

GADTs meet their Match

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KU LEUVEN

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Checking Pattern Matches

Property I:

Exhaustiveness

```
zip []      []      = []
zip (a:as) (b:bs) = (a,b) : zip as bs
```

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*Main> zip [] [1]
```

*** Exception:

Non-exhaustive patterns in function zip

Property I: Exhaustiveness

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```
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```

*** Exception:

Non-exhaustive patterns in function zip

Runtime crash



Property I:

Exhaustiveness

`zip [] [] = []`

`zip (a:as) (b:bs) = (a,b) : zip as bs`

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zip []      []      = []
zip (a:as) (b:bs) = (a,b) : zip as bs
```

```
*Main> :l zip.hs
[1 of 1] Compiling Main
```

Warning:

Pattern match(es) are non-exhaustive

In an equation for ‘zip’:

Patterns not matched:

[] (_ : _)
(_ : _) []

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[1 of 1] Compiling Main
```

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In an equation for ‘zip’:

Patterns not matched:

[] (_ : _)
(_ : _) []

Compiler warning



GADT Problem I

```
data Vect :: Nat -> * -> * where
  VN :: Vect Zero a
  VC :: a -> Vect n a -> Vect (Succ n) a

vzip :: Vect n a -> Vect n b -> Vect n (a,b)
vzip VN      VN      = VN
vzip (VC x xs) (VC y ys) = VC (x,y) (vzip xs ys)
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(VC)VN

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Pattern match(es) are non-exhaustive
In an equation for ‘vzip’:

Patterns not matched:

VN (VC)
(VC)VN



Property 2: Redundancy

<code>len []</code>	$=$	0
<code>len (x:xs)</code>	$=$	$1 + \text{len } xs$
<code>len 1</code>	$=$	42

Property 2: Redundancy

`len [] = 0`

`len (x:xs) = 1 + len xs`

`len l = 42`



dead
code

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dead
code

Warning:

Pattern match(es) are overlapped

In an equation for ‘len’: `len l = ...`

Compiler warning



GADT Problem 2

vzip VN VN

= VN

vzip (VC x xs) (VC y ys) = VC (x,y) (vzip xs ys)

vzip _ _ = error "unreachable"

GADT Problem 2

vzip VN VN	=	VN
vzip (VC x xs) (VC y ys)	=	VC (x,y) (vzip xs ys)
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suppresses bogus
non-exhaustiveness
warnings

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dead
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No compiler
warning



Quick Summary

Exhaustiveness }
Redundancy }

broken
for GADTs

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Exhaustiveness }
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broken
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But there is more ...

Challenge I

```
data F a where
  F1 :: F Int
  F2 :: F Bool
data G a where
  G1 :: G Int
  G2 :: G Char
```

```
h :: F a -> G a -> Int
h F1 G1 = 1
h _ _ = 2
```

Challenge I

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data F a where  
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```
> h F2 undefined  
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redundant?

```
> h F2 undefined  
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```

```
> h F2 undefined  
*** Exception
```

Challenge 2

```
g :: Bool -> Bool -> Int
g _    False = 1
g True False = 2
g _    _      = 3
```

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g :: Bool -> Bool -> Int  
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> g undefined True  
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Challenge 2

‘useful’
pattern

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g :: Bool -> Bool -> Int
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g True False = 2
g _ _ = 3
```

```
> g undefined True
*** Exception
```

```
> g undefined True
3
```

Challenge 2

‘useful’
pattern

```
g :: Bool -> Bool -> Int  
g _ False = 1  
g True False = 2  
g _ _ = 3
```

unreachable
RHS

```
> g undefined True  
*** Exception
```

```
> g undefined True  
3
```

The Guard Challenge

```
abs2 :: Int -> Int  
abs2 x | x < 0 = -x  
       | x >= 0 = x
```

exhaustive?

The Guard Challenge

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abs2 :: Int -> Int  
abs2 x | x < 0 = -x  
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exhaustive?

```
append xs ys  
| []      <- xs = ys  
| (p:ps) <- xs = p : append ps ys
```

exhaustive?

What we need:

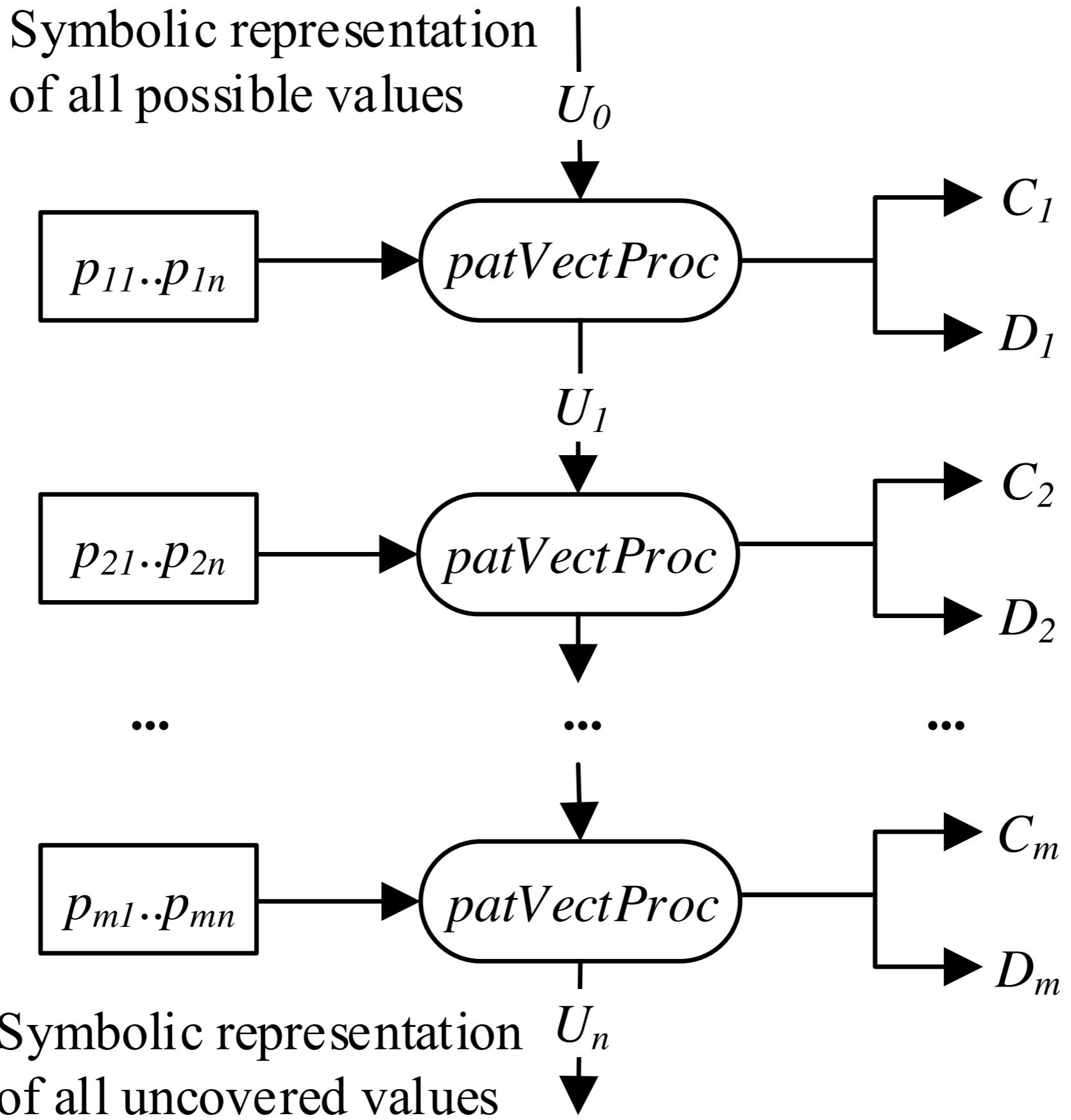
A uniform approach to deal with:

GADTs

Laziness

Guards

Our Approach



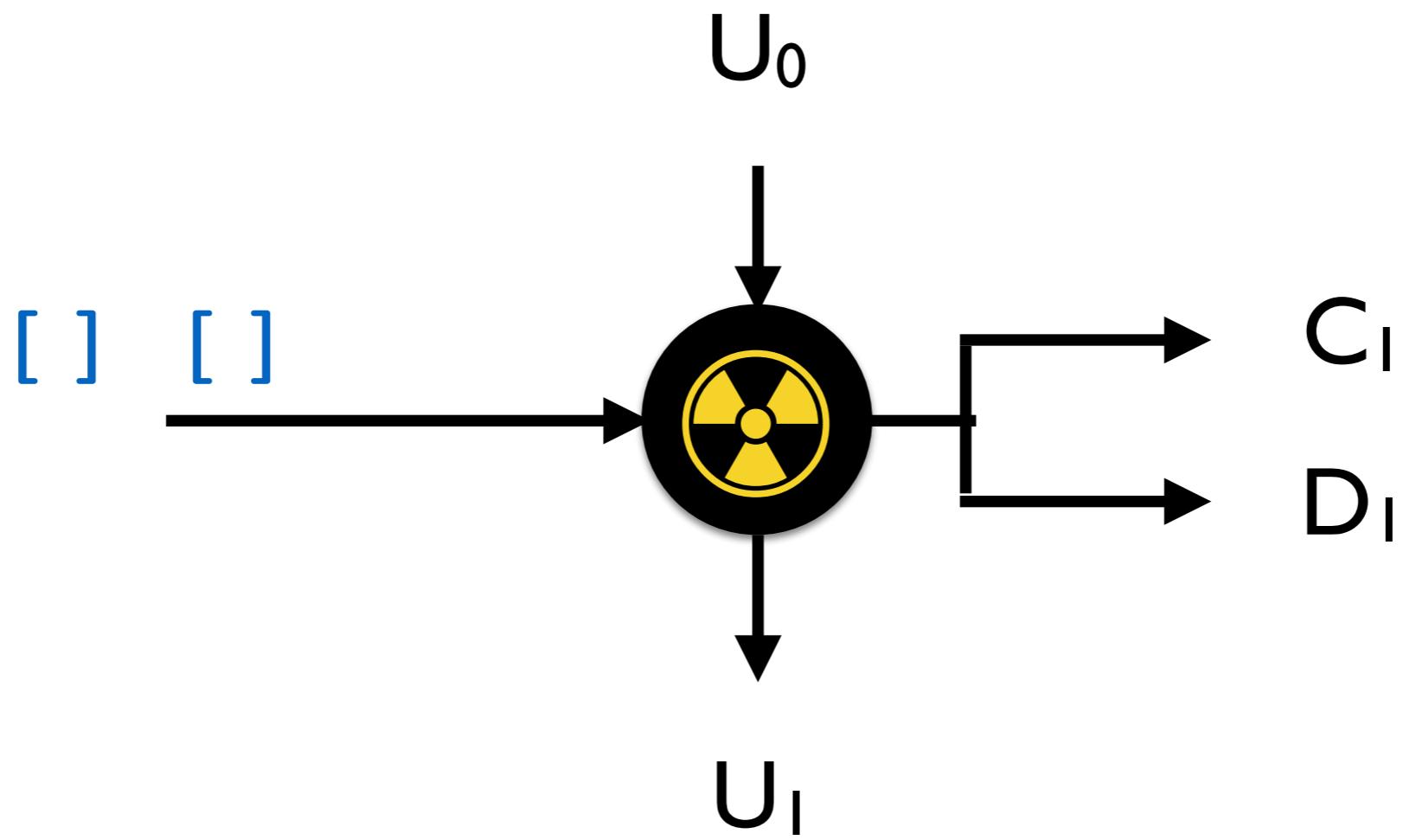
Basic ADT Support

```
zip [] [] = ...
zip (x:xs) (y:ys) = ...
```

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zip [] [] = ...
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```

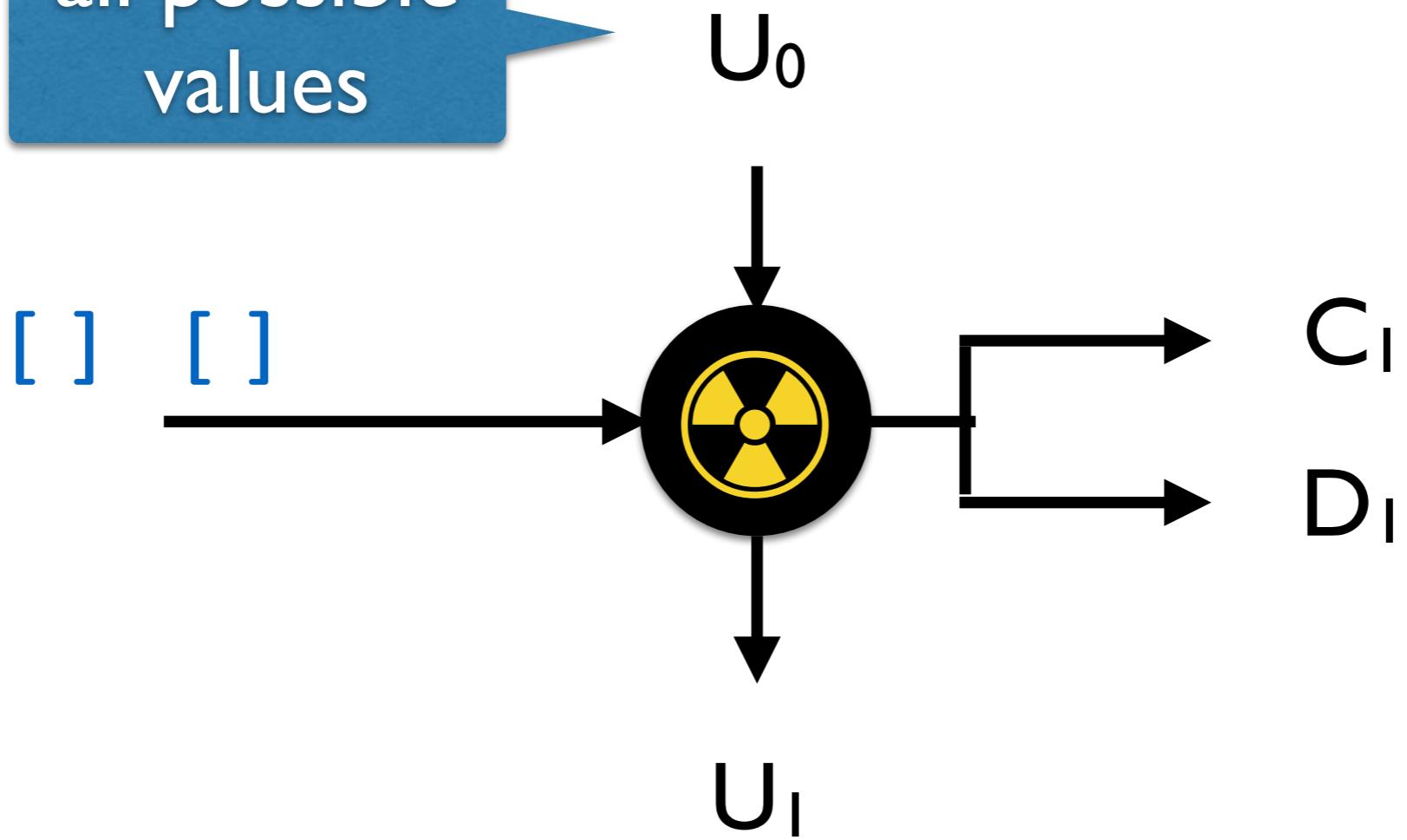
[] []

(x : xs) (y : ys)



(x:xs) (y:ys)

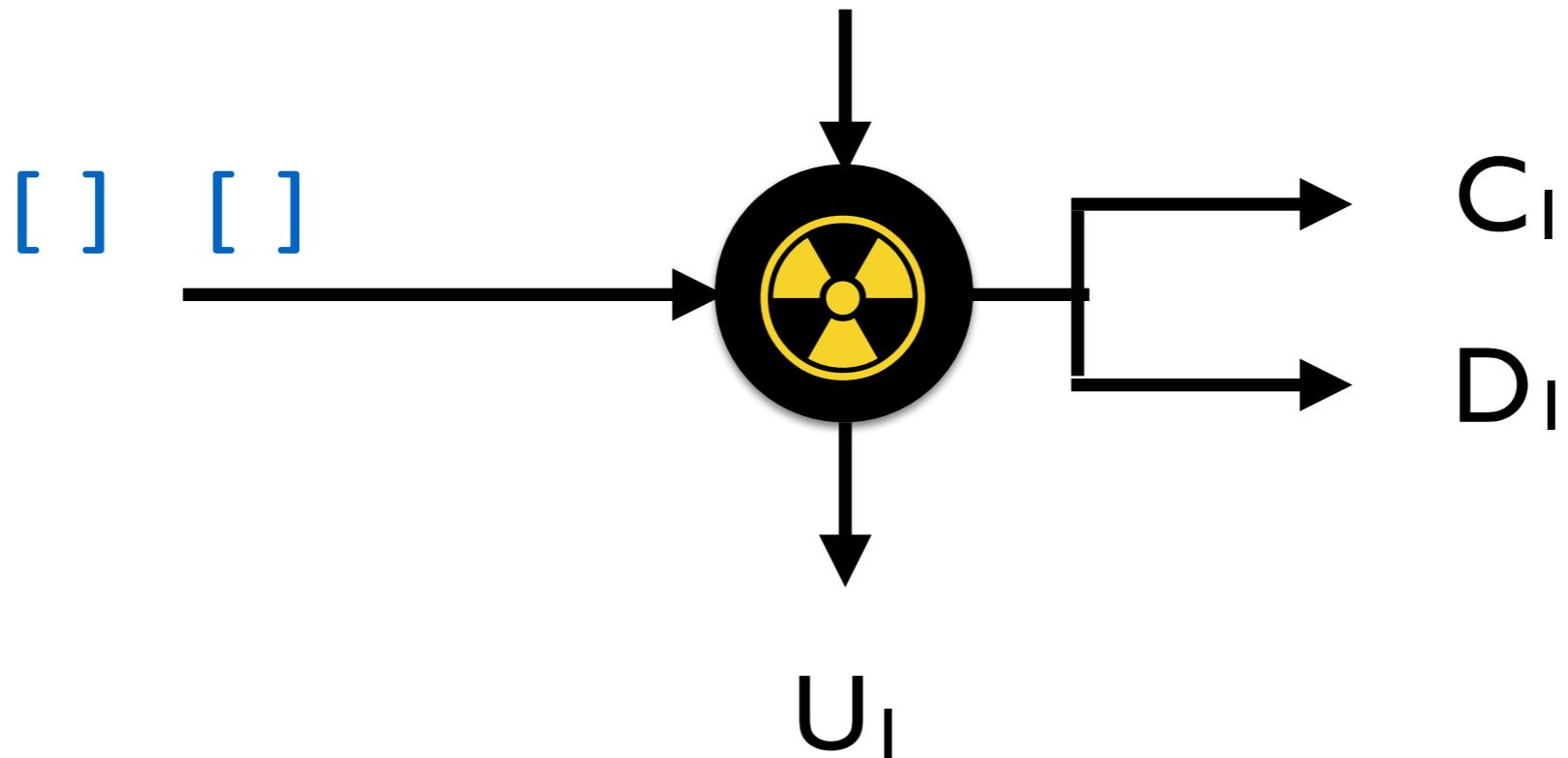
all possible
values



($x:xs$) ($y:ys$)

all possible
values

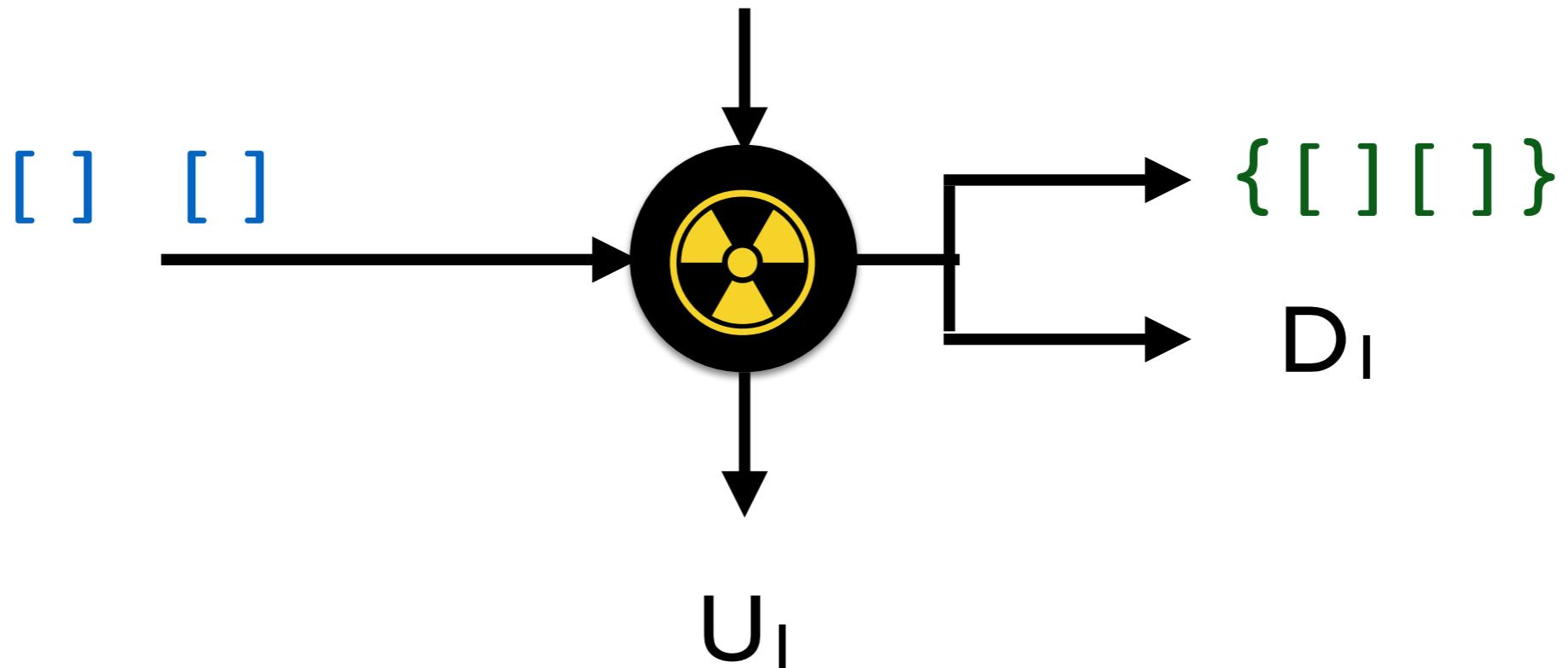
{ _ _ }



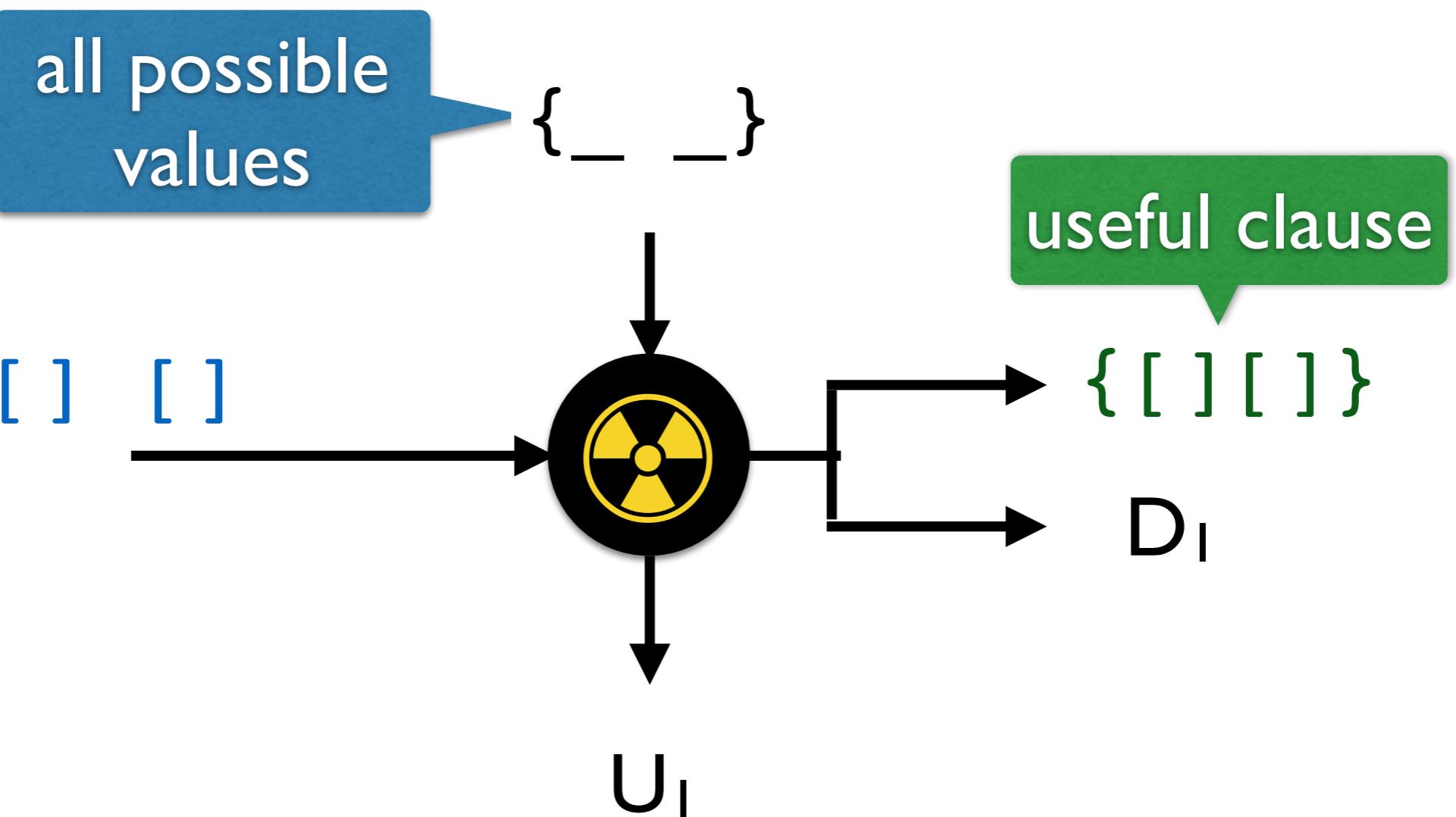
(x:xs) (y:ys)

