

COMPUTER ORGANIZATION

Homework Two

(100 points equally distributed across all questions)

1. Consider this simple ISA:

Opcode	Register 1	Register 2
8 bits	4 bits	4 bits

What is the size of the register file (or how many registers can be addressed)? If each register's size is the same as the instruction's size and the memory address size, what is the size of the memory in bytes? How do you update the program counter in this case?

There are 16 addressable registers. The size of memory is 2 bytes. The Program Counter (PC) is updated by adding '2' to the current PC.

2. Complete this table with the content of the register file after each operation, considering the memory's state as shown bellow. This is a two-operand ISA and for auto decrement, d=2:

Register file:

	Initial state	Sub R0, R0	Add R2, #2	Add R3, (4)	Add R4, 10(R2)	Add R5, (R4+R3)	Add R6, @ (R5)	Add R7, -(R2)
R0	1	0	0	0	0	0	0	0
R1	1	1	1	1	1	1	1	1
R2	1	1	3	3	3	3	3	1
R3	1	1	1	2	2	2	2	2
R4	1	1	1	1	3	3	3	3
R5	1	1	1	1	1	4	4	4
R6	1	1	1	1	1	1	5	5
R7	1	1	1	1	1	1	1	3

Memory content:

Address	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Content	5	4	3	2	1	3	9	8	7	11	5	4	7	2	1	0

3. Write a procedure *Reduce* that takes a vector with n elements, adds all the values, and stores the result $\$v0$. The address of the original vector is in $\$a1$ and the number of elements is stored in $\$a0$.

Reduce:

move $\$t0, \$a0$	-- move n length into temp reg.
move $\$t1, \$a1$	-- move vector address into temp reg.
lw $\$t2, (\$t1) +$	-- load value at current pointer, increment
sub $\$t0, \$t0, \#1$	-- decrement number of vecs. left in array.
add $\$v0, \$v0, \$t2$	-- add loaded value to total
bne $\$t0, 0, \text{Reduce}$	-- repeat until done with whole array.
jr $\$ra$	-- return to caller (necessary?)

4. The procedure below uses 3 numbers stored in $\$a0$, $\$a1$ and $\$a2$.

addi $\$sp, \$sp, -4$	--moving by a word on the stack (alloc): $SP \leftarrow SP - 4$
sw $\$t0, 0(\$sp)$	--store value in $\$t0$ at the stack pointer: $Mem[0 + SP] \leftarrow T0$
add $\$t0, \$a0, \$a1$	--add $\$a0$ and $\$a1$ and put result in $\$t0$: $T0 \leftarrow A0 + A1$
add $\$v0, \$t0, \$a2$	--add $\$t0$ and $\$a2$ and put result in $\$v0$: $V0 \leftarrow T0 + A2$
lw $\$t0, 0(\$sp)$	--retrieve old $\$t0$ from stack: $T0 \leftarrow Mem[0 + SP]$
addi $\$sp, \$sp, 4$	--de-allocate mem on stack: $SP \leftarrow SP + 4$
jr $\$ra$	--return to caller: (no RTN?)

- a. Write comments by the side of each instruction using RTN that describe its functionality.
- b. What is the code doing?

This code first is allocating a word on the stack and storing some value from $\$t0$ in that new allocation. The code then adds the values in $\$a0$ and $\$a1$, using $\$t0$ as a temporary location. This half result is then added to $\$a2$ and this final result is put in $\$v0$. The value that was in $\$t0$ and stored on the stack is now retrieved from the stack and placed back in $\$t0$. The stack is reset and the procedure ends.

- c. Write a procedure *Add4* that calls this one as a procedure to add 4 numbers. Use QTSPIM to test it and show the screen captures.

```

1      .globl main
2  main:
3      jal Add4
4
5  Add4: move $v1, $ra      # save return address
6      li $a0, 3           # load 4 random numbers to add together
7      li $a1, 5
8      li $a2, 7
9      li $a3, 10
10     jal Sumn            # jump and link to given procedure
11     add $s1, $v0, $a3   # add fourth value since the procedure only adds three
12     jr $v1
13
14
15  Sumn: addi $sp, $sp, -4  # moving by a word on the stack (alloc): SP <- SP - 4
16       sw $t0, 0($sp)    # store value in $t0 at the stack pointer: Mem[0 + SP] <- T0
17       add $t0, $a0, $a1  # add $a0 and $a1 and put result in $t0: T0 <- A0 + A1
18       add $v0, $t0, $a2  # add $t0 and $a2 and put result in $v0: V0 <- T0 + A2
19       lw $t0, 0($sp)    # retrieve old $t0 from stack: T0 <- Mem[0 + SP]
20       addi $sp, $sp, 4   # de-allocate word on stack: SP <- SP + 4
21       jr $ra            # return to caller: (no RTN?)
22
23

