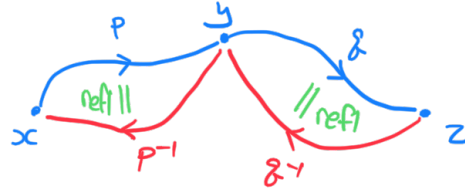


# Worksheet 3

HoTTEST Summer School 2022

The HoTTEST TAs, and **FRANK**  
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1 (★)



Let  $A$  be a type, and  $x, y, z : A$ . Show that path inversion distributes over concatenation, i.e., construct a term of the following type:

$$\prod_{p:x=y} \prod_{q:y=z} (p \cdot q)^{-1} = \underbrace{q^{-1}}_{z=y} \cdot \underbrace{p^{-1}}_{y=x}$$

$$(refl_x \cdot refl_y)^{-1} \\ \doteq refl_x^{-1} \doteq refl_z$$

$$refl_y^{-1} \cdot refl_x^{-1} \\ = refl_z \cdot refl_y \doteq refl_z$$

2 (★★)

Let  $A$  and  $B$  be types. We can define the product type as  $A \times B := \Sigma_A B$ , where  $B$  is considered a constant type family over  $A$ . The resulting elimination principle is:

$$\frac{\Gamma, z : A \times B \vdash D(z) \text{ type} \quad \Gamma, a : A, b : B \vdash d : D(a, b)}{\Gamma, z : A \times B \vdash \text{ind}_x(d, z) : D(z)} \times\text{-Elim}$$

Define the two projections  $\text{pr}_1 : A \times B \rightarrow A$  and  $\text{pr}_2 : A \times B \rightarrow B$  using the elimination principle.

$$\frac{\Gamma, z : A \times B \vdash A \text{ type} \quad \Gamma, a : A, b : B \vdash a : A}{\Gamma, z : A \times B \vdash \text{pr}_1(z) : A} \quad \frac{\Gamma, z : A \times B \vdash B \text{ type} \quad \Gamma, a : A, b : B \vdash b : B}{\Gamma, z : A \times B \vdash \text{pr}_2(z) : B} \\ \text{ind}_x(a, z) \quad \text{ind}_x(b, z)$$

Now use the elimination principle to give a term of the following type:

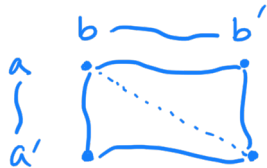
$$\prod_{z:A \times B} z = (\text{pr}_1 z, \text{pr}_2 z)$$

$$\text{proof } (a, b) = \text{refl}_{(a, b)}$$

### 3 (★)

Let  $A$  and  $B$  be types. Give an informal construction of a term of the following type:

$$\prod_{a,a':A} \prod_{b,b':B} (a =_A a') \times (b =_B b') \rightarrow ((a, b) =_{A \times B} (a', b'))$$



(We will later see that this map is an *equivalence*, and this map *characterises* paths in product types as pairs of paths between the components.)

### 4 (★★)

Let  $A$  be a type and  $a : A$ . Show that  $\text{refl}_a$  is unique among paths starting at  $a$  but with the other endpoint free. That is, for any  $z : \Sigma_{x:A} (a = x)$  construct a term

$$(a, \text{refl}_a) =_{\Sigma_{x:A} (a=x)} z$$

$\underbrace{(a', a = a')}_{\text{refl}_a}$

(We emphasize that this does not mean that  $\text{refl}_a$  is unique as a loop at  $a$ !)

(This is Prop. 5.5.1 in Egbert's book.)

### 5 (★★★)

Let  $A$  be a type. In the third lecture, a function **concat** of the following type was constructed:

$$\text{concat} : \prod_{x,y,z:A} (x = y) \rightarrow ((y = z) \rightarrow (x = z))$$

Call the above for  $\text{concat}_1$ . Define two different terms  $\text{concat}_2$  and  $\text{concat}_3$  of the same type.

Choose your favorite numbers  $i, j \in \{1, 2, 3\}$  ( $i \neq j$ ). Give a term of the following type:

$$\prod_{x,y,z:A} \prod_{p:x=y} \prod_{q:y=z} \text{concat}_i(p, q) = \text{concat}_j(p, q)$$

**6**     $(\star \star \star)$ 

Let  $A$  be a type with an element  $a : A$ . Can you construct a term of the following type? If not, what goes wrong?

$$\prod_{p:(a=Aa)} p =_{(a=Aa)} \text{refl}_a$$

No, this implies that every proof of identity is propositionally equal to  $\text{refl}_a$ , so it implies UIP, but UIP is not derivable.