/CD CSMA with Collision Detection CSMA Carrier Sense Multiple Access CTS Clear To Send CW Contention Window DAMA Demand-Assigned Multiple Access

DA Destination Address
DBPSK Differential Binary Phase Shift Keying DCF Distributed Coordination Function

DH Diffie-Hellman
DNS Domain Name System
DQPSK Differential Quadr Quadrature Phase Shift Keying DSDV Destination Sequenced Distance Vector
DSRC Dedicated Short Range Communica-DSR Dynamic Source Routing DSSS Direct Sequence Spread Spectrum DS Differentiated Service DS Distribution System DTIM Delivery Traffic Indication Map DoS Denial of Service EAP-TLS TLS over EAP EAPOL TLS OVER LAN EAP Extensible Authentication Protocol EDCA Enhanced Distributed Channel Ac-EHF Extra High Frequency EPC Electronic Product Code ESP Encapsulating Security Payload ESP_{info} Contains SPI

FA Foreign Agent

DECT Digital Enhanced Cordless Telecom-

munications

DHCP Dynamic Host Configuration Proto-

 $ESP_{transform}$ Supported crypto suites ESS Extended Service Set
FAMA Floor Acquisition Multiple Access

FQDN Fully Qualified Domain Name MACA-BI MACA By Invitation GFSK Gaussian Frequency Shift Keying GMK Group Master Key GPRS General Packet Radio Service GSM Global System for Mobile Communi-HA Home Agent HCCA HCF Controlled Channel Access HCF Hybrid Coordination Function HF High Frequency HIP Host Identity Protocol HIT Host Identity Tag HI Host Identifier HMIP Hierarchical Mobile IP HSPDA High Speed Downlink Packet Ac-ICMP Internet Control Message Protocol IFS Inter Frame Spacing IHL Internet Header Length IKE Internet Key Exchange IMSI International Mobile Subscriber ISI InterSymbol Interference KISS Keep It Simple and Stupid LDPC Low Density Parity Check

MACA Multiple Access with Collision Avoidance (RTS-CTS(+ACK)) MAC Message Authentication Code MAHO Mobile Assisted Handover MAP Mobility Anchor Point MD Mobile Device MF Medium Frequency MH Mobile Host MIB Management Information Base MIC Message Integrity Code MN Mobile Node MSC Mobile service Switching Center MTSO Mobile Telecommunications Switching Office NAASS Normalized Average Anonymity NAT Network Address Translation
NAV Net Allocation Vector
OFDMA Orthogonal Frequency-Division
Multiple Access OLSR Optimized Link- State Routing OTP One-Time Password
PCF Point Coordination Function
PEAP Protected EAP
PEP Performances Enhancing Proxies

LFSR Linear Feedback Shift Register

PMD Col Physical Medium Dependent PMK Pairwise Master Kev PN Pseudo-random Noise PSTN Public Switched Telephone Network PTK Pairwise Transient Key QoS Quality of Service RADIUS Remote Authentication Dial-In RA Receiver Address
RERR Route ERRor
RFID Radio Frequency Identification RREP Route REPly RREQ Route REQuests RSN Robust Security Network RTCP Real Time Control Protocol RTM Retransmission TimeOut RTP Real Time Protocol RTS Request To Send RVS Rendez-Vous Server RWND Receiver Window SACK Selective ACKnowledgment SA Security Association SA Source Address SDMA Space Division Multiple Access

SHF Super High Frequency

SIFS Short Inter Frame Spacing

SIM Subscriber Identity Module

PIN Personal Identification Number PLCP Physical Layer Convergence Proto-

STresh Slow Start Threshold STA STAtion STA Station IA Transmitter Address ICP Transmitter Address ICP Transmitter Department of the Depar TDMA Time Division Multiple Access TIM Traffic Indication Map TKIP Temporal Key Integrity Protocol TLS Transport Laver Security TMSI Temorary Mobile Subscriber Iden-TOS Type Of Service TSF Timing Synchronisation Function TTL Time To Live UHF Ultra High Frequency UMTS Universal Mobile Telecommunica-

UV Ultraviolet Light VANET Vehicular Ad-hoc NETwork VHF Very High Frequency VLF Very Low Frequency WAP Wireless Access Point WEP Wired Equivalent Privacy WIAN Wireless Local Area Network WMN Wireless Mesh Network WPAN Wireless Personal Area Network WPA WiFi Protected Access