all possible
values

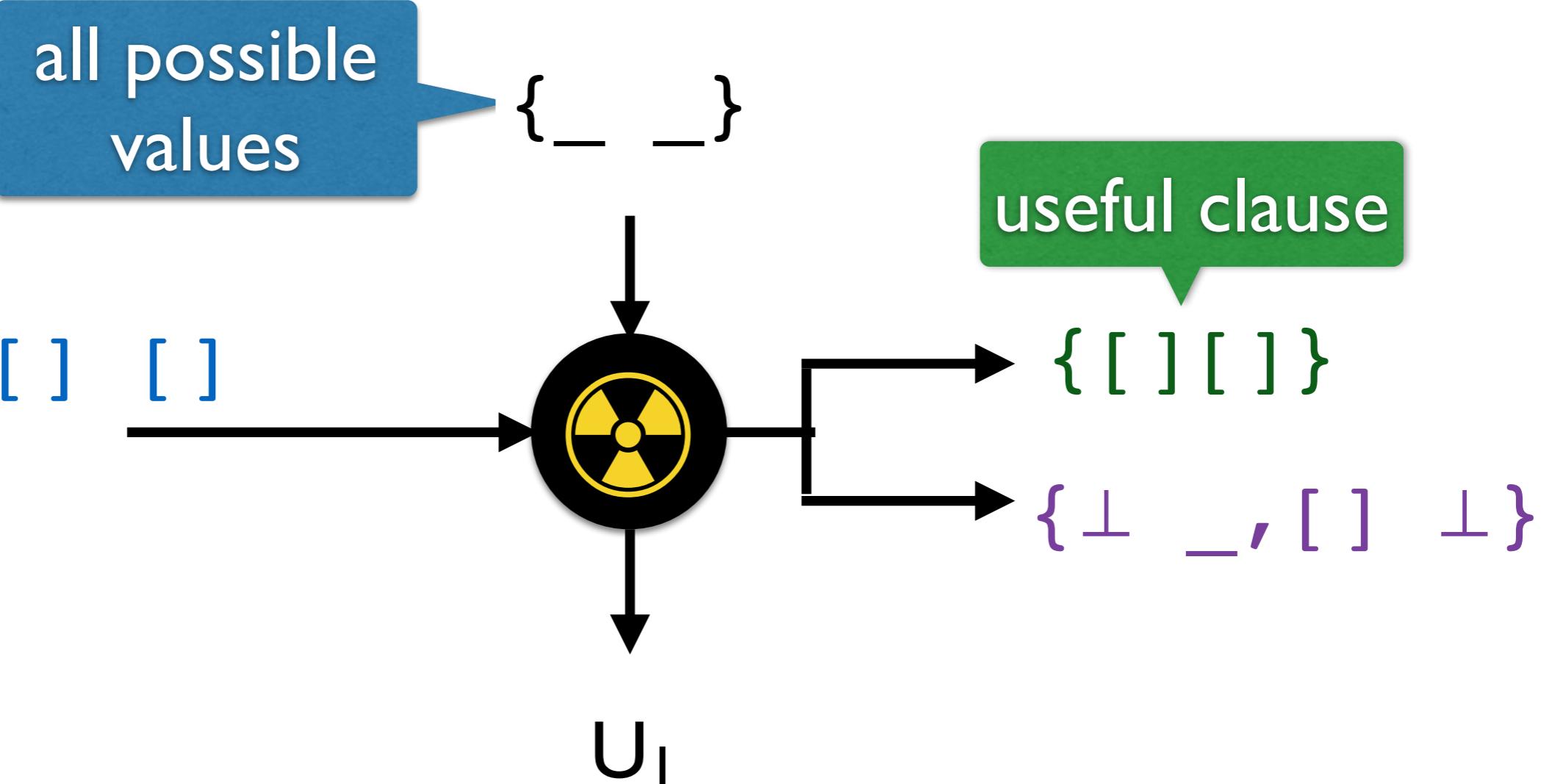
{ _ _ }



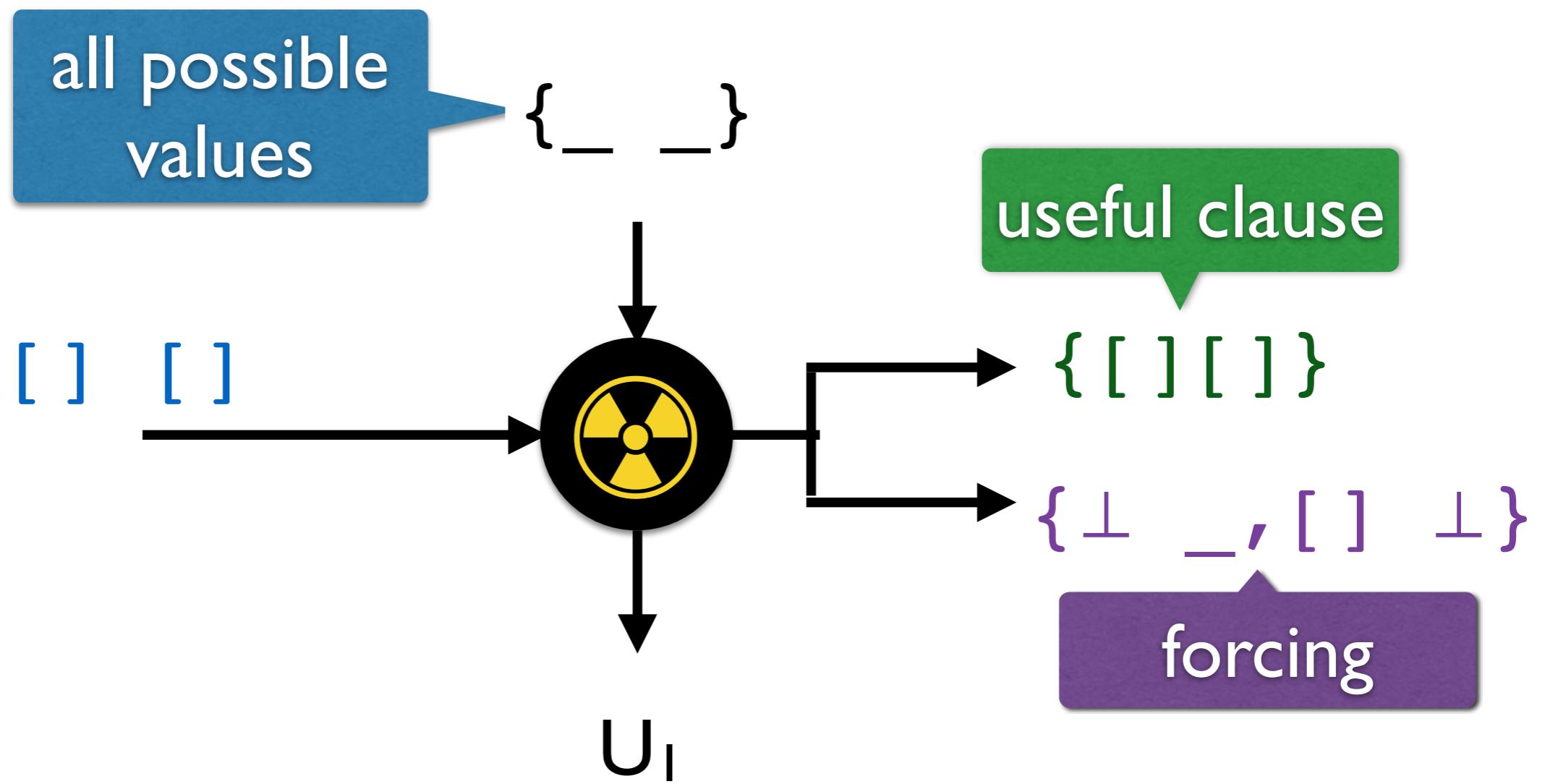
(x:xs) (y:ys)



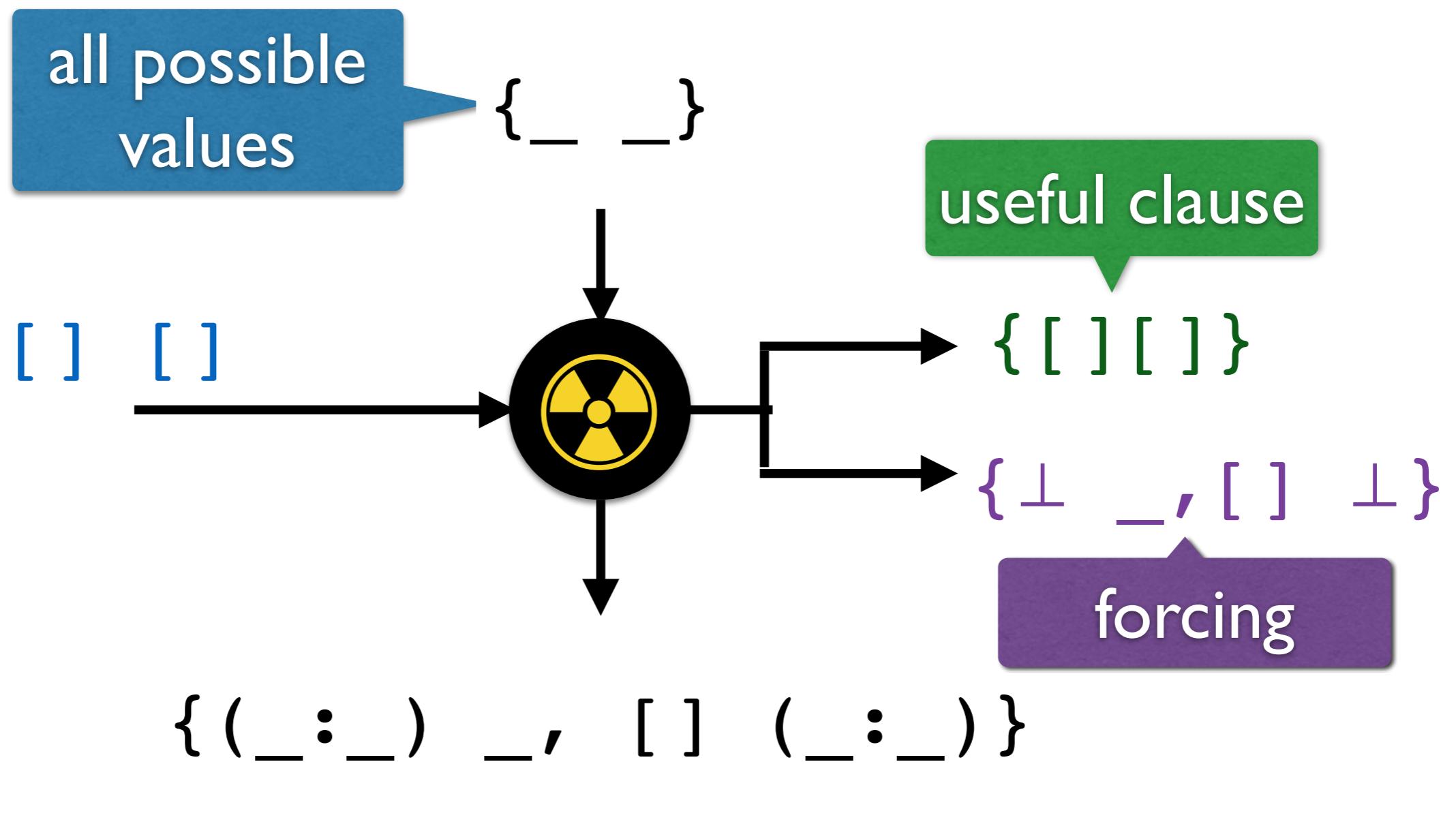
$(x:xs) \quad (y:ys)$



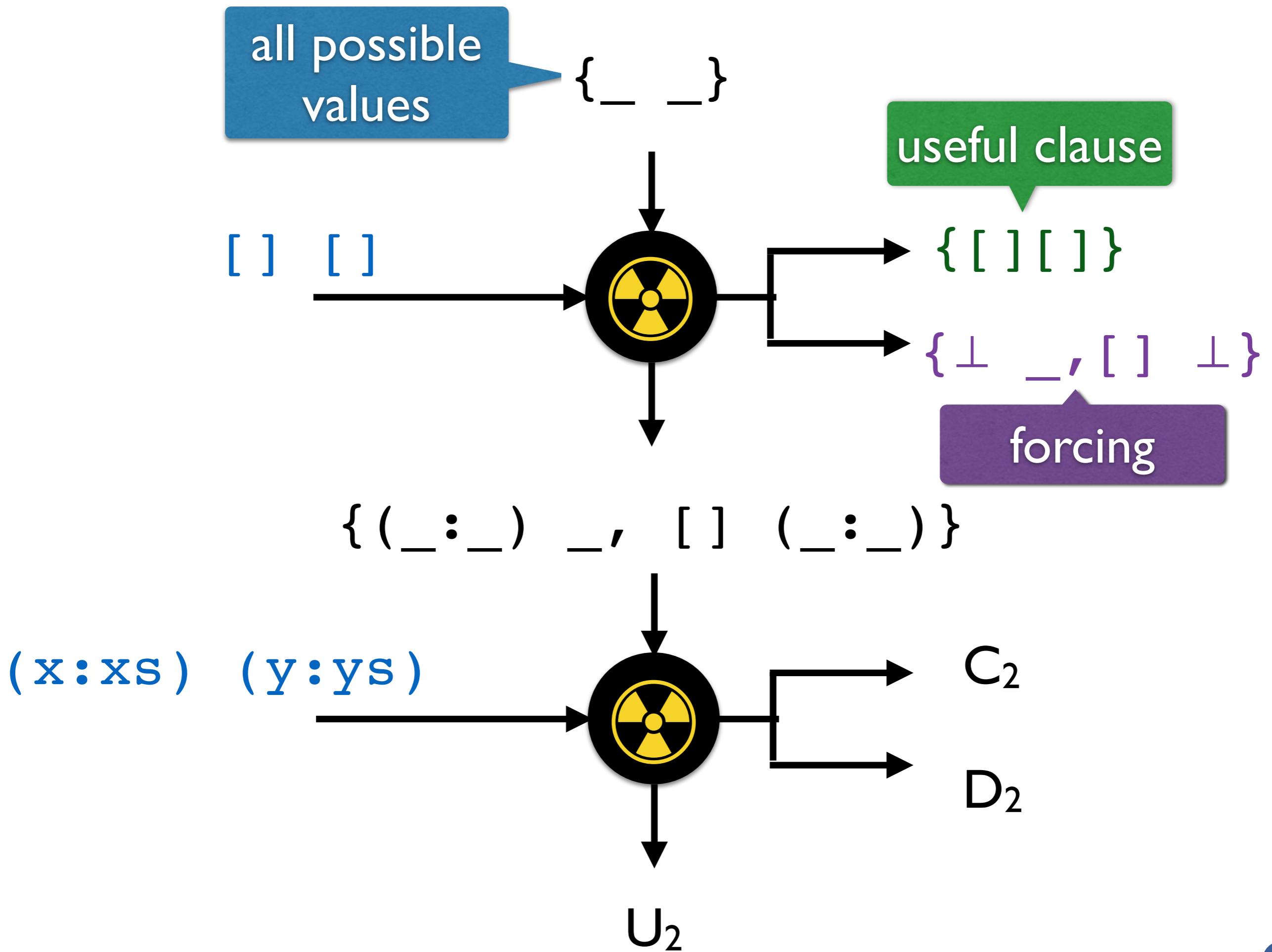
(x:xs) (y:ys)

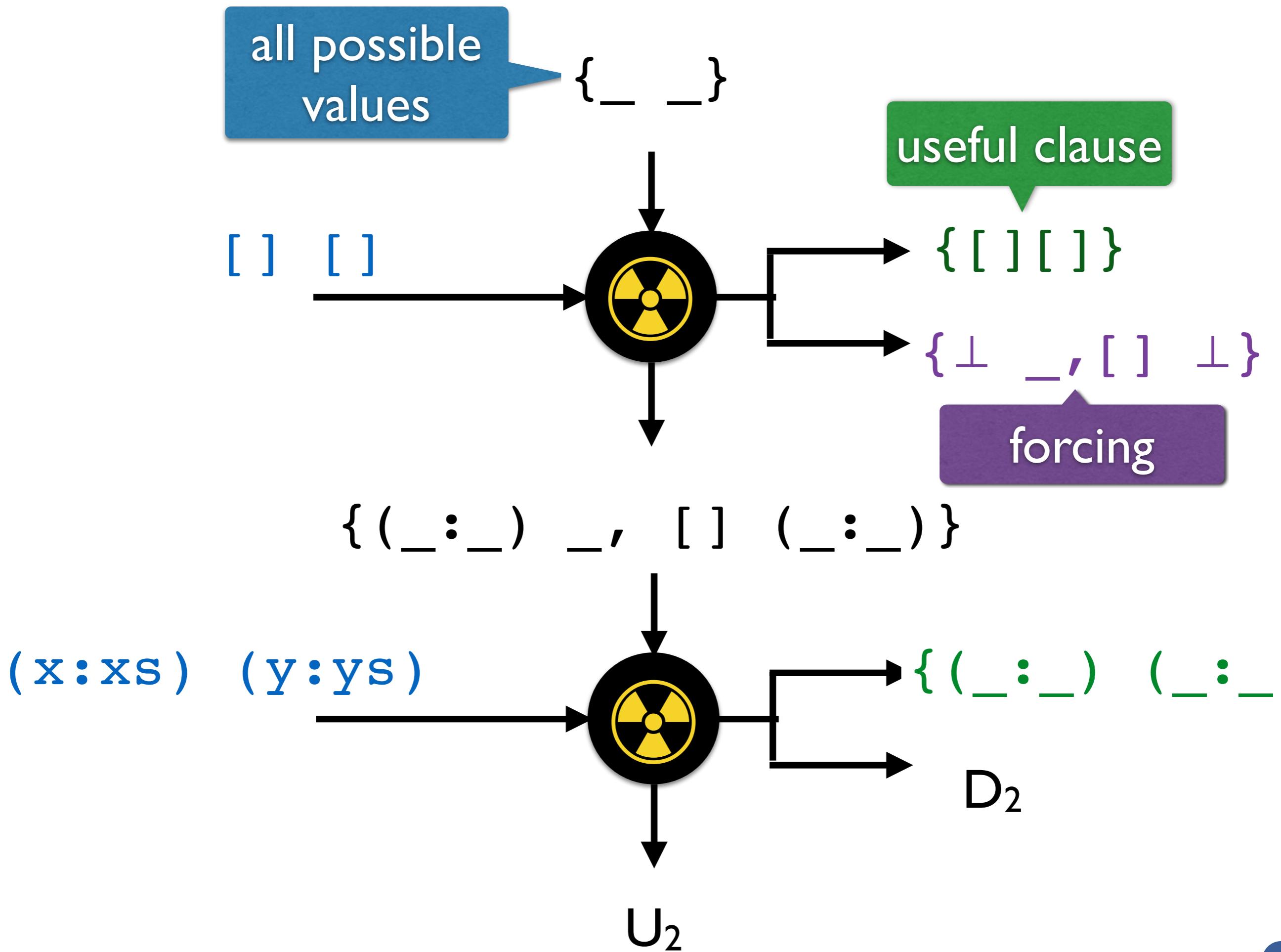


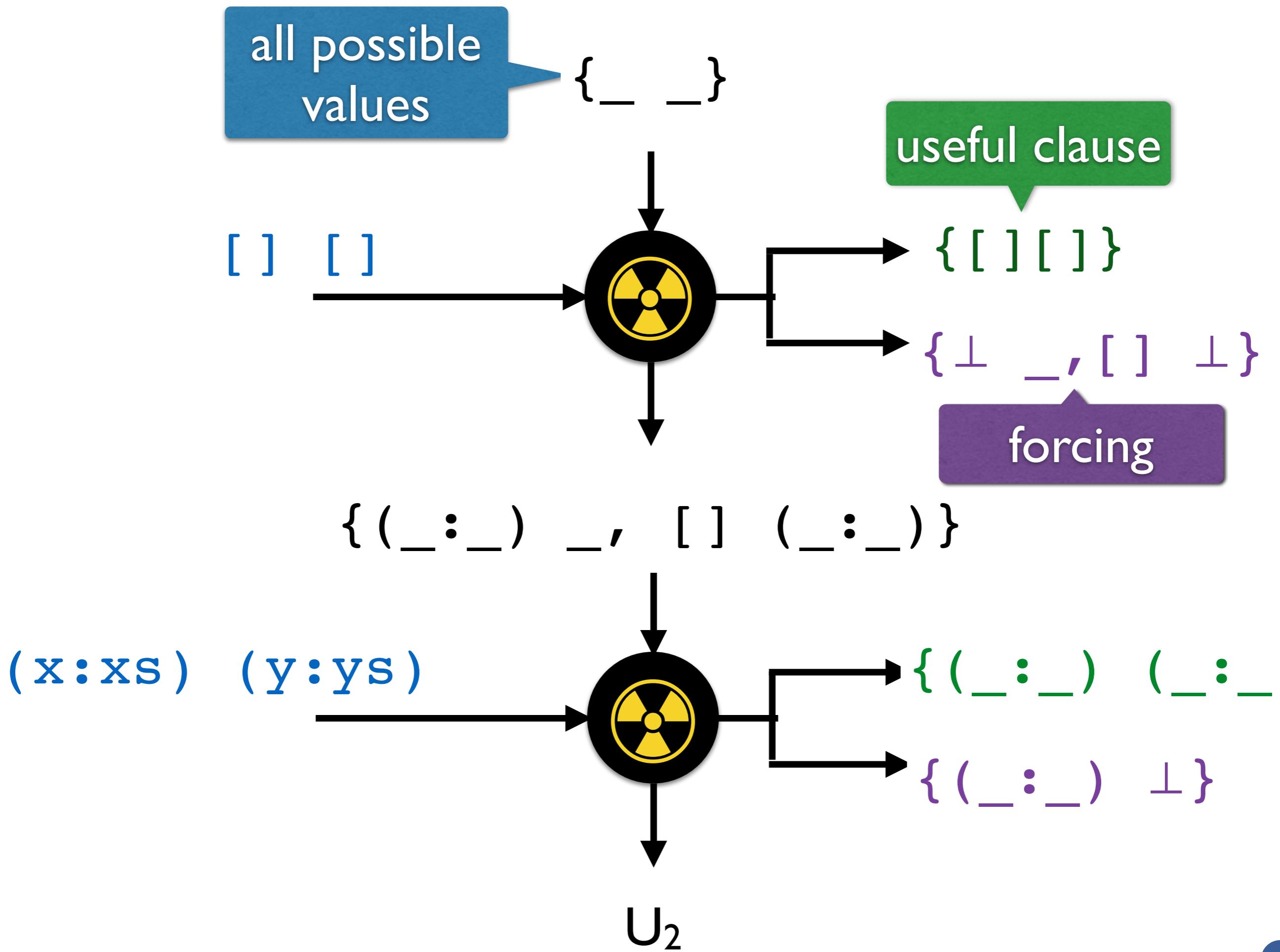
(x:xs) (y:ys)

