Lab 3 for uC/OS-II: Ceiling Priority Protocol

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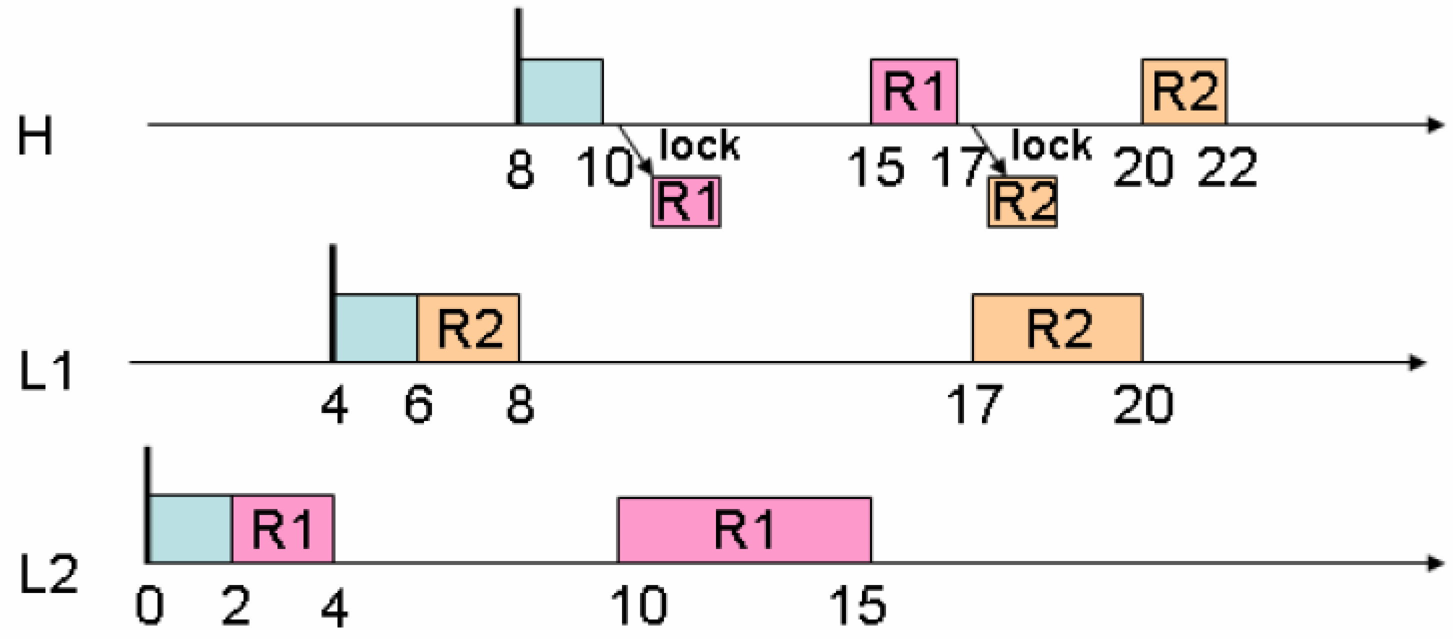
Objective

- To implement Ceiling Priority Protocol for ucOS's mutex locks

ucOS Mutex Locks

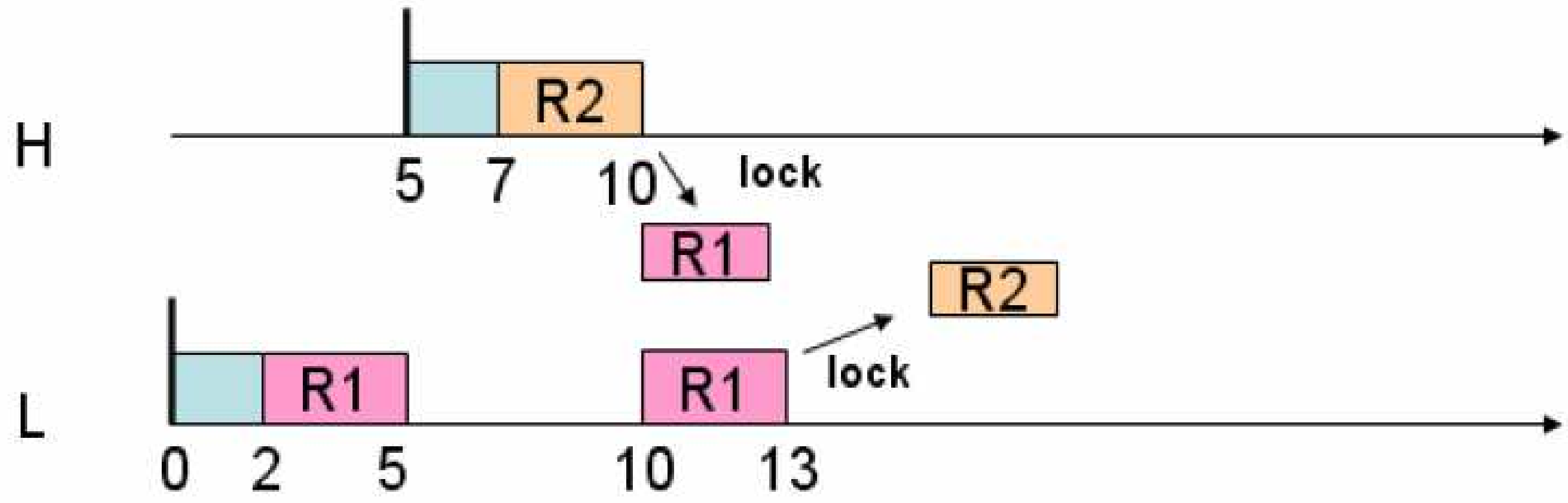
- Mutex lock is assigned to a priority upon creation
  - Its priority is higher than the priorities of all the tasks sharing the mutex lock
  - E.g., the lock's priority will be 2 if T_3 and T_4 share the lock
  - When a task blocks another task via a mutex lock, the task inherits the priority of the lock

