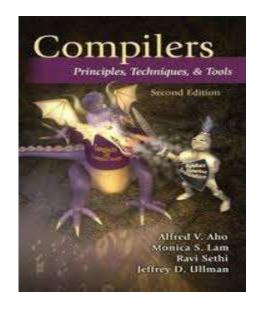
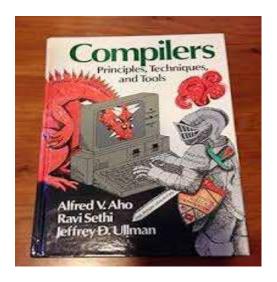


# Addis Ababa Institute of Technology - AAiT School of Information Technology and Engineering - SiTE

### **Compiler Design**

Henock Mulugeta (PhD)





# <u>Outline</u>

- Introduction
- Programs related to compiler
- ☐ The translation process
  - Analysis
    - Lexical analysis
    - Syntax analysis
    - Semantic analysis

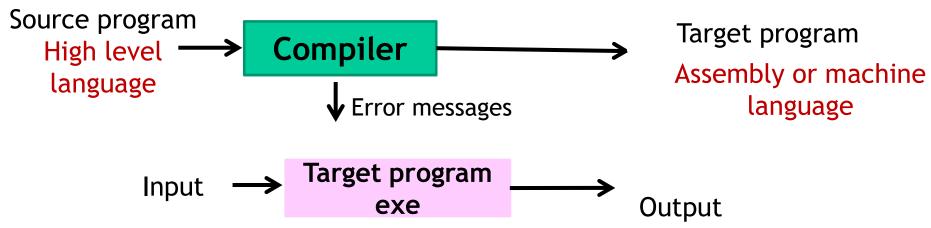
#### Synthesis

- IC generator
- IC optimizer
- Code generator
- Code optimizer
- Major data and structures in a compiler
- Types of compilers
- Compiler construction tools

# **Introduction**

#### What is a compiler?

- □ a program that reads a program written in one language (the source language) and translates it into an equivalent program in another language (the target language).
- Why we design compiler?
- Why we study compiler construction techniques?
  - Compilers provide an essential interface between applications and architectures
  - Compilers embody a wide range of theoretical techniques



### Introduction...

Using a high-level language for programming has a large impact on how fast programs can be developed.

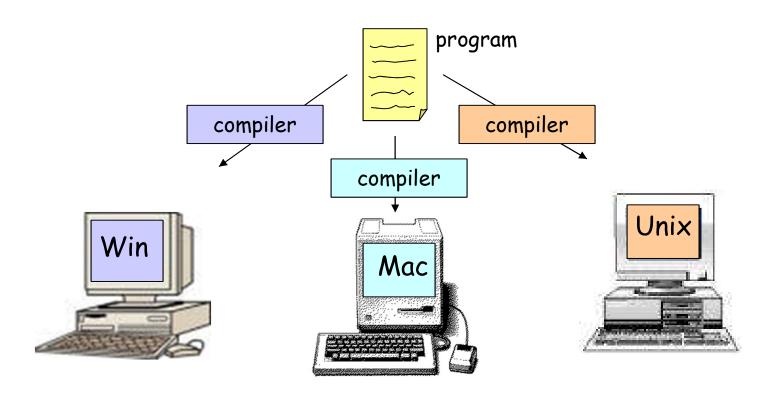
#### The main reasons for this are:

- Compared to machine language, the notation used by programming languages is closer to the way humans think about problems.
- The compiler can spot some obvious programming mistakes.
- O Programs written in a high-level language tend to be shorter than equivalent programs written in machine language.
- The same program can be compiled to many different machine languages
  - and, hence, be brought to run on many different machines.



### Introduction...

□ Since different platforms, or hardware architectures along with the operating systems (Windows, Macs, Unix), require different machine code, you must compile most programs separately for each platform.