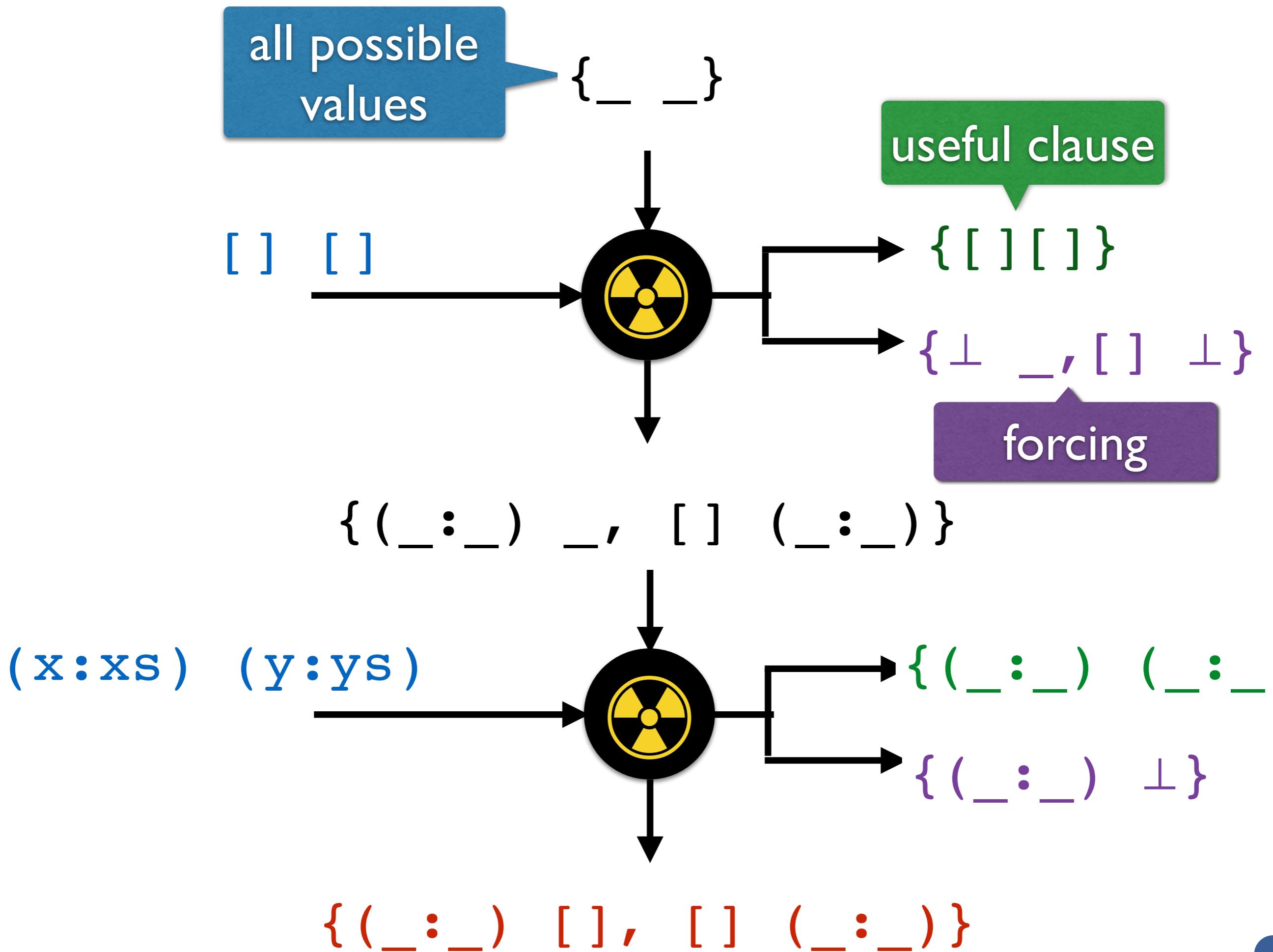


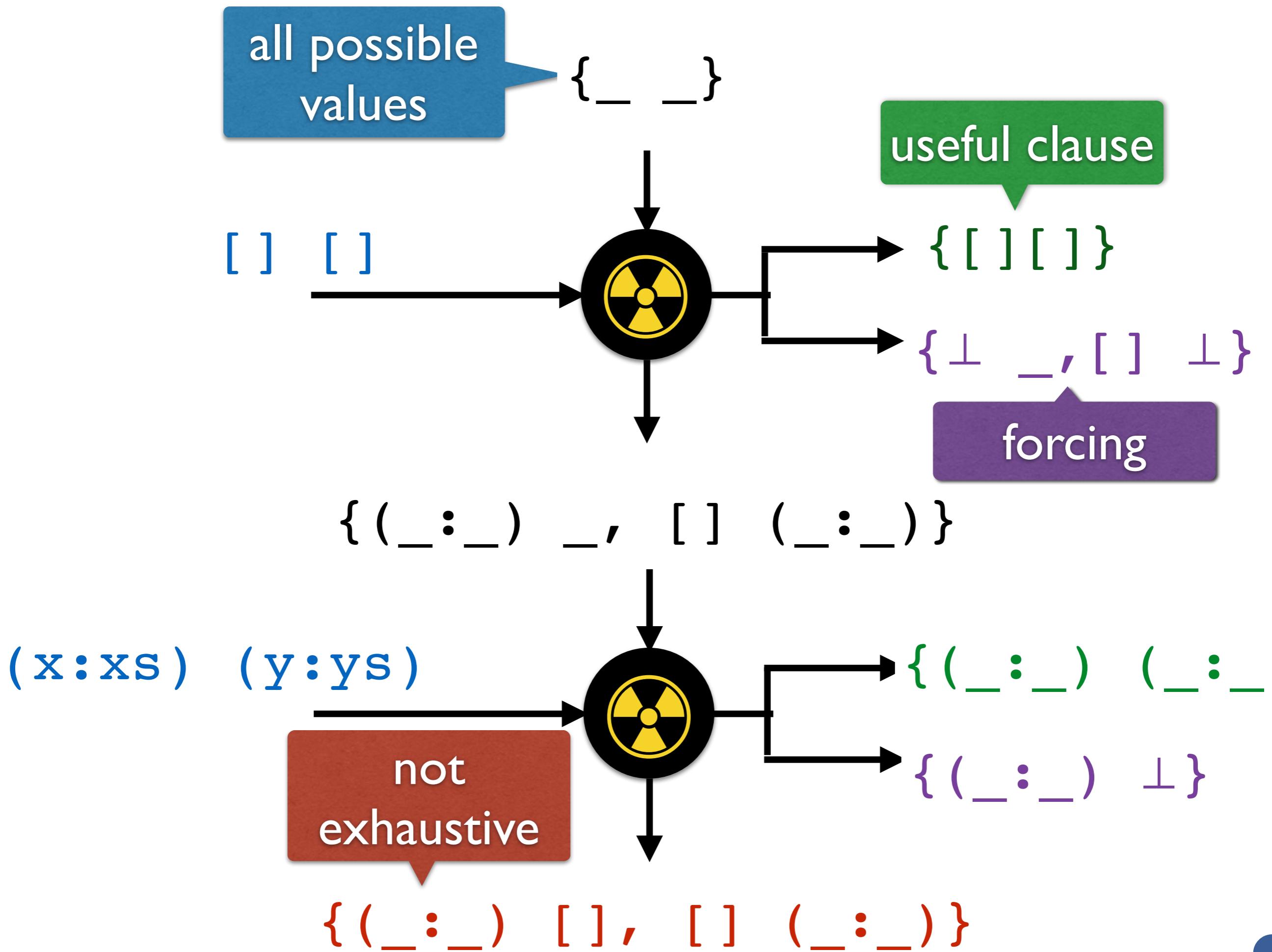
$(x:xs) \ (y:ys)$











Basic Syntax

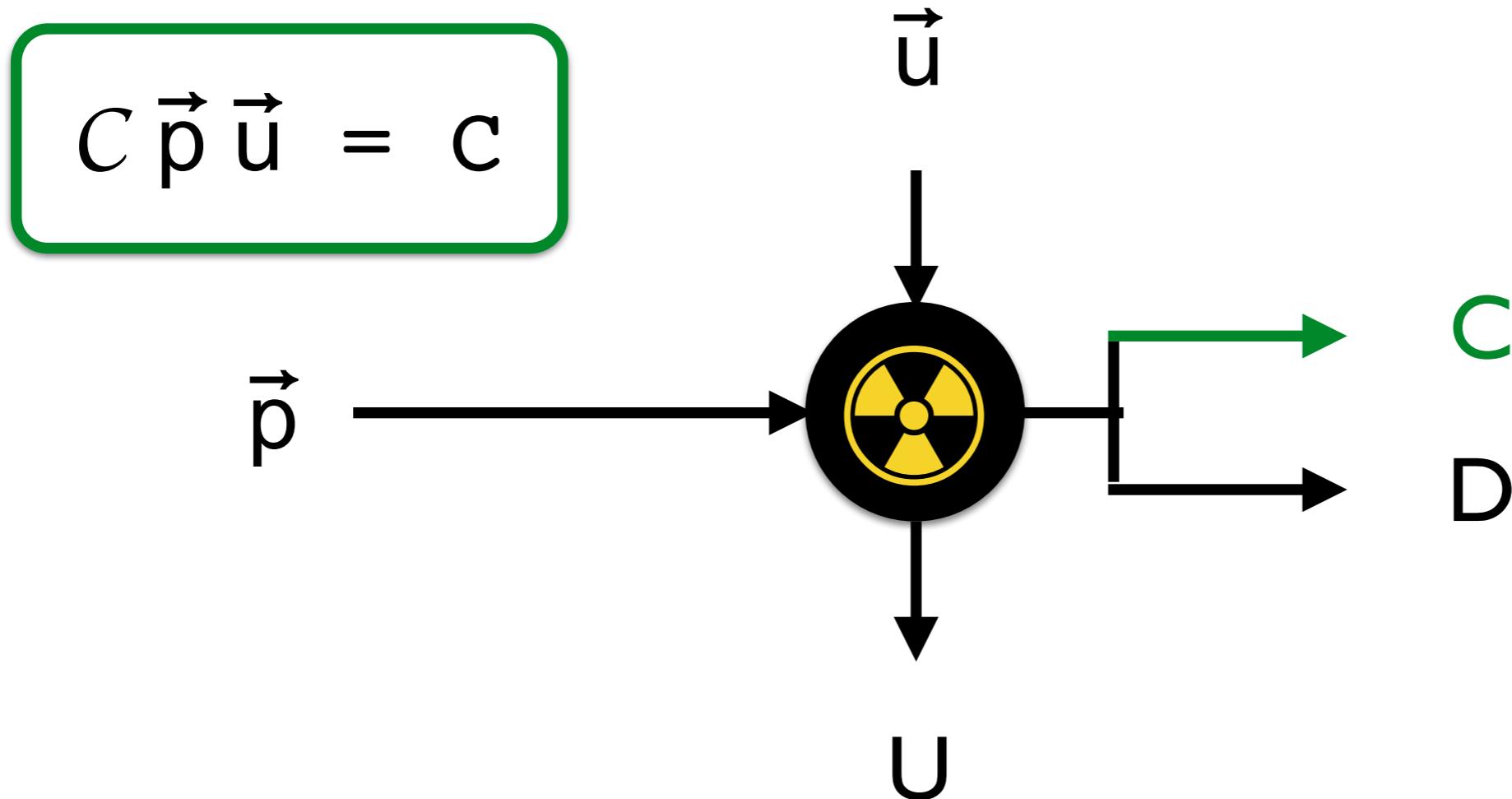
Pattern

$$p ::= \underline{\quad}$$
$$\mid K \vec{p}$$

Value Abstraction

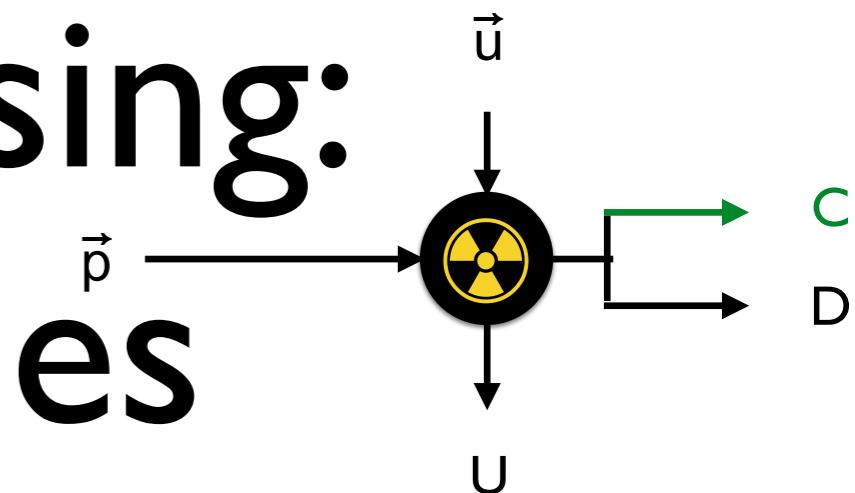
$$u ::= \underline{\quad}$$
$$\mid K \vec{u}$$

Clause Processing: Covered Values



Clause Processing: Covered Values

$$C \vec{p} \vec{u} = c$$

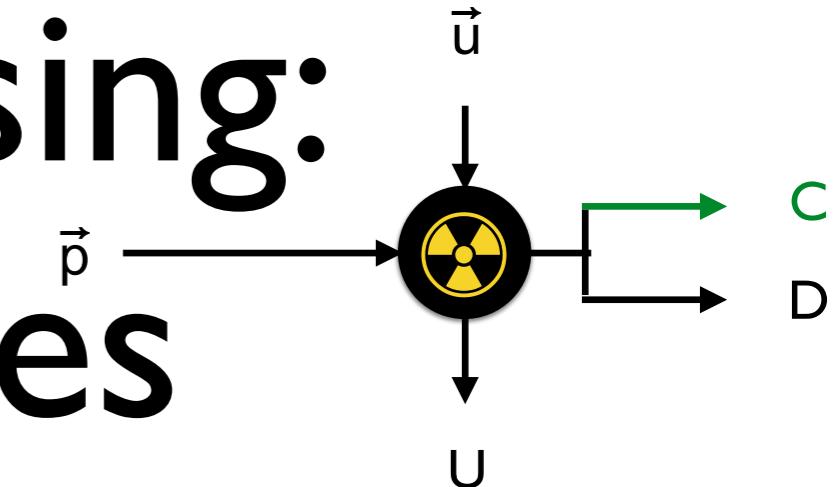


$$C \varepsilon \quad \varepsilon = \{\varepsilon\}$$

Clause Processing:

$$C \vec{p} \vec{u} = c$$

Covered Values



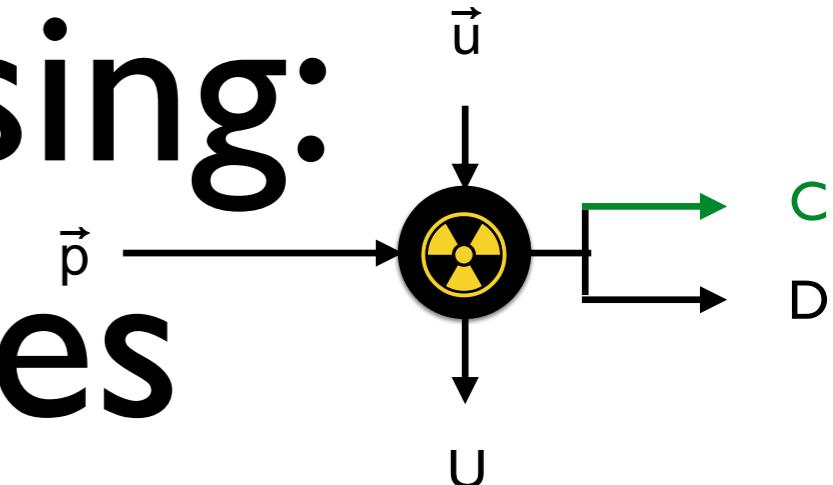
$$C \varepsilon \quad \varepsilon = \{\varepsilon\}$$

$$C (_ \vec{p}) (u \vec{u}) = \{u \vec{w} \mid \vec{w} \leftarrow C \vec{p} \vec{u}\}$$

Clause Processing:

Covered Values

$$C \vec{p} \vec{u} = C$$



$$C \varepsilon \quad \varepsilon = \{\varepsilon\}$$

$$C (_ \vec{p}) (\vec{u} \vec{u}) = \{\vec{u} \vec{w} \mid \vec{w} \leftarrow C \vec{p} \vec{u}\}$$

$$C ((K_i \vec{q}) \vec{p}) ((K_j \vec{w}) \vec{u})$$

$$\mid K_i == K_j = \{(K_i \vec{w}') \vec{u}'$$

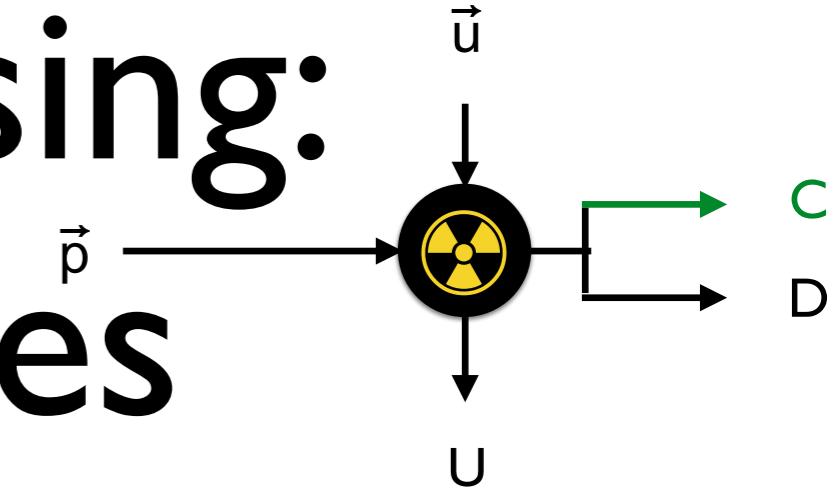
$$\mid \vec{w}' \vec{u}' \leftarrow C (\vec{q} \vec{p}) (\vec{w} \vec{u})\}$$

$$\mid \text{otherwise} = \{\}$$

Clause Processing:

Covered Values

$$C \vec{p} \vec{u} = C$$



$$C \varepsilon \quad \varepsilon = \{\varepsilon\}$$

$$C (_ \vec{p}) (u \vec{u}) = \{u \vec{w} \mid \vec{w} \leftarrow C \vec{p} \vec{u}\}$$

$$C ((K_i \vec{q}) \vec{p}) ((K_j \vec{w}) \vec{u})$$

$$\begin{aligned} | K_i == K_j &= \{ (K_i \vec{w}') \vec{u}' \\ &\quad | \vec{w}', \vec{u}' \leftarrow C (\vec{q} \vec{p}) (\vec{w} \vec{u}) \} \end{aligned}$$

$$| \text{otherwise} = \{ \}$$

$$C ((K_i \vec{q}) \vec{p}) (_ \vec{u})$$

$$= C ((K_i \vec{q}) \vec{p}) ((K_i _ _) \vec{u})$$

Clause Processing:

Covered Values

$$C \vec{p} \vec{u} = C$$

Uncovered

$$U \vec{p} \vec{u} = U$$

Diverging

$$\mathcal{D} \vec{p} \vec{u} = D$$

$$= \{\varepsilon\}$$

$$= \{u \vec{w} \mid \vec{w} \leftarrow C \vec{p} \vec{u}\}$$

$$(K_j \vec{w}) \vec{u})$$

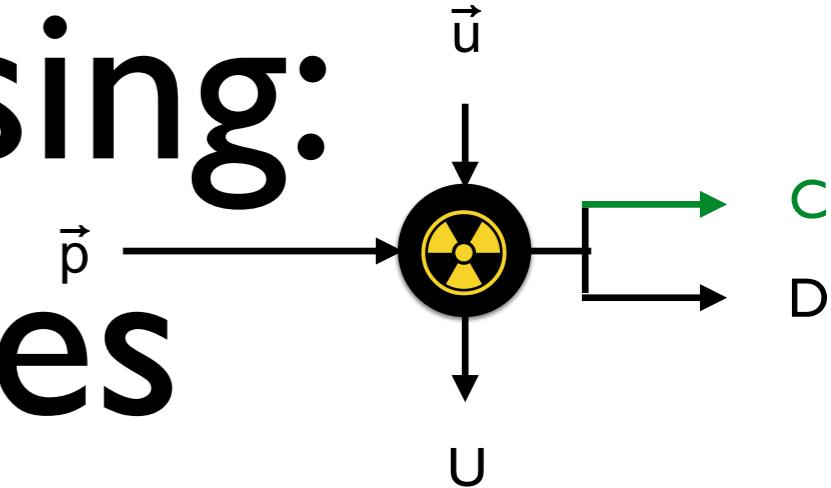
$$\{(K_i \vec{w}') \vec{u}'$$

$$| \vec{w}' \vec{u}' \leftarrow C (\vec{q} \vec{p}) (\vec{w} \vec{u})\}$$

$$= \{\}$$

$$C ((K_i \vec{q}) \vec{p}) (\underline{\quad} \vec{u})$$

$$= C ((K_i \vec{q}) \vec{p}) ((K_i \underline{\dots}) \vec{u})$$



GADTs

GADTs and Type Constraints

```
data Vec :: Nat -> * -> * where
  VN ::                                     Vec Z a
  VC :: a -> Vec m a -> Vec (S m) a
```

GADTs and Type Constraints

```
data Vec :: Nat -> * -> * where
  VN ::                                     Vec z a
  VC :: a -> Vec m a -> Vec (S m) a
```



```
data Vec :: Nat -> * -> * where
  VN :: n ~ z =>                                     Vec n a
  VC :: n ~ S m => a -> Vec m a -> Vec n a
```

GADTs and Type Constraints

```
data Vec :: Nat -> * -> * where
  VN ::                                     Vec z a
  VC :: a -> Vec m a -> Vec (S m) a
```



equivalent

```
data Vec :: Nat -> * -> * where
  VN :: n ~ z =>                                     Vec n a
  VC :: n ~ S m => a -> Vec m a -> Vec n a
```

type constraints

Syntax with Type Constraints

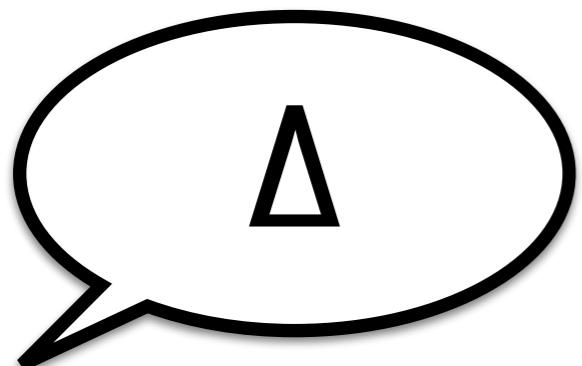
Value Abstraction

$$v ::= \Gamma \vdash \vec{u} \triangleright \Delta$$
$$u ::= _ \mid K \vec{u}$$

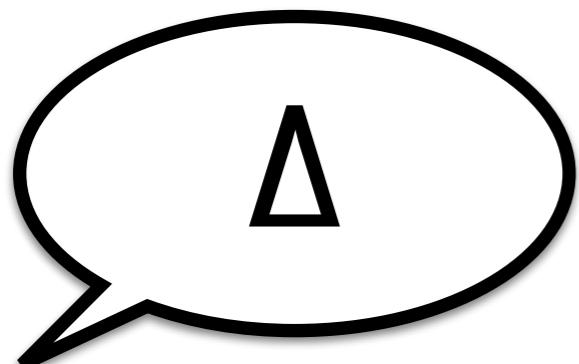
refined by
type constraints

$$\Gamma ::= \varepsilon \mid \Gamma, a$$
$$\Delta ::= \varepsilon \mid \Delta \cup \Delta \mid \tau \sim \tau$$

