#### **Routing Protocols**



- Setting all routes in large size networks is an unfeasible solution
- There exists several routing algorithm developed for Linux-based OSs (BGP, OSPF, OLSR, AODV, etc.)
- In this seminary we will focus on routing protocol for wireless multi-hop networks (specifically for MANETs and WMNs)
  - AODV: reactive routing protocol
  - OLSR: proactive routing protocol



#### **Routing Protocols**



- Reactive Protocols (AODV)
  - Generally involve large delays between the request and first packet delivery
  - Incur low overhead in low traffic scenarios
- Proactive Protocols (OLSR)
  - Packets are immediately delivered as paths are already established
  - Results in high path maintenance overhead since the paths are kept regardless of traffic Patterns
- Hybrid Protocols
  - Operate midway of delay and overhead performance



#### **Trade-Off**



- Latency of route discovery
  - Proactive protocols may have lower latency since routes are maintained at all times
  - Reactive protocols may have higher latency because a route from X to Y will be found only when X attempts to send to Y
- Overhead of route discovery/maintenance
  - Reactive protocols may have lower overhead since routes are determined only if needed
  - Proactive protocols can (but not necessarily) result in higher overhead due to continuous route updating
- Which approach achieves a better trade-off depends on the traffic and mobility patterns





## Reactive Routing Protocols

#### **Route Discovery**



- Reactive Routing Protocols discover reactively the route towards a destination D
  - When a packet needs to be sent to D and the route to D is unknown
- When node S wants to send a packet to node D, but does not know a route to D, node S initiates a route discovery
  - Source node S floods Route Request (RREQ)
- Destination D on receiving the first RREQ, sends back a Route Reply (RREP)



#### **Route Discovery**



- The route used to forward data packets is
  - stored in the packet header (DSR)
  - derived dinamically by using routing tables at each node (AODV)
- Storing routes in packet header results in a large overhead (particularly when data contents of a packet are small)
- AODV improves DSR by maintaining routing tables at the nodes
- Using AODV, routes are maintained only between nodes which need to communicate



## Ad Hoc On-Demand Distance Vector Routing (AODV)

- When a source node S desires a route to a destination D for which it does not already have a route, it broadcasts a route request (RREQ) packet across the network
- When a node re-broadcasts a RREQ, it sets up a reverse path pointing towards the source S
  - AODV assumes symmetric (bi-directional) links
- When the intended destination receives a Route Request, it replies by sending a Route Reply (RREP)
- RREP travels along the reverse path set-up when Route Request is forwarded



#### Route Request (RREQ)





## Route Request (RREQ)



Field	Size (bit)	<b>Description</b>
Type	8	Message Identifier (1 = RREQ)
Flags	5	See Flags description (J, R, G, D, U)
Reserved	11	Ignored
Hop Count	8	The number of hops from the Originator IP Address to the node handling the request
RREQ ID	32	A sequence number uniquely identifying the particular RREQ
Destination IP	32	The IP address of the destination for which a route is desired
Dst Seq Number	32	The latest sequence number received in the past by the originator for any route towards the destination
Originator IP Address	32	The IP address of the node which originated the Route Request
Originator Sequence Number	32	The current sequence number to be used in the route entry pointing towards the originator of the route rquest

#### RREQ Flags (5)



- J: Join flag; reserved for multicast.
- R: Repair flag; reserved for multicast.
- **G**: **Gratuitous** RREP flag; indicates whether a gratuitous RREP should be unicast to the node specified in the Destination IP Address field
- D: Destination only flag; indicates only the destination may respond to this RREQ
- U: Unknown sequence number; indicates the destination sequence number is unknown



