#### The C10K Problem and Solutions

#### References:

- •http://www.kegel.com/c10k.html
- http://www.monkey.org/~provos/libevent
- •http://byteliu.com/2019/05/08/LINUX-%E2%80%93-IO-MULTIP LEXING-%E2%80%93-SELECT-VS-POLL-VS-EPOLL/
- •A Scalable and Explicit Event Delivery Mechanism for UNIX, <a href="http://static.usenix.org/event/usenix99/full\_papers/banga/banga.pdf">http://static.usenix.org/event/usenix99/full\_papers/banga/banga.pdf</a>
- •Acknowledgement: Modified from the slides of Che-Yi Lin and Hao-Yun Liu.

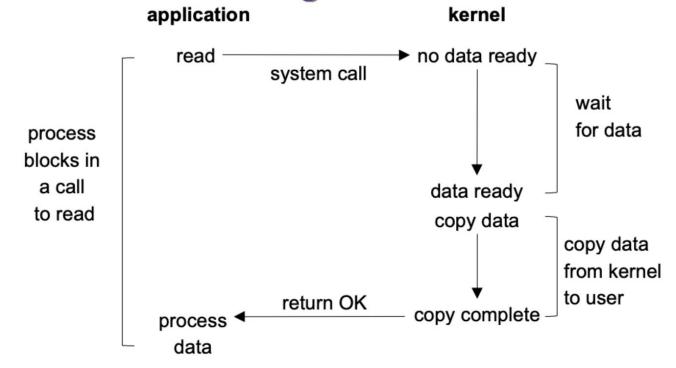
#### Outline

- Background
  - Types of I/O
  - Problem of select()
- The C10K problem
- Asynchronous I/O (AIO)
- Design of networking software
- Using libevent
- Benchmarks

#### Background: Types of I/O

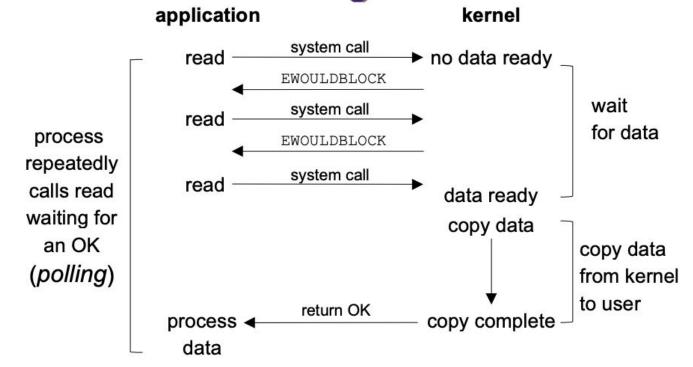
Wait until data is ready

### **Blocking I/O Model**



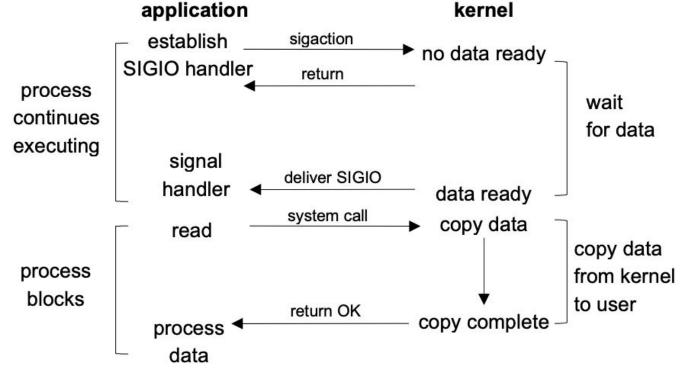
Return an error if data was not ready

#### Nonblocking I/O Model

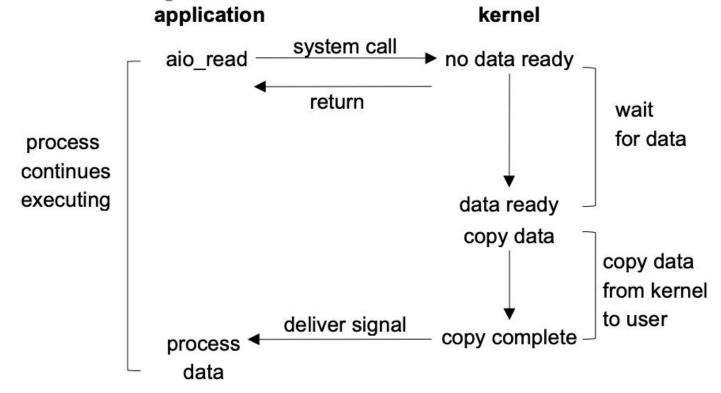


- Block until the given file descriptors are ready to perform I/O
- Block select() or poll(), instead of blocking I/O
  - select() I/O Multiplexing Model poll() epoll() application kernel system call select no data ready process blocks wait waiting for for data one of many fds return readable data ready system call copy data read copy data process from kernel blocks to user return OK copy complete process data