Scenario #2
Multiple blocking in ucOS2 PIP



Task H is in turn blocked by task L1 and task L2

Scenario #2: Deadlock in uCOS-2 PIP



A deadlock

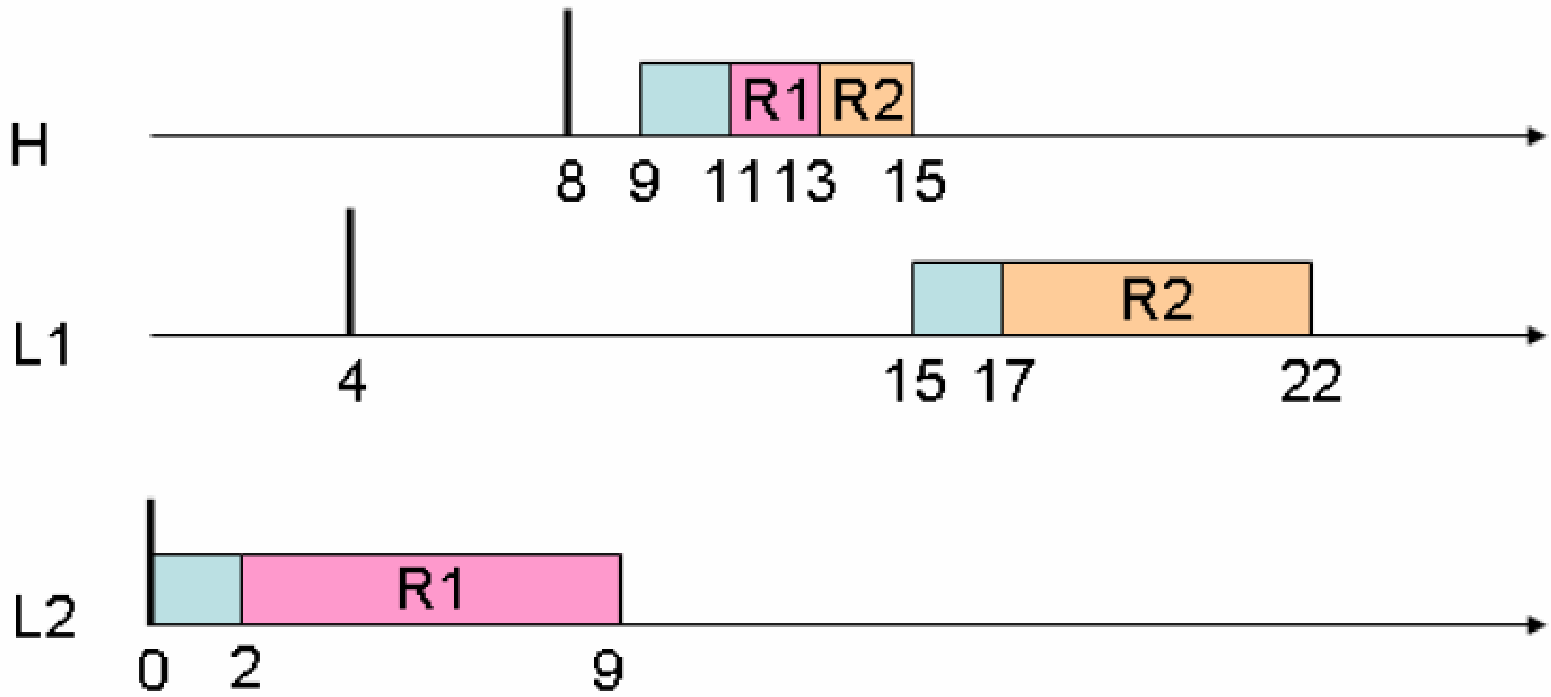
Disadvantages of PIP

- The “PIP” avoids uncontrolled priority inversion, but it has two disadvantages
  - A high priority task can be blocked multiple times
  - Deadlocks are possible