#### Route Reply (RREP)





## Route Reply (RREP)



Field	Size (bit)	<b>Description</b>
Туре	8	Message Identifier (2 = RREP)
Flags	2	See Flags description (R, A)
Reserved	9	Ignored
Prefix Sz	5	refix Size specifies that the indicated next hop may be used for any nodes with the same routing prefix
Hop Count	8	The number of hops from the Originator IP Address to the Destination IP Address
Destination IP Address	32	The IP address of the destination for which a route is supplied
Dst Seq Number	32	The destination sequence number associated to the route
Originator IP Address	32	The IP address of the node which originated the RREQ for which the route is supplied
Lifetime	32	The time in milliseconds for which nodes receiving the RREP consider the route to be valid

## RREP Flags (2)



- Flags of Route Replay:
  - R: Repair flag; used for multicast.
  - A: Acknowledgment required.



#### Route Error (RERR)





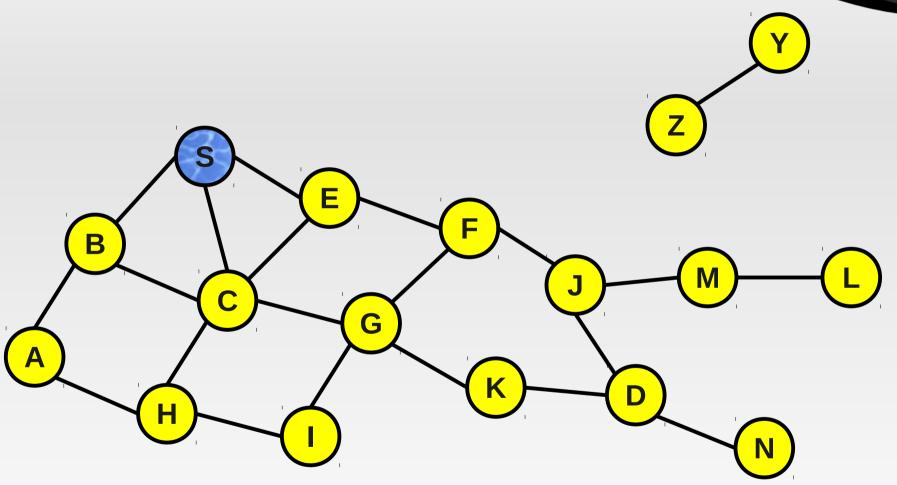
## Route Error (RREP)



Field	Size (bit)	Description
Туре	8	Message Identifier (3 = RERR)
N flag	1	No delete flag; set when a node has performed a local repair of a link
Reserved	15	Ignored
DestCount	8	The number of unreachable destinations included in the message; MUST be at least 1
Unreachable Destination IP	32	The IP address of the destination that has become unreachable due to a link break
Unreachable Dst Seq Number	32	The sequence number in the route table entry for the destination listed in the previous Unreachable Destination IP Address field