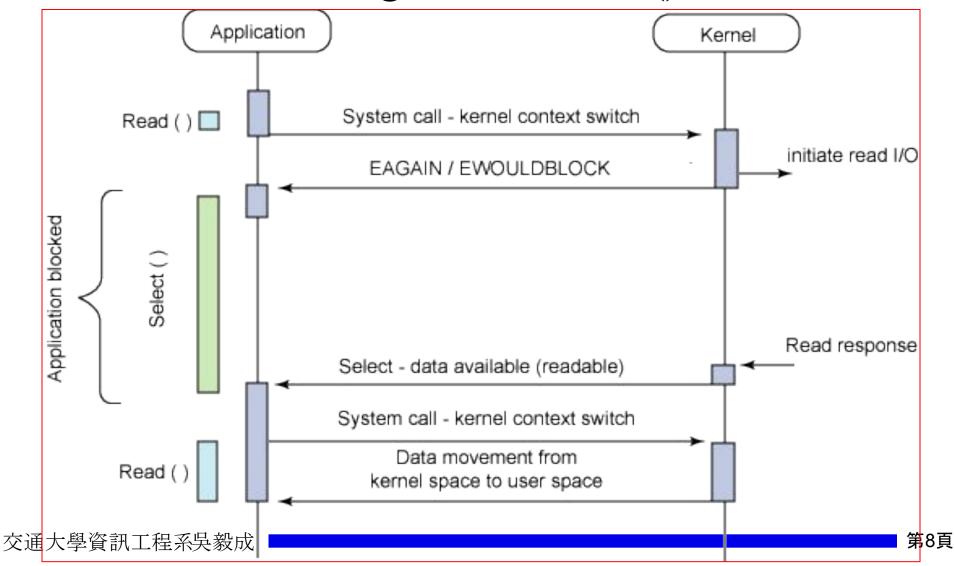
# Signal Driven I/O Model



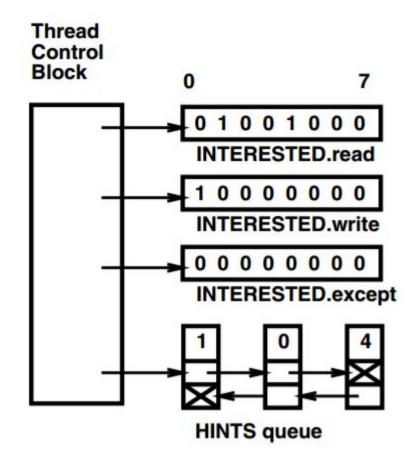
## Asynchronous I/O Model



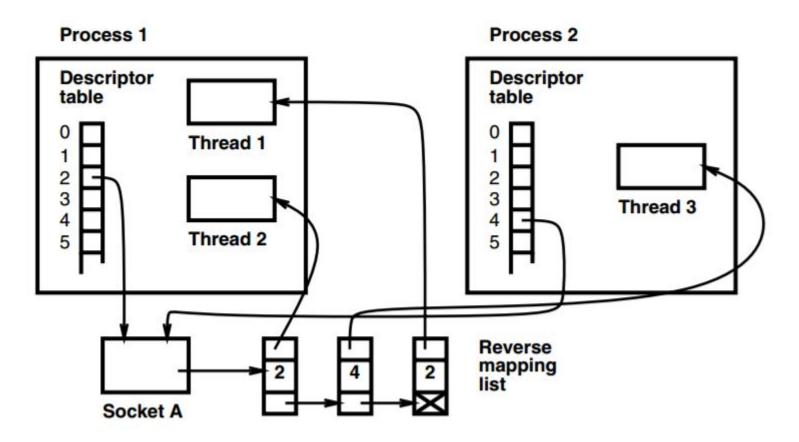
## Background: select()



