

My Experience

Monday, December 24, 2012

Segment Trees and lazy propagation

In this topic i will explain a very interesting data structure that can be used to solve a specific set of problems. I will start by explaining its definition and the proceeding with an example problem to solve with it.

Table of contents:

- What is segment trees?
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What is segment trees?

Segment Trees is a Tree data structure for storing intervals, or segments. It allows querying which of the stored segments contain a given point. It is, in principle, a static structure; that is, its content cannot be modified once the structure is built. It only uses $O(N \lg(N))$ storage.

A segment trees has only three operations: `build_tree`, `update_tree`, `query_tree`.

Building tree: To init the tree segments or intervals values

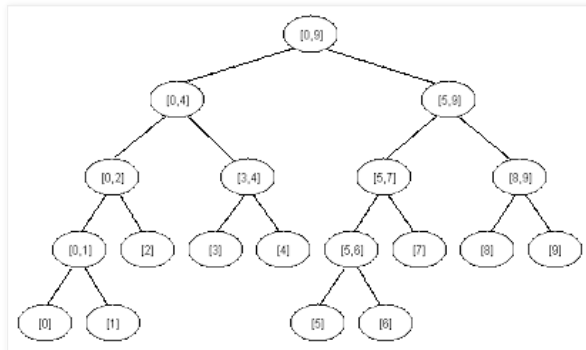
Update tree: To update value of an interval or segment

Query tree: To retrieve the value of an interval or segment

Example Segment Tree:

- The first node will hold the information for the interval $[i, j]$
- If $i < j$ the left and right son will hold the information for the intervals $[i, (i+j)/2]$ and $[(i+j)/2+1, j]$

Notice that the height of a segment tree for an interval with N elements is $\lceil \log N \rceil + 1$. Here is how a segment tree for the interval $[0, 9]$ would look like:



Order of growth of segment trees operations

- **build_tree:** $O(N \lg(N))$
- **update_tree:** $O(\lg(N + k))$
- **query_tree:** $O(\lg(N + k))$

K = Number of retrieved intervals or segments

Show me your code

```
1  /**
2   * In this code we have a very large array called arr, and very large set
3   * Operation #1: Increment the elements within range [i, j] with value va
```

Clustermaps

About Me

[se7so](#)

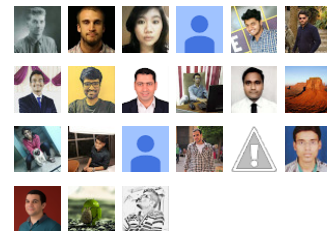
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```

4  * Operation #2: Get max element within range [i, j]
5  * Build tree: build_tree(1, 0, N-1)
6  * Update tree: update_tree(1, 0, N-1, i, j, value)
7  * Query tree: query_tree(1, 0, N-1, i, j)
8  */
9
10 #include<iostream>
11 #include<algorithm>
12 using namespace std;
13
14 #include<string.h>
15 #include<math.h>
16
17 #define N 20
18 #define MAX (1+(1<<6)) // Why? :D
19 #define inf 0x7fffffff
20
21 int arr[N];
22 int tree[MAX];
23
24 /**
25  * Build and init tree
26  */
27 void build_tree(int node, int a, int b) {
28     if(a > b) return; // Out of range
29
30     if(a == b) { // Leaf node
31         tree[node] = arr[a]; // Init value
32         return;
33     }
34
35     build_tree(node*2, a, (a+b)/2); // Init left child
36     build_tree(node*2+1, 1+(a+b)/2, b); // Init right child
37
38     tree[node] = max(tree[node*2], tree[node*2+1]); // Init root value
39 }
40
41 /**
42  * Increment elements within range [i, j] with value value
43  */
44 void update_tree(int node, int a, int b, int i, int j, int value) {
45
46     if(a > b || a > j || b < i) // Current segment is not within range
47         return;
48
49     if(a == b) { // Leaf node
50         tree[node] += value;
51         return;
52     }
53
54     update_tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
55     update_tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right child
56
57     tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root value
58 }
59
60 /**
61  * Query tree to get max element value within range [i, j]
62  */
63 int query_tree(int node, int a, int b, int i, int j) {
64
65     if(a > b || a > j || b < i) return -inf; // Out of range

```