## Programs related to compilers

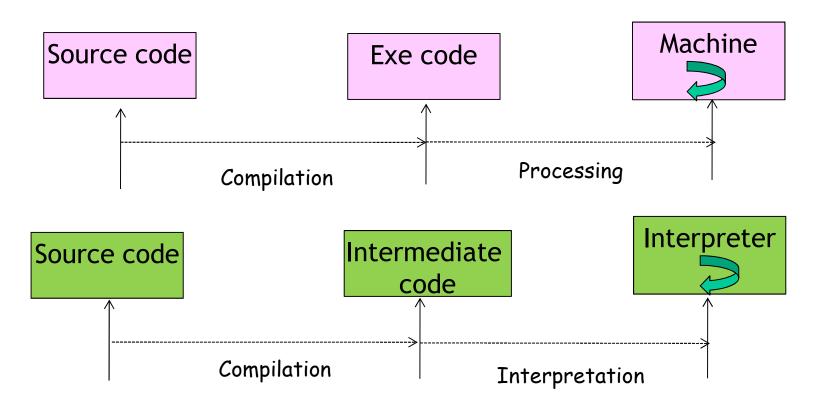
#### Interpreter

- Is a program that reads a source program and executes it
- Works by analyzing and executing the source program commands one at a time
- Does not translate the whole source program into object code

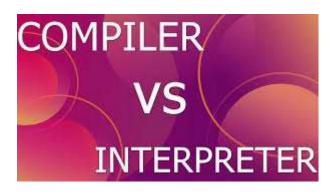
#### Interpretation is important when:

- Programmer is working in interactive mode and needs to view and update variables
- Running speed is not important
- Commands have simple formats, and thus can be quickly analyzed and executed
- Modification or addition to user programs is required as execution proceeds

# Programs related to compilers...



Interpreter and compiler



#### Interpreter and compiler differences



- Interpreter takes one statement then translates it and executes it and then takes another statement.
- Interpreter will stop the translation after it gets the first error.
- Interpreter takes less time to analyze the source code.

Over all execution speed is less.

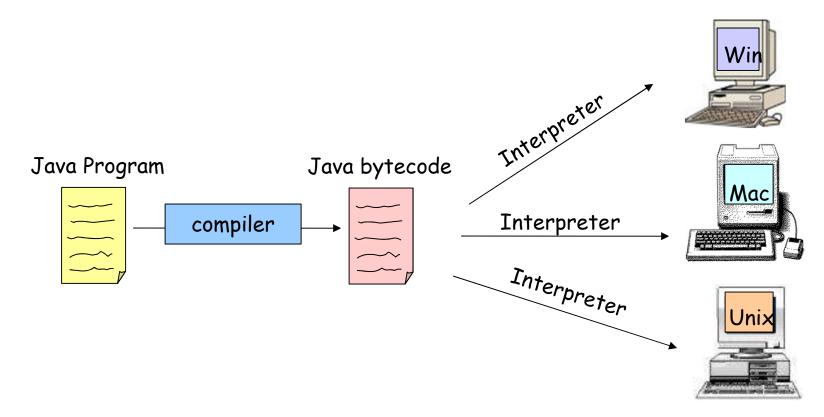
- While compiler translates the entire program in one go and then executes it.
- Compiler generates the error report after the translation of the entire program.
- Compiler takes a large amount of time in analyzing and processing the high level language code.
- Overall execution time is faster.

### Programs related to compilers...

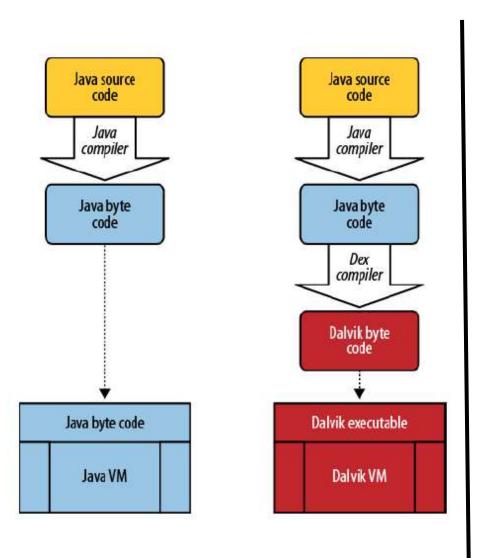
- Interpreter...
- Well-known examples of interpreters:
  - Basic interpreter, Lisp interpreter, UNIX shell command interpreter, SQL interpreter, java interpreter...
- In principle, any programming language can be either interpreted or compiled:
  - Some languages are designed to be interpreted, others are designed to be compiled
- Interpreters involve large overheads:
  - Execution speed degradation can vary from 10:1 to 100:1
  - Substantial space overhead may be involved