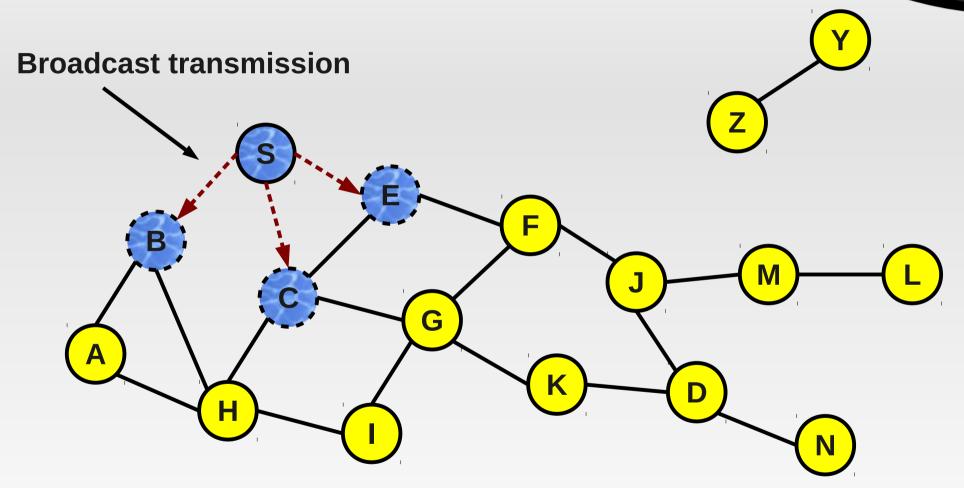




Represents a node that has received RREQ for D from S

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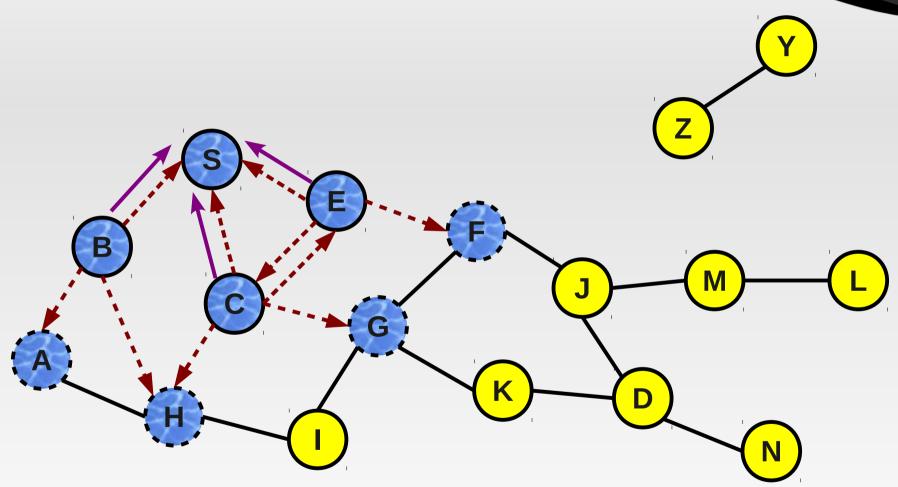