```

66
67     if(a >= i && b <= j) // Current segment is totally within range [
68         return tree[node];
69
70     int q1 = query_tree(node*2, a, (a+b)/2, i, j); // Query left child
71     int q2 = query_tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right child
72
73     int res = max(q1, q2); // Return final result
74
75     return res;
76 }
77
78 int main() {
79     for(int i = 0; i < N; i++) arr[i] = 1;
80
81     build_tree(1, 0, N-1);
82
83     update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
84     update_tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12
85     update_tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100
86
87     cout << query_tree(1, 0, N-1, 0, N-1) << endl; // Get max element
88 }

```

segment_tree.cpp hosted with ❤ by GitHub

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Lazy Propagation

Sometimes a segment tree operation wouldn't survive if the problem constraints is too large, here it come lazy propagation along with the segment tree.

In the current version when we update a range, we branch its childs even if the segment is covered within range. In the lazy version we only mark its child that it needs to be updated and update it when needed.

```

1  /**
2   * In this code we have a very large array called arr, and very large segment tree
3   * Operation #1: Increment the elements within range [i, j] with value value
4   * Operation #2: Get max element within range [i, j]
5   * Build tree: build_tree(1, 0, N-1)
6   * Update tree: update_tree(1, 0, N-1, i, j, value)
7   * Query tree: query_tree(1, 0, N-1, i, j)
8   */
9
10 #include<iostream>
11 #include<algorithm>
12 using namespace std;
13
14 #include<string.h>
15 #include<math.h>
16
17 #define N 20
18 #define MAX (1+(1<<6)) // Why? :D
19 #define inf 0x7fffffff
20
21 int arr[N];
22 int tree[MAX];
23 int lazy[MAX];
24
25 /**
26  * Build and init tree
27  */

```

```

28 void build_tree(int node, int a, int b) {
29     if(a > b) return; // Out of range
30
31     if(a == b) { // Leaf node
32         tree[node] = arr[a]; // Init value
33         return;
34     }
35
36     build_tree(node*2, a, (a+b)/2); // Init left child
37     build_tree(node*2+1, 1+(a+b)/2, b); // Init right child
38
39     tree[node] = max(tree[node*2], tree[node*2+1]); // Init root val
40 }
41
42 /**
43  * Increment elements within range [i, j] with value value
44  */
45 void update_tree(int node, int a, int b, int i, int j, int value) {
46
47     if(lazy[node] != 0) { // This node needs to be updated
48         tree[node] += lazy[node]; // Update it
49
50         if(a != b) {
51             lazy[node*2] += lazy[node]; // Mark child as lazy
52             lazy[node*2+1] += lazy[node]; // Mark child as lazy
53         }
54
55         lazy[node] = 0; // Reset it
56     }
57
58     if(a > b || a > j || b < i) // Current segment is not within range
59         return;
60
61     if(a >= i && b <= j) { // Segment is fully within range
62         tree[node] += value;
63
64         if(a != b) { // Not leaf node
65             lazy[node*2] += value;
66             lazy[node*2+1] += value;
67         }
68
69         return;
70     }
71
72     update_tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
73     update_tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right child
74
75     tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root
76 }
77
78 /**
79  * Query tree to get max element value within range [i, j]
80  */
81 int query_tree(int node, int a, int b, int i, int j) {
82
83     if(a > b || a > j || b < i) return -inf; // Out of range
84
85     if(lazy[node] != 0) { // This node needs to be updated
86         tree[node] += lazy[node]; // Update it
87
88         if(a != b) {
89             lazy[node*2] += lazy[node]; // Mark child as lazy

```