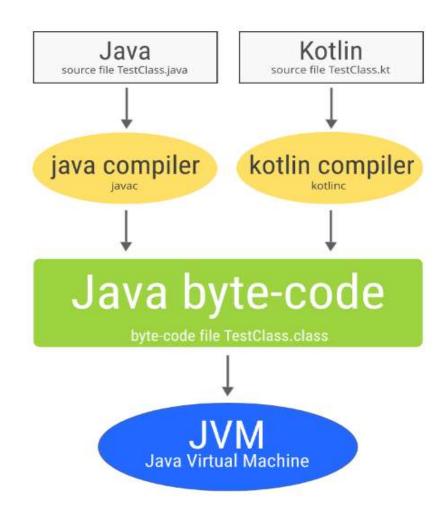
## E.g., Compiling Java Programs

- The Java compiler produces bytecode not machine code
- Bytecode is converted into machine code using a Java Interpreter
- You can run bytecode on any computer that has a Java Interpreter installed



# Android: Java and Kotlin





# Programs related to compiler...

#### Assemblers

- Translator for the assembly language.
- Assembly code is translated into machine code
- Output is relocatable machine code.

#### Linker

- Links object files separately compiled or assembled
- Links object files to standard library functions
- Generates a file that can be loaded and executed

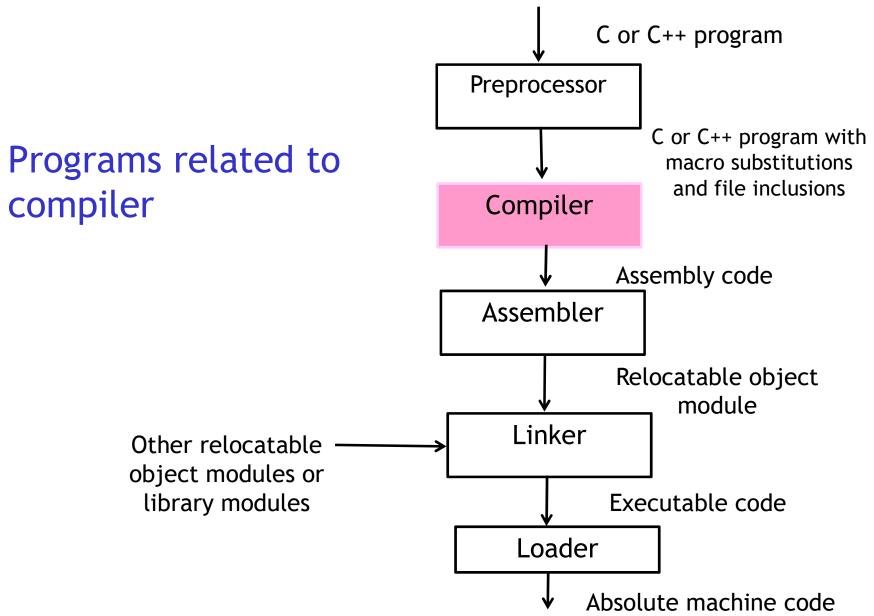
### Programs related to compiler...

#### Loader

 Loading of the executable codes, which are the outputs of linker, into main memory.

#### Pre-processors

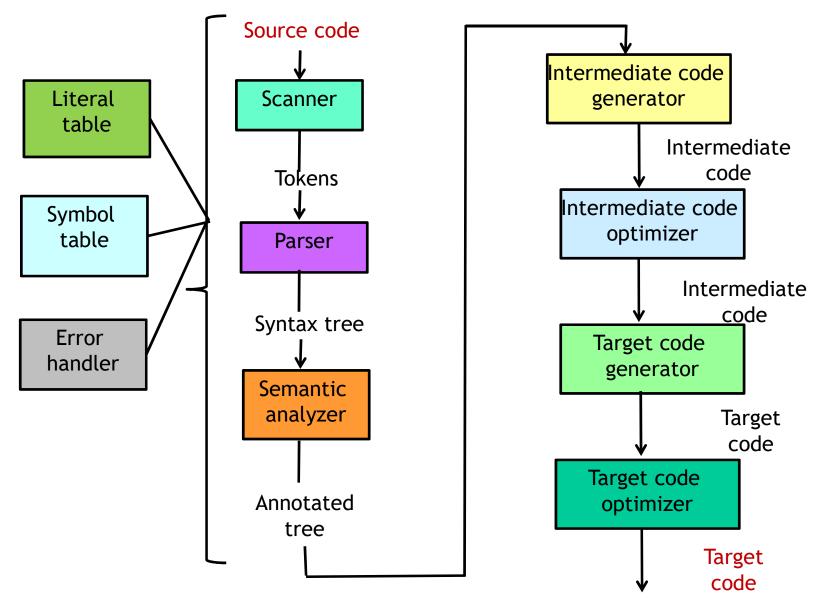
- A pre-processor is a separate program that is called by the compiler before actual translation begins.
- Such a pre-processor:
  - Produce input to a compiler
  - can delete comments,
  - Macro processing (substitutions)
  - include other files...