The Oracle

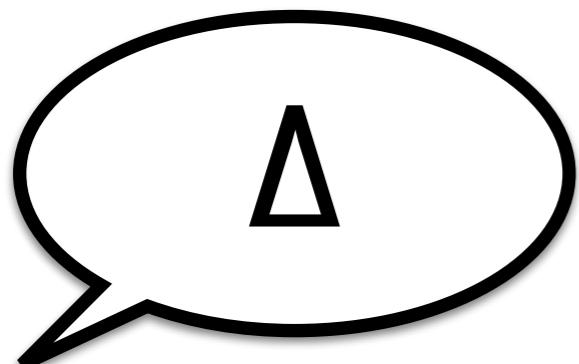


The Oracle



Not
Satisfiable

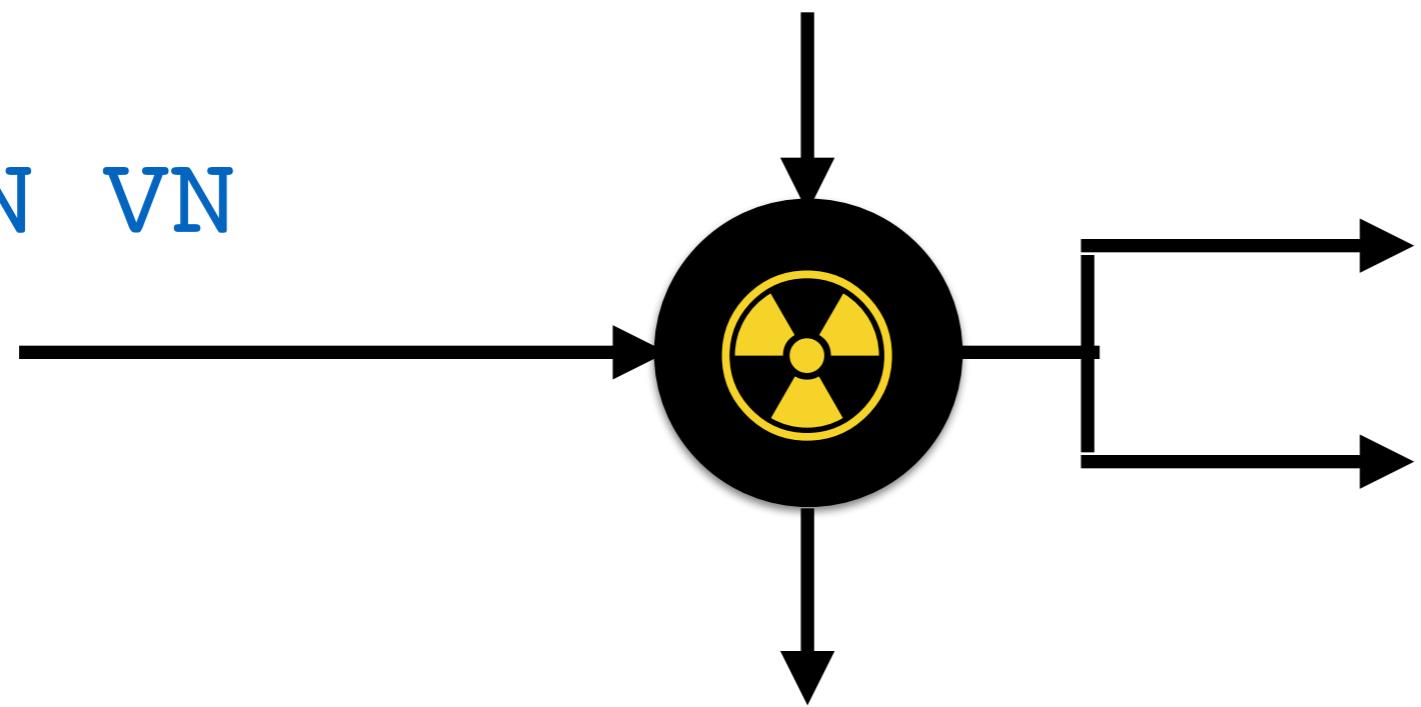
The Oracle

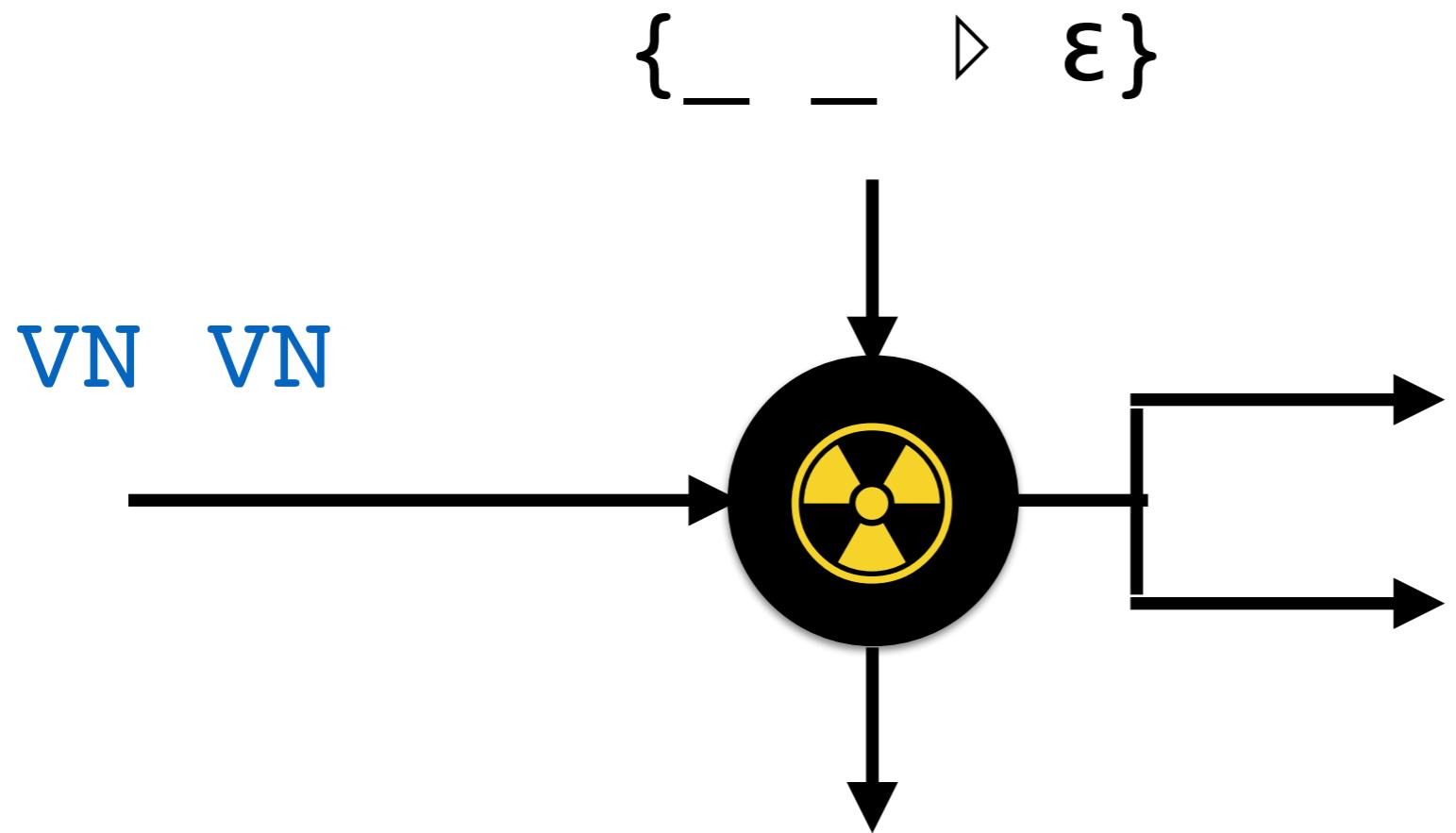


Not
Satisfiable

Maybe
Satisfiable

VN VN

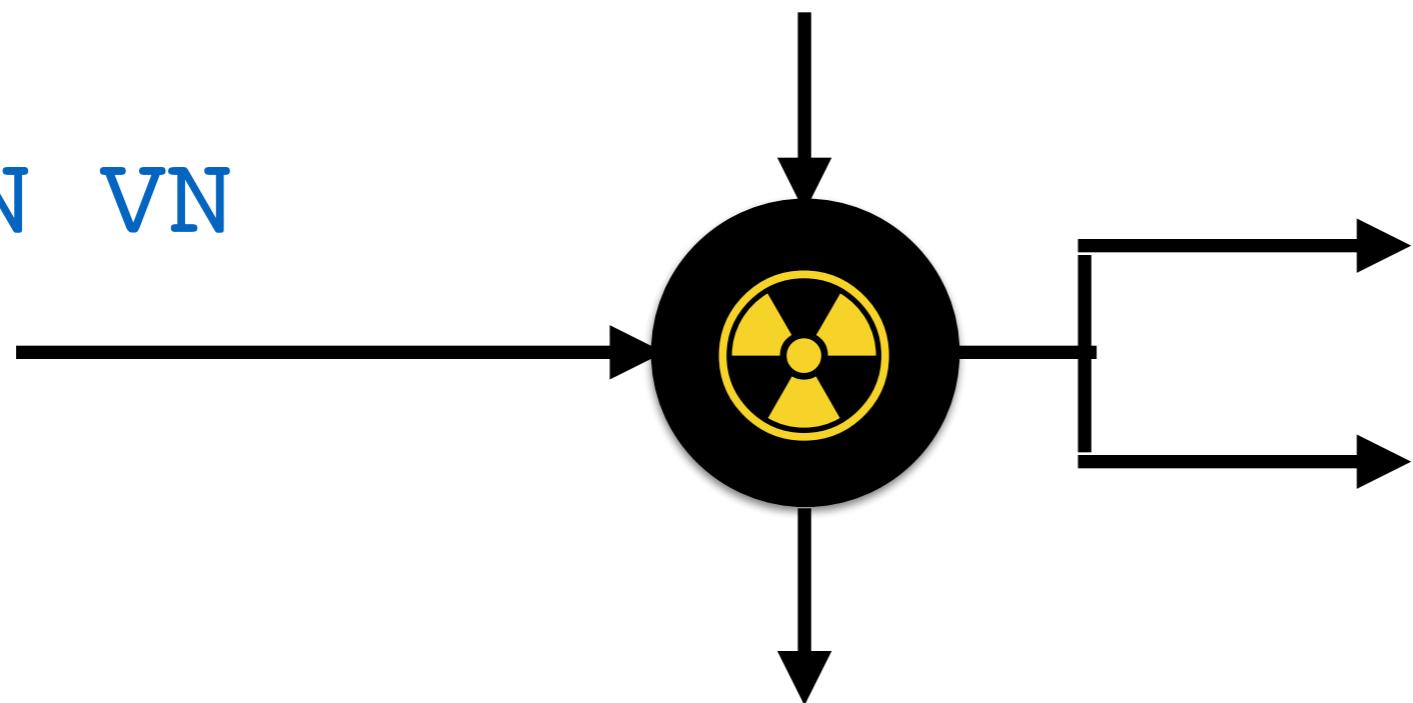


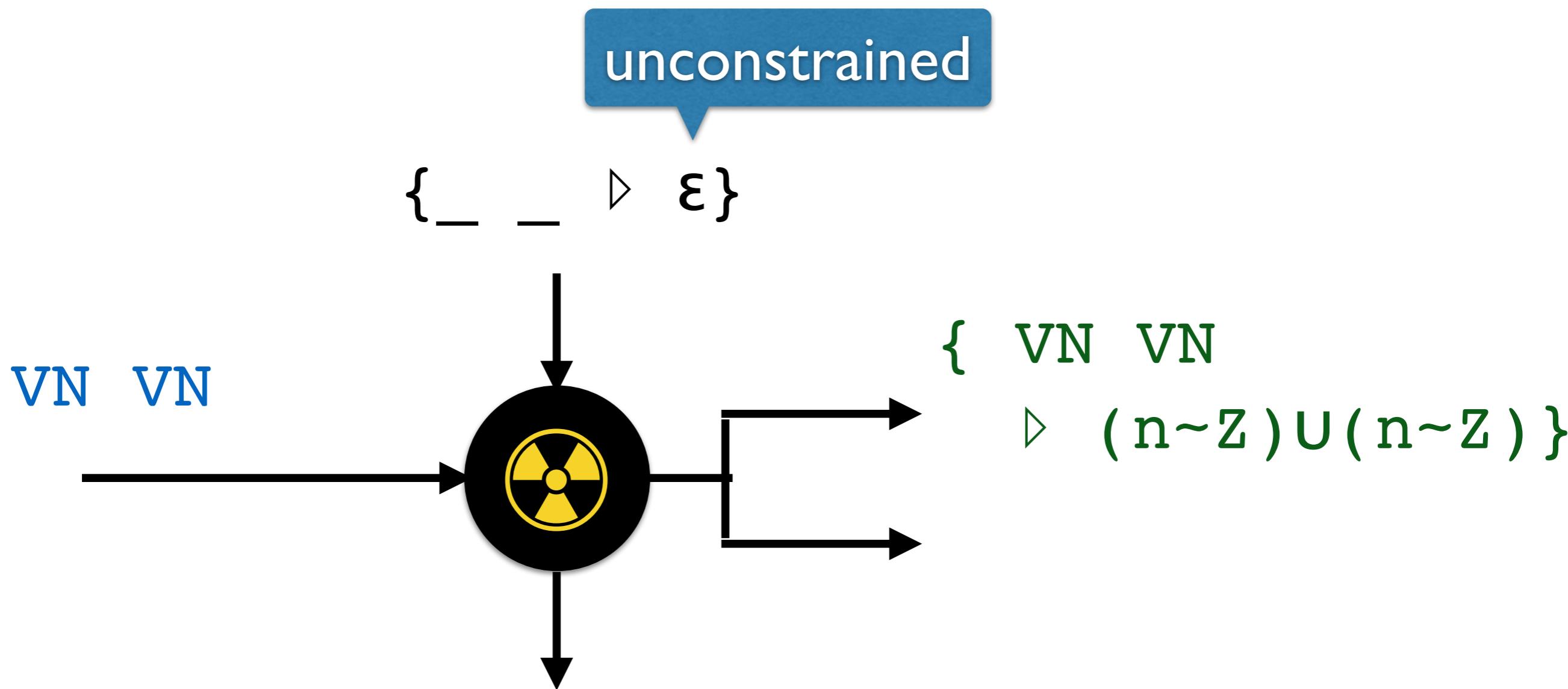


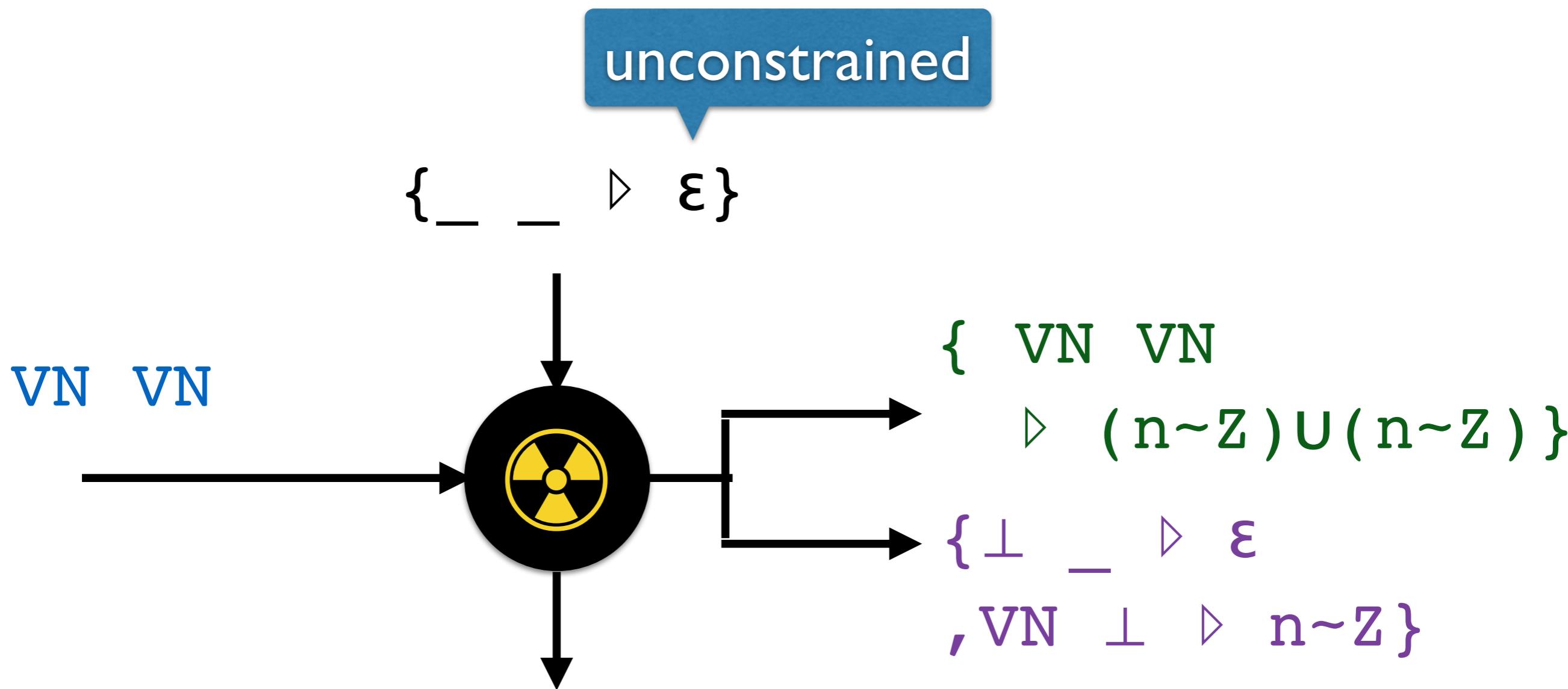
unconstrained

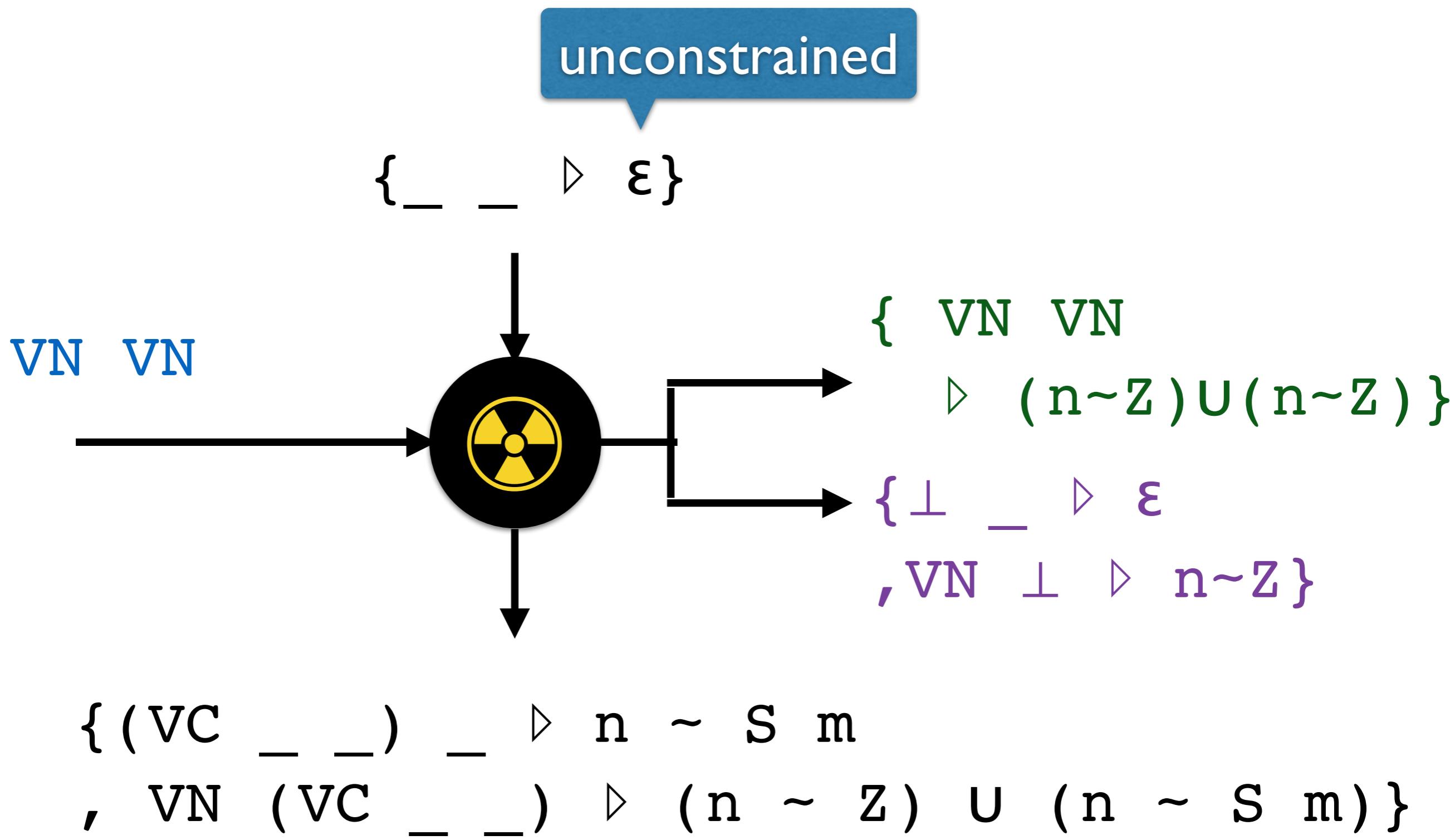
{ _ _ \triangleright ε }

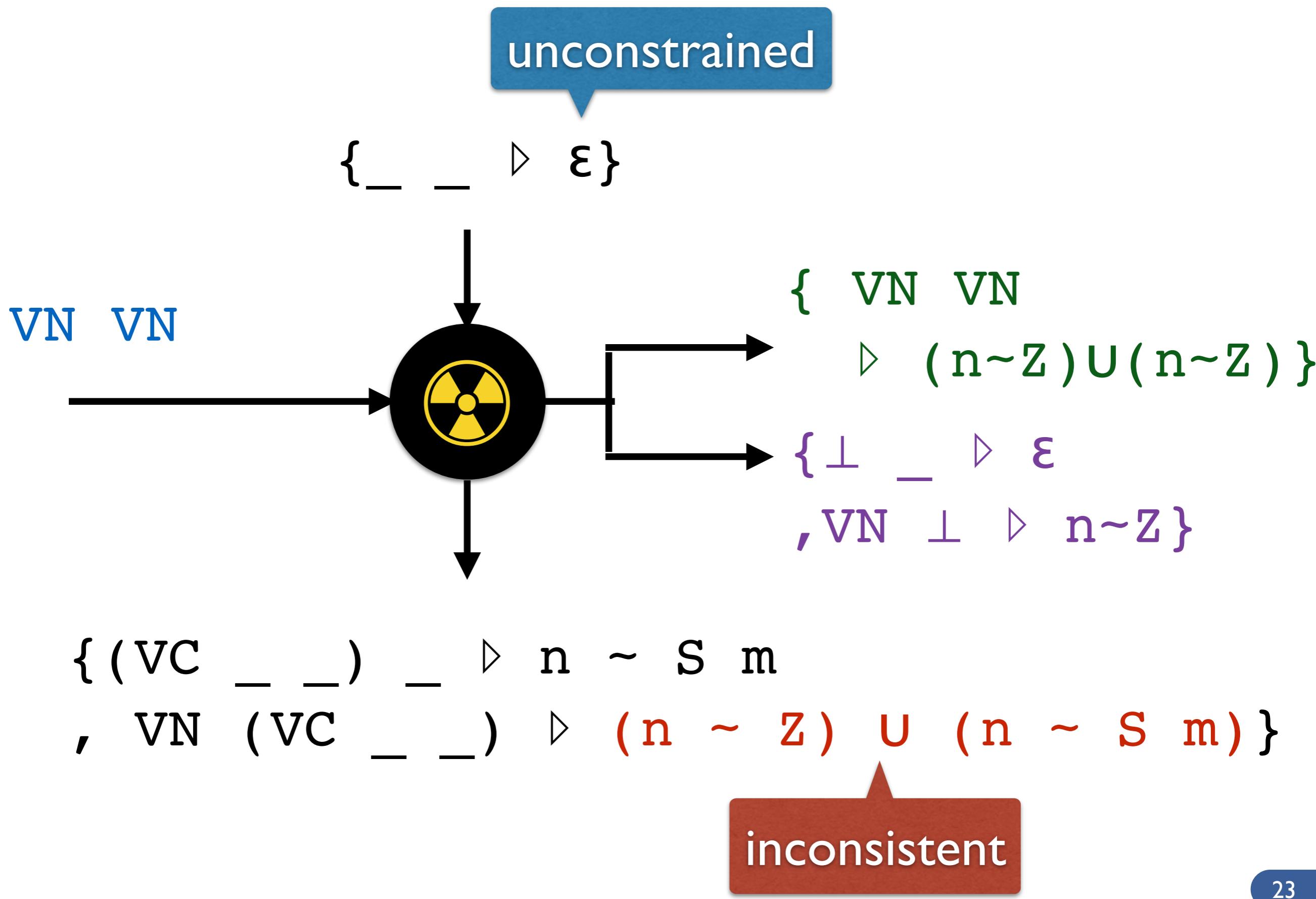
VN VN







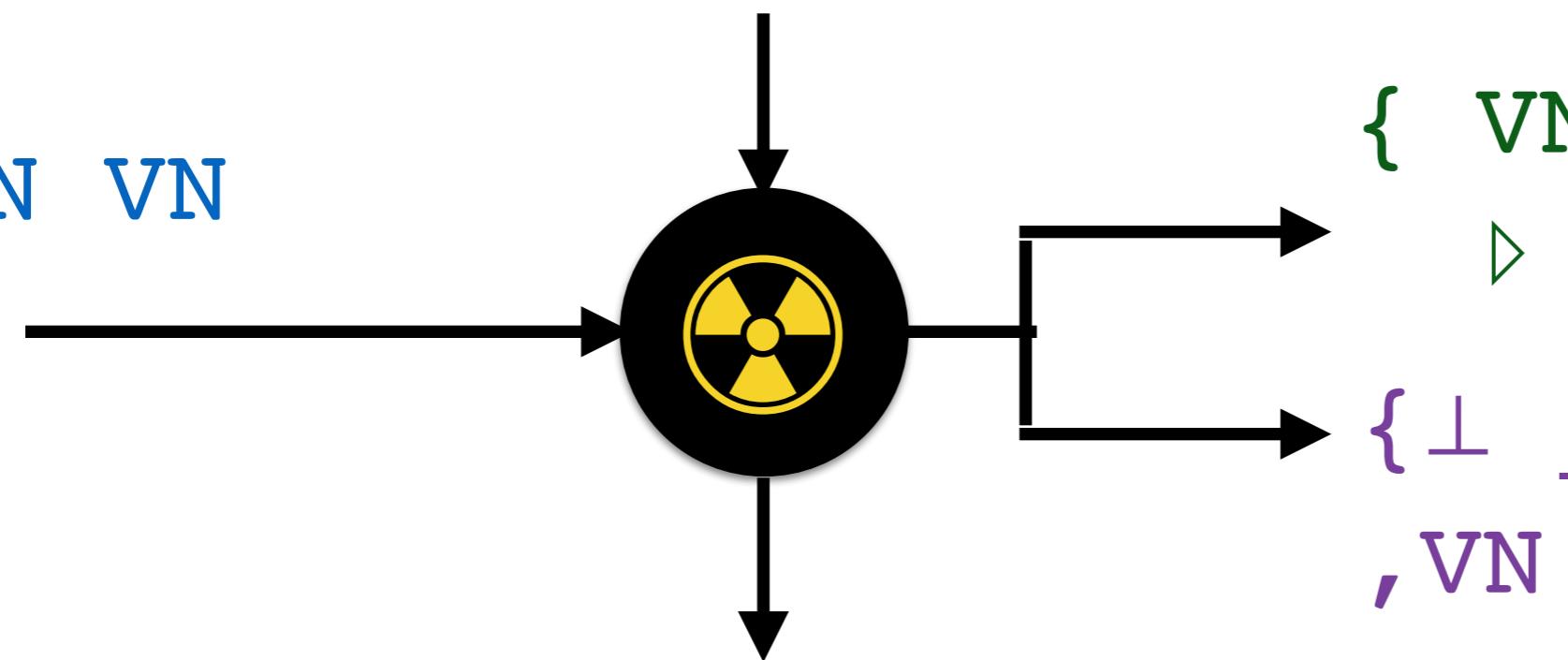




unconstrained

{ } > Σ

VN VN



VN VN

▷ $(n-Z) \cup (n-Z)$

{ ⊥ ▷ ε }

, VN \perp \triangleright n~Z }

$$\{ (\text{VC } \quad \quad \quad) \quad \quad \triangleright n \sim S \text{ } m$$