```

90         lazy[node*2+1] += lazy[node]; // Mark child as lazy
91     }
92
93     lazy[node] = 0; // Reset it
94 }
95
96 if(a >= i && b <= j) // Current segment is totally within range
97     return tree[node];
98
99 int q1 = query_tree(node*2, a, (a+b)/2, i, j); // Query left child
100 int q2 = query_tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right child
101
102 int res = max(q1, q2); // Return final result
103
104 return res;
105 }
106
107 int main() {
108     for(int i = 0; i < N; i++) arr[i] = 1;
109
110     build_tree(1, 0, N-1);
111
112     memset(lazy, 0, sizeof lazy);
113
114     update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
115     update_tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12
116     update_tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100
117
118     cout << query_tree(1, 0, N-1, 0, N-1) << endl; // Get max element
119 }

```

lazy_segment_tree.cpp hosted with ♥ by GitHub

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Note: Read my solution for problem [Can you answer these queries I](#) in this [article](#).

Sample Problems to try

- [Quadrant Queries](#)
- [D-Query](#)
- [Can You answer these queries I](#)

References

- [Wiki](#)
- [Topcoder tutorials](#)

Posted by [se7so](#) at 4:58 AM

58 comments:



[donngghi](#) December 24, 2012 at 12:25 PM

In lazy's version, i thinks it's better if you replace update tree[2*node] and tree[2*node+1] in 49th, 50th, 80th and 81th line by lazy[2*node] and lazy[2*node+1].

Its reason is your query is not really come down to higher level, so lazy[] should be updated

[Reply](#)



[Hussein El-Sayed](#) December 24, 2012 at 10:44 PM

I can't understand you :)

[Reply](#)

**donngghi** December 25, 2012 at 3:03 AM

So, in line 80th:
`tree[node*2] += lazy[node]; // Mark child as lazy`
`tree[node*2+1] += lazy[node]; // Mark child as lazy`

=> replaced by:
`lazy[node*2] += lazy[node];`
`lazy[node*2+1] += lazy[node];`

It's correct ?

[Reply](#)

**Hussein El-Sayed** December 25, 2012 at 3:29 AM

Yes you are totally right :). thanks for correcting me ;).

[Reply](#)

**Hussein El-Sayed** December 25, 2012 at 3:30 AM

The same at line 49 and 50, updated check it now and tell me :)

[Reply](#)

[Replies](#)

**Max Li** April 17, 2015 at 9:07 AM

f

**Max Li** April 17, 2015 at 9:10 AM

At lines 72 and 73, shouldn't this be replaced by a `lazy[node]`? sorry, I am a beginner.

[Reply](#)

**Sandipan Manna** December 31, 2012 at 12:13 PM

`update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5`
`update_tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12`
`update_tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100`

So, The maximum element in the whole array should be 101
 and your final array becomes
`[6 6 6 6 6 6 13 13 13 13 13 101 101 101 101 101 101 101]`

but your program gives output as 117 !!!

[Reply](#)

**Hussein El-Sayed** January 1, 2013 at 3:06 AM

No it should be 113, however the size of the array needs to be $(1+(1<<6))$ as it should be $2^{(1+\lg N)}$.

Also there was some checks needed to be added in the lazy version.. please check it and get back to me.

[Reply](#)

**Sandipan Manna** January 11, 2013 at 11:19 PM

Yes your segment tree size should be

`int x = (int)(ceil(log2(N)))+1;`
`size = (1 << x);`
 This one!

[Reply](#)

**InfiniteComplexity** February 16, 2013 at 3:27 AM

Wow! Thanks for this post, it was very helpful! However, I'm trying to implement another `update_tree_val` function that sets the values from a range to one value. e.g. `update_tree_val(3,7,4)` would set the range `[3,7]` to the value of 4. How can I do this using lazy propagation on your tree?