#### Per-Thread Data Structure



#### Per-Socket Data Structure



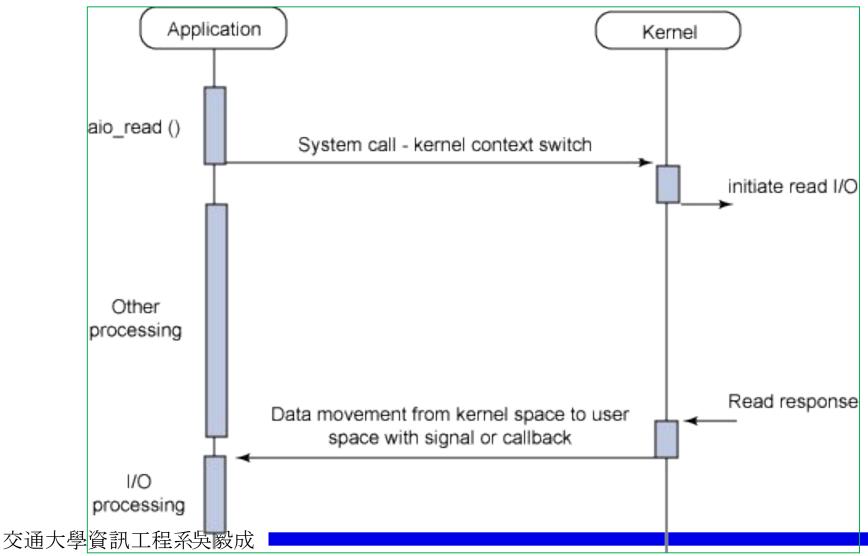
## The C10K problem

- Web servers have to handle ten thousand clients simultaneously
- Web is a big place now
- Hardware is no longer the bottleneck

http://www.kegel.com/c10k.html

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### Asynchronous I/O (AIO)



#### AIO in Linux

Non-blocking Blocking Read/wirte Synchronous Read/write (O\_NONBLOCK) i/O multiplexing Asynchronous AIO (select/poll)

#### AIO in Linux (cont.)

- Introduced in Linux kernel 2.6 (released at 2008) and also available in 2.4 if patched.
- The completion of I/O can be notified by two method.
  - Signal.
  - Register a completion handler function to create a new thread.
- API:
  - aio\_read
  - aio error
  - aio return
  - aio\_write
  - aio\_suspend
  - aio cancel
  - lio\_listio

### Design of application

#### • Use fork()

defect: High overhead for each connection

solution: Return to accept() and child process die automatically

example: Apache 1.3

#### • Use pthread create()

defect: Thread-safe and Memory-leak problems

solution: Use Thread-safe library and Garbage collection library

example: Apache 2.0 Thread MPM

### Design of application (cont.)

#### Event-based process

advantage: Without overhead of create process or thread, no need to use Share Memory or Mutex for process / thread

#### hard to implement:

- BSD = kqueue(), Linux = epoll(), Solaris = /dev/poll
   None of these are Standard!
- Buffering of nonblocking I/O

Solution: libevent library

### Using libevent

- libevent by Niels Provos

  <a href="http://www.monkey.org/~provos/libevent/">http://www.monkey.org/~provos/libevent/</a>
- Is a lightweight C I/O framework
- Support kqueue(), epoll(), /dev/poll, and the traditional select(), poll()
- Under 3-clause BSD license!!

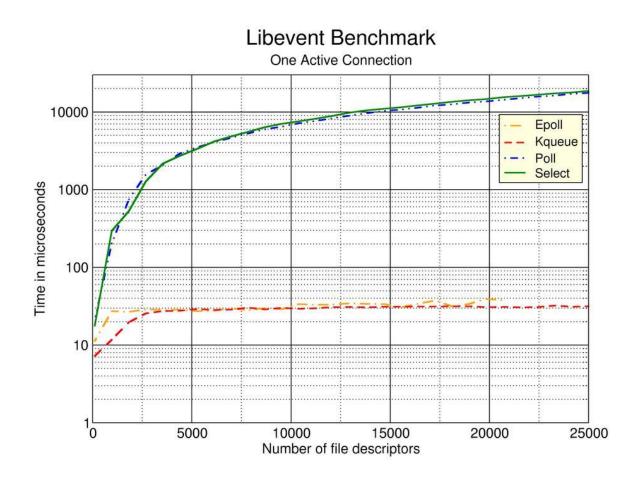
### Using libevent (cont.)

- Execute a callback function when a specific event occurs on a file descriptor or after a timeout has been reached
- Replace the event loop found in event driven network servers
- An application just needs to call event\_dispatch() and then add or remove events dynamically without having to change the event loop
- · Can also be used for multi-threaded applications

### Using libevent (cont.)

```
/* Initial libevent. */
event init();
/* Create event. */
struct event ev;
event_set(&ev, sfd, EV READ | EV PERSIST,
  connection accept, &ev);
/* Add event. */
event add(&ev, NULL);
event dispatch();
```

#### Benchmarks of libevent



#### Benchmarks (cont.)