# The translation process

- A compiler consists of internally of a number of steps, or phases, that perform distinct logical operations.
- □ The phases of a compiler are shown in the next slide, together with three auxiliary components that interact with some or all of the phases:
  - The symbol table,
  - the literal table,
  - and error handler.
- There are two important parts in compilation process:
  - Analysis and
  - Synthesis.

# The translation process...



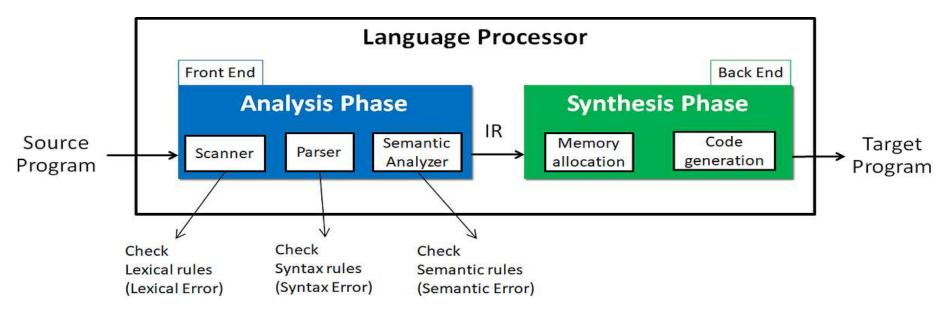
# **Analysis and Synthesis**

#### Analysis (front end)

- Breaks up the source program into constituent pieces and
- Creates an intermediate representation of the source program.
- During analysis, the operations implied by the source program are determined and recorded in hierarchical structure called a tree.

#### Synthesis (back end)

 The synthesis part constructs the desired program from the intermediate representation.



# Analysis of the source program

- Analysis consists of three phases:
  - Linear/Lexical analysis
  - Hierarchical/Syntax analysis
  - Semantic analysis

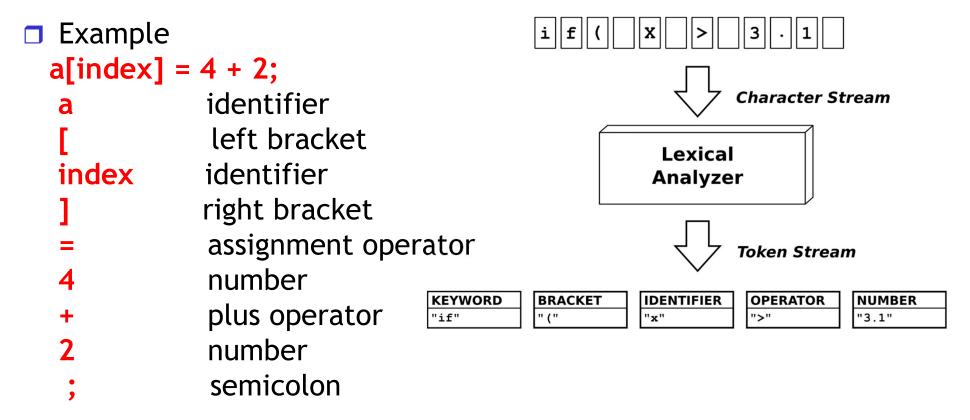
# 1. Lexical analysis or Scanning

- A lexical analyzer, also called a lexer or a scanner,
  - receives a stream of characters from the source program and
  - groups them into tokens.
- A token is a sequence of characters having a collective meaning.
- Examples:
  - Identifiers



- Keywords
- Symbols (+, -, ...)
- Numbers ...
- Blanks, new lines, tabulation marks will be removed during lexical analysis.

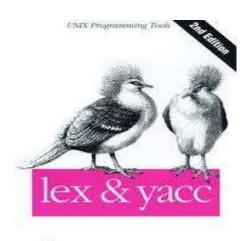
# Lexical analysis or Scanning...



- A scanner may perform other operations along with the recognition of tokens.
  - It may inter identifiers into the symbol table, and
  - It may inter literals into literal table.

## Lexical Analysis Tools

- □ There are tools available to assist in the designing of lexical analyzers.
  - lex produces C source code (UNIX/linux).
  - o flex produces C source code (gnu).
  - JfLex produces Java source code.