```

QtSpim

File Simulator Registers Text Segment Data Segment Window Help

FP Regs Int Regs [10] Data Text

Int Regs [10]

PC = 4194372
EPC = 0
Cause = 0
BadVAddr = 0
Status = 805371664

HI = 0
LO = 0

R0 [r0] = 0
R1 [at] = 0
R2 [v0] = 15
R3 [v1] = 4194368
R4 [a0] = 3
R5 [a1] = 5
R6 [a2] = 7
R7 [a3] = 10
R8 [t0] = 0
R9 [t1] = 0
R10 [t2] = 0
R11 [t3] = 0
R12 [t4] = 0
R13 [t5] = 0
R14 [t6] = 0
R15 [t7] = 0
R16 [s0] = 0
R17 [s1] = 25
R18 [s2] = 0
R19 [s3] = 0
R20 [s4] = 0
R21 [s5] = 0
R22 [s6] = 0
R23 [s7] = 0
R24 [t8] = 0
R25 [t9] = 0
R26 [k0] = 0
R27 [k1] = 0

result

00400000 8fa40000 lw \$4, 0(\$29) ; 183: lw \$a0 0(\$sp) # argv
00400004 27a50004 addiu \$5, \$29, 4 ; 184: addiu \$a1 \$sp 4 # argv
00400008 24a60004 addiu \$6, \$5, 4 ; 185: addiu \$a2 \$a1 4 # envp
0040000c 00041080 sll \$2, \$4, 2 ; 186: sll \$v0 \$a0 2
00400010 00c23021 addu \$6, \$6, \$2 ; 187: addu \$a2 \$a2 \$v0
00400014 0c100009 jal 0x00400024 [main] ; 188: jal main
00400018 00000000 nop ; 189: nop
0040001c 3402000a ori \$2, \$0, 10 ; 191: li \$v0 10
00400020 0000000c syscall ; 192: syscall # syscall 10 (exit)
00400024 0c10000a jal 0x00400028 [Add4] ; 3: jal Add4
00400028 001f1821 addu \$3, \$0, \$31 ; 5: move \$v1, \$ra # save return address
0040002c 34040003 ori \$4, \$0, 3 ; 6: li \$a0, 3 # load 4 random numbers to add together
00400030 34050005 ori \$5, \$0, 5 ; 7: li \$a1, 5
00400034 34060007 ori \$6, \$0, 7 ; 8: li \$a2, 7
00400038 3407000a ori \$7, \$0, 10 ; 9: li \$a3, 10
0040003c 0c100012 jal 0x00400048 [Sumn] ; 10: jal Sumn # jump and link to given procedure
00400040 00479820 add \$17, \$2, \$7 ; 11: add \$s1, \$v0, \$a3 # add fourth value since the procedure only adds three
00400044 00600009 jr \$2 ; 12: jr \$v1
00400048 23bdfffc addi \$29, \$29, -4 ; 15: addi \$sp, \$sp, -4 # moving by a word on the stack (alloc): SP
0040004c afa30000 sw \$0, 0(\$29) ; 16: sw \$t0, 0(\$sp) # store value in \$t0 at the stack pointer: Mem[0 + SP]
00400050 00854020 add \$0, \$4, \$5 ; 17: add \$t0, \$a0, \$a1 # add \$a0 and \$a1 and put result in \$t0: T0
00400054 01061020 add \$2, \$0, \$6 ; 18: add \$v0, \$t0, \$a2 # add \$t0 and \$a2 and put result in \$v0: V0
00400058 8fa30000 lw \$0, 0(\$29) ; 19: lw \$t0, 0(\$sp) # retrieve old \$t0 from stack: T0
0040005c 23bd0004 addi \$29, \$29, 4 ; 20: addi \$sp, \$sp, 4 # de-allocate word on stack: SP
00400060 03e00008 jr \$31 ; 21: jr \$ra # return to caller: (no RTN?)
Kernel Text Segment [80000000]..[80010000]
80000180 0001d821 addu \$27, \$0, \$1 ; 90: move \$k1 \$at # Save \$at
80000184 3c019000 lui \$1, -28672 ; 92: sw \$v0 \$1 # Not re-entrant and we can't trust \$sp
80000188 ac220200 sw \$2, 512(\$1)
8000018c 3c019000 lui \$1, -28672 ; 93: sw \$a0 \$2 # But we need to use these registers
80000190 ac240204 sw \$4, 516(\$1)
80000194 401a6800 mfc0 \$26, \$13 ; 95: mfc0 \$k0 \$13 # Cause register
80000198 001a2082 srl \$4, \$26, 2 ; 96: srl \$a0 \$k0 2 # Extract ExcCode Field

Memory and registers cleared

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Memory and registers cleared

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