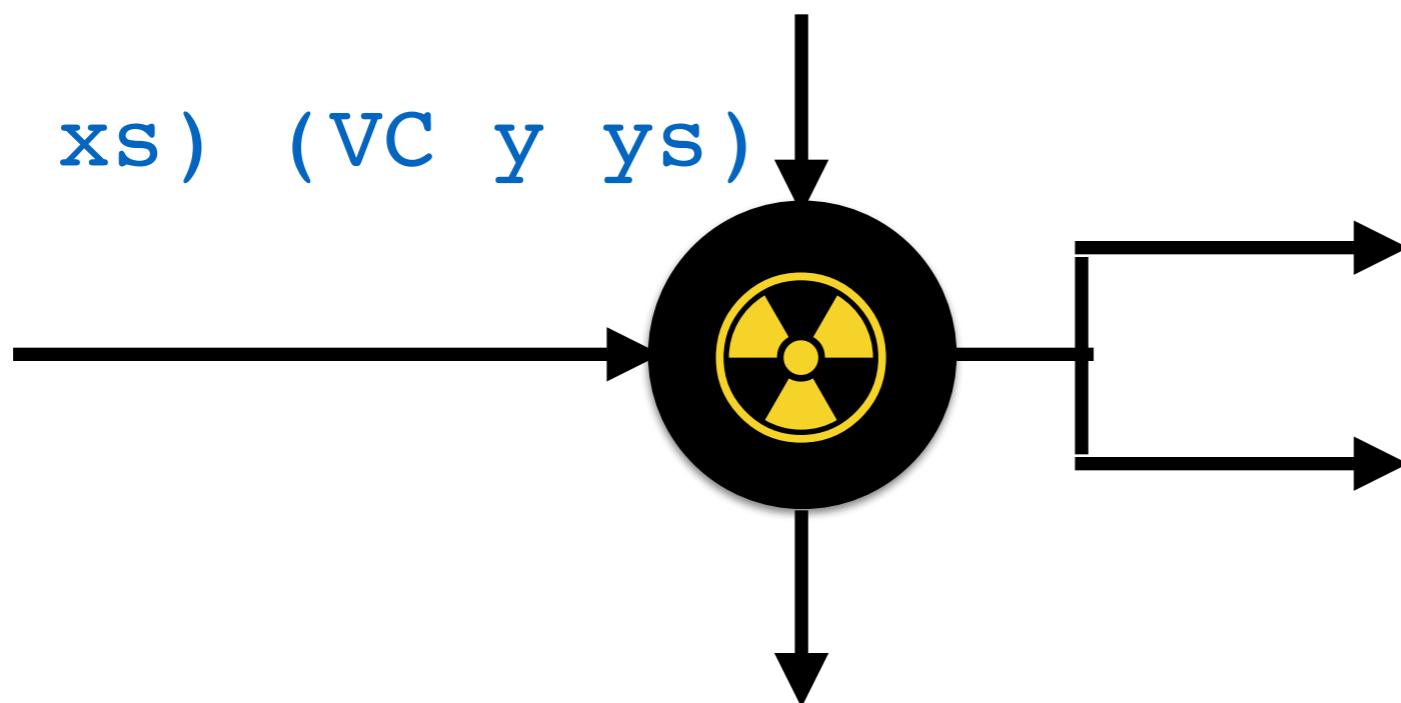
, ~~VN (VC _____)~~ \Rightarrow ~~(n - z) U (n - s - m)~~ }

inconsistent

continued

{ (VC _ _) _ \triangleright n ~ S m }

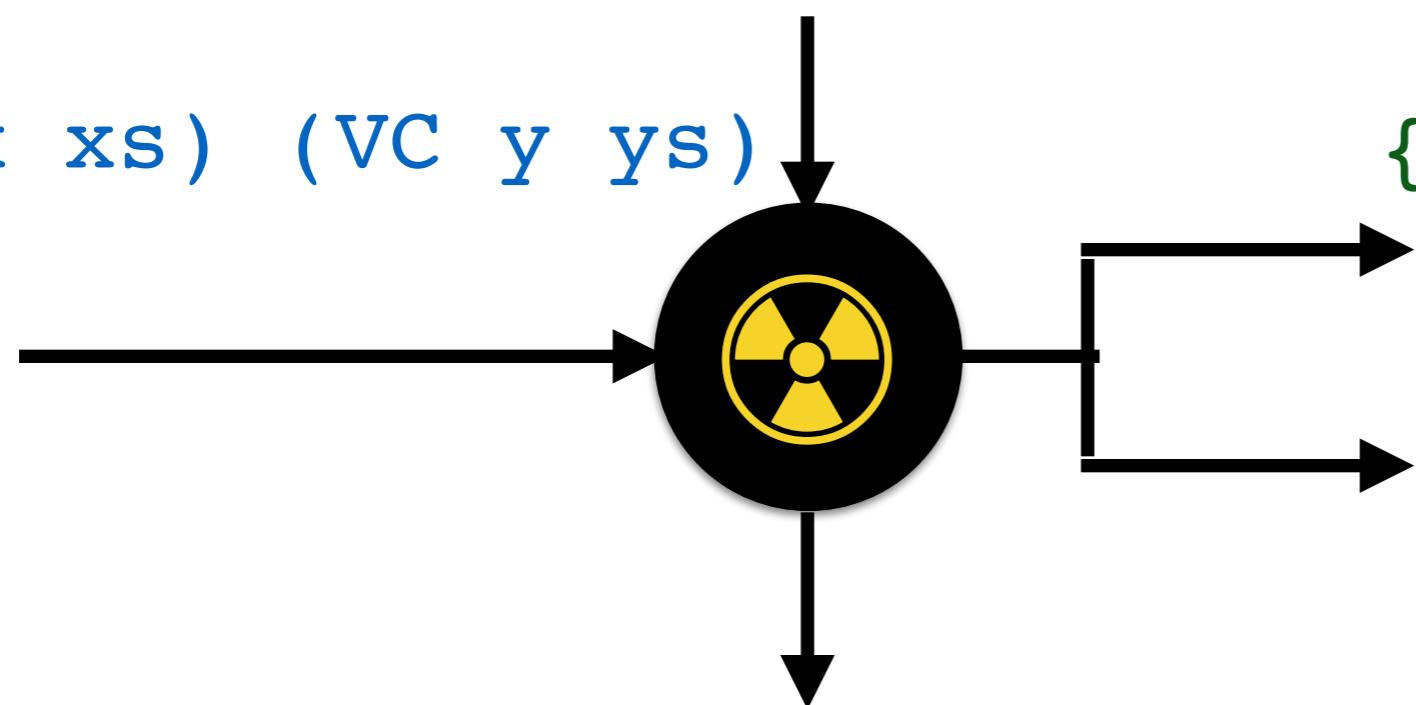
(VC x xs) (VC y ys)



continued

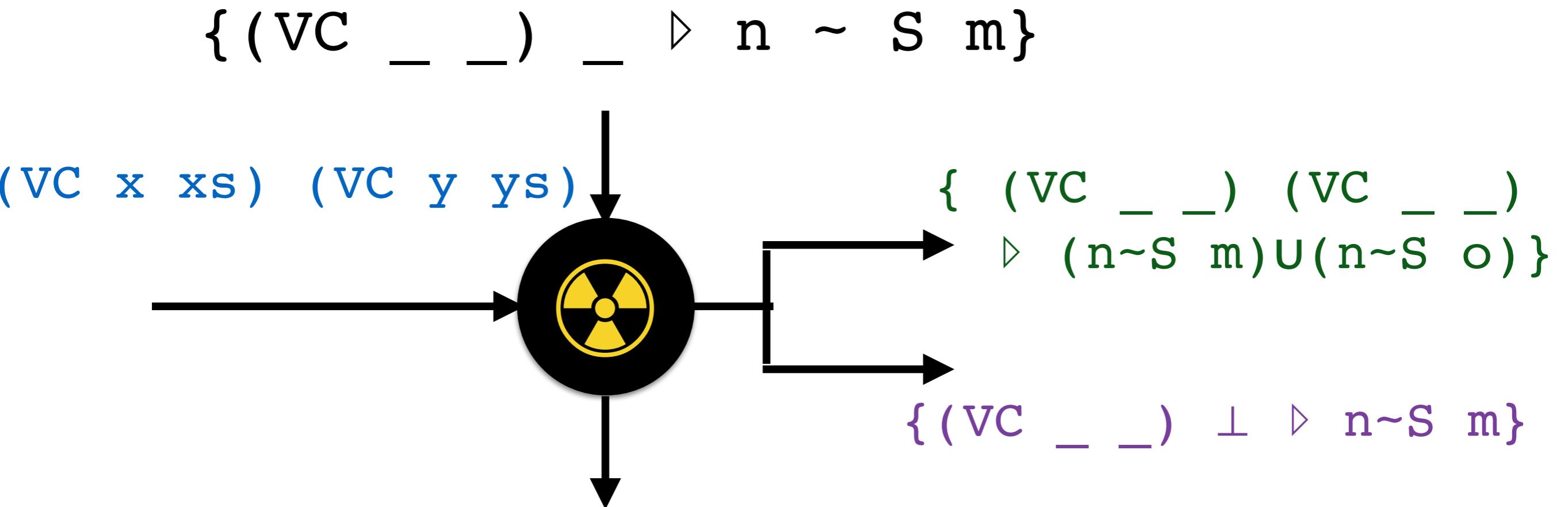
{ (VC _ _) _ } $\triangleright n \sim S m$

(VC x xs) (VC y ys)



{ (VC _ _) (VC _ _)
 $\triangleright (n \sim S m) \cup (n \sim S o)$ }

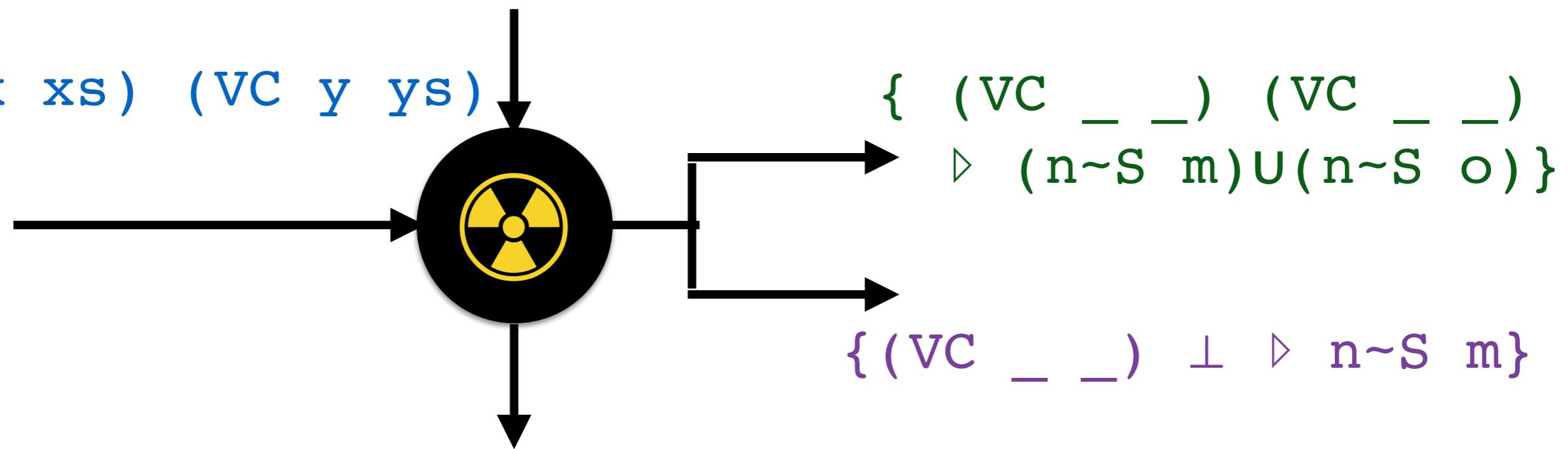
continued



continued

$$\{ (\text{VC } _ _) _ _ \triangleright n \sim S \ m \}$$

$$(\text{VC } x \ x_s) \ (\text{VC } y \ y_s)$$

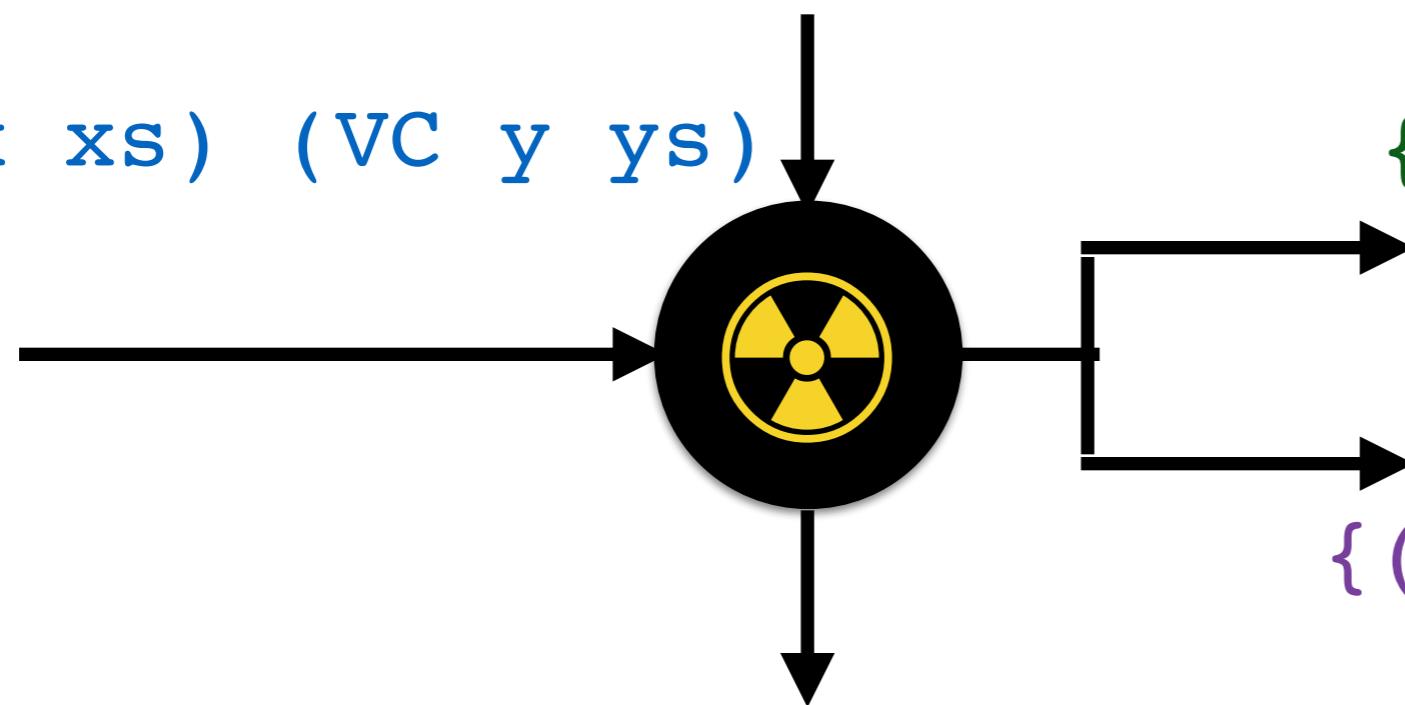


$$\{ (\text{VC } _ _) \text{ VN} \triangleright (n \sim S \ m) \cup (n \sim Z) \}$$

continued

$\{ (\text{VC } _ _) _ _ \triangleright n \sim S \ m \}$

$(\text{VC } x \ x_s) \ (\text{VC } y \ y_s)$



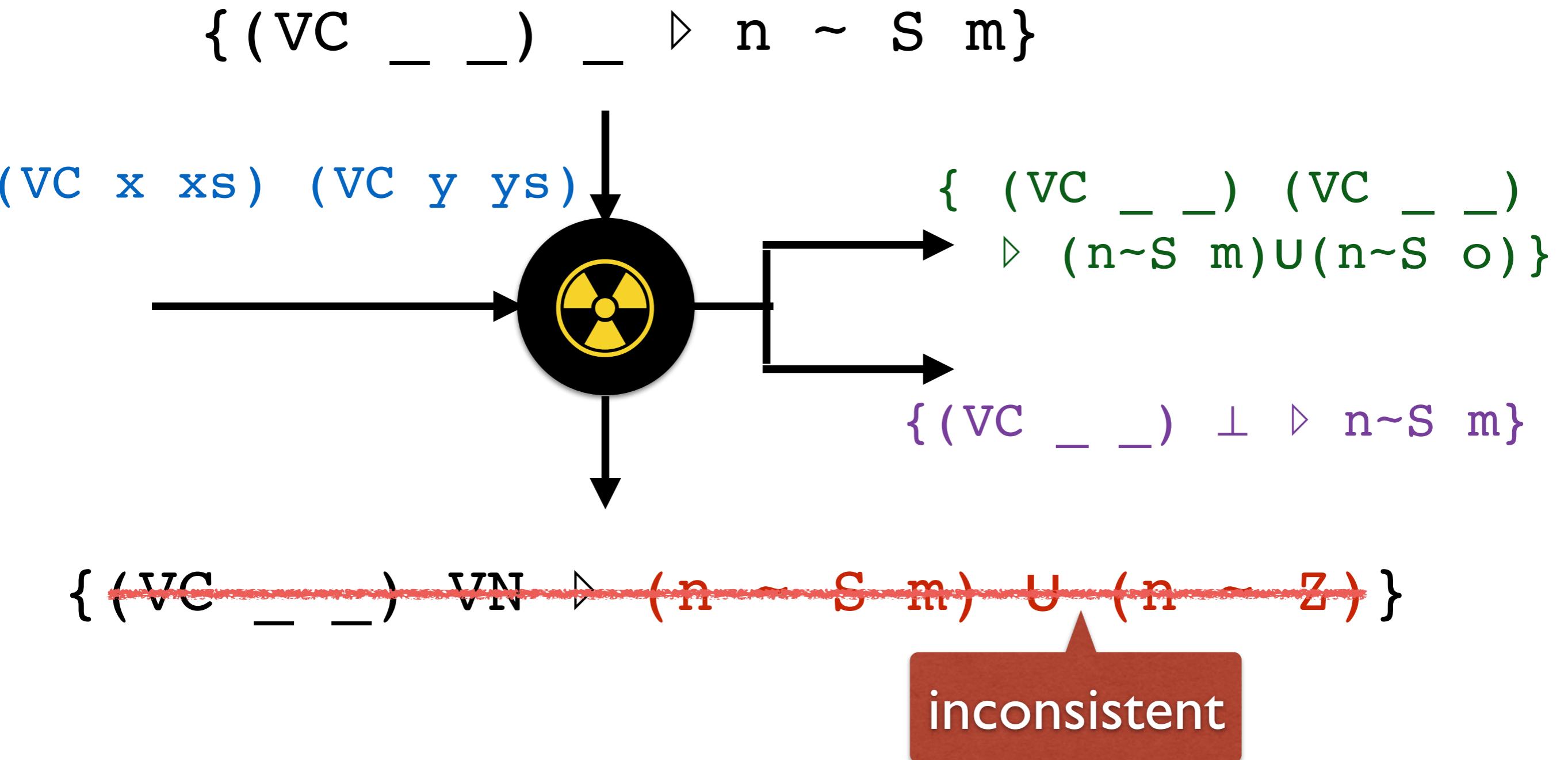
$\{ (\text{VC } _ _) (\text{VC } _ _) \triangleright (n \sim S \ m) \cup (n \sim S \ o) \}$