■ We will use Lex.





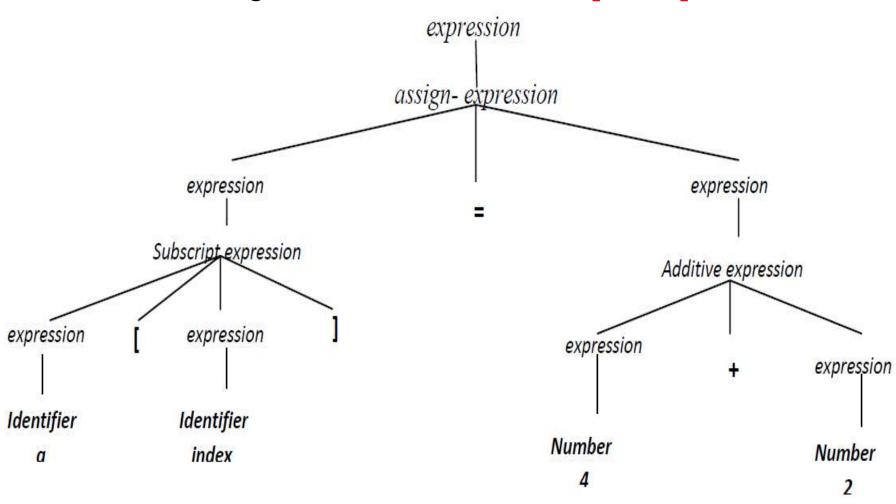
### 2. Syntax analysis or Parsing

- The parser receives the source code in the form of tokens from the scanner and performs syntax analysis.
- The results of syntax analysis are usually represented by a parse tree or a syntax tree.
  - Syntax tree → each interior node represents an operation and the children of the node represent the arguments of the operation.
- The syntactic structure of a programming language is determined by context free grammar (CFG).



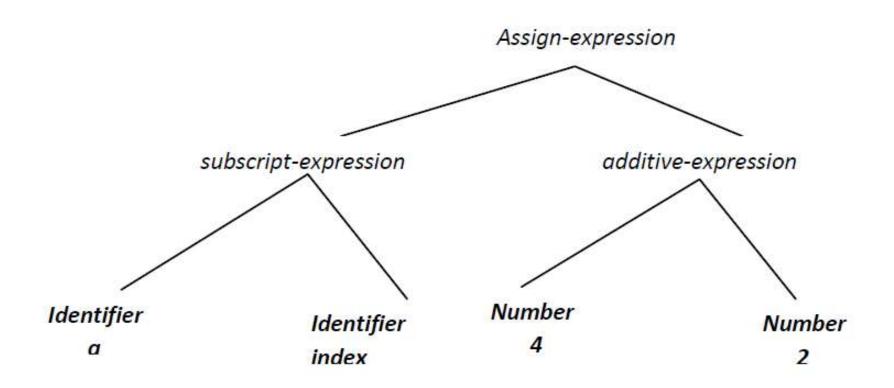
### Syntax analysis or Parsing...

Ex. Consider again the line of C code: a[index] = 4 + 2



### Syntax analysis or Parsing...

□ Sometimes syntax trees are called **abstract syntax trees**, since they represent a further abstraction from parse trees. Example is shown in the following figure.



# Syntax Analysis Tools

- There are tools available to assist in the writing of parsers.
  - yacc produces C source code (UNIX/Linux).
  - bison produces C source code (gnu).
  - CUP produces Java source code.
- ☐ We will use yacc.