Represents transmission of RREQ



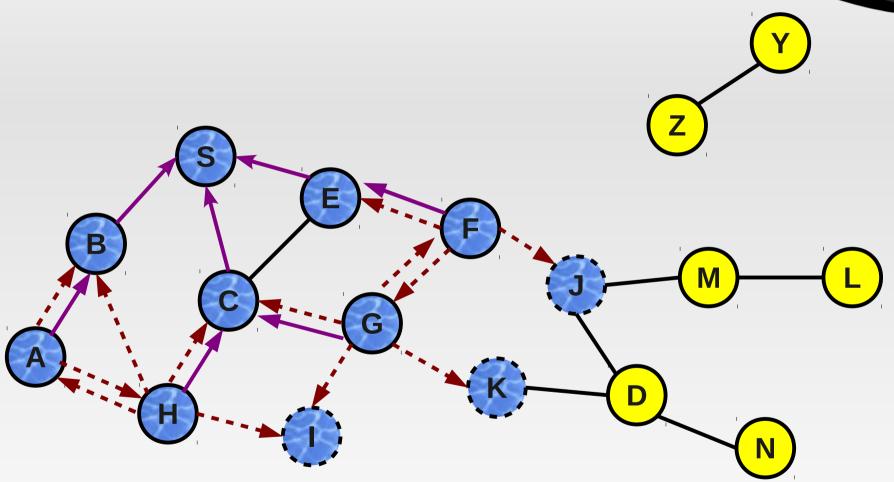




Represents links on Reverse Path



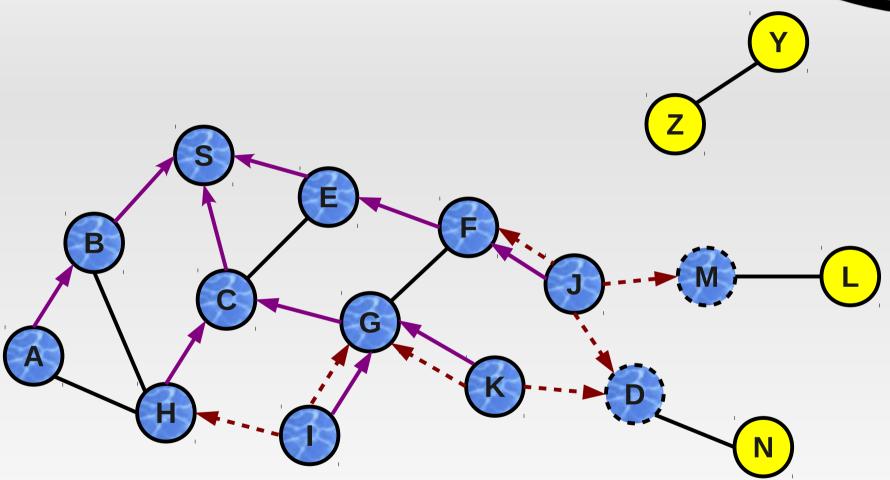




Node C receives RREQ from G and H, but does not forward it again, because node C has already forwarded RREQ once