Ceiling Priority Protocol

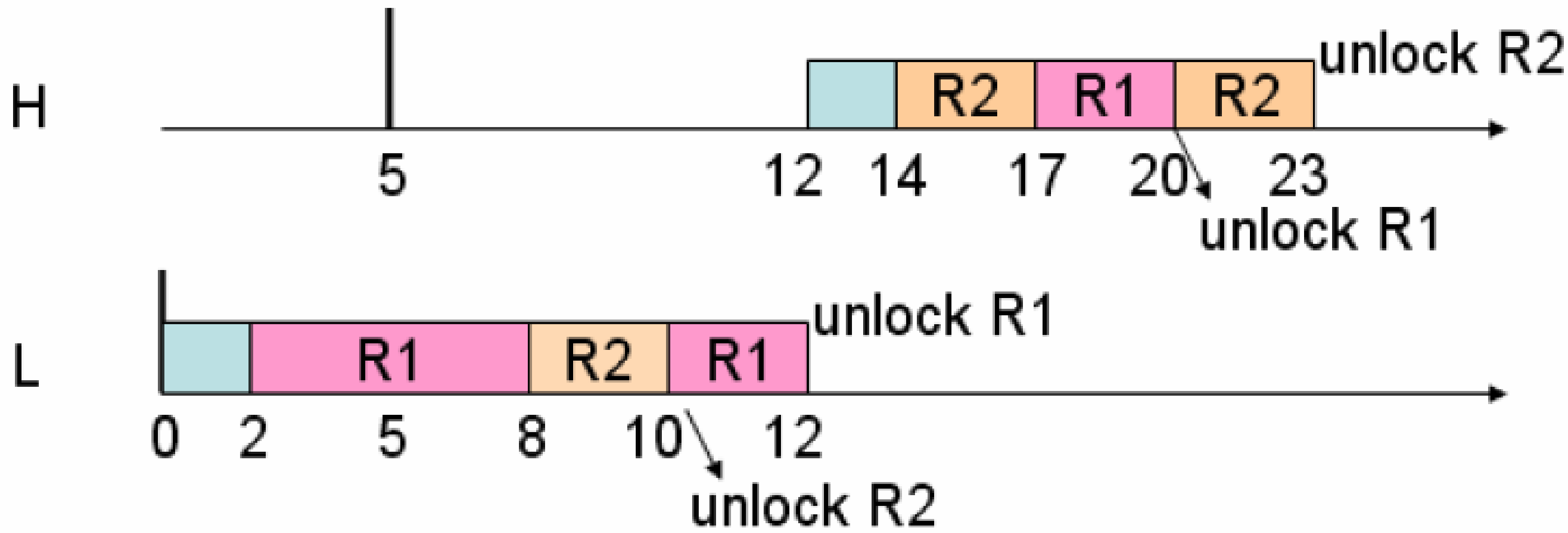
- Highest Locker's Priority Protocol
- The priority of the mutex lock is higher than the priorities of all the tasks sharing the lock
- Differently, when a task acquires a mutex lock, the task's priority boosts to the priority of the lock

CPP: Removing Multiple Blockings



The result of applying CPP

CPP: Avoiding Deadlocks



Deadlocks are avoided by using CPP

Implementation

- Reuse your code of Lab 0 (not Lab 1)

- Modify the two functions
  - OSMutexPend
    - Boost the locker's priority to the mutex priority
  - OSMutexPost
    - Restore the original priority of the locker

- Do not use OSTaskChangePrio to boost tasks' priorities
  - It calls OS_Sched() and results in unexpected behaviors

Implementation

- All tasks should add proper OSTimeDly() at their beginning to emulate their arrival times

- Emulate durations of CPU execution and resource use with your code from Lab 0
  - 2 ticks → lock R1→ 6 ticks → lock R2 → 2 ticks → unlock R2 → 2 ticks → unlock R1

Output

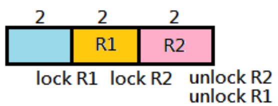
- Similar to those in Labs 0 and 1, but add lock/unlock events
- Output the results of using CPP for Scenarios 1 and 2

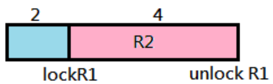
Output Example

Priority initialization:
R1: 1
R2: 2
Task1: 3
Task2: 4
Task3: 5

Task arrival time:
Task1: 8
Task2: 4
Task3: 0

Task execution time and resource used:

Task1: 

Task2: 

Task3: 


```
 20      lock    R1   (Prio=5 changes to=1)
 90    unlock    R1   (Prio=1 changes to=5)
 90  complete          5          3
110      lock    R1   (Prio=3 changes to=1)
130      lock    R2   (Prio=1 changes to=1)
150    unlock    R2   (Prio=1 changes to=1)
150    unlock    R1   (Prio=1 changes to=3)
150  complete          3          4
170      lock    R2   (Prio=4 changes to=2)
220    unlock    R2   (Prio=2 changes to=4)
220  complete          4          19
```