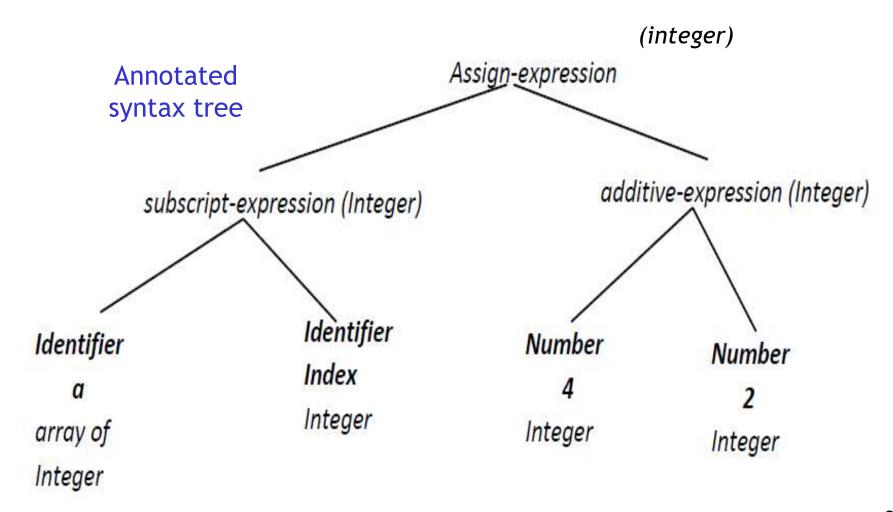


# 3. Semantic analysis

- The semantics of a program are its meaning as opposed to syntax or structure
- ☐ The semantics consist of:
  - Runtime semantics behavior of program at runtime
  - Static semantics checked by the compiler
- Static semantics include:
  - Declarations of variables and constants before use
  - Calling functions that exist (predefined in a library or defined by the user)
  - Passing parameters properly
  - Type checking.
- The semantic analyzer does the following:
  - Checks the static semantics of the language
  - Annotates the syntax tree with type information

# Semantic analysis...

Ex. Consider again the line of C code: a[index] = 4 + 2



# Synthesis of the target program

- Intermediate code generator
- Intermediate code optimizer
- ☐ The target code generator
- ☐ The target code optimizer

# Code Improvement

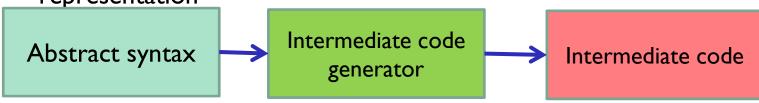
- Code improvement techniques can be applied to:
  - Intermediate code independent of the target machine
  - Target code dependent on the target machine
- □ Intermediate code improvement include:
  - Constant folding
  - Elimination of common sub-expressions



- Improving loops
- Improving function calls
- □ Target code improvement include:
  - Allocation and use of registers
  - Selection of better (faster) instructions and addressing modes

# Intermediate code generator

- Comes after syntax and semantic analysis
- Separates the compiler front end from its backend
- Intermediate representation should have 2 important properties:
  - Should be easy to produce
  - Should be easy to translate into the target program
- Intermediate representation can have a variety of forms:
  - Three-address code, P-code for an abstract machine, Tree or DAG representation



☐ Three address code for the original C expression a[index]=4+2 is:

$$t_1=2$$
  
 $t_2 = 4 + t_1$   
 $a[index] = t_2$ 

# IC optimizer

- ☐ An IC optimizer reviews the code, looking for ways to reduce:
  - the number of operations and
  - the memory requirements.
- A program may be optimized for speed or for size.
- This phase changes the IC so that the code generator produces a faster and less memory consuming program.
- The optimized code does the same thing as the original (non-optimized) code but with
  - less cost in terms of CPU time and memory space.



# IC optimizer...

- ☐ There are several techniques of optimizing code and they will be discussed in the forthcoming chapters.
- Ex. Unnecessary lines of code in loops (i.e. code that could be executed outside of the loop) are moved out of the loop.

```
for (i=1; i<10; i++) {
    x = y+1;
    z = x+i; }

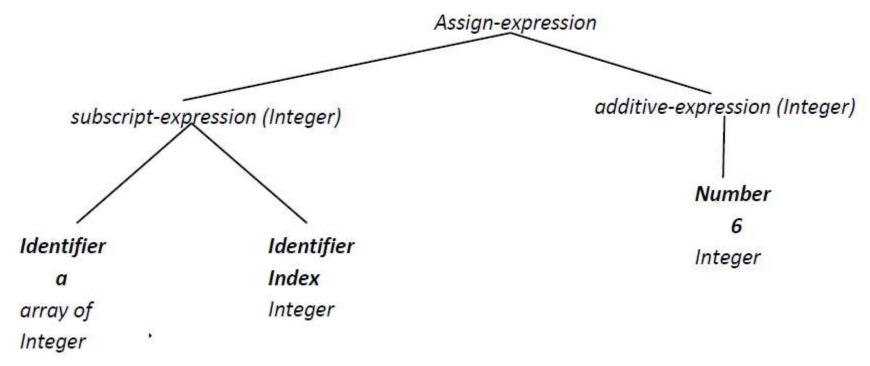
x = y+1;
for (i=1; i<10, i++)
    z = x+i;</pre>
```

## IC optimizer...

In our previous example, we have included an opportunity for source level optimization; namely, the expression 4 + 2 can be recomputed by the compiler to the result 6(This particular optimization is called **constant folding**).