Thanks

[Reply](#)

Replies

**Hussein El-Sayed** February 16, 2013 at 11:33 PM

It would be the same but without incrementing..

**Aditya Choudhary** November 9, 2014 at 8:58 AM

This means in line 48 ,I have to modify it to `tree[node] = lazy[node];`
 and in line 62, `tree[node] = value; ?`

Reply

**PRASHANTH SOUNDAR** February 19, 2013 at 9:37 AM

This comment has been removed by the author.

Reply

**Tilak Raj Singh** July 26, 2013 at 6:08 PM

in line 86 it should be
`tree[node] += (b-a+1)*lazy[node];`

Reply

Replies

**ravi** November 2, 2014 at 8:22 PM

I think you forgot that query returns maximum from given range
 not the sum of elements of given range

**Sourabh Khandelwal** January 31, 2017 at 8:45 AM

This is true only when your query is range sum and not max range.

Reply

**rl** September 15, 2013 at 10:15 AM

In order to do the following:

- 1.) Add x to `A[i], A[i+1], ..., A[j]`
- 2.) Output the sum of `A[i], A[i+1], ..., A[j]`

I simply replaced `max()` with `sum()` however i am not getting the correct answer.

Reply

Replies

**ravi** November 2, 2014 at 8:25 PM

also replace line 48 and 86 by `tree[node] += (b - a + 1) * lazy[node];`
 and line 62 by `tree[node] += (b - a + 1) * value;`

**Matheus dall rosa** October 16, 2015 at 6:36 PM

This comment has been removed by the author.

Reply

**হুমতিয়াজ** November 13, 2013 at 11:37 AM

This comment has been removed by the author.

Reply

**Gautam Singh** May 24, 2014 at 12:55 AM

in `query_tree()` method we could have updated the lazy value to the tree before we look for the out
 of range condition.... it would result in better performance...

am i right about it....please correct me if I am wrong!!

Reply

Replies

**Unknown** August 25, 2018 at 4:02 AM

yes whenever u reach a node that has pending updates ..be it in the update or query function u should always update that node.this should be a good practise.

[Reply](#)**VIPUL JAIN** August 1, 2014 at 7:16 AM

Can you please elaborate how to implement DQUERY?

[Reply](#)[Replies](#)**Gaurav Singh** March 12, 2015 at 1:27 PM

The best link is:
<http://apps.topcoder.com/forums/?sessionId=C7BE6D64D8F4953865BCE8B7945FA2F6?module=Thread&threadID=627423&start=0&mc=13#1060242>

[Reply](#)**Ashish Tilokani** September 12, 2014 at 1:36 AM

<http://discuss.codechef.com/questions/50866/segment-trees-doubt>

[Reply](#)**Ashu Pachauri** September 25, 2014 at 10:11 PM

Very helpful

[Reply](#)**Rahul Kumar** December 18, 2014 at 1:23 AM

we have to take array size of 2^n for make segment tree of n element array,then how to make segment tree of 30 or greater elements,because $2^{30} = 1073741824$
 then how to take array of this larger size for make segment tree,how to implement ?

[Reply](#)**Hussein El-Sayed** December 18, 2014 at 1:27 AM

It should be $2^{(1+\lg N)}$ not 2^N

[Reply](#)**ম্যাট্রিক্স.কোড** January 6, 2015 at 2:36 PM

This comment has been removed by the author.

[Reply](#)**ম্যাট্রিক্স.কোড** January 6, 2015 at 2:43 PM

Very good and nice blog..
 I am following your code structure in my SG tree implementation :)

[Reply](#)**Jayaram Prabhu Durairaj** January 7, 2015 at 7:49 PM

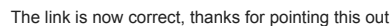
This comment has been removed by the author.

[Reply](#)**Anonymous** January 18, 2015 at 11:33 PM

Hey Note: Read my solution for problem Can you answer these queries I in this article.

The article doesn't point to the solution. It's pointing to the problem itself. Please correct it.

[Reply](#)[Replies](#)**Hussein El-Sayed** January 19, 2015 at 4:12 AM



Gaurav Singh March 12, 2015 at 1:36 PM

Reply

Anonymous March 12, 2015 at 11:38 PM

nice post [good writing service for scholarship essays](#)