$\{ (\text{VC } _ _) \perp \triangleright n \sim S \ m \}$

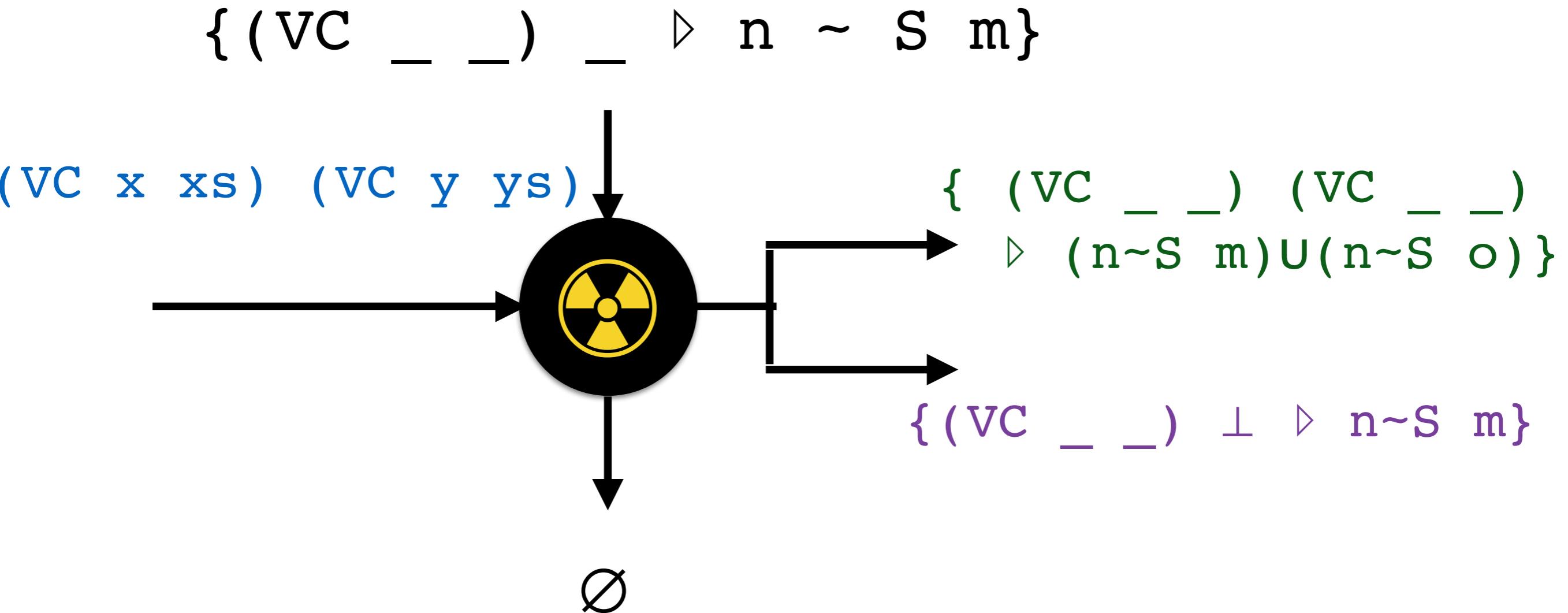
$\{ (\text{VC } _ _) \text{ VN } \triangleright (n \sim S \ m) \cup (n \sim Z) \}$

inconsistent

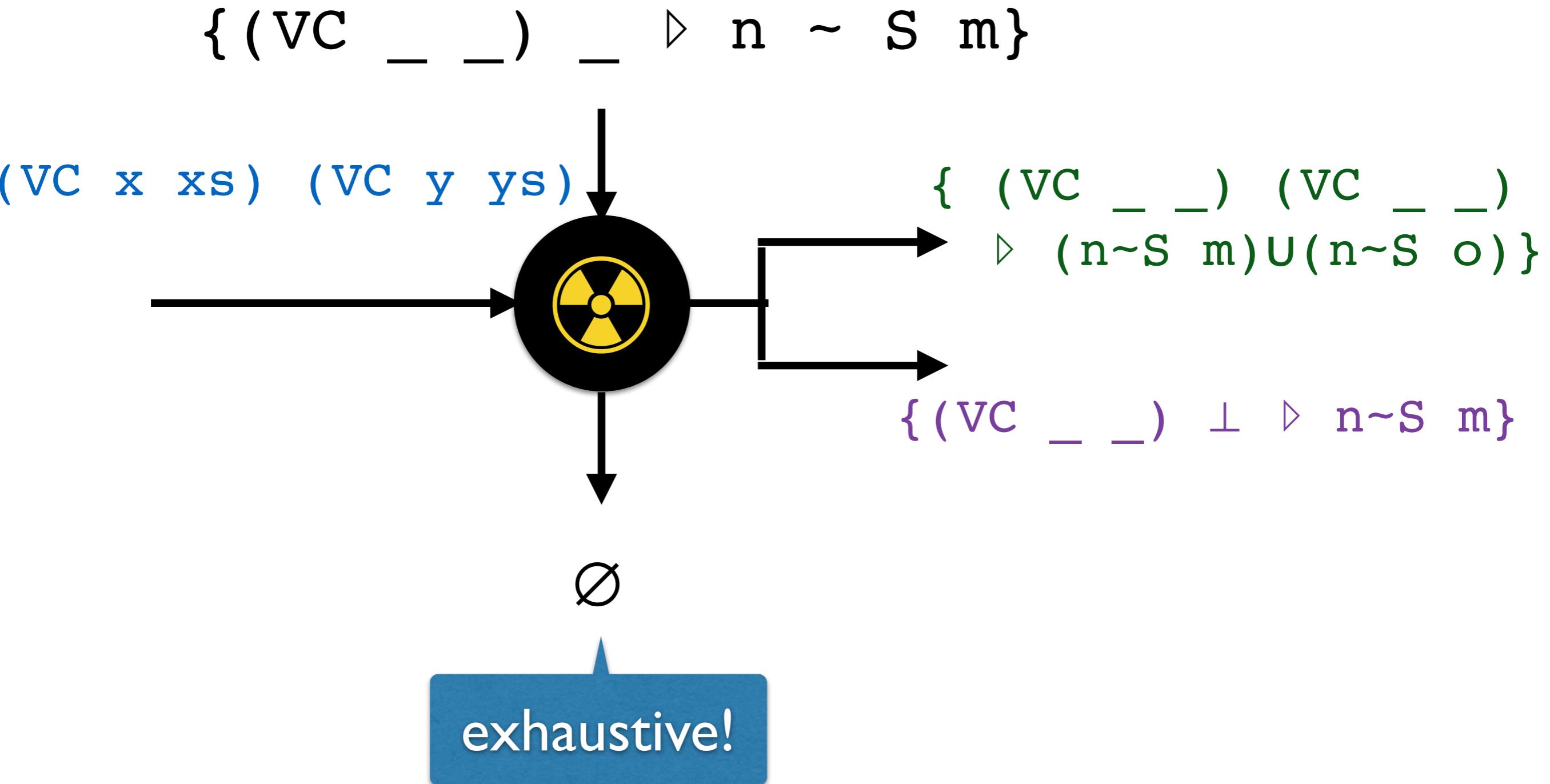
continued



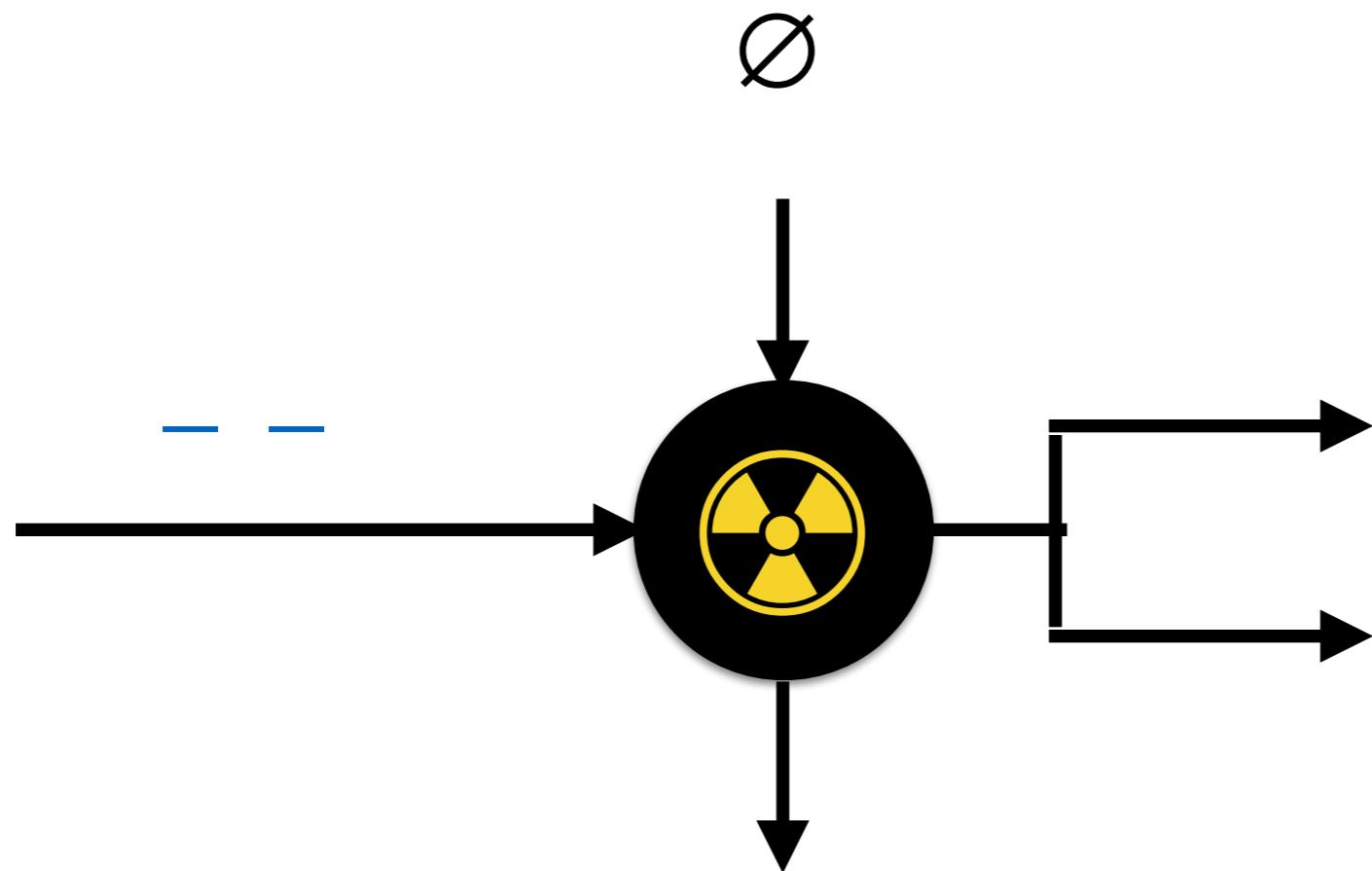
continued



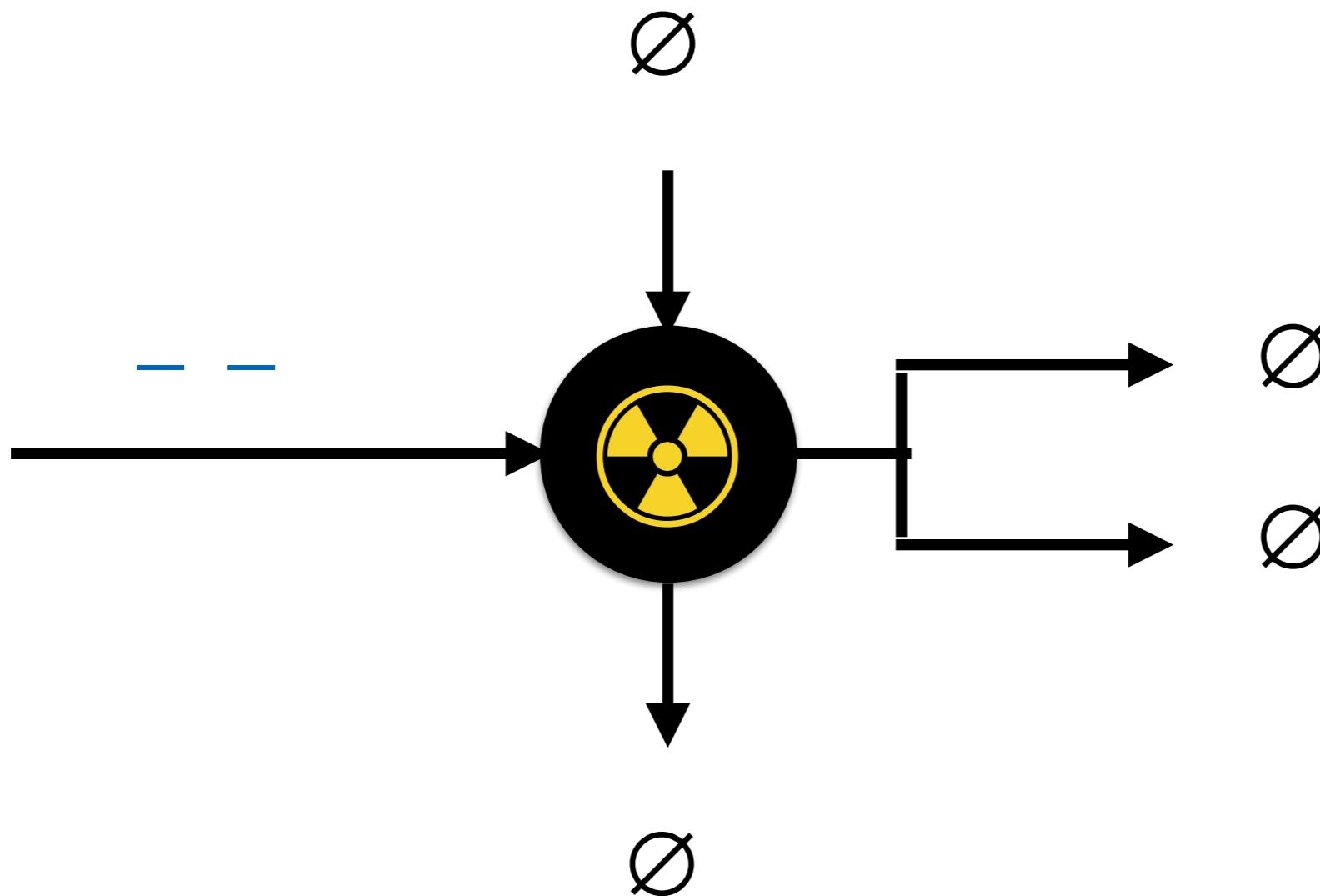
continued



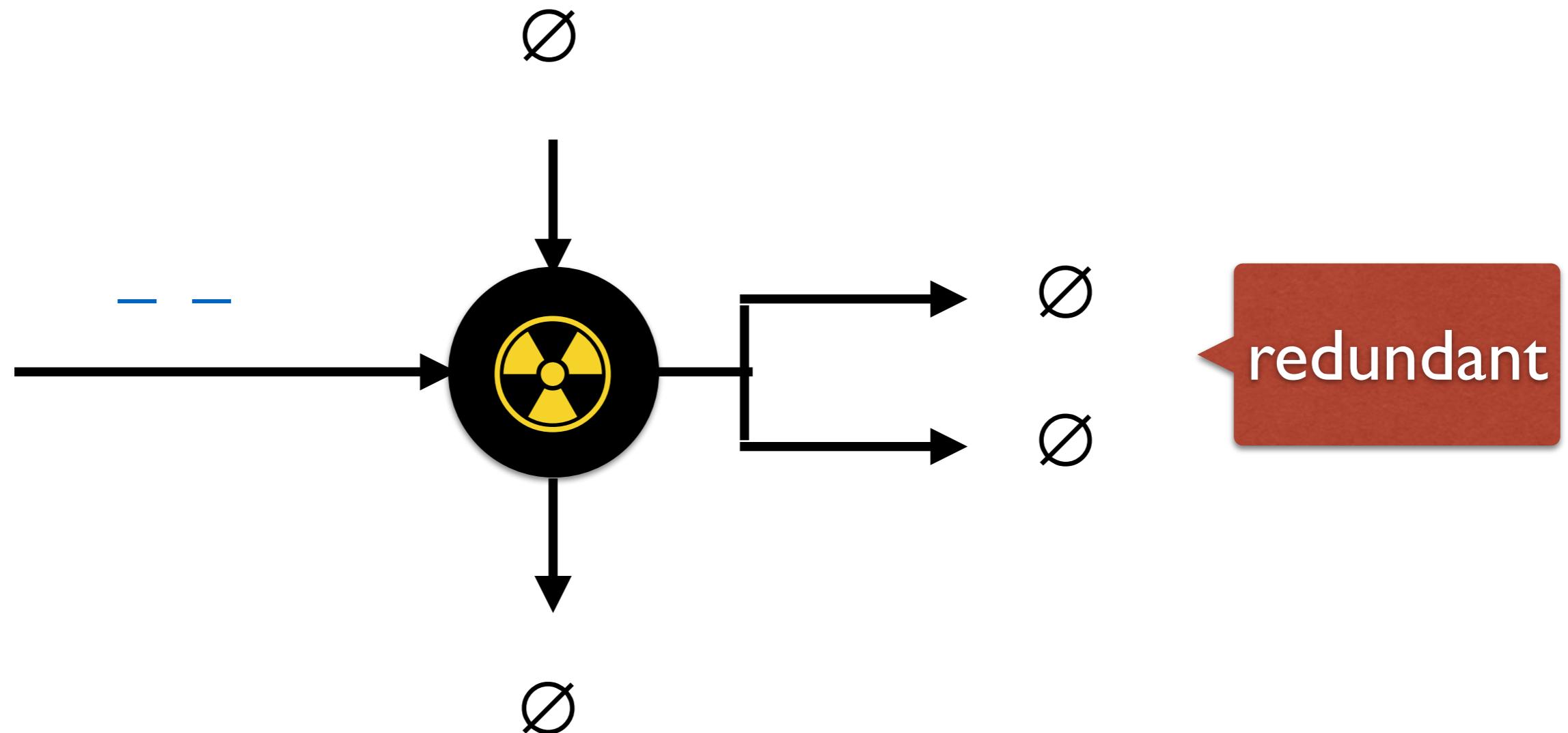
suppose one more catch-all clause



suppose one more catch-all clause



suppose one more catch-all clause



Guards

Guards

$\text{abs } x$	$ $	$x < 0$	$=$	\dots
	$ $	$x \geq 0$	$=$	\dots

Guard Core Syntax

Pattern

```
p ::= x
    | K p
    | p <- e
```

variable names pattern guards

```
abs x (True <- x < 0) = ...
abs x (True <- x >= 0) = ...
```

More Desugaring

literal patterns

```
f :: Num a => a -> [b] -> ...
f 0 [] = ...
```

More Desugaring

literal patterns

```
f :: Num a => a -> [b] -> ...
```

```
f 0 [] = ...
```



```
f x (True <- x == fromInteger 0) [] = ...
```

More Desugaring

literal patterns

```
f :: Num a => a -> [b] -> ...
```

```
f 0 [] = ...
```



```
f x (True <- x == fromInteger 0) [] = ...
```

view patterns

```
g (toRad -> (theta,r)) = ...
```

More Desugaring

literal patterns

```
f :: Num a => a -> [b] -> ...
```

```
f 0 [] = ...
```



```
f x (True <- x == fromInteger 0) [] = ...
```

view patterns

```
g (toRad -> (theta,r)) = ...
```



```
g x ((theta,r) <- toRad x) = ...
```

Guarded Value Abstraction

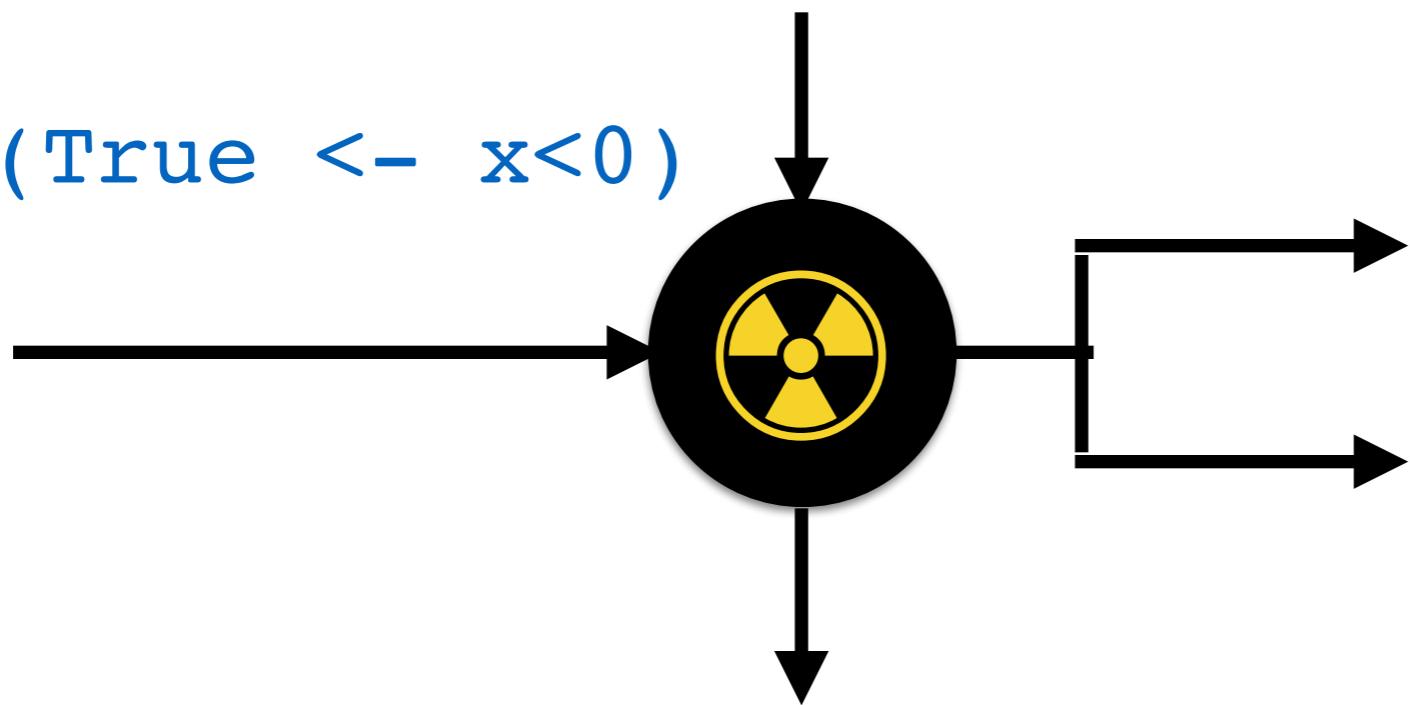
Value Abstraction

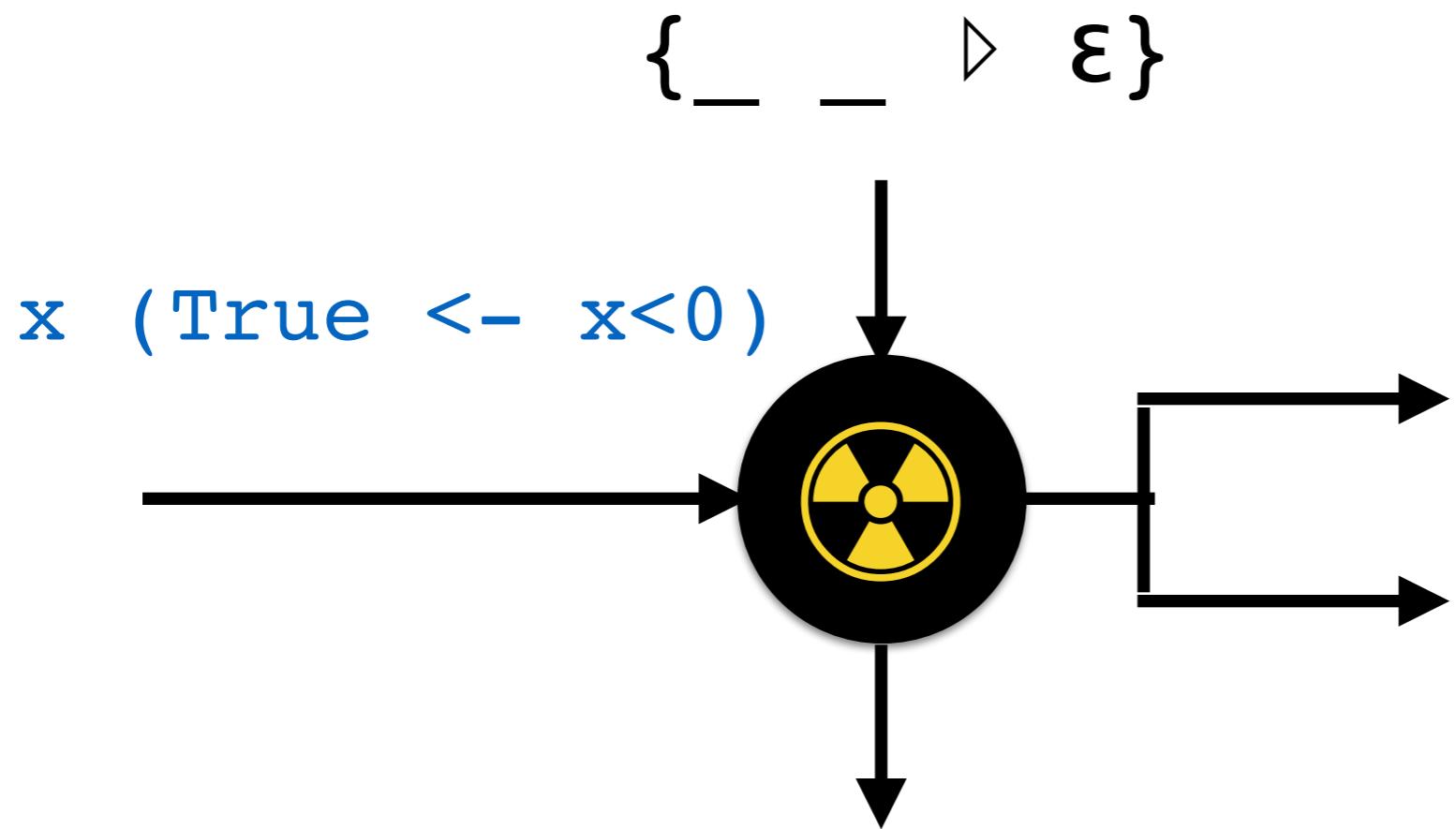
$$v ::= \Gamma \vdash \vec{u} \triangleright \Delta$$
$$u ::= x \mid K \vec{u}$$
$$\Gamma ::= \varepsilon \mid \Gamma, a \mid \Gamma, x:\tau$$
$$\Delta ::= \varepsilon \mid \Delta \cup \Delta \mid \tau \sim \tau$$
$$\begin{array}{l} | \quad x \approx e \\ | \quad x \approx \perp \end{array}$$

term constraint

strictness constraint

x (True \leftarrow $x < 0$)