This optimization can be performed directly on the syntax tree as shown below.



# IC optimizer...

- Many optimizations can be performed directly on the tree.
- However, in a number of cases, it is easier to optimize a linearized form of the tree that is closer to assembly code.
- A standard choice is **Three-address code**, so called because it contains the addresses of up to **three locations in memory**.
- In our example, three address code for the original C expression might look like this:

$$t_1=2$$
 $t_2 = 4 + t_1$ 
a[index] =  $t_2$ 

□ Now the optimizer would improve this code in two steps, first computing the result of the addition

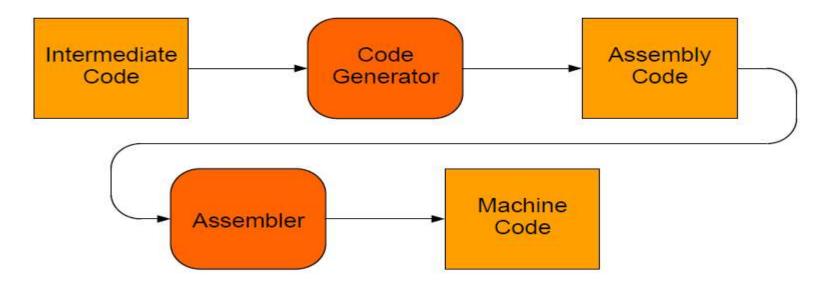
And then replacing t by its value to get the three-address statement a[index] = 6

### Code generator

- The machine code generator receives the (optimized) intermediate code, and then it produces either:
  - Machine code for a specific machine, or
  - Assembly code for a specific machine and assembler.

#### Code generator

- Selects appropriate machine instructions
- Allocates memory locations for variables
- Allocates registers for intermediate computations



# Code generator...

- The code generator takes the IR code and generates code for the target machine.
- ☐ Here we will write target code in assembly language: a[index]=6

```
MOV R0, index
;; value of index -> R0
MUL R0, 2
;; double value in R0
MOV R1, &a
;; address of a -> R1
ADD R1, R0
;; add R0 to R1
MOV *R1, 6
;; constant 6 -> address in R1
```

- ta -the address of a (the base address of the array)
- \*R1-indirect registers addressing (the last instruction stores the value 6 to the address contained in R1)

## The target code optimizer

- ☐ In this phase, the compiler attempts to improve the target code generated by the code generator.
- ☐ Such improvement includes:
  - Choosing addressing modes to improve performance
  - Replacing slow instruction by faster ones
  - Eliminating redundant or unnecessary operations
- □ In the sample target code given, use a shift instruction to replace the multiplication in the second instruction.
- ☐ Another is to use a more powerful addressing mode, such as indexed addressing to perform the array store.
- With these two optimizations, our target code becomes:

```
MOV R0, index ;; value of index -> R0

SHL R0 ;; double value in R0

MOV &a [R0], 6 ;; constant 6 -> address a + R0
```

## Major Data and Structures in a Compiler

#### Token

- Represented by an integer value or an enumeration literal
- Sometimes, it is necessary to preserve the string of characters that was scanned
- For example, name of an identifiers or value of a literal

### Syntax Tree

- Constructed as a pointer-based structure
- Dynamically allocated as parsing proceeds
- Nodes have fields containing information collected by the parser and semantic analyzer

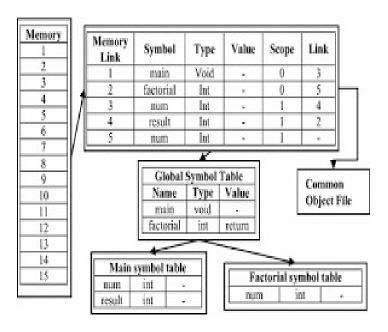
## Major Data and Structures in a Compiler...

#### Symbol Table

- Keeps information associated with all kinds of tokens:
  - Identifiers, numbers, variables, functions, parameters, types, fields, etc.
- Tokens are entered by the scanner and parser
- Semantic analyzer adds type information and other attributes
- Code generation and optimization phases use the information in the symbol table

#### Performance Issues

- Insertion, deletion, and search operations need to be efficient because they are frequent
- More than one symbol table may be used



## Major Data and Structures in a Compiler...

#### Literal Table

- Stores constant values and string literals in a program.
- One literal table applies globally to the entire program.
- Used by the code generator to:
  - Assign addresses for literals.
- Avoids the replication of constants and strings.
- Quick insertion and lookup are essential.