Reply



Can you do range multiplication with segment trees?

Reply

Replies



of course. its same as update value with $k * x$ (x is elements in the range)

Reply



How to initialize values in a range i.e $A[i]=x$ (some integer) using update function in lazy propagation?

Reply



If I change $(2 * \text{node})$ with $(2 * \text{node} + 1)$ and $(2 * \text{node} + 1)$ with $(2 * \text{node} + 2)$ and start the node position with 0 instead of 1. Will there be any changes in the result? Why have you started node of segment tree with 1?

[Reply](#)

in 46 line when will $a > b$ condition becomes true?.

Reply



well i still cant understand the lazy part.....what does the lazy[] array stores ?.....i mean for a particular node

Reply



I cannot say if I understood Lazy propagation part.

[Reply](#)

Great Explanation Of Segment trees

please explain more the lazy propagation part i cantnt understand it.

[Reply](#)

Sky Blue August 13, 2015 at 4:35 AM



i don't understand

[Reply](#)**Thirupathi Reddy** August 24, 2015 at 5:11 PM

thank u for your post.....really helped me.....

[Reply](#)**Pradyumna Bang** September 7, 2015 at 3:49 AM

If you want to learn this in a real easy way, watch tushar roy's youtube video on segment tree along with this blog post.

[Reply](#)**P.Vijay Kumar** October 15, 2015 at 5:49 PM

D-Query can be easily solved using fast io + MO's algorithm...But using BIT is safer to avoid TLE sometimes...Nice Post...!!!

[Reply](#)**Ashish Yadav** January 22, 2016 at 11:36 PM

This comment has been removed by the author.

[Reply](#)**Suman Patel** June 21, 2016 at 12:55 AM

the above content you are posted is very nice.

[kapil sharma new show online](#)[kapil sharma show hit](#)[bigg boss season 10 promo](#)[bigg boss season 10 episodes](#)[one night stand sexy pics](#)[Photo of Sunny Leone](#)[good morning images download free](#)[best good morning images](#)[Reply](#)**Unknown** September 14, 2016 at 11:21 PM

what if i have to update each nodes with different values means i m not updating whole range with unique value

[Reply](#)**Rahul Mehlawat** January 9, 2017 at 10:11 PM

Wow,, This is great blog... Thanks for sharing this informative Content here....

We offers [Tree Trimming in Sacramento](#), Contact us now for [Tree Pruning in Sacramento](#)Source: Best Tree Company in Sacramento www.cisnerostreecare.com[Reply](#)**Sarwar Jahan** September 27, 2017 at 12:08 AM

nice :)

really helpful one.

[Reply](#)**crapperGenius11** November 12, 2017 at 9:13 AM

Hey, this is a really intuitive post with which i learnt how a segment tree is coded!

Could you point to me some links that will explain how the worst case complexity of updating intervals is $O(\log(N+k))$?[Reply](#)**charlie** December 17, 2018 at 11:14 PM

The article is praiseworthy. It talks about Segment Trees and lazy propagation. The article will let you understand what are segment trees? According to the article, 'Segment Trees is a Tree data structure for storing intervals or segments. It allows querying which of the stored segments contain

a given point. A segment trees have only three operations such as build tree, update tree, and query tree.'

Thanks,
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Unknown March 25, 2019 at 4:19 AM

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Admin April 30, 2019 at 4:30 AM

This will help you to find detail about download festival:
https://calendar.prattlibrary.org/event/community_poetree#.XMgsVegzaMq

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usama habib noor July 19, 2019 at 3:39 AM

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