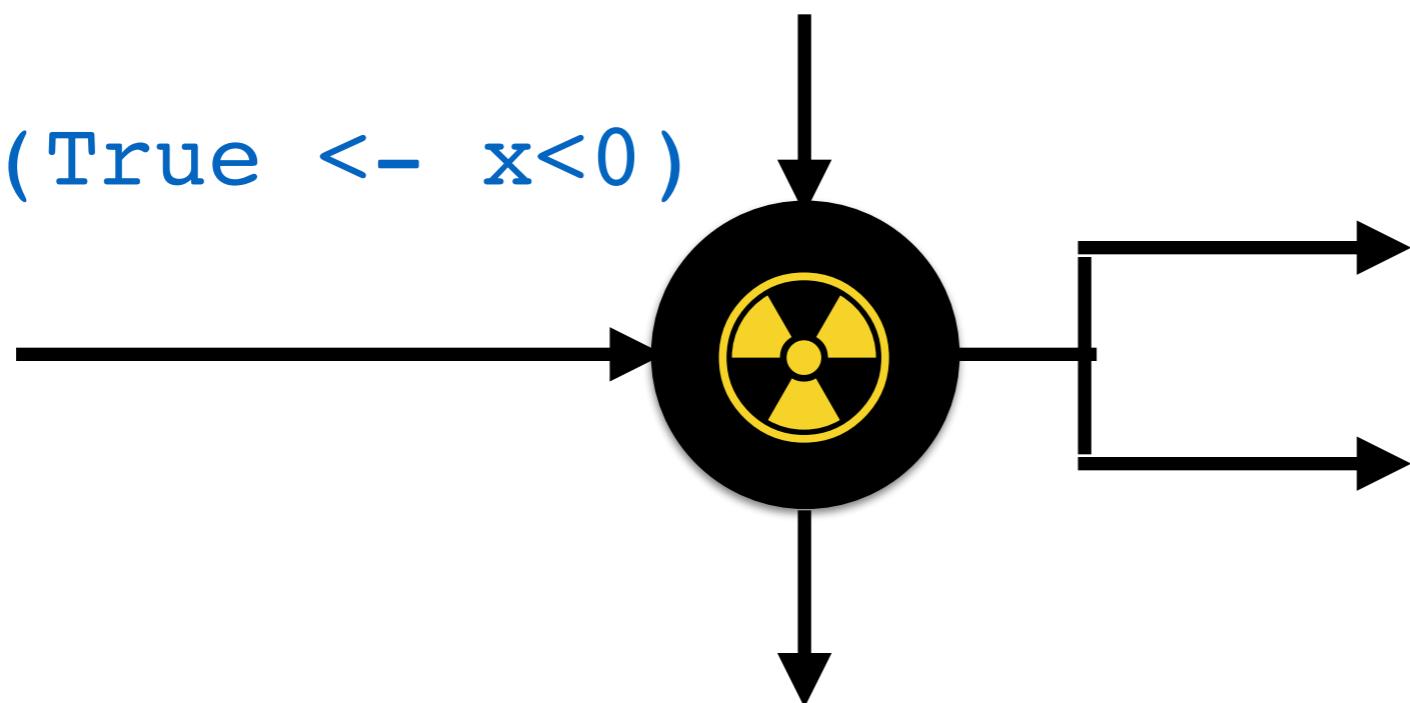




unconstrained

{ _ _ $\triangleright \varepsilon$ }

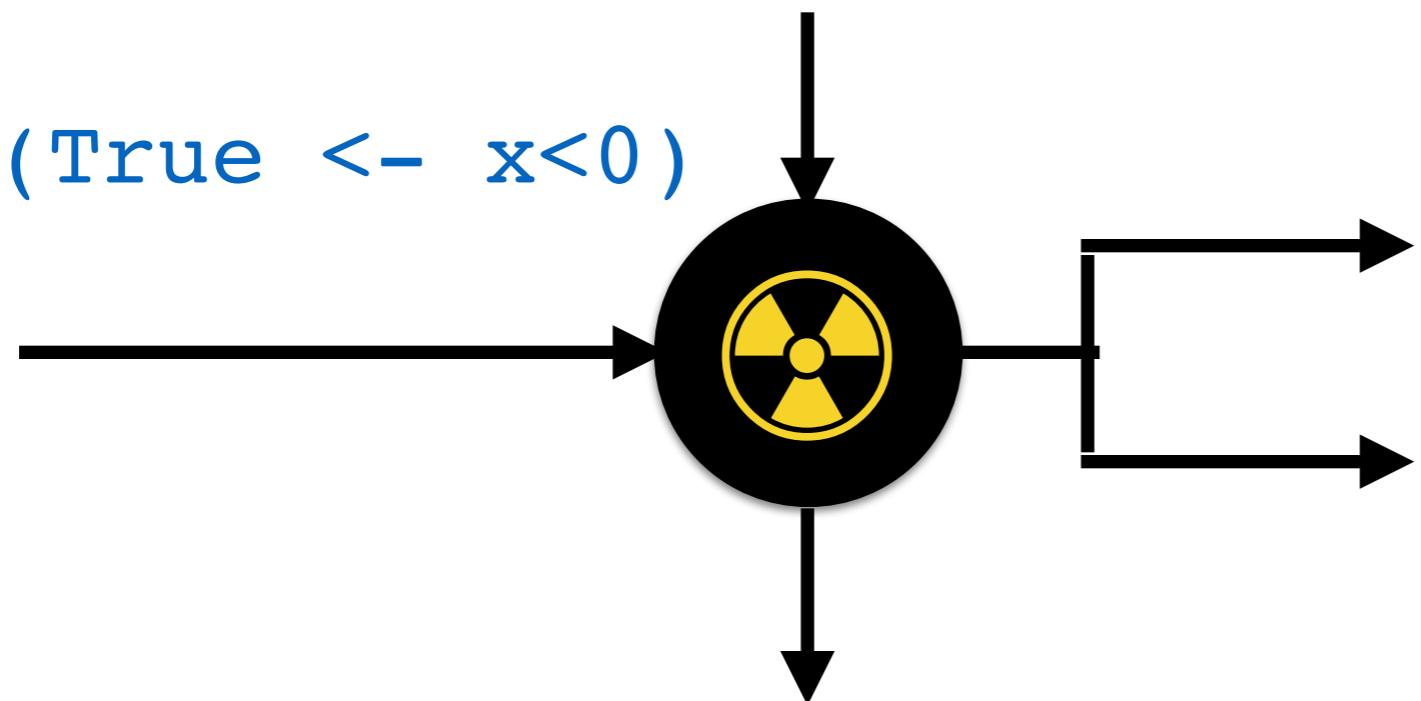
x (True $\leftarrow x < 0$)



unconstrained

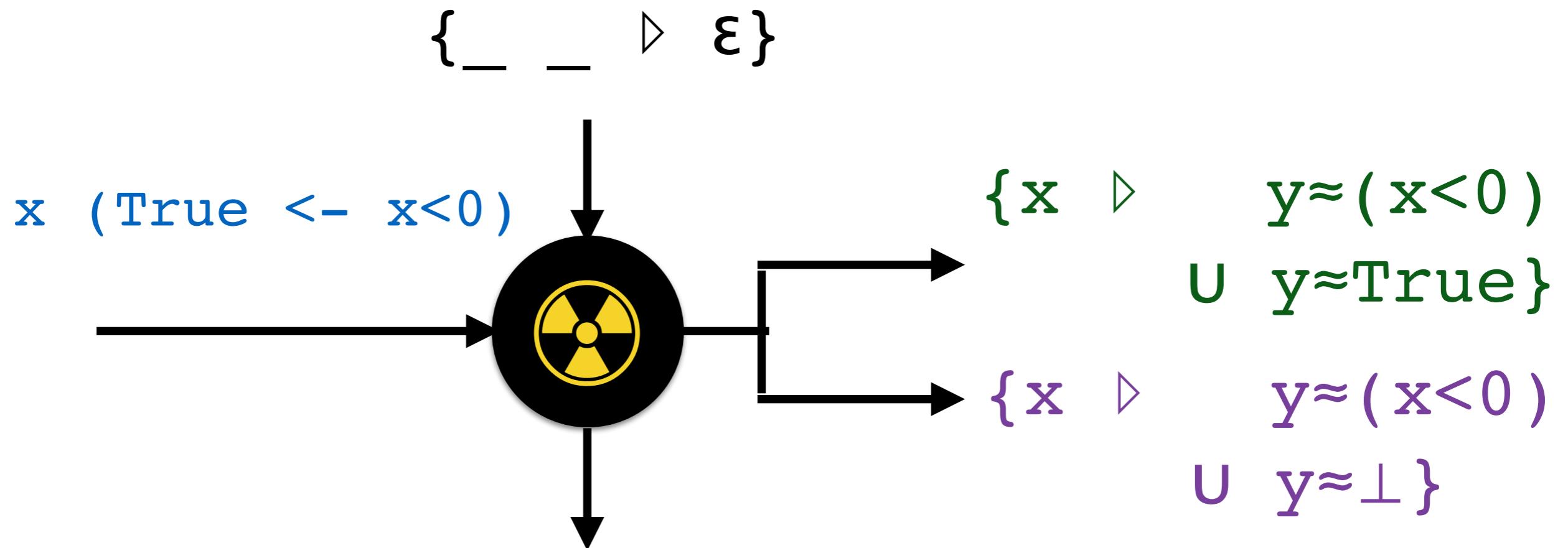
$$\{ _ _ \triangleright \varepsilon \}$$

x (True <- x<0)