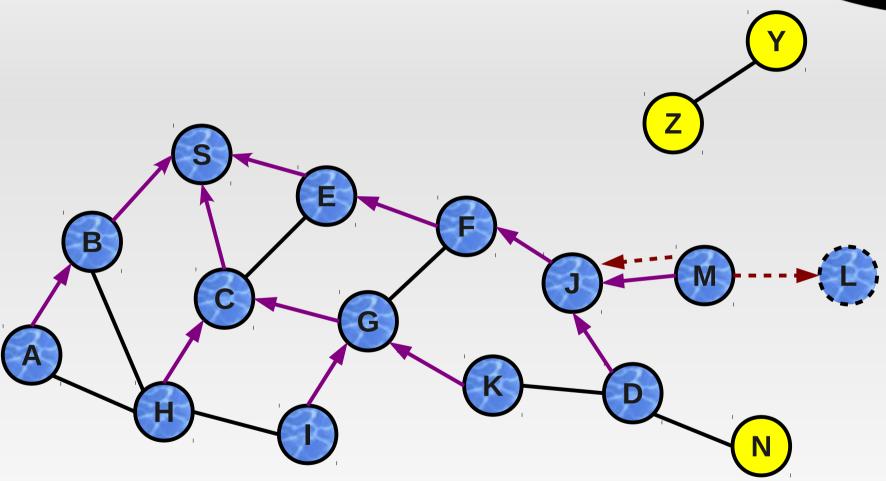
ubuntu





#### Route Replies in AODV



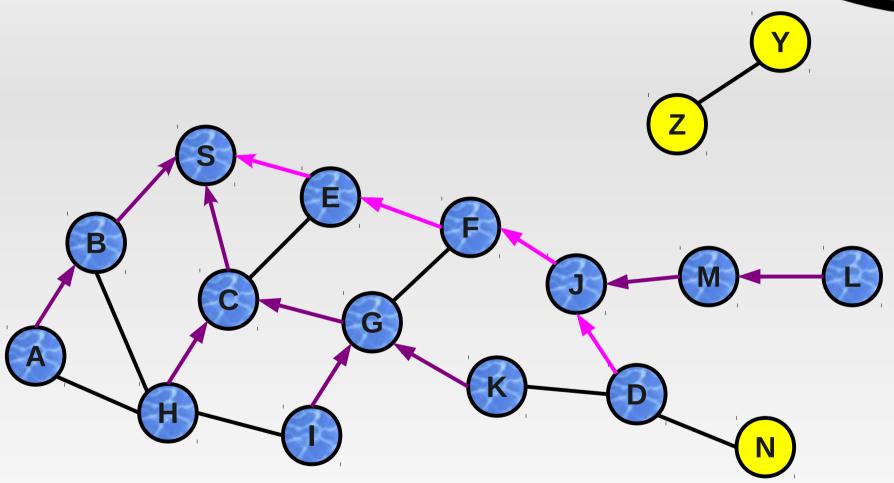


Node D does not forward RREQ, because node D is the intended target of the RREQ



### Route Replies in AODV





Represents links on path taken by RREP

ubuntu

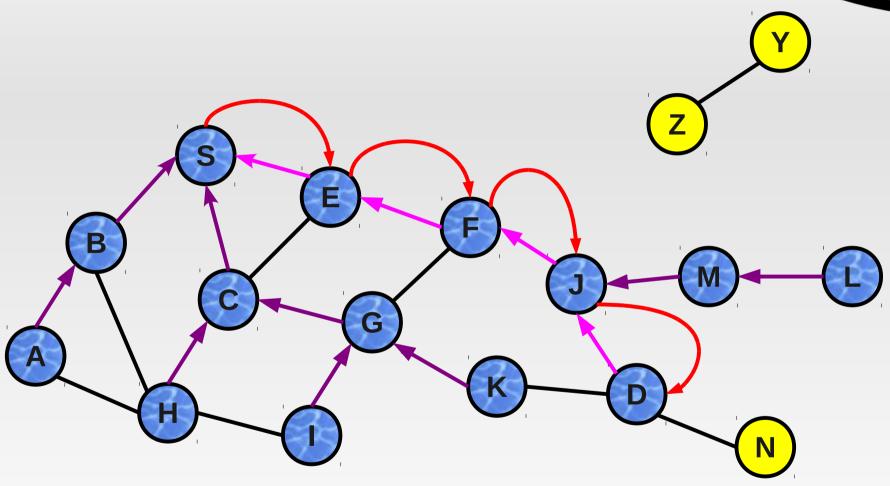
#### Route Reply in AODV



- An intermediate node (not the destination) may also send a Route Reply (RREP) provided that it knows a more recent path than the one previously known to sender S
- To determine whether the path known to an intermediate node is more recent, destination sequence numbers are used
  - A new RREQ by node S for a destination is assigned a higher destination sequence number.
     An intermediate node which knows a route, but with a smaller sequence number, cannot send RREP







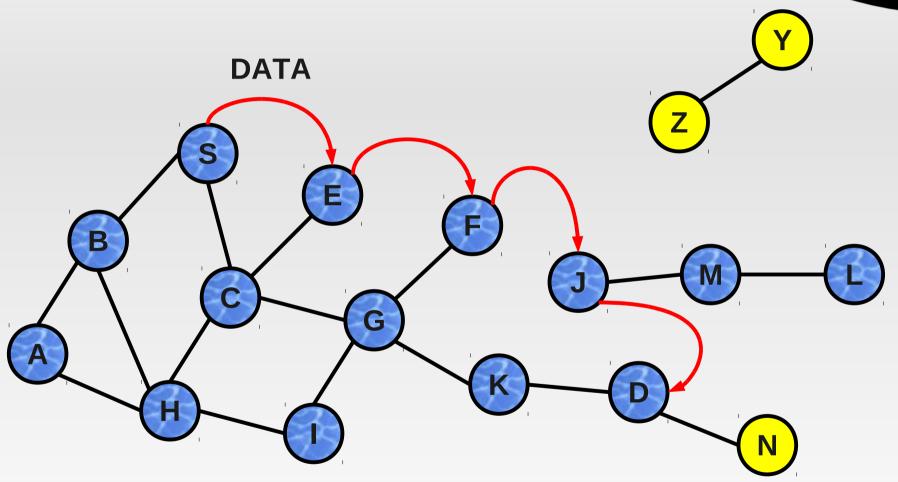


Forward links are setup when RREP travels along the reverse path

Represents a link on the forward path

ubuntu





Routing table entries used to forward data packet. Route is **not** included in packet header