#### Literal Table

Name of the Literal '4'	Value of the Literal 4	Address of Usage of Symbol 0003,0007	Address of Defination of Symbol 0014 0015
-------------------------------	------------------------------	---	---

## Types of Compilers

- ☐ Based on the **number of passes**:
- □ Single-pass compilers: These compilers read the source code and generate the target code in a single pass.
  - They are generally faster than multi-pass compilers but can be less efficient in terms of code quality. Pascal
- Two-pass compilers: These compilers read the source code twice.
  - The first pass gathers information about the code, and the second pass generates the target code.
  - They can produce more efficient code than single-pass compilers but are slower. C, C++, and Java
- Multi-pass compilers: These compilers read the source code multiple times.
  - They can produce the most efficient code but are also the slowest. Ada

## Types of Compilers...

- Based on the target platform:
- Cross compilers: These compilers run on one platform but generate code for a different platform.
  - They are used to develop software for embedded systems or other platforms where a compiler is not available.
  - SDCC (Small Device C Compiler)
- Native compilers: These compilers run or a platform and generate code for the same platform.
  - They are the most common type of compiler.
  - C# program on Windows using the .NET framework's



## Types of Compilers...

- Based on the type of code generated
- □ Just-in-time (JIT) compilers: These compilers compile code at runtime, rather than ahead of time.
  - They are used in Java and other languages where code is often downloaded from the internet.
  - Java, JavaScript, .NET languages (C#, VB.NET, F#), Python, Ruby, PHP
- □ Ahead-of-Time (AOT) Compiler: It converts the entire source code into machine code before the program runs.
  - C and C++, Fortran, Go, Rust, Swift, Java (with GraalVM), .NET (with .NET Native)
- Hybrid Approaches: Modern runtime environments often use a combination of AOT and JIT techniques to optimize performance.

## Other types of compilers

- □ Source-to-source compilers: These compilers translate code from one high-level language to another high-level language.
  - Haxe (pronounced "hex") is an open-source high-level multiplatform programming language and compiler that can produce code for many different target platforms from a single code-base.
  - O Haxe can be compiler to C++, C#, Java, Lua, Python
- □ Decompilers: These programs attempt to reverse the compilation process, converting machine code back into a higher-level language.
  - O Java (e.g., JAD) and .NET (e.g., .NET Reflector)

# Key Considerations When Choosing a Compiler

- □ Compilation speed: How quickly can the compiler translate source code into executable code?
- ☐ Code efficiency: How well does the compiler optimize the generated code for performance?
- ☐ Error detection: How effectively does the compiler identify and report errors in the source code?
- □ Portability: Can the compiler generate code for multiple platforms?
- ☐ Features: Does the compiler offer advanced features, such as debugging support or code optimization tools?

### Popular Compiler Examples by Language

- □ C/C++: GCC, G++ (GNU Compiler Collection), Clang/LLVM, Microsoft Visual C++ Compiler
- □ Java: OpenJDK, Oracle JDK, IBM SDK
- C#: Microsoft Visual Basic Compiler (VBC)
- □ Python: While Python is primarily interpreted, there are also compilers like MyPy (for type checking) and Cython (for performance optimization)
- ☐ Go: Go Compiler (gc)
- JavaScript: V8 (used in Chrome and Node.js), SpiderMonkey (used in Firefox)

## Compiler construction tools

□ Various tools are used in the construction of the various parts of a compiler.

#### Scanner generators

- Ex. Lex, flex, JLex
- These tools generate a scanner /lexical analyzer/ if given a regular expression.

#### Parser Generators

- Ex. Yacc, Bison, CUP
- These tools produce a parser /syntax analyzer/ if given a Context Free Grammar (CFG) that describes the syntax of the source language.

## Compiler construction tools...

#### Syntax directed translation engines

- O Ex. Cornell Synthesizer Generator
- It produces a collection of routines that walk the parse tree and execute some tasks.

#### Automatic code generators

 Take a collection of rules that define the translation of the IC to target code and produce a code generator.

#### ■ Data-Flow Analysis Engines

 It supports code optimization by analyzing the flow of values throughout different parts of the program.

#### □ Compiler Construction Toolkits:

 It provides integrated routines to facilitate the construction of various compiler components.

## This completes our brief description of the compiler Design.

Thank u