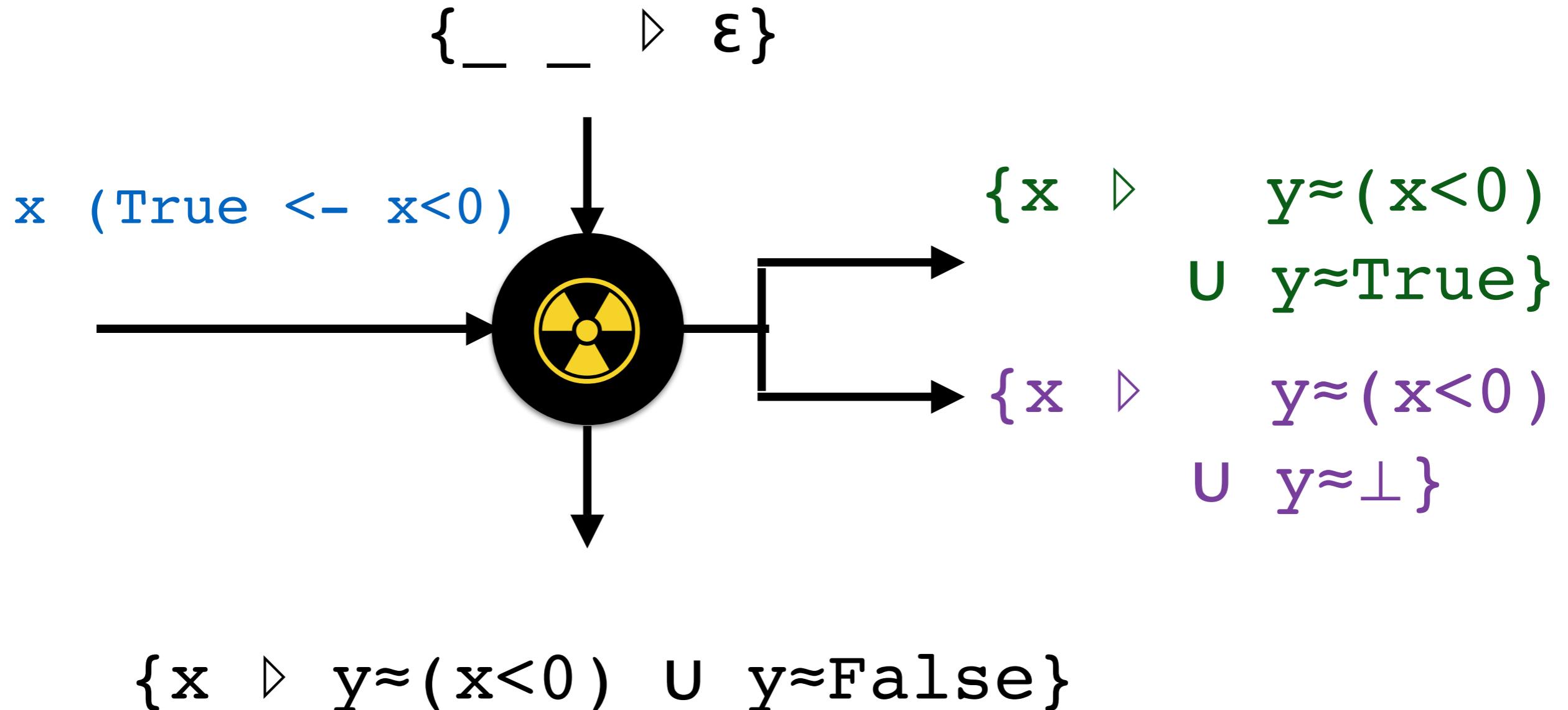


$\{ x \triangleright y \approx (x < 0) \cup y \approx \text{True} \}$

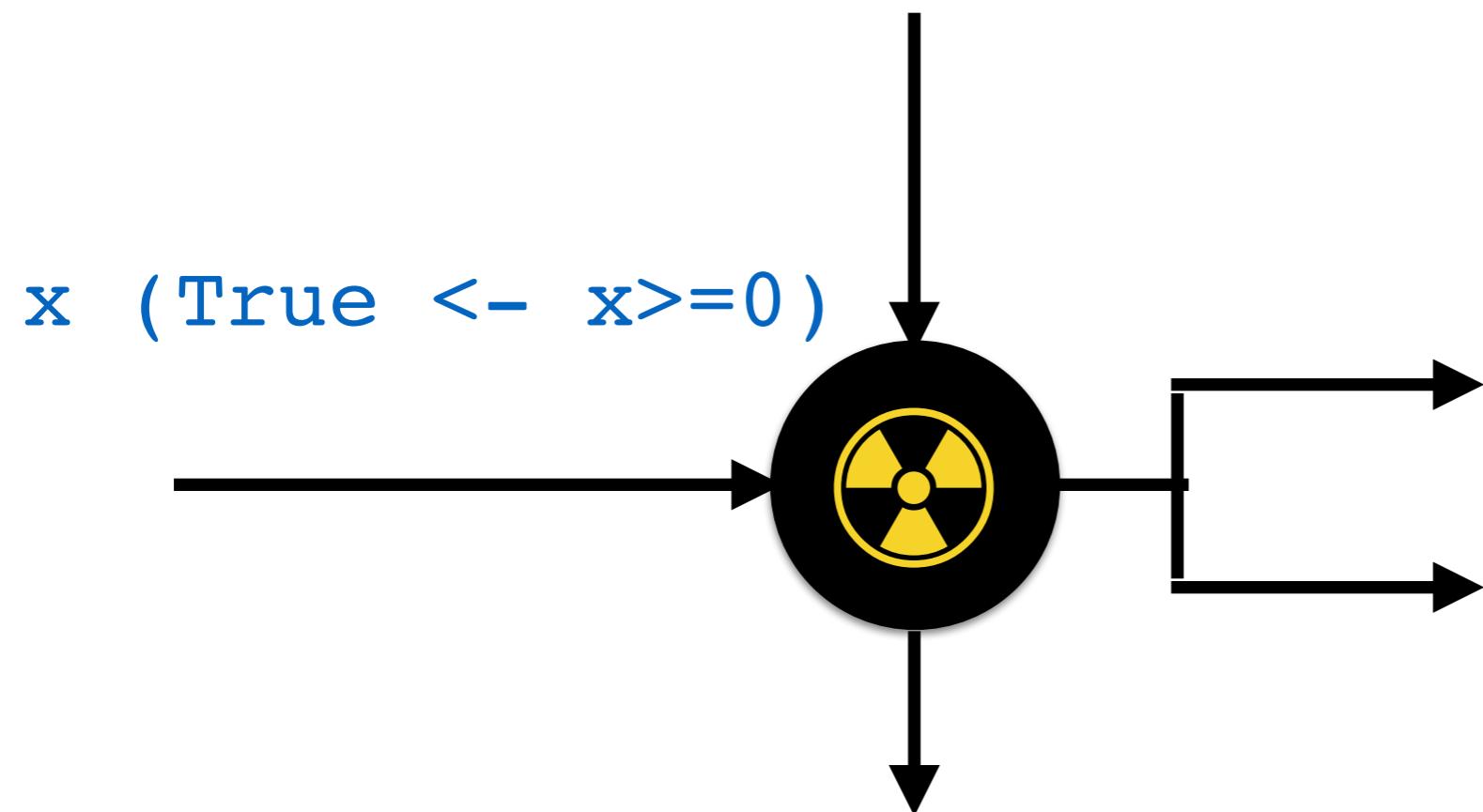
unconstrained



unconstrained

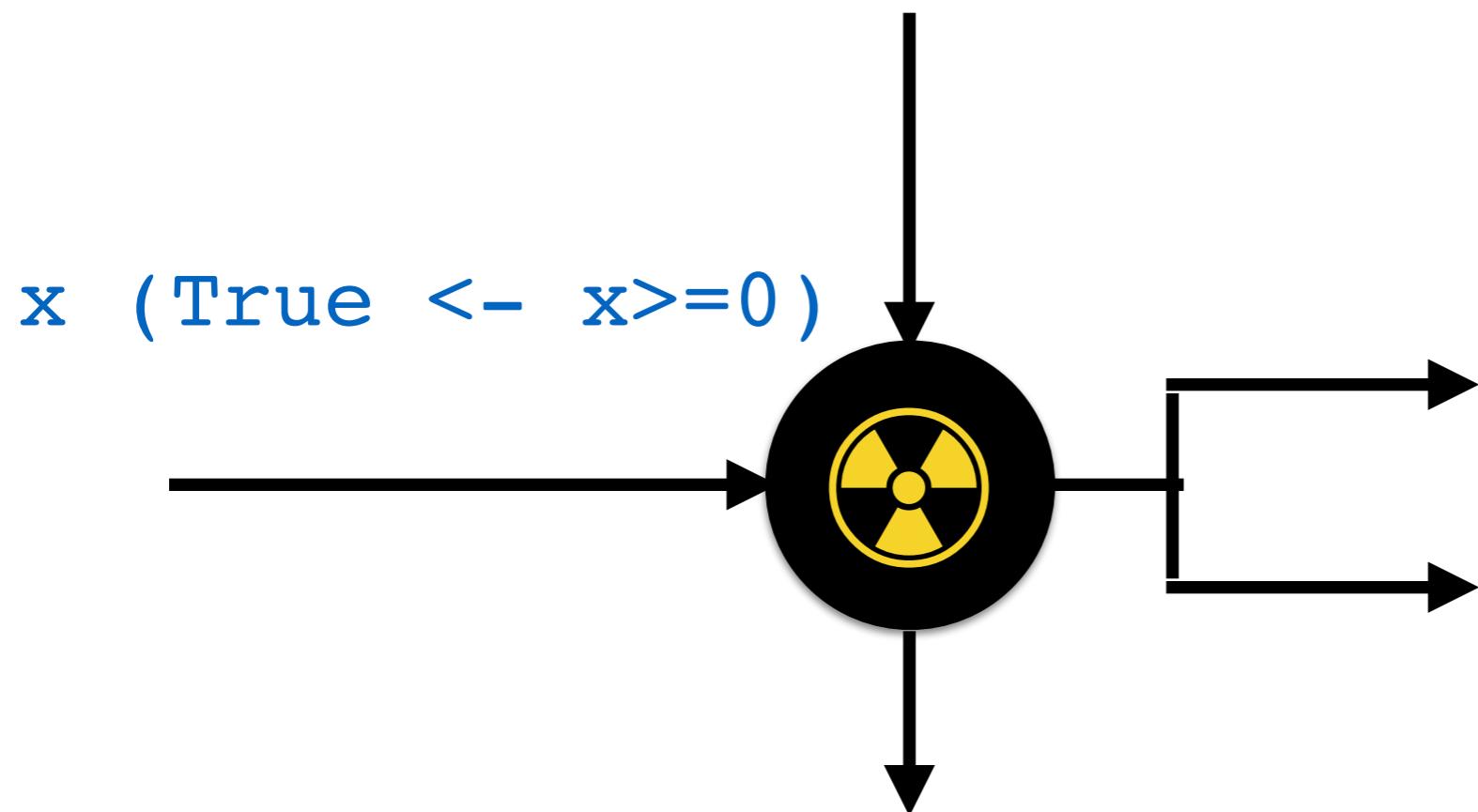


continued

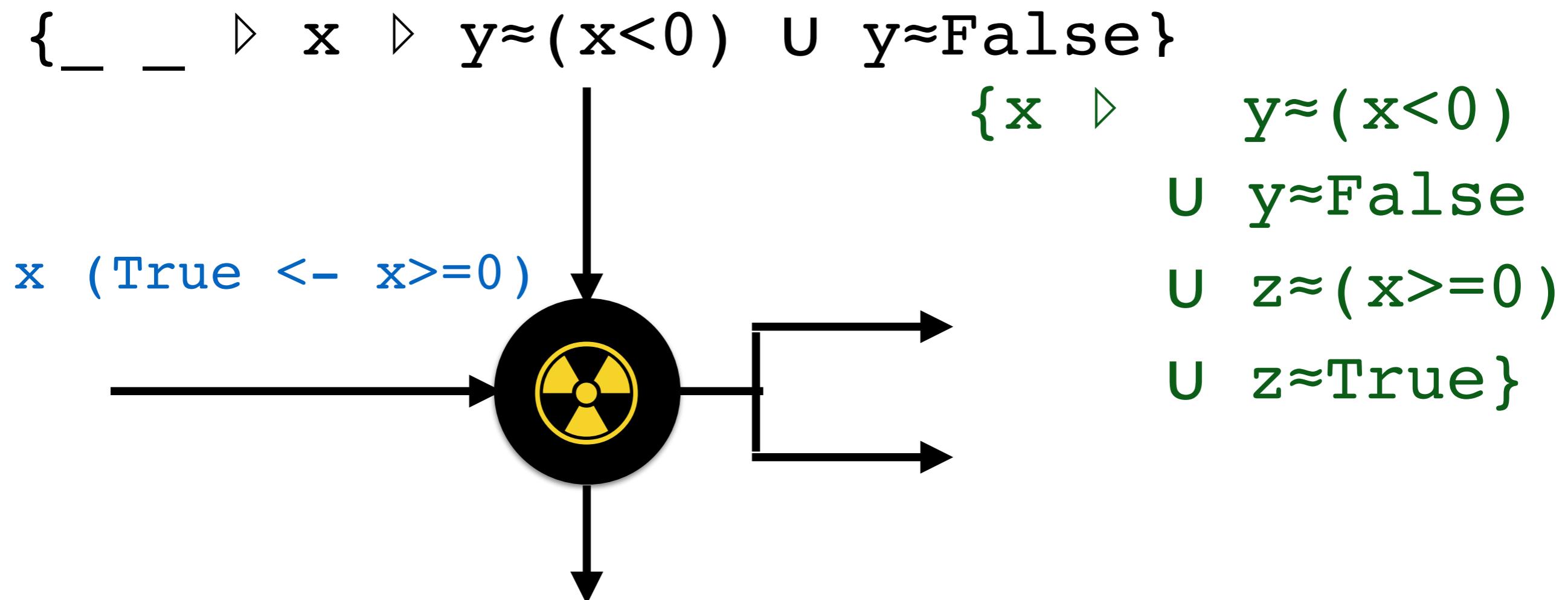


continued

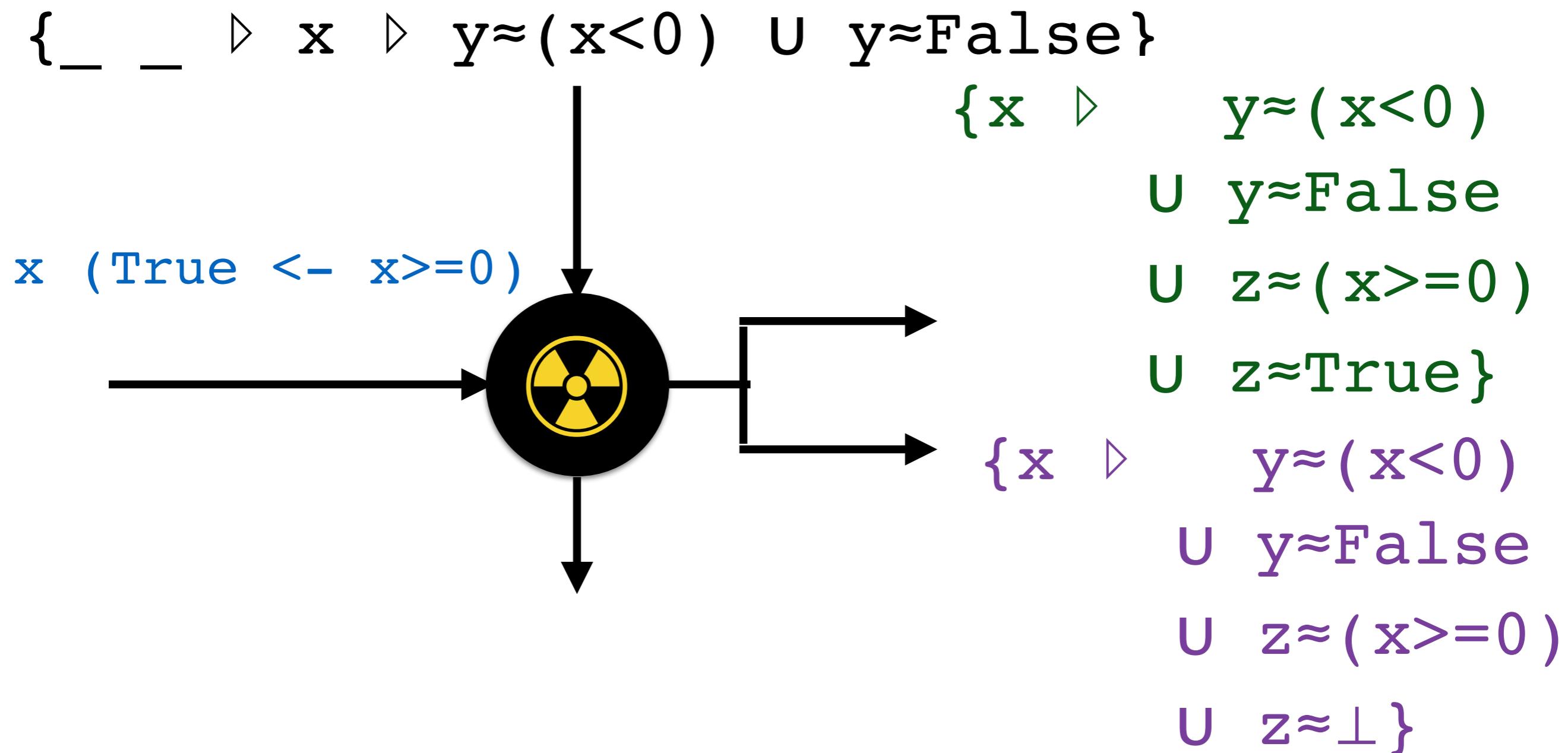
{ _ _ > x > y≈(x<0) U y≈False }



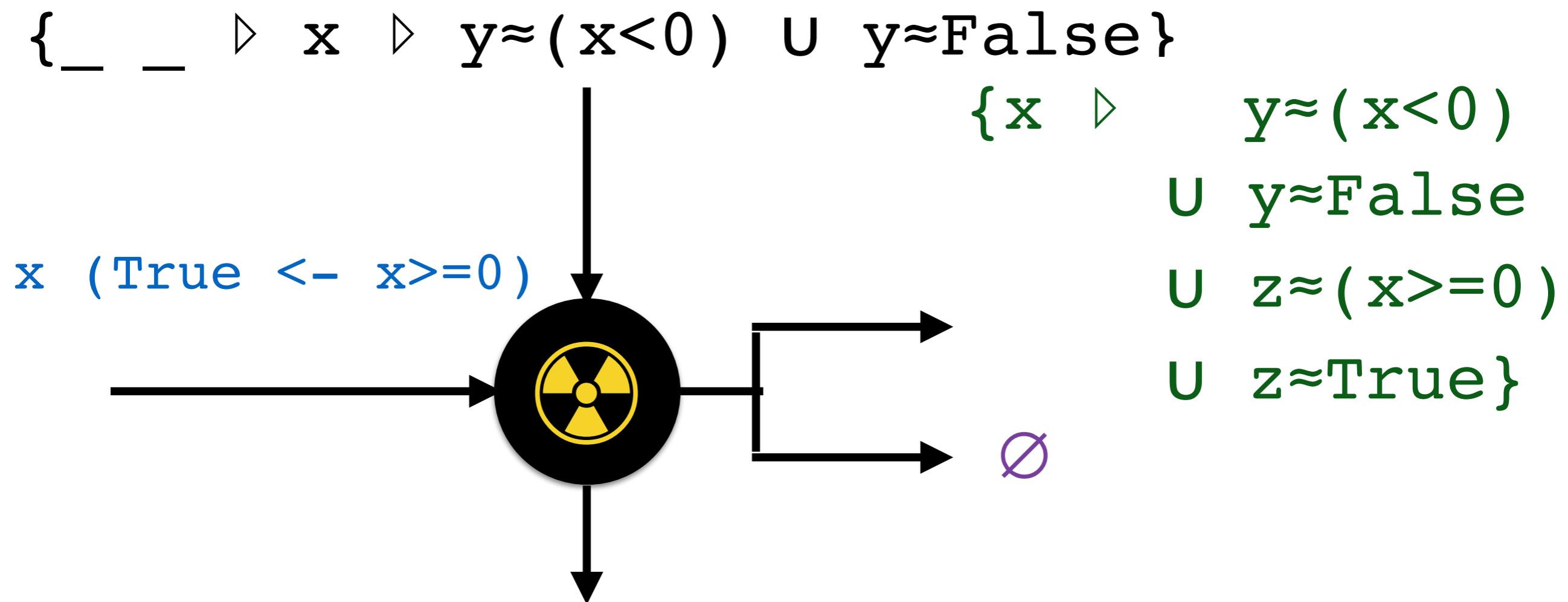
continued



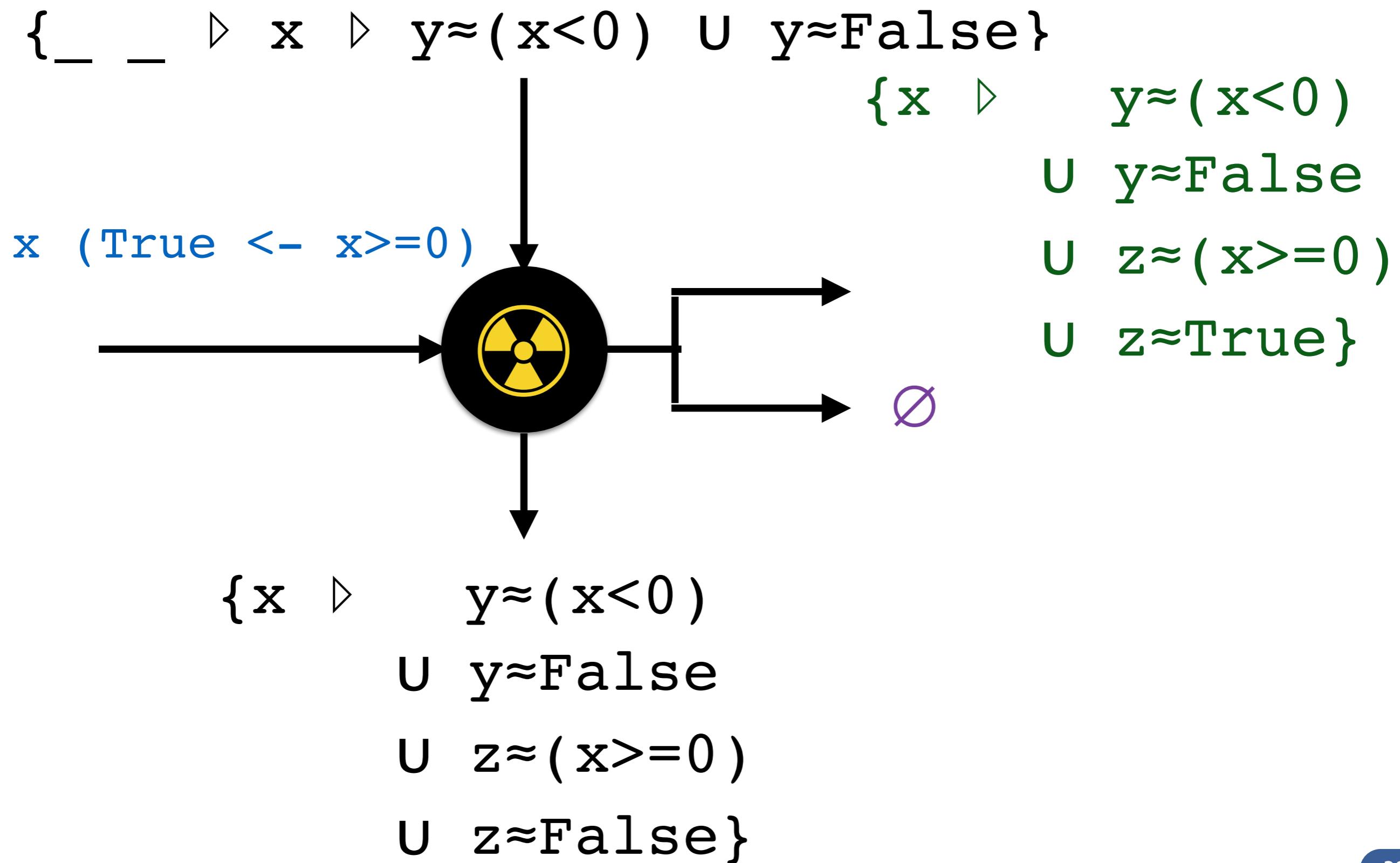
continued



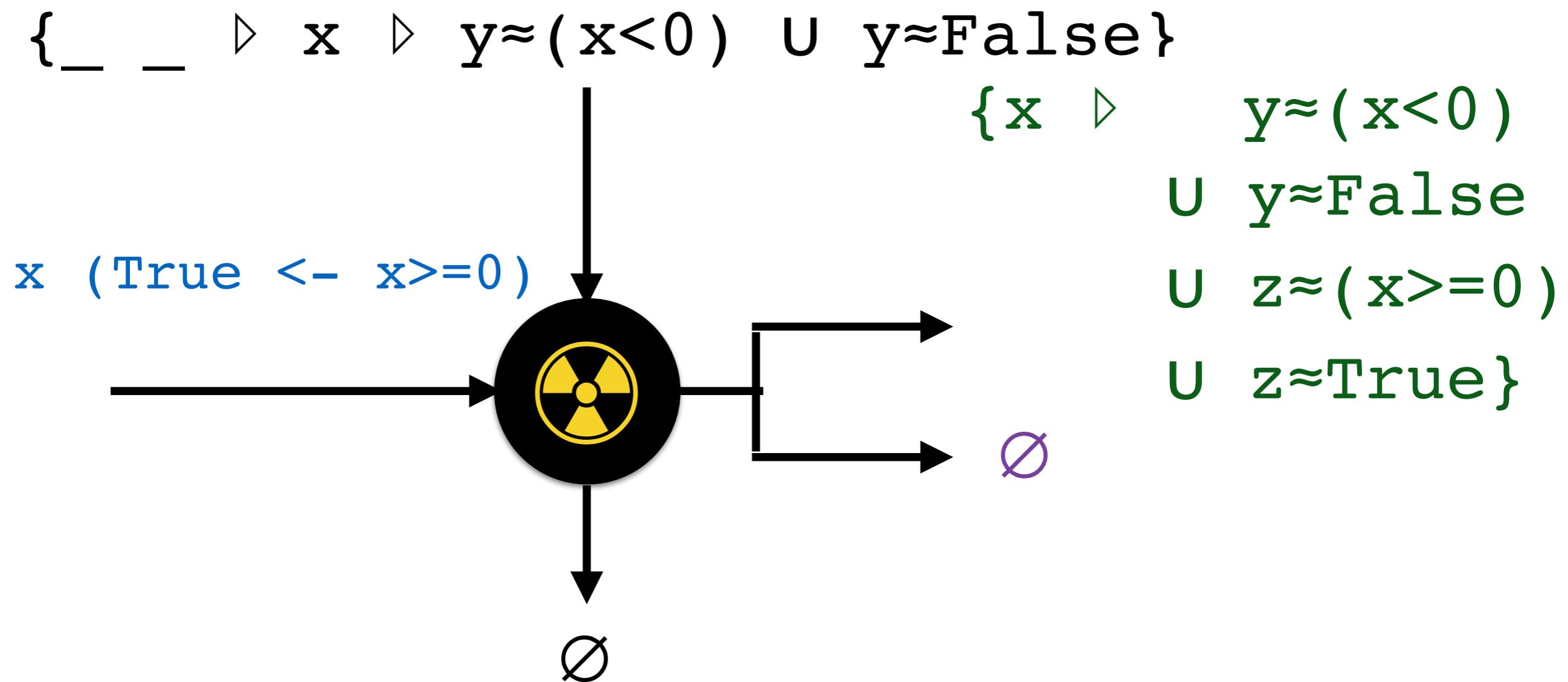
continued



continued



continued



Evaluation

Prototype Implementation

- GHC branch
- 504 LoC (vs. old 588 LoC)
- Type oracle: GHC type checker
- Term oracle:
for now only $(y \approx \text{False} \cup y \approx \text{True})$

Hackage Packages		LoC	M	R	M	R	M	R
		GHC-1		GHC-2		New		
<i>accelerate</i>	11,393	11	0	9	0	8	14	
<i>ad</i>	1,903	2	0	0	0	0	6	
<i>boolsimplifier</i>	256	10	0	0	0	0	0	
<i>d-bus</i>	2,753	45	0	42	0	16	1	
<i>generics-sop</i>	1,008	0	0	0	0	0	3	
<i>hoopl</i>	2,147	33	0	0	0	0	3	
<i>json-sop</i>	393	0	0	0	0	0	2	
<i>lens-sop</i>	280	2	0	0	0	0	2	
<i>pretty-sop</i>	27	0	0	0	0	0	1	

Additional tests		LoC	M	R	M	R	M	R
		GHC-1		GHC-2		New		
<i>lists</i>	66	1	0	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	0	1

		LoC	M	R	M	R	M	R		
		Original		GHC		GHC-1		GHC-2		New
Hackage Packages										
<i>accelerate</i>	11			0	9	0	8	14		
<i>ad</i>	1			0	0	0	0	6		
<i>boolsimplifier</i>	256	10	0	0	0	0	0	0		
<i>d-bus</i>	2,753	45	0	42	0	16	1			
<i>generics-sop</i>	1,008	0	0	0	0	0	0	3		
<i>hoopl</i>	2,147	33	0	0	0	0	0	3		
<i>json-sop</i>	393	0	0	0	0	0	0	2		
<i>lens-sop</i>	280	2	0	0	0	0	0	2		
<i>pretty-sop</i>	27	0	0	0	0	0	0	1		
Additional tests		LoC	M	R	M	R	M	R		
<i>lists</i>	66	1	0	0	0	0	0	3		
<i>heterogeneous lists</i>	38	0	0	0	0	0	0	2		
<i>heaps</i>	540	3	0	0	0	0	0	1		

Hackage Packages	LoC	M	R	GHC-1	GHC-2	New	
				Original	Adhoc	M	R
<i>accelerate</i>	11	0	0	0	0	8	14
<i>ad</i>	1	0	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,008	0	0	0	0	0	3
<i>hoopl</i>	2,147	33	0	0	0	0	3
<i>json-sop</i>	393	0	0	0	0	0	2
<i>lens-sop</i>	280	2	0	0	0	0	2
<i>pretty-sop</i>	27	0	0	0	0	0	1

Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

		LoC	M	R	M	R	M	R	
		Original GHC			Adhoc Patch		Our Approach		
Hackage Packages		I	S	M	R	M	R	M	R
<i>accelerate</i>	11								
<i>ad</i>	1								
<i>boolsimplifier</i>		256	10	0	0	0	0	0	0
<i>d-bus</i>		2,753	45	0	42	0	16	1	
<i>generics-sop</i>		1,008	0	0	0	0	0	0	3
<i>hoopl</i>		2,147	33	0	0	0	0	0	3
<i>json-sop</i>		393	0	0	0	0	0	0	2
<i>lens-sop</i>		280	2	0	0	0	0	0	2
<i>pretty-sop</i>		27	0	0	0	0	0	0	1
Additional tests		LoC	M	R	M	R	M	R	
<i>lists</i>		66	1	0	0	0	0	0	3
<i>heterogeneous lists</i>		38	0	0	0	0	0	0	2
<i>heaps</i>		540	3	0	0	0	0	0	1

Hackage Packages		LoC	M	R	M	R	M	R
<i>accelerate</i>		11,393	11	0	9	0	8	14
<i>ad</i>		1,903	2	0	0	0	0	6
<i>boolsimplifier</i>		256	10	0	0	0	0	0
<i>d-bus</i>		2,753	45	0	42	0	16	1
<i>generics-sop</i>				0	0	0	0	3
<i>hoopl</i>				0	0	0	0	3
<i>json-sop</i>				0	0	0	0	2
<i>lens-sop</i>				0	0	0	0	2
<i>pretty-sop</i>				0	0	0	0	1
libraries with checking problems								
Additional tests			R	M	R	M	R	
<i>lists</i>		66	1	0	0	0	0	3
<i>heterogeneous lists</i>		38	0	0	0	0	0	2
<i>heaps</i>		540	3	0	0	0	0	1

		GHC-1		GHC-2		New	
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,003	0	0	0	0	0	3
<i>hoopl</i>	7	33	0	0	0	0	3
<i>json-sop</i>	8	0	0	0	0	0	2
<i>lens-sop</i>	0	2	0	0	0	0	2
<i>pretty-sop</i>	7	0	0	0	0	0	1
clauses erroneously reported Missing							
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

		GHC-1		GHC-2		New	
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,003	0	0	0	0	0	3
<i>hoopl</i>	7,003	33	0	0	0	0	3
<i>json-sop</i>	3,003	0	0	0	0	0	2
<i>lens-sop</i>	1,003	2	0	0	0	0	2
<i>pretty-sop</i>	7,003	0	0	0	0	0	1
clauses erroneously reported Missing							
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

		GHC-1		GHC-2		New	
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,003	0	0	0	0	0	3
<i>hoopl</i>	7,337	33	0	0	0	0	3
<i>json-sop</i>	3,033	0	0	0	0	0	2
<i>lens-sop</i>	1,000	2	0	0	0	0	2
<i>pretty-sop</i>	7,337	0	0	0	0	0	1
clauses erroneously reported Missing							
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

			GHC-1	GHC-2	New		
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	1
<i>ad</i>	1,903	2	0	0	0	0	0
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,003	0	0	0	0	0	0
<i>hoopl</i>	7	33	0	0	0	0	0
<i>json-sop</i>	8	0	0	0	0	0	2
<i>lens-sop</i>	0	2	0	0	0	0	2
<i>pretty-sop</i>	7	0	0	0	0	0	1
clauses erroneously reported Missing							
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

better
term
oracle
needed

		GHC-1	GHC-2	New			
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,888	8	0	0	0	0	3
<i>hoopl</i>			0	0	0	0	3
<i>json-sop</i>			0	0	0	0	2
<i>lens-sop</i>			0	0	0	0	2
<i>pretty-sop</i>			0	0	0	0	1
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

clauses
reported
Redundant

		GHC-1	GHC-2	New			
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,888	8	0	0	0	0	3
<i>hoopl</i>			clauses	0	0	0	3
<i>json-sop</i>			reported	0	0	0	2
<i>lens-sop</i>			Redundant	0	0	0	2
<i>pretty-sop</i>				0	0	0	1
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

		GHC-1	GHC-2	New			
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,888	8	0	0	0	0	3
<i>hoopl</i>			0	0	0	0	3
<i>json-sop</i>			0	0	0	0	2
<i>lens-sop</i>			0	0	0	0	2
<i>pretty-sop</i>			0	0	0	0	1
clauses reported Redundant							
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	0	0	0	2
<i>heaps</i>	540	3	0	0	0	0	1

		GHC-1	GHC-2	New			
Hackage Packages	LoC	M	R	M	R	M	R
<i>accelerate</i>	11,393	11	0	9	0	8	14
<i>ad</i>	1,903	2	0	0	0	0	6
<i>boolsimplifier</i>	256	10	0	0	0	0	0
<i>d-bus</i>	2,753	45	0	42	0	16	1
<i>generics-sop</i>	1,888	8	0	0	0	0	3
<i>hoopl</i>			0	0	0	0	3
<i>json-sop</i>			0	0	0	0	2
<i>lens-sop</i>			0	0	0	0	2
<i>pretty-sop</i>			0	0	0	0	1
Additional tests	LoC	M	R	M	R	M	R
<i>lists</i>	66	1	0	0	0	0	3
<i>heterogeneous lists</i>	38	0	0	f	—	= ...	2
<i>heaps</i>	540	3	0	0	0	0	1

clauses
reported
Redundant

only
suppressing
catch-all
clauses

f — = ...

Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
10 – 99	181	2.04
100 – 2813	5	0.06

Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
10 – 99	181	2.04
100 – 2813	5	0.06

maximum size of
Covered/Uncovered
sets per match

Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
10 – 99	181	2.04
100 – 2813	5	0.06

Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
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Space Explosion?

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Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
10 – 99	181	2.04
100 – 2813	5	0.06

all in package ad

```
data T = A | B | C  
f A A = True  
f B B = True  
f C C = True
```

Space Explosion?

Maximum size of C/U	Pattern Matches	(%)
1 – 9	8702	97.90
10 – 99	181	2.04
100 – 2813	5	0.06

54 x 54
constructors

all in package ad

```
data T = A | B | C
f A A = True
f B B = True
f C C = True
```

Wrapping Up

Summary

A uniform approach to deal with:

- ❖ GADTs
- ❖ Laziness
- ❖ Guards and extensions
 - pattern guards, view patterns, literal patterns, ...

Coming soon
in GHC release

More in the Paper

- Full Algorithm
- Optimisation Opportunities
- Meta-Theory
- Related Work

GADTs meet their match
Pattern-matching warnings that account for GADTs, guards, and laziness.

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Abstract
For ML and Haskell, accurate warnings when a function definition contains redundant or missing patterns are mission critical. But today's compilers generate bogus warnings when the programmer uses guards (even simple ones), GADTs, pattern guards, or view patterns. We give the first algorithm that handles all these cases in a single, unified framework, together with an implementation in GHC, and evidence of its utility in practice.

Categories and Subject Descriptors CR-number [subcategory]: third-level
Keywords Haskell, pattern matching, Generalized Algebraic Data Types, OUTSIDER

and exhaustiveness warnings when these features are used. Certainly our own compiler, GHC, does not, and nor does OCaml. In this paper we solve the problem. Our contributions are these:

- We characterise the challenges of generating accurate warnings (Section 2). The problem goes beyond GADTs! There are subtle issues concerning how to view patterns, guards, and laziness; the latter at least has never been treated before.
- We give the first type-aware algorithm for determining missing or redundant patterns (Sections 3 and 4). The algorithm is parameterised over an oracle that can compute both type constraints and type annotations for variables. Existing type checkers are an oracle for this purpose. Extending the oracle allows us to accommodate type system extensions or improve the precision of the reported warnings without affecting the oracle.

The central abstraction in this algorithm is the compact symbolic representation of a set of values by a triple $(\Gamma \vdash u \triangleright \Delta)$ consisting of an environment Γ , a syntactic value abstraction u and a set of constraints Δ . It is key to note that we do not include the constraints Δ to refine the set of values; for example $(x:\text{Int} \vdash \text{Just } x \triangleright x>0)$ is the set of all applications of `Just` to values of type `x` that are greater than zero.

No, it is not: the call `(zip [] True)` will fail, because neither equation matches the call. Good compilers will report missing pattern warnings for such cases, but they will not warn about partially-defined. They will also warn about completely-overlapped, and hence redundant, equations. Although technically optional for soundness, these features are an important part of Haskell, especially when the program is refactored (i.e. throughout its active life) (Section 2).

But what about this function?

```
zip :: Vect n a -> Vect n b -> Vect n (a,b)
zip [] [] = []
zip (Vect x xs) (Vect y ys) = Vect (zip x y) (vzip xs ys)
```

where the type `Vect n a` represents lists of length n with elements of type `a`. `Vect` is a Generalised Algebraic Data Type (GADT):

```
data Vect :: Nat -> * -> * where
  Nil :: Vect Zero a
  (:>) :: a -> Vect (Succ n) a
```

Unlike `zip`, function `vzip` is fully defined: a call with arguments of unequal length, such as `(vzip Nil (Vect True Vect))`, is simply ill-typed. Comparing `zip` and `vzip`, it should be clear that only a type system that can reason about pattern guards can tell if the pattern-matches of a function definition are exhaustive.

Despite the runaway popularity of GADTs, and other pattern-matching features such as view patterns and pattern guards, no production compiler known to us gives accurate pattern-match overlap and exhaustiveness warnings when these features are used. Certainly our own compiler, GHC, does not, and nor does OCaml. In this paper we solve the problem. Our contributions are these:

- We characterise the challenges of generating accurate warnings (Section 2). The problem goes beyond GADTs! There are subtle issues concerning how to view patterns, guards, and laziness; the latter at least has never been treated before.
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- We characterise the challenges of generating accurate warnings (Section 2).
- The challenge of laziness (Section 2.3).
- The challenge of guards (Section 2.4).

To our knowledge, none of these difficulties are systematically addressed in the literature, let alone in a real compiler. In this paper we show how to tackle all three, in a single unified framework.

The End