#### **Timeout**



- A routing table entry maintaining a reverse path is purged after a timeout interval
  - timeout should be long enough to allow RREP to come back
- A routing table entry maintaining a forward path is purged if not used for a active\_route\_timeout interval
  - if no data is being sent using a particular routing table entry, that entry will be deleted from the routing table (even if the route may actually still be valid)



#### Link Failure Reporting



- A neighbor of node X is considered active for a routing table entry, if the neighbor sent a packet within active\_route\_timeout interval which was forwarded using that entry
- When the next hop link in a routing table entry breaks, all active neighbors are informed
- Link failures are propagated by means of Route Error (RERR) messages, which also update destination sequence numbers



#### Route Error



- When node X is unable to forward packet P (from node S to node D) on link (X,Y), it generates a RERR message
- Node X increments the destination sequence number for D cached at node X
- The incremented sequence number N is included in the RERR
- When node S receives the RERR, it initiates a new route discovery for D using destination sequence number at least as large as N
- When node D receives the RREQ with destination sequence number N, node D will set its sequence number to N, unless it is already larger than N



#### **Link Failure Detection**



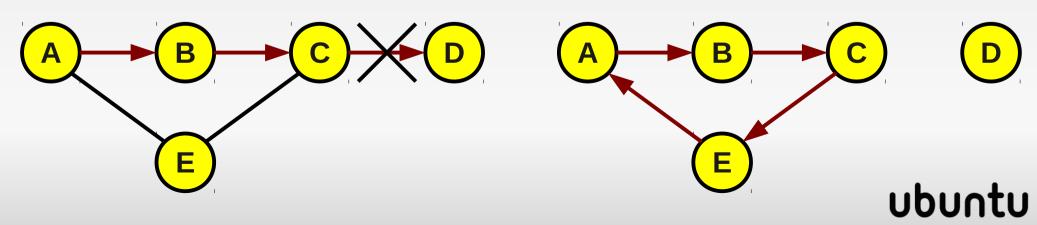
- Hello messages: Neighboring nodes periodically exchange hello message
- Absence of hello message is used as an indication of link failure
- Alternatively, failure to receive several MAClevel acknowledgement may be used as an indication of link failure



## Why Sequence Numbers in AODV



- To avoid using old/broken routes
  - To determine which route is newer
- To prevent formation of loops
  - Assume that A does not know about failure of link CD because RERR sent by C is lost
  - Now C performs a route discovery for D. Node A receives the RREQ (say, via path C-E-A)
  - Node A will reply since A knows a route to D via node B
  - Results in a loop (for instance, C-E-A-B-C )



# Optimization: Expanding Ring Search



- RREQs are initially sent with small Time-to-Live (TTL) field, to limit their propagation
  - DSR also includes a similar optimization
- If no Route Reply is received, then larger TTL tried





## **AODV for Linux**

#### **AODV for Linux**



- There exists several Linux implementations of the AODV routing protocols
- In this course we use AODV-UU, the version developed at Uppsala University
  - Most stable version
  - Available at http://sourceforge.net/projects/aodvuu/
- Drawback: it does not support correctly multiinterface devices
- We will force multi-hop communication setting filtering rules with iptables

ubuntu

#### **AODV-UU**



- AODV-UU implements the version 13 of the RFC 3561 (AODV)
- AODV-UU is composed of
  - A user-space process (aodvd) which implements the routing algorithm
  - A kernel-space loadable module (kaodv) which captures the data packets
- The software must be cross-compiled on the host machine
- Pre-compiled modules are available on the website of the course



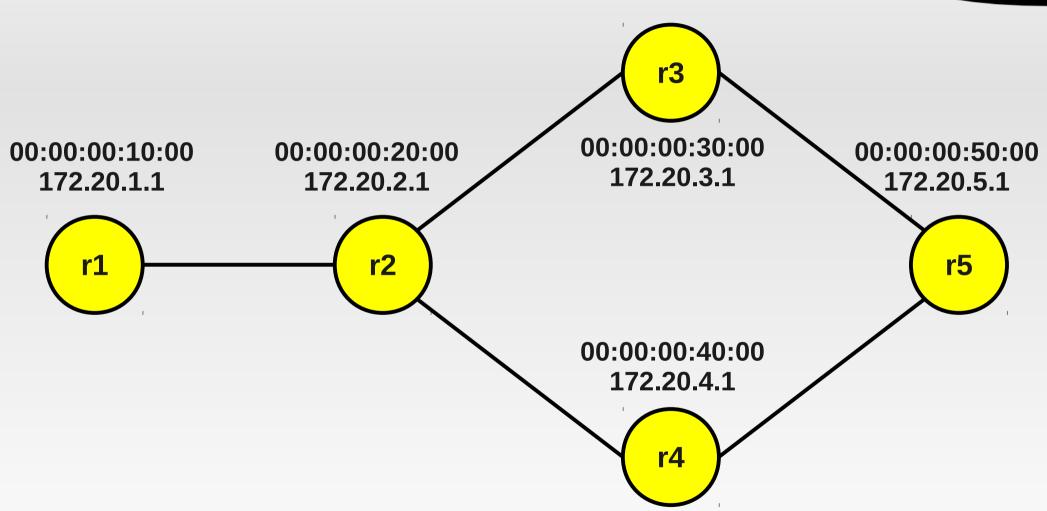
#### **AODV-UU**



AODV-UU File Edit View Terminal Help stefano@bender:~\$ ./aodvd -h Usage: aodvd [-dghjlouwxLDRV] [-i if0,if1,..] [-r N] [-n N] [-q THR] -d, --daemon Daemon mode, i.e. detach from the console. -g, --force-gratuitous Force the gratuitous flag to be set on all RREQ's. -h. --help This information. -i, --interface Network interfaces to attach to. Defaults to first wireless interface. -j, --hello-jitter Toggle hello jittering (default ON). -l, --log Log debug output to /var/log/aodvd.log. -o, --opt-hellos Send HELLOs only when forwarding data (experimental). -r, --log-rt-table Log routing table to /var/log/aodvd.rtlog every N secs. Receive N hellos from host before treating as neighbor. -n, --n-hellos -u, --unidir-hack Detect and avoid unidirectional links (experimental). -w, --gateway-mode Enable experimental Internet gateway support. -x, --no-expanding-ring Disable expanding ring search for RREQs. -D, --no-worb Disable 15 seconds wait on reboot delay. -L, --local-repair Enable local repair. Enable link layer feedback. -f, --llfeedback -R, --rate-limit Toggle rate limiting of RREQs and RERRs (default ON). -q, --quality-threshold Set a minimum signal quality threshold for control packets. -V, --version Show version. Erik Nordstrûm, <erik.nordstrom@it.uu.se> stefano@bender:~\$

### **Network Topology**







#### **Network Topology**



- Multi-hop communication is forced defining a white-list:
  - List of MAC addresses from which frames are accepted
  - Any other frame containing a different src MAC address is discarded
- The routing algorithm, which operates at the network layer, does not see routing messages



#### **Network Topology**



- Example of white-list configuration
  - Accept all frames containig as source MAC address 00:00:00:20:00 (r1)
    - iptables -t mangle -A PREROUTING -m mac \\
      --mac-source 00:00:00:20:00 -j ACCEPT
  - Drop all frames
     iptables -t mangle -A PREROUTING -i eth0 -j DROP
- The order of the two commands is important, since it defines the triggering order of the rules



#### Run AODV protocol



 After having set up the network interface, we start the sw daemon implementing the AODV routing algorithm

aodvd -I -r 3 -i eth0 -d

- Log both the aodvd messages (-I) and the routing table updates (-r)
  - The routing table is sampled every 3 seconds
- Attach the daemon on te network interface eth0
- Run in